A Choreographic View of Smart Contracts

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A tutorial @ FORTE 2025, Lille

Work partly supported by the PRIN 2022 PNRR project DeLiCE (F53D23009130001)

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What's up doc?

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Prologue . . . . An inspiring initiative

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Prologue .... An inspiring initiative

Act I .... A coordination framework

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Act II .... Some tool support

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Epilogue . . . . Work in progress

– Prologue –

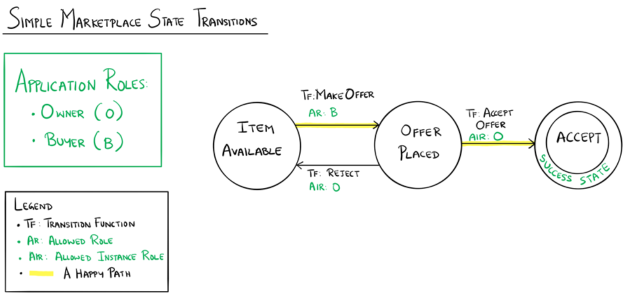
[An inspiring initiative]

– Prologue –

[An inspiring initiative]

A nice sketch! [6, 7]

A smart contract among Owners and Buyers



initially buyers can make offers
then
  either an owner can accept an offer and the protocol stops
  or the offer is rejected and the protocol restarts

A nice sketch! [6, 7]

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What did we just see?

A smart contract looks like

a choreographic model
global specifications determine the enabled actions along the evolution of the protocol

a typestate
In OOP, "can reflects how the legal operations on imperative objects can change at runtime as their internal state changes." [3]

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So, we saw an interesting model where

distributed components coordinate through a global specification

which specifies how actions are enabled along the computation

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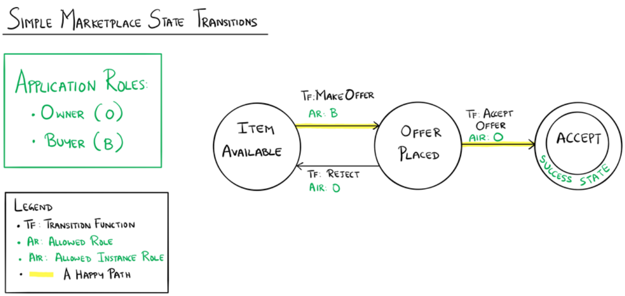

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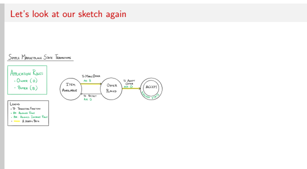

The diagram specifies a lot...

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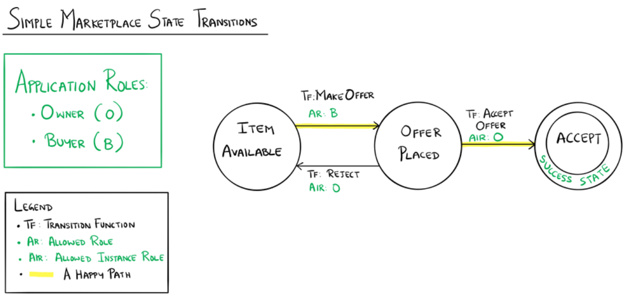
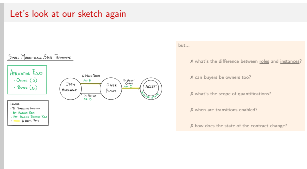


but...

- ✗ what's the difference between roles and instances?
- ✗ can buyers be owners too?
- ✗ what's the scope of quantifications?
- ✗ when are transitions enabled?
- ✗ how does the state of the contract change?



The diagram specifies a lot...

1. is the sketch giving semantics to roles and instances?
2. not forbidden...however what if we wanted to separate the roles?
3. from [7]: "The transitions between the Item Available and the Offer Placed states can continue until the owner is satisfied with the offer made." so, after a rejection, the new offer must be from the original buyer or a new one?
4. ok
5. should the price of the item remain unchanged when the owner rejects offers?

Let's go formal!

Our first attempt was to "look for into our toolbox", but

- ✗ are known notions of well-formedness suitable?
- ✗ data-awareness is crucial
- ✓ we got roles okay, but
- ✗ limitations on instances of roles
- ✗ instances can have one role only

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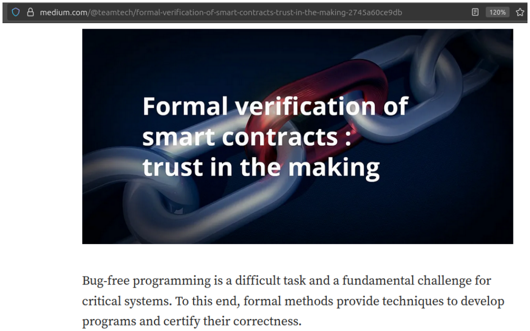
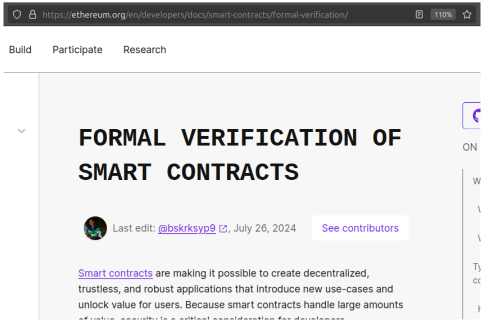


https://medium.com/@teamtech/formal-verification-of-smart-contracts-trust-in-the-making-2745a60ce9db



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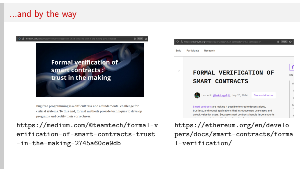

– Act I –

[A coordination framework]

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Basic concepts and notation

Participants $\mathsf{p}, \mathsf{p}', \ldots$

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states of the coordinator determine which operations each roles is entitled to invoke

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We assume that sorts can be inferred; **TRAC** instead requires to assign sorts explicitly

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Expressions are standard but for state variables occurring in rhs e must have the old _ qualifier; this concept will be used in the definition of (progress for) well-formedness

We adapt the mechanism based on the old keyword from the Eiffel language [5] which, as explained in [4] is necessary to render assignments into logical formulae since e.g., $x = x+1 \iff$ False.

