Algorithms and Analysis

Lesson 31: *Understand Time Complexity*
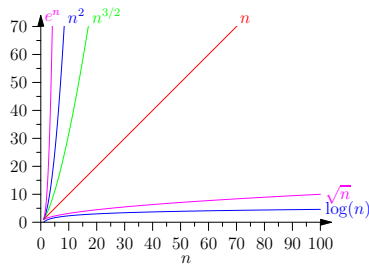
Theta, Big-O, little-o, Big-Omega, little-omega

Outline

1. **Time Complexity Classes**
   - Theta—Θ
   - Big O
   - Little o
   - Big Omega—Ω
   - Little omega—ω

2. Computing Time Complexity

Recap

- We have seen many algorithms taking times of order 1, $\log(n)$, n, $n\log(n)$, n^2, etc

- Sometimes these are worst time, average time or best time results

- We have lots of different notations, e.g. $O(1)$, $\Theta(\log(n))$, $\Omega(n^2)$, etc.

- What does it all mean?

Complexity Class Sets

- The correct way to think about complexity classes is in terms of sets

- Suppose we have an algorithm which takes an input of size n and computes an output in $f(n)$ operations

- E.g. $f(n) = 4n^2 + 2n + 3$

- We can partition all run times into sets by considering only the leading order term and ignoring the constant term

- We denote these sets by $\Theta(g(n))$

  ⋆ $4n^2 + 2n + 3 \in \Theta(n^2)$
  ⋆ $5n\log(n) + 3n + 2 \in \Theta(n\log(n))$

Defining $\Theta(g(n))$

- A function $f(n) \in \Theta(g(n))$ if

$$\lim_{n\to\infty} \frac{f(n)}{g(n)} = c \qquad\qquad 0 < c < \infty$$

- E.g.

$$\lim_{n\to\infty} \frac{4n^2 + 2n + 3}{n^2} = 4$$

$$\lim_{n\to\infty} \frac{5n\log(n) + 3n + 2}{n\log(n)} = 5$$

Ignoring the Constant

- Does an algorithm that uses $4n^2$ operations run faster than one that uses $7n^2$ operations?

- Answer depends on the operations which might depend on the programming language or the machine architecture

- However asymptotically (i.e. for sufficiently large n) an order $n\log(n)$ algorithm will always run faster than an order n^2 algorithm even when they are run on different machines

- The constant is important in practice (if there are two algorithms A and B that are both $n\log(n)$, but algorithm A runs twice as fast as algorithm B, which one should you use?)

- Nevertheless, ignoring the constant is often essential to make analysis of algorithms doable

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Ordering Complexity Classes

- We can define the relation $\Theta(f(n)) < \Theta(g(n))$ if

$$\lim_{n\to\infty} \frac{f(n)}{g(n)} = 0$$

- Informally if algorithm A has time complexity $\Theta(f(n))$ and algorithm B has time complexity $\Theta(g(n))$ then if $\Theta(f(n)) < \Theta(g(n))$ algorithm A is faster for sufficiently large n
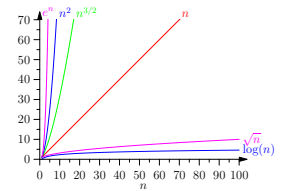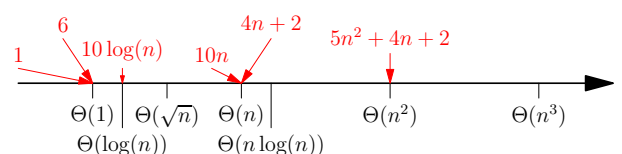
- The relation defines a **complete ordering**

The Complexity Line

- We can order all complexity classes. E.g.

$$\Theta(1) < \Theta(\log(n)) < \Theta(\sqrt{n}) < \Theta(n) < \Theta(n^2)$$

- We can depict this as a complexity line

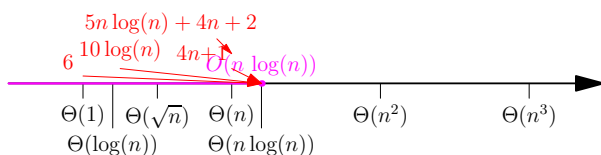- The line is dense (i.e. there are an uncountable infinity of complexity classes)

Complexity Dependent on Inputs

- The run time of many algorithms depends on the input

- In this case we can define different time complexities
  - ⋆ Worst case time complexity (the longest time an algorithm will take)
  - ⋆ Average complexity (the expected time averaged over all possible inputs)
  - ⋆ Best case time complexity (the shortest time an algorithm will take)—usually not very interesting

- Every algorithm will have a Θ complexity class for the worst, average and best time complexity

Unknown Time Complexity

- Algorithms are often rather complicated and knowing the exact time complexity (for either worst, average or best cases) might not be known

