

# Lecture 4: Games with Infinite Strategy Spaces

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## The Road Ahead

1. Introduction to Infinite Strategy Spaces
2. Games Without Payoff Matrices
3. Existence and Non-existence of Equilibria
4. Second Price Auctions
5. The Median Voter Theorem
6. Applications and Extensions

## Introduction to Infinite Strategy Spaces

**So far:** All games had **finite** number of pure strategies → could represent with matrices or game trees

**Reality:** Many strategic situations involve **infinite** strategy choices:

- Firms choosing production levels (any real number  $\geq 0$ )
- Politicians choosing policy positions (continuum of possibilities)
- Auction bidders choosing bid amounts (any positive value)
- Pricing decisions (continuous price range)

**Key insight:** Matrices and trees become **impractical** or **impossible** with infinite strategies

**New challenge:** How to find equilibria without drawing payoff matrices?

## Games Without Payoff Matrices

**Game definition remains the same:**

- Set of players
- Strategy spaces (now possibly infinite)
- Order of moves
- Payoff functions

**New solution method** for infinite games:

1. Consider a single pure strategy from one player
2. Find the other player's best responses to that strategy
3. Check whether the original strategy is a best response to any of those best responses
4. If yes → equilibrium found; if no → try next strategy
5. Repeat for all strategies

**Advantage:** No need to draw massive (impossible) matrices!

## Example: Simple Infinite Game

**Setup:** Players 1 and 2 simultaneously select whole numbers between 1 and 100. Each player's payoff = product of the two numbers.

**Analysis using new method:**

**Step 1:** Suppose Player 2 chooses 1

- Player 1's payoff =  $x \times 1 = x \rightarrow$  best response: choose 100
- But Player 2's best response to 100 is 100 (not 1)  $\rightarrow$  No equilibrium

**Step 2:** Suppose Player 2 chooses any  $k < 100$

- Player 1's payoff =  $x \times k \rightarrow$  best response: choose 100
- But Player 2's best response to 100 is 100 (not  $k$ )  $\rightarrow$  No equilibrium

**Step 3:** Suppose Player 2 chooses 100

- Player 1's payoff =  $x \times 100 \rightarrow$  best response: choose 100
- Player 2's best response to 100 is 100  $\checkmark \rightarrow$  **Equilibrium: (100, 100)**

## Equilibrium Existence in Infinite Games

**Nash's theorem:** Every **finite** game has at least one Nash equilibrium

**Warning:** Infinite games may have **no equilibrium!**

**Counter-example:** Players select any number  $> 0$ , payoff = product of numbers

**Analysis:**

- Suppose Player 1 chooses any value  $x > 0$
- Player 2's best response: choose something slightly larger than  $x$
- But then Player 1 wants to choose something even larger
- **Result:** No equilibrium exists (infinite regress)

**Key lesson:** Existence theorems don't apply to infinite games

## When Do Infinite Games Have Equilibria?

**Sufficient conditions** for equilibrium existence:

1. **Bounded strategy spaces** (e.g.,  $[0, 100]$  instead of all positive numbers)
2. **Continuous payoff functions**
3. **Compact strategy sets**
4. **Quasi-concave payoff functions**

**Example of bounded game:** Choose number in  $[1, 100]$

- Upper bound prevents infinite regress
- Guarantees equilibrium exists

**Example without bounds:** Choose any positive number

- No upper limit  $\rightarrow$  players keep trying to choose larger numbers
- No equilibrium possible

## Applications: Second Price Auctions

### Setup:

- $n$  bidders compete for single item
- Each bidder has private valuation  $v_i$  for the item
- Sealed bid auction: highest bidder wins, pays **second-highest bid**

**Strategy space:** Each bidder chooses bid  $b_i \in [0, \infty)$

### Payoff function:

- Winner:  $v_i - (\text{second highest bid})$
- Losers: 0

**Key question:** What is the optimal bidding strategy?



## Second Price Auction Analysis

**Claim:** Bidding your true valuation ( $b_i = v_i$ ) is a **dominant strategy**

**Proof idea:**

**Case 1:** Your bid wins

- You pay second-highest bid (independent of your exact bid)
- Bidding higher than  $v_i$ : same outcome, but risk paying more than value
- Bidding lower than  $v_i$ : might lose auction you could have won profitably

**Case 2:** Your bid loses

- You pay 0 regardless of exact bid amount
- No incentive to deviate from truth-telling

**Result:** Truth-telling is optimal regardless of others' strategies → **dominant strategy equilibrium**

## Applications: The Median Voter Theorem

### Setup:

- Two candidates choose policy positions on a line (e.g., left-right spectrum)
- Voters distributed along this line with single-peaked preferences
- Each voter supports candidate closest to their ideal position
- Candidates want to maximize vote share

**Strategy space:** Choose any position  $x \in \mathbb{R}$

**Key insight:** Both candidates will converge to the **median voter's position**

### Intuition:

- If you're away from median, opponent can move slightly toward median
- Opponent captures majority of voters between your positions
- Only stable outcome: both at median position

## Median Voter Theorem: Formal Analysis

### Assumptions:

1. Voters have single-peaked preferences
2. Candidates care only about winning (not ideology)
3. Voters vote for candidate closest to their ideal point
4. Odd number of voters (to avoid ties)

**Equilibrium prediction:** Both candidates locate at median voter's ideal point

### Real-world implications:

- Explains convergence to political center
- Why extreme candidates often lose general elections
- Strategic moderation in two-party systems

**Extensions:** Result may fail with multiple dimensions, strategic voting, or candidate ideology

## Summary: Key Insights

### Main takeaways from infinite strategy games:

1. **Solution method:** Use best-response analysis instead of matrices
2. **Existence:** Infinite games may have no equilibrium (unlike finite games)
3. **Bounded sets:** Help guarantee equilibrium existence
4. **Dominant strategies:** Still powerful solution concept (second-price auctions)
5. **Convergence:** Competition can lead to clustering (median voter theorem)

### Applications beyond this lecture:

- Cournot competition (firms choosing quantities)
- Bertrand competition (firms choosing prices)
- Public goods provision
- Arms races and conflict models

## Practice Problems

**Problem 1:** Two firms choose production levels  $x_1, x_2 \in [0, 10]$ . Profit functions:  $\pi_1 = x_1(12 - x_1 - x_2)$ ,  $\pi_2 = x_2(12 - x_1 - x_2)$ . Find Nash equilibrium.

**Problem 2:** In a first-price auction (highest bidder wins and pays their bid), is truth-telling still optimal? Why or why not?

**Problem 3:** Consider median voter theorem with three candidates instead of two. What happens to the equilibrium prediction?

**Problem 4:** Players choose numbers in  $[0, 1]$ . Player 1's payoff =  $x_1 - x_1^2$ , Player 2's payoff =  $x_2 - 2x_1x_2$ . Find best responses and equilibrium.

## Takeaway Points

1. Infinite strategy spaces require new solution methods beyond payoff matrices.
2. Nash's existence theorem doesn't apply to infinite games - equilibria may not exist.
3. Bounded strategy sets and continuous payoffs help ensure equilibrium existence.
4. Truth-telling is dominant in second-price auctions but not in first-price auctions.
5. The median voter theorem explains convergence to the center in two-candidate elections.