

Task 1: Deriving the Private Key

Here we can see the specified p, q, and e, followed by the d they are meant to generate using the extended euclidean algorithm.

```
[04/30/19]seed@VM:~/.../hw3$ gcc bignum2.c -lcrypto
[04/30/19]seed@VM:~/.../hw3$ ./a.out
p = F7E75FDC469067FFDC4E847C51F452DF
q = E85CED54AF57E53E092113E62F436F4F
n = E103ABD94892E3E74AFD724BF28E78366D9676BCCC70118BD0AA1968DBB143D1
e = 0D88C3
phi = E103ABD94892E3E74AFD724BF28E78348D52298BD687C44DEB3A81065A7981A4
d = 3587A24598E5F2A21DB007D89D18CC50ABA5075BA19A33890FE7C28A9B496AEB
Public Key: e=0D88C3, n=E103ABD94892E3E74AFD724BF28E78366D9676BCCC70118BD0AA1968DBB143D1
Private Key: d = 3587A24598E5F2A21DB007D89D18CC50ABA5075BA19A33890FE7C28A9B496AEB
```

Task 2: Encrypting a Message

Here the plaintext m = "4120746f702073656372657421" is encrypted with the supplied n and e values, then successfully decrypted with d to produce an identical m.

```
[05/01/19]seed@VM:~/.../hw3$ ./a.out
Public Key: e = 010001, n = DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5
Private Key: d = 74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D

m = 4120746f702073656372657421
c = 6FB078DA550B2650832661E14F4F8D2CFAEF475A0DF3A75CACDC5DE5CFC5FADC
m decrypted = 4120746f702073656372657421
```

Task 3: Decrypting a Message

The same d is used to decrypt the supplied ciphertext. When converted back to ASCII, it reads "Password is dees".

```
[05/01/19]seed@VM:~/.../hw3$ ./a.out
Public Key: e = 010001, n = DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5
Private Key: d = 74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D

c = 8C0F971DF2F3672B28811407E2DABBE1DA0FEBBBD7C7DCB67396567EA1E2493F
m decrypted = 50617373776F72642069732064656573
[05/01/19]seed@VM:~/.../hw3$ echo 50617373776F72642069732064656573 | xxd -r -p
Password is dees[05/01/19]seed@VM:~/.../hw3$
```

Task 4: Signing a Message

Here we generate a signature by applying the private key to two plaintexts: "I owe you \$2000." and "I owe you \$3000.". As we can see, changing only one character entirely changes the signature.

```
[05/01/19]seed@VM:~/.../hw3$ ./a.out
Public Key: e = 010001, n = DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5
Private Key: d = 74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D

m = 49206F776520796F752024323030302E // I owe you $2000.
c = 3A759CBF53901AC41373EEC603955A8E6AF8D3BCD5E9F6DD62C873CBB675051E
m decrypted = 49206F776520796F752024323030302E
[05/01/19]seed@VM:~/.../hw3$ gcc bignum2.c -lcrypto
[05/01/19]seed@VM:~/.../hw3$ ./a.out
Public Key: e = 010001, n = DCBFFE3E51F62E09CE7032E2677A78946A849DC4CDDE3A4D0CB81629242FB1A5
Private Key: d = 74D806F9F3A62BAE331FFE3F0A68AFE35B3D2E4794148AACBC26AA381CD7D30D

m = 49206F776520796F752024333030302E // I owe you $3000.
c = D06908047527906C724937169FA68CE0AC442FEB99D1880438D331A88F44B074
m decrypted = 49206F776520796F752024333030302E
```

Task 5: Verifying a Signature

Here we verify the signature s of a message. By applying the public key, we accurately create the message “Launch a missile.”. However, by changing only a single hex character, the result after verifying is completely unreadable.

```
[05/01/19]seed@VM:~/.../hw3$ gcc bignum2.c -lcrypto  
[05/01/19]seed@VM:~/.../hw3$ ./a.out  
s = 643D6F34902D9C7EC90CB0B2BCA36C47FA37165C0005CAB026C0542CBDB6802F  
verify = 4C61756E63682061206D697373696C652E  
[05/01/19]seed@VM:~/.../hw3$ echo 4C61756E63682061206D697373696C652E | xxd -r -p  
Launch a missile.  
[05/01/19]seed@VM:~/.../hw3$  
[05/01/19]seed@VM:~/.../hw3$ gcc bignum2.c -lcrypto  
[05/01/19]seed@VM:~/.../hw3$ ./a.out  
s = 643D6F34902D9C7EC90CB0B2BCA36C47FA37165C0005CAB026C0542CBDB6803F  
verify = 91471927C80DF1E42C154FB4638CE8BC726D3D66C83A4EB6B7BE0203B41AC294  
[05/01/19]seed@VM:~/.../hw3$ echo 91471927C80DF1E42C154FB4638CE8BC726D3D66C83A4EB6B7BE0203B41AC294 | xxd  
-r -p  
00,00c000rm=f0:N000000000 05/01/19]seed@VM:~/.../hw3$
```


Here are the certificates we downloaded from a real web server, the certificate snippet on the left is the web server's certificate we downloaded from. The one on the right is the intermediate CA.

