

## Experiment 2

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### Problem 1.2.1: Two Sum

1. **Aim-** Given an array of integers `nums` and an integer `target`, return the indices of the two numbers such that they add up to `target`. Each input has exactly one solution, and you cannot use the same element twice.
2. **Algorithm –**
  1. Initialize an empty hash map (dict).
  2. Iterate through the `nums` array:
    - a. For each element `num`, calculate the complement: `complement = target - num`.
    - b. Check if the complement exists in the hash map:
      - i. If it does, return the indices of the complement and the current number.
      - ii. If it doesn't, add the current number and its index to the hash map.
  3. Return the indices of the two numbers that add up to the target.

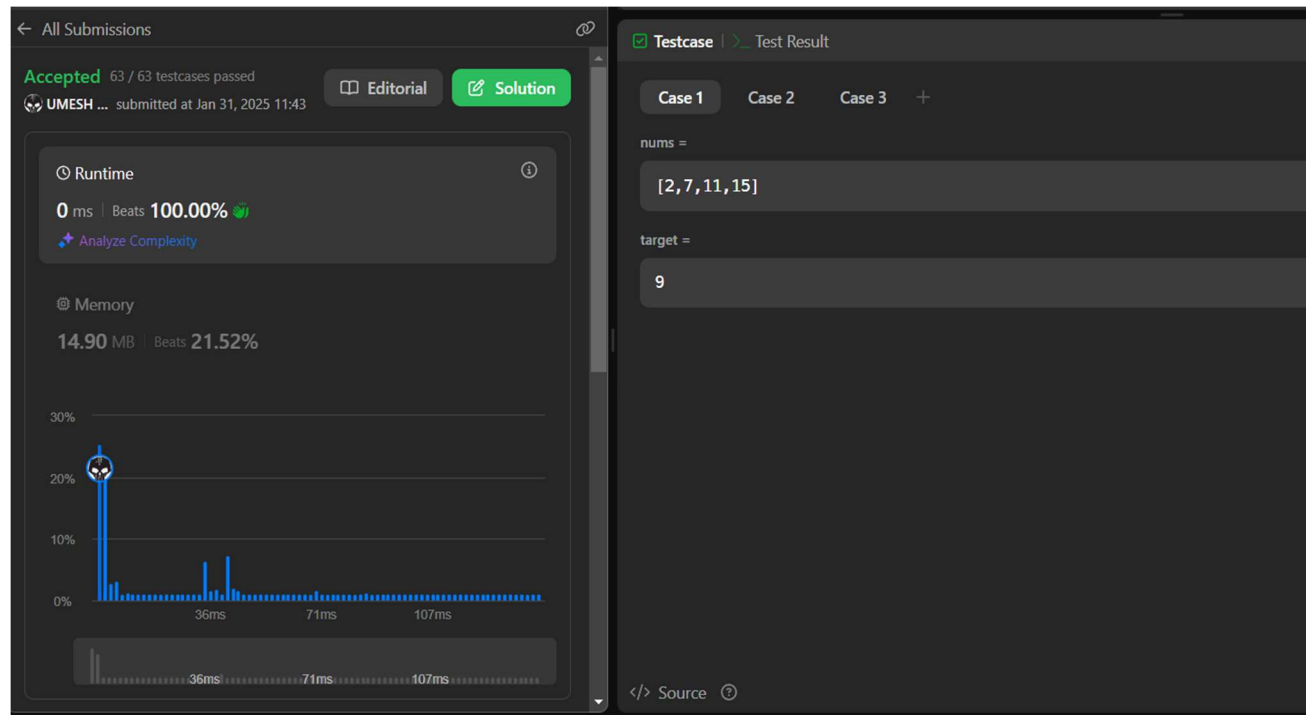
### 3. Code-

```
</> Code
C++ v Auto

1 #include <vector>
2 #include <unordered_map>
3 using namespace std;
4 class Solution {
5 public:
6     vector<int> twoSum(std::vector<int>& nums, int target) {
7         unordered_map<int, int> numMap;
8         for (int i = 0; i < nums.size(); i++) {
9             int complement = target - nums[i];
10            if (numMap.find(complement) != numMap.end()) {
11                return {numMap[complement], i};
12            }
13            numMap[nums[i]] = i;
14        }
15        return {};
16    }
17 };

Saved Ln 7, Col 40
```

## 4. Output –



## 5. Time Complexity –

- Time Complexity:  $O(n)$ , where  $n$  is the number of elements in the `nums` array. This is because we only iterate through the array once.
- Space Complexity:  $O(n)$ , since we store each number and its index in the hash map.

## Problem 1.2.2: Jump Game II

**1. Aim -** You are given a 0-indexed array `nums` of length `n`. You are initially positioned at `nums[0]`. Each element `nums[i]` represents the maximum length of a forward jump from index `i`. Return the minimum number of jumps to reach `nums[n - 1]`.

