#### Quantum Error Correction

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> Physics 160 Harvard University

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#### Introduction

"To be an Error and to be Cast out is part of God's Design."

William Blake

- Noise as a longstanding problem in information processing systems
  - e.g., classical computers, modems, CD players, etc.
  - Noise is still a problem in quantum information
- Key idea: to protect a message against noise, encode the message by adding redundant information; even if some information is corrupted, redundancy allows us to decode and recover the original message

## Project Framework

- Goals:
  - to implement various quantum error-correcting codes
    - we chose the 3-qubit, 9-qubit, 7-qubit codes
  - to analyze and compare their performances
    - when are they effective?
    - when should we use error-correcting codes?
- Tools:
  - Python's Qiskit package
  - IBM's quantum machines

## 3-Qubit Codes: Classical Inspiration

#### **Classical Error Correction**

Encoding by repetition codes:

$$0 \rightarrow 000$$
  
 $1 \rightarrow 111$ .

• Decoding by majority voting:

*Ex.:* 
$$001 \rightarrow 0$$
.

• Analysis: Let p be the probability that a bit is flipped. This method fails when 2 or more bits are flipped, which occurs with probability  $3p^2(1-p)+p^3$ , so the probability of error is  $p_e=3p^2-2p^3$ . Then this method is preferred when  $p_e < p$ , or p < 1/2.

## Noisy Channels: The Bit Flip Channel

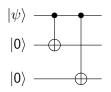
- One model for noise is the *bit flip channel* (analogous to classical channel).
- The bit flip channel flips qubits with probability p and leaves them untouched with probability 1 p.
- Equivalent to applying X gate with probability p.
- We protect qubits from this channel with the bit flip code.

# 3-Qubit Bit Flip Code: Encoding Logical Bits

- The goal is to correct bit flip errors.
- Encoding:

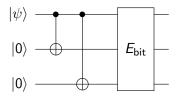
$$|0\rangle \rightarrow |0_L\rangle \equiv |000\rangle$$
  
 $|1\rangle \rightarrow |1_L\rangle \equiv |111\rangle$ .

• Encoding circuit for 3-qubit bit flip code:



#### 3-Qubit Bit Flip Code: Detecting Errors

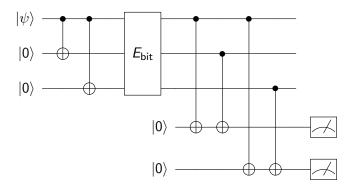
Suppose there is a bit flip error after encoding:



- Error Detection (or syndrome diagnosis):
  - we would like to determine which, if any, of the qubits have been corrupted
  - four error syndromes: no error, bit flip on qubit one, bit flip on qubit two, bit flip on qubit three

#### 3-Qubit Bit Flip Code: Detecting Errors

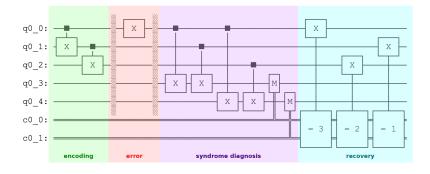
• We can diagnose the syndrome using two ancillary qubits:



• Based on measurement results, we know where the error occured.

# 3-Qubit Bit Flip Code: Correcting Errors

• Complete circuit for error correction (or *recovery*):



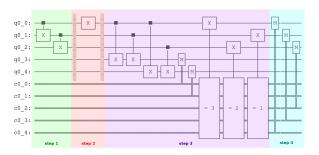
#### Analyzing the Bit Flip Code: Simulation

- Let's look at the performance of the 3-qubit bit flip code against bit flip channels of varying error probabilities *p*.
- Setup:
  - **1** encode a single qubit in state  $|0\rangle$  into a logical state  $|0_L\rangle = |000\rangle$

  - o run error correcting code
  - measure final state

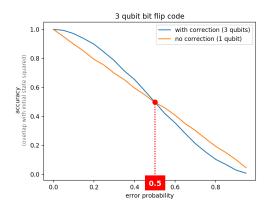
#### Analyzing the Bit Flip Code: Simulation

- Let's look at the performance of the 3-qubit bit flip code against bit flip channels of varying error probabilities *p*.
- Setup:
  - **1** encode a single qubit in state  $|0\rangle$  into a logical state  $|0_L\rangle = |000\rangle$
  - 2 create a bit flip channel which adds X gates with probability p
  - o run error correcting code
  - measure final state



#### Analyzing the Bit Flip Code: Simulation

- Ran tests on Qiskit's simulator
- Probability p ranging from 0 to 1; 10000 trials for each p



- Observe crossover point at p = 0.5.
- For p < 0.5, error correcting code performs better than a single qubit with no correction.

