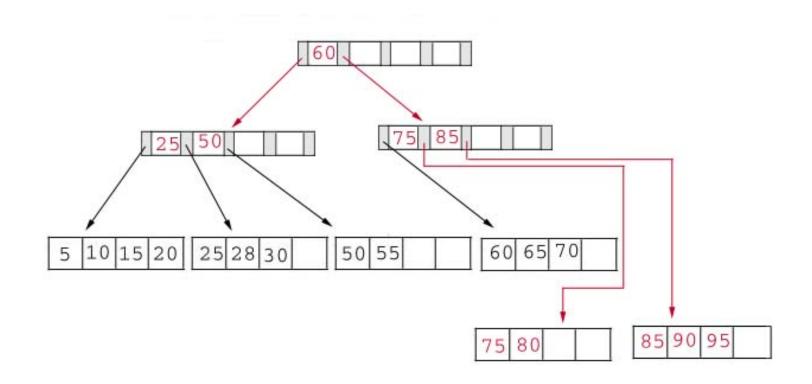
# **Algorithms and Analysis**

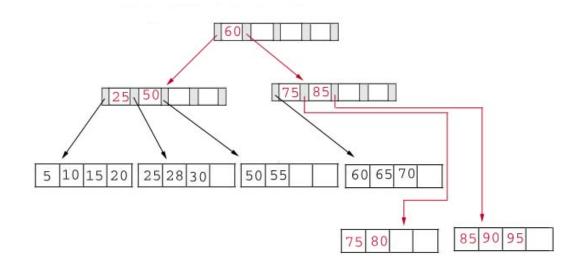
#### Lesson 9: Sometimes It Pays Not to Be Binary



B-Trees, Tries, Suffix Trees

#### **Outline**

- 1. B-Trees
- 2. Tries
- 3. Suffix Tree



- B-trees are balanced trees for fast search, finding successors and predecessors, insert, delete, maximum, minimum, etc.
- Not to be confused with binary trees
- They are designed to keep related data close to each other in (disk) memory to minimise retrieval time
- Important when working with large amount of data that is stored on secondary storage (e.g. disks)
- Used extensively in databases

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- An underlying assumption of Big-O is that all elementary operations take roughly the same amount of time
- This just isn't true of disk look-up
- The typical time of an elementary operation on a modern processor is  $10^{-9}$  seconds
- But a typical hard disk might do 7 200 revolutions per minute or 120 revolutions per second
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- In database applications we want to store data as large sets
- Storing data in binary trees is disastrous as we typically need around  $\log_2(n)$  disk accesses before we locate our data
- It is not unusual in databases for  $n=10\,000\,000$  so that  $\log_2(n)\approx 24$
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• To remedy this we can use M-way trees so that the access time is

$$\log_M(n) = \frac{\log_2(n)}{\log_2(M)}$$

- In practice we might use  $M \approx 200 \approx 2^8$  so we can reduce the depth of the tree by around a factor of 8
- The basic data structures for doing this is the B-tree
- There are many variants of B-tree, all trying to squeeze a bit more performances from the basic structure

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- A pretty basic implementation would obey the following rules
  - 1. The data items are stored at leaves
  - 2. The non-leaf nodes store up to M-1 keys to guide the search: key i represents the smallest key in subtree i+1
  - 3. The root is either a leaf or has between 2 and M children
  - 4. All non-leaf nodes except the root have between  $\lceil M/2 \rceil$  and M children
  - 5. All leaves are at the same depth and have between  $\lceil L/2 \rceil$  and L data entries

