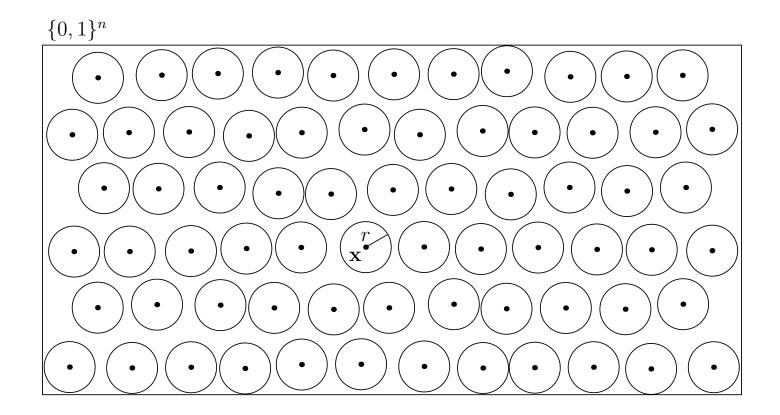
Advanced Machine Learning

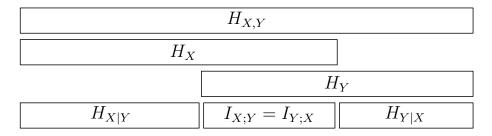
Information Theory



Information, KL-divergence, Minimum Description Length

Outline

- 1. Information Theory
- 2. KL-Divergence
- 3. Minimum Description Length
- 4. Variational Auto-Encoders



Communicating Via a Noisy Channel

Information theory considers communicating down a (noisy) channel

$$X \sim \mathbb{P}(X) \xrightarrow{\text{noisy channel}} Y \sim \mathbb{P}(Y \mid X)$$

- We send a message X (with probability $\mathbb{P}(X)$) and receive a message Y with probability $\mathbb{P}(Y\mid X)$
- ullet The uncertainty of the message sent, given we received a message y is

$$H_{X|Y=y} = -\sum_{x \in \mathcal{X}} \mathbb{P}(X = x \mid Y = y) \log(\mathbb{P}(X = x \mid Y = y))$$

• The expected uncertainty in the message sent is

$$H_{X|Y} = \sum_{y \in \mathcal{Y}} \mathbb{P}(Y = y) H_{X|Y=y} = -\sum_{x,y} \mathbb{P}(X = x, Y = y) \log(\mathbb{P}(X = x \mid Y = y))$$

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We can define the joint entropy

$$H_{X,Y} = -\sum_{x,y} P_{X,Y}(x,y) \log(P_{X,Y}(x,y))$$

- If the message we receive is independent of the message that is sent then $H_{X,Y} = H_X + H_Y$ (we saw this in the last lecture)
- $H_{X,Y} \neq H_X + H_Y$ if X and Y are correlated
- Since $\mathbb{P}(X,Y) = \mathbb{P}(Y|X)\mathbb{P}(X) = \mathbb{P}(X|Y)\mathbb{P}(Y)$ if follows

$$H_{X,Y} = H_X + H_{Y|X} = H_Y + H_{X|Y}$$

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Mutual Information

- The amount of uncertainty about the message being sent, X, before receiving the message is $H_X = -\mathbb{E}_X[\log \mathbb{P}(X)]$
- Shannon define the $mutual\ information$ to be the expected loss in uncertainty when we receive a message

$$I_{X;Y} = H_X - H_{X|Y}$$

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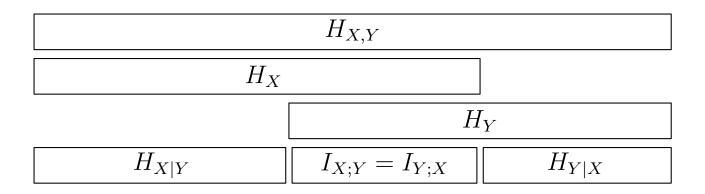
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Channel Capacity