## Noisy Channels: Phase Flip Channel

- Another quantum channel is the *phase flip* error model.
- With probability p the relative phase of states  $|0\rangle$  and  $|1\rangle$  is flipped, with probability 1-p it is left alone.
- Equivalent to applying Z operator with probability p.
- We fight this channel with the *phase flip code*.

## 3-Qubit Phase Flip Code

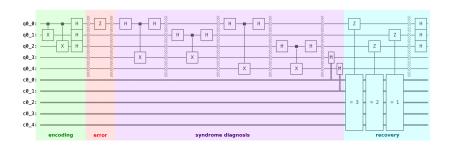
- No classical analog, but it is easy to turn the phase flip channel into a bit flip channel.
- Use x-basis for encoding:

$$\begin{split} |0\rangle &\rightarrow |0_L\rangle \equiv |+++\rangle \\ |1\rangle &\rightarrow |1_L\rangle \equiv |---\rangle \,. \end{split}$$

• Phase flip Z acts as bit flip for this encoding!

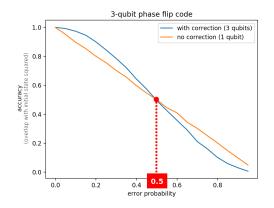
# 3-Qubit Phase Flip Code

#### Complete 3-qubit phase flip code:



#### Analyzing the Phase Flip Code: Simulation

- Ran tests on Qiskit's simulator
- Probability p ranging from 0 to 1; 10000 trials for each p



- Observe crossover point at p = 0.5.
- For p < 0.5, error correcting code performs better than a single qubit with no correction.
- Nearly identical result to the bit flip code.

#### 3-Qubit Codes on IBM's Machines

Adjusting the circuit for IBM's machines (which do not have feed-forward capability)

#### The Shor 9-Qubit Code

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- Yes! → The Shor code

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• Then encode each of those qubits with the bit flip code:

$$|+\rangle \rightarrow (|000\rangle + |111\rangle)/\sqrt{2}, \quad |-\rangle \rightarrow (|000\rangle - |111\rangle)/\sqrt{2}.$$

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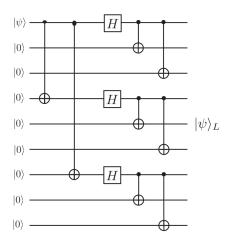
$$|+\rangle \rightarrow (|000\rangle + |111\rangle)/\sqrt{2}, \quad |-\rangle \rightarrow (|000\rangle - |111\rangle)/\sqrt{2}.$$

• The result is a 9-qubit code with codewords

$$\begin{split} |0\rangle \rightarrow |0_L\rangle &\equiv \frac{(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)}{2\sqrt{2}} \\ |1\rangle \rightarrow |1_L\rangle &\equiv \frac{(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)(|000\rangle - |111\rangle)}{2\sqrt{2}}. \end{split}$$

# The Shor 9-Qubit Code: Encoding

#### Encoding circuit for 9-qubit code:



# The Shor 9-Qubit Code: Correcting Errors

#### Bit Flip Error Correction

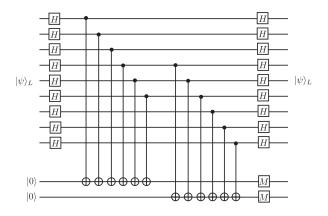
• On each block of three (i.e. qubits 0-2, 3-5, and 6-8), the 3-qubit circuit is utilized to correct for bit flips.

#### **Phase Flip Error Correction**

- The phase of the first two blocks of three (qubits 0-2 and 3-5) and the second two blocks of three (qubits 3-5 and 6-8) are compared to correct for phase flips.
- The phase correction necessitates two ancillary qubits. Thus, we need 8 ancilla: 6 for bit flip correction, and 2 for phase flip correction.

# The Shor 9-Qubit Code: Correcting Phase Errors

• The phase correction circuit, shown below, converts the qubits from the x-basis to the z-basis and checks parity of each block of two.



# The Shor 9-Qubit Code: Correcting Phase Errors

 The following corrections are performed depending on the measured ancilla for phase flip correction:

$$10 
ightarrow \sigma_z$$
 on block  $1$ 

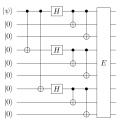
$$01 \rightarrow \sigma_z$$
 on block 3

$$11 \rightarrow \sigma_z$$
 on block 2.

