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1 Contest
                                                              1
2 Data structures
3 Matching
4 Math
Contest (1)
template.cpp
#include <bits/stdc++.h>
using namespace std;
#define int long long
#define endl '\n'
#define rep(i, a, b) for(int i = (a); i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
#define debug(var) cerr << #var << ": " << var << endl
#define pb push back
#define eb emplace_back
typedef long long 11;
typedef pair<int, int> pii;
typedef vector<int> vi;
void solve() {
int32 t main() {
  ios_base::sync_with_stdio(0); cout.tie(0); cin.tie(0);
  int t = 1;
  while(t--) solve();
 return 0;
.bashrc
                                                          3 lines
alias c='q++ -Wall -Wconversion -Wfatal-errors -q -std=c++17 \
  -fsanitize=undefined,address
xmodmap -e 'clear lock' -e 'keycode 66=less greater' \#caps = \Leftrightarrow
.vimrc
                                                           4 lines
set cin ai is ts=4 sw=4 nu rnu
" Select a region and then type : Hash
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \
\| md5sum \| cut -c-6
hash.sh
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed.
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum |cut -c-6
troubleshoot.txt
Write a few simple test cases if sample is not enough.
Are time limits close? If so, generate max cases.
Is the memory usage fine?
Could anything overflow?
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Make sure to submit the right file.

```
Wrong answer:
Print your solution! Print debug output, as well.
Are you clearing all data structures between test cases?
Can your algorithm handle the whole range of input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Have you understood the problem correctly?
Any uninitialized variables?
Any overflows?
Confusing N and M, i and j, etc.?
Are you sure your algorithm works?
What special cases have you not thought of?
Are you sure the STL functions you use work as you think?
Add some assertions, maybe resubmit.
Create some testcases to run your algorithm on.
Go through the algorithm for a simple case.
Go through this list again.
Explain your algorithm to a teammate.
Ask the teammate to look at your code.
Go for a small walk, e.g. to the toilet.
Is your output format correct? (including whitespace)
Rewrite your solution from the start or let a teammate do it.
Runtime error:
Have you tested all corner cases locally?
Any uninitialized variables?
Are you reading or writing outside the range of any vector?
Any assertions that might fail?
Any possible division by 0? (mod 0 for example)
Any possible infinite recursion?
Invalidated pointers or iterators?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see Various).
Time limit exceeded:
Do you have any possible infinite loops?
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (References)
How big is the input and output? (consider scanf)
Avoid vector, map. (use arrays/unordered_map)
What do your teammates think about your algorithm?
Memory limit exceeded:
What is the max amount of memory your algorithm should need?
Are you clearing all data structures between test cases?
Data structures (2)
SegTree.h
Description: Iterative SegTree Can be changed by modifying Spec
Time: \mathcal{O}(\log N)
                                                     0386e5, 37 lines
