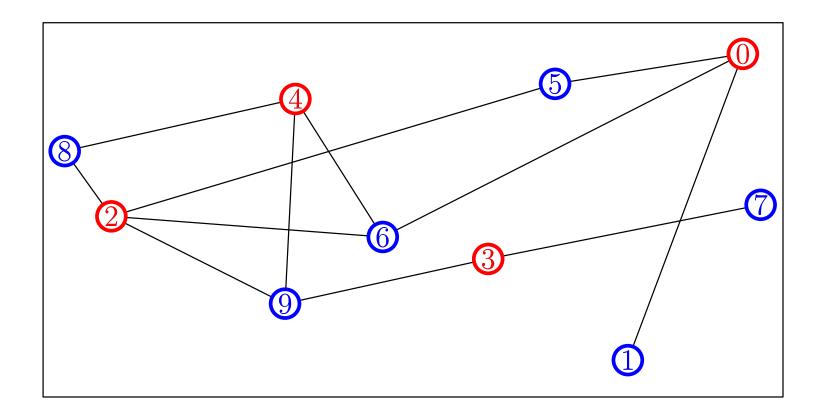
### **Algorithms and Analysis**

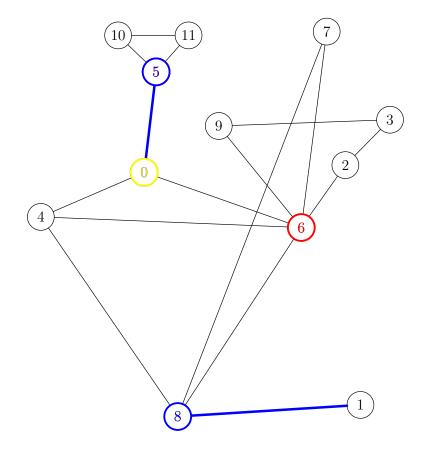
# Lesson 20: Learn to Traverse Graphs



Breadth First Search, Depth First Search, Topological Sort

#### **Outline**

- BFS applications
- 2. Depth First Search
  - DFS applications
- 3. Topological Sort



- Graphs provide an abstraction for a huge number of real world processes: social networks, compute network, road networks, etc.
- Increasing applications focus on very large (sparse) graphs (usually implemented using an adjacency list)
- Require (near) linear time algorithms
- Basic building block are graph traversal algorithms
  - \* Breadth First Search
  - ⋆ Depth First Search

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- This introduces a new level of generics which makes the algorithms very powerful
- Increases the steepness of the learning curve to use these algorithms
- Once you get familiar with using these algorithms this level of generics starts to pay off
- Libraries which does this include Boost Graph Library and LEDA in C++; JDSL and JGraphT in Java

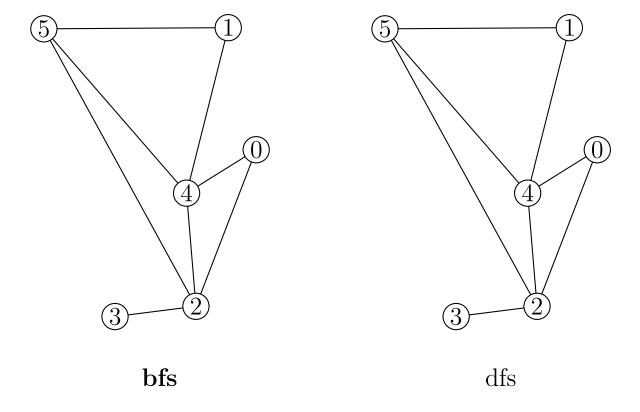
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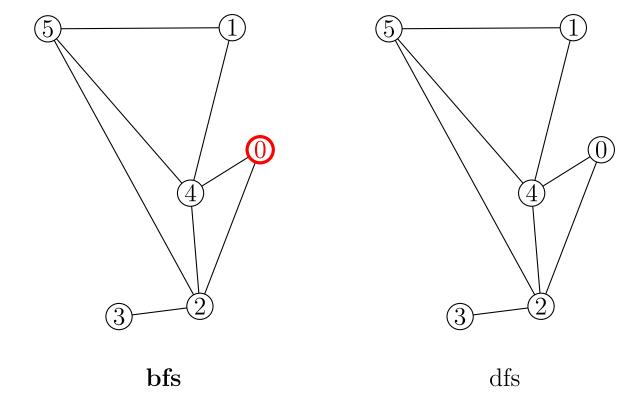
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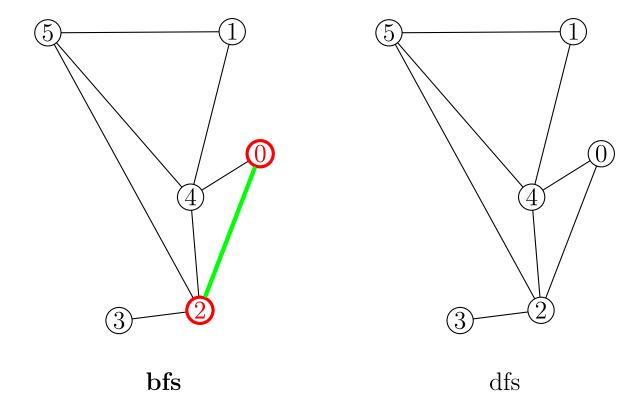
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- We the follow edges to create a tree



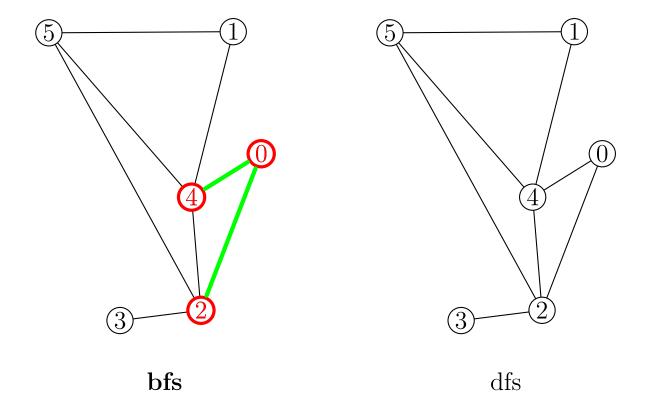
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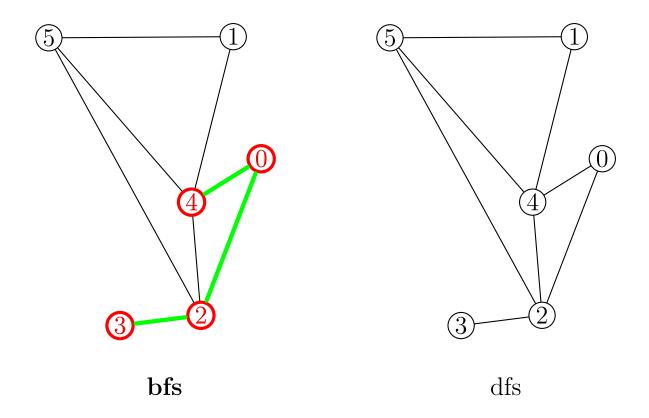
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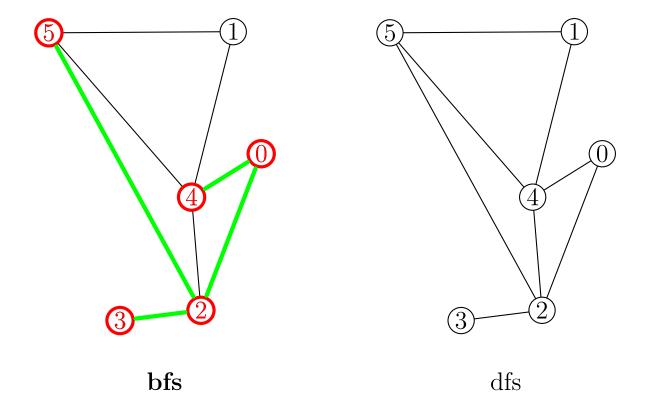
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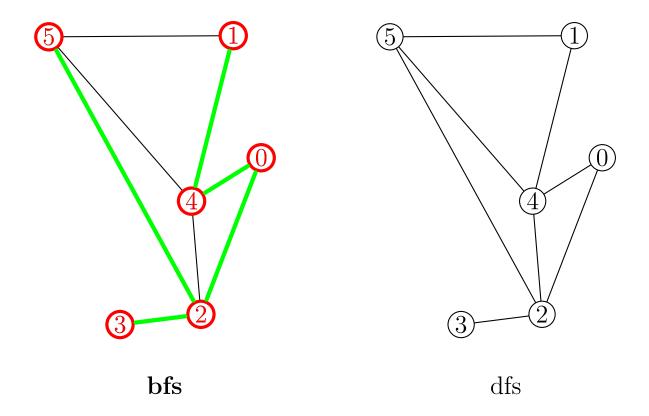
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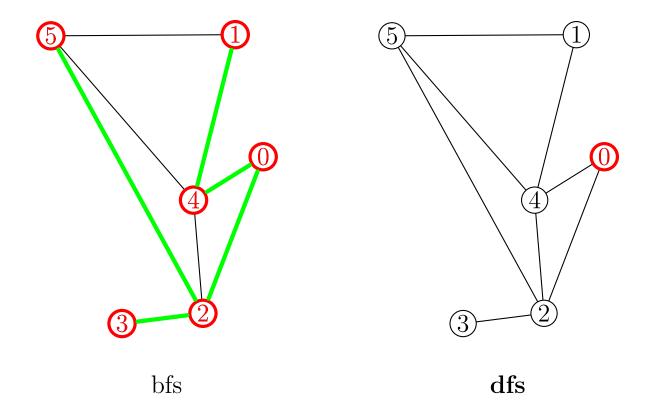
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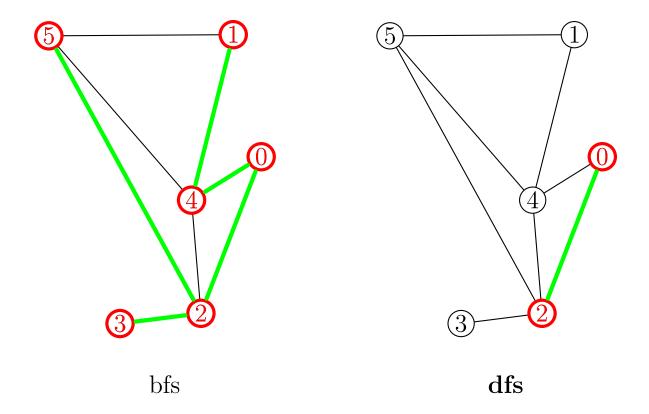
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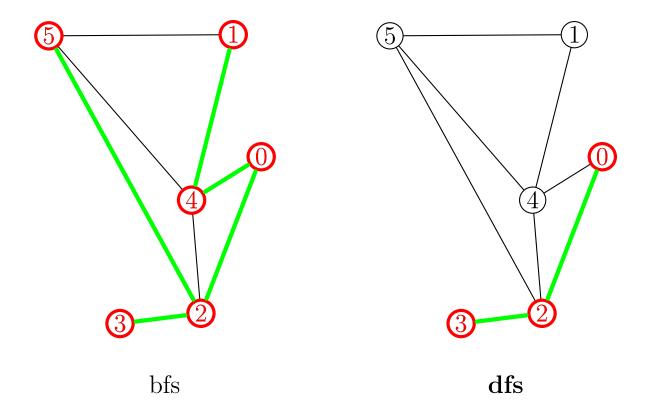
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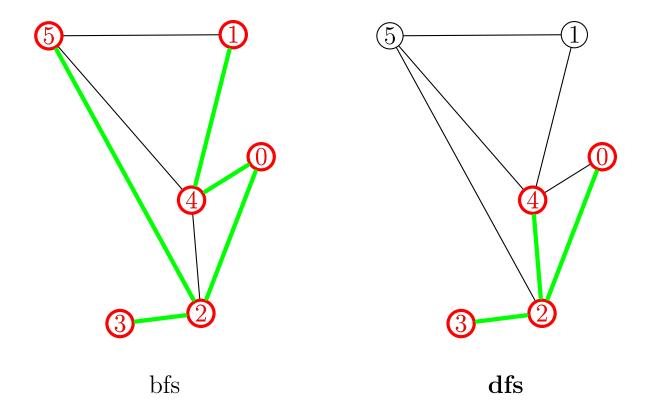
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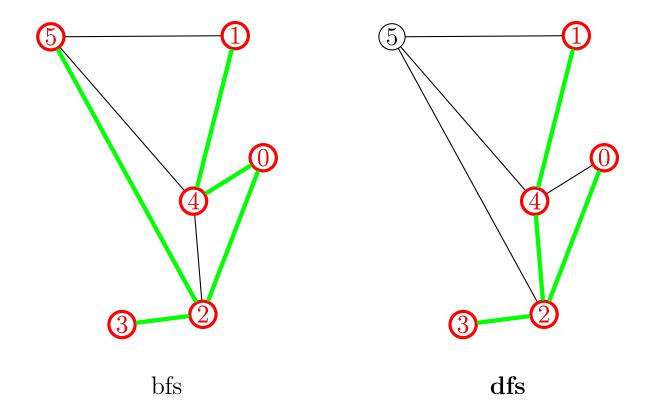
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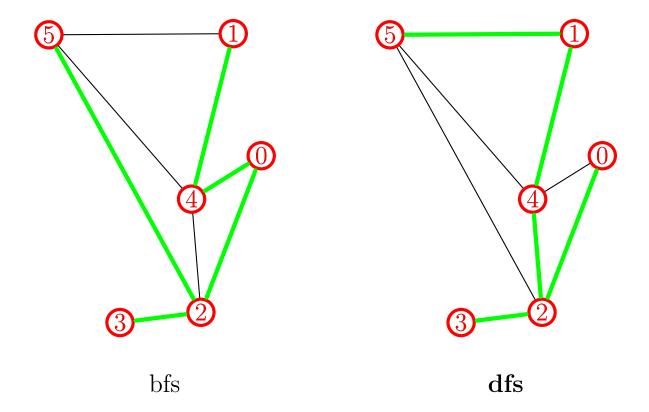
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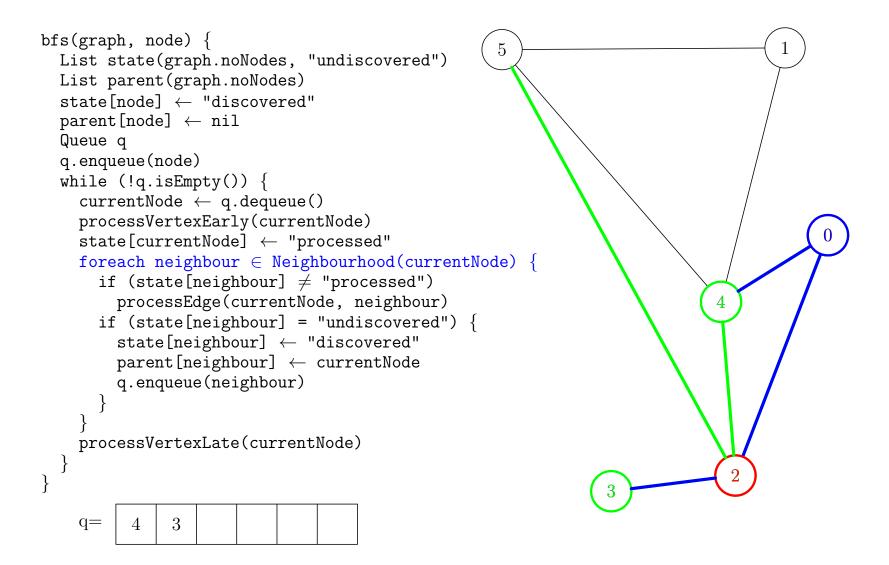
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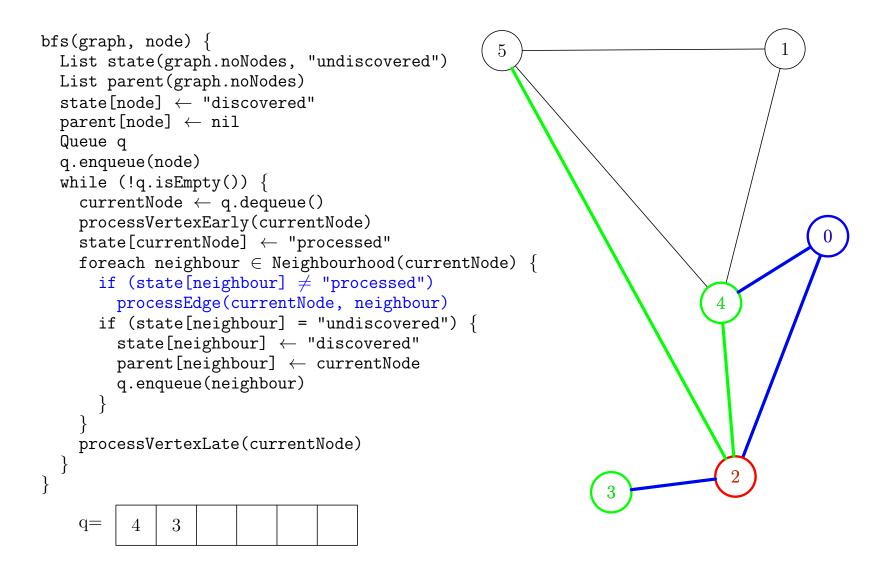
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    foreach neighbour ∈ Neighbourhood(currentNode) {
      if (state[neighbour] \neq "processed")
        processEdge(currentNode, neighbour)
      if (state[neighbour] = "undiscovered") {
        state[neighbour] = "discovered"
        parent[neighbour] = currentNode
        q.enqueue(neighbour)
   processVertexLate(currentNode)
    q=
              3
```

