Advanced Algorithms Plagiarism Detection using an dynamic LCS algorithm

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Introduction

- Pre-processing
- Summary of Algorithms Implemented
 - LCS
 - Printing Neatly
- Algorithms
 - Branch and Bound
- Complexity
 - Longest common sub-sequence
 - Printing Neatly
- Plagiarism Detection techniques
- Outliers
- Demo of UI
- Project planning and Execution
- Lessons Learned and Challenges
- Conclusions

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Pre-processing

- Four different techniques :
 - Light Pre-processing
 - Advanced Pre-processing
 - Stop Word Removal
 - Word Ordering

Light Pre-processing

- A basic pre-processing technique which simply removes or replaces given symbol such as:
 - Directional quotes -> Straight quotes
 - Extra newline characters are removed
 - Extra spaces are removed

Advanced Pre-processing

- A more advanced technique which remove all non-ASCII characters, and all symbols that are not periods or comma.
 It makes everything in lower case.
- This pre-processing technique allow the LCS to be more robust to changes in punctuation.

Advanced Pre-processing

Text before the pre-process:

 This kind of relationship can be visualised as a tree structure, where 'student' would be the more general root node and both 'postgraduate' and 'undergraduate' would be more specialised extensions of the 'student' node (or the child nodes).

Text after:

 this kind of relationship can be visualised as a tree structure, where student would be the more general root node and both postgraduate and undergraduate would be more specialised extensions of the student node or the child nodes.

Stop word removal

Text before the pre-process:

 this kind of relationship can be visualised as a tree structure, where student would be the more general root node and both postgraduate and undergraduate would be more specialised extensions of the student node or the child nodes.

Text after:

 kind relationship visualised tree structure, student would general root node postgraduate undergraduate specialised extensions student node child nodes

Word Ordering

- Word Ordering is a technique that allow us to see if a text has been copied but rearranged. In order to see that we order every sentence of the text.
- The technique reordered words in alphabetical order.
 However, we did not find this technique to be an improvement on the plagiarism score, so we did not include it in the final plagiarism detector.

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LCS Forward

LCS Forward on « ABBA » and « ABCD » -> 2

	A	В	В	A
A	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

	A	В	В	A
Α	D	L	L	L
В	U	D	L	L
С	U	D	U	U
D	U	U	U	U

Time Complexity : O(m*n)

Space Complexity : O((m*n)*2)

LCS Backward

LCS Forward on « ABBA » and « ABCD » -> 2

	A	В	В	A
A	2	1	0	0
В	1	1	0	0
С	1	1	0	0
D	0	0	0	0

	A	В	В	A
Α	U	L	L	L
В	L	U	L	L
С	L	U	L	L
D	L	L	L	D

Time Complexity : O(m*n)

Space Complexity : O((m*n) *2)

LCS Forward on « ABBA » and « ABCD » -> 2

Dynamic

	A	В	В	A
Α	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

Space Efficient

	A	В	В	A
-	0	0	0	0
A	1	1	1	1

LCS Forward on « ABBA » and « ABCD » -> 2

Dynamic

	A	В	В	A
Α	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

Space Efficient

	A	В	В	A
A	1	1	1	1
В	1	2	2	2

LCS Forward on « ABBA » and « ABCD » -> 2

Dynamic

	A	В	В	A
Α	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

Space Efficient

	A	В	В	Α
В	1	2	2	2
С	1	2	2	2

LCS Forward on « ABBA » and « ABCD » -> 2

Dynamic

	A	В	В	A
Α	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

Space Efficient

	A	В	В	Α
С	1	2	2	2
D	1	2	2	2

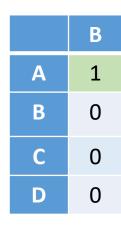
Space Efficient Forwards

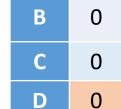
	A	В	В	-
-	0	0	0	-
A	1	1	1	-

Space Efficient Backwards

	-	-	В	A
D	-	-	0	0
-	-	-	0	0

Array for result of column B





	В
A	1
В	2
С	2
D	2

Time Complexity : O(m*n)

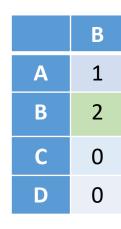
Space Complexity: O(2n)