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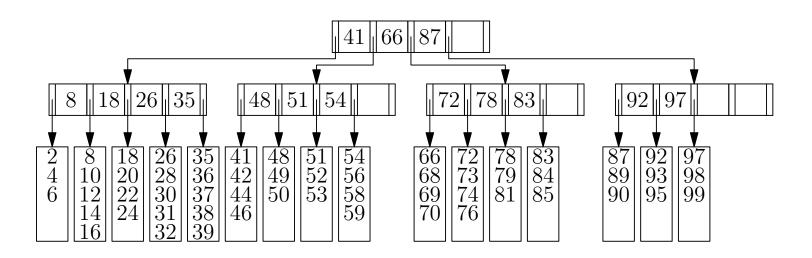
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- In the examples below we consider tiny M=L=5 which is unrealistic, but drawable

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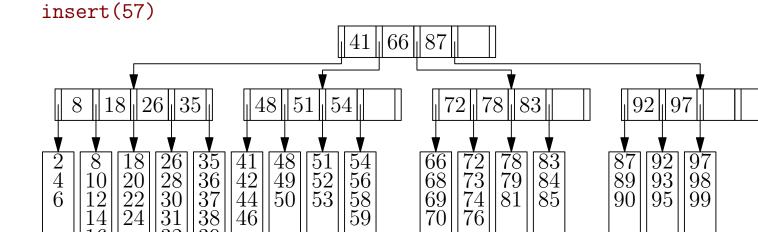
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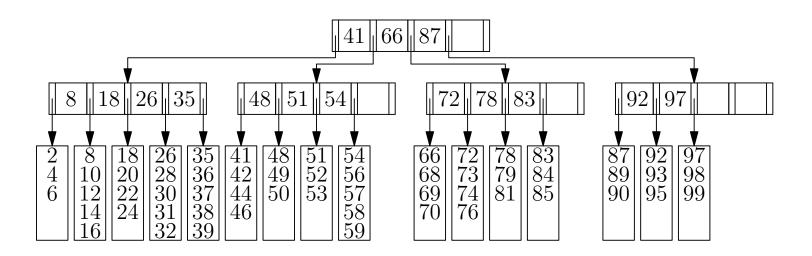
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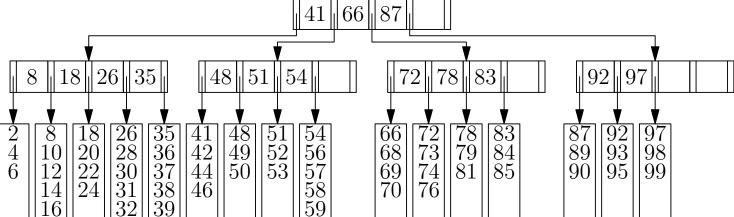
• M = 5, L = 5

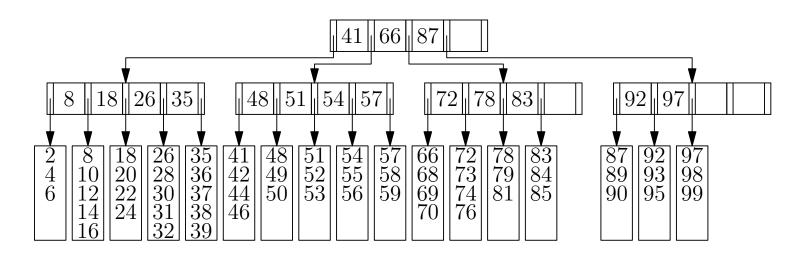


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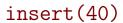


