

THEORY PROBLEMS FOR DATA STRUCTURES AND ALGORITHMS (COMP1009)

1 Consider the program (valid for inputs $n \geq 1$)

```
foo(int n) {
    bar();
    if (n==1)
        return;
    foo(n-1);
    foo(n-1);
    foo(n-1);
}
```

- (a) Let $C(n)$ be the number of times the function `bar()` is called when we call `foo(n)`. Write down a recurrence relation for $C(n)$. (2 marks)

$$C(n) =$$

- (b) Write down the boundary condition for the recurrence relation. (1 marks)

$$C(1) =$$

- (c) Using the recurrence relation to compute $C(2)$, $C(3)$ and $C(4)$. (3 marks)

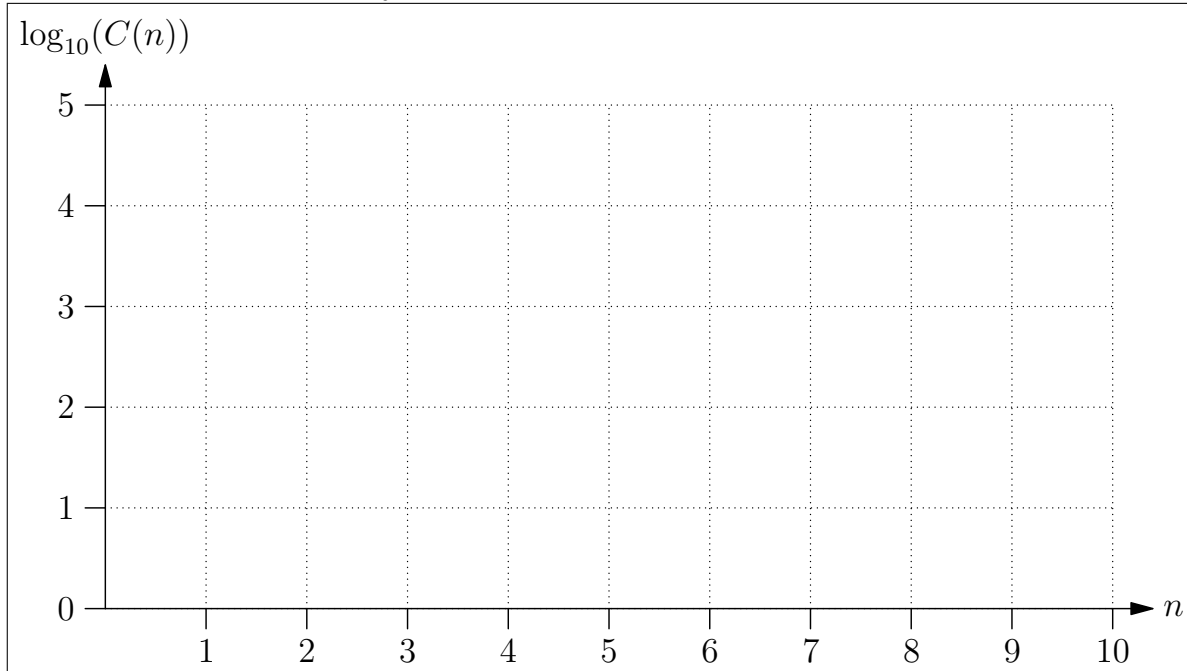
$$C(2) =$$
$$C(3) =$$
$$C(4) =$$

- (d) Prove by induction that $f(n) = \frac{3^n - 1}{2}$ satisfies the recurrence relation for $C(n)$.
(5 marks)

[illegible]

(e) Sketch the curve $\log_{10}(C(n))$ on the graph below.

(2 marks)



(f) Assume that the most time consuming operation is calling function `bar()` then, if it takes 100s to compute `foo(5)` approximately how long will it take to compute `foo(10)`? (2 marks)

2 Consider the program (valid for inputs $n \geq 1$)

```
foo(int n) {
    for(i=1; i<=2n-1; i++)
        bar();
    if (n==1)
        return;
    foo(n-1);
}
```

- (a) Let $C(n)$ be the number of times the function `bar()` is called when we call `foo(n)`. Write down a recurrence relation for $C(n)$. (2 marks)

$C(n) =$

- (b) Write down the boundary condition for the recurrence relation. (1 marks)

$C(1) =$

- (c) Using the recurrence relation to compute $C(2)$, $C(3)$ and $C(4)$. (3 marks)

