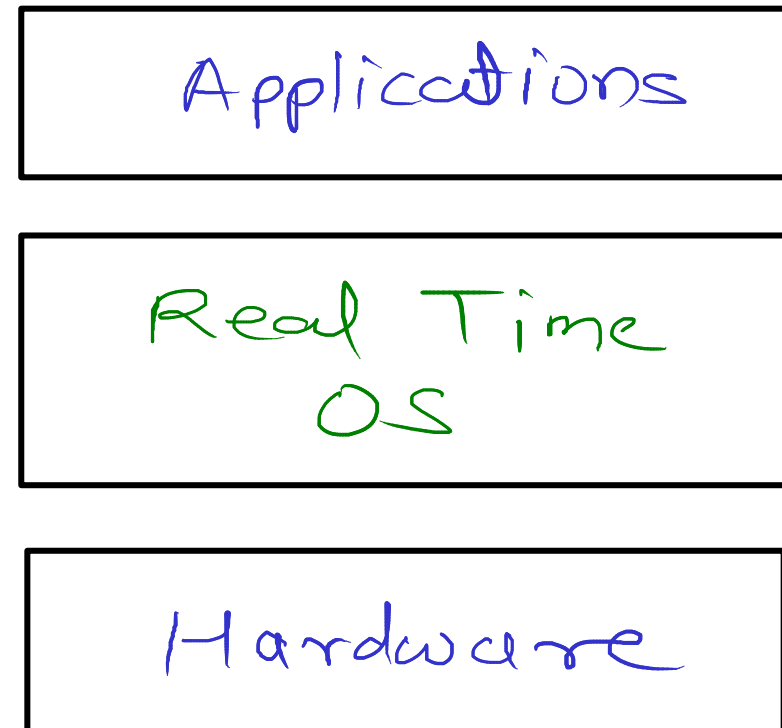


# Real Time Operating System



\* result should be accurate  
\* result should be calculated in minimal/deterministic

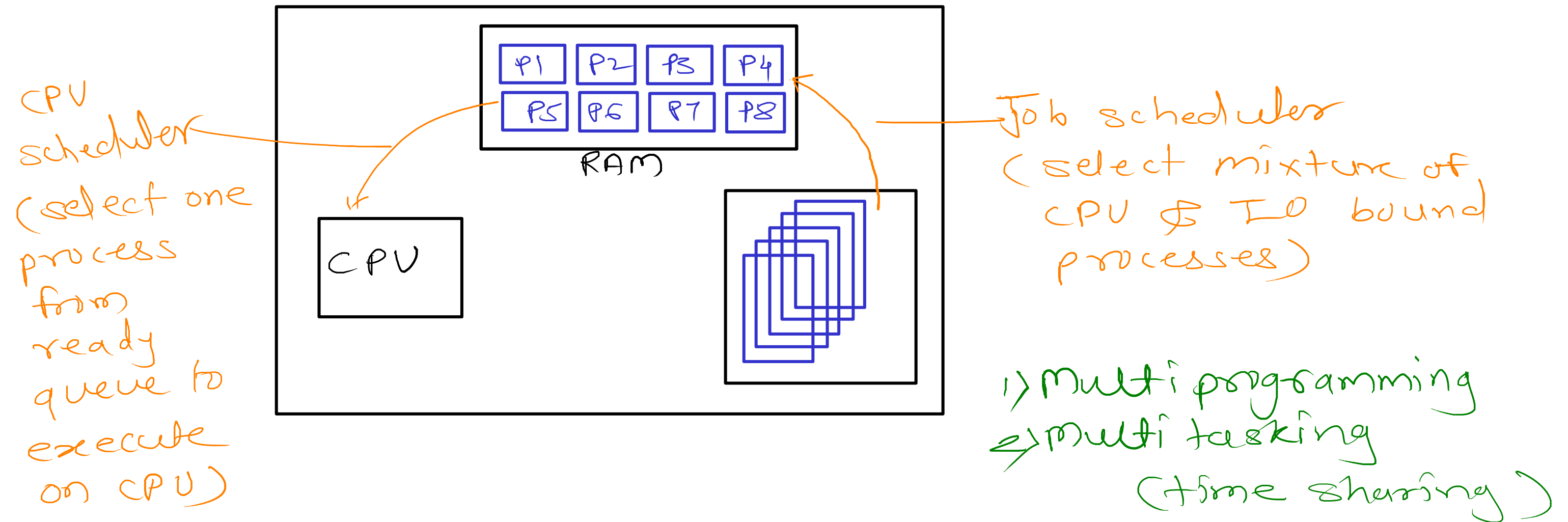
## Types of RTOS

1. Hard Real Time OS
2. Soft Real Time OS
3. Firm Real Time OS

## Functions of RTOS

1. Hardware Abstraction
2. Task Management
3. Memory Management
4. CPU Scheduling
5. Interrupt Handling

# CPU Scheduling

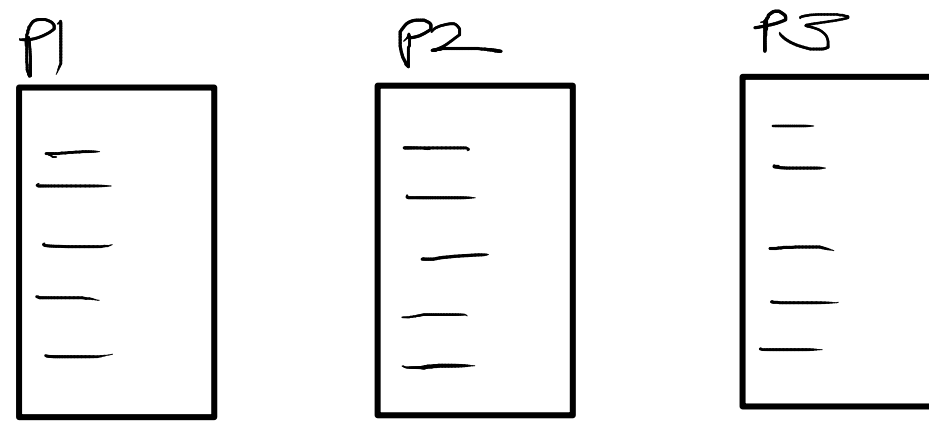


## OS data structures

- 1) Job queue/process list
- 2) Ready queue ←
- 3) Waiting queues

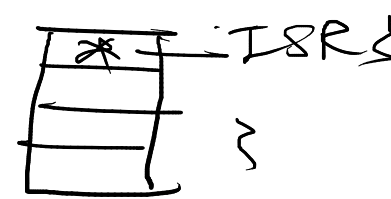
## Algorithms:

- 1) FCFS
- 2) SJF
- 3) Priority
- 4) RR
- 5) Fair Share



interrupt\_handler() {

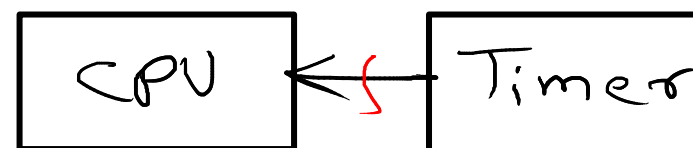
- 1) Save execution context of current save into its PCB
- 2) find address of ISR from IVT
- 3) call ISR
- 4) pid = cpu\_scheduler();
- 5) cpu\_dispatcher(pid);



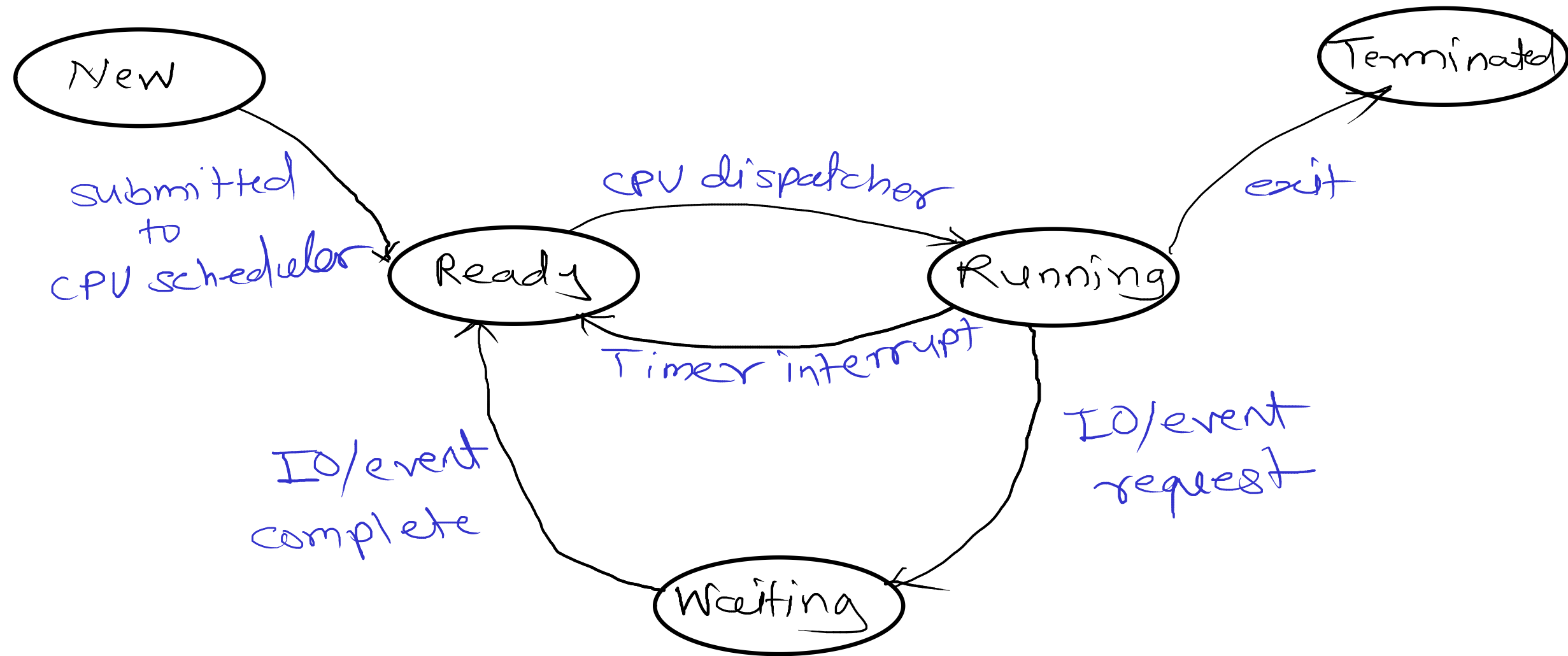
```
int cpu_scheduler() {
    if (remaining > 0)
        select some process
        to be executed next
    else
        select new process
        to be executed next
    return pid;
}
```

```
cpu_dispatcher() {
    restore execution
    context of
    selected process
}
```

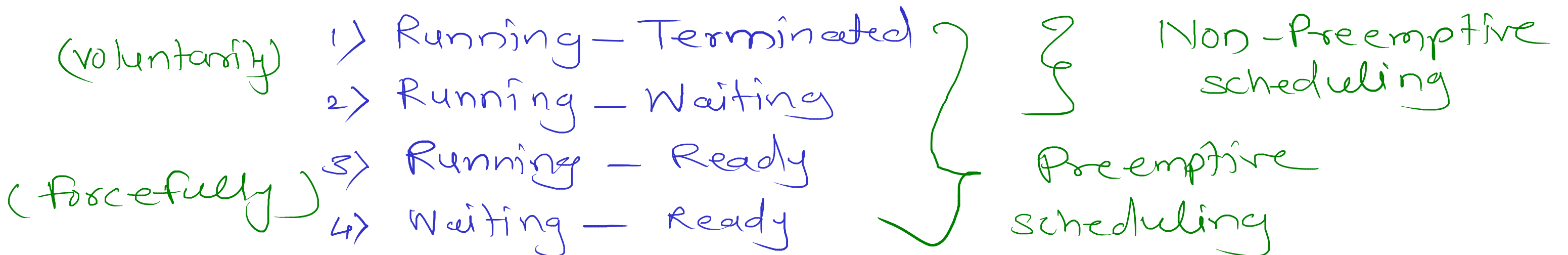
}



## Process Life Cycle



## Types of Scheduling



# CPU Scheduling Criterias

## 1. CPU Utilization (Ideal : Max)

- desktop system - 70%
- server system - 90%

## 2. Throughput (Ideal : Max)

- Amount of work done in unit time

## 3. Waiting Time (Ideal : Min)

- waiting for CPU
- total time spent by process into ready queue

