Algorithms and Analysis

Lesson 5: *Use Arrays*

Variable length arrays, implementing stacks

Outline

1. **Why Arrays?**

2. Variable Length Arrays

3. Programming Language

4. Implementing Stacks

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  ```
  int *array = new int[20]
  ```

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  `int *array = new int[20]`

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  int *array = **new int**[20]

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  `int *array = new int[20]`

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  `int *array = new int[20]`

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  `int *array = new int[20]`

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance

Use Arrays

- An array is a contiguous chunk of memory

- In C we can create arrays using
  ```
  int *array = new int[20]
  ```

- The array has an access time of $\Theta(1)$

- The constant factor is small (i.e. access time ≈ 1 time step)

- Arrays provide a very efficient use of memory

- 95% of the time using arrays is going to give you the best performance, although never use raw arrays!

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

  - ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

  ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

  ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

  ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

  ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are

Disadvantages of Arrays

- Arrays have a fixed length

- Very often we don't know how big an array we want

    ⋆ E.g. reading words from a file

- Adding or deleting elements from the middle of an array is costly

- Sorted arrays are expensive to maintain

- Arrays don't know how big they are—annoying

Outline

1. Why Arrays?

2. **Variable Length Arrays**

3. Programming Language

4. Implementing Stacks

Variable Length Arrays

- We want a variable length array

- Initially a variable length array would have length zero

- We should be able to

  ⋆ Add an element to an array
  ⋆ Access any element in the array
  ⋆ Change an element
  ⋆ Delete elements
  ⋆ Know how many elements we have

Variable Length Arrays

- We want a variable length array

- Initially a variable length array would have length zero

- We should be able to

  ⋆ Add an element to an array
  ⋆ Access any element in the array
  ⋆ Change an element
  ⋆ Delete elements
  ⋆ Know how many elements we have

Variable Length Arrays

- We want a variable length array

- Initially a variable length array would have length zero

- We should be able to

  ⋆ Add an element to an array
  ⋆ Access any element in the array
  ⋆ Change an element
  ⋆ Delete elements
  ⋆ Know how many elements we have

ADT for a List

- What do we want of a list of **int**s?

ADT for a List

- What do we want of a list of **int**s?

  ⋆ **void** push_back(**int** value)

ADT for a List

- What do we want of a list of **int**s?

  ⋆ **void** push_back(**int** value)
  ⋆ random access array[i]

ADT for a List

- What do we want of a list of **int**s?

  ⋆ **void** push_back(**int** value)
  ⋆ random access array[i]
  ⋆ **int** size()

ADT for a List

- What do we want of a list of **int**s?

  - ⋆ **void** push_back(**int** value)
  - ⋆ random access array[i]
  - ⋆ **int** size()

- It would be useful if it resized

ADT for a List

- What do we want of a list of **int**s?

  ⋆ **void** push_back(**int** value)
  ⋆ random access array[i]
  ⋆ **int** size()

- It would be useful if it resized

- It would be great to have some algorithms (e.g. sort) that can be
  run on a list

Implementation

- How should we implement a list?

Implementation

- How should we implement a list?

- Use an array, of course!

Implementation

- How should we implement a list?

- Use an array, of course!

- We need to distinguish between

  ★ the number of elements in the list `size()`
  ★ the number of elements in the array `capacity()`

Implementation

- How should we implement a list?

- Use an array, of course!

- We need to distinguish between

  ⋆ the number of elements in the list `size()`
  ⋆ the number of elements in the array `capacity()`

- If the number of elements grows larger than the capacity then we need to increase the capacity

---

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use cn bits of memory where we want to keep c small

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use cn bits of memory where we want to keep c small

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use $c\,n$ bits of memory where we want to keep c small

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use cn bits of memory where we want to keep c small

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use $c\,n$ bits of memory where we want to keep c small

Initial Capacity

- We could prevent resizing arrays by using a huge initial capacity

- However, how big is big enough?

- What happens when we have an array of arrays?

- Memory like time is resource we should care about

- In an analogy with **time complexity** we also care about **space complexity** (i.e. how much memory we need)

- If we want to store n elements it is reasonable to expect that we use $c\,n$ bits of memory where we want to keep c small
