Algorithms and Analysis

Lesson 14: *Use Heaps!*

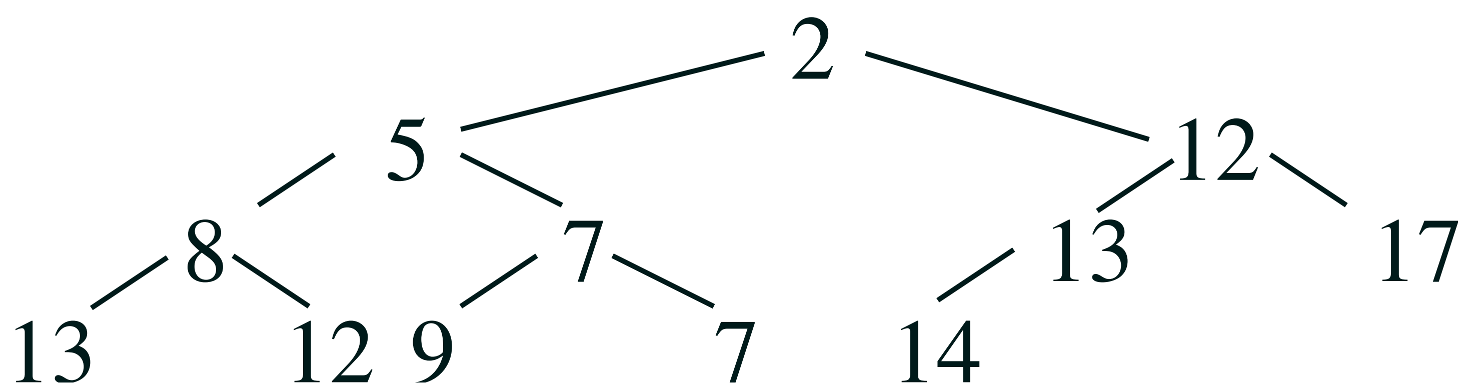
Heaps, Priority queues, Heap Sort, Other heaps

Outline

1. **Heaps**
2. Priority Queues
   - Array Implementation
3. Heap Sort
4. Other Heaps

Heaps

- A (min-)heap is (from one perspective) a binary tree
- It is a binary tree satisfying two constraints
  - ⋆ It is a **complete** tree
  - ⋆ Each child has a value ‘greater than or equal to’ its parent



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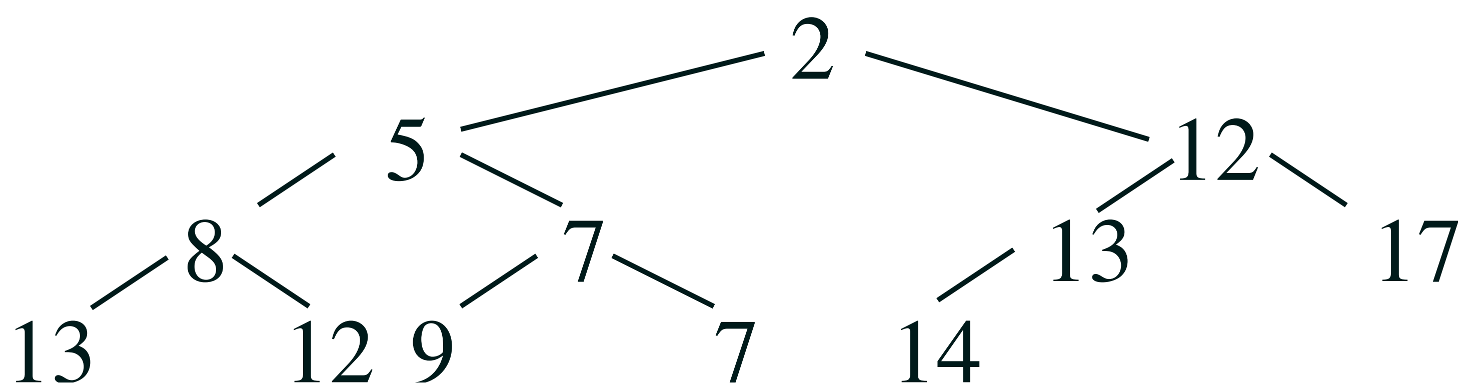

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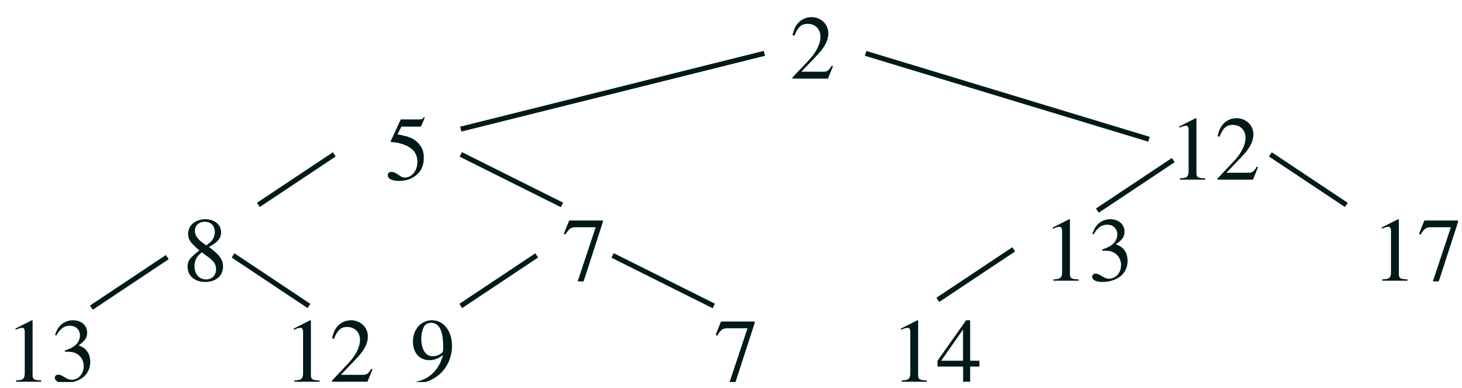

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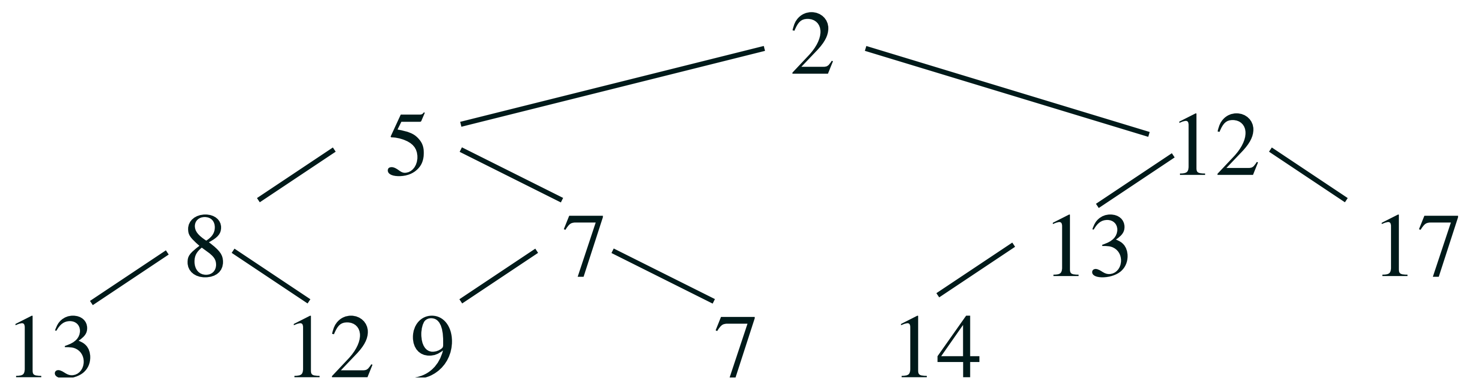

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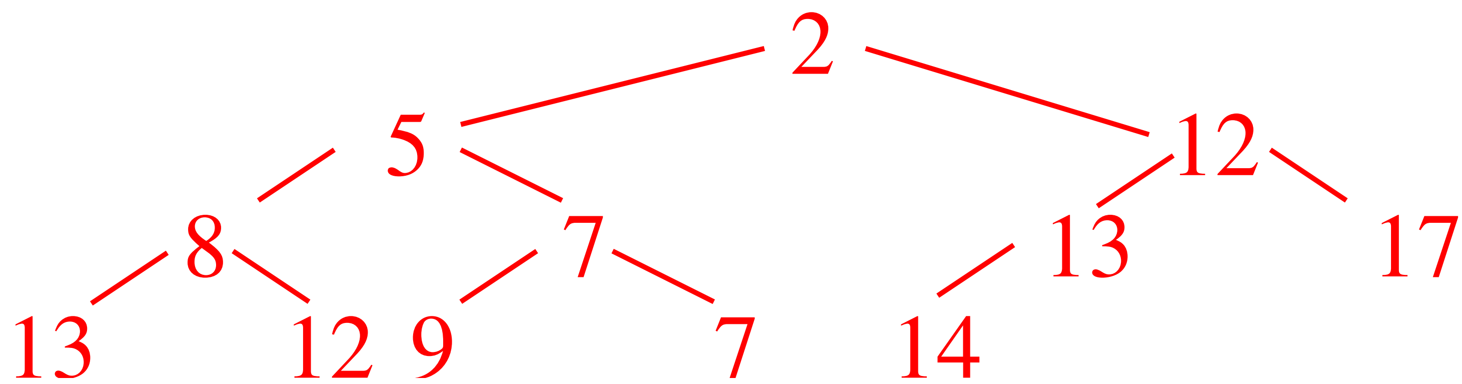

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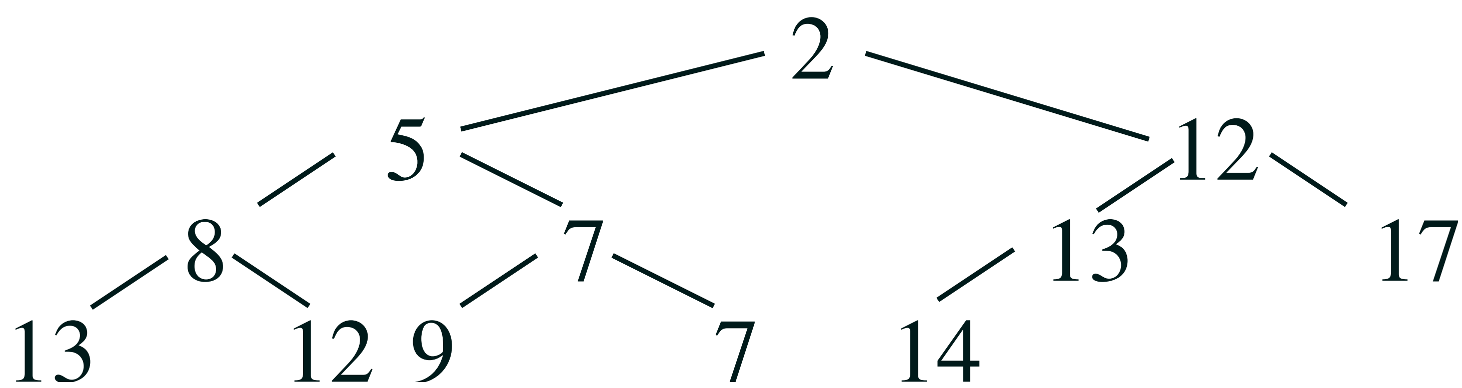

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Priority Queues

- One of the prime uses of heaps is to implement a priority queue
- A Priority Queue is a queue with priorities
- That is, we assign a priority to each element we add
- The head of the queue is the element with highest priority (smallest number)
- Used, for example, in simulating real time events
- Used to implement "greedy algorithms"

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Priority Queue

- A simple Priority Queue might include
  - `unsigned size()` returning the the number of elements
  - `bool empty()` returns true if empty
  - `void push(T element, int priority)` adds an element
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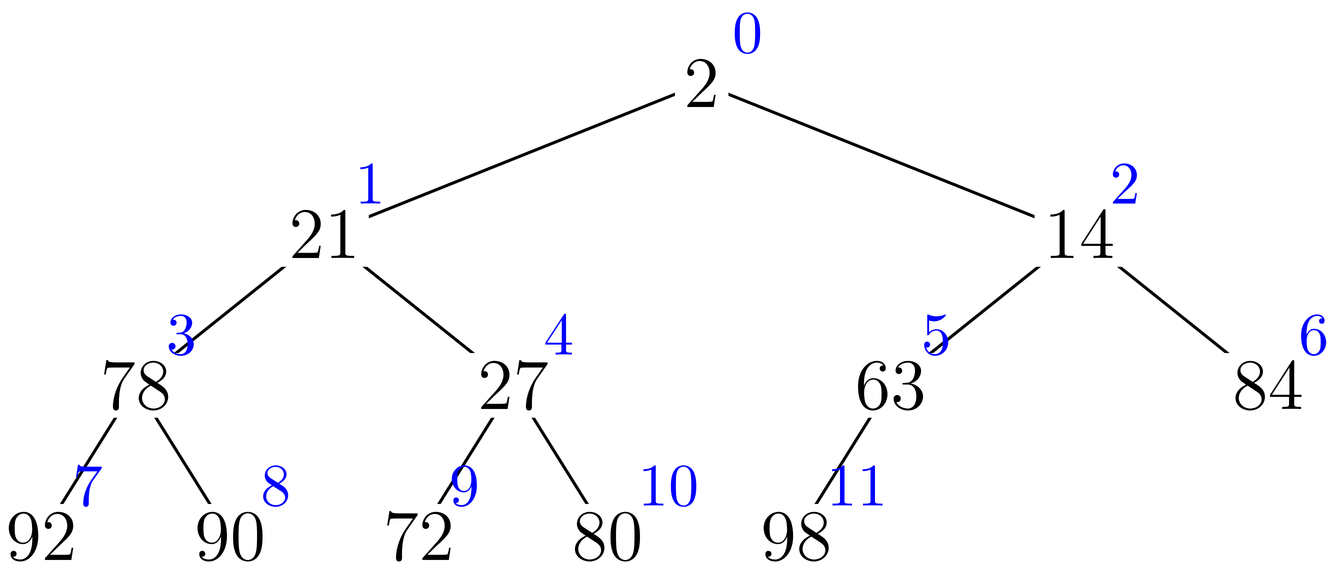

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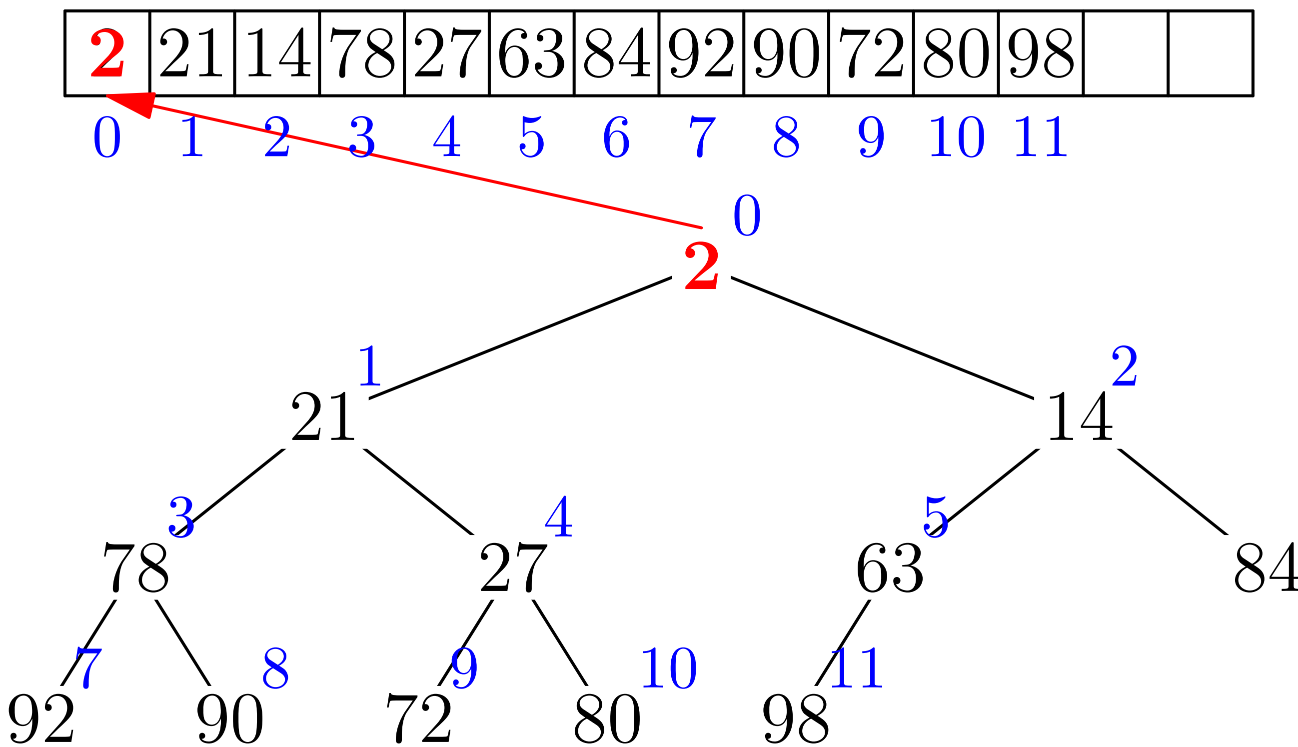

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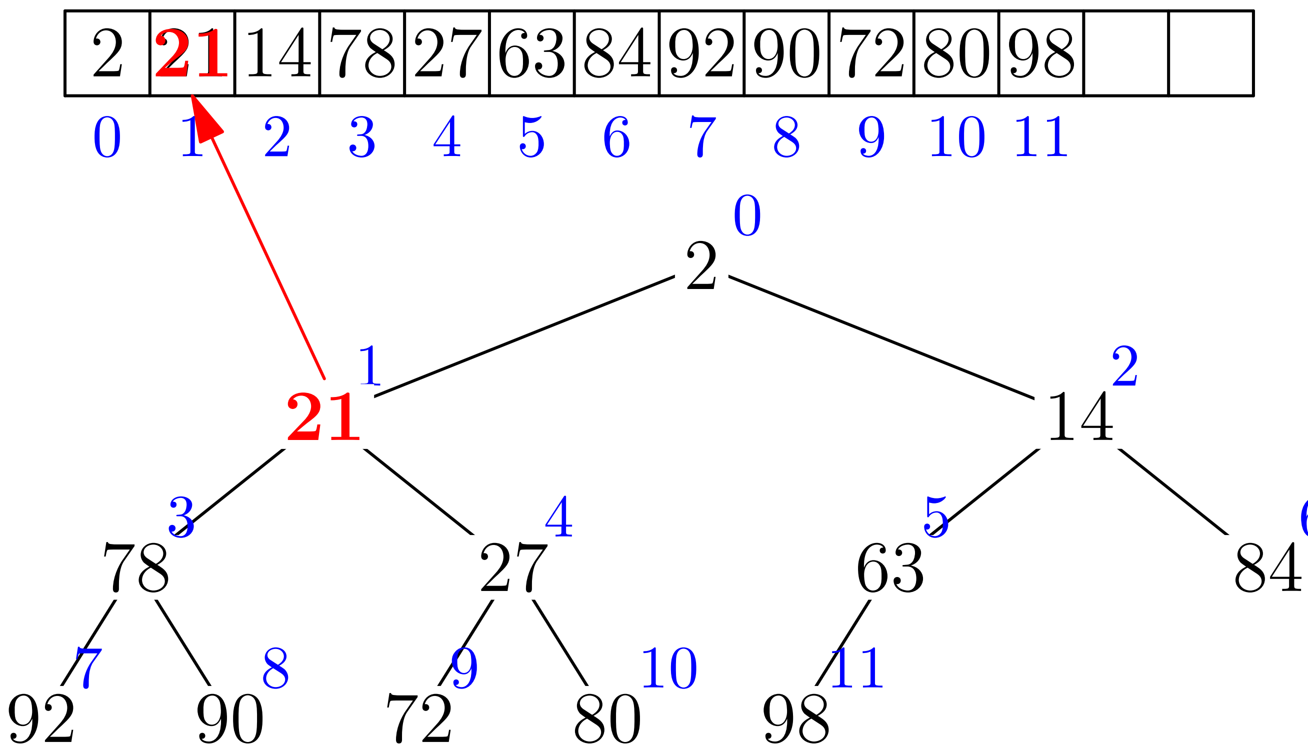

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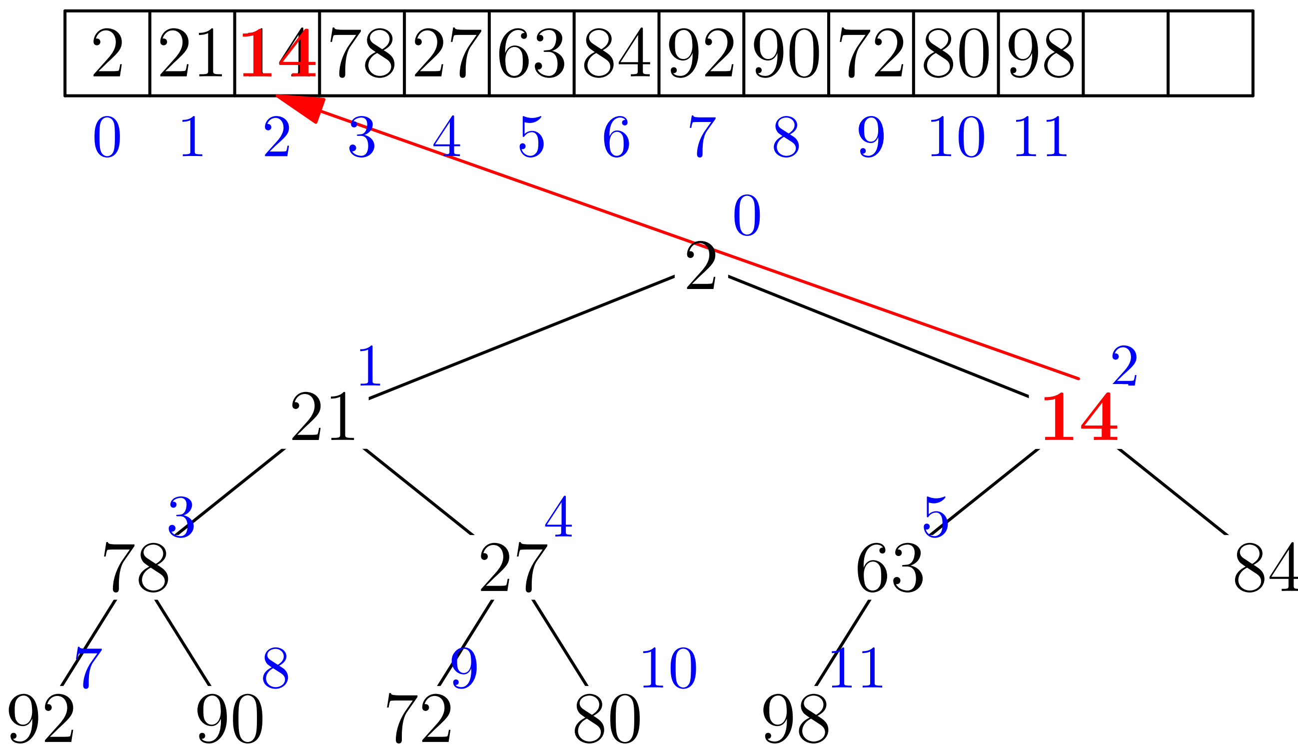

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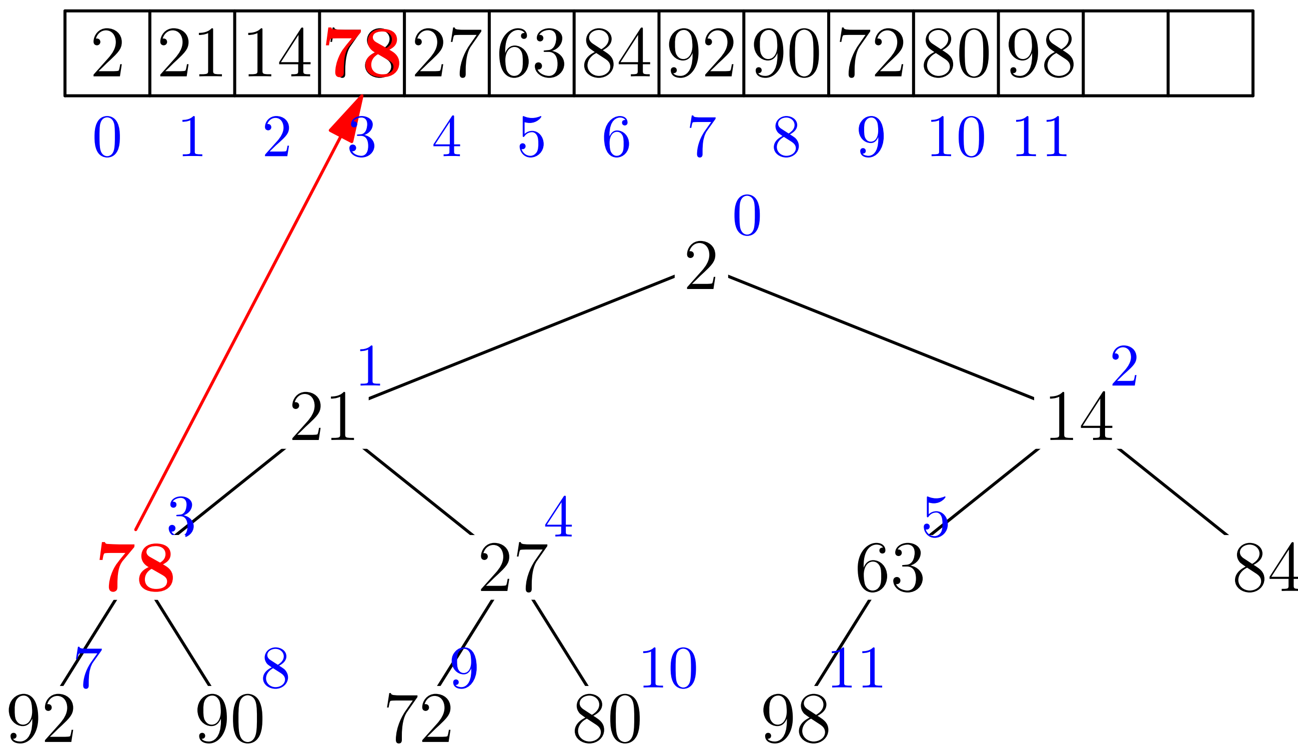

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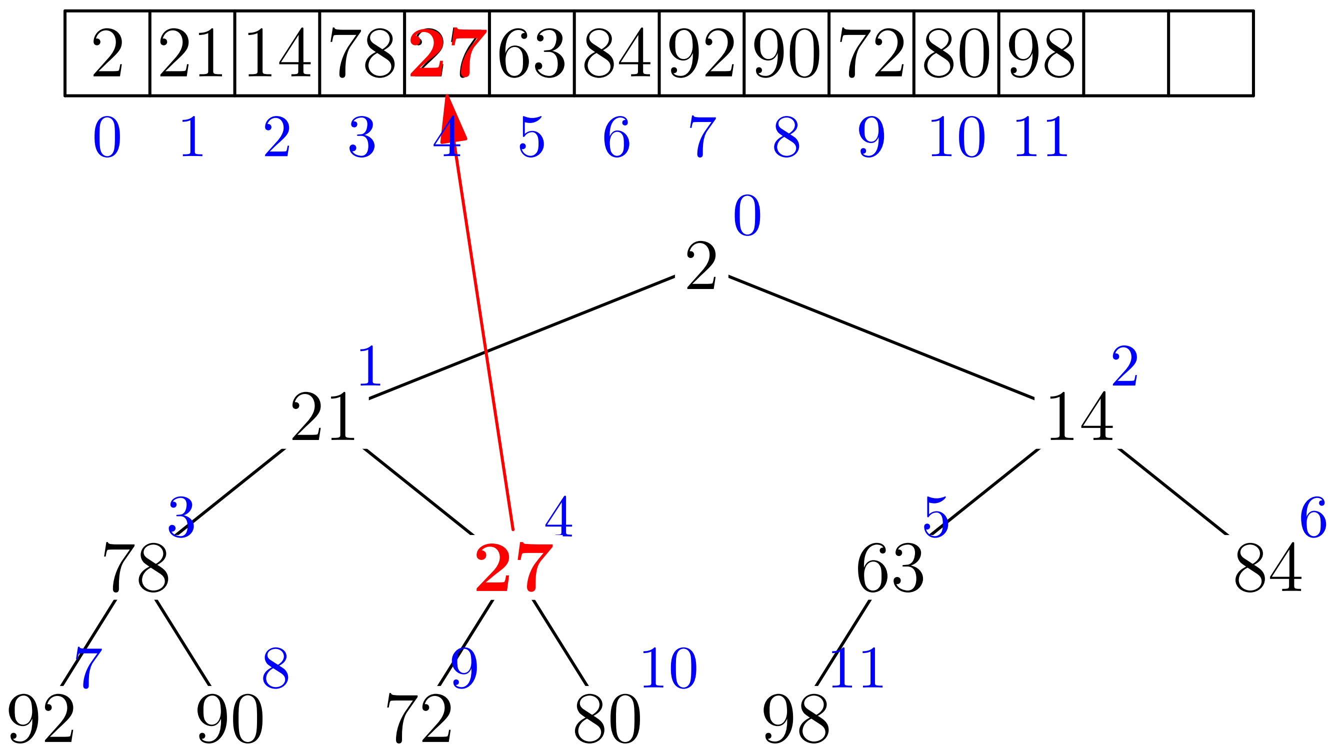

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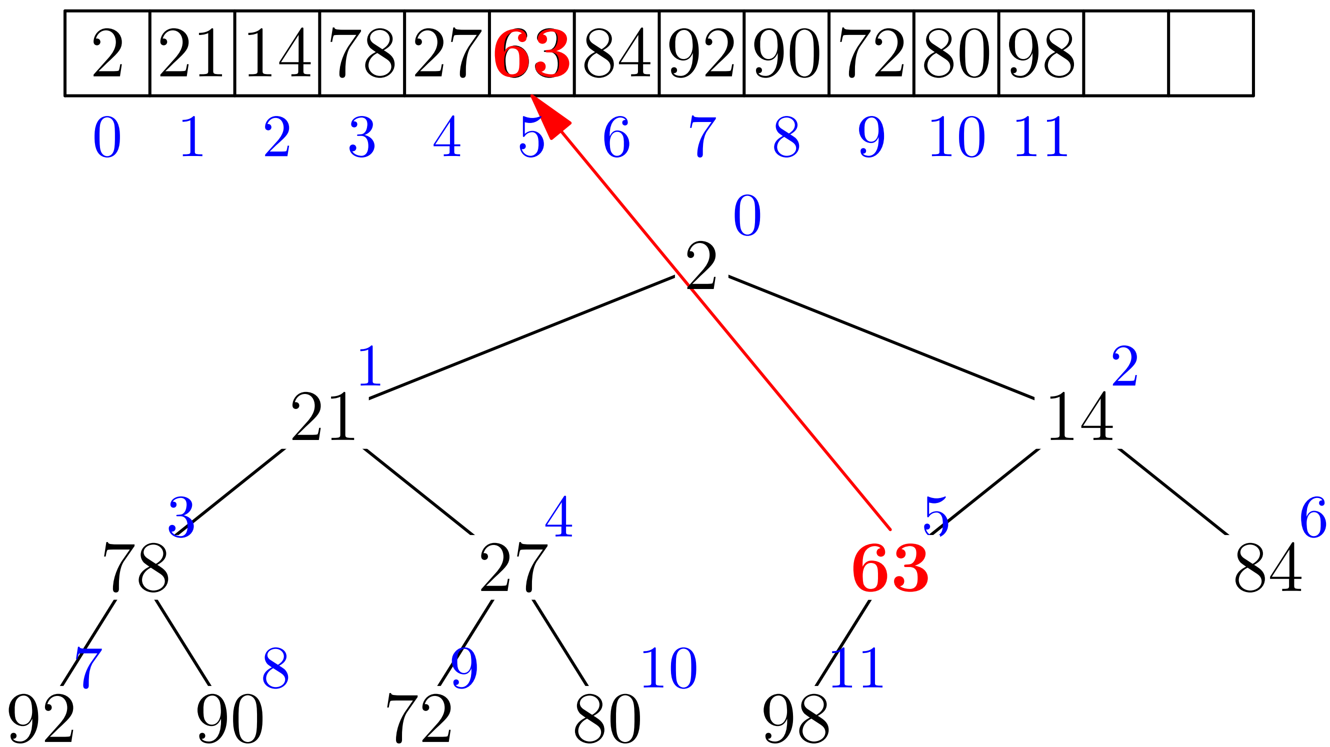

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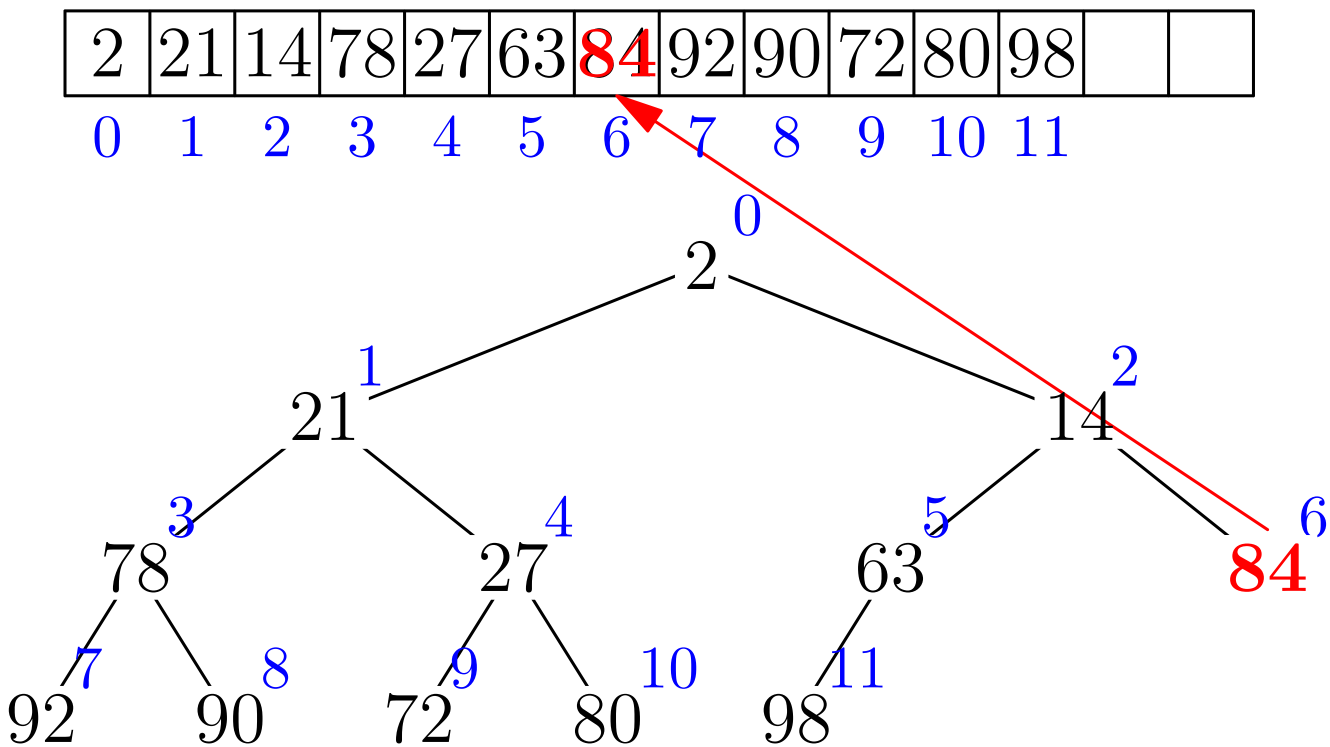

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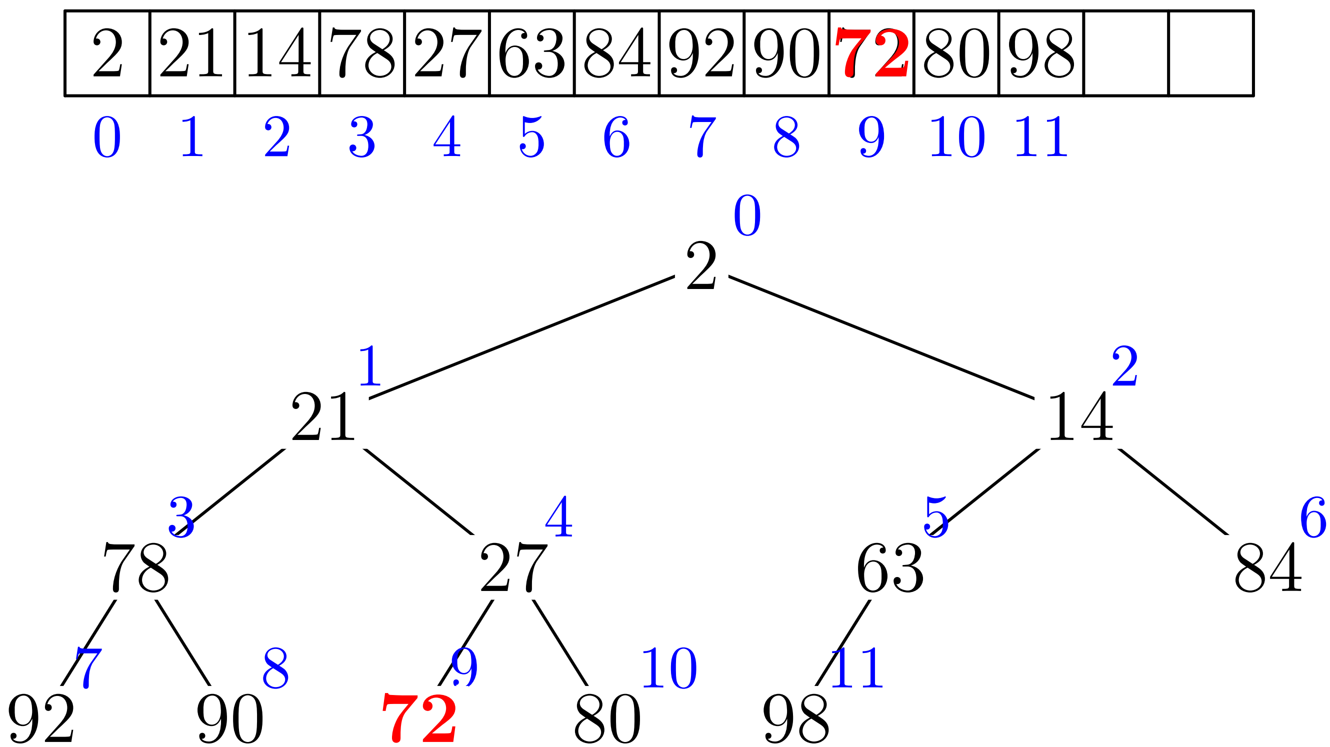

