

VIENNA UNIVERSITY OF TECHNOLOGY

105.625 PR ADVANCED ECONOMICS PROJECT

Double Sided Matching

9702170, Karin Leithner
0725439, Florin Bogdan Balint
1025735, Clemens Proyer
1027143, Mattias Haberbusch
1027433, Thomas Solich

April 3, 2016

Contents

1	Introduction	2
1.1	Motivation and Problem Description	2
1.2	Aim and Objectives	2
1.3	Structure	3
2	Theory	4
2.1	The Double Sided Matching Problem	4
2.2	Gale and Shapley Algorithm	4
2.3	Relevance to Game Theory	5
2.3.1	General	5
2.3.2	Example for the application of game theory	7
2.3.3	Double Sided Matching Problem and Game Theory	8
3	Prototype	9
3.1	Subsection	9
4	Prototype Evaluation: College Applications in Austria	10
4.1	Subsection	10
5	Prototype Evaluation: Labor Market in the European Union	11
5.1	Subsection	11
6	Summary	12
6.1	Subsection	12
7	Example	13
7.1	Heading on level 2 (subsection)	14
7.1.1	Heading on level 3 (subsubsection)	14
7.2	Example for list (3*itemize)	14
7.3	Example for list (enumerate)	15
	List of Figures	16
	List of Tables	17
	References	18

Introduction

Double sided matching is about finding a match between two sets of elements. The first algorithm which should solve this problem was described by David Gale and Lloyd S. Shapley in 1962 [1]. The focus will lie on the examination of this algorithm and the creation of a simulation model for it, which can be applied to different economic issues. The selected economic issues will be presented in the following chapters.

add one
sentence
about
something

Motivation and Problem Description

In many daily live situations there are hidden double sided matching problems. For example students usually have the opportunity to apply to a number of schools. In this case students and schools have to be matched together. Both parties have preferences which need to be taken into consideration. In our everyday life we usually don't have the time to find an appropriate solution.

In some cases these problems even go by without solving them. In other cases however, it is crucial to solve them, e.g. if we consider medical school graduates and open positions in hospitals. Due to multiple reasons (e.g. different preferences, high number of elements in the two sets) this problem cannot be solved intuitively or in linear time.

Aim and Objectives

The aim of this paper is to develop a simulation model prototype which can solve similar problems, like the ones discussed above. With this model we could answer questions like:

- Are there "stable" solutions for this kind of problem?
- How can it be applied to all prospective students and universities in Austria?
- Is there also a solution for the current labor supply and demand in the European Union?

Structure

In the second chapter a theoretical overview is provided regarding the double sided matching problem and what stability means in this context. The focus lies on the Gale and Shapley Algorithm, basic information about the game theory and the relationship between them.

The third chapter describes the prototype of the simulation model, which implements the previously discussed theory into practice. As a simulation environment Netlogo will be used.

The fourth chapter deals with the examination of data from Austrian universities and their integration into the prototype. As a result, a solution for this double sided matching problem will be provided.

Chapter five deals with the labor market in the European Union, trying to match the labor demand and labor supply.

Finally the most important findings and results are summarized and compared in the last chapter. Additionally there is a short outlook into further research.

rework
shortly
before
we are
finished,
because
structure
changes
most time

Theory

The Double Sided Matching Problem

The double sided matching problem describes a problem, which has two sets of elements. These sets are disjoint and may be equal in size, but this is not mandatory. The purpose of the double sided matching algorithm is to find a solution, so every element of the first set has a corresponding counterpart in the other set. This can be best described in an example: Assume that the elements are men and women and that there are only heterosexual preferences for the individuals. The two sets are now male and female. The double sided matching algorithm shall create couples, which consists of a man and woman.

The solutions of the matchings can have different properties, which might be very important (e.g. stability). The output of the algorithm of men and women is not allowed to be unstable, i.e. there are two men α and β who are assigned to women A and B although β prefers A to B and A prefers α to β . If this does not occur, the assignment is called *stable*. If there is more than one stable solution the *optimal* one can be of particular interest.

Optimality means that a solution has the best outcome for an individual or set. In this case optimality can be achieved from the perspective of the first set of elements or from the second set. There can also be a trade-off from the perspective of the both sets, in this case optimality is achieved by weighting the expected value of both sets. These properties have been described in the paper by Gale and Shapley [1].

Generally in economics this is also known as pareto efficiency [2, p. 46]. Pareto efficiency is a characteristic for a solution where any change which would improve one property on one hand, but will worsen another one on the other hand.