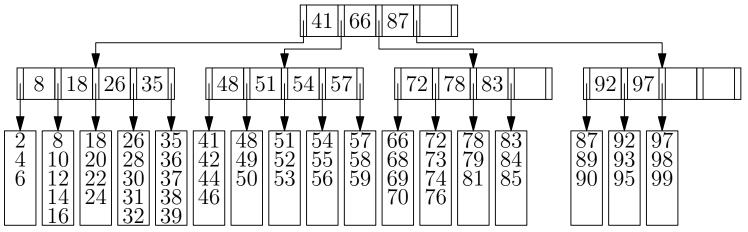




### **B-Tree Example**

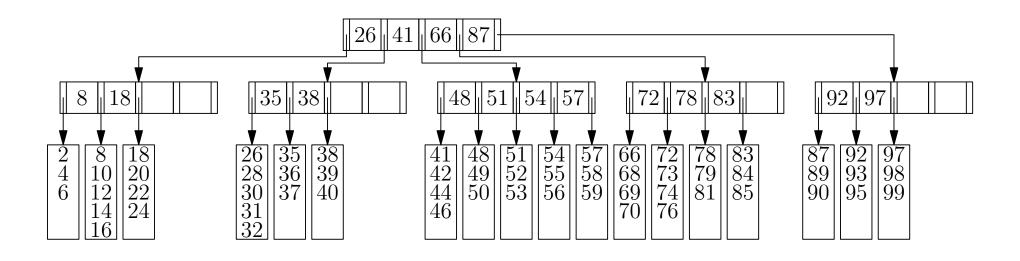
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## **B-Tree Example**

• M = 5, L = 5



- If the root is full then it can be split into two and a new root created
- B-trees also have to allow the removal of records without losing its structure
- There are a number of variant strategies (e.g. neighbouring nodes can adopt a child if the current node cannot expand any more)
- The actual implementation of B-trees is tricky because there are many special cases

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- B-trees are an important data structure for databases where reducing the number of disk searches is vital
- They tend to be much more complex than the other data structures we have seen
- The problem of disk access can be improved by replacing disk memory with solid-state drives (still slow compared to memory)
- For massive databases new data structures have been developed to allow faster (although less flexible) information access (e.g. NOSQL, MongoDB, Neo4j)

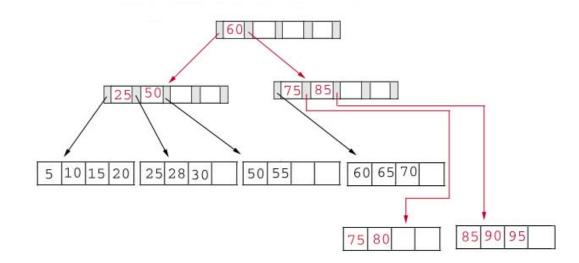
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### **Outline**

- 1. B-Trees
- 2. Tries
- 3. Suffix Tree



- A Trie (pron. 'try') or digital tree is a multiway tree often used for storing large sets of words
- They are trees with a possible branch for every letter of an alphabet
- Their names comes from the word retrieval
- Tries usually compactify the edges in the tree
- All words end with a special letter "\$"

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### **Trie**

 $\{a\$, abba\$, ba\$, baa\$, baba\$, cab\$\}$ b\* ab\* ba\* ca\* b abb\* bab\* baa\* cab\* abba\* baba\* baa\$ abba\$ baba\$

## Trie

{a, abba, ba, baa, baba, cab}

a b cab\$

cab\$

\$bab\$

a\$ baba\$

ba\$

baa\$

baba\$

## Trie

{a, abba, ba, baa, baba, cab}

ba

cab\$

cab\$

\$ba\*

ba\$

\$ba\$

baba\$

baba\$

baba\$

- Tries are yet another way of implementing sets
- They provide quick insertion, deletion and find
- Typically considerably quicker than binary trees and hash tables
- They are particularly good for spell checkers, completion algorithms, longest-prefix matching, hyphenation
- Each search finds the longest match between the words in the set and the query

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Т	THAT\$	IT\$
U		
V		
W	WITH\$	
X		
Y		
Z		

add("HIS")

	0	1	2
\$		I\$	
A			HAD\$
В			
С			
D			
E			
F			
G			
Н	H* (2)		
I	I* (1)		HIS\$
J			
K			
L			
M			
N	NOT\$	IN\$	
О			
P			
Q			
R			
S			
Т	THAT\$	IT\$	
U			
V			
W	WITH\$		
X			
Y			
Z			

add("HAVE")

