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1 Contest
2 Data structures
3 Math
4 String
                                                              3
5 Graph
6 DP
Contest (1)
template.cpp
                                                         26 lines
#include <bits/stdc++.h>
using namespace std;
#define int long long
#define endl '\n'
#define rep(i, a, b) for(int i = (a); i < (b); ++i)
#define all(x) begin(x), end(x)
#define sz(x) (int)(x).size()
#define debug(var) cerr << #var << ": " << var << endl
#define pb push_back
#define eb emplace_back
typedef long long 11;
typedef pair<int, int> pii;
typedef vector<int> vi;
void solve() {
int32_t main() {
 ios_base::sync_with_stdio(0); cout.tie(0); cin.tie(0);
 int t = 1;
  while (t--) solve();
  return 0;
.bashrc
alias c='q++ -Wall -Wconversion -Wfatal-errors -q -std=c++17 \
 -fsanitize=undefined,address'
xmodmap -e 'clear lock' -e 'keycode 66=less greater' \#caps = \diamondsuit
.vimrc
set cin ai is ts=4 sw=4 nu rnu
" Select a region and then type :Hash
ca Hash w !cpp -dD -P -fpreprocessed \| tr -d '[:space:]' \
\| md5sum \| cut -c-6
hash.sh
# Hashes a file, ignoring all whitespace and comments. Use for
# verifying that code was correctly typed.
cpp -dD -P -fpreprocessed | tr -d '[:space:]' | md5sum |cut -c-6
```

```
troubleshoot.txt
                                                          52 lines
Pre-submit:
Write a few simple test cases if sample is not enough.
Are time limits close? If so, generate max cases.
Is the memory usage fine?
Could anything overflow?
Make sure to submit the right file.
Wrong answer:
Print your solution! Print debug output, as well.
Are you clearing all data structures between test cases?
Can your algorithm handle the whole range of input?
Read the full problem statement again.
Do you handle all corner cases correctly?
Have you understood the problem correctly?
Any uninitialized variables?
Any overflows?
Confusing N and M, i and i, etc.?
Are you sure your algorithm works?
What special cases have you not thought of?
Are you sure the STL functions you use work as you think?
Add some assertions, maybe resubmit.
Create some testcases to run your algorithm on.
Go through the algorithm for a simple case.
Go through this list again.
Explain your algorithm to a teammate.
Ask the teammate to look at your code.
Go for a small walk, e.g. to the toilet.
Is your output format correct? (including whitespace)
Rewrite your solution from the start or let a teammate do it.
Runtime error:
Have you tested all corner cases locally?
Any uninitialized variables?
Are you reading or writing outside the range of any vector?
Any assertions that might fail?
Any possible division by 0? (mod 0 for example)
Any possible infinite recursion?
Invalidated pointers or iterators?
Are you using too much memory?
Debug with resubmits (e.g. remapped signals, see Various).
Time limit exceeded:
Do you have any possible infinite loops?
What is the complexity of your algorithm?
Are you copying a lot of unnecessary data? (References)
How big is the input and output? (consider scanf)
Avoid vector, map. (use arrays/unordered_map)
What do your teammates think about your algorithm?
Memory limit exceeded:
What is the max amount of memory your algorithm should need?
Are you clearing all data structures between test cases?
Data structures (2)
SegTree.h
Description: Iterative SegTree Can be changed by modifying Spec
Time: \mathcal{O}(\log N)
                                                    607842, 25 lines
template<typename LS>
struct SegTree {
  using S = typename LS::S;
  using K = typename LS::K;
  int n;
  vector<S> seg;
  SegTree(int n)
```

