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MODEL

杂项

最大化最小值(二分值最小越小越容易满足条件,求最大值)

```
区间长度为1时的写法:
解的范围为
// 计算区间为[lb, rb]
while( rb > lb ) // 区间长度为1 时终止循环
   // 防止溢出
   int m = lb + (rb - lb + 1) / 2; // 由于是区间长度为1 时终止循环,所以要加1
   if(ok(m))lb = m;
   else rb = m - 1;
// 跳出循环时 Lb == rb
区间长度为2时的写法:
解的范围为
while( rb - lb > 1 ) // 区间长度为2 时终止循环
   // 防止溢出
   int m = lb + (rb - lb) / 2; // 由于是区间长度为2 时终止循环,所以不用加1(不会死循环)
   if(ok(m)) lb = m;
   else rb = m;
// 跳出循环时 Lb + 1 == rb
// 答案为 Lb
最小化最大值(二分值越大越容易满足条件,求最小值)
区间长度为1时的写法:
解的范围为
while( rb > lb )
   // 防止溢出
   int m = 1b + (rb - 1b) / 2;
                          // 这里虽然区间长度为1,但不需要加1(不会死循环)
   if(ok(m)) rb = m;
   else 1b = m + 1;
// 跳出循环时 Lb == rb
区间长度为2时的写法:
解的范围为
while (rb - lb > 1)
   // 防止溢出
   int m = 1b + (rb - 1b) / 2;
   if(ok(m)) rb = m;
```

```
else lb = m:
// 跳出循环时 Lb + 1 == rb
// 答案为 rb
浮点数的二分,100次循环可以达到2-100(約为10-30)的精度范围
以最大化最小值为例(即小于该数的解均满足条件)
for( int i = 0; i < 100; ++ i )</pre>
    double m = (1b + rb) / 2;
   if(check(m)) lb=m;
    else rb=m;
// 跳出循环时 Lb 与 rb 近似相等,所以都可作为答案
莫队
//动态
struct unit {
        int 1, r, id;
  bool operator<(const unit &x) const {</pre>
   return 1 / sq == x.1 / sq
              ? (r == x.r ? 0 : ((1 / sq) & 1) ^ (r < x.r))
              : 1 < x.1;//莫队奇偶优化
};
int totq=0,totxiu=0,sq,ans[N],ans1=0,n,m,a[N],sum[N],l=1,r=0,now=0;
int main()
    scanf("%d%d",&n,&m);
    for (int i=1;i<=n;i++) scanf("%d",&a[i]);</pre>
    for (int i=1;i<=m;i++)</pre>
       char s[2];
       int x,y;
       scanf("%s%d%d",&s,&x,&y);
       if (s[0]=='0')
           totq++;
           q[totq].l=x; q[totq].r=y;
           q[totq].x=totxiu;
           q[totq].id=totq;
       else
           totxiu++;
           xiu[totxiu].id=x;
           xiu[totxiu].cl=y;
    sq=sqrt(n*1.0);
    sort(q+1,q+1+totq,cmp);
        for(int i=1;i<=totq;i++)</pre>
```

```
while (1>q[i].1)1--,ins(1);
        while(r<q[i].r)r++,ins(r);</pre>
                  while(1<q[i].1)del(1),1++;
                  while(r>q[i].r)del(r),r--;
                  while(now<q[i].x)change(now+1),now++;</pre>
                  while(now>q[i].x)change(now),now--;
                  ans[q[i].id]=ans1;
         for(int i=1;i<=totq;i++)</pre>
         printf("%d\n",ans[i]);
}
//基本
struct node {
    11 1,r,id,A,B;
}Q[maxn];
11 be[maxn],a[maxn],sq,num,l=1,r=0,sum[maxn];
ll ans1=0;
ll gcd(ll a,ll b)
 {
         if(b==0)return a;
         else return gcd(b,a%b);
bool cmp1(const node& A,const node&B) {return be[A.1]==be[B.1]?A.r<B.r:A.1<B.1;}</pre>
bool cmp2(const node& A,const node&B){return A.id<B.id;}</pre>
void ins(int x) {ans1=ans1+2*sum[x];sum[x]++;}
void del(int x) {ans1=ans1-2*sum[x]+2;sum[x]--;}
int main()
{
        scanf("%11d%11d",&n,&m);
        sq=sqrt(n);
        for(ll i=1;i<=n;i++)
        scanf("%d",&a[i]);
        for(ll i=1;i<=m;++i)</pre>
            scanf("%d%d",&Q[i].1,&Q[i].r);
            Q[i].id=i;
        for(ll i=1;i<=n;++i){
         be[i]=i/sq+1;
                  sort(Q+1,Q+1+m,cmp1);
        for(int i=1;i<=m;i++){</pre>
                      while(1<Q[i].1)del(a[1]),++1;
            while (1>Q[i].1)--1, ins (a[1]);
            while(r<Q[i].r)++r,ins(a[r]);</pre>
            while(r>Q[i].r)del(a[r]),--r;
            if(Q[i].l==Q[i].r){Q[i].A=0;Q[i].B=1;continue;}
                           Q[i].B=ans1;Q[i].A=(Q[i].r-Q[i].l+1)*(Q[i].r-Q[i].1);
        sort(Q+1,Q+1+m,cmp2);
        for(ll i=1;i<=m;i++){</pre>
         if(Q[i].A==0)printf("0/1\n");
                  int g=gcd(Q[i].A,Q[i].B);
             printf("%11d/%11d\n",Q[i].B/g,Q[i].A/g);
```

```
return 0;
三分
// 整数
while(l+1<r)
    int lm=(l+r)>>1, rm=(lm+r)>>1;
    if(judge(lm)>judge(rm))
        r=rm;
    else
        1=1m;
// double
while(l+eps<r)</pre>
    double lm=(l+r)/2, rm=(lm+r)/2;
    if(judge(lm)>judge(rm))
        r=rm;
    else
        1=1m;
数位 dp
int dfs(int pos, int lim1, int lim2, bool zero)
    if (pos== -1) return 1;
    if (dp[pos][lim1][lim2] != -1)
        return dp[pos][lim1][lim2];
    int up1 = lim1 ? a[pos] : 1;
    int up2 = lim2 ? b[pos] : 1;
    int tmp1 = 0, tmp2 = 0;
    for (int i = 0; i <= up1; i++)</pre>
        for (int j = 0; j <= up2; j++)
            if (i & j)
                continue:
            int tmp = dfs(pos - 1, lim1 && (i == up1), lim2 && (j == up2), zero | | (i ^
j));
            if (!zero && (i ^ j))
                tmp1 = (tmp1 + tmp) \% mod;
            tmp2 = (tmp2 + tmp) \% mod;
    ans = (ans + tmp1 * (pos + 1) % mod) % mod;
    dp[pos][lim1][lim2] = tmp2;
    return tmp2:
void cw(int x, int y)//拆位
    int pos1 = 0, pos2 = 0;
    while (x)
        a[pos1++] = x & 1;
        x >>= 1;
```

```
while (y)
       b[pos2++] = y \& 1;
       y >>= 1;
    for (int i = pos1 - 1; i >= pos2; i--) b[i] = 0;
   dfs(pos1 - 1, 1, 1, 0);
枚举子集
for(int now=S;S!=0;now=(now-1)&S)
        tmp=S^now;//now 为子集, tmp 为 now 的补集
SOSDP (高维前缀和)
/*
f[i]代表i 所代表的所有子集之和
for(int j=0;j<n;j++)//枚举每一位
        for(int i=0;i<1<<n;i++)//枚举每个可能的集合
                if(i>>j&1)//该位为1
                         f[i]+=f[i^{(1<< j)}];//代表i 中缺少了任意一个物品的集合
随机数
unsigned seed=std::chrono::system_clock::now().time_since_epoch().count();
        mt19937 rand(seed);
    uniform int distribution<int> dis(0,1000000000);
   dis(rand);//最后生成的随机数
###
字符串
Manacher
vector<int> d1(n);
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i);
  while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k]) {
   k++;
  d1[i] = k--;
  if (i + k > r) {
   1 = i - k;
   r = i + k;
vector<int> d2(n);
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
  while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[i + k]) {
   k++;
  d2[i] = k--;
```

```
if (i + k > r) {
   1 = i - k - 1;
   r = i + k;
字典树
#include <cstdio>
const int N = 500010;
char s[60];
int n, m, ch[N][26], tag[N], tot = 1;
int main() {
        scanf("%d", &n);
        for (int i = 1; i <= n; ++i) {</pre>
                 scanf("%s", s + 1);
                 int u = 1;
                 for (int j = 1; s[j]; ++j) {
                         int c = s[j] - 'a';
                         if (!ch[u][c]) ch[u][c] = ++tot;
                          u = ch[u][c];
                 tag[u] = 1;
        scanf("%d", &m);
        while (m--) {
                 scanf("%s", s + 1);
                 int u = 1;
                          for (int j = 1; s[j]; ++j) {
                                  int c = s[j] - 'a';
                                  u = ch[u][c];
                                  if (!u) break; // 不存在对应字符的出边说明名字不存在
                 if (tag[u] == 1) {
                          tag[u] = 2;
                          puts("OK");
                 else if (tag[u] == 2)
                          puts("REPEAT");
                 else
                          puts("WRONG");
        return 0;
后缀数组
#include <algorithm>
#include <cstdio>
#include <cstring>
#include <iostream>
using namespace std;
const int N = 1000010;
char s[N];
int n, sa[N], rk[N << 1], oldrk[N << 1], id[N], cnt[N];</pre>
int main() {
        int i, m, p, w;
        scanf("%s", s + 1);
        n = strlen(s + 1);
        m = max(n, 300);
```

```
for (i = 1; i <= n; ++i) ++cnt[rk[i] = s[i]];</pre>
        for (i = 1; i \leftarrow m; ++i) cnt[i] += cnt[i - 1];
        for (i = n; i >= 1; --i) sa[cnt[rk[i]]--] = i;
        for (w = 1; w < n; w <<= 1) {
                 memset(cnt, 0, sizeof(cnt));
                 for (i = 1; i <= n; ++i) id[i] = sa[i];</pre>
                 for (i = 1; i <= n; ++i) ++cnt[rk[id[i] + w]];</pre>
                 for (i = 1; i \leftarrow m; ++i) cnt[i] += cnt[i - 1];
                         for (i = n; i >= 1; --i) sa[cnt[rk[id[i] + w]]--] = id[i];
                 memset(cnt, 0, sizeof(cnt));
                 for (i = 1; i <= n; ++i) id[i] = sa[i];</pre>
                 for (i = 1; i <= n; ++i) ++cnt[rk[id[i]]];</pre>
                 for (i = 1; i <= m; ++i) cnt[i] += cnt[i - 1];</pre>
                 for (i = n; i >= 1; --i) sa[cnt[rk[id[i]]]--] = id[i];
                 memcpy(oldrk, rk, sizeof(rk));
                 for (p = 0, i = 1; i <= n; ++i) {
                         if (oldrk[sa[i]] == oldrk[sa[i - 1]] &&
                                  oldrk[sa[i] + w] == oldrk[sa[i - 1] + w]) {
                                  rk[sa[i]] = p;
                         else {
                                  rk[sa[i]] = ++p;
后缀自动机的各种应用
struct state {
        int len, link;
        std::map<char, int> next;//可以考虑少一个Log
};
// SAM 本身将会存储在一个 state 结构体数组中。我们记录当前自动机的大小 sz 和变量 Last ,当前整
个字符串对应的状态。
const int MAXLEN = 100000;
state st[MAXLEN * 2];
int sz, last;
//我们定义一个函数来初始化 SAM (创建一个只有初始状态的 SAM)。
void sam init() {
        st[0].len = 0;
        st[0].link = -1;
        SZ++;
        last = 0:
//最终我们给出主函数的实现:给当前行末增加一个字符,对应地在之前的基础上建造自动机。
void sam extend(char c) {
        int cur = sz++;
        st[cur].len = st[last].len + 1;
        int p = last:
        while (p != -1 && !st[p].next.count(c)) {
                 st[p].next[c] = cur;
                 p = st[p].link;
        if (p == -1) {
                 st[cur].link = 0;
        else {
                 int q = st[p].next[c];
                 if (st[p].len + 1 == st[q].len) {
```

```
st[cur].link = q;
                else {
                         int clone = sz++;
                         st[clone].len = st[p].len + 1;
                         st[clone].next = st[q].next;
                         st[clone].link = st[q].link;
                         while (p != -1 && st[p].next[c] == q) {
                                 st[p].next[c] = clone;
                                 p = st[p].link;
                         st[q].link = st[cur].link = clone;
        last = cur;
}
广义后缀自动机(多个串中不同子串个数)//理解什么的再说
#include <bits/stdc++.h>
using namespace std;
const int MAXN = 2000000; // 双倍字符串长度
const int CHAR_NUM = 30; // 字符集个数,注意修改下方的 (-'a')
struct exSAM {
 int len[MAXN]:
                           // 节点长度
  int link[MAXN];
                            // 后缀链接, Link
  int next[MAXN][CHAR NUM]; // 转移
  int tot:
                           // 节点总数: [0, tot)
                           //初始化函数
  void init() {
   tot = 1:
   link[0] = -1;
  int insertSAM(int last, int c) { // last 为父 c 为子
   int cur = next[last][c];
    if (len[cur]) return cur;
    len[cur] = len[last] + 1;
    int p = link[last];
    while (p != -1) {
     if (!next[p][c])
       next[p][c] = cur;
     else
       break;
     p = link[p];
    if (p == -1) {
     link[cur] = 0;
     return cur:
    int q = next[p][c];
    if (len[p] + 1 == len[q]) {
     link[cur] = q;
     return cur;
    int clone = tot++;
    for (int i = 0; i < CHAR NUM; ++i)</pre>
     next[clone][i] = len[next[q][i]] != 0 ? next[q][i] : 0;
    len[clone] = len[p] + 1:
    while (p != -1 \&\& next[p][c] == q) {
```

```
next[p][c] = clone;
     p = link[p];
   link[clone] = link[q];
   link[cur] = clone;
   link[q] = clone;
   return cur;
  int insertTrie(int cur, int c) {
   if (next[cur][c]) return next[cur][c]; //已有该节点 直接返回
   return next[cur][c] = tot++;
                                           //无该节点 建立节点
 void insert(const string &s) {
   int root = 0;
   for (auto ch : s) root = insertTrie(root, ch - 'a');
  void insert(const char *s, int n) {
   int root = 0;
   for (int i = 0; i < n; ++i)
     root =
         insertTrie(root, s[i] - 'a'); //一边插入一边更改所插入新节点的父节点
  void build() {
   queue<pair<int, int>> q;
   for (int i = 0; i < 26; ++i)</pre>
     if (next[0][i]) q.push({i, 0});
   while (!q.empty()) { //广搜遍历
     auto item = q.front();
     q.pop();
     auto last = insertSAM(item.second, item.first);
     for (int i = 0; i < 26; ++i)
       if (next[last][i]) q.push({i, last});
} exSam;
char s[1000100];
int main() {
 int n;
 cin >> n;
 exSam.init();
 for (int i = 0; i < n; ++i) {</pre>
   cin >> s;
   int len = strlen(s);
   exSam.insert(s, len);//有string 版本的多态
  exSam.build();
 long long ans = 0;
 for (int i = 1; i < exSam.tot; ++i) {</pre>
   ans += exSam.len[i] - exSam.len[exSam.link[i]];
 cout << ans << endl;</pre>
```

