1. Extensive Games with Perfect Information

An Extensive game with perfect information consists of the following;

- Set of players
- Set of Sequences (terminal histories) with proper entry that no sequence is a proper subhistory of any other sequence.
- A player function that assigns a player to every sequence that is a proper subhistory of some terminal history.
- Prefernces for each player over the set of terminal histories.

Examples

Game 1: Simultaneous move Duoploy

An enterant firm decides to enter an Industry or not while the icumbent firm in that industry decides to wether to accommodate or fight the new enterant firm.

- Set of players: $N = \text{Enterant}(\mathbf{E}), \text{Incumbent}(\mathbf{I})$
- Action Set of Each Player: $T = \{ \text{StayOut}, (\text{Enter}, \text{Fight}), (\text{Enter}, \text{Acommodate}) \}$
- Player Function:

$$\mathcal{P}(\phi) = E$$

 $\mathcal{P}(Enter) = I$

• Payoffs of both players:

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U_E(\text{Stay out}) = 0

U_E(\text{Enter,Fight}) = -50

U_E(\text{Enter,Accomodate}) = 50

U_I(\text{Stay out}) = 100

U_I(\text{Stay out}) = 0

U_I(\text{Stay out}) = 50
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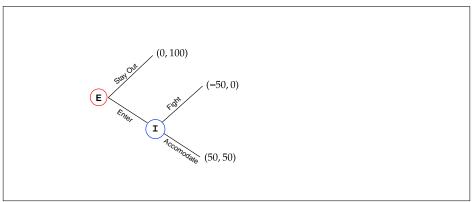


Figure 1: Simultaneous move Duopoly

Nash Equilibrium, $NE = \{(Stayout, Fight), (Enter, Accomodate)\}$ Subgame Perfect Equilibrium, $SPE = \{(Enter, Accomodate)\}$

Note that (Stayout, Fight) is a Nash equilbrium but not a subgame perfect equilibrium because it is not a Nash Equilibrium of every subgame in this game.

Game 2: @3866, Piazza Contributing Efforts game

Consider the following simultaneous-move effort-contributions' game of n players:

- Set of players: $\{1, 2, \ldots, n\}$
- Action set of each player : $A_i = [0, 1]$
- Payoff of player: $u_i(e_1, e_2, ..., e_n) = n \min(e_1, e_2, ..., e_n) e_i$

Let us solve this game for n = 3 players;

- Strategy of player 3: $S_3: [0,1] \times [0,1], S_3(e_1,e_2)$
- Strategy of Player 2: S_2 : $S_2(e_1)$
- Strategy of Player 1: $S_1 \in [0,1]$

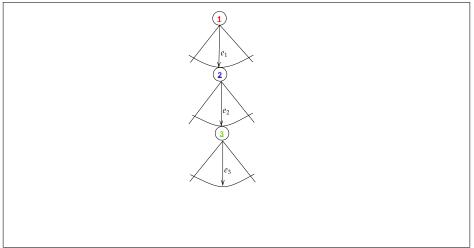


Figure 2: Efforts Contribution Game

Solving the game using backward induction, We solve for the player 3 first; Player 3 will solve the following utility maximization problem given (e_1, e_2) ;

$$\max_{0 \le e_3 \le 1} 3 \min(e_1, e_2, e_3) - e_3$$

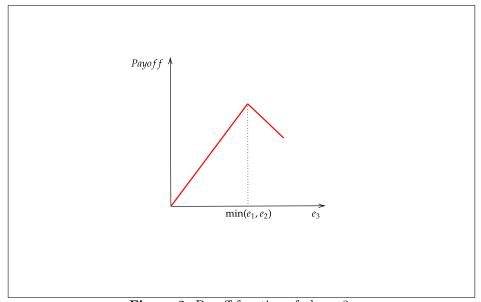


Figure 3: Payoff function of player 3

The solution to this problem is; $e_3 = S_3(e_1, e_2) = \min(e_1, e_2)$ Folding the game backward again, Player 2 will now solve his following utility maximization problem given e_1 ;

$$\max_{0 \le e_2 \le 1} 3 \min(e_1, e_2, e_3) - e_2$$
s.t. $e_3 = S_3(e_1, e_2) = \min(e_1, e_2)$

