

1 Theory

1.1 Control flow graphs

In figure 1 I assume that the code is:

```
for ( a ; b ; c ) {}  
d;  
e;
```

not:

```
for ( a ; b ; c )  
  d;  
e;
```

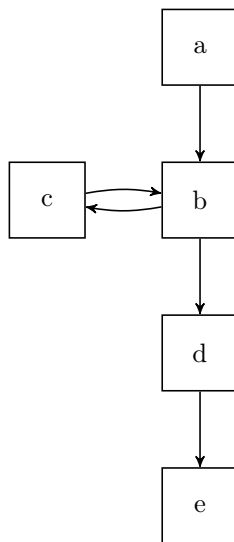


Figure 1: `for (a ; b ; c) d ; e ;`

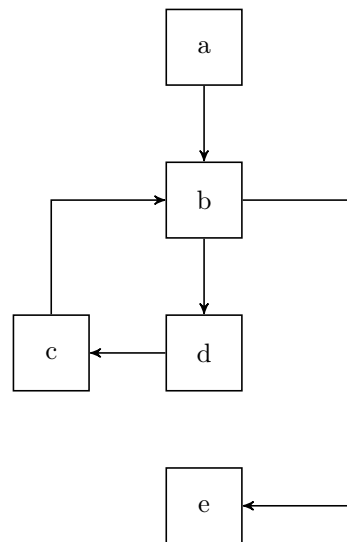


Figure 2: `a ; while (b) { d ; c ; } e ;`

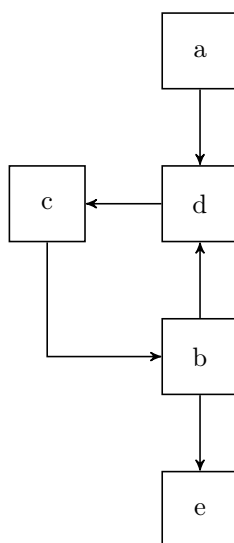


Figure 3: `a ; do { d ; c ; } while (b); e ;`

1.2 Optimizations

1.2.1 Copy propagation

1. $a=1$
2. $b=2$
3. $c=3$
4. $d=a+x$
5. $e=b+c$
6. $f=e$
7. $g=f$
8. $g=d+y$
9. $a=b+c$

Figure 4: Before

1. $a=1$
2. $b=2$
3. $c=3$
4. $d=a+x$
5. $e=b+c$
6. $f=e$
7. $g=e$
8. $g=d+y$
9. $a=b+c$

Figure 5: After

1.2.2 Common subexpression elimination

1. $a=1$
2. $b=2$
3. $c=3$
4. $d=a+x$
5. $e=b+c$
6. $f=e$
7. $g=f$
8. $g=d+y$
9. $a=b+c$

Figure 6: Before

1. $a=1$
2. $b=2$
3. $c=3$
4. $d=a+x$
5. $t1=b+c$
6. $e=t1$
7. $f=e$
8. $g=f$
9. $g=d+y$
10. $a=t1$

Figure 7: After

1.2.3 Constant propagation

Assuming only constant propagation, not constant folding.

1. $a = 1$
2. $b = 2$
3. $c = 3$
4. $\mathbf{d = a + x}$
5. $\mathbf{e = b + c}$
6. $f = e$
7. $g = f$
8. $g = d + y$
9. $\mathbf{a = b + c}$

Figure 8: Before

1. $\mathbf{a=1}$
2. $\mathbf{b=2}$
3. $\mathbf{c=3}$
4. $\mathbf{d = 1 + x}$
5. $\mathbf{e = 2 + 3}$
6. $\mathbf{f=e}$
7. $\mathbf{g=f}$
8. $\mathbf{g=d+y}$
9. $\mathbf{a = 2 + 3}$

Figure 9: After