A Choreographic View of Smart Contracts

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This slides

What's up doc?

Prologue ..... An inspiring initiative

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Prologue . . . . . An inspiring initiative

Act I . . . . . A coordination framework

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Act I ..... A coordination framework

Act II ..... Some tool support

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Act III ..... A little exercise

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Act II ..... Some tool support

Act III ..... A little exercise

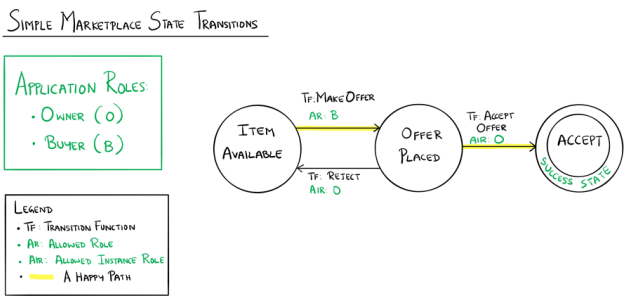
Epilogue ..... Work in progress

– Prologue –

[An inspiring initiative]

A nice sketch! [5, 6]

A smart contract among Owners and Buyers



initially buyers can make offers

then

either an owner can accept an offer and the protocol stops

or the offer is rejected and the protocol restarts

What did we just see?

A smart contract looks like

a choreographic model

global specifications determine the enabled actions along the evolution of the protocol

a typestate

In OOP, "can reflects how the legal operations on imperative objects can change at runtime as their internal state changes." [2]

A new coordination model

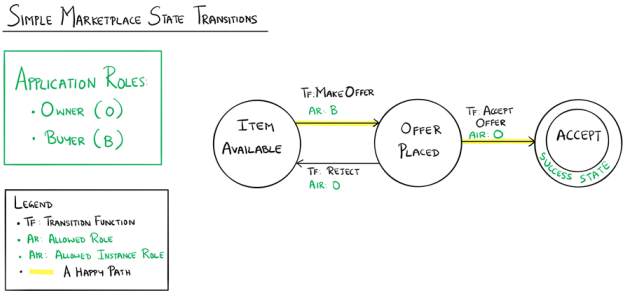
So, we saw an interesting model where

distributed components coordinate through a global specification

which specifies how actions are enabled along the computation

"without forcing" components to be cooperative!

Let's look at our sketch again



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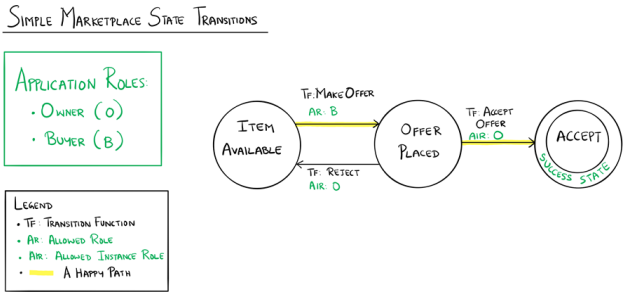

but...

- ✗ what's the difference between roles and instances?
- ✗ can buyers be owners too?
- ✗ what's the scope of quantifications?
- ✗ when are transitions enabled?
- ✗ how does the state of the contract change?

Let's go formal!

Our first attempt was to "look for into our toolbox", but

- ✗ are known notions of well-formedness suitable?
- ✗ data-awareness is crucial
- ✓ we got roles okay, but
- ✗ limitations on instances of roles
- ✗ instances can have one role only

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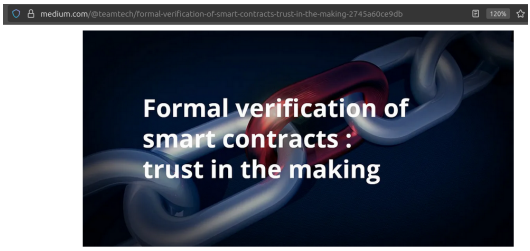
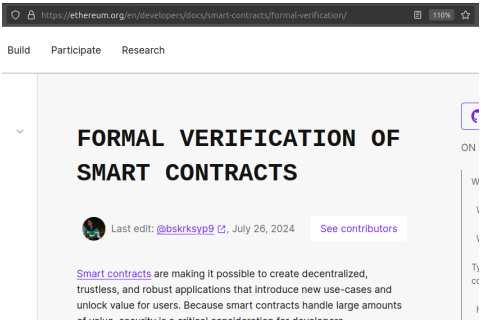
So we had to came up with some new behavioural types.