	0	1	2
\$		I\$	
A			HAD\$
В			
С			
D			
E			
F			
G			
Н	H* (2)		
I	I* (1)		HIS\$
J			
K			
L			
M			
N	NOT\$	IN\$	
О			
P			
Q			
R			
S			
Т	THAT\$	IT\$	
U			
V			
W	WITH\$		
X			
Y			
Z			

add("HAVE")

			_	
	0	1	2	3
\$		I\$		
A			HA* (3)	
В				
C				
D				HAD\$
E				
F				
G				
H	H* (2)			
I	I* (1)		HIS\$	
J				
K				
L				
M				
N	NOT\$	IN\$		
O				
P				
Q				
R				
S				
$\overline{T}$	THAT\$	IT\$		
U				
V				HAVE\$
W	WITH\$			
X				
Y				
Z				

add("A")

	0	1	2	3
\$		I\$		
A			HA* (3)	
В				
C				
D				HAD\$
Е				
F				
G				
H	H* (2)			
I	I* (1)		HIS\$	
J				
K				
L				
M				
N	NOT\$	IN\$		
O				
P				
Q				
R				
S				
Т	THAT\$	IT\$		
U				
V				HAVE\$
W	WITH\$			
X				
Y				
Z				

add("A")

			_	
	0	1	2	3
\$		I\$		
A	A\$		HA* (3)	
В				
C				
$\overline{}$ D				HAD\$
E				
F				
G				
H	H* (2)			
I	I* (1)		HIS\$	
J				
K				
L				
M				
N	NOT\$	IN\$		
О				
P				
Q				
R				
S				
Т	THAT\$	IT\$		
U				
V				HAVE\$
W	WITH\$			
X				
Y				
Z				

add("WHICH")

			_	
	0	1	2	3
\$		I\$		
A	A\$		HA* (3)	
В				
C				
$\overline{}$ D				HAD\$
E				
F				
G				
H	H* (2)			
I	I* (1)		HIS\$	
J				
K				
L				
M				
N	NOT\$	IN\$		
О				
P				
Q				
R				
S				
Т	THAT\$	IT\$		
U				
V				HAVE\$
W	WITH\$			
X				
Y				
Z				

add("WHICH")

					1
	0	1	2	3	4
\$		I\$			
A	A\$		HA* (3)		
В					
$\mathbf{C}$					
D				HAD\$	
E					
F					
G					
H	H* (2)				WHICH\$
I	I* (1)		HIS\$		WITH\$
J					
K					
L					
M					
N	NOT\$	IN\$			
О					
P					
Q					
R					
S					
T	THAT\$	IT\$			
U					
V				HAVE\$	
W	W* (4)				
X					
Y					
Z					

add("HER")

	0	1	2	3	4
\$		I\$			
A	A\$		HA* (3)		
В					
С					
D				HAD\$	
Е					
F					
G					
Н	H* (2)				WHICH\$
I	I* (1)		HIS\$		WITH\$
J					
K					
L					
M					
N	NOT\$	IN\$			
O					
P					
Q					
R					
S					
T	THAT\$	IT\$			
U					
V				HAVE\$	
W	W* (4)				
X					
Y					
Z					

add("HER")

	0	1	2	3	4
\$		I\$			
A	A\$		HA* (3)		
В					
С					
D				HAD\$	
E			HER\$		
F					
G					
Н	H* (2)				WHICH\$
I	I* (1)		HIS\$		WITH\$
J					
K					
L					
M					
N	NOT\$	IN\$			
О					
P					
Q					
R					
S					
T	THAT\$	IT\$			
U					
V				HAVE\$	
W	W* (4)				
X					
Y					
Z					

add("AT")

	0	1	2	3	4
\$		I\$			
A	A\$		HA* (3)		
В					
С					
D				HAD\$	
E			HER\$		
F					
G					
Н	H* (2)				WHICH\$
I	I* (1)		HIS\$		WITH\$
J					
K					
L					
M					
N	NOT\$	IN\$			
О					
P					
Q					
R					
S					
T	THAT\$	IT\$			
U					
V				HAVE\$	
W	W* (4)				
X					
Y					
Z					

add("AT")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
C						
D				HAD\$		
E			HER\$			
F						
G						
Н	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				
O						
P						
Q						
R						
S						
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("IS")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
C						
D				HAD\$		
E			HER\$			
F						
G						
Н	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				
O						
Р						
Q						
R						
S						
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("IS")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
C						
				HAD\$		
E			HER\$			
F						
G						
Н	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				
О						
Р						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("AND")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
C						
D				HAD\$		
E			HER\$			
F						
G						
Н	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				
O						
P						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("AND")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
$\overline{}$ C						
D				HAD\$		
E			HER\$			
F						
G						
H	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				AND\$
O						
P						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
<u>Y</u>						
Z						

add("BUT")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В						
$\overline{}$ C						
D				HAD\$		
E			HER\$			
F						
G						
H	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				AND\$
O						
P						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
<u>Y</u>						
Z						

add("BUT")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В	BUT\$					
$\overline{}$ C						
D				HAD\$		
E			HER\$			
F						
G						
H	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				AND\$
O						
P						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("THE")

	0	1	2	3	4	5
\$		I\$				A\$
A	A* (5)		HA* (3)			
В	BUT\$					
C						
D				HAD\$		
E			HER\$			
F						
G						
H	H* (2)				WHICH\$	
I	I* (1)		HIS\$		WITH\$	
J						
K						
L						
M						
N	NOT\$	IN\$				AND\$
О						
P						
Q						
R						
S		IS\$				
T	THAT\$	IT\$				AT\$
U						
V				HAVE\$		
W	W* (4)					
X						
Y						
Z						

add("THE")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
C							
D				HAD\$			
E			HER\$				THE\$
F							
G							
H	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
O							
P							
Q							
R							
S		IS\$					
Т	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	W* (4)						
X							
Y							
Z							

add("FOR")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
C							
D				HAD\$			
E			HER\$				THE\$
F							
G							
H	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
O							
P							
Q							
R							
S		IS\$					
Т	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	W* (4)						
X							
Y							
Z							

add("FOR")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
C							
D				HAD\$			
E			HER\$				THE\$
F	FOR\$						
G							
Н	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
О							
P							
Q							
R							
S		IS\$					
T	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	W* (4)						
X							
Y							
$\mathbf{Z}$						·	

add("ON")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
C							
D				HAD\$			
E			HER\$				THE\$
F	FOR\$						
G							
H	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
О							
P							
Q							
R							
S		IS\$					
Т	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	$W^*(4)$						
X							
Y							
Z							

add("ON")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
$\overline{}$ C							
D				HAD\$			
E			HER\$				THE\$
F	FOR\$						
G							
Н	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
О	ON\$						
P							
Q							
R							
S		IS\$					
Т	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	$W^*(4)$						
X							
Y							
Z							

add("HE")

	0	1	2	3	4	5	6
\$		I\$				A\$	
A	A* (5)		HA* (3)				THAT\$
В	BUT\$						
C							
D				HAD\$			
E			HER\$				THE\$
F	FOR\$						
G							
H	H* (2)				WHICH\$		
I	I* (1)		HIS\$		WITH\$		
J							
K							
L							
M							
N	NOT\$	IN\$				AND\$	
O	ON\$						
P							
Q							
R							
S		IS\$					
Т	TH* (6)	IT\$				AT\$	
U							
V				HAVE\$			
W	W* (4)						
X							
Y							
Z							

add("HE")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)				THAT\$	
В	BUT\$							
$\overline{C}$								
D				HAD\$				
E			HE* (7)				THE\$	
F	FOR\$							
G								
H	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R								HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

#### add("WAS")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)				THAT\$	
В	BUT\$							
C								
$\overline{D}$				HAD\$				
E			HE* (7)				THE\$	
F	FOR\$							
G								
Н	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R								HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

#### add("WAS")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)		WAS\$		THAT\$	
В	BUT\$							
C								
D				HAD\$				
Е			HE* (7)				THE\$	
F	FOR\$							
G								
H	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R								HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

add("ARE")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)		WAS\$		THAT\$	
В	BUT\$							
$\overline{C}$								
D				HAD\$				
E			HE* (7)				THE\$	
F	FOR\$							
G								
Н	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R								HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

#### add("ARE")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)		WAS\$		THAT\$	
В	BUT\$							
C								
D				HAD\$				
E			HE* (7)				THE\$	
F	FOR\$							
G								
H	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R						ARE\$		HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

add("BY")

	0	1	2	3	4	5	6	7
\$		I\$				A\$		HE\$
A	A* (5)		HA* (3)		WAS\$		THAT\$	
В	BUT\$							
C								
D				HAD\$				
E			HE* (7)				THE\$	
F	FOR\$							
G								
H	H* (2)				WHICH\$			
I	I* (1)		HIS\$		WITH\$			
J								
K								
L								
M								
N	NOT\$	IN\$				AND\$		
O	ON\$							
P								
Q								
R						ARE\$		HER\$
S		IS\$						
T	TH* (6)	IT\$				AT\$		
U								
V				HAVE\$				
W	W* (4)							
X								
Y								
Z								

add("BY")

	0	1	2	3	4	5	6	7	8
\$		I\$				A\$		HE\$	
A	A* (5)		HA* (3)		WAS\$		THAT\$		
В	B* (8)								
$\overline{C}$									
D				HAD\$					
E			HE* (7)				THE\$		
F	FOR\$								
G									
H	H* (2)				WHICH\$				
I	I* (1)		HIS\$		WITH\$				
J									
K									
L									
M									
N	NOT\$	IN\$				AND\$			
O	ON\$								
P									
Q									
R						ARE\$		HER\$	
S		IS\$							
T	TH* (6)	IT\$				AT\$			
U									BUT\$
V				HAVE\$					
W	W* (4)								
X									
Y									BY\$
Z									

add("OR")

	0	1	2	3	4	5	6	7	8
\$		I\$				A\$		HE\$	
A	A* (5)		HA* (3)		WAS\$		THAT\$		
В	B* (8)								
$^{-}$ C									
D				HAD\$					
E			HE* (7)				THE\$		
F	FOR\$								
G									
H	H* (2)				WHICH\$				
I	I* (1)		HIS\$		WITH\$				
J									
K									
L									
M									
N	NOT\$	IN\$				AND\$			
O	ON\$								
P									
Q									
R						ARE\$		HER\$	
S		IS\$							
T	TH* (6)	IT\$				AT\$			
U									BUT\$
V				HAVE\$					
W	W* (4)								
X									
Y									BY\$
Z									

add("OR")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$					
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
О	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$								
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

add("AS")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$					
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
О	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$								
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

add("AS")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
C										
D				HAD\$						
Е			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$					
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
0	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

#### add("THIS")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$					
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
O	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

#### add("THIS")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
О	O* (9)									
Р										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
T	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

#### add("YOU")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
О	O* (9)									
Р										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
T	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y									BY\$	
Z										

#### add("YOU")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
O	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
T	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y	YOU\$								BY\$	
Z										

#### add("BE")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$			
F	FOR\$									
G										
Н	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
О	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y	YOU\$								BY\$	
Z						·				

#### add("BE")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$		BE\$	
F	FOR\$									
G										
H	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
O	O* (9)									
Р										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
Т	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y	YOU\$								BY\$	
Z										

#### add("FROM")

	0	1	2	3	4	5	6	7	8	9
\$		I\$				A\$		HE\$		
A	A* (5)		HA* (3)		WAS\$		THAT\$			
В	B* (8)									
$\overline{C}$										
D				HAD\$						
E			HE* (7)				THE\$		BE\$	
F	FOR\$									
G										
Н	H* (2)				WHICH\$					
I	I* (1)		HIS\$		WITH\$		THIS\$			
J										
K										
L										
M										
N	NOT\$	IN\$				AND\$				ON\$
O	O* (9)									
P										
Q										
R						ARE\$		HER\$		OR\$
S		IS\$				AS\$				
T	TH* (6)	IT\$				AT\$				
U									BUT\$	
V				HAVE\$						
W	W* (4)									
X										
Y	YOU\$								BY\$	
Z										