```
: n(_n), seg(2*n, LS::id()) {}
  void update(int no, K val) {
    no += n;
    seg[no] = LS::update(val, seg[no]);
    while (no > 1) no /= 2, seg[no] = LS::op(seg[no*2], seg[no
         *2+11);
 S query(int 1, int r) { // [l, r)
    S vl = LS::id(), vr = LS::id();
    for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
      if (1 & 1) v1 = LS::op(v1, seg[1++]);
      if (r & 1) vr = LS::op(seg[--r], vr);
    return LS::op(vl, vr);
};
LazySeg.h
Description: Iterative Lazy SegTree Can be changed by modifying Spec
Time: \mathcal{O}(\log N)
                                                     ee5763, 96 lines
template<typename Spec>
struct LazySeq {
  using LS = Spec;
  using S = typename LS::S;
  using K = typename LS::K;
  int n;
  vector<S> seq:
  vector<K> lazv;
  vector<bool> has_lazy;
  // vector<int> lx, rx; // Aditional info
    LazySeg(vector<S> & v) : n(sz(v)), seg(2*n) , lazy(n),
        has lazv(n) {
    rep(no, 0, n) seg[no+n] = v[no];
    for (int no = n-1; no >= 1; no--) pull(no);
    // Aditional info, n must be power of two
    lx.assign(2*n, 0); rx.assign(2*n, 0);
    lx/1/ = 0; rx/1/ = n;
    rep(no, 1, n) {
      int \ mid = (lx[no] + rx[no])/2;
      lx[no*2] = lx[no]; rx[no*2] = mid;
      lx[no*2+1] = mid; rx[no*2+1] = rx[no];
  S query(int 1, int r) { // [l, r)
   1 += n;
    push_to(1); push_to(r-1);
    S vl = LS::id(), vr = LS::id();
    while(1 < r) {
     if (1 & 1) v1 = LS::op(v1, seg[1++]);
      if (r & 1) vr = LS::op(seg[--r], vr);
      1 >>= 1; r >>= 1;
    return LS::op(vl, vr);
  void update(int 1, int r, K val) {
   1 += n;
    r += n;
```

 $push_to(1)$; $push_to(r-1)$;

int lo = 1, ro = 1;

```
while(1 < r) {
      if (1 & 1) lo = max(lo, 1), apply(l++, val);
      if (r \& 1) ro = max(ro, r), apply(--r, val);
     1 >>= 1; r >>= 1;
    pull_from(lo);
    pull_from(ro-1);
  void apply(int no, K val) {
    seg[no] = LS::update(val, seg[no]);
    // seg[no] = LS::update(val, seg[no], lx[no], rx[no]);
    if (no < n) {
      if (has_lazy[no]) lazy[no] = LS::compose(val, lazy[no]);
      else lazy[no] = val;
      has_lazy[no] = true;
  void pull_from(int no) {
   while(no > 1) no >>= 1, pull(no);
  void pull(int no) {
   seq[no] = LS::op(seq[no*2], seq[no*2+1]);
  void push_to(int no) {
   int h = 0; int p2 = 1;
   while (p2 < no) p2 \star= 2, h++;
    for (int i = h; i >= 1; i--) push(no >> i);
  void push (int no) {
    if (has_lazy[no]) {
      apply(no*2, lazy[no]);
      apply (no *2+1, lazy[no]);
      has lazv[no] = false;
};
struct Spec {
  using S = int;
  using K = int;
  static S op(S a, S b) { return max(a, b); }
  static S update(K f, S a) { return f + a; }
  static K compose(const K f, const K g) { return f + g; }
  static S id() { return 0; }
RecSeg.h
Description: Recursive generic persistent dinamic SegTree Can be changed
by modifying Spec, queries are inclusive exclusive
Memory: \mathcal{O}(Q * \log N)
Time: \mathcal{O}(\log N)
                                                       99dd6a, 36 lines
template<typename LS>
struct Node {
  using S = typename LS::S;
  using K = typename LS::K;
  Node<LS> *1 = 0, *r = 0;
  int lo, hi;
  S val = 0;
  Node (int lo, int hi): lo(lo), hi(hi), val(LS::id()) {}
```