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- $A, A', \ldots$ range over finite sets of assignments where each variable can be assigned at most once

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Data-Aware FSMs

A DAFSM c on roles $R_1, \ldots R_m$ and state variables $u_1, \ldots, u_n$ is a finite-state machine "instantiated" by a participant p whose transitions are decorated as follows[1]

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$$\xrightarrow{\text{new } \mathsf{p} : \mathsf{R} \triangleright \mathsf{start}(\mathsf{c}, \cdots, T_i\ \mathsf{x}_i, \cdots)\ \{\cdots \mathsf{u}_j := \mathsf{e}_j \cdots\}} \bigcirc$$

c is freshly created by p which also initilise state variables u_j with expressions e_j which are built on state variables and parameters x_i

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start is a "built-in" (and pleonastic) function name

each state variable is declared and initialises with type-consistent expressions on state variables and parameters x_i

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$$\circ \xrightarrow{\{\gamma\}\ \pi \triangleright \text{f}(\cdots, T_i\ x_i, \cdots)\ \{A\}} \circ$$

where γ is a guard (ie a boolean expression) and

$$\pi \ ::= \ \text{new } p : R \ \big| \ \text{any } p : R \ \big| \ p$$

is a qualified participant calling f with parameters x_i
state variables are reassigned according to A if the invocation is successful

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γ predicates over state variables and formal parameters of its transition; guards have to be satisfied for the invocation to succeed: an invocation that makes the guard false is rejected

- new p: R specifies that p must be a fresh participant with role R
- any p: R qualifies p as an existing participant with role R
- p refers to a participant in the scope of a binder
- invocations from non-suitable callers are rejected

the variables occurring in the right-hand side of assignments in A are either state variables or parameters of the invocation

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$$\circ \xrightarrow{\ell} \circledcirc$$

accepting states are denoted as usual

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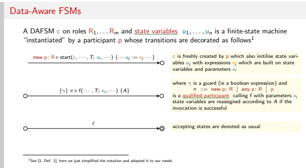
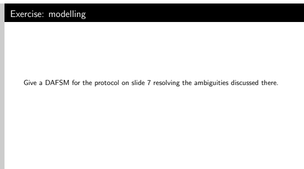
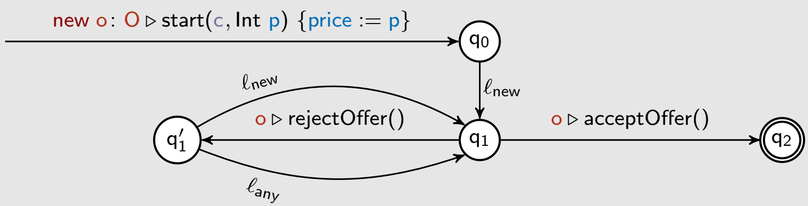


Exercise: modelling

Give a DAFSM for the protocol on slide 7 resolving the ambiguities discussed there.

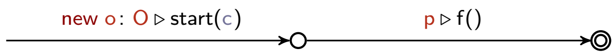


Let $\ell_{\mathsf{new}} = \{\mathsf{newOffer} > 0\}$ new b: B ▷ makeOffer(Int newOffer) $\{\mathsf{offer} := \mathsf{newOffer}\}$
and $\ell_{\mathsf{any}} = \{\mathsf{newOffer} > 0\}$ any b: B ▷ makeOffer(Int newOffer) $\{\mathsf{offer} := \mathsf{newOffer}\}$



A new participant o acts as owner O for a coordinator c assigning an initial value p to the state variable price in the initial state q_0 where the only enabled function is makeOffer(Int offer). The first buyer b invoking this function with an actual parameter newOffer, satisfying the guard $\mathsf{newOffer} > 0$, moves the protocol to state q_1 while recording the new offer in the coordinator state with the assignment $\mathsf{offer} := \mathsf{newOffer}$. Contextually, the state of the coordinator records that the caller b plays role B.