```
bfs(graph, node) {
 List state(graph.noNodes, "undiscovered")
 List parent(graph.noNodes)
 state[node] ← "discovered"
 parent[node] \leftarrow nil
 Queue q
 q.enqueue(node)
 while (!q.isEmpty()) {
    currentNode = q.dequeue()
    processVertexEarly(currentNode)
    state[currentNode] = "processed"
    foreach neighbour ∈ Neighbourhood(currentNode) {
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        q.enqueue(neighbour)
   processVertexLate(currentNode)
    q=
              3
```

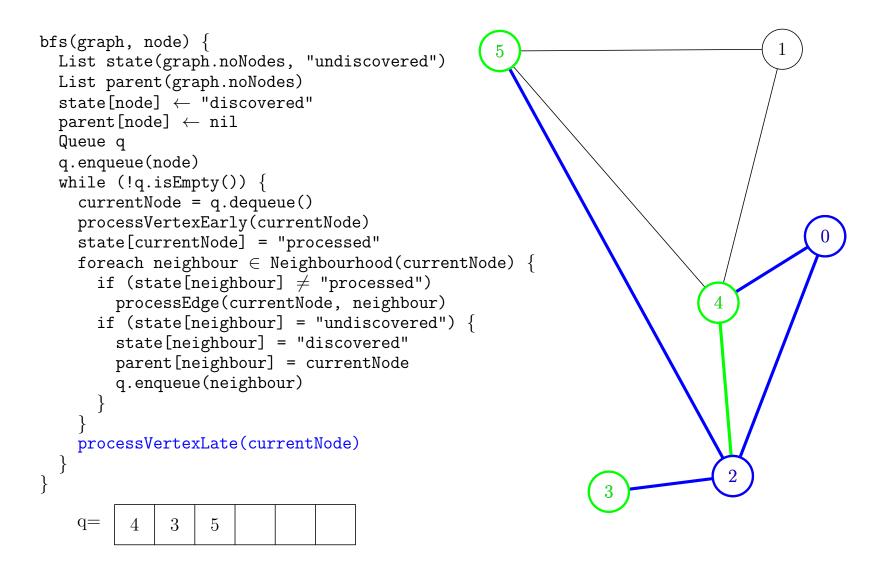


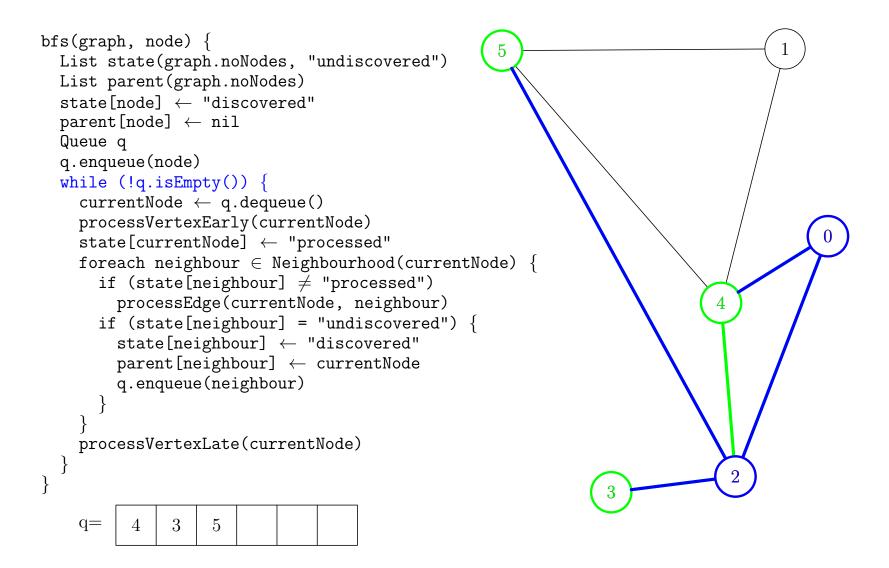


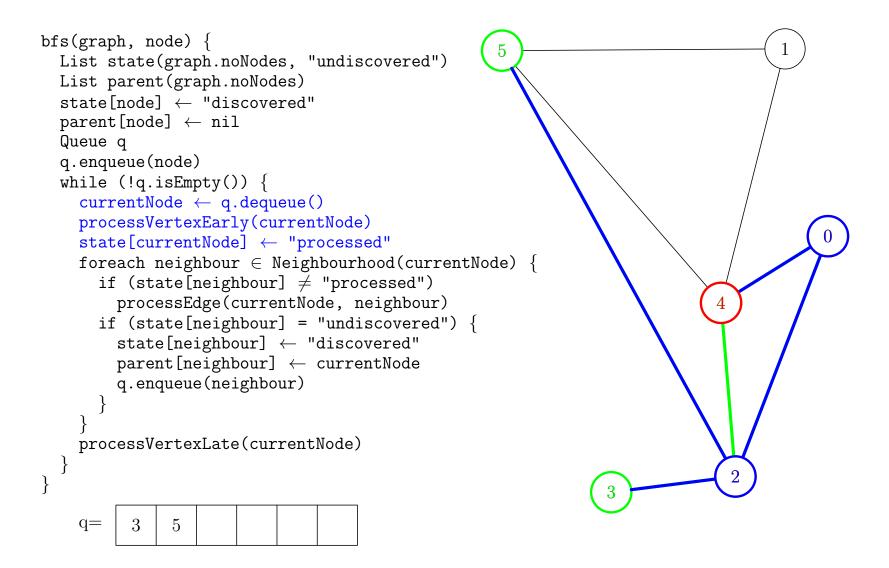
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      if (state[neighbour] = "undiscovered") {
        state[neighbour] 

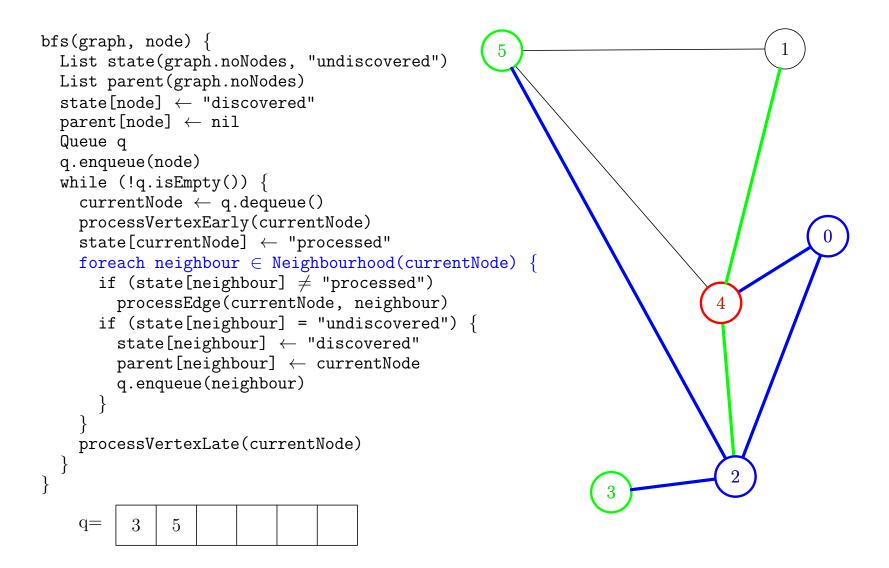
- "discovered"
        parent[neighbour] ← currentNode
        q.enqueue(neighbour)
    processVertexLate(currentNode)
    q =
              3
```