```
S query(int L, int R) { // [L, R)
    if (R <= lo || hi <= L) return LS::id();</pre>
    if (L <= lo && hi <= R) return val;</pre>
    return LS::op(l->query(L, R), r->query(L, R));
  Node<LS>* update(int idx, K x) {
    if (hi <= idx || idx < lo) return this;</pre>
    Node<LS>* me = new Node(lo, hi);
    push(); me->1 = 1; me->r = r;
    if (hi - lo == 1) me->val = LS::update(x, val);
      me \rightarrow l = l \rightarrow update(idx, x), me \rightarrow r = r \rightarrow update(idx, x);
      me \rightarrow val = LS::op(me \rightarrow l \rightarrow val, me \rightarrow r \rightarrow val);
    return me;
  void push() {
    if (!1) {
      int mid = lo + (hi - lo)/2;
      1 = new Node(lo, mid); r = new Node(mid, hi);
};
Description: Multidimensional Psum Requires Abelian Group S (op, inv,
Memory: \mathcal{O}\left(N^{D}\right)
Time: \mathcal{O}(1)
                                                            65f259, 29 lines
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct Psum{ using T = typename S::T;
  vector<Psum<D-1, S>> v;
  MAs Psum(int s, As... ds):n(s+1),v(n,Psum<D-1, S>(ds...)){}
  MAs void set(T x, int p, As... ps){v[p+1].set(x, ps...);}
  void push(Psum& p) {rep(i, 1, n)v[i].push(p.v[i]);}
  void init() {rep(i, 1, n)v[i].init(),v[i].push(v[i-1]);}
  MAs T query (int 1, int r, As... ps) {
    return S::op(v[r+1].query(ps...),S::inv(v[1].query(ps...)))
};
template<class S>
struct Psum<0, S>{ using T = typename S::T;
  T val=S::id;
  void set(T x) {val=x;}
  void push(Psum& a) {val=S::op(a.val,val);}
  void init(){}
  T query() {return val;}
struct G{
  using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return a+b;}
  static T inv(T a) {return -a;}
MultiDSegTree.h
Description: Multidimensional SegTree Requires Monoid S (op, id)
Memory: \mathcal{O}(N^D)
Time: \mathcal{O}\left((\log N)^D\right)
                                                           53621d, 37 lines
```

```
//#pragma once
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct SegTree{ using T = typename S::T;
  vector<SegTree<D-1, S>> seg;
  MAs SegTree(int s, As... ds):n(s),seg(2*n, SegTree<D-1, S>(ds
       . . . ) ) { }
  MAs T get(int p, As... ps) {return seg[p+n].get(ps...);}
  MAs void update (T x, int p, As... ps) {
    p+=n; seg[p].update(x, ps...);
    for(p>>=1;p>=1;p>>=1)
    seg[p].update(S::op(seg[2*p].get(ps...), seg[2*p+1].get(ps
          ...)), ps...);
  MAs T query(int 1, int r, As... ps) {
    T lv=S::id, rv=S::id;
    for(l+=n,r+=n+1;l<r;l>>=1,r>>=1){
      if (1&1) lv = S::op(lv, seg[1++].query(ps...));
      if (r&1)rv = S::op(seg[--r].query(ps...),rv);
    return S::op(lv,rv);
};
template < class S>
struct SegTree<0, S>{ using T = typename S::T;
 T val=S::id;
  T get() {return val;}
  void update(T x) {val=x;}
  T query() {return val;}
struct M{ //monoid
  using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return max(a,b);}
SparseTable.h
Description: Multidimensional Sparse Table Requires Idempotent Monoid
S (op, inv, id)
Memory: \mathcal{O}\left(\left(n\log n\right)^D\right)
Time: \mathcal{O}(1) query, \mathcal{O}((n \log n)^D) build
                                                        c900f0, 39 lines
\#define MAs template < class...As > //multiple arguments
template<int D, class S>
struct SpTable{ using T = typename S::T;
  using isp = SpTable<D-1, S>;
  inline int lg(signed x) {return __builtin_clz(1) -__builtin_clz
       (x);}
  int n;
  vector<vector<isp>> tab;
  MAs SpTable(int s, As... ds):n(s),
  tab(1+lg(n), vector < isp > (n, isp(ds...))) {}
  MAs void set(T x, int p, As... ps){tab[0][p].set(x, ps...);}
  void join(SpTable& a, SpTable& b) {
    rep(i, 0, 1+lg(n))rep(j, 0, n)
      tab[i][j].join(a.tab[i][j], b.tab[i][j]);
  void init(){
    rep(i, 0, n)tab[0][i].init();
    rep(i, 0, lg(n)) rep(j, 0, n-(1 << i))
      tab[i+1][j].join(tab[i][j], tab[i][j+(1<<i)]);
  MAs T query(int 1, int r, As... ps) {
```

vi s(sz(0));