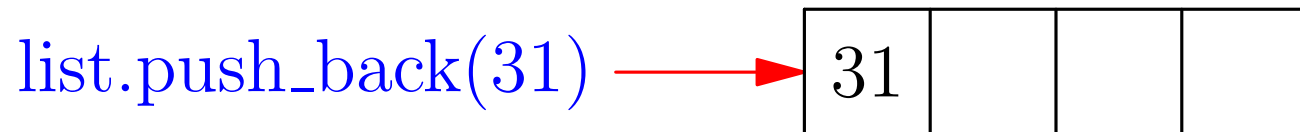
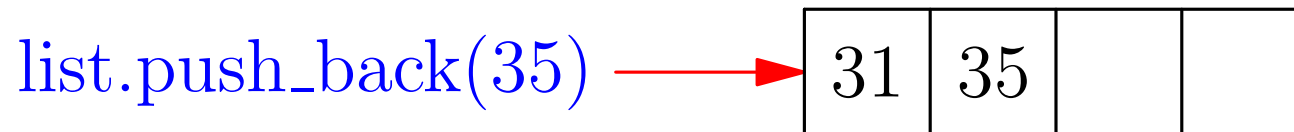# Resizing Memory

• We start with some reasonable capacity
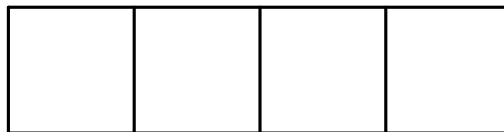
Resizing Memory

- We start with some reasonable capacity

- We can add elements

$$\text{list.push\_back}(31) \longrightarrow \boxed{31} \quad \boxed{} \quad \boxed{} \quad \boxed{}$$

Resizing Memory

- We start with some reasonable capacity

- We can add elements

$$\text{list.push\_back}(35) \longrightarrow \boxed{31 \mid 35 \mid \mid }$$

Resizing Memory

- We start with some reasonable capacity

- We can add elements

$\text{list.push\_back}(85)$ ⟶ | 31 | 35 | 85 | |

Resizing Memory

- We start with some reasonable capacity

- We can add elements

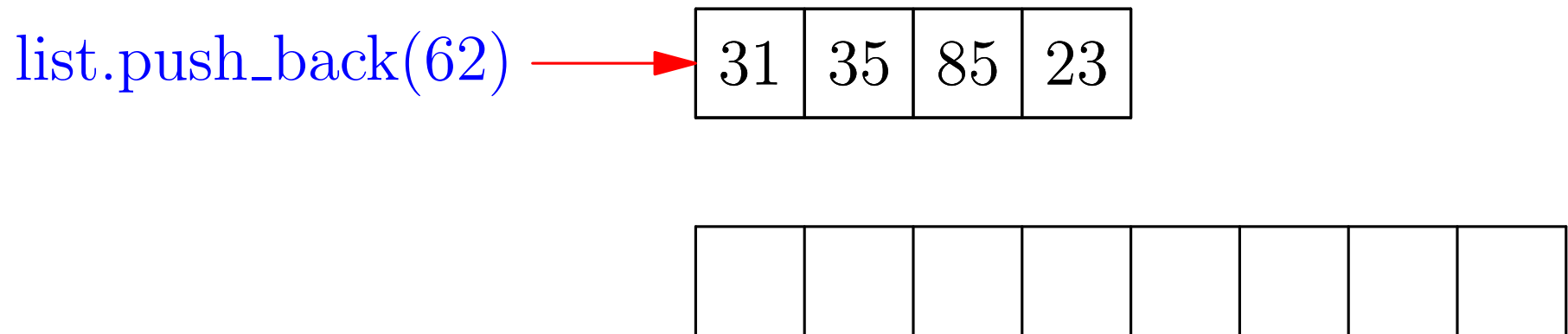list.push_back(23) ⟶ | 31 | 35 | 85 | 23 |

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

$$\text{list.push\_back}(62) \longrightarrow \boxed{\begin{array}{c|c|c|c} 31 & 35 & 85 & 23 \end{array}}$$

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array

list.push_back(62) $\longrightarrow$

31	35	85	23

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array
  ⋆ copy the existing elements from the old array to the new array

list.push_back(62) ⟶ | 31 | 35 | 85 | 23 |

| 31 | | | | | | | |

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array
  ⋆ copy the existing elements from the old array to the new array

list.push_back(62) ⟶ | 31 | 35 | 85 | 23 |

| 31 | 35 | | | | | | |

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array
  ⋆ copy the existing elements from the old array to the new array

list.push_back(62) ⟶ | 31 | 35 | 85 | 23 |

| 31 | 35 | 85 | | | | | |

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  - ⋆ create a new array with double the capacity of the old array
  - ⋆ copy the existing elements from the old array to the new array

list.push_back(62) ⟶ | 31 | 35 | 85 | 23 |

| 31 | 35 | 85 | 23 | | | | |

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array
  ⋆ copy the existing elements from the old array to the new array

$\text{list.push\_back}(62)$ $\longrightarrow$

31	35	85	23				

Resizing Memory

- We start with some reasonable capacity

- We can add elements until we reach the capacity

- A simple method for resizing memory is

  ⋆ create a new array with double the capacity of the old array
  ⋆ copy the existing elements from the old array to the new array

$\text{list.push\_back}(62)$ ⟶

31	35	85	23	62			

Amortised Time Analysis

- How efficient is resizing?

- Most `push_back(elem)` operations are $\Theta(1)$

- When we are at full capacity we have to copy all elements

- Adding to a full array is slow but it is **amortised** by other quick adds

  amortised: effect of a single operations 'deadened' by other operations

Amortised Time Analysis

- How efficient is resizing?

- Most `push_back(elem)` operations are $\Theta(1)$

- When we are at full capacity we have to copy all elements

- Adding to a full array is slow but it is **amortised** by other quick adds

    amortised: effect of a single operations 'deadened' by other operations

Amortised Time Analysis

- How efficient is resizing?

- Most `push_back(elem)` operations are $\Theta(1)$

- When we are at full capacity we have to copy all elements

- Adding to a full array is slow but it is **amortised** by other quick adds

  amortised: effect of a single operations 'deadened' by other operations

Amortised Time Analysis

- How efficient is resizing?

- Most `push_back(elem)` operations are $\Theta(1)$

- When we are at full capacity we have to copy all elements

- Adding to a full array is slow but it is **amortised** by other quick adds

  amortised: effect of a single operations 'deadened' by other operations

Amortised Time Analysis

- How efficient is resizing?

- Most `push_back(elem)` operations are $\Theta(1)$

- When we are at full capacity we have to copy all elements

- Adding to a full array is slow but it is **amortised** by other quick adds

    amortised: effect of a single operations 'deadened' by other operations

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

    ⋆ adds: 100

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

  - ⋆ adds: 100
  - ⋆ copies: 10

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

    ⋆ adds: 100
    ⋆ copies: 10+20

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

  ⋆ adds: 100
  ⋆ copies: 10+20+40

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

  ★ adds: 100
  ★ copies: 10+20+40+80

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

    ⋆ adds: 100
    ⋆ copies: 10+20+40+80=150

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

  - ⋆ adds: 100
  - ⋆ copies: 10+20+40+80=150
  - ⋆ **new int**[]: 4

Example

- If we have an initial capacity of 10 and add 100 elements then the number of operations needed is

  - ⋆ adds: 100
  - ⋆ copies: 10+20+40+80=150
  - ⋆ **new int**[]: 4

- 250 adds and copies operations $+$ 4 **new** operations

General Time Analysis

- If we perform N adds with an initial capacity of n
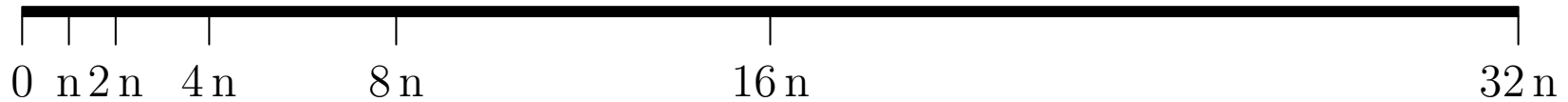
General Time Analysis

- If we perform N adds with an initial capacity of n

- We must perform m copies where

```
├──┬─┬──┬─────┬─────────────┬───────────────────────────────┤
0  n 2n  4n      8n            16n                             32n
```
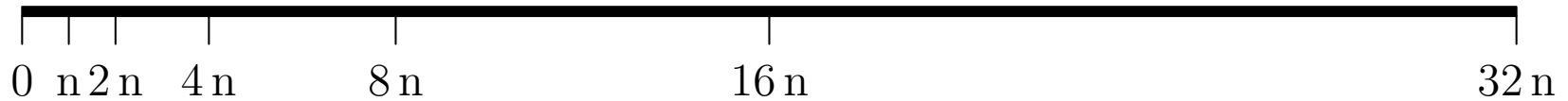
$$n \times 2^{m-1} < N \leq n \times 2^m$$

General Time Analysis

- If we perform N adds with an initial capacity of n

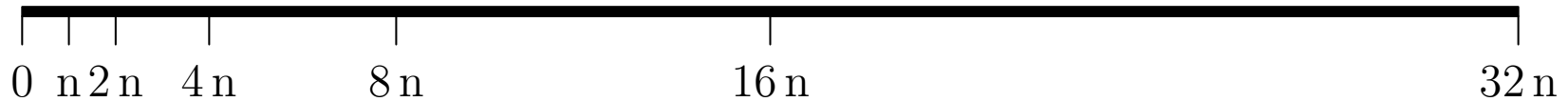- We must perform m copies where

```
 0  n 2n  4n        8n                16n                                32n
```

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2 \left(\frac{N}{n} \right) \right\rceil$$

General Time Analysis

- If we perform N adds with an initial capacity of n
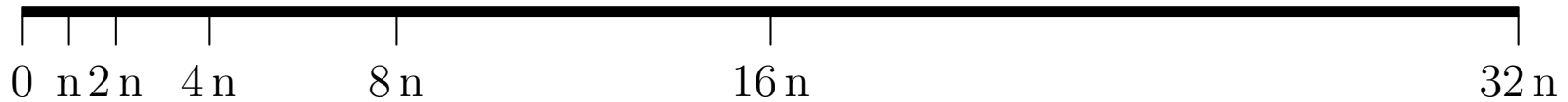
- We must perform m copies where

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2\left(\frac{N}{n}\right) \right\rceil$$

- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1}\,n$$

General Time Analysis

- If we perform N adds with an initial capacity of n

- We must perform m copies where

```
0  n 2n   4n        8n              16n                                    32n
```

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2 \left(\frac{N}{n} \right) \right\rceil$$

- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1} n = n(1 + 2 + \cdots + 2^{m-1})$$

General Time Analysis

- If we perform N adds with an initial capacity of n
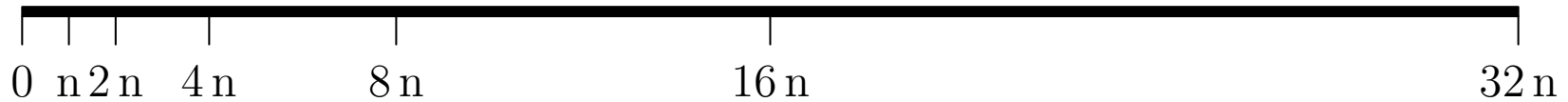
- We must perform m copies where

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2 \left(\frac{N}{n} \right) \right\rceil$$

- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1} \, n = n(1 + 2 + \cdots + 2^{m-1}) = n \left(2^m - 1 \right)$$

General Time Analysis

- If we perform N adds with an initial capacity of n
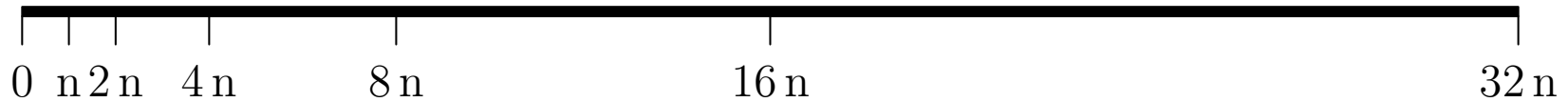
- We must perform m copies where

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2\left(\frac{N}{n}\right) \right\rceil$$

- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1}\,n = n(1 + 2 + \cdots + 2^{m-1}) = n\left(2^m - 1\right)$$

- Total number of operations is (using $\lceil \log(a) \rceil < \log(a) + 1$)

$$N + n\left(2^m - 1\right)$$

General Time Analysis

- If we perform N adds with an initial capacity of n

- We must perform m copies where

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2\left(\frac{N}{n}\right) \right\rceil$$

- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1}\,n = n(1 + 2 + \cdots + 2^{m-1}) = n\,(2^m - 1)$$

- Total number of operations is (using $\lceil \log(a) \rceil < \log(a) + 1$)

$$N + n\,(2^m - 1) = N + n2^{\left\lceil \log_2\left(\frac{N}{n}\right) \right\rceil} - n$$

Algorithms and Analysis

General Time Analysis

- If we perform N adds with an initial capacity of n

- We must perform m copies where

$$0 \quad n\,2\,n \quad 4\,n \qquad 8\,n \qquad\qquad 16\,n \qquad\qquad\qquad\qquad 32\,n$$

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2\left(\frac{N}{n}\right) \right\rceil$$
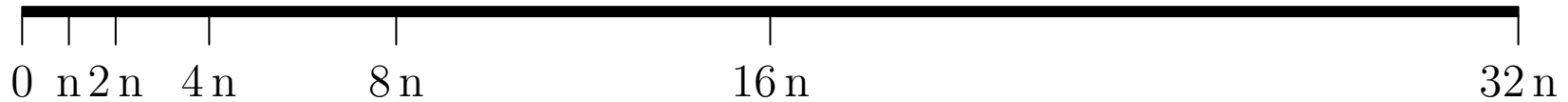
- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1}\,n = n(1 + 2 + \cdots + 2^{m-1}) = n\left(2^m - 1\right)$$

- Total number of operations is (using $\lceil \log(a) \rceil < \log(a) + 1$)

$$N + n\left(2^m - 1\right) = N + n 2^{\left\lceil \log_2\left(\frac{N}{n}\right)\right\rceil} - n < N + 2N - n$$

General Time Analysis

- If we perform N adds with an initial capacity of n

- We must perform m copies where

$$n \times 2^{m-1} < N \leq n \times 2^m \qquad \text{i.e.} \quad m = \left\lceil \log_2 \left(\frac{N}{n} \right) \right\rceil$$

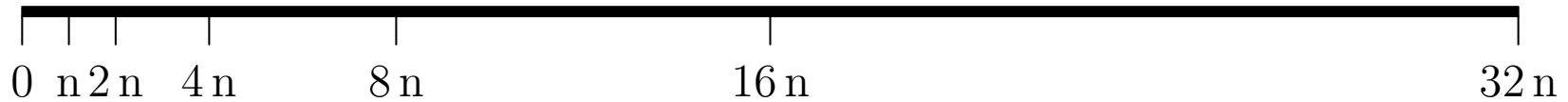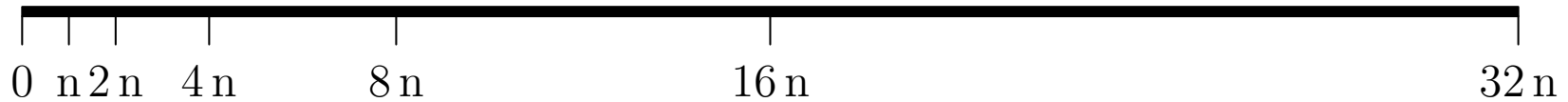- The number of elements copied is

$$n + 2n + 4n + \cdots + 2^{m-1} n = n(1 + 2 + \cdots + 2^{m-1}) = n \left(2^m - 1 \right)$$

- Total number of operations is (using $\lceil \log(a) \rceil < \log(a) + 1$)

$$N + n \left(2^m - 1 \right) = N + n 2^{\left\lceil \log_2 \left(\frac{N}{n} \right) \right\rceil} - n < N + 2N - n < 3N$$

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order

- Linked lists allow you to splice in a sublist into a list in constant time

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order

- Linked lists allow you to splice in a sublist into a list in constant time

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order

- Linked lists allow you to splice in a sublist into a list in constant time

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order (binary trees will do that for you much more efficiently)

- Linked lists allow you to splice in a sublist into a list in constant time

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order (binary trees will do that for you much more efficiently)

- Linked lists allow you to splice in a sublist into a list in constant time

Insertion and Deletion

- `vector<T>` is very useful and very fast for lots of things

- But if you try to insert or delete an element anywhere other than the end then you have to shove all the subsequent elements one space forward

- This is not the right data structure if you want to keep elements in order (binary trees will do that for you much more efficiently)

- Linked lists allow you to splice in a sublist into a list in constant time although linked lists have a lot of drawbacks

Outline

1. Why Arrays?

2. Variable Length Arrays

3. **Programming Language**

4. Implementing Stacks

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C)

- **Java** was designed to be vary safe (avoiding lots of bugs)

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

---

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C)

- **Java** was designed to be vary safe (avoiding lots of bugs)

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C), it pays the price of allowing bugs that hard to detect

- **Java** was designed to be vary safe (avoiding lots of bugs)

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C), it pays the price of allowing bugs that hard to detect

- **Java** was designed to be vary safe (avoiding lots of bugs)

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C), it pays the price of allowing bugs that hard to detect

- **Java** was designed to be vary safe (avoiding lots of bugs), but is not fast and a bit long winded

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C), it pays the price of allowing bugs that hard to detect

- **Java** was designed to be vary safe (avoiding lots of bugs), but is not fast and a bit long winded

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code

Computer Languages

- Different computer languages are designed for different roles and have different advantages and disadvantages

- **C++** was designed to be fast (as fast as C), it pays the price of allowing bugs that hard to detect

- **Java** was designed to be vary safe (avoiding lots of bugs), but is not fast and a bit long winded

- **Python** was designed so you can rapidly write powerful programmes with a small amount of code, but it is not fast or safe

Problems with C++

- Amongst a number of issues that make C++ dangerous are

  - ⋆ Memory management
  - ⋆ Writing to parts of memory that you should not
  - ⋆ Multiple inheritance

- However, by using existing data structures (STL) and following established programming patterns these don't have to be an issue

Problems with C++

- Amongst a number of issues that make C++ dangerous are

  ⋆ Memory management
  ⋆ Writing to parts of memory that you should not
  ⋆ Multiple inheritance, although you seldom need to do this

- However, by using existing data structures (STL) and following established programming patterns these don't have to be an issue

Problems with C++

- Amongst a number of issues that make C++ dangerous are

    ⋆ Memory management
    ⋆ Writing to parts of memory that you should not
    ⋆ Multiple inheritance, although you seldom need to do this

- However, by using existing data structures (STL) and following established programming patterns these don't have to be an issue

Memory Management

- Most programming languages have two types of memory

    The Stack: is the area of memory controlled by compiler for local variables, function calls, etc.

    The Heap: is area that the programmer (you) can request

Memory Management

• Most programming languages have two types of memory

The Stack: is the area of memory controlled by compiler for local variables, function calls, etc.

The Heap: is area that the programmer (you) can request

Memory Management

- Most programming languages have two types of memory

  The Stack: is the area of memory controlled by compiler for local variables, function calls, etc.

  The Heap: is area that the programmer (you) can request, which is nice

Memory Management

- Most programming languages have two types of memory

  The Stack: is the area of memory controlled by compiler for local variables, function calls, etc.

  The Heap: is area that the programmer (you) can request, which is nice

- In C++ you are given the **right** to ask for memory