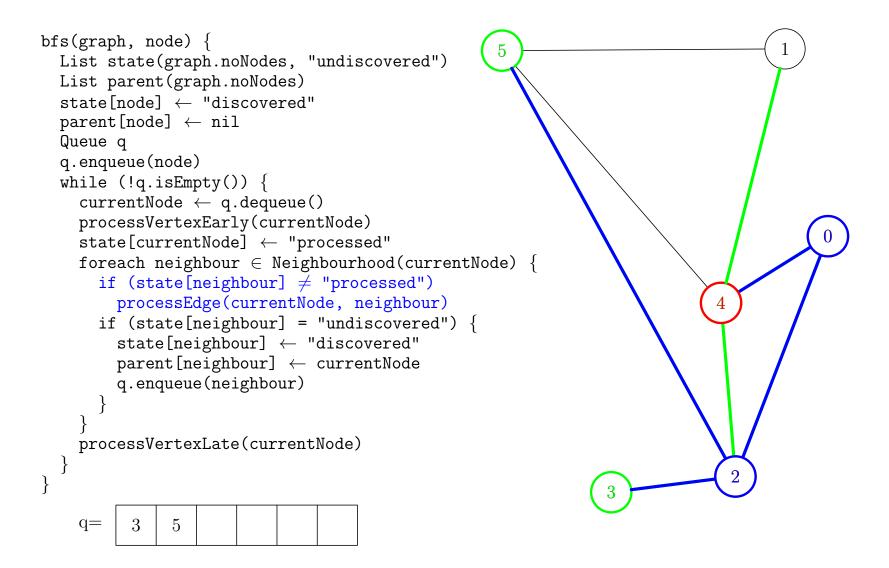
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       processEdge(currentNode, neighbour)
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       parent[neighbour] ← currentNode
       q.enqueue(neighbour)
   processVertexLate(currentNode)
   q =
             3
                 5
```

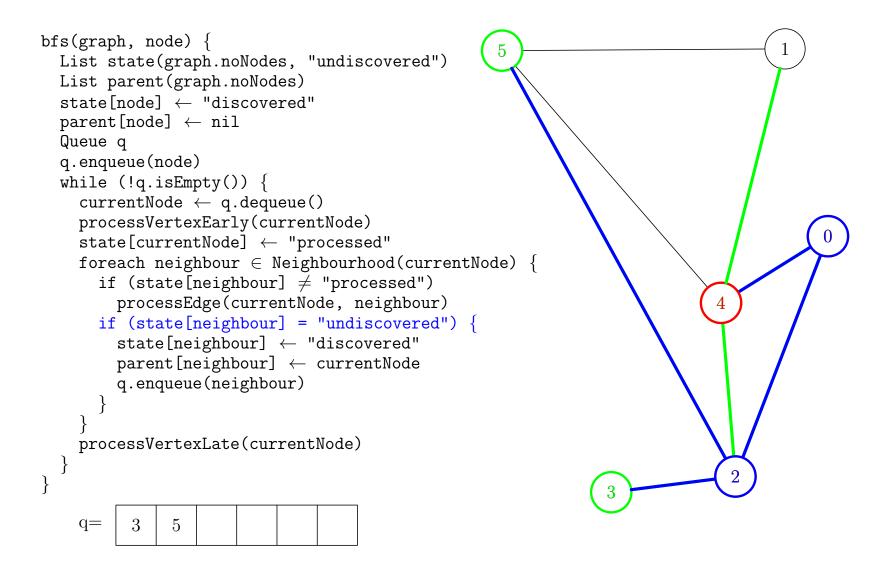


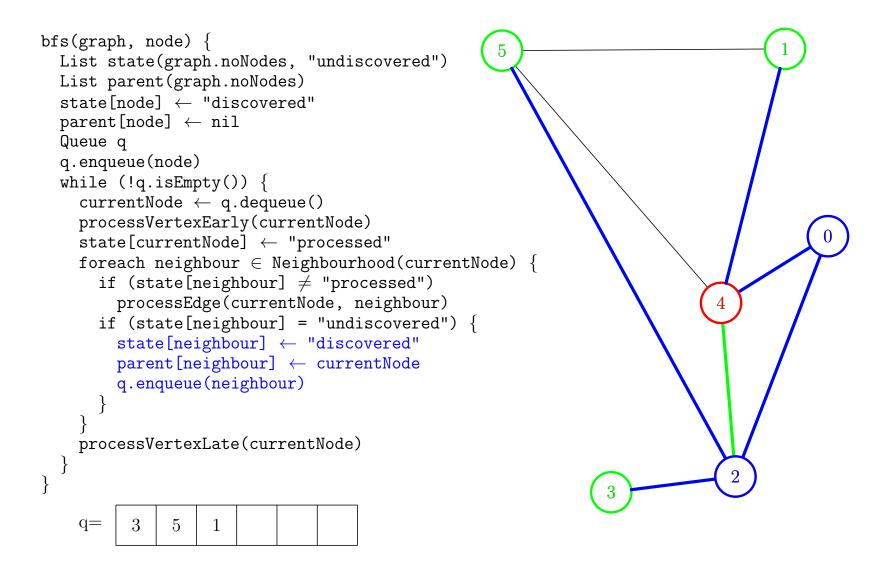


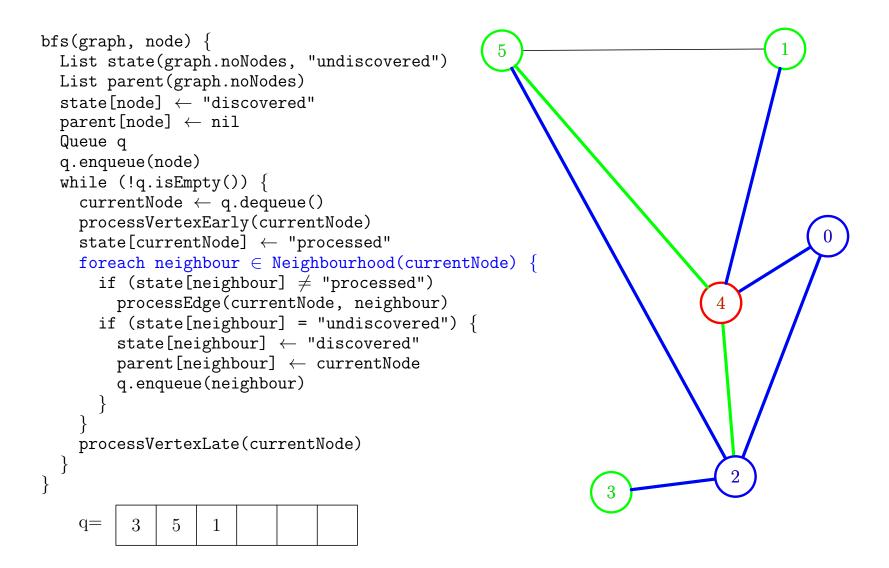


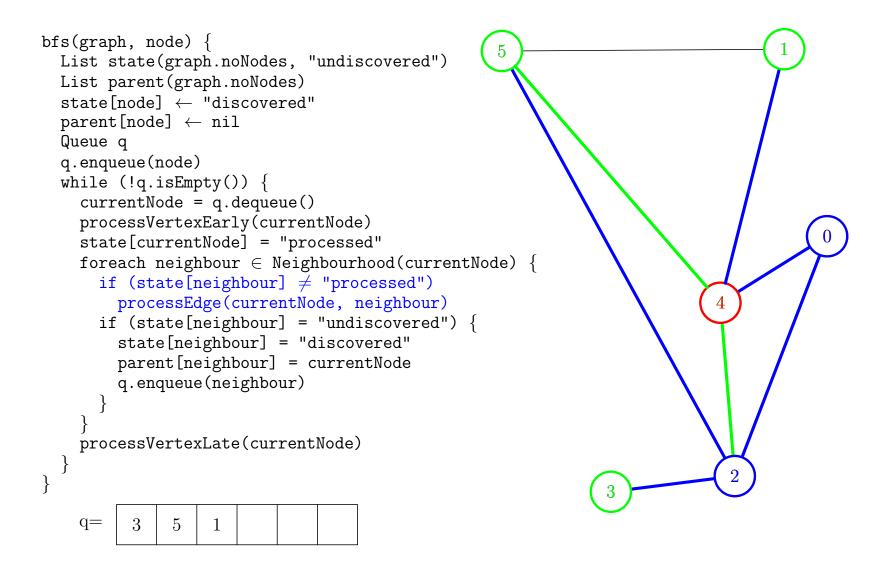


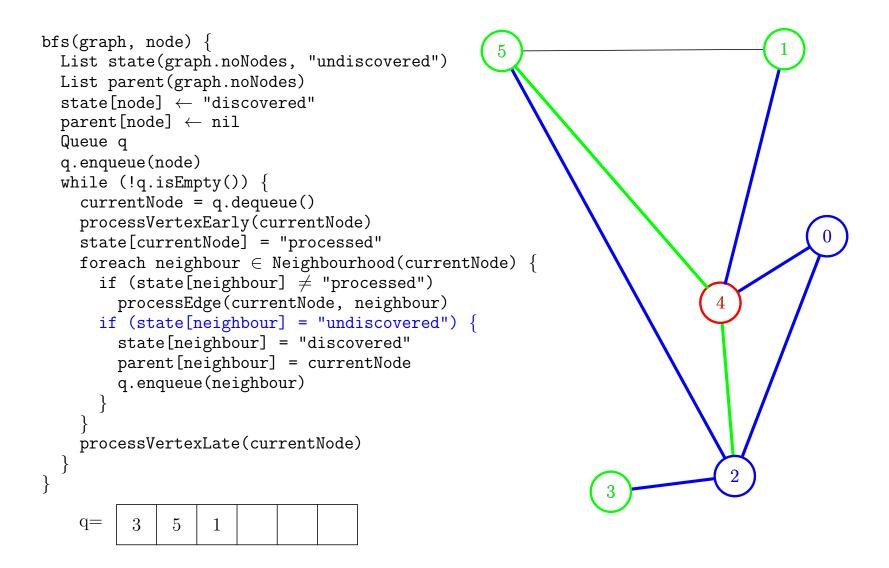


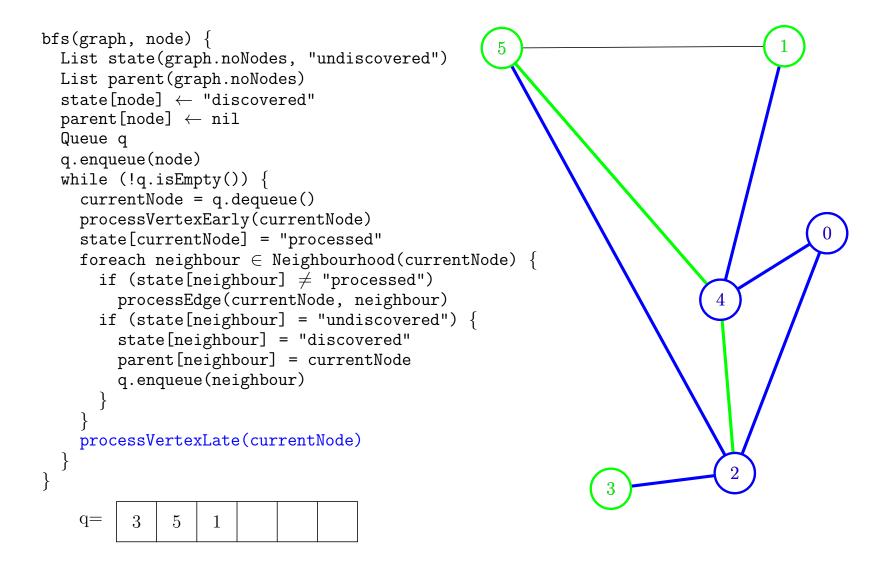


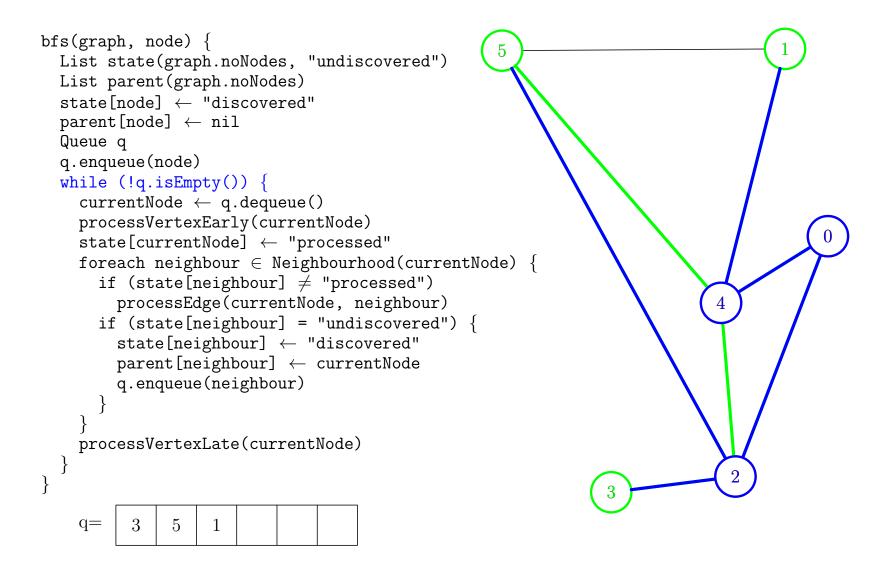


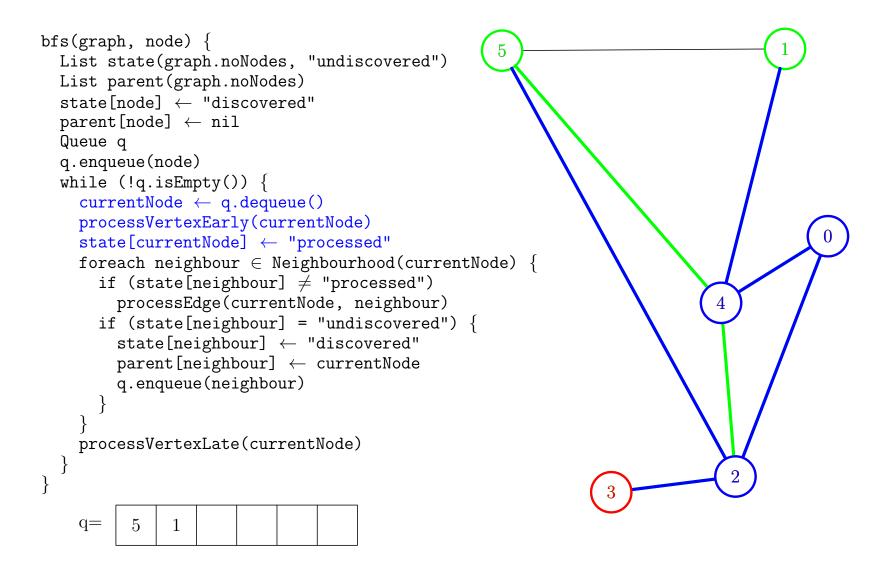


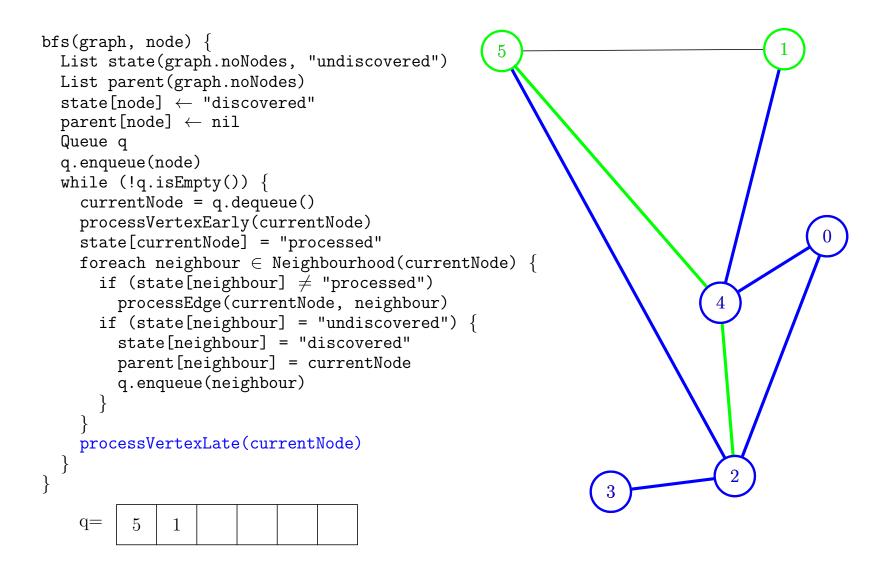


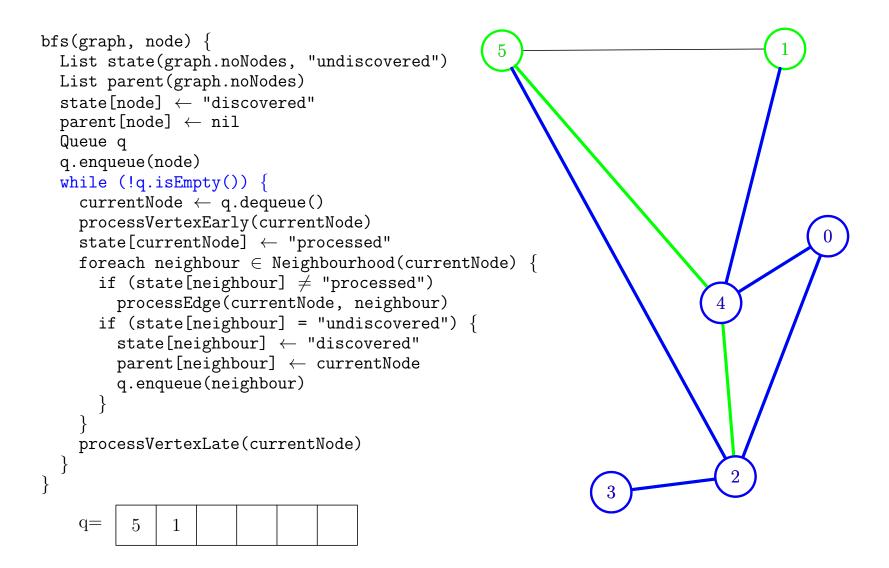


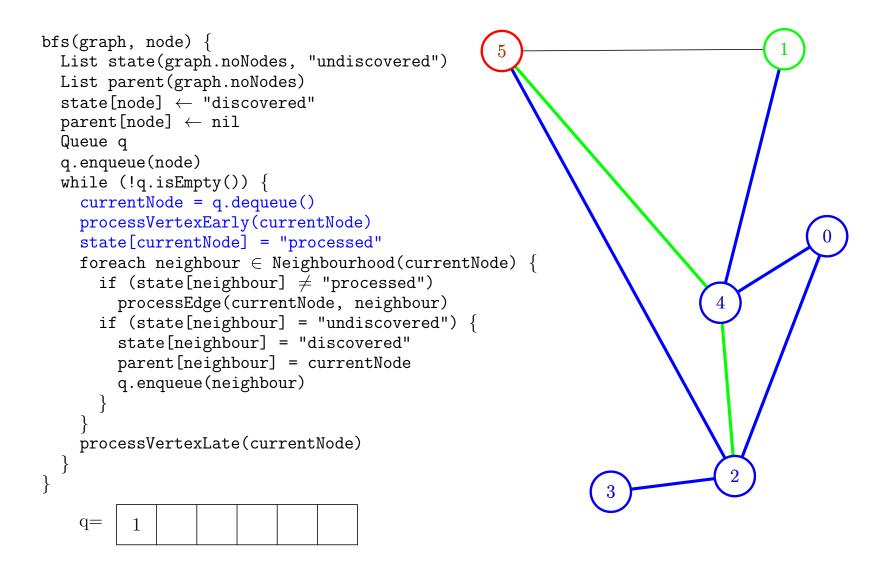


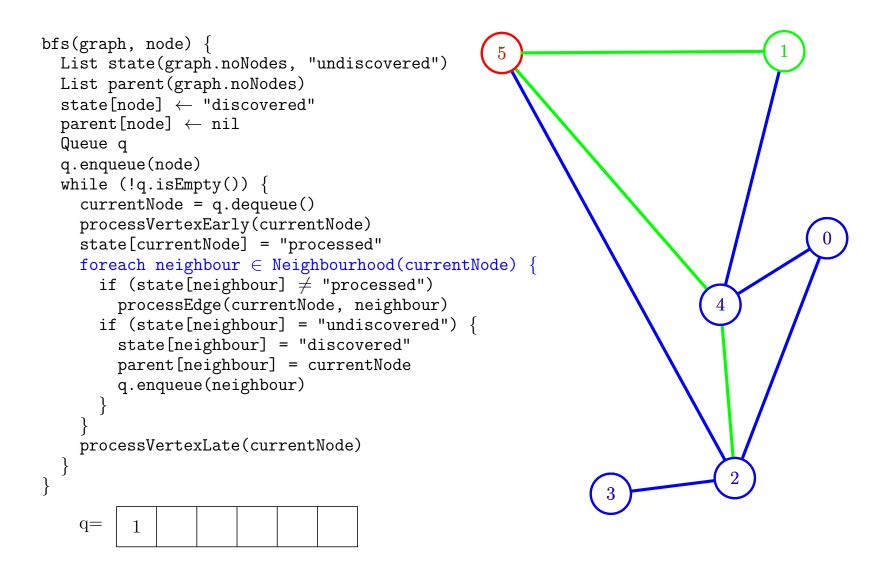


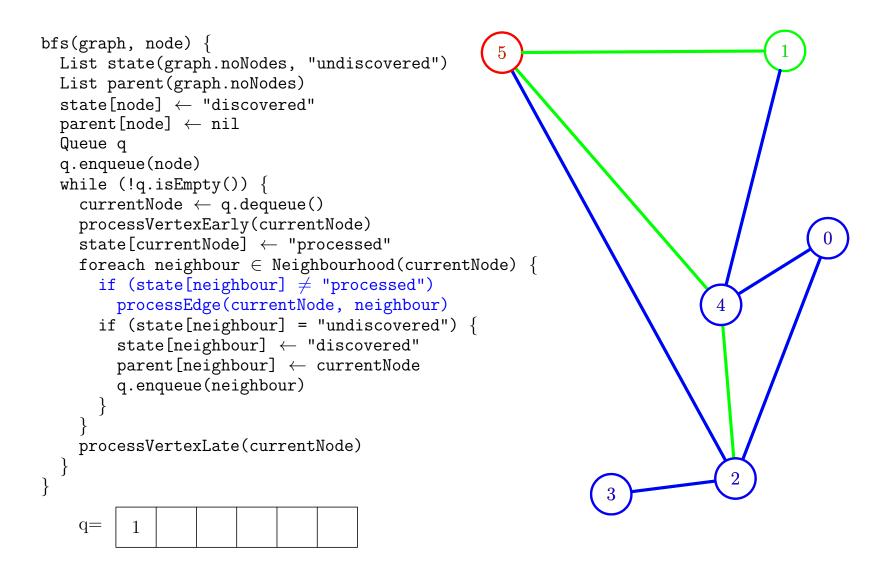


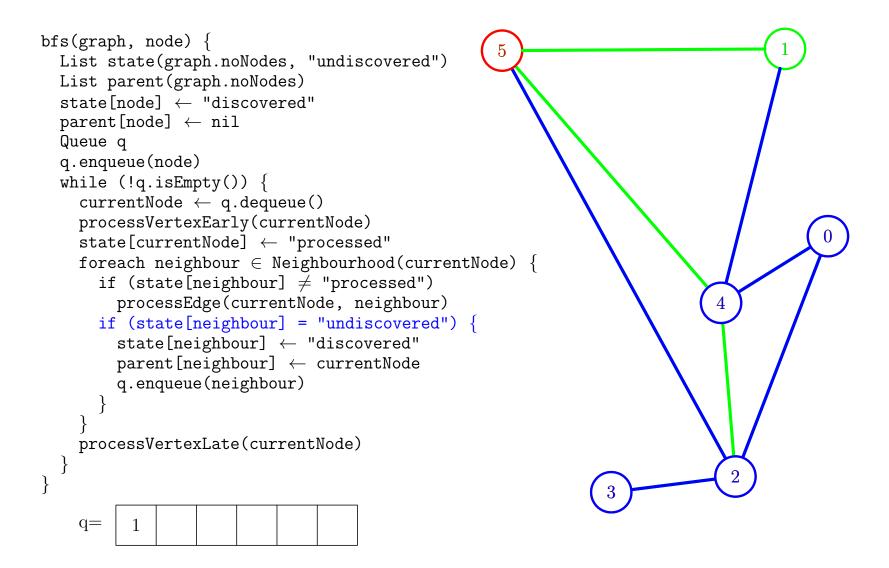


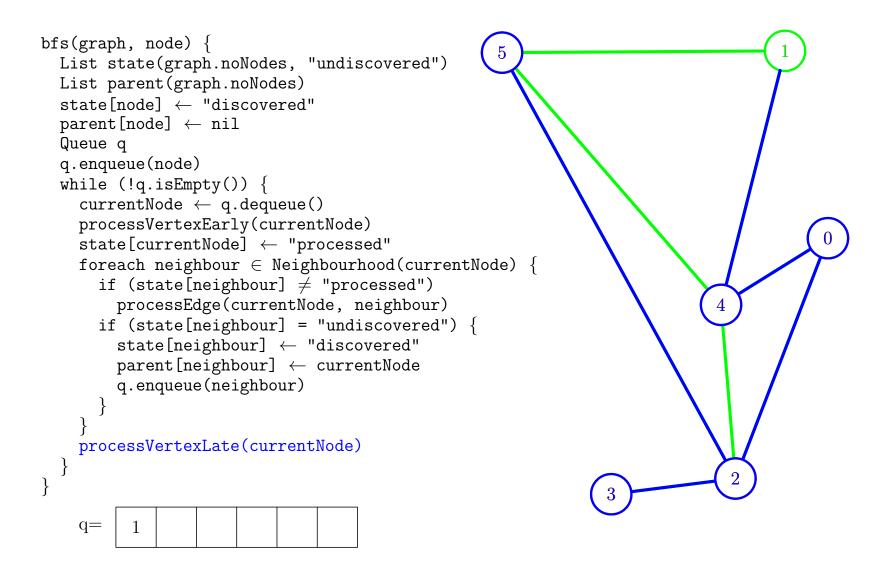


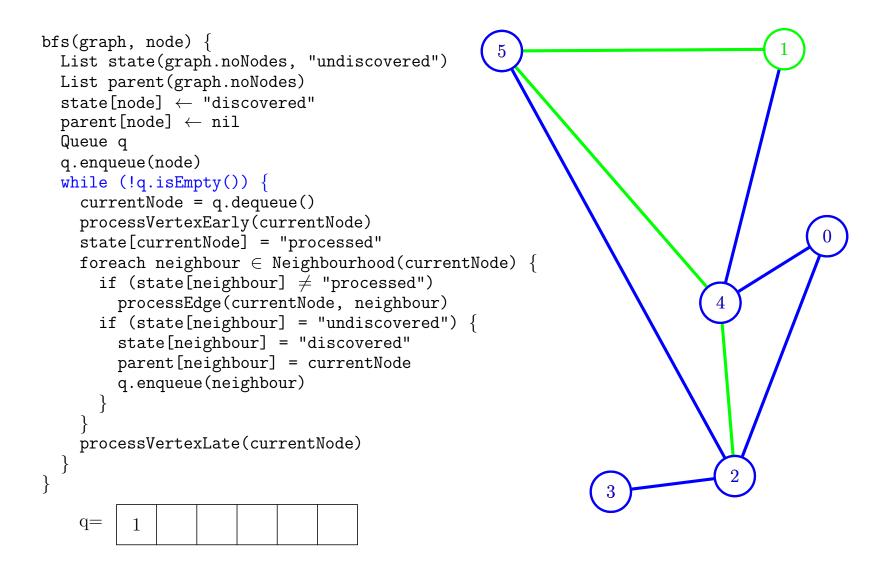








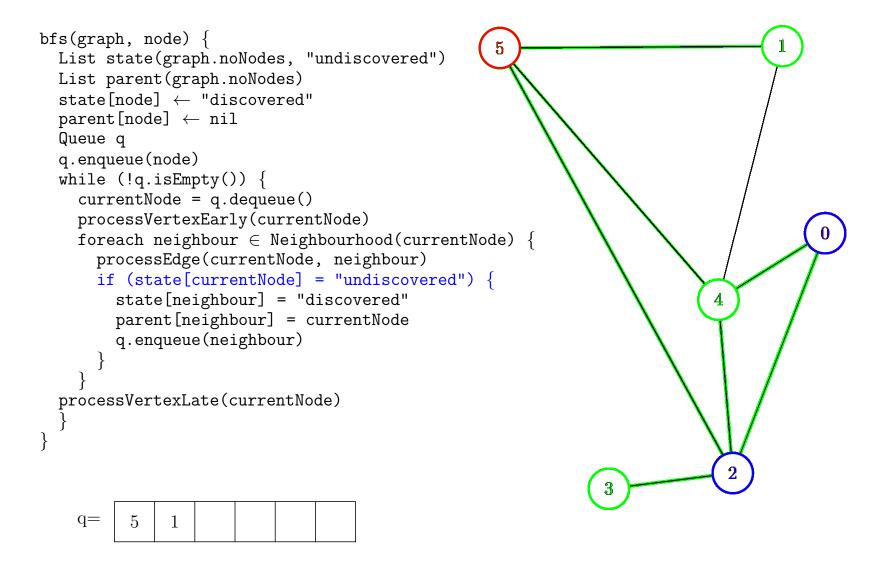




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bfs(graph, node) {
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 Queue q