Space Efficient Forwards

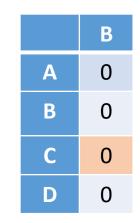
	A	В	В	-
A	1	1	1	-
В	1	2	2	-

Space Efficient Backwards

	-	-	В	A
С	-	-	0	0
D	-	-	0	0



Array for result of column B



Sum of forwards + Backwards

	В
Α	1
В	2
С	2
D	2

Space Efficient Forwards

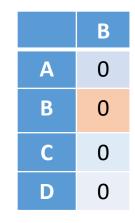
	A	В	В	-
В	1	2	2	-
С	1	2	2	-

Space Efficient Backwards

	-	-	В	A
В	-	-	0	0
С	-	-	0	0

	В
Α	1
В	2
С	2
D	0

Array for result of column B



Sum of forwards + Backwards

	В
A	1
В	2
С	2
D	2

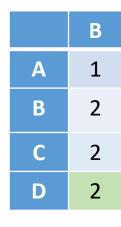
Time Complexity : O(m*n)O(2n) Space Complexity:

Space Efficient Forwards

	A	В	В	A
С	1	2	2	-
D	1	2	2	-

Space Efficient Backwards

	-	-	В	A
Α	-	-	0	0
В	-	-	0	0



Array for result of column B

	В
A	0
В	0
С	0
D	0

Sum of forwards + Backwards

	В
A	1
В	2
С	2
D	2

	Α	В	В	A
A	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

Time Complexity : O(m*n)

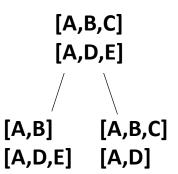
Space Complexity: O(2n)

	Α	В	В	A
Α	1	1	1	1
В	1	2	2	2
С	1	2	2	2
D	1	2	2	2

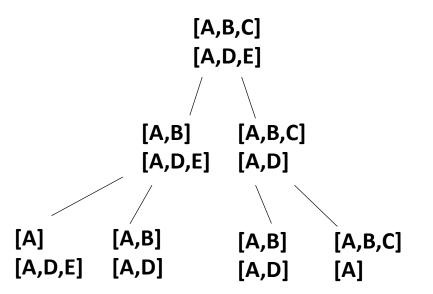
Time Complexity : O(m*n)

Space Complexity : O(2n)

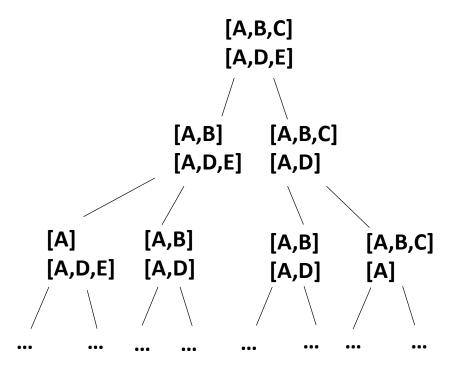
 Possible solutions are explored recursively, with the best results being taken.



 Possible solutions are explored recursively, with the best results being taken.

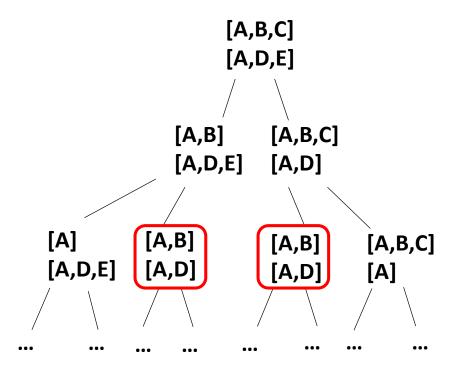


 Possible solutions are explored recursively, with the best results being taken.



 Possible solutions are explored recursively, with the best results being taken.

 Duplicate items in tree demonstrate why dynamic approach is suitable for this problem



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Branch and Bound: Bounds

- Upper Bound
 - Computed as the sum of the common letters between two sequences.
- Lower Bound
 - Computed as the maximum number of common letters between the two sequences.
- E.g. For Letters [A,B,C,A] & [B,C,A,A]
 - Common letters are A:2, B:1 & C:1
 - The Upper bound = 2+1+1 = 4
 - The Lower bound = 2

Branch and Bound: Implementation

- Implemented with a LIFO queue
- The last element added to the queue is the branch with the largest bound (the most potential)
- Items are only added to the queue if their upper bound is evaluated to be higher than the current best
- Elements are popped off the queue, if their upper bound is greater than the current best, they are explored further.

Branch and Bound: Index and Tracker

Index
ABCA
BCAA
Tracker

Branch and Bound: Index and Tracker

Index

ABC A

BCAA

Tracker

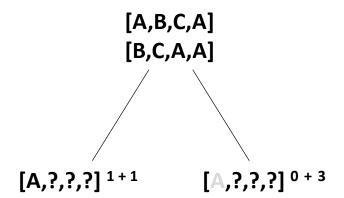
Branch and Bound: Index and Tracker

Index
A B C A
B C A A
Tracker

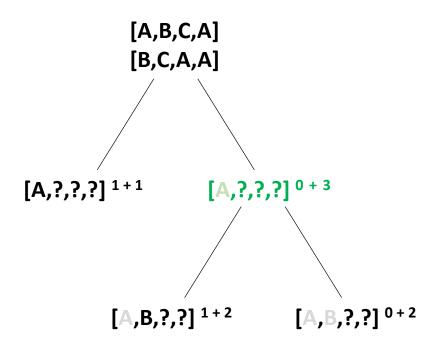
[A,B,C,A] [B,C,A,A]

(Solution, bound, index, tracker)

LIFO QUEUE (represented as a stack)

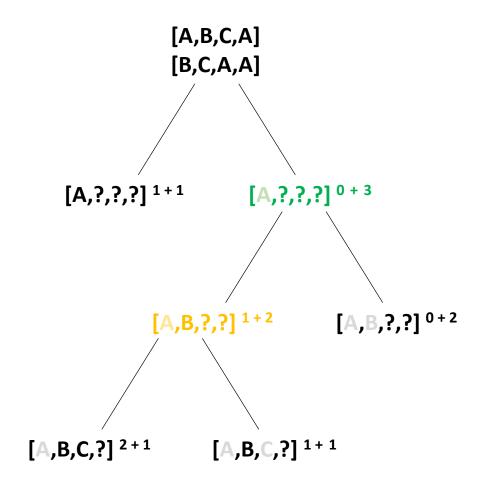


LIFO QUEUE (represented as a stack)



```
([A,B,?,?], 3, 2, 1)
([A,B,?,?], 2, 2, 0)
([A,?,?,?], 3, 1, 0)
([A,?,?,?], 2, 1, 3)
(Solution, bound, index, tracker)
```

LIFO QUEUE (represented as a stack)



```
([A,B,C,?], 3, 3, 2)

([A,B,C,?], 2, 3, 2)

<del>([A,B,?,?], 3, 2, 1)</del>

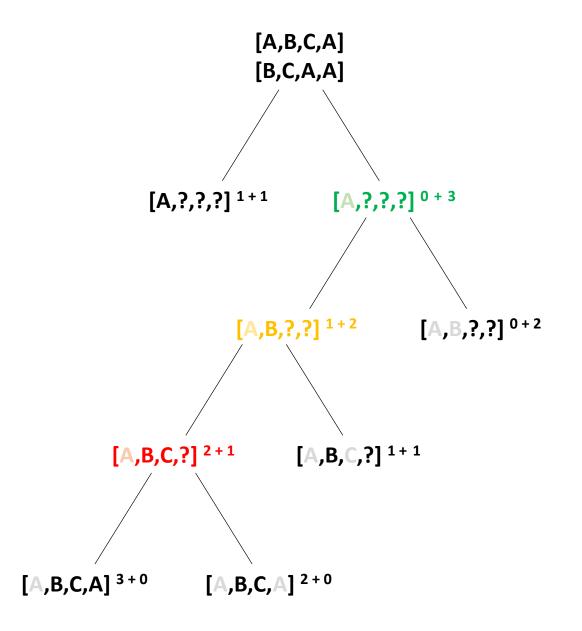
([A,B,?,?], 2, 2, 0)

<del>([A,?,?,?], 3, 1, 0)</del>

([A,?,?,?], 2, 1, 3)

(Solution, bound, index, tracker)
```