- In reality it will have some run time (e.g. $f(n) = 3n^2 \log(n) + 2n^2 - n + 3$) and will belong to a Θ time complexity set (e.g. $\Theta(n^2 \log(n))$) but we might not be able to calculate it

- However, we can usually bound the run times of algorithms
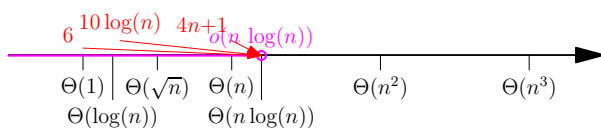
Big-O

- Big-O is an upper bound on the time complexity

- If an algorithm is $O(g(n))$ then its time complexity is no more than $\Theta(g(n))$

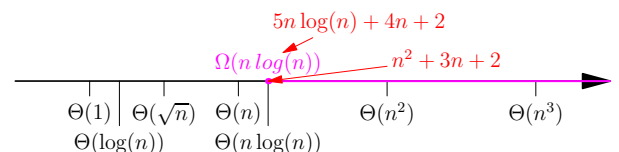- I.e. $\Theta(f(n)) \leq \Theta(g(n))$ implies $f(n) \in O(g(n))$

Upper Bounding Time Complexity

- Consider a program

```
// define stuff
for(int i=0; i<n; i++) {
  // do something
  if (/* some condition */) {
    for (int j=0; j<n; j++) {
      // do other stuff
    }
  }
}
// clean up
```

- If the `if` statements is never true this is a $\Theta(n)$ algorithm if it is always true it is a $\Theta(n^2)$ algorithm

- If we don't know then we can at least say that the run time is in $O(n^2)$—we assume the worst, but the worst may never happen

Little-o

- Sometimes we want to say that the time complexity for an algorithm $\Theta(f(n))$ is **strictly less** than a known time complexity $\Theta(g(n))$

- I.e. $\Theta(f(n)) < \Theta(g(n))$ implies $f(n) \in o(g(n))$

Lower Bounds—Ω

- It is often easy to obtain a lower bound on a particular algorithm although getting a tight general lower bound is often very difficult

- I.e. $\Theta(f(n)) \geq \Theta(g(n))$ implies $f(n) \in \Omega(g(n))$

Lower Bounding Time Complexity

- Returning to the program

```
// define stuff
for(int i=0; i<n; i++) {
  // do something
  if (/* some condition */) {
    for (int j=0; j<n; j++) {
      // do other stuff
    }
  }
}
// clean up
```

- We might not know how frequently the `if` statement is true, but we know in all cases the first `for` loop iterates over n

- Thus we know this algorithm is in $\Omega(n)$—we assume the best, but the best may never happen
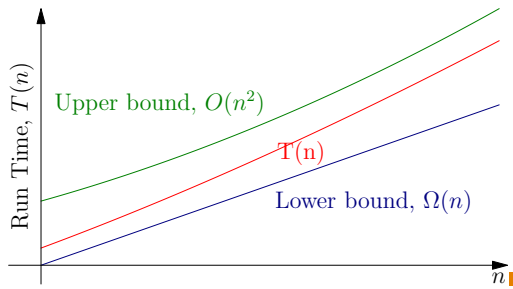
Little Omega—ω

- It is sometimes useful to talk about a strict lower bound

- I.e. $\Theta(f(n)) > \Theta(g(n))$ implies $f(n) \in \omega(g(n))$
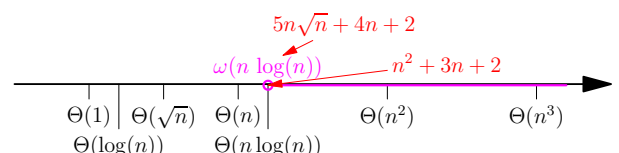
Bounding Run Time Complexity

- When we are given an algorithm to analyse we want to compute $\Theta(n)$

- This may be difficult, however, it is often easy to find bounds

Proving Asymptotic Time Complexity

- If we know an algorithm is
  - $T(n) \in O(f(n))$
  - $T(n) \in \Omega(f(n))$

- Then $T(n) \in \Theta(f(n))$

- This is a common proof strategy

Meaning of Time Complexity

- Insertion sort has time complexity $\Theta(n^2)$

- Because it consists of two `for` loops

- It takes 2 seconds to sort $100\,000$ items

- How long does it take to sort $1\,000\,000$ items?

- n increases by 10, time complexity increases by $10^2 = 100$

- Time taken is approximately 200 seconds or around 3.5 minutes

Exponential Time Complexity

- When we talk about exponential time complexity we usually mean that

$$\log(T(n)) \in \Theta(n)$$

- This is true if
  - $T(n) = 2^n$ $\log(T(n)) = n\log(2)$
  - $T(n) = 6.1\,\mathrm{e}^{0.003n}$ $\log(T(n)) = 0.03\,n + \log(6.1)$
  - $T(n) = n\,10^n$ $\log(T(n)) = n\log(10) + \log(n)$

- Note that none of these are in complexity class $\Theta(\mathrm{e}^n)$

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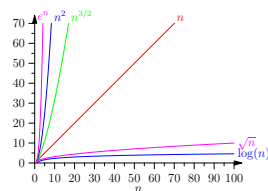
Counting For Loops

- How long does the following code take?

```
for(int i=0; i<n; i++)  {
   // prepare stuff
}
for(int i=0; i<n; i++)  {
   // do something
   for (int j=0; j<n; j++) {
      // do other stuff
   }
}
```

- The first `for` loop takes $\Theta(n)$ operations the second double `for` loop takes $\Theta(n^2)$

- Answer $\Theta(n^2)$

Recursion

- Determining time complexity is harder when we use recursion

- Consider Euclid's algorithm for determining the greatest common divisor

```
long gcd(long m, long n)
{
   while(n!=0) {
      long rem = m%n;
      m = n;
      n = rem;
   }
   return m;
}
```

```
long gcd(long m, long n)
{
   if (n==0)
      return m;
   else
      return gcd(n, m%n);
}
```

- This doesn't even look like a recursion

Example of gcd

- Example of Euclid's algorithm `gcd(1989,1590)`

- Sequence of remainders is 399, 393, 6, 3, 0

- The greatest common divisor is 3

- How long does it take compute `gcd(n,m)` with $n > m$

- This is subtle as could depend in a complex way on the pair n and m
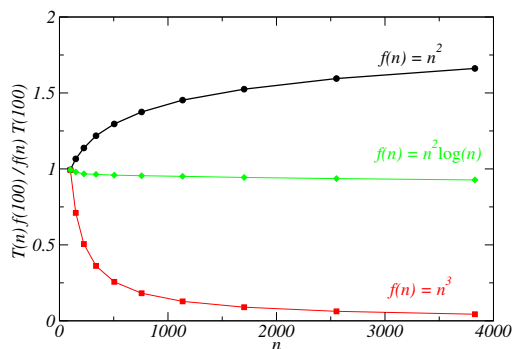
Recursive Formulae

- An observation which makes the analysis relatively simple is that the remainder is reduced by at least 2 after two iterations∎

- To prove
  - ⋆ Using the recursion (assuming $m, n < 0$)

$$\gcd(m, n) = \gcd(n, \text{rem}(m, n)) = \gcd(\text{rem}(m, n), \text{rem}(n, \text{rem}(m, n)))$$

  - ⋆ The proof follows by showing that $\text{rem}(n, \text{rem}(m, n)) < n/2$∎

- Thus $T(n) < T(n/2) + 2$∎

Solving Recursions

- To show that $T(n) \in O(\log(n))$ we observe
  - ⋆ Note that $T(1) = 1$∎

$$T(n) < T(2^{-1}n) + 2 < T(2^{-2}n) + 4 < \cdots < T(2^{-t}n) + 2t$$∎

  - ⋆ Choose $t = \lceil \log_2(n) \rceil$
  - ⋆ then $2^{-t}n = 2^{-\lceil \log_2(n) \rceil}n \leq 2^{-\log_2(n)}n = \frac{n}{n} = 1$∎
  - ⋆ Thus $T(2^{-t}n) < T(1) = 1$∎
  - ⋆ $T(n) < 1 + 2t = 1 + 2\lceil \log_2(n) \rceil \in O(\log(n))$∎

- A huge calculation shows the the average number of iterations is about $(12 \log(2) \log(n))/\pi^2 + 1.47$∎

Probability of Relative Primes

- Consider the following program to compute the probability of relative primes for all numbers up to n

```java
public static double probRelPrime(n)
{
  int rel=0, tot=0;
  for(int i=1; i<=n; i++)
    for(int j=i+1; j<=n; j++) {
      tot++;
      if (gcd(i,j)==1)
        rel++;
    }
  return (double) rel / tot;
}
```

- What is the time complexity?∎

Time Complexity

- Program involves two nested loops of size $O(n)$∎

- Then we need to calculate `gcd(i, j)` at each iteration∎

- Time complexity is $n \times n \times \log(n) = n^2 \log(n)$∎

- How could we provide empirical support for this calculation?∎

Testing Hypothesis

- We can test our hypothesis by scaling the run time by the complexity

Conclusions

- You should understand the difference between Θ, O, o, Ω and ω∎

- You need to be able to compute time complexity by loop counting∎

- To compute time complexity for recursive functions you need to be able to obtain recurrence equations∎

- You should be able to solve simple recurrence equations and sum up simple series∎

- You should be able to prove more complicated results using proof by induction∎

- Thank you for attending the course∎