

1) a)  $2n^3 - 2n + 1 \in O(n^3)$  doğrudur

$$2n^3 - 2n + 1 \leq cn^3$$

$$2n^3 - 2n + 1 \leq 2n^3 + 2n^3 + n^3 \leq 5n^3 \quad \begin{matrix} c=5 \\ n_0=1 \end{matrix} \quad \forall n \geq n_0$$

b)  $2n^2 - n + 1 \in O(n^3)$  Yanlış ✗

$$c_1 g(n) \leq f(n) \leq c_2 g(n) \quad \forall n \geq n_0$$

$$c_1 \cdot n^3 \stackrel{?}{\leq} 2n^2 - n + 1 \stackrel{\checkmark}{\leq} c_2 \cdot n^3 \quad c_2 > 0$$

$$c_1 n^3 \leq 2n^2 - n + 1 \quad \alpha \quad \text{Yanlış}$$

↳ Hiçbir sabit sonucu sıfırlanmaz.

c)  $2n^3 - n + 1 \in O(n^4)$  doğrudur

$$f(n) \leq g(n)$$

$$2n^3 - n + 1 \leq cn^4$$

$$2n^3 - n + 1 \leq cn^3 \leq cn^4 \quad \checkmark$$

$$2n^3 + n^3 + n^3$$

d)  $2n^3 - 3n + 1 \in \Omega(n^2)$  doğrudur ✓

$$f(n) \geq c g(n)$$

$$2n^3 - 3n + 1 \geq cn^2$$

↳ Sonuçta daha hızlı gider

2) a)  $x(n) = x(n+1) + 7, n \geq 1, x(1) = 0$

$$x(1) = x(2) + 7$$

$$x(2) = x(3) + 7$$

$$\vdots$$

$$x(n-2) = x(n-1) + 7$$

$$+ \quad x(n-1) = x(n) + 7$$

$$x(1) = x(n) + (n-1) \cdot 7 \Rightarrow$$

$$0 = x(n) + 7n - 7$$

$$\boxed{x(n) = 7 - 7n}$$

b)  $x(n) = 3x(n-1) + 4$ ,  $n > 1$ ,  $x(1) = 1$

$$x(n-1) = 3x(n-2) + 4$$

$$x(n) = 3[3x(n-2) + 4] + 4$$

$$x(n-2) = 3x(n-3) + 4$$

$$x(n) = 3^2[3x(n-3) + 4] + 3 \cdot 4 + 3 \cdot 4 + 4$$

$$x(n) = 3^3 x(n-3) + 3^2 \cdot 4 + 3 \cdot 4 + 4$$

$$x(n) = 3^k x(n-k) + 3^{k-1} \cdot 4 + 3^{k-2} \cdot 4 + \dots + 3 \cdot 4 + 4$$

$k = n-1$  Sei! med!

$$x(n) = 3^{n-1} x(1) + 3^{n-2} \cdot 4 + 3^{n-3} \cdot 4 + \dots + 3 \cdot 4 + 4$$

$$= 3^{n-1} + 4 \cdot (3^{n-2} + \dots + 1)$$

$r = 3$

$$= 3^{n-1} + 4 \cdot \left[ \frac{3^{n-1} - 1}{3 - 1} \right] = 3^{n-1} + 2 \cdot 3^{n-1} - 2$$

$$x(n) = 3^n - 2$$

c)  $x(n) = x(n/3) + 2$ ,  $n > 1$ ,  $x(1) = 1$  ( $n = 3^k$  ist angenommen)

$$n = 3^k \quad x(3^k) = x(3^{k-1}) + 2$$

$$x(3^{k-1}) = x(3^{k-2}) + 2$$

⋮

$$n = 3 \quad x(3) = x(1) + 2$$

$$x(3^k) = x(1) + 2 \cdot k = 2k + 1$$

$$x(n) = 2 \cdot \log_3 n + 1$$

$$k = \log_3 n$$

3) Özyinelemeli Ağac Yöntemi ile çözünüz

$$T(n) = 3T(n/2) + n^2$$

$$T(1) = 1$$

$$T(n) = n^2 + 3\left(\frac{n}{2}\right)^2 + 3^2\left(\frac{n}{2^2}\right)^2 + \dots + 3^h\left(\frac{n}{2^h}\right)^2$$

