

M1-IRELE

**ELEC-H401 Modulation and coding**

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# DVB-C project

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## Introduction

This report aims to complete the code that simulates a DVB-C transmission chain in matlab. It provides additional information from the theoretical part of the project.

The first part builds the transmission chain and link the blocks together such that the received signal is the same as the one transmitted (in a noiseless case).

Second part to be explained later

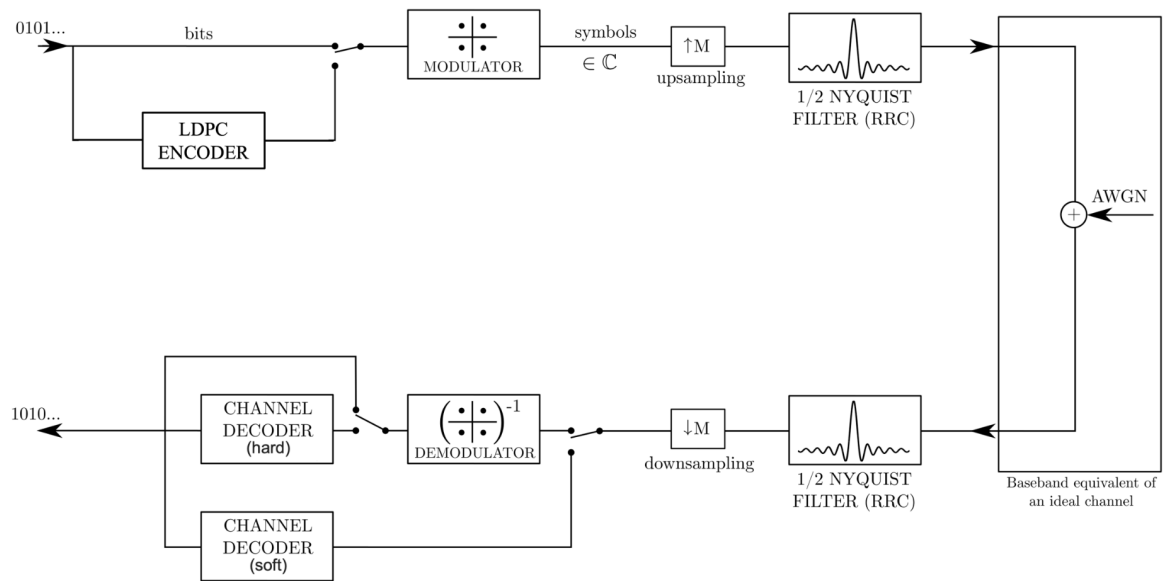


Figure 1: DVB-C transmission chain

### Baseband representation

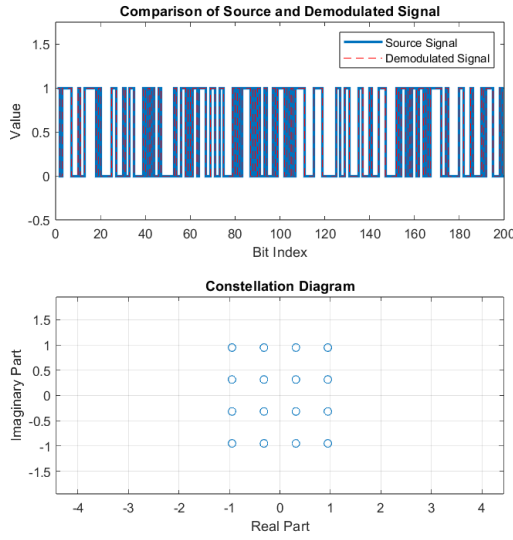
By looking at the block diagram of the transmission chain 1, one can see we never move the baseband signal to the carrier frequency. As the simulation runs on a computer, using the bandpass representation of the signal would require much more samples as the sampling frequency would need to be at least twice the carrier frequency. By simulating the chain in baseband, the minimal sampling frequency is reduced to the symbol rate in order to have at least one sample per symbol.

Because the signal is oversampled, the sampling frequency is then equal to the symbol rate multiplied by the oversampling factor.