template<typename Spec>
struct SegTree {
  using LS = Spec;
  using S = typename LS::S;
  using K = typename LS::K;
  int n;
  vector<S> seq;
  SegTree(const vector<S> & v)
    : n(sz(v)), seq(2*n) {
    rep(i, 0, n) seg[i+n] = v[i];
    for (int i = n-1; i \ge 1; i--) seq[i] = LS::op(seq[i*2], seq
         [i*2+1]);
  void update(int no, K val) {
```

```
no += n;
    seg[no] = LS::update(val, seg[no]);
    while (no > 1) no /= 2, seg[no] = LS::op(seg[no\star2], seg[no
         *2+1]);
 S query(int 1, int r) { // [l, r)
    S vl = LS::id(), vr = LS::id();
    for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
      if (1 & 1) v1 = LS::op(v1, seg[1++]);
      if (r \& 1) vr = LS::op(seq[--r], vr);
    return LS::op(vl, vr);
};
struct Spec {
 using S = int;
 using K = int;
  static S op(S a, S b) { return max(a, b); }
  static S update(K f, S a) { return f + a; }
  static S id() { return 0; }
LazvSeg.h
Description: Iterative Lazy SegTree Can be changed by modifying Spec
Time: \mathcal{O}(\log N)
template<typename Spec>
struct LazySeq {
  using LS = Spec;
  using S = typename LS::S;
  using K = typename LS::K;
  int n;
  vector<S> seg:
  vector<K> lazv:
  vector<bool> has lazy;
  // vector<int> lx, rx; // Aditional info
    LazySeg(vector<S> & v) : n(sz(v)), seg(2*n), lazy(n),
        has_lazy(n) {
    rep(no, 0, n) seg[no+n] = v[no];
    for (int no = n-1; no >= 1; no--) pull(no);
    // Aditional info, n must be power of two
    lx.assign(2*n, 0); rx.assign(2*n, 0);
    lx/1/ = 0; rx/1/ = n;
    rep(no, 1, n) {
      int \ mid = (lx[no] + rx[no])/2;
      lx[no*2] = lx[no]; rx[no*2] = mid;
      lx[no*2+1] = mid; rx[no*2+1] = rx[no];
 S query(int 1, int r) { // [l, r)
   1 += n;
    r += n;
    push_to(1); push_to(r-1);
    S vl = LS::id(), vr = LS::id();
    while(1 < r) {
      if (1 & 1) v1 = LS::op(v1, seg[1++]);
      if (r & 1) vr = LS::op(seg[--r], vr);
     1 >>= 1; r >>= 1;
    return LS::op(vl, vr);
```

```
void update(int 1, int r, K val) {
   1 += n;
    r += n;
    push_to(1); push_to(r-1);
    int lo = 1, ro = 1;
    while(1 < r)  {
     if (1 & 1) lo = max(lo, 1), apply(l++, val);
     if (r \& 1) ro = max(ro, r), apply(--r, val);
     1 >>= 1; r >>= 1;
    pull_from(lo);
   pull_from(ro-1);
  void apply(int no, K val) {
    seg[no] = LS::update(val, seg[no]);
    // seg[no] = LS::update(val, seg[no], lx[no], rx[no]);
    if (no < n) {
      if (has_lazy[no]) lazy[no] = LS::compose(val, lazy[no]);
      else lazy[no] = val;
     has_lazy[no] = true;
  void pull_from(int no) {
    while(no > 1) no >>= 1, pull(no);
  void pull(int no) {
    seg[no] = LS::op(seg[no*2], seg[no*2+1]);
  void push_to(int no) {
   int h = 0; int p2 = 1;
    while (p2 < no) p2 \star= 2, h++;
    for (int i = h; i >= 1; i--) push(no >> i);
  void push(int no) {
    if (has_lazy[no]) {
      apply(no*2, lazy[no]);
      apply(no*2+1, lazy[no]);
      has_lazy[no] = false;
};
struct Spec {
  using S = int;
  using K = int;
  static S op(S a, S b) { return max(a, b); }
  static S update(K f, S a) { return f + a; }
  static K compose(const K f, const K q) { return f + q; }
  static S id() { return 0; }
Description: Multidimensional Psum Requires Abelian Group S (op, inv,
Memory: \mathcal{O}\left(N^{D}\right)
Time: \mathcal{O}(1)
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct Psum{ using T = typename S::T;
```