...and by the way



https://medium.com/@teamtech/formal-verification-of-smart-contracts-trust-in-the-making-2745a60ce9db



https://ethereum.org/en/developers/docs/smart-contracts/formal-verification/

– Act I –

[A coordination framework]

Basic concepts and notation

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- state variables u or old u (representing the value of u before the assignment) [3, 4]

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  - function parameters
  - state variables u or old u (representing the value of u before the assignment) [3, 4]
- $A, A', \ldots$ range over finite sets of assignments where each variable can be assigned at most once

Data-Aware FSMs

A DAFSM c on roles $R_1, \ldots R_m$ and state variables $u_1, \ldots, u_n$ is a finite-state machine "instantiated" by a participant p whose transitions are decorated as follows[1]

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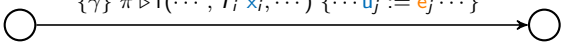

where γ is a guard (ie a boolean expression) and

$$\pi \ ::= \ \text{new R p} \ \mid \ \text{any R p} \ \mid \ \text{p}$$

is a qualified participant calling f with parameters x_i
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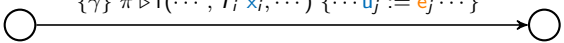

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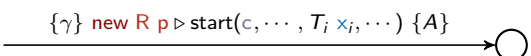

c is freshly created by p which also initialises state variables u_j with expressions e_j which are built on state variables and parameters x_i

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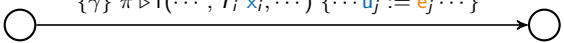

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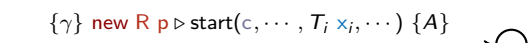

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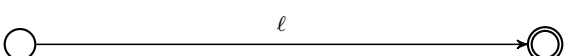

accepting states are denoted as usual

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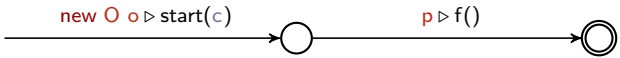
Exercise: modelling

Give a DAFSM for the protocol on slide 8 resolving the ambiguities discussed there.

Not all DAFSMs “make sense”



names’ freeness

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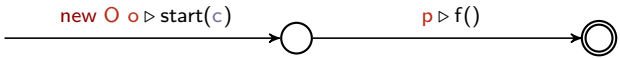

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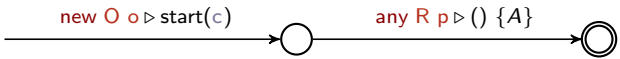

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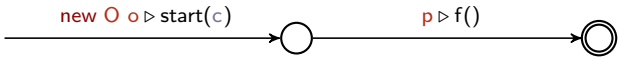

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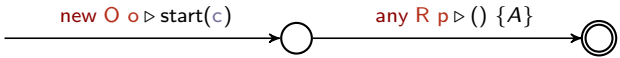

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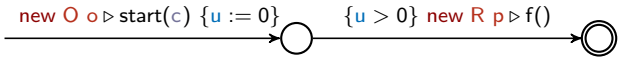

no progress

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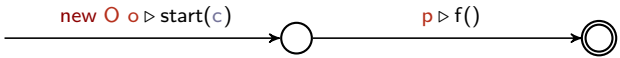

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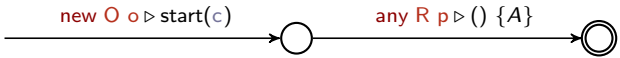

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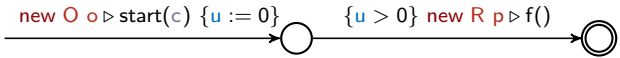

no progress

Save names' freeness, the other properties are undecidable in general, so we'll look for sufficient conditions to rule out nonsensical DAFSMs

Closed DAFSMs

Binders: parameter declarations in function calls (with scope local to the transition), new R p, and any R p (with scope along paths)

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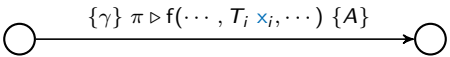
p is bound in 

if, for some role R,

$\pi = \text{new R p}$ or $\pi = \text{any R p}$ or there is i s.t. $x_i = \text{p}$ and $\text{T} = \text{R}_i$