M1H5Q3CJBIGagAIwIgt7YqCWYQn9bVoiPvZlK7ZF2HAnBgkqkIc9w0BAQsFADB1
M1JW3CQVDDQVQGEVJZUPEtVMBGAIUECHMRGlnAu1NcnJq3w5JMRKwFvG0VQLEXB3
d3CuzcJlNaWNLcnqUy2TmTQmWgVYDQDExTEawdpQ2YdCBTSEeYIEV4dGvCuzGVk
IFzhbZkXyRxpB24G2U2YmVmyEINBMB4QDYE4MH4UDDWADMBwFoXDRtIMWDYwMzEY
MDAwMFowggczHTABNgBMB8ABFByaXzhdGUGT3JnyYSpemF0aHwXMRREYQLWkYB
BAGCnzwCAQMTALVtMkrhFWyLkYwBAGCnzwCAQITCER1bGF3JYLmRAAdgVDDVQF
EGe1MTU3nTQ1W03mCQYDQVQGEWJvUzETMBEGA1UECHMRKQ2f3aZwYcM5pYTEHMBQ
A1UEBXMkRwFTEZUyZyYw5jXaXjBzEVNBMBGA1UECHMR210SHVlLjBjBmNmRwMEYQY
VQDDWpnaxRxdUyIUy29tMT1B81JAnBgkqhK1G9w0BAQEFAAOCAQAMIBTCGKCAQE
jxyq8jYXDDRBTyitcnB90865tWBzPHSbIndG/XqYQkzFMBLXmqkzCzFdBTRByneZ
wSPzXWQVLF474JW6LScnzCE2F0XcEqL0JCuz9jPaqBr7u0rNLghwG3Y1Ydqbg6S0j
/4x+ogEG3dr/USY7UvF658DKyEM56v0eY9vMDvTJustJkgw+o5ZAKFVYMLUe
SMwFtoTDJFnvF6Jk10wXsp1xwQ/MPQK1cyq00tYgVylGd3NfY6CEPF22AX8A1Q
xbcaI+GwFqL1F87j3yJ+kJME19/LRgVQ7o2R1XN5XpFCBWu+XfE6pewP/3JrbkM
fZ2beYfSAUkZmh1LwPePSYIDAA0840IdETCA3UyHwYDVR0JBGwFAUPNdQpdag
re7ZSNFvM2dMh1Pj41g8WbQIDA080BBYEFnMCn2Fmnnv7JfomZq84mqhJ6k1pMCUQ
A1UEdQE0MByCcMndpdGh1Yi5jb22Cnd3dySnaXRodUyUy29tMA4GA1UdDwEB/wQE
A1F0dADBgNBHVSUEfYAubggrBGEFBCQDAQYTKwYb9QAwHawtdQYDVR0fBg4wboDA
oDKgMTyUAHR0EDoV2LnybBMuZGlnaWNLcnqUy29tL3NoYTI2XZYTzXZyTvyLWcy
LmNybDA0EDoDKgMTyUAHR0EDoV2LnybBMuZGlnaWNLcnqUy29tL3NoYTI2XZYTzVY
LmNybLWcyLmNybDBLNgBVNSAEADBBQDCGCHCAGGAG/wCAATAqMCGCCsAGCFBwB1F
FhoxdHRwcZov13d3dy5kaidpY2YyGd3j2b0v1BMTACGBwEBDAEMIBGgrBgEgF
BCQBAQRHMowJA1YkYwBBQUHACGGGH0dH6Ly9v3NwLRnqUy29tL2XJZJ0LmNBwTBS
BgrBgEgFBCwAozGahr0EDoV2LnyhY2VdHMuZGlnaWNLcnqUy29tLRpZ2L2DZXJ0
U0BhMkZ4dGVuZGVkVmw5aSRhWdGlnbLNLcnqUy29tL2XJZJ0LmNBwTBSBgrBgEgF
M1BfGyKKYwBkVMEQIEA5GaC4EggfGyAlHgAdgckuqmd0tBHyfIE7H6L2M3AKPQY
Pbk37j3Jd809yA3CEAAANBNBY0K5AAEAwBHMEUCIQRZp38cTWSH2GD8Dpe/Uf
Wns+M4EC2+d1c5yKZyG1Q5Gy6vyaazXk2BNkM2mgmUeFNS52pARBM1pJfG
USUADgBwFAaal9F9cn174b1Esj7HRNa5vJkRXMDv13vY1onQ3QAAMNBYGv0d0TAAE
AwBHMEUCIQC17omxYb2LMD0b2LobTEhrAaY1Io7n6J3drtYdmPUWJZtYgVdWIAZ51
vK9EN1NbG27f0mYb8TLVND005217eXRCr21AdG4C7d+8H4pCzt0U1SeqkthOFEv
CqTS6bqQ1mQ2j3h7RQA0AANBNY3fAAEAwBHMEUCIqChzdTKUu2N+XcqCk00JY9N
8EYtN1oVxh04yPk6d3EPgIGDNH5uR83Uc1QV489y0e0w2040ndREGTKUyUExpG
e0QWdQ3KvJ2IhvcNAQELBQADgGEBAHApPannPwH/p2o3J5yqAH8mqGfaunucVE+v
ca+881kDK/LkdZ1f86kThZ1YcLtkfzcG93hWkBS214NRNNP91aQodnK1275vN
HnXVUGw+xyjJMLGqgkpeOnZ2RB14kcTOGpA45AJuuaMwXmCo7jYUwPwLE1NULVB
Kqg6LKH0c4K0sZnxEFHhF1z29pW2AVWJrMEC/2z5z8hdvnxVUYy10e67XINk
myQKc+Y5SBZ2YLnSFWFHR3u5dcaQGGAR426Yd4L3B8d4d0iBaM+FXSXB1pUf
WJbST4VXmdao17uzFm0JA4zkQXDZAVFSXgJlAFadfySna/teik=

