CMPT 210: Probability and Computing

Lecture 14

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Recap

Probability density function (PDF): Let R be a r.v. with codomain V. The probability density function of R is the function $PDF_R: V \to [0,1]$, such that $PDF_R[x] = Pr[R = x]$ if $x \in Range(R)$ and equal to zero if $x \notin Range(R)$.

Cumulative distribution function (CDF): The cumulative distribution function of R is the function $CDF_R : \mathbb{R} \to [0,1]$, such that $CDF_R[x] = Pr[R \le x]$.

Importantly, neither PDF_R nor CDF_R involves the sample space of an experiment.

Example: If we flip three coins, and C counts the number of heads, then $PDF_C[0] = Pr[C=0] = \frac{1}{8}$, and $CDF_C[2.3] = Pr[C \le 2.3] = Pr[C=0] + Pr[C=1] + Pr[C=2] = \frac{7}{8}$.

Recap

A **distribution** can be specified by its probability density function (PDF) (denoted by f).

Bernoulli Distribution: $f_p(0) = 1 - p$, $f_p(1) = p$. Example: When tossing a coin such that Pr[heads] = p, random variable R is equal to 1 if we get a heads (and equal to 0 otherwise). In this case, R follows the Bernoulli distribution i.e. $R \sim Ber(p)$.

Uniform Distribution: If $R: S \to V$, then for all $v \in V$, f(v) = 1/|V|. Example: When throwing an n-sided die, random variable R is the number that comes up on the die. $V = \{1, 2, \ldots, n\}$. In this case, R follows the Uniform distribution i.e. $R \sim \text{Uniform}(1, n)$.

Binomial Distribution: $f_{n,p}(k) = \binom{n}{k} p^k (1-p)^{n-k}$. Example: When tossing n independent coins such that $\Pr[\text{heads}] = p$, random variable R is the number of heads in n coin tosses. In this case, R follows the Binomial distribution i.e. $R \sim \text{Bin}(n,p)$.

Pr(heads) = p for each coin

Geometric Distribution: $f_p(k) = (1-p)^{k-1}p$. Example: When repeatedly tossing a coin such that $\Pr[\text{heads}] = p$, random variable R is the number of tosses needed to get the first heads. In this case, R follows the Geometric distribution i.e. $R \sim \text{Geo}(p)$.

Q: It is known that disks produced by a certain company will be defective with probability 0.01 independently of each other. The company sells the disks in packages of 10 and offers a money-back guarantee that at most 1 of the 10 disks is defective (the package can be returned if there is more than 1 defective disk). What proportion of packages is returned? If someone buys three packages, what is the probability that exactly one of them will be returned? It X be the number of disks which are defective in the packet

$$Pr(X > 1) = 1 - pr(X <=1)$$

Think of a disk being defective as a coin coming up I

 $X \sim Bin(10, 0.01)$

Event of interest: package is returned

$$Pr(X <= 1) = Pr(X = 0) + Pr(X = 1) \\ Pr(X <= 1) = 0.99^10 + 0.99^9 * 0.01 * 10 = 0.99^9(0.99 + 0.01) = 0.99 \\ pr(X > 1) = 1 - line above = approximately 0.05 \\ Pr(package returned) = 0.05 \\ Let P be the amount of packages returned \\ P \sim Bin(3, 0.05)$$

 $Pr(P = 1) (3c1)0.05 * 0.95^2$

Find pr(X >= 1)

Q: It is known that disks produced by a certain company will be defective with probability 0.01 independently of each other. The company sells the disks in packages of 10 and offers a money-back guarantee that at most 1 of the 10 disks is defective (the package can be returned if there is more than 1 defective disk). What proportion of packages is returned? If someone buys three packages, what is the probability that exactly one of them will be returned?