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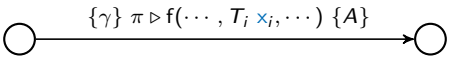
p is bound in 

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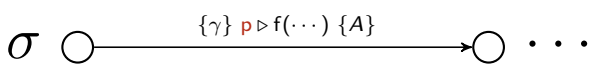
The occurrence of p is bound in a path



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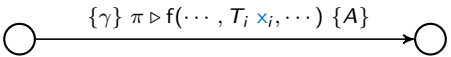
p is bound in 

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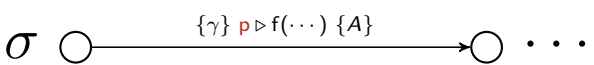
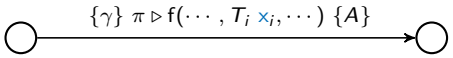


if p is bound in a transition of σ

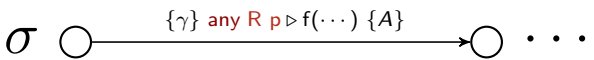
A DAFSM is closed if all occurrences of variables are bound in the paths of the DAFSM they occur on

Role emptyness

A transition 

expands role R if
π = new R p or there is i s.t. x_i = p and T_i = R

Role R is expanded in a path

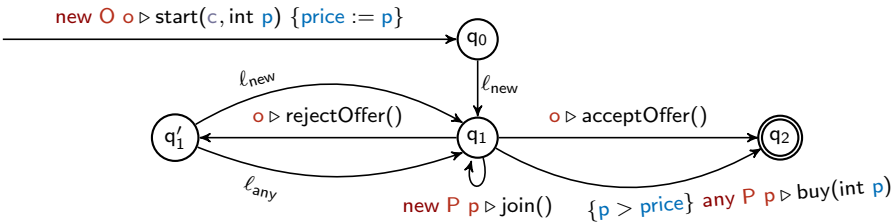


if a transition in σ expands R

A DAFSM expands R if R is expanded on all the paths it occurs on
A DAFSM is (strongly) empty-role free if it expands all its roles

Exercise: Role emptyness

Is the DAFSM below empty-role free?



where

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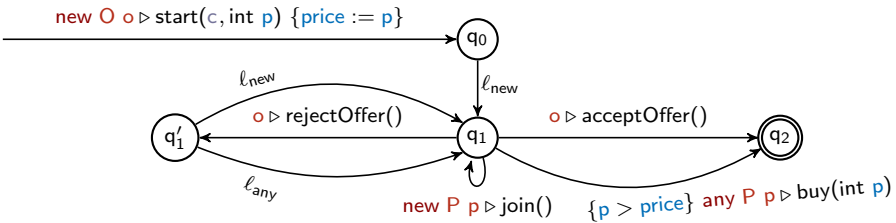

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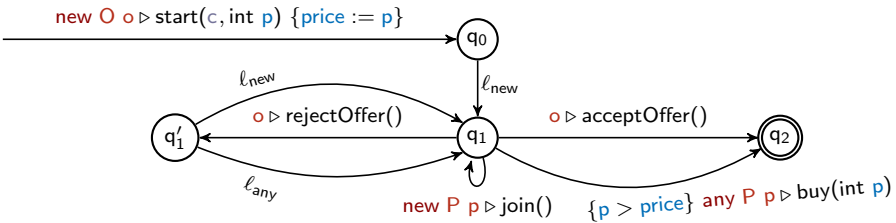

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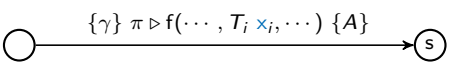
Progress

A DAFSM with state variables $u_1, \ldots, u_n$ is consistent if

for each 

$\mathbb{V}_U\, \mathbb{H}_X\, (\gamma\{\text{old } u_1, \ldots, \text{old } u_n / u_1, \ldots, u_n\} \wedge \gamma_A \implies \bigvee_{1 \le j \le m} \mathbb{H}_{Y_j}\, \gamma_j)$ is satisfiable

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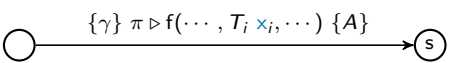
for each 

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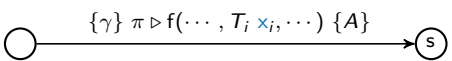
for each 

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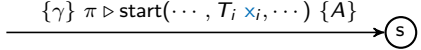
and 

is such that $\mathbb{H}_X \, \gamma$ is satisfiable

where

$U = \{u_i, \text{old } u_i\}_{1 \le i \le n}$

$X = \{x \mid \exists i : x = x_i\}$

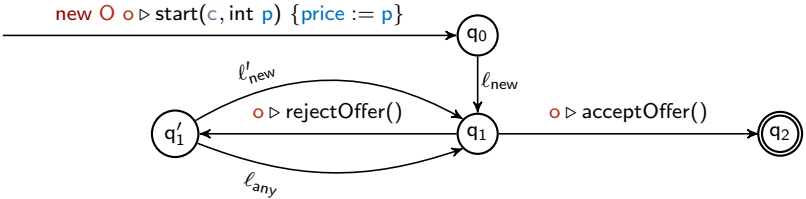
$\gamma_A = \bigwedge_{u := e \in A} u = e \wedge \bigwedge_{u \notin A} u = \text{old } u$

$Y_j = \{x \mid x \text{ is a parameter of the } j^{\text{th}} \text{ outgoing transition of s}\}$

$\gamma_j = \begin{cases} \text{the guard of the } j^{\text{th}} \text{ outgoing transitions of s} & \text{if s not accepting} \\ \text{True} & \text{otherwise} \end{cases}$

Exercise: Consistency

Is the DAFSM below consistent?



where

$\ell_{\text{new}} = \{\text{newOffer} > 0\}$ new B b ▷ makeOffer(int newOffer) {offer := newOffer},

$\ell'_{\text{new}} = \{\text{newOffer} \geq \textit{price} * 1.05\}$ new B b ▷ makeOffer(int newOffer) {offer := newOffer}, and

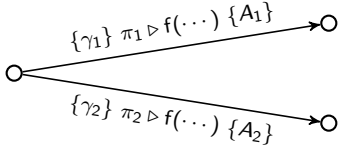
$\ell_{\text{any}} = \{\text{newOffer} \geq \textit{price} * 1.05\}$ any B b ▷ makeOffer(int newOffer) {offer := newOffer}

Determinism

Let $_\#_$ be the least binary symmetric relation s.t.

$$\text{new R p}\#\text{p}' \quad \text{and} \quad \text{new R p}\#\text{any R}' \text{ p}' \quad \text{and} \quad \text{R} \neq \text{R}' \implies \text{any R p}\#\text{any R}' \text{ p}'$$

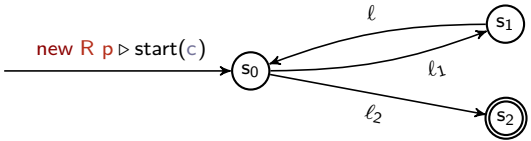
A DAFSM is deterministic if

whenever 

then $\gamma_1 \wedge \gamma_2 \implies \pi_1 \# \pi_2$

Exercise: Determinism

The DAFSM $\mathcal{S} =$



is deterministic or not, depending on the labels ℓ_1 and ℓ_2.

1. Is it the case that $\mathcal{S}$ is not deterministic whenever $\ell_1 = \ell_2$?
2. Find two labels ℓ_1 and ℓ_2 that make $\mathcal{S}$ deterministic
3. Find two labels $\ell_1 \neq \ell_2$ that make $\mathcal{S}$ non-deterministic

Well-formedness

A DAFSM is well-formed when it is

closed,

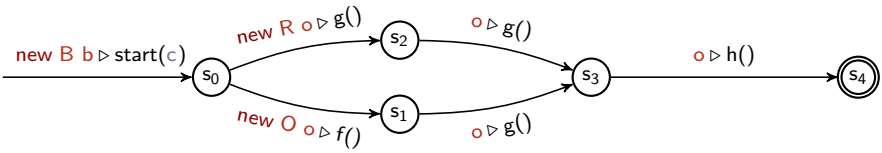
empty-role free,

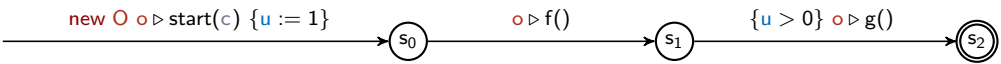
consistent, and

deterministic

Exercise: Well-formedness

Which of the following DAFSM is well-formed?





– Act II –

[A tool]

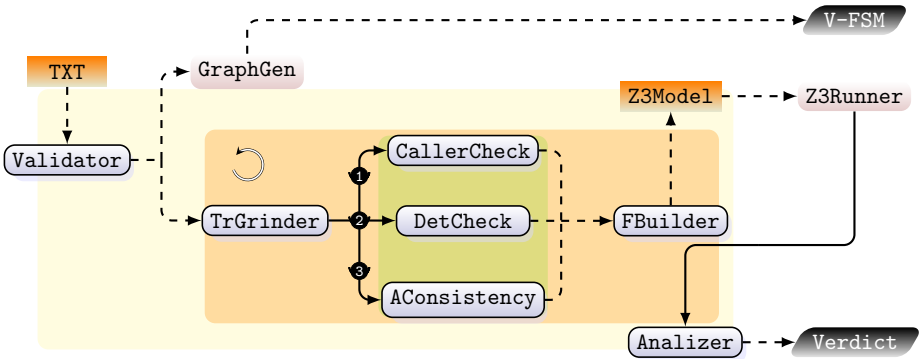
Verification

Checking well-formedness by hand is laborious and cumbersome (and boring)

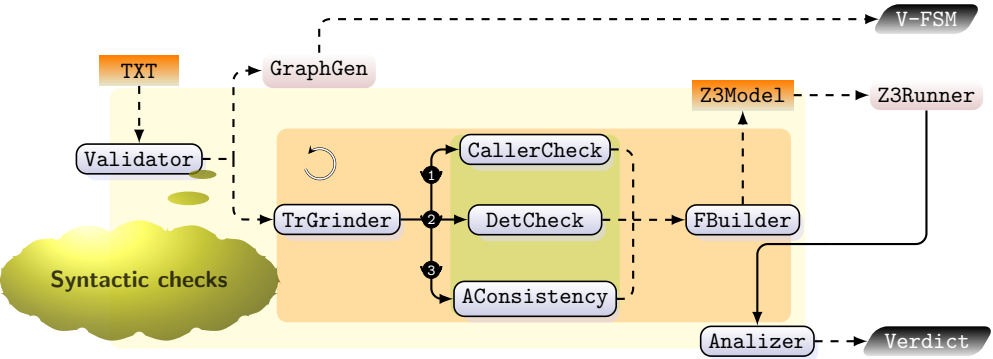
So we implemented a tool which

- ✓ features a DSL to specify DAFSMs
- ✓ verifies well-formedness (relying on the SMT solver Z3)
- ✓ it's efficient enough
- ✗ but cannot handle roles and inter-contract interactions