```
int k = lq(r-l+1); r+=1-(1<< k);
    return S::op(tab[k][1].query(ps...),tab[k][r].query(ps...))
};
template<class S>
struct SpTable<0, S>{ using T = typename S::T;
 T val=S::id;
  void set(T x) {val=x;}
  void join(SpTable& a, SpTable& b) {val=S::op(a.val,b.val);}
  void init(){}
 T query() {return val;}
struct IM{
 using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return max(a, b);}
BIT.h
Description: Multidimensional BIT/Fenwick Tree Requires Abelian Group
"S" (op, inv, id)
Memory: \mathcal{O}\left(N^{D}\right)
Time: \mathcal{O}\left((\log N)^D\right)
                                                        778135, 31 lines
#define MAs template<class... As> //multiple arguments
template<int D, class S>
struct BIT{ using T = typename S::T;
  vector<BIT<D-1, S>> bit;
  MAs BIT(int s, As... ds):n(s),bit(n+1, BIT<D-1, S>(ds...)){}
  inline int lastbit(int x) {return x&(-x);}
  MAs void add(T x, int p, As... ps) {
    for (p++; p<=n; p+=lastbit (p) ) bit [p] .add (x, ps...);</pre>
  MAs T query(int 1, int r, As... ps) {
    T lv=S::id, rv=S::id; r++;
    for(;r>=1;r-=lastbit(r))rv=S::op(rv,bit[r].query(ps...));
    for(; l>=1; l-=lastbit(l)) lv=S::op(lv,bit[l].query(ps...));
    return S::op(rv,S::inv(lv));
template < class S>
struct BIT<0, S>{ using T = typename S::T;
 T val=S::id;
  void add(T x) {val=S::op(val,x);}
 T query() {return val;}
struct AG{ //abelian group analogous to int addition
  using T = int;
  static constexpr T id = 0;
  static T op(T a, T b) {return a+b;}
  static T inv(T a) {return -a;}
MoQueries.h
Description: Solve queries offline Can be changed by modifying Spec
Time: \mathcal{O}\left(n*\sqrt{(q)}\right)
template<typename LS>
void mo(LS & v, vector<pii> Q) { // Queries in Q are [l, r)
  int L = 0, R = 0, blk = 350; // N/sqrt (Q)
```

```
auto K = [&](pii x) {return pii(x.first/blk, x.second ^ -(x.
      first/blk & 1)); };
  iota(all(s), 0);
  sort(all(s), [\&](int s, int t){ return K(Q[s]) < K(Q[t]); });
  for (int qi : s) {
   pii q = Q[qi];
    while (L > q.first) v.add(--L, 0);
    while (R < q.second) v.add(R++, 1);
    while (L < q.first) v.del(L++, 0);
    while (R > q.second) v.del(--R, 1);
    v.calc(gi);
}
struct Spec {
 using S = 11;
 void add(int ind, int end) {}
 void del(int ind, int end) {}
 void calc(int idx) {}
Math (3)
Linear Diophantine Equation.h
Description: Find a solution to equation a^*x + b^*y = c
Time: \mathcal{O}(log(a))
                                                      538f05, 14 lines
array<11, 3> exgcd(11 a, 11 b) {
 if (a == 0) return {0, 1, b};
  auto [x, y, g] = exgcd(b % a, a);
  return {y - b / a * x , x, g};
array<11, 4> find_any_solution(11 a, 11 b, 11 c) {
  assert(a != 0 || b != 0);
  auto[x, y, g] = exgcd(a, b);
 if (c % q) return {false, 0, 0, 0};
 x *= c / g;
 y *= c / g;
 return {true, x, y, g};
String (4)
Hash.h
Description: String hashing
Memory: \mathcal{O}(n)
Time: \mathcal{O}(1) query, \mathcal{O}(n) build
<sys/time.h>
                                                      2d1ad6, 43 lines
 Arithmetic mod two primes and 2°32 simultaneously.
  C can be initilialize to a number less than MOD or random
  timeval tp;
  gettimeofday(\&tp, 0);
  C = (int)tp.tv\_usec; // (less than modulo)
  assert((ull)(H(1)*2+1-3) == 0);
typedef uint64_t ull;
template<int M, class B>
struct A {
  int x; B b; A(int x=0) : x(x), b(x) {}
  A(int x, B b) : x(x), b(b) {}
  A operator+(A o) const {int y = x+o.x; return{y - (y>=M) *M, b
```