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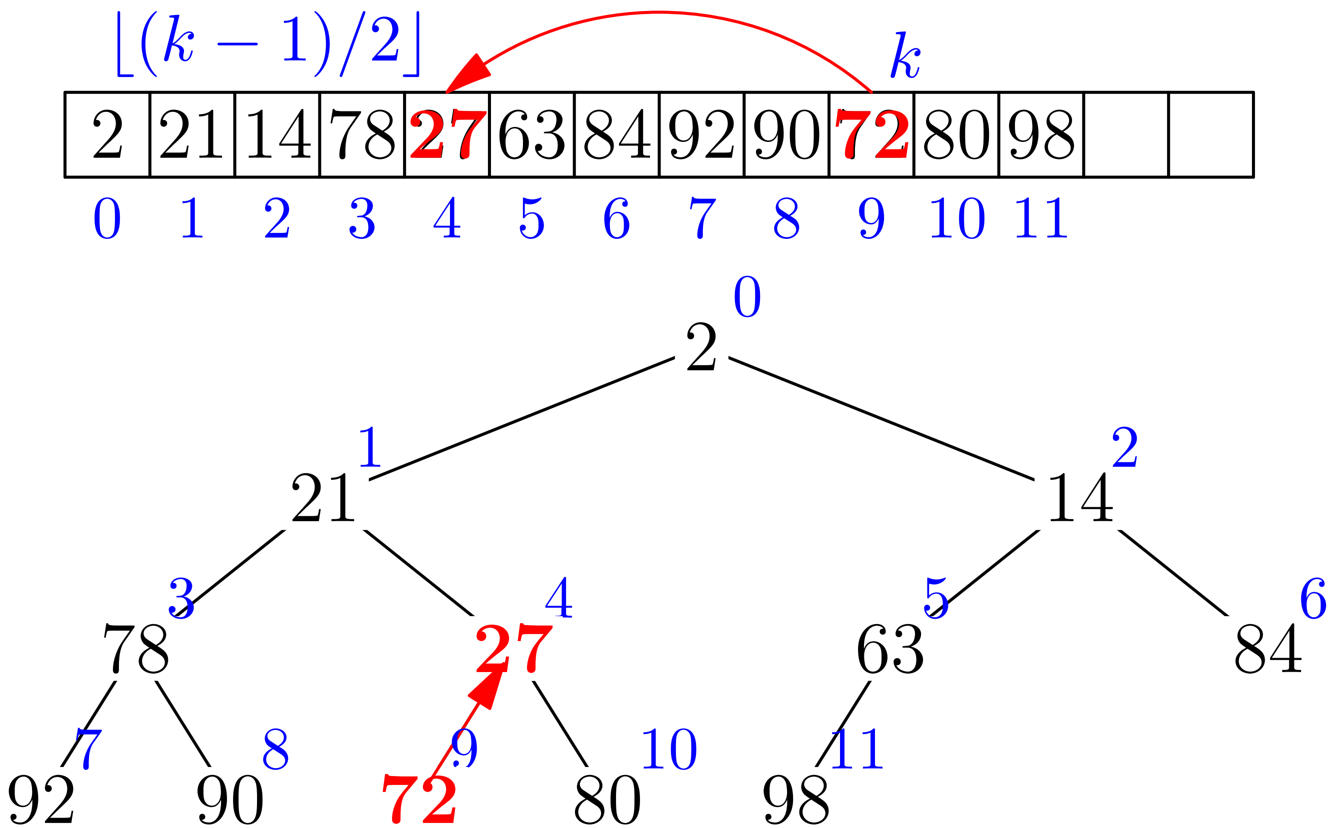

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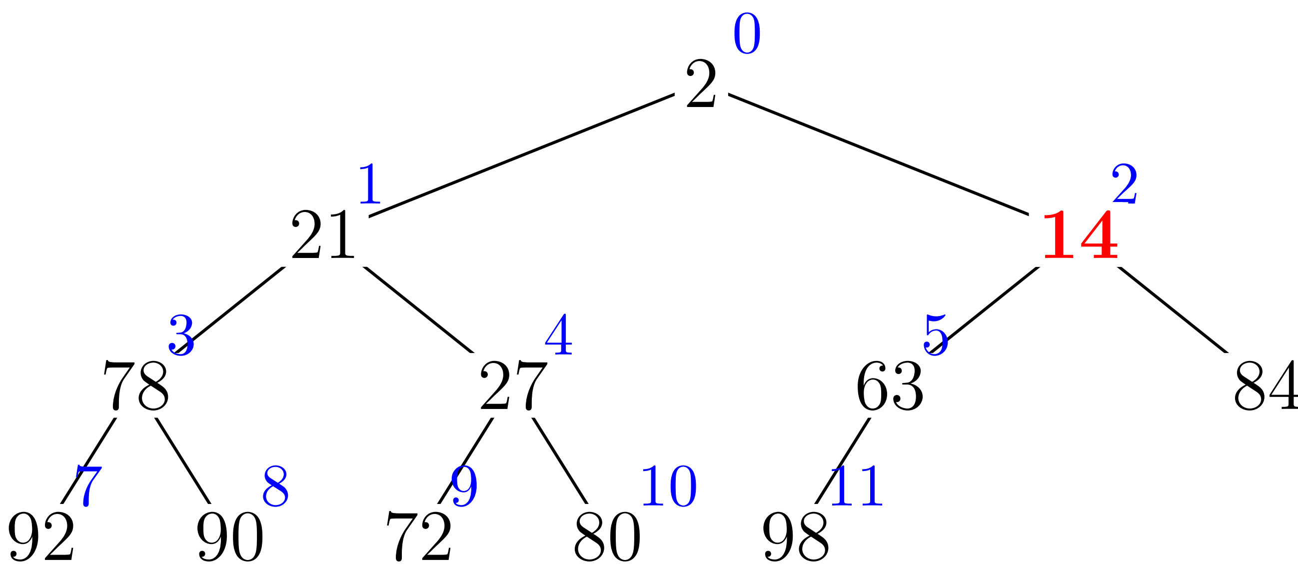

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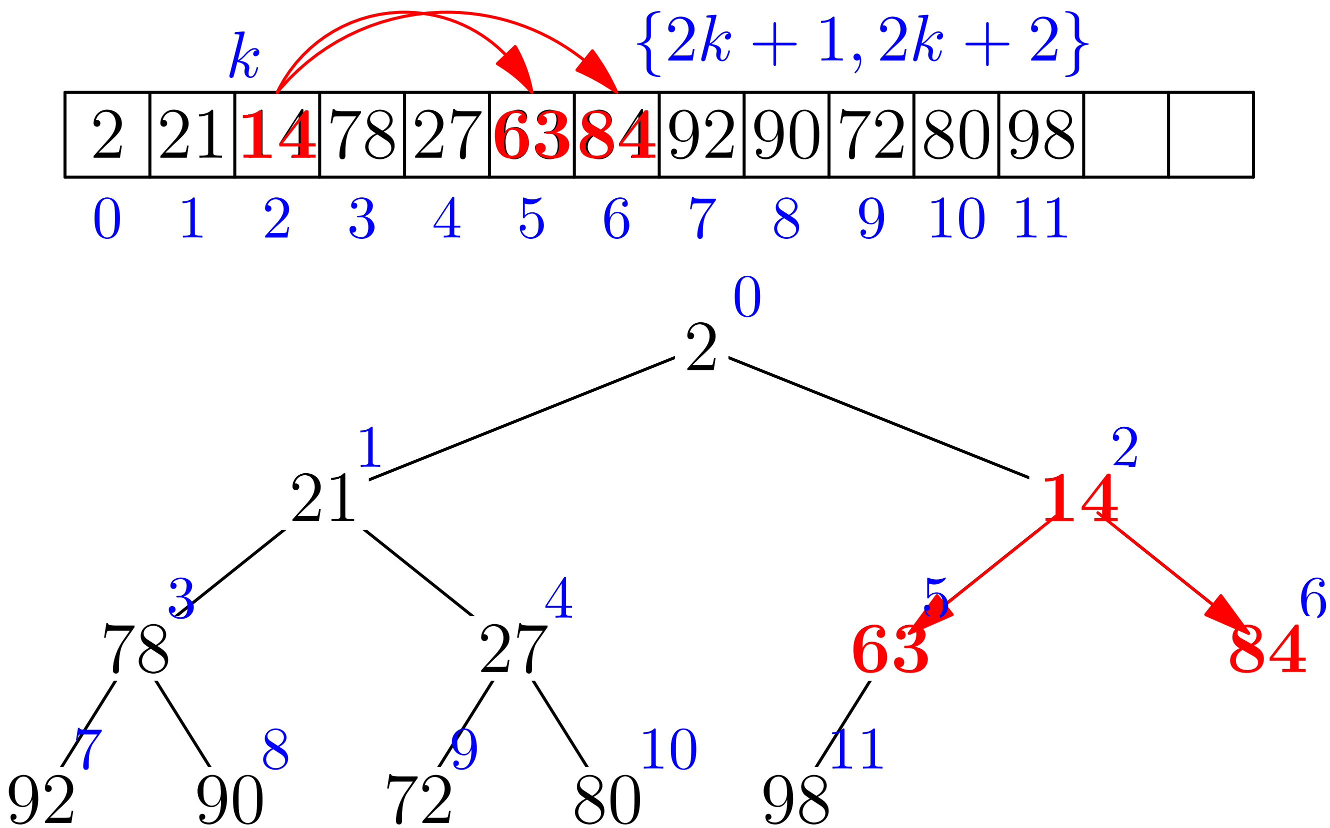

Code for a Priority Queue

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using namespace std;

template <typename T, typename P>
class heapPQ {
private:
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  const T& top() {return array[0].first;}
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Adding an Element

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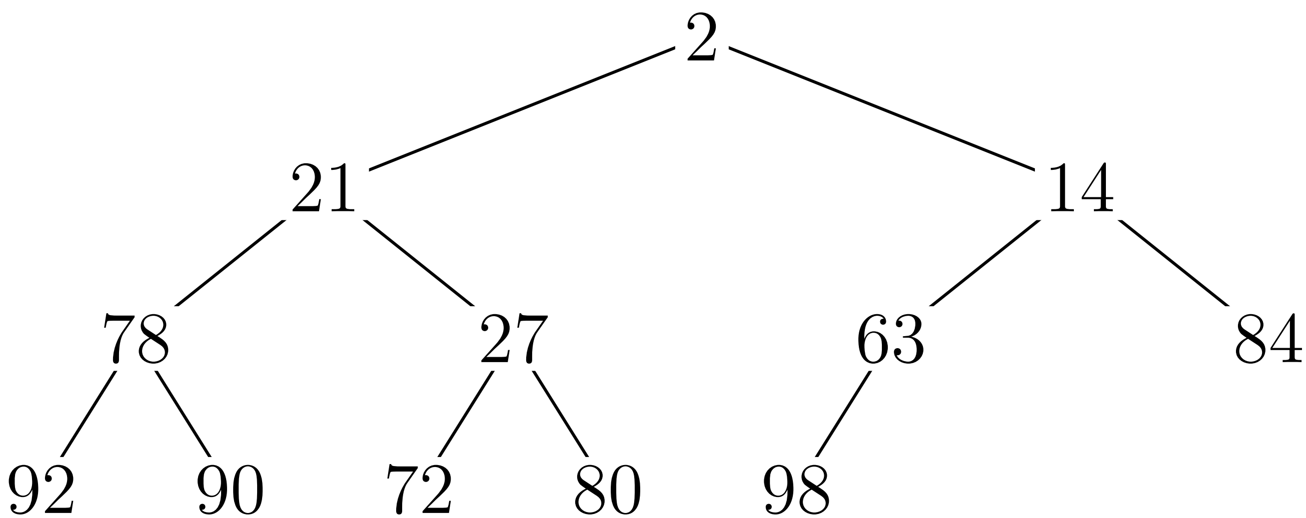
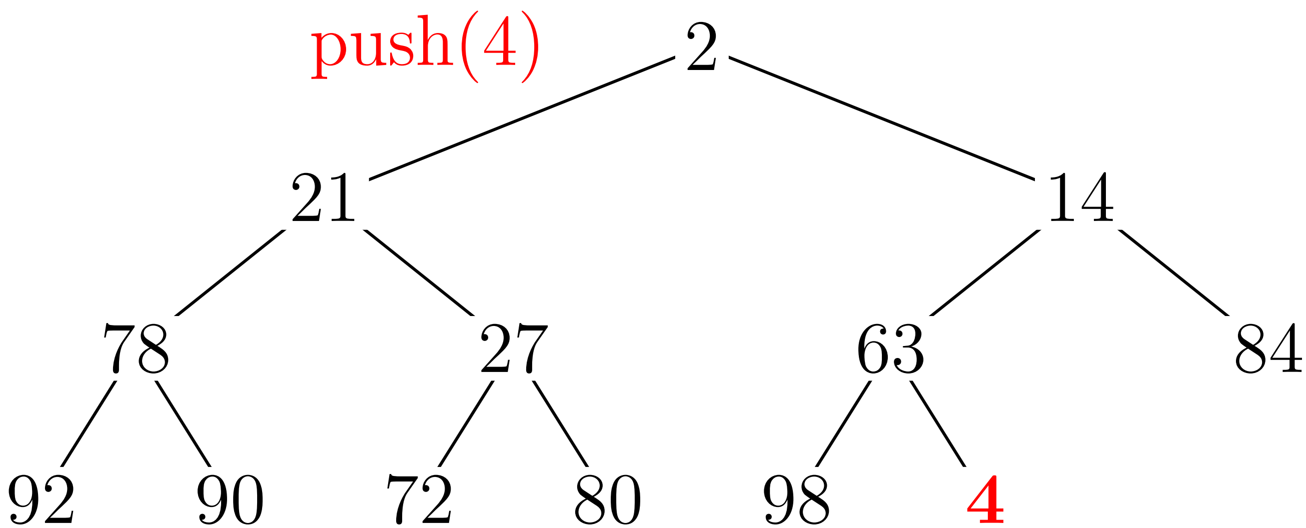


Adding an Element

2	21	14	78	27	63	84	92	90	72	80	98	**4**	



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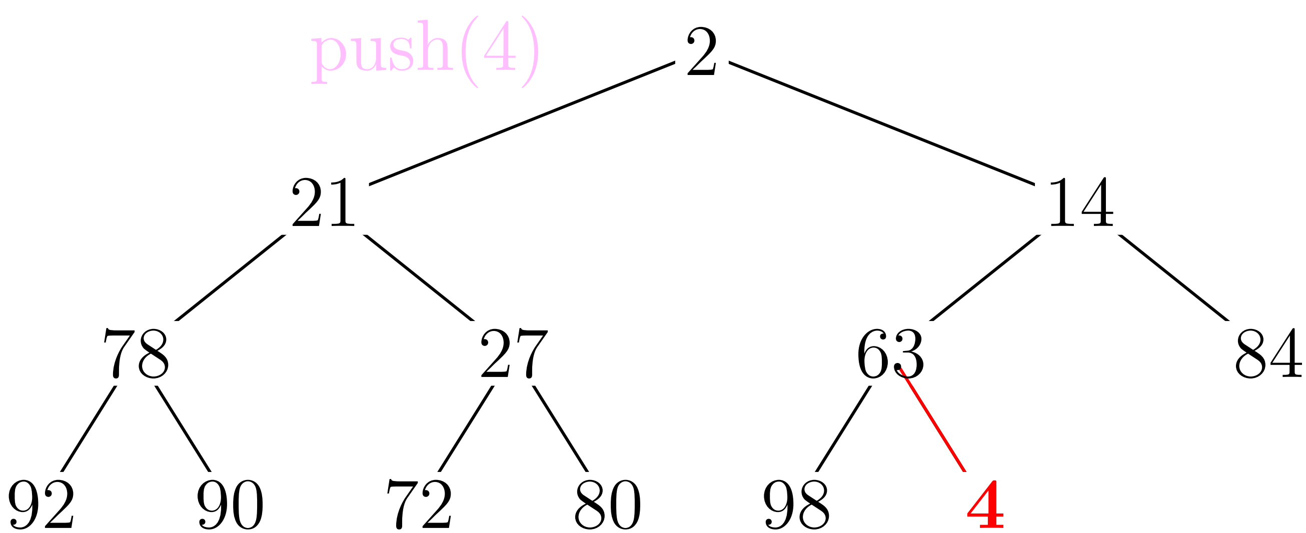

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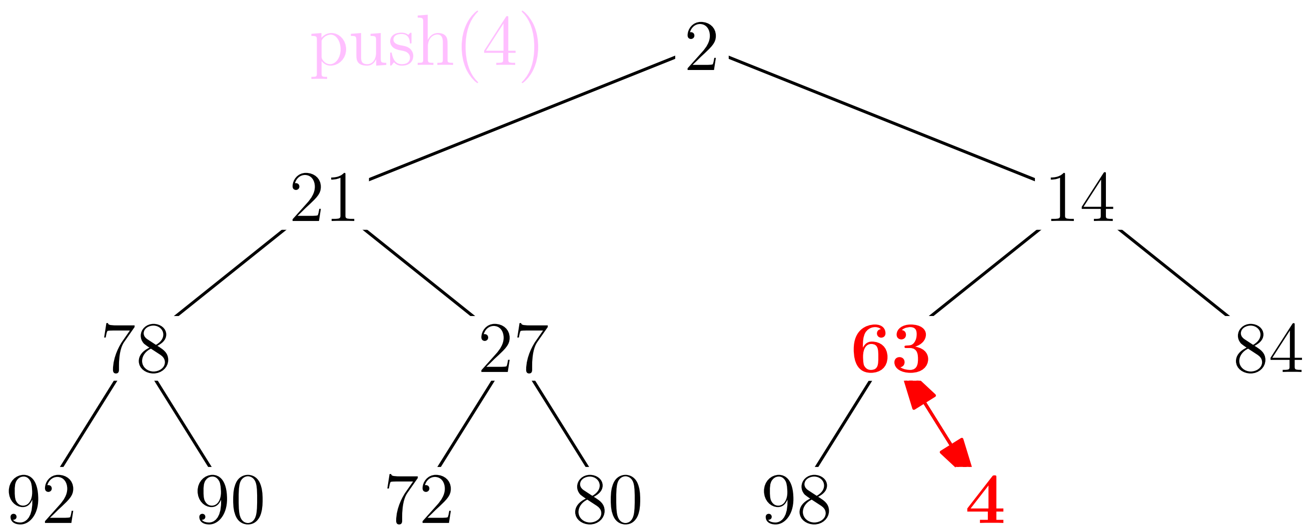

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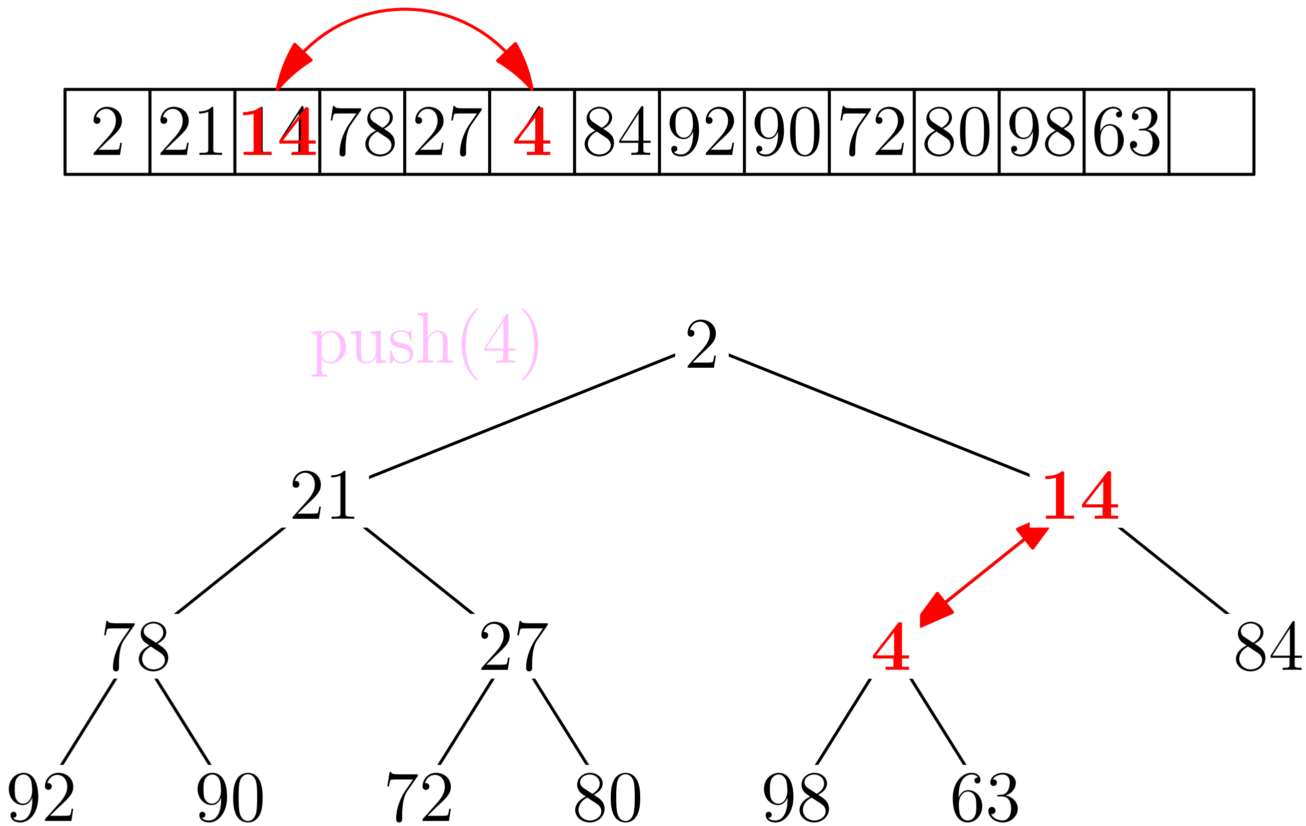

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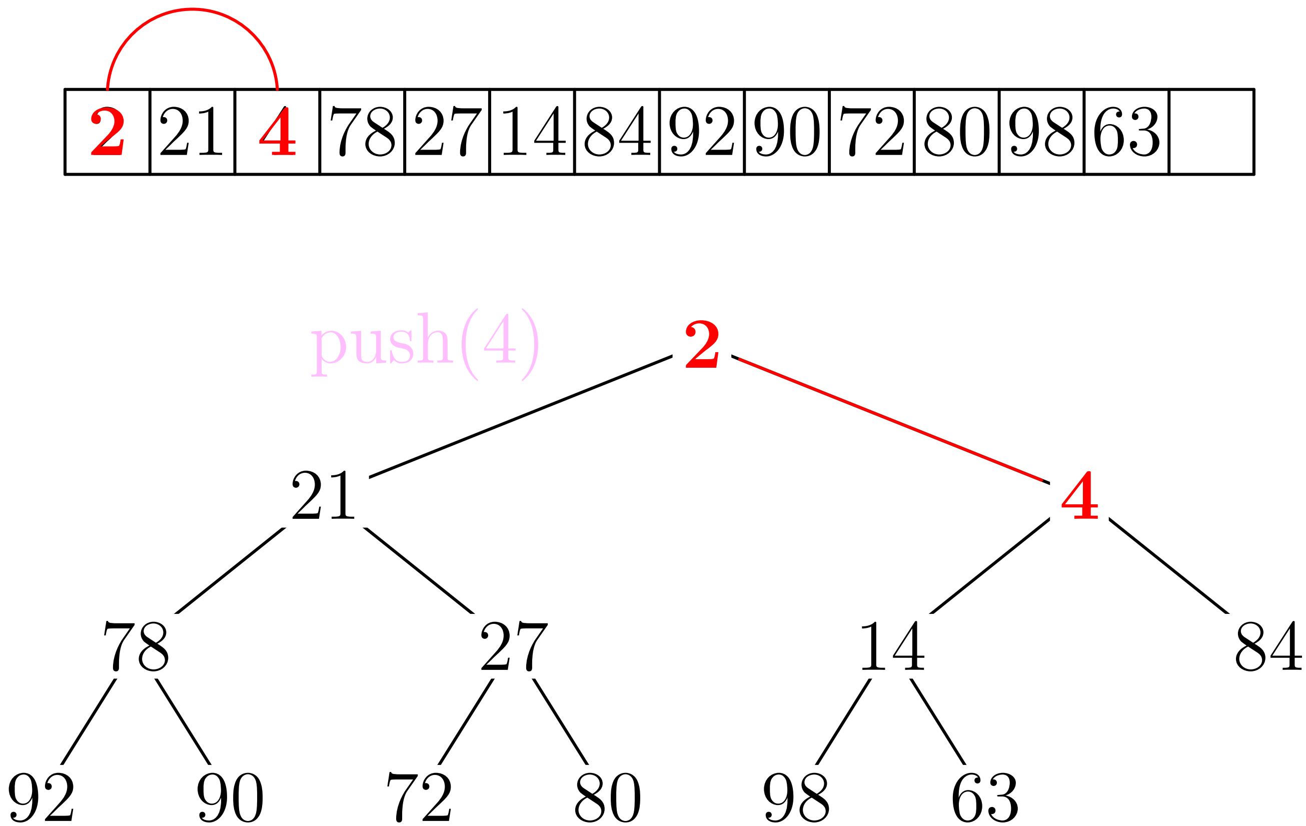

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void push(T value, P priority) {
  pair<T,P> tmp(value, priority);
  array.push_back(tmp);
  unsigned child = size() - 1;

  /* Percolate Up */
  while(child!=0) {
    unsigned parent = (child-1)>>1;  // floor((child-1)/2)
    if (array[parent].second < array[child].second)
      return;
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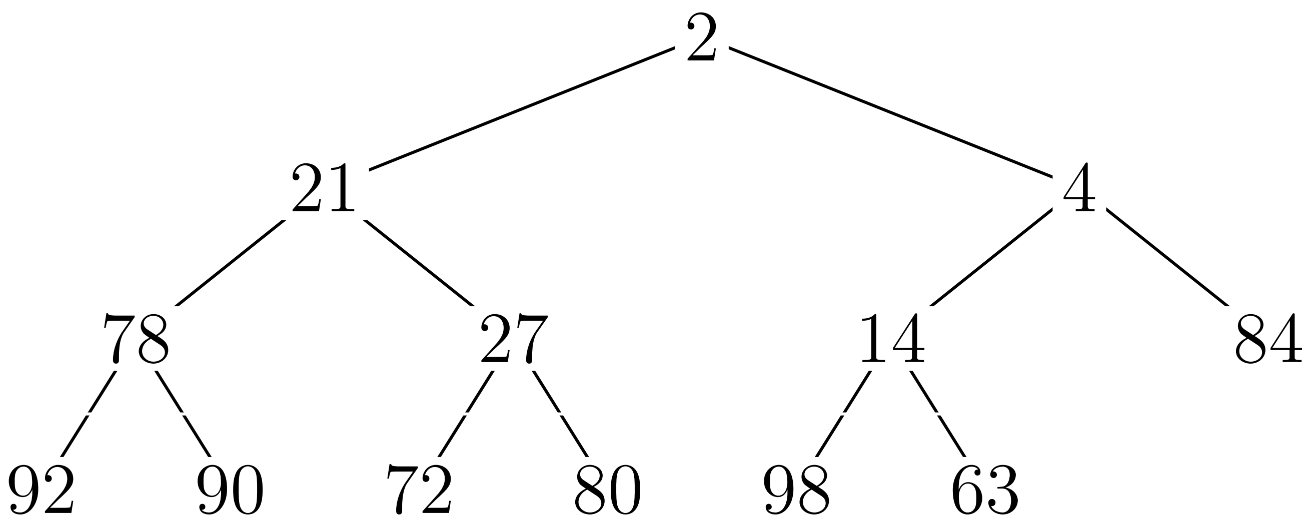
2	21	4	78	27	14	84	92	90	72	80	98	63	



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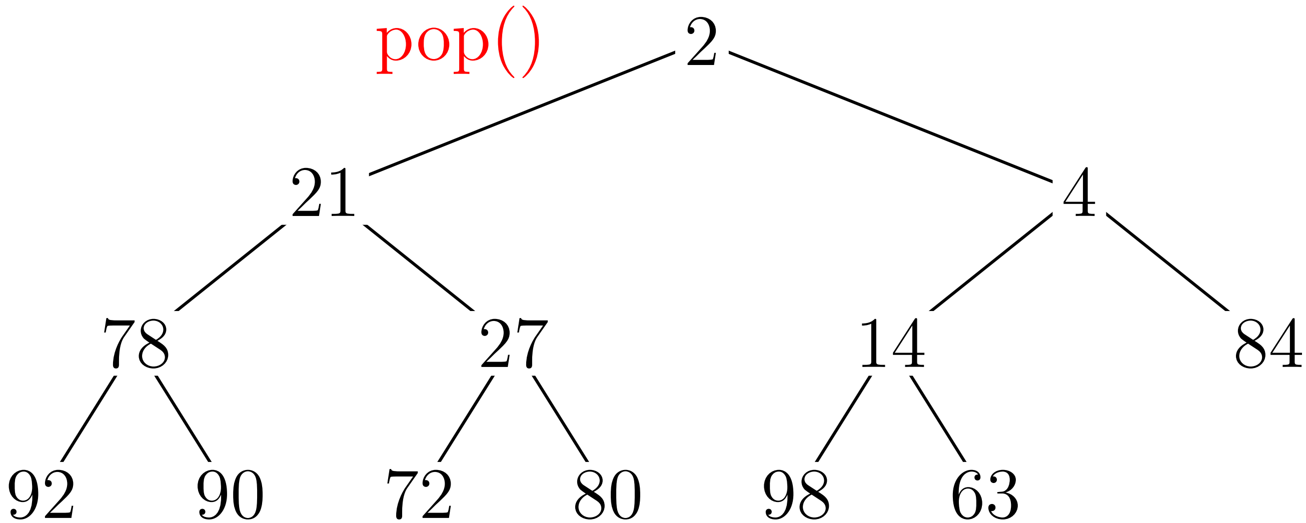

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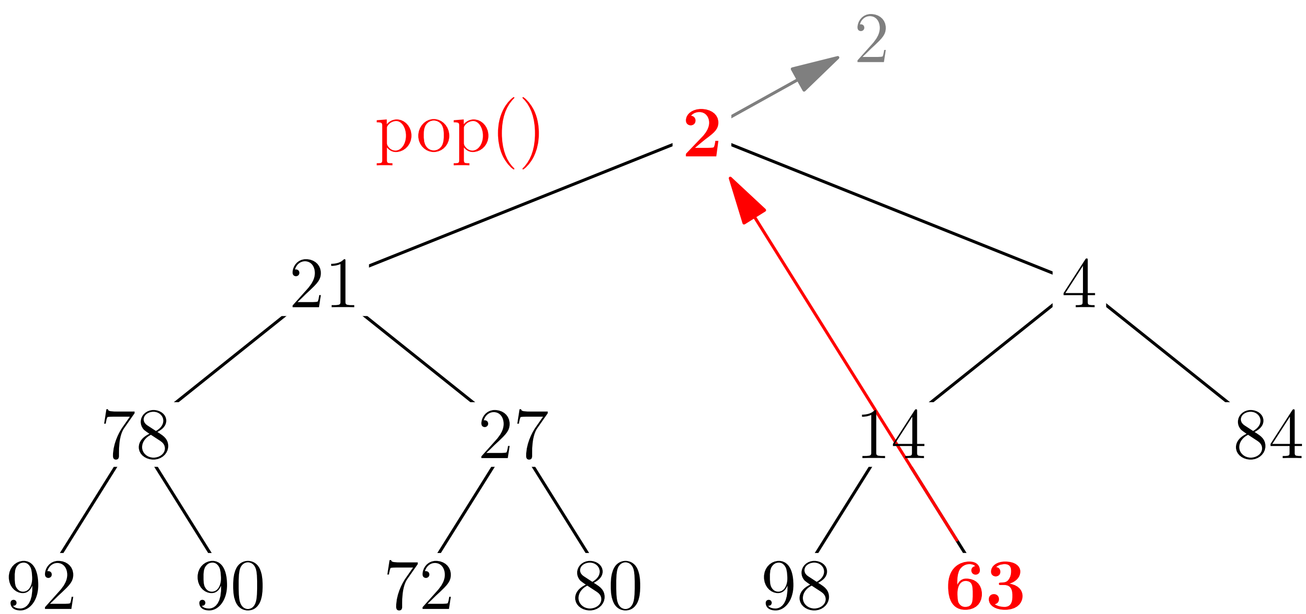
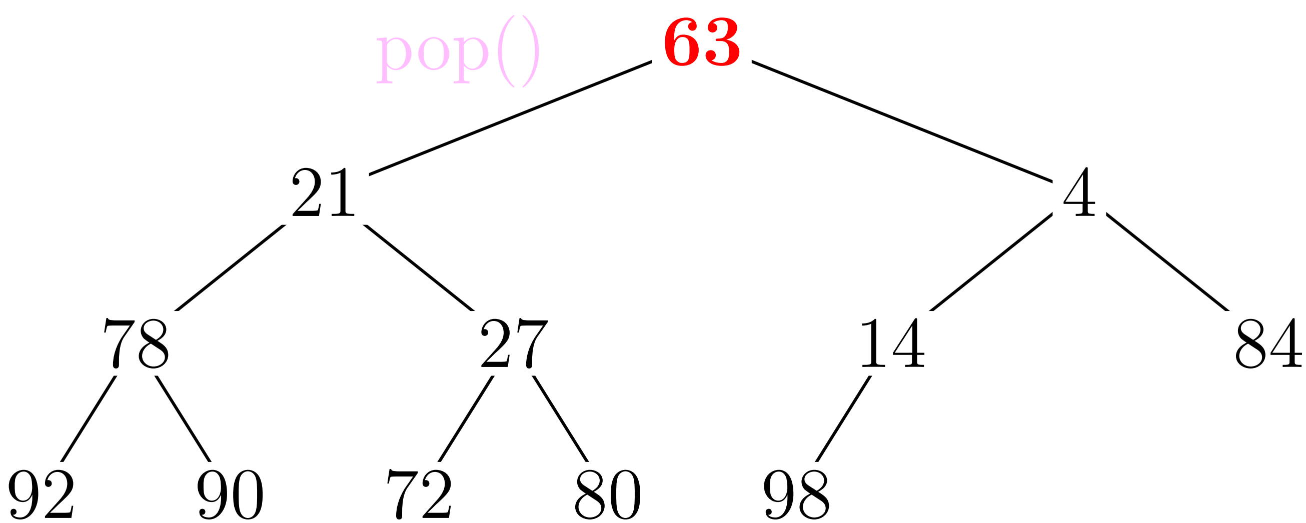
2	21	4	78	27	14	84	92	90	72	80	98	**63**	



Popping the Top

63	21	4	78	27	14	84	92	90	72	80	98		



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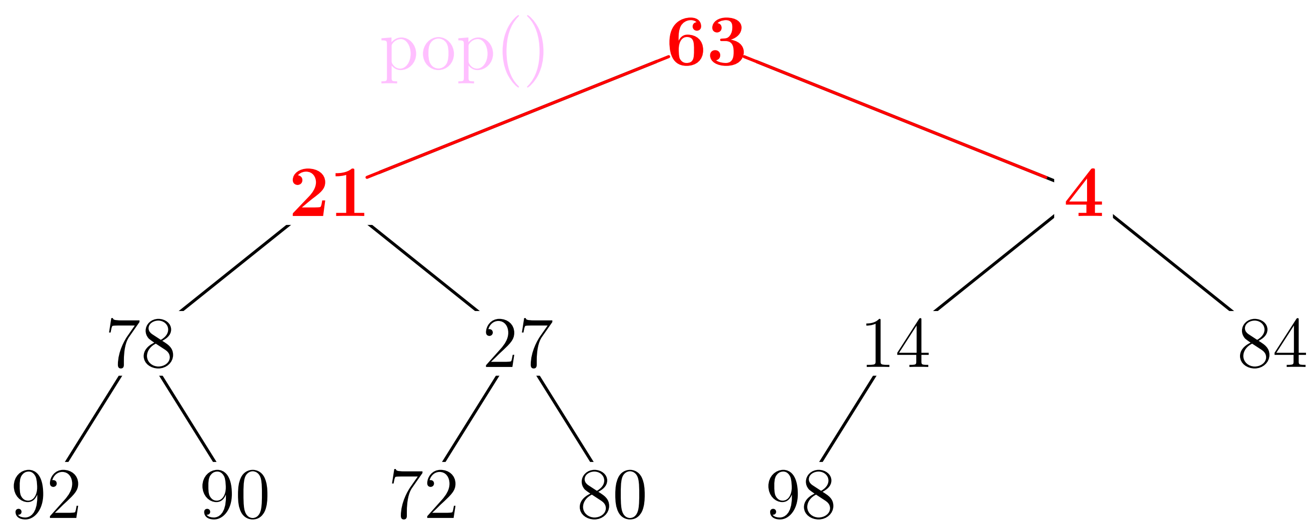

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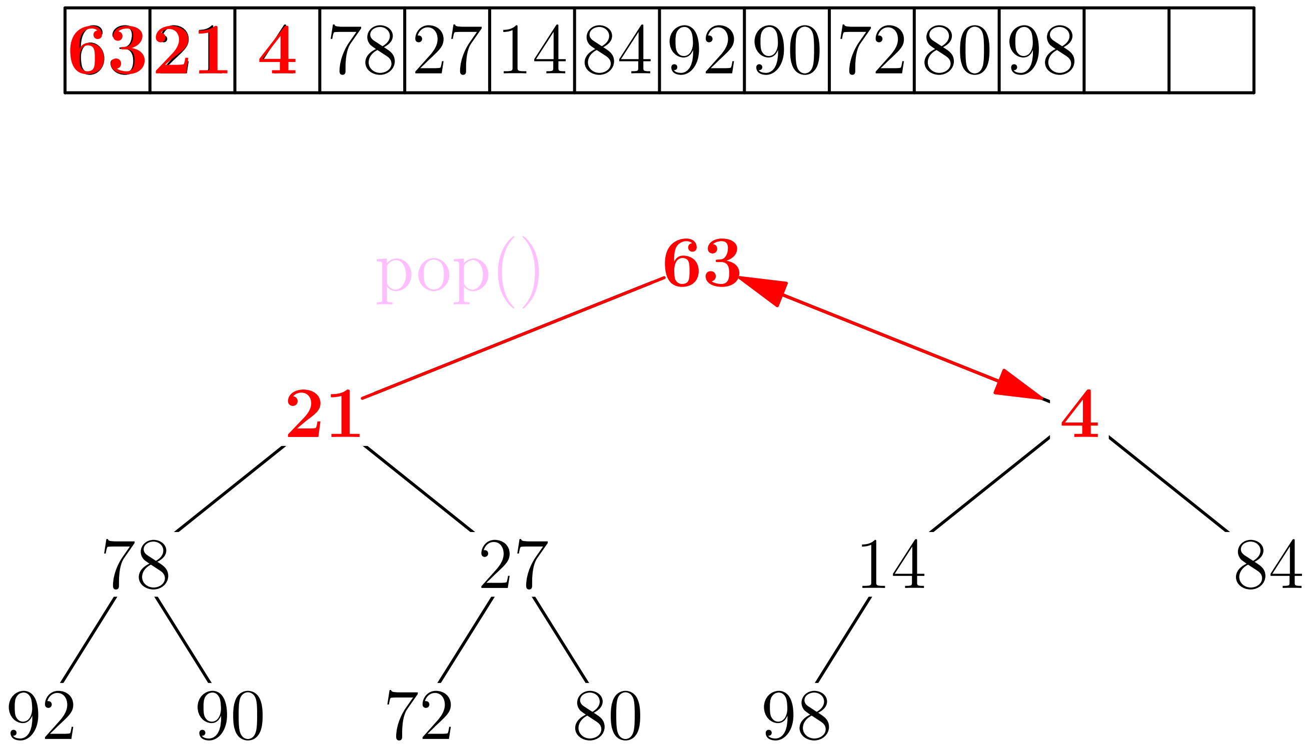

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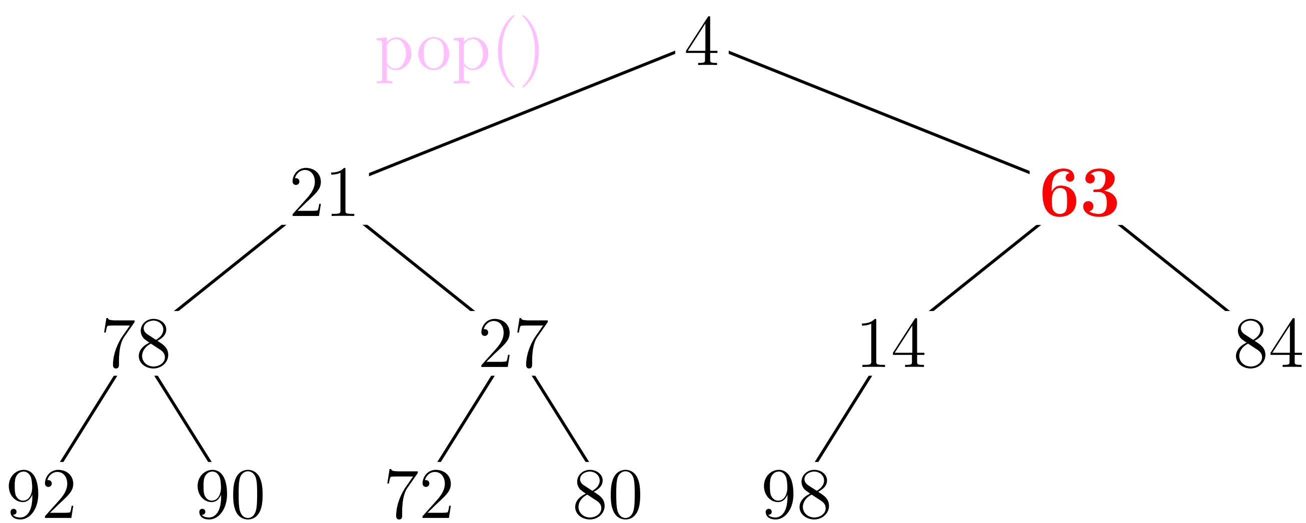