We can summarise these relationships diagrammatically



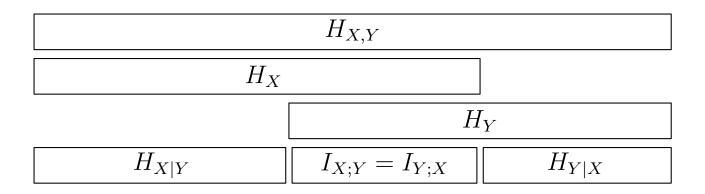
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Independent Noise

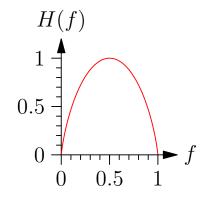
ullet The simplest model of a noisy channel is a binary channel where each symbol is corrupted independently with a probability f

$$\mathbb{P}(X = 1 | Y = 0) = \mathbb{P}(X = 0 | Y = 1) = f$$



$$H_{X_i|Y_i} = -(1-f)\log(1-f) - f\log(f) = H(f)$$





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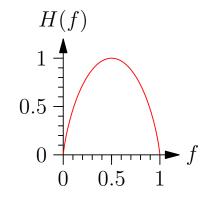
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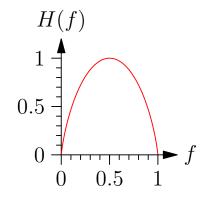
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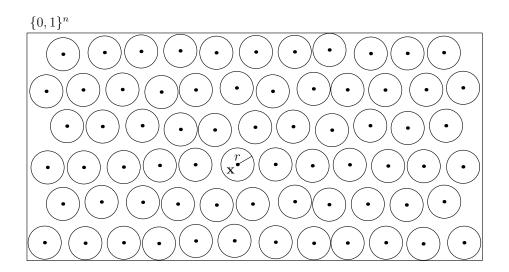
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Error Correcting Codes

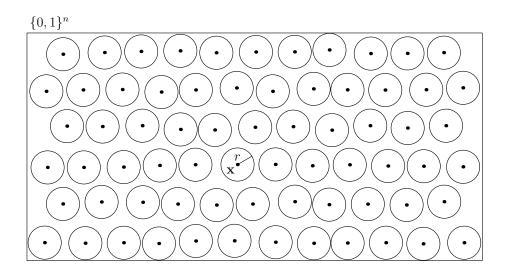
- To reduce the chance of misinterpreting a message we need to build an error correcting code
- We can do this dividing the space of binary messages into a set of Hamming balls



• A Hamming ball B(x,r) is the set of strings that differ from n-dimensional binary string, x, by at most r digits

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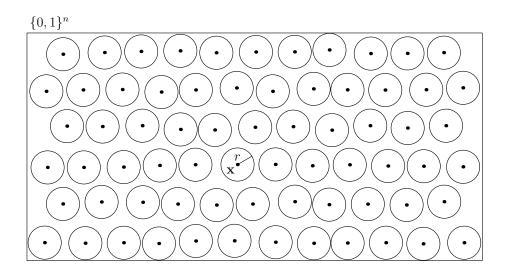
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- For sufficiently large n we would expect all errors are smaller than $(f+\epsilon)n$ (for $\epsilon>0$)
- If we make the radius of the Hamming ball $r = (f + \epsilon)n$ ($\epsilon > 0$) then we would expect no error for sufficiently large n
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 m bits}$ so we need codes of just over twice as long to communicate accurately over a noisy channel with a 10% corruption rate
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 - ★ Wikipedia mentions 14 applications
- Suppose we want to align two sets of images through some non-linear transformations
- One way of doing this is to choose the non-linear transformations that maximise the mutual information (or normalised mutual information) between the two sets of images

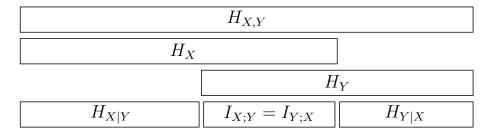
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KL-Divergence