```
-0.h}:
  A operator* (A o) const { return \{(int)((11)x*o.x % M), b*o.b\}
  explicit operator ull() const { return x ^ (ull) b << 21; }</pre>
 bool operator==(A o) const { return (ull) *this == (ull) o; }
 bool operator<(A o) const { return (ull) *this < (ull) o; }</pre>
typedef A<1000000007, A<1000000009, unsigned>> H;
static int C;
struct HashInterval {
 int n:
 vector<H> ha, pw;
 template<typename S>
 HashInterval(const S & str) : n(sz(str)), ha(n+1), pw(n+1) {
   pw[0] = 1;
   rep(i,0,n)
      ha[i+1] = ha[i] * C + str[i],
      pw[i+1] = pw[i] * C;
 H query (int a, int b) { // hash (a, b)
    return ha[b] - ha[a] * pw[b - a];
 H queryI(int a, int b) {
    return query (n - b, n - a);
DinamicHashing.h
Description: Dinamic string hashing
Memory: \mathcal{O}(n)
Time: \mathcal{O}(1) query, \mathcal{O}(n) build
                                                       cf8a3c, 41 lines
typedef uint64_t ull;
template<int M, class B>
struct A {
 int x; B b;
 constexpr A(int x=0) : x(x), b(x) {}
 constexpr A(int x, B b) : x(x), b(b) {}
 A operator+(A o) const {int y = x+o.x; return{y - (y>=M) *M, b
 A operator-(A o) const {int y = x-o.x; return{y + (y < 0) *M, b
 A operator*(A o) const { return {(int)(1LL*x*o.x % M), b*o.b}
  explicit operator ull() const { return x ^ (ull) b << 21; }</pre>
 bool operator==(A o) const { return (ull) *this == (ull) o; }
 bool operator<(A o) const { return (ull) *this < (ull) o; }</pre>
typedef A<989831283, A<912391239 , unsigned>> H;
const static int C = 12312;
struct DinamicHash {
 int n;
 vector<int> s;
 vector<H> p;
 SeaTree sea:
 DinamicHash(const vector<int> & v) : n(v.size()), s(v), p(n
       +1) {
    p[0] = 1;
    vector<H> values(n);
    rep(i, 0, n) {
      values[i] = p[i] * s[i];
      p[i+1] = p[i] * C;
```

A operator-(A o) const {int y = x-o.x; return{y + (y< 0)*M, b

seg = SegTree(values);

void update(int idx, int v) {

s[idx] = v;

H query(int 1, int r) const $\{ // [l, r) \}$

return seg.query(1, r) * p[n-r];

seg.update(idx, p[idx] * s[idx]);

Automata KMP Kosaraju TwoSat

```
};
Automata.h
Description: Suffix automata
Memory: \mathcal{O}(n*26)
Time: \mathcal{O}(n) build
                                                      92d90c, 49 lines
struct Automata {
  int saID = 1, last = 1;
  int n;
  vector<int> len, lnk;
  vector<array<int,27>> to;
  vector<int> occ, fpos;
  vector<int> states;
  Automata(const string & s, const char a = 'a')
    : n(s.size()), len(2*n+2), lnk(2*n+2), to(2*n+2, {0}), occ
         (2*n+2), fpos(2*n+2) {
    for (const auto & c: s) push(c-a);
    states.assign(saID, 0);
    iota(all(states), 1);
    sort(all(states), [&](const auto & u, const auto & v) {
         return len[u] > len[v]; });
    for (auto st: states) {
     occ[lnk[st]] += occ[st];
  void push (int c) {
    int a = ++saID;
    int p = last;
    last = a:
    len[a] = len[p] + 1;
    occ[a] = 1;
    fpos[a] = len[a] - 1;
    for (; p > 0 && !to[p][c]; p = lnk[p]) to[p][c] = a;
    int q = to[p][c];
    if (p == 0) {
     lnk[a] = 1;
    else if (len[p] + 1 == len[q]) {
     lnk[a] = q;
    else {
      int clone = ++saID;
     lnk[clone] = lnk[q];
     to[clone] = to[q];
      fpos[clone] = fpos[q];
     len[clone] = len[p] + 1;
     lnk[a] = lnk[q] = clone;
      for (; to[p][c] == q; p = lnk[p]) to[p][c] = clone;
};
```

