

Rod Cutting

Given a rod of length n and prices of piece of rod length i , where $1 \leq i \leq n$. Find the optimal way to cut the given rod into smaller rods in order to maximize profit.

e.g.

Given a rod of length 4m and the prices of different lengths of rods

Length (l)	1	2	3	4	5	6	7	8	9	10
Price (pi)	1	5	8	9	10	17	17	20	24	30

Possibilities

- We don't cut the given rod of length 4 (9)
- We cut the rod into 4 pieces (3 cuts), each of length 1 ($1 + 1 + 1 + 1 = 4$)
- We cut the rod into 2 pieces (1 cut), one of length 1 and other of length 3 ($1 + 8 = 9$)
- We cut the rod into 3 pieces (2 cuts), two pieces of length 1 and other of length 2 ($1 + 1 + 5 = 7$)
- We cut the rod into 2 pieces (1 cut), each of length 2 ($5 + 5 = 10$)

e.g.

Length of the Rod is 5m. $n = 5m$, Prices = {2, 5, 7, 3}

To get to the solution to the actual problem (given rod of length 5m and prices of rod length 1m(2\$), 2m(5\$), 3m(7\$), 4m(3\$), what is the maximum profit we can get?

We will first solve the subproblems.

e.g.

- Given rod of length 1m and when we can only make cut of 1m, what is the maximum profit we can get (answer is 2\$ {1m})
- Given rod of length 4m and when we can only make cut of 1m, what is the maximum profit we can get (answer is 8\$ {1m, 1m, 1m, 1m})
- Given rod of length 2m and when we can only make cut of 1m and 2m, what is the maximum profit we can get (answer is 5\$ {2m})
- Given rod of length 4m and when we can make cut of 1m, 2m, 3m and 4m, what is the maximum profit we can get (answer is 10\$ {2m, 2m})

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0						
1						
2						
3						
4						

When we make a cut of 0 meter, the maximum profit we can make is 0\$.

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1						
2						
3						
4						

Similarly when the rod length is 0 meter regardless of the different cuts we are allowed to make, the maximum profit we can make is 0\$.

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0					
2	0					
3	0					
4	0					

When we can only make cuts of 1m

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0					
3	0					
4	0					

When we can only make cuts of 1m and 2m

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0	2	5	7	10	12
3	0					
4	0					

When we can only make cuts of 1m, 2m and 3m

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0	2	5	7	10	12
3	0	2	5	7	10	12
4	0					

When we can make cuts of 1m, 2m, 3m and 4m

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0	2	5	7	10	12
3	0	2	5	7	10	12
4	0	2	5	7	10	12

When we can't make i^{th} meter cut from the rod of length j meter , because $i > j$

$$\text{dpTable}[i][j] = \text{dpTable}[i-1][j]$$

When we can make i^{th} meter cut from the rod of length j meter , because i is not greater than j ($i \leq j$)

$$\text{dpTable}[i][j] = \text{Max}(\text{dpTable}[i-1][j], \text{prices}[i-1] + \text{dpTable}[i][j-i])$$

$n = 5\text{m}$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0	2	5	7	10	12
3	0	2	5	7	10	12
4	0	2	5	7	10	12

$\text{dpTable}[2][1]$ = Since rod length is only 1m and we can't make cut of 2m from a rod of length just 1 m,
 $\text{dpTable}[2][1] = \text{dpTable}[1][1] = 2$

$\text{dpTable}[2][2] = \text{Max}(\text{dpTable}[1][2], 5 + \text{dpTable}[2][2-2]) = \text{Max}(4, 5) = 5$

$\text{dpTable}[2][3] = \text{Max}(\text{dpTable}[1][3], 5 + \text{dpTable}[2][3-2]) = \text{Max}(6, 7) = 7$

$\text{dpTable}[2][4] = \text{Max}(\text{dpTable}[1][4], 5 + \text{dpTable}[2][4-2]) = \text{Max}(8, 10) = 10$

$\text{dpTable}[2][5] = \text{Max}(\text{dpTable}[1][5], 5 + \text{dpTable}[2][5-2]) = \text{Max}(10, 12) = 12$

$\text{dpTable}[3][3] = \text{Max}(\text{dpTable}[2][3], 7 + \text{dpTable}[3][3-3]) = \text{Max}(7, 7) = 7$

$\text{dpTable}[3][4] = \text{Max}(\text{dpTable}[2][4], 7 + \text{dpTable}[3][4-3]) = \text{Max}(10, 9) = 10$

$\text{dpTable}[4][5] = \text{Max}(\text{dpTable}[3][5], 3 + \text{dpTable}[4][5-4]) = \text{Max}(12, 5) = 12$

But what cuts we made?

$n = 5m$, Prices = {2, 5, 7, 3}

Cut/Length	0(m)	1(m)	2(m)	3(m)	4(m)	5(m)
0	0	0	0	0	0	0
1	0	2	4	6	8	10
2	0	2	5	7	10	12
3	0	2	5	7	10	12
4	0	2	5	7	10	12

Cuts : 2m, 2m, 1m

Implementation

```
public class App {  
    public static void main(String[] args) {  
        int[] prices = { 2, 5, 7, 3 };  
        int rodLength = 5;  
        RodCutting rodCutting = new RodCutting(prices, rodLength);  
        int maximumProfit = rodCutting.solve();  
        System.out.println("Maximum Profit : " + maximumProfit);  
        rodCutting.showCuts();  
    }  
}
```

```

public class RodCutting {
    int [] prices;
    int rodLength;
    int[][] dpTable;
    public RodCutting(int[] prices, int rodLength) {
        this.prices = prices;
        this.rodLength = rodLength;
        dpTable = new int[prices.length + 1][rodLength + 1];
    }

    public int solve() {
        ....
    }

    public void showCuts() {
        ....
    }
}

```

```

public int solve() {
    for (int i = 1; i <= prices.length; i++) {
        for (int j = 1; j <= rodLength; j++) {
            if (i > j) {
                dpTable[i][j] = dpTable[i - 1][j];
            } else {
                dpTable[i][j] = Math.max(dpTable[i - 1][j],
                                         prices[i - 1] + dpTable[i][j - i]);
            }
        }
    }
    return dpTable[prices.length][rodLength];
}

```

```

public void showCuts() {
    for(int i= prices.length, j=rodLength; i > 0; ) {
        if(dpTable[i][j] != 0 && dpTable[i][j] != dpTable[i-1][j]) {
            System.out.println("Make a cut of " + i+"m");
            j = j -i;
        }else {
            i--;
        }
    }
}

```

