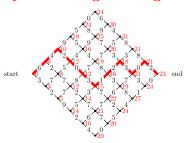
Algorithms and Analysis

Outline

Lesson 23: Dynamic Programming



 $\label{eq:continuous} Dynamic\ programming,\ line\ breaking,\ edit\ distance,\ Dijkstra,\ TSP$

AICE1005

Algorithms and Analysi

Dynamic Programming

- One of the most powerful strategies for solving optimisation problems is **dynamic programming**!
 - ★ Build a set of optimal partial solutions
 - ⋆ Increase the size of the partial solutions until you have a full solution
 - ★ Each step uses the set of optimal partial solutions found in the previous step!
- Developed by Richard Bellman in the early 1950's
- The name is unfortunate as it doesn't have much to do with programming!

AICE1005

Algorithms and Analys

Brute Force

- The obvious brute force strategy is to try every path
- For a problem with n steps we require n/2 to be diagonally up and n/2 to be diagonally down!
- The total number of paths is

$$\binom{n}{n/2} \approx \sqrt{\frac{2}{\pi\,n}}\,2^n {\rm I}$$

- \bullet For the problem shown above with n=8 there are 70 paths!
- \bullet For a problem with n=100 there are 1.01×10^{29} paths!

AICE1005

Algorithms and Analys

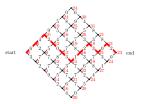
Forward Algorithm

- ullet Suppose we know the optimal costs for all the edge in column i
- ullet Our task is to find the optimal cost at column i+1
- If we consider the sites in the lattice then the optimal cost will be

$$c_{(i+1,j)} = \min \left(c_{(i,j+1)} + w_{(i,j+1)(i+1,j)}, \ c_{(i,j-1)} + w_{(i,j-1)(i+1,j)} \right) \mathbf{I}$$

- This is the defining equation in dynamic programming
- We have to treat the boundary sites specially, but this is just book-keeping

- 1. Dynamic Programming
- 2. Applications
 - Line Breaks
 - Edit Distance
 - Dijkstra's Algorithm
- 3. Limitation

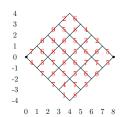


AICE1005

Igorithms and Analysis

A Toy Problem

 Consider the problem of find a minimum cost path from point (0,0) to (8,0) on the lattice



- The costs of traversing each link is shown in red
- The cost of a path is the sum of weights on each link!

AICE100

gorithms and Analysis

Building a Solution

- We can solve this problem efficiently using dynamic programming by considering optimal paths of shorter length
- \bullet Let $c_{(i,j)}$ denote the cost of the optimal path to node (i,j)
- \bullet We denote the weights between two points on the lattice by $w_{(i,j)(i+1,j\pm 1)}$

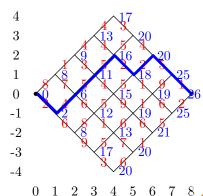


• Clearly $c_{(0,0)} = 0$

AICE1005

Algorithms and Analysis

Example



Backward Algorithm

- \bullet Having found the optimal costs $c_{(i,j)}$ we can find the optimal path starting from $(n,0) {\rm I\!I}$
- At each step we have a choice of going up or down
- We choose the direction which satisfies the constraint

$$c_{(i,j)} = c_{(i-1,j\pm 1)} + w_{(i-1,j\pm 1)(i,j)}$$

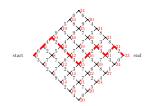
 If both directions satisfy the constraint we have more than one optimal path

AICE1005

Algorithms and Analysi

Outline

- 1. Dynamic Programming
- 2. Applications
 - Line Breaks
 - Edit Distance
 - Dijkstra's Algorithm
- 3. Limitation



AICE1005

Ilgorithms and Analysis

Using Dynamic Programming

- The challenge is recognising that you can use dynamic programming and representing the problem right
- Learn this from examples
- Consider writing a word processor that splits paragraphs up into lines
- You want to choose the line breaks so that the lines are all roughly the same length!
- This is a global optimisation task
 minimise the total number of spaces left at the end of each
 line

AICE1005

Algorithms and Analys

Real Word Breaking

- In advanced word processing you care about hyphenation, large gaps at the end of lines, etc.
- These all affect the way you would assign costs
- Dynamic programming is used in LATEX to produce nice line breaks
- A similar algorithm is used to produce nice page breaks

Time Complexity

- \bullet In our dynamic programming solution we had to compute the cost $c_{(i,j)}$ at each lattice point!
- There were $(\frac{n+1}{2})^2$ lattice point
- It took constant time to compute each cost so the total time to perform the forward algorithm was $\Theta(n^2)$
- ullet The time complexity of the backward algorithm was $\Theta(n)$
- This compares with $\exp(\Theta(n))$ for the brute force algorithm

AICE1005 Algorithms at

Applications of Dynamic Programming

- Dynamic programming is used in a vast number of applications
 - ★ String matching algorithms
 - ⋆ Shape matching in images
 - ★ Dynamical time-warping in speech
 - ⋆ Hidden Markov Models in machine learning
- Unlike greedy algorithms the idea is readily extended to many different applications

AICE1005 Algorithms and Analysis 12

Using Dynamic Programming

 \bullet I have a dream that one day this nation will rise up and live out the. . . \blacksquare