图论

最小生成树

```
Boruvka
```

```
求最小森林 O(ElogV)
#include<bits/stdc++.h>
using namespace std;
const int MaxN = 5000 + 5, MaxM = 2000000 + 5;
int N, M;
int U[MaxM], V[MaxM], W[MaxM];
bool used[MaxM];
int par[MaxN], Best[MaxN];
void init() {
    //scanf("%d %d", &N, &M);
    cin>>N>>M;
    for (int i = 1; i \leftarrow M; ++i)
        cin>>U[i]>>V[i]>>W[i];
        //scanf("%d %d %d", &U[i], &V[i], &W[i]);
void init dsu() {
    for (int i = 1; i \leftarrow N; ++i)
        par[i] = i;
int get_par(int x) {
    if (x == par[x]) return x;
    else return par[x] = get_par(par[x]);
// 比较统一连通块的出边边权
inline bool Better(int x, int y) {
    if (y == 0) return true;
    if (W[x] != W[y]) return W[x] < W[y];</pre>
    return x < y;
void Boruvka() {
    init_dsu();
    int merged = 0, sum = 0:
    bool update = true;
    while (update) {
        update = false;
        memset(Best, 0, sizeof Best);
        for (int i = 1; i <= M; ++i) {
            if (used[i] == true) continue;
            int p = get_par(U[i]), q = get_par(V[i]);
            if (p == q) continue;
```

```
if (Better(i, Best[p]) == true) Best[p] = i;
           if (Better(i, Best[q]) == true) Best[q] = i;
       for (int i = 1; i \leftarrow N; ++i)
           if (Best[i] != 0 && used[Best[i]] == false) {
               update = true;
               merged++; sum += W[Best[i]];
               used[Best[i]] = true;
               // 合并连通块
               par[get_par(U[Best[i]])] = get_par(V[Best[i]]);
   if (merged == N - 1) //printf("%d\n", sum);
       cout<<sum<<"\n";</pre>
   else cout<<"orz\n";</pre>
int main() {
    ios::sync with stdio(false);
   init();
   Boruvka();
   return 0;
二分图
二分图带权最大匹配-KM
template <typename T>
struct hungarian { // km
 int n;
 vector<int> matchx; // 左集合对应的匹配点
 vector<int> matchy; // 右集合对应的匹配点
 vector<int> pre; // 连接右集合的左点
 vector<bool> visx; // 拜访数组 左
 vector<bool> visy; // 拜访数组 右
  vector<T> lx;
 vector<T> ly;
 vector<vector<T> > g;
 vector<T> slack;
 T inf;
 T res;
 queue<int> q;
  int org_n;
 int org m;
  hungarian(int _n, int _m) {
   org_n = n;
   org m = m;
   n = max(\underline{n}, \underline{m});
   inf = numeric limits<T>::max();
   res = 0:
   g = vector<vector<T> >(n, vector<T>(n));
```

```
matchx = vector<int>(n, -1);
 matchy = vector<int>(n, -1);
 pre = vector<int>(n);
 visx = vector<bool>(n);
 visy = vector<bool>(n);
 lx = vector<T>(n, -inf);
 ly = vector<T>(n);
 slack = vector<T>(n);
void addEdge(int u, int v, int w) {
 g[u][v] = max(w, 0); // 负值还不如不匹配 因此设为 0 不影响
 // 最小权匹配改为
     // g[u][v]=w;
bool check(int v) {
 visy[v] = true;
 if (matchy[v] != -1) {
   q.push(matchy[v]);
   visx[matchy[v]] = true; // in S
   return false;
 // 找到新的未匹配点 更新匹配点 pre 数组记录着"非匹配边"上与之相连的点
 while (v != -1) {
   matchy[v] = pre[v];
   swap(v, matchx[pre[
 return true;
void bfs(int i) {
 while (!q.empty()) {
   q.pop();
 q.push(i);
 visx[i] = true;
 while (true) {
   while (!q.empty()) {
     int u = q.front();
     q.pop();
     for (int v = 0; v < n; v++) {
       if (!visy[v]) {
        T delta = lx[u] + ly[v] - g[u][v];
         if (slack[v] >= delta) {
          pre[v] = u;
           if (delta) {
            slack[v] = delta;
           } else if (check(v)) { // delta=0 代表有机会加入相等子图 找增广路
                                 // 找到就 return 重建交错树
            return;
   // 没有增广路 修改顶标
   Ta = inf;
```

```
for (int j = 0; j < n; j++) {
     if (!visy[j]) {
       a = min(a, slack[j]);
   for (int j = 0; j < n; j++) {</pre>
     if (visx[j]) { // 5
       lx[j] -= a;
     if (visy[j]) { // T
       ly[j] += a;
     } else { // T'
       slack[j] -= a;
   for (int j = 0; j < n; j++) {</pre>
     if (!visy[j] && slack[j] == 0 && check(j)) {
       return;
 II \lambda \Box
void solve() {
 // 初始顶标
 for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < n; j++) {
     lx[i] = max(lx[i], g[i][j]);
 for (int i = 0; i < n; i++) {</pre>
   fill(slack.begin(), slack.end(), inf);
   fill(visx.begin(), visx.end(), false);
   fill(visy.begin(), visy.end(), false);
   bfs(i);
 // custom
 for (int i = 0; i < n; i++) {
   if (g[i][matchx[i]] > 0) {
     res += g[i][matchx[i]];
   } else {
     matchx[i] = -1;
      // 最小权
      // q 的初始化要改为-inf
      for(int i=0;i<n;++i){</pre>
               if(g[i][matchx[i]]!=-inf)
                        res+=g[i][matchx[i]];
               else
                        matchx[i]=-1;
      */
```

```
cout << res << "\n";
    for (int i = 0; i < org_n; i++) {</pre>
      cout << matchx[i] + 1 << " ";
    cout << "\n";</pre>
};
点分治
#include <cstdio>
#include <algorithm>
#include <vector>
#include <cstring>
using namespace std;
const int N=1e5+5;
struct node {
   int v, 1;
};
vector<node> g[N];
int n, k, Size, s[N], f[N], root, d[N], K, ans;
vector<int> dep;
bool done[N];
void getroot(int now, int fa) {
    int u:
    s[now] = 1; f[now] = 0;
    for (int i=0; i<g[now].size(); i++)</pre>
        if ((u = g[now][i].v) != fa && !done[u]) {
            getroot(u, now);
            s[now] += s[u];
            f[now] = max(f[now], s[u]);
    f[now] = max(f[now], Size-s[now]);
    if (f[now] < f[root]) root = now;</pre>
void getdep(int now, int fa) {
    int u;
    dep.push_back(d[now]);
    s[now] = 1;
    for (int i=0; i<g[now].size(); i++)</pre>
        if ((u = g[now][i].v) != fa && !done[u]) {
            d[u] = d[now] + g[now][i].1;
            getdep(u, now);
            s[now] += s[u];
int calc(int now, int init) {
    dep.clear(); d[now] = init;
    getdep(now, 0);
    sort(dep.begin(), dep.end());
    int ret = 0;
    for (int l=0, r=dep.size()-1; l<r; )</pre>
        if (dep[1] + dep[r] <= K) ret += r-1++;
        else r--:
    return ret;
void work(int now) {
```

```
int u;
    ans += calc(now, ∅);
    done[now] = true;
    for (int i=0; i<g[now].size(); i++)</pre>
        if (!done[u = g[now][i].v]) {
            ans -= calc(u, g[now][i].1);
            f[0] = Size = s[u];
            getroot(u, root=0);
            work(root);
signed main() {
    while (scanf("%d%d", &n, &K)) {
        if (n == 0 && K == 0) break;
        for (int i=0; i<=n; i++) g[i].clear();</pre>
        memset(done, false, sizeof(done));
        int u, v, 1;
        for (int i=1; i<n; i++) {</pre>
            scanf("%d%d%d", &u, &v, &l);
            g[u].push_back(node(v, 1));
            g[v].push back(node(u, 1));
        f[0] = Size = n;
        getroot(1, root=0);
        ans = 0;
        work(root);
        printf("%d\n", ans);
    return 0;
Floyd
for (k = 1; k \le n; k++) {
         for (i = 1; i <= n; i++) {
                  for (j = 1; j <= n; j++) {
                          f[i][j] = min(f[i][j], f[i][k] + f[k][j]);
}
Tarjan
 强连通分量
const int MAXN = 1e5 + 10;
struct Edge{
         int to, next, dis;
}edge[MAXN << 1];</pre>
int head[MAXN], cnt, ans;
bool inStack[MAXN];
                          //判断是否在栈中
//dfn 第一次访问到该节点的时间(时间戳)
//Low[i] Low[i]能从哪个点(最早时间戳)到达这个点的。
int dfn[MAXN], low[MAXN], tot;
stack<int> stc;
void add_edge(int u, int v, int dis) {
         edge[++cnt].to = v;
```

```
edge[cnt].next = head[u];
        head[u] = cnt;
void Tarjan(int x) {
        dfn[x] = low[x] = ++tot;
        stc.push(x);
        inStack[x] = 1;
        for(int i = head[x]; i; i = edge[i].next) {
                 int to = edge[i].to;
                 if ( !dfn[to] ) {
                          Tarjan(to);
                         low[x] = min(low[x], low[to]);
                 } else if (inStack[to]){
                         low[x] = min(low[x], dfn[to]);
                 }
        //cout << x << " " << low[x] << " " << dfn[x] << endl;
        if(low[x] == dfn[x]) { //发现是整个强连通分量子树里 的最小根。
                 //int cnt = 0;
                 ans++; //强连通分量计数器
                 while(1) {
                          int top = stc.top();
                          stc.pop();
                          //cnt ++;
                         inStack[top] = 0;
                          //cout << top << " "; 每个强连通分量内的点
                         if(top == x) break;
void init() {
        cnt = 1:
        tot = 0;
        ans = 0;
        memset(inStack, 0, sizeof(inStack));
        memset(head, 0, sizeof(head));
        memset(dfn, 0, sizeof(dfn));
        memset(low, 0, sizeof(low));
        while(!stc.empty()) stc.pop();
int main () {
        std::ios::sync with stdio(false);
        cin.tie(0);
        int n, m;
        while(cin >> n >> m && (n \mid \mid m))\{
                 init():
                 int x, y;
                 for(int i = 1; i <= m; ++i) {</pre>
                          cin >> x >> y;
                          add edge(x, y, 0);
                                                   //有向图求强连通
                 for(int i = 1; i <= n; ++i) {</pre>
                         if( !dfn[i] )
                                  Tarjan(i);
        return 0;
```

```
缩点
const int MAXN = 5e3 + 20;
const int MAXM = 1e6 + 10;
int head[MAXN], cnt, tot, dfn[MAXN], low[MAXN], color[MAXN], col;
bool vis[MAXN];
int degree[MAXN];
stack<int> stc;
int n, m;
struct Edge {
        int to, next, dis;
}edge[MAXM << 1];</pre>
void add edge(int u, int v, int dis) {
        edge[++cnt].to = v;
        edge[cnt].next = head[u];
        head[u] = cnt;
void Tarjan(int x) {
        vis[x] = 1;
        dfn[x] = low[x] = ++tot;
        stc.push(x);
        for(int i = head[x]; i; i = edge[i].next) {
                 int to = edge[i].to;
                 if (!dfn[to]) {
                          Tarjan(to);
                          low[x] = min(low[x], low[to]);
                 } else if( vis[to] ) {
                          low[x] = min(low[x], dfn[to]);
        if(dfn[x] == low[x]) {
                 col ++;
                 while(true) {
                          int top = stc.top();
                          stc.pop();
                          color[top] = col; //颜色相同的点缩点
                          vis[top] = 0;
                          cout << top << " ":</pre>
                 //
                          if(top == x) break;
                 //cout << endl;</pre>
void solve(){
        for(int i = 1; i <= n; ++i) {
                 if(!dfn[i])
                          Tarjan(i);
        for(int x = 1; x <= n; ++x) {
                                          //遍历 n 个节点
                 for(int i = head[x]; i; i = edge[i].next) { //缩点后 每个点的出度
                          int to = edge[i].to;
                          if(color[x] != color[to]) {
                                   degree[color[x]] ++;
```

```
void init () {
       cnt = 1;
       tot = 0;
       col = 0:
       memset(vis, 0, sizeof(vis));
       memset(head, 0, sizeof(head));
       memset(dfn, 0, sizeof(dfn));
       memset(low, 0, sizeof(low));
       memset(degree, 0, sizeof(degree));
       memset(color, 0, sizeof(color));
       while(!stc.empty()) stc.pop();
int main () {
       std::ios::sync with stdio(false);
       cin.tie(0);
       while(cin >> n && n) {
               cin >> m;
               init();
               int x, y;
               for(int i = 1; i <= m; ++i) {</pre>
                       cin >> x >> y;
                       add_edge(x, y, 0);
               solve();
       return 0;
割点
#include <bits/stdc++.h>
using namespace std;
int n, m; // n: 点数 m: 边数
int num[100001], low[100001], inde, res;
// num: 记录每个点的时间戳
// Low: 能不经过父亲到达最小的编号, inde: 时间戳, res: 答案数量
bool vis[100001], flag[100001]; // flag: 答案 vis: 标记是否重复
vector<int> edge[100001]; // 存图用的
void Tarjan(int u, int father) { // u 当前点的编号, father 自己爸爸的编号
       vis[u] = true; // 标记
       low[u] = num[u] = ++inde; // 打上时间戳
       int child = 0; // 每一个点儿子数量
       for (auto v : edge[u]) { // 访问这个点的所有邻居 (C++11)
               if (!vis[v]) {
                       child++; // 多了一个儿子
                       Tarjan(v, u); // 继续
                       low[u] = min(low[u], low[v]); // 更新能到的最小节点编号
                       if (father != u && low[v] >= num[u] &&!flag[u]) // 主要代码
                               // 如果不是自己,且不通过父亲返回的最小点符合割点的要求,并
且没有被标记过
                               // 要求即为: 删了父亲连不上去了, 即为最多连到父亲
                       {
                               flag[u] = true;
                               res++; // 记录答案
```

```
else if (v != father)
                        low[u] =
                        min(low[u], num[v]); // 如果这个点不是自己,更新能到的最小节点编号
        if (father == u && child >= 2 &&
                !flag[u]) { // 主要代码,自己的话需要 2 个儿子才可以
                flag[u] = true;
                res++; // 记录答案
int main() {
        cin >> n >> m: // 读入数据
        for (int i = 1; i <= m; i++) { // 注意点是从 1 开始的
                int x, y;
                cin >> x >> y;
                edge[x].push back(y);
                edge[y].push back(x);
        } // 使用 vector 存图
        for (int i = 1; i <= n; i++) // 因为 Tarjan 图不一定连通
                if (!vis[i]) {
                        inde = 0; // 时间戳初始为 0
                        Tarjan(i, i); // 从第 i 个点开始, 父亲为自己
        cout << res << endl;</pre>
        for (int i = 1; i <= n; i++)
                if (flag[i]) cout << i << " "; // 输出结果
        return 0;
点双连通分量
#include<bits/stdc++.h>
using namespace std;
typedef long long 11:
const 11 mod=998244353;
const int maxn=1e6+50;
const 11 inf=0x3f3f3f3f3f3f3f3f3f1LL;
struct Edge{
    int u,v;
};
///割顶 bccno 无意义
int pre[maxn],iscut[maxn],bccno[maxn],dfs_clock,bcc_cut;
vector<int>G[maxn],bcc[maxn];
stack<Edge>S;
int dfs(int u,int fa){
    int lowu = pre[u] = ++dfs clock;
   int child = 0;
   for(int i = 0; i < G[u].size(); i++){</pre>
       int v =G[u][i];
       Edge e = (Edge)\{u,v\};
       if(!pre[v]){ ///没有访问过
           S.push(e);
           child++;
           int lowv = dfs(v, u);
           lowu=min(lowu, lowv);
                                              ///用后代更新
           if(lowv >= pre[u]){
```

```
iscut[u]=true;
                bcc cut++;bcc[bcc cut].clear(); ///注意 bcc 从1 开始
                for(;;){
                    Edge x=S.top();S.pop();
                    if(bccno[x.u] != bcc_cut){bcc[bcc_cut].push_back(x.u);bccno[x.u]=bcc
_cut;}
                    if(bccno[x.v] != bcc_cut){bcc[bcc_cut].push_back(x.v);bccno[x.v]=bcc
_cut;}
                    if(x.u==u&&x.v==v)break:
        else if(pre[v] < pre[u] && v !=fa){</pre>
            S.push(e);
            lowu = min(lowu,pre[v]);
    if(fa < 0 && child == 1) iscut[u] = 0;
    return lowu;
void find_bcc(int n){
    memset(pre, 0, sizeof(pre));
    memset(iscut, 0, sizeof(iscut));
    memset(bccno, 0, sizeof(bccno));
    dfs clock = bcc cut = 0;
    for(int i = 0; i < n; i++)
        if(!pre[i])dfs(i,-1);
边双连通分量
去除所有桥 dfs 即可
int low[MAXN], dfn[MAXN], iscut[MAXN], dfs clock;
bool isbridge[MAXN];
vector<int> G[MAXN];
int cnt bridge:
int father[MAXN];
void tarjan(int u, int fa) {
        father[u] = fa;
        low[u] = dfn[u] = ++dfs clock;
        for (int i = 0; i < G[u].size(); i++) {</pre>
                 int v = G[u][i];
                 if (!dfn[v]) {
                          tarjan(v, u);
                          low[u] = min(low[u], low[v]);
                          if (low[v] > dfn[u]) {//主要
                                   isbridge[v] = true;
                                   ++cnt_bridge;
                 else if (dfn[v] < dfn[u] && v != fa) {</pre>
                          low[u] = min(low[u], dfn[v]);
```

```
2-SAT
#include<bits/stdc++.h>
using namespace std;
#define int long long
const int maxn=1e6+5;
int a[maxn<<1];</pre>
vector<int>g[maxn<<1];</pre>
int tot;
int dfn[maxn<<1],low[maxn<<1];</pre>
stack<int>sta;
int insta[maxn<<1];</pre>
int scccnt;
int color[maxn<<1];</pre>
int n,m;
void tarjan(int u)
    dfn[u]=low[u]=++tot;
    sta.push(u);
    insta[u]=1;
    for(auto v:g[u])
        if(!dfn[v])
             tarjan(v);
             low[u]=min(low[u],low[v]);
        else if(insta[v])
             low[u]=min(low[u],dfn[v]);
    if(dfn[u]==low[u])
        ++scccnt:
        do{
             color[u]=scccnt;
             u=sta.top();sta.pop();
             insta[u]=0;
        }while(low[u]!=dfn[u]);
signed main()
    ios::sync_with_stdio(false);
    cin>>n>>m:
    for(int i=1;i<=m;i++)</pre>
        int a, va, b, vb; //a 为 va 或者 b 为 vb
        cin>>a>>va>>b>>vb;
        if(va&vb)
             g[a+n].push_back(b);
             g[b+n].push_back(a);
        else if(!va&vb)
```

```
g[a].push back(b);
            g[b+n].push back(a+n);
        else if(va&!vb)
            g[a+m].push_back(b+n);
            g[b].push_back(a);
        else if(!va&!vb)
            g[a].push_back(b+n);
            g[b].push_back(a+n);
    /*for (int i = 0; i < m; ++i) {
    int a = read(), va = read(), b = read(), vb = read();
   g[a + n * (va & 1)].push_back(b + n * (vb ^ 1));
    a[b + n * (vb & 1)].push_back(a + n * (va ^ 1));*/
    for(int i=1;i<=(n<<1);i++)</pre>
        if(!dfn[i])
            tarjan(i);
    for(int i=1;i<=n;i++)</pre>
        if(color[i]==color[i+n])
            cout<<"IMPOSSIBLE\n";</pre>
            return 0;
    cout<<"POSSIBLE\n";</pre>
    for(int i=1;i<=n;i++)</pre>
        int tmp=(color[i]<color[i+n])?1:0;</pre>
        cout<<tmp<<" ";
找环
void tarjan(int u) {
   low[u] = dfn[u] = ++dfsClock;
    stk.push(u); ins[u] = true;
    for (const auto &v : g[u]) {
        if (!dfn[v]) tarjan(v), low[u] = std::min(low[u], low[v]);
        else if (ins[v]) low[u] = std::min(low[u], dfn[v]);
   if (low[u] == dfn[u]) {
        ++sccCnt:
        do {
            color[u] = sccCnt;
            u = stk.top(); stk.pop(); ins[u] = false;
        } while (low[u] != dfn[u]);
   }
// 笔者使用了 Tarjan 找环,得到的 color[x] 是 x 所在的 scc 的拓扑逆序。
for (int i = 1; i <= (n << 1); ++i) if (!dfn[i]) tarjan(i);</pre>
for (int i = 1; i <= n; ++i)</pre>
```

```
if (color[i] == color[i + n]) { // x 与 -x 在同一强连通分量内,一定无解
       puts("IMPOSSIBLE");
       exit(0);
puts("POSSIBLE");
for (int i = 1; i \le n; ++i)
   print((color[i] < color[i + n])), putchar(' '); // 如果不使用 Tarjan 找环,请改成大于号
网络流
最大流-Dinic
template <class T>
struct Dinic
    struct Edge
       int v, next;
       T flow;
       Edge() {}
       Edge(int v, int next, T flow) : v(v), next(next), flow(flow) {}
    } e[N * 30];
    int head[N], tot;
    int cur[N]; // 当前弧优化
   int dep[N];
   void init(int siz)
       memset(head, -1, sizeof(head)*(siz+2));
       tot = 0;
   void adde(int u, int v, T w, T rw = 0)
       e[tot] = Edge(v, head[u], w);
       head[u] = tot++;
       cur[u]=head[u];
       e[tot] = Edge(u, head[v], rw);
       head[v] = tot++;
       cur[v]=head[v];
   bool BFS(int _S, int _T)
       memset(dep, 0, sizeof(dep));
       queue<int> q;
       q.push(_S);
       dep[S] = 1;
       while (!q.empty())
           int u = q.front();
           q.pop();
           for (int i = head[u]; ~i; i = e[i].next)
               int v = e[i].v;
               if (!dep[v] && e[i].flow > 0)
                   dep[v] = dep[u] + 1;
                   q.push(v);
       }
```

```
return dep[ T] != 0;
   T dfs(int _S, int _T, T a)
        T flow = 0, f:
        if (_S == _T || a == 0)
            return a;
        for (int i = cur[ S]; ~i; i = e[i].next)
            cur[_S]=i;
            int v = e[i].v;
            if (dep[v] != dep[ S] + 1)
                continue;
            f = dfs(v, _T, min(a, e[i].flow));
            if (f)
                e[i].flow -= f;
                e[i ^ 1].flow += f;
                flow += f:
                a -= f;
                if (a == 0)
                    break;
        if (!flow)
            dep[S] = -1;
        return flow;
    T dinic(int S, int T)
       T max_flow = 0;
        while (BFS(_S, _T))
            max_flow += dfs(_S, _T, INF);
        return max_flow;
};
```