  ```
  int *storage = new int[n];
  ```

Memory Management

- Most programming languages have two types of memory

  The Stack: is the area of memory controlled by compiler for local variables, function calls, etc.

  The Heap: is area that the programmer (you) can request, which is nice

- In C++ you are given the **right** to ask for memory

  ```cpp
  int *storage = new int[n];
  ```

- You have **responsibility** to free the memory

  ```cpp
  delete[] storage;
  ```

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete[] `storage;`

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an *undefined behaviour*, but often the programme will crash

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete[] `storage;`

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an *undefined behaviour*, but often the programme will crash

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete[] storage;

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an *undefined behaviour*, but often the programme will crash

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete`[] storage;`

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an $undefined$ $behaviour$, but often the programme will crash

---

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete[] storage;

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an *undefined behaviour*, but often the programme will crash

Trouble with Memory Management

- If you don't release memory acquired with `new` using `delete` you cause a **memory leak**

- Often memory leaks are no concern, but in large programs memory leaks will rapid exhaust the computer's memory, slowing down the code and eventually leading to the programme crashing

- To release a block of memory we can use:
  delete`[] storage;`

- Now `storage` is a **dangling pointer** and must not be used as it is no longer valid

- If we accidentally delete the storage twice we get an *undefined behaviour*, but often the programme will crash

---

Resource Acquistion is Initialisation (RAII)

- Java and Python use garbage collectors which automatically checks whether memory can be accessed and if not it is removed

- In C++ this is your responsibility

- But there is a standard **programming pattern** to elevate the problem known as **Resource Acquistion is Initialisation (RAII)**

  Wrap all resources in classes. Request the resources in the constructor and release the resource in the destructor

- When the object goes out of scope (you leave a `for` loop, function call, etc.) the destructor is called and the resource is safely released

---

Resource Acquistion is Initialisation (RAII)

- Java and Python use garbage collectors which automatically checks whether memory can be accessed and if not it is removed

- In C++ this is your responsibility

- But there is a standard **programming pattern** to elevate the problem known as **Resource Acquistion is Initialisation (RAII)**

  Wrap all resources in classes. Request the resources in the constructor and release the resource in the destructor

- When the object goes out of scope (you leave a `for` loop, function call, etc.) the destructor is called and the resource is safely released

---

Resource Acquistion is Initialisation (RAII)

- Java and Python use garbage collectors which automatically checks whether memory can be accessed and if not it is removed

- In C++ this is your responsibility

- But there is a standard **programming pattern** to elevate the problem known as **Resource Acquistion is Initialisation (RAII)**

  Wrap all resources in classes. Request the resources in the constructor and release the resource in the destructor

- When the object goes out of scope (you leave a `for` loop, function call, etc.) the destructor is called and the resource is safely released

---

Resource Acquistion is Initialisation (RAII)

- Java and Python use garbage collectors which automatically checks whether memory can be accessed and if not it is removed

- In C++ this is your responsibility

- But there is a standard **programming pattern** to elevate the problem known as **Resource Acquistion is Initialisation (RAII)**

  Wrap all resources in classes. Request the resources in the constructor and release the resource in the destructor

- When the object goes out of scope (you leave a `for` loop, function call, etc.) the destructor is called and the resource is safely released

---

Resource Acquistion is Initialisation (RAII)

- Java and Python use garbage collectors which automatically checks whether memory can be accessed and if not it is removed

- In C++ this is your responsibility

- But there is a standard **programming pattern** to elevate the problem known as **Resource Acquistion is Initialisation (RAII)**

    Wrap all resources in classes. Request the resources in the constructor and release the resource in the destructor

- When the object goes out of scope (you leave a `for` loop, function call, etc.) the destructor is called and the resource is safely released

---

Writing over Memory

- In C++ the following will compile and run

```cpp
int *array = new int[4];
int *a = new int[2];
double *darray = new double[4];
array[4] = 4;
```

- However `array[4]` has not been assigned (unlike `array[0]`, `array[1]`, `array[2]` and `array[3]`)

- The memory on the heap corresponding to the address of `array[4]` might have been assigned to `a[0]` in which case you may inadvertently have set `a[0]` to 4 leading to the program not doing what you want

- It might be that you have put an `int` into `darray[0]` which will then crash the system when you read `darray[0]`

Writing over Memory

- In C++ the following will compile and run

```cpp
int *array = new int[4];
int *a = new int[2];
double *darray = new double[4];
array[4] = 4;
```

- However array[4] has not been assigned (unlike array[0], array[1], array[2] and array[3])

- The memory on the heap corresponding to the address of array[4] might have been assigned to a[0] in which case you may inadvertently have set a[0] to 4 leading to the program not doing what you want

- It might be that you have put an int into darray[0] which will then crash the system when you read darray[0]

Writing over Memory

- In C++ the following will compile and run

```
int *array = new int[4];
int *a = new int[2];
double *darray = new double[4];
array[4] = 4;
```

- However array[4] has not been assigned (unlike array[0], array[1], array[2] and array[3])

- The memory on the heap corresponding to the address of array[4] might have been assigned to a[0] in which case you may inadvertently have set a[0] to 4 leading to the program not doing what you want

- It might be that you have put an int into darray[0] which will then crash the system when you read darray[0]

Writing over Memory

- In C++ the following will compile and run

```
int *array = new int[4];
int *a = new int[2];
double *darray = new double[4];
array[4] = 4;
```

- However `array[4]` has not been assigned (unlike `array[0]`, `array[1]`, `array[2]` and `array[3]`)

- The memory on the heap corresponding to the address of `array[4]` might have been assigned to `a[0]` in which case you may inadvertently have set `a[0]` to 4 leading to the program not doing what you want

- It might be that you have put an `int` into `darray[0]` which will then crash the system when you read `darray[0]`

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Guarding Against Mistakes

- These are really hard problems to debug because where the program goes wrong or crashes can be very far from the assignment that caused the error

- Java takes the approach that it always tests whether you are writing in valid memory

- By default C++ doesn't even for data structures—making this check slows down random access

- Checks can also make pipeline optimisations harder to make

- The onus is on the user to use the memory correctly

Follow Programming Idioms

- Using common data structures and following common idioms will prevent most errors

```cpp
int n = 5;
vector<int> array(n);

for(int i=0; i<array.size(); ++i) {
  array[i] = i;
}

for(auto pt=array.begin(); pt != array.end(); ++pt){
  *pt *= 2
}

for(int& element: array) {
  element += 2;
}
```

Follow Programming Idioms

- Using common data structures and following common idioms will prevent most errors

```cpp
int n = 5;
vector<int> array(n);

for(int i=0; i<array.size(); ++i) {
  array[i] = i;
}

for(auto pt=array.begin(); pt != array.end(); ++pt){
  *pt *= 2
}

for(int& element: array) {
  element += 2;
}
```

Follow Programming Idioms

- Using common data structures and following common idioms will prevent most errors