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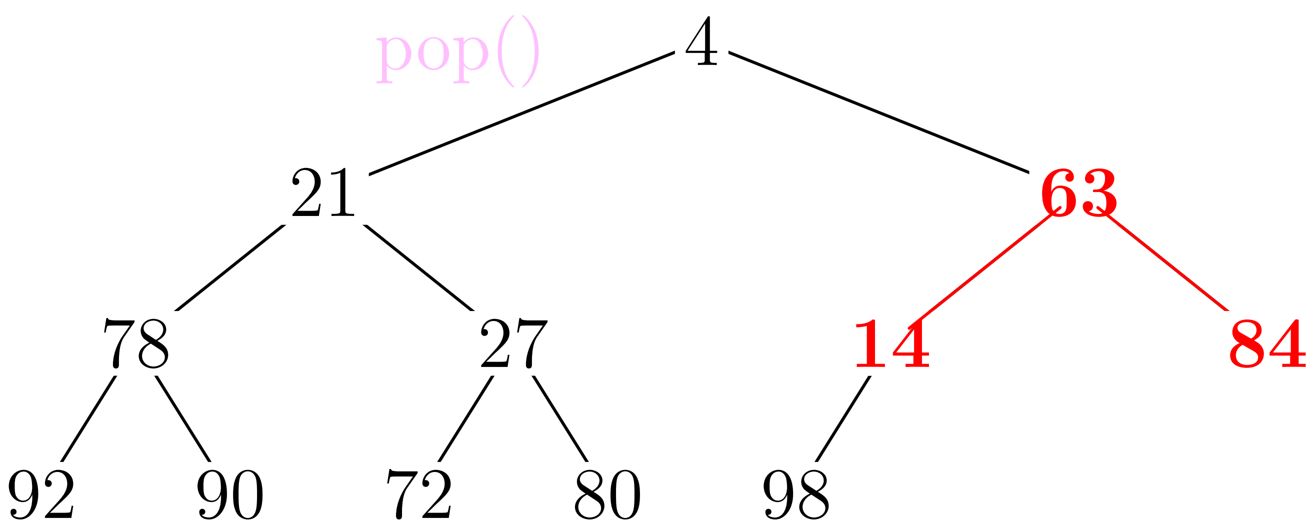

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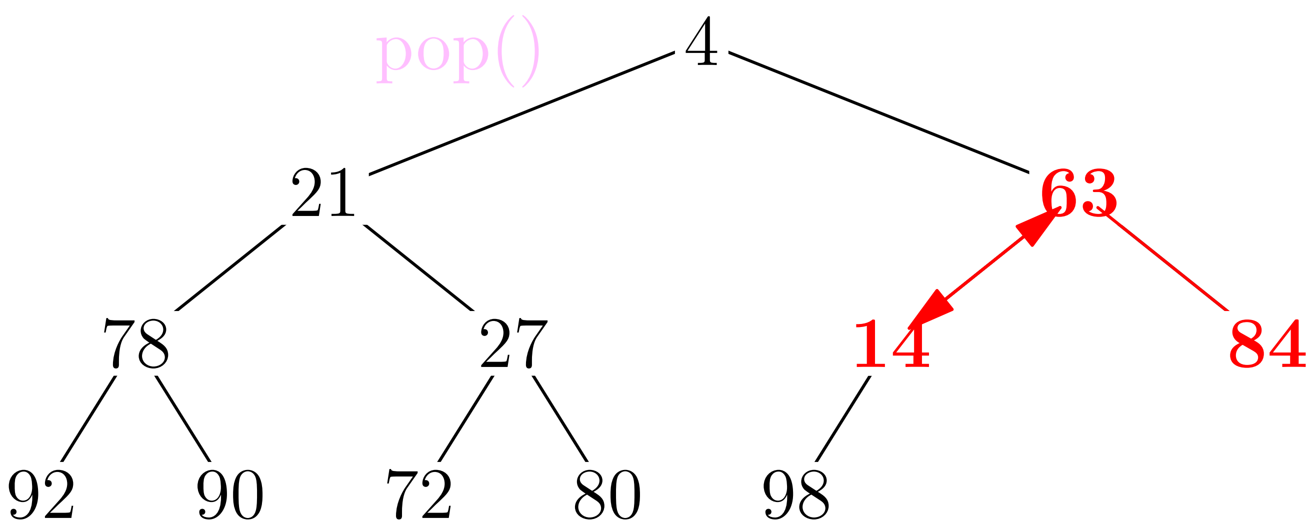

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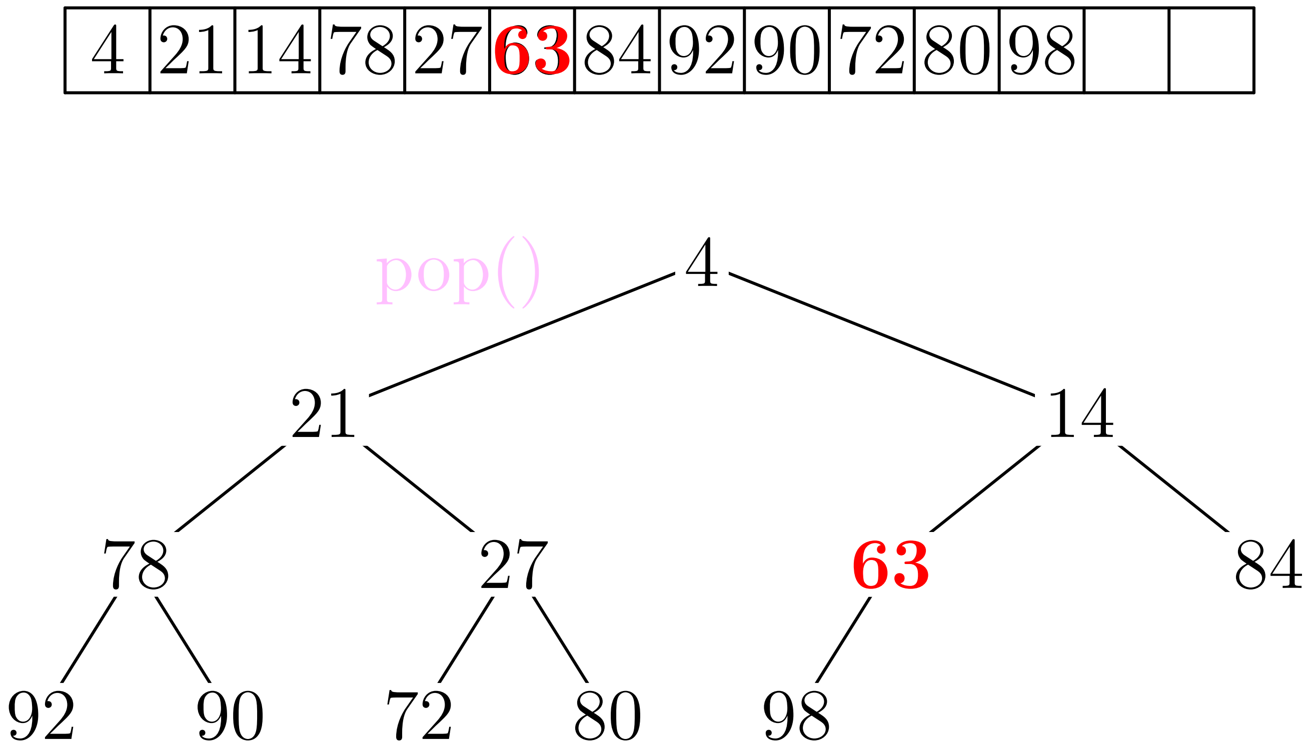

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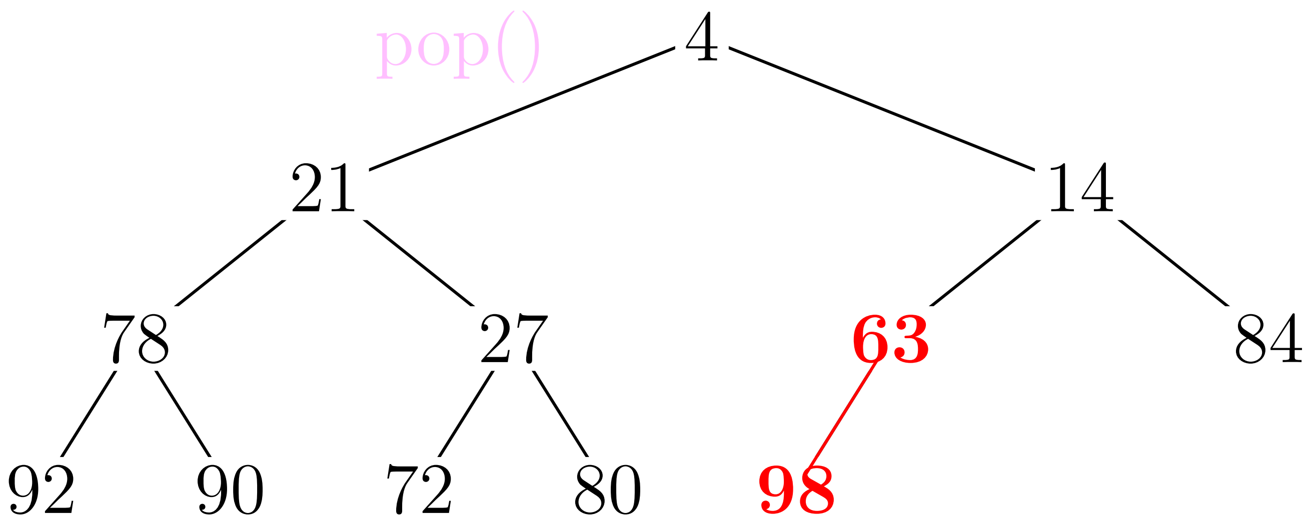

