
Name: Ernesto Diaz

Panther-ID: $x \times x \times - 0534$

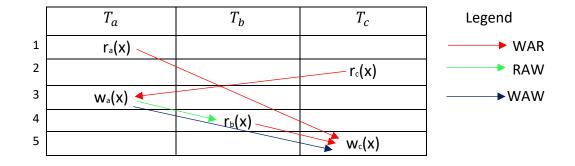
Course: COP 4722
Assignment#: 2

Due: Wed, Feb 21, 2018

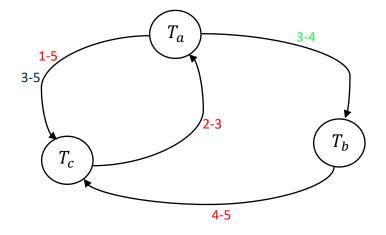
I hereby certify that this work is my own and none of it is the work of any other person.

Signature: _____

 S_1 : $r_a(x)$; $r_c(x)$; $w_a(x)$; $r_b(x)$; $w_c(x)$;



Conflict Precedence Graph:



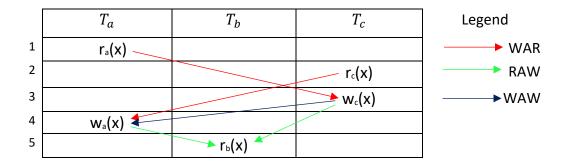
Since the precedence graph has the following cycles:

$$T_a \to T_b \to T_c \to T_a$$

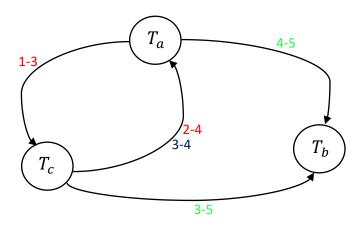
$$T_c \to T_a \to T_c$$

this schedule is not conflict serializable. So, there is no conflict equivalent serial schedule for this schedule.

$$S_2$$
: $r_a(x)$; $r_c(x)$; $w_c(x)$; $w_a(x)$; $r_b(x)$;



Conflict Precedence Graph:



Since the precedence graph has the following cycle:

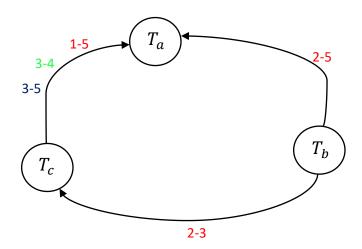
$$T_a \to T_c \to T_a$$

this schedule is not conflict serializable. So, there is no conflict equivalent serial schedule for this schedule.

 S_3 : $r_c(x)$; $r_b(x)$; $w_c(x)$; $r_a(x)$; $w_a(x)$;

	T_a	T_b	T_c	Legend
1			r _c (x)	→ WAR
2				→ RAW
3			W _c (x)	→WAW
4	r₃(x)			
5	W _a (x)			

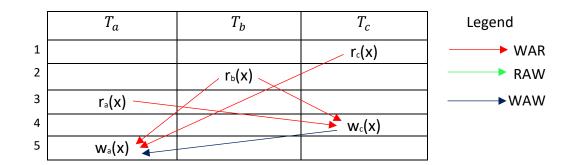
Conflict Precedence Graph:



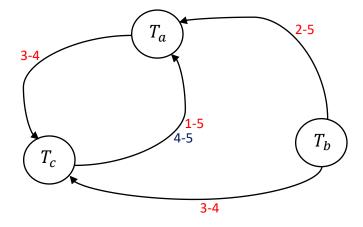
Since the precedence graph has no cycle this schedule is conflict serializable. So, the conflict equivalent serial schedule for this schedule is:

$$T_b \to T_c \to T_a$$

 S_4 : $r_c(x)$; $r_b(x)$; $r_a(x)$; $w_c(x)$; $w_a(x)$;



Conflict Precedence Graph:

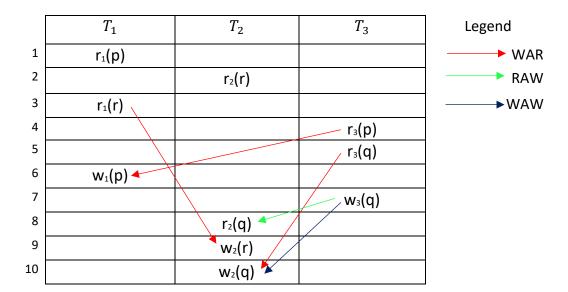


Since the precedence graph has the following cycle:

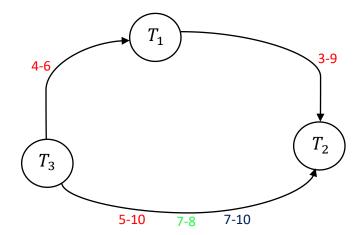
$$T_a \to T_c \to T_a$$

this schedule is not conflict serializable. So, there is no conflict equivalent serial schedule for this schedule.

 $S_5: r_1(p); \, r_2(r); \, r_1(r); \, r_3(p); \, r_3(q); \, w_1(p); \, w_3(q); \, r_2(q); \, w_2(r); \, w_2(q);$



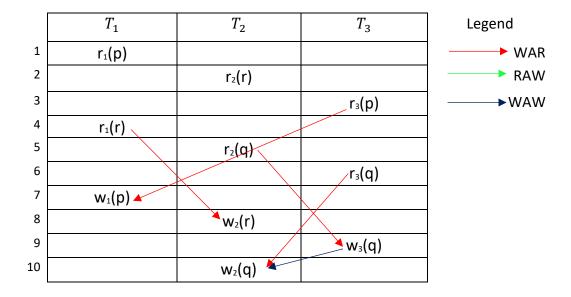
Conflict Precedence Graph:



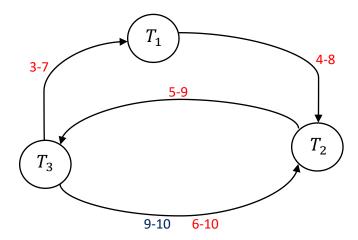
Since the precedence graph has no cycle this schedule is conflict serializable. So, the conflict equivalent serial schedule for this schedule is:

$$T_3 \rightarrow T_1 \rightarrow T_2$$

 $S_6 \colon r_1(p); \ r_2(r); \ r_3(p); \ r_1(r); \ r_2(q); \ r_3(q); \ w_1(p); \ w_2(r); \ w_3(q); \ w_2(q);$



Conflict Precedence Graph:



Since the precedence graph has the following cycles:

$$T_1 \rightarrow T_2 \rightarrow T_3 \rightarrow T_1$$

$$T_2 \rightarrow T_3 \rightarrow T_2$$

this schedule is not conflict serializable. So, there is no conflict equivalent serial schedule for this schedule.