• As two bit-flip errors occur with probability  $3p^2$  per block for 3 blocks, and two phase-flip errors occur on separate blocks with probability  $3p^2$ , the total error  $\sim 12p^2$ . Thus, this error code is efficient for  $12p^2 .$ 

# The Shor 9-Qubit Code: Error Correction Methodology

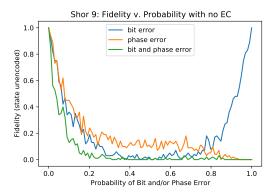
 We only consider error that occurs between the encoding step and the correcting step, thus simulating a memory error.



- Specifically, we consider a complete phase flip and/or bit flip (i.e. X or Z) that occurs independently on each of the 9 physical qubits with probability p.
- After the error, we measure the ancilla and apply the appropriate error correcting steps. Finally, we run the encoding circuit in reverse and measure the output to determine fidelity.

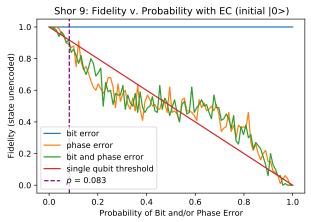
# The Shor 9-Qubit Code: Simulation Performance with No Error Correction

- Initial state:
  - $|0\rangle \rightarrow |0_L\rangle \equiv \frac{1}{\sqrt{8}}(|000\rangle + |111\rangle)(|000\rangle + |111\rangle)(|000\rangle + |111\rangle).$
- Fidelity of un-encoded state measured against |000000000).



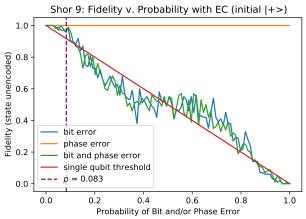
# The Shor 9-Qubit Code: Simulation Performance with Error Correction

- Initial state:  $|0\rangle \rightarrow |0_L\rangle$
- Fidelity of un-encoded state measured against |000000000).



# The Shor 9-Qubit Code: Simulation Performance with Error Correction

- Initial state:  $|+\rangle \rightarrow \frac{1}{\sqrt{2}}(|0_L\rangle + |1_L\rangle)$
- Fidelity of un-encoded state measured against |000000000).



#### 7-Qubit Code

#### Encodes 1 logical qubit using 7 physical qubits:

$$\begin{split} |0_L\rangle &= \frac{|0000000\rangle + |1010101\rangle + |0110011\rangle + |1100110\rangle + |0001111\rangle + |1011010\rangle + |0111100\rangle + |1101001\rangle}{\sqrt{8}} \\ |1_L\rangle &= \frac{|1111111\rangle + |0101010\rangle + |1001100\rangle + |0011001\rangle + |1110000\rangle + |0100101\rangle + |1000011\rangle + |0010110\rangle}{\sqrt{8}} \end{split}$$

$$\begin{split} H^{\otimes 7} \left| \mathbf{0}_L \right\rangle &= \frac{\left| \mathbf{0}_L \right\rangle + \left| \mathbf{1}_L \right\rangle}{\sqrt{2}} \\ H^{\otimes 7} \left| \mathbf{1}_L \right\rangle &= \frac{\left| \mathbf{0}_L \right\rangle - \left| \mathbf{1}_L \right\rangle}{\sqrt{2}} \end{split}$$

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$$H^{\otimes 7} |0_L\rangle = \frac{|0_L\rangle + |1_L\rangle}{\sqrt{2}}$$

$$H^{\otimes 7} |1_L\rangle = \frac{|0_L\rangle - |1_L\rangle}{\sqrt{2}}$$

- Of the 16 bit strings above, any two differ by  $\geq$  3 bits
- Intuition: therefore a single bit flip can be recovered
  - X error flips bit in  $|0_L\rangle$ ,  $|1_L\rangle$
  - ullet Z error flips bit in  $H^{\otimes 7}\ket{0_L}, H^{\otimes 7}\ket{1_L}$

## Example recovery for X error in qubit 3

$$|0_L\rangle = \frac{|0000000\rangle + |1010101\rangle + |0110011\rangle + |1100110\rangle + |0001111\rangle + |1011010\rangle + |0111100\rangle + |1110101\rangle}{\sqrt{8}} \\ X^{(3)} |0_L\rangle = \frac{|0010000\rangle + |1000101\rangle + |0100011\rangle + |1110110\rangle + |0011111\rangle + |1001010\rangle + |0101100\rangle + |1111001\rangle}{\sqrt{8}}$$

$$\begin{pmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 \end{pmatrix} \begin{pmatrix} \\ \\ \\ \\ \\ \end{pmatrix} = \begin{pmatrix} 1 \\ \\ 1 \\ \\ \end{pmatrix} = |3\rangle \text{ (in binary)}$$

In fact:

$$\begin{pmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 \end{pmatrix} X^{(i)} \ket{0_L} = \ket{i} \text{ (in binary) for all } i = 1, \dots, 7$$

Also works for logical state 1, and for phase flips.