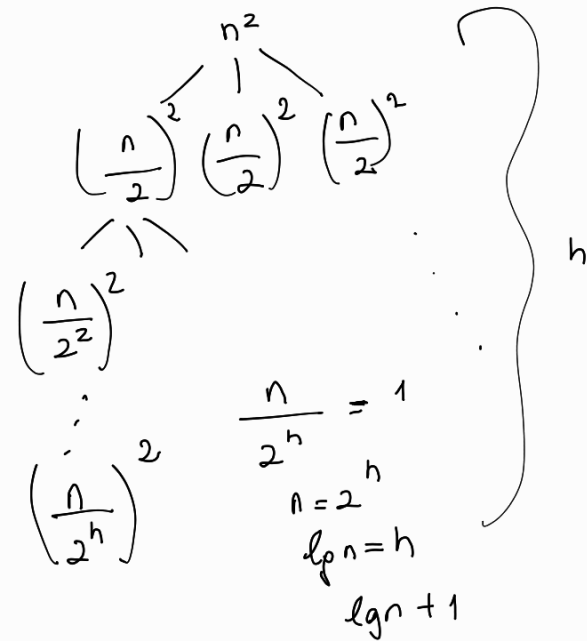
$$T(n) = n^2 \left[ 1 + 3 \cdot \frac{1}{4} + 3^2 \cdot \frac{1}{4^2} + \dots + 3^h \cdot \frac{1}{2^{2h}} \right]$$

$$T(n) = n^2 \cdot \left[ \frac{1 - \left(\frac{3}{4}\right)^{h+1}}{\frac{1}{4}} \right] = 4 \cdot n^2 \cdot \left[ 1 - \left(\frac{3}{4}\right)^{h+1} \right]$$

$$= 4n^2 - 4n^2 \cdot \left(\frac{3}{4}\right)^{h+1}$$

$$= 4n^2 - \frac{4n^{h/2}}{4^{h+1}} \cdot 3^{h+1} = \boxed{4n^2 - 3^{\lg n + 1}}$$

~~$\frac{\lg n}{4} \cdot 4$~~



4) Master Teoremiyle çözünüz ( $T(1)=1$ )

a)  $T(n) = 3T(n/3) + \underline{n^2 \lg^3 n}$

$$f(n) = n^2 \lg^3 n$$

$$\begin{matrix} a=3 \\ b=3 \end{matrix} \quad n^{\lg_3 3} = \underline{n^2} \text{ ile karşılaştır}$$

Durum 1

$$n^2 \lg^3 n = O(n^{2-\epsilon})$$

$$\cancel{n^2 \lg^3 n} = \cancel{n^2} \cdot \frac{1}{n^\epsilon}$$

$$\lg^3 n > \frac{1}{n^\epsilon}$$

Durum 1 olamaz

Durum 2

$$n^2 \lg^3 n \stackrel{?}{=} \Omega(n^{2+\epsilon})$$

$$\cancel{n^2 \lg^3 n} = \cancel{n^2} \cdot \left(\frac{\epsilon}{n}\right)$$

$$n^\epsilon > \lg n$$

Durum 2 olamaz

↳ Alt sınır doğrudur

Durum 3

$$\boxed{n^2 \lg^3 n = \Theta(n^2 \lg^k n)}$$

$$\boxed{k=3}$$

$$\forall k \geq 3$$

$$\boxed{T(n) = \Theta(n^2 \lg^4 n)}$$

b)  $T(n) = 10T(n/3) + n^2 \lg^5 n$   
 $f(n) = n^2 \lg^5 n \rightarrow a=10, b=3, n^{\log_3 10} > n^2$

$$n^2 \lg^5 n = O(n^{\log_3 10 - \epsilon})$$

Durum 1

$$T(n) = O(n^{\log_3 10})$$

c)  $T(n) = 7T(n/3) + n^2 \lg^5 n$

$f(n) = n^2 \lg^5 n \rightarrow a=7, b=3, n^{\log_3 7} < n^2 < n^2 \lg^5 n$

$\rightarrow$  Durum 1 değil

$f(n) = n^2 \lg^5 n = \Omega(n^{\log_3 7 + \epsilon})$  Alt sınırı, yani

$\rightarrow n^2 \lg^5 n$  daha büyük ✓

Rizgönük kasusı seçtiyorsa mı?

a.  $f(\frac{n}{b}) \leq c f(n), 0 < c < 1$  ise

$7 \cdot \left(\frac{n}{3}\right)^2 \cdot \cancel{\lg^5\left(\frac{n}{3}\right)} \leq c \cdot \cancel{n^2 \lg^5 n}$

$7 \cdot \left(\frac{n}{3}\right)^2 \leq c n^2$

$7 \cdot \frac{n^2}{9} \leq c n^2$

$\frac{7}{9} \leq c < 1$

$$T(n) = O(f(n)) = O(n^2 \lg^5 n)$$