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### **Applications of Breadth First Search**

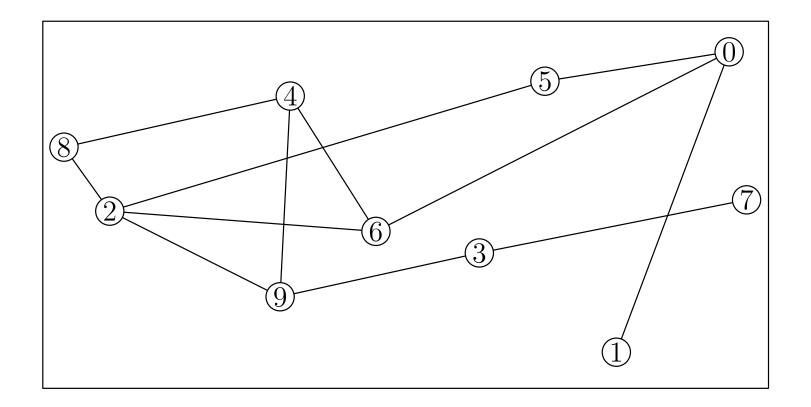
- Breadth first search can be used to find the shortest path from a source node to a destination node for an unweighted graph
  - \* Run bfs (graph, source)
  - Use parent information to find path from destination back to source
- BFS (as well as DFS) can be used to find connected components
  - ★ Use process VertexEarly to mark vertices connected to the current connected component
  - $\star$  Run bfs from all vertices that are not labelled

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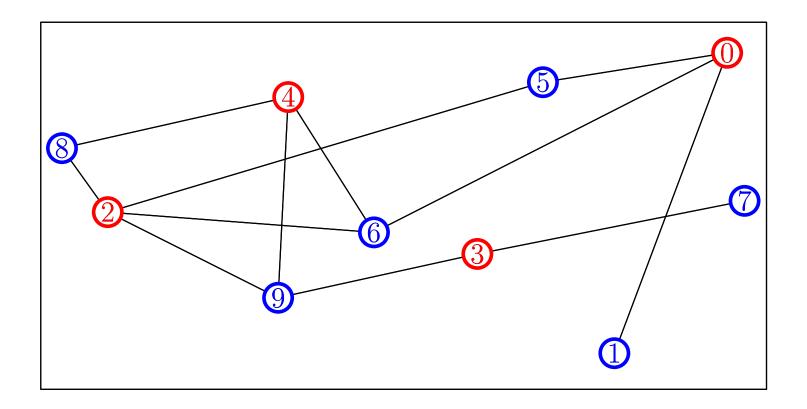
### **Bipartite Graphs**

 Bipartite graphs are graphs where the vertices can be split into two sets so that there are no edges between vertices in the same graph



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```
isBipartite(graph) {
   colour = List(graph.noNodes(), "uncoloured");
   bipartite = true;
   foreach node in graph {
      if (colour[node] == "uncoloured") {
         colour[node] = white;
         bfs(graph, node);
   return bipartite;
processEdge(node1, node2) {
   if (colour[node1] == colour[node2])
       bipartite = false;
   colour[node2] = (colour[node1] == "white")? "black": "white";
```