We have met the Kullback-Leibler divergence

$$KL(p||q) = \mathbb{E}_{X \sim p(X)} \left[\log \left(\frac{p(X)}{q(X)} \right) \right]$$
$$= -\mathbb{E}_{X \sim p(X)} [\log(q(X))] - H_X$$

- Recall $-\log(q(X=x))$ is the length of code need to send a message x with a probability q(X=x)
- Thus $-\mathbb{E}_{X\sim p(X)}[\log(q(X))]$ is the expected length of message needed to code $X\sim p(X)$ using the optimal code for the distribution q(X) that than p(X)
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Variational Approximation

- Recall we use MCMC in Bayesian inference because the posterior distribution is too complicated to write down in closed form
- In the variational approximation we approximate the posterior distribution by a simpler (typically factored distribution), e.g.

$$f(\boldsymbol{\theta} \mid \mathcal{D}) \approx g(\boldsymbol{\theta} \mid \boldsymbol{\phi}) = \prod_{i} g(\theta_i \mid \phi_i)$$

 The standard method for solving this is to maximise the variational free energy

$$\Phi(\phi) = -\int g(\boldsymbol{\theta} \mid \phi) \log \left(\frac{g(\boldsymbol{\theta} \mid \phi)}{f(\boldsymbol{\theta}, \mathcal{D})} \right) d\boldsymbol{\theta}$$

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- If we maximise $\Phi(\phi)$, we end up minimising the KL divergence between g and f so that $g \approx f$ and $\Phi(\phi) \approx \log(f(\mathbf{D}))$
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Put Another Way

• We can rewrite the variational free energy as $\Phi(\phi) = L_q(\phi) + H_q(\phi)$ where

$$L_q(\boldsymbol{\phi}) = \int g(\boldsymbol{\theta} \mid \boldsymbol{\phi}) \left(\log(f(\mathcal{D}|\boldsymbol{\theta})) + \log(f(\boldsymbol{\theta})) \right) d\boldsymbol{\phi}$$

acts like an expected posterior term that is maximised when the data is well modelled (we put the probability density, $g(\theta \mid \phi)$ where the $f(\theta, \mathcal{D})$ is large)

The second term is an entropy

$$H_q(\boldsymbol{\phi}) = -\int g(\boldsymbol{\theta} \mid \boldsymbol{\phi}) \log(g(\boldsymbol{\theta} \mid \boldsymbol{\phi})) d\boldsymbol{\phi}$$

That is, we maximise the uncertainty of the distribution $g(\theta \mid \phi)$

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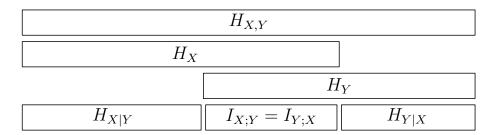
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Outline

- 1. Information Theory
- 2. KL-Divergence
- 3. Minimum Description Length
- 4. Variational Auto-Encoders



- Outside of the Bayesian framework it is difficult to do model selection
- When is it better to accept a more complex model for a better fit and when are we just over-fitting?
- Usually we answer this using a validation set, but this is not always possible
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- Suppose Alice has data $\mathcal{D}=\{(\boldsymbol{x}_i,y_i)\mid i=1,2,...,m\}$ while Bob has only the feature vectors $\{\boldsymbol{x}_i\mid i=1,2,...,m\}$
- Alice wants to communicate y_i to Bob as efficiently as possible
- ullet We suppose Alice & Bob have available a model $\hat{f}(oldsymbol{x}|oldsymbol{ heta})$
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$$\delta_i = y_i - \hat{f}(\boldsymbol{x}_i|\boldsymbol{\theta})$$

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• The **description length** for $\{y_i \mid i=1,2,...,m\}$ is then the cost of transmitting θ plus the cost of transmitting the errors

$$L = \sum_{k=1}^{n} \ell(\theta_k) - \sum_{i=1}^{m} \left(\log \left(p_{\delta} \left(y_i - \hat{f}(\boldsymbol{x}_i | \boldsymbol{\theta}) \right) \right) + \log(\Delta) \right)$$

where $\ell(\theta_k)$ is the number of bits need to communicate θ_k (we get to choose the accuracy if is worth encoding the parameters)

