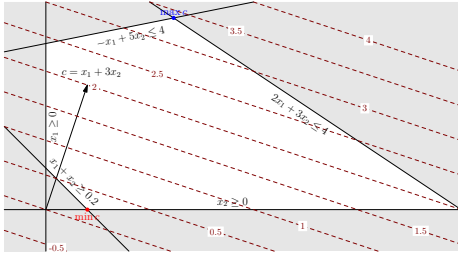


Lesson 26: Use Linear Programmings



linear programming, applications

AICE1005

Algorithms and Analysis

1

Going Shopping

- Suppose we have a number of food stuffs which we label with indices $f \in \mathcal{F}$
- The price of food stuff f per kilogram we denote p_f
- We are interested in buying a selection of foods $x = (x_f | f \in \mathcal{F})$ where x_f is the quantity (in kg) of food f
- We want to minimise the total price $\sum_f p_f x_f = p \cdot x$
- However we want to ensure that the food has enough vitamins

AICE1005

Algorithms and Analysis

3

Optimisation Problem

- We can write the food shopping problem as

$$\min_x p \cdot x \quad \text{subject to} \quad Ax \geq b \quad \text{and} \quad x \geq 0$$

- Note that the inequalities involving vectors means that each component must be satisfied, i.e.

$$\begin{aligned} Ax \geq b &\Rightarrow \forall v \in \mathcal{V} \quad \sum_{f \in \mathcal{F}} A_{vf} x_f \geq b_v \\ x \geq 0 &\Rightarrow \forall f \in \mathcal{F} \quad x_f \geq 0 \end{aligned}$$

- This is an example of a “linear program”

AICE1005

Algorithms and Analysis

5

Constraints

- Each factory can only produce a certain overall tonnage of commodities

$$\sum_{c \in \mathcal{C}} x_{cf} \leq b_f \quad \forall f \in \mathcal{F}$$

where b_f is the maximum production capacity of factory f

- The total demand for each commodity is d_c so

$$\sum_{f \in \mathcal{F}} x_{cf} = d_c \quad \forall c \in \mathcal{C}$$

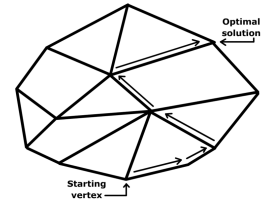
- We can only produce positive amounts, i.e. $x_{cf} \geq 0$

AICE1005

Algorithms and Analysis

7

1. Examples
2. Linear Programs
3. Properties of Solution
4. Normal Form



AICE1005

Algorithms and Analysis

2

Nutrition

- We consider the set of vitamins \mathcal{V}
- Let A_{vf} be the quantity of vitamin v in food stuff f
- Let b_v be the minimum daily requirement of vitamin v
- We therefore require

$$\forall v \in \mathcal{V} \quad \sum_{f \in \mathcal{F}} A_{vf} x_f \geq b_v$$

AICE1005

Algorithms and Analysis

4

Transportation

- We consider a set of factories \mathcal{F} producing a set of commodities \mathcal{C}
- The amount of commodity c produced by factory f we denote by x_{cf}
- The shipping cost of commodity c from factory f to the retailer of c we denote by p_{cf}
- We want to choose x_{cf} to minimise the transportation costs

$$\sum_{c \in \mathcal{C}, f \in \mathcal{F}} p_{cf} x_{cf}$$

- However, we have constraints. . .

AICE1005

Algorithms and Analysis

6

Linear Program

- We can write the full problem as

$$\min_x \sum_{c \in \mathcal{C}, f \in \mathcal{F}} p_{cf} x_{cf}$$

subject to

$$\sum_{c \in \mathcal{C}} x_{cf} \leq b_f \quad \forall f \in \mathcal{F}$$

$$\sum_{f \in \mathcal{F}} x_{cf} = d_c \quad \forall c \in \mathcal{C}$$

$$x_{cf} \geq 0 \quad \forall c \in \mathcal{C}, \quad \forall f \in \mathcal{F}$$