## 4. Response Time (Ideal : Min)

- time from arrival of process into ready queue upto first time getting scheduled

## 5. Turn Around Time (Ideal : Min)

- total time of process spent into memory
- time from creation to termination of process

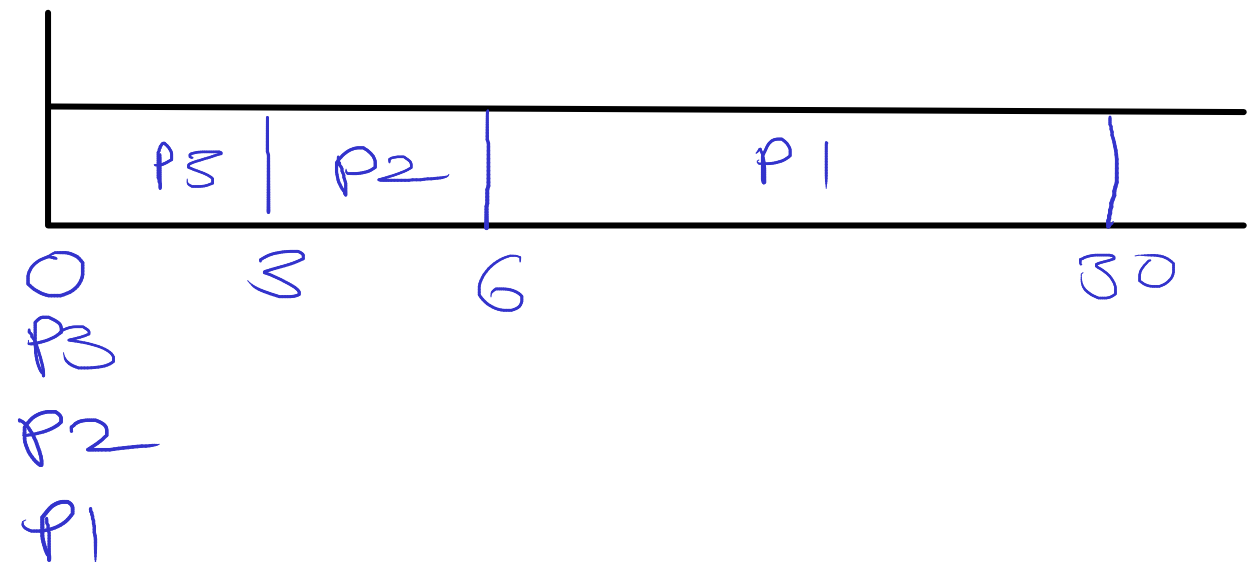
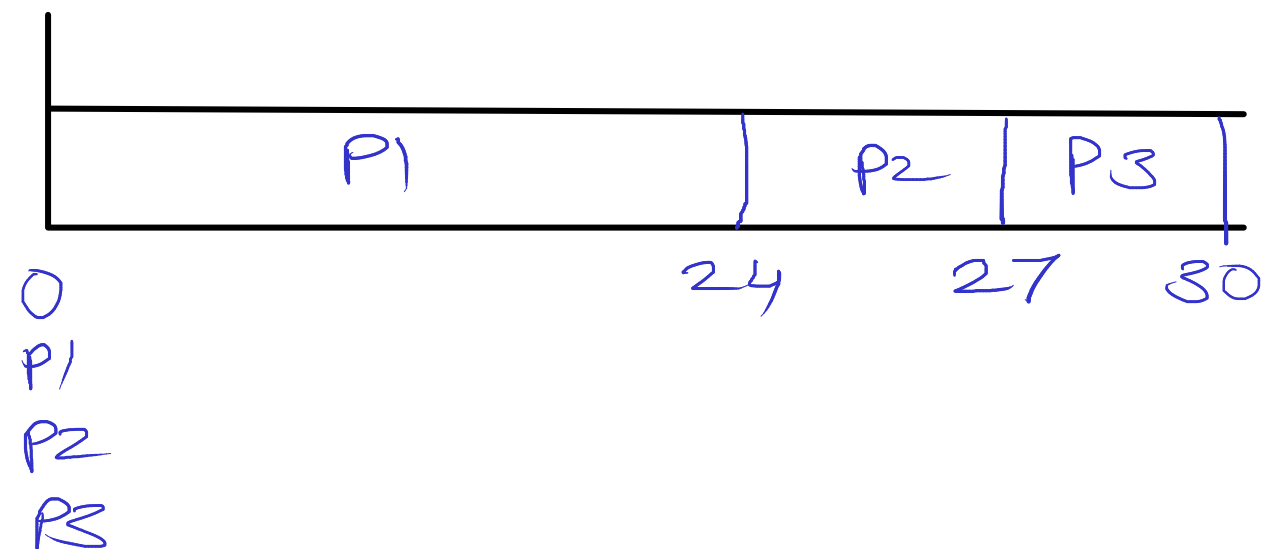
$$TAT = \text{CPU waiting time} + \text{CPU burst} + \text{IO waiting time} + \text{IO burst}$$