vector<Psum<D-1, S>> v;

```
MAs Psum(int s, As... ds):n(s+1),v(n,Psum<D-1, S>(ds...)){}
  MAs void set (T x, int p, As... ps) {v[p+1].set(x, ps...);}
  void push(Psum& p) {rep(i, 1, n)v[i].push(p.v[i]);}
  void init(){rep(i, 1, n)v[i].init(),v[i].push(v[i-1]);}
 MAs T query(int 1, int r, As... ps) {
    return S::op(v[r+1].query(ps...), S::inv(v[1].query(ps...)))
};
template<class S>
struct Psum<0, S>{ using T = typename S::T;
 T val=S::id;
 void set(T x) {val=x;}
  void push(Psum& a) {val=S::op(a.val,val);}
  void init(){}
 T query() {return val;}
};
struct G{
 using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return a+b;}
 static T inv(T a) {return -a;}
MultiDSegTree.h
Description: Multidimensional SegTree Requires Monoid S (op, id)
Memory: \mathcal{O}\left(N^{D}\right)
Time: \mathcal{O}\left((\log N)^D\right)
                                                       53621d, 37 lines
//#pragma once
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct SegTree{ using T = typename S::T;
  vector<SegTree<D-1, S>> seg;
 MAs SegTree(int s, As... ds):n(s),seg(2*n, SegTree<D-1, S>(ds
      ...)){}
  MAs T get(int p, As... ps){return seg[p+n].get(ps...);}
 MAs void update (T x, int p, As... ps) {
    p+=n; seg[p].update(x, ps...);
    for (p>>=1; p>=1; p>>=1)
    seg[p].update(S::op(seg[2*p].get(ps...),seg[2*p+1].get(ps
         ...)), ps...);
 MAs T query(int 1, int r, As... ps) {
    T lv=S::id,rv=S::id;
    for (1+=n, r+=n+1; 1<r; 1>>=1, r>>=1) {
      if (1&1) lv = S::op(lv, seq[l++].query(ps...));
      if (r&1) rv = S::op(seg[--r].query(ps...), rv);
    return S::op(lv,rv);
};
template<class S>
struct SegTree<0, S>{ using T = typename S::T;
 T val=S::id;
 T get() {return val;}
 void update(T x) {val=x;}
 T query() {return val;}
struct M{ //monoid
  using T = int;
  static constexpr T id = 0;
```

```
static T op(T a, T b) {return max(a,b);}
SparseTable.h
Description: Multidimensional Sparse Table Requires Idempotent Monoid
S (op, inv, id)
Memory: \mathcal{O}\left((n\log n)^D\right)
Time: \mathcal{O}(1) query, \mathcal{O}((n \log n)^D) build
                                                         c900f0, 39 lines
#define MAs template<class...As> //multiple arguments
template<int D, class S>
struct SpTable{ using T = typename S::T;
 using isp = SpTable<D-1, S>;
 inline int lg(signed x) {return __builtin_clz(1) -__builtin_clz
  int n;
  vector<vector<isp>> tab;
  MAs SpTable(int s, As... ds):n(s),
  tab(1+lg(n), vector < isp > (n, isp(ds...))) {}
  MAs void set(T x, int p, As... ps){tab[0][p].set(x, ps...);}
  void join(SpTable& a, SpTable& b) {
    rep(i, 0, 1+lg(n))rep(j, 0, n)
      tab[i][j].join(a.tab[i][j], b.tab[i][j]);
  void init(){
    rep(i, 0, n)tab[0][i].init();
    rep(i, 0, lg(n)) rep(j, 0, n-(1 << i))
      tab[i+1][j].join(tab[i][j], tab[i][j+(1<<i)]);
  MAs T query(int 1, int r, As... ps) {
    int k = lq(r-l+1); r+=1-(1<< k);
    return S::op(tab[k][1].query(ps...),tab[k][r].query(ps...))
};
template<class S>
struct SpTable<0, S>{ using T = typename S::T;
 T val=S::id;
  void set(T x){val=x;}
  void join(SpTable& a, SpTable& b) {val=S::op(a.val,b.val);}
 void init(){}
 T query() {return val;}
};
struct IM{
 using T = int;
 static constexpr T id = 0;
 static T op(T a, T b) {return max(a, b);}
};
Description: Multidimensional BIT/Fenwick Tree Requires Abelian Group
"S" (op, inv, id)
Memory: \mathcal{O}\left(N^{D}\right)
Time: \mathcal{O}\left((\log N)^D\right)
                                                         778135, 31 lines
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct BIT{ using T = typename S::T;
 vector<BIT<D-1, S>> bit;
 MAs BIT(int s, As... ds):n(s), bit(n+1, BIT<D-1, S>(ds...)){}
 inline int lastbit(int x){return x&(-x);}
 MAs void add(T x, int p, As... ps) {
    for (p++; p<=n; p+=lastbit (p) ) bit [p] .add (x, ps...);</pre>
```

```
MAs T query(int 1, int r, As... ps) {
   T lv=S::id, rv=S::id; r++;
    for(;r>=1;r-=lastbit(r))rv=S::op(rv,bit[r].query(ps...));
    for(; 1>=1; 1-=lastbit(1)) lv=S::op(lv,bit[1].query(ps...));
    return S::op(rv,S::inv(lv));
};
template<class S>
struct BIT<0, S>{ using T = typename S::T;
 T val=S::id;
  void add(T x) {val=S::op(val,x);}
 T query() {return val;}
struct AG{ //abelian group analogous to int addition
  using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return a+b;}
  static T inv(T a) {return -a;}
MoQueries.h
Description: Solve queries offline Can be changed by modifying Spec
Time: \mathcal{O}\left(n*\sqrt{(q)}\right)
vi mo(vector<pii> Q, vector<int> V) { // Queries in Q are \lceil l, r
  int L = 0, R = 0, blk = 350; // N/sqrt (Q)
  vi s(sz(Q)), res = s;
  auto K = [&](pii x) {return pii(x.first/blk, x.second ^ -(x.
       first/blk & 1)); };
  iota(all(s), 0);
  sort(all(s), [\&](int s, int t){ return K(Q[s]) < K(Q[t]); });
  int sum = 0;
  auto add = [&](int ind, int end) { sum += V[ind]; }; // add
       a \ | \ ind \ | \ (end = 0 \ or \ 1)
  auto del = [&](int ind, int end) { sum -= V[ind]; }; //
       remove a [ ind ]
  auto calc = [&]() { return sum; };
       compute current answer
  for (int qi : s) {
   pii q = Q[qi];
    while (L > q.first) add(--L, 0);
    while (R < q.second) add(R++, 1);</pre>
   while (L < q.first) del(L++, 0);
    while (R > q.second) del(--R, 1);
    res[qi] = calc();
  return res;
```