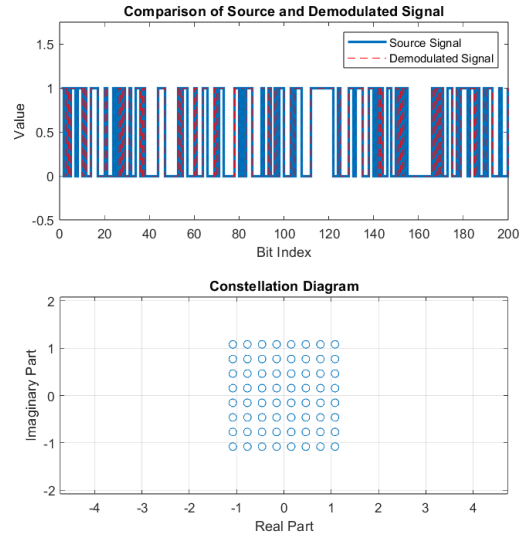
### Modulation and Demodulation

After generating  $N$  random bits, they are modulated. This allows to send fewer symbols than the number of bits. We chose QAM modulation as it combines ASK and PSK. Depending on the number of bits per symbol ( $N_{\text{bps}}$ ), the number of bits sent ( $N$ ) had to be chosen such that  $N/N_{\text{bps}} \in \mathbb{N}$ .

Figure 2 compares the constellation diagrams obtained for QAM-16 and QAM-64. As the constellations points are more spaced on the left, QAM-16 is less prone to a wrong demodulation (when noise will be added). This comes at the cost of a lower bitrate: for the same symbol rate, QAM-64 will send 6 bits while QAM-16 only send 4. It clearly shows a compromise between reliability and capacity.



(a) QAM-16 modulation



(b) QAM-64 modulation

Figure 2: Comparison of QAM modulations, where the mean square error is computed between the transmitted and received bitstream

## Optimal demodulator and detector

First, it is important to remind that the transmitted signal is represented by a set of coefficients which results from the projection of the signal on an orthonormal basis related to the chosen modulation. Once transmitted, the signal is affected by noise (AWG noise here). In the general case, this noise moves the signal, resulting in a change in its coefficient in the basis, as can be seen on Figure 3

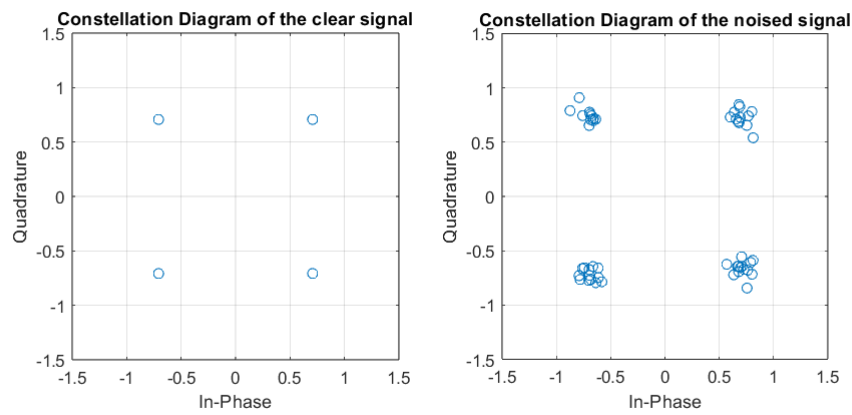


Figure 3: Effect of noise on the constellation diagram

To construct an optimal demodulator, 2 criteria should be taken into account. The first one is the sufficient statistic criteria. It is proven that once the received signal is projected on the sub-space defined by the previous basis functions, the noise component outside of the sub-space is independent from the pro-

jected signal. It means that there is no information loss when projecting the received signal on the original sub-space and the optimal decision can therefore be made using the projection of the incoming signal.

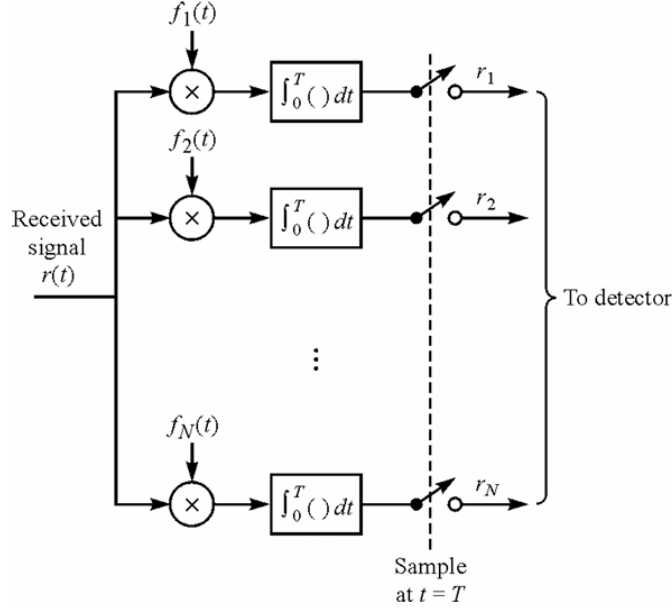


Figure 4: Projection on basis sub-space

$$\begin{aligned}
 E[n'(t)r_k] &= E[n'(t)]s_{mk} + E[n'(t)n_k] \\
 &= E[n'(t)n_k] \\
 &= E\left[\left(n(t) - \sum_{l=1}^K n_l f_l(t)\right) n_k\right] \\
 &= \int_{t'=-\infty}^{\infty} E[n(t)n(t')] f_k(t') dt' - \sum_{l=1}^K E[n_l n_k] f_l(t) \\
 &= \frac{N_0}{2} f_k(t) - \frac{N_0}{2} f_k(t) = 0
 \end{aligned}$$

Figure 5: Criteria of sufficient statistics

The second criteria is the usage of matched filters. The demodulator used to achieve the sufficient statistic property is composed of a bank of correlators (projection on a basis function). Instead of using a bank of correlators, we can use a bank of filters matched to the basis functions of the chosen modulation. It is proven that such filters at the demodulator gives a maximized SNR (minimize the power of the noise at the exit of the demodulator).

In conclusion, by using a bank of filters matched on the orthonormal basis function set by the choice of the modulation, we can construct a optimal demodulator which will ensure an optimal decision based on the received signal and ensure a maximum SNR at the output of this demodulator.

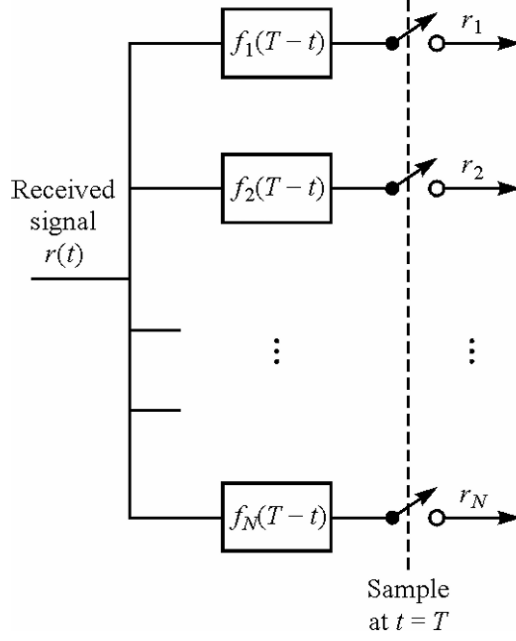


Figure 6: Bank of filters matching basis functions

$$\text{SNR} := \frac{\left[ \int_{\tau=-\infty}^{\infty} s(\tau) h(-\tau) d\tau \right]^2}{\frac{N_0}{2} \int_{\tau=-\infty}^{\infty} h^2(-\tau) d\tau}$$

Cauchy-Schwartz inequality:

$$\left[ \int_{t=-\infty}^{\infty} g_1(t) g_2(t) dt \right]^2 \leq \int_{t=-\infty}^{\infty} g_1^2(t) dt \int_{t=-\infty}^{\infty} g_2^2(t) dt$$

Equality if  $g_1(t) = k g_2(t)$  for any arbitrary constant  $k$

Figure 7: Maximum SNR demonstration

At the output of the demodulator, we still need to make sure the optimal choice of the  $M$  possible  $s_m(t)$  is made based on the received signal. To achieve this, the maximum likelihood criteria is used. It is equivalent to the maximum a posteriori criteria (general criteria) in the case of equiprobable symbols  $s_m(t)$ . The criteria leads to the following result : the optimizal  $s_m(t)$  choice is found by taking the minimum euclidian distance between the observable received signal  $r(t)$  and all the possible modulated signal  $s_m(t)$ .

## Pulse shaping

With modulation only, the bandwidth of the transmitted signal is infinite. This is problematic as it could interfere with neighboring channels. A filtering is applied to resolve this but the chosen filter must respect two other constraints: it must cancel inter-symbol interference (ISI) and must maximize the SNR.

The raised cosine filter is chosen as it limits the bandwidth and cancels ISI. To maximize the SNR, it is applied as a matched filter by using the square root of it at the transmitter and at the receiver.

The time domain and frequency domain representation of the raised cosine filter is shown in Figure 8. Figure 9 shows how the signal is shaped in the time domain and how there is indeed no ISI. Finally, the spectrum of the transmitted signal is plotted in figure 10 where the frequency band is limited to  $[-3, 3]$  MHz.

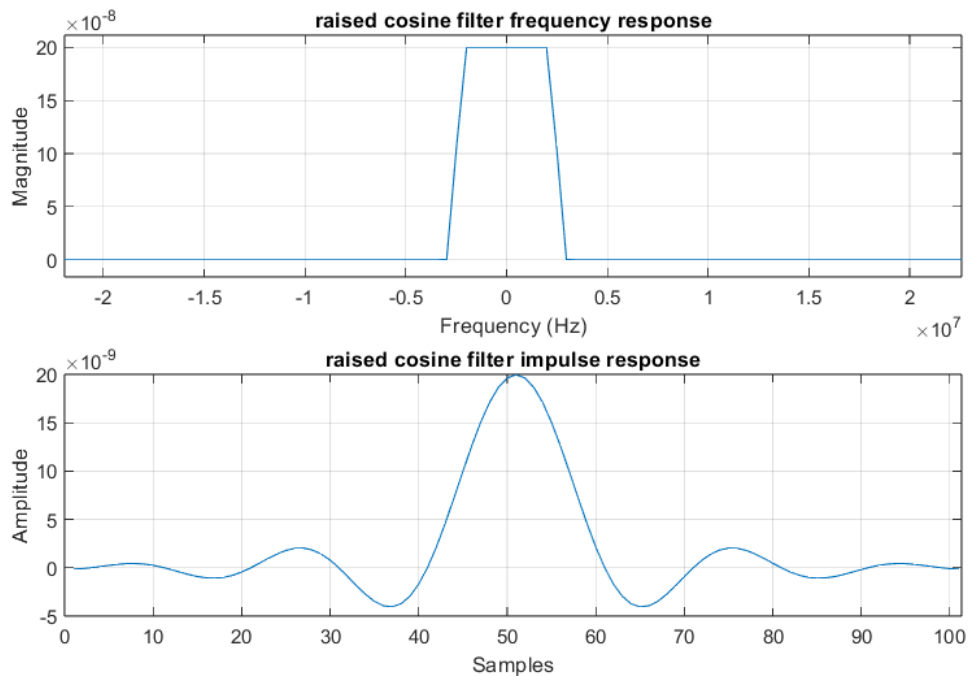


Figure 8: Time and frequency domain representation of the raised cosine filter



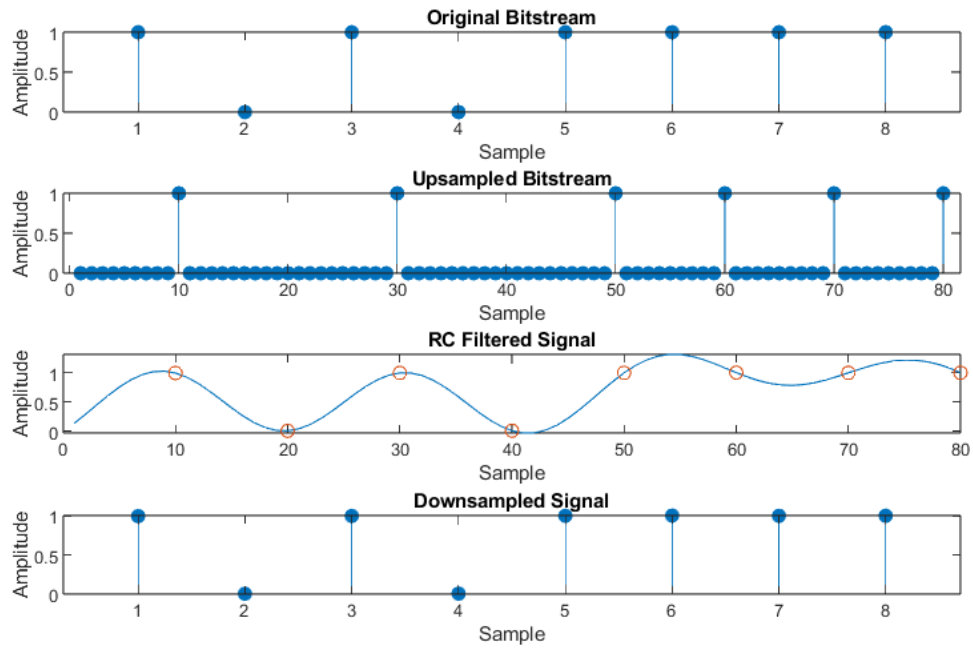


Figure 9: Pulse shaping with a raised cosine filter

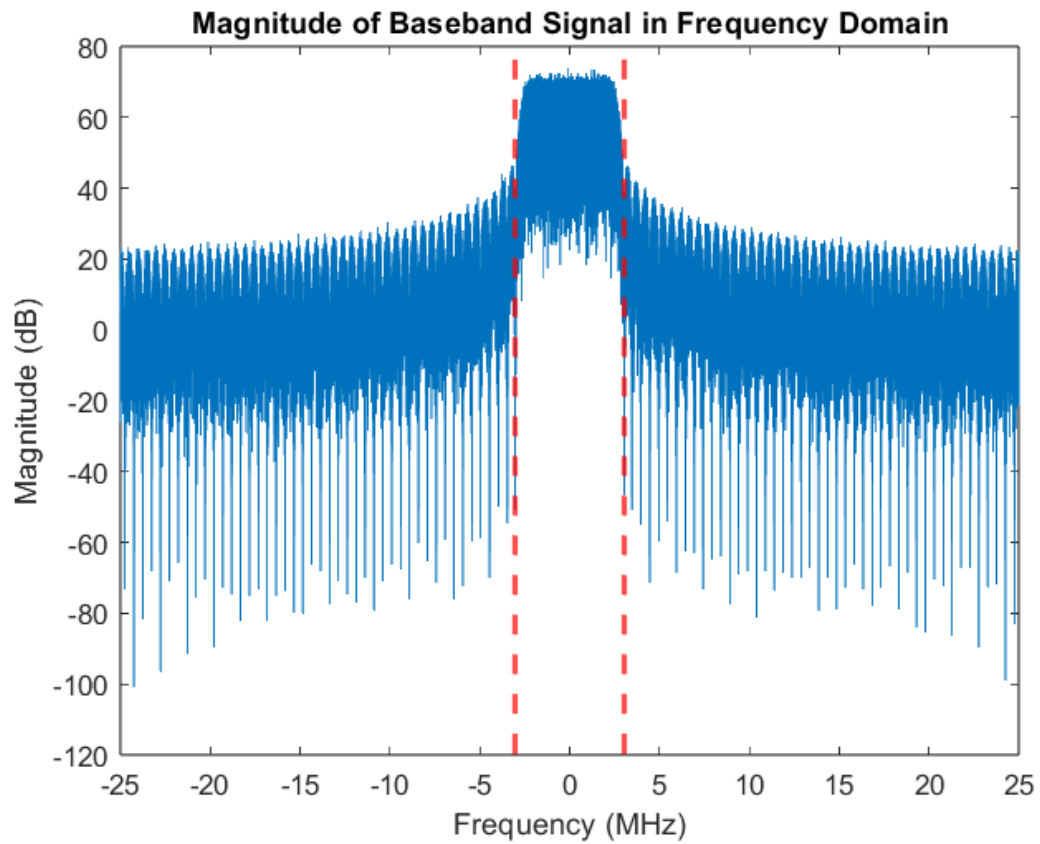


Figure 10: Spectrum of the transmitted signal after pulse shaping

## Noise addition

The last building block is a noise source. It generates additive white Gaussian noise in baseband. When the signal is too much corrupted, the demodulation can fail. The BER curves are plotted in figure 11 and they show the impact of the noise power  $N_0$  on the bit error rate. The compromise between reliability and capacity is again visible: in the same conditions (same  $E_b/N_0$ ), a modulation with lower capacity will have a smaller BER.