树形 DP 求树的最小支配集,最小点覆盖,最大独立集

一:最小支配集

考虑最小支配集,每个点有两种状态,即属于支配集合或者不属于支配集合,其中不属于支配集合时此点还需要被覆盖,被覆盖也有两种状态,即被子节点覆盖或者被父节点覆盖.总结起来就是三种状态,现对这三种状态定义如下:

1):dp[i] [0],表示点 i 属于支配集合,并且以点 i 为根的子树都被覆盖了的情况下支配集中所包含最少点的个数.

2):dp[i] [1],表示点 i 不属于支配集合,且以 i 为根的子树都被覆盖,且 i 被其中不少于一个子节点覆盖的情况下支配集所包含最少点的个数.

3):dp[i] [2],表示点 i 不属于支配集合,且以 i 为根的子树都被覆盖,且 i 没被子节点覆盖的情况下支配集中所包含最少点的个数.即 i 将被父节点覆盖.

对于第一种状态,dp[i] [0]含义为点 i 属于支配集合,那么依次取每个儿子节点三种状态中的最小值,再把取得的最小值全部加起来再加 1,就是 dp[i] [0]的值了.即只要每个以 i 的儿子为根的子树都被覆盖,再加上当前点 i,所需要的最少的点的个数,DP 转移方程如下:

dp[i] [0] = 1 + ∑(u 取 i 的子节点)min(dp[u] [0], dp[u] [1], dp[u] [2])

对于第三种状态,dp[i] [2]含义为点 i 不属于支配集合,且 i 被其父节点覆盖.那么说明点 i 和点 i 的儿子节点都不属于支配集合.所以点 i 的第三种状态之和其儿子节点的第二种状态有关.方程为:

dp[i] [2] = Σ(u 取 i 的子节点)dp[u] [1]

对于第二种状态,略有些复杂.首先如果点i没有子节点那么 dp[i] [1]应该初始化为 INF;否则为了保证它的每个以i的儿子为根的子树被覆盖,那么要取每个儿子节点的前两种状态的最小值之和,因为此时点i不属于支配集,不能支配其子节点,所以子节点必须已经被支配,与子节点的第三种状态无关,如果当前所选状态中每个儿子都没被选择进入支配集即在每个儿子的前两种状态中,第一种状态都不是所需点最小的,那么为了满足第二种状态的定义(因为点i的第二种状态必须被其子节点覆盖,即其子节点必须有一个属于支配集,如果此时没有,就必须强制使一个子节点的状态为状态一),需要重新选择点i的一个儿子节点为第一种状态,这时取花费最少的一个点,即取 min(dp[u] [0] - dp[u] [1])的儿子节点 u,强制取其第一种状态,其他的儿子节点取第二种状态,DP 转移方程为:

if(i 没有子节点) dp[i] [1] = INF

else dp[i] [1] = ∑(u 取 i 的子节点)min(dp[u] [0], dp[u] [1]) + inc

其中对于 inc 有:

if(上面式子中的 Σ (u 取 i 的子节点)min(dp[u] [0], dp[u] [1]) 中包含某个 dp[u] [0], 即存在一个所选的最小值为状态一的儿子节点) inc = 0

else inc = min(dp[u] [0] - dp[u] [1]) (其中 u 取点 i 的儿子节点)

代码:

```
1 void DP(int u, int p) {// p 为 u 的父节点
       dp[u][2] = 0;
 3
       dp[u][0] = 1;
       bool s = false;
       int sum = 0, inc = INF;
       for(int k = head[u]; k != -1; k = edge[k].next) {
 6
 7
           int to = edge[k].to;
 8
           if(to == p) continue;
 9
           DP(to, u);
10
           dp[u][0] += min(dp[to][0], min(dp[to][1], dp[to][2]));
11
           if(dp[to][0] <= dp[to][1]) {
12
               sum += dp[to][0];
13
               s = true;
14
          }
15
           else {
16
               sum += dp[to][1];
17
               inc = min(inc, dp[to][0] - dp[to][1]);
18
19
           if(dp[to][1] != INF && dp[u][2] != INF) dp[u][2] += dp[to][1];
           else dp[u][2] = INF;
20
21
```

二:最小点覆盖

对于最小点覆盖,每个点只有两种状态,即属于点覆盖或者不属于点覆盖:

1):dp[i] [0]表示点 i 属于点覆盖,并且以点 i 为根的子树中所连接的边都被覆盖的情况下点覆盖集中所包含最少点的个数.

2):**dp[i]** [1]表示点 i 不属于点覆盖,且以点 i 为根的子树中所连接的边都被覆盖的情况下点覆盖集中所包含最少点的个数.

对于第一种状态 dp[i] [0],等于每个儿子节点的两种状态的最小值之和加 1,DP 转移方程如下:

dp[i] [0] = 1 + Σ(u 取 i 的子节点)min(dp[u] [0], dp[u] [1])

对于第二种状态 dp[i] [1],要求所有与 i 连接的边都被覆盖,但是点 i 不属于点覆盖,那么点 i 的所有子节点就必须属于点覆盖,即对于点 i 的第二种状态与所有子节点的第一种状态有关,在数值上等于所有子节点第一种状态的和,DP 转移方程如下:

dp[i] [1] = ∑(u 取 i 的子节点)dp[u] [0]

代码:

```
1 void DP(int u, int p) {// p 为 u 的父节点
      dp[u][0] = 1;
       dp[u][1] = 0;
       for(int k = head[u]; k != -1; k = edge[k].next) {
 5
          int to = edge[k].to;
          if(to == p) continue;
 6
7
          DP(to, u);
 8
          dp[u][0] += min(dp[to][0], dp[to][1]);
9
          dp[u][1] += dp[to][0];
10
11 }
```

三:最大独立集

对于最大独立集,每个点也只有两种状态,即属于点 i 属于独立集或者不属于独立集两种情况:

- 1):dp[i] [0]表示点 i 属于独立集的情况下,最大独立集中点的个数.
- 2):dp[i] [1]表示点 i 不属于独立集的情况下.最大独立集中点的个数.

对于第一种状态 dp[i] [0],由于 i 点属于独立集,所以它的子节点都不能属于独立集,所以对于点 i 的第一种状态,只和它的子节点的第二种状态有关,等于其所有子节点的第二种状态的值加 1,DP 转移方程如下:

dp[i] [0] = 1 + ∑(u 取 i 的子节点) dp[u] [1]

对于第二种状态 dp[i] [1],由于点 i 不属于独立集,所以子节点可以属于独立解,也可以不属于独立集,所取得子节点状态应该为数值较大的那个状态,DP 转移方程:

dp[i] [1] = ∑(u 取 i 的子节点)max(dp[u] [0], dp[u] [1])

```
代码:
```

```
void DP(int u, int p) {// p 为 u 的父节点
    dp[u][0] = 1;
    dp[u][1] = 0;
    for(int k = head[u]; k != -1; k = edge[k].next) {
        int to = edge[k].to;
        if(to == p) continue;
        DP(to, u);
        dp[u][0] += dp[to][1];
        dp[u][1] += max(dp[to][0], dp[to][1]);
    }
}
数据结构
```

```
并杳集
int parent[maxn], rk[maxn];
void init(int n)
{
   for(int i=0;i<n;i++)</pre>
       parent[i]=i;
       rk[i]=0; // 初始树的高度为0
// 合并x 和 v 所属的集合
int fid(int x)
                  //查找 x 元素所在的集合, 回溯时压缩路径
   if (x != parent[x])
       parent[x] = fid(parent[x]);
                                   //回溯时的压缩路径
            //从x 结点搜索到祖先结点所经过的结点都指向该祖先结点
   return parent[x];
void unite(int x,int y)
   x=fid(x);
   y=fid(y);
   if(x==y) return ;
   if(rk[x]<rk[y])</pre>
       parent[x]=y; // 合并是从 rank 小的向 rank 大的连边
   else
       parent[y]=x;
       if(rk[x]==rk[y]) rk[x]++;
```

```
树状数组(区间查询区间修改)
#include<iostream>
#include<stdio.h>
#include<string.h>
#include<algorithm>
using namespace std;
const int maxn=1e5+10;
int sum1[maxn];
int sum2[maxn];
int a[maxn];
int n,m;
int lowbit(int x){
  return x&(-x);
void update(int x,int w){//更新效果: 把x 位置后面所有的数的值+w
  for (int i=x;i<=n;i+=lowbit(i)){</pre>
     sum1[i]+=w;//维护前缀和c[i]
     sum2[i]+=w*(x-1);//维护前缀和c[i]*(n-1)
void range_update(int l,int r,int val)//更新效果: 把 L 位置到 r 位置所有的数的值+w
    update(1,val);
    update(r+1,-val);
int sum(int x){//求1-x的和
 int ans=0;
  for (int i=x;i>0;i-=lowbit(i)){
    ans+=x*sum1[i]-sum2[i];
  return ans;
int range_ask(int l,int r){//求し-r 的和
   return sum(r)-sum(l-1);
int main(){
  while(~scanf("%d%d",&n,&m)){
   for (int i=1;i<=n;i++){</pre>
       scanf("%d",&a[i]);
       update(i,a[i]-a[i-1]);//维护差分数组
  return 0;
主席树(非工程版)
#include <algorithm>
#include <cstdio>
#include <cstring>
using namespace std;
const int maxn = 1e5; // 数据范围
int tot, n, m;
int sum[(maxn << 5) + 10], rt[maxn + 10], ls[(maxn << 5) + 10],
    rs[(maxn << 5) + 10];
int a[maxn + 10], ind[maxn + 10], len;
inline int getid(const int &val) { // 离散化
  return lower_bound(ind + 1, ind + len + 1, val) - ind;
```

```
int build(int l, int r) { // 建树
 int root = ++tot;
 if (1 == r) return root;
 int mid = 1 + r \gg 1;
 ls[root] = build(1, mid);
  rs[root] = build(mid + 1, r);
 return root; // 返回该子树的根节点
int update(int k, int l, int r, int root) { // 插入操作
 int dir = ++tot;
 ls[dir] = ls[root], rs[dir] = rs[root], sum[dir] = sum[root] + 1;
 if (1 == r) return dir;
 int mid = 1 + r \gg 1;
  if (k <= mid)</pre>
   ls[dir] = update(k, 1, mid, ls[dir]);
   rs[dir] = update(k, mid + 1, r, rs[dir]);
  return dir;
int query(int u, int v, int l, int r, int k) { // 查询操作
  int mid = 1 + r \gg 1,
     x = sum[ls[v]] - sum[ls[u]]; // 通过区间减法得到左儿子的信息
  if (1 == r) return 1;
  if (k <= x) // 说明在左儿子中
   return query(ls[u], ls[v], l, mid, k);
  else // 说明在右儿子中
   return query(rs[u], rs[v], mid + 1, r, k - x);
inline void init() {
  scanf("%d%d", &n, &m);
 for (int i = 1; i <= n; ++i) scanf("%d", a + i);</pre>
  memcpy(ind, a, sizeof ind);
  sort(ind + 1, ind + n + 1);
  len = unique(ind + 1, ind + n + 1) - ind - 1;
  rt[0] = build(1, len);
  for (int i = 1; i <= n; ++i) rt[i] = update(getid(a[i]), 1, len, rt[i - 1]);</pre>
int 1, r, k;
inline void work() {
 while (m--) {
   scanf("%d%d%d", &l, &r, &k);
   printf("%d\n", ind[query(rt[l - 1], rt[r], 1, len, k)]); // 回答询问
int main() {
 init();
 work();
 return 0;
线段树
        simple 线段树
初始化可以传vector
vector 从1-n
```

```
区间修改 modify(L,r,k) 默认区间加
单点修改 modify(p,k)
区间查询 modify(l,r)
单点查询 modify(p)
*/
using i64=long long ;
#pragma GCC optimize(2)
#define lson rt<<1
#define rson rt<<1|1
template<class Info.
    class Merge = plus<Info>>
struct SegmentTree{
    const int n;
    const Merge merge:
    vector<Info> info;
    SegmentTree(int n) :n(n) ,merge(Merge()),info(4 << lg(n)) {}</pre>
    SegmentTree(vector<Info> init) : SegmentTree(init.size()-1) {
        function<void(int,int,int)> build = [&](int rt,int l,int r) {
            if (r == 1) {
               info[rt] = init[l];
                return:
            int mid = (1 + r) / 2;
            build(lson.l.mid);
           build(rson,mid+1,r);
            pull(rt);
        build(1,1,n);
    void pull(int rt) { info[rt] = merge(info[lson],info[rson]); }
    void push(int rt)
        pushDown(info[rt],info[lson]);
        pushDown(info[rt],info[rson]);
        cleanLazy(info[rt]);
    void pushDown(Info &a,Info &b){
        if(a.lazy){
           b.s += a.lazy*(b.r-b.l+1);
            b.lazy += a.lazy;
    void cleanLazy(Info &a){ a.lazy=0; }
    void modify(int rt,int l,int r,int L,int R,const int &v) {
        if (L <= 1 && r <= R) {
            info[rt] = info[rt] + v;
            return;
        push(rt):
        int mid = (1 + r) / 2;
        if(L > mid)
            modifv(rson,mid+1,r,L,R,v);
        else if (R <= mid)
            modify(lson, 1,mid,L,R,v);
        else
```

```
modify(lson,l,mid,L,mid,v),modify(rson,mid+1,r,mid+1,R,v);
        pull(rt);
    void modify(int l,int r,const int &v) { modify(1,1,n,1,r,v); }
    void modify(int p,const int & v) {modify(1,1,n,p,p,v);}
    Info rangeQuery(int rt,int l,int r,int L,int R){
        if(R < 1 || r < L )
            return Info(1,r,0);
        if(L <= 1 && r <= R)
            return info[rt];
        push(rt);
        int \ mid = (1 + r) / 2;
        return merge(rangeQuery(lson,1,mid,L,R),rangeQuery(rson,mid+1,r,L,R));
    Info rangeQuery(int 1,int r)
    {
        return rangeQuery(1,1,n,1,r);
};
struct Info {
    int 1,r;
    i64 s, lazy;
    Info(): l(0),r(0),s(0),lazy(0) {}
    Info(int x, i64 val) : l(x), r(x), s(val), lazy(0) {}
    Info(i64 val) : l(0), r(0), s(val), lazy(0) {}
    Info(int L,int R,i64 val) : l(L),r(R),s(val),lazy(0) {}
    Info(int L, int R, i64 val, i64 lz) : l(L), r(R), s(val), lazy(lz) {}
};
Info operator+ (const Info &a,const int& b) // Lazy 下标的运算符重载
    return Info(a.l,a.r,a.s+b*(a.r-a.l+1),a.lazy+b);
Info operator+ (const Info &a, const Info &b) { // 区间合并的运算符重载
    return Info(a.1,b.r,a.s+b.s);
树链剖分
#include<bits/stdc++.h>
using namespace std;
#define int long long
#define pii pair<int,int>
using namespace std;
const int maxn=1e5+10;
struct Node{
    int sum,lazy,l,r,ls,rs;
}node[2*maxn];
int rt,n,m,r,p,a[maxn],cnt,f[maxn],d[maxn],siz[maxn],son[maxn],rk[maxn],top[maxn],id[max
n];
vector<int>g[maxn];
int mod(int a,int b)
    return (a+b)%p;
```

```
inline void add edge(int x,int y)
                  g[x].push_back(y);
void dfs1(int u,int fa,int depth)
                  f[u]=fa;
                  d[u]=depth;
                  siz[u]=1;
                  for(auto &v:g[u])
                                    if(v==fa)
                                                       continue;
                                    dfs1(v,u,depth+1);
                                    siz[u]+=siz[v];
                                    if(siz[v]>siz[son[u]])
                                                      son[u]=v;
void dfs2(int u,int t)
                  top[u]=t;
                  id[u]=++cnt;
                  rk[cnt]=u;
                  if(!son[u])
                                    return;
                  dfs2(son[u],t);
                  for(auto &v:g[u])
                                    if(v!=son[u]\&\&v!=f[u])
                                                      dfs2(v,v);
void pushup(int x)
                  node \lceil x \rceil. sum = (node \lceil node \lceil x \rceil.1s \rceil. sum + node \lceil node \lceil x \rceil.rs \rceil. sum + node \lceil x \rceil.1azy*(node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \lceil x \rceil.r - node \rceil sum + node \lceil node \rceil sum + node \lceil node \rceil sum + node \rceil sum + node \lceil node \rceil sum + node \rceil sum + node \lceil node \rceil sum + node \rceil sum + node \rceil sum + node \lceil node \rceil sum + node
[x].l+1))%p;
void build(int li,int ri,int cur)
                  if(li==ri)
                                    node[cur].l=node[cur].r=li;
                                    node[cur].sum=a[rk[li]];
                                    return;
                  int mid=(li+ri)>>1;
                  node[cur].ls=cnt++;
                  node[cur].rs=cnt++;
                  build(li,mid,node[cur].ls);
                  build(mid+1,ri,node[cur].rs);
                  node[cur].l=node[node[cur].ls].l;
                  node[cur].r=node[node[cur].rs].r;
                  pushup(cur);
void update(int li,int ri,int c,int cur)
                  if(li<=node[cur].1&&node[cur].r<=ri)</pre>
```

```
node[cur].sum=mod(node[cur].sum,c*(node[cur].r-node[cur].l+1));
        node[cur].lazy=mod(node[cur].lazy,c);
        return;
    int mid=(node[cur].l+node[cur].r)>>1;
   if(li<=mid)</pre>
        update(li,ri,c,node[cur].ls);
   if(mid<ri)</pre>
        update(li,ri,c,node[cur].rs);
   pushup(cur);
int query(int li,int ri,int cur)
   if(li<=node[cur].1&&node[cur].r<=ri)</pre>
        return node[cur].sum;
    int tot=node[cur].lazy*(min(node[cur].r,ri)-max(node[cur].l,li)+1)%p;
    int mid=(node[cur].l+node[cur].r)>>1;
   if(li<=mid)</pre>
        tot=mod(tot,query(li,ri,node[cur].ls));
   if(mid<ri)</pre>
        tot=mod(tot,query(li,ri,node[cur].rs));
    return tot%p;
int sum(int x,int y)
    int ans=0:
   int fx=top[x],fy=top[y];
   while(fx!=fy)
        if(d[fx]>=d[fy])
            ans=mod(ans,query(id[fx],id[x],rt));
            x=f[fx],fx=top[x];
        else
            ans=mod(ans,query(id[fy],id[y],rt));
            y=f[fy],fy=top[y];
    if(id[x]<=id[y])</pre>
        ans=mod(ans,query(id[x],id[y],rt));
   else
        ans=mod(ans,query(id[y],id[x],rt));
    return ans%p;
void updates(int x,int y,int c)
    int fx=top[x],fy=top[y];
    while(fx!=fy)
        if(d[fx]>=d[fy])
            update(id[fx],id[x],c,rt);
            x=f[fx],fx=top[x];
        else
            update(id[fy],id[y],c,rt);
```

```
y=f[fy],fy=top[y];
    if(id[x]<=id[v])</pre>
        update(id[x],id[y],c,rt);
    else
        update(id[y],id[x],c,rt);
signed main()
    ios::sync_with_stdio(false);
    cin>>n>>m>>r>>p;
    for(int i=1;i<=n;i++)</pre>
        cin>>a[i];
    for(int i=1;i<n;i++)</pre>
        int x,y;
        cin>>x>>y;
        add edge(x,y);
        add_edge(y,x);
    cnt=0;
    dfs1(r,0,1);
    dfs2(r,r);
    cnt=0;
    rt=cnt++;
    build(1,n,rt);
    for(int i=1;i<=m;i++)</pre>
        int op,x,y,z;
        cin>>op;
        if(op==1)
            cin>>x>>y>>z;
            updates(x,y,z);
        else if(op==2)
            cin>>x>>y;
            cout<<sum(x,y)<<'\n';</pre>
        else if(op==3)
            cin>>x>>z;
            //子树也有连续区间的性质
            update(id[x],id[x]+siz[x]-1,z,rt);
        else if(op==4)
            \verb|cout|<<| query(id[x],id[x]+siz[x]-1,rt)<<'\n';
    return 0;
```

FHQ-Treap

```
FHQ-Treap-ptr
#include<random>
using namespace std;
uniform int distribution < unsigned > u(1, 100000000);
random device rd:
mt19937 e(rd());
struct fhqTreap
    struct Node {
        int val, rnd, siz, num;
        Node* rs, * ls;
        Node() {}
        Node(int x) : val(x), ls(nullptr), rs(nullptr), rnd(u(e)), siz(1), num(1) {}
        Node(int x, int y) :val(\emptyset), ls(nullptr), rs(nullptr), rnd(\emptyset), siz(\emptyset), num(\emptyset) {}
    Node* nullnode = new Node(0, 0);
    Node* root = nullnode:
    int n;
    fhqTreap() :n(0) {};
    Node* newNode(int x) {
        Node* ret = new Node(x);
        ret->ls = nullnode;
        ret->rs = nullnode;
        return ret:
    void pushup(Node* x) {
        if (x == nullnode) return;
        x->siz = x->ls->siz + x->rs->siz:
        x->siz += x->num;
    void split(Node* rt, int k, Node** x, Node** y) {
        if ((rt) == nullnode) { *x = *y = nullnode; return; }
        if ((rt)->val <= k) *x = rt, split(rt->rs, k, &(rt->rs), y);
        else *y = rt, split(rt->ls, k, x, &(rt->ls));
        pushup(rt);
    Node* merge(Node* x, Node* y) {
        if (x == nullnode | | y == nullnode) {
            return x == nullnode ? y : x;
        if (x->rnd < y->rnd) {
            x->rs = merge(x->rs, y);
            pushup(x);
            return x;
        else {
            y->1s = merge(x, y->1s);
            pushup(y);
            return y;
    void Delete(int k)
        Node* x, * y, * z;
        split(root, k, &x, &y);
        split(x, k - 1, &x, &z);
```

```
if (z->num == 1) {
        delete z;
        root = merge(x, y);
    else {
        z->num--, z->siz--;
        root = merge(merge(x, z), y);
void insert(int k)
    Node* x, * y, * z;
    x = y = z = nullnode;
    split(root, k, &x, &y);
    split(x, k - 1, &x, &z);
    // z=new Node(k);
    if (z != nullnode) z->num++, z->siz++;
    else z = newNode(k);
    root = merge(merge(x, z), y);
Node* find_kth(Node* rt, int k)
    while (1) {
        if (rt->ls->siz >= k) rt = rt->ls;
        else if (rt->ls->siz < k \&\& k <= rt->ls->siz + rt->num) return rt;
        else k \rightarrow rt \rightarrow ls \rightarrow siz + rt \rightarrow num, rt = rt \rightarrow rs;
Node* kth(int k)
    return find_kth(root, k);
int kth val(int k)
    return kth(k)->val;
Node* max(Node* rt) {
    if (rt == nullnode) return rt;
    while (rt->rs != nullnode) {
        rt = rt->rs:
    return rt;
Node* max() { return max(root); }
Node* min(Node* rt) {
    if (rt == nullnode) return rt;
    while (rt->ls != nullnode) rt = rt->ls:
    return rt;
Node* min() { return min(root); }
Node* pre(Node** rt, int k)
    Node* x, * y;
    split(*rt, k - 1, &x, &y);
    Node* ret = max(x);
    *rt = merge(x, y);
    return ret:
Node* pre(int k) { return pre(&root, k); }
```

```
Node* sub(Node** rt, int k)
        Node* x, * y;
        split(*rt, k, &x, &y);
        Node* ret = min(y);
        *rt = merge(x, y);
        return ret;
    Node* sub(int k) { return sub(&root, k); }
    int rank(Node** rt, int k)
        Node* x, * y;
        split(*rt, k - 1, &x, &y);
        int ret = x->siz;
        *rt = merge(x, y);
        return ret + 1;
    int rank(int k) { return rank(&root, k); }
};
区间翻转(可以部分替代 splay)
# include<iostream>
# include<cstdio>
# include<cstring>
# include<cstdlib>
using namespace std;
const int MAX=1e5+1;
int n,m,tot,rt;
struct Treap{
    int pos[MAX],siz[MAX],w[MAX];
    int son[MAX][2];
    bool fl[MAX];
    void pus(int x)
        siz[x]=siz[son[x][0]]+siz[son[x][1]]+1;
    int build(int x)
        w[++tot]=x,siz[tot]=1,pos[tot]=rand();
        return tot;
    void down(int x)
        swap(son[x][0], son[x][1]);
        if(son[x][0]) fl[son[x][0]]^=1;
        if(son[x][1]) fl[son[x][1]]^=1;
        fl[x]=0;
    int merge(int x,int y)
        if(!x||!y) return x+y;
        if(pos[x]<pos[y])</pre>
            if(fl[x]) down(x);
            son[x][1]=merge(son[x][1],y);
            pus(x);
```

```
return x;
        if(fl[y]) down(y);
        son[y][0] = merge(x, son[y][0]);
        pus(y);
        return y;
   void split(int i,int k,int &x,int &y)
        if(!i)
            x=y=0;
            return;
        if(fl[i]) down(i);
        if(siz[son[i][0]]<k)
        x=i, split(son[i][1], k-siz[son[i][0]]-1, son[i][1], y);
        y=i, split(son[i][0], k, x, son[i][0]);
        pus(i);
    void coutt(int i)
        if(!i) return;
        if(fl[i]) down(i);
        coutt(son[i][0]);
        printf("%d ",w[i]);
        coutt(son[i][1]);
}Tree;
int main()
    scanf("%d%d",&n,&m);
    for(int i=1;i<=n;i++)</pre>
      rt=Tree.merge(rt,Tree.build(i));
    for(int i=1;i<=m;i++)</pre>
          int 1,r,a,b,c;
          scanf("%d%d",&1,&r);
          Tree.split(rt,l-1,a,b);
        Tree.split(b,r-l+1,b,c);
        Tree.fl[b]^=1;
        rt=Tree.merge(a,Tree.merge(b,c));
   Tree.coutt(rt);
    return 0;
可持久化
#include<cstdio>
#include<cctype>
#include<cstring>
#include<cstdlib>
#include<ctime>
#include<utility>
#include<algorithm>
using namespace std;
typedef pair<int,int> Pair;
```

```
int read() {
    int x=0, f=1;
    char c=getchar();
    for (;!isdigit(c);c=getchar()) if (c=='-') f=-1;
    for (;isdigit(c);c=getchar()) x=x*10+c-'0';
    return x*f;
const int maxn=5e4+5;
const int nlogn=1.3e7+5;
struct node {
    int x,hp,l,r,sum,size;
    bool rev:
    void clear() {
        x=hp=l=r=sum=size=rev=0;
};
struct TREAP
    int pool[nlogn];
    int pooler;
    node t[nlogn];
    int now,all;
    int root[maxn];
    TREAP ():now(0),pooler(1) {
        for (int i=1;i<nlogn;++i) pool[i]=i;</pre>
        root[now]=pool[pooler++];
    int newroot() {
        int ret=pool[pooler++];
        return ret;
    int newnode(int x) {
        int ret=pool[pooler++];
        t[ret].hp=rand();
        t[ret].size=1;
        t[ret].x=t[ret].sum=x;
        return ret;
    void delnode(int x) {
        t[x].clear();
        pool[--pooler]=x;
    void next() {
        root[++all]=newroot();
        t[root[all]]=t[root[now]];
        now=all;
    void back(int x) {
        now=x;
    void update(int x) {
        t[x].sum=t[x].x+t[t[x].1].sum+t[t[x].r].sum;
        t[x].size=t[t[x].1].size+t[t[x].r].size+1;
    void pushdown(int x) {
        if (!t[x].rev) return;
        if (t[x].1) {
            int tx=newnode(t[t[x].1].x);
            t[tx]=t[t[x].1];
            t[tx].rev^=true;
```

```
t[x].1=tx;
    if (t[x].r) {
        int tx=newnode(t[t[x].r].x);
        t[tx]=t[t[x].r];
       t[tx].rev^=true;
        t[x].r=tx;
    swap(t[x].1,t[x].r);
    t[x].rev=false;
int merge(int x,int y) {
    if (!x) return y;
    if (!y) return x;
    int now:
    if (t[x].hp <= t[y].hp) {
        now=newnode(t[x].x);
        t[now]=t[x];
        pushdown(now);
        t[now].r=merge(t[now].r,y);
        now=newnode(t[y].x);
        t[now]=t[y];
        pushdown(now);
        t[now].1=merge(x,t[now].1);
    update(now);
    return now;
Pair split(int x,int p) {
    if (t[x].size==p) return make_pair(x,0);
    int now=newnode(t[x].x);
    t[now]=t[x];
    pushdown(now);
    int l=t[now].1, r=t[now].r;
    if (t[1].size>=p) {
        t[now].1=0;
        update(now);
        Pair g=split(l,p);
        now=merge(g.second,now);
        return make_pair(g.first,now);
    } else if (t[l].size+1==p) {
        t[now].r=0;
        update(now);
        return make_pair(now,r);
    } else {
       t[now].r=0;
        update(now);
        Pair g=split(r,p-t[1].size-1);
        now=merge(now,g.first);
        pushdown(now);
        return make_pair(now, g. second);
void rever(int 1,int r) {
    ++1,++r;
    Pair g=split(root[now],l-1);
    Pair h=split(g.second,r-l+1);
    int want=h.first;
```

```
int here=newnode(t[want].x);
       t[here]=t[want];
       t[here].rev^=true;
       int fi=merge(g.first,here);
       int se=merge(fi,h.second);
       root[now]=se;
   int query(int 1,int r) {
       ++1,++r;
       Pair g=split(root[now],l-1);
       Pair h=split(g.second,r-l+1);
       int want=h.first;
       int ret=t[want].sum;
        int fi=merge(g.first,want);
       int se=merge(fi,h.second);
       root[now]=se;
       return ret;
   void insert(int x) {
       int k=newnode(x);
        root[now]=merge(root[now],k);
} Treap;
int main() {
#ifndef ONLINE JUDGE
    freopen("test.in","r",stdin);
    freopen("my.out","w",stdout);
#endif
    srand(time(0));
    int n=read(), m=read();
   for (int i=1;i<=n;++i) {</pre>
        int x=read();
       Treap.insert(x);
   while (m--) {
       int op=read();
       if (op==1) {
            Treap.next();
            int l=read(),r=read();
            Treap.rever(1,r);
       } else if (op==2) {
            int l=read(),r=read();
            int ans=Treap.query(1,r);
            printf("%d\n",ans);
       } else if (op==3) {
            Treap.back(read());
    return 0;
#include <bits/stdc++.h>
using namespace std;
typedef long long 11;
const int N = 100010;
int m,n;
```

```
struct Splay
   int rt, tot, fa[N], ch[N][2], val[N], cnt[N], sz[N]; // cnt 权值出现次数
   int maxVal=-0x3f3f3f3f;
   bool rev[N];
   void pushdown(int x)
       if (rev[x])
           swap(ch[x][0], ch[x][1]);
           rev[ch[x][0]] ^= 1;
           rev[ch[x][1]] ^= 1;
           rev[x] = 0;
   void maintain(int x)
       sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + cnt[x];
   // 右儿子返回1 左儿子返回0
   bool get(int x) { return x == ch[fa[x]][1]; }
   void clear(int x)
       ch[x][0] = ch[x][1] = fa[x] = val[x] = sz[x] = cnt[x] = 0;
   // 旋转操作
   void rotate(int x)
       int y = fa[x], z = fa[y], chk = get(x);
       // x 是左儿子右旋,右儿子左旋
       ch[y][chk] = ch[x][chk ^ 1];
       if (ch[x][chk ^ 1])
           fa[ch[x][chk ^ 1]] = y;
       ch[x][chk ^ 1] = y;
       fa[y] = x;
       fa[x] = z;
       if (z)
           ch[z][y == ch[z][1]] = x;
       maintain(x);
       maintain(y);
   // SLpay 操作
   void splay(int x, int goal = 0)
       for (int f = fa[x]; f = fa[x], f != goal; rotate(x))
           if (fa[f] != goal)
               rotate(get(x) == get(f) ? f : x);
       if (!goal)
           rt = x;
   // 插入
   void ins(int k)
       maxVal=max(maxVal,k);
       // 树空
```

```
if (!rt)
        val[++tot] = k;
        cnt[tot]++;
       rt = tot;
       maintain(rt);
        return;
   int cur = rt, f = 0;
   while (1)
       if (val[cur] == k)
           cnt[cur]++;
           maintain(cur);
           maintain(f);
           splay(cur);
           break;
       f = cur;
        cur = ch[cur][val[cur] < k];</pre>
       if (!cur)
           val[++tot] = k;
           cnt[tot]++;
            fa[tot] = f;
           ch[f][val[f] < k] = tot;
           maintain(tot);
           maintain(f);
            splay(tot);
           break;
// 查询排名 等价于find
int rk(int k)
   int res = 0, cur = rt;
   while (1)
        if (k < val[cur])</pre>
            cur = ch[cur][0];
       else
           res += sz[ch[cur][0]];
           if (k == val[cur])
                splay(cur);
                return res + 1;
           res += cnt[cur];
           cur = ch[cur][1];
// 查询第k 大 索引
```

```
int kth idx(int k)
    int cur = rt;
    while (1)
        pushdown(cur);
        if (ch[cur][0] && k <= sz[ch[cur][0]])</pre>
           cur = ch[cur][0];
        else
           k -= cnt[cur] + sz[ch[cur][0]];
           if (k <= 0)
                splay(cur);
               return cur;
           cur = ch[cur][1];
}
// 查询第k 大 值
int kth_val(int k) { return val[kth_idx(k)]; }
// 查询前驱
int pre()
    int cur = ch[rt][0];
    if (!cur)
        return cur;
    while (ch[cur][1])
       cur = ch[cur][1];
    splay(cur);
    return cur;
// 查询后继
int nxt()
    int cur = ch[rt][1];
    if (!cur)
       return cur;
    while (ch[cur][0])
        cur = ch[cur][0];
    splay(cur);
    return cur;
// pre 封装
int q_pre(int x)
    ins(x);
    int ret = val[pre()];
    del(x);
    return ret;
// nxt 封装
int q_nxt(int x)
   ins(x);
```

```
int ret = val[nxt()];
   del(x);
   return ret;
// 删除
void del(int k)
   rk(k);
   if (cnt[rt] > 1)
        cnt[rt]--;
        maintain(rt);
        return;
   if (!ch[rt][0] && !ch[rt][1])
        clear(rt);
        rt = 0;
        return;
   if (!ch[rt][0])
        int cur = rt;
        rt = ch[rt][1];
        fa[rt] = 0;
        clear(cur);
        return:
   if (!ch[rt][1])
        int cur = rt;
        rt = ch[rt][0];
        fa[rt] = 0;
        clear(cur);
        return;
   int cur = rt, x = pre();
   fa[ch[cur][1]] = x;
   ch[x][1] = ch[cur][1];
   clear(cur);
   maintain(rt);
void reverse(int 1, int r)
   int x = kth_idx(1), y = kth_idx(r + 2);
   splay(x);
   splay(y, x);
    rev[ch[y][0]] ^= 1;
// 打印索引为x 的节点及其子树
void output(int x)
   pushdown(x);
   if (ch[x][0])
        output(ch[x][0]);
   if (val[x] && val[x] <= n)</pre>
        cout << val[x] <<" ";
```

```
if (ch[x][1])
            output(ch[x][1]);
    //打印整颗树
    void print_tree()
        output(rt);
    // 已知序列建树
    int build(int 1,int r)
        int x=++tot;
        int mid=(1+r)/2;
        cnt[tot]++;
        // val[tot]=xxxx;
        if(l==r){
            maintain(x);
            return x;
        ch[x][0]=build(1,mid-1);
        ch[x][1]=build(mid+1,r);
        maintain(x);
        return x;
} tree;
LCA
// 倍增方法:
#include<bits/stdc++.h>
#define MXN 50007
using namespace std;
std::vector<int> v[MXN];
std::vector<int> w[MXN];
int fa[MXN][31], cost[MXN][31], dep[MXN];
int n, m;
int a, b, c;
void dfs(int root, int fno) {
        fa[root][0] = fno;
        dep[root] = dep[fa[root][0]] + 1;
        for (int i = 1; i < 31; ++i) {
                 fa[root][i] = fa[fa[root][i - 1]][i - 1];
                 cost[root][i] = cost[fa[root][i - 1]][i - 1] + cost[root][i - 1];
        int sz = v[root].size();
        for (int i = 0; i < sz; ++i) {</pre>
                 if (v[root][i] == fno) continue;
                 cost[v[root][i]][0] = w[root][i];
                 dfs(v[root][i], root);
int lca(int x, int y) {
        if (dep[x] > dep[y]) swap(x, y);
        int tmp = dep[y] - dep[x], ans = 0;
        for (int j = 0; tmp; ++j, tmp >>= 1)
                 if (tmp \& 1) ans += cost[y][j], y = fa[y][j];
        if (y == x) return ans;
        for (int j = 30; j >= 0 && y != x; --j) {
                 if (fa[x][j] != fa[y][j]) {
```

```
ans += cost[x][j] + cost[y][j];
                          x = fa[x][j];
                          y = fa[y][j];
         ans += cost[x][0] + cost[y][0];
         return ans;
int main() {
    ios::sync_with_stdio(false);
         memset(fa, 0, sizeof(fa));
         memset(cost, 0, sizeof(cost));
                  memset(dep, 0, sizeof(dep));
         // scanf("%d", &n);
    cin>>n:
        for (int i = 1; i < n; ++i) {
                 // scanf("%d %d %d", &a, &b, &c);
        cin>>a>>b>>c;
                 ++a, ++b;
                 v[a].push_back(b);
                 v[b].push back(a);
                 w[a].push_back(c);
                 w[b].push back(c);
        dfs(1, 0);
        // scanf("%d", &m);
    cin>>m:
        for (int i = 0; i < m; ++i) {
                 // scanf("%d %d", &a, &b);
                 cin>>a>>b;
        ++a, ++b;
                  // printf("%d\n", lca(a, b));
        cout<<lca(a,b)<<"\n";</pre>
        return 0;
// Tarjan 方法:
#include <algorithm>
#include <iostream>
using namespace std;
class Edge {
public:
         int toVertex, fromVertex;
         int next;
        int LCA;
         Edge(): toVertex(-1), fromVertex(-1), next(-1), LCA(-1) \{\};
        Edge(int u, int v, int n) : fromVertex(u), toVertex(v), next(n), LCA(-1) {};
};
const int MAX = 100;
int head[MAX], queryHead[MAX];
Edge edge[MAX], queryEdge[MAX];
int parent[MAX], visited[MAX];
int vertexCount, edgeCount, queryCount;
void init() {
        for (int i = 0; i <= vertexCount; i++) {</pre>
                 parent[i] = i;
```

```
int find(int x) {
        if (parent[x] == x) {
                  return x;
        else {
                  return find(parent[x]);
void tarjan(int u) {
        parent[u] = u;
        visited[u] = 1;
        for (int i = head[u]; i != -1; i = edge[i].next) {
                 Edge& e = edge[i];
                 if (!visited[e.toVertex]) {
                          tarjan(e.toVertex);
                          parent[e.toVertex] = u;
        for (int i = queryHead[u]; i != -1; i = queryEdge[i].next) {
                 Edge& e = queryEdge[i];
                 if (visited[e.toVertex]) {
                          queryEdge[i ^ 1].LCA = e.LCA = find(e.toVertex);
int main() {
        memset(head, 0xff, sizeof(head));
        memset(queryHead, 0xff, sizeof(queryHead));
        cin >> vertexCount >> edgeCount >> queryCount;
        int count = 0;
        for (int i = 0; i < edgeCount; i++) {</pre>
                 int start = 0, end = 0;
                 cin >> start >> end;
                          edge[count] = Edge(start, end, head[start]);
                 head[start] = count;
                 count++;
                 edge[count] = Edge(end, start, head[end]);
                 head[end] = count;
                 count++;
        count = 0;
        for (int i = 0; i < queryCount; i++) {</pre>
                 int start = 0, end = 0;
                 cin >> start >> end;
                 queryEdge[count] = Edge(start, end, queryHead[start]);
                  queryHead[start] = count;
                 count++;
                 queryEdge[count] = Edge(end, start, queryHead[end]);
                 queryHead[end] = count;
                 count++;
        init();
        tarian(1):
        for (int i = 0; i < queryCount; i++) {</pre>
                 Edge& e = queryEdge[i * 2];
                 cout << "(" << e.fromVertex << "," << e.toVertex << ") " << e.LCA << en
dl:
```

```
return 0;
大数
//注: 可以直接把BigInt 和一样用 cin cout 都行,就是高精乘为了速度才用了 FFT 降低了精度,有需要可
#include <cstdio>
#include <iostream>
#include <cmath>
#include <string>
#include <cstring>
#include <vector>
#include <algorithm>
using namespace std;
const double PI = acos(-1.0);
struct Complex{
    double x,y;
    Complex(double _x = 0.0, double _y = 0.0){
       x = x;
       y = _y;
   Complex operator-(const Complex &b)const{
        return Complex(x - b.x,y - b.y);
    Complex operator+(const Complex &b)const{
        return Complex(x + b.x, y + b.y);
   Complex operator*(const Complex &b)const{
        return Complex(x*b.x - y*b.y,x*b.y + y*b.x);
void change(Complex y[],int len){
    int i,j,k;
    for(int i = 1, j = len/2; i < len-1; i++){}
        if(i < j) swap(y[i],y[j]);
       k = len/2;
       while(j >= k){
           j = j - k;
            k = k/2;
       if(j < k)
                    j+=k;
void fft(Complex y[],int len,int on){
    change(y,len);
    for(int h = 2;h <= len;h<<=1){
        Complex wn(cos(on*2*PI/h),sin(on*2*PI/h));
       for(int j = 0; j < len; j += h){
            Complex w(1,0);
            for(int k = j; k < j + h/2; k++){
               Complex u = y[k];
               Complex t = w*y[k + h/2];
               y[k] = u + t;
               y[k + h/2] = u - t;
               w = w*wn;
```

```
if(on == -1){
        for(int i = 0; i < len; i++){}
            y[i].x /= len;
class BigInt
#define Value(x, nega) ((nega) ? -(x) : (x))
#define At(vec, index) ((index) < vec.size() ? vec[(index)] : 0)</pre>
    static int absComp(const BigInt &lhs, const BigInt &rhs)
        if (lhs.size() != rhs.size())
            return lhs.size() < rhs.size() ? -1 : 1;</pre>
        for (int i = lhs.size() - 1; i >= 0; --i)
            if (lhs[i] != rhs[i])
                return lhs[i] < rhs[i] ? -1 : 1;
        return 0;
    using Long = long long;
    const static int Exp = 9;
    const static Long Mod = 10000000000;
    mutable std::vector<Long> val;
    mutable bool nega = false;
    void trim() const
    {
        while (val.size() && val.back() == 0)
            val.pop_back();
        if (val.empty())
            nega = false;
    int size() const { return val.size(); }
    Long &operator[](int index) const { return val[index]; }
    Long &back() const { return val.back(); }
    BigInt(int size, bool nega) : val(size), nega(nega) {}
    BigInt(const std::vector<Long> &val, bool nega) : val(val), nega(nega) {}
public:
    friend std::ostream &operator<<((std::ostream &os, const BigInt &n)</pre>
        if (n.size())
            if (n.nega)
                putchar('-');
            for (int i = n.size() - 1; i >= 0; --i)
                if (i == n.size() - 1)
                    printf("%lld", n[i]);
                else
                    printf("%0*1ld", n.Exp, n[i]);
        else
            putchar('0');
        return os;
    friend BigInt operator+(const BigInt &lhs, const BigInt &rhs)
```

```
BigInt ret(lhs);
        return ret += rhs;
    friend BigInt operator-(const BigInt &lhs, const BigInt &rhs)
        BigInt ret(lhs);
        return ret -= rhs;
    BigInt(Long x = 0)
        if (x < 0)
            x = -x, nega = true;
        while (x >= Mod)
            val.push_back(x % Mod), x /= Mod;
        if (x)
            val.push back(x);
    BigInt(const char *s)
        int bound = 0, pos;
        if (s[0] == '-')
            nega = true, bound = 1;
        Long cur = 0, pow = 1;
        for (pos = strlen(s) - 1; pos >= Exp + bound - 1; pos -= Exp, val.push_back(cur),
 cur = 0, pow = 1)
            for (int i = pos; i > pos - Exp; --i)
                cur += (s[i] - '0') * pow, pow *= 10;
        for (cur = 0, pow = 1; pos >= bound; --pos)
            cur += (s\lceil pos \rceil - '0') * pow, pow *= 10;
        if (cur)
            val.push_back(cur);
    BigInt &operator=(const char *s){
        BigInt n(s);
        *this = n;
        return n;
    BigInt &operator=(const Long x){
        BigInt n(x);
        *this = n:
        return n;
    friend std::istream &operator>>(std::istream &is, BigInt &n){
        string s;
        is >> s;
        n=(char*)s.data();
        return is;
    BigInt &operator+=(const BigInt &rhs)
        const int cap = std::max(size(), rhs.size()) + 1;
        val.resize(cap);
        int carry = 0;
        for (int i = 0; i < cap - 1; ++i)
            val[i] = Value(val[i], nega) + Value(At(rhs, i), rhs.nega) + carry, carry =
0;
            if (val[i] >= Mod)
                val[i] -= Mod, carry = 1;
```

```
else if (val[i] < 0)</pre>
            val[i] += Mod, carry = -1;
    if ((val.back() = carry) == -1) //assert(val.back() == 1 or 0 or -1)
        nega = true, val.pop_back();
        bool tailZero = true;
        for (int i = 0; i < cap - 1; ++i)
            if (tailZero && val[i])
                val[i] = Mod - val[i], tailZero = false;
                val[i] = Mod - 1 - val[i];
    trim();
    return *this;
friend BigInt operator-(const BigInt &rhs)
    BigInt ret(rhs);
    ret.nega ^= 1;
    return ret;
BigInt &operator -= (const BigInt &rhs)
    rhs.nega ^= 1;
    *this += rhs;
    rhs.nega ^= 1;
    return *this;
friend BigInt operator*(const BigInt &lhs, const BigInt &rhs)
    int len=1;
    BigInt ll=lhs,rr=rhs;
    11.nega = lhs.nega ^ rhs.nega;
    while(len<2*lhs.size()||len<2*rhs.size())len<<=1;</pre>
    11.val.resize(len),rr.val.resize(len);
    Complex x1[len],x2[len];
    for(int i=0;i<len;i++){</pre>
        Complex nx(ll[i],0.0),ny(rr[i],0.0);
        x1[i]=nx;
        x2[i]=ny;
    fft(x1,len,1);
    fft(x2,len,1);
    for(int i = 0 ; i < len; i++)</pre>
        x1[i] = x1[i] * x2[i];
    fft( x1 , len , -1 );
    for(int i = 0 ; i < len; i++)</pre>
        11[i] = int(x1[i].x + 0.5);
    for(int i = 0 ; i < len; i++){</pre>
       ll[i+1]+=ll[i]/Mod;
        ll[i]%=Mod;
    11.trim();
    return 11:
friend BigInt operator*(const BigInt &lhs, const Long &x){
```

```
BigInt ret=lhs;
    bool negat = (x < 0);
    Long xx = (negat) ? -x : x;
    ret.nega ^= negat;
    ret.val.push back(0);
    ret.val.push_back(0);
    for(int i = 0; i < ret.size(); i++)</pre>
        ret[i]*=xx;
    for(int i = 0; i < ret.size(); i++){</pre>
        ret[i+1]+=ret[i]/Mod;
        ret[i] %= Mod;
    ret.trim();
    return ret;
BigInt &operator*=(const BigInt &rhs) { return *this = *this * rhs; }
BigInt &operator*=(const Long &x) { return *this = *this * x; }
friend BigInt operator/(const BigInt &lhs, const BigInt &rhs)
    static std::vector<BigInt> powTwo{BigInt(1)};
    static std::vector<BigInt> estimate;
    estimate.clear();
    if (absComp(lhs, rhs) < 0)</pre>
        return BigInt();
    BigInt cur = rhs;
    int cmp:
    while ((cmp = absComp(cur, lhs)) <= 0)</pre>
        estimate.push back(cur), cur += cur;
        if (estimate.size() >= powTwo.size())
            powTwo.push_back(powTwo.back() + powTwo.back());
    if (cmp == 0)
        return BigInt(powTwo.back().val, lhs.nega ^ rhs.nega);
    BigInt ret = powTwo[estimate.size() - 1];
    cur = estimate[estimate.size() - 1];
    for (int i = estimate.size() - 1; i >= 0 && cmp != 0; --i)
        if ((cmp = absComp(cur + estimate[i], lhs)) <= 0)</pre>
            cur += estimate[i], ret += powTwo[i];
    ret.nega = lhs.nega ^ rhs.nega;
    return ret;
friend BigInt operator/(const BigInt &num,const Long &x){
    bool negat = (x < 0);
    Long xx = (negat) ? -x : x;
    BigInt ret;
    Long k = 0;
    ret.val.resize( num.size() );
    ret.nega = (num.nega ^ negat);
    for(int i = num.size() - 1; i >= 0; i--){
        ret[i] = (k * Mod + num[i]) / xx;
        k = (k * Mod + num[i]) % xx;
    ret.trim():
    return ret;
bool operator == (const BigInt &rhs) const
    return nega == rhs.nega && val == rhs.val;
```

```
bool operator!=(const BigInt &rhs) const { return nega != rhs.nega || val != rhs.val;
   bool operator>=(const BigInt &rhs) const { return !(*this < rhs); }</pre>
   bool operator>(const BigInt &rhs) const { return !(*this <= rhs); }</pre>
   bool operator<=(const BigInt &rhs) const</pre>
       if (nega && !rhs.nega)
           return true;
       if (!nega && rhs.nega)
           return false;
       int cmp = absComp(*this, rhs);
       return nega ? cmp >= 0 : cmp <= 0;
   bool operator<(const BigInt &rhs) const</pre>
       if (nega && !rhs.nega)
           return true;
       if (!nega && rhs.nega)
           return false;
       return (absComp(*this, rhs) < 0) ^ nega;</pre>
   void swap(const BigInt &rhs) const
       std::swap(val, rhs.val);
       std::swap(nega, rhs.nega);
BigInt ba,bb;
int main(){
   cin>>ba>>bb;
   cout << ba + bb << '\n';//和
   cout << ba - bb << '\n';//差
   cout << ba * bb << '\n';//积
   BigInt d;
   cout << (d = ba / bb) << '\n';//商
   cout << ba - d * bb << '\n';//余
   return 0;
带修改整体二分
#include <bits/stdc++.h>
using namespace std:
const int maxn = 1e5 + 5;
struct query1
   int type;
   int 1, r, k;
   int id;
   // 对于询问, type = 1, l, r 表示区间左右边界, k 表示询问第 k 小
   // 对于修改,type = 0,L 表示修改位置,r 表示修改后的值,
   // k 表示当前操作是插入(1)还是擦除(-1), 更新树状数组时使用,
   // id 记录每个操作原先的编号, 因二分过程中操作顺序会被打散
};
struct num1
```

```
int p, x;
};
vector<int> ans(100000);
int n, m;
int t[maxn];
int a[maxn];
int sum(int p)
    int ans = 0;
    while (p)
        ans += t[p];
        p -= p \& (-p);
    return ans;
void add(int p, int x)
    while (p <= n)
        t[p] += x;
        p += p \& (-p);
void solve(int 1, int r, vector<query1> &q)
    if (q.size() == 0)
        return;
    if (1 == r)
        for (auto i : q)
            if (i.type == 1)
                ans[i.id] = 1;
        return;
    int mid = (1 + r) >> 1;
    vector<query1> q1;
    vector<query1> q2;
    for (auto i : q)
        if (i.type == 1)
            int t = sum(i.r) - sum(i.l - 1);
            if (i.k <= t)
                q1.push_back(i);
            else
                i.k -= t;
                q2.push_back(i);
        else
            if (i.r <= mid)</pre>
                add(i.1, i.k);
                q1.push_back(i);
```

```
else
                q2.push_back(i);
   for (auto i : q)
        if(i.type==0)
            if (i.r <= mid)</pre>
                add(i.1, -i.k);
    solve(1, mid, q1);
   solve(mid + 1, r, q2);
signed main()
   ios::sync with stdio(false);
   int tmp1, tmp2, tmp3;
    vector<query1> query;
   cin >> n >> m;
    for (int i = 1; i <= n; i++)</pre>
        cin >> a[i];
        query.push_back({0, i, a[i], 1, 0});
   char op;
   int 1, r, k;
   int cnt = 0;
   int x, y;
    for (int i = 1; i <= m; i++)</pre>
        cin >> op;
        if (op == 'Q')
            cin >> 1 >> r >> k;
            query.push_back({1, 1, r, k, ++cnt});
        else if (op == 'C')
            cin >> x >> y;
            query.push_back({0, x, a[x], -1, 0});
            a[x]=y;
            query.push_back({0, x, y, 1, 0});
    solve(0, 1e9, query);
   for (int i = 1; i <= cnt; i++)
        cout << ans[i] << "\n";</pre>
```

数学

```
线形基
#include<bits/stdc++.h>
using namespace std;
// #pragma GCC optimize(2)
#define fi first
#define se second
typedef pair<int, int> pii;
typedef long long ll;
typedef unsigned long long ull;
typedef long double ld;
void io() { ios::sync_with_stdio(false); cin.tie(0); cout.tie(0); }
template<typename T>
inline void debug(T const& x) { cout << x << "\n"; }</pre>
struct LinearBase {
    const int siz=64:
    int MN;
    vector<11>p, tmp;
    bool flag = false;
    LinearBase(){
        p.resize(siz);
        tmp.resize(siz);
        MN=siz-1;
    void clear()
       // siz = MN = 0;
        p.clear(); tmp.clear();
        flag = false;
    void resize(int size)
        p.resize(siz);
        tmp.resize(siz);
        MN = siz - 1;
        flag = false;
    void insert(ll x)
        for (int i = MN; ~i; --i) {
            if (x & (111 << i)) {
                if (!p[i]) {
                    p[i] = x;
                    return;
                else
                    x ^= p[i];
        flag = true;
    bool check(ll x)
```

```
for (int i = MN; ~i; i--) {
            if (x & (111 << i)) {
                if (!p[i]) return false;
                else x ^= p[i];
        return true;
    11 Qmax()
        11 res = 011;
        for (int i = MN; \sim i; --i) {
            res = max(res, res ^ p[i]);
        return res;
    11 Qmin()
        if (flag) return 0;
        for (int i = 0; i <= MN; ++i)</pre>
            if (p[i]) p[i];
    }
    // void rebuild()
    // {
    //
           int cnt=0,top=0;
    //
           for(int i=MN;~i)
    // }
    11 Qnth_element(l1 k)
        11 \text{ res} = 0;
        int cnt = 0;
        k -= flag;
        if (!k) return 0;
        for (int i = 0; i <= MN; ++i) {</pre>
            for (int j = i - 1; ~j; j--) {
                if (p[i] & (1ll << j)) p[i] ^= p[j];</pre>
            if (p[i]) tmp[cnt++] = p[i];
        if (k >= (111 << cnt)) return -1;
        for (int i = 0; i < cnt; ++i)
            if (k & (1ll << i))
                res ^= tmp[i];
        return res;
};
int main()
```

```
io();
    int n;
    cin >> n;
    LinearBase lb;
    for (int i = 0; i < n; ++i) {
        11 _;
        cin >> ;
        lb.insert( );
    cout << lb.Qmax() << "\n";</pre>
    return 0;
pollard's rho
#include <bits/stdc++.h>
#define sz(x) int((x).size())
#define all(x) begin(x), end(x)
using namespace std;
template<class T>
using vc = vector<T>;
using ull = unsigned long long;
using 11 = long long;
ull modmul(ull a, ull b, ull M) {
         11 ret = a * b - M * ull(1.L / M * a * b);
         return ret + M * (ret < 0) - M * (ret >= (11)M);
ull modpow(ull b, ull e, ull mod) {
         ull ans = 1:
         for (; e; b = modmul(b, b, mod), e /= 2)
                  if (e & 1) ans = modmul(ans, b, mod);
         return ans:
}
ull Qpow(ull b, int e) {
  ull res = 1;
  for (; e; b *= b, e /= 2) if (e & 1) res *= b;
  return res:
bool isPrime(ull p) {
         if (p == 2) return true;
         if (p == 1 || p % 2 == 0) return false;
         ull s = p - 1;
         while (s \% 2 == 0) s /= 2;
  for (int i = 0; i < 15; ++i) {</pre>
                 ull a = rand() \% (p - 1) + 1, tmp = s;
                  ull mod = modpow(a, tmp, p);
                  while (tmp != p - 1 && mod != 1 && mod != p - 1) {
                          mod = modmul(mod, mod, p);
                          tmp *= 2:
                  if (mod != p - 1 && tmp % 2 == 0) return false;
         return true;
}
```

```
ull pollard(ull n) {
        auto f = [n](ull x) { return modmul(x, x, n) + 1; };
        ull x = 0, y = 0, t = 30, prd = 2, i = 1, q;
        while (t++ % 40 || __gcd(prd, n) == 1) {
                 if (x == y) x = ++i, y = f(x);
                 if ((q = modmul(prd, max(x,y) - min(x,y), n))) prd = q;
                 x = f(x), y = f(f(y));
        return __gcd(prd, n);
}
vector<ull> factor(ull n) {
        if (n == 1) return {};
        if (isPrime(n)) return {n};
        ull x = pollard(n);
        auto 1 = factor(x), r = factor(n / x);
        1.insert(l.end(), all(r));
        return 1;
}
int main() {
#ifdef LOCAL
  freopen("in.txt", "r", stdin);
#endif
        cin.tie(nullptr)->sync with stdio(false);
  int n; cin >> n;
  while (n--) {
    ull x; cin >> x;
    auto fac = factor(x);
    map<ull, int> mp;
    for (auto e: fac) {
      ++mp[e];
    ull ans = 1;
    for (auto p: mp) {
     ans *= Opow(p.first, p.second / 3);
    cout << ans << '\n';</pre>
数论和杂项
扩欧求逆元
inline void exgcd(LL a, LL b, LL &x, LL &y)
    if (!b)
        x = 1:
        y = 0;
        return;
    exgcd(b, a % b, y, x);
    y -= a / b * x;
inline LL inv(LL a, LL mo)
```

```
LL x, y;
    exgcd(a, mo, x, y);
    return x >= 0 ? x : x + mo;
NTT
#include<bits/stdc++.h>
using namespace std;
inline int read() {
  int x = 0, f = 1;
  char ch = getchar();
  while (ch < '0' || ch > '9') {
    if (ch == '-') f = -1;
    ch = getchar();
  while (ch <= '9' && ch >= '0') {
   x = 10 * x + ch - '0';
    ch = getchar();
  return x * f;
void print(int x) {
  if (x < 0) putchar('-'), x = -x;
  if (x >= 10) print(x / 10);
  putchar(x % 10 + '0');
const int N = 300100, P = 998244353;
inline int qpow(int x, int y) {
  int res(1);
  while (v)
   if (y & 1) res = 111 * res * x % P;
   x = 111 * x * x % P;
    y >>= 1;
  return res;
int r[N];
void ntt(int *x, int lim, int opt) {
  register int i, j, k, m, gn, g, tmp;
  for (i = 0; i < lim; ++i)
    if (r[i] < i) swap(x[i], x[r[i]]);
  for (m = 2; m <= lim; m <<= 1) {</pre>
    k = m \gg 1;
    gn = qpow(3, (P - 1) / m);
    for (i = 0; i < lim; i += m) {
      g = 1;
      for (j = 0; j < k; ++j, g = 111 * g * gn % P) {
        tmp = 111 * x[i + j + k] * g % P;
       x[i + j + k] = (x[i + j] - tmp + P) \% P;
        x[i + j] = (x[i + j] + tmp) \% P;
   }
```

```
if (opt == -1) {
    reverse(x + 1, x + lim);
    register int inv = qpow(lim, P - 2);
    for (i = 0; i < \lim; ++i) x[i] = 111 * x[i] * inv % P;
int A[N], B[N], C[N];
char a[N], b[N];
int main() {
  register int i, lim(1), n;
  scanf("%s", &a);
  n = strlen(a);
  for (i = 0; i < n; ++i) A[i] = a[n - i - 1] - '0';
  while (lim < (n << 1)) lim <<= 1;
  scanf("%s", &b);
  n = strlen(b);
  for (i = 0; i < n; ++i) B[i] = b[n - i - 1] - '0';
  while (lim < (n << 1)) lim <<= 1;
  for (i = 0; i < \lim; ++i) r[i] = (i & 1) * (\lim >> 1) + (r[i >> 1] >> 1);
  ntt(A, lim, 1);
  ntt(B, lim, 1);
  for (i = 0; i < lim; ++i) C[i] = 111 * A[i] * B[i] % P;</pre>
  ntt(C, lim, -1);
  int len(0);
  for (i = 0; i < lim; ++i) {</pre>
   if (C[i] >= 10) len = i + 1, C[i + 1] += C[i] / 10, C[i] \% = 10;
   if (C[i]) len = max(len, i);
  while (C[len] >= 10) C[len + 1] += C[len] / 10, C[len] %= 10, len++;
  for (i = len; ~i; --i) putchar(C[i] + '0');
  puts("");
  return 0;
拉格朗日插值
#include <algorithm>
#include <cstdio>
#include <cstring>
const int maxn = 2010;
using 11 = long long;
11 \mod = 998244353;
11 n, k, x[maxn], y[maxn], ans, s1, s2;
11 powmod(l1 a, l1 x) {
        11 \text{ ret} = 111, \text{ nww} = a;
         while (x)
                  if (x & 1) ret = ret * nww % mod;
                           nww = nww * nww % mod;
                  x >>= 1;
         return ret;
ll inv(ll x) { return powmod(x, mod - 2); }
int main() {
```

```
scanf("%11d%11d", &n, &k);
        for (int i = 1; i <= n; i++) scanf("%lld%lld", x + i, y + i);</pre>
        for (int i = 1; i <= n; i++) {
                  s1 = v[i] \% mod;
                  s2 = 111;
                  for (int j = 1; j <= n; j++)
                          if (i != j)
                                   s1 = s1 * (k - x[j]) % mod, s2 = s2 * ((x[i] - x[j]) %
mod) % mod) % mod;
                  ans += s1 * inv(s2) % mod;
                  ans = (ans + mod) \% mod;
        printf("%1ld\n", ans);
         return 0;
高斯消元
void Gauss() {
        for(int i = 0; i < n; i ++) {</pre>
                  r = i;
                  for(int j = i + 1; j < n; j ++)
                          if(fabs(A[j][i]) > fabs(A[r][i])) r = j;
                  if(r != i) for(int j = 0; j <= n; j ++) std :: swap(A[r][j], A[i][j]);
                  for(int j = n; j >= i; j --) {
                          for(int k = i + 1; k < n; k ++)
                                   A[k][j] -= A[k][i] / A[i][i] * A[i][j];
        for(int i = n - 1; i >= 0; i --) {
                  for(int j = i + 1; j < n; j ++)</pre>
                          A[i][n] -= A[j][n] * A[i][j];
                  A[i][n] /= A[i][i];
        }
组合数学
```