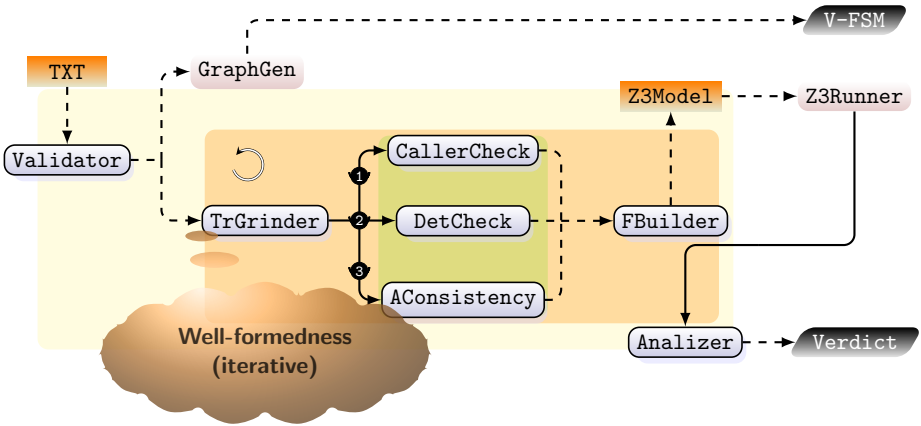
The architecture of TRAC



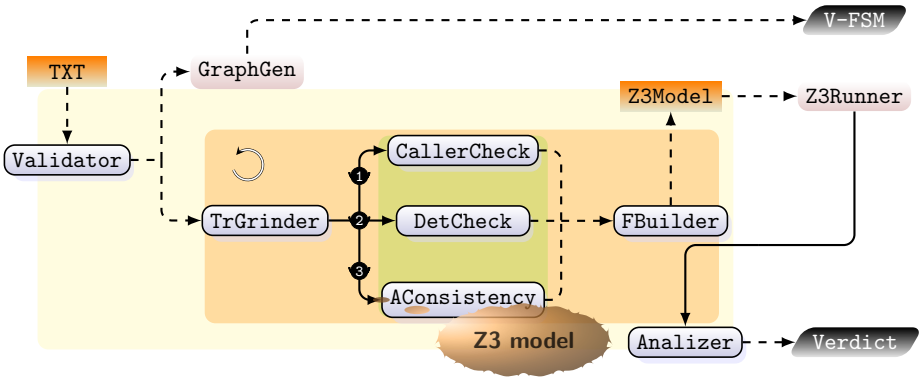
The architecture of TRAC



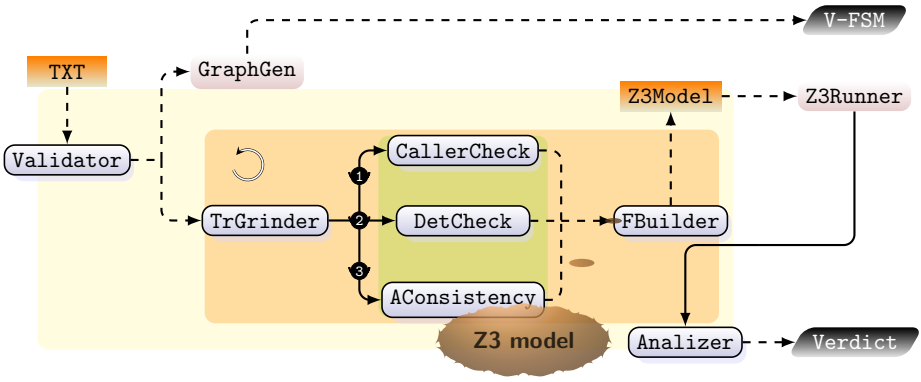
The architecture of TRAC



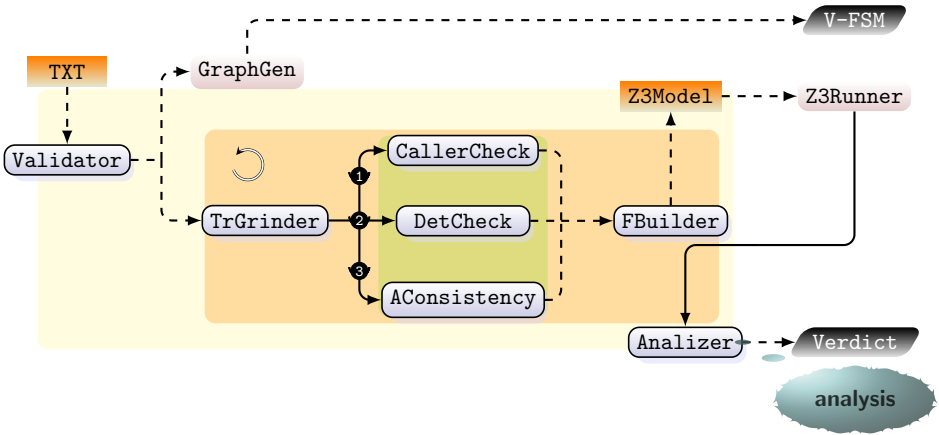
The architecture of TRAC



The architecture of TRAC



The architecture of TRAC



Getting **TRAC**

In the cloud: `https://trac-5sy1.onrender.com/`

or

In your hands: `https://github.com/loctet/TRAC/tree/TRAC_v1/`

Dependencies: GraphViz and Python 3.7 or later

Installation instructions in the `README.md`

Concrete syntax (I)

$\langle pars \rangle \;::=\; \varepsilon \;\mid\; \langle dcl \rangle (, \langle dcl \rangle)^{\star}$

