

Further Mathematics and Algorithms

Lesson 10: *Make Friends with Trees*



Binary trees, binary search trees, sets, tree iterators

Outline

1. **Trees**
2. Binary Trees
 - Implementing Binary Trees
3. Binary Search Trees
 - Definition
 - Implementing a Set
4. Tree Iterators



Trees

- Trees are one of the major ways of structuring data
- They are used in a vast number of data structures
 - ★ Binary search trees
 - ★ B-trees
 - ★ splay trees
 - ★ heaps
 - ★ tries
 - ★ suffix trees
- We shall cover most of these

Trees

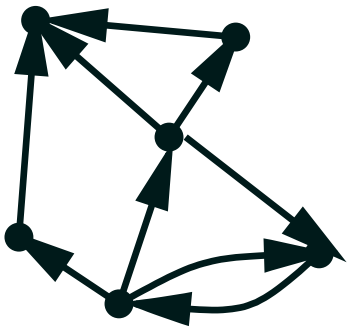
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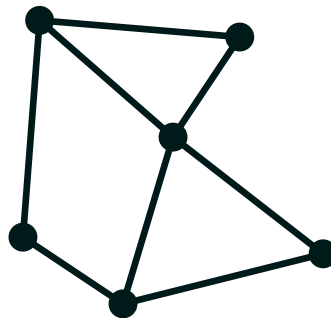
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Defining Trees

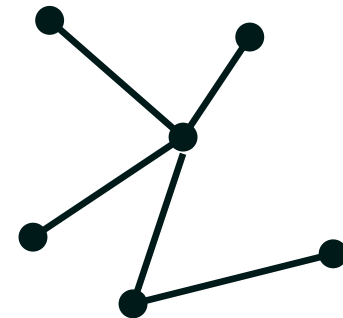
- Mathematically a tree is an **acyclic undirected graph**
 - ★ **graph**: a structure consisting of **nodes** or **vertices** joined by **edges**
 - ★ **undirected**: the edges goes both ways
 - ★ **acyclic**: there are no cycles in the graph



graph



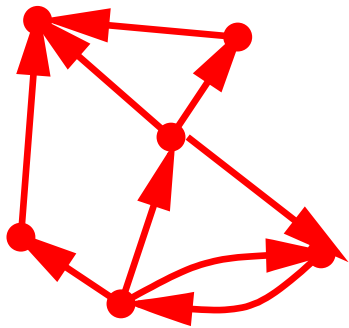
undirected graph



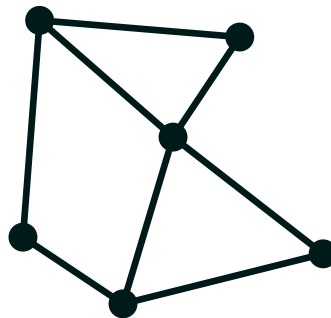
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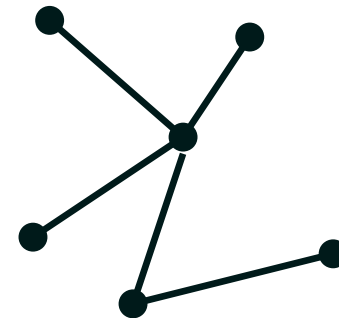
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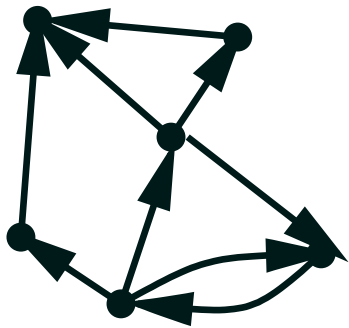
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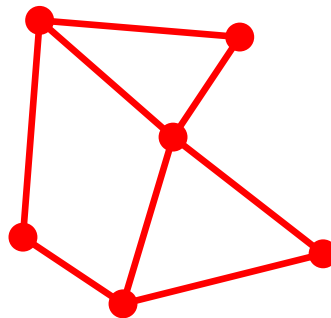
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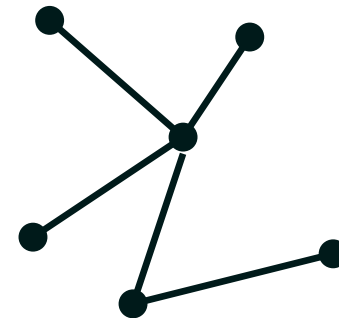
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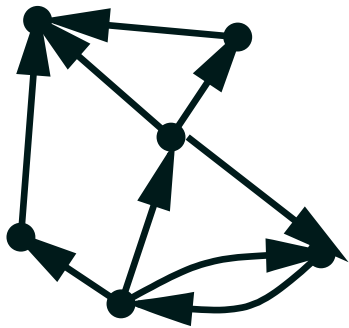
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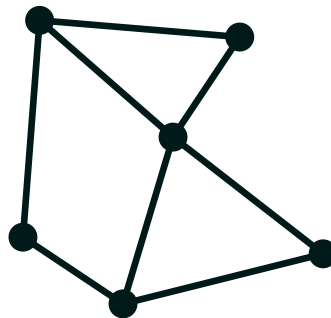
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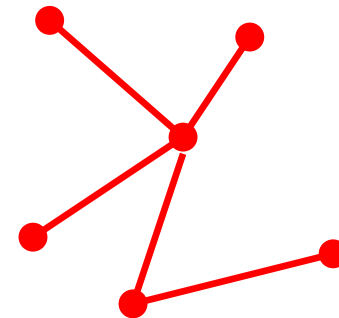
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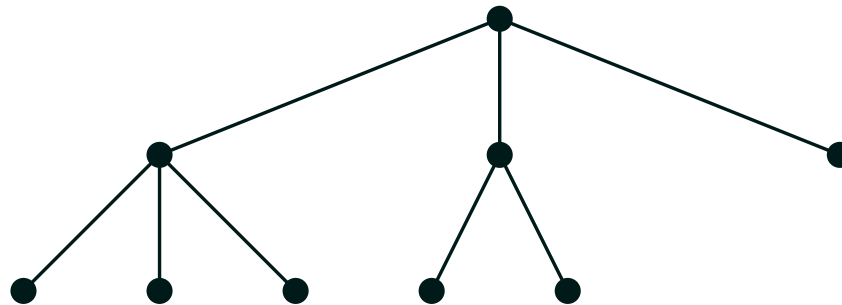
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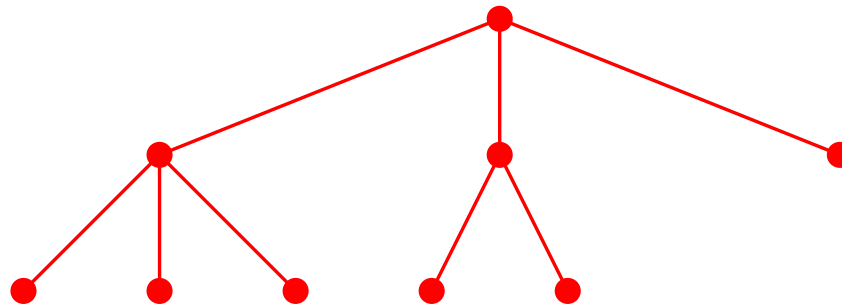
Borrowing from Nature

- We often impose an ordering on the nodes (or a direction on the edges)—known as a rooted tree
- Borrowing from nature, we recognise one node as the **root** node
- Nodes have **children** nodes living beneath them
- Each child has a **parent** node above them except the root
- Nodes with no children are **leaf** nodes



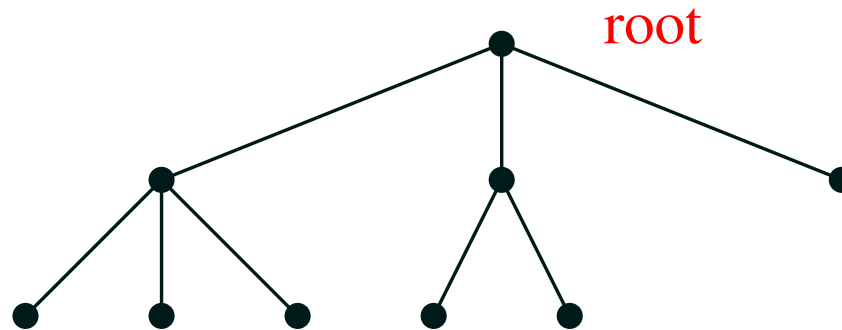
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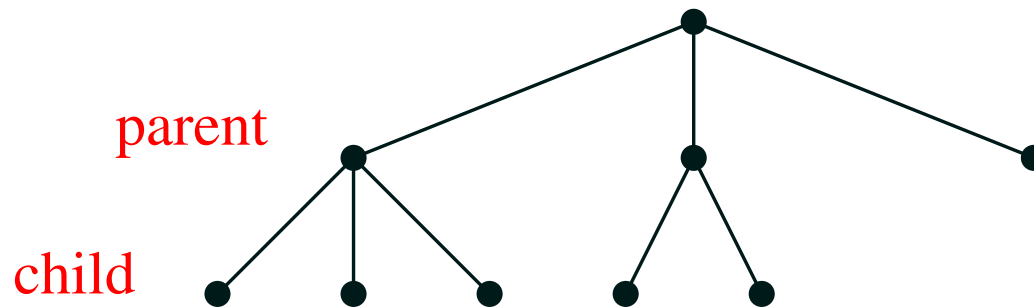
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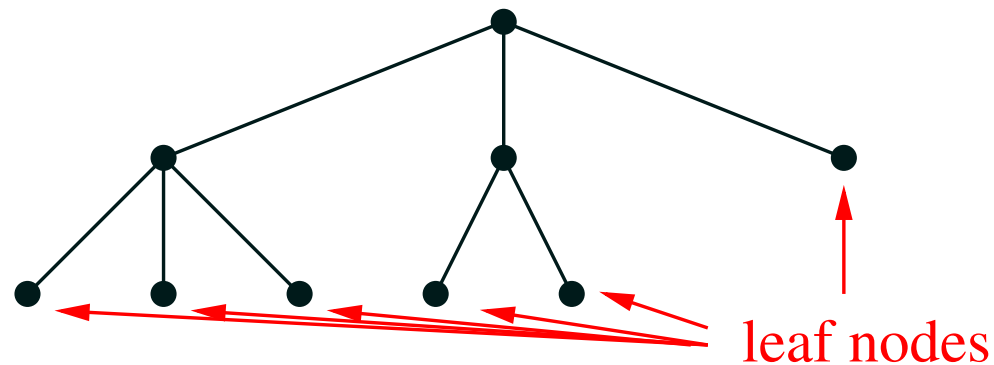
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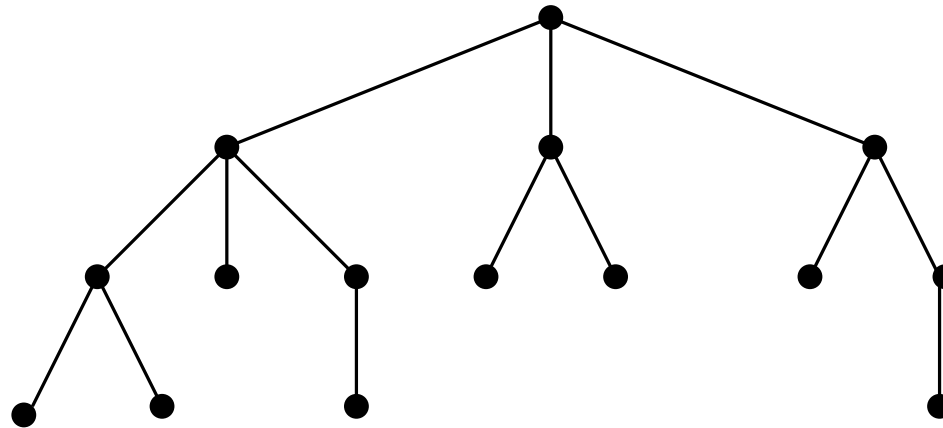


Spot the Error

- One small biological inconsistency

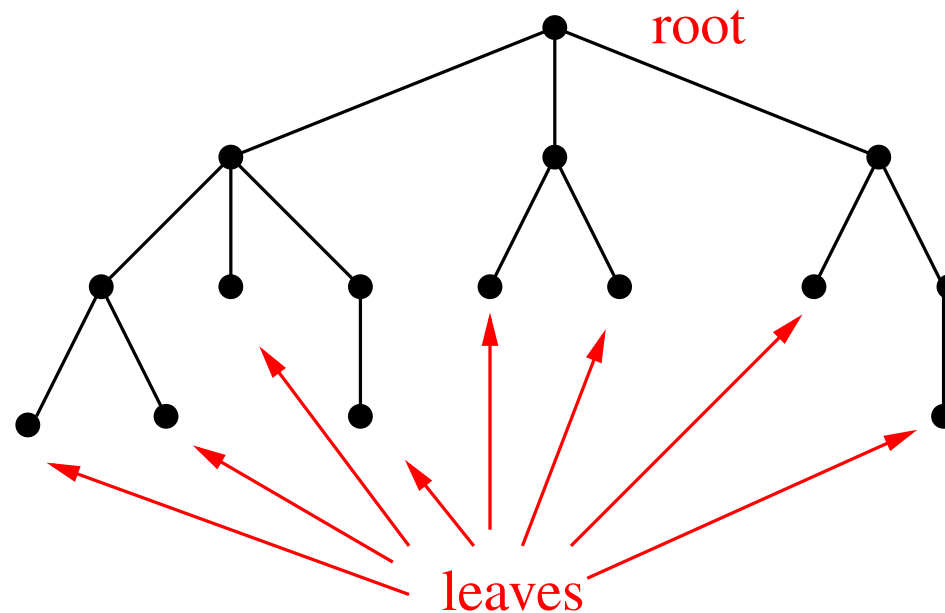
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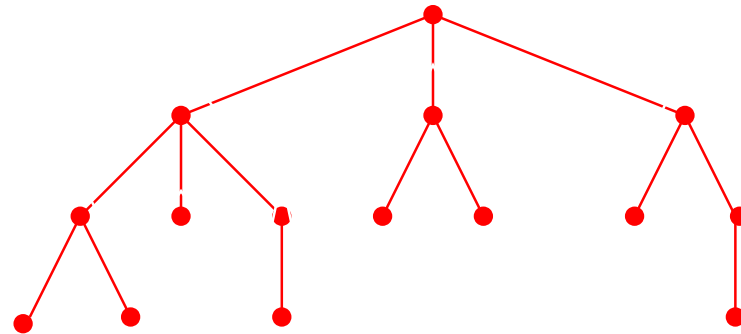
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- Yep!, computer scientists draw there trees upside down
 - ★ root at the top
 - ★ leaves at the bottom



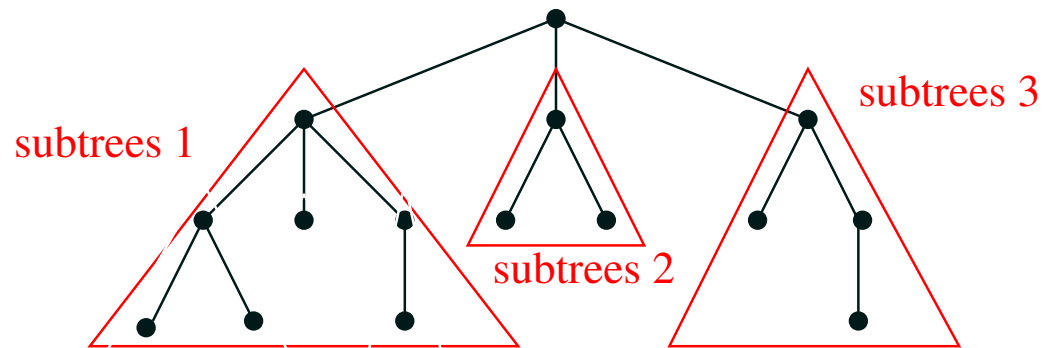
Subtrees

- We can think of the tree made up of **subtrees**



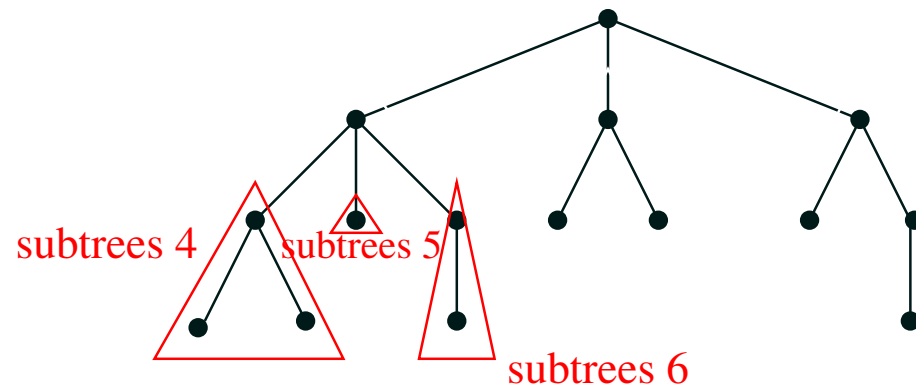
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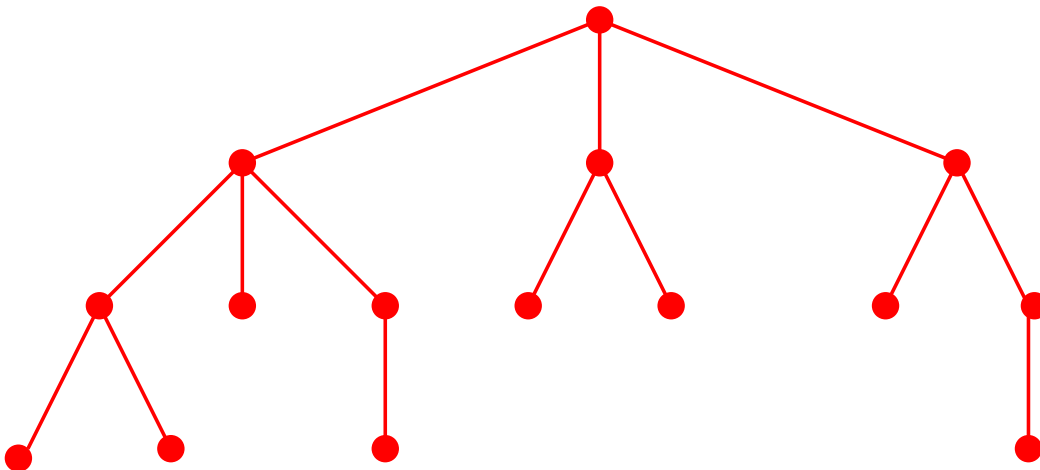
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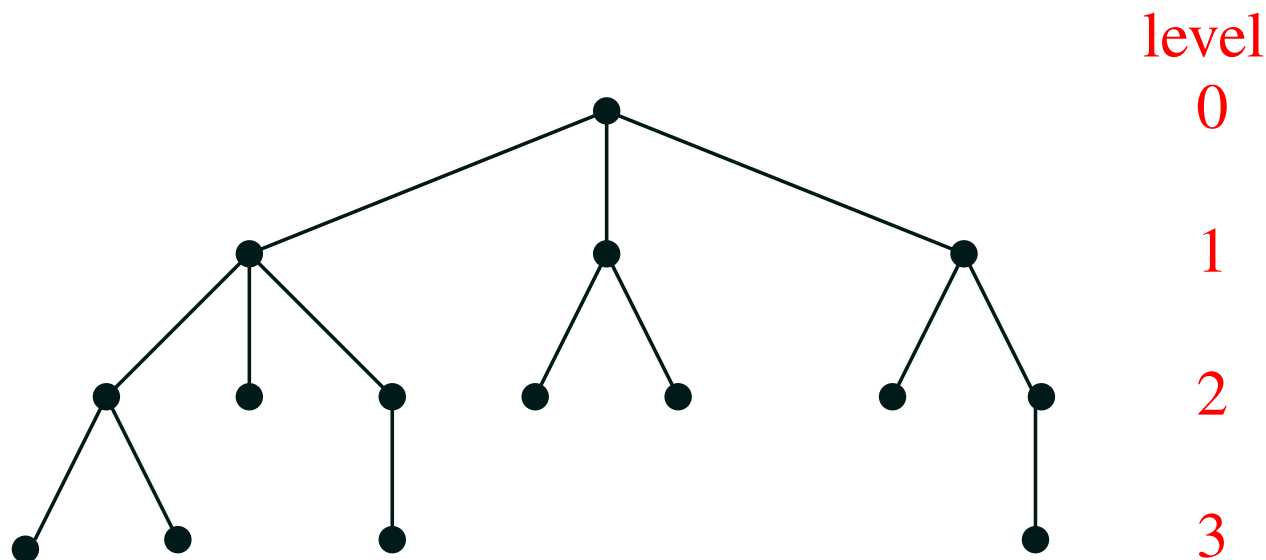
Level of Nodes

- It is useful to label different levels of the tree
- We take the **level** of a node in a tree as its distance from the root
- We take the **height** of a tree to be the number of levels



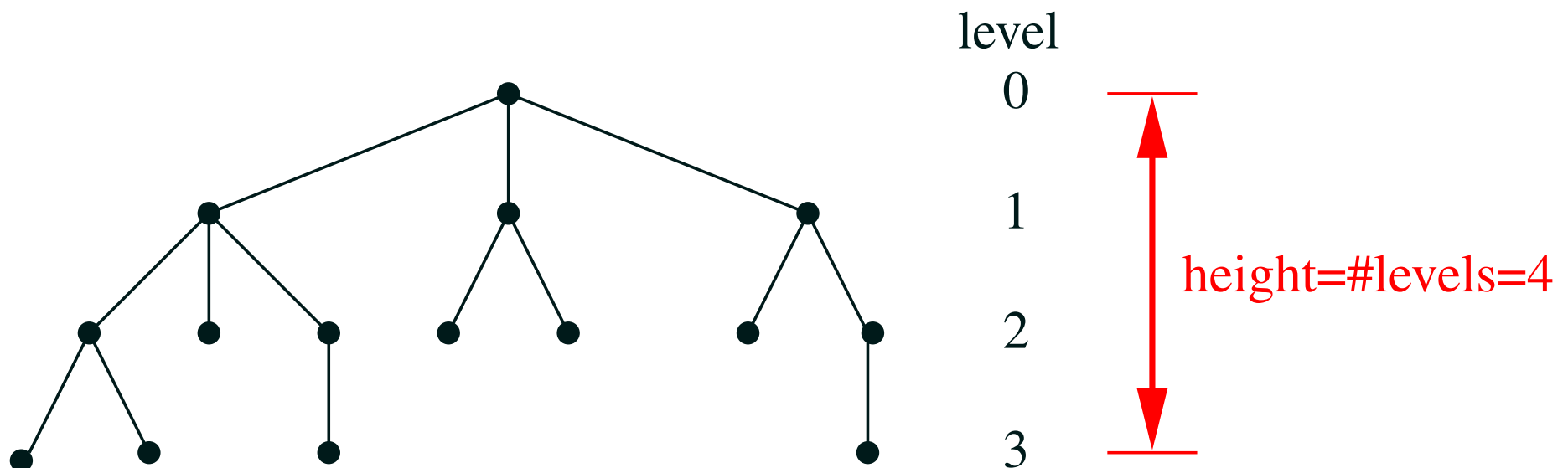
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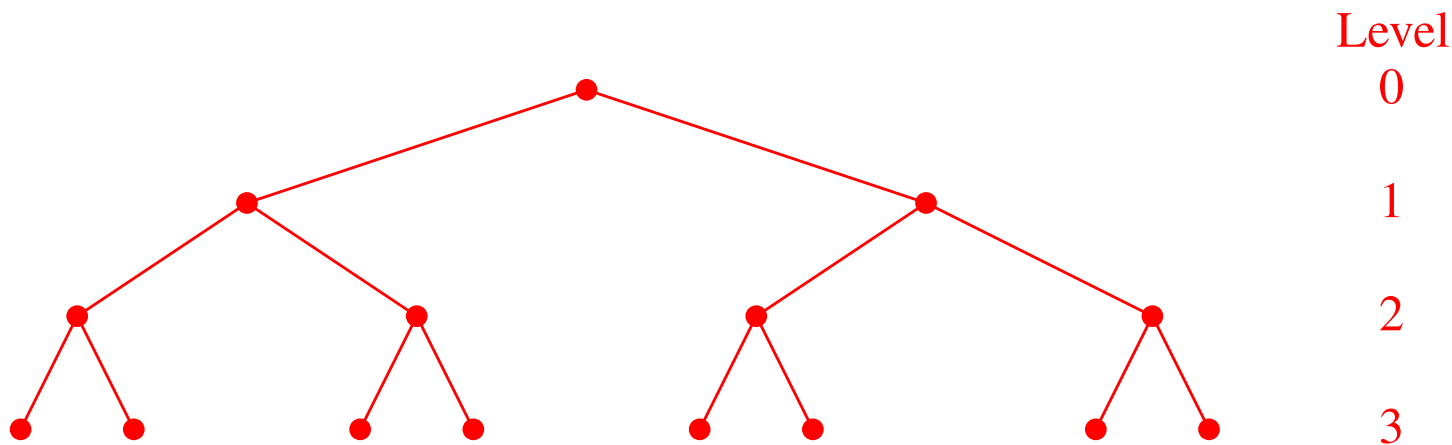
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2. **Binary Trees**
 - Implementing Binary Trees
3. Binary Search Trees
 - Definition
 - Implementing a Set
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Binary Trees

- A **binary tree** is a tree where each node can have zero, one or two children
- The total number of possible nodes at level l is 2^l
- The total number of possible nodes of a tree of height h is

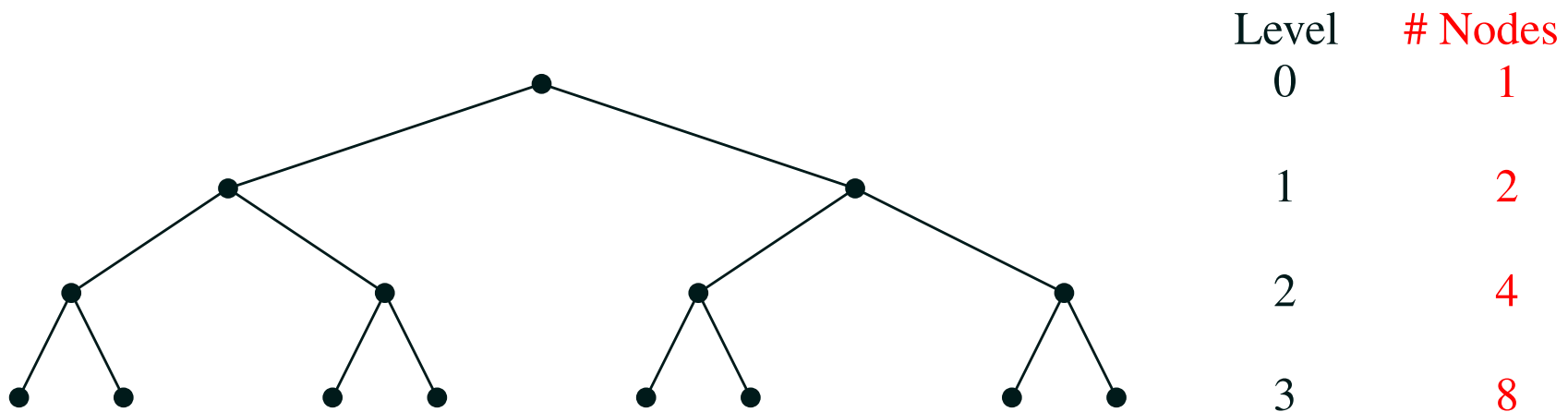
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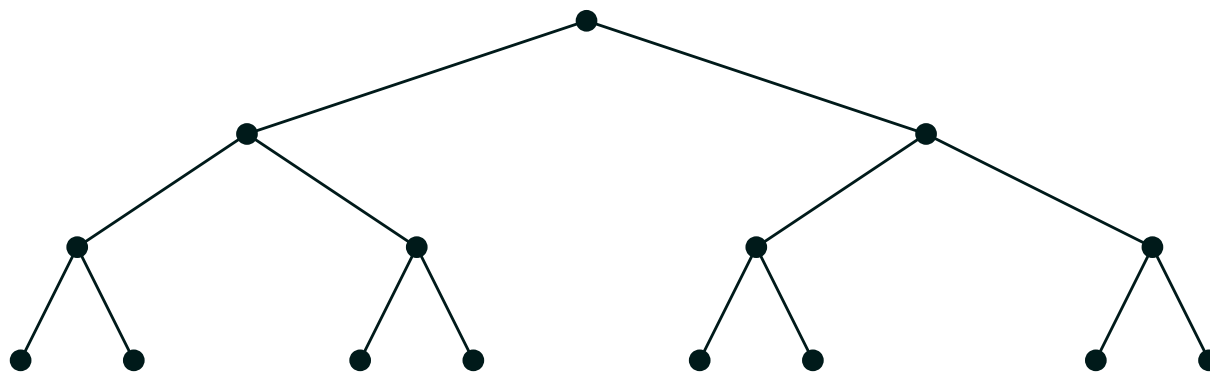
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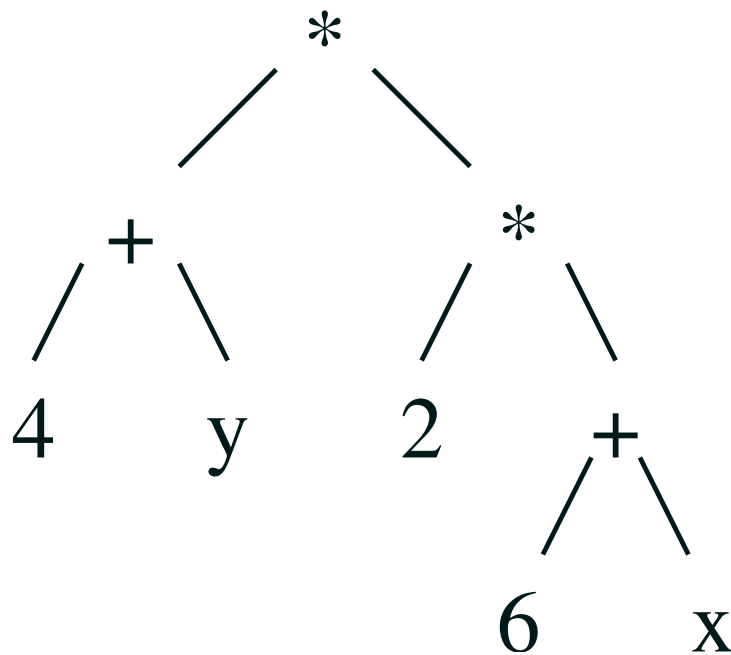
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Level	# Nodes
0	1
1	2
2	4
3	8
	<u>15</u>

Uses of Binary Trees

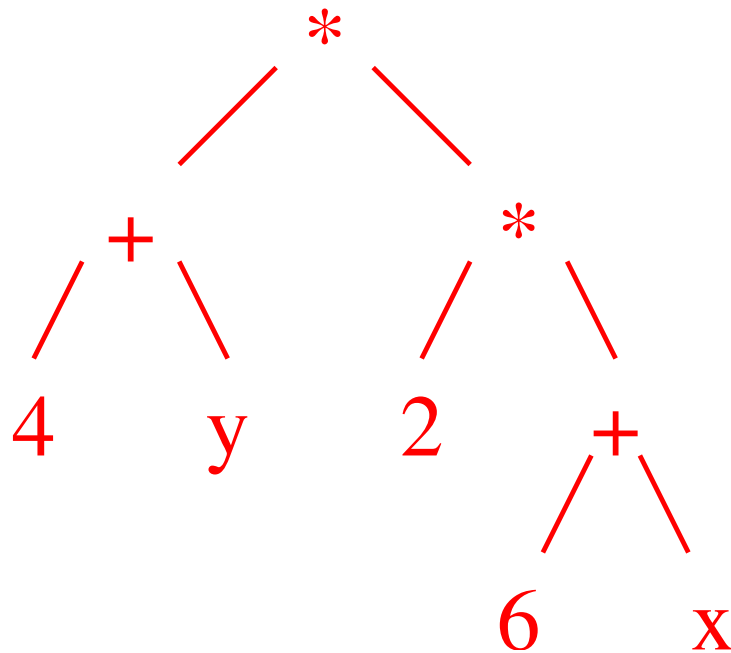
- Binary trees have a huge number of applications
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Implementation

- We wish to build a generic binary tree class with each node housing an element
- Again we use a `Node<T>` class as the building block for our data structure—in this case a node of the tree
- The `Node<T>` class will contain a pointer to left and right children
- To help navigate the tree each node will contain a pointer to its parent

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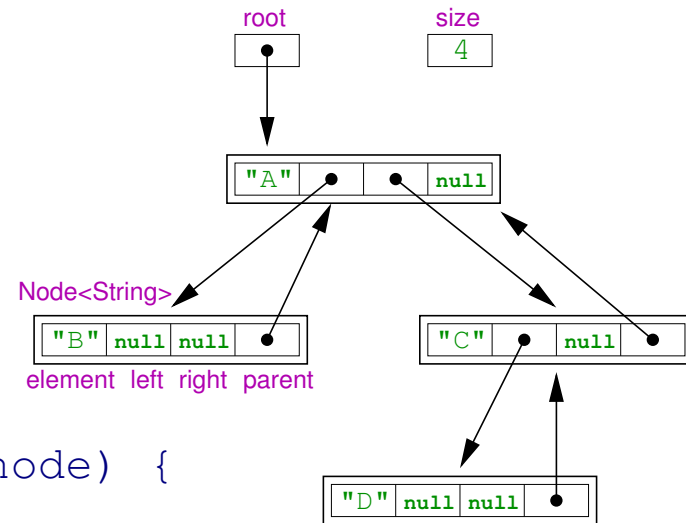
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C++ Code

```
template <typename T>
class binary_tree {
private:
    class Node {
    public:
        T element;
        Node* parent;
        Node* left = 0;
        Node* right = 0;

        Node(const T& value, Node* parent_node) {
            element = value;
            parent = parent_node;
        }
    };
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```
unsigned no_elements = 0;
Node* root = 0;
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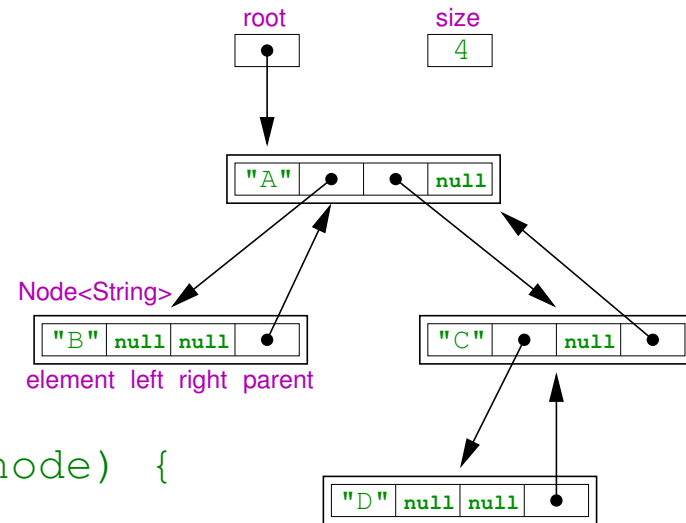


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Binary Search Trees

- We will concentrate on one of the most important binary trees, namely the **binary search tree**
- The binary search tree keeps the elements ordered
- We can define a binary search tree recursively
 1. Each element in the left subtree is less than the root element
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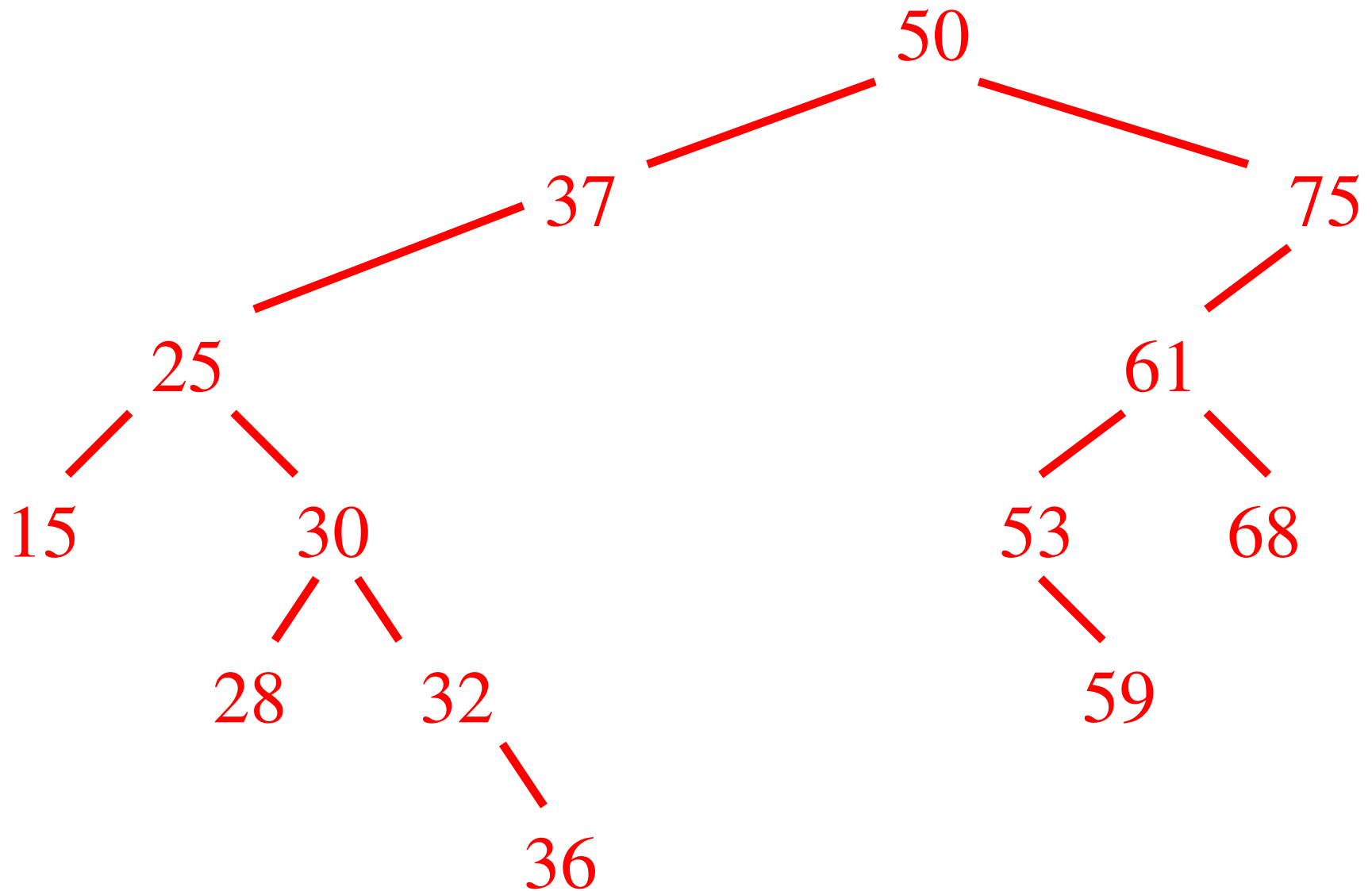
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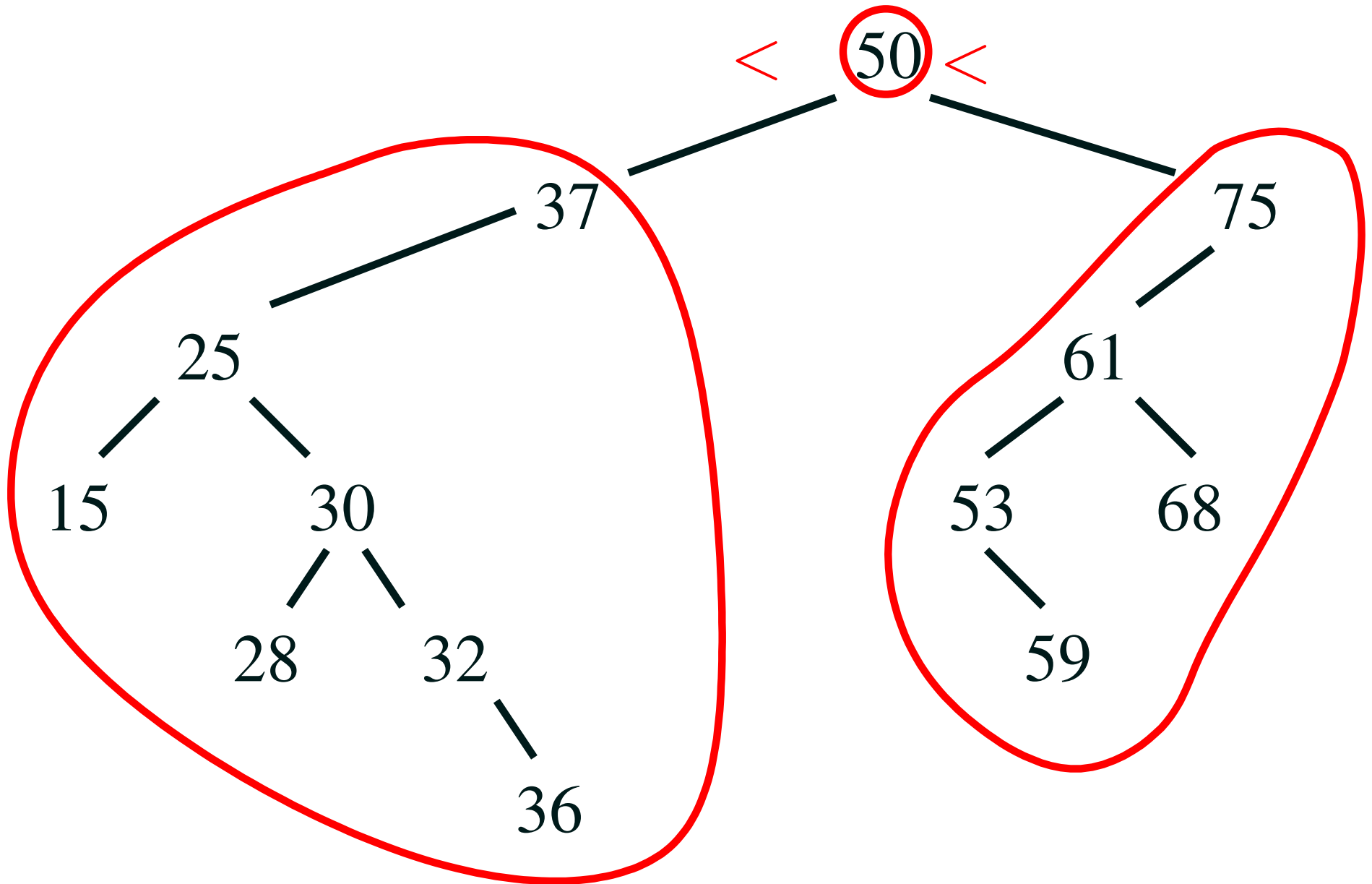
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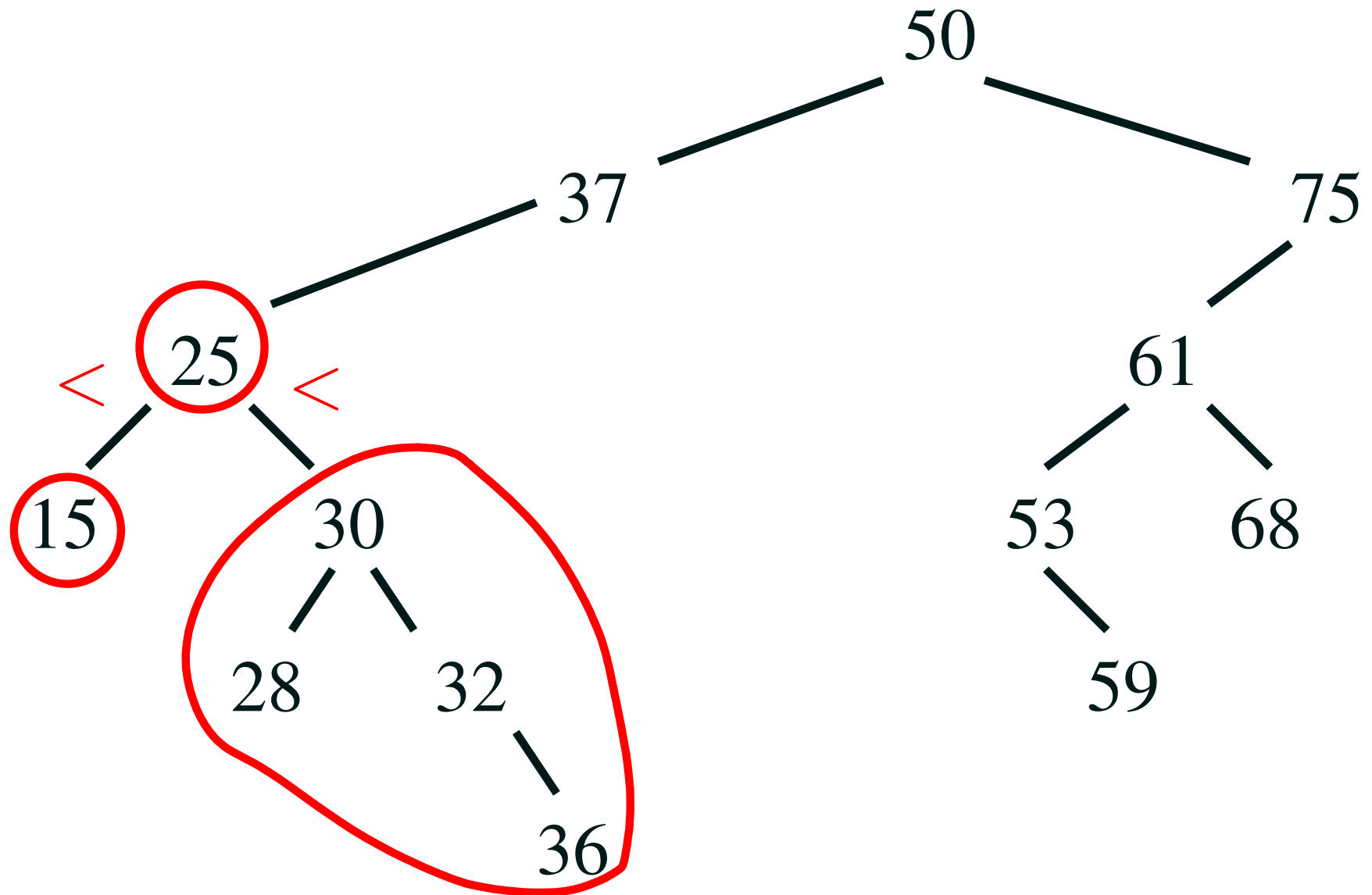
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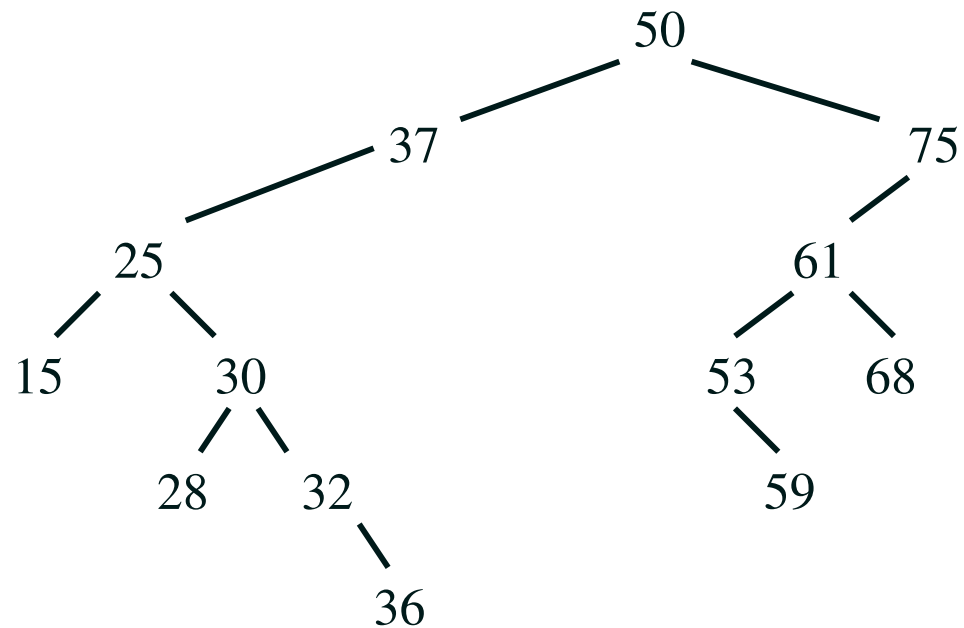


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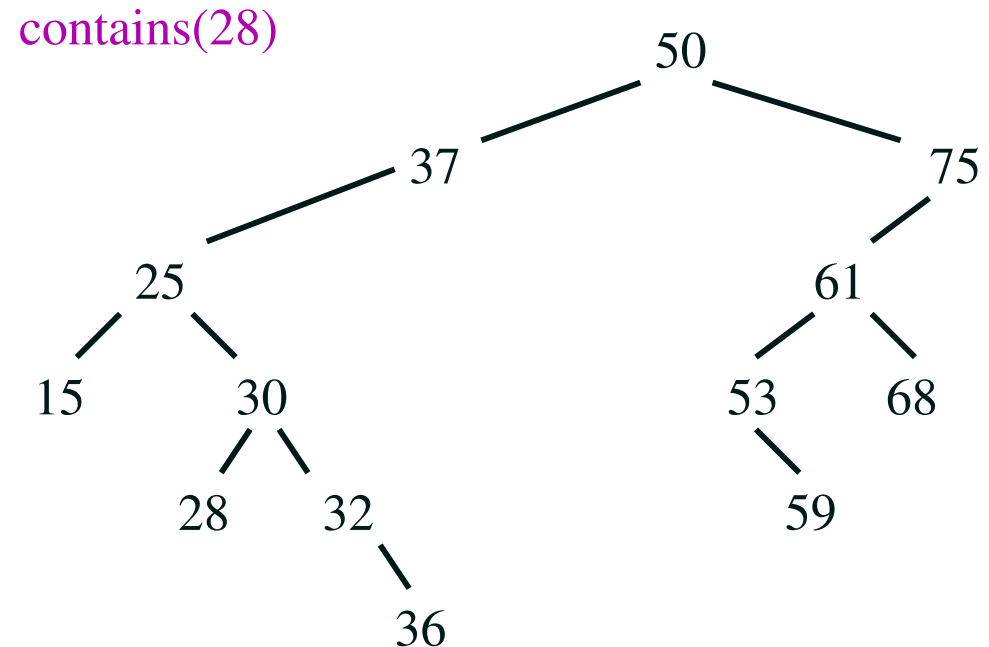
Searching A Binary Search Tree

- Searching a binary search tree is easy
- Start at the root
- Compare with element
 - ★ If less than element go left
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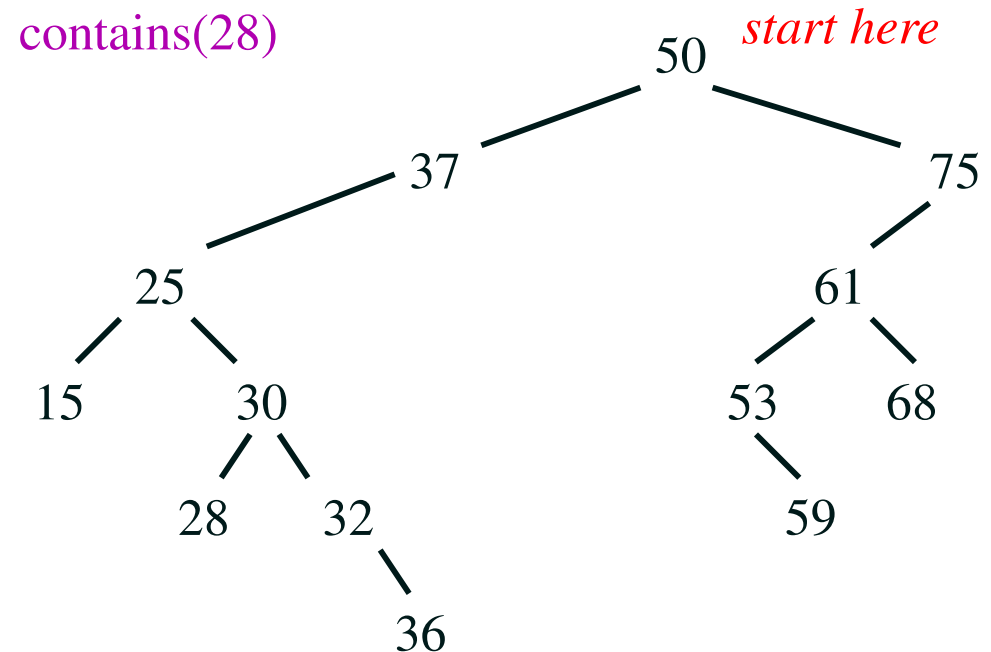
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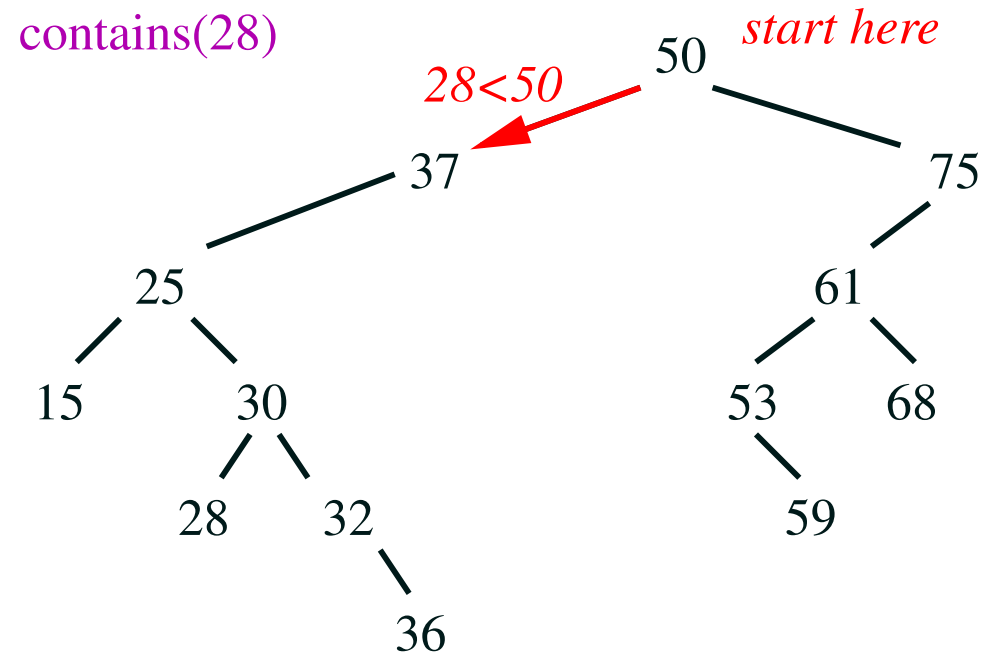
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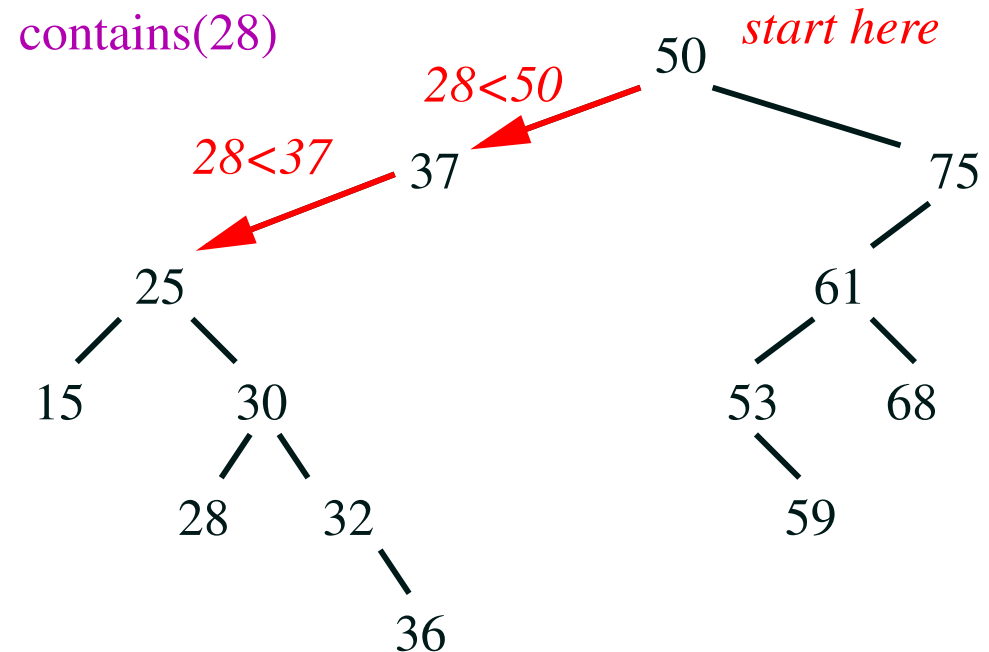
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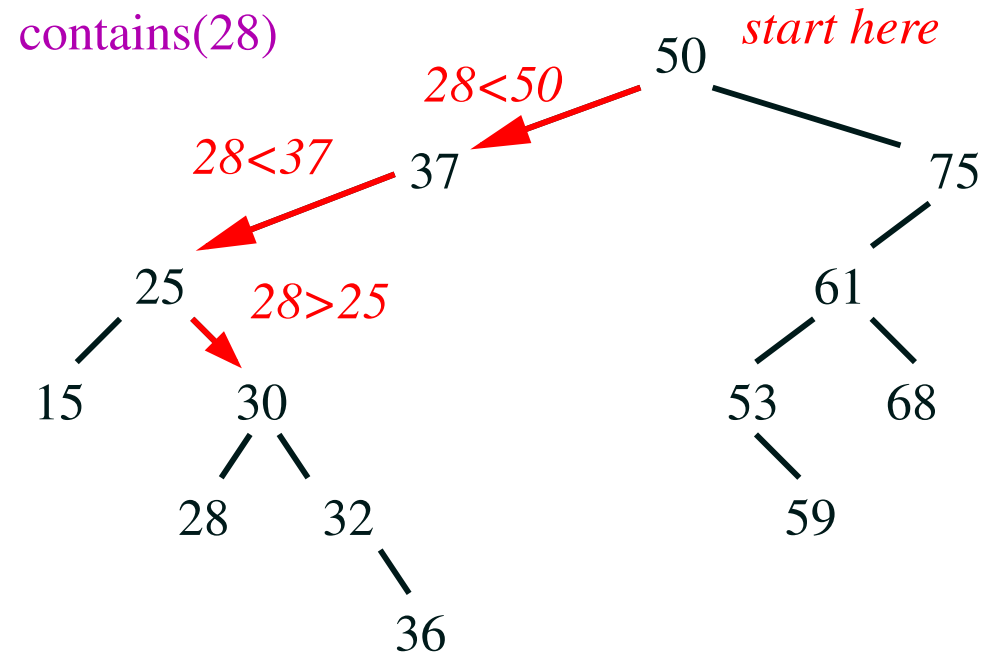
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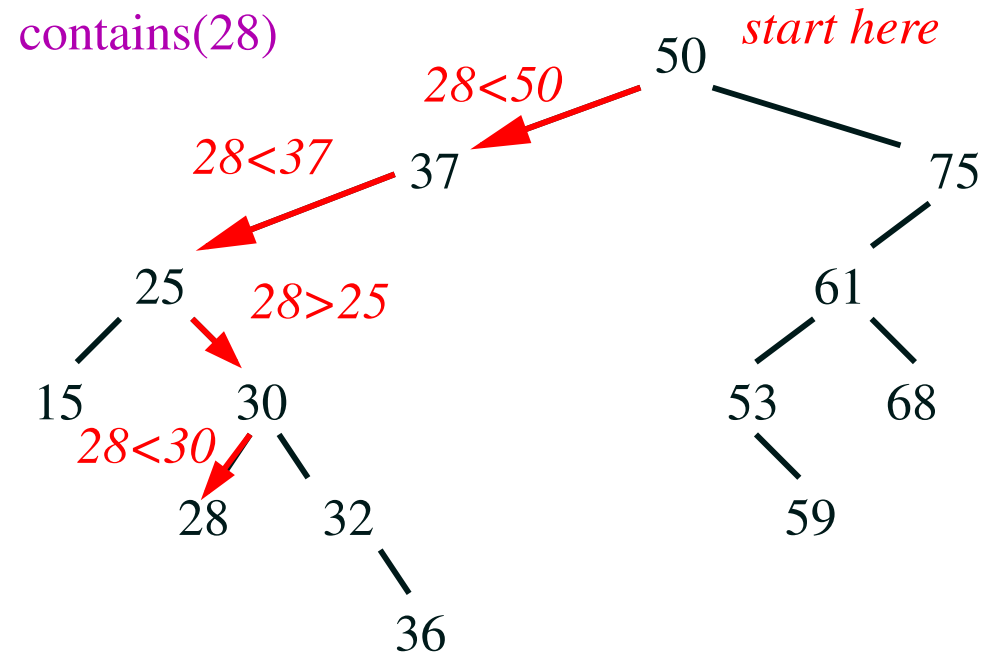
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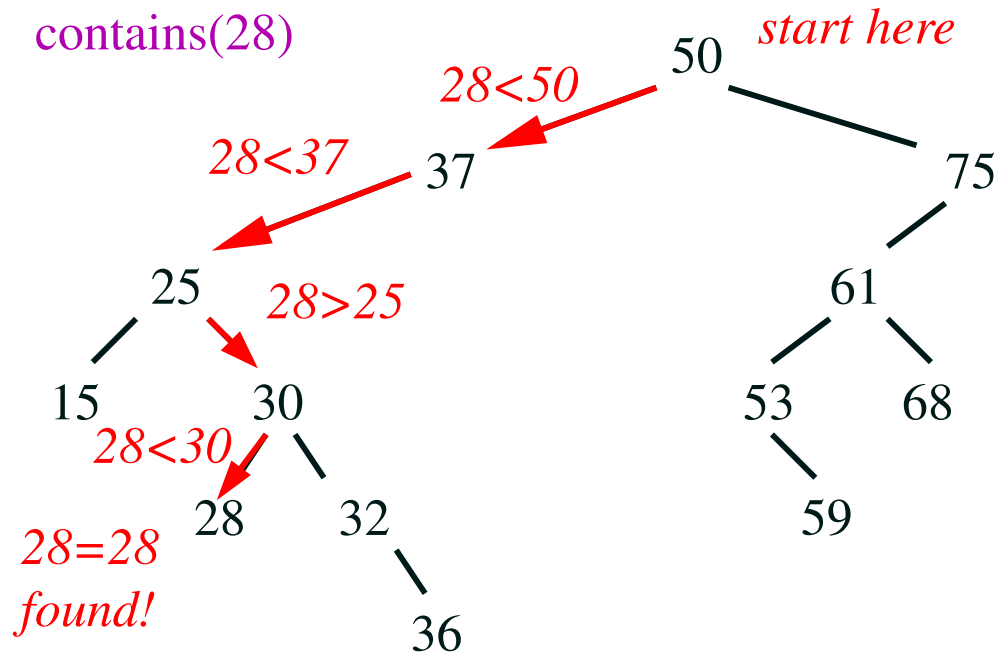
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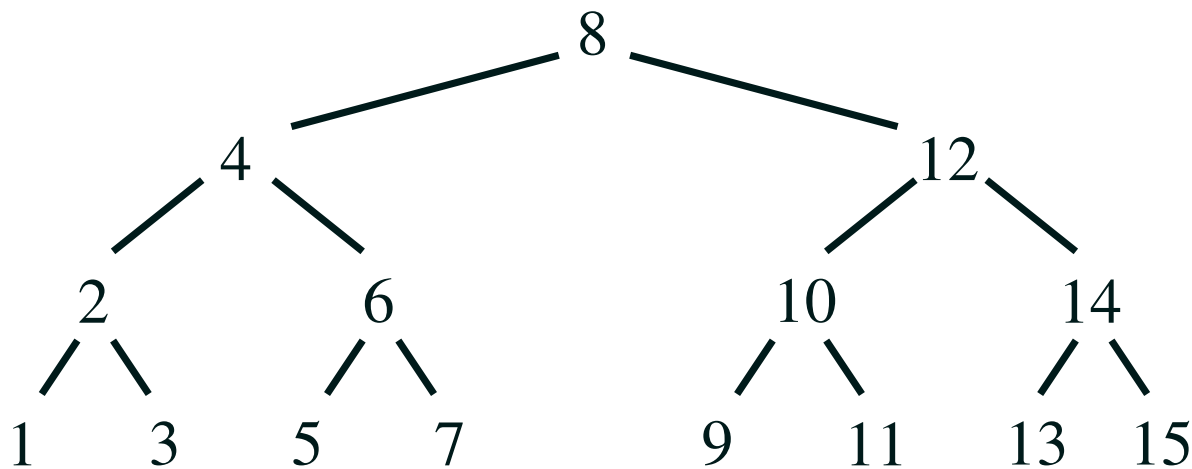
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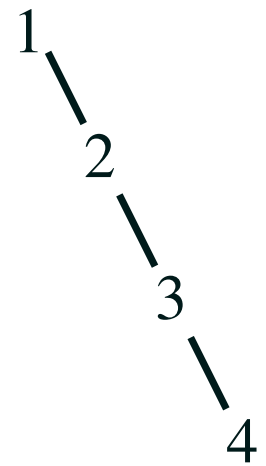


Speed of Search

- The number of comparisons necessary to find an element in a binary tree depends on the level of the node in the tree
- The worst case number of comparisons is therefore the height of the tree
- This depends on the density of the tree



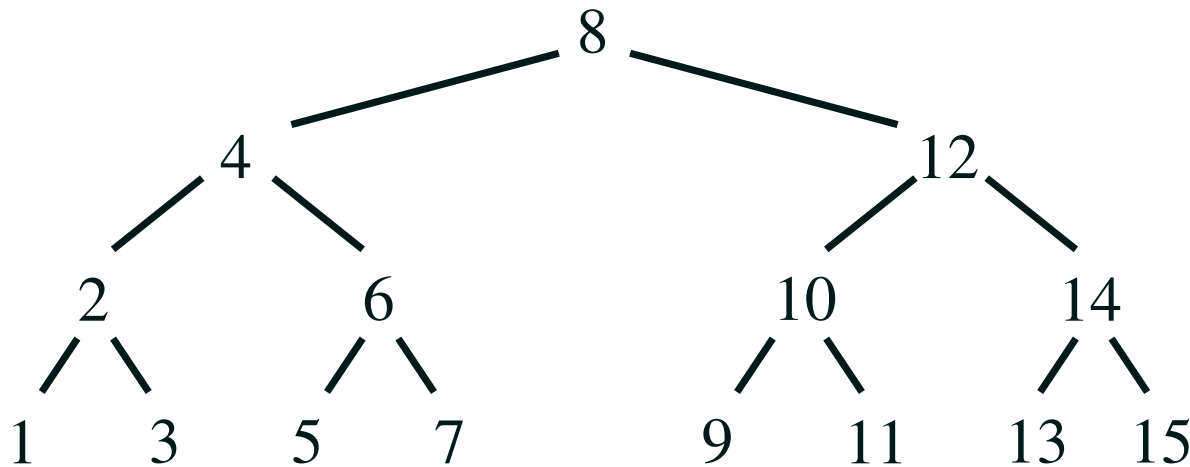
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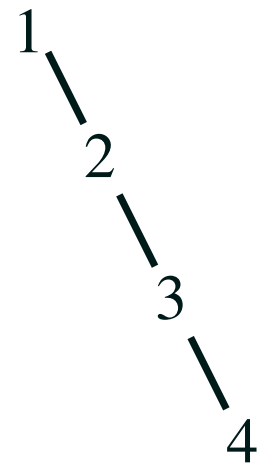
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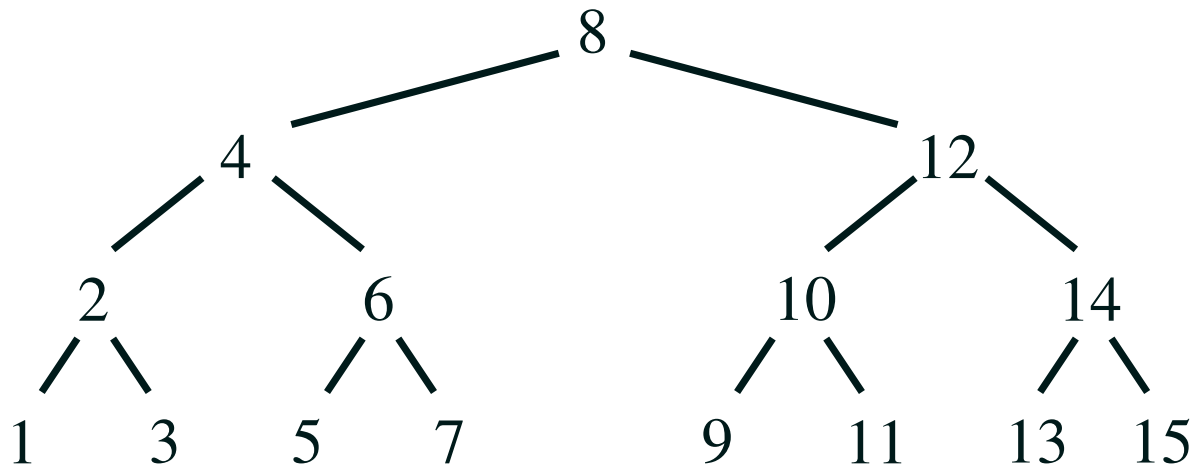
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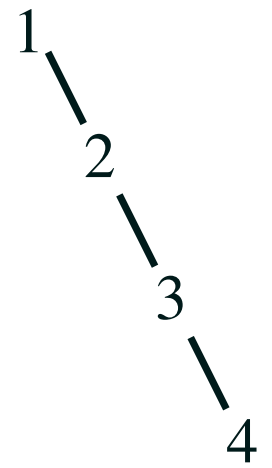
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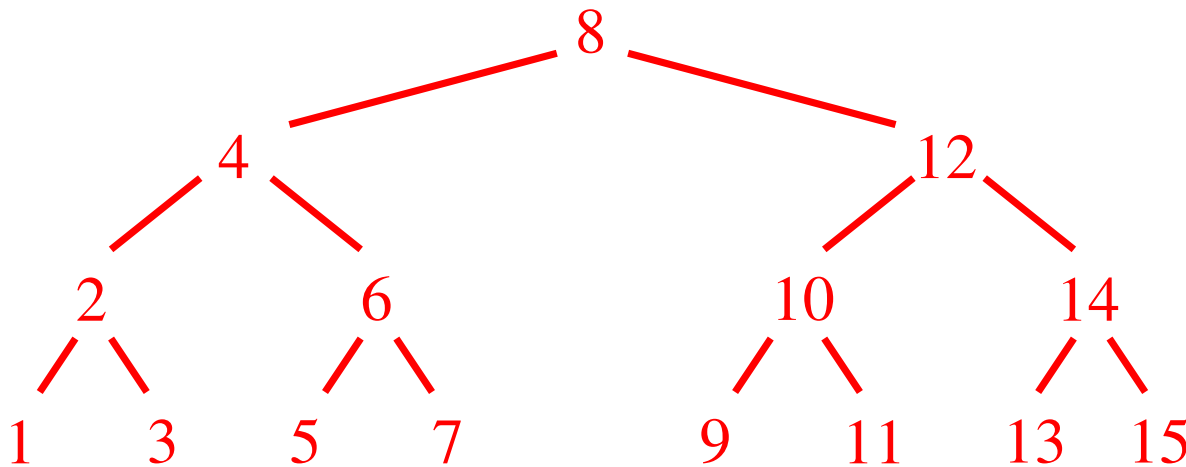
full tree



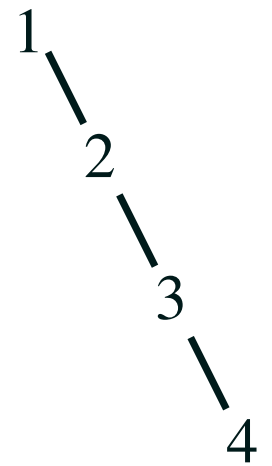
sparse tree

Speed of Search

- The number of comparisons necessary to find an element in a binary tree depends on the level of the node in the tree
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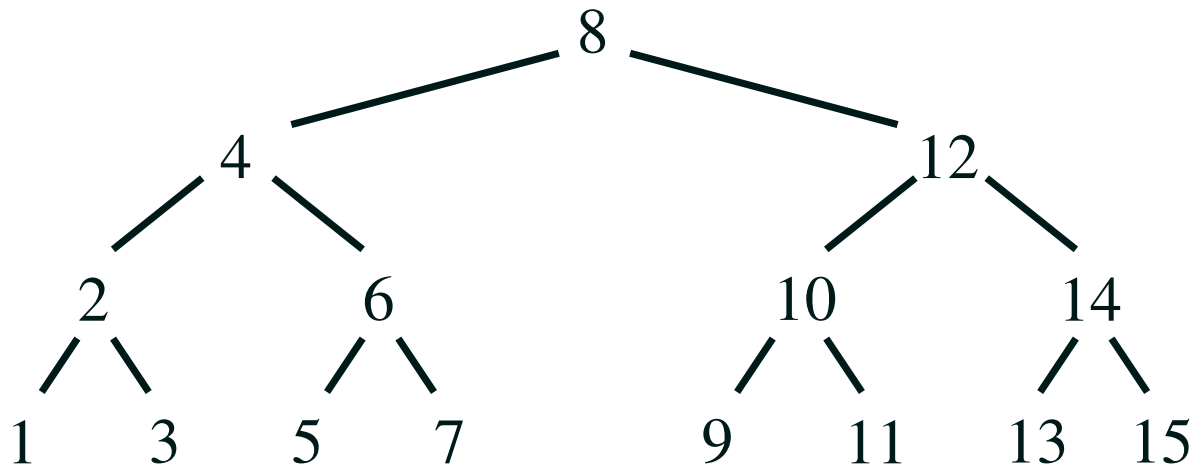
full tree



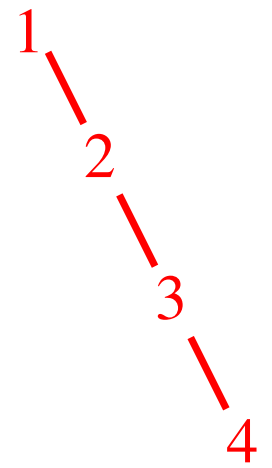
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full tree



sparse tree

Implementing a Set

- A set is a fundamental **abstract data type**
- It is a collection of things with no repetition and no order
- Ironically because order doesn't matter we can order the elements

$$\{1, 3, 5, 5, 3, 4\} = \{5, 3, 4, 1\} = \{1, 3, 4, 5\}$$

- This allows rapid search—a feature we care about
- Binary trees are one of the efficient ways of implementing a set

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Fitting In

- The standard template library provides a class `std::set<T>`
- This contains many functions like
 - ★ Constructors
 - ★ `size()`
 - ★ `insert(T o)`
 - ★ `find(Object o)`
 - ★ `erase(Object o)`
 - ★ `begin()` and `end()`

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Comparable

- To sort any objects they must be comparable
- In the STL the set implementation has a second template parameter: `std::set<T, Compare = less<T> >`
- by default this is defined to be `less<T>` (which is a function already defined for most common types) which you can define
- If you have a set of complex objects you will have to define `Compare`

```
bool MyCompare(MyObject left, MyObject right) {  
    return something  
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```

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mySet = set<MyObject, MyCompare>;
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Find an Element

- One of the core operations of a binary tree is to find a node

```
iterator find(const T& element) {  
    Node* current = root;  
    while (current!=0) {  
        if (current->element == element) {  
            return iterator(current);  
        }  
        if (element < current->element) {  
            current = current->left;  
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}
```

Add an Element

```
pair<iterator, bool> insert(const T& element) {
    if (no_elements==0) {
        root = new Node(element, 0);
        ++no_elements;
        return pair<iterator, bool>(iterator(root), true);
    }
    Node* parent = 0;
    Node* current = root;
    while(current != 0) {
        if (current->element == element) {
            return pair<iterator, bool>(iterator(0), false);
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        parent = current;
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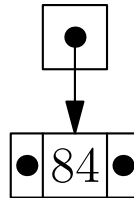
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current = new Node(element, parent);  
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++no_elements;  
return pair<iterator, bool>(iterator(current), true);  
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```

Tree in Action

add(84)

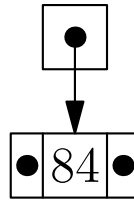


Tree in Action

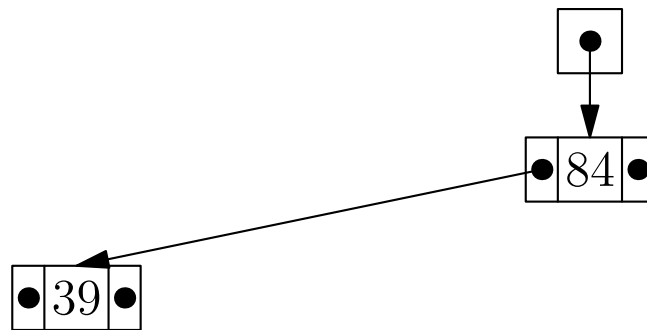


Tree in Action

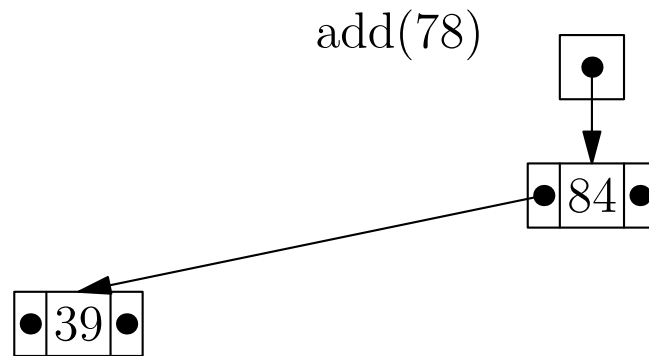
add(39)



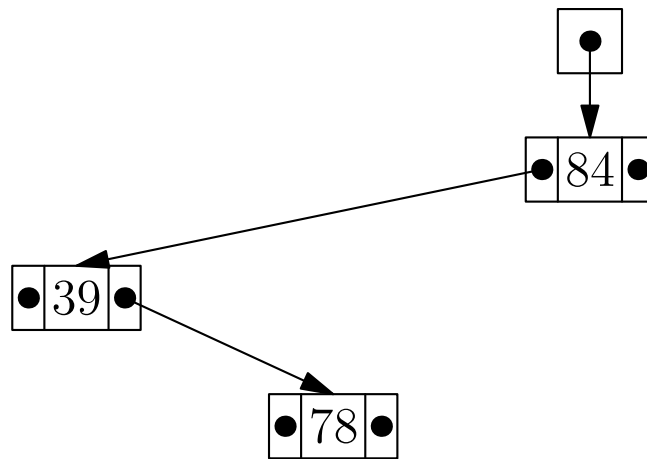
Tree in Action



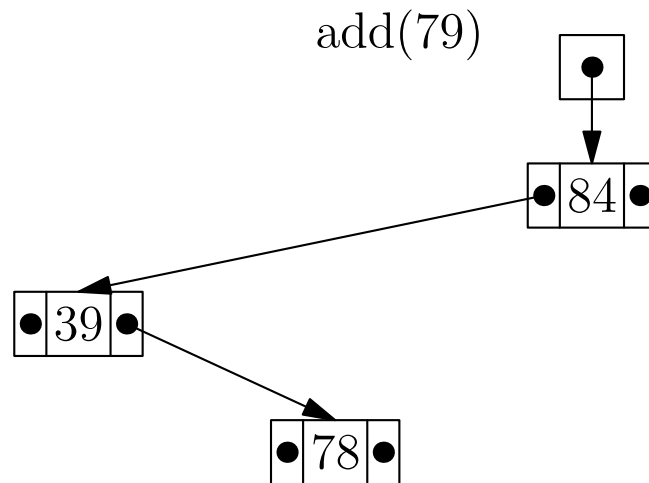
Tree in Action



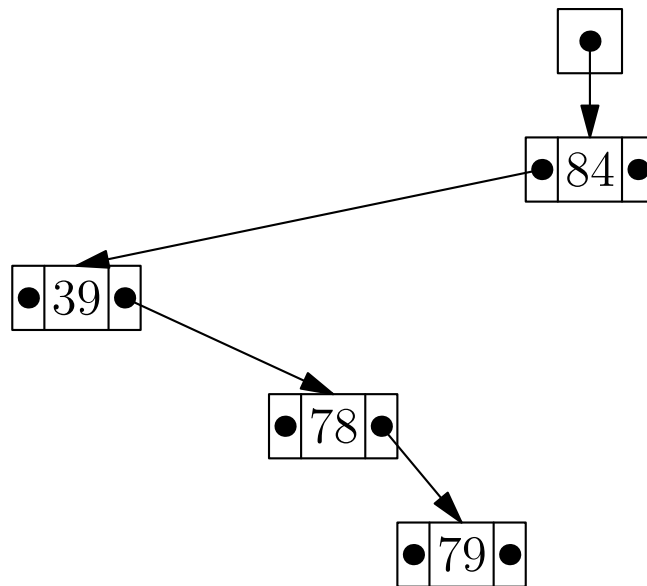
Tree in Action



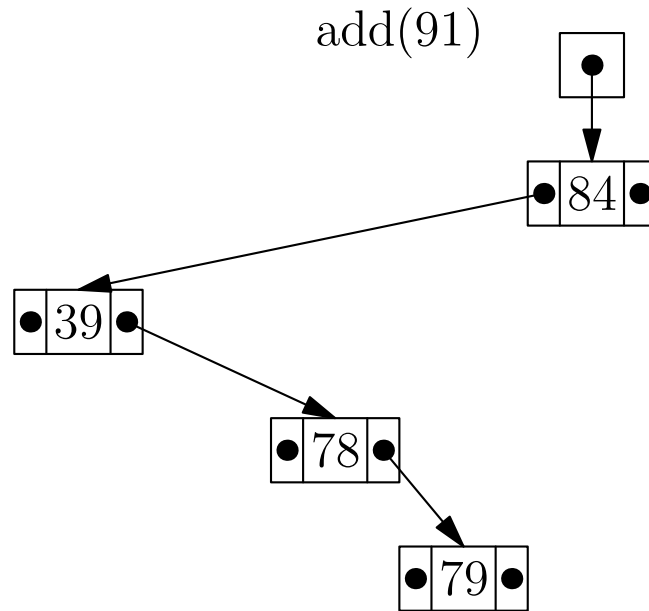
Tree in Action



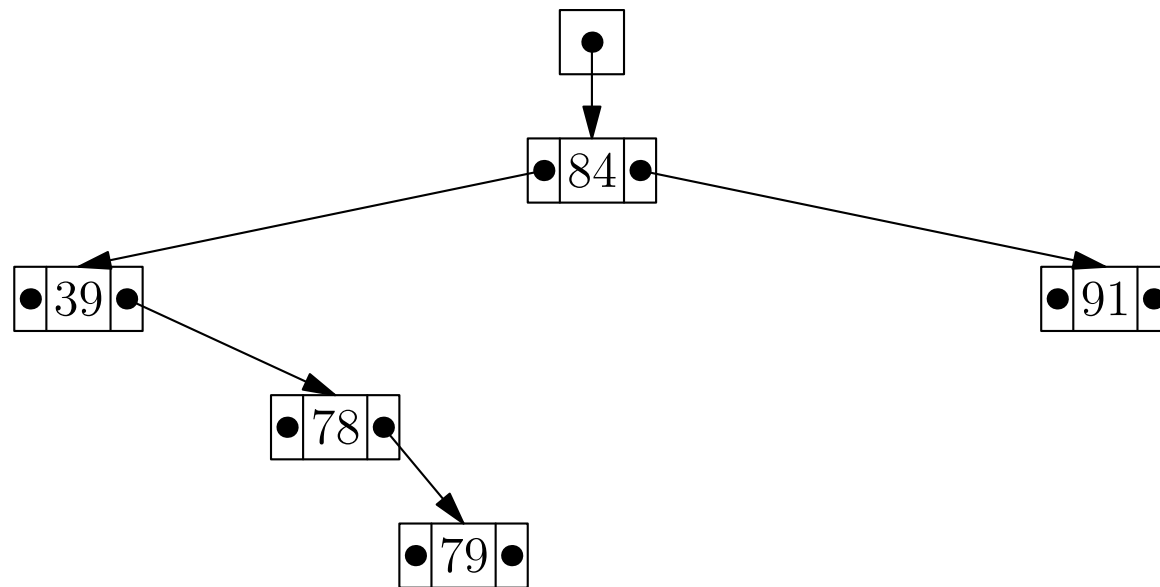
Tree in Action



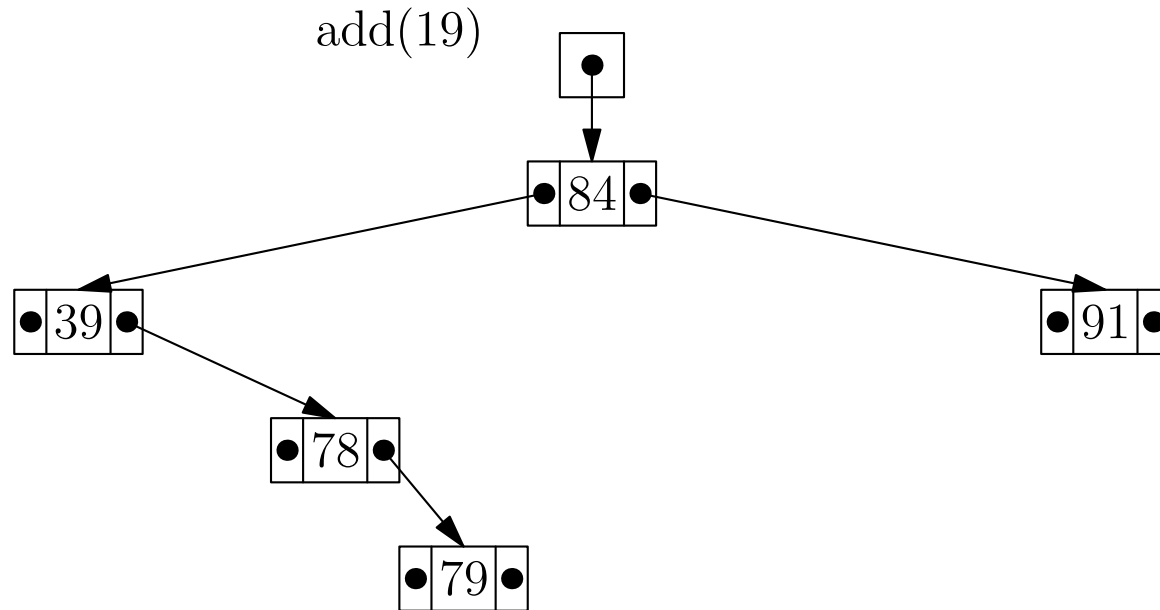
Tree in Action



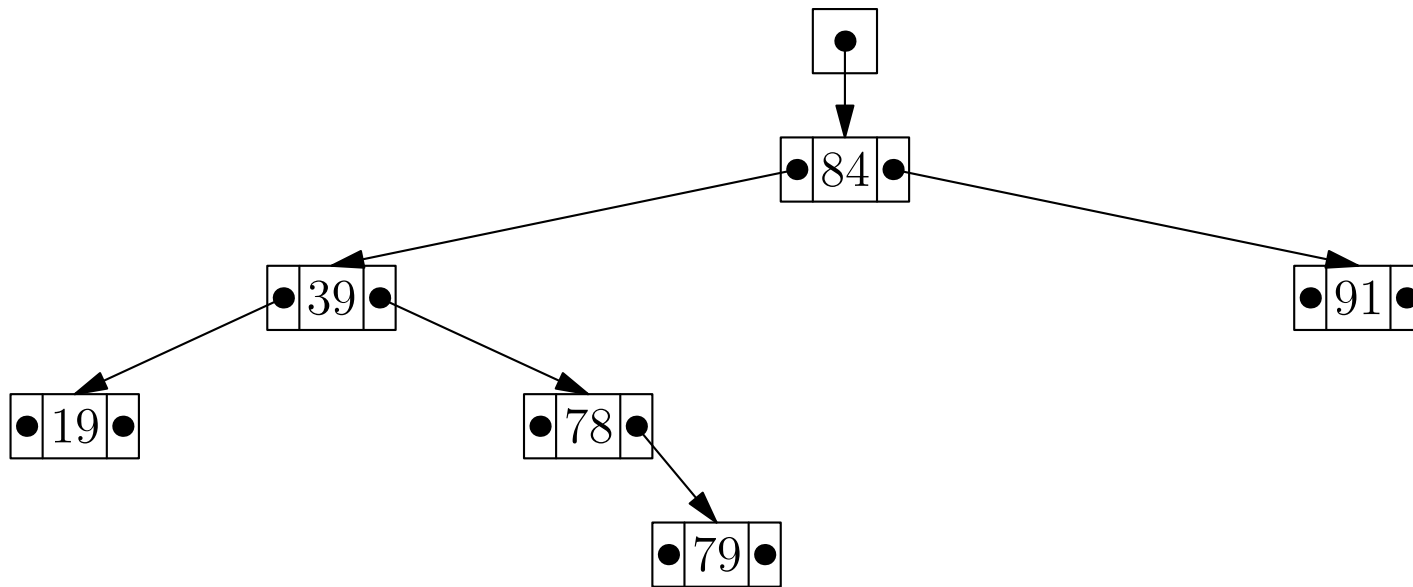
Tree in Action



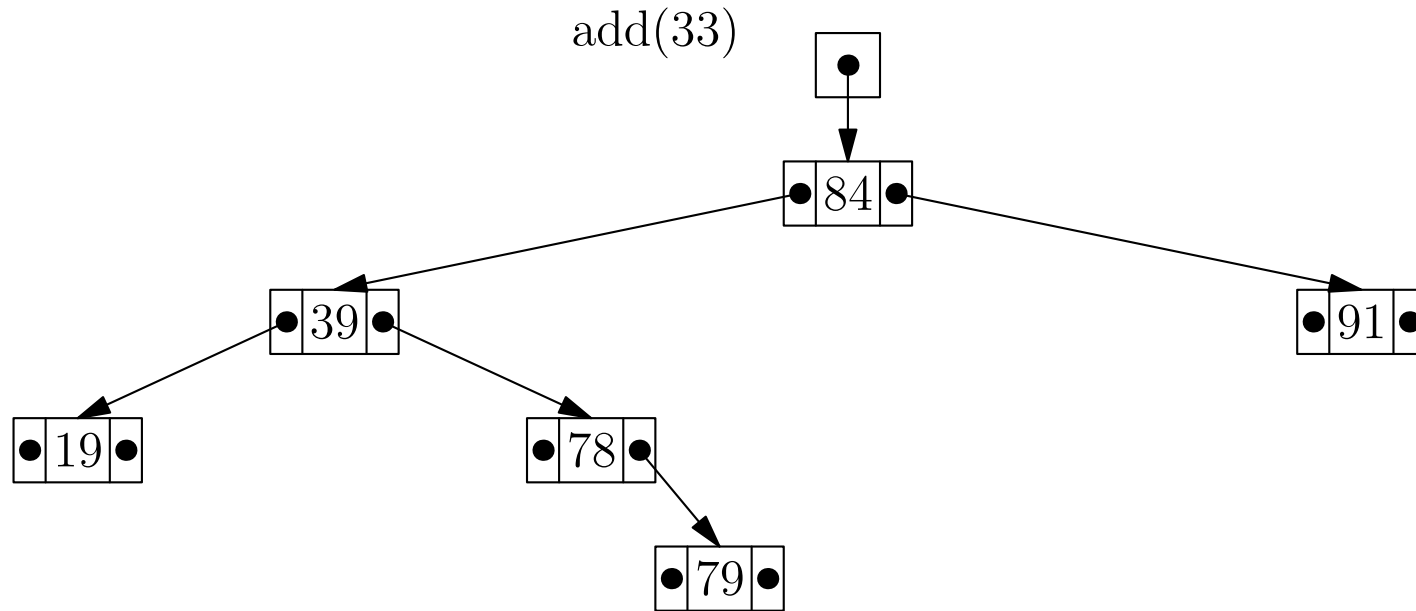
Tree in Action



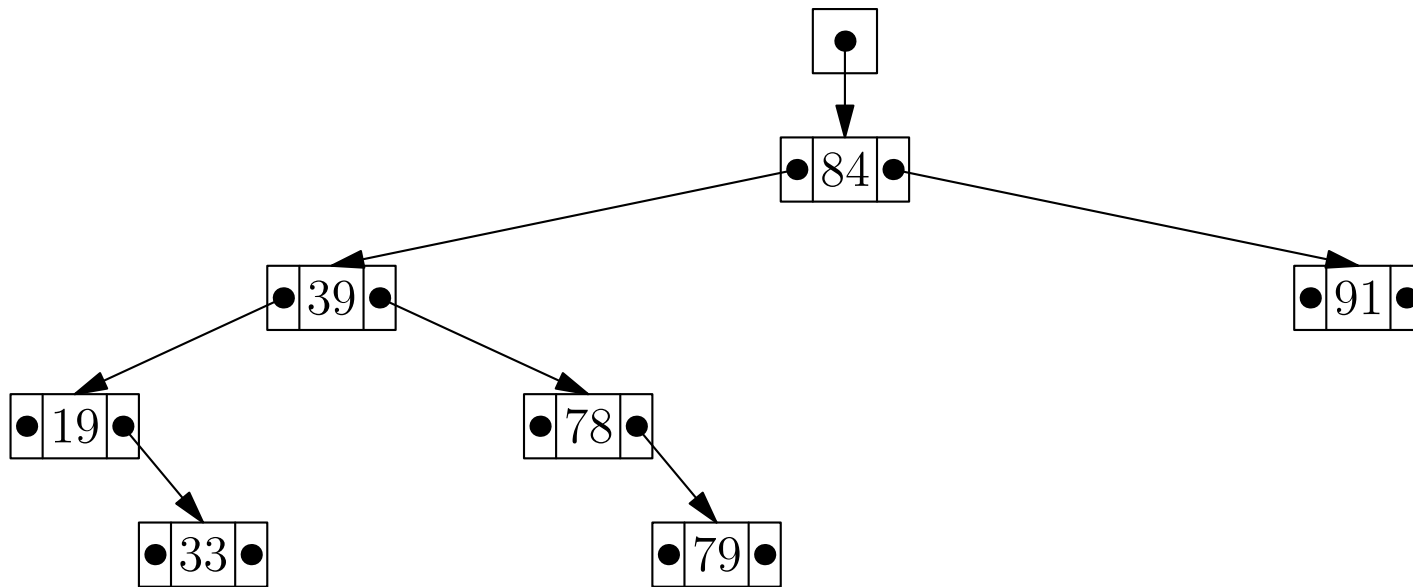
Tree in Action



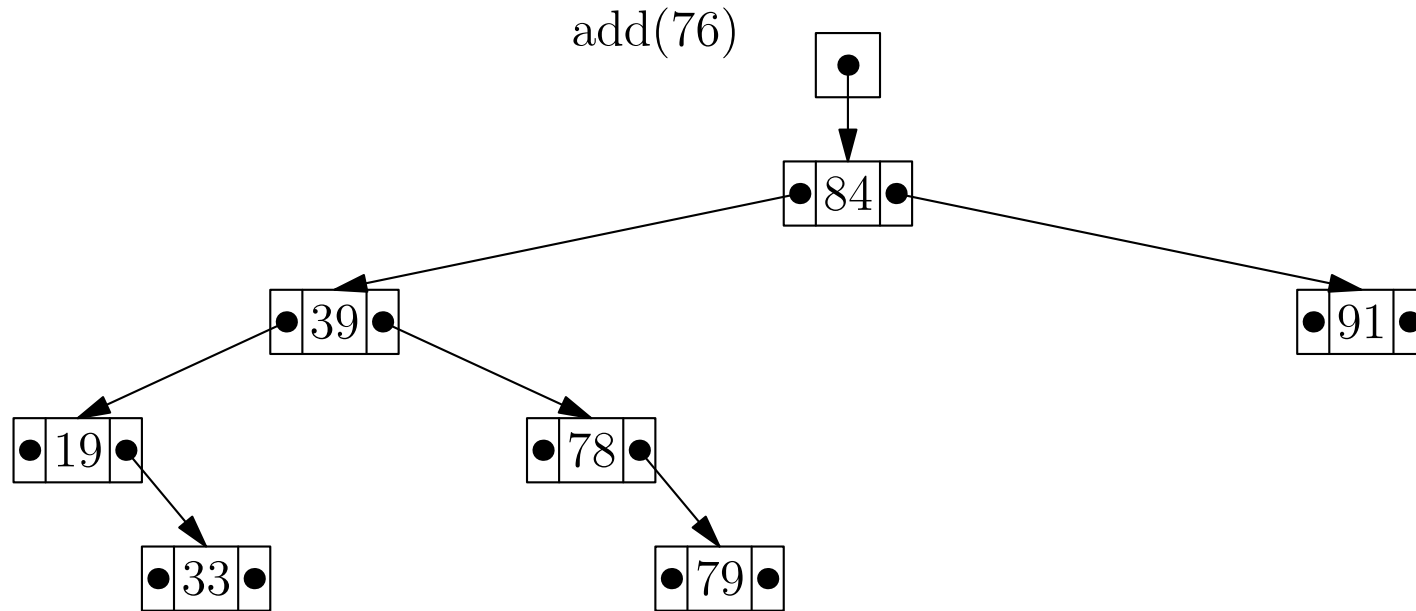
Tree in Action



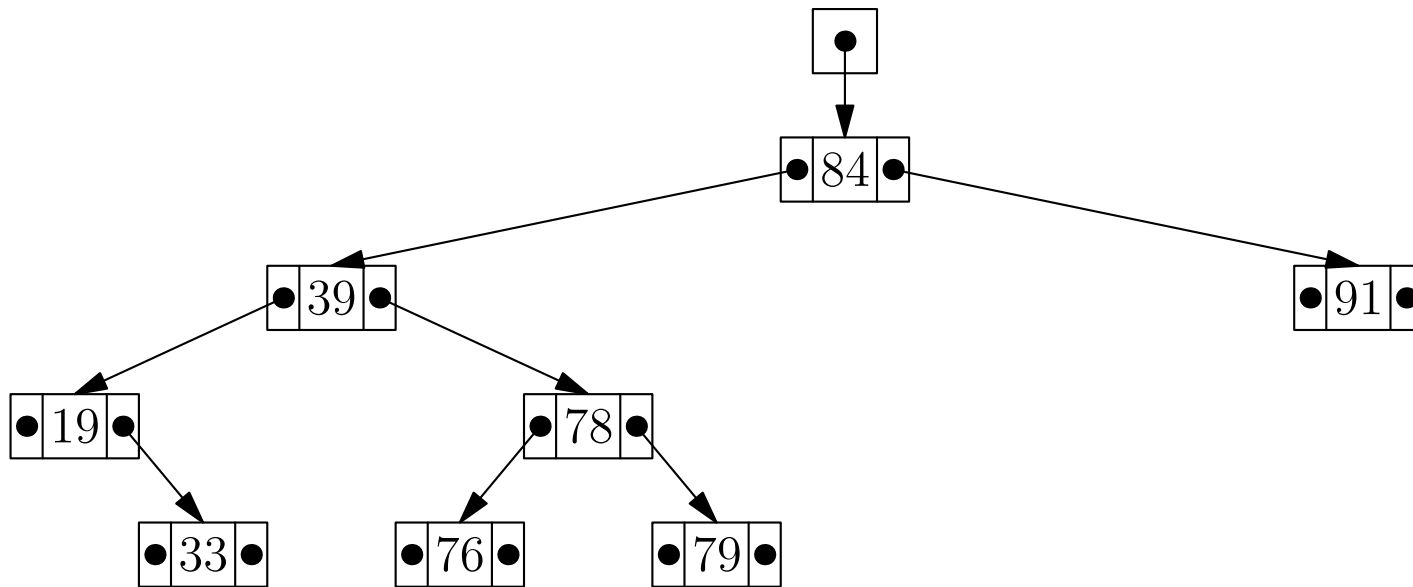
Tree in Action



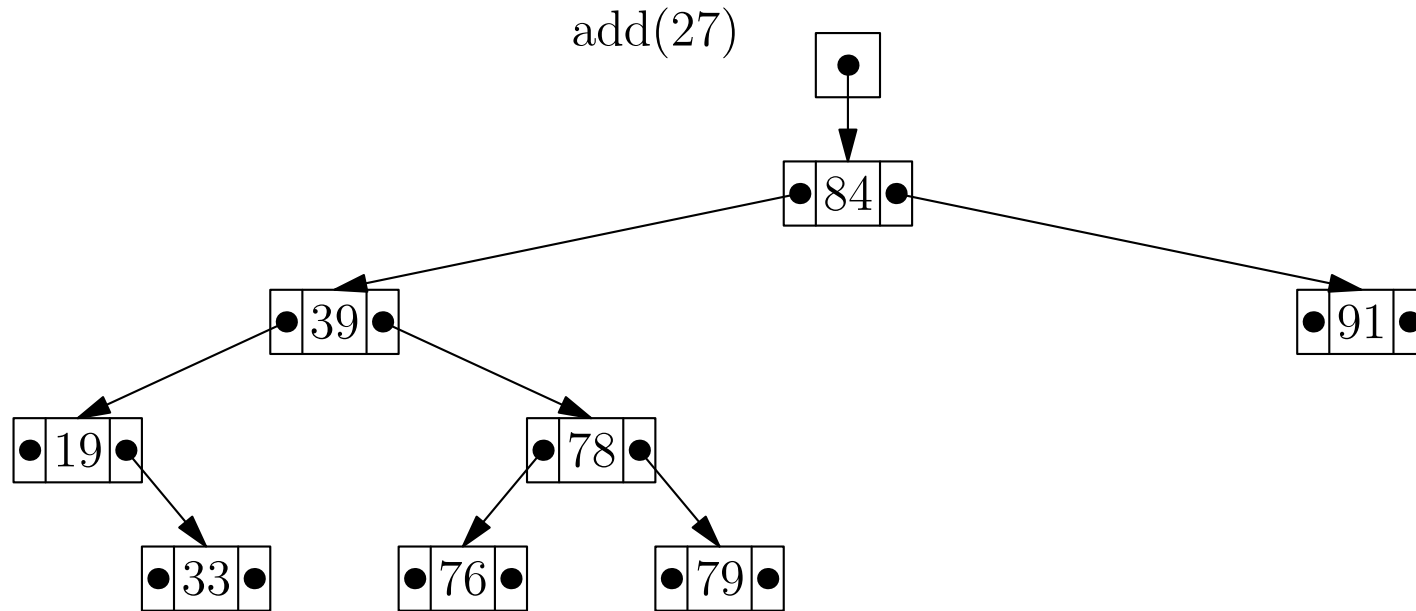
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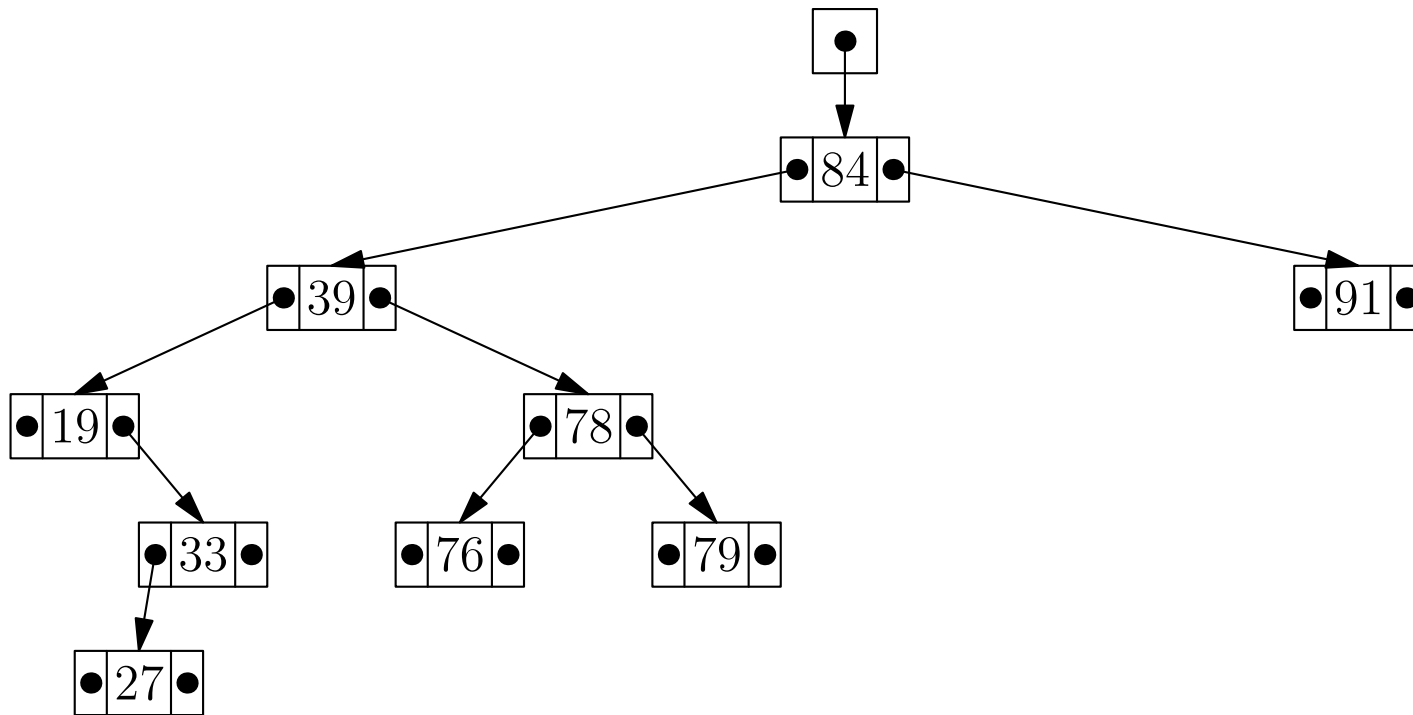
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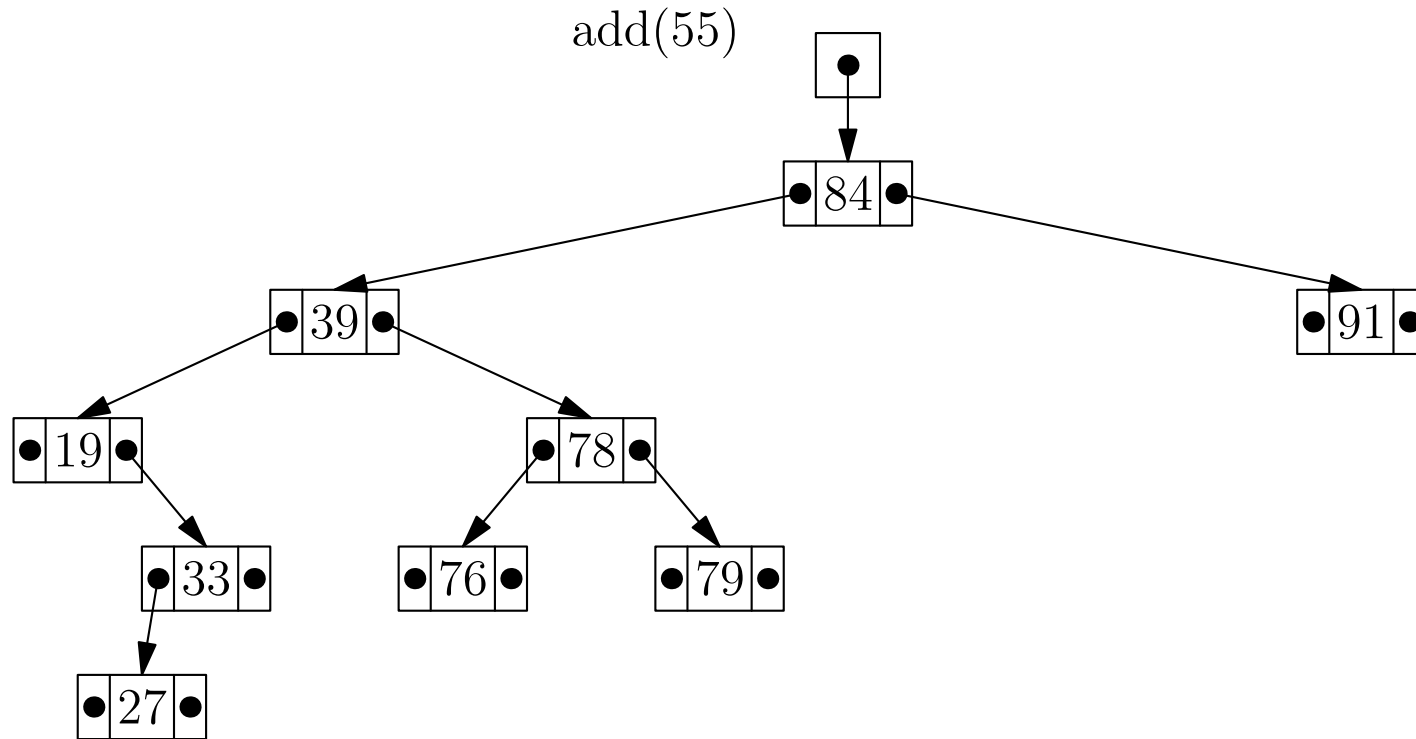
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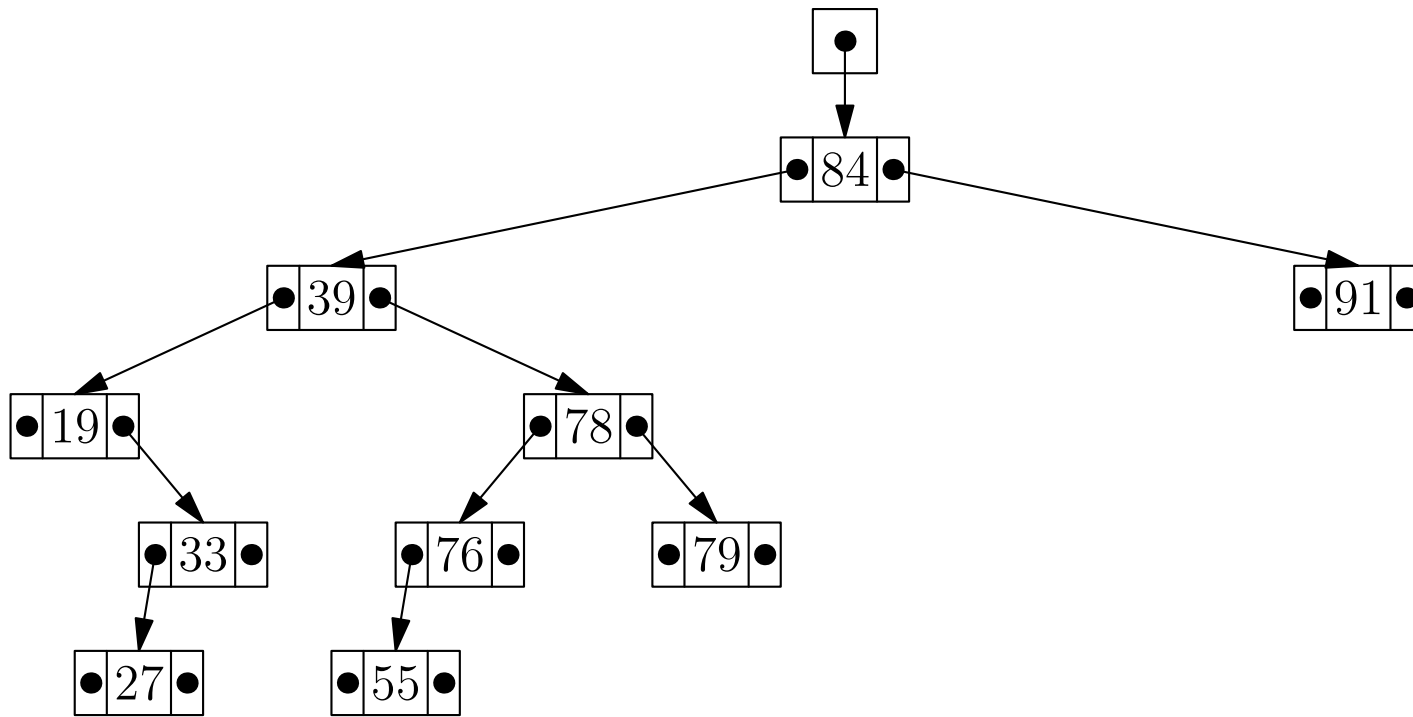
Tree in Action



Tree in Action



Tree in Action



Shape of Tree

- The structure of the tree depends on the order in which we add elements to it
- Suppose we add

*To be, or not to be: that is the question:
Whether 'tis nobler in the mind to suffer
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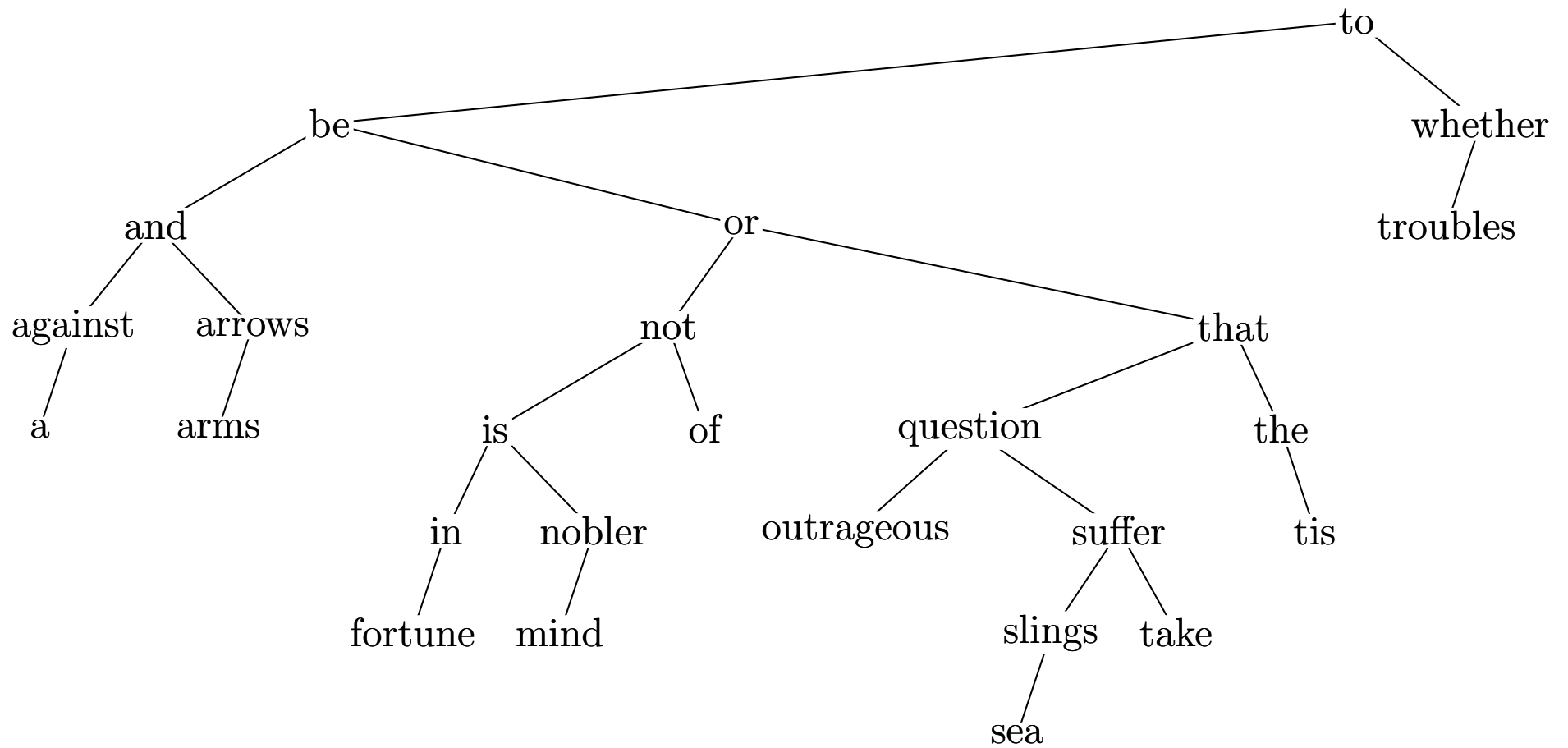
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Hamlet



Outline

1. Trees
2. Binary Trees
 - Implementing Binary Trees
3. Binary Search Trees
 - Definition
 - Implementing a Set
4. **Tree Iterators**



Tree Iterators

- As with most container classes it is very useful to define iterators
- `begin()` should return a “pointer” to the start of the tree
- `end()` provides a “pointer” past the end
- `operator*()` returns the element
- `operator++()` increments the “pointer”
- `operator!=(lhs, rhs)` is used to compare iterators

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set<int> mySet;  
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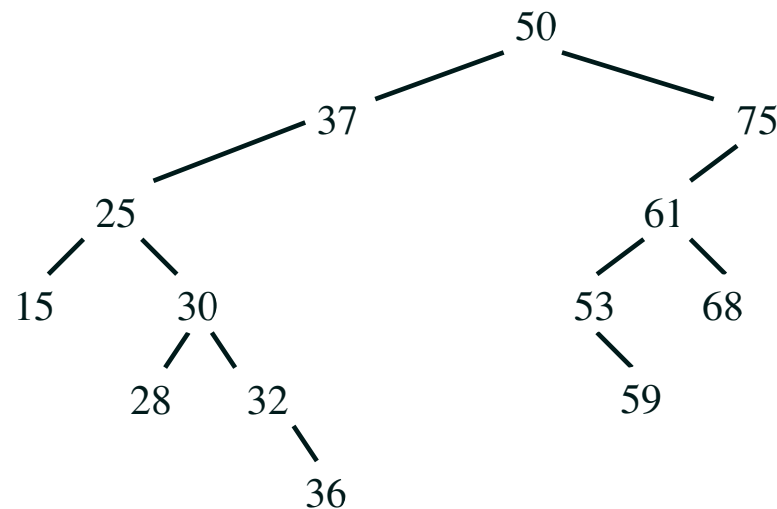
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Successor

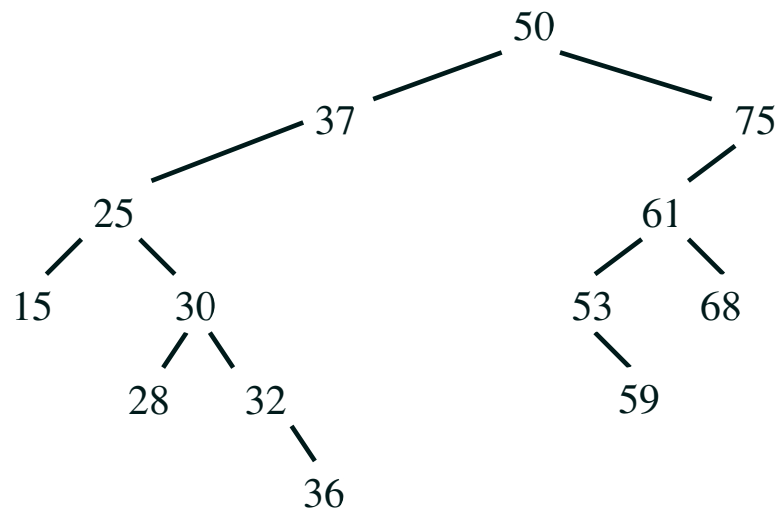
- To find the successor we first start in the left most branch
- We follow two rules
 1. **If** right child exist **then** move right once and then move as far left as possible
 2. **else** go *up* to the left as far as possible and then move up right



{15 25 28 30 32 36 37 50 53 59 61 68 75}

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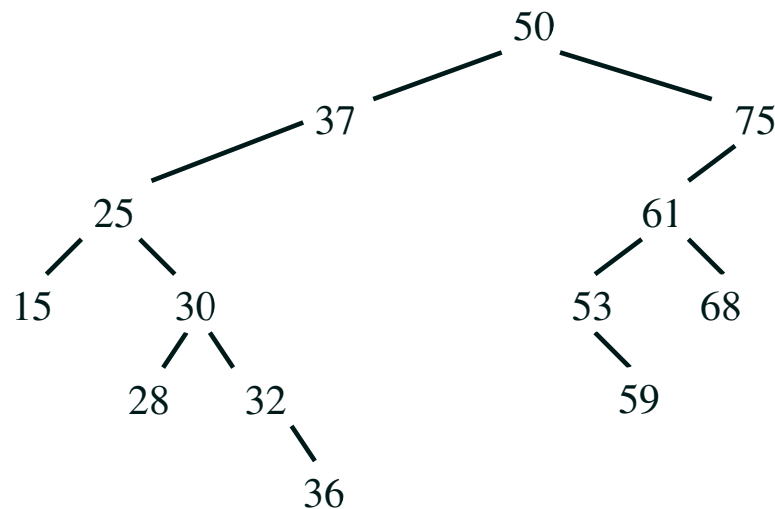
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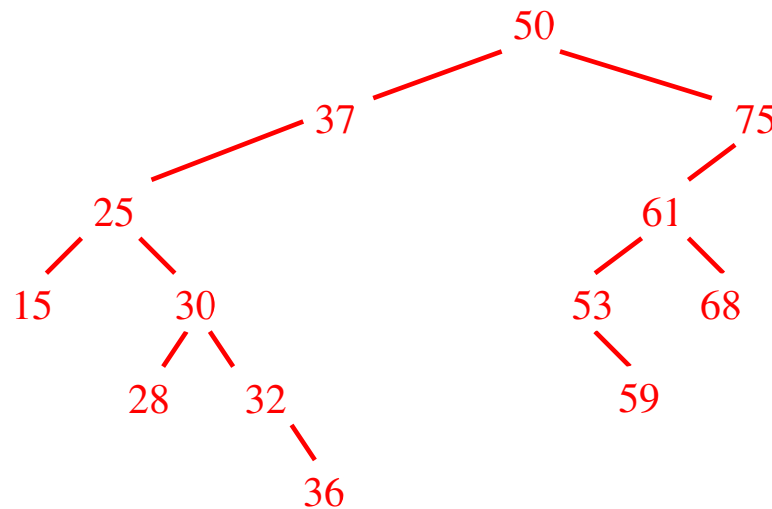
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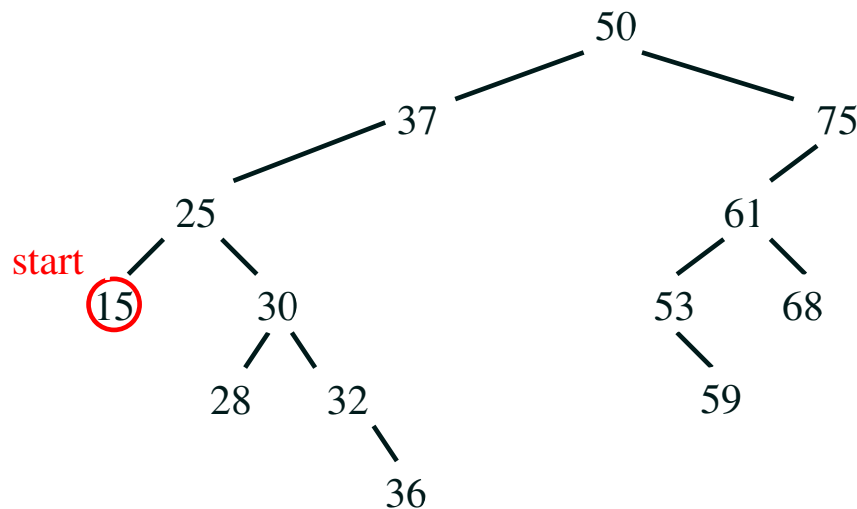
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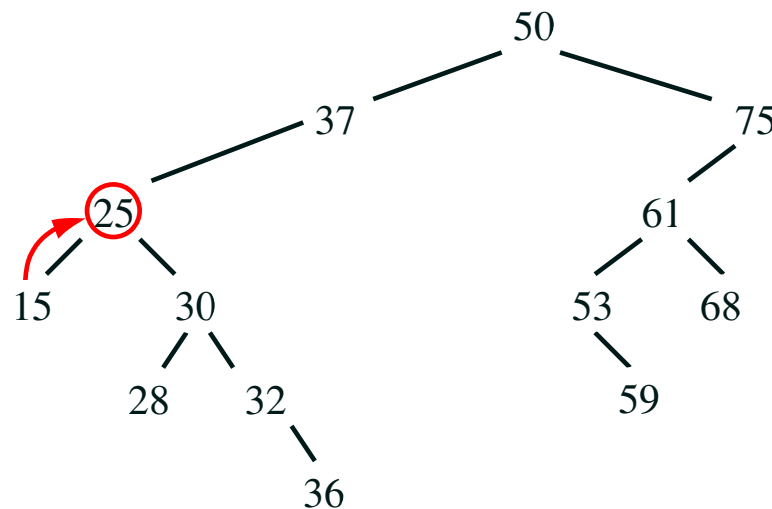
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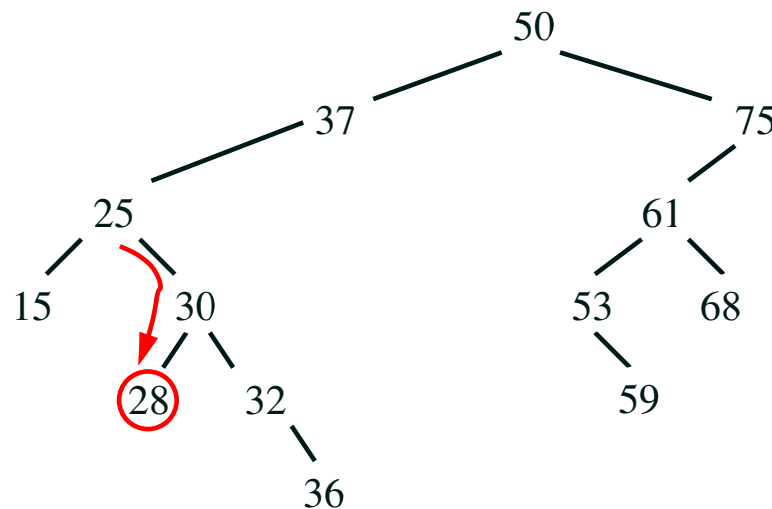
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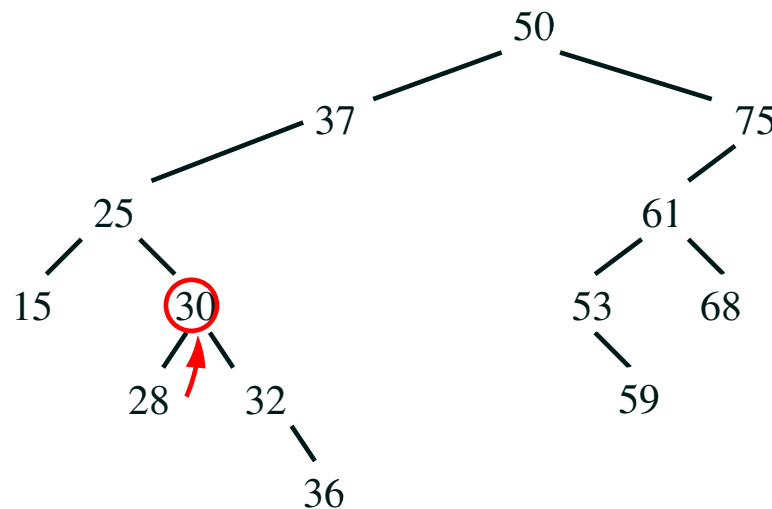
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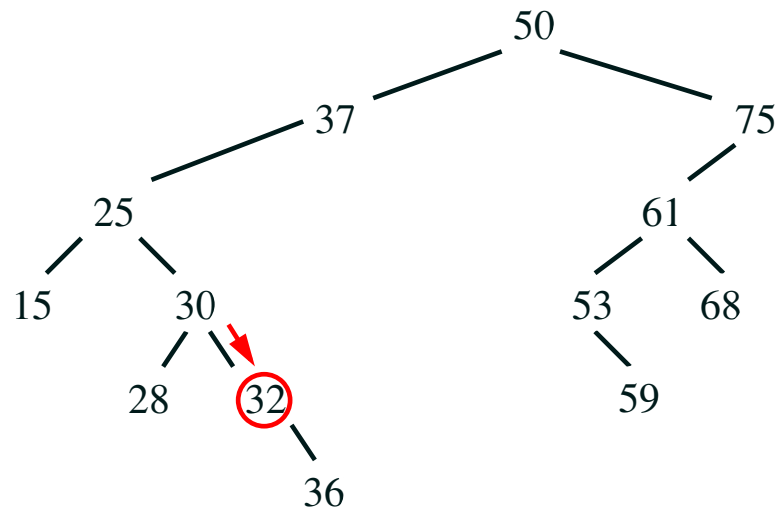
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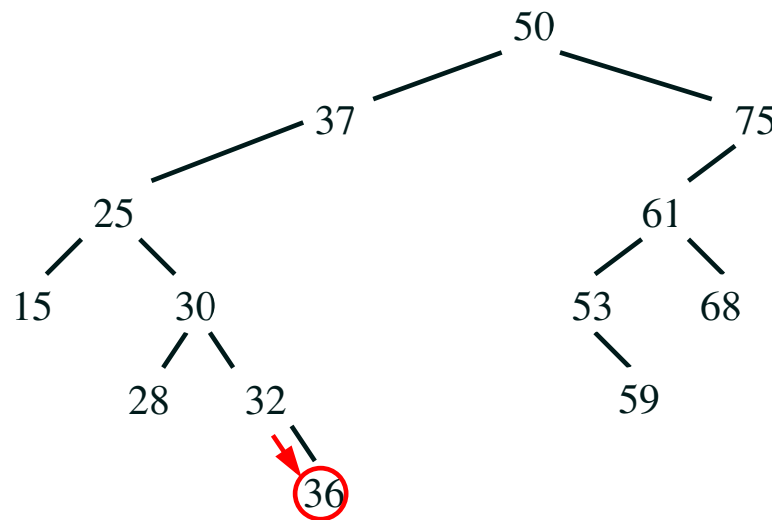
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Successor

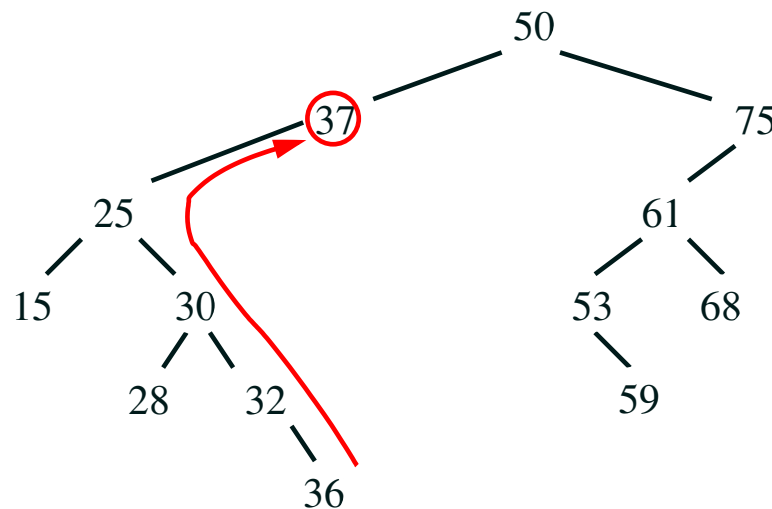
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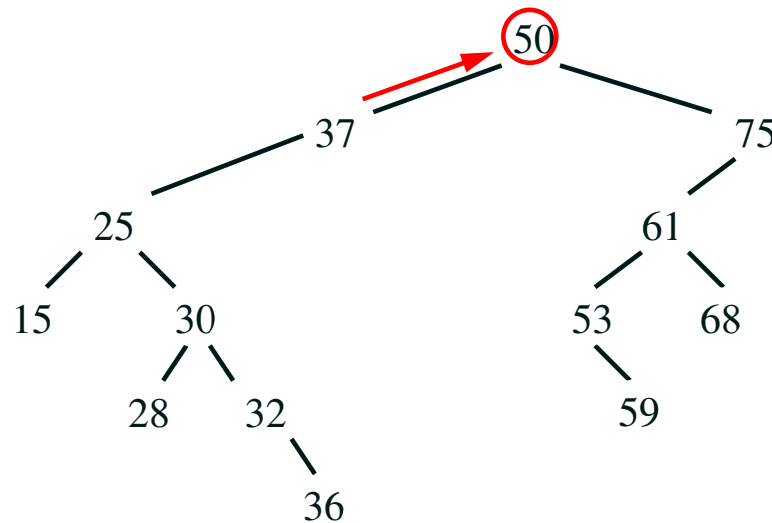
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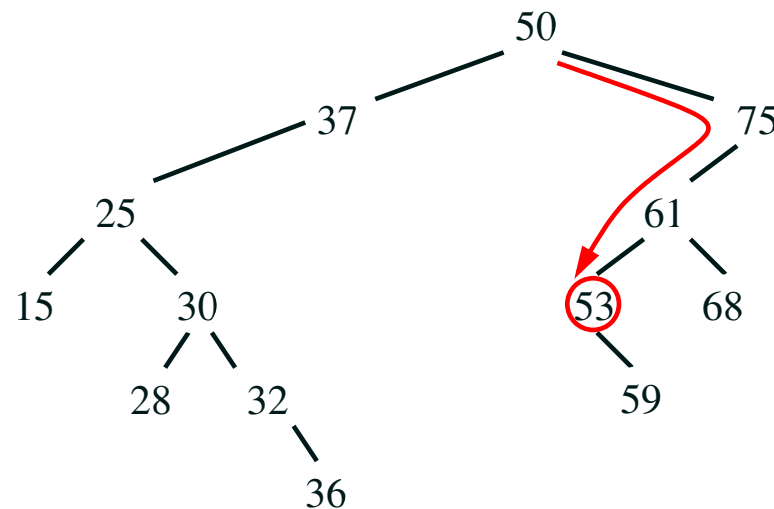
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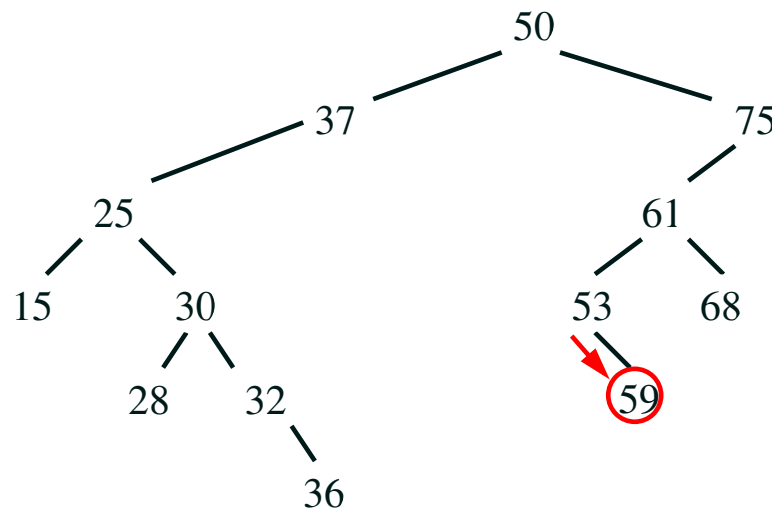
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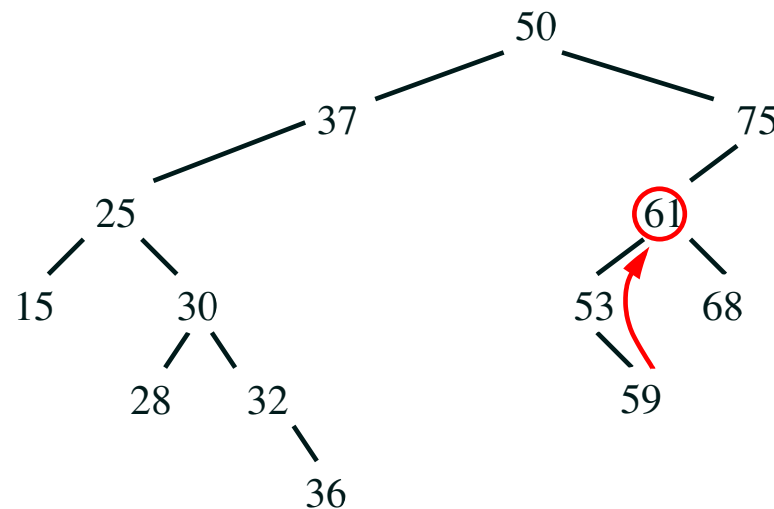
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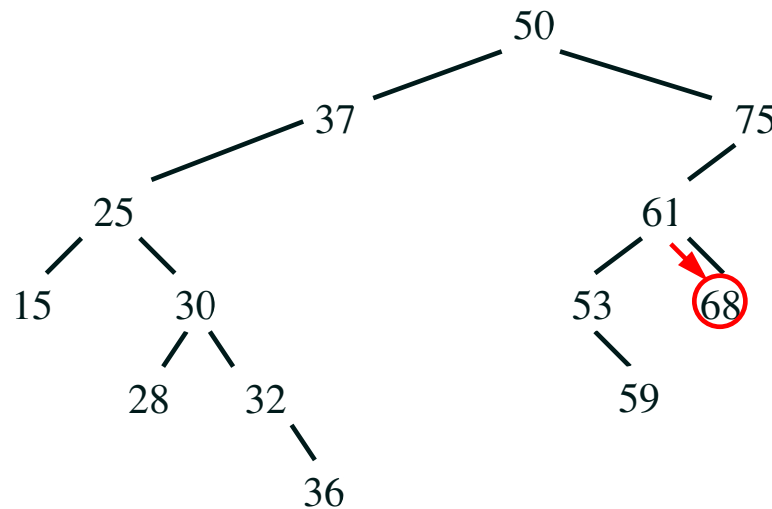
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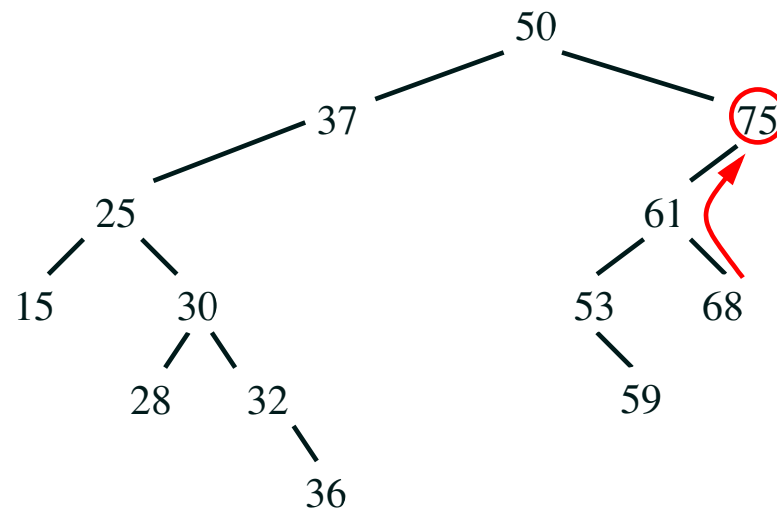
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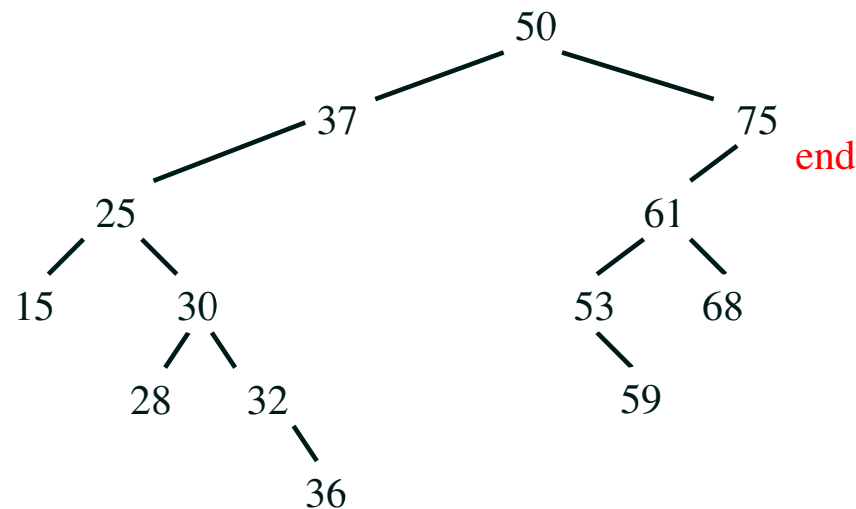
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Lessons

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- Conceptually they are quite simple
- However, there are a lot of details that need to be understood
- Coding even simple trees needs great care
- As we will see things get more complicated

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