Not all DAFSMs “make sense”



names’ freeness

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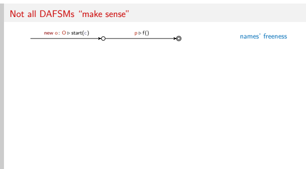

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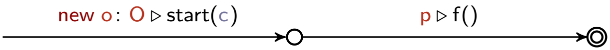

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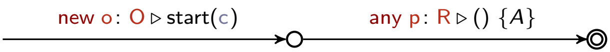
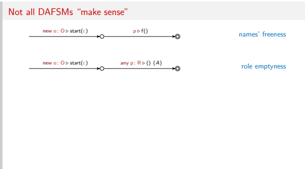


role emptyness



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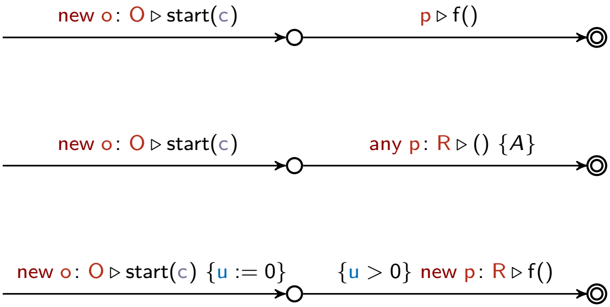

names’ freeness

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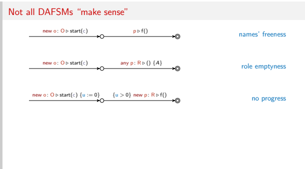
no progress



Not all DAFSMs "make sense"

$$\xrightarrow{\text{new } o \colon O \triangleright \text{start}(c)} \bigcirc \xrightarrow{p \triangleright f()} ◎$$

names' freeness

$$\xrightarrow{\text{new } o \colon O \triangleright \text{start}(c)} \bigcirc \xrightarrow{\text{any } p \colon R \triangleright () \ \{A\}} ◎$$

role emptyness

$$\xrightarrow{\text{new } o \colon O \triangleright \text{start}(c) \ \{u := 0\}} \bigcirc \xrightarrow{\{u > 0\} \ \text{new } p \colon R \triangleright f()} ◎$$

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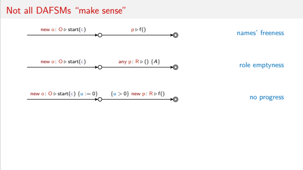

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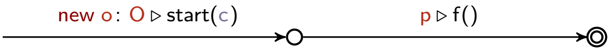

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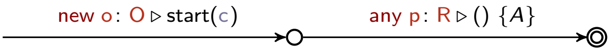
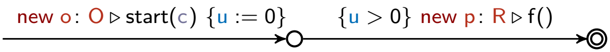
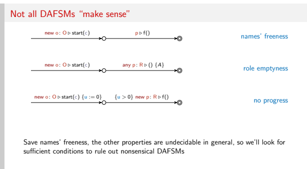


role emptyness



no progress

Save names' freeness, the other properties are undecidable in general, so we'll look for sufficient conditions to rule out nonsensical DAFSMs



Closed DAFSMs

Binders: parameter declarations in function calls, new $\mathsf{p} : \mathsf{R}$, and any $\mathsf{p} : \mathsf{R}$

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Closed DAFSMs

Binders: parameter declarations in function calls, new p: R, and any p: R

p is bound in $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright \mathsf{f}(\cdots,\ T_i\ \mathsf{x}_i, \cdots)\ \{A\}} \circ$ if, for some role R,

$$\pi = \mathsf{new}\ \mathsf{p}\colon \mathsf{R} \quad \text{or} \quad \pi = \mathsf{any}\ \mathsf{p}\colon \mathsf{R} \quad \text{or} \quad \text{there is } i \text{ s.t. } \mathsf{x}_i = \mathsf{p} \text{ and } T_i = \mathsf{R}$$

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The occurrence of p is bound in a path

$$\sigma \quad \circ \xrightarrow{\{\gamma\}\ \mathsf{p} \triangleright \mathsf{f}(\cdots)\ \{A\}} \circ \cdots$$

if p is bound in a transition of σ

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$\pi =$ new p: R or $\pi =$ any p: R or there is i s.t. $\mathsf{x}_i = \mathsf{p}$ and $T_i = \mathsf{R}$
The occurrence of p is bound in a path
$\sigma \quad \circ \xrightarrow{\{\gamma\}\ \mathsf{p} \triangleright \mathsf{f}(\cdots)\ \{A\}} \circ \cdots$
if p is bound in a transition of σ

Closed DAFSMs

Binders: parameter declarations in function calls, new p: R, and any p: R

p is bound in $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright \mathsf{f}(\cdots, T_i\ \mathsf{x}_i, \cdots)\ \{A\}} \circ$ if, for some role R,

$$\pi = \text{new p: R} \quad \text{or} \quad \pi = \text{any p: R} \quad \text{or} \quad \text{there is } i \text{ s.t. } \mathsf{x}_i = \mathsf{p} \text{ and } T_i = \mathsf{R}$$

The occurrence of p is bound in a path

$$\sigma \quad \circ \xrightarrow{\{\gamma\}\ \mathsf{p} \triangleright \mathsf{f}(\cdots)\ \{A\}} \circ \quad \cdots$$

if p is bound in a transition of σ

A DAFSM is closed if all occurrences of participant variables are bound in the paths of the DAFSM they occur on

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$\sigma \quad \circ \xrightarrow{\{\gamma\}\ \mathsf{p} \triangleright \mathsf{f}(\cdots)\ \{A\}} \circ \quad \cdots$

if p is bound in a transition of σ

A DAFSM is closed if all occurrences of participant variables are bound in the paths of the DAFSM they occur on

Role emptyness

A transition $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright \mathsf{f}(\cdots, T_i\ \mathsf{x}_i, \cdots)\ \{A\}} \circ$ <u>expands</u> role R if $\pi = \mathsf{new}\ \mathsf{p} \colon \mathsf{R}$ or there is i s.t. $\mathsf{x}_i = \mathsf{p}$ and $T_i = \mathsf{R}$