```cpp
int n = 5;
vector<int> array(n);

for(int i=0; i<array.size(); ++i) {
  array[i] = i;
}

for(auto pt=array.begin(); pt != array.end(); ++pt){
  *pt *= 2
}

for(int& element: array) {
  element += 2;
}
```

Outline

1. Why Arrays?

2. Variable Length Arrays

3. Programming Language

4. **Implementing Stacks**

Stacks

- Lets look at implementing a stack

- Remember a stack has methods

  - ⋆ `push(Object )`
  - ⋆ `pop()`
  - ⋆ `top()`
  - ⋆ `empty()`

Stacks

- Lets look at implementing a stack

- Remember a stack has methods

  ⋆ `push(Object )`
  ⋆ `pop()`
  ⋆ `top()`
  ⋆ `empty()`

Implementation of Stack

```cpp
template <typename T>
class MyStack
{
private:
  std::vector<T> stack;

public:
  void push(const T& obj) {stack.push_back(obj);}

  T top() const {return stack.back();}

  T pop() {
    T tmp = stack.back();
    stack.pop_back();
    return tmp;
  }

  T empty() {return stack.size()==0;}
};
```

Implementation of Stack

```cpp
template <typename T>
class MyStack
{
private:
  std::vector<T> stack;

public:
  void push(const T& obj) {stack.push_back(obj);}

  T top() const {return stack.back();}

  T pop() {
    T tmp = stack.back();
    stack.pop_back();
    return tmp;
  }

  T empty() {return stack.size()==0;}
};
```

Implementation of Stack

```cpp
template <typename T>
class MyStack
{
private:
  std::vector<T> stack;

public:
  void push(const T& obj) {stack.push_back(obj);}

  T top() const {return stack.back();}

  T pop() {
    T tmp = stack.back();
    stack.pop_back();
    return tmp;
  }

  T empty() {return stack.size()==0;}
};
```

Implementation of Stack

```cpp
template <typename T>
class MyStack
{
private:
  std::vector<T> stack;

public:
  void push(const T& obj) {stack.push_back(obj);}

  T top() const {return stack.back();}

  T pop() {
    T tmp = stack.back();
    stack.pop_back();
    return tmp;
  }

  T empty() {return stack.size()==0;}
};
```

Implementation of Stack

```cpp
template <typename T>
class MyStack
{
private:
  std::vector<T> stack;

public:
  void push(const T& obj) {stack.push_back(obj);}

  T top() const {return stack.back();}

  T pop() {
    T tmp = stack.back();
    stack.pop_back();
    return tmp;
  }

  T empty() {return stack.size()==0;}
};
```

Notes on Implementation

- I don't need to write a constructor as C++ generates a default constructor that will initialise the stack correctly

- I don't need to write a desctuctor because by default the destructor for `vector<T>` will be called which releases memory

- I've written the `pop` command, that I like, but if I run
  ```
  stack<Widget> widget_stack;
  Widget w;
  widget_stack.push(w);
  Widget w1(widget.pop());
  ```
  if the last command throws an exception then the last term on the stack is lost for ever

Notes on Implementation

- I don't need to write a constructor as C++ generates a default constructor that will initialise the stack correctly

- I don't need to write a desctuctor because by default the destructor for `vector<T>` will be called which releases memory

- I've written the `pop` command, that I like, but if I run

```
stack<Widget> widget_stack;
Widget w;
widget_stack.push(w);
Widget w1(widget.pop());
```

  if the last command throws an exception then the last term on the stack is lost for ever

Notes on Implementation

- I don't need to write a constructor as C++ generates a default constructor that will initialise the stack correctly

- I don't need to write a desctuctor because by default the destructor for `vector<T>` will be called which releases memory

- I've written the `pop` command, that I like, but if I run