Let X be the random variable corresponding to the number of defective disks in a package. Let E be the event that the package is returned. We wish to compute $\Pr[E] = \Pr[X > 1]$. X follows the Binomial distribution Bin(10, 0.01). Hence,

$$Pr[E] = Pr[X > 1] = 1 - Pr[X \le 1] = 1 - Pr[X = 0] - Pr[X = 1]$$
$$= 1 - {10 \choose 0} (0.99)^{10} - {10 \choose 1} (0.99)^{9} (0.01)^{1} \approx 0.05$$

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 ${f Q}$: It is known that disks produced by a certain company will be defective with probability 0.01 independently of each other. The company sells the disks in packages of 10 and offers a money-back guarantee that at most 1 of the 10 disks is defective (the package can be returned if there is more than 1 defective disk). If someone buys three packages, what is the probability that exactly one of them will be returned?

Let F be the event that someone bought 3 packages and exactly one of them is returned.

Answer 1: Let E_i be the event that package i is returned. From the previous question, we know that $\Pr[E_i] = \Pr[\text{Package } i \text{ has more than 1 defective disk}] \approx 0.05$.

$$F = (E_1 \cap E_2^c \cap E_3^c) \cup (E_1^c \cap E_2^c \cap E_3) \cup (E_1^c \cap E_2 \cap E_3^c)$$

$$\Pr[F] = \Pr[E_1](1 - \Pr[E_2])(1 - \Pr[E_3]) + (1 - \Pr[E_1])(1 - \Pr[E_2])\Pr[E_3] + \dots$$

$$\Pr[F] \approx 3 \times (0.05)(0.95)(0.95) \approx 0.15.$$

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$$Pr[F] \approx 3 \times (0.05)(0.95)(0.95) \approx 0.15.$$

Answer 2: Let Y be the random variable corresponding to the number of packages returned. Y follows the Binomial distribution Bin(3,0.05) and we wish to compute $Pr[F] = Pr[Y = 1] \approx \binom{3}{1}(0.05)^1(0.95)^2 \approx 0.15$.

Q: You are randomly and independently throwing darts. The probability that you hit the bullseye in throw i is p. Once you hit the bullseye you win and can go collect your reward. (a) What is the probability that you win after exactly k throws? (b) What is the probability you win in less than k throws?

Geometric distribution

Let D be the random variable measuring the number of darts thrown.

$$G \sim Geo(k, p)$$

$$Pr(G = k) = (1 - p)^{k} (k - 1)^{*} p$$

Alternate: sum over from $(1 - p)^0p$ to $(1-p)^k - 2p$

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(a) The number of throws (T) to hit the bullseye and win follows a geometric distribution Geo(p) and we wish to compute Pr[T=k]. Using the PDF for the Geometric distribution, this is equal to $(1-p)^{k-1}p$.

Review these slides

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- (a) The number of throws (T) to hit the bullseye and win follows a geometric distribution Geo(p) and we wish to compute Pr[T=k]. Using the PDF for the Geometric distribution, this is equal to $(1-p)^{k-1}p$.
- (b) **Answer 1**: If *E* is the event that we win in less than *k* throws, $\Pr[E] = \Pr[T < k] = \sum_{i=1}^{k-1} \Pr[T = i] = p \sum_{i=1}^{k-1} (1-p)^{i-1} = 1 (1-p)^{k-1}$.

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- (b) **Answer 1**: If E is the event that we win in less than k throws,

$$\Pr[E] = \Pr[T < k] = \sum_{i=1}^{k-1} \Pr[T = i] = p \sum_{i=1}^{k-1} (1-p)^{i-1} = 1 - (1-p)^{k-1}.$$

Answer 2:

$$Pr[E] = 1 - Pr[E^c] = 1 - Pr[do \text{ not hit the bullseye in } k - 1 \text{ throws}] = 1 - (1 - p)^{k-1}$$
.

Q: We have two envelopes. Each contains a distinct number in $\{0, 1, 2, \dots, 100\}$. To win the game, we must determine which envelope contains the larger number. We are allowed to peek at the number in one envelope selected at random. Can we devise a winning strategy?