Role R is <u>expanded</u> in a path

$$\sigma \quad \circ \xrightarrow{\{\gamma\}\ \mathsf{any}\ \mathsf{p} \colon \mathsf{R} \triangleright \mathsf{f}(\cdots)\ \{A\}} \circ \quad \cdots$$

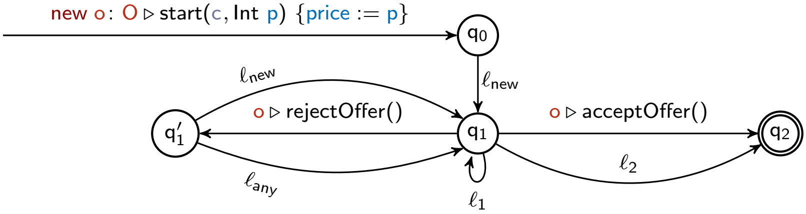
if a transition in σ expands R

A DAFSM <u>expands</u> R if all its paths expand R and is <u>(strongly) empty-role free</u> if it expands all its roles

'expands' means register a new participant with that role in the protocol (the participant might already be registered with a different role)

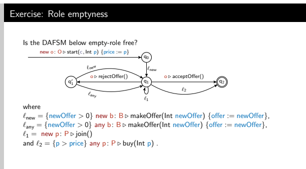
Exercise: Role emptyness

Is the DAFSM below empty-role free?



where

$\ell_{new} = \{\text{newOffer} > 0\}$ new b: B ▷ makeOffer(Int newOffer) {offer := newOffer},
$\ell_{any} = \{\text{newOffer} > 0\}$ any b: B ▷ makeOffer(Int newOffer) {offer := newOffer},
$\ell_1 =$ new p: P ▷ join()
and $\ell_2 = \{\text{p} > \text{price}\}$ any p: P ▷ buy(Int p) .



No, because of the paths excluding the self-loop on the role P.
We can fix the problem by adding adding a parameter of type p to the start transition.

Progress

A DAFSM with state variables $\mathsf{u}_1, \ldots, \mathsf{u}_n$ is consistent if it is closed and the following implication holds for each transition $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright \mathsf{f}(\cdots, T_i\ \mathsf{x}_i, \cdots)\ \{A\}} (\mathsf{s})$

$$\forall\!\!\!\forall_U\, \exists\!\!\!\exists_X \left(\gamma\{\mathsf{old}\ \mathsf{u}_1, \ldots, \mathsf{old}\ \mathsf{u}_n / \mathsf{u}_1, \ldots, \mathsf{u}_n\} \ \wedge\ \gamma_A \implies \gamma_{\mathsf{s}}\right)$$

where

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$$\forall\!\!\!\forall_U\, \exists\!\!\!\exists_X \left(\gamma\{\mathsf{old}\ \mathsf{u}_1, \ldots, \mathsf{old}\ \mathsf{u}_n / \mathsf{u}_1, \ldots, \mathsf{u}_n\} \ \wedge\ \gamma_A \implies \gamma_{\mathsf{s}}\right)$$

where

for a finite set of symbols Z, $\forall\!\!\!\forall_Z\,(\_)$ and $\exists\!\!\!\exists_Z\,(\_)$ are the universal and existential closures of a logical formula on the symbols in Z

Progress

A DAFSM with state variables $u_1, \ldots, u_n$ is consistent if it is closed and the following implication holds for each transition $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright f(\cdots, T_i\ x_i, \cdots)\ \{A\}} (s)$

$$\forall\!\!\forall_U\ \exists\!\!\exists_X\ \left(\gamma\{\text{old } u_1, \ldots, \text{old } u_n / u_1, \ldots, u_n\} \ \wedge\ \gamma_A \implies \gamma_s\right)$$

where

$$U = \{u_i, \text{old } u_i\}_{1 \le i \le n}$$
$$X = \{x \mid \exists i :\ x = x_i \text{ or } x \text{ is a parameter of an outgoing transition of } s\}$$

$$\gamma_s = \begin{cases} \text{the disjunction of guards of the outgoing transitions of } s & \text{is not accepting} \\ \text{True} & \text{otw} \end{cases}$$

$$\gamma_A = \bigwedge_{u := e \in A} u = e \ \wedge \bigwedge_{u \notin A} u = \text{old } u$$

Progress
A DAFSM with state variables $u_1, \ldots, u_n$ is consistent if it is closed and the following implication holds for each transition $\circ \xrightarrow{\{\gamma\}\ \pi \triangleright f(\cdots, T_i\ x_i, \cdots)\ \{A\}} (s)$
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for a finite set of symbols Z, $\forall\!\!\forall_Z\,(\_)$ and $\exists\!\!\exists_Z\,(\_)$ are the universal and existential closures of a logical formula on the symbols in Z

$u \notin A$
iff
for all $v := e \in A$, $u \neq v$ and $\text{old } u$ does not occur in e

Exercise: Consistency

todo

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todo

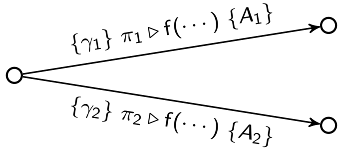
Determinism

Let $\_\#\_$ be the least binary symmetric relation s.t.