Popping the Top

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void pop() {
unsigned parent = 0;
pair<T, P> tmp = array.back();
array[0] = tmp;
array.pop_back();
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while(child<size()) {
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    ++child;
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  child = 2*parent + 1;
 }
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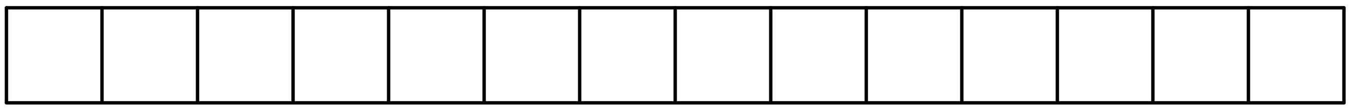
```

Heaps in Action

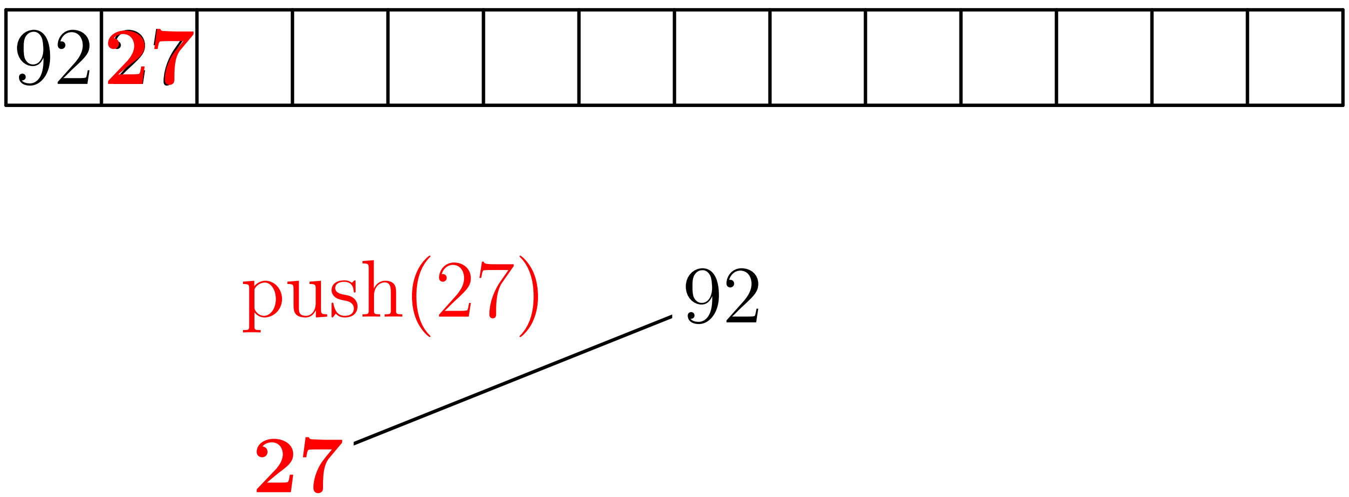


Heaps in Action

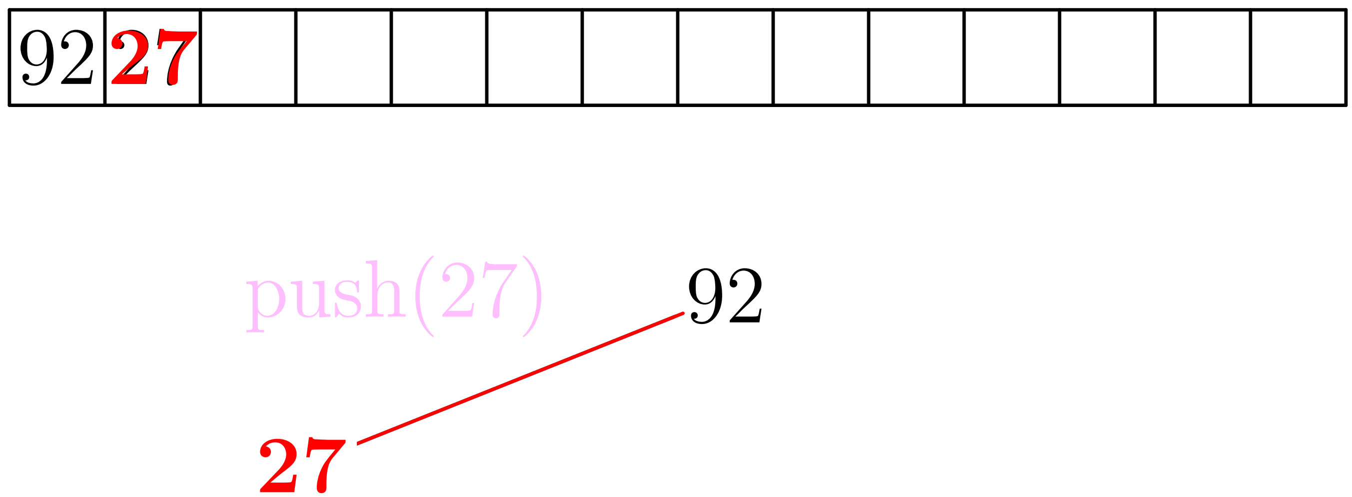
92													

push(92) **92**

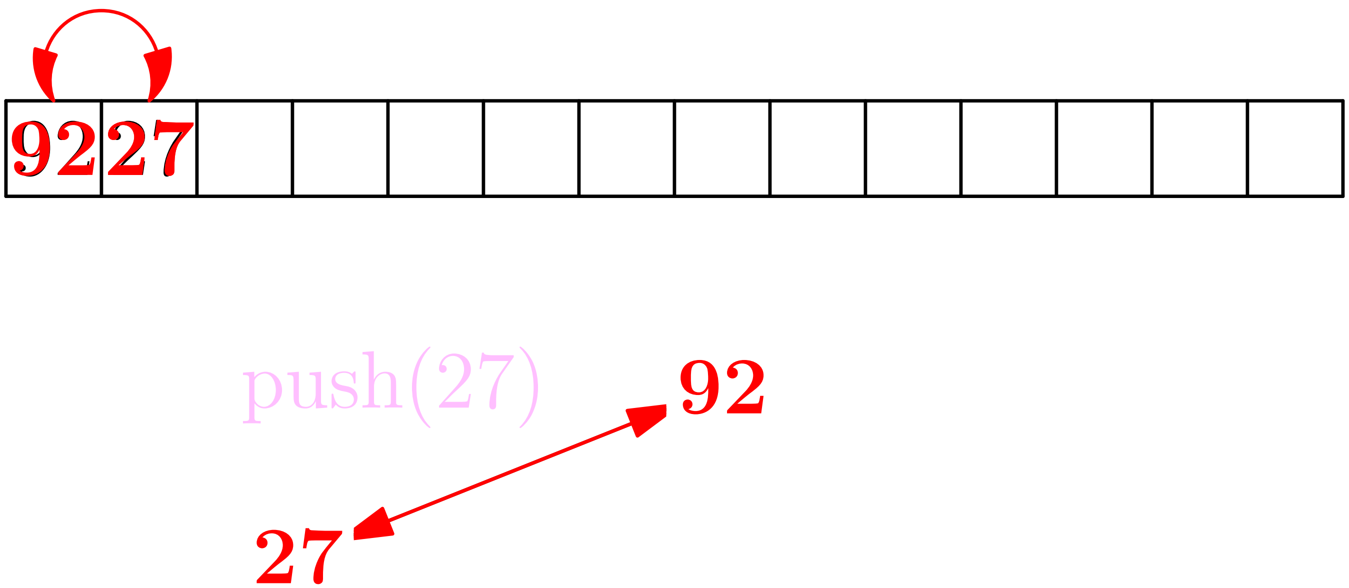
Heaps in Action



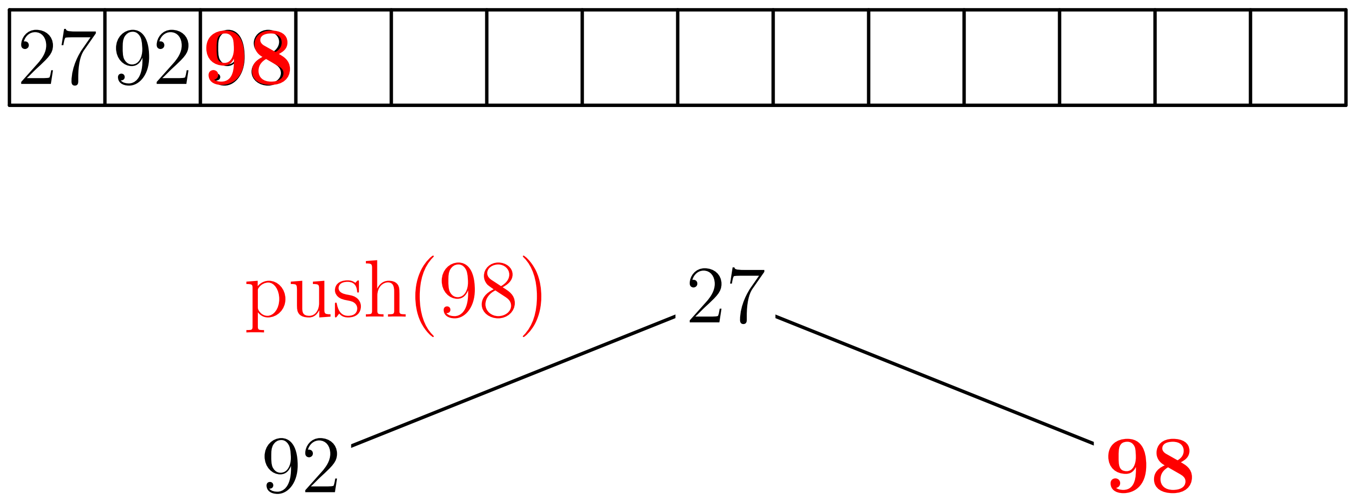
Heaps in Action



Heaps in Action



Heaps in Action



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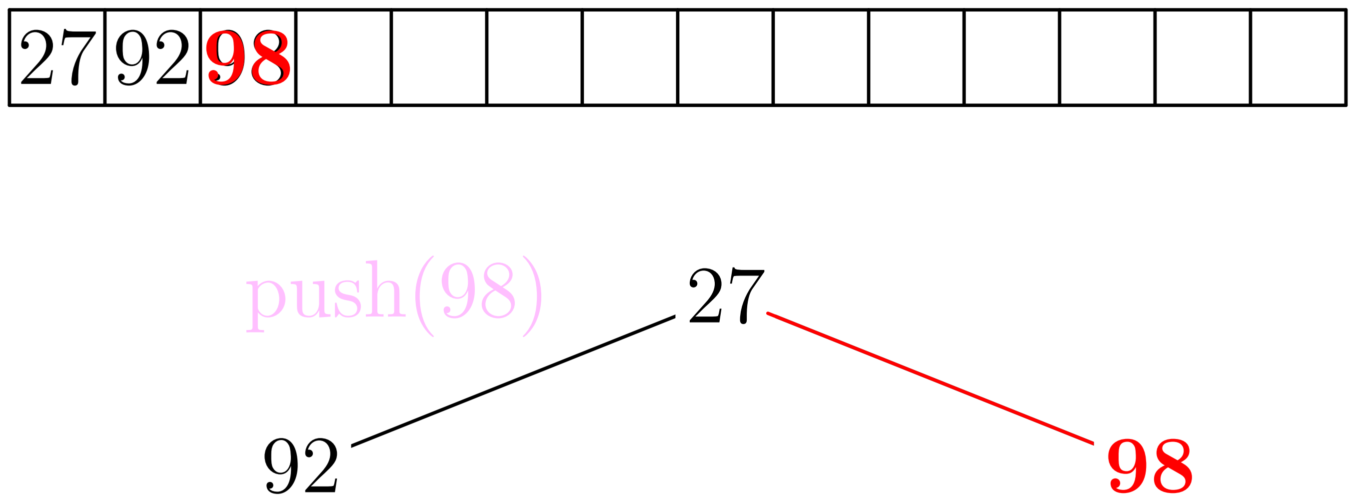

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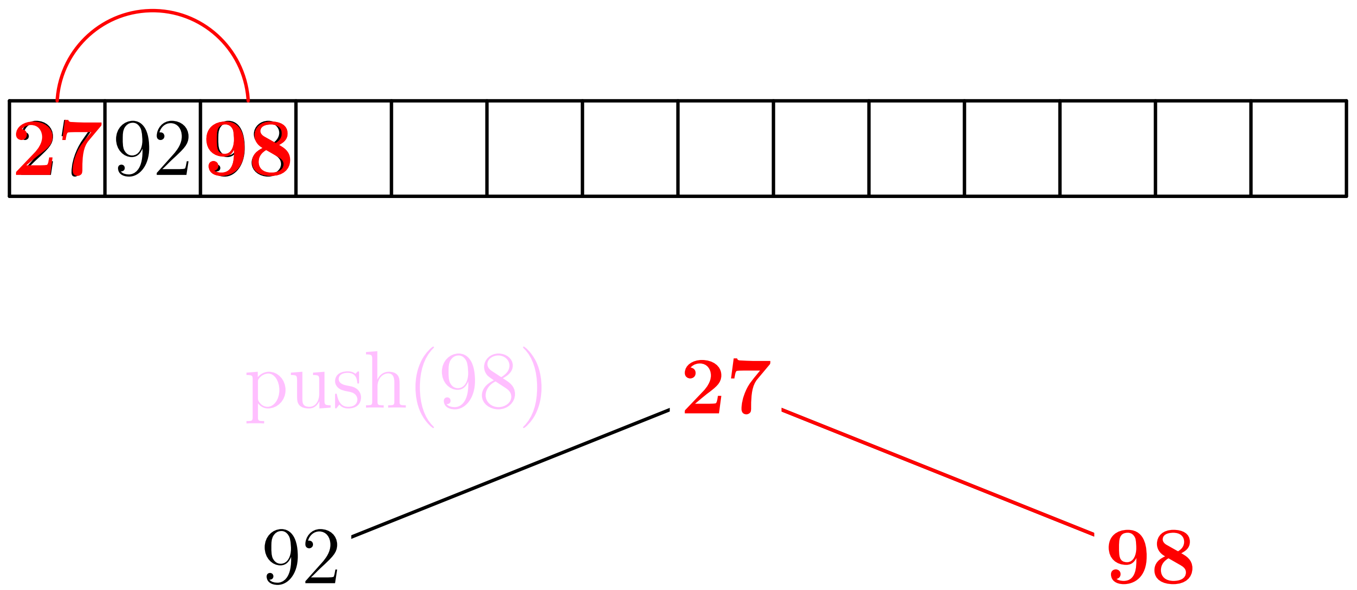

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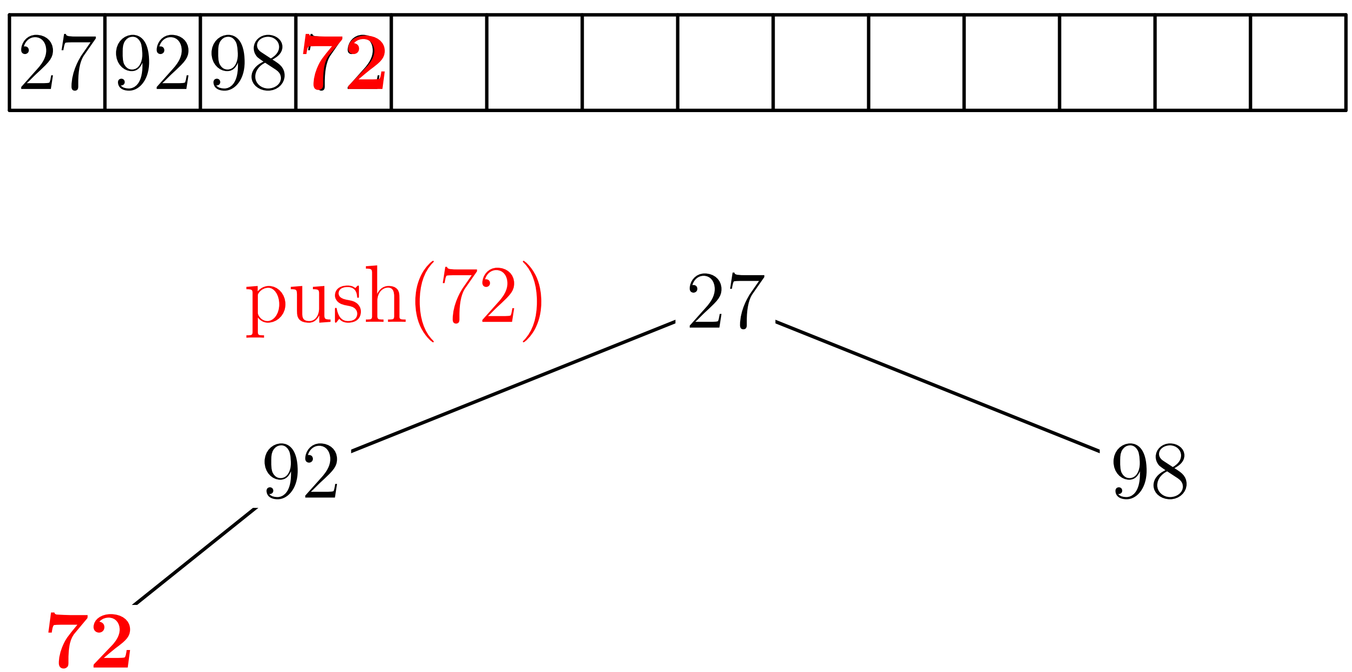

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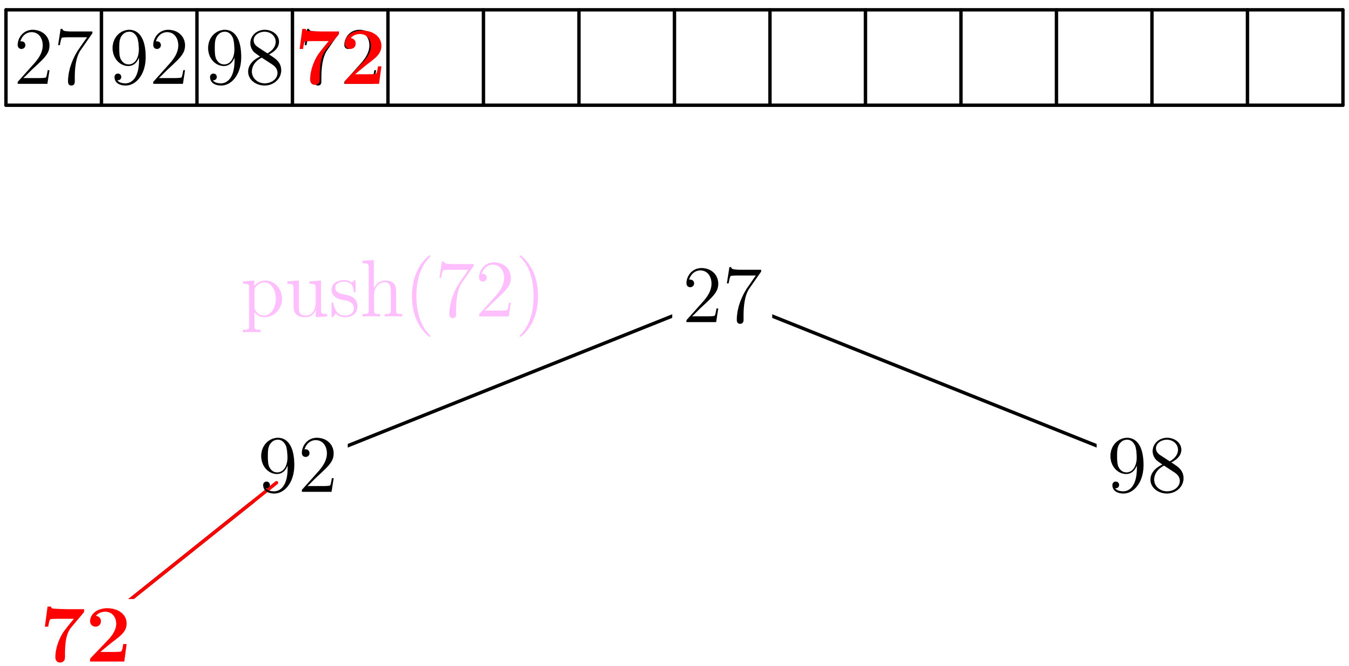

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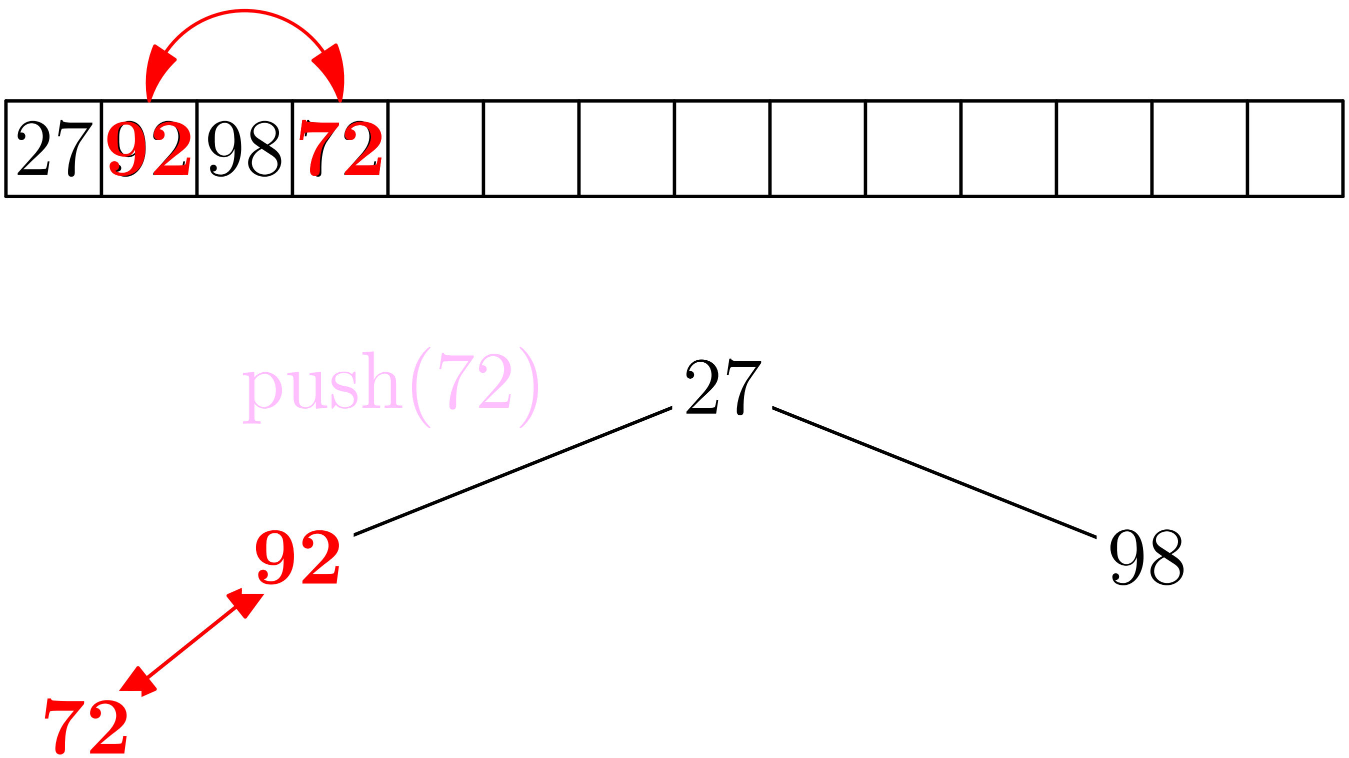

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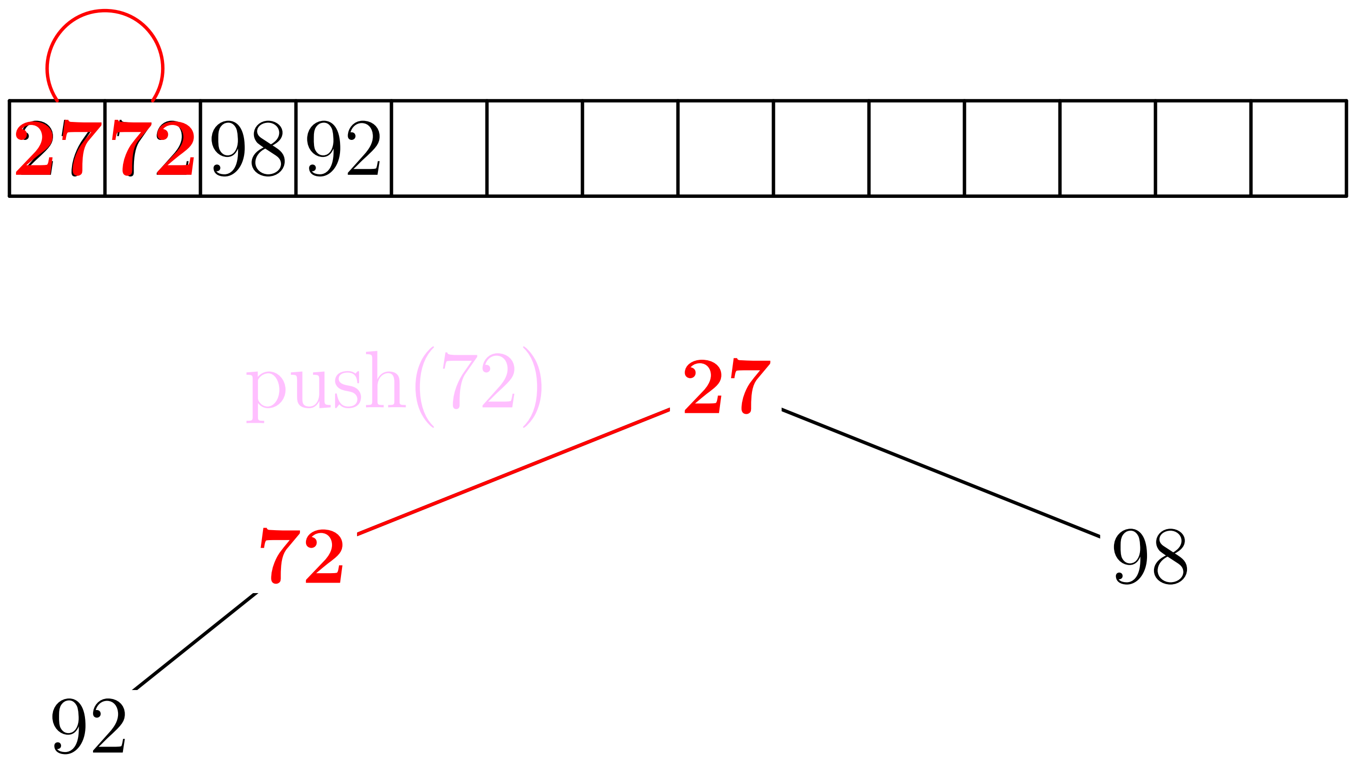
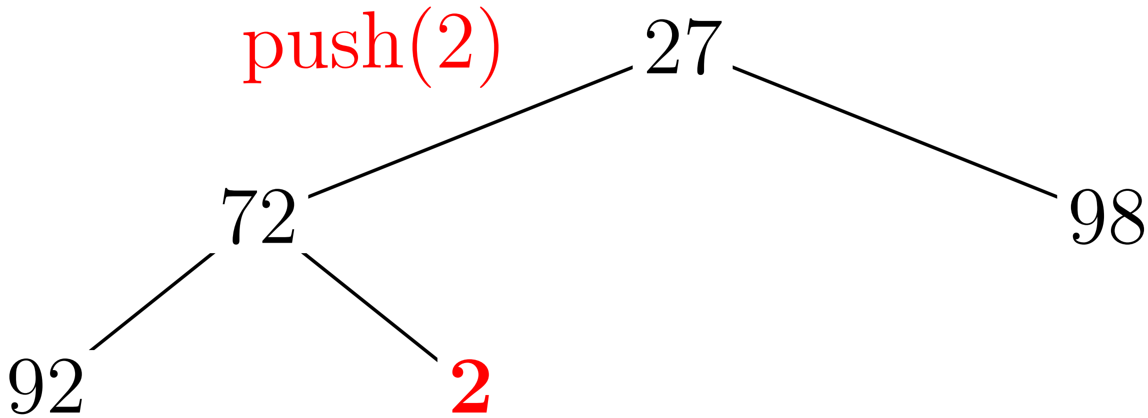


Heaps in Action

27	72	98	92	**2**									



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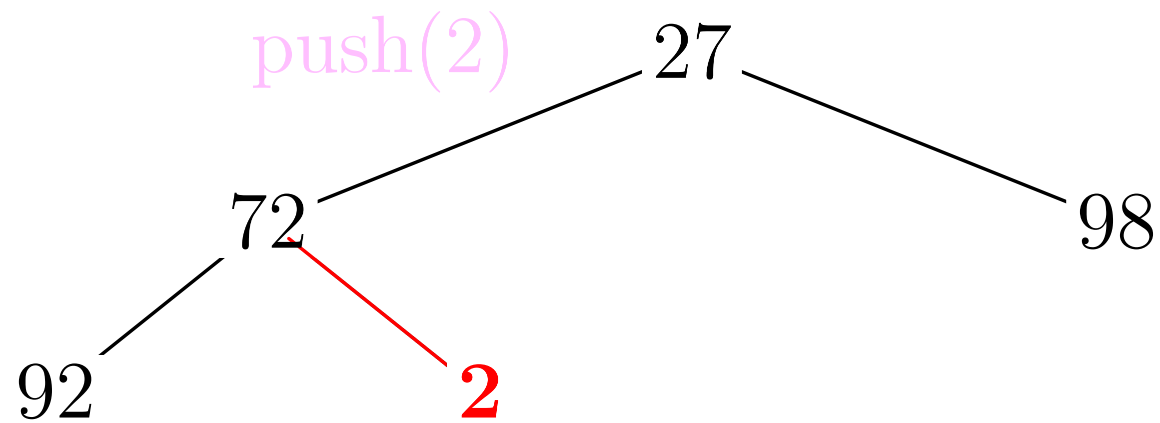

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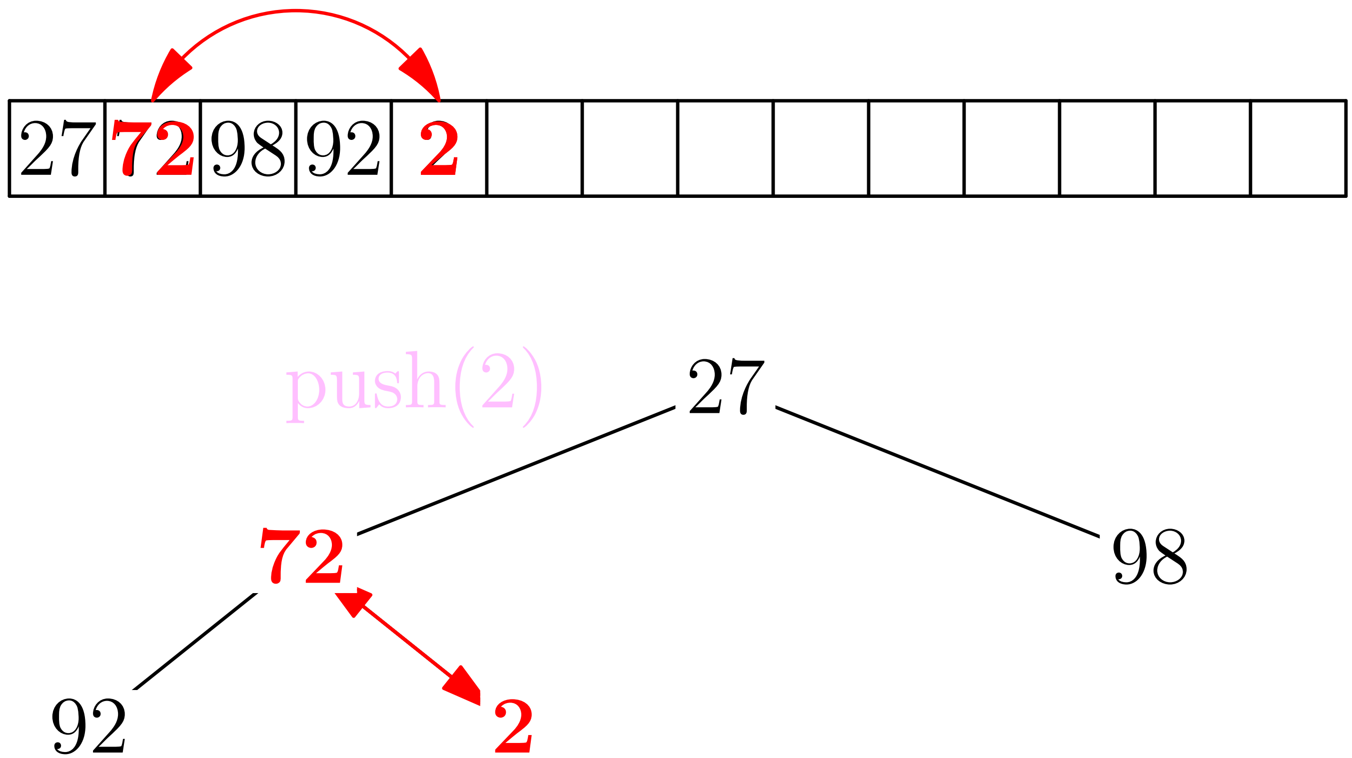

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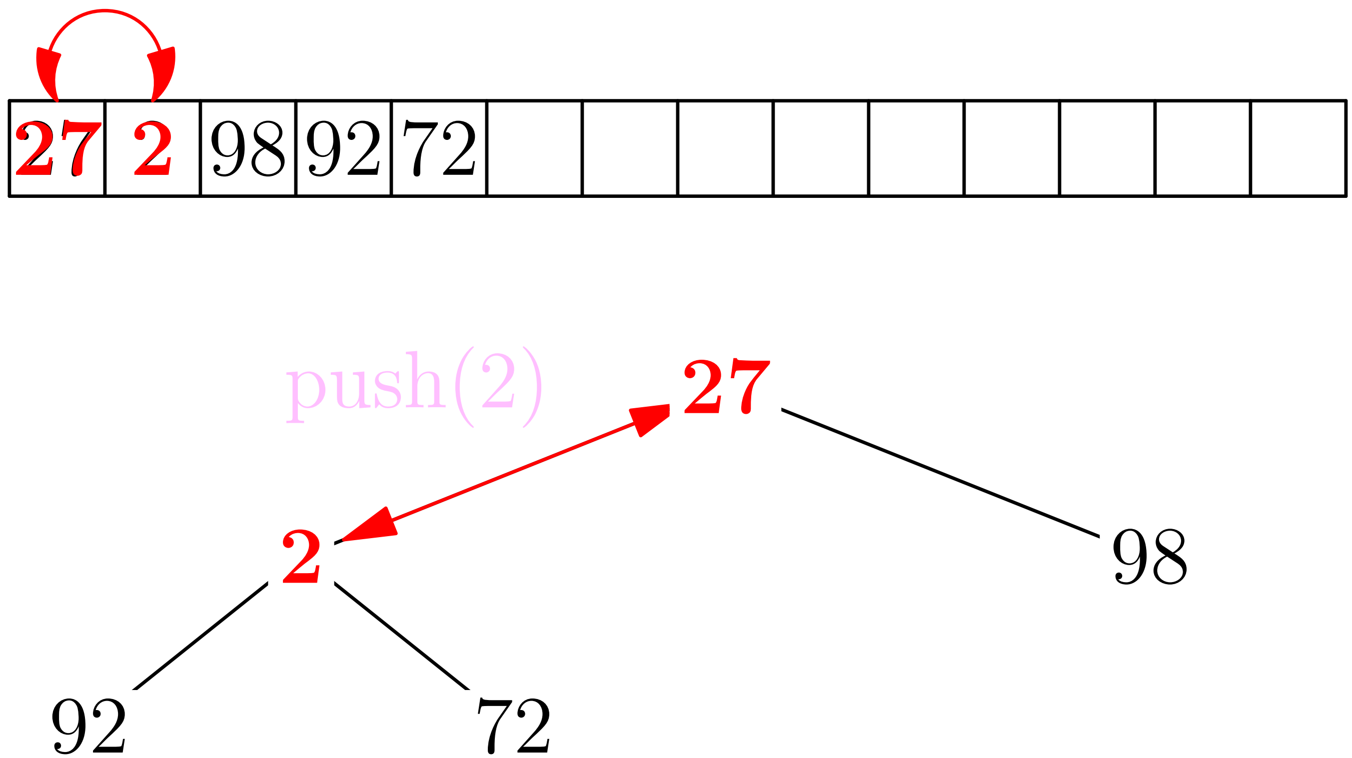

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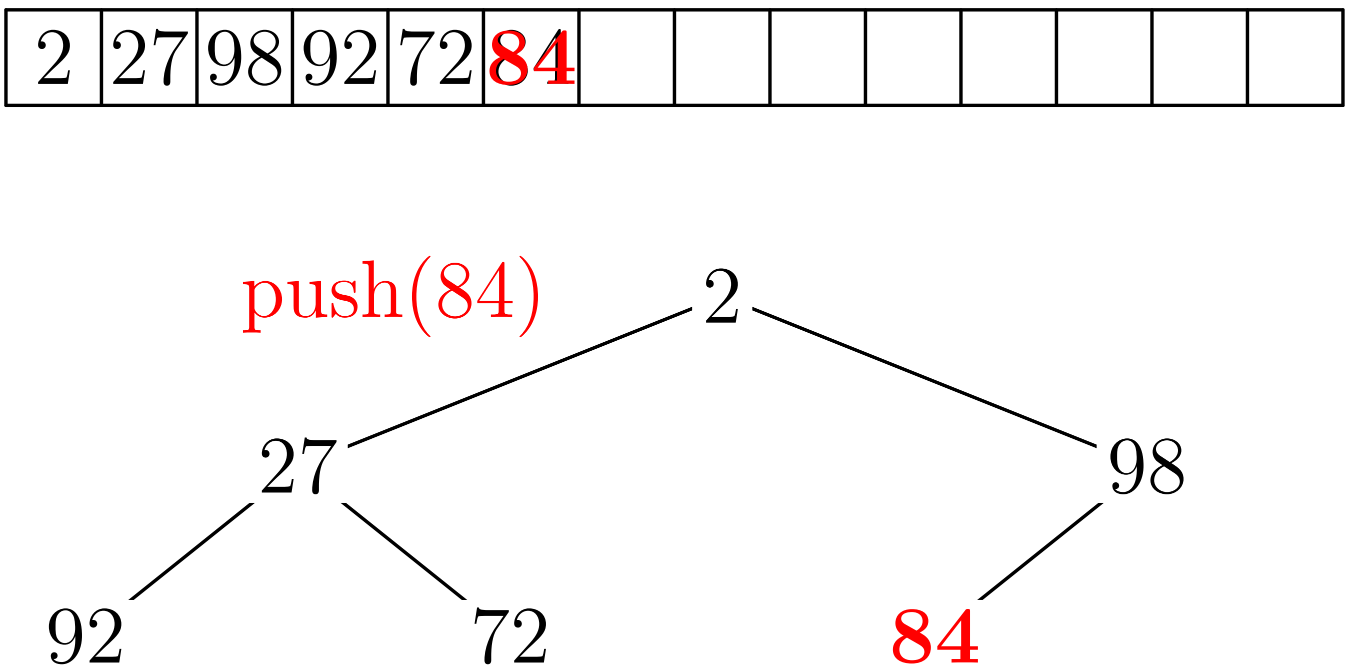

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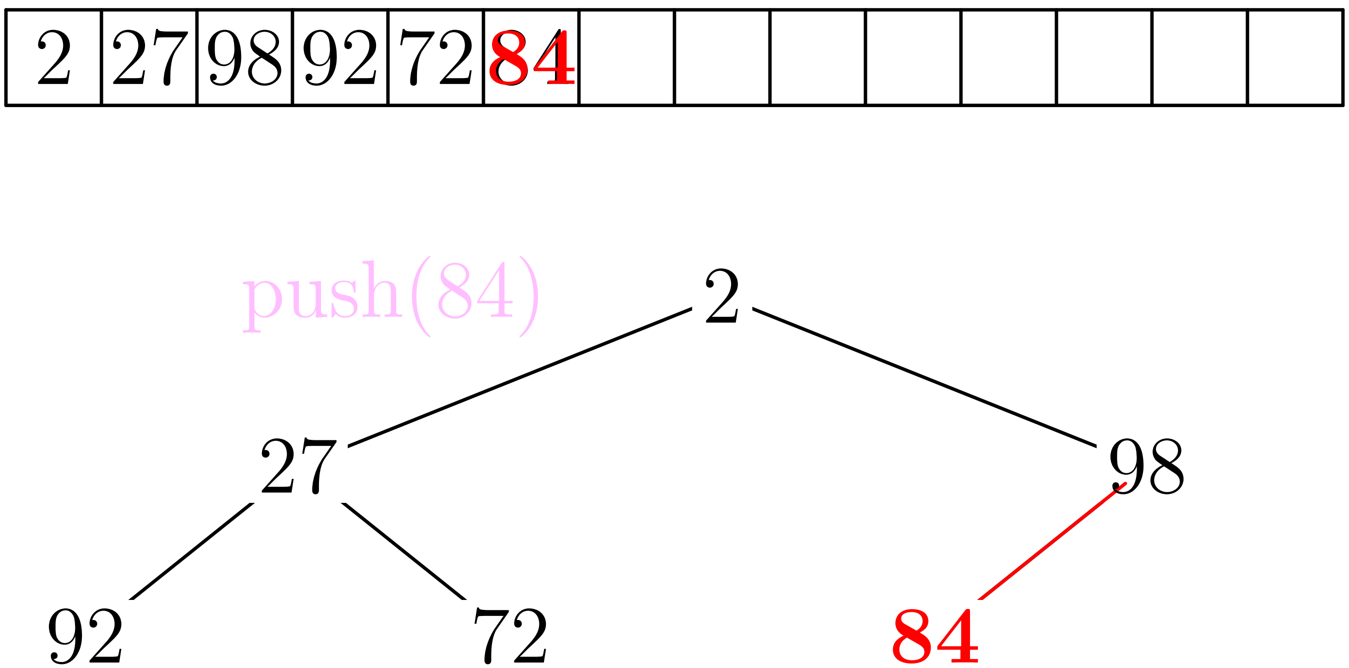

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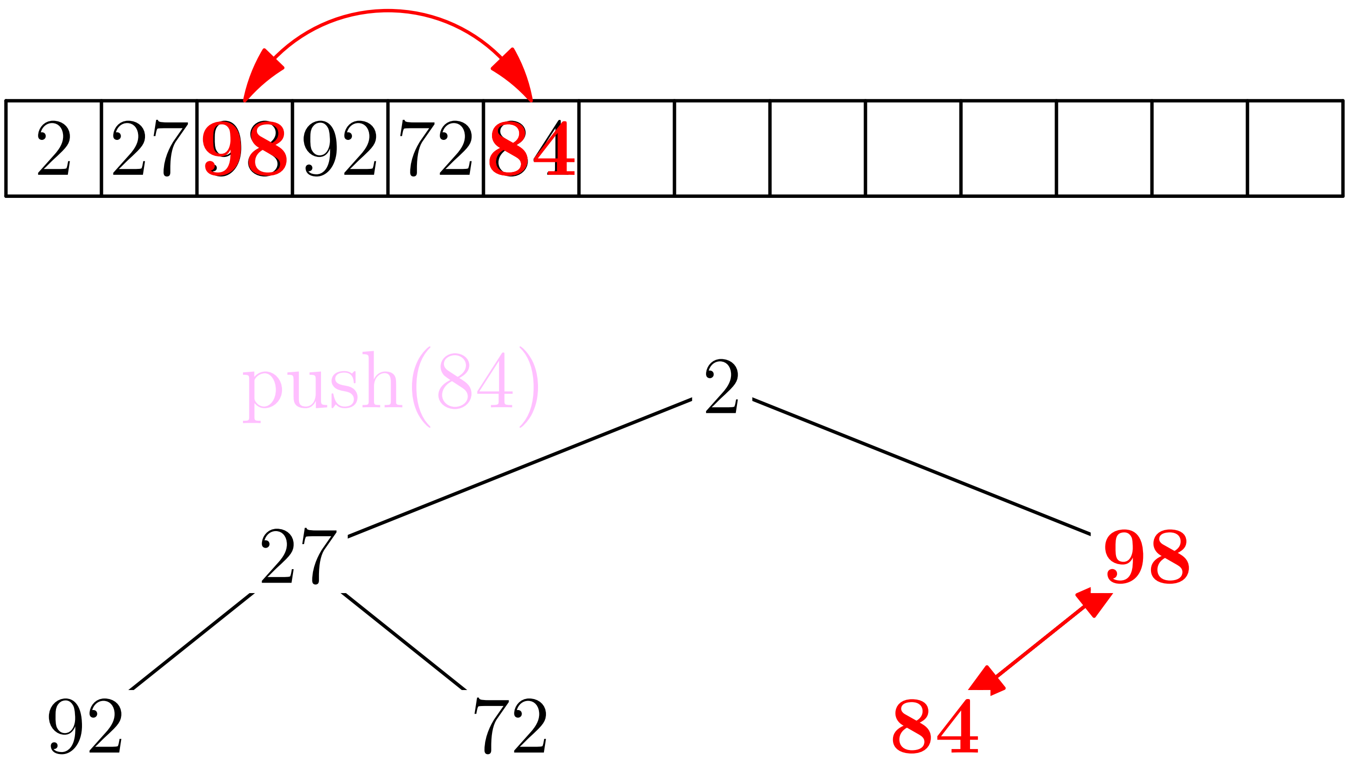

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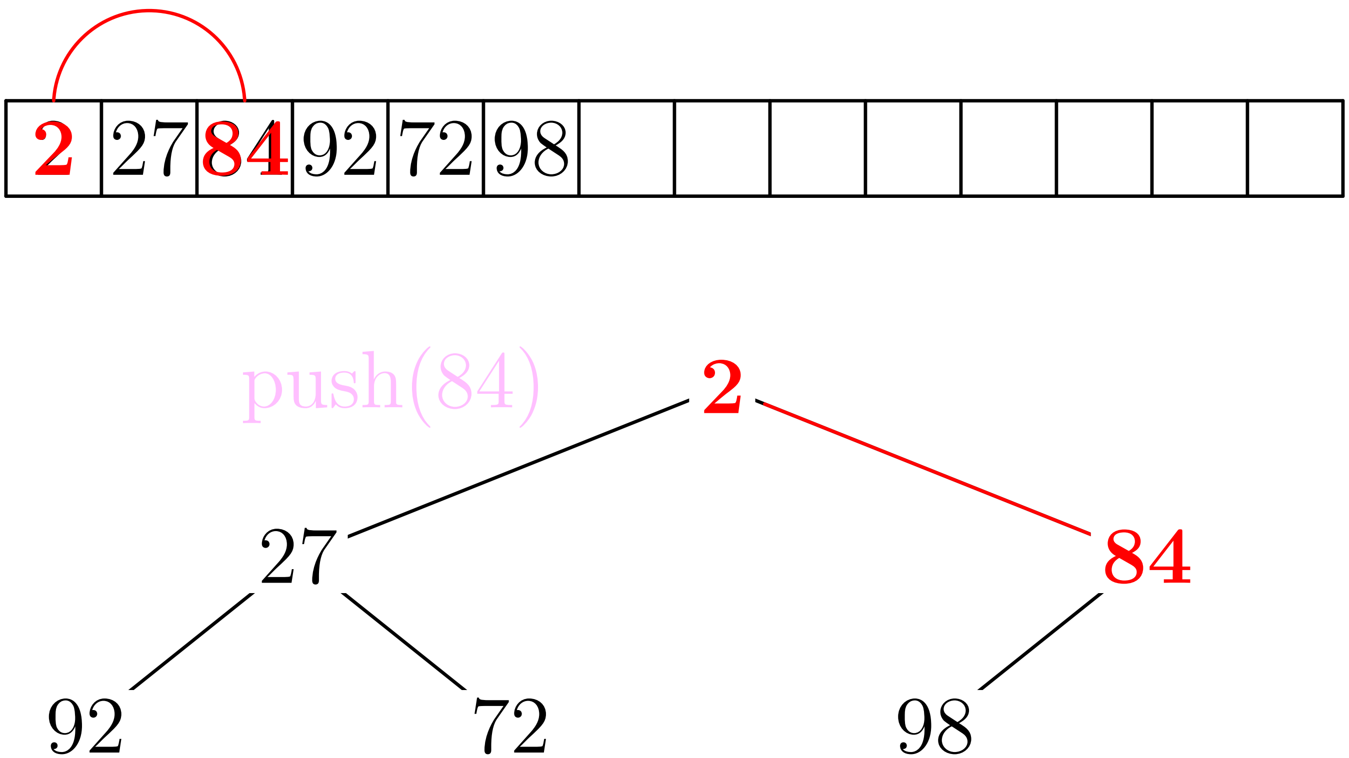
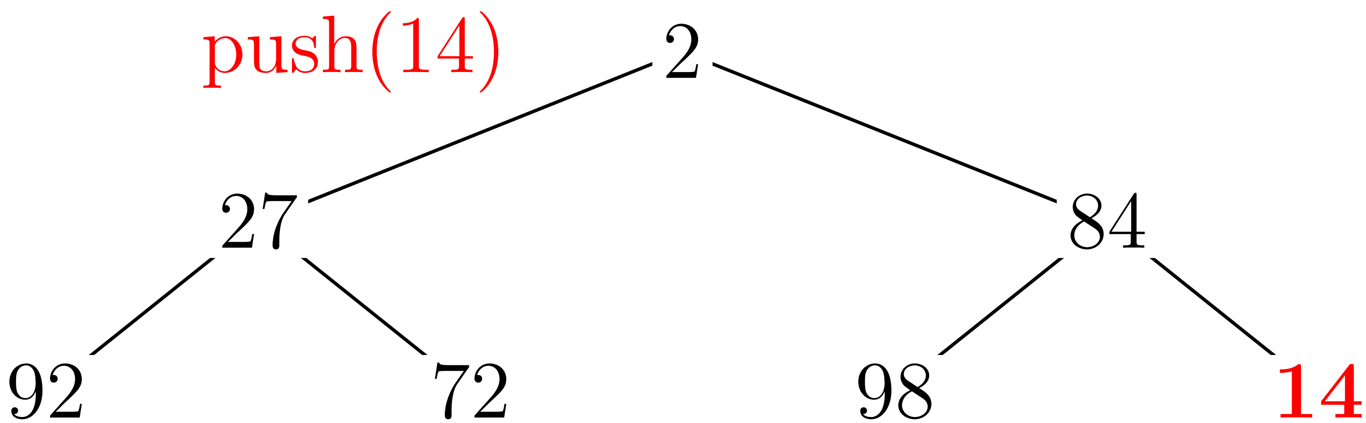


Heaps in Action

2	27	84	92	72	98	**14**							



Heaps in Action

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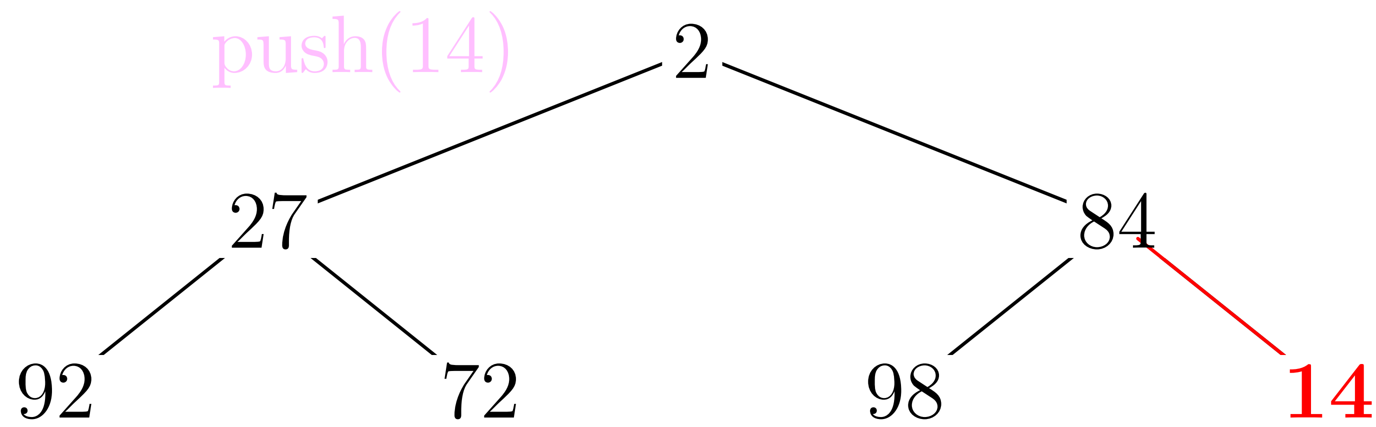
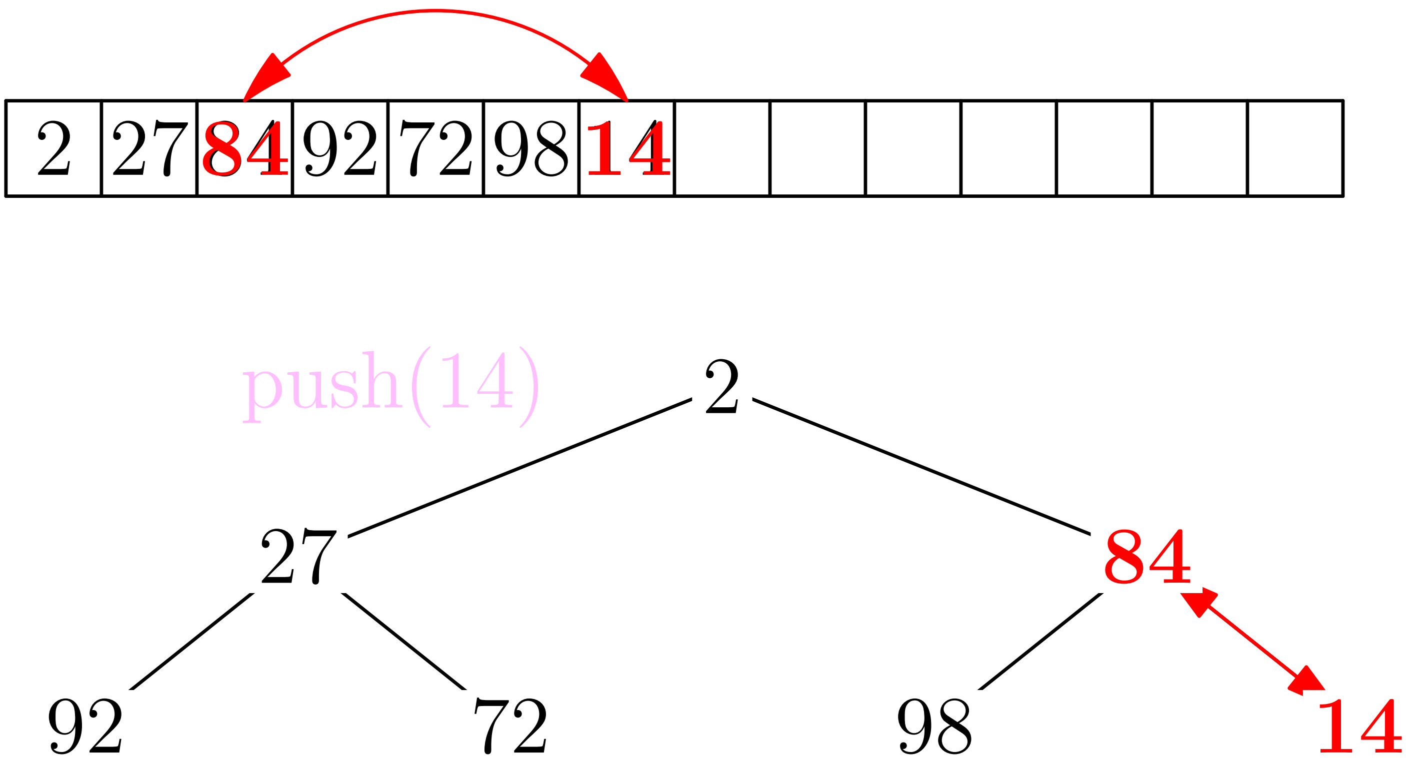


Heaps in Action



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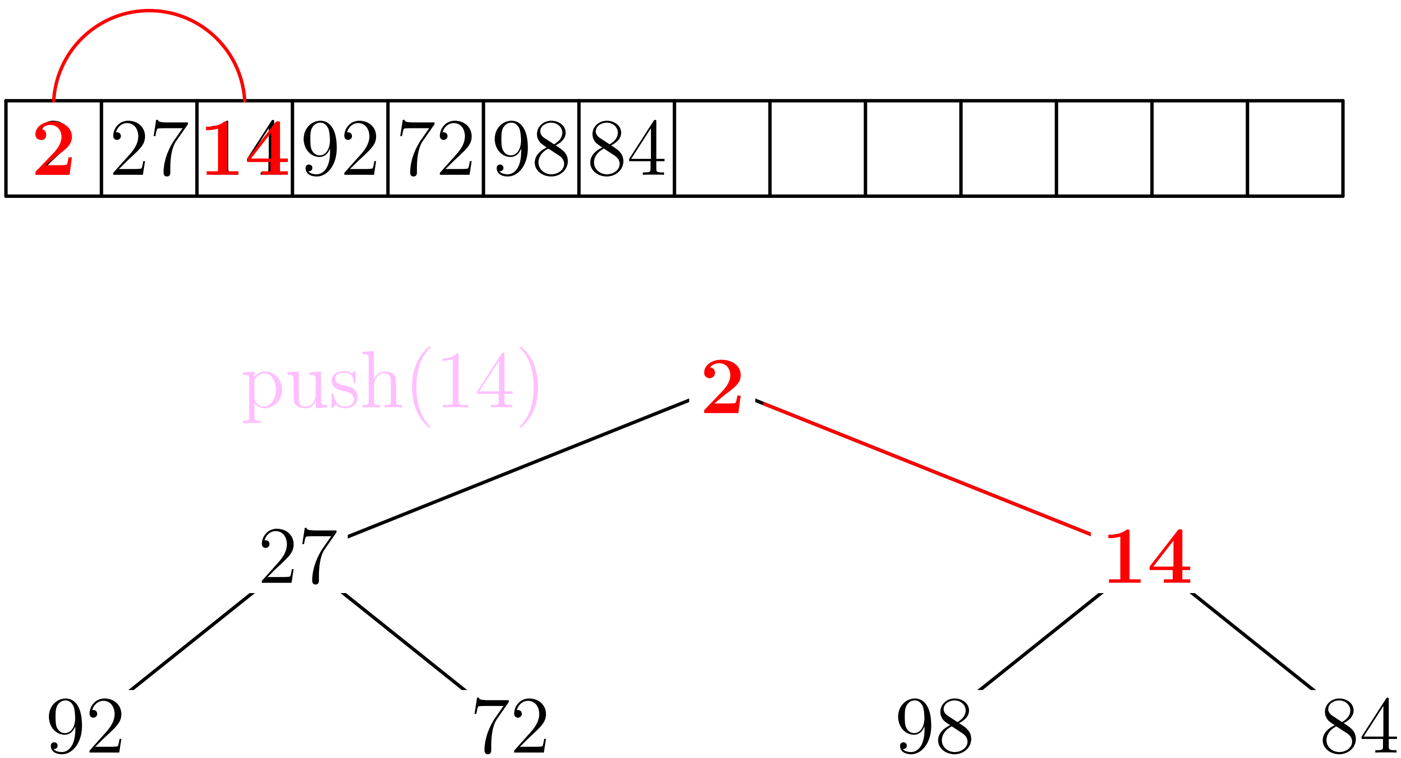
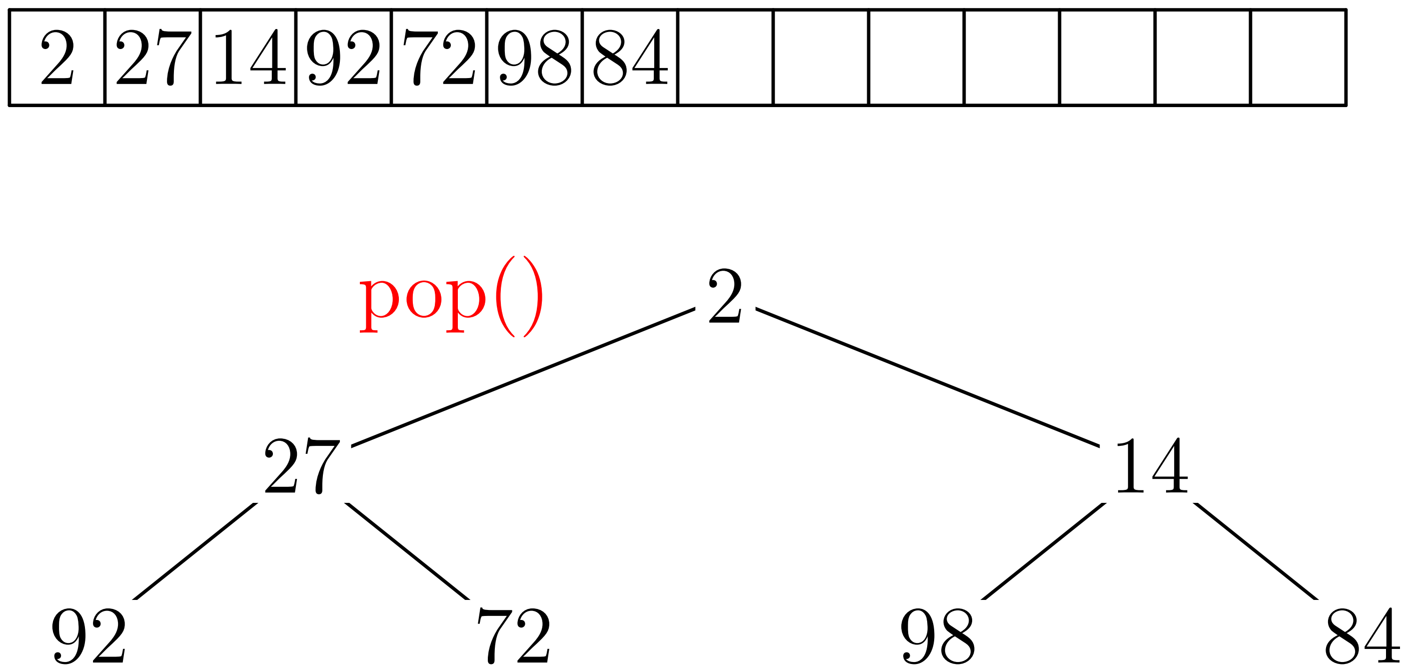


Heaps in Action



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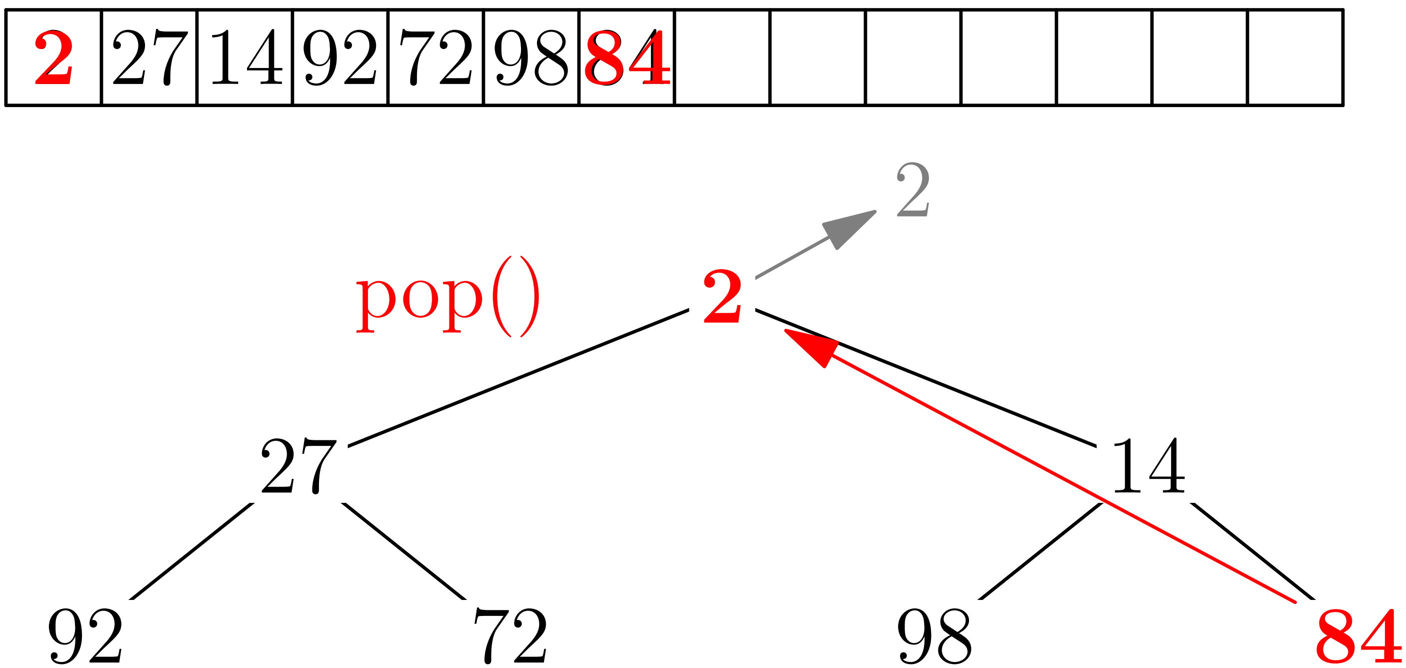
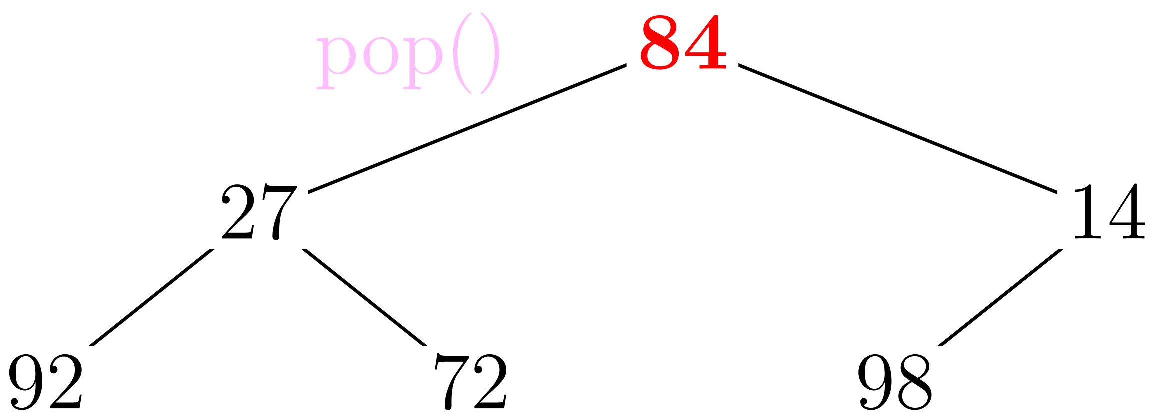


Heaps in Action

84	27	14	92	72	98								



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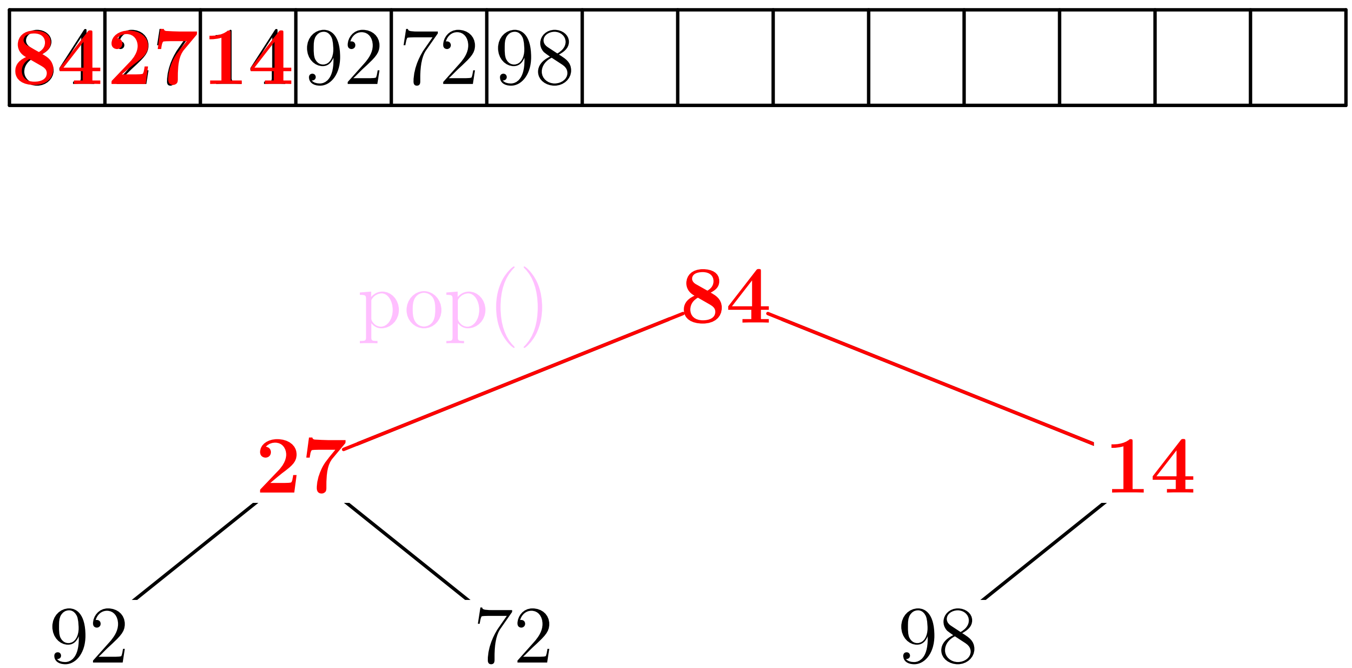

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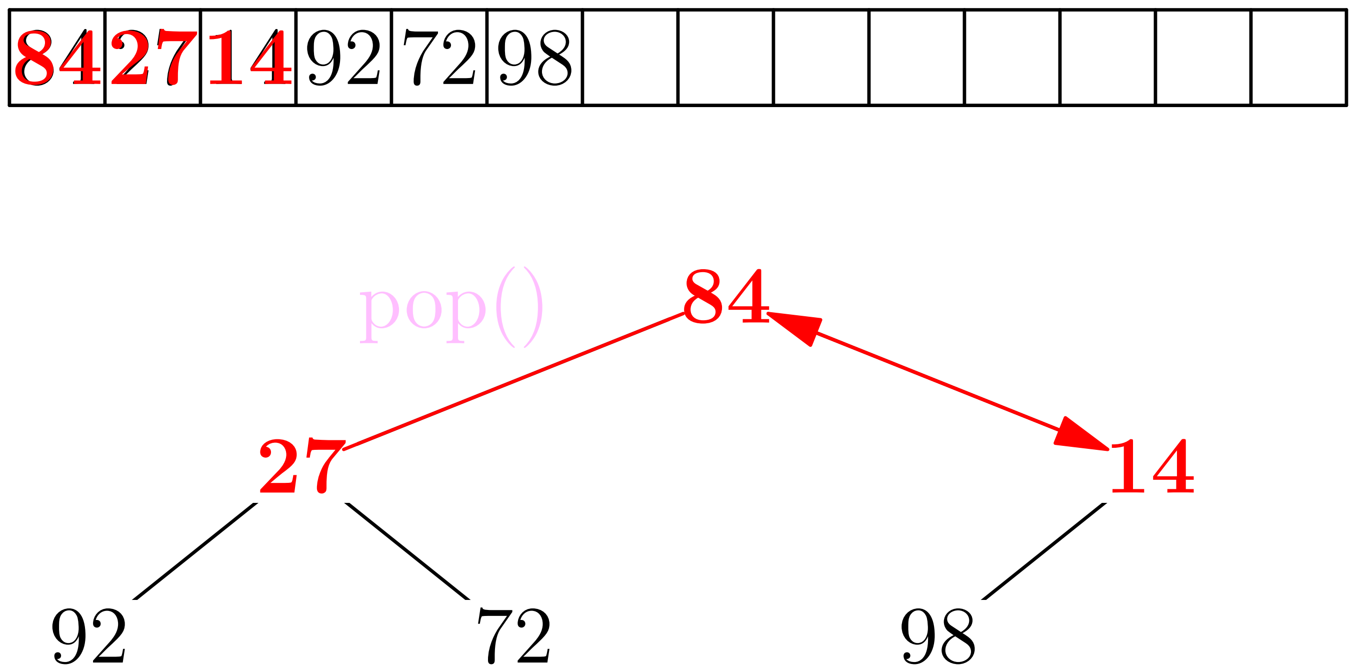
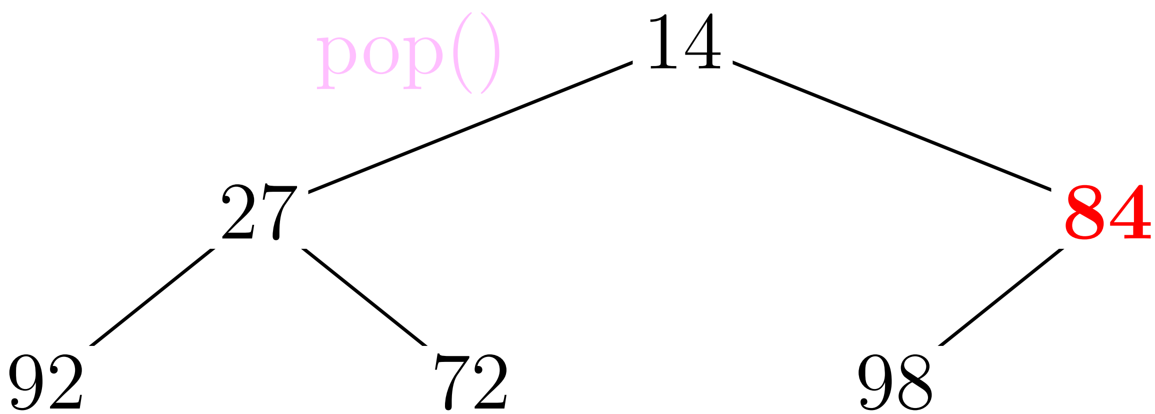


Heaps in Action

14	27	**84**	92	72	98								



Heaps in Action

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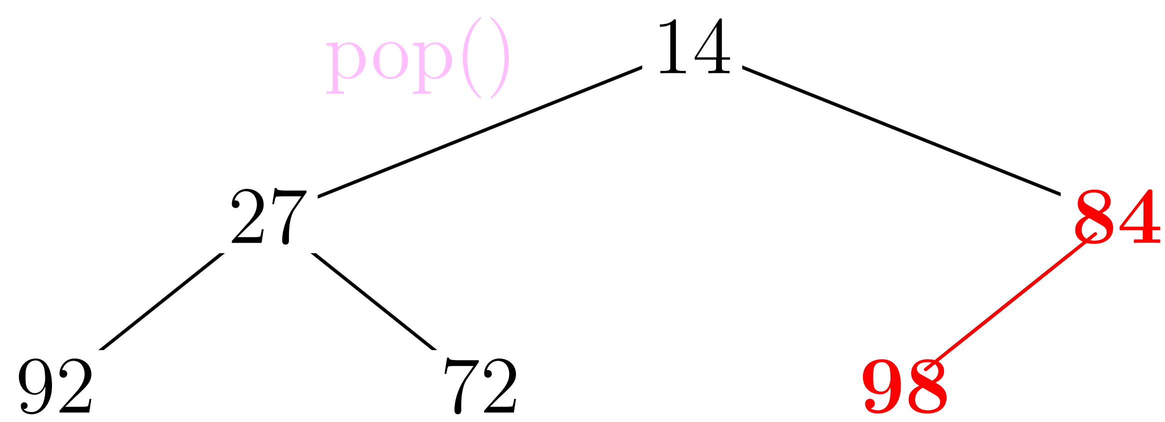

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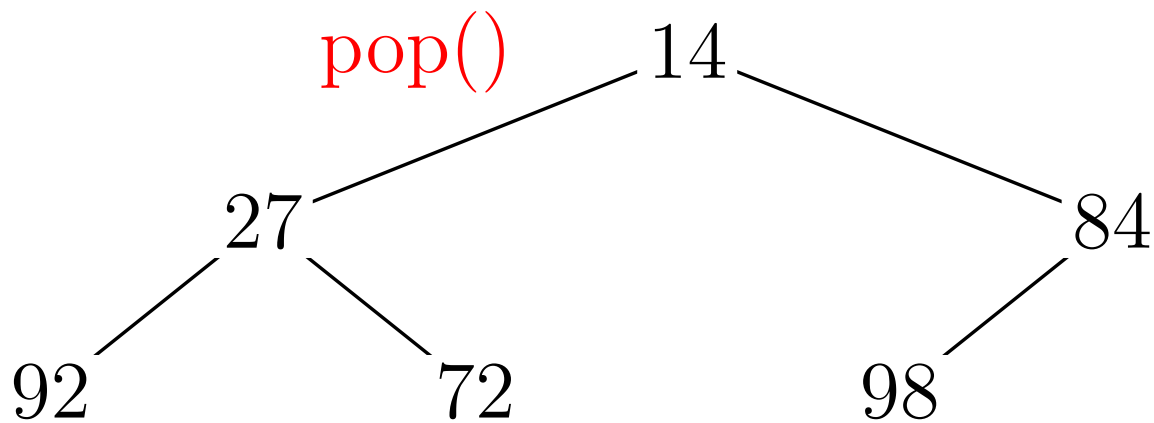

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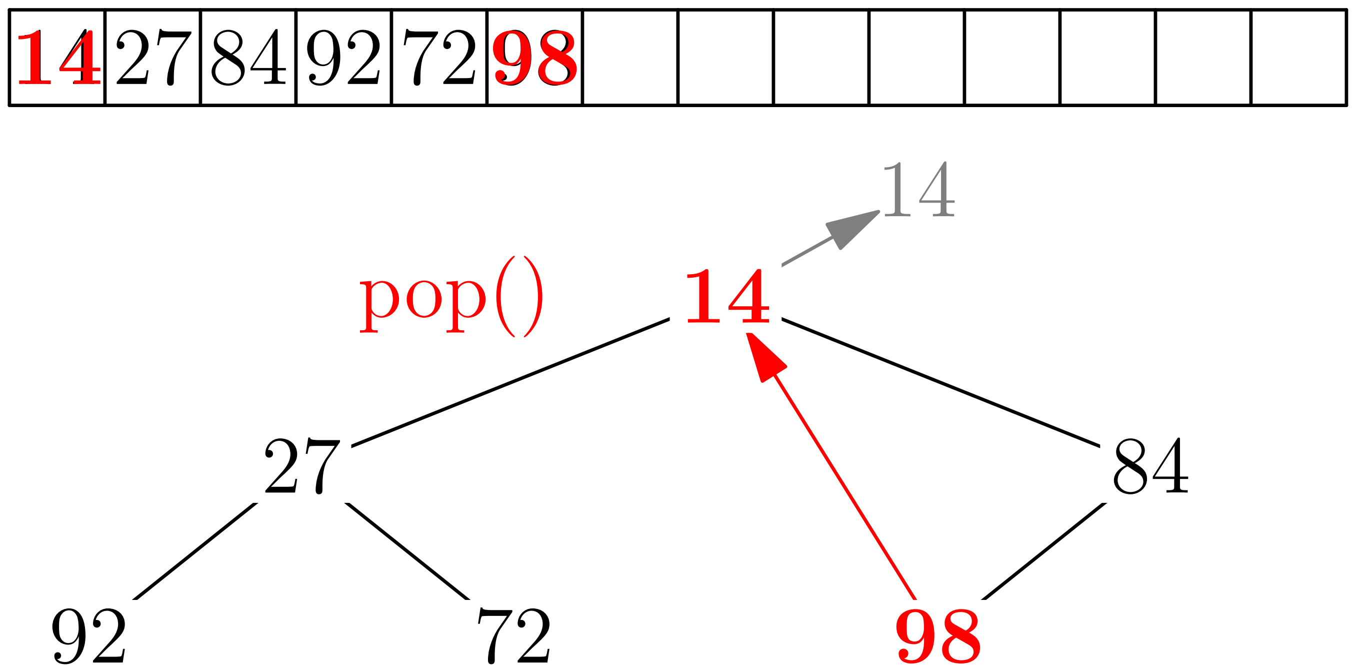

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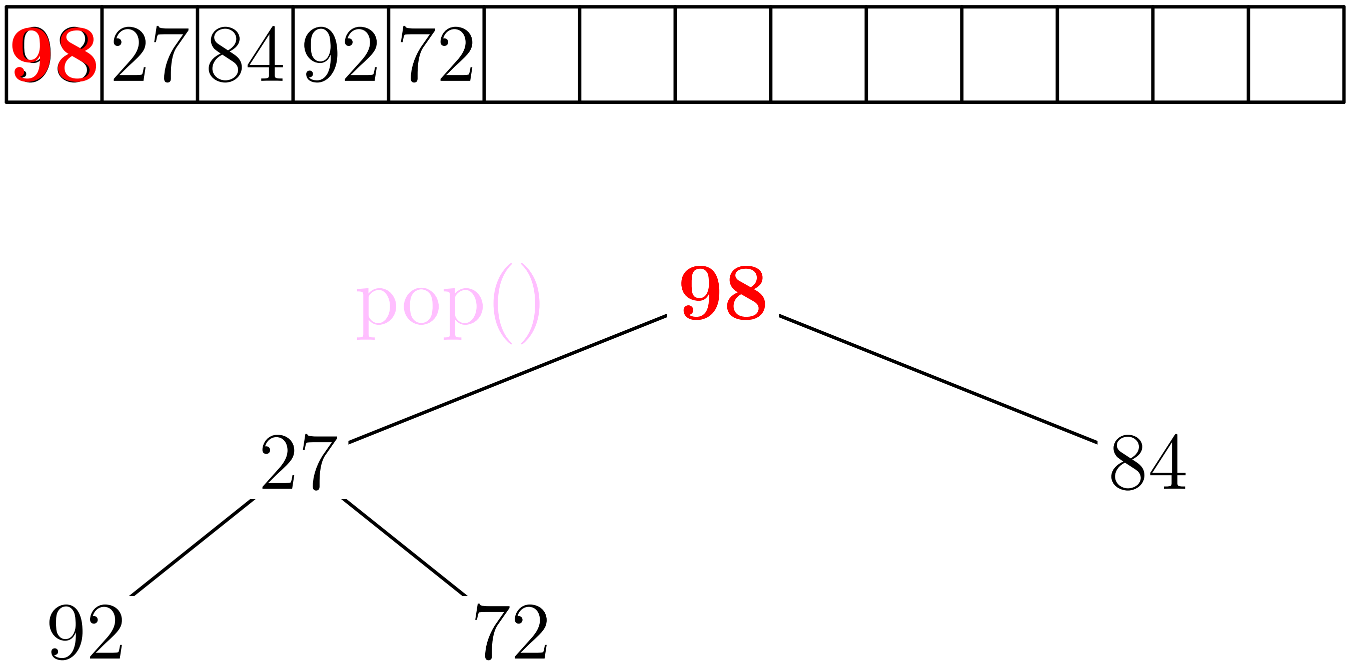

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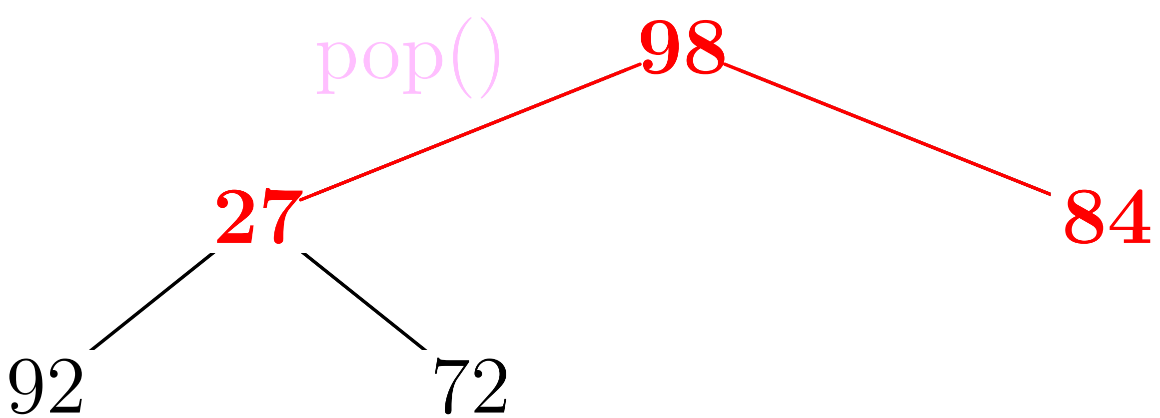
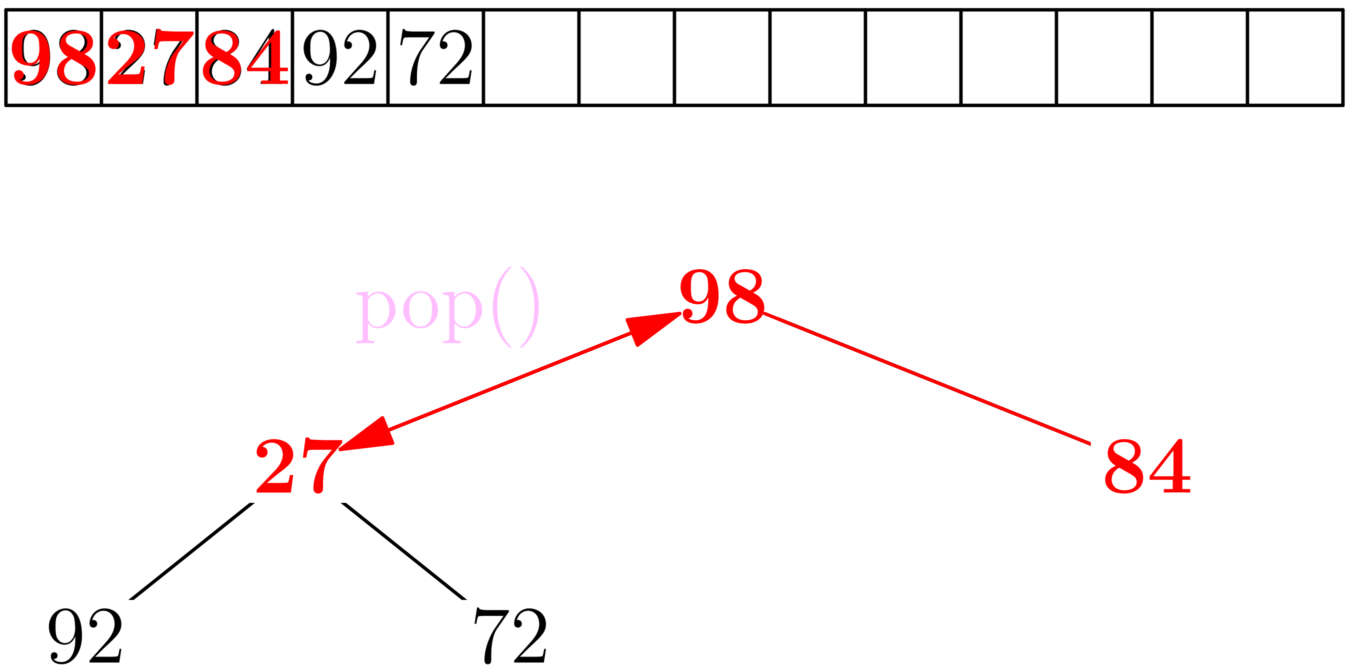


Heads in Action



Heads in Action

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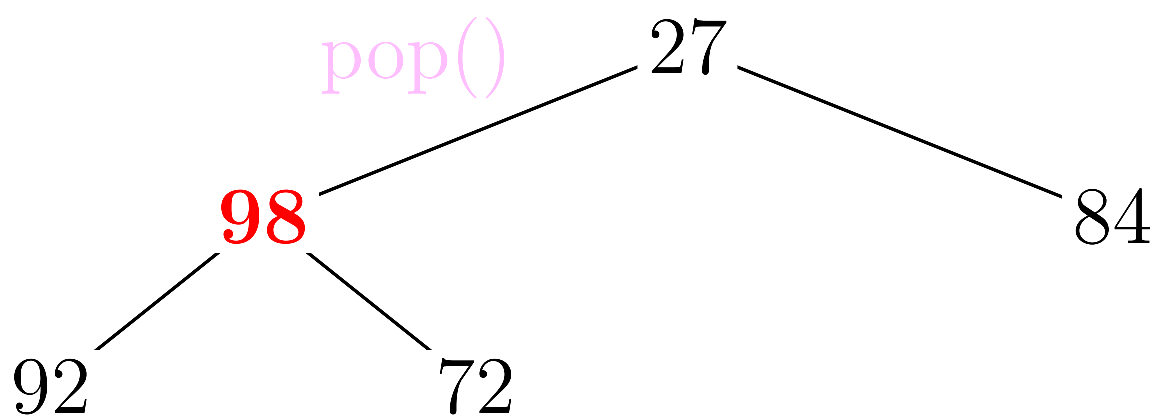

Heaps in Action

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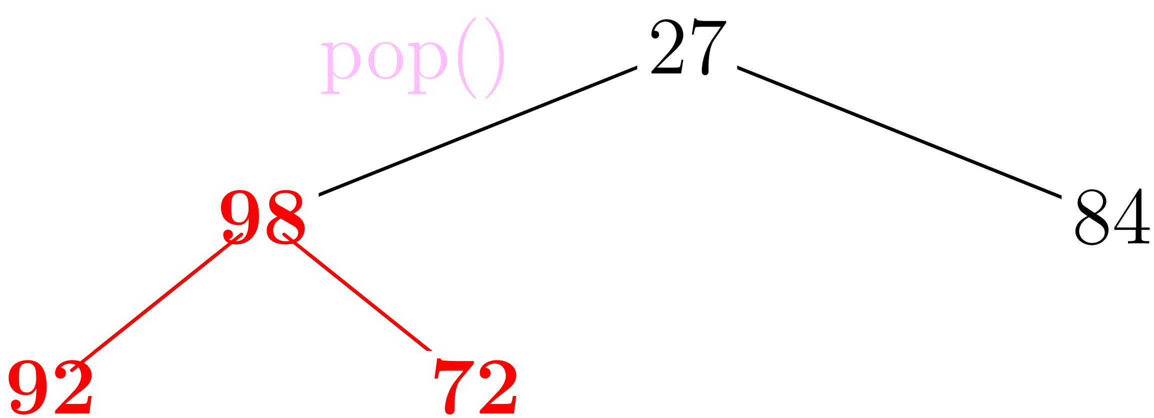

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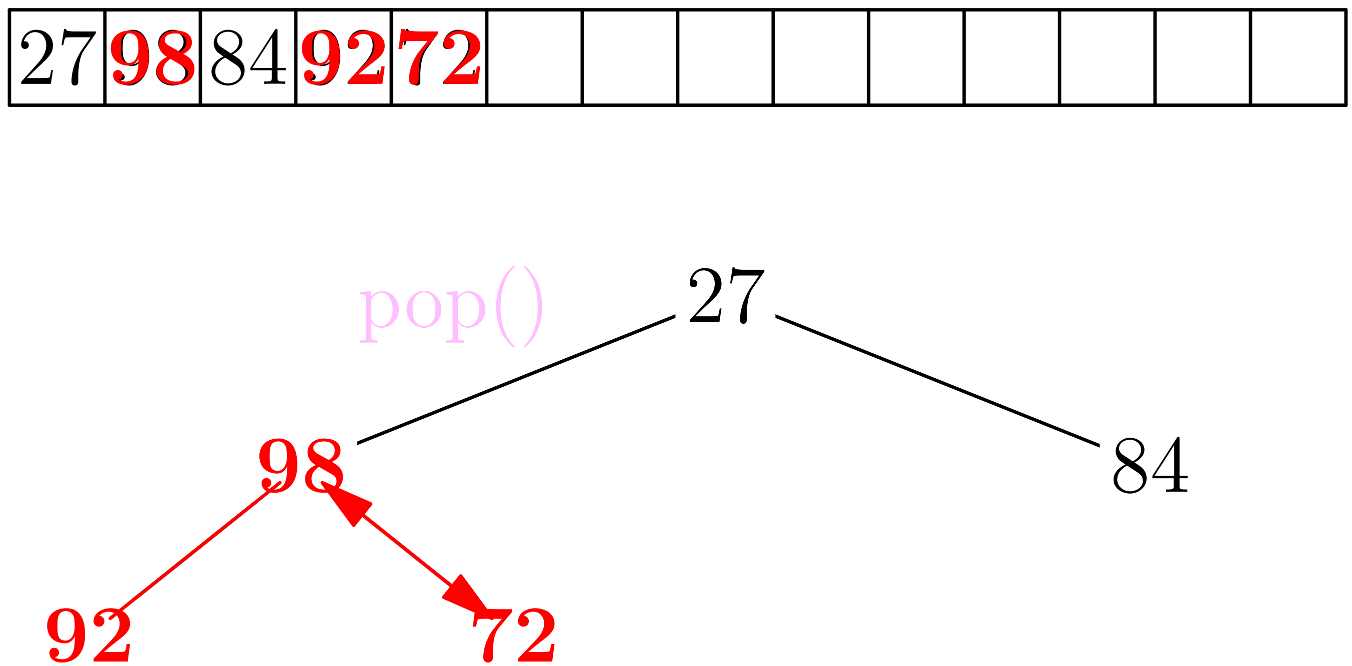

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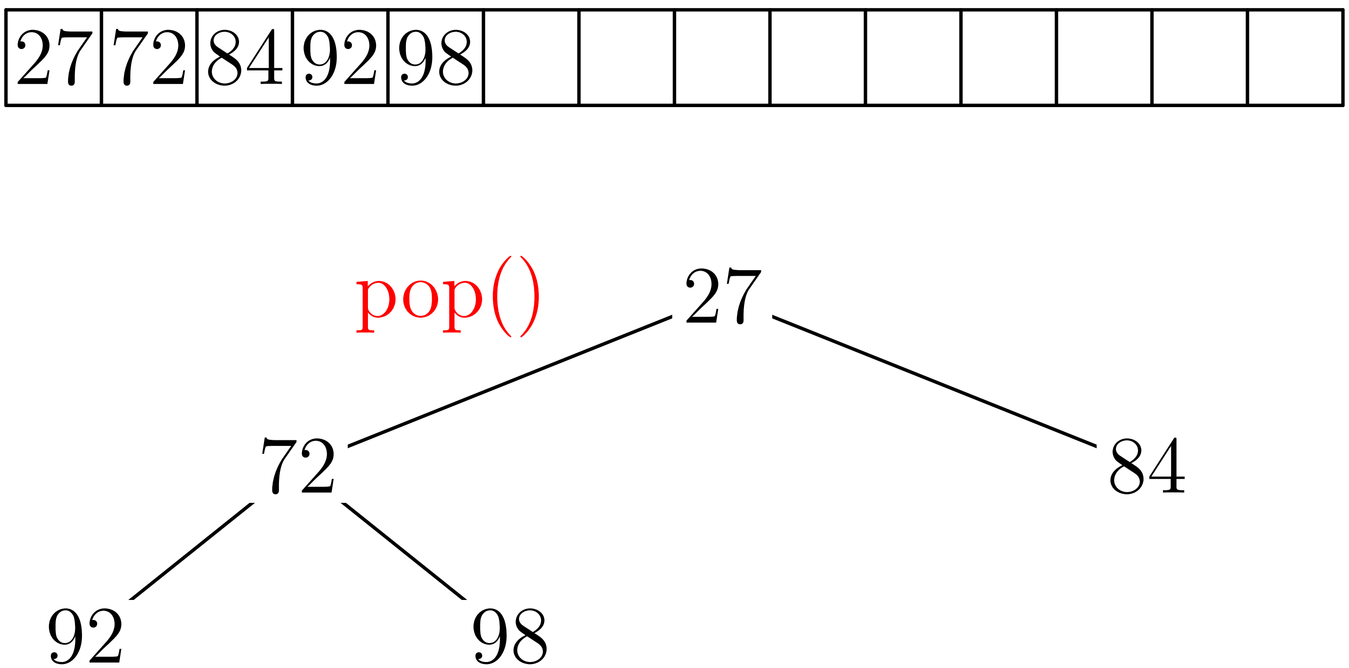

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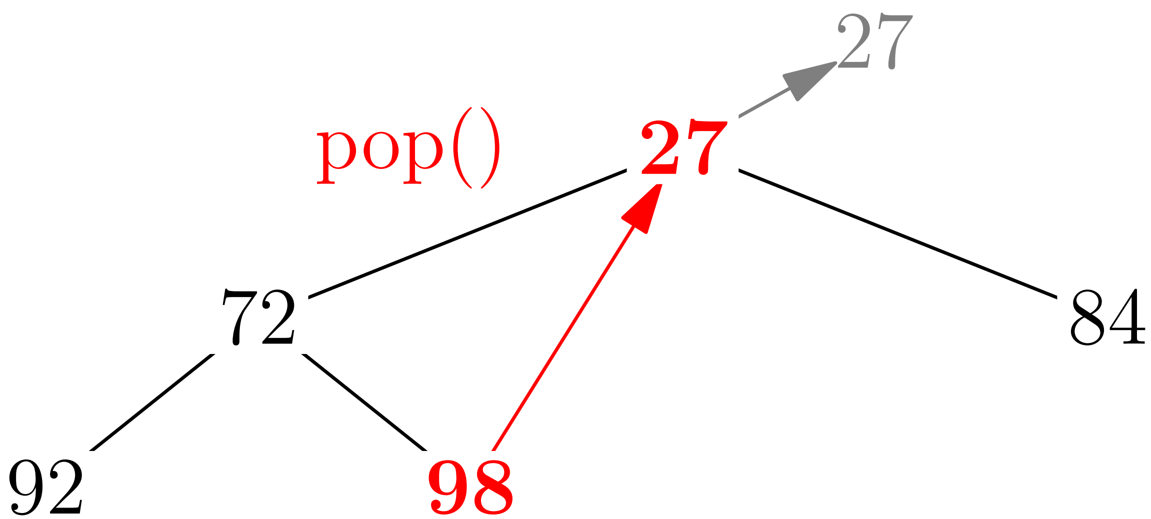

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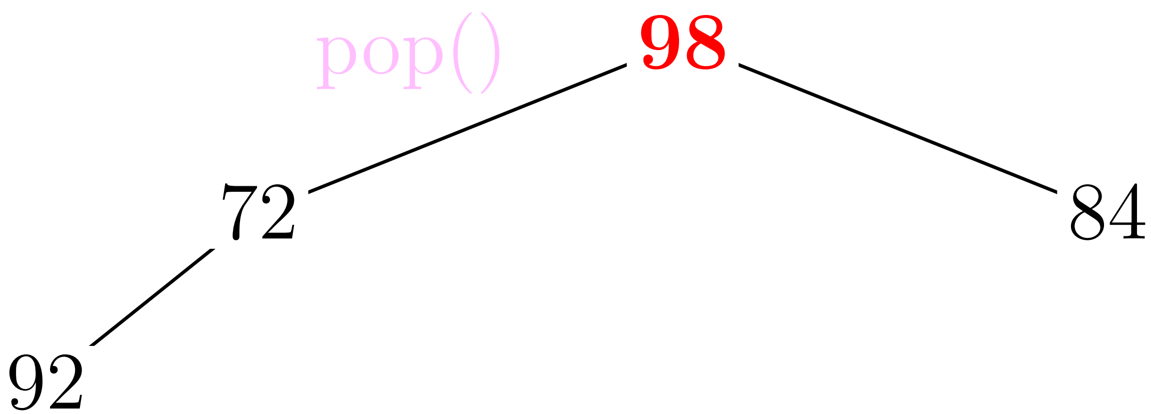