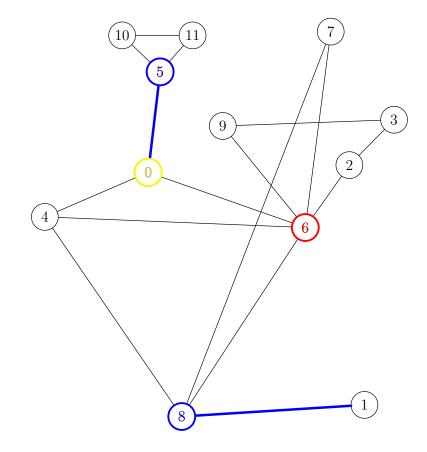
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#### **Outline**

- 1. Breadth First Search
  - BFS applications
- 2. Depth First Search
  - DFS applications
- 3. Topological Sort



- Depth first search is essentially like breadth first search except we replace the queue by a stack
- In practice it is often implemented using recursion rather than a stack
- It proves useful to keep a record of the traversal time for each vertex
  - \* the clock ticks each time a vertex is entered or exited

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- In practice it is often implemented using recursion rather than a stack
- It proves useful to keep a record of the traversal time for each vertex
  - \* the clock ticks each time a vertex is entered or exited

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=0
    node=0
```

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dfs(graph, node) {
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    } else if (state[neighbour] \neq "processed") {
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    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=0
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```

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dfs(graph, node) {
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  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=3
    node=2
                               time=2
```

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     processEdge(node, neighbour)
    if ("finished") return
 processVertexLate(currentNode)
  state[currentNode] ← "processed"
 time \leftarrow time + 1
              neighbour=nil time=3
    node=3
```

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    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=4
    node=3
```

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  time \leftarrow time + 1
 processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
     parent[neighbour] ← node
     processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
     processEdge(node, neighbour)
    if ("finished") return
 processVertexLate(currentNode)
  state[currentNode] ← "processed"
 time \leftarrow time + 1
              neighbour=nil time=5
    node=4
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
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  state[currentNode] ← "processed"
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  state[currentNode] ← "processed"
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    node=4
```

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      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=1
    node=4
                                time=5
```

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      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=1
    node=4
                                time=5
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      dfs(graph, neighbour)
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      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=1
    node=4
                                time=5
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      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=1
    node=4
                                time=5
```

```
dfs(graph, node) {
  if ("finished") return
  time \leftarrow time + 1
 processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
     parent[neighbour] ← node
     processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
     processEdge(node, neighbour)
    if ("finished") return
 processVertexLate(currentNode)
  state[currentNode] ← "processed"
 time \leftarrow time + 1
              neighbour=nil time=6
    node=1
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=1
                                time=6
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=1
                                time=6
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
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      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=1
                                time=6
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
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      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=1
                                time=6
```

```
dfs(graph, node) {
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     parent[neighbour] ← node
     processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
     processEdge(node, neighbour)
    if ("finished") return
 processVertexLate(currentNode)
  state[currentNode] ← "processed"
 time \leftarrow time + 1
              neighbour=nil time=7
    node=5
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=5
                                time=7
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
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  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=5
                                time=7
```

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dfs(graph, node) {
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    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=5
                                time=7
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  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=5
                                time=7
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    if ("finished") return
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  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=2
    node=5
                                time=7
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dfs(graph, node) {
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      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=4
    node=5
                                time=7
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
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    if ("finished") return
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  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=4
    node=5
                                time=7
```

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  state[currentNode] ← "processed"
  time \leftarrow time + 1
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    node=5
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               neighbour=4
    node=5
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  state[currentNode] ← "processed"
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    node=5
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      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=8
    node=5
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=1
                                time=8
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
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      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=9
    node=1
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=1
    node=4
                                time=9
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
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      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=4
                                time=9
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
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    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=4
                                time=9
```

```
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  time \leftarrow time + 1
               neighbour=5
    node=4
                                time=9
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      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=10
    node=4
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=4
    node=2
                                time=10
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=2
                                time=10
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
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    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=2
                                time=10
```

```
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  \texttt{time} \leftarrow \texttt{time} + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] = "undiscovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    } else if (state[neighbour] \neq "processed") {
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=5
    node=2
                                time=10
```

```
dfs(graph, node) {
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  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=nil time=11
    node=2
```

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      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
               neighbour=4
                                time=11
    node=0
```

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dfs(graph, node) {
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               neighbour=nil time=12
    node=0
```

## **Applications of DFS**

- Depth first search has many applications
- Suppose we want to check if the graph is a tree (i.e. has no cycles)
- The only edges that are allowed are parent edges

```
processEdges(node1, node2) {
   if (parent[node1] ≠ node2) {
      isTree ←false
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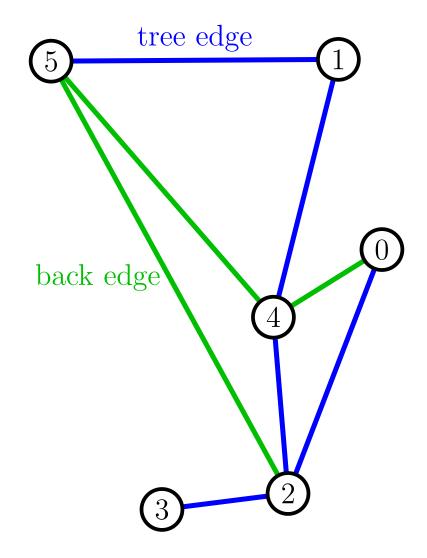
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- A brute force method for identifying articulation vertices is to remove each and check for connectivity—this would take O(n(m+n)) time. Can we do this any faster?

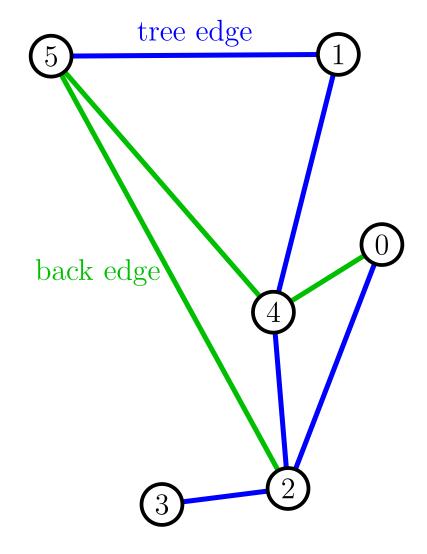
# Single Pass Algorithm

- In DFS we divide the edges into tree edges that define the search tree (edges between nodes and parents) and back edges which take us back to vertices we have already seen
- Without the back edges we have a tree where all non-leaf nodes are articulation nodes
- The back edges secure the edges to the rest of the tree



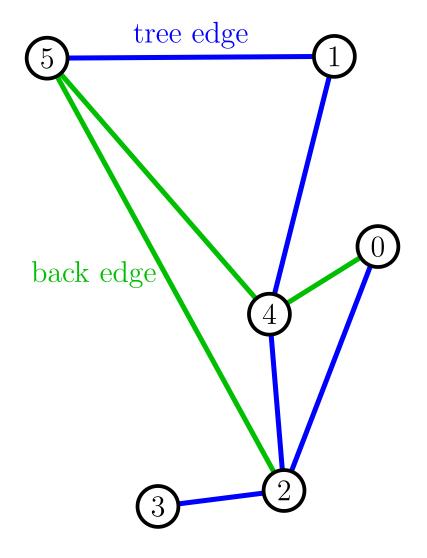
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- We maintain an array noting the reachable ancestors of all nodes
- This is initialised in the processVertexEarly method to the node itself
- In the processEdge method, if the edge is a
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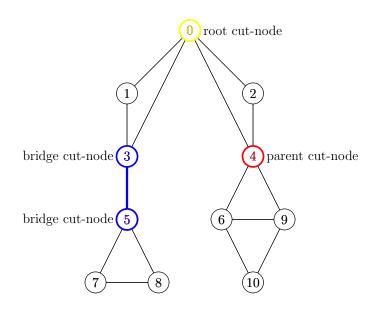
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 Key is to recognise that articulated vertices only occur in three version

