Algorithm Library

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Algorithm Library by palayutm

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1 图论

```
1.1 connected graph
1.1.1 割点
void dfs(int u, int fa)
    int low[u] = pre[u] = ++dfs_block;
    int child = 0;
    for(int i = 0; i < G[u].size(); i++)</pre>
        int v = G[u][i];
        if(!pre[v])
        {
             child++;
             dfs(v, u);
             low[u] = min(low[u], low[v]);
             if(low[v] >= pre[u])
                 iscut[u] = true;
        }
        else if(v != fa)
             low[u] = min(low[u], pre[v]);
    }
    if(fa < 0 && child == 1) iscut[u] = false;</pre>
}
1.1.2 割边
struct node
{
    int fir, sec;
    bool operator < (node a) const {</pre>
        if(fir != a.fir) return fir < a.fir;</pre>
        return sec < a.sec;</pre>
    }
};
void dfs(int u, int fa)
    low[u] = pre[u] = ++dfs_block;
    for(int i = 0; i < G[u].size(); i++)</pre>
    {
        int v = G[u][i];
        if(!pre[v])
```

{

```
dfs(v, u);
            low[u] = min(low[u], low[v]);
            if(low[v] > pre[u])
                all[co++] = node\{u, v\};
        }
        else if(v != fa)
            low[u] = min(low[u], low[v]);
    }
}
1.1.3 强连通分量
void dfs(int u)
    low[u] = dfn[u] = ++dfs_block;
    S.push(u);
    for(int i = 0; i < G[u].size(); i++)</pre>
    {
        int v = G[u][i];
        if(!dfn[v])
            dfs(v);
            low[u] = min(low[u], low[v]);
        else if(!sccno[v])
            low[u] = min(low[u], low[v]);
    }
    if(low[u] == dfn[u])
    {
        ++scc_cnt;
        while(1)
            int x = S.top(); S.pop();
            sccno[x] = scc_cnt;
            if(x == u) break;
        }
    }
}
1.1.4 强连通分量缩点
void dfs(int u, int fa)
{
    low[u] = dfn[u] = ++dfs_block;
```

```
S.push(u);
    bool flag = false;
    for(int i = 0; i < G[u].size(); i++)</pre>
        int v = G[u][i];
        if(!dfn[v])
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            if(low[v] > dfn[u])
                 bridge++;
        else if(v != fa || flag)
            low[u] = min(low[u], low[v]);
2
        if(v == fa) flag = true;
    if(low[u] == dfn[u])
        ++scc_cnt;
        while(1)
        {
            int x = S.top(); S.pop();
            sccno[x] = scc_cnt;
            if(x == u) break;
        }
    }
}
void find_scc()
{
    init();
    for(int i = 1; i <= n; i++)
        if(!dfn[i]) dfs(i, -1);
}
for(int i = 1; i <= n; i++)</pre>
    for(int j = 0; j < G[i].size(); j++)</pre>
    {
        int v = G[i][j];
        if(sccno[i] != sccno[v])
        {
            M[sccno[i]].push_back(sccno[v]);
            M[sccno[v]].push_back(sccno[i]);
        }
    }
void bfs(int s)
```

```
{
    queue<int> q;
    q.push(s);
    memset(vis, 0, sizeof(vis));
    vis[s] = 1; dis[s] = 0; last_node = s; ret = 0;
    while(!q.empty())
        int u = q.front();
        q.pop();
        for(int i = 0; i < M[u].size(); i++)</pre>
            int v = M[u][i];
            if(!vis[v])
            {
                vis[v] = 1;
                dis[v] = dis[u] + 1;
                if(dis[v] > ret)
                     ret = dis[v];
                     last_node = v;
                }
                q.push(v);
            }
        }
    }
}
    bfs(1);
    bfs(last_node);
```

```
String
                                     top = 0;
                                     head = last = &a[0];
2.1 KMP
                                  void add(int x)
 * Args:
                                    node *p = &a[++top], *mid;
 * s[]: string
                                     p->len = last->len + 1;
 * Return:
                                     mid = last, last = p;
 * fail[]: failure function
                                     for (; mid && !mid->chd[x];
 */

→ mid = mid->link)
int fail[N];
                                     \rightarrow mid->chd[x] = p;
void getfail(char s[])
                                     if (!mid) p->link = head;
                                     else{
  fail[0] = -1;
                                       if (mid->len + 1 ==
  int p = -1;
                                       \rightarrow mid->chd[x]->len) {
  for (int i = 0; s[i]; i ++)
                                         p->link = mid->chd[x];
  → {
                                       } else {
    while (p!=-1 \&\&
                                         node *q = mid->chd[x],
    \rightarrow s[i]!=s[p]) p =
                                         \rightarrow *r = &a[++top];
    _{\hookrightarrow} \quad \texttt{fail[p];}
                                         *r = *q, q->link =
    fail[i+1] = ++p;
                                         \rightarrow p->link = r;
  }
                                         r->len = mid->len + 1;
}
                                         for (; mid &&
                                          \rightarrow mid->chd[x] == q;
2.2 Suffix Automaton

    mid = mid->link)

/*
                                          \rightarrow mid->chd[x] = r;
* 1 call init()
                                       }
* 2 call add(x) to add every
                                     }
}
 * Args:
 * Return:
   an automaton
   link: link path pointer
     len: maximum length
 */
struct node{
  node* chd[26], *link;
  int len;
}a[3*N], *head, *last;
int top;
void init()
  memset(a, 0, sizeof(a));
```