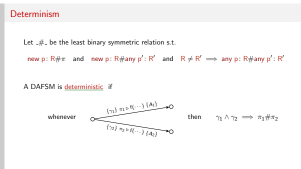
$$\text{new } p : R \# \pi \quad \text{and} \quad \text{new } p : R \# \text{any } p' : R' \quad \text{and} \quad R \neq R' \implies \text{any } p : R \# \text{any } p' : R'$$

A DAFSM is deterministic if

whenever

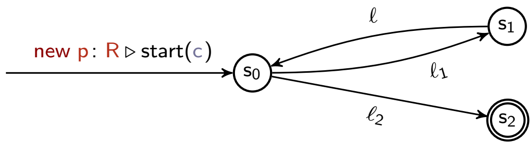


then $\quad \gamma_1 \wedge \gamma_2 \implies \pi_1 \# \pi_2$



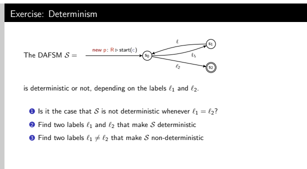
transitions from the same source state and calling the same function

The DAFSM $\mathcal{S}$ =



is deterministic or not, depending on the labels ℓ_1 and ℓ_2.

1. Is it the case that $\mathcal{S}$ is not deterministic whenever $\ell_1 = \ell_2$?
2. Find two labels ℓ_1 and ℓ_2 that make $\mathcal{S}$ deterministic
3. Find two labels $\ell_1 \neq \ell_2$ that make $\mathcal{S}$ non-deterministic



1. no: eg for $\ell_1 = \ell_2 =$ new p: R $\mathcal{S}$ is deterministic
2. $\ell_1 = \ell_2 =$ new p: R ▷ f($\cdots$, T_i x_i, $\cdots$) make $\mathcal{S}$ deterministic because the next state is unambiguously determined by the caller which is fresh on both transitions
3. $\ell_1 = \{x \leq 0\}$ p ▷ f(Int x) and $\ell_2 = \{x \geq -1\}$ p ▷ f(Int x) make $\mathcal{S}$ non-deterministic because the guards of ℓ_1 and of ℓ_2 are not disjoint therefore the next state is not determined by the caller

Well-formedness

A DAFSM is well-formed when it is

empty-role free

consistent, and

deterministic

Well-formedness

A DAFSM is well-formed when it is

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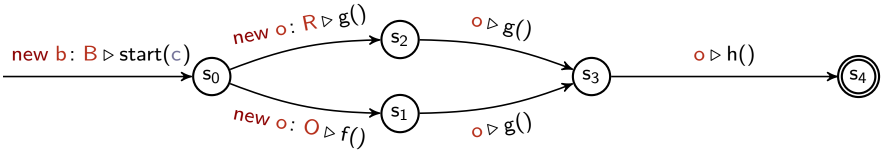
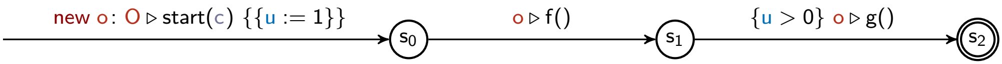
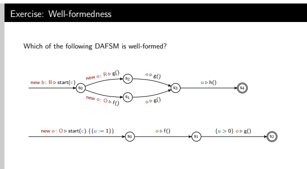
Exercise: Well-formedness

Which of the following DAFSM is well-formed?







yes: o is defined on paths it occurs on and the DAFSM is deterministic.

no: the transition from s_0 violates consistency since True does not imply $u > 0$ hinting that the protocol could get stuck in state s_1. However, this never happens because u is initially set to 1 and never changed, hence the transition from s_1 would be enabled when the protocol lands in s_1.

– Act II –

[A tool]

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[A tool]

Verification

Checking well-formedness by hand is laborious and cumbersome (and boring)

So we implemented **TRAC**, which

✓ features a DSL to specify DAFSMs

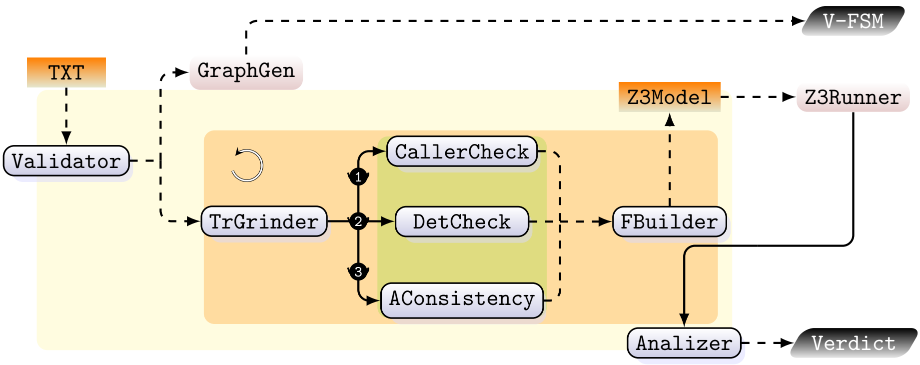
✓ verifies well-formedness condition relying on the SMT solver Z3

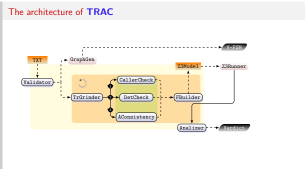
✓ it's efficient enough

✗ but cannot handle roles and inter-contract interactions

Verification
Checking well-formedness by hand is laborious and cumbersome (and boring)
So we implemented TRAC, which
✓ features a DSL to specify DAFSMs
✓ verifies well-formedness condition relying on the SMT solver Z3
✓ it's efficient enough
✗ but cannot handle roles and inter-contract interactions