Gale and Shapley Algorithm

In the paper of Gale and Shapley [1] an algorithm was proposed for solving the double sided matching problem. Let's assume there are two sets of elements: one contains men and the other one women. The size of the two sets is equal, the preferences are heterosexual and there are no polygamies (i.e. one element is only allowed to be matched with one element from the other set). The aim of the algorithm is to find a stable way of marrying men and women. Each man and

women has their own preference for their marriage partner.

First Iteration

In the first iteration each man proposes to the woman who he ranked first. Afterwards a woman has a list of men that proposed to her. This list might be zero, if not she rejects every men, except her highest preferred man on this list. If a man is not rejected he and the woman create a temporary couple. All man who are not in a couple are single.

Second Iteration

In the second interaction each man who is single, proposes to his highest ranked woman, who he didn't already proposed to. Each woman now chooses the man which she ranks the highest out of her received proposals. If a man is not rejected he and the woman create a temporary couple (Note: after the second iteration some men, who have been in a couple from the first iteration might be now single).

Further Iterations

The third iteration is the same as the second one. This continues until no men is single anymore. If this condition is satisfied a stable solution is found for the double sided matching problem. The stability of the solution of this algorithm has been proven by contradiction [1].

Relevance to Game Theory

General

"Game theory is about what happens when people - or genes, or nations - interact" [3, p. 1]. An important aspect for participating parties is to anticipate how the opposite party will react on certain actions. Mathematics shall help to analyze, understand and estimate outcomes of such games. Depending on the information participants have, they decide how to act basing on rules contained in their strategy. Game theory is widely applied in economics. Companies use game theory to estimate e.g. reactions of competitors or behavior of employees. The

major advantages of game theory are its precision and its universal applicability to all kind of games. [3, pp. 1-3]

In addition there are some assumptions which are made during preparation and execution by every individual who is interacting with another one:

- Well-specified choices
- Well-defined end-state
- Specified pay-off
- Perfect knowledge
- Rationality

However, these perfect preconditions will never be met in real scenarios. In the context of double sided matching, it is important to have well-specified choices defined. Nevertheless, there might be chances to get into a nearly worst-case scenario instead. But furthermore, as a consequence of these listed assumptions above, there are two different kind of games:

- Static ones
- Dynamic ones

The two items can be combined with the following kind of knowledge

- Incomplete
- Complete (there is no private information)

Depending on the nature of the game, whether it is static or dynamic and the knowledge is complete or incomplete four different combinations can be made. For a better understanding, following examples only consist of two players (every combination can be used for a higher amount of players too):

- Static and (in)complete information

As of complete information, the players action set can be presented in an $N \times M$ matrix. Therefore all information (also private) is known to every player and each player is able to easily eliminate the pay-offs of the other players. The pay-off is the outcome of a chosen action of a player which depends on the state or previous actions of the game. The pay-off can change dynamically while the game is in progress. As of incomplete information players are forced to be uncertain about the other player's actions and pay-offs. Besides that, there are two basic steps

1. Player 1 and 2 are choosing actions out of a set, which were predefined by them before.
2. After their actions, they receive their pay-offs.

- Dynamic and complete information

When basic assumptions are made on complete information, players will create a strategy to ensure to get the highest return of each action. Although the strategy might be changed during the game after each turn. Incomplete information leads to communication of uninformed parties (i.e. closed private information) which might end up to under- or overestimate the others pay-offs. In general there will be a Nash Equilibrium between those parties. This equilibrium is created by players because they choose a strategy which do not depend on the other players. There are three basic steps:

1. Player 1 chooses an action out of his/her predefined set.
2. Player 2 observes player 1's action and chooses an action based on his/her observation.
3. After their actions, they receive their pay-offs.

The previously described theory is based on information of the paper "An introduction to applicable game theory" [4].

Example for the application of game theory

Game theory is also very popular in the field of economics. Many companies primarily apply game theory to estimate the reactions of their competitors. In the last auction for broadband internet frequencies in Germany, all major providers hired game theorists out of two reasons. Firstly, they wanted to win the auction

for the desired frequency and secondly, they also did not want to overpay for the frequency. The bidding behavior of a provider indicated the interest rate for a frequency. As a consequence, other providers retained their bids for parts of this frequency so the interested provider did not need to overpay for it. The positive result for the providers was, that the 700-MHz frequency parts were sold for the minimum bidding price. [5]

Double Sided Matching Problem and Game Theory

Like described in chapter 2, the goal of the three main problems, which will be simulated in the course of this project, is to bring different parties together. In the original marriage problem discussed by Shapley, the goal was to match men and women so the overall situation is stable. This idea can also be applied to other areas like the labor market or university applications. In each of these three problems the parties need to have a strategy how to react to moves of the counterpart. If, for example, a university offers a place to a student and this is not the student's favorite one, he/she has to decide whether to accept the place (to be on the safe side) or still hope to get accepted from the favorite university (and as a consequence accept to be on the waiting list there). Not only the student needs to have a strategy how to deal with these situations but also universities have to incorporate different reactions of students into their application process (e.g. how many students to invite) [1].