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- Note that the accuracy Δ will lead to the same cost, $-m\log(\Delta)$, for all models so doesn't affect which model is selected

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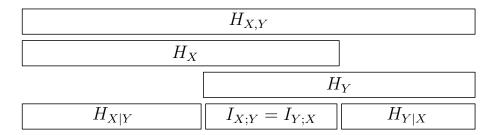
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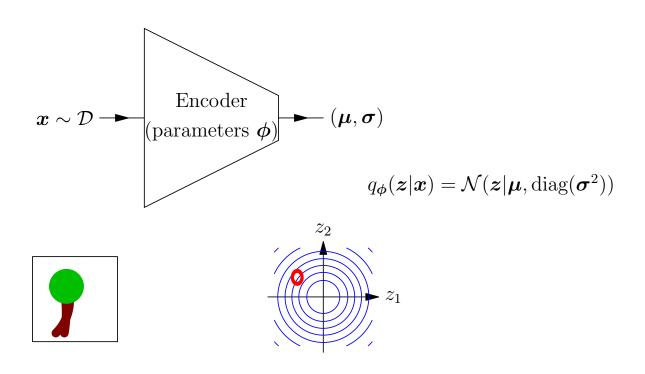
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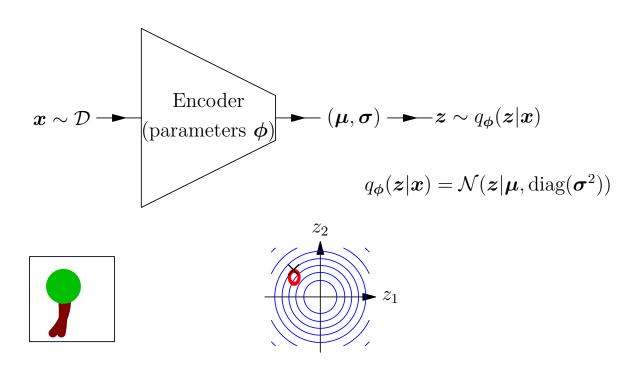
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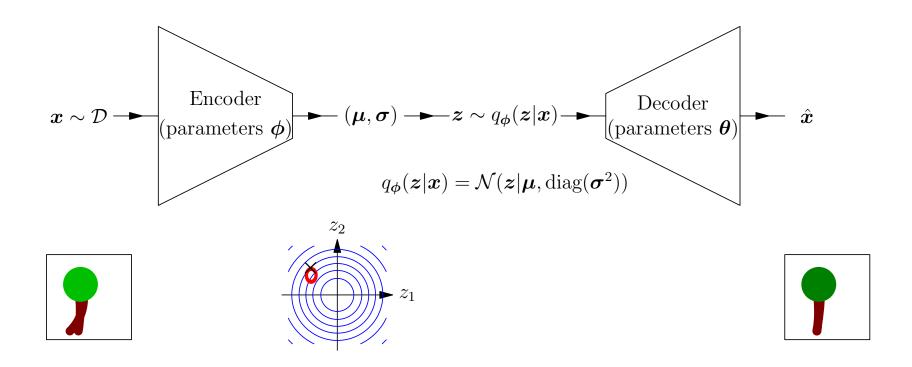