**Root cut-nodes** Occur when the root has more than one child

Bridge cut-nodes Occurs when the earliest reachable vertex (not including the tree edge to the parent) is the vertex itself. The parent will be an articulation node as will be the node itself if it is not a leaf node

Parent cut-nodes If the earliest reachable vertex is it parent then the parent is an articulation node

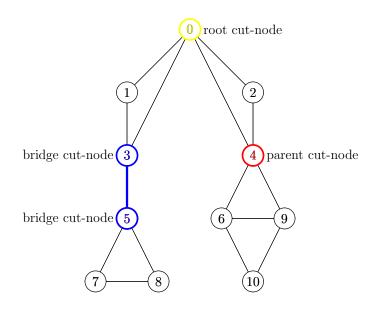


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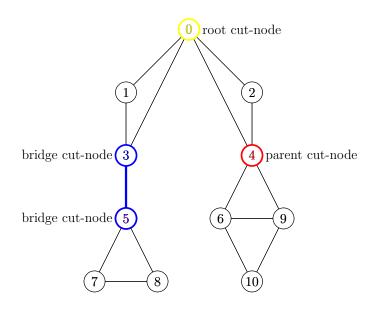


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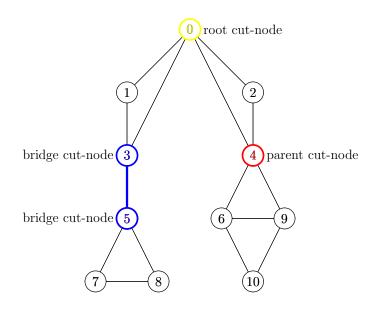


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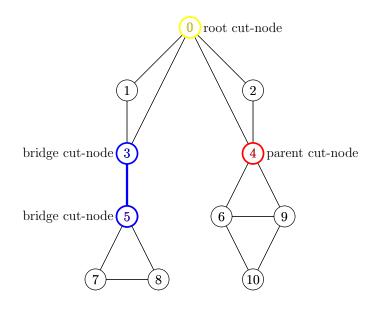


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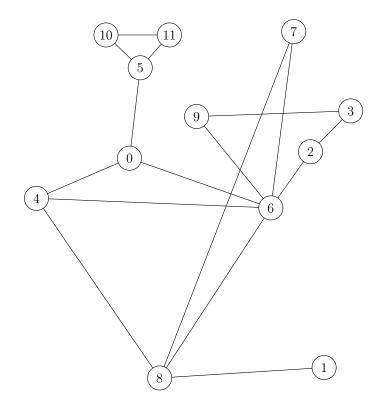
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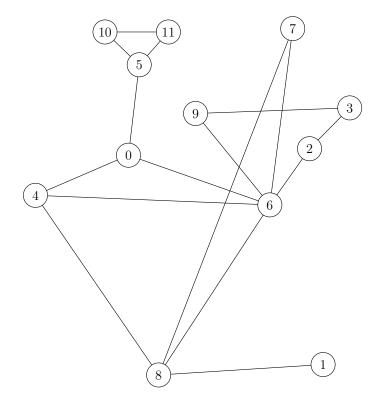
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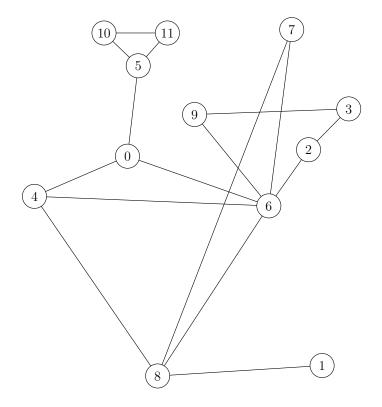
- Algorithmic details are not too important
- One pass (O(n+m)) algorithm
- Uses processVertexEarly, procesEdge and processVerexLate methods
- Bridge cuts also shows articulation edges



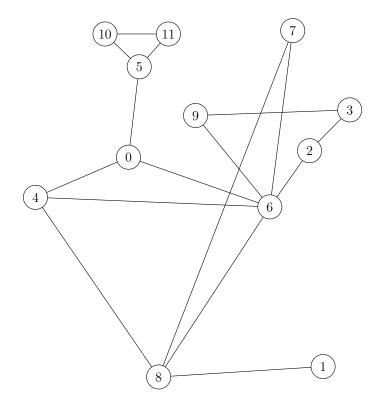
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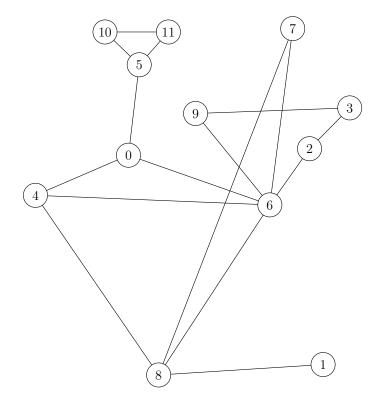


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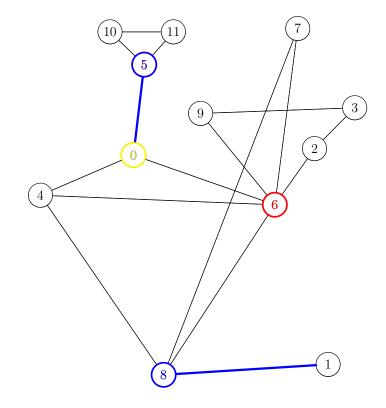
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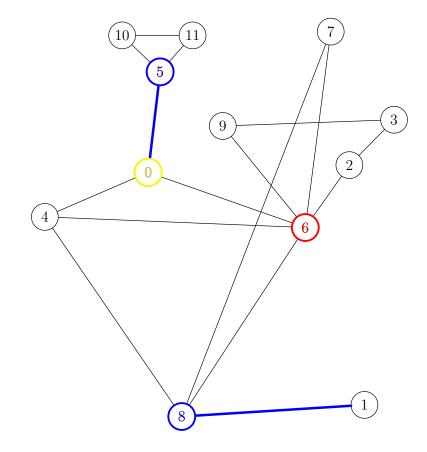
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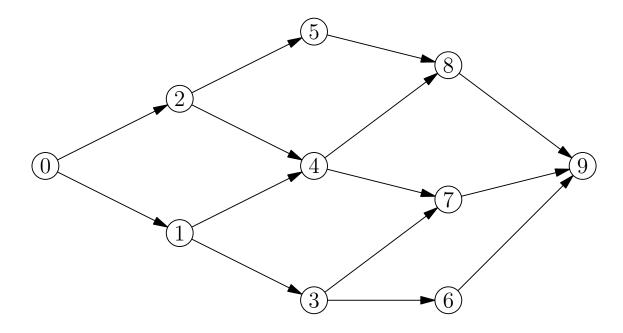
#### **Outline**

- 1. Breadth First Search
  - BFS applications
- 2. Depth First Search
  - DFS applications
- 3. Topological Sort



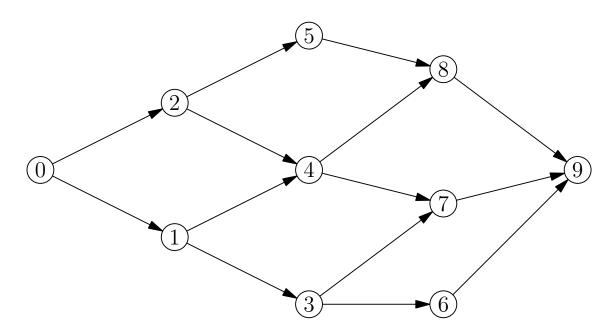
#### **DAGs**

- Directed acyclic graphs or DAGs are directed graphs without cycles
- They are often used to represent complex processes
  - Vertices are processes
  - $\star$  Directed edge (i,j) indicates process i needs to occur before process j



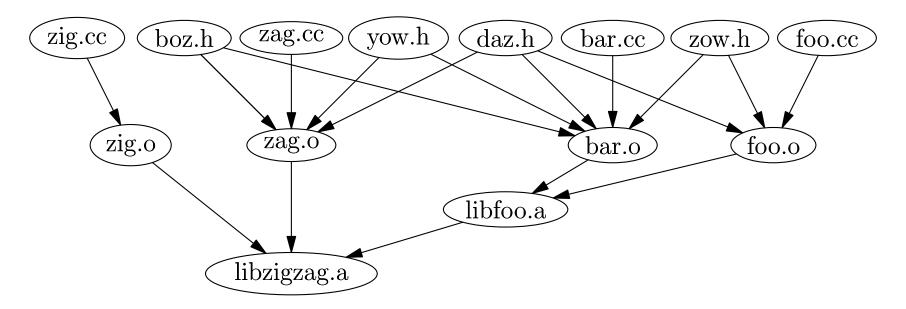
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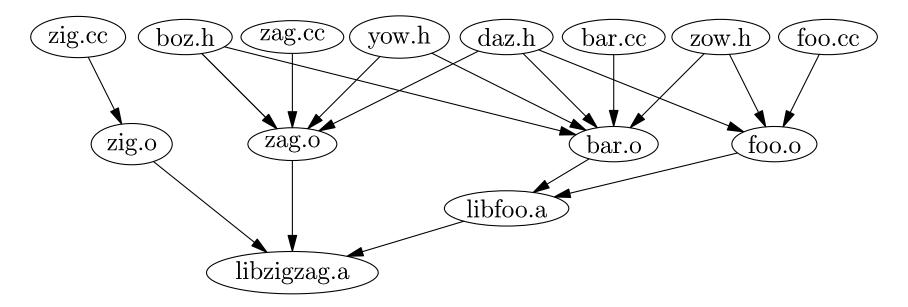
## **Program Compilation**

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- Some programs depend on other programs so they need compiling first



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  - ★ What order should you compile the classes?
- In taking a degree various modules have other modules as prerequisites
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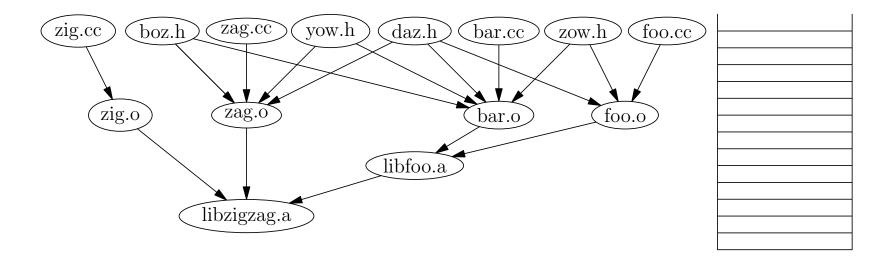
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- If the graphs were not acyclic it would be impossible process them!

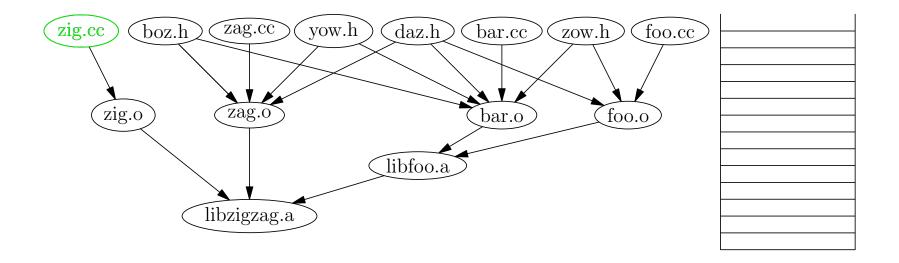
- Given a DAG a topological sort outputs an ordered list of vertices which respects the ordering imposed by the edges
- That is, for each edge (i, j), vertex i will occur before vertex j
- Any DAG will have at least one topological sort, but most DAGs will have many topological sorts
- Topological sort is not a "sort", but it is a useful algorithm for some applications