LIFO QUEUE (represented as a stack)



```
([A,B,C,A], 3, 4, 3)

([A,B,C,A], 2, 4, 3)

([A,B,C,?], 3, 3, 2)

([A,B,C,?], 2, 3, 2)

([A,B,?,?], 3, 2, 1)

([A,B,?,?], 3, 2, 1)

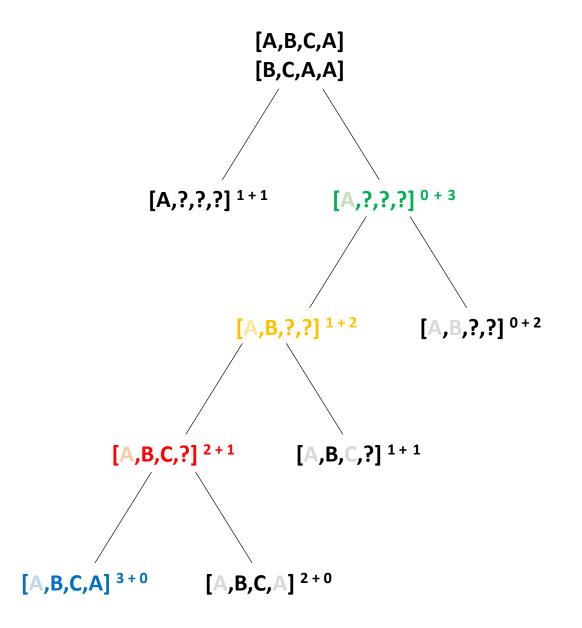
([A,B,?,?], 3, 1, 0)

([A,?,?,?], 3, 1, 0)

([A,?,?,?], 2, 1, 3)

(Solution, bound, index, tracker)
```

LIFO QUEUE (represented as a stack)



```
([A,B,C,A],3,4,3)

([A,B,C,A],2,4,3)

([A,B,C,?],3,3,2)

([A,B,C,?],2,3,2)

([A,B,C,?],3,2,1)

([A,B,?,?],3,2,1)

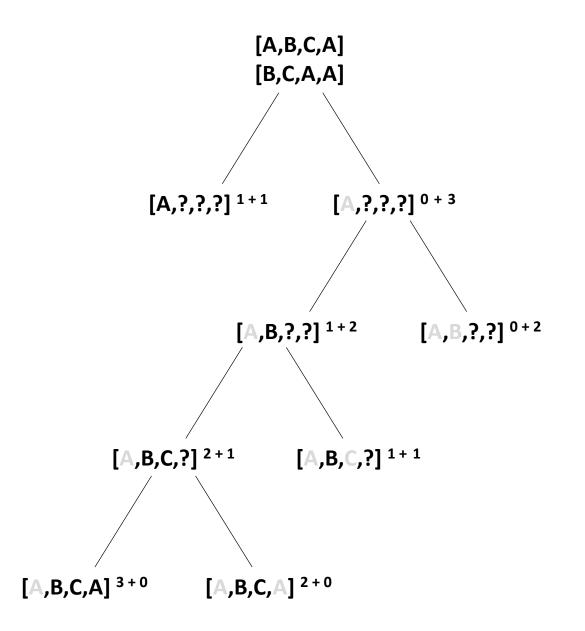
([A,B,?,?],2,2,0)

([A,?,?,?],3,1,0)

([A,?,?,?],2,1,3)

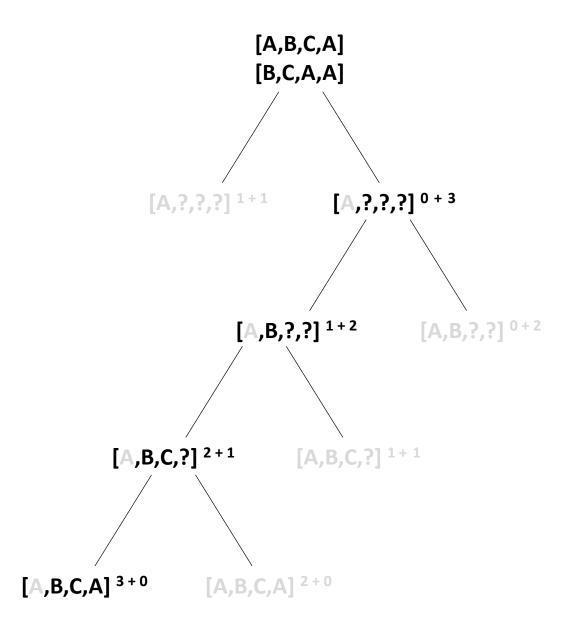
(Solution, bound, index, tracker)
```

