

## ② Fast Fourier Transform.

- DFT takes  $N^2$  operations to calculate values

- FFT takes  $N \log(N)$  steps.

$x[n]$  -  $n$ th sample of signal  $x$ .

$X_k$  - DFT coefficient ( $k = 0 : N-1$ )

$$X_k = \sum_{n=0}^{N-1} x[n] \exp\left(\frac{-2\pi j kn}{N}\right)$$

Let  $k = \cancel{N} \cdot \frac{N}{2} k_1 + k_2$   
 $k_1 = 0/1$   
 $k_2 \Rightarrow 0 \leq k_2 \leq \frac{N}{2} - 1$

$$X_k = \sum_{n=0}^{N-1} x[n] \exp\left(\frac{-2\pi j n}{N} \left(\frac{N}{2} k_1 + k_2\right)\right)$$

$$= \sum_{n=0}^{N-1} x[n] \exp\left(\frac{-2\pi j n}{N} \left(\frac{N}{2} k_1\right)\right) \cdot \exp\left(\frac{-2\pi j n}{N} k_2\right)$$

$$= \sum_{n=0}^{N-1} x[n] \exp(-\pi j n k_1) \exp\left(\frac{-2\pi j n k_2}{N}\right)$$

$$= \sum_{n=0}^{N-1} x[n] (-1)^{nk_1} \exp\left(\frac{-2\pi j k_2 n}{N}\right)$$

Now, split 'n' into even and odd indices  
(2n) and (2n+1)

$$X_k = \sum_{n=0}^{N/2-1} x[2n] (-1)^{k_1 2n} \exp\left(\frac{-2\pi j k_2 (2n)}{N}\right)$$

$$+ \sum_{n=0}^{N/2-1} x[2n+1] (-1)^{k_1 (2n+1)} \exp\left(\frac{-2\pi j k_2 (2n+1)}{N}\right)$$

$$= \sum_{n=0}^{N/2-1} x[2n] (1) \exp\left(\frac{-2\pi j k_2 (2n)}{N}\right)$$

$$+ \sum_{n=0}^{N/2-1} x[2n+1] \exp\left(\frac{-2\pi j k_2 n}{N}\right) \exp\left(\frac{-2\pi j k_2}{N}\right) \times (-1)^{k_1}$$

$$= \sum_{n=0}^{N/2-1} x[2n] \exp\left[\frac{-2\pi j k_2 (2n)}{N}\right]$$

$$+ (-1)^{k_1} \exp\left[\frac{-2\pi j k_2}{N}\right] \sum_{n=0}^{N/2-1} x[2n+1] \exp\left[\frac{-2\pi j k_2 n}{N}\right]$$

$$* \sum_{n=0}^{N/2-1} x[2n+1] \exp\left(\frac{-2\pi j k_2 n}{N}\right)$$

$$= F_0 + QF_1$$



When  $k_1 = 0$

①  $X_{k_2} = F_0 + QF_1 \rightarrow$  FFT of  $N/2$  elements  
FFT of  $N/2$  elements

②  $k_1 = 1$   
 $X_{N/2 + k_2} = F_0 - QF_1$

$F_0 =$  DFT of  $N/2$  even terms.

$F_1 =$  DFT of  $N/2$  odd terms

To calculate DFT of an  $N$  sample signal, we break it down (recursively) into parts of  $N/2$  each time.

$$\therefore T_N = 2T_{N/2} + 2N \rightarrow \frac{N}{2} \text{ multiplications}$$

$\frac{N}{2}$  additions for

$$k_1 = 0/1$$

$$\therefore T_N = O(N \log N)$$

$\rightarrow$  Master's theorem