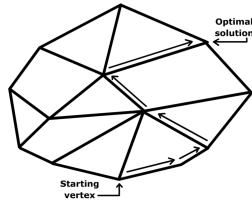
AICE1005

Algorithms and Analysis

8

Outline

1. Examples
2. **Linear Programs**
3. Properties of Solution
4. Normal Form



AICE1005

Algorithms and Analysis

9

Maximising

- We can also maximise rather than minimise
- Whether we want to maximise or minimise will depend on the application
- Note that

$$\max_x c \cdot x \equiv \min_x (-c) \cdot x$$

- We can thus always reformulate a maximisation problem as a minimisation problem and vice versa

AICE1005

Algorithms and Analysis

11

Key Features

- There are three key features of linear programs
 1. The cost (objective function) is linear in x_i ($c \cdot x$)
 2. The constraints are linear in x_i (e.g. $A_1 x \leq b_1$)
 3. The component of x are non-negative (i.e. $x_i \geq 0$)
- These are very special features, very often they don't apply, but a surprising large number of problems can be formulated as linear programming problems

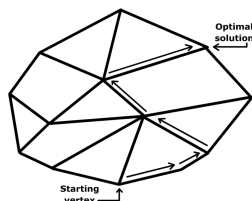
AICE1005

Algorithms and Analysis

13

Outline

1. Examples
2. Linear Programs
3. **Properties of Solution**
4. Normal Form



AICE1005

Algorithms and Analysis

15

General Linear Programs

- Linear programs are problems that can be formulated as follows

$$\begin{aligned} & \min_x c \cdot x \\ & \text{subject to} \\ & A^{\leq} x \leq b^{\leq}, \quad A^{\geq} x \geq b^{\geq}, \quad A^{\text{=}} x = b^{\text{=}}, \quad x \geq 0 \end{aligned}$$

- Note in the previous example it was convenient to use two indices c and f to denote the components x_{cf} , however, it still has this structure

AICE1005

Algorithms and Analysis

10

Linear Program Applications

- A huge number of problems can be mapped to linear programming problems
- Or modelled as linear (even when they're not, e.g. oil extraction)
- Realistic problems might have many more constraints and large number of variables
- State of the art solvers can deal with problems with hundreds of thousands or even millions of variables

AICE1005

Algorithms and Analysis

12

History

- Linear programming was "invented" by Leonid Kantorovich in 1939 to help Soviet Russia maximise its production
- It was kept secret during the war, but was finally made public in 1947 when George Dantzig published the **simplex method** which still today is a standard method for solving linear programs
- John von Neumann developed the idea of duality (you can turn a maximisation problem for a set of variables x into a minimisation problem for a dual set of variables λ associated with each constraint)
- von Neumann used this idea as the basis for "game theory"

AICE1005

Algorithms and Analysis

14

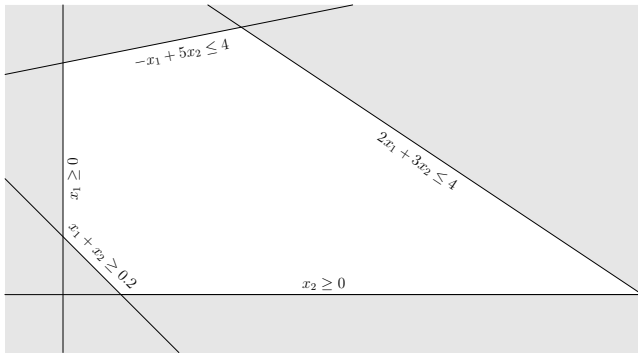
Structure of Linear Programs

- Before we go into the details of solving linear programs its useful to consider the structure of the solutions
- The set of x that satisfy all the constraints is known as the set of **feasible solutions**
- The set of feasible solutions may be empty in which case it is impossible to satisfy all the constraints
- This is rather disappointing, but usually doesn't happen if we have formulated a sensible problem

AICE1005

Algorithms and Analysis

16



AIICE1005

Algorithms and Analysis

17

- The space of feasible solutions is a polyhedra or polytope
- The maximum or minimum solution will always lie at a vertex of the polytope
- Our solution policy will be to start at a vertex and move to a neighbouring vertex that gives the best improvement in cost
- When this isn't possible then we are finished
- However, there is still a lot of work to realise this solution strategy