Matching (3)

OnlineMatching.h

int t = 0;

Description: Modified khun developed for specific question able to run $2 * 10^6$ queries, in $2 * 10^6$ x 10^6 graph in 3 seconds codeforces Time: $\mathcal{O}\left(confia\right)$

```
struct OnlineMatching {
   int n = 0, m = 0;
   vector<int> vis, match, dist;
   vector<vector<int>> q;
  vector<int> last;
```

6ac539, 42 lines

```
OnlineMatching(int n_, int m_) : n(n_), m(m_),
   vis(n, 0), match(m, -1), dist(n, n+1), g(n), last(n, -1)
    {}
    void add(int a, int b) {
        g[a].pb(b);
   bool kuhn(int a) {
   vis[a] = t;
   for(int b: g[a]) {
       int c = match[b];
       if (c == -1) {
       match[b] = a;
       return true;
       if (last[c] != t || (dist[a] + 1 < dist[c]))</pre>
       dist[c] = dist[a] + 1, last[c] = t;
    for (int b: q[a]) {
     int c = match[b];
     if (dist[a] + 1 == dist[c] && vis[c] != t && kuhn(c)) {
       match[b] = a;
        return true;
        return false;
 bool can_match(int a) {
    last[a] = t;
    dist[a] = 0;
    return kuhn(a);
Math (4)
Linear Diophantine Equation.h
Description: Find a solution to equation a^*x + b^*y = c
Time: \mathcal{O}(log(a))
                                                     538f05, 14 lines
array<11, 3> exgcd(11 a, 11 b) {
 if (a == 0) return {0, 1, b};
 auto [x, y, q] = exgcd(b % a, a);
 return {y - b / a * x , x, g};
array<11, 4> find_any_solution(11 a, 11 b, 11 c) {
 assert(a != 0 || b != 0);
 auto[x, y, g] = exgcd(a, b);
 if (c % g) return {false, 0, 0, 0};
 x *= c / q;
 y *= c / g;
 return {true, x, y, g};
```

Techniques (A)

techniques.txt

Combinatorics

159 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Scheduling Max contiquous subvector sum Invariants Huffman encoding Graph theory Dynamic graphs (extra book-keeping) Breadth first search Depth first search * Normal trees / DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Konig's theorem and vertex cover Min-cost max flow Lovasz toggle Matrix tree theorem Maximal matching, general graphs Hopcroft-Karp Hall's marriage theorem Graphical sequences Floyd-Warshall Euler cycles Flow networks * Augmenting paths * Edmonds-Karp Bipartite matching Min. path cover Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Edge coloring * Trees Vertex coloring * Bipartite graphs (=> trees) * 3^n (special case of set cover) Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Coin change Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag Dynprog over intervals Dynprog over subsets Dynprog over probabilities Dynprog over trees 3^n set cover Divide and conquer Knuth optimization Convex hull optimizations RMQ (sparse table a.k.a 2^k-jumps) Bitonic cycle Log partitioning (loop over most restricted)

Computation of binomial coefficients Pigeon-hole principle Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts Divisibility Euclidean algorithm Modular arithmetic * Modular multiplication * Modular inverses * Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Ouadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping Game theory Combinatorial games Game trees Mini-max Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search Unimodality and convex functions Binary search on derivative Numerical methods Numeric integration Newton's method Root-finding with binary/ternary search Golden section search Matrices Gaussian elimination Exponentiation by squaring Sorting Radix sort Geometry Coordinates and vectors * Cross product * Scalar product Convex hull Polygon cut Closest pair Coordinate-compression Ouadtrees KD-trees All segment-segment intersection Sweeping Discretization (convert to events and sweep) Angle sweeping Line sweeping Discrete second derivatives Strings Longest common substring Palindrome subsequences

Knuth-Morris-Pratt Tries Rolling polynomial hashes Suffix array Suffix tree Aho-Corasick Manacher's algorithm Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Best-first (A*) Bidirectional search Iterative deepening DFS / A* Data structures LCA (2^k-jumps in trees in general) Pull/push-technique on trees Heavy-light decomposition Centroid decomposition Lazy propagation Self-balancing trees Convex hull trick (wcipeg.com/wiki/Convex_hull_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree