

# APPM 5610 - Homework 4

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1. Show that Jacobi's method for finding eigenvalues of a real symmetric matrix is ultimately quadratically convergent. Assume that all off-diagonal elements of the matrix  $\mathbf{A}^k$  are  $\mathcal{O}(\epsilon)$ , where  $k$  enumerates Jacobi sweeps. Show that then all rotations of the next Jacobi sweep are of the form

$$1$$

Then demonstrate that this implies that, after the sweep, all off-diagonal elements of  $\mathbf{A}_{k+1}$  are  $\mathcal{O}(\epsilon^2)$ . Assume that all eigenvalues are non-zero and distinct.

Let  $\mathbf{A}$  be a given  $n \times n$  real, symmetric matrix whose eigenvalues are non-zero and distinct.

We denote Jacobi rotation matrices by  $\mathbf{J}$ , and we see that  $\mathbf{J}$  is identical to the  $n \times n$  identity matrix  $\mathbf{I}$  except for in four entries:  $\mathbf{J}_{pp} = c$ ,  $\mathbf{J}_{pq} = s$ ,  $\mathbf{J}_{qp} = -c$ , and  $\mathbf{J}_{qq} = s$ , where  $p$  and  $q$  are the row and column, respectively, of  $\mathbf{A}$  such that  $\mathbf{A}_{pq}$  is the largest off-diagonal entry in magnitude;  $c = \cos(\theta)$ ; and  $s = \sin(\theta)$ . Here,  $\theta$  is the angle by which we rotate the two-dimensional subspace spanned by columns  $p$  and  $q$  of the  $n$ -dimensional subspace spanned by the columns of  $\mathbf{A}$ .

From *Matrix Computations*, the values  $c$  and  $s$  are determined from the following equations:

$$c = \frac{1}{\sqrt{1 + t_{\min}^2}},$$
$$s = t_{\min} c,$$

where

$$\tau = \frac{a_{qq} - a_{pp}}{2a_{pq}},$$
$$t_{\min} = \begin{cases} \frac{1}{\tau + \sqrt{1 + \tau^2}} & \text{if } \tau \geq 0 \\ \frac{1}{\tau - \sqrt{1 + \tau^2}} & \text{if } \tau < 0 \end{cases}.$$

To show that Jacobi's method for finding eigenvalues is ultimately quadratically convergent, we suppose that all off-diagonal elements of  $\mathbf{A}^k$  are  $\mathcal{O}(\epsilon)$ . From the definition of  $\tau$ , we see that

$$\begin{aligned} \tau &= \frac{a_{qq} - a_{pp}}{2a_{pq}} \\ &\sim \frac{a_{qq} - a_{pp}}{2\mathcal{O}(\epsilon)} \quad \text{since the off-diagonal entries of } \mathbf{A}^k \text{ are } \mathcal{O}(\epsilon) \\ &\sim \mathcal{O}\left(\frac{1}{\epsilon}\right). \end{aligned}$$

From the definition of  $t_{\min}$ , we see that

$$\begin{aligned}
t_{\min} &= \begin{cases} \frac{1}{\tau + \sqrt{1 + \tau^2}} & \text{if } \tau \geq 0 \\ \frac{1}{\tau - \sqrt{1 + \tau^2}} & \text{if } \tau < 0 \end{cases} \\
&\sim \begin{cases} \frac{1}{\mathcal{O}(\frac{1}{\epsilon}) + \sqrt{1 + \mathcal{O}(\frac{1}{\epsilon^2})}} & \text{if } \tau \geq 0 \\ \frac{1}{\mathcal{O}(\frac{1}{\epsilon}) - \sqrt{1 + \mathcal{O}(\frac{1}{\epsilon^2})}} & \text{if } \tau < 0 \end{cases} \\
&\sim \begin{cases} \frac{1}{\mathcal{O}(\frac{1}{\epsilon}) + \sqrt{\mathcal{O}(\frac{1}{\epsilon^2})}} & \text{if } \tau \geq 0 \\ \frac{1}{\mathcal{O}(\frac{1}{\epsilon}) - \sqrt{\mathcal{O}(\frac{1}{\epsilon^2})}} & \text{if } \tau < 0 \end{cases} \\
&\sim \begin{cases} \mathcal{O}(\epsilon) & \text{if } \tau \geq 0 \\ \frac{1}{\mathcal{O}(\epsilon)} & \text{if } \tau < 0 \end{cases}.
\end{aligned}$$

Taking the smaller of the two values, we see that  $t_{\min} = \mathcal{O}(\epsilon)$ . From the definition of  $c$ , we see that

$$\begin{aligned}
c &= \frac{1}{\sqrt{1 + t_{\min}^2}} \\
&\sim 1 - \frac{t_{\min}^2}{2} + \dots \\
&\sim 1 - \mathcal{O}(\epsilon^2),
\end{aligned}$$

as desired. Then, from the definition of  $s$ , we see that

$$\begin{aligned}
s &= t_{\min} c \\
&\sim \mathcal{O}(\epsilon) (1 - \mathcal{O}(\epsilon^2)) \\
&\sim \mathcal{O}(\epsilon) - \mathcal{O}(\epsilon^3) \\
&\sim \mathcal{O}(\epsilon),
\end{aligned}$$

as desired. Thus,  $c \sim 1 - \mathcal{O}(\epsilon^2)$  and  $s \sim \mathcal{O}(\epsilon)$ .

To show that the Jacobi eigenvalue algorithm, when applied to  $\mathbf{A}^k$ , results in  $\mathbf{A}^{k+1}$  having off-diagonal entries that are  $\mathcal{O}(\epsilon^2)$ , we note that the first iteration of the algorithm

- (a) zeroes the  $a_{pq}$  and  $a_{qp}$  entries of  $\mathbf{A}^k$ ,
- (b) leaves the diagonal entries of  $\mathbf{A}^k$  unaffected, and
- (c) leaves the other off-diagonal entries of  $\mathbf{A}^k$  unaffected, so they remain  $\mathcal{O}(\epsilon)$ .

In the next iteration of the algorithm, we select columns  $r$  and  $s$  of  $\mathbf{A}^k$  so as to zero out  $a_{rs}$  and  $a_{sr}$ . After doing so, we see that

- (a) the diagonal entries of  $\mathbf{A}^k$  remain unaffected,
- (b) the  $(p, q)$  and  $(q, p)$  entries from earlier are now  $\mathcal{O}(\epsilon^2)$ ,
- (c) the  $(r, s)$  and  $(s, r)$  entries are zeroed out, and
- (d) rows and columns other than  $r$  and  $s$  are unaffected, so they remain either  $\mathcal{O}(\epsilon)$  or  $\mathcal{O}(\epsilon^2)$ .

As we continue with the iterations of the sweep, the above pattern continues, and by the end of the sweep, the diagonal entries of  $\mathbf{A}^{k+1}$  are  $\mathcal{O}(\epsilon^2)$ .

2. Show that  $\mathbf{A}$  is diagonalizable if and only if there is a positive-definite self-adjoint (Hermitian) matrix  $\mathbf{H}$  such that  $\mathbf{H}^{-1}\mathbf{A}\mathbf{H}$  is normal.

( $\implies$ ) Suppose that  $\mathbf{A}$  is diagonalizable, i.e., there exists an invertible matrix  $\mathbf{P}$  and a diagonal matrix  $\mathbf{D}$  such that  $\mathbf{A} = \mathbf{P}\mathbf{D}\mathbf{P}^{-1}$ . We can then use  $\mathbf{P}$ 's polar decomposition to write  $\mathbf{P} = \mathbf{S}\mathbf{U}$ , where  $\mathbf{S}$  is positive-definite, Hermitian matrix and  $\mathbf{U}$  is unitary. Here, the positive-definiteness of  $\mathbf{S}$  is guaranteed since  $\mathbf{P}$  is invertible. Thus,

$$\begin{aligned}\mathbf{A} = \mathbf{P}\mathbf{D}\mathbf{P}^{-1} &\implies \mathbf{A} = \mathbf{S}\mathbf{U}\mathbf{D}\mathbf{U}^*\mathbf{S}^{-1} \\ &\implies \mathbf{S}^{-1}\mathbf{A}\mathbf{S} = \mathbf{U}\mathbf{D}\mathbf{U}^*.\end{aligned}$$

Let  $\mathbf{B} = \mathbf{U}\mathbf{D}\mathbf{U}^*$ . We see that

$$\begin{aligned}\mathbf{B}^*\mathbf{B} &= \mathbf{U}\mathbf{D}^*\mathbf{U}^*\mathbf{U}\mathbf{D}\mathbf{U}^* \\ &= \mathbf{U}\mathbf{D}^*\mathbf{D}\mathbf{U}^* \\ &= \mathbf{U}\mathbf{D}\mathbf{U}^*\mathbf{U}\mathbf{D}^*\mathbf{U}^* \\ &= \mathbf{U}\mathbf{D}\mathbf{D}^*\mathbf{U}^* \\ &= \mathbf{B}\mathbf{B}^*\end{aligned}$$

since diagonal matrices commute. Thus, there exists a positive-definite self-adjoint matrix  $\mathbf{S}$  such that  $\mathbf{S}^{-1}\mathbf{A}\mathbf{S}$  is normal.

( $\impliedby$ ) Suppose that there exists a positive-definite self-adjoint matrix  $\mathbf{H}$  such that  $\mathbf{H}^{-1}\mathbf{A}\mathbf{H}$  is normal. By earlier results, we know that all normal matrices are unitarily diagonalizable, i.e.,

$$\mathbf{H}^{-1}\mathbf{A}\mathbf{H} = \mathbf{U}\mathbf{D}\mathbf{U}^*,$$

where  $\mathbf{D}$  is a diagonal matrix and  $\mathbf{U}$  is a unitary matrix. Then, we have that

$$\begin{aligned}\mathbf{H}^{-1}\mathbf{A}\mathbf{H} = \mathbf{U}\mathbf{D}\mathbf{U}^* &\implies \mathbf{A} = \mathbf{H}\mathbf{U}\mathbf{D}\mathbf{U}^*\mathbf{H}^{-1} \\ &\implies \mathbf{A} = \mathbf{P}\mathbf{D}\mathbf{P}^{-1},\end{aligned}$$

where  $\mathbf{P} = \mathbf{H}\mathbf{U}$ . Thus, we see that  $\mathbf{A}$  is similar to a diagonal matrix, i.e., we see that  $\mathbf{A}$  is diagonalizable.