Discussion 12: November 15, 2023

Getting Started

Someone from your group should join Discord.

To get help from a TA, send a message to the discuss-queue channel with the @discuss tag and your discussion group number.

If you have only 1 or 2 people in your group, you can join the other group in the room with you.

Q1: Ice Breaker

Everybody say your name, and then figure out who has the most elaborate Thanksgiving tradition. They're your group's facilitator for the day.

Facilitator: If you ever don't want to be facilitator anymore, just ask: "Would someone else be facilitator?" You can even do that now.

Macros

A macro is a code transformation that is created using define-macro and applied using a call expression. A macro call is evaluated by:

- 1. Binding the formal paramters of the macro to the **unevaluated** operand expressions of the macro call.
- 2. Evaluating the body of the macro, which returns an expression.
- 3. Evaluating the expression returned by the macro in the frame of the original macro call.

```
scm> (define-macro (twice expr) (list 'begin expr expr))
twice
scm> (twice (+ 2 2)) ; evaluates (begin (+ 2 2) (+ 2 2))
4
scm> (twice (print (+ 2 2))) ; evaluates (begin (print (+ 2 2)) (print (+ 2 2)))
4
4
4
```

Q2: Mystery Macro

Figure out what mystery-macro does.

```
(define-macro (mystery-macro expr old new)
        (mystery-helper expr old new))

(define (mystery-helper e o n)
      (if (pair? e)
            (cons (mystery-helper (car e) o n) (mystery-helper (cdr e) o n))
            (if (eq? e o) n e)))
```

Here are some example uses of mystery-macro that could help you understand what it does and how it might be used.

```
scm> (define a-very-long-symbol 5)
a-very-long-symbol
scm> (mystery-macro (+ (* x x) (* 2 x) 1) x a-very-long-symbol)
36
scm> (mystery-macro (+ (* x x) (* 2 x) 1) x (+ a-very-long-symbol 1))
49
scm> (mystery-macro '(+ (* x x) (* 2 x) 1) x y)
(+ (* y y) (* 2 y) 1)
scm> (mystery-macro (> (x) (> (y) (+ x y))) > lambda)
(lambda (x) (lambda (y) (+ x y)))
scm> (mystery-macro (begin e e e) e (print 5))
5
5
5
```

Pro tip: Please don't just look at the hint right away. Try to figure out what it does just by talking about the code as a group and discussing examples. Hints are for when you get stuck.

Facilitator: Ask your group, "What does mystery-macro do?" Once you all agree, work together to come up with an interesting use of this macro. Then send a message to the #discuss-queue channel with the @discuss tag, your discussion group number, and the message "We made some changes!". A member of the course staff will join your voice channel to hear you explain your use case and give feedback. The facilitator can decide who will present your explanation (or ask for volunteers).

Q3: Repeat

Define repeat, a macro that is called on a number n and an expression expr. Calling it evaluates expr in a local frame n times, and its value is the final result.

```
For example, (repeat (+ 2 3) (print 1)) is equivalent to: (repeated-call (+ 2 3) (lambda () (print 1)))
```

```
(define-macro (repeat n expr)
   `(repeated-call ,n (lambda () ,expr)))

; Call zero-argument procedure f n times and return the final result.
(define (repeated-call n f)
   (if (= n 1) (f) (begin (f) (repeated-call (- n 1) f))))

(repeat (+ 1 2) (print 'three))    ; This should print three 3 times
(repeat 2 (repeat 2 (print 'four)))  ; This should print four 4 times
```

The repeated-call procedure takes a zero-argument procedure, so (lambda () ___) must appear in the blank. The body of the lambda is expr, which must be unquoted.

Call f on no arguments with (f). If n is 1, just call f. If n is greater than 1, first call f and then call (repeated-call (- n 1) f).

Discussion time: The repeat macro from lecture (*repeated below*) included (eval n), while this version just has n in the body of the macro. Why does n not need to be evaluated in this implementation of repeat? If you talk about it for more than 5 minutes and still aren't sure, click the explanation below to find out.

In the version of repeat with lambda, n is part of the call to repeated-call returned from the macro, and so it will be evaluated automatically when that call expression is evaluated. For example, (repeat (+ 2 3) (print 1)) returns (repeated-call (+ 2 3) (lambda () (print 1))), and when this call is evaluated, the operand (+ 2 3) is evaluated to 5, which is passed to repeated-call.

In the lecture version of repeat, the repeated-expr procedure is being called within the macro, and so n will be passed as an expression unless it is evaluated using (eval n).

Version of repeat from lecture that contains (eval n):

```
; Evaluate expr n times and return the last value.
(define-macro (repeat n expr)
  (cons 'begin (repeated-expr (eval n) expr)))

; Return a list containing expr n times.
; scm> (repeated-expr 4 '(print 2))
; ((print 2) (print 2) (print 2) (print 2))
(define (repeated-expr n expr)
  (if (zero? n) nil (cons expr (repeated-expr (- n 1) expr))))
```

Facilitator: Ask your group, "What do you think of macros so far?" Together, write a sentence that starts with "Macros are ..." in your group's Discord channel's text chat.

Q4: Multiple Assignment

In Scheme, the expression returned by a macro procedure is evaluated in the same environment in which the macro was called. Therefore, it's possible to return a **define** expression from a macro and have it affect the environment in which the macro was called. This differs from a regular scheme procedure that contains a **define** expression, which would only affect the procedure's local frame.

In Python, we can bind two names to values in one line as follows:

```
>>> x, y = 1 + 1, 3  # now x is bound to 2 and y is bound to 3
>>> x, y = y, x  # swap the values of x and y
>>> x
3
>>> y
2
```

Implement the assign Scheme macro, which takes in two symbols sym1 and sym2 as well as two expressions expr1 and expr2. It should bind sym1 to the value of expr1 and sym2 to the value of expr2 in the environment from which the macro was called.

```
scm> (assign x y (+ 1 1) 3) ; now x is bound to 2 and y is bound to 3
scm> (assign x y y x) ; swap the values of x and y
scm> x
3
scm> y
2
```

Make sure that expr2 is evaluated before sym1 is changed. Assume that expr1 and expr2 do not have side effects (and so do not contain define or assign expressions).

```
(define-macro (assign sym1 sym2 expr1 expr2)
  `(begin
     (define ,sym1 ,expr1)
     (define ,sym2 ',(eval expr2))))
(assign x y (+ 1 1) 3)
(assign x y y x)
(expect x 3)
(expect y 2)
; Extra tests
(define z 'x)
                  ; z is bound to the symbol x
(assign v w 2 z) \,; now v is bound to 2 and w is bound to the symbol x
(assign v w w v)
                 ; swap the values of v and w
(expect v x)
(expect w 2)
```

Call eval on expr2 so that its value is included in the define expression created by assign: ,(eval expr2). That way, the define for expr1 won't affect the value of expr2, because expr2 will already have been evaluated.

For an **optional extra challenge**, paste in these additional tests that make sure **assign** works correctly even when the value of **expr2** is not a number, but instead a symbol.

In order to ensure that the value of expr2 is not evaluated a second time, quote the result of evaluating it.

For example, (assign v w 2 z) should be equivalent to:

```
(begin
 (define v 2)
 (define w (quote x)))
```

In this begin expression, (quote x) comes from first evaluating z and then quoting the result.

Document the Occasion

Please all fill out the attendance form (one submission per person per week).