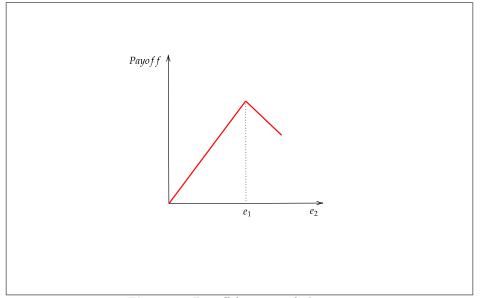


Figure 4: Payoff function of player 2

Solving the above problem by substituting for e_3 into the objective function we can rewrite our problem as

$$\max_{e_1, e_2} 3 \min(e_1, e_2) - e_2$$

solving this we get, $e_2 = S_2(e_1) = e_1$

And then finally player 1 solves his utilty maximization problem;

$$\max_{0 \le e_1 \le 1} 3 \min(e_1, e_2, e_3) - e_1$$
s.t. $e_3 = S_3(e_1, e_2) = \min(e_1, e_2)$

$$e_2 = e_1$$

Solving the above gives $e_1 = 1$

Therefore $e_1^* = e_2^* = e_3^* = \min(e_1, e_2) = 1$

And the resulting terminal History where players move according to the equilibrium strategy (1, 1, 1) is the Subgame Perfect Equilibrium outcome and the SPE strategy is;

$$e_1 = S_1 = 1$$

 $e_2 = S_2 = e_1$
 $e_2 = S_3(e_1, e_2) = \min(e_1, e_2)$

Game 3: Stackelberg Duopoly Using Isoprofit Curves

1. Set of Players;

$$N = \{1, 2\}$$

2. Action set of players;

$$T = \{ (q_1, q_2) \in \mathbb{R}_+ \times \mathbb{R}_+ \}$$

3. Player function;

$$\mathcal{P}(\phi) = 1$$

 $\mathcal{P}(q_1) = 2, \forall q_1 \in \mathbb{R}_+$

4. Payoff of palyer $i \ \forall i \in \{1, 2\};$ $\pi_i(q_1, q_2) = q_i \max(15 - q_1 - q_2, 0)$

Solving this game Using Backward Induction;

First firm 2(follower) solves its profit maximization problem taking output of firm 1, q_1 as given;

$$\max_{q_2 \ge 0} q_2(\max(15 - q_1 - q_2, 0))$$

Solution to which gives us the best response function of player 2 given q_1 , $BR_2(q_1)$

$$BR_2(q_1) \in \begin{cases} \{\frac{15-q_1}{2}\} & \text{if } q_1 < 15\\ \mathbb{R}_+ & \text{if } q_1 \ge 15 \end{cases}$$

So,

$$q_2(q_1) = \begin{cases} \frac{15 - q_1}{2} & \text{if } q_1 < 15\\ 0 & \text{if } q_1 \ge 15 \end{cases}$$

Then the firm 1(leader) solve its profit maximization problem subject to the best response function of the follower;

$$\max_{q_1 \ge 0} q_1(\max(15 - q_1 - q_2, 0))$$

$$s.t. \quad q_2(q_1) = \begin{cases} \frac{15 - q_1}{2} & \text{if } q_1 < 15\\ 0 & \text{if } q_1 \ge 15 \end{cases}$$

The above problem can be rewritten as;

$$\max_{0 \le q_1 \le 15} q_1 \left(\frac{15 - q_1}{2}\right)$$

solving which gives us the following; $q_1^*=7.5, q_2^*=\frac{7.5}{2}$

The above problem can also be solved using the isoprofit curves approach as follow;

$$q_1(15 - q_1 - q_2) = \overline{\pi}$$
differntiating both sides w.r.t q_1

$$\Rightarrow \frac{dq_1}{dq_2} = \frac{15 - 2q_1 - q_2}{q_1}$$
But,
$$\frac{dq_2}{dq_1} = \frac{-1}{2}$$

$$\Rightarrow \frac{15 - 2q_1 - q_2}{q_1} = \frac{-1}{2}$$

 $\implies q_1 = 7.5, q_2 = 3.75$

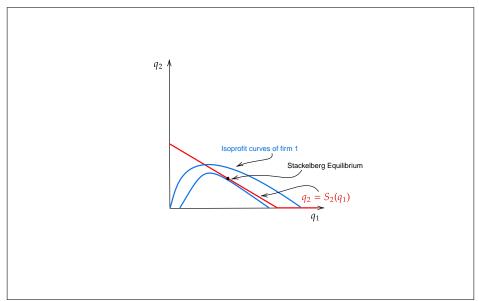


Figure 5: Stackelberg Equilibrium

Game 4: Stackelberg Duopoly with fixed costs

•
$$\pi_i = q_i \max(16 - q_1 - q_2,) - F_i$$
; if $q_i > 0$

$$\bullet \ F_i(q_i) = \begin{cases} 25 & \text{if } q_i > 0\\ 0 & \text{if } q_i = 0 \end{cases}$$