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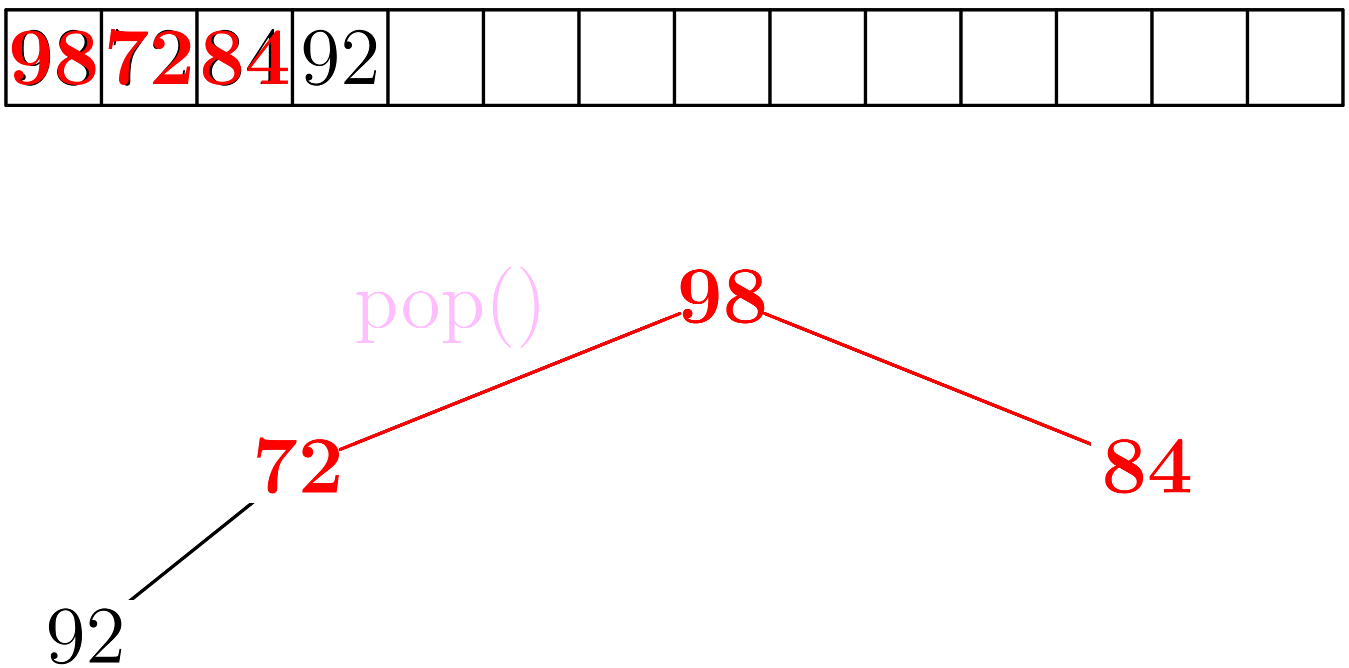

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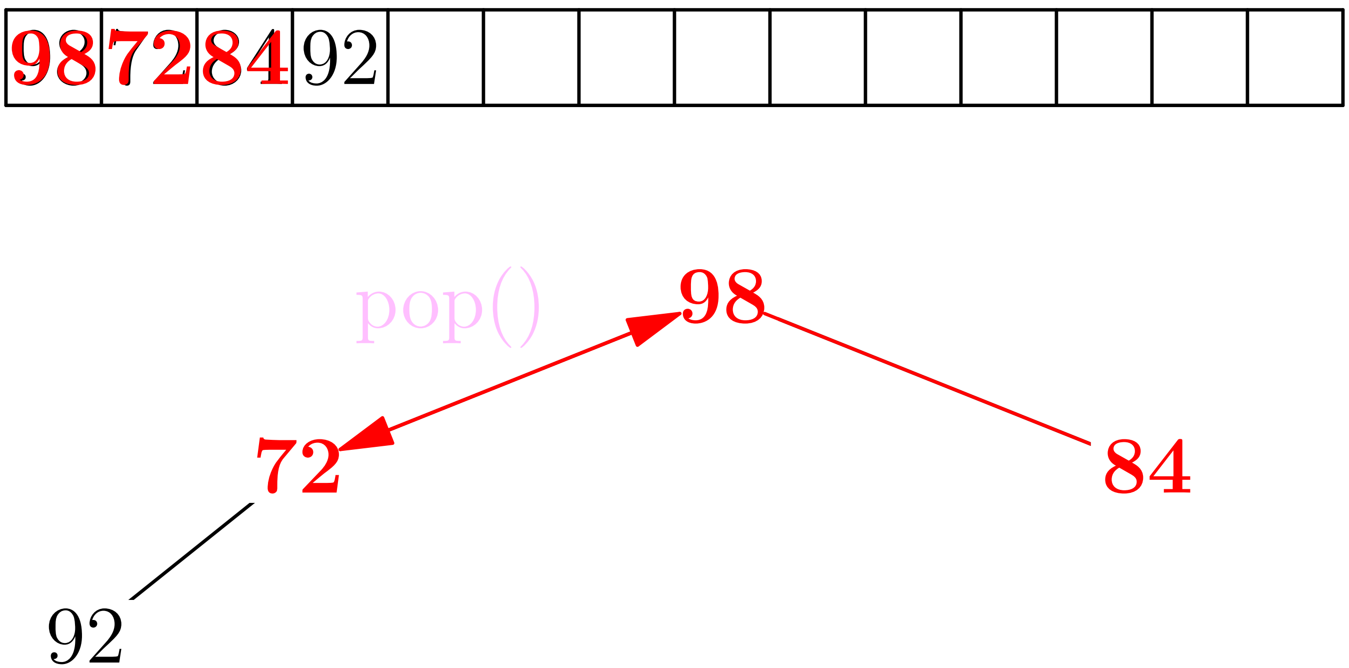

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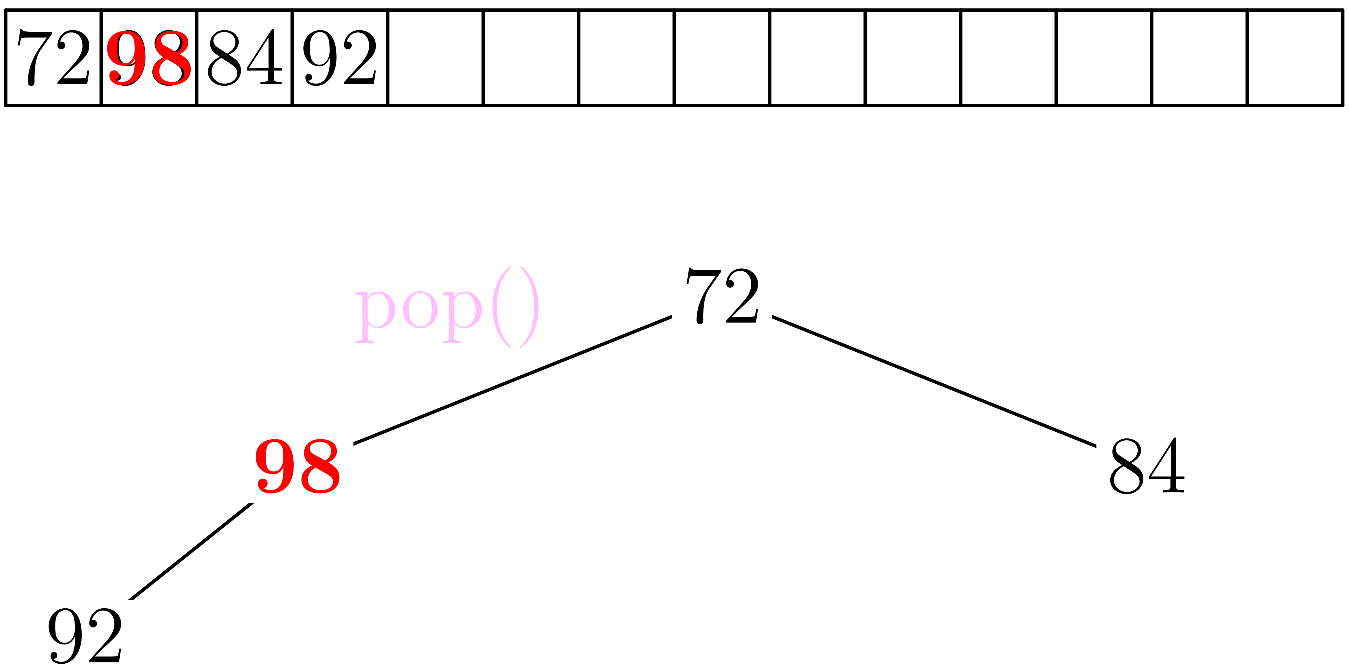

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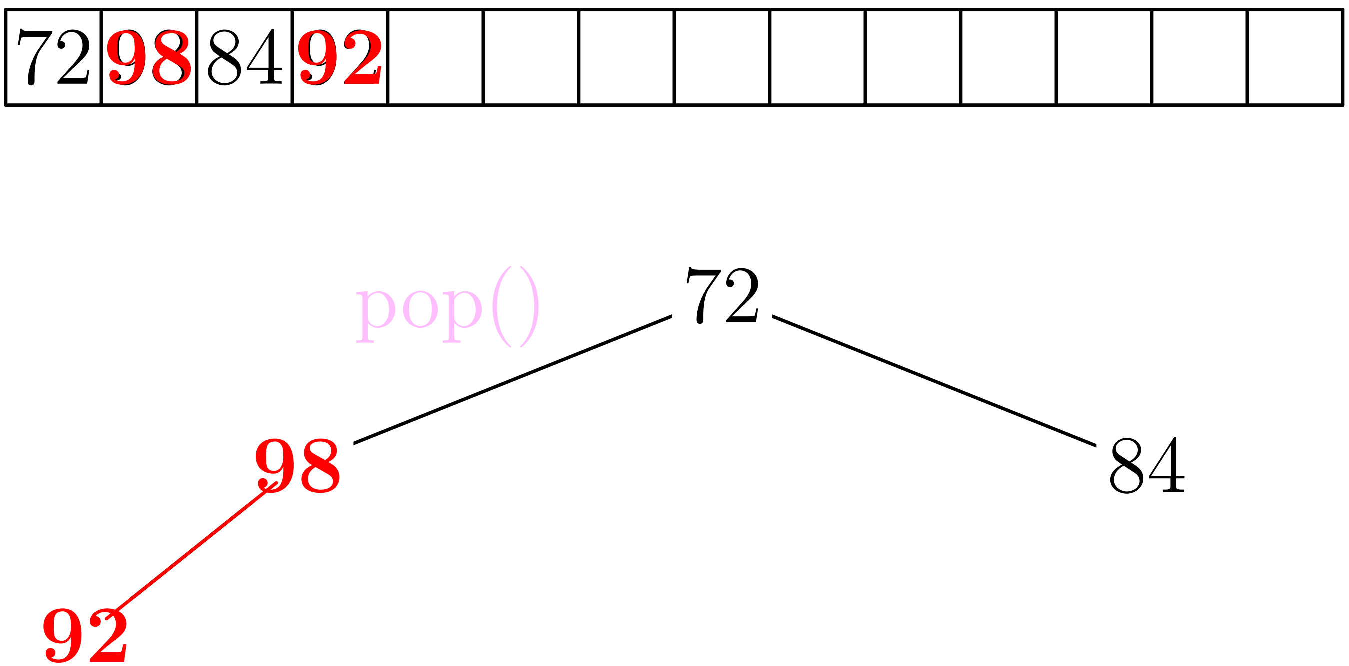

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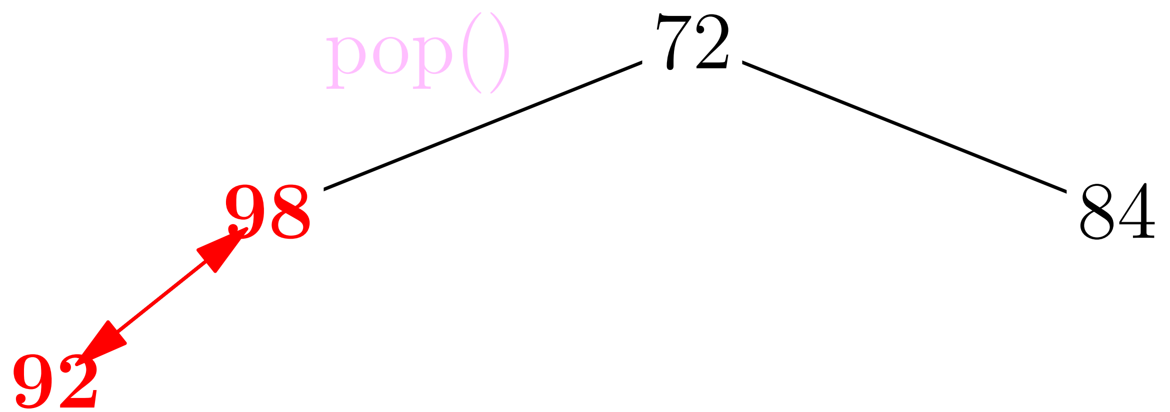

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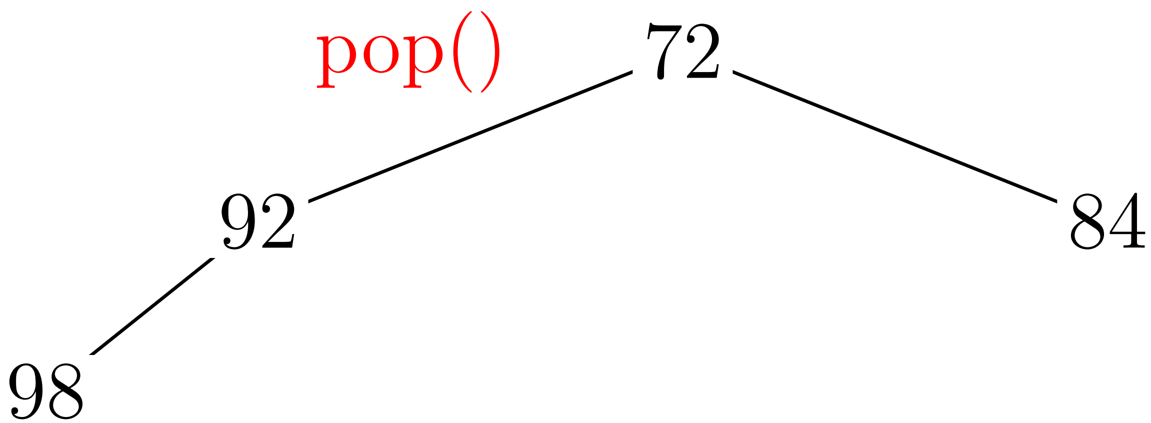

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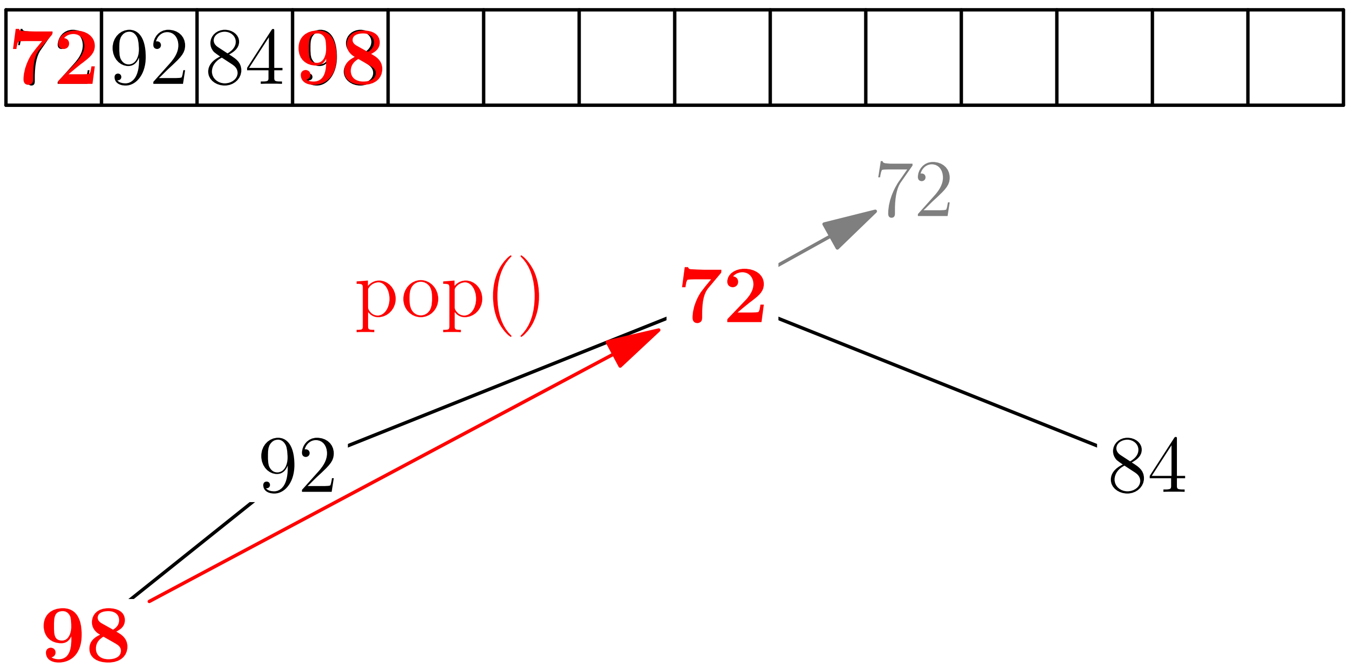

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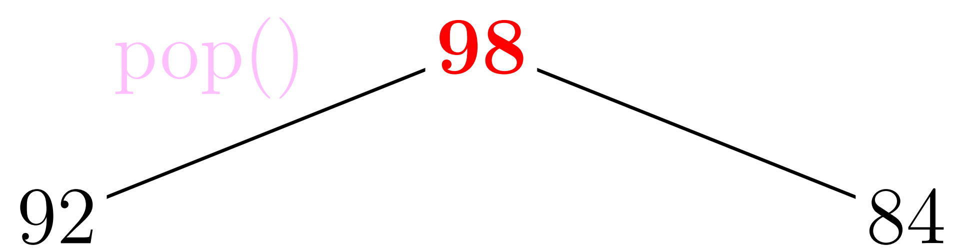

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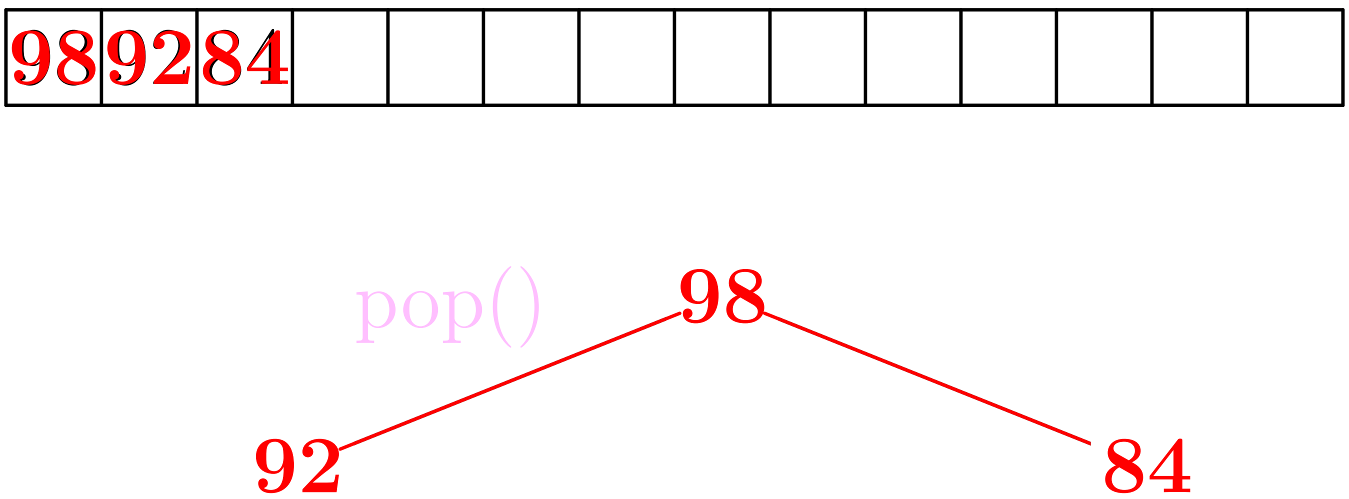

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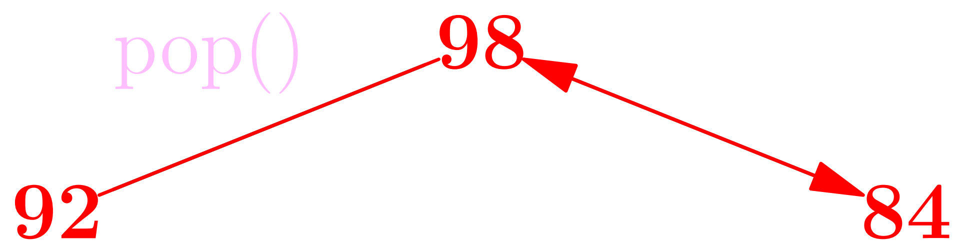

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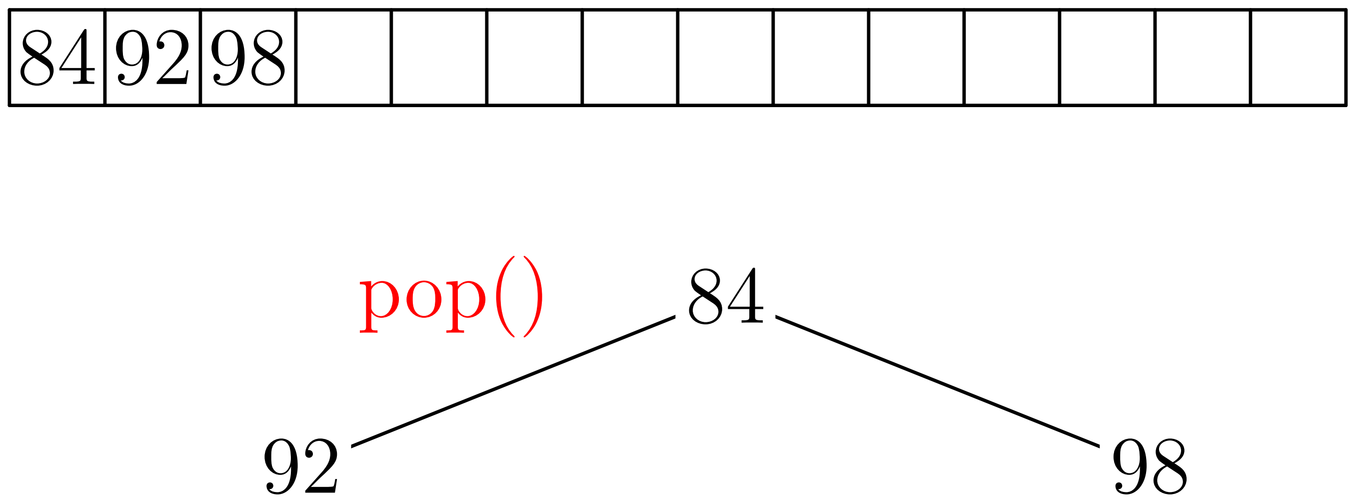

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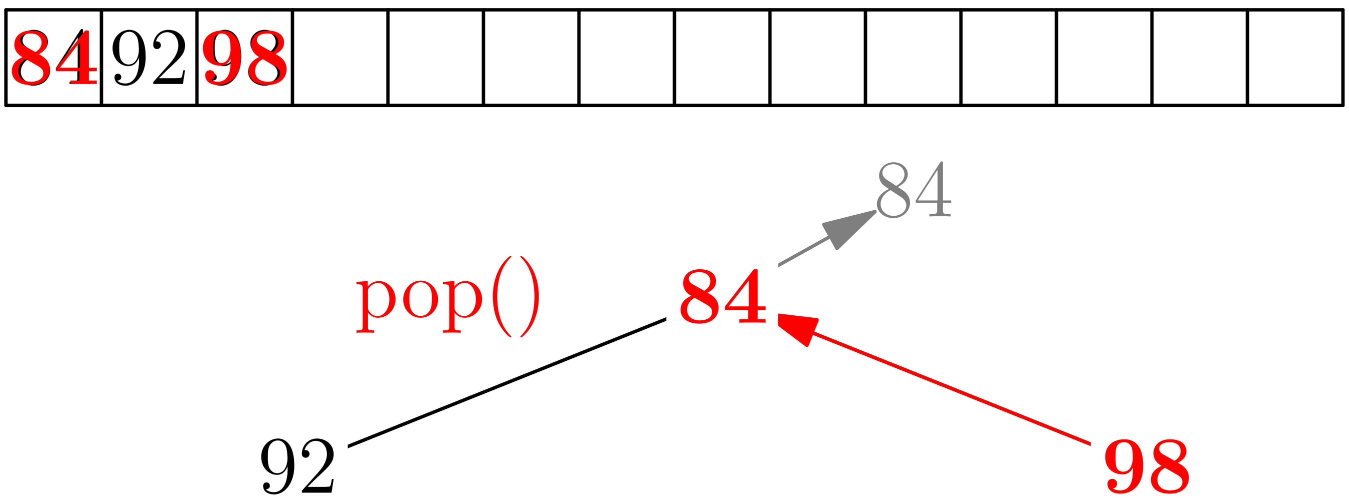

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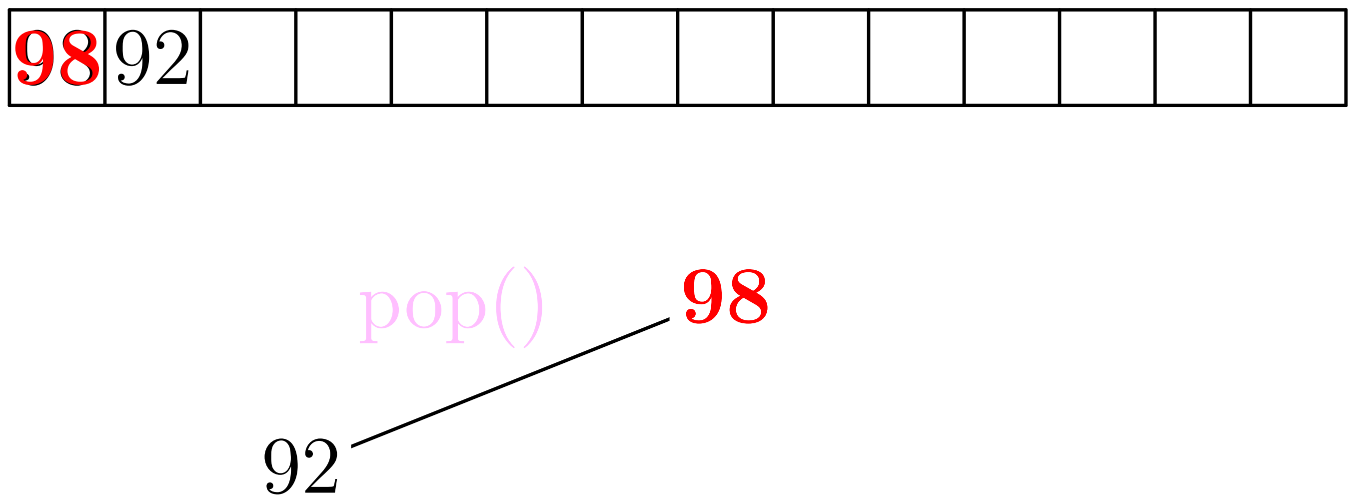

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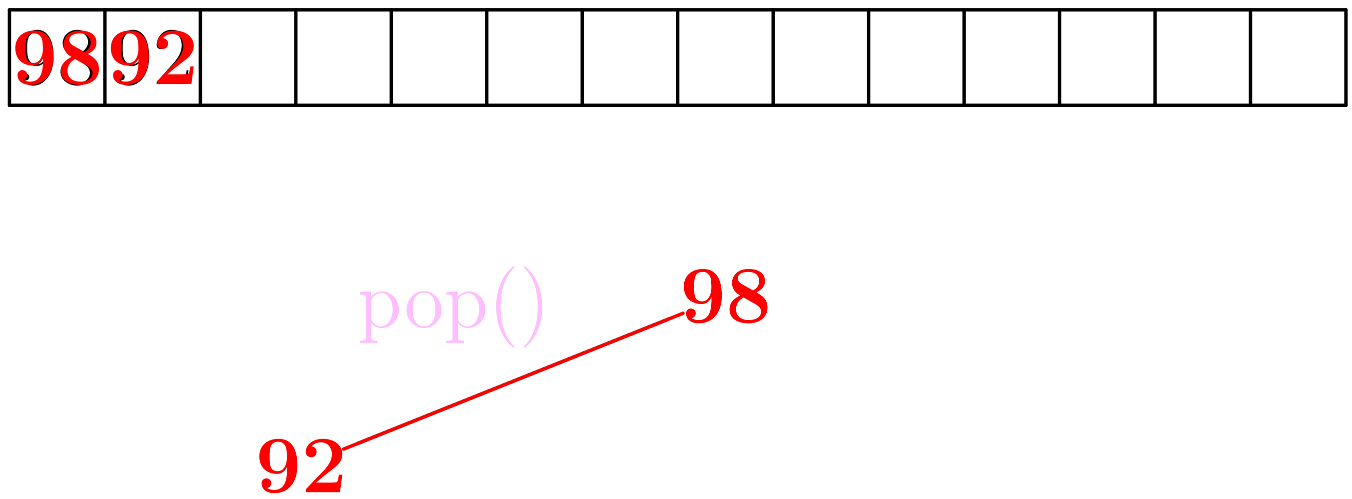

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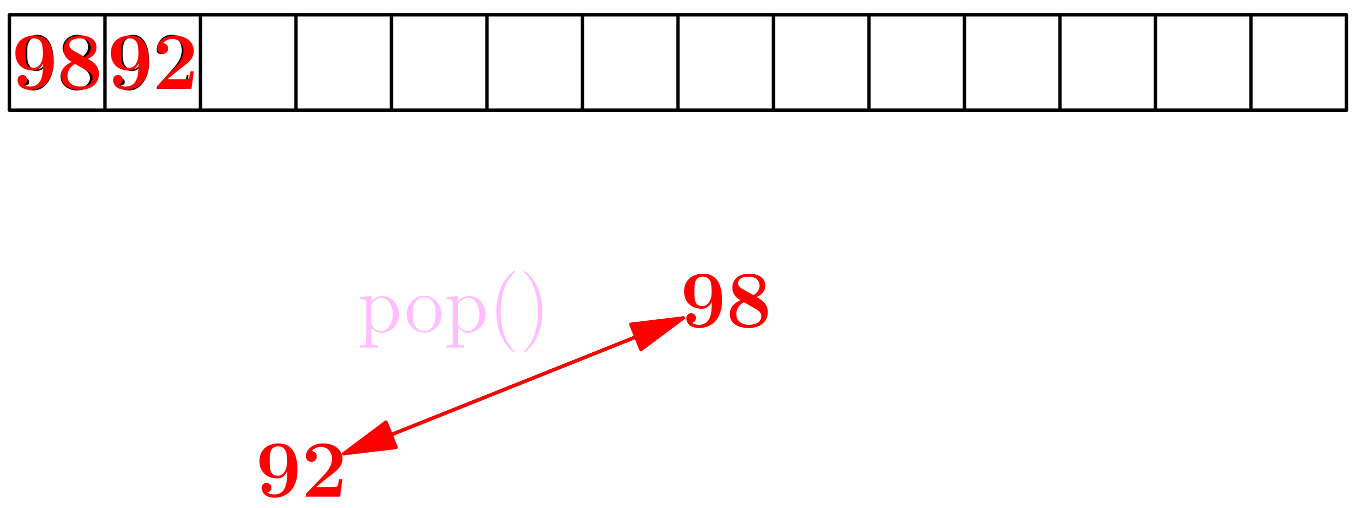

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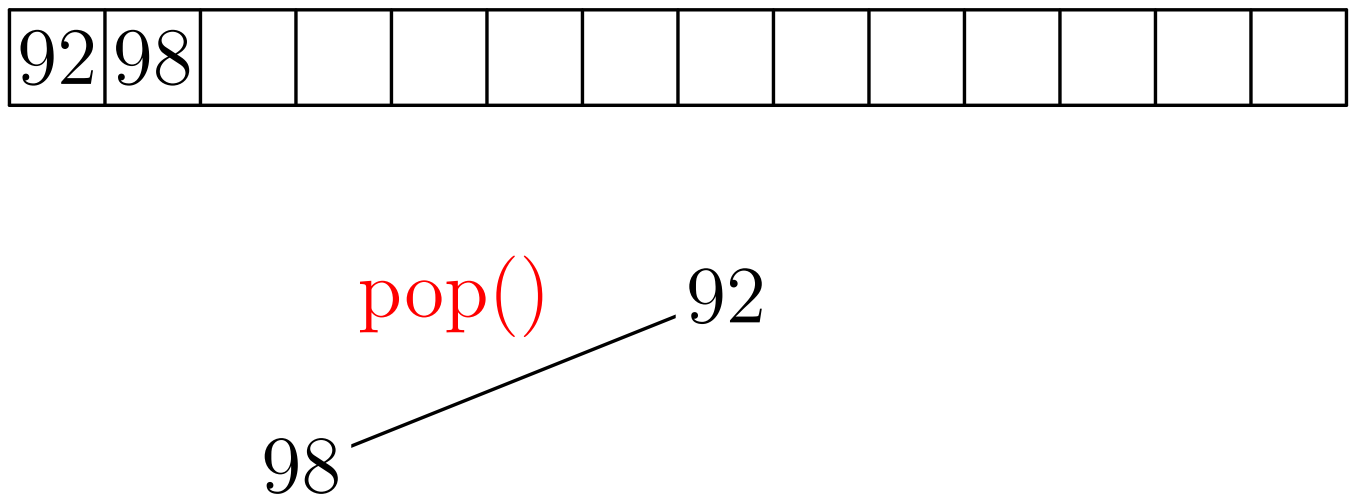

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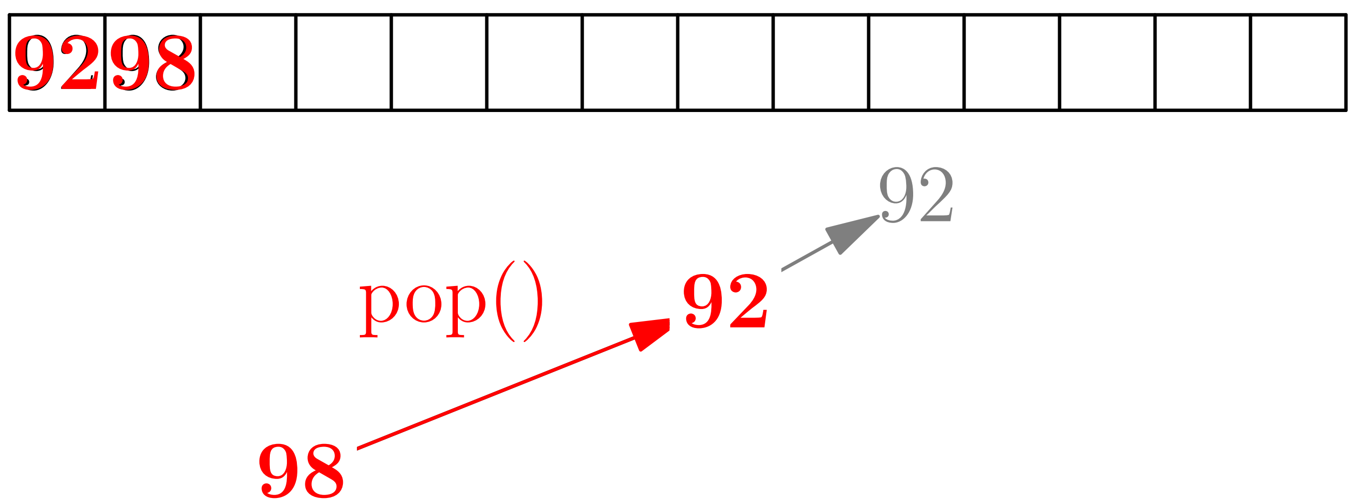

Heaps in Action

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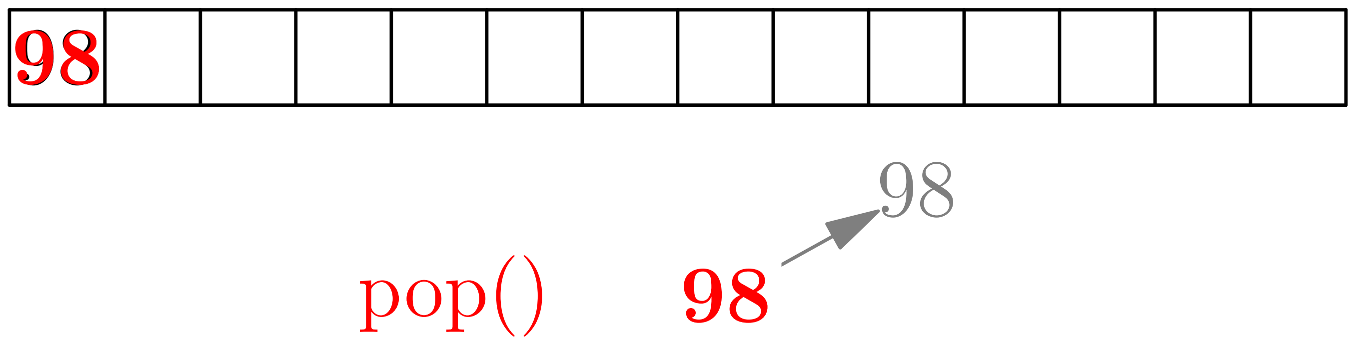
Heaps in Action



Time Complexity of Heaps

- The two important operation are `add` and `removeMin`
- These both either percolating an element up the tree or percolating an element down the tree
- The number of operations depends on the depth of the tree which is $\Theta(\log(n))$
- Thus `add` and `removeMin` are $O(\log(n))$
- Except `add` could also require resizing the array, but the amortised cost of this is low

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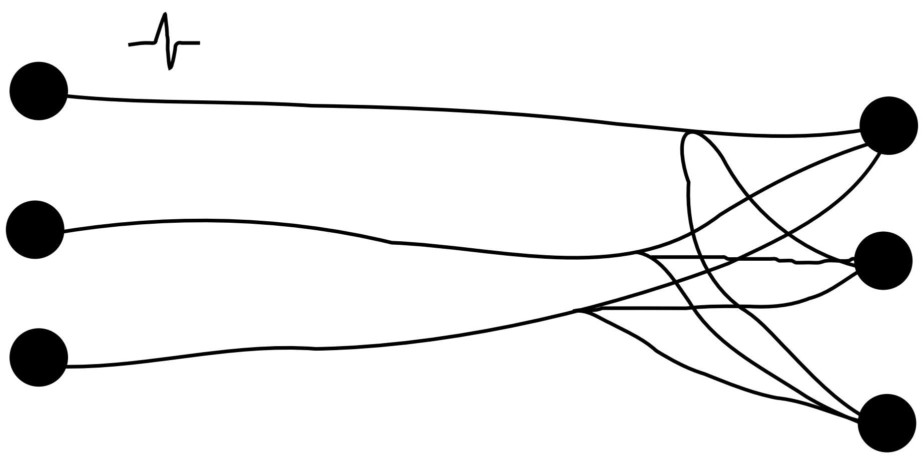
Real Time Simulation

- A nice application of priority queues is to perform real time simulations
- I was once modelling a neural network where neuron fired an impulse which would then be received by other neurons



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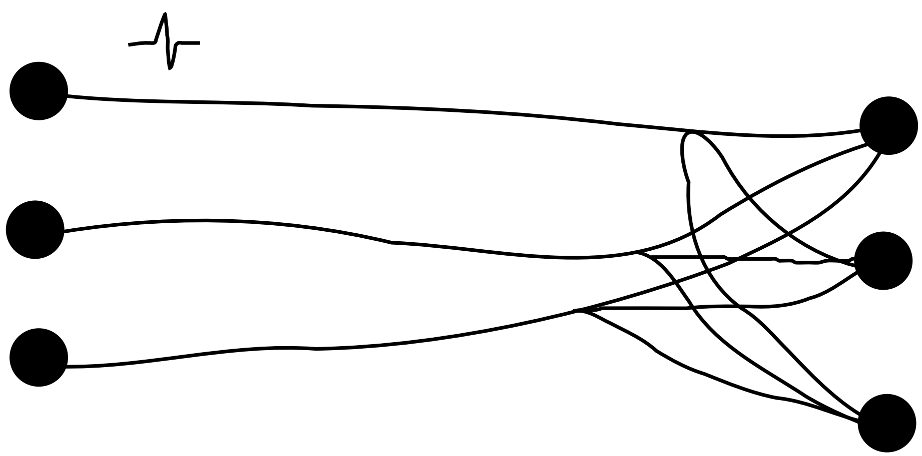

Synchronised Firing

- We wanted to show that if a group of neurons fired together they could make another group of neurons fire in synchrony despite the fact that it would take different times for the receiving neurons to feel the pulse (due to the different lengths of the axons)

- A famous Israeli group had "proved" this couldn't happen

- Using a priority queue we modelled the neurons
  - ⋆ When a neuron fired the receiving neurons would be put on a priority queue according to when they received the pulse
  - ⋆ If the receiving neurons received enough pulse in a short enough time they would then fire

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  - ⋆ When a neuron fired the receiving neurons would be put on a priority queue according to when they received the pulse
  - ⋆ If the receiving neurons received enough pulse in a short enough time they would then fire

Efficient Modelling

- Using a priority queue meant we knew when the next event would happen
- We did not have to run a clock where most of the time nothing happened
- This allowed us to perform a very large simulation efficiently
- The simulation showed that the pulse of neurons synchronised despite the “proof” that this wouldn’t happen

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Outline

1. Heaps
2. Priority Queues
   - Array Implementation
3. **Heap Sort**
4. Other Heaps

Heap Sort

- A priority queue suggests a very simple way of performing sort

- We simply add elements to a heap and then take them off again