### 2. Algorithm-

1. Initialize `jumps` as 0 (the number of jumps needed), `current_end` as 0 (the farthest index you can reach in the current jump), and `farthest` as 0 (the farthest index you can reach so far).
2. Loop through each index in `nums` (except the last index):
  - a. Update `farthest` to the maximum of `farthest` and `i + nums[i]`.
  - b. If the current index `i` is the last index of the current jump (`i == current_end`):
    - i. Increment the jumps.
    - ii. Set `current_end` to `farthest`.
3. Return `jumps` when `current_end` reaches or exceeds the last index (`n - 1`).

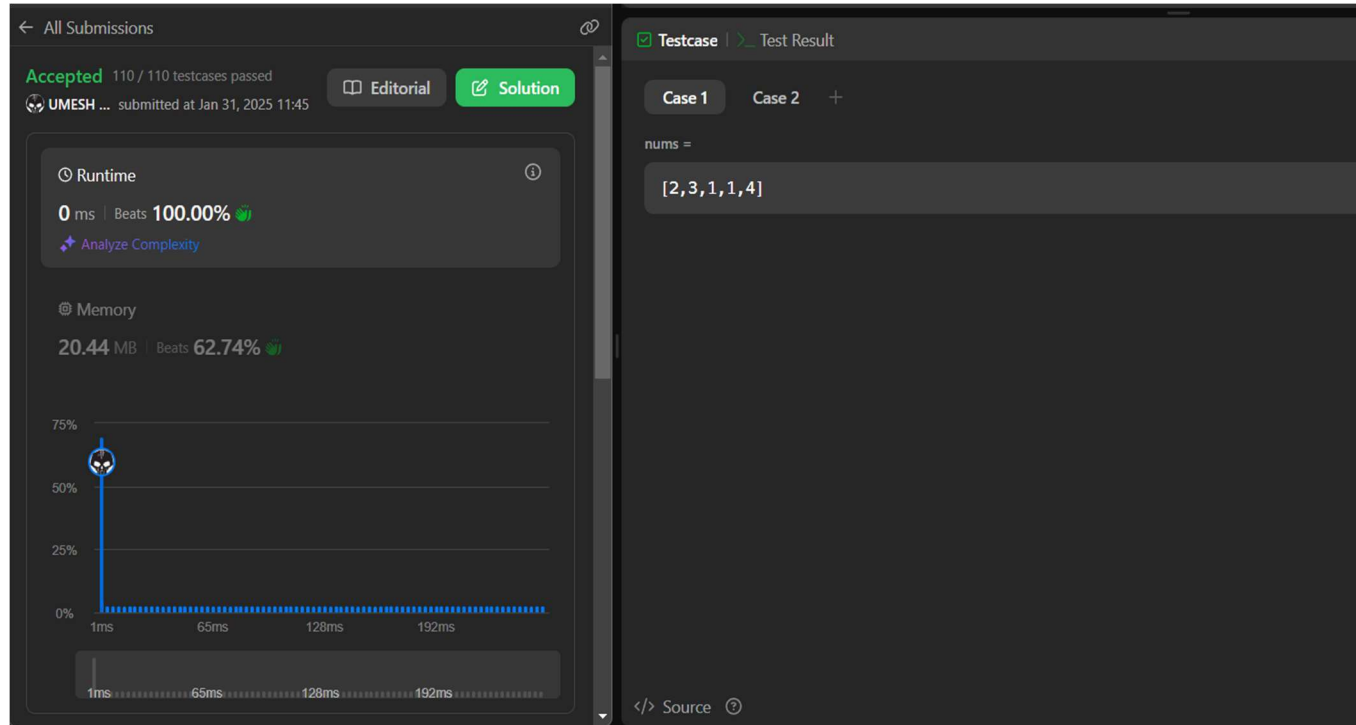
### 3. Code –

```
</> Code
C++ v Auto

1  #include <vector>
2  using namespace std;
3
4  class Solution {
5  public:
6      int jump(vector<int>& nums) {
7          int n = nums.size();
8          if (n <= 1) return 0;
9
10         int jumps = 0, curEnd = 0, farthest = 0;
11
12         for (int i = 0; i < n - 1; i++) {
13             farthest = max(farthest, i + nums[i]);
14
15             if (i == curEnd) {
16                 jumps++;
17                 curEnd = farthest;
18                 if (curEnd >= n - 1) break;
19             }
20         }
21
22         return jumps;
23     }
24 };

Saved Ln 8, Col 31
```

## 4. Output-



## 5. Time Complexity –

- Time Complexity:  $O(n)$ , where  $n$  is the number of elements in the `nums` array. We iterate through the array once.
- Space Complexity:  $O(1)$ , since we use only a few extra variables.

## 6. Learning Outcomes -

- Using a hash map (`unordered_map` in C++ / `HashMap` in Java) to store visited numbers for  $O(1)$  average-time complexity lookup.
- Understanding how to store values vs indices to return correct results
- Understanding Greedy Strategy-Keep track of the farthest reachable index. Iterate through the array, updating the maximum reach.
- Using a Single-pass vs Two-pass Approach -Two-pass-Store elements first, then find pairs. Single-pass to Store and check simultaneously.