LIFO QUEUE (represented as a stack)



([A,B,C,A], 3, 4, 3) ([A,B,C,A], 2, 4, 3) ([A,B,C,?], 3, 3, 2) ([A,B,C,?], 2, 3, 2) ([A,B,?,?], 3, 2, 1) ([A,B,?,?], 2, 2, 0) ([A,?,?,?], 3, 1, 0) ([A,?,?,?], 2, 1, 3) (Solution, bound, index, tracker)

LIFO QUEUE (represented as a stack)

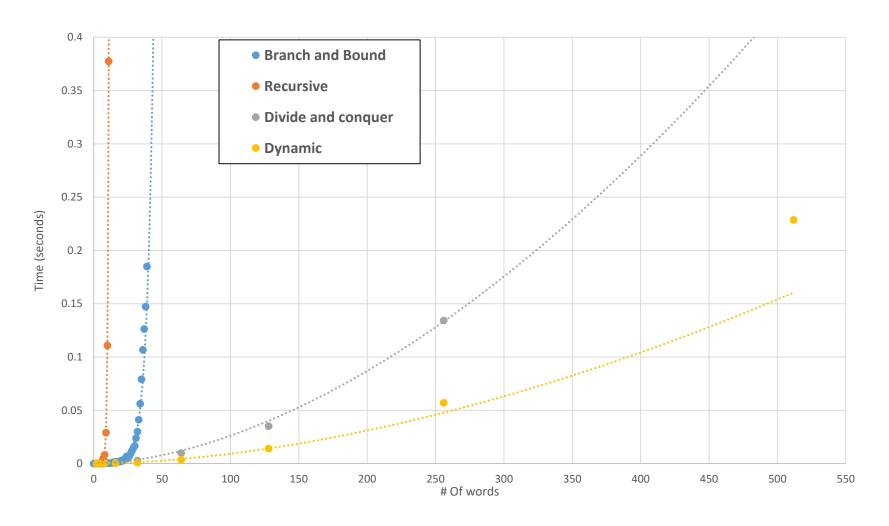


LIFO QUEUE (represented as a stack)

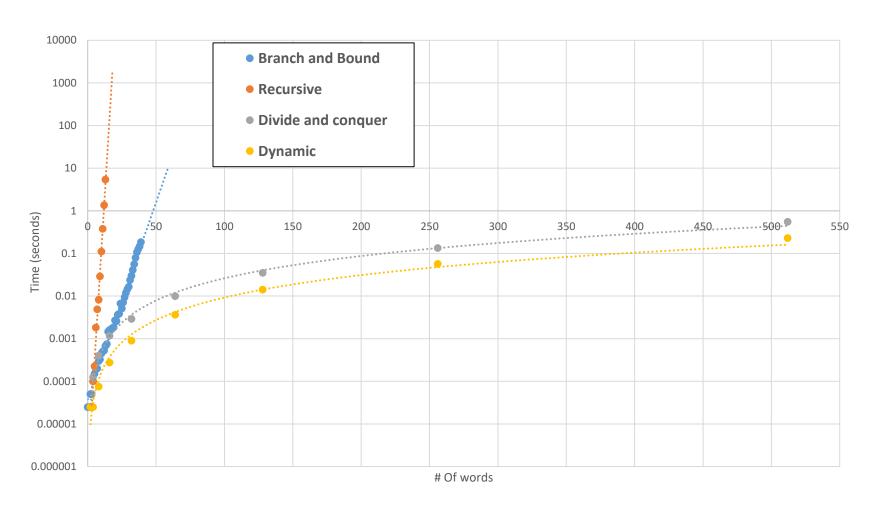
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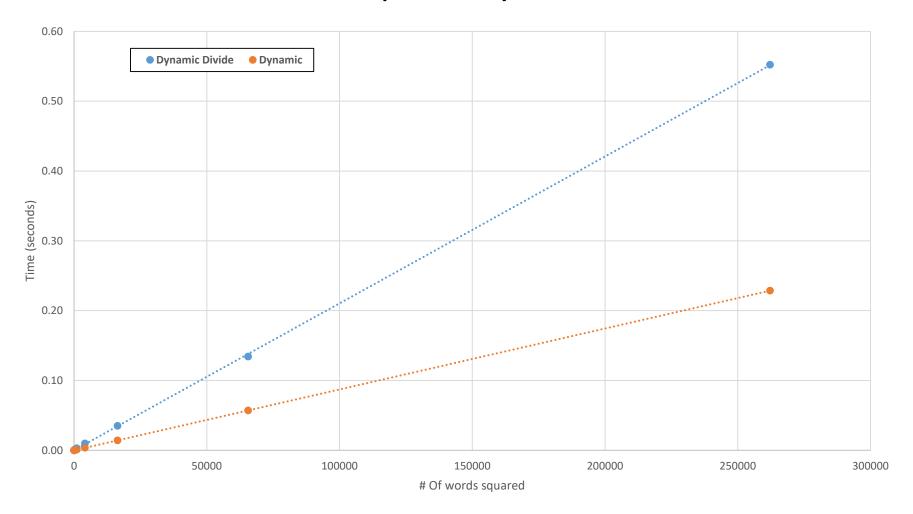
LCS: Time complexity