```
template <typename T>
void sort(vector<T> aList)
{
    HeapPQ<T> aHeap = new HeapPQ<T>(aList.size());
    for (T element: aList)
      aHeap.push(element, element);

    aList.clear();
    while(aHeap.size() > 0) {
      aList.push_back(aHeap.top());
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- Note that this is not an in-place sort algorithm (i.e. it uses lots of memory)

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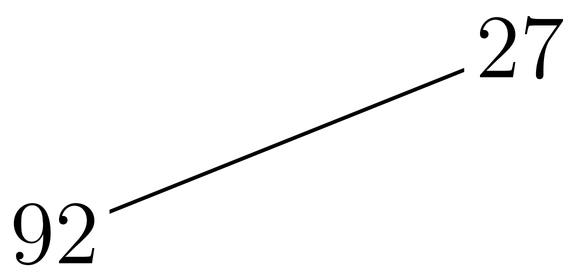

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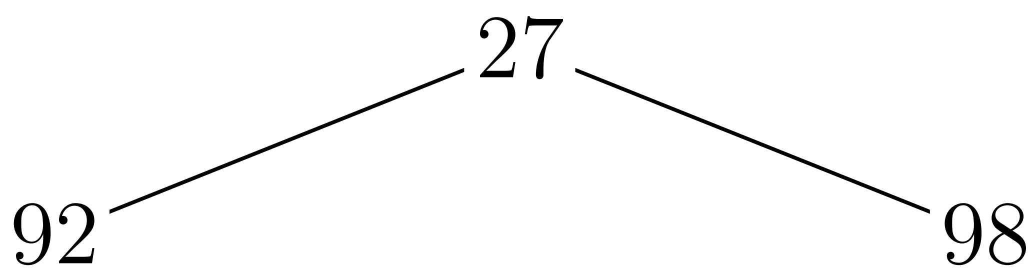

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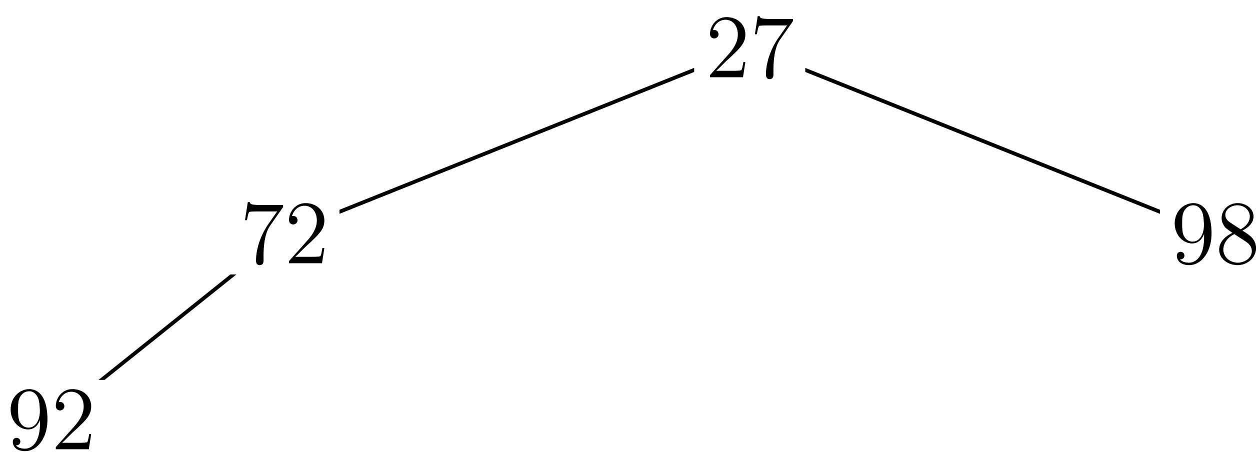

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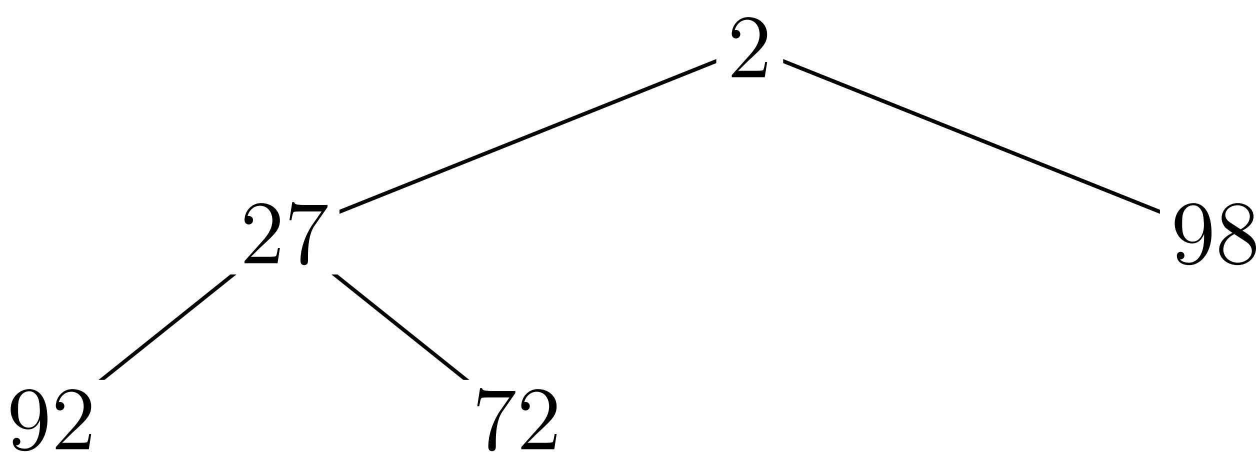

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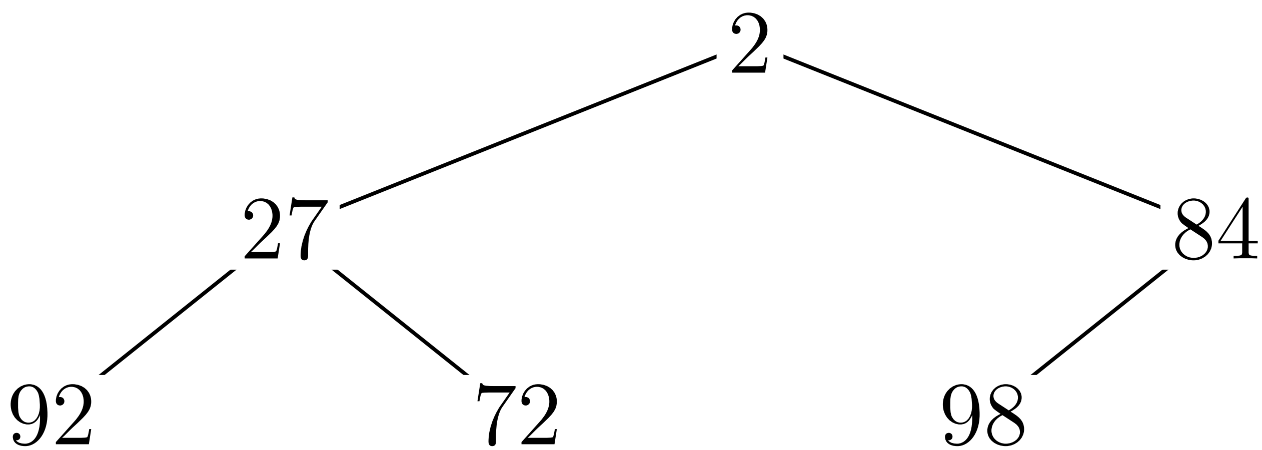