```
KMP.h
Description: KMP automaton
Memory: \mathcal{O}(26n)
Time: \mathcal{O}(n) build, \mathcal{O}(1) query
                                                        32c2a7, 22 lines
struct KMP {
 const int inic = 'a';
 int n; vi pi; vector<vi> aut;
 KMP(string s):n(s.size()),pi(n),aut(n+1, vi(26,0)){
    rep(i, 1, n) {
      int j = pi[i-1];
      while (j > 0 \text{ and } s[j]!=s[i])j = pi[i-1];
      pi[i] = j + (s[i] = s[j]);
   s.pb('#');
    rep(i, 0, n+1){
      rep(c, 0, 26){
        if (i > 0 and c+inic!=s[i])aut[i][c] = aut[pi[i-1]][c];
        else aut[i][c] = i + (c+inic == s[i]);
 pair<bool, int> nxt(int cur, char c) {
   cur = aut[cur][c];
    return {cur==n, cur};
};
```

Graph (5)

Description: Kosaraju

Kosaraju.h

```
Memory: \mathcal{O}(n)
Time: \mathcal{O}(n+m) query, \mathcal{O}(n+m) build
                                                       eb04c1, 43 lines
struct Kosaraju {
 int n:
 vector<vector<int>> g, rg;
 vector<bool> vis;
 vector<int> id;
 vector<vector<int>> dag, comp;
 int cc = 0;
 vector<int> S;
 Kosaraju(int _n)
   : n(_n), g(n), rg(n), vis(n), id(n) {}
 void add_edge(int a, int b) {
    g[a].eb(b);
    rg[b].eb(a);
 void dfs(int a) {
    vis[a] = true;
    for (auto b: g[a]) if (!vis[b]) dfs(b);
    S.eb(a);
 void scc(int a, int c) {
   vis[a] = true;
   id[a] = c;
    for (auto b: rg[a]) if (!vis[b]) scc(b, c);
 void run() {
    rep(a, 0, n) if (!vis[a]) dfs(a);
    fill(all(vis), 0);
    reverse(all(S));
    for (auto a: S) if (!vis[a]) scc(a, cc++);
```

dag.resize(cc); comp.resize(cc);

```
vector<pair<int, int>> edges;
    rep(a, 0, n) {
      comp[id[a]].eb(a);
      for (int b: g[a]) if (id[a] != id[b]) edges.eb(id[a], id[
           b]);
    sort (all (edges));
    edges.erase(unique(all(edges)), edges.end());
    for (const auto & [a, b]: edges) dag[a].eb(b);
};
TwoSat.h
Description: Two sat
Memory: \mathcal{O}(n)
Time: \mathcal{O}(n+m) query, \mathcal{O}(n+m) build
                                                     0b8692, 64 lines
struct TwoSat{
 int n:
 vector<vector<int>> g, gi;
 vector<bool> vis;
  vector<int> vars, comp;
  vector<int> top;
 TwoSat(int _n)
   : n(_n), q(2*n), qi(2*n), vis(2*n), vars(n, -1), comp(2*n)
 int neg(int a) {
   if (a >= n) return a-n;
    return a + n;
 void add_or(int a, int b) {
    g[neg(a)].eb(b);
    q[neq(b)].eb(a);
    gi[b].eb(neg(a));
    gi[a].eb(neg(b));
 void add_imp(int a, int b) {
    add_or(neg(a), b);
  void add_either(int a, int b) {
    add_or(a, b);
    add_or(neg(a), neg(b));
  void dfs(int a) {
   vis[a] = true;
    for (auto b: q[a]) if (!vis[b]) dfs(b);
    top.eb(a);
 void idfs(int a, int c){
   vis[a] = true;
    comp[a] = c;
    for (auto b: gi[a]) if (!vis[b]) idfs(b, c);
 bool sat() {
    int c = 0;
    rep(a, 0, 2*n) if (!vis[a]) dfs(a);
    fill(all(vis), 0);
    reverse(all(top));
    for(int a : top) if (!vis[a]) idfs(a, c++);
```