Being adverserial:

choose two numbers in the envelopes based on the strategy

If you use randomness, you can overcome the adversarial challenge.

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Strategy 1: We pick an envelope at random and guess that it contains the larger number (without even peeking at the number).

Q: What is the probability that we win with this strategy?

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Strategy 1: We pick an envelope at random and guess that it contains the larger number (without even peeking at the number).

Q: What is the probability that we win with this strategy?

Strategy 2: We peek at the number and if its below 50, we choose the other envelope.

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Strategy 1: We pick an envelope at random and guess that it contains the larger number (without even peeking at the number).

Q: What is the probability that we win with this strategy?

Strategy 2: We peek at the number and if its below 50, we choose the other envelope.

But the numbers in the envelopes need not be random! The numbers are chosen "adversarially" in a way that will defeat our guessing strategy. For example, to "beat" Strategy 2, the two numbers can always be chosen to be below 50Text

Q: Can we do better than 50% chance of winning?

If we want to win with better odds, we need more randomness.

X is in (L, H)

Suppose that we somehow knew a number x that was in between the numbers in the envelopes. If we peek in one envelope and see a number. If it is bigger than x, we know its the higher number and choose that envelope. If it is smaller than x, we know that is the smaller number and choose the other envelope.

If T > x, then we pick T

If T < x, then we pick the other envelope

Suppose that we somehow knew a number x that was in between the numbers in the envelopes. If we peek in one envelope and see a number. If it is bigger than x, we know its the higher number and choose that envelope. If it is smaller than x, we know that is the smaller number and choose the other envelope.

Of course, we do not know such a number x. But we can guess it!

Strategy 3: Choose a random number x from $\{0.5, 1.5, 2.5, \dots n - 1/2\}$ according to the uniform distribution i.e. $\Pr[x = 0.5] = \Pr[1.5] = \dots = 1/n$. Then we peek at the number (denoted by T) in one envelope, and if T > x, we choose that envelope, else we choose the other envelope.

We choose a number in the fractional multiples since we cannot have that X = L or X = H Uniform distribution for n.

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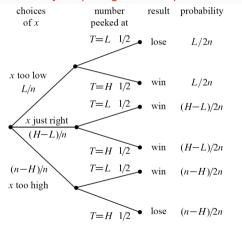
The advantage of such a randomized strategy is that the adversary cannot easily "adapt" to it.

Q: But does it have better than 50% chance of winning?
Why is it that increasing the probability slightly is a tremendous benefit?

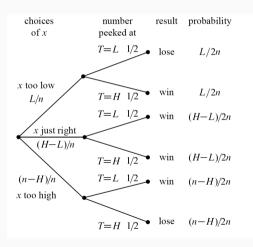
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Let the numbers in the two envelopes be L (lower number) and H (the higher number).

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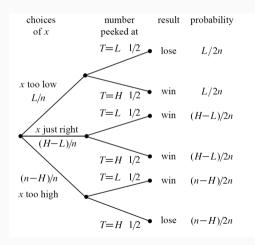
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$$\Pr[\text{win}] = \frac{L}{2n} + \frac{H - L}{2n} + \frac{H - L}{2n} + \frac{n - H}{2n}$$
$$= \frac{1}{2} + \frac{H - L}{2n} \ge \frac{1}{2} + \frac{1}{2n} > \frac{1}{2}$$

Hence our strategy has a greater than 50% chance of winning! If n = 10, $Pr[win] \ge 0.55$, for n = 100, $Pr[win] \ge 0.505$.

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$$= \frac{1}{2} + \frac{H - L}{2n} \ge \frac{1}{2} + \frac{1}{2n} > \frac{1}{2}$$

Hence our strategy has a greater than 50% chance of winning! If n = 10, $Pr[win] \ge 0.55$, for n = 100, $Pr[win] \ge 0.505$.

Q: For n = 100, if L = 23 and H = 54, compute Pr[guessing too low | we win]