```
stack<Widget> widget_stack;
Widget w;
widget_stack.push(w);
Widget w1(widget.pop());
```

if the last command throws an exception then the last term on the stack is lost for ever

---

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory

Why not use a `vector`

- Surely it is mad to use `MyStack<T>` as I could just use the more powerful `vector<T>`

- I can make `MyStack<T>` as efficient as `vector<T>` by inlining function calls

- But why would I want to lose functional?

- By using `MyStack<T>` I am **declaring my intention** of using this data structure as a stack

- I'm not going to do something weird like modify an element inside the stack

- My code becomes self-explanatory—I don't need to write comments as it is clear what I am doing

Using `MyStack`

• Implementing a stack using a dynamically re-sizable array is trivial

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial
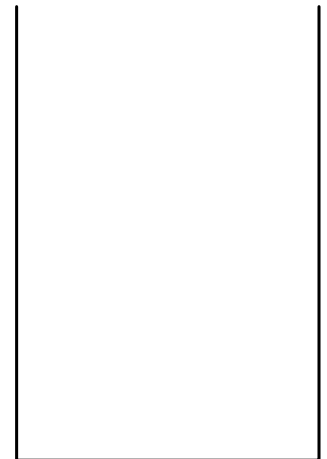
- Stacks have many applications

Using **MyStack**

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

you can cage a swallow can't you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file
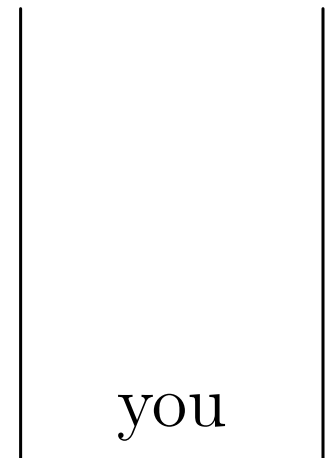
can cage a swallow can't you

you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

cage a swallow can't you

can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

a swallow can't you

cage
can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

swallow can't you

$$\begin{array}{|c|} \hline \\ \\ \\ a \\ cage \\ can \\ you \\ \hline \end{array}$$

Using `MyStack`

• Implementing a stack using a dynamically re-sizable array is trivial

• Stacks have many applications

• E.g. suppose we want to write a program to reverse the order of strings in a file

can't you

swallow
a
cage
can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

<div align="center">
you
</div>

$$\begin{array}{|c|}
\hline
\text{can't} \\
\text{swallow} \\
\text{a} \\
\text{cage} \\
\text{can} \\
\text{you} \\
\hline
\end{array}$$

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

$$
\begin{array}{c}
\text{you} \\
\text{can't} \\
\text{swallow} \\
\text{a} \\
\text{cage} \\
\text{can} \\
\text{you}
\end{array}
$$

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

you

$$\begin{array}{c} \text{can't} \\ \text{swallow} \\ \text{a} \\ \text{cage} \\ \text{can} \\ \text{you} \end{array}$$

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file
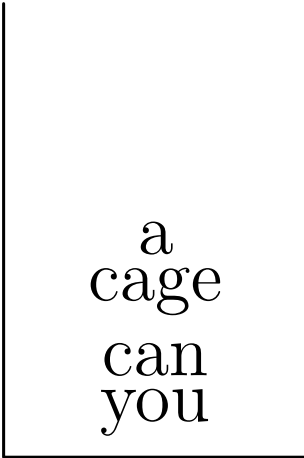
you can't

swallow
a
cage
can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

you can't swallow

```
   a
 cage
  can
  you
```

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

you can't swallow a

cage
can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file
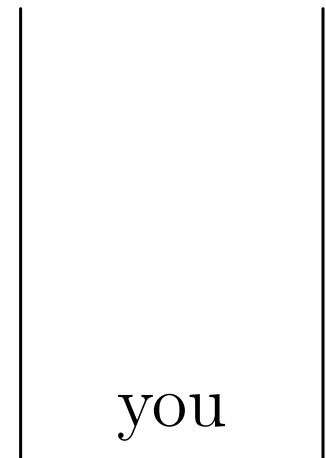
you can't swallow a cage

can
you

Using `MyStack`

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file

you can't swallow a cage can

you

Using **MyStack**

- Implementing a stack using a dynamically re-sizable array is trivial

- Stacks have many applications

- E.g. suppose we want to write a program to reverse the order of strings in a file
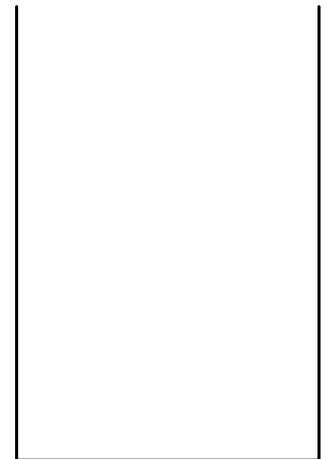
you can't swallow a cage can you

Reversing Strings in File

```cpp
#include <stack>
#include <iostream>
#include <fstream>
using namespace std;

int main(int argc, char *argv[]) {
  ifstream in(argv[1]);

  stack<string> stack;

  string word;
  while (in >> word)
    stack.push(word);

  while(!stack.empty()) {
    cout << stack.top() << ' ';
    stack.pop();
  }
}
```

Lessons

- Arrays are very efficient both in space (memory) and access time

- Resizing an array is not that costly

- insertion and deletion are expensive, $O(n)$

- Arrays are often the simplest way to implement many other data structures, e.g. stacks

- Use (dynamically re-sizable) arrays (vector¡T¿) frequently!

Lessons

- Arrays are very efficient both in space (memory) and access time

- Resizing an array is not that costly

- insertion and deletion are expensive, $O(n)$

- Arrays are often the simplest way to implement many other data structures, e.g. stacks

- Use (dynamically re-sizable) arrays (vector¡T¿) frequently!

Lessons

- Arrays are very efficient both in space (memory) and access time

- Resizing an array is not that costly

- insertion and deletion are expensive, $O(n)$

- Arrays are often the simplest way to implement many other data structures, e.g. stacks

- Use (dynamically re-sizable) arrays (vector¡T¿) frequently!

Lessons

- Arrays are very efficient both in space (memory) and access time

- Resizing an array is not that costly

- insertion and deletion are expensive, $O(n)$

- Arrays are often the simplest way to implement many other data structures, e.g. stacks

- Use (dynamically re-sizable) arrays (vector¡T¿) frequently!

Lessons

- Arrays are very efficient both in space (memory) and access time

- Resizing an array is not that costly

- insertion and deletion are expensive, $O(n)$

- Arrays are often the simplest way to implement many other data structures, e.g. stacks

- Use (dynamically re-sizable) arrays (vector¡T¿) frequently!