卡特兰数

$$K_n = \frac{C_{2n}^n}{n+1}$$

$$K_0 = K_1 = 1$$

$$K_n = \sum_{i=1}^n K_{i-1} K_{n-i} (n >= 2)$$

$$K_n = \frac{K_{n-1} (4n-2)}{n+1}$$

$$K_n = C_{2n}^n - C_{2n}^{n-1}$$

斯特林数

第二类斯特林数

将 n 个两两不同的元素划分为 k 个互不区分的非空子集方案数

$$S(n,k) = S(n-1,k-1) + k * S(n-1,k)$$

$$S(n,0) = [n=0]$$

$$S(n,m) = \sum_{i=0}^{m} \frac{(-1)^{m-i}i^{n}}{i!(m-i)!}$$

$$S[0][0] = 1;$$

 $FOR (i, 1, N)$
 $FOR (j, 1, i + 1) S[i][j] = (S[i - 1][j - 1] + j * S[i - 1][j]) % MOD;$

反演

二项式反演

fn 为至多, gn 为恰好的方案数

$$f_n = \sum_{i=0}^n C_n^i g_i$$

$$g_n = \sum_{i=0}^{n} (-1)^{n-i} C_n^i f_i$$

或者:

fn 为恰好, gi 为至少

$$g_k = \sum_{i=k}^n C_n^i f$$

$$g_{k} = \sum_{i=k}^{n} C_{n}^{i} f_{i}$$

$$f_{k} = \sum_{i=0}^{n} (-1)^{i-k} C_{i}^{k} g_{i}$$

min_max 反演

min{S}为集合最小值, max{S}为集合最大值

$$\max\{S\} = \sum_{T \subseteq S, T = 0} (-1)^{|T|-1} \min\{T\}$$

牛成函数

普通牛成函数

常用生成函数的开放,收敛转化:

$$1 + px + p^{2}x^{2} + \dots = \frac{1}{1 - px}$$

$$\sum_{n=0}^{\infty} C_{m}^{n} x^{n} = (1 + x)^{m}$$

$$\sum_{n=0}^{\infty} C_{n+m}^{n} x^{n} = \frac{1}{(1 - x)^{m+1}}$$

自然数幂和表

```
MP = {
    0:"1 1 0",
    1:"2 1 1 0",
    2:"6 2 3 1 0",
    3:"4 1 2 1 0 0",
    4:"30 6 15 10 0 -1 0",
    5:"12 2 6 5 0 -1 0 0",
    6:"42 6 21 21 0 -7 0 1 0",
    7:"24 3 12 14 0 -7 0 2 0 0"
    8: "90 10 45 60 0 -42 0 20 0 -3 0",
    9:"20 2 10 15 0 -14 0 10 0 -3 0 0",
    10:"66 6 33 55 0 -66 0 66 0 -33 0 5 0",
    11:"24 2 12 22 0 -33 0 44 0 -33 0 10 0 0",
    12:"2730 210 1365 2730 0 -5005 0 8580 0 -9009 0 4550 0 -691 0",
    13: "420 30 210 455 0 -1001 0 2145 0 -3003 0 2275 0 -691 0 0",
    14:"90 6 45 105 0 -273 0 715 0 -1287 0 1365 0 -691 0 105 0",
    15: "48 3 24 60 0 -182 0 572 0 -1287 0 1820 0 -1382 0 420 0 0",
    16:"510 30 255 680 0 -2380 0 8840 0 -24310 0 44200 0 -46988 0 23800 0 -3617 0".
    17:"180 10 90 255 0 -1020 0 4420 0 -14586 0 33150 0 -46988 0 35700 0 -10851 0 0",
    18:"3990 210 1995 5985 0 -27132 0 135660 0 -529074 0 1469650 0 -2678316 0 2848860 0
-1443183 0 219335 0".
    19:"840 42 420 1330 0 -6783 0 38760 0 -176358 0 587860 0 -1339158 0 1899240 0 -14431
83 0 438670 0 0",
    20:"6930 330 3465 11550 0 -65835 0 426360 0 -2238390 0 8817900 0 -24551230 0 4476780
0 0 -47625039 0 24126850 0 -3666831 0"
}
预处理组合数
LL C[M][M];
void init C(int n) {
        FOR (i, 0, n) {
                 C[i][0] = C[i][i] = 1;
                 FOR (j, 1, i)
                          C[i][j] = (C[i - 1][j] + C[i - 1][j - 1]) \% MOD;
```

OTHER

```
BM 线性递推
#include<bits/stdc++.h>
using namespace std;
#define rep(i,a,n) for (int i=a;i<n;i++)</pre>
#define per(i,a,n) for (int i=n-1;i>=a;i--)
#define pb push back
#define mp make pair
#define all(x) (x).begin(),(x).end()
#define fi first
#define se second
#define SZ(x) ((int)(x).size())
typedef vector<int> VI;
typedef long long 11;
typedef pair<int,int> PII;
const 11 mod=10000000007;
ll powmod(ll a,ll b) {ll res=1;a%=mod; assert(b>=0); for(;b;b>>=1){if(b&1)res=res*a%mod;
a=a*a%mod;}return res;}
// head
11 n:
namespace linear_seq {
    const int N=10010:
    11 res[N],base[N],_c[N],_md[N];
    vector<int> Md;
    void mul(ll *a,ll *b,int k) {
         rep(i,0,k+k) _c[i]=0;
         rep(i,0,k) if (a[i]) rep(j,0,k) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
         for (int i=k+k-1; i>=k; i--) if (c[i])
              rep(j,0,SZ(Md)) \ \_c[i-k+Md[j]] = (\_c[i-k+Md[j]] - \_c[i] * \_md[Md[j]]) %mod;
         rep(i,0,k) a[i]=c[i];
    int solve(ll n, VI a, VI b) { // a 系数 b 初值 b[n+1]=a[0]*b[n]+...
         11 ans=0, pnt=0;
         int k=SZ(a);
         assert(SZ(a)==SZ(b));
         rep(i,0,k) _md[k-1-i]=-a[i];_md[k]=1;
         Md.clear();
         rep(i,0,k) if (_md[i]!=0) Md.push_back(i);
         rep(i,0,k) res[i]=base[i]=0;
         res[0]=1;
         while ((111<<pnt)<=n) pnt++;</pre>
         for (int p=pnt;p>=0;p--) {
              mul(res,res,k);
              if ((n>>p)&1) {
                   for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
                  \label{eq:continuous} \begin{split} \text{rep}(\texttt{j},\emptyset,\mathsf{SZ}(\texttt{Md})) \ \ \text{res}[\texttt{Md}[\texttt{j}]] = & (\texttt{res}[\texttt{Md}[\texttt{j}]] - \texttt{res}[\texttt{k}] * \_\texttt{md}[\texttt{Md}[\texttt{j}]]) \% \texttt{mod}; \end{split}
         rep(i,0,k) ans=(ans+res[i]*b[i])%mod;
         if (ans<0) ans+=mod;</pre>
         return ans;
    VI BM(VI s) {
         VI C(1,1), B(1,1);
```

int L=0, m=1, b=1;

```
rep(n,0,SZ(s)) {
            11 d=0;
            rep(i,0,L+1) d=(d+(ll)C[i]*s[n-i])%mod;
            if (d==0) ++m;
            else if (2*L<=n) {
               VI T=C;
                11 c=mod-d*powmod(b,mod-2)%mod;
                while (SZ(C) < SZ(B) + m) C.pb(0);
                rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
                L=n+1-L; B=T; b=d; m=1;
            } else {
                11 c=mod-d*powmod(b,mod-2)%mod;
                while (SZ(C)<SZ(B)+m) C.pb(∅);</pre>
                rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
                ++m;
        return C;
    int gao(VI a,ll n) {
       VI c=BM(a);
       c.erase(c.begin());
       rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;
       return solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
};
int main() {
   /*push back 进去前 8~10 项左右、最后调用 qao 得第 n 项*/
   vector<int>v:
   v.push back(3);
   v.push_back(9);
   v.push back(20);
   v.push back(46);
    v.push_back(106);
    v.push back(244);
   v.push back(560);
   v.push back(1286);
   v.push_back(2956);
   v.push back(6794);
    int nCase:
    scanf("%d", &nCase);
    while(nCase--){
       scanf("%11d", &n);
       printf("%lld\n",1LL * linear seg::gao(v,n-1) % mod); ///求第n 项
常用宏及函数与快读
// v2021.5.22 主席树更面向对象
#include <ext/pb ds/tree policy.hpp>
#include <ext/pb ds/assoc container.hpp>
__gnu_pbds::tree<int, __gnu_pbds::null_type, std::less<int>, __gnu_pbds::rb_tree_tag, __
gnu_pbds::tree_order_statistics_node_update> TTT;
```

```
// 函数不返回值可能会 RE
// 少码大数据结构,想想复杂度更优的做法
// 小数 二分/三分 注意 break 条件
// 浮点运算 sqrt(a^2-b^2) 可用 sqrt(a+b)*sqrt(a-b) 代替,避免精度问题
// Long double -> %Lf 别用C11 (C14/16)
// 控制位数 cout << setprecision(10) << ans;
// reverse vector 注意判空 不然会re
// 分块注意维护块上标记 来更新块内数组 a[]
// vector+Lower bound 常数 < map/set/(unordered map)
// map.find 不会创建新元素 map[]会 注意空间
// 别对指针用 memset
// 用位运算表示 2^n 注意加 LL 1LL<<20
// 注意递归爆栈
// 注意边界
// 注意memset 多组会T
// Lambda
// sort(p + 1, p + 1 + n,
              [](const point &x, const point &y) -> bool { return x.x < y.x; });
// append l1 to l2 (l1 unchanged)
// L2.insert(L2.end(), L1.begin(), L1.end());
// append L1 to L2 (elements appended to L2 are removed from L1)
// (general form ... TG gave form that is actually better suited
// for your needs)
// L2.splice(L2.end(), L1, L1.begin(), L1.end());
//位运算函数
//int builtin ffs (unsigned int x 最后一位1 的是从后向前第几位,1110011001000 返回4
//int builtin clz (unsigned int x)前导0 个数
//int __builtin_ctz (unsigned int x)末尾0 个数
//int builtin popcount (unsigned int x) 1 的个数
//此外,这些函数都有相应的 usigned Long 和 usigned Long 版本,只需要在函数名后面加上 L 或 LL
就可以了,比如int builtin clzll。
//java 大数
//import java.io.*;
//import java.math.BigInteger;
//import java.util.*;
//public class Main {
    public static void main(String args[]) throws Exception {
         Scanner cin=new Scanner(System.in);
//
         BigInteger a:
//
         BigInteger b;
         a = cin.nextBigInteger();
        b = cin.nextBigInteger();
//
//
         System.out.println(a.add(b));
//}
//
```

```
//生成超过32767 的随机数
//unsigned seed = std::chrono::system_clock::now().time_since_epoch().count();
// mt19937 rand_num(seed); // 大随机数
// uniform_int_distribution<long long> dist(0, 1000000000); // 给定范围
// cout << dist(rand_num) << endl;
```

常见博弈

巴什博弈

只有一堆 \mathbf{n} 个物品,两个人轮流从这堆物品中取物,规定每次至少取一个,最多 \mathbf{n} 个。最后取光者得胜。

n% (m+1) ==0 必败, 否则必胜

威佐夫博弈

有两堆各若干个物品,两个人轮流从任意一堆中取出至少一个或者同时从两堆中取出同样多的物品,规定每次至少取一个,至多不限,最后取光者胜利。设两堆分别为 n 和 m

较小堆*两堆之差为(黄金分割比+1)时必败,否则比胜

nim 博弈

有若干堆各若干个物品,两个人轮流从某一堆取任意多的物品,规定每次至少取一个,多者不限,最后 取光者得胜。

各堆物品异或为0时必败,否则必胜

anti-nim 博弈

有若干堆各若干个物品,两个人轮流从某一堆取任意多的物品,规定每次至少取一个,多者不限,最后 取光者得败。

先手胜当且仅当 ①所有堆石子数都为 1 且游戏的 SG 值为 0 (即有偶数个孤单堆-每堆只有 1 个石子数); ②存在某堆石子数大于 1 且游戏的 SG 值不为 0.

阶梯博弈

地面表示第 0 号阶梯。每次都可以将一个阶梯上的石子向其左侧移动任意个石子,没有可以移动的空间时(及所有石子都位于地面时)输。

阶梯博弈等效为奇数号阶梯的尼姆博弈

对拍

check.cpp (Windows) #include <windows.h> #include <bits/stdc++.h> using namespace std; int main() system("g++ data.cpp -o data --std=c++17"); system("g++ std.cpp -o std --std=c++17"); system("g++ test.cpp -o test --std=c++17"); int t = 10000: while (t--) system("data.exe > data.txt"); clock t st=clock(); system("test.exe < P3372_8.in > test.out"); clock t end=clock(); system("std.exe < data.txt > std.txt"); if (system("fc P3372 8.out test.out")) t=-1: break: cout<<"TIME: "<<end-st<<" ms\n\n";</pre> **if** (t == 0) cout << "Accepted!" << endl;</pre> cout << "Wrong Answer!" << endl;</pre> return 0; } check.cpp (Linux) 重点!数据比较器 #include <bits/stdc++.h> using namespace std; int main() system("g++ ./data.cpp -o data --std=c++17"); system("g++ ./std.cpp -o std --std=c++17"); system("g++ ./test.cpp -o test --std=c++17"); int t = 10000: while (t--) system("./data.exe > ./data.txt"); clock_t st = clock(); system("./test.exe < ./data.txt > ./test.txt"); clock t end = clock(); system("./std.exe < ./data.txt > ./std.txt"); if (system("diff ./std.txt ./test.txt")) t=-1; break: cout << "TIME: " << end - st << " ms\n\n";</pre> **if** (t == 0) cout << "Accepted!" << endl;</pre> else

100000000000000049 100000000000000177

- cout << "Wrong Answer!" << endl;
 return 0;
 }</pre>
 - data.cpp 生成数据
 - std.cpp 暴力程序
 - test.cpp 需确认正确性

质数表

1e2	1e3	1e4	1e5	1e6
101	1009	10007	100003	1000003
103	1013	10009	100019	1000033
107	1019	10037	100043	1000037
109	1021	10039	100049	1000039
113	1031	10061	100057	1000081
127	1033	10067	100069	1000099
131	1039	10069	100103	1000117
137	1049	10079	100109	1000121
139	1051	10091	100129	1000133
149	1061	10093	100151	1000151
151	1063	10099	100153	1000159

1e7	1e8	1e10	1e11	1e12
10000019	100000007	1000000007	1000000019	100000000003
10000079	100000037	1000000009	10000000033	100000000019
10000103	100000039	1000000021	10000000061	100000000057

1e13	1e14	1e15	1e16
1000000000037	100000000000031	1000000000000037	100000000000000061
10000000000051	100000000000067	10000000000000091	10000000000000069
10000000000099	100000000000097	1000000000000159	10000000000000079
1000000000129	100000000000099	1000000000000187	10000000000000099
1000000000183	100000000000133	1000000000000223	10000000000000453

1e17	1e18
100000000000000003	10000000000000000003
100000000000000013	10000000000000000009
100000000000000019	10000000000000000031
1000000000000000021	1000000000000000079