$C(2) =$

$C(3) =$

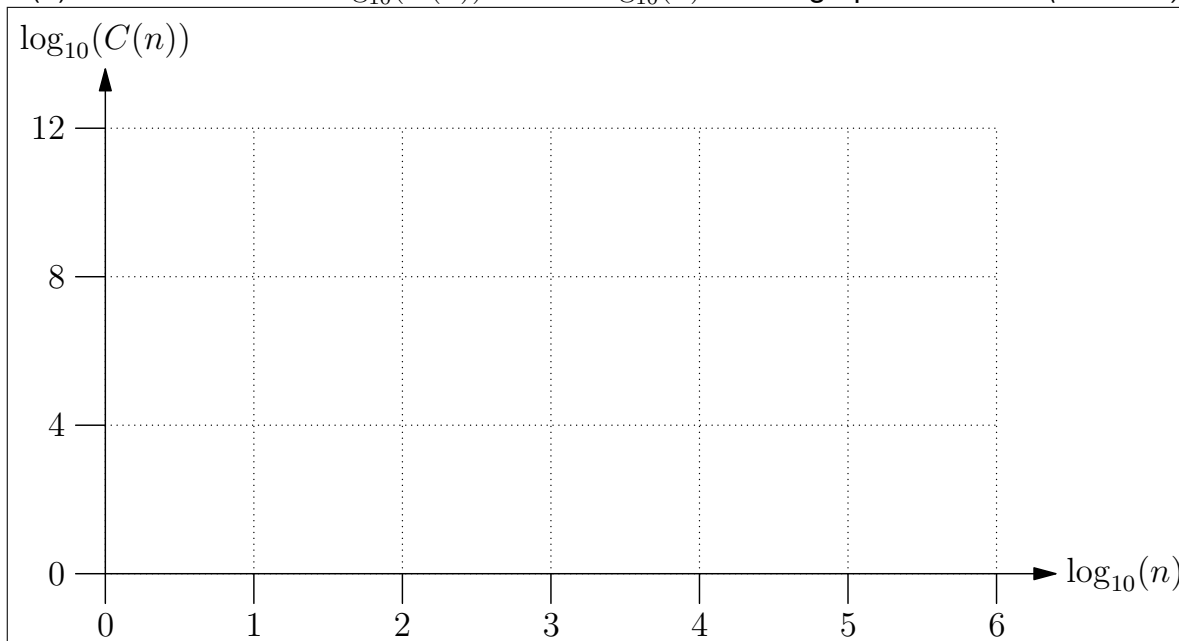
$C(4) =$

- (d) Guess the solution for $C(n)$ and prove by induction that it satisfies the recurrence relation for $C(n)$. (5 marks)

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TURN OVER

(e) Sketch the curve $\log_{10}(C(n))$ versus $\log_{10}(n)$ on the graph below. (2 marks)



(f) If it takes 100s to compute $f_{oo}(1000)$ approximately how long will it take to compute $f_{oo}(2000)$? (2 marks)

3 Consider the program (valid for inputs $n \geq 1$)

```
foo(int n) {
    bar();
    if (n==1)
        return;
    int m = (int) n/2
    foo(m);
}
```

where $(\text{int})\ n/2$ returns the greatest integer less than or equal to $n/2$ (i.e. $\lfloor n/2 \rfloor$).

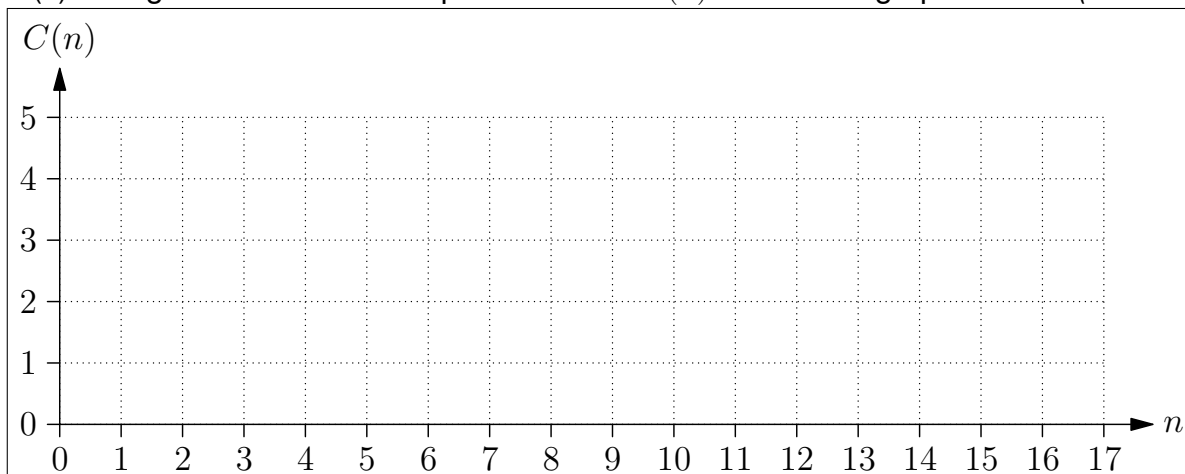
- (a) Let $C(n)$ be the number of times the function `bar()` is called when we call `foo(n)`. Write down a recurrence relation for $C(n)$. (2 marks)

$$C(n) =$$

- (b) Write down the boundary condition for the recurrence relation. (1 marks)

$$C(1) =$$

- (c) Using the recurrence compute values of $C(n)$ to draw the graph below. (4 marks)



TURN OVER

- (d) Prove by induction that $f(n) = \lfloor \log_2(n) \rfloor + 1$ satisfies the recurrence relation for $C(n)$. This is simplified if we perform the inductive step over the set of integers $S_m = \{2^m, 2^m + 1, \dots, 2^{m+1} - 1\}$. (6 marks)

- (e) If it takes 100s to compute $f_{oo}(512)$ approximately how long will it take to compute $f_{oo}(1024)$? (2 marks)

END OF PAPER