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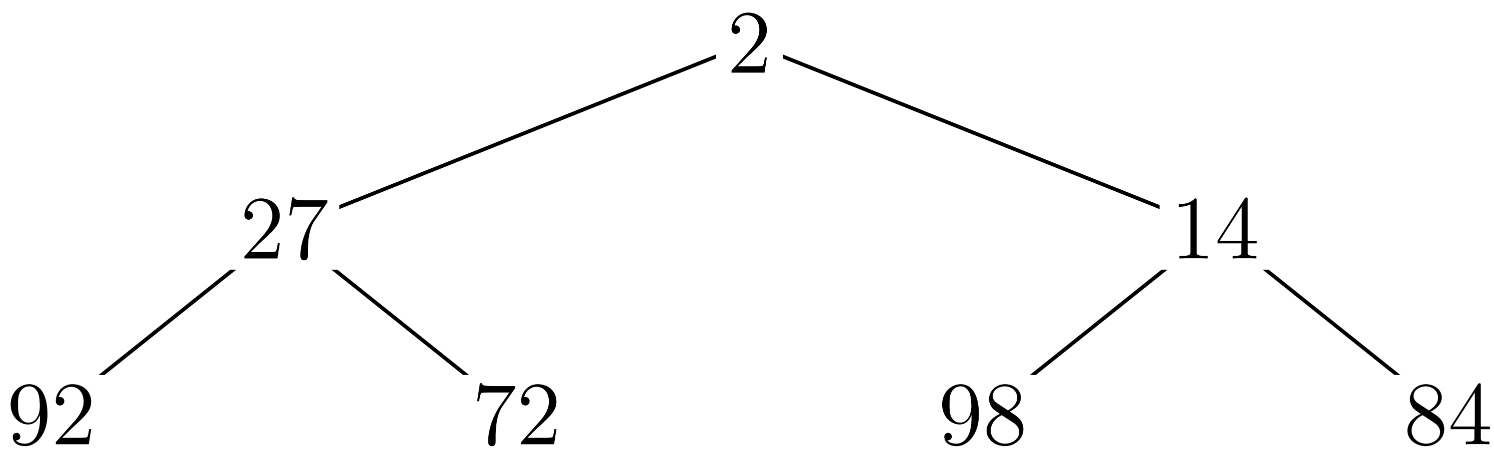

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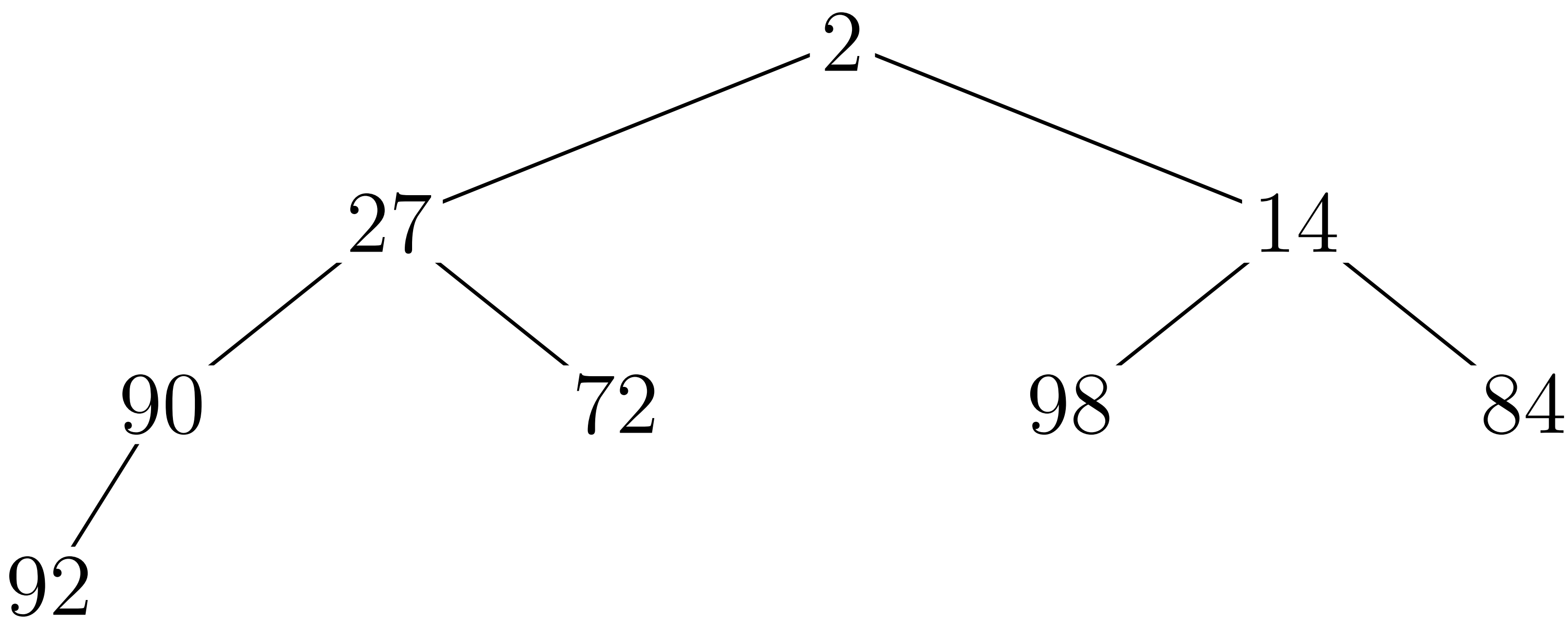

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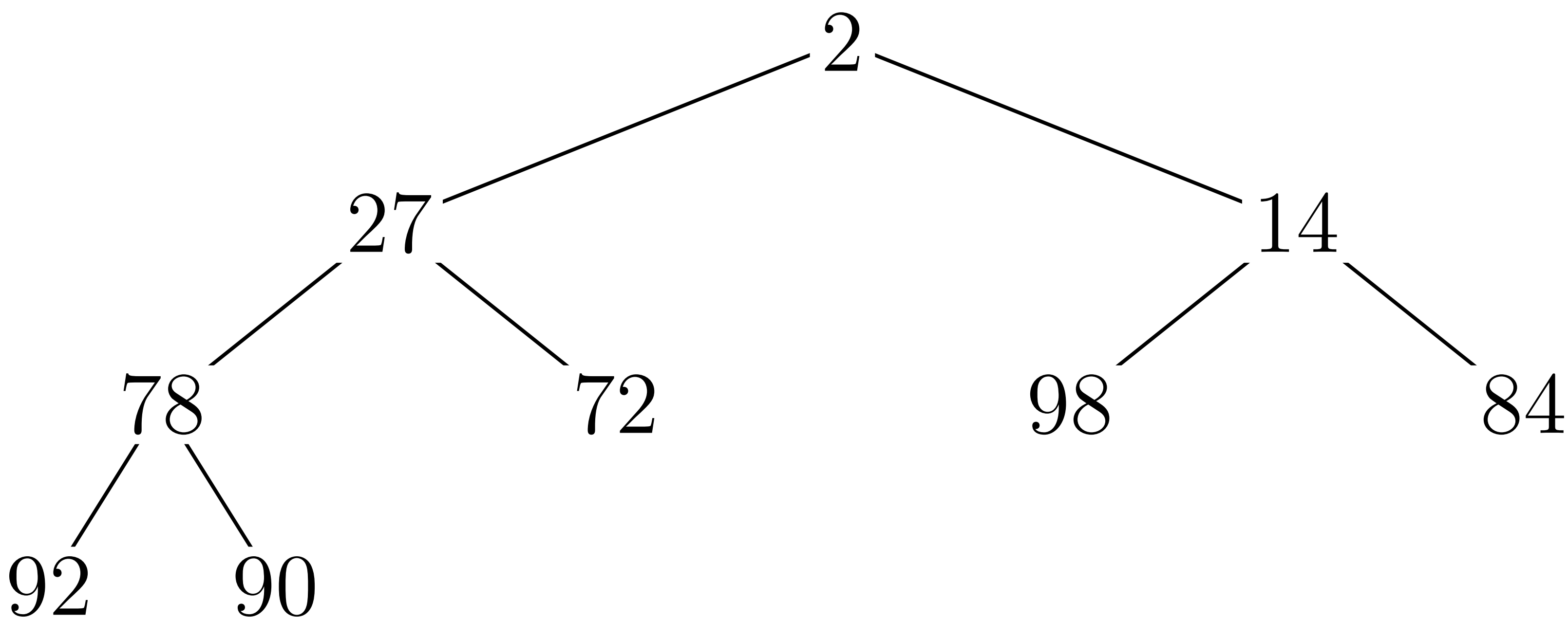

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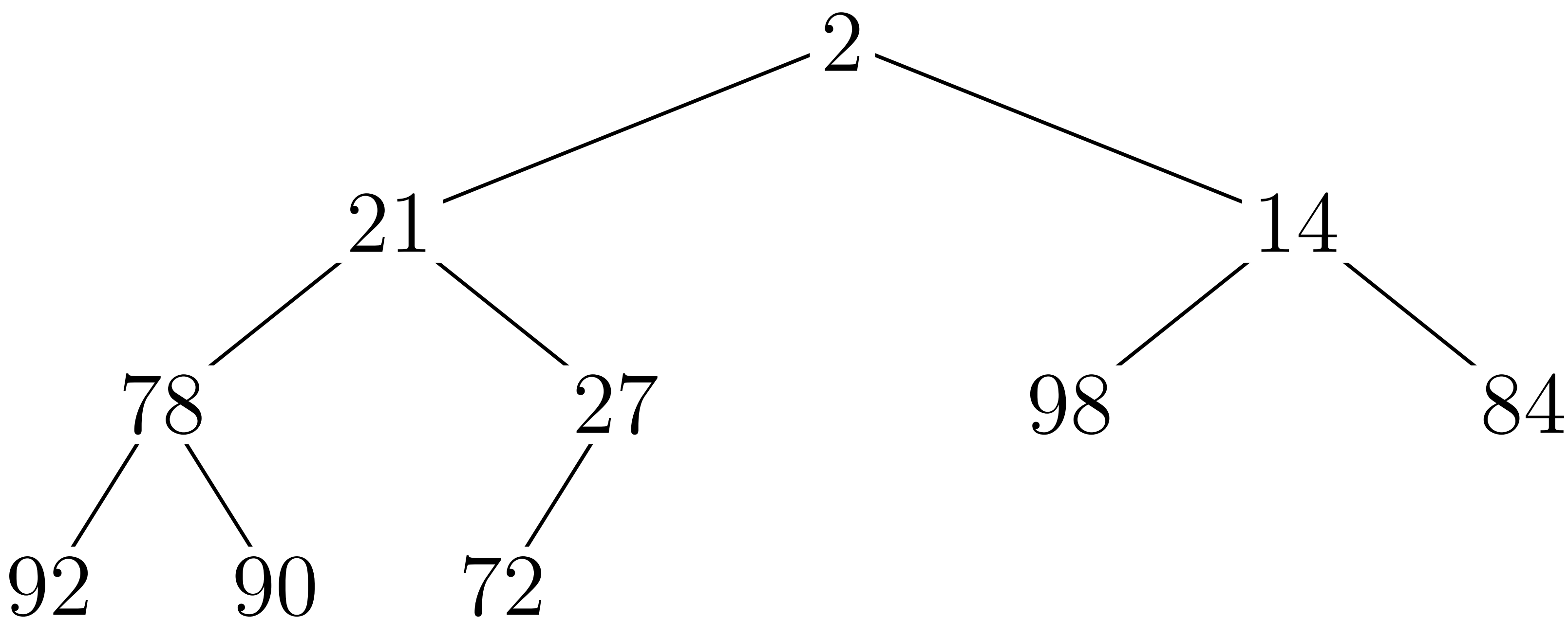

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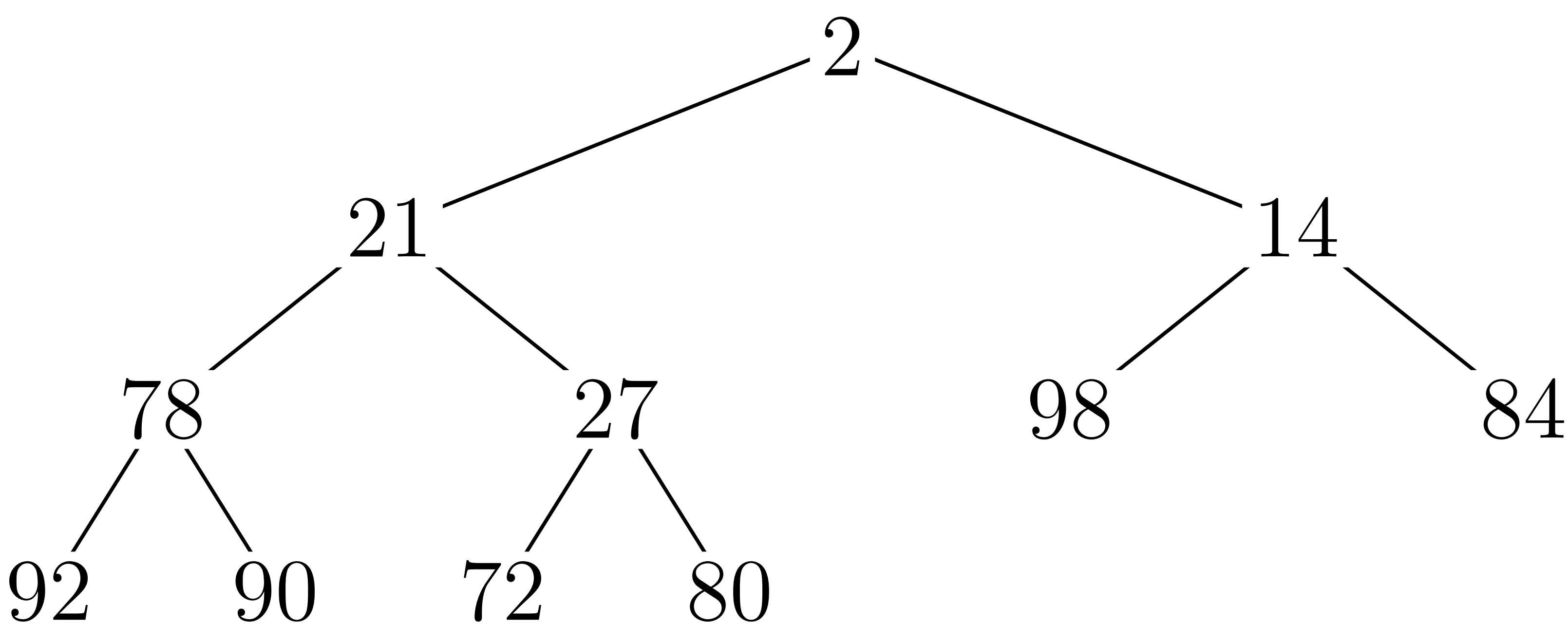

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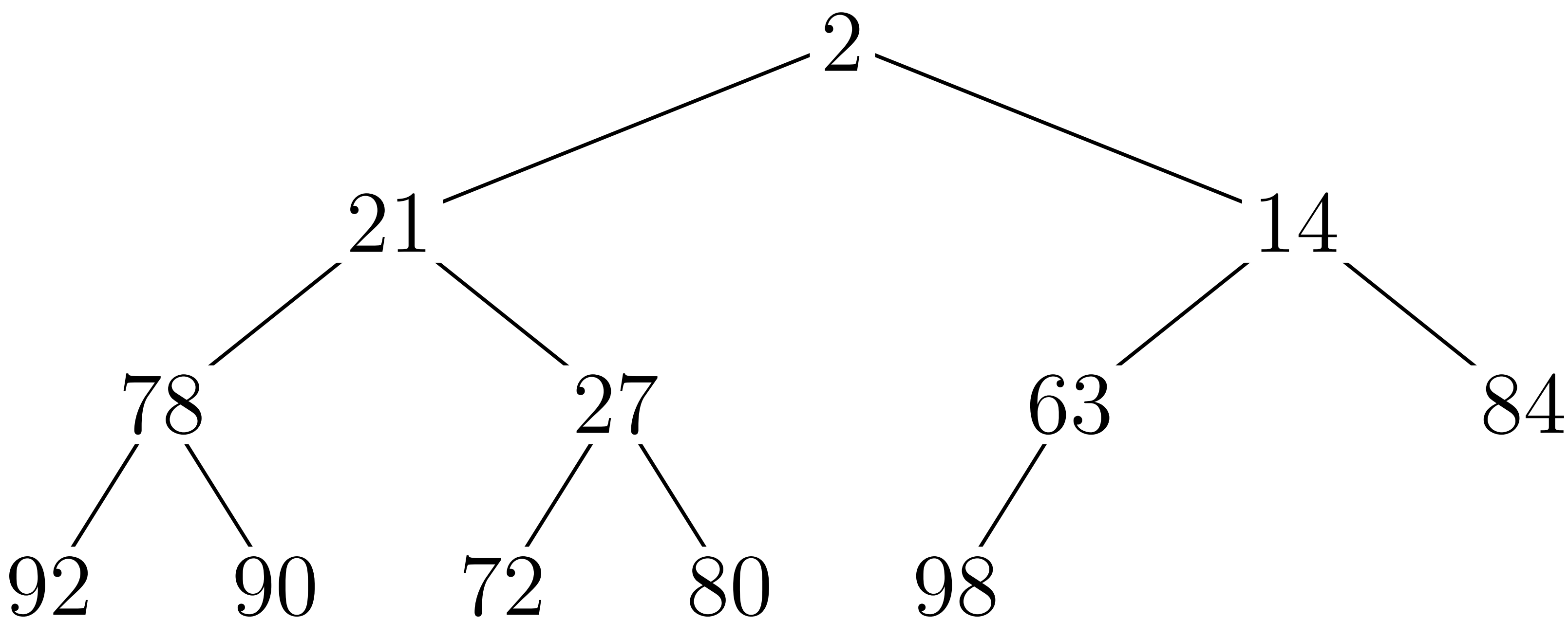

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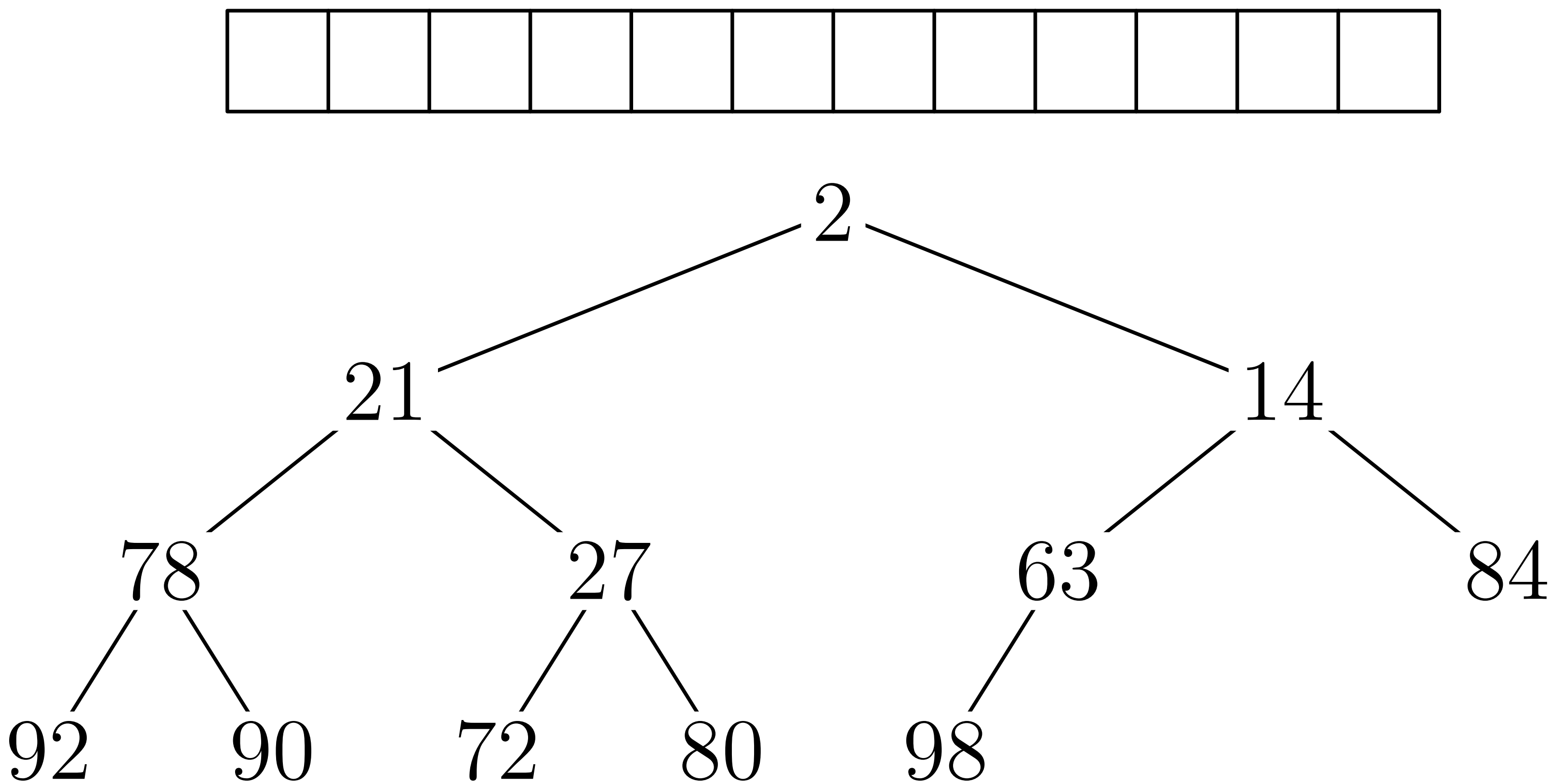

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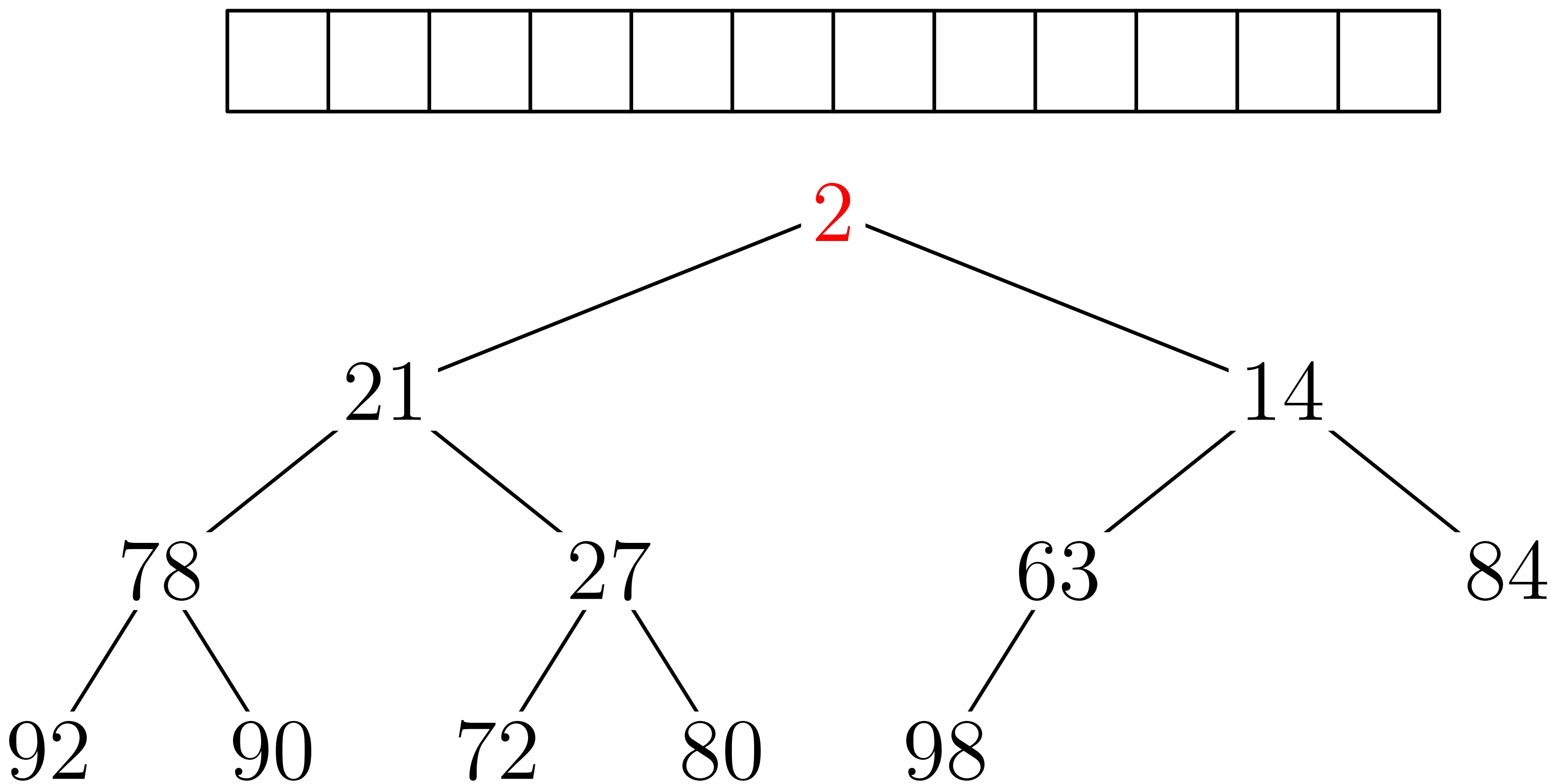

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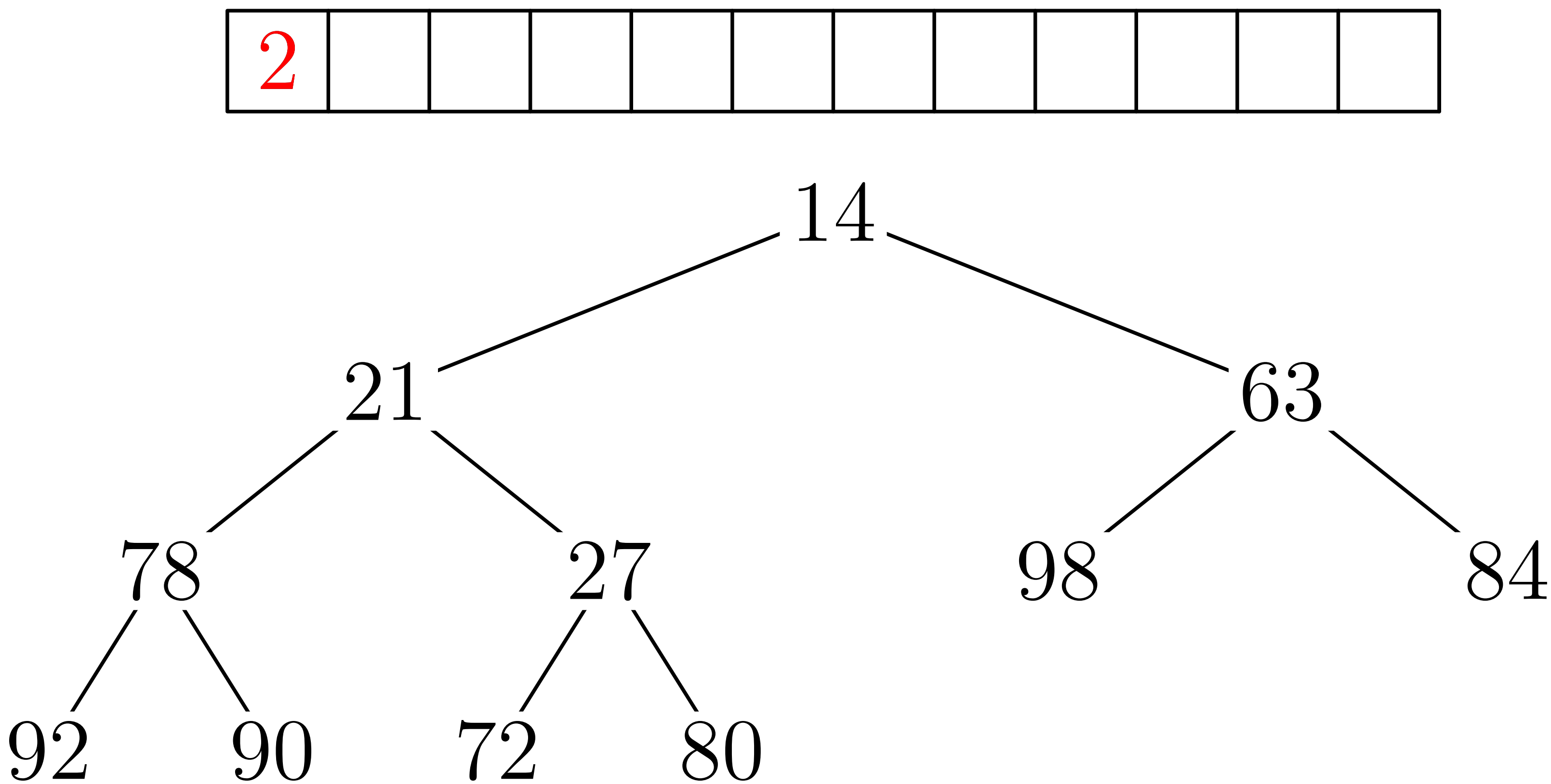
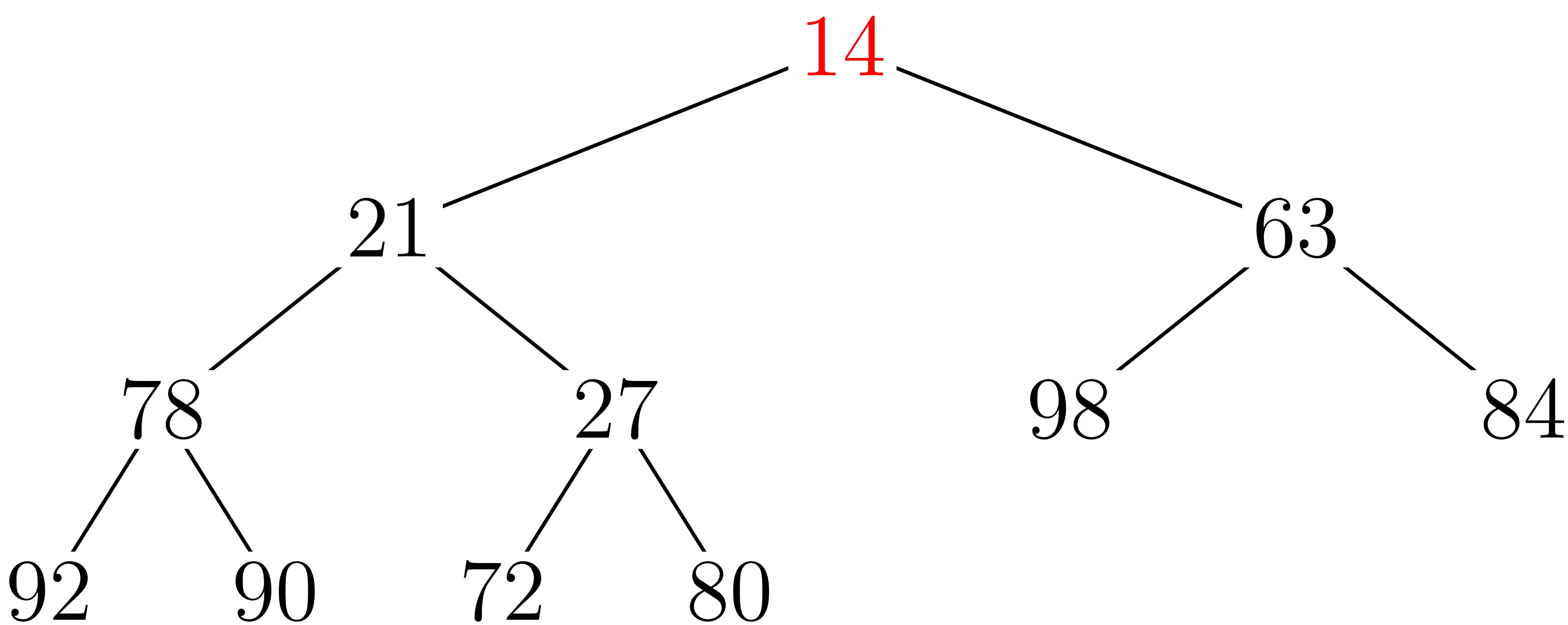


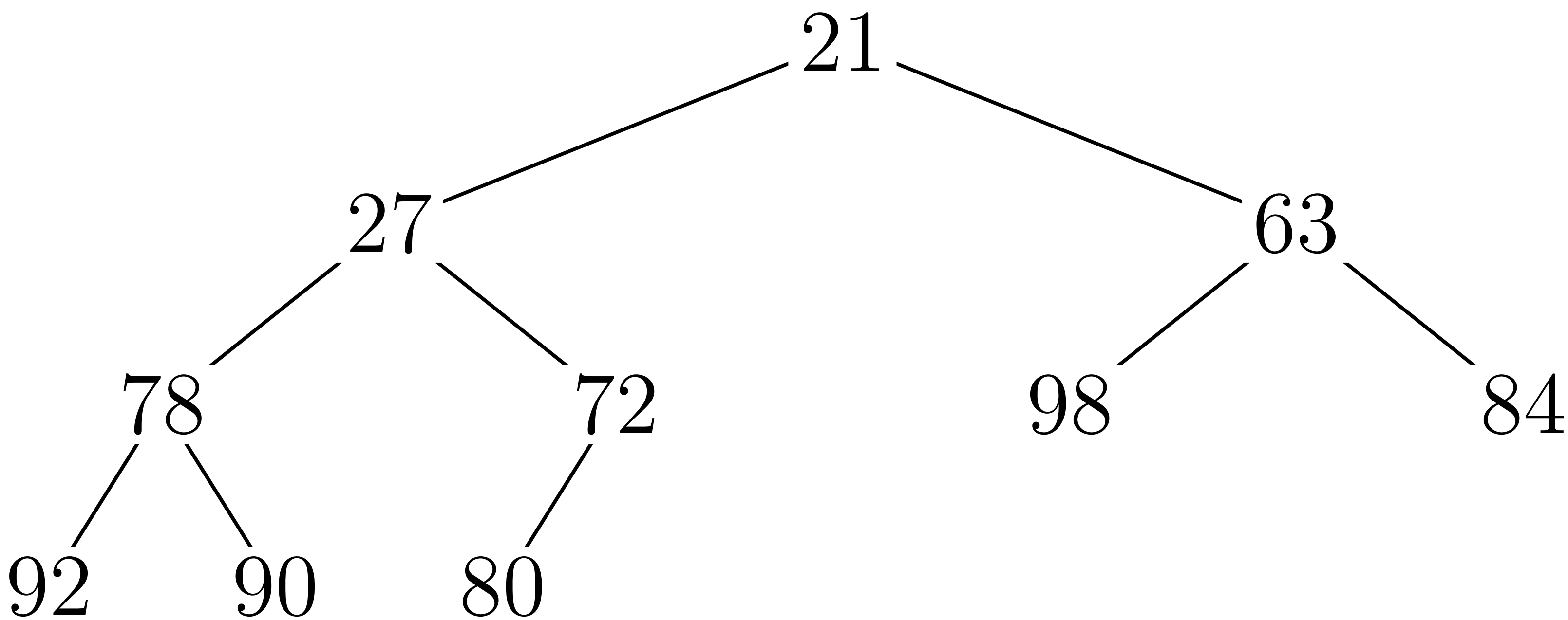
Example of Heap Sort

2											



Example of Heap Sort

2	14										



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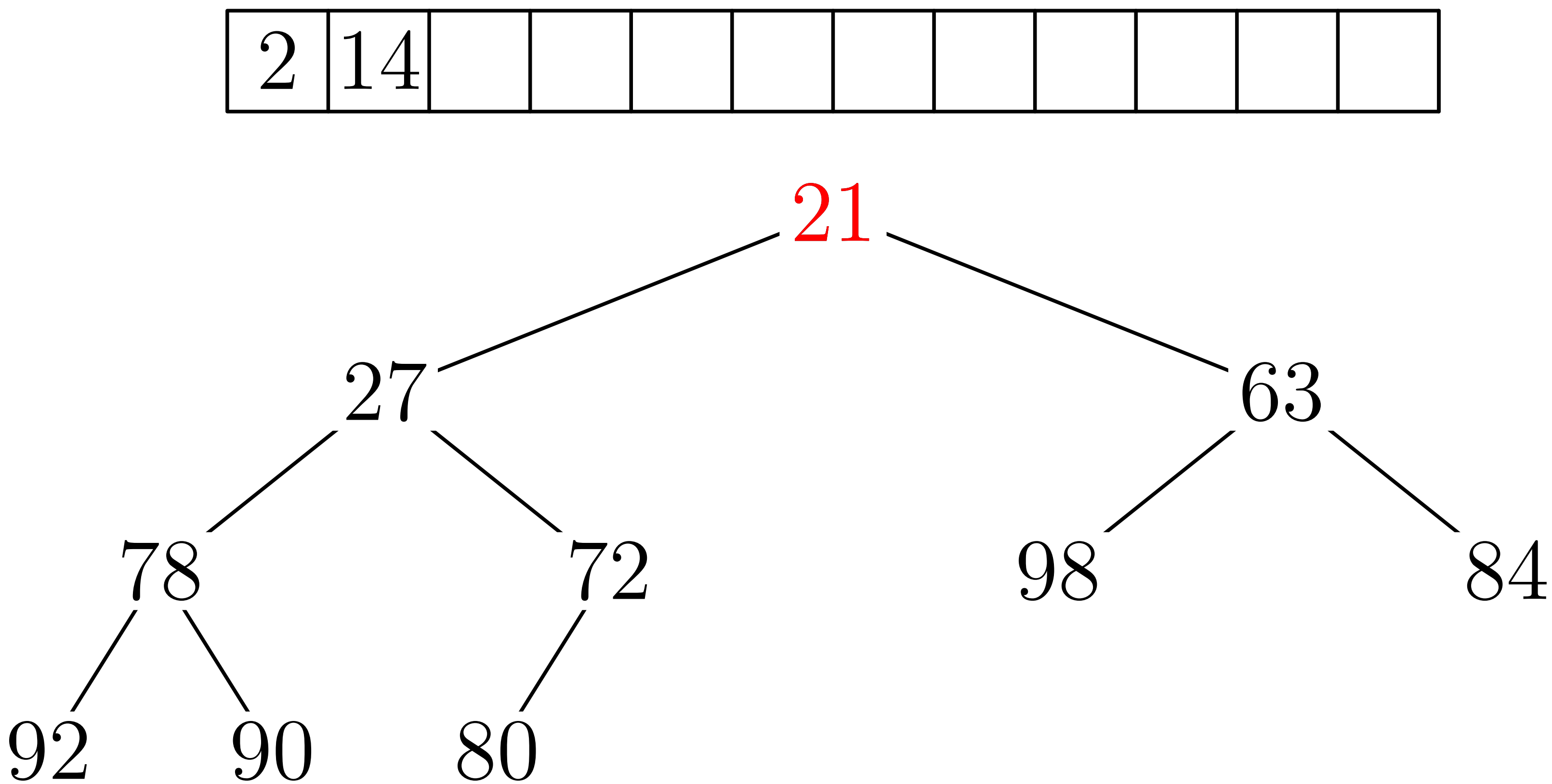
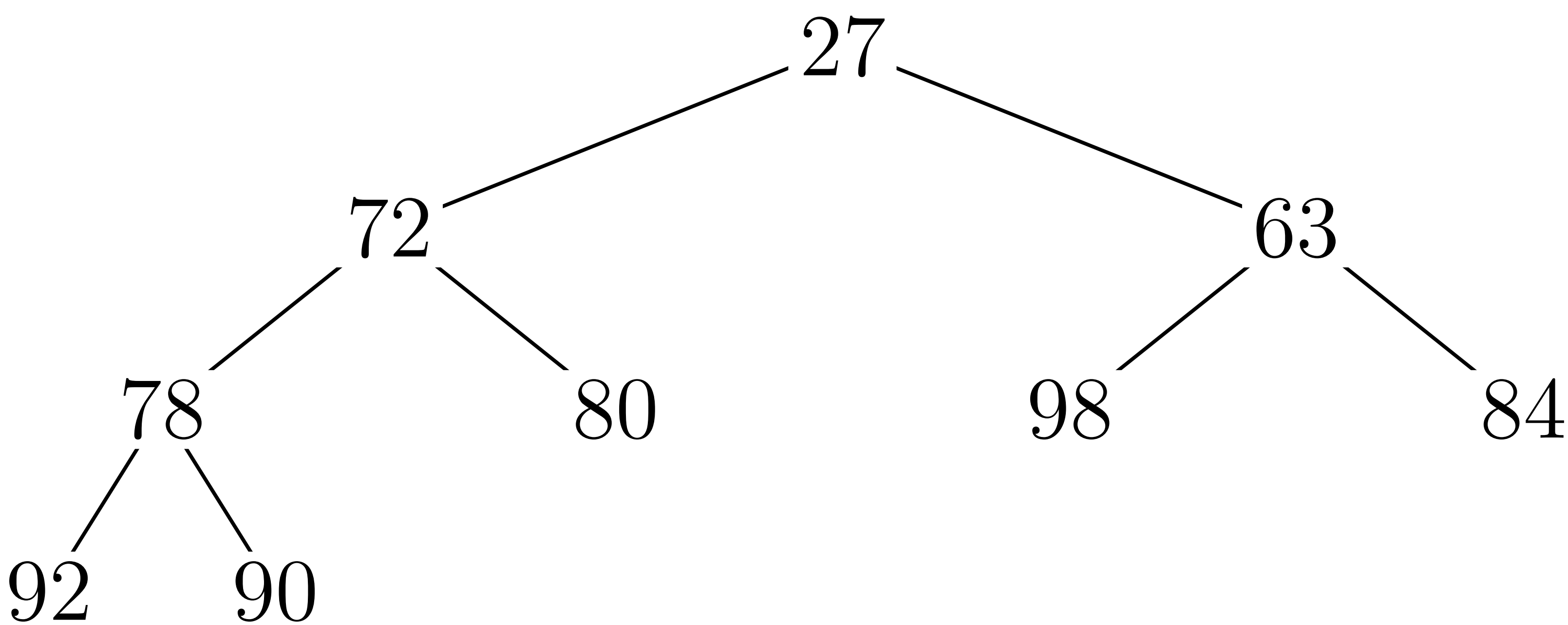


