# **Lecture 2: Extensive Form Games**

**Instructor**: Fei Tan

@econdojo Decondojo Decond

Saint Louis University

Date: September 4, 2025

## The Road Ahead

- 1. From Simultaneous to Sequential Games
- 2. Subgame Perfect Equilibrium
- 3. Backward Induction
- 4. Multiple Subgame Perfect Equilibria
- 5. Making Threats Credible
- 6. Commitment Problems
- 7. Backward Induction with N Players
- 8. Forward Induction

## From Simultaneous to Sequential Games

**Simultaneous Move Games**: players move at the same time or cannot observe each other's moves

- Prisoner's dilemma: prisoners in separate interrogation rooms
- Football: offense and defense call plays simultaneously
- Meeting coordination: couple choosing locations without communication

Sequential Games: strategic interactions that flow over time in specific steps

- Military invasion followed by response decision
- Police officer requesting search permission, then deciding thoroughness
- Chess: white moves, black responds, cycle repeats
- Firm entry followed by incumbent's competitive response

# **Introducing Selten's Game**

Firm 1 considers entering Firm 2's monopoly market. If Firm 1 enters, Firm 2 must decide whether to accommodate entry or wage a price war.

## **Payoff Structure:**

- Price war: both firms earn 0 (profits eliminated)
- Accommodation: Firm 1 earns 3, Firm 2 earns 1
- No entry: Firm 1 earns 2 (saves investment), Firm 2 earns 2 (maintains monopoly)

	Firm 2: Accommodate	Firm 2: War
Firm 1: Enter	3, 1	0, 0
Firm 1: Stay Out	2, 2	2, 2

## **Extensive Form Games**

**Game Tree Representation** of the firm entry game:

## **Key Elements**:

- **Decision nodes**: where players make choices
- **Terminal nodes**: end points with payoffs
- Branches: represent available strategies
- Sequential structure: order of play is explicit

# Multiple Nash Equilibria in Strategic Form

Recall Selten's game has multiple Nash equilibria:

### 1. Pure strategy equilibria:

- (Enter, Accommodate): (3, 1)
- (Stay Out, War): (2, 2)

### 2. Mixed strategy equilibria:

- Firm 1 plays Stay Out with probability 1
- $\circ$  Firm 2 plays War with probability  $\geq rac{1}{3}$

Multiple equilibria make prediction difficult. Which one will actually occur? **Sequential structure** helps us resolve this ambiguity!

# **Analyzing Sequential Play**

If Firm 1 enters, Firm 2 faces this decision:

```
Firm 2
/ \
Accommodate War
/ (3, 1) (0, 0)
```

Firm 2's analysis: 1 > 0, so Firm 2 prefers Accommodate over War

Firm 1's analysis (knowing Firm 2 will accommodate):

- Enter → Firm 2 accommodates → payoff = 3
- Stay Out → payoff = 2
- Since 3 > 2, Firm 1 should enter

Unique sequential solution: (Enter, Accommodate) with payoffs (3, 1)

## **Subgame Perfect Equilibrium**

A **Subgame Perfect Equilibrium (SPE)** is a strategy profile where each player's strategy constitutes a Nash equilibrium in every subgame.

Key insight: SPE ensures that threats are credible

- Firm 2's threat to wage war is NOT credible
- Once Firm 1 enters, Firm 2 has no incentive to follow through
- Firm 1 recognizes this and enters anyway

Refinement: SPE eliminates Nash equilibria that rely on non-credible threats

- All SPE are Nash equilibria
- Not all Nash equilibria are SPE
- SPE is the gold standard for extensive form games

## Games with Simultaneous Moves in Extensive Form

Some extensive form games include simultaneous moves, e.g. matching pennies

```
Player 1

Heads Tails

Player 2 --- Player 2 <- information set

Heads Tails Heads Tails

(1,-1) (-1,1) (-1,1) (1,-1)
```

Information set: dashed line shows Player 2 cannot observe Player 1's choice

Solution method: convert to matrix form and solve as simultaneous game

# **Constructing Games with Simultaneous Moves**

Critical rules for information sets:

- 1. Identical strategy sets: same available actions at all nodes in the information set
- 2. **Irrelevance of player order**: whoever plays first, the resulting payoff matrix must be the same

### **Example of violation:**

```
If Player 1 chooses Heads → Player 2 chooses {A, B}
If Player 1 chooses Tails → Player 2 chooses {C, D}
```

**Problem**: Player 2 can infer Player 1's move from available actions, violating simultaneity

## **Backward Induction**

**Method**: solve extensive form games by working backwards from the end:

- 1. Start at the final decision nodes and determine optimal actions
- 2. Work backwards using optimal future play to determine current optimal actions
- 3. Continue until reaching the initial node

Backward induction always finds the subgame perfect equilibrium

**Note**: complete strategy must specify actions at ALL decision nodes, even those not reached in equilibrium

# **Example: Escalation Game**

Two countries are on the brink of war:

```
Player 1: Accept —— (0, 0)

Threaten — Player 2: Concede —— (1, -2)

Escalate — Player 1: War —— (-1, -1)

Back Down —— (-2, 1)
```

**Step 1**: Player 1 chooses between War (-1) and Back Down (-2). Since -1 > -2, choose War.

**Step 2**: Player 2 chooses between Concede (-2) and Escalate (leading to War: -1). Since -1 > -2, choose Escalate.

Step 3: Player 1 chooses between Accept (0) and Threaten (leading to Escalate → War: -1). Since 0 > -1, choose Accept.

SPE: ((Accept, War), Escalate) with outcome (0, 0)

# **Example: Ultimatum Game**

Setup: Player 1 has a good worth 2 and must bargain with Player 2 over division

- **Split**: Offer equal division (1, 1)
- Take: Attempt to take everything (2, 0)
- Player 2 can Accept or Reject any proposal
- Rejection leads to (0, 0) for both

# Multiple SPE in Ultimatum Game

After Split: Player 2 prefers Accept (payoff 1) over Reject (payoff 0)

After Take: Player 2 is indifferent between Accept (payoff 0) and Reject (payoff 0)

Multiple Equilibria Arise from Player 2's indifference:

- 1. Player 2 always Accepts: Player 1 takes (payoff 2 > 1), so SPE is \(\tag{Take, (Accept, Accept})\)
- 2. **Player 2 always Rejects**: Player 1 splits (payoff 1 > 0), so SPE is <Split, (Accept, Reject)>
- 3. Player 2 mixes: Any probability p of accepting after "Take"
  - $\circ$  If p > 1/2: Player 1 takes
  - ∘ If p < 1/2: Player 1 splits
  - $\circ$  If p = 1/2: Player 1 indifferent, can mix

# Making Threats Credible: Burning Bridges

The Classic Story: Two armies fight over an island. Each has a bridge for access.

#### **Strategic Problem:**

- Island is valuable but not worth fighting for
- Each army prefers to concede rather than fight
- First army occupies island, second decides whether to invade

#### Game Tree:

```
Army 1: Burn Bridge?

Burn Don't Burn

Army 2 Army 2 invades

Don't Invade Fight Retreat

(1,0) (-1,-1) Army 1 (0,1)

Fight Retreat

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```

# **Bridge Burning Analysis**

#### **Backward Induction:**

**Step 1** - If bridge not burned and Army 2 invades:

Army 1 chooses Fight(-1) vs Retreat(0) → Choose Retreat

**Step 2** - Army 2's decision if bridge not burned:

Don't Invade(0) vs Invade(1, since Army 1 retreats) → Choose Invade

**Step 3** - Army 2's decision if bridge burned:

Don't Invade(0) vs Invade(-1, since Army 1 must fight) → Choose Don't Invade

Step 4 - Army 1's initial decision:

Burn(1) vs Don't Burn(0) → Choose Burn

**SPE**: 〈(Burn, Retreat), (Don't Invade, Invade)〉

Key Insight: Burning bridges makes the threat to fight credible!

# **Tying Hands**

**Story**: Boss wants to stop employee theft but prefers keeping employee over firing **Without Warning**:

```
Employee steals \rightarrow Boss: Fire(-1) vs Ignore(0) \rightarrow Boss ignores \rightarrow Employee steals
```

## With Warning (Boss values reputation):

```
Employee steals after warning → Boss: Fire(-1) vs Ignore(-2) → Boss fires → Employee doesn't steal
```

## **Analysis:**

- Warning changes Boss's payoffs (reputation cost for not following through)
- Makes threat to fire credible
- Employee responds by not stealing

SPE: ((Warning, Ignore, Follow Through), (Steal, Don't Steal))

## **Commitment Problems**

**Definition**: Situations where the ability to commit to future actions provides strategic advantage

### **Examples**:

- 1. Burning Bridges: Remove option to retreat, making fight threat credible
- 2. **Tying Hands**: Make deviation from threat extremely costly
- 3. Reputation Building: Invest in credibility of future threats

Paradox: Reducing your own options can improve your payoff

### Applications:

- Military strategy and deterrence
- Business competition and market entry
- Labor negotiations and strikes
- International relations and sanctions

# The Pirate Game (5 Pirates, 10 Gold Coins)

#### Rules:

- Hierarchy: Nash > Pirate 2 > Pirate 3 > Pirate 4 > Pirate 5
- Senior-most pirate proposes allocation
- Majority vote required (including proposer's vote)
- If proposal fails, proposer walks the plank, next pirate proposes
- All pirates are rational, greedy, but prefer life over death

Question: What allocation will Nash (most senior) propose?

## **Pirate Game: Backward Induction**

If only Pirate 4 and 5 remain: Pirate 4 proposes (10, 0) and wins with his own vote If Pirates 3, 4, 5 remain:

- Pirate 3 needs 2 votes total (including his own)
- Pirate 4 rejects any offer (can get 10 coins if Pirate 3 dies)
- Pirate 5 gets 0 if Pirate 3 dies, so accepts any positive offer
- Pirate 3's optimal proposal: (9, 0, 1)

### If Pirates 2, 3, 4, 5 remain:

- Pirate 2 needs 2 votes total
- Pirate 3 gets 9 if Pirate 2 dies (expensive to buy)
- Pirate 4 gets 0 if Pirate 2 dies (cheap to buy needs only 1 coin)
- Pirate 2's optimal proposal: (9, 0, 1, 0)

## **Pirate Game Solution**

### Nash's Decision (all 5 pirates):

- Needs 3 votes total (including his own)
- Pirate 2 gets 9 if Nash dies (too expensive)
- Pirates 3 and 5 get 0 if Nash dies (each needs only 1 coin)
- Pirate 4 gets 1 if Nash dies (needs 2 coins)

Nash's Optimal Strategy: Buy Pirates 3 and 5 with 1 coin each

Final Allocation: Nash(8), Pirate 2(0), Pirate 3(1), Pirate 4(0), Pirate 5(1)

Votes: Nash ✓, Pirate 2 ×, Pirate 3 ✓, Pirate 4 ×, Pirate 5 ✓

**Key Insight**: Even the most senior player cannot extract everything - must share just enough to create winning coalition

## Nim: Game of Perfect Information

Setup: 21 chips, players alternate taking 1 or 2 chips, last player to take chips wins

### **Backward Induction Logic:**

- 1 chip left: Player must take it and wins
- 2 chips left: Player takes both and wins
- 3 chips left: Player loses (opponent will win regardless)
- 4 chips left: Player takes 1, leaving 3 for opponent → Player wins
- 5 chips left: Player takes 2, leaving 3 for opponent → Player wins
- 6 chips left: Player loses (opponent can force win)

Pattern: Player loses with 3, 6, 9, 12, 15, 18, 21 chips remaining

21-chip game: Player 1 starts with 21 chips and is in a losing position!

Optimal strategy: Leave opponent with multiple of 3 chips

## **Forward Induction**

**Definition**: Forward induction uses the assumption that all previous play was rational to make inferences about opponents' private information or strategies

## **Key Difference from Backward Induction:**

- Backward Induction: Assumes optimal play in the future
- Forward Induction: Assumes rational play in the past

**Application**: When players observe actions that seem suboptimal, forward induction helps explain why those actions might actually be rational

Requirement: Extremely sophisticated thinking and strong rationality assumptions

# Forward Induction Example: Battle of Sexes with Burning Money

Modified Battle of the Sexes: Player 1 can burn money before choosing

Standard Battle of the Sexes:

	Ballet	Fight
Ballet	1, 2	0, 0
Fight	0, 0	2, 1

With Money Burning: Player 1 can burn \$1 before the game, changing his payoffs to (0,1) and (1,0)

#### **Forward Induction Logic:**

- 1. If Player 1 burns money, he must be planning to go to Fight (only way to get positive payoff)
- 2. Knowing this, Player 2 should also choose Fight
- 3. This gives Player 1 payoff of 1 (better than mixing in original game) Fei Tan | Made on Earth by humans.
  - 4. Therefore, Player 1 burns money and both go to Fight

# **Takeaway Points**

- 1. **Subgame Perfect Equilibrium** requires credible threats eliminates Nash equilibria based on non-credible threats
- 2. **Backward Induction** systematically finds SPE by analyzing optimal play from the end of the game backwards
- 3. **Multiple SPE** can exist when players are indifferent between actions at some decision nodes
- 4. **Commitment Strategies** (burning bridges, tying hands) can improve payoffs by making threats credible
- 5. **Sequential games** often have unique predictions even when the corresponding simultaneous game has multiple equilibria
- 6. **Forward induction** uses rationality assumptions about past play to make inferences about opponents' strategies
- 7. **Complete strategies** must specify actions at all decision nodes, not just those

# **Applications and Extensions**

#### **Real-World Applications:**

- Business Strategy: Entry deterrence, capacity investment, price wars
- Military Strategy: Nuclear deterrence, alliance formation
- Political Science: Legislative bargaining, coalition formation
- Economics: Auction design, contract theory, mechanism design

#### **Advanced Topics** (beyond this course):

- Perfect Bayesian Equilibrium (incomplete information)
- Repeated games and reputation
- Bargaining theory
- Evolutionary game theory

**Next Lecture**: Advanced Strategic Form Games - dominance, rationalizability, and equilibrium refinements