AICE1005 Algorithms and Analysis 1

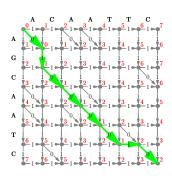
Inexact Matching

- A second example of dynamic programming is to find inexact matches
- The edit distance between two strings is the number of changes needed to move from one string to another
- The exact metric depends on the application, but might include number of substitutions, insertions and deletions.
- This has many applications, e.g. in genomics to see what DNA strings (or proteins) are related

AICE1005 Algorithms and Analysis 15 AICE1005 Algorithms and Analysis

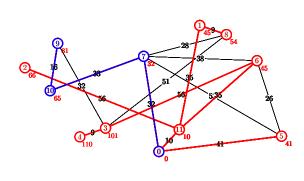
Edit Distance

 What is the minimum edit distance between ACAATTC and AGCAATCI



AICE1005 Algorithms and Analysis

Going from Node 0 to Node 9



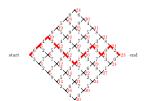
AICE1005 Algorithms and Analysis

- The use of transposition tables and refutation tables in computer chess
- The Viterbi algorithm (used for hidden Markov models)
- The Earley algorithm (a type of chart parser)
- The Needleman–Wunsch and other algorithms used in bioinformatics, including sequence alignment, structural alignment, RNA structure prediction
- Floyd's all-pairs shortest path algorithm
- Optimizing the order for chain matrix multiplication
- Pseudo-polynomial time algorithms for the subset sum and knapsack and partition problems
- The dynamic time warping algorithm for computing the global distance between two time series
- The Selinger (a.k.a. System R) algorithm for relational database query optimization
- De Boor algorithm for evaluating B-spline curves
- Duckworth–Lewis method for resolving the problem when games of cricket are interrupted

AICE1005 Algorithms and Analysis 2:

Outline

- 1. Dynamic Programming
- 2. Applications
 - Line Breaks
 - Edit Distance
 - Dijkstra's Algorithm
- 3. Limitation



Dijkstra's Algorithm

- We saw early Dijkstra's algorithm for find the minimum distance between a source and destination node!
- We grouped this with the greedy algorithms as we choose the next node to add to the minimum-distance spanning tree to be the closest node to the source we could access!
- However, we should perhaps more rightly identify it as using a dynamic programming strategy as we are building up the cost of getting of the partial solution to reach a nodel
- We use the greedy strategy to ensure that we always find the shorter paths first

AICE1005 Algorithms and Analysis

Uses of Dynamic Programming

- · Recurrent solutions to lattice models for protein-DNA binding
- Backward induction as a solution method for finite-horizon discrete-time dynamic optimization problems
- Method of undetermined coefficients can be used to solve the Bellman equation in infinite-horizon, discrete-time, discounted, time-invariant dynamic optimization problems
- Many string algorithms including longest common subsequence, longest increasing subsequence, longest common substring, Levenshtein distance (edit distance)
- Many algorithmic problems on graphs can be solved efficiently for graphs of bounded treewidth or bounded clique-width by using dynamic programming on a tree decomposition of the graph.
- The Cocke–Younger–Kasami (CYK) algorithm which determines whether and how a given string can be generated by a given context-free grammar
- Knuth's word wrapping algorithm that minimizes raggedness when word wrapping text

AICE1005 Algorithms and Analysis 2

- The value iteration method for solving Markov decision processes
- Some graphic image edge following selection methods such as the "magnet" selection tool in Photoshop
- Some methods for solving interval scheduling problems
- Some methods for solving word wrap problems
- Some methods for solving the travelling salesman problem, either exactly (in exponential time) or approximately (e.g. via the bitonic tour)
- Recursive least squares method
- Beat tracking in music information retrieval
- Adaptive-critic training strategy for artificial neural networks
- Stereo algorithms for solving the correspondence problem used in stereo vision
- Seam carving (content aware image resizing)
- \bullet The Bellman–Ford algorithm for finding the shortest distance in a graph
- Some approximate solution methods for the linear search problem
- Kadane's algorithm for the maximum subarray problem

AICE1005 Algorithms and Analysis 22

When You Can't Use It

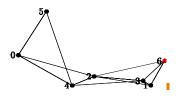
- Not all problems can be split neatly to make dynamic programming possible
- Dynamic programming works on problems with some natural ordering
- We need this to build up a list of optimum cost of partial solutions—these have to depend on the cost of previous partial solutions
- Sometime no natural ordering exists

CE1005 Algorithms and Analysis 23 AICE1005 Algorithms and Analysis

ī

Travelling Salesman Problem

- For TSP we can find a dynamic programming solution by considering optimal sub-tours (paths) involving sets of k cities
- If we know the optimal sub-tour through all sets of cities of size k (starting and finishing at each possible pair in the set) then we can quickly compute the optimal sub-tour between set of size k+1



AICE1005 Algorithms and Analysis 25

Conclusions

- Dynamic programming is one of the most powerful strategies for solving hard optimisation problems
- It works by iteratively building up costs for partial solutions using the costs of smaller partial solutions
- When it works it is great and there are hosts of practical algorithms which use DPI
- However, it doesn't always work

AICE1005

Algorithms and Analys

Travelling Salesman Problem

- The problem is there are $\binom{n}{k}$ subsets consisting of k cities out of a possible $n\mathbf{l}$
- The total number of subsets that need to be considered is

$$\sum_{k=0}^{n} \binom{n}{k} = 2^{n}$$

• The time complexity of the DP solution is $n^2\,2^n$ which is better than n! and is currently the fastest known exact algorithm for TSP, but it ain't very useful in practice!

AICE1005 Algorithms and Analysis 2