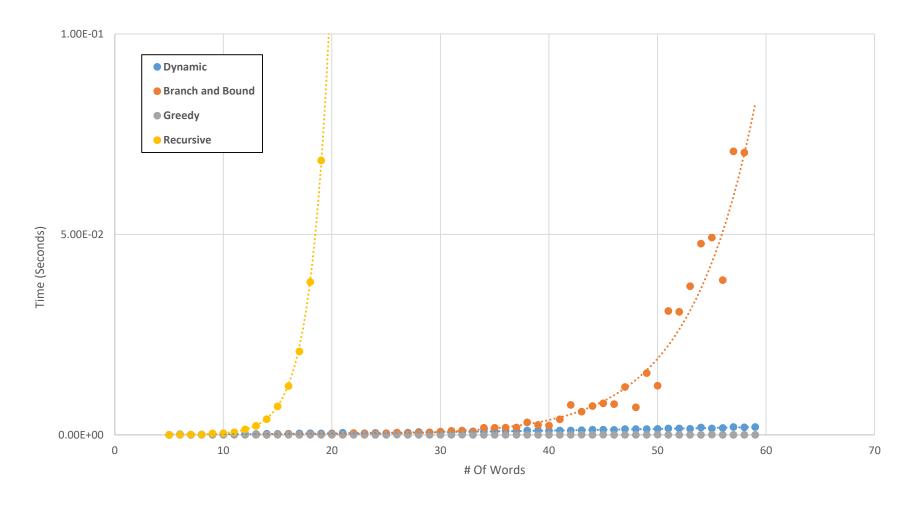
LCS: Time complexity



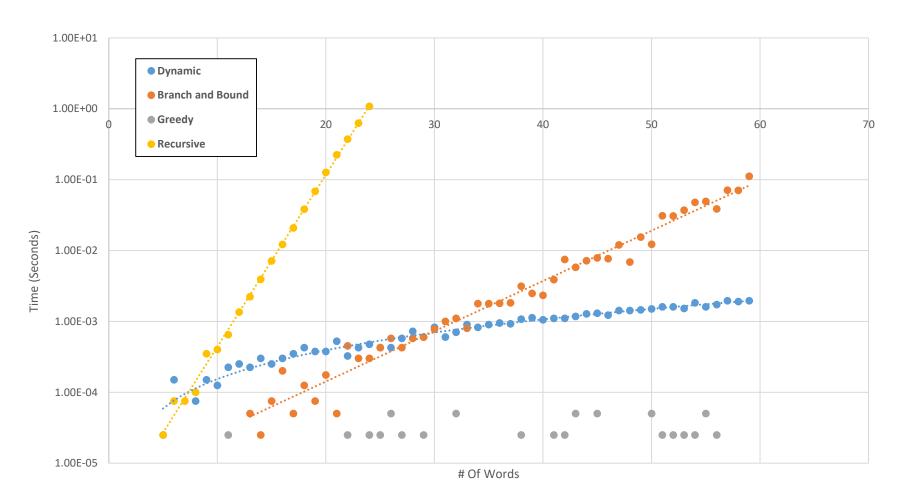
LCS: Time complexity



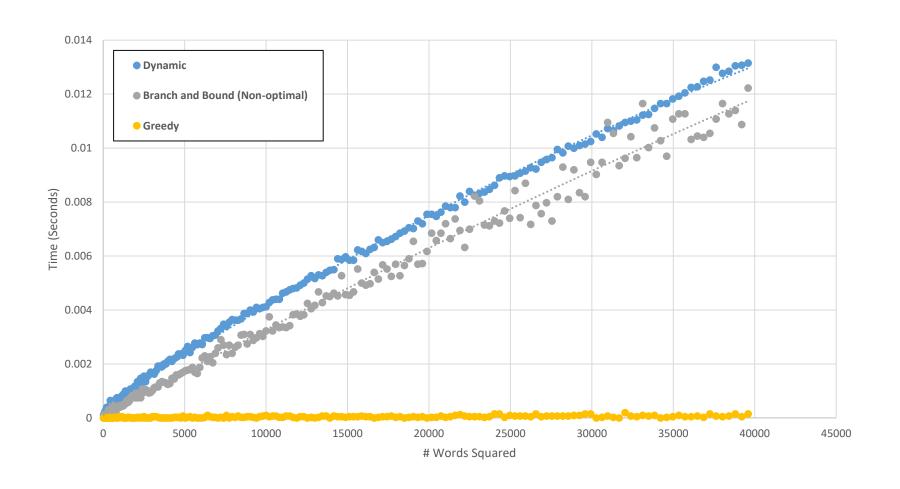
Printing Neatly: Time complexity



Printing Neatly: Time complexity



Printing Neatly: Time complexity



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Plagiarism Detection Techniques

- Several methods were explored:
 - Comparing ratio of # words in LCS with # words in corpus example
 - Generating the ratio above without stop words
 - Looking at number of plagiarised sentences in the body of text
 - Looking at the grouping of words and weighting large groups of words higher than small groups of words.