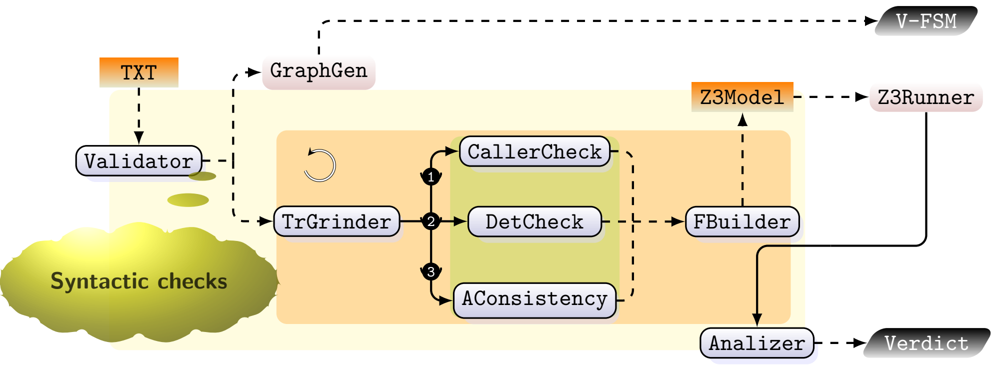
The architecture of **TRAC**

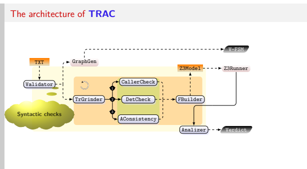




the architecture of **TRAC** is compartmentalised into two principal modules:
parsing and visualisation (yellow box) and
TRAC's core (orange box). The latter module implements well-formedness check (green box).
Solid arrows represent calls between components while dashed arrows data IO.

The architecture of **TRAC**

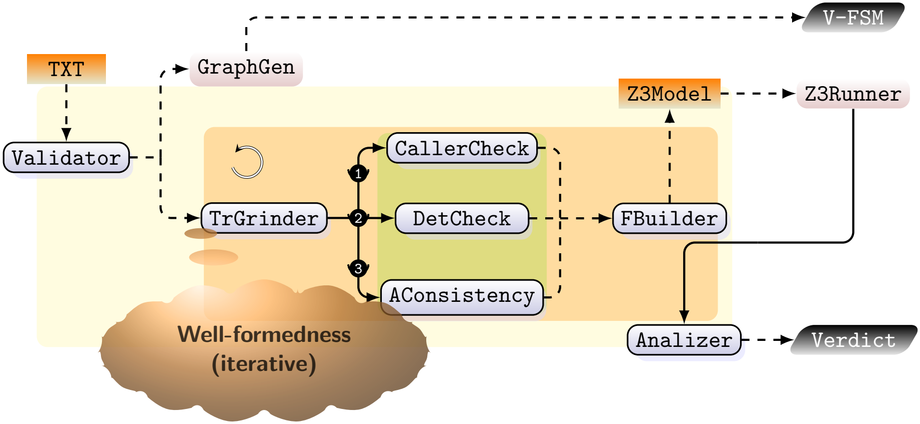




basic syntactic checks on a DSL representation of DAFSMs and transforming the input in a format that simplifies the analysis of the following phases:

- passed to `GraphGen` for visual representation of DAFSMs (`V-FSM` output)
- passed to the `TrGrinder` component (orange box) for well-formedness checking.

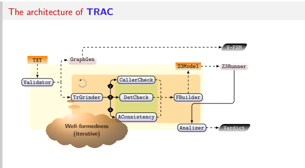
The architecture of **TRAC**



2025-05-26

A Choreographic View of Smart Contracts

└The architecture of **TRAC**



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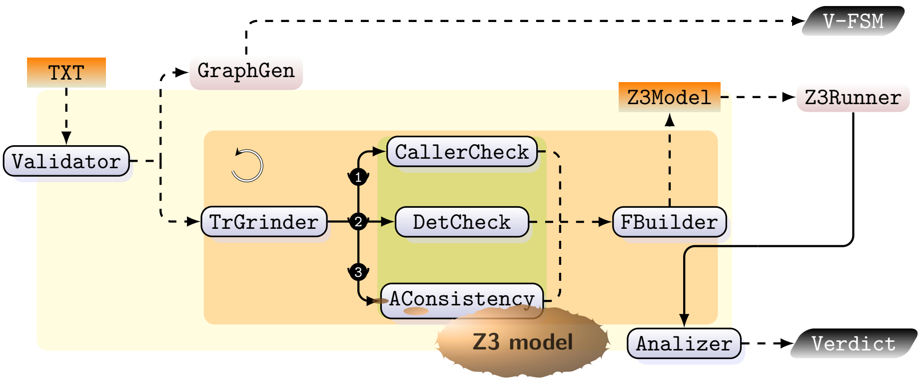
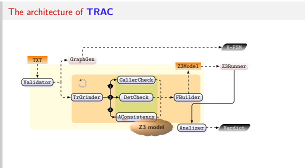
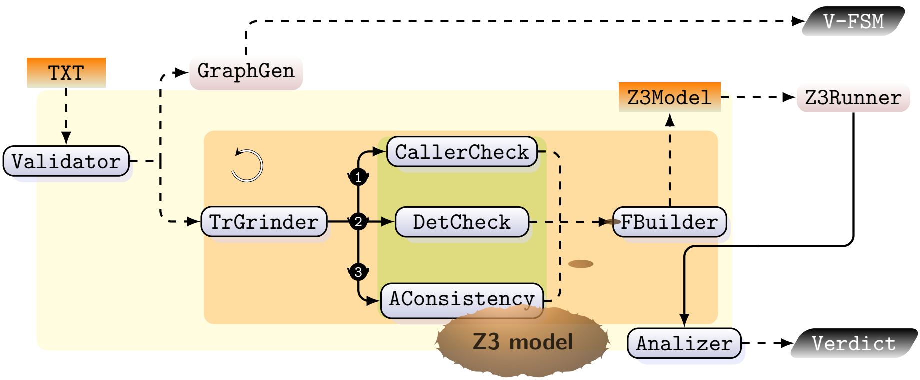




`AConsistency` (arrow ❸) to generate a Z3 formula which holds if, and only if, the transtion is consistent.

