

# 1 Properties of Relations

- Reflexivity:  $\forall x \in S : (x, x) \in \mathcal{R}$
- Completeness:  $\forall x, y \in S : x \neq y \Rightarrow (x, y) \in \mathcal{R} \text{ or } (y, x) \in \mathcal{R}$
- Transitivity:  $\forall x, y, z \in S : ((x, y) \in \mathcal{R} \text{ and } (y, z) \in \mathcal{R}) \Rightarrow (x, z) \in \mathcal{R}$
- Symmetry:  $\forall x, y \in S : (x, y) \in \mathcal{R} \Rightarrow (y, x) \in \mathcal{R}$
- Anti-symmetry:  $\forall x, y \in S : ((x, y) \in \mathcal{R} \text{ and } (y, x) \in \mathcal{R}) \Rightarrow x = y$
- Asymmetry:  $\forall x, y \in S : (x, y) \in \mathcal{R} \Rightarrow (y, x) \notin \mathcal{R}$
- Negative transitivity:  $\forall x, y, z \in S : ((x, y) \notin \mathcal{R} \text{ and } (y, z) \notin \mathcal{R}) \Rightarrow (x, z) \notin \mathcal{R}$
- Equivalence: Relation which is symmetric, reflexive and transitive.

## 1.1 The indifference Relation

We know that the preference relation  $\succsim$  on  $A$  is Reflexive, Transitive and Complete. Define the indifference relation on  $A$  as follows;

$$\sim = \{(a, b) \in A \times A \mid (a, b) \in \succsim \wedge (b, a) \in \succsim\}$$

Now we can check for the reflexivity, transitivity, completeness and symmetry of the indifference relation  $\sim$ ;

- Reflexivity: By the definition of the indifference relation  $\sim$  on  $A$  we have,

$$\begin{aligned} \sim \text{ on } A &= \{(a, a) \in A \times A \mid (a, a) \in \succsim \wedge (a, a) \in \succsim\} \\ &\implies (a, a) \in \sim \end{aligned}$$

- Construct an example by yourself to show that  $\sim$  is not complete.
- Transitivity: We want to show,  $(\forall A)(\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\sim \text{ relation derived in the above manner is transitive})$

or,

$$(\forall a, b, c \in A) ((a, b) \in \sim \wedge (b, c) \in \sim) \implies ((a, c) \in \sim)$$

Proof; Consider any arbitrary  $A$ , and any arbitrary  $\succsim$  on  $A$  that is transitive, reflexive and complete.

Now consider the indifference relation

$$\sim = \{(a, b) \in A \times A \mid (a, b) \in \succsim \wedge (b, a) \in \succsim\}$$

and any arbitrary  $a, b, c \in A$  such that  $(a, b) \in \sim$  and  $(b, c) \in \sim$   
Now since  $(a, b) \in \sim$  and  $(b, c) \in \sim$  then by the definition of the indifference relation,

$$(a, b) \in \succsim \wedge (b, a) \in \succsim \quad (1)$$

$$(b, c) \in \succsim \wedge (c, b) \in \succsim \quad (2)$$

Then from the transitivity of  $\succsim$  and from (1), (2) we get,

$$(a, c) \in \succsim \text{ and } (c, a) \in \succsim$$

so  $(a, a) \in \succsim$  or by the definition of  $\sim$ ,  $(a, c) \in \sim$

Hence proved that the indifference relation  $\sim$  is transitive.

- Symmetry: We want to show,  $(\forall A) (\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\sim \text{ relation derived in the above manner is symmetric})$

Proof; Consider any arbitrary  $A$ , and any arbitrary  $\succsim$  on  $A$  that is transitive, reflexive and complete.

Now consider the indifference relation

$$\sim = \{(a, b) \in A \times A \mid (a, b) \in \succsim \wedge (b, a) \in \succsim\}$$

and the symmetric relation  $\mathcal{R}$  on  $A$  defined as follows,

$$(\forall a, b \in A) ((a, b) \in \mathcal{R} \implies (b, a) \in \mathcal{R})$$

Now consider arbitrary  $a, b$  such that  $(a, b) \in \sim$ ,

Then by the definition of the indifference relation we get,

$$(a, b) \in \succsim \wedge (b, a) \in \succsim$$

which can be rewritten as

$$(b, a) \in \succsim \wedge (a, b) \in \succsim$$

and therefore  $(b, a) \in \sim$

So we have shown that if the preference relation  $\succsim$  is reflexive, transitive and complete, then the indifference relation  $\sim$  derived from it is reflexive, transitive and symmetric.