Example of Heap Sort

2	14	21									



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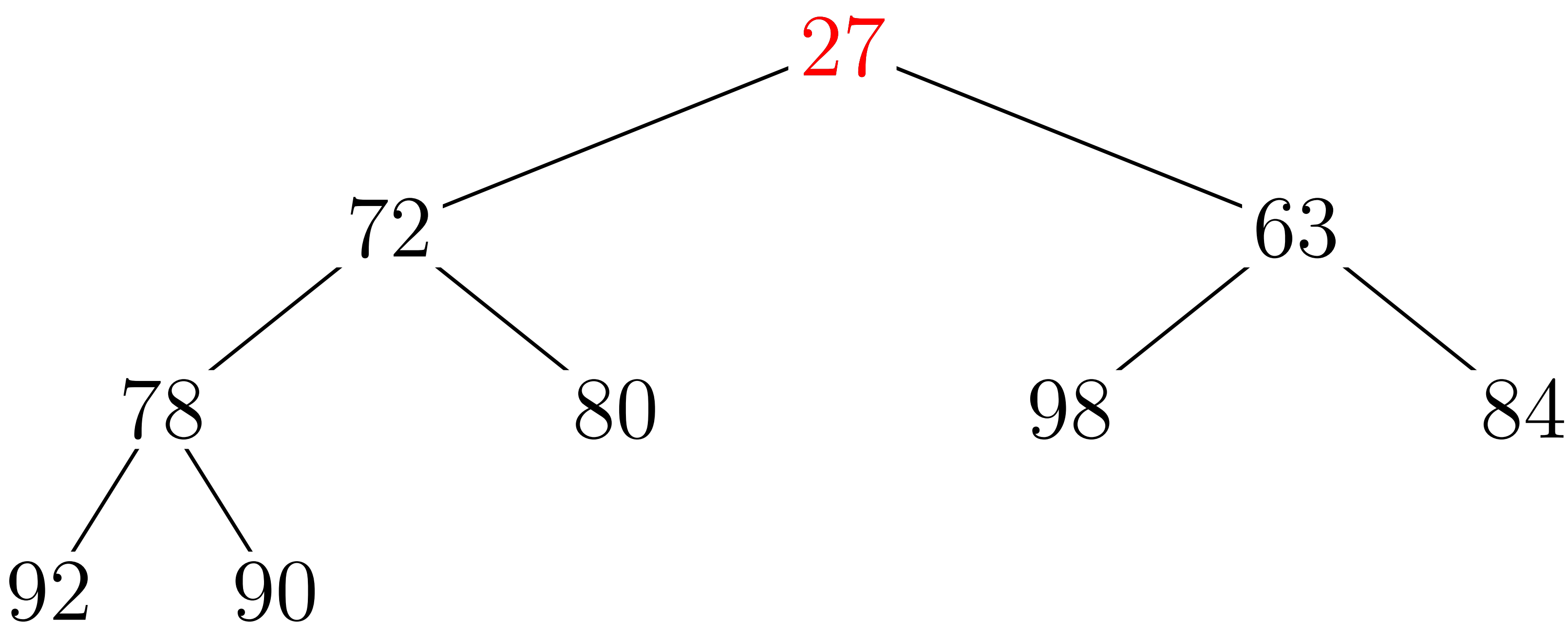
2	14	21									



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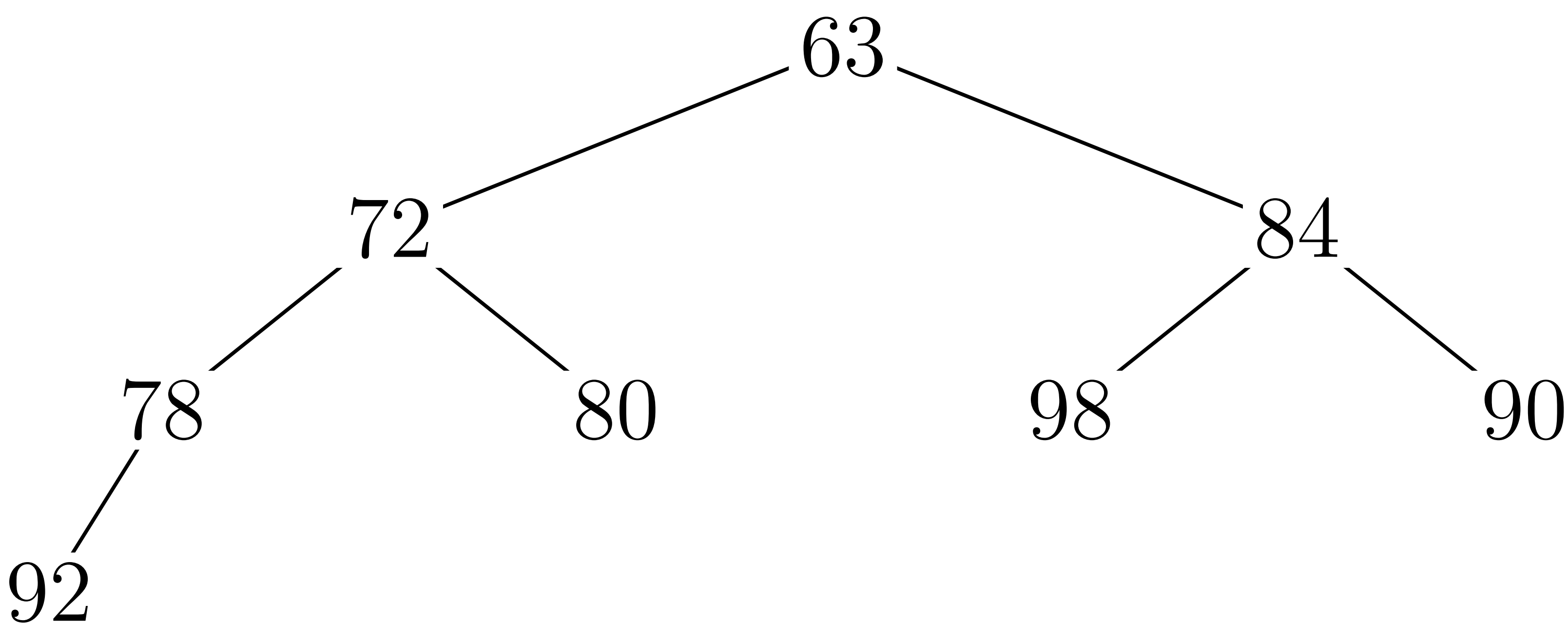
2	14	21	27								



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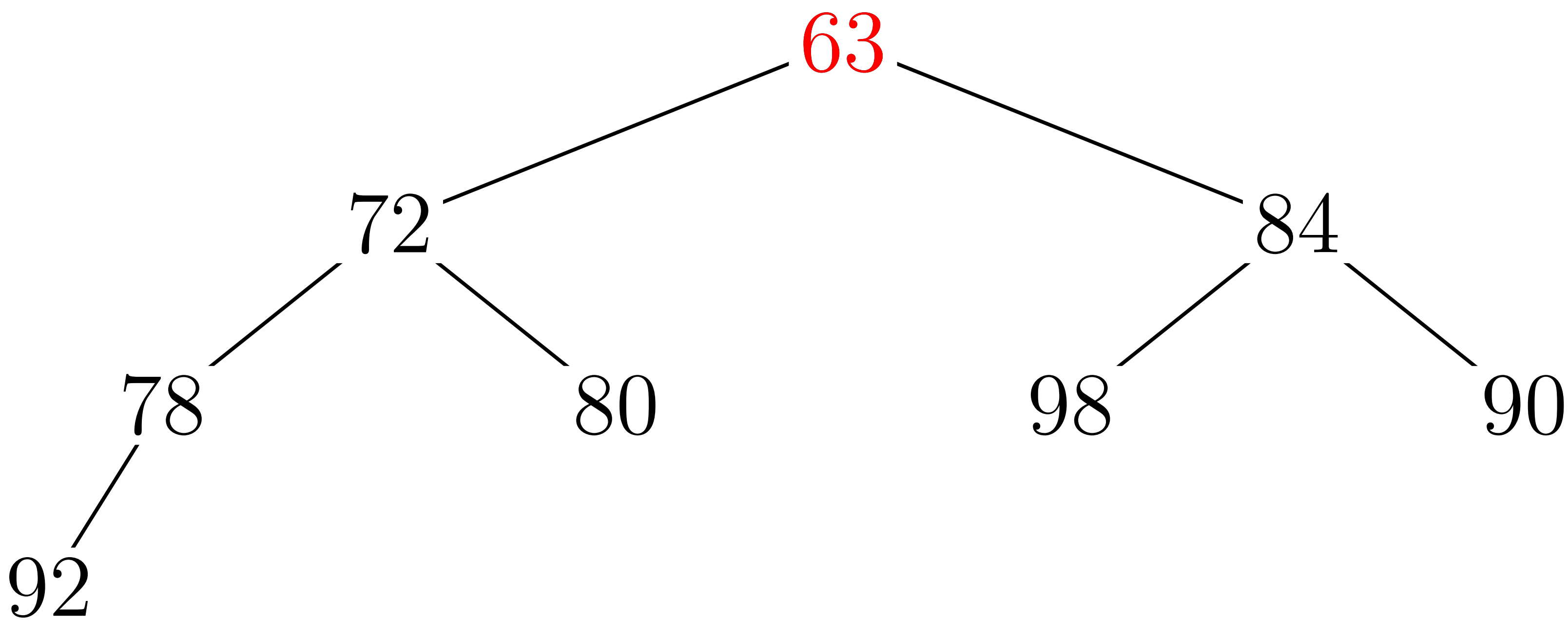

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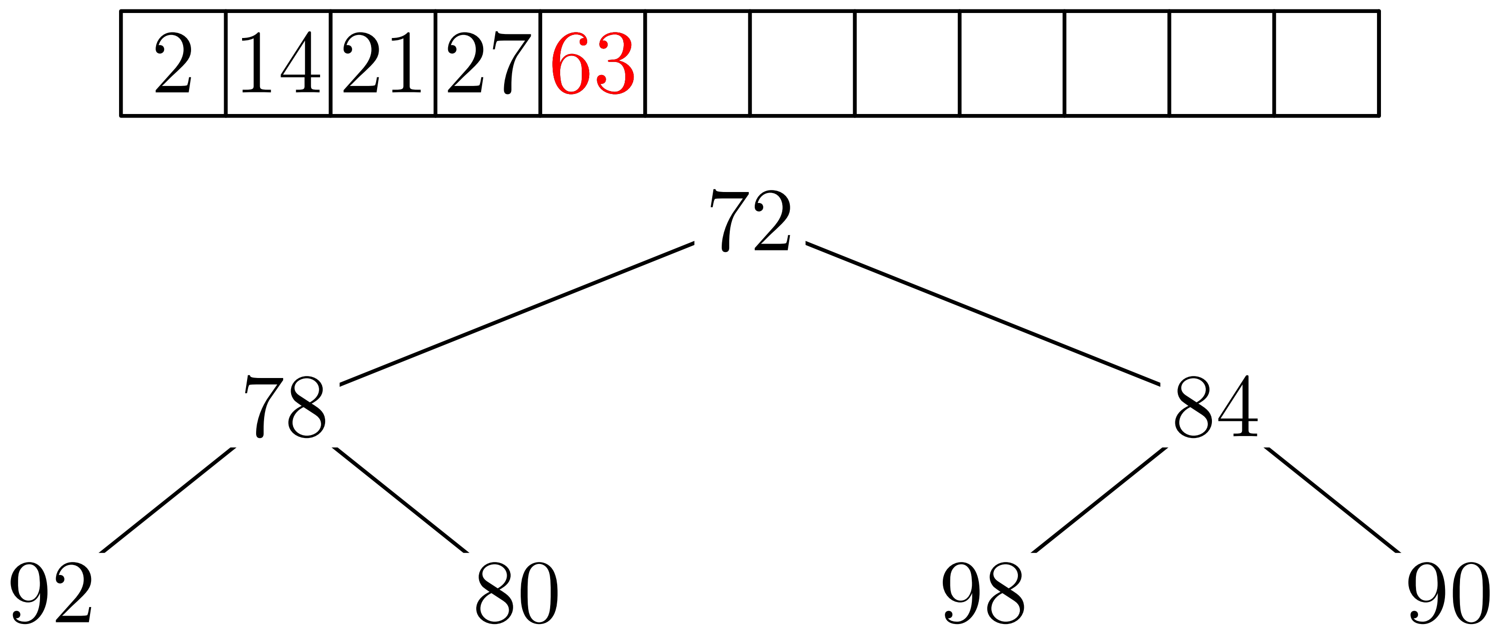

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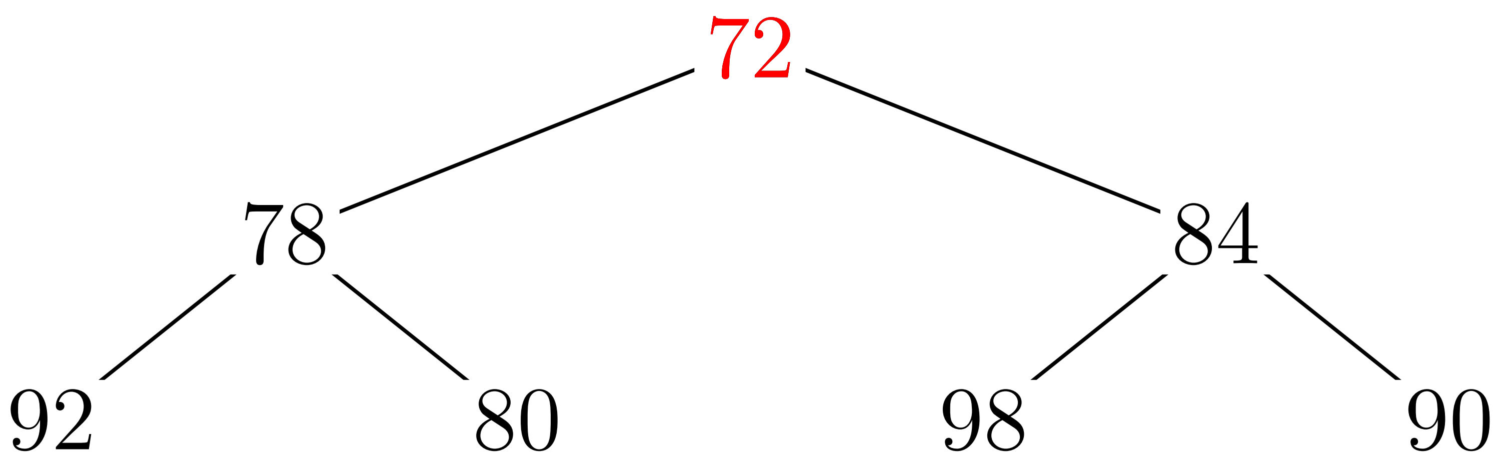
2	14	21	27	63							



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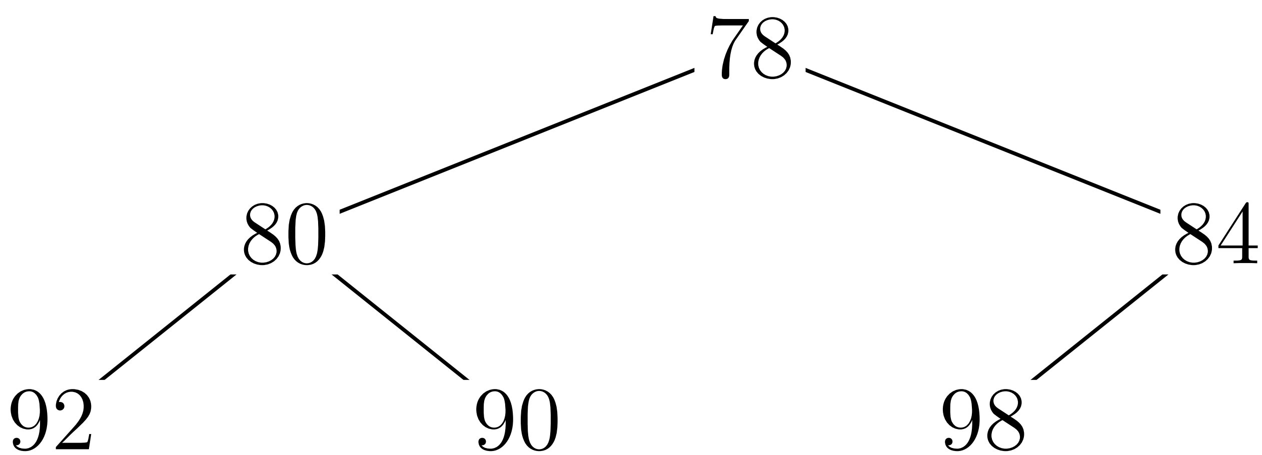
2	14	21	27	63	72						



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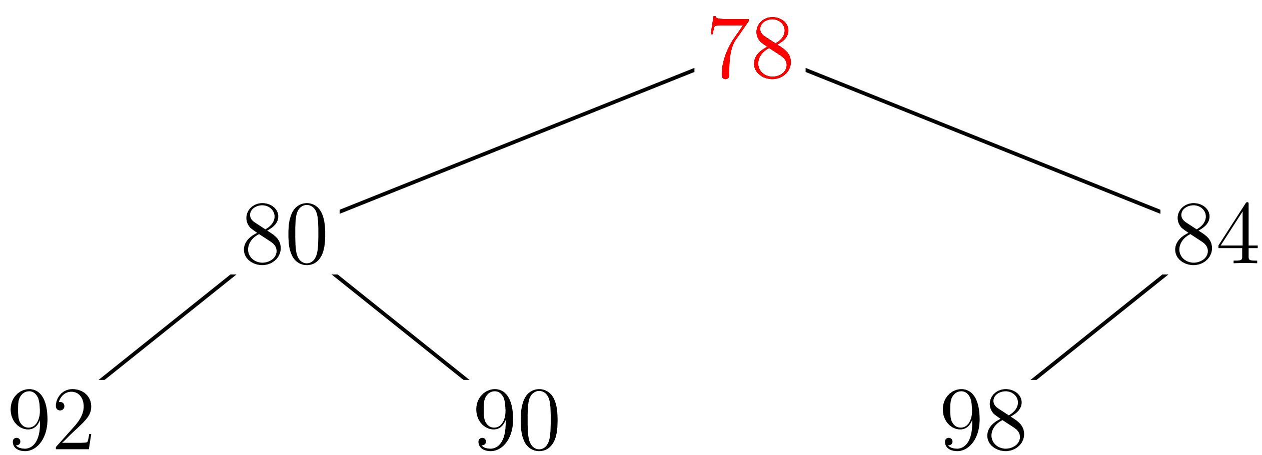

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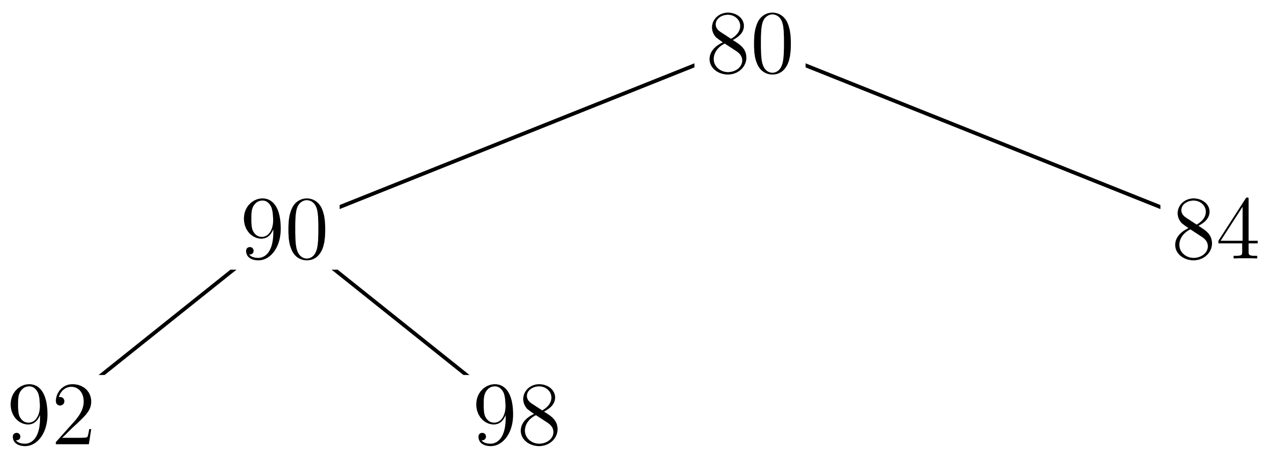

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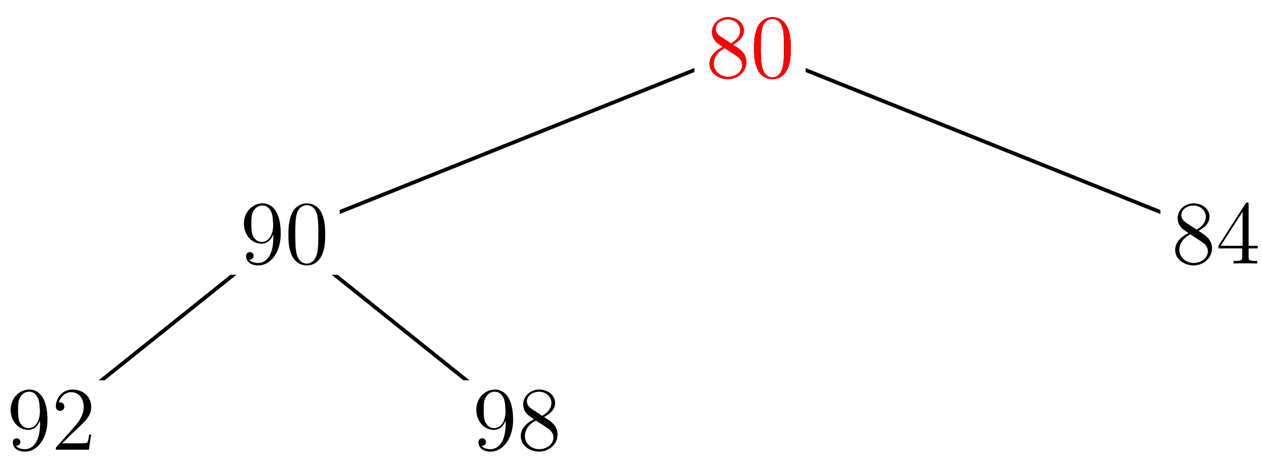

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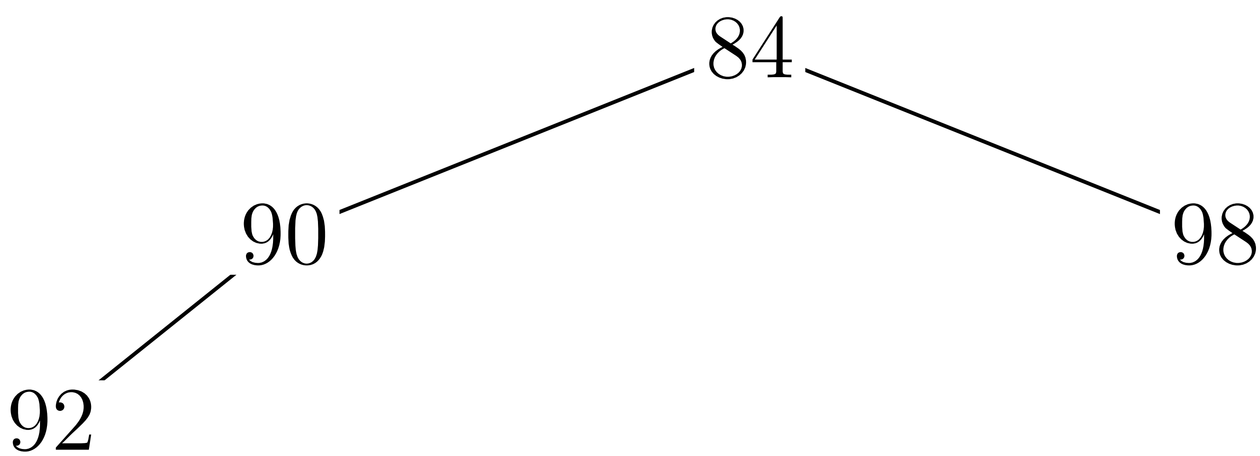
2	14	21	27	63	72	78	80				



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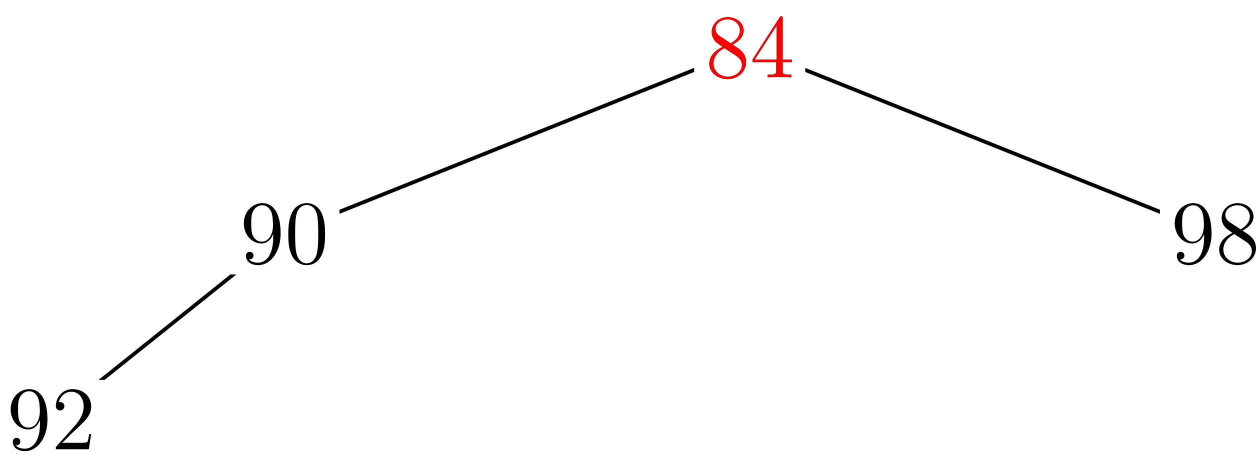
2	14	21	27	63	72	78	80				



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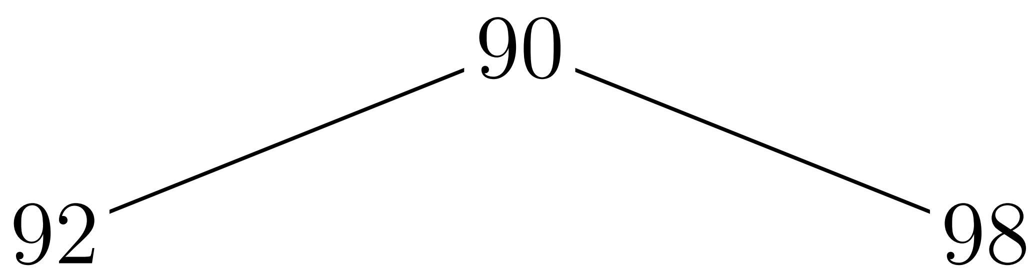
2	14	21	27	63	72	78	80	84			



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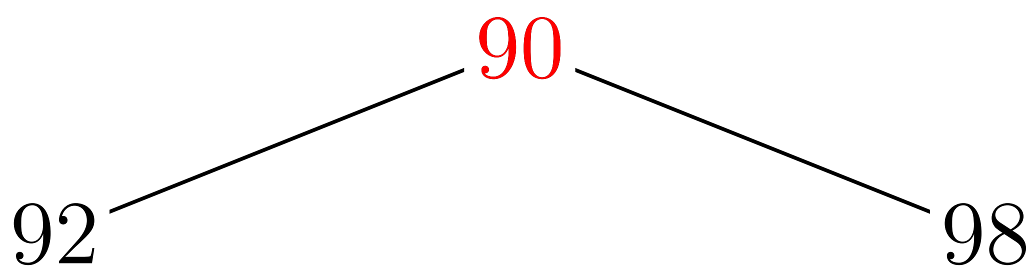

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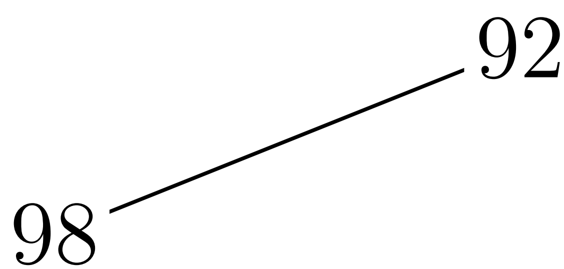

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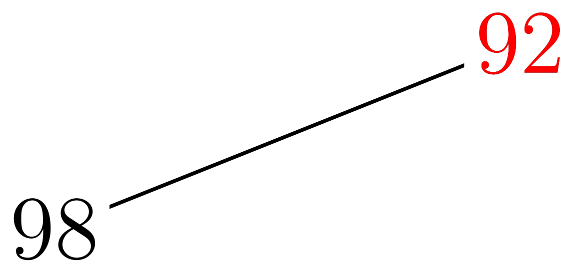

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98

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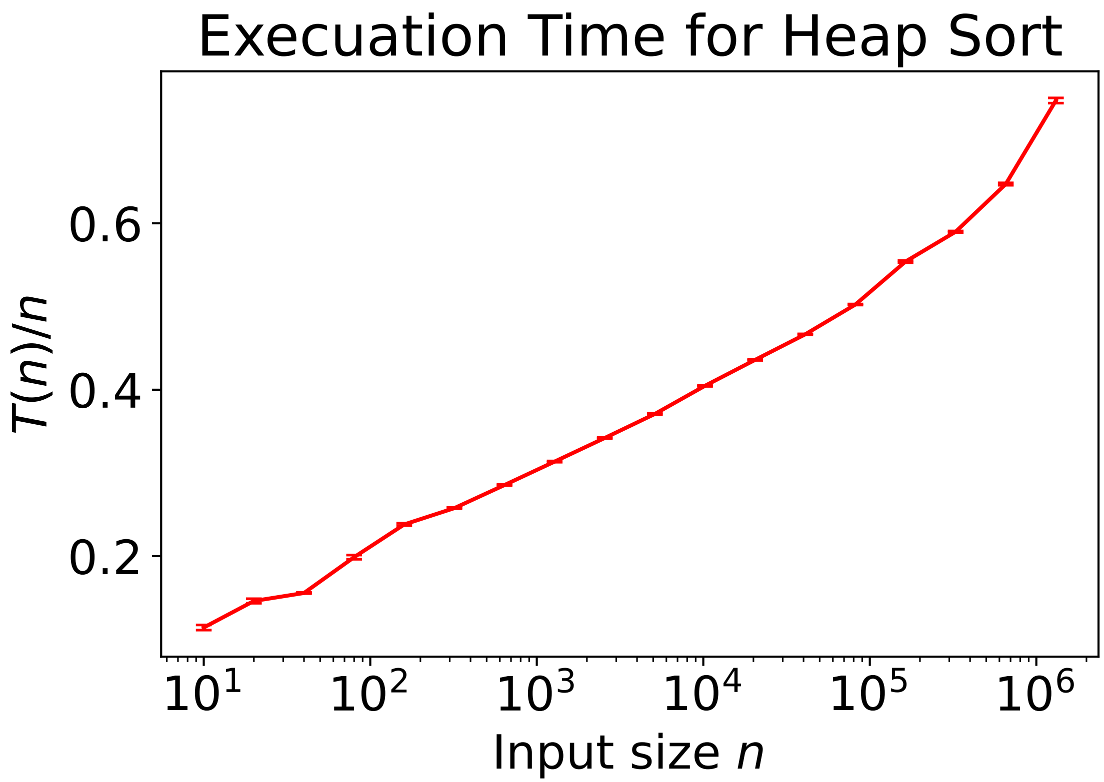
Complexity of Heap Sort

- As we have to add n elements and then remove n elements the time complexity is log-linear, i.e. $O(n\log(n))$



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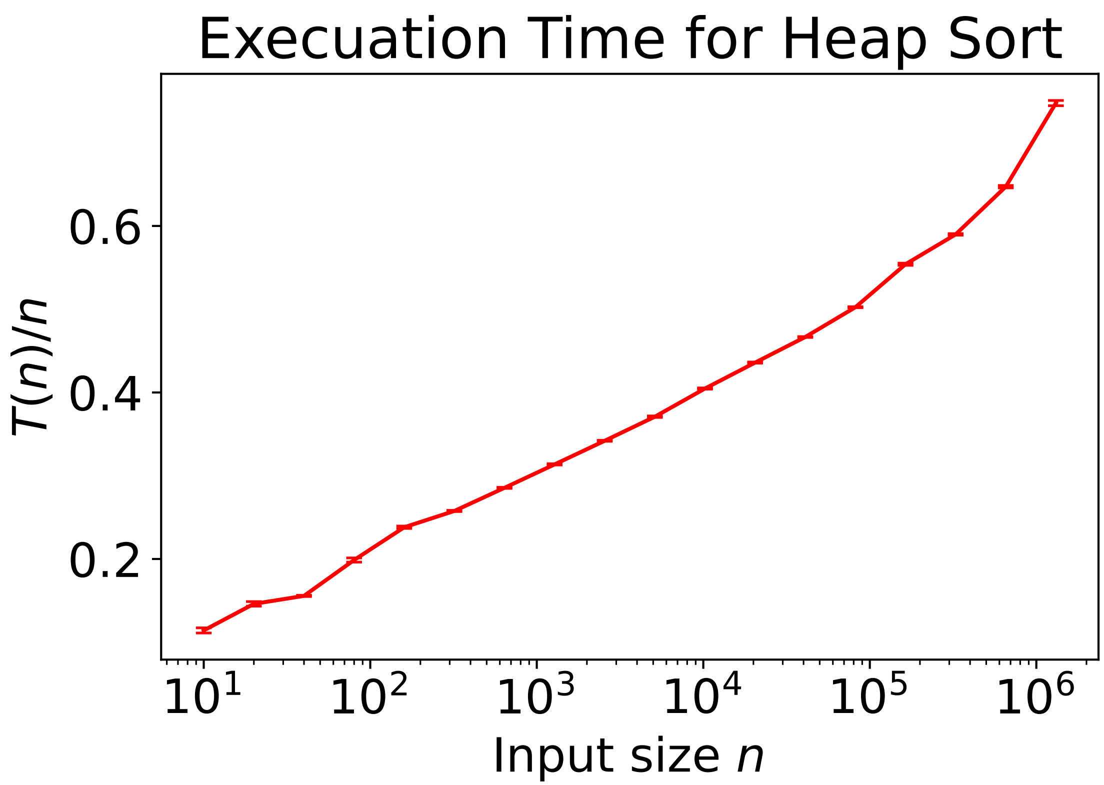

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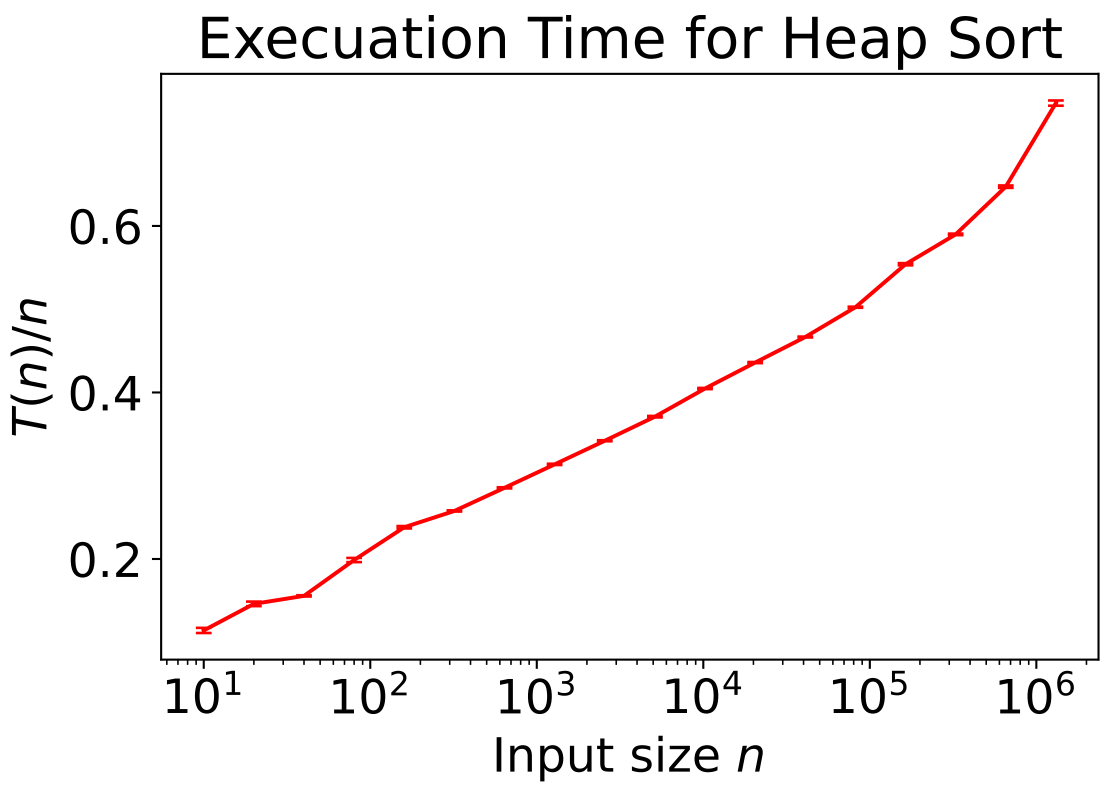

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Other Heaps

- Binary Heaps are so useful that other types of heaps have been developed
- The simplest enhancement is to combine a binary heap with a map which maintains a pointer to each element
- The map has to be updated every time elements are moved in the heap (fortunately only $O(\log(n))$ elements are move each time the heap is updated)
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Merging Heaps

- One common demand on a heaps is to merge two heaps
- Unfortunately binary heaps are not efficiently merged
- There are a variety of different heaps (leftist heaps, skew heaps, binomial queues, . . .) designed to be merged
- All these heaps are real binary trees (i.e. represented by pointers rather than put in an array)
- They are slower than a binary heap because indexing is slightly slower as is creating node objects
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Other Operations

- All other operations are achieved by merging
  - ⋆ Adding an element is achieved by merging the current heap with a heap of one element
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- Heaps are binary trees that can be implemented as arrays
- Priority queue have many (often surprising uses)
  - ⋆ They are used when you need a queue with priorities, e.g. in operating systems
  - ⋆ They can be used to perform pretty efficient sort
  - ⋆ They are often used for implementing greedy type algorithms
  - ⋆ One important application is in real time simulations
- There exists many extensions of heaps

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- Priority queue have many (often surprising uses)
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  - ⋆ They can be used to perform pretty efficient sort
  - ⋆ They are often used for implementing greedy type algorithms
  - ⋆ One important application is in real time simulations
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