# FCFS (First Come First Serve) (Non-preemptive)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	24	0	0	24
P2	0	3	24	24	27
P3	0	3	27	27	30

Process	Arrival	CPU Burst	WT	RT	TAT
P3	0	3	0	0	3
P2	0	3	3	3	6
P1	0	24	6	6	30

## Gantt's chart



## Convoy Effect

- due to arrival of longer process early, all other processes has to wait for longer time to get CPU access

- we can not control the sequence of process arrival into ready queue

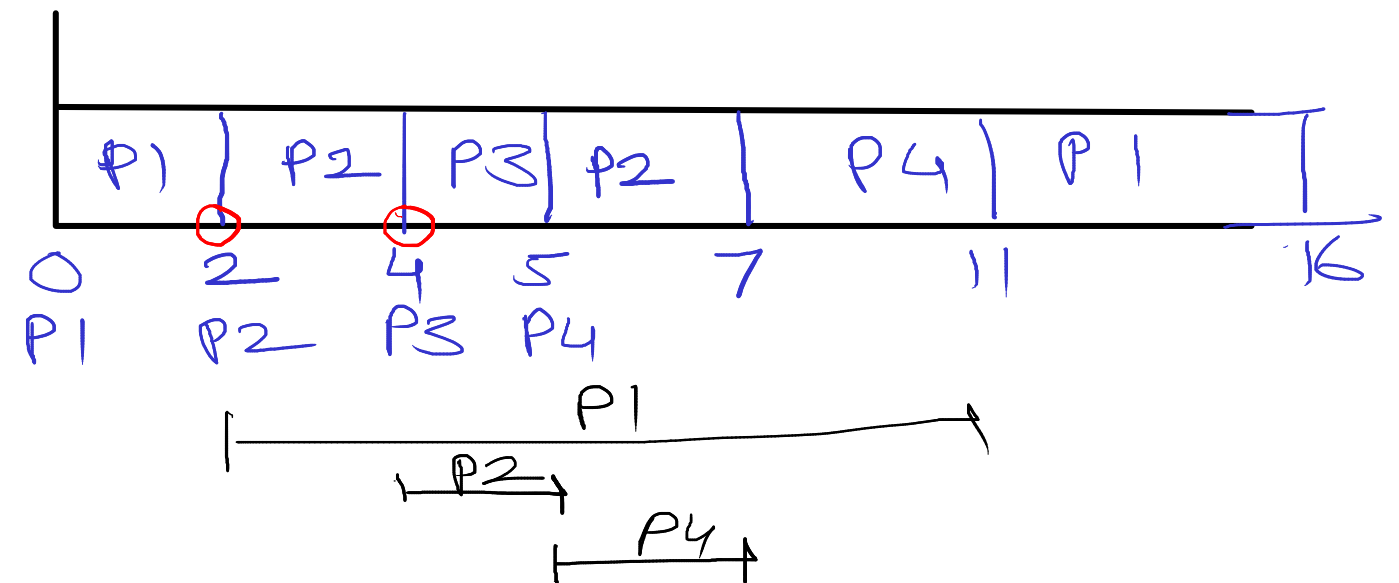
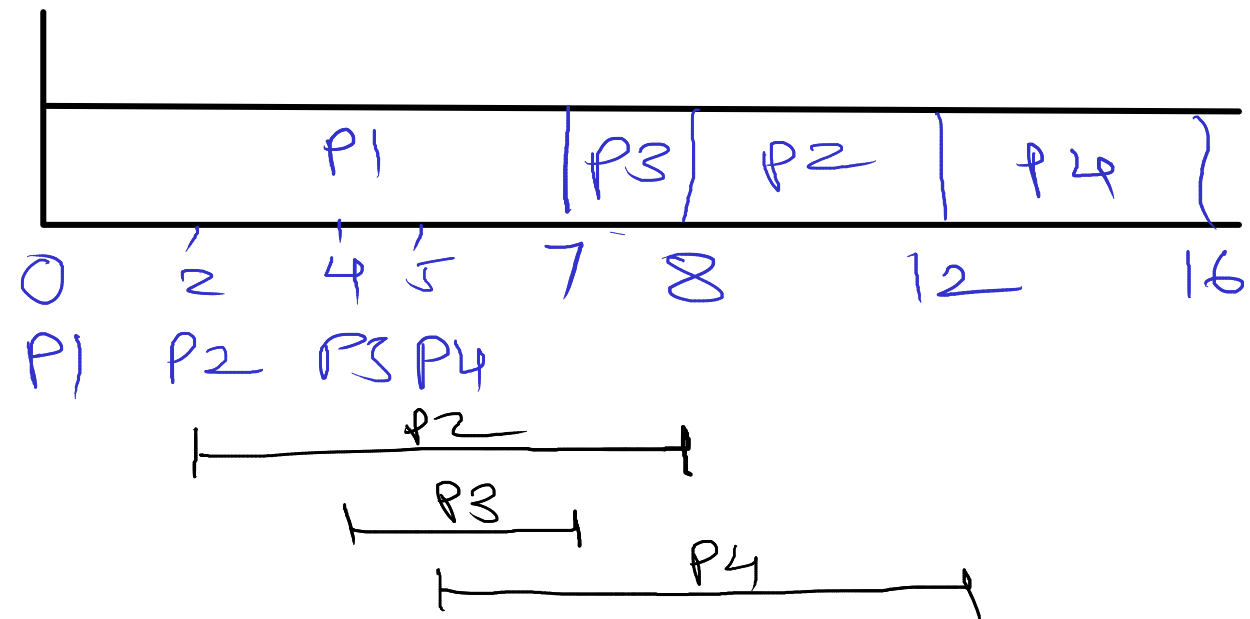
(Non-Preemptive)

**SJF** (Shortest Job First)  
(Preemptive)

(Shortest Remaining Time First)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	7	0	0	7
P2	2	4	6	0	10
P3	4	1	3	0	4
P4	5	4	7	0	11

Process	Arrival	CPU Burst	Remain time	WT	RT	TAT
P1	0	7	5	9	0	16
P2	2	4	2	1	0	5
P3	4	1	0	0	0	1
P4	5	4	4	2	2	6



Priority

(Non-Preemptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	0	1	1(H)
P3	0	2	4(L)
P4	0	5	2

WT	RT	TAT
6	6	16
0	0	1
16	16	18
1	1	6

(Preemptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	1	1	1
P3	3	2	4
P4	0	5	2

WT	RT	TAT
6	6	16
0	0	1
13	13	15
4	0	6