The architecture of **TRAC**



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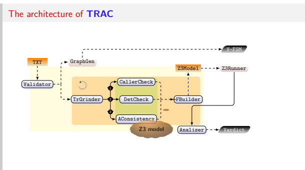
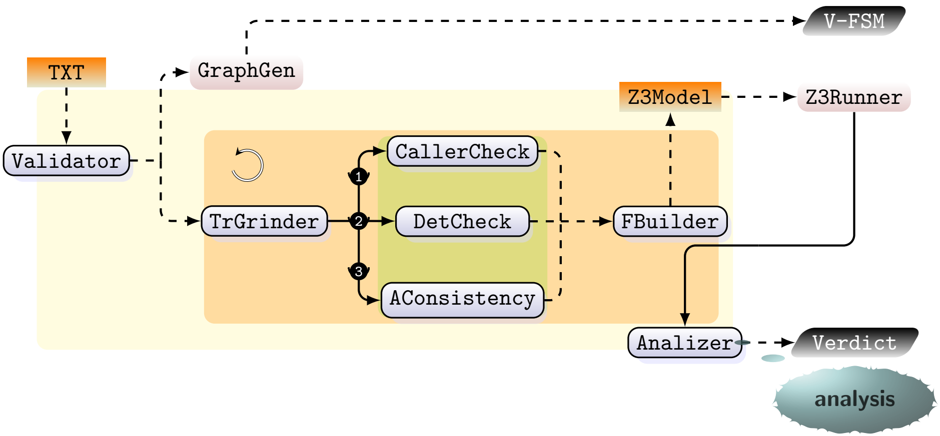
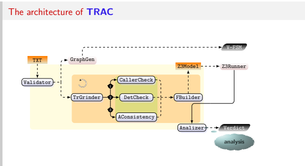


computes the z3 f.la equivalent to the conjunction of the outputs which is then passed to a Z3 engine to check its satisfiability

The architecture of **TRAC**





Finally, the `Analizer` component that diagnoses the output of Z3 and produces a `Verdict` which reports (if any) the violations of well-formedness of the DAFSM in input.

Installation

Detailed instructions at `https://github.com/loctet/TRAC`

Dependencies: Java RE (to render DAFSM graphically) & Python 3.6 or later

```
$ pip install z3-solver matplotlib networkx
```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Concrete syntax (I)

$\langle \mathit{pars} \rangle ::= \varepsilon \mid \langle \mathit{dcl} \rangle(,\langle \mathit{dcl} \rangle)^\star$ $\langle \mathit{dcl} \rangle ::= \langle \mathit{str} \rangle\ \langle \mathit{str} \rangle$

roles $\langle \mathit{str} \rangle^+$ . . . . role declaration

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  ⋮
  $\langle \mathit{dcl} \rangle = e$; # state variables with initial assignment (if any)
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A Choreographic View of Smart Contracts

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  ⋮
  if γ # initial guard (this clause can be omitted)
}

A Choreographic View of Smart Contracts

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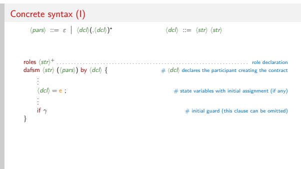

recall that e and γ are SMT-Lib2 syntax for expressions and boolean expressions respectively

Concrete syntax (I)

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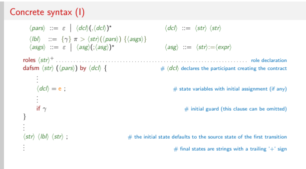
$\langle \mathit{lbl} \rangle \ ::= \ \{\gamma\} \ \pi > \langle \mathit{str} \rangle (\langle \mathit{pars} \rangle) \ \{\langle \mathit{asgs} \rangle\}$

$\langle \mathit{asgs} \rangle \ ::= \ \varepsilon \ \mid \ \langle \mathit{asg} \rangle (; \langle \mathit{asg} \rangle)^\star \qquad \langle \mathit{asg} \rangle \ ::= \ \langle \mathit{str} \rangle := \langle \mathit{expr} \rangle$

```
roles ⟨str⟩+ ........................................................ role declaration
dafsm ⟨str⟩ (⟨pars⟩) by ⟨dcl⟩ {       # ⟨dcl⟩ declares the participant creating the contract
    ⋮
    ⟨dcl⟩ = e ;                         # state variables with initial assignment (if any)
    ⋮
    if γ                                # initial guard (this clause can be omitted)
}
⋮
⟨str⟩ ⟨lbl⟩ ⟨str⟩ ;          # the initial state defaults to the source state of the first transition
⋮                                       # final states are strings with a trailing '+' sign
```