The double sided matching problem is a dynamic game. The players of the two sets have incomplete information. The algorithm uses the information of both sets, therefore it has complete information.

An important aspect of the game or matching is that the rules have to be clear, which was mentioned by Roth and Sotomayor. The way agents are matched to each other influences the analysis of the problem. A possible rule might be that individuals like a student and university are only brought together if both parties agree to the matching. Other norms of a game might contain the way an individual proposes to another one or whether there exists a moderating individual [6, p. 492].

Prototype

Subsection

Text...

Prototype Evaluation: College Applications in Austria

Subsection

Text...

Prototype Evaluation: Labor Market in the European Union

Subsection

Text...

Summary

Subsection

Text...

Example

Lorem ipsum dolor sit amet, consectetur adipiscing elit. Aenean commodo ligula eget dolor. Aenean massa. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Donec quam felis, ultricies nec, pellentesque eu, pretium quis, sem. In enim justo, rhoncus ut, imperdiet a, venenatis vitae, justo. Nullam dictum felis eu pede mollis pretium. Integer tincidunt. Cras dapibus. Vivamus elementum semper nisi. Aliquam lorem ante, dapibus in, viverra quis, feugiat a, tellus:

$$\begin{aligned}(x+y)^3 &= (x+y)^2(x+y) \\ &= (x^2 + 2xy + y^2)(x+y) \\ &= (x^3 + 2x^2y + xy^2) + (x^2y + 2xy^2 + y^3) \\ &= x^3 + 3x^2y + 3xy^2 + y^3\end{aligned}\tag{7.1}$$

Phasellus viverra nulla ut metus varius laoreet. Quisque rutrum. Aenean imperdiet. Etiam ultricies nisi vel augue. Curabitur ullamcorper ultricies

Figure 7.1: One angry bird.



Heading on level 2 (subsection)

Lorem ipsum dolor sit amet, consectetur adipiscing elit.

$$A = \begin{bmatrix} A_{11} & A_{21} \\ A_{21} & A_{22} \end{bmatrix} \quad (7.2)$$

Aenean commodo ligula eget dolor. Aenean massa. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Donec quam felis, ultricies nec, pellentesque eu, pretium quis, sem.

Heading on level 3 (subsubsection)

Nulla consequat massa quis enim. Donec pede justo, fringilla vel, aliquet nec, vulputate eget, arcu. In enim justo, rhoncus ut, imperdiet a, venenatis vitae, justo. Nullam dictum felis eu pede mollis pretium. Integer tincidunt. Cras dapibus. Vivamus elementum semper nisi. Aenean vulputate eleifend tellus. Aenean leo ligula, porttitor eu, consequat vitae, eleifend ac, enim.

Heading on level 4 (paragraph) Lorem ipsum dolor sit amet, consectetur adipiscing elit. Aenean commodo ligula eget dolor. Aenean massa. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus. Donec quam felis, ultricies nec, pellentesque eu, pretium quis, sem. Nulla consequat massa quis enim.

Example for list (3*itemize)

- First item in a list
 - First item in a list
 - * First item in a list
 - * Second item in a list
 - Second item in a list
- Second item in a list

Example for list (enumerate)

1. First item in a list
2. Second item in a list
3. Third item in a list

List of Figures

7.1	One angry bird.	13
-----	-------------------------	----

List of Tables

References

- [1] David Gale and Lloyd S. Shapley. College admissions and the stability of marriage. *The American Mathematical Monthly*, 69(1):9–15, 1962.
- [2] Nicholas Barr. *Economics of the Welfare State*. Oxford University Press, 2012.
- [3] Colin Camerer. *Behavioral game theory: Experiments in strategic interaction*. Princeton University Press, 2003.
- [4] Robert Gibbons. An introduction to applicable game theory. *The Journal of Economic Perspectives*, 11:127–149, 1997.
- [5] Benedikt Fuest. Frequenzauktion droht zum mini-geschaeft zu werden, 2015.
- [6] Alvin E Roth and Marilda A Oliveira Sotomayor. *Two-sided matching: A study in game-theoretic modeling and analysis*. Number 18. Cambridge University Press, 1992.