```
for(int a: top) {
      if (comp[a] == comp[neq(a)]) return false;
      bool is_neg = a >= n;
      if (is_neq) a = neq(a);
      if (vars[a] == -1) vars[a] = is_neg;
    return true:
};
OnlineMatching.h
Description: Modified khun developed for specific question able to run
2 * 10^6 queries, in 2 * 10^6 x 10^6 graph in 3 seconds codeforces
Time: \mathcal{O}(confia)
                                                       6ac539, 42 lines
struct OnlineMatching {
    int n = 0, m = 0;
    vector<int> vis, match, dist;
    vector<vector<int>> q;
  vector<int> last;
  int t = 0;
    OnlineMatching(int n , int m ) : n(n ), m(m ),
    vis(n, 0), match(m, -1), dist(n, n+1), g(n), last(n, -1)
    void add(int a, int b) {
        g[a].pb(b);
    bool kuhn(int a) {
    vis[a] = t;
    for(int b: q[a]) {
        int c = match[b];
        if (c == -1) {
        match[b] = a;
        return true;
        if (last[c] != t || (dist[a] + 1 < dist[c]))</pre>
        dist[c] = dist[a] + 1, last[c] = t;
    for (int b: q[a]) {
      int c = match[b];
      if (dist[a] + 1 == dist[c] && vis[c] != t && kuhn(c)) {
        match[b] = a;
        return true;
        return false;
  bool can_match(int a) {
   t++;
    last[a] = t;
    dist[a] = 0;
    return kuhn(a);
};
FunctGraph.h
Description: Functional Graph
Memory: \mathcal{O}(n)
Time: \mathcal{O}(n)
```

```
struct FunctGraph{
 int n;
  vi head;
```

```
vector<vi> gr, comps;
  FunctGraph (vi& fn):
   n(sz(fn)), head(n, -1), gr(n) {
    rep(i, 0, n)gr[fn[i]].pb(i);
   vi visited(n, 0);
    auto dfs = [&](auto rec, int v, int c) -> void{
     head[v] = c; visited[v] = 1;
     for(int f : gr[v])if (head[f]!=f)rec(rec, f, c);
    };
    rep(i, 0, n){
     if (visited[i])continue;
      int l=fn[i], r=fn[fn[i]];
      while(1!=r) l=fn[1], r=fn[fn[r]];
      vi cur = \{r\};
      for(l=fn[1]; 1!=r; l=fn[1]) cur.pb(1);
      for(int x : cur) dfs(dfs, x, x);
      comps.pb(cur);
};
Hierholzer.h
Description: Eulerian path/cycles if existing
Memory: \mathcal{O}(V+E)
Time: \mathcal{O}(E)
                                                     a85ad9, 29 lines
vi hierholzer(int n, vector<pii>& edges, int inic) {
 vi ans; int m = sz(edges);
  auto check = [&]()->bool{return true;};
  if (not check()){
    //a function should be created to check conditions
    //acording to type of graph and problem restrictions on
    //the path type and enpoints
    //base conditions: edge connectivity and vertex degree
    return ans; //empty vector if impossible
  vector<vi> q(n);
  rep(i, 0, m){
    auto [a, b] = edges[i];
    g[a].pb(i); g[b].pb(i); //remove the latter if it's
  vi used(m, false), st = {inic};
  while(not st.empty()){
    int v = st.back();
    while(not g[v].empty() and used[g[v].back()])g[v].pop_back
    if (g[v].empty())st.pop_back(), ans.pb(v);
      int idx = q[v].back(); q[v].pop_back();
      auto [a, b] = edges[idx]; used[idx] = true;
     st.pb((v==a ? b : a));
 reverse(all(ans));
 return ans;
DominatorTree.h
```

Description: Dominator Tree, creates the graph tree, where all ancestors of a u in the tree are necessary in the path from the root to u Memory: $\mathcal{O}(n)$

```
Time: \mathcal{O}((n+m)log(n)) build
```

7eb669, 56 lines

```
struct DominatorTree {
 vector<vector<int>> g, gt, tree, bucket, down;
 vector<int> S;
```