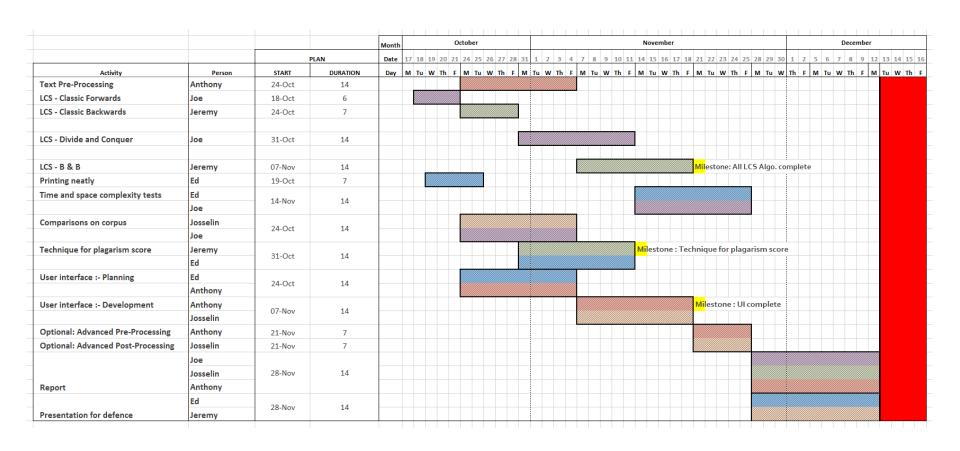
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Project Planning

- At the start of the project, a project plan was made and all team members agreed.
- The following are the changes made to the schedule:
 - The original plan did not include LCS linear space, LCS recursive or printing neatly branch and bound
 - It was noted that the implementation of the linear space algorithm was taking longer than expected, so another team-member was moved onto the LCS branch & bound algorithm.
 - The complexity tests and final technique for plagiarism took longer than expected to finalise.

Project Planning: Plan



Project Planning: Plan