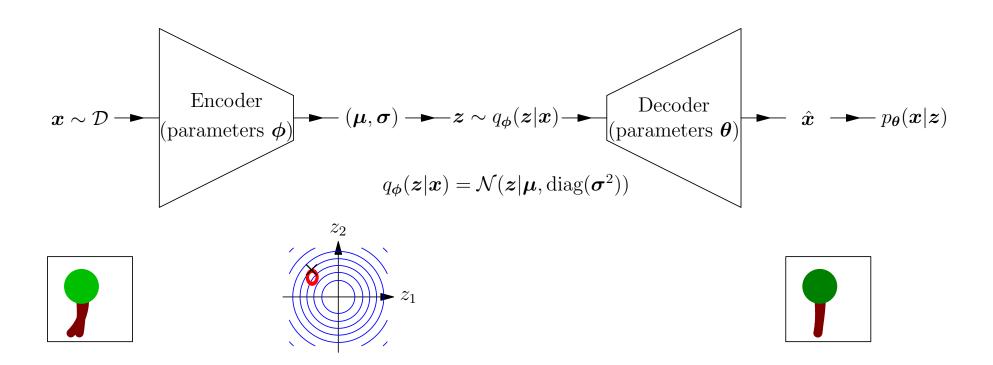
 $oldsymbol{x} \sim \mathcal{D}$

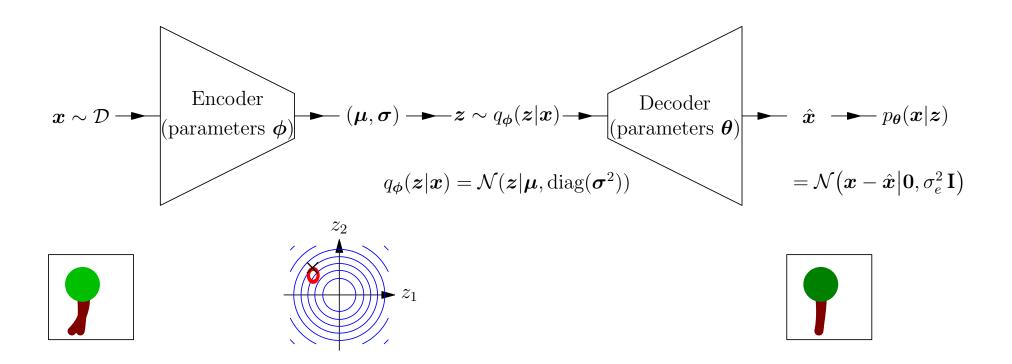


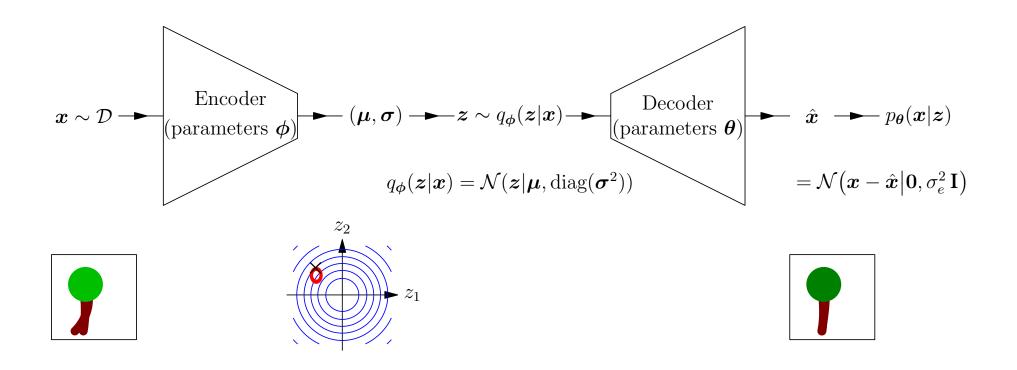




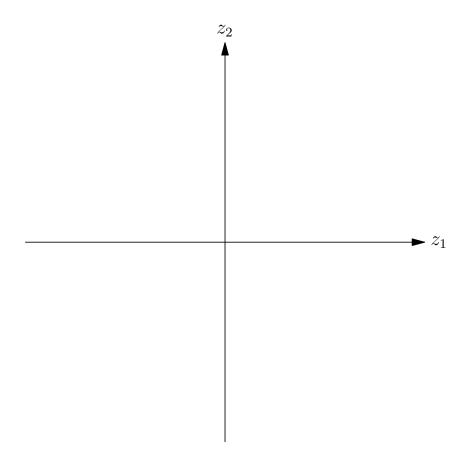


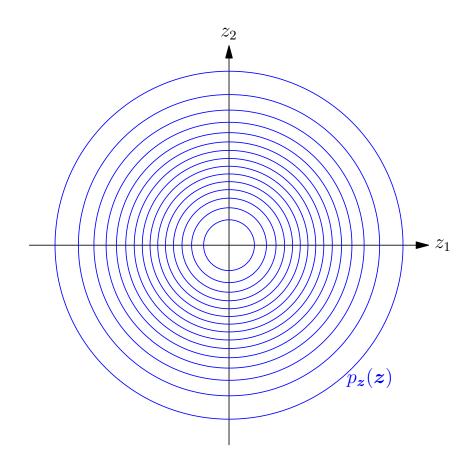


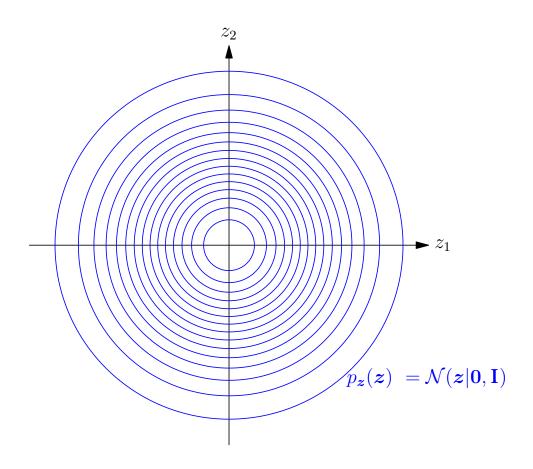


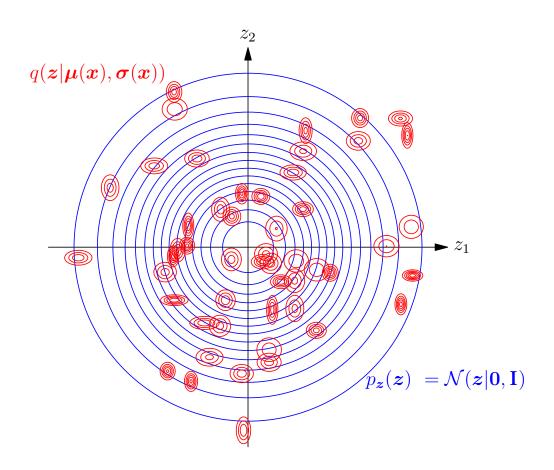


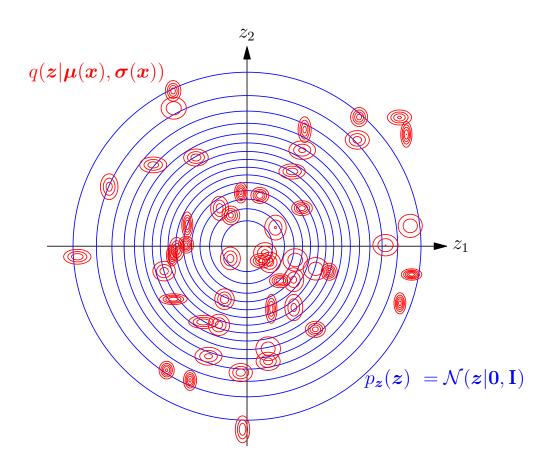
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$$\mathrm{KL} ig(q_{m{ heta}}(m{z}|m{x}) ig| \mathcal{N}(m{0},m{I}) ig)$$

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- It has a very natural explanation in terms of minimum description length
- Alice wants to communicate the images to Bob
- Alice uses the encoder to derive a (latent) code $q(\boldsymbol{z}|\boldsymbol{x})$ which she communicates to Bob
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- \star The cost of communicating the code $\mathrm{KL}ig(q_{m{ heta}}(m{z}|m{x})ig\|\mathcal{N}(\mathbf{0},\mathbf{I})ig)$
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