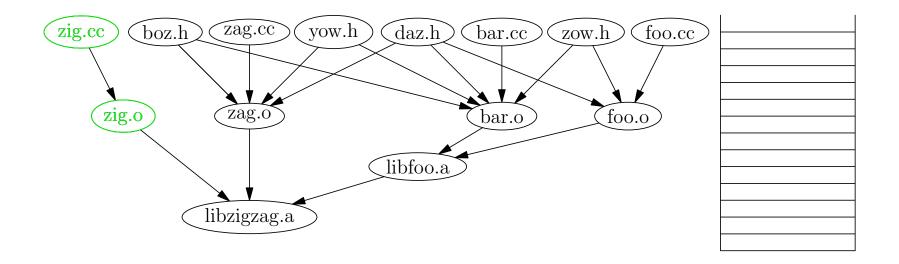
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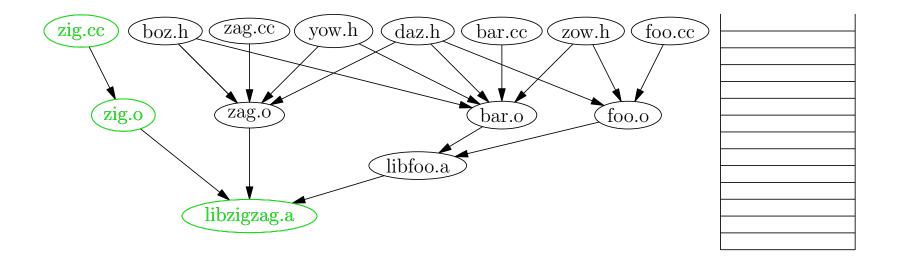
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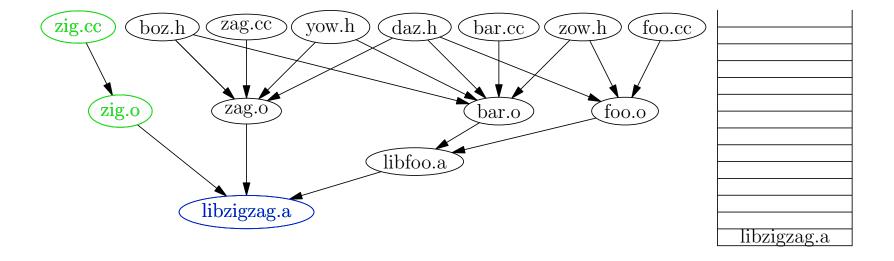
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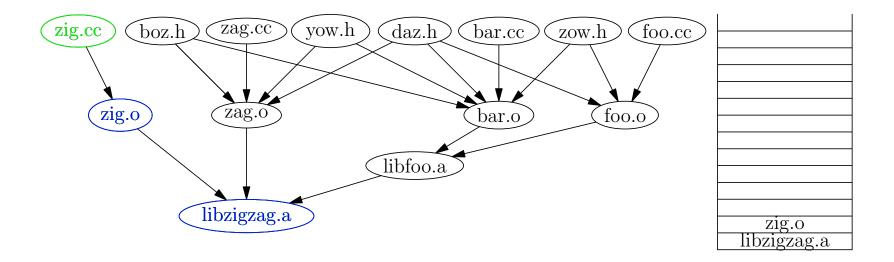