[illegible]

-----END CERTIFICATE-----

Step 2-4:

Below is the value of n that is extracted from the x509 certificate using the flag “modulus”. And by printing out all the fields, we found the value of the exponent ‘e’.

Since there was no specific *openssl* command to extract the signature field, we needed to print out all the fields, and copy/paste the specific signature block into a file, but since we needed a hex-string version of the signature, we used the *-d* flag to delete the “:” and spaces from the data.

To take advantage of the ASN.1 encoding used on X.509 certificates, we used the built-in ASN.1 parser in *openssl* called *asn1parse*. Once parsed, we use the starting offset of 4 and *-strparse* flag to only retrieve the body of the certificate without including the signature block. And with the both of the certificate, we calculated the hash, which is notated below as *hashed body*.

```
mod
D753A40451F899A616484B6727AA9349D039ED0CB0B00087F1672886858C8E63
DABCB14038E2D3F5ECA50518B83D3EC5991732EC188CFAF10CA6642185CB0710
34B052882B1F689BD2B18F12B0B3D2E7881F1FEF387754535F80793F2E1AAAA8
1E4B2B0DABB763B935B77D14BC594BDF514AD2A1E20CE29082876AAEEAD764D6
9855E8FDAF1A506C54BC11F2FD4AF29D8B7F0EF4D5BE8E16891255D8C07134EE
F6DC2DECC48725868DD821E4B04D0C89DC392617DDF6D79485D80421709D6F6F
FF5CBA19E145CB5657287E1C0D4157AAB7B827BBB1E4FA2AEF2123751AAD2D9B
86358C9C77B573ADD8942DE4F30C9DEEC14E627E17C0719E2CDEF1F910281933
```

```
exp
10001
```

```
sig
700f5a96a758e5bf8a9da827982b007f26a907daba7b82544faf69cfbcf25903
2bf2d5745825d81ea42076626029732ad7dccc6f77856bca6d24f83513473fd2
e2690a9d342d7b7b9bcd1e75d5506c3ecb1ca330b1aa9207a93a767645bd7891
c4ce1a9e22e40b89bae68cc17982a3b8d4c0fc1f2ded4d5255412aa83a2cad07
72ae0ad2c667c44f07171899f765a95760155a344c11cff6cf6b213680efc6f1
5463263539eebbc483649b240a73eca0481673c8b9d7485556987af7bb975c69
a406180478dafa9876be222f7f0777874e88199af855ec5c122a5948db493e15
5e675aa25eecc53288c0e33931403640bc5e5780994015a75fc929dafed7a29
```

```
hashed body
85088f934d3e58e3673ea5be32c7c8cf6965e4ab93fbed4fff634723f46d5693
```

Below is finally using the CA's public key/signature and the body of the server's certificate and running our program, the underlines verify that the program works. There is SHA-256 padding that we can ignore. Looking at the last bit of characters we can see that both hashes are indeed the same.