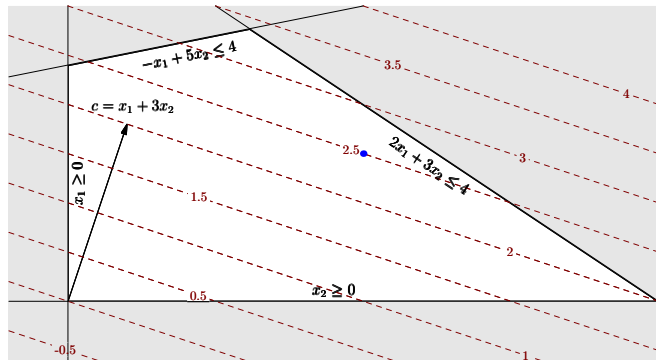
AIICE1005

Algorithms and Analysis

18

Optimal Solution

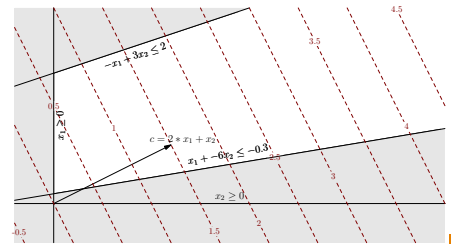
Unbounded Solutions



AIICE1005

Algorithms and Analysis

19



- But usually this would not happen because of the problem definition

AIICE1005

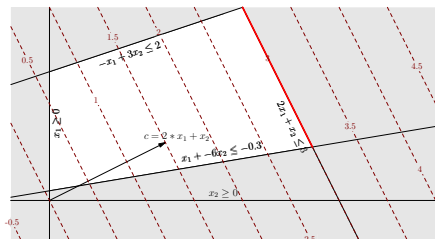
Algorithms and Analysis

20

Multiple Solutions

Outline

- You can also get multiple solutions if a constraint is orthogonal to the objective function



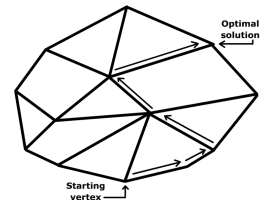
- Nevertheless the optimal will be at a vertex

AIICE1005

Algorithms and Analysis

21

1. Examples
2. Linear Programs
3. Properties of Solution
4. **Normal Form**



AIICE1005

Algorithms and Analysis

22

Converting Linear Programs

Slack Variables

- Solving full linear programs is difficult
- However, it is much easier to solve linear programs in **normal form**
- This is basically a form where we get rid of all inequalities and rewriting the equalities
- Fortunately its rather easy to convert linear programs to normal form

- We can change an inequality into an equality by introducing a new "slack" variable

- E.g.

$$\begin{aligned} a_1 \cdot x &\geq 0 &\Rightarrow & a_1 \cdot x - z_1 = 0 & z_1 &\geq 0 \\ a_2 \cdot x &\leq 0 &\Rightarrow & a_2 \cdot x + z_2 = 0 & z_2 &\geq 0 \end{aligned}$$

z_1 (the excess) and z_2 (the deficit) are known as slack variables

- We eliminate inequalities at the expense of increasing the number of variables
- We can treat the slack variables on an equivalent footing to the normal variables (they just provide a different way of describing the original problem)

AIICE1005

Algorithms and Analysis

23

AIICE1005

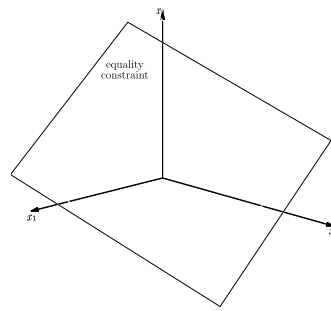
Algorithms and Analysis

24

Normal Form

- A linear program with only equality constraints is said to be in **normal form**
- We will find in the next lecture that this is a convenient form for solving linear programs
- An equality constraint restricts the solutions to a subspace (some lower dimensional space)

Solving Linear Programming



- The basic feasible points for LP problems with n variables and m constraints have at least $n - m$ zero variables
- Typical number of basic feasible solutions is $\binom{n}{m} \geq \left(\frac{n}{m}\right)^m$
- Simplex algorithm organises iterative search for global solutions

Lessons

- There are a huge number of problems that can be set up as linear programs
- They are particularly useful in resource allocation where the resources are all positive
- The solution to linear programming problems is at the vertex of the feasible space (intersection of constraints)
- We can search for solutions by moving from vertex to vertex
- We can transform inequality constraints to equality constraints using slack variables