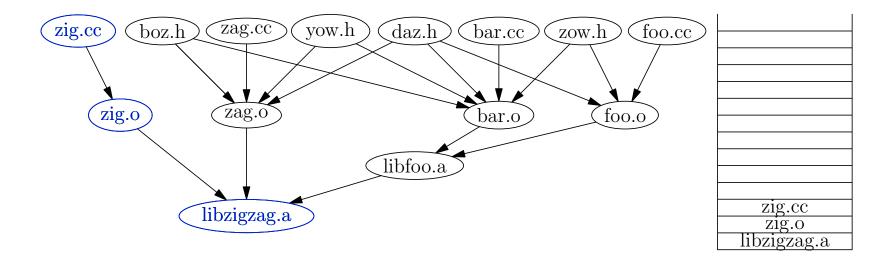


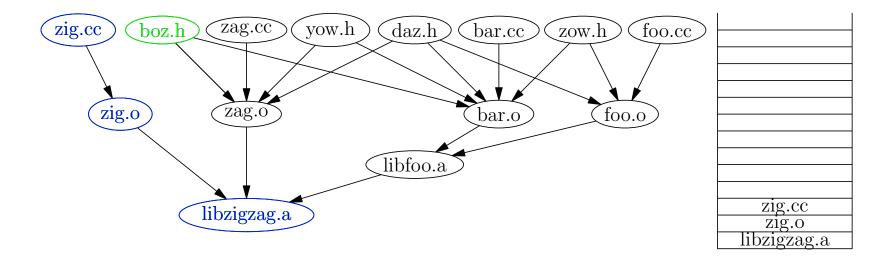


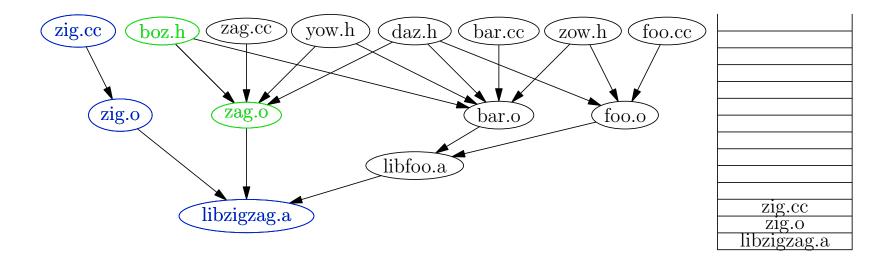


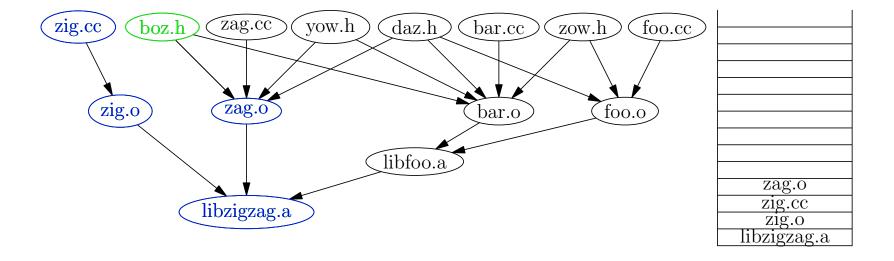


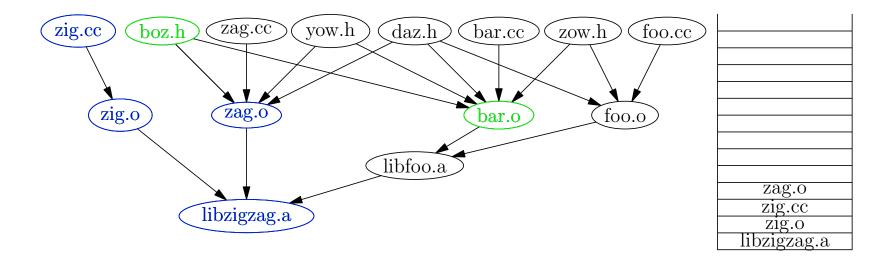


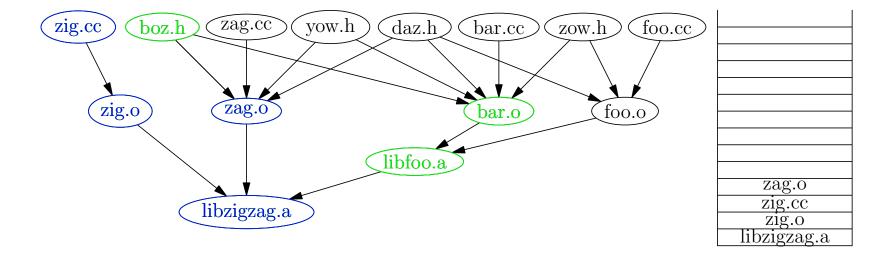


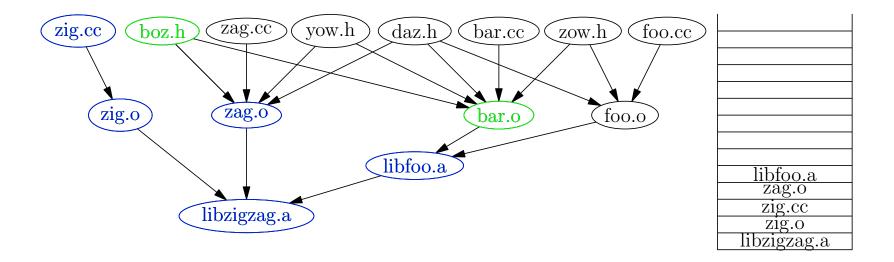


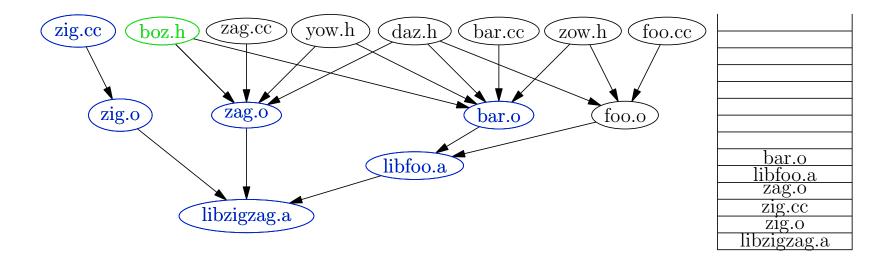


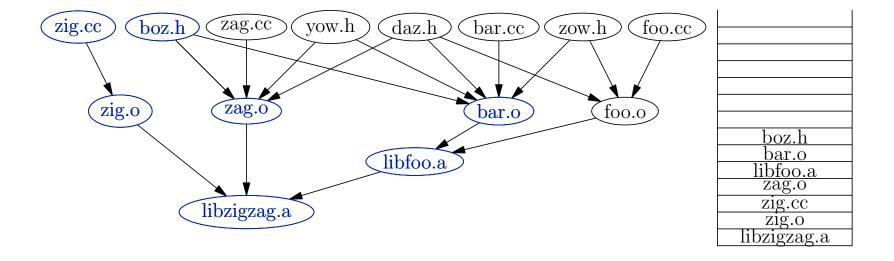


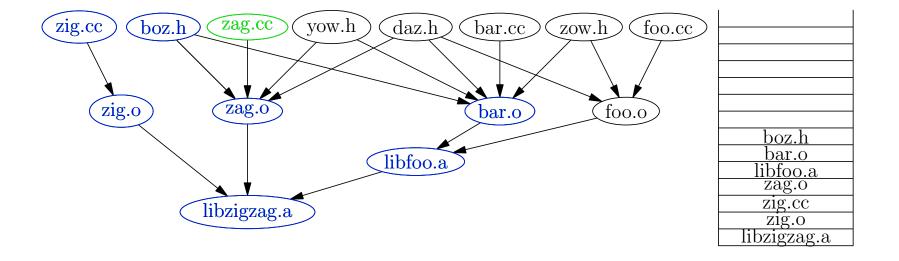


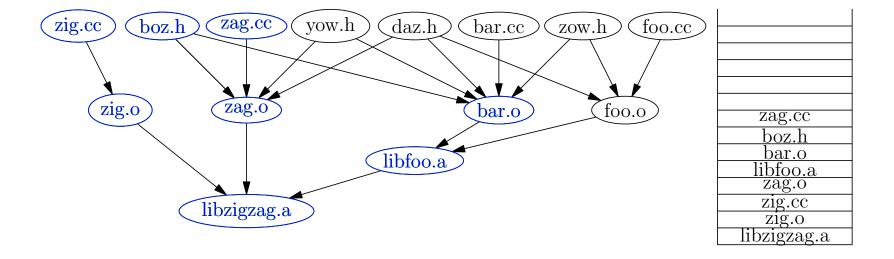


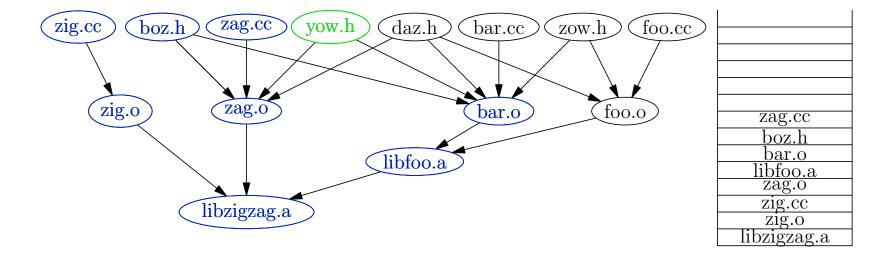


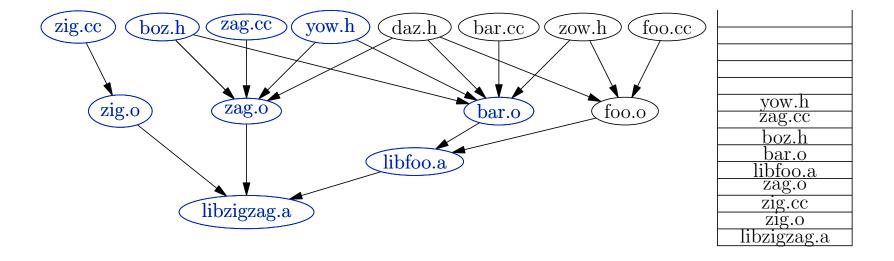


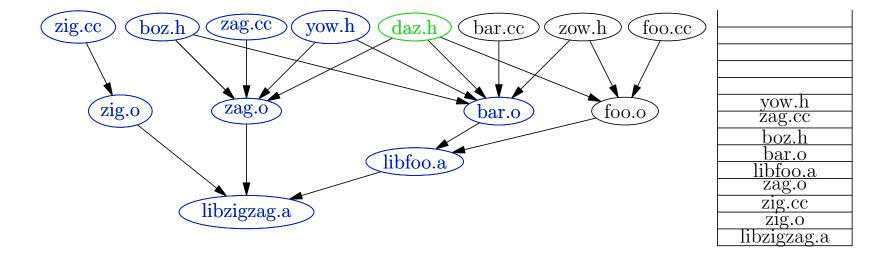


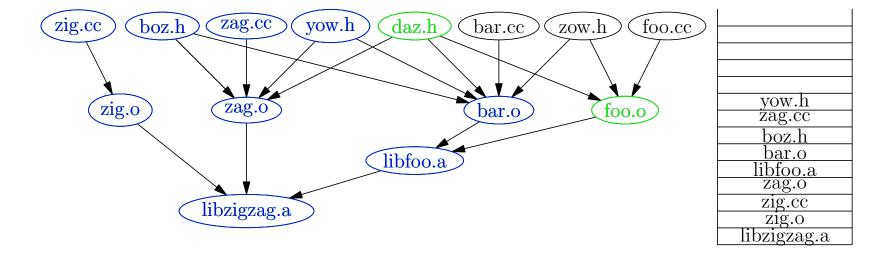


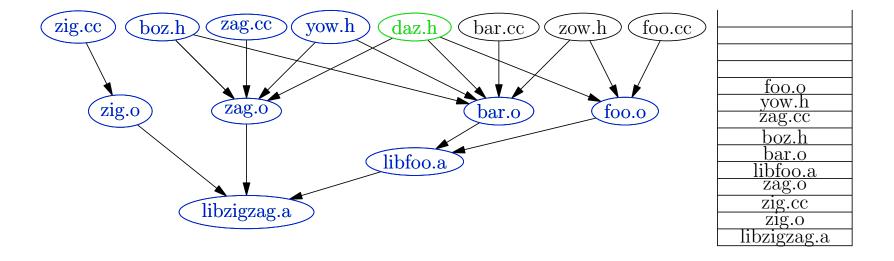


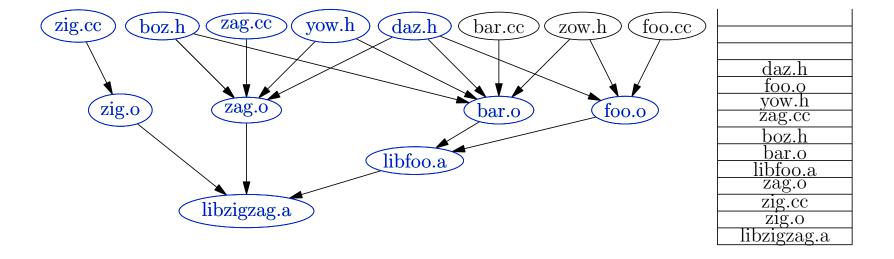


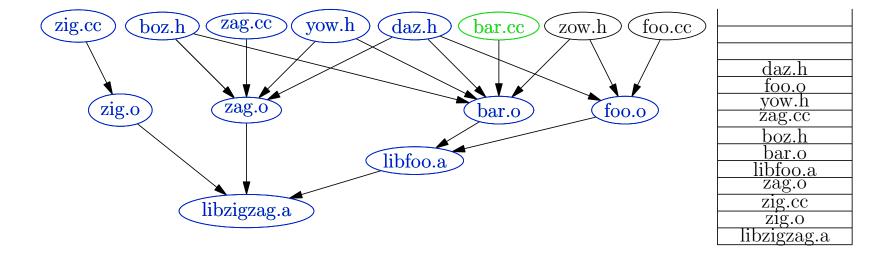


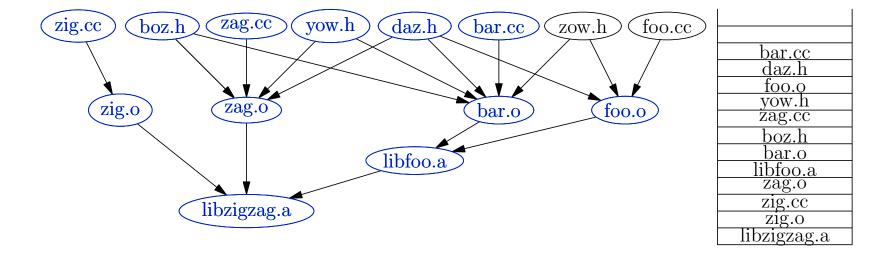


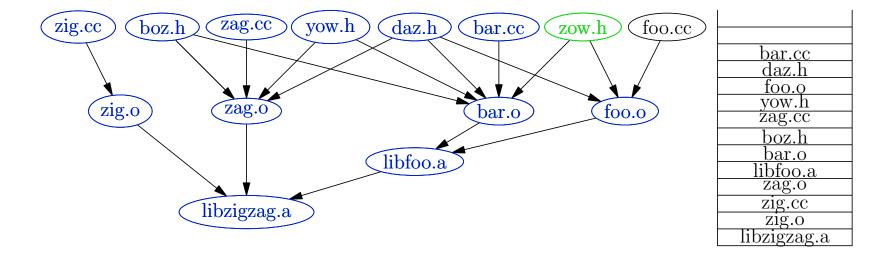


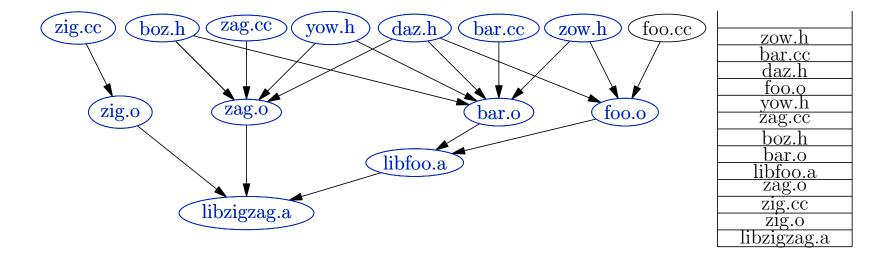


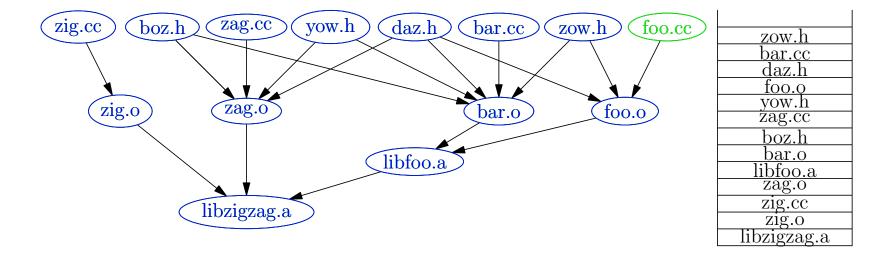


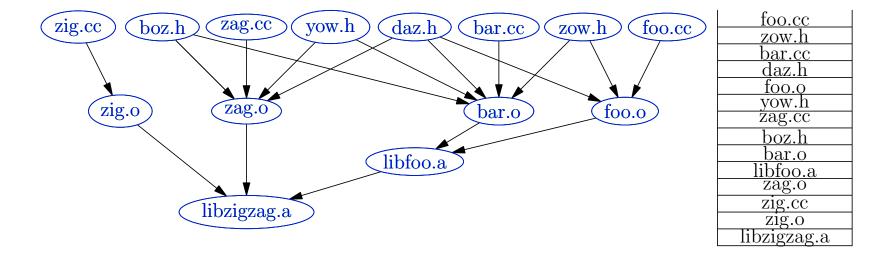












#### **DFS** on Digraphs

```
DFS for undirected graphs
dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  time \leftarrow time + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] \( \neq \) "discovered") \( \{ \)
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    \} else if (state[neighbour] \neq "processed") \{
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  \texttt{time} \leftarrow \texttt{time} + 1
```

#### **DFS** on Digraphs

#### DFS for directed graphs

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dfs(graph, node) {
  if ("finished") return
  state[node] ← "discovered"
  time \leftarrow time + 1
  processVertexEarly(node)
  foreach neighbour ∈ Neighbourhood(node) {
    if (state[neighbour] \( \neq \) "discovered") {
      parent[neighbour] ← node
      processEdge(node, neighbour)
      dfs(graph, neighbour)
    \} else if (state[neighbour] \neq "processed" \lor graph is directed) \{
      processEdge(node, neighbour)
    if ("finished") return
  processVertexLate(currentNode)
  state[currentNode] ← "processed"
  time \leftarrow time + 1
```

### Implementing Topological Sort

• Given our DFS programme we define

```
topologicalSort(graph) {
   Stack stack
   for node ∈ graph.vertexSet()
      if (¬discovered[node])
          dfs(graph, node)

List topSortList
   while (¬stack.isEmpty())
      topSortList.add(stack.pop())

return topSortList
}
```

#### **Enhance DFS**

• Requires us to define a couple of helper function

```
processVertexLast(node) {
    stack.push(node)
}

processEdge(currentNode, neighbour) {
    if (state[neighbour] == "processed") {
        print "error:_graph_not_a_DAG"
        finished = true
    }
}
```

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- Most awkward part of the implementation is that the topologicalSort algorithm needs access to dfs structures (discovered[])
- processVertexLast (node) needs access to the stack
- Need to be able to redefine processVertexFirst, processEdge and processVertexLast
- Different languages and libraries cope with this differently
  - ⋆ Java: JDSL, JGraphT
  - ★ C++: Boost Graph Library, LEDA

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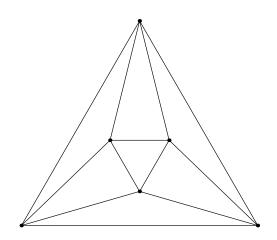
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• DFS is used for many other classic problems

Euler Cycles

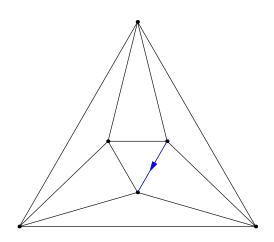
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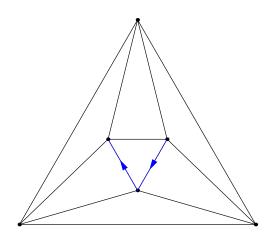
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Euler Cycles



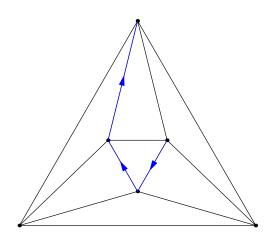
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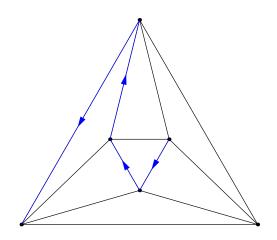
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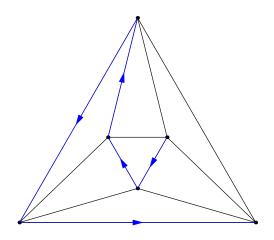
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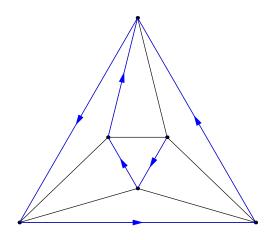
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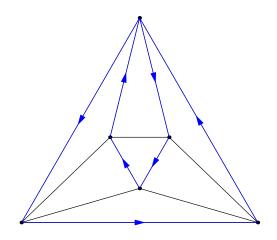
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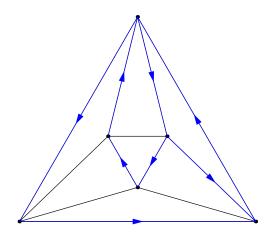
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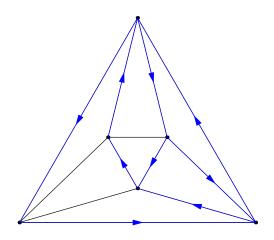
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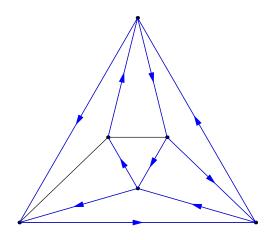
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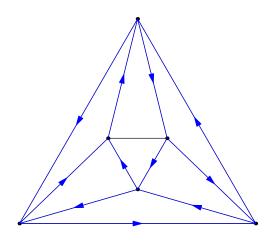
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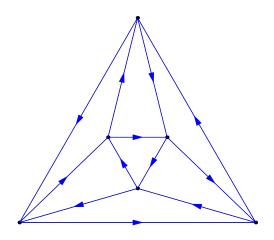
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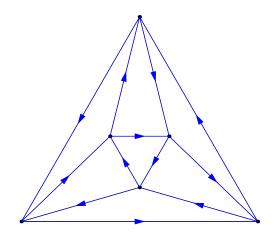
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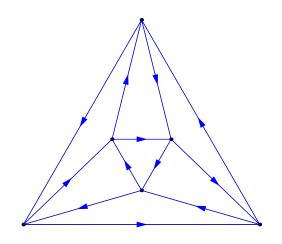
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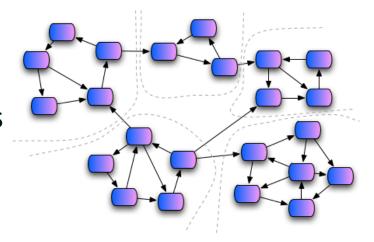
Euler Cycles



DFS is used for many other classic problems

Euler Cycles





- Breadth first and depth first search are different methods for traversing graphs
- They are used as part of many specific algorithms for discovering graph properties
- Breadth first search is particularly important for finding shortest paths in unweighted graphs
- Depth first search is used in a whole host of applications (finding articulation points, Euler cycles, strongly connected components)
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