### SEMESTER 2 EXAMINATION 2008/2009

#### DATA STRUCTURES AND ALGORITHMS

Duration: 120 mins

You must enter your Student ID and your ISS login ID (as a cross-check) on this page. You must not write your name anywhere on the paper.

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| Question | Marks |
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| 1        |       |
| 2        |       |
| 3        |       |
| 4        |       |
| Total    |       |

Answer THREE questions out of FOUR.

This examination is worth 85%. The tutorials were worth 15%.

University approved calculators MAY be used.

Each answer must be completely contained within the box under the corresponding question. No credit will be given for answers presented elsewhere.

You are advised to write using a soft pencil so that you may readily correct mistakes with an eraser.

You may use a blue book for scratch—it will be discarded without being looked at.

**Question 1** Below we show part of a binary search tree class for storing integers

```
public class BinarySearchTree {
   Node root;
   int number_of_elements;

   private static class Node {
      int element;
      Node left;
      Node right;
      Node parent;
   }
}
```

(a) Write a recursive method sum() which sums all the elements stored in the tree. Approximately correct code is sufficient. (5 marks)

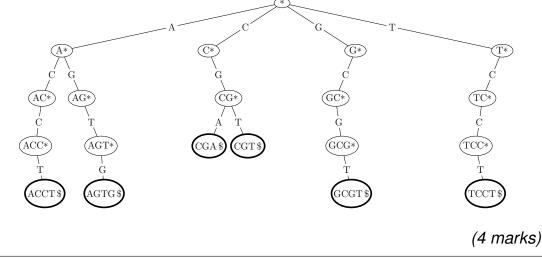
| (b) Show how you would make the BinarySearchTree into a generic class. (5 marks) |
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**TURN OVER** 

| (c) Write a method getNode(E obj) which returns the node with element objectists otherwise it returns null. You may assume E extends the Comparainterface.  (8 ma |  |
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(d) A digital tree or trie is a multi-way tree which can also be used to implement a set. An illustration of a trie that could be used for storing DNA sequences is shown below. This is a 4-way trie storing the sequences "ACCT\$", "AGTG\$", "CGA\$", "CGT\$", "GCGT\$" and "TCCT\$". Describe how the tree can be modified to be more efficient (turned into a Patricia trie).



(e) Give the retrieval time of an element from a TreeSet, a HashSet and a Trie. For full marks give a brief explanations of your answer. (6 marks)

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Pseudo code for a possible hashing function for DNA sequences is shown below

```
hash = 17;
foreach base b in a sequence
  hash = 31*hash + integer(b)
return hash
```

where integer(b) turns the bases, A, C, G, T, into integers 0, 1, 2, 3.

| (f) | Explain, why a HashS | et may not be | any faster than | the Trie at | inserting or |
|-----|----------------------|---------------|-----------------|-------------|--------------|
|     | retrieving data.     |               |                 |             | (5 marks)    |

End of question 1

Q1: (a) 
$$\frac{}{5}$$
 (b)  $\frac{}{5}$  (c)  $\frac{}{8}$  (d)  $\frac{}{4}$  (e)  $\frac{}{6}$  (f)  $\frac{}{5}$  Total  $\frac{}{33}$ 

## **Question 2** The Towers of Hanoi problem is solved by the program

```
hanoi(n, A, B, C)
{
   if (n>0) {
     hanoi(n-1, A, C, B);
     move(A, C);
     hanoi(n-1, B, A, C);
   }
}
```

(a) Let T(n) be the number of times move is called to solve the Hanoi problem of size n. Write down a recurrence relation for T(n). (4 marks)

```
T(n) =
```

(b) Write down the boundary condition T(1) and use the recurrence relation to compute T(2), T(3), and T(4) (4 marks)

| T(1) = |  |
|--------|--|
| T(2) = |  |
| T(3) = |  |
| T(4) = |  |

(c) Prove by induction that,  $f(n)=2^n-1$  satisfies the recurrence relation in part (a) (8 marks)