Solving using backward induction follower firm 2 solves its profit maximization problem first taking q_1 as given;

$$\max_{q_2>0} q_2(\max(16-q_1-q_2,0))-25$$

Solving which gives us,

$$BR_2(q_1) \in \begin{cases} \{0\} & \text{if } q_1 > 6\\ \{\frac{16 - q_1}{2}\} & \text{if } q_1 < 6\\ \{0, \frac{16 - q_1}{2}\} & \text{if } q_1 = 6 \end{cases}$$

and,

$$q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 6\\ 0 & \text{if } q_1 \ge 6 \end{cases}$$

Then the leader firm 1 solves it's profit maximization problem subject to followers's best response;

$$\max_{q_1 \ge 0} \quad q_1(\max(16 - q_1 - q_2, 0)) - 25$$

$$s.t. \quad q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 6\\ 0 & \text{if } q_1 \ge 6 \end{cases}$$

The solution to this problem is that firm 1 will produce the monoply outcome of $q_1 = 8$ given follower's strategy and $q_2 = 0$. This game is also known as Natural Monopoly since the monopoly output is enough to deter the entry of the follower firm.

Game 5: Stackelberg Duopoly with fixed costs

•
$$\pi_i = q_i \max(16 - q_1 - q_2) - F_i$$
; if $q_i > 0$

$$\bullet \ F_i(q_i) = \begin{cases} 9 & \text{if } q_i > 0 \\ 0 & \text{if } q_i = 0 \end{cases}$$

Solving using backward induction follower firm 2 solves its profit maximization problem first taking q_1 as given;

$$\max_{q_2>0} q_2(\max(16-q_1-q_2,0)) - 9$$

Solving which gives us,

$$BR_2(q_1) \in \begin{cases} \left\{ \frac{16 - q_1}{2} \right\} & \text{if } q_1 < 10\\ \mathbb{R}_+ & \text{if } q_1 \ge 10 \end{cases}$$

and,

$$q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 10\\ 0 & \text{if } q_1 \ge 10 \end{cases}$$

Then the leader firm 1 solves it's profit maximization problem subject to followers's best response;

$$\max_{q_1 \ge 0} \quad q_1(\max(16 - q_1 - q_2, 0)) - 9$$

$$s.t. \quad q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 10\\ 0 & \text{if } q_1 \ge 10 \end{cases}$$

Unlike the previous game the monopoly outcome $q_1 = 8$ is not enough to deter the entry and hence the leader firm produces $q_1 = 10$ now and $q_2 = 0$. This case is known as the Entry deterrence case.

Game 6: Stackelberg Duopoly with fixed costs

•
$$\pi_i = q_i \max(16 - q_1 - q_2) - F_i$$
; if $q_i > 0$

•
$$\pi_i = q_i \max(16 - q_1 - q_2) - F_i$$
; if $q_i > 0$
• $F_i(q_i) = \begin{cases} 1 & \text{if } q_i > 0 \\ 0 & \text{if } q_i = 0 \end{cases}$

Solving using backward induction follower firm 2 solves its profit maximization problem first taking q_1 as given;

$$\max_{q_2>0} \ q_2(\max(16-q_1-q_2,0))-1$$

Solving which gives us,

$$BR_2(q_1) \in \begin{cases} \{\frac{16-q_1}{2}\} & \text{if } q_1 < 14\\ \mathbb{R}_+ & \text{if } q_1 \ge 14 \end{cases}$$

and,

$$q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 14\\ 0 & \text{if } q_1 \ge 14 \end{cases}$$

Then the leader firm 1 solves it's profit maximization problem subject to followers's best response;

$$\max_{q_1>0} q_1(\max(16-q_1-q_2,0))-1$$

$$\max_{q_1 \ge 0} \quad q_1(\max(16 - q_1 - q_2, 0)) - 1$$

$$s.t. \quad q_2(q_1) = \begin{cases} \frac{16 - q_1}{2} & \text{if } q_1 < 14\\ 0 & \text{if } q_1 \ge 14 \end{cases}$$

In this case the fixed costs are so low that leader firm decides to accomodate the entry and produces $q_1 = 8$ and the follower produces $q_2 = 8$.