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$$\langle pars \rangle ::= \varepsilon \mid \langle dcl \rangle(,\langle dcl \rangle)^\star \qquad \langle dcl \rangle ::= \langle str \rangle\ \langle str \rangle$$

$$\langle lbl \rangle ::= \{\gamma\}\ \pi > \langle str \rangle(\langle pars \rangle)\ \{\langle asgs \rangle\}$$

$$\langle asgs \rangle ::= \varepsilon \mid \langle asg \rangle(;\langle asg \rangle)^\star \qquad \langle asg \rangle ::= \langle str \rangle{:=}\langle expr \rangle$$

roles $\langle str \rangle^+$ . . . . role declaration

dafsm $\langle str \rangle$ ($\langle pars \rangle$) by $\langle dcl \rangle$ { # $\langle dcl \rangle$ declares the participant creating the contract

    ⋮
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⋮

$\langle str \rangle$ $\langle lbl \rangle$ $\langle str \rangle$; # the initial state defaults to the source state of the first transition

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Exercise: TRAC usage (I)

Edit a `.trac` file for the contract specified at
`https://github.com/Azure-Samples/blockchain/blob/master/blockchain-workbench/application-and-smart-contract-samples/basic-provenance/readme.md`

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```
roles Owner Conterparty
dafsm basicProvenance(Owner o) by cp : Conterparty {}
q0 cp > TransferResponsibility(Conterparty cp) {} q1
q1 any cp: Conterparty > TransferResponsibility(Conterparty cp) {} q1
q1 o > Complete() {} q2+
```

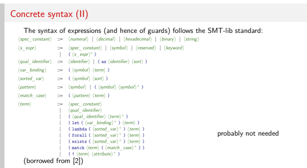
Concrete syntax (II)

The syntax of expressions (and hence of guards) follows the SMT-lib standard:

```
⟨spec_constant⟩   ::=  ⟨numeral⟩ | ⟨decimal⟩ | ⟨hexadecimal⟩ | ⟨binary⟩ | ⟨string⟩
⟨s_expr⟩          ::=  ⟨spec_constant⟩ | ⟨symbol⟩ | ⟨reserved⟩ | ⟨keyword⟩
                   |   ( ⟨s_expr⟩* )
⟨qual_identifier⟩ ::=  ⟨identifier⟩ | ( as ⟨identifier⟩ ⟨sort⟩ )
⟨var_binding⟩     ::=  ( ⟨symbol⟩ ⟨term⟩ )
⟨sorted_var⟩      ::=  ( ⟨symbol⟩ ⟨sort⟩ )
⟨pattern⟩         ::=  ⟨symbol⟩ | ( ⟨symbol⟩ ⟨symbol⟩+ )
⟨match_case⟩      ::=  ( ⟨pattern⟩ ⟨term⟩ )
⟨term⟩            ::=  ⟨spec_constant⟩
                   |   ⟨qual_identifier⟩
                   |   ( ⟨qual_identifier⟩ ⟨term⟩+ )
                   |   ( let ( ⟨var_binding⟩+ ) ⟨term⟩ )
                   |   ( lambda ( ⟨sorted_var⟩+ ) ⟨term⟩ )
                   |   ( forall ( ⟨sorted_var⟩+ ) ⟨term⟩ )
                   |   ( exists ( ⟨sorted_var⟩+ ) ⟨term⟩ )
                   |   ( match ⟨term⟩ ( ⟨match_case⟩+ ) )
                   |   ( ! ⟨term⟩ ⟨attribute⟩+ )
```

probably not needed

(borrowed from [2])

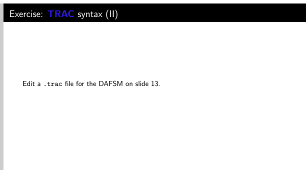


https://smt-lib.org/papers/smt-lib-reference-v2.6-r2021-05-12.pdf

http://smtlib.github.io/jSMTLIB/SMTLIBTutorial.pdf

Exercise: **TRAC** syntax (II)

Edit a `.trac` file for the DAFSM on slide 13.



TODO

– Act III –

[A little exercise]

– Act III –
[A little exercise]

A Choreographic View of Smart Contracts

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https://github.com/blockchain-unica/rosetta-smart-contracts/tree/main/contracts/vesting

– Epilogue –

[Work in progress]

– Epilogue –

[Work in progress]

Work in progress

Work in progress

Thank you

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[1] J. Afonso, E. Konjoh Selabi, M. Murgia, A. Ravara, and E. Tuosto. TRAC: A tool for data-aware coordination - (with an application to smart contracts).
In I. Castellani and F. Tiezzi, editors, *Coordination Models and Languages - 26th IFIP WG 6.1 International Conference, COORDINATION 2024, Held as Part of the 19th International Federated Conference on Distributed Computing Techniques, DisCoTec 2024, Groningen, The Netherlands, June 17-21, 2024, Proceedings*, volume 14676 of *LNCS*, pages 239–257. Springer, 2024.

[2] and and. *The SMT-LIB Standard*, version 2.7 edition.

[3] R. Garcia, E. Tanter, R. Wolff, and J. Aldrich. Foundations of typestate-oriented programming.
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[4] B. Meyer. *Introduction to the Theory of Programming Languages*.
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A Choreographic View of Smart Contracts

2025-05-26

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[5] B. Meyer. *Eiffel: The Language*.
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[6] Microsoft. The blockchain workbench.
https://github.com/Azure-Samples/blockchain/tree/master/blockchain-workbench, 2019.

[7] Microsoft. Simple marketplace sample application for azure blockchain workbench.
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[5] B. Meyer. *Eiffel: The Language*.
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https://github.com/Azure-Samples/blockchain/tree/master/blockchain-workbench/application-and-smart-contract-samples/simple-marketplace, 2019.