### Question

Suppose  $A = \{a, b, c\}$  and  $\mathcal{R} = \{(a, b), (b, a), (a, a)\}$ .

Is  $\mathcal{R}$  a valid indifference relation?

### Answer

Note that  $\mathcal{R}$  is not reflexive, not complete and neither transitive and an indifference relation is always reflexive, transitive and symmetric so therefore,  $\mathcal{R}$  is not a valid indifference relation.

## 1.2 The strict preference relation

We know that the preference relation  $\succsim$  on  $A$  is reflexive, transitive and complete. From this we can define the strict preference relation as follows;

$$\succ = \{(a, b) \in A \times A \mid (a, b) \in \succsim \wedge (b, a) \notin \succsim\}$$

So,  $(a, b) \in \succ \wedge \neg(b, a) \in \succ$

Now we can check for the reflexivity, transitivity, completeness and symmetry of the strict preference relation  $\succ$  as follows;

- Reflexivity: We want to show,  $(\forall A)(\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\succ \text{ relation derived in the above manner is reflexive})$

for the proof, try to negate the above proposition and you can see that the strict preference relation is not reflexive.

- Completeness: try to find an example to show that the strict preference relation is not complete.

Consider  $A = \{a, b, c\}$  and the preference relation  $\succsim = A \times A$ .

Now because of the given preference the relation the strict preference relation will be an empty set or  $\succ = \phi$  which is neither reflexive and nor complete.

- Transitivity: We want to show,  $(\forall A)(\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\succ \text{ relation derived in the above manner is transitive})$

Proof; Consider any arbitrary  $A$  and an arbitrary  $\succsim$  satisfying reflexivity, transitivity and completeness and also consider any arbitrary  $a, b, c$  in  $A$  such that  $(a, b) \in \succ$  and  $(b, c) \in \succ$ .

Now because  $(a, b) \in \succ$  and  $(b, c) \in \succ$  we get,

$$(a, b) \in \succsim \wedge (b, a) \notin \succsim \quad (3)$$

$$(b, c) \in \succsim \wedge (c, b) \notin \succsim \quad (4)$$

and because the preference relation  $\succsim$  is transitive and from (3), (4) we get that  $(a, c) \in \succsim$  and now we want to show that  $(c, a) \notin \succsim$

We will do so by the way of contradiction:-

Suppose  $(c, a) \in \succsim$ , given that  $(a, b) \in \succ$  from, the previous assumption, we get that  $(c, b) \in \succsim$  because the preference relation is transitive.

But this is a contradiction because we know that  $(c, b) \notin \succsim$  from (4) and therefore,  $(c, a) \notin \succsim$

So we have  $(a, c) \in \succsim \wedge (c, a) \notin \succsim$ , which clearly implies that the strict preference relation is transitive.

So we have shown that if the preference relation  $\succsim$  is reflexive, transitive and complete, then the strict preference relation  $\succ$  derived from it is not reflexive, nor complete but transitive.

### Examples

$$A = \{a, b, c\}$$

$$\mathcal{R} = \{(a, a), (a, b), (b, b)\}$$

Note that  $\mathcal{R}$  is Anti-symmetric but not Asymmetric and it is also not negative transitive because  $(a, c) \notin \mathcal{R}$  and  $(c, b) \notin \mathcal{R}$  but  $(a, b) \in \mathcal{R}$

How to negate  $\neg((\forall x), (\forall y), (P(x) \implies Q(y)))$

$$\begin{aligned} \neg((\forall x), (\forall y), (P(x) \implies Q(y))) &\equiv ((\exists x) (\exists y) (\neg(P(x) \implies Q(y)))) \\ &\equiv ((\exists x) (\exists y) (P(x) \wedge \neg Q(y))) \end{aligned}$$

**Propositon 1**

$(\forall A) (\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\succ \text{ relation derived from } \succsim \text{ is asymmetric})$

Proof; Consider any arbitrary  $a, b \in A$  and  $(a, b) \in \succ$   
Now because  $(a, b) \in \succ$  we get

$$(a, b) \in \succsim \wedge (b, a) \notin \succsim$$

but this implies that  $(b, a) \notin \succ$   
and therefore the strict prefernce relation  $\succ$  is assymetric.

