

Answer Key

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|-------------|--------------|--------------|
| 1. D | 8. D | 15. E |
| 2. B | 9. A | 16. D |
| 3. E | 10. B | 17. E |
| 4. D | 11. A | 18. A |
| 5. B | 12. C | 19. C |
| 6. C | 13. C | 20. B |
| 7. B | 14. A | 21. B |

Answer Explanations

- (D)** Tail recursion is when the recursive call of a method is made as the last executable step of the method. Divide-and-conquer algorithms like those used in merge sort or quicksort have recursive calls *before* the last step. Thus, statement II is false.
- (B)** Code segment I is wrong because there is no base case. Code segment III is wrong because, besides anything else, `sum(n)` prevents the method from terminating—the base case `n == 1` will not be reached.
- (E)** When `stringRecur` is invoked, it calls itself irrespective of the length of `s`. Since there is no action that leads to termination, the method will not terminate until the computer runs out of memory (run-time error).
- (D)** The base case is `s.length() ≥ 15`. Since `s` gets longer on each method call, the method will eventually terminate. If the original length of `s` is ≥ 15 , the method will terminate without output on the first call.
- (B)** Letting R denote the method result, we have

$$\begin{aligned}
 R(5) &= 2 * R(4) \\
 &= 2 * (2 * (R(3))) \\
 &= \dots \\
 &= 2 * (2 * (2 * (2 * R(1)))) \\
 &= 2^5 \\
 &= 32
 \end{aligned}$$

- (C)** For `result(n)` there will be $(n - 1)$ recursive calls before `result(1)`, the base case, is reached. Adding the initial call gives a total of n method calls.
- (B)** This method returns the n th term of an arithmetic sequence with first term a and common difference d . Letting M denote method `mystery`, we have

$$\begin{aligned}
 M(3, 2, 6) &= 6 + M(2, 2, 6) \\
 &= 6 + (6 + M(1, 2, 6)) \quad (\text{base case}) \\
 &= 6 + 6 + 2 \\
 &= 14
 \end{aligned}$$

- (D)** Here are the recursive calls that are made, in order: $f(6, 8) \rightarrow f(6, 2) \rightarrow f(4, 2) \rightarrow f(2, 2)$, base case. Thus, 2 is returned.