#### add("FROM")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)										
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z		·						·			

#### add("OF")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)										
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	$W^*(4)$										
X											
Y	YOU\$								BY\$		
Z		·						·			

#### add("OF")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("YOU")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("YOU")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
$\overline{D}$				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
$\overline{R}$						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("YOU") True

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
-C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("THE")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("THE")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
С											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("THE")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
$\overline{}$ D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z				·							

#### search("THE") True

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
Е			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
0	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z				·							

#### search("NO")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("NO")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("NO") False

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
$\overline{\mathbf{D}}$				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z								·			

#### search("THEM")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("THEM")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z								·			

#### search("THEM")

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
$\overline{C}$											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("THEM") False

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
С											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
$\overline{D}$				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
H	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M										 	
N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
P											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
E			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	$H^*(2)$				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
K											
L											
M											
N	NOT\$	IN\$				AND\$				ON\$	
О	O* (9)										FOR\$
Р											
Q											
R						ARE\$		HER\$		OR\$	FROM\$
S		IS\$				AS\$					
T	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
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F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
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L											
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N	NOT\$	IN\$				AND\$				ON\$	
O	O* (9)										FOR\$
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S		IS\$				AS\$					
Т	TH* (6)	IT\$				AT\$					
U									BUT\$		
V				HAVE\$							
W	W* (4)										
X											
Y	YOU\$								BY\$		
Z											

#### search("HERE") False

	0	1	2	3	4	5	6	7	8	9	10
\$		I\$				A\$		HE\$			
A	A* (5)		HA* (3)		WAS\$		THAT\$				
В	B* (8)										
C											
D				HAD\$							
Е			HE* (7)				THE\$		BE\$		
F	F* (10)									OF\$	
G											
Н	H* (2)				WHICH\$						
I	I* (1)		HIS\$		WITH\$		THIS\$				
J											
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N	NOT\$	IN\$				AND\$				ON\$	
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- Table-based tries typically waste large amounts of memory
- Often table-based tries are used for the first few layers, while lower levels use a less memory intensive data structure
- These days memory is less of a problem so table-based tries are acceptable for some applications
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## **Binary Tries**

- One extreme (though not uncommon) solution to address memory issues is to build a bit-level trie so the data-structure is a binary tree
- It differs from a binary tree in that the decisions to go left or right depends on the current bit
- Although you lose the advantage of a multiway tree (of reducing the depth) it does find the longest match and it speeds up finds which fail

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- Tries are slightly specialist and tend to get used in very particular applications
  - ★ Finding longest matches
  - ★ Completion, spell checking, etc.
- A basic trie is not too complicated, however, . . .
- There are many implementation which try to overcome the difficulty of wasting too much memory

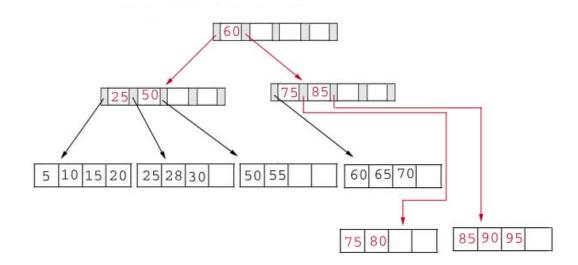
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### **Outline**

