

Push-Pull Signal-Function Functional Reactive Programming

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Functional Reactive Programming (FRP)

- *Functional*: First-class and higher-order functions.
- *Reactive*: Behavior changes in response to temporal inputs.
- Basic abstractions:
 - Signals - Functions of time¹.
 - Events - Sequences of temporally ordered and labeled discrete values.
- “Classic” FRP - Signals and events are first-class.

¹Often also called behaviors.

Signal-Function FRP (SF-FRP)

- Signal functions are transformers of events and signals.
- Signal functions are first-class in Signal-Function FRP.
- Signals and events are not first-class in Signal-Function FRP.
- This approach is more composable than first-class signals and events (since input may be transformed)².

²Courtney and Elliott, “Genuinely functional user interfaces”.

Push vs. Pull Evaluation

- When to evaluate what?
- Pull evaluation (“demand-driven”):
Evaluate when output is needed.
- Push evaluation (“data-driven”):
Evaluate when input is available.
- Ideally, an FRP system uses push for events and pull for signals³.
- This has been achieved for “classic” FRP (first class signals and events) but not SF-FRP⁴.

³Elliott, “Push-pull functional reactive programming”.

⁴Ibid.

Separating Events and Signals

```
class (Category a) => Arrow a where  
  arr :: (b -> c) -> a b c  
  ...
```

- Traditional Signal-Function FRP encodes signal functions as a Haskell Arrow.

```
data SF a b = ...
```

```
instance Arrow SF where  
  ...
```

Separating Events and Signals

- The abstraction must lift any function to a corresponding signal function, without input/output type annotation.
- Events encoded as optional signals:
`type Event a = Maybe a`
- Combined signals and events encoded as one signal with a pair type:

Separating Events and Signals

```
class (Category a) => Arrow a where
  ...
  first :: a b c -> a (b, d) (c, d)
  ...

capture :: SF (a, Event b) (Event a)

ghci> :t first . arr
first . arr
  :: (Arrow a) => (b -> c) -> a (b, d) (c, d)
```


Type Signal Functions with Signal Vectors

- Described by Sculthorpe⁵.
- Describe separation of individual signals and events.
- In Haskell (with `-XEmptyDataDecls`):

```
data SVEmpty
data SVSignal a
data SVEvent a
data SVAppend svl svr
type ^: svl svr = SVAppend svl svr
```
- Would like to use `-XDataKinds` but it's not working well yet.

⁵Sculthorpe, "Towards Safe and Efficient Functional Reactive Programming".

Type Signal Functions with Signal Vectors

- Two combinator examples:

- Lifting a pure function to transform a signal:

```
pureSignal ::      (a -> b)
              -> SF (SVSignal a) (SVSignal b)
```

- Passing through input on the right:

```
first ::      SF svIn svOut
          -> SF (svIn :^: svRight) (svOut :^: svRight)
```

- Composing these leads to:

```
ghci> :t first . pureSignal
first . pureSignal
  :: (a -> b) -> SF (SVSignal a :^: sv)
                  (SVSignal b :^: sv)
```

Partial Representations of Signal Vectors

- To evaluate signals and events differently we must represent them separately.
- Signal vectors enable this by distinguishing them in the types.
- We can construct several representations of a signal vector.
 - Represent signal leaves, event leaves, or both.
 - Represent one leaf, a subset of leaves, or all applicable leaves.
 - Transform the type at the leaf, or don't.

Signal Representation

- Represent an entire signal sample (for initializing a signal function at time zero):

data SVSample where

```
SVSample      ::      a -> SVSample (SVSignal a)
SVSampleEvent ::      SVSample (SVEvent a)
SVSampleEmpty ::      SVSample SVEempty
SVSampleBoth  ::      SVSample svLeft
                  -> SVSample svRight
                  -> SVSample (svLeft :^: svRight)
```

Signal Representation

- Represent a signal delta
(replacement values for a subset of signals):

```
data SVDelta where
    SVDelta      :: a -> SVDelta (SVSignal a)
    SVDeltaNothing :: SVDelta sv
    SVDeltaBoth   :: SVDelta svLeft
                  -> SVDelta svRight
                  -> SVDelta (svLeft :^: svRight)
```


Event Representation

- Represent an event occurrence:

```
data SVOccurrence where
  SVOccurrence :: a -> SVOccurrence (SVEvent a)
  SVOccLeft    :: SVOccurrence svLeft
                -> SVOccurrence
                (svLeft :^: svRight)
  SVOccRight   :: SVOccurrence svRight
                -> SVOccurrence
                (svLeft :^: svRight)
```

Signal Function Representation

- Separate continuations for time advancement and event occurrences:

```
data Initialized
```

```
data NonInitialized
```

```
data SF init svIn svOut where
```

```
  SF ::      (SVSample svIn
```

```
             -> (SVSample svOut,
```

```
                 SF Initialized svIn svOut))
```

```
      -> SF NonInitialized svIn svOut
```

```
...
```

Signal Function Representation

- Separate continuations for time advancement and event occurrences.

...

```
SFInit :: (Double
  -> SVDelta svIn
  -> (SVDelta svOut,
      [SVOccurrence svOut],
      SF Initialized svIn svOut))
-> (SVOccurrence svIn
  -> ([SVOccurrence svOut],
      SF Initialized svIn svOut)
-> SF Initialized svIn svOut
```

Evaluation

- Yampa/AFRP: Supply SF and input/output actions to an evaluation loop (`reactimate`).
- Here: initialize an evaluation state with:
 - A signal function.
 - Initial values for all input signals.
 - Handlers for all outputs.
- Then, the evaluation monad carries this state and provides the actions:
 - `push` Push an event (which will be immediately reacted to).
 - `update` Update the value of an input signal (with no immediate effect).
 - `step` Update the time and evaluate new values of signals.

Conclusion

- Signal vectors fix an abstraction leak in Signal-Function FRP.
- Representing events distinctly from signals in the type permits heterogeneous (push-pull) evaluation of signal functions.

Further Work

- Dynamic collections: Dynamically switch between collections of signal functions.
- Demonstrations and performance comparisons.
- Semantics/correctness proof (especially for event merging).
- Time-independence optimization.

Questions?

Time-independence optimization

- Inefficiency: running every time continuation every sample step.
- This happens even if the signal delta is empty (no change)
- Currently this is necessary because the *time* delta might cause some output.
- Optimization:
 - Mark time-independent SFs using a separate constructor.
 - "Smart" composition-routing to make a composite SF time-independent if all of its components are.
 - Time-independent time continuation has no time input, only delta.
 - Only called if non-empty delta.

Event Merging (Semantics Idea)

- Push-based events are not tied to samples.
- We thus know only what the last sample time was for an event, not the actual time.
- Can we hide event merging in this loss of precision?