The theoretical BER curves are plotted on figure 11 and are compared with the simulation results. They stay close to each other until the BER reaches  $10^{-4}$ . This limit could go even lower by increasing the number of bits sent but we limited it to  $10^6$  in order for the code to run quite fast.

To impose a value of  $E_b/N_0$ , we start by computing the energy of the transmitted signal before adding the noise. The power of the noise is then chosen as  $N_0 = E_b / (E_b/N_0)_{\text{desired}}$ .

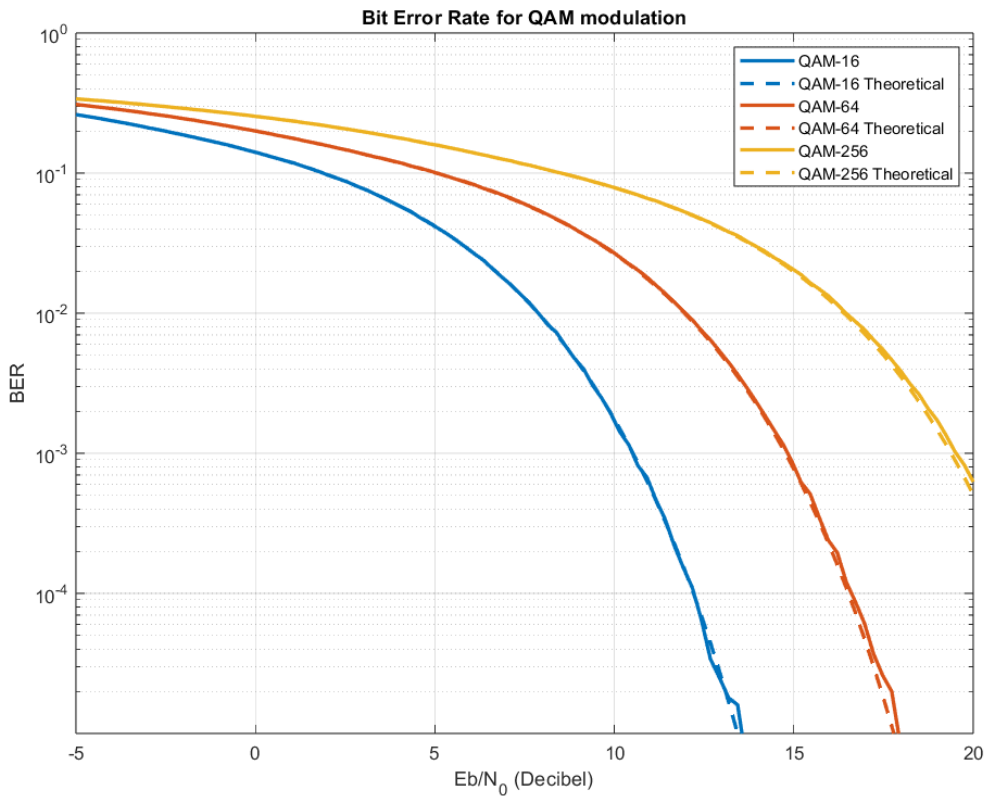


Figure 11: BER curves for different QAM modulations

## Bit rate

Considering the following characteristics :

- (a) Physical bandwidth of  $B_{\text{physical}} = 6 \text{ MHz}$
- (b) Roll-off factor  $\beta = 0.2$
- (c) QAM 16 modulation  $\rightarrow \text{Nbps} = 4$

We can derive the symbol rate from the physical bandwidth

$$f_{\text{symbol}} = \frac{B_{\text{physical}}}{1 + \beta} = 5 \text{ MHz}$$

$$\text{Bit-rate} = f_{\text{symbol}} * \text{Nbps} = 5 \text{ MHz} * 4 = 20 \text{ MBps}$$