Here is an algorithm for performing sort

```
public static void Sort(double[] data) {
    for (int i = 1; i < data.length; i++) {
        if (data[i] < data[i-1]) {
            double temp = data[i];
            int j = i;
            do {
                data[j] = data[j-1];
                j--;
            } while (j>0 && data[j-1] > temp);
            data[j] = temp;
        }
    }
}
```

(d) Give upper and lower bounds on the time complexity of the algorithm? Explain your reasoning. (9 marks)

(e) Briefly explain how binary search works and explain its time complexity (8 marks)

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End of question 2

Q2: (a)  $\frac{}{4}$  (b)  $\frac{}{4}$  (c)  $\frac{}{8}$  (d)  $\frac{}{9}$  (e)  $\frac{}{8}$  Total  $\frac{}{33}$ 

# Question 3

| (a) Briefly describe the quick sort algorithm.   | (5 marks)                               |
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| (b) Explain why choosing a pivot for quick sort is crucial to the post of the algorithm and explain a common method for choosing a | performance<br>good pivot.<br>(6 marks) |
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**TURN OVER** 

| (c) What is the worst case and average case time complexity of quick a justification for your answer. | sort. Give<br>(6 marks) |
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| (d) Explain why no algorithm using binary comparisons can sort an than $\Theta(n  \log(n))$ operation on average. | array in less<br>(10 marks) |
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**TURN OVER** 

| (f) What is the time complexity of radix sort? Why does this not contradict part (c) (2 marks)  1 | (e) | Describe very briefly how radix sort works on a set of integers. | (4 marks)  |
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| part (c) (2 marks)  |     |  |            |
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| part (c) (2 marks)  | (6) | NAME   |            |
|   | (†) | What is the time complexity of radix sort? Why does this not     | contradict |
|   |     | part (c)   | (2 marks)  |
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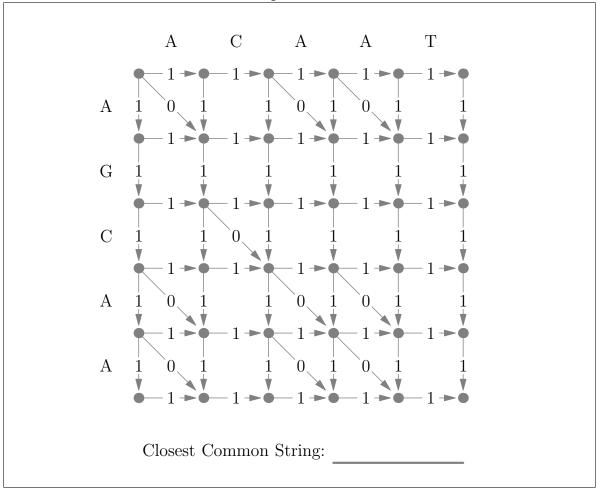
End of question 3

Q3: (a)  $\frac{}{5}$  (b)  $\frac{}{6}$  (c)  $\frac{}{6}$  (d)  $\frac{}{10}$  (e)  $\frac{}{4}$  (f)  $\frac{}{2}$  Total  $\frac{}{33}$ 

# Question 4

| (a) Explain the dynamic programming strategy. | (5 marks) |
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(b) The edit distance between two strings is the number of edit operations needed to go from one string to another. In the simplest case we can consider the only edit operation to be insertion. In this case, the edit distance between the string "AAC" and "ACT" is two as each string would have to have an insertion to make them identical (i.e. both strings could be turned into "AACT"). The edit distance can be computed by finding the lowest cost path on a edit graph starting from the top left hand corner and ending at the bottom right hand corner. The edit graph for two strings "ACAAT" and "AGCAA" is shown below. Using dynamic programming compute the costs of reaching each node (write the costs next to the nodes). Using the backwards algorithm to find the optimal path (show this on the graph). Write down the closest common string.



(16 marks)

(c) Briefly describe the following strategies, giving examples of their use. 1) Brute force, 2) Branch and bound, 3) Divide and conquer and 4) Greedy method. (12 marks)

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End of question 4

Q4: (a)  $\frac{}{5}$  (b)  $\frac{}{16}$  (c)  $\frac{}{12}$  Total  $\frac{}{33}$ 

**END OF PAPER**