Magic bytes (padding) for digest algorithms

```
[05/02/19]seed@VM:~/.../hw3$ ./a.out  
verify = 01FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF  
FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF  
FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF  
FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF  
FFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFF  
1300d06096086480165030402010500042085088f934D3E58E3673EA5BE32C7C8CF6965E4AB93FBED4FFF634723F46D5693  
[05/02/19]seed@VM:~/.../hw3$ ./a.out | tail -c 65  
85088f934D3E58E3673EA5BE32C7C8CF6965E4AB93FBED4FFF634723F46D5693  
[05/02/19]seed@VM:~/.../hw3$ sha256sum c0 body.bin  
85088f934d3e58e3673ea5be32c7c8cf6965e4ab93fbcd4fff634723f46d5693    c0 body.bin
```

Code:

```
// Eduard Klimenko
// Textbook RSA (without padding).
#include <stdio.h>
#include <openssl/bn.h>

int main () {
    BN_CTX *ctx = BN_CTX_new();           // struct used to hold BIGNUM
    BIGNUM *one = BN_new();                // the number 1
    BIGNUM *p = BN_new();                  // p
    BIGNUM *p_minus_1 = BN_new();          // p-1
    BIGNUM *q = BN_new();                  // q
    BIGNUM *q_minus_1 = BN_new();          // q-1
    BIGNUM *n = BN_new();                  // n = p*q
    BIGNUM *e = BN_new();                  // e
    BIGNUM *e_mult_inverse = BN_new();     // e^-1
    BIGNUM *phi = BN_new();                // phi = (p-1)*(q-1)
    BIGNUM *d = BN_new();                  // d
    BIGNUM *m_in = BN_new();               // message to be encrypted
    BIGNUM *m_out = BN_new();              // message produced from decryption
    BIGNUM *c = BN_new();                  // ciphertext

    // the number 1
    BN_dec2bn(&one, "1");

    // should be randomly generated
    BN_dec2bn(&p, "7");
    BN_dec2bn(&q, "11");
    BN_dec2bn(&e, "7");

    // message
    BN_dec2bn(&m_in, "25");

    // key generation:
    // n = p*q
    BN_mul(n, p, q, ctx);

    // p-1 and q-1
    BN_sub(p_minus_1, p, one);
    BN_sub(q_minus_1, q, one);

    // phi(n) = (p-1)*(q-1)
    BN_mul(phi, p_minus_1, q_minus_1, ctx);

    // e*x+phi(n)*y=1
    BN_mod_inverse(e_mult_inverse, e, phi, ctx);

    // d = e^-1 mod phi(n)
    BN_mod(d, e_mult_inverse, phi, ctx);

    // print out p, q, e, n, d, m
    printf("p = %s, q = %s\n\n", BN_bn2dec(p), BN_bn2dec(q));
    printf("Public Key: e = %s, n = %s\n\nPrivate Key: d = %s\n\n", BN_bn2dec(e), BN_bn2dec(n),
    BN_bn2dec(d));
    printf("m = %s\n", BN_bn2dec(m_in));

    // encryption: c = m^e mod n
    BN_mod_exp(c, m_in, e, n, ctx);
    printf("c = %s\n\n", BN_bn2dec(c));

    // decryption: m = c^d mod n
    BN_mod_exp(m_out, c, d, n, ctx);
    printf("c decrypted = %s\n", BN_bn2dec(m_out));

    // free struct
    BN_CTX_free(ctx);

    return 0;
}
```

Key Generation

$p = 7, q = 11$
 $n = 77 \quad \phi(n) = (7 - 1)(11 - 1) = 60$
 $e = 7$
 $d = 7^{-1} \bmod 60 = 43$
 $\text{pub: } \{7, 77\} \quad \text{priv: } \{43, 7, 11\}$

Encrypt

$M = 25$
 $C = M^e \bmod n = 25^7 \bmod 77 = 53$

Decrypt

$M = C^d \bmod n = 53^{43} \bmod 77$

$C_p = 53^{43} \bmod 7$
 $43 \bmod 6 = 1$
 $53 \bmod 7 = 4$
 $C_p = 4^1 \bmod 7 = 4$

$C_q = 53^{43} \bmod 11$
 $43 \bmod 10 = 3$
 $53 \bmod 11 = 9$
 $C_q = 9^3 \bmod 11 = 3$

$M = 4(11 * 11^{-1} \bmod 7) + 3(7 * 7^{-1} \bmod 11) \bmod 77 = 25$

The program was built off of this small example for correctness as there was no way to verify the algorithm at first. For the completion of the lab, the only variables that were changed were the inputs such as p, n, e, d and the message.