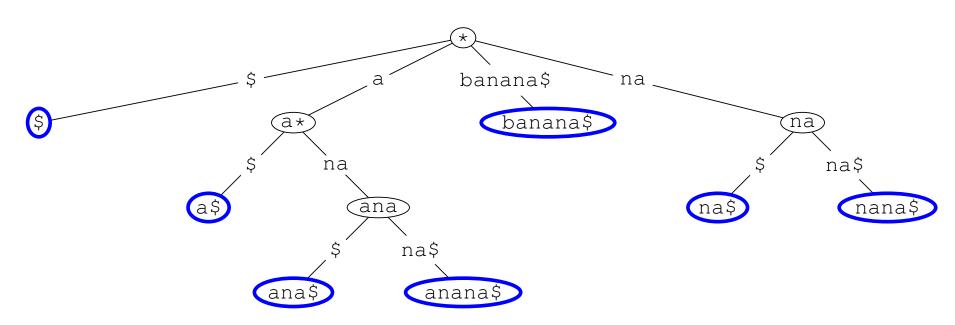
- 1. B-Trees
- 2. Tries
- 3. Suffix Tree



### **Suffix Tree**

- Suffix tree is a trie of all suffixes of a string
- E.g. banana

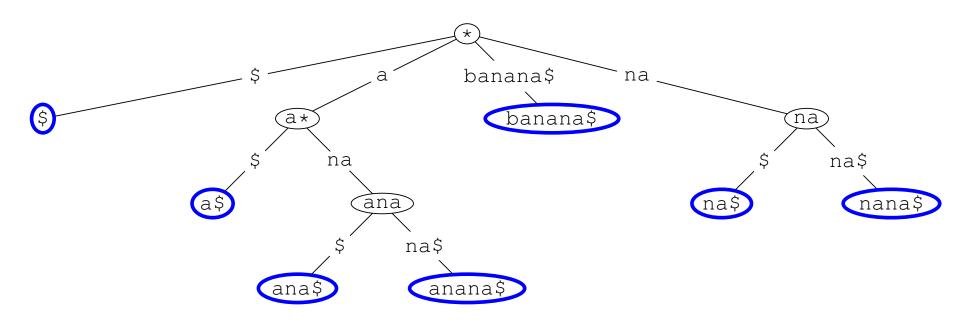
 $\{\$, a\$, na\$, ana\$, nana\$, anana\$, banana\$\}$ 



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- It really only got implemented after this
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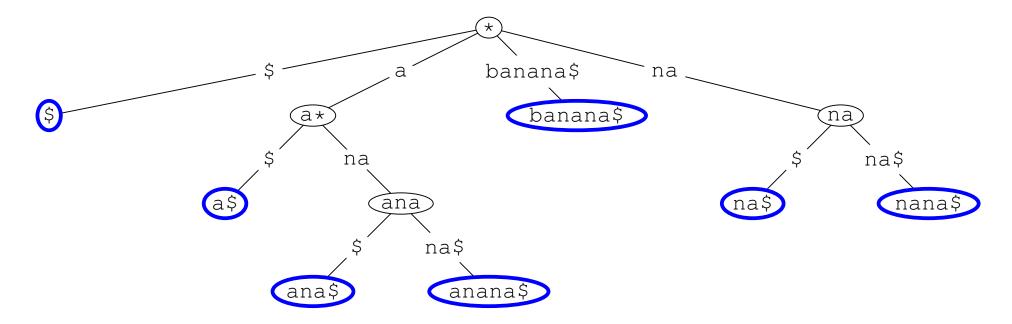
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# **String Matching**

• To find a match of a query string, Q, in a text, T, we can first construct the suffix tree of the string T we then simple look up the query, Q, using the trie

{\$, a\$, na\$, ana\$, nana\$, anana\$, banana\$}



- Using a regular trie for a suffix tree would typically use far too much memory to be useful
- However, by using pointers to the original text it is possible to build a suffix tree using O(n) memory where n is the length of the text
- Furthermore (and rather incredibly) there is a linear time (O(n)) algorithm to construct the trie
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- Exact word matching is in itself a very important application
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- But they can give the best performance—sometimes performance matters enough to make it worthwhile implementing multiway trees