$\langle dcl \rangle \;::=\; \langle str \rangle \; \langle str \rangle$

roles $\langle str \rangle^{+}$ // set the roles

dafsm $\langle str \rangle (\langle pars \rangle)$ by $\langle dcl \rangle$ { // $\langle dcl \rangle$ declares the participant creating the contract

Concrete syntax (I)

$\langle pars \rangle \ ::= \ \varepsilon \ \mid \ \langle dcl \rangle (, \langle dcl \rangle)^{\star}$ $\qquad$ $\langle dcl \rangle \ ::= \ \langle str \rangle \ \langle str \rangle$

roles $\langle str \rangle^{+}$ // set the roles
dafsm $\langle str \rangle (\langle pars \rangle)$ by $\langle dcl \rangle$ { // $\langle dcl \rangle$ declares the participant creating the contract
⋮
$\langle dcl \rangle$:=e ; // state variables (if any) with their initial assignment

Concrete syntax (I)

$\langle pars \rangle \;::=\; \varepsilon \;\mid\; \langle dcl \rangle (, \langle dcl \rangle)^{\star}$

$\langle dcl \rangle \;::=\; \langle str \rangle \; \langle str \rangle$

roles $\langle str \rangle^{+}$ // set the roles
dafsm $\langle str \rangle (\langle pars \rangle)$ by $\langle dcl \rangle$ { // $\langle dcl \rangle$ declares the participant creating the contract
⋮
$\langle dcl \rangle$:=e ; // state variables (if any) with their initial assignment
⋮
if γ // initial guard (this clause can be omitted)
}

Concrete syntax (I)

$\langle pars \rangle \ ::= \ \varepsilon \ \mid \ \langle dcl \rangle (,\langle dcl \rangle)^\star$ $\qquad$ $\langle dcl \rangle \ ::= \ \langle str \rangle \ \langle str \rangle$

$\langle lbl \rangle \ ::= \ \{\gamma\} \ \pi \ > \ \langle str \rangle (\langle pars \rangle) \ \{\langle asgs \rangle\}$

$\langle asgs \rangle \ ::= \ \varepsilon \ \mid \ \langle asg \rangle (;\langle asg \rangle)^\star$ $\qquad$ $\langle asg \rangle \ ::= \ \langle str \rangle :=$ e

roles $\langle str \rangle^+$ // set the roles

dafsm $\langle str \rangle (\langle pars \rangle)$ by $\langle dcl \rangle$ { // $\langle dcl \rangle$ declares the participant creating the contract

⋮

$\langle dcl \rangle$:=e ; // state variables (if any) with their initial assignment

⋮

if γ // initial guard (this clause can be omitted)

}

⋮

$[\langle str \rangle]$ $\langle lbl \rangle$ $[\langle str \rangle]$; // the initial state defaults to the source state of the first transition

⋮ // final states are strings with a trailing '+' sign

Exercise: TRAC usage (I)

Edit a `.trac` file for the contract specified at
`https://github.com/Azure-Samples/blockchain/blob/master/blockchain-workbench/application-and-smart-contract-samples/basic-provenance/readme.md`

Concrete syntax (II)

The syntax of non-logical expressions is as in python while logical expressions are of the form

- Not(e)
- And($e_1, \ldots, e_n$)
- Or($e_1, \ldots, e_n$)
- Imply(e_1, e_2)

See `https://ericpony.github.io/z3py-tutorial/guide-examples.htm` for examples

Exercise: TRAC syntax (II)

Edit a `.trac` file for the DAFSM on slide 14.

– Act III –

[A little exercise]

A non-trivial contract

Use **TRAC** to specify a well-formed DAFSM for a contract that helps to raise funds.

The instantiating participant plays the role of the owner of the contract.

Once instantiated with a goal (the amount of money to raise), the contract handles a fundraising campaign whereby contributors can deposit funds if the campaign is active.

Once the goal is met, the owner triggers an inspection phase done by two agents.

At the end of the inspection one of the agents closes the campaign so that the owner can finally withdraw the funds.

– Epilogue –

[Work in progress]

Work in progress

DAFSMs

Refine well-formedness

Extend the model

Projections & code generation

TRAC

Other verification approaches (model checking)

Integration with other tools

Keep improving usability

Thank you

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https://github.com/Azure-Samples/blockchain/tree/master/blockchain-workbench/application-and-smart-contract-samples/simple-marketplace, 2019.