```
vector<int> dsu, label, sdom, idom, id;
  int dfstime =0;
 DominatorTree(vector<vector<int>> & _g, int root)
     : n(sz(_g)), g(_g), gt(n), tree(n), bucket(n), down(n),
      S(n), dsu(n), label(n), sdom(n), idom(n), id(n) {
    prep(root); reverse(all(S));
    for(int u : S) {
      for(int v : gt[u]) {
       int w = fnd(v);
        if(id[ sdom[w] ] < id[ sdom[u] ])</pre>
          sdom[u] = sdom[w];
      gt[u].clear();
      if(u != root) bucket[ sdom[u] ].push_back(u);
      for(int v : bucket[u]) {
       int w = fnd(v);
        if(sdom[w] == sdom[v]) idom[v] = sdom[v];
        else idom[v] = w;
      bucket[u].clear();
      for(int v : down[u]) dsu[v] = u;
      down[u].clear();
    reverse(S.begin(), S.end());
    for(int u : S) if(u != root) {
      if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
      tree[ idom[u] ].push_back(u);
  void prep(int u) {
    S[dfstime] = u;
    id[u] = ++dfstime;
    label[u] = sdom[u] = dsu[u] = u;
    for(int v : g[u]){
      if(!id[v])
        prep(v), down[u].push_back(v);
      gt[v].push back(u);
  int fnd(int u, int flag = 0){
   if(u == dsu[u]) return u;
    int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
    if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
     label[u] = b;
    dsu[u] = v;
    return flag ? v : label[u];
};
```

DP(6)

```
sos.h
Description: Sos DP
Time: \mathcal{O}(n*2^n)
```

```
5063f0, 19 lines
```

```
//iterative version
for(int mask = 0; mask < (1<<N); ++mask){</pre>
  dp[mask][-1] = A[mask]; //handle base case separately (leaf
       states)
  for(int i = 0; i < N; ++i){</pre>
    if(mask & (1<<i))
      dp[mask][i] = dp[mask][i-1] + dp[mask^(1<<i)][i-1];
      dp[mask][i] = dp[mask][i-1];
```

UnB

```
}
F[mask] = dp[mask][N-1];
}

//memory optimized, super easy to code.
for(int i = 0; i < (1 < N); ++i)
   F[i] = A[i];
for(int i = 0; i < N; ++i) for(int mask = 0; mask < (1 << N); ++
        mask) {
        if(mask & (1 << i))
            F[mask] += F[mask^(1 << i)];
}</pre>
```

6

Techniques (A)

techniques.txt

117 lines

Recursion Divide and conquer Finding interesting points in N log N Algorithm analysis Master theorem Amortized time complexity Greedy algorithm Max contiguous subvector sum Invariants Graph theory DP, com cyclo no dikstra reverso Breadth first search Depth first search DFS trees Dijkstra's algorithm MST: Prim's algorithm Bellman-Ford Min-cost max flow Flovd-Warshall Euler cycles Flow networks Bipartite matching Topological sorting Strongly connected components Cut vertices, cut-edges and biconnected components Vertex coloring Bipartite graphs Diameter and centroid K'th shortest path Shortest cycle Dynamic programming Knapsack Longest common subsequence Longest increasing subsequence Number of paths in a dag Shortest path in a dag over intervals over subsets over probabilities over trees 3^n set cover Divide and conquer Convex hull optimizations RMQ (sparse table a.k.a 2^k-jumps) Combinatorics Inclusion/exclusion Catalan number Pick's theorem Number theory Integer parts (School's excursion) Divisibility Euclidean algorithm Modular inverses Modular exponentiation by squaring Chinese remainder theorem Fermat's little theorem Euler's theorem Phi function Frobenius number Quadratic reciprocity Pollard-Rho Miller-Rabin Hensel lifting Vieta root jumping

Game theory Combinatorial games Game trees Mini-max Nim Games on graphs Games on graphs with loops Grundy numbers Bipartite games without repetition General games without repetition Alpha-beta pruning Probability theory Optimization Binary search Ternary search (Convex functions) Binary search on derivative Numerical methods Newton's method Root-finding with binary/ternary search Matrices Gaussian elimination Exponentiation by squaring Geometry Cross product Scalar product Convex hull Polygon cut Closest pair (Distance functions) Hull diameter (Distance functions) Sweeping Discretization (convert to events and sweep) Angle sweeping Line sweeping Strings Longest common substring Knuth-Morris-Pratt Tries Rolling polynomial hashes Aho-Corasick Letter position lists Combinatorial search Meet in the middle Brute-force with pruning Data structures LCA (2^k-jumps in trees in general) Centroid decomposition SegTree, LazySeg Convex hull trick (wcipeg.com/wiki/Convex_hull_trick) Monotone queues / monotone stacks / sliding queues Sliding queue using 2 stacks Persistent segment tree