**Propositon 2**

$(\forall A) (\forall \succsim \text{ on } A \text{ that are transitive, reflexive and complete}) (\succ \text{ relation derived from } \succsim \text{ is anti-symmetric})$

this is a corollary of the above proof because

$(\forall \mathcal{R} \text{ on } A) (\mathcal{R} \text{ is asymmetric} \implies \mathcal{R} \text{ is anti-symmetric})$

**Propositon 3**

$(\forall A) (\forall \succsim$  on  $A$  that are transitive, reflexive and complete) ( $\succ$  relation derived from  $\succsim$  is Negative transitive)

Proof. Consider any arbitrary  $A$ , and any weak preference relation  $\succsim$  on set  $A$  that is reflexive, complete and transitive. Now consider arbitrary  $a, b, c \in A$  such that  $(a, b) \notin \succ$  and  $(b, c) \notin \succ$ . We will try and show that  $(a, c) \notin \succ$ .

Suppose, by way of contradiction,  $(a, c) \in \succ$ .

Step 1. By completeness of  $\succsim$ , we also have  $(c, b) \in \succsim$ . By transitivity of  $\succsim$ . This implies that  $(a, b) \in \succsim$  and we also have  $(b, a) \in \succsim$  by completeness of  $\succsim$ . So, we get  $(a, b) \in \sim$ .

Step 2. In the similar fashion as above, by completeness of  $\succsim$ , we also have  $(b, a) \in \succsim$ . By transitivity of  $\succsim$ . This implies that  $(b, c) \in \succsim$  and we also have  $(c, b) \in \succsim$  by completeness of  $\succsim$ . So, we get  $(b, c) \in \sim$ .

So, by step 1 and 2, we get that  $(a, c) \in \sim$  (by transitivity of  $\succsim$ ). This is a contradiction.

Therefore,  $(a, c) \notin \succ$ .

Alternative Proof (Direct Proof). Consider any arbitrary  $A$ , and any weak preference relation  $\succsim$  on set  $A$  that is reflexive, complete and transitive. Now consider arbitrary  $a, b, c \in A$  such that  $(a, b) \notin \succ$  and  $(b, c) \notin \succ$ . By completeness of  $\succsim$   $(b, a) \in \succsim$  and  $(c, b) \in \succsim$ . By transitivity of  $\succsim$   $(c, a) \in \succsim$ . Therefore,  $(a, c) \notin \succ$ .

**Question**

Consider a binary relation  $\succsim$  on a set  $A$ . Suppose  $\succsim$  is transitive. Define relations  $\succ$  and  $\sim$  on  $A$  by: for  $x, y \in A$ ,

$x \succ y$  if and only if  $x \succsim y$  and not  $y \succsim x$

$x \sim y$  if and only if  $x \succsim y$  and  $y \succsim x$

Prove the following. (i) If  $x \succ y$  and  $y \succ z$  then  $x \succ z$

(ii) If  $x \sim y$  and  $y \sim z$  then  $x \sim z$

(iii) If  $x \succ y$  and  $y \succsim z$  then  $x \succ z$

Proof. Let  $\succsim$  be transitive, then we want to prove that

$$(x \succ y \wedge y \succsim z) \implies x \succ z$$

Consider any arbitrary  $x, y, z \in A$  such that;

$x \succ y \wedge y \succsim z$  holds

Suppose by the way of contradiction  $x \succ z$  does not hold  
then by definition of strict preference,

$$\begin{aligned} & \neg(x \succ z \wedge \neg(z \succ x)) \\ \implies & \neg(x \succ z) \vee z \succ x \end{aligned}$$

But then by the transitivity of  $\succsim$

$$z \succ x \wedge y \succsim z \implies y \succ x$$

But this contradicts the fact that  $x \succ y$  holds. Hence proved!