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In fact:

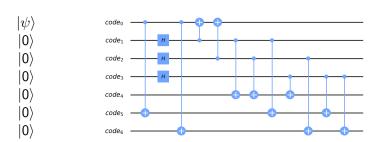
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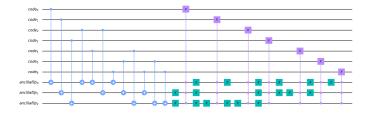
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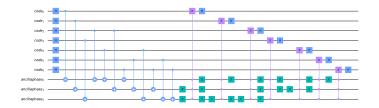


 $|\psi_{\mathsf{L}}\rangle$ 

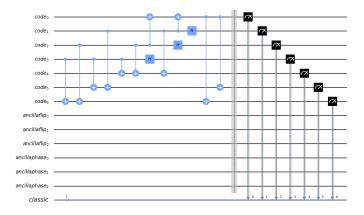
#### 7-qubit code: Flip correction



### 7-qubit code: Phase correction

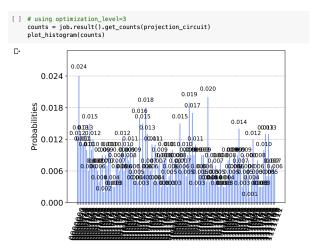


## 7-qubit code: Decoding and measurement



# 7-qubit code: Simulation vs running on quantum computers?

- The states should be clearly defined, but noise dominates the system



## 7-qubit code: Fidelity of X Gate under Depolarization

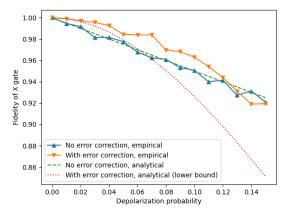
Encode  $|0_L\rangle$ 

Χ

Correct flip

Correct phase

Decode, measure



# 7-qubit code: Fidelity of X Gate under Depolarization

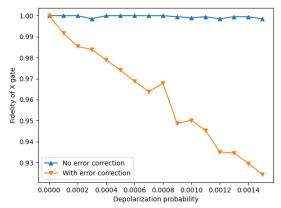
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- Example:  $\langle 0_L | X^{(i)} | 0_L \rangle = 0$ 
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- Example:  $\langle 0_L | X^{(i)} | 0_L \rangle = 0$ 
  - Single bit flip to  $|0_L\rangle$  gives fidelity 0, but can be corrected!
- 7-qubit code reduces single gate error rate for depolarization rates lower than detectable by our qiskit simulations ( $p \lesssim 10^{-4}$ )

## 7-qubit code: Extending to Multiple Gates

• Here, we assumed that the  $|0\rangle_L$  state was perfectly encoded, before performing up to 200 pairs of X gates on the seven qubits simultaneously

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- We compared the fidelity of the error correcting algorithm to a single qubit that goes through up to 200 X rotations

## 7-qubit code: Useful with lower error probabilities

Encode  $|\overline{0}\rangle$ 

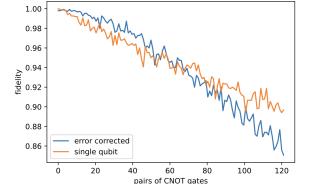


Correct flip

Correct phase

Decode, measure

Comparison of fidelities with and without Steane Code, X-Fidelity=0.999



# 7-qubit code: Adding Error Correction at different Timesteps

- Using a noise model with (p=0.001) noise in the  $\times$  and z-direction per gate

Encode  $|\overline{0}\rangle$ 

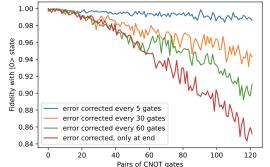


Correct flip

Correct phase

Decode, measure

Comparison of fidelities with varying error correction schedules, X-Fidelity=0.999



# 7-qubit code: Adding Error Correction at different Timesteps

- Using a noise model with (p=0.01) noise in the  $\times$  and z-direction per gate

Comparison of fidelities with varying error correction schedules. Gate-Fidelity=0.99 1.0 0.9 idelity with |0> state 0.8 corrected every 5 gates corrected every 30 gates corrected every 60 gates single gubit (control) 0.6 0.5 20 n 40 60 100 120

Pairs of CNOT gates

#### 7-qubit code: Generalizing to Error Correction

- Ultimately, error correction is useful, but must be intelligently deployed - here, error-correction is heavily sensitive to the noise in each gate
- It may be useful to characterize the thresholds before using error-correcting codes or consider using error-correcting codes continuously throughout a computation