Supplementary Notes to KARL Library

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I. SOTA ALGORITHM

SOTA [5] adopts the multi-step approach [7], [1], [2] for evaluating the bound functions (LB and UB), combining with existing indexing structures to support kernel density classification (query type I-τ). We extend their algorithm to support different other types of machine learning models, including approximate kernel density estimation (query type I-ϵ), 1-class SVM (query type II-τ) and 2-class SVM (query type III-τ). We name this algorithm as Multi-Step Kernel Prediction (MSKP).

Algorithm 1 Multi-Step Kernel Prediction (MSKP)

1: **procedure** MSKP(query $\mathbf{q}$, weights $\{w_1,...,w_n\}$, tree T, threshold τ)
2:     Create a max-heap H
3:     $e \leftarrow T.R_{root}$
4:     $\widehat{lb} \leftarrow LB(\mathbf{q}, e), \widehat{ub} \leftarrow UB(\mathbf{q}, e)$
5:     enheap e to H
6:     **while** $H \neq \emptyset$ **do**
7:         **if** $\widehat{lb} \geq \tau$ **then**
8:            return 1
9:         **if** $\widehat{ub} < \tau$ **then**
10:           return -1
11:         $R \leftarrow$ deheap an entry in H
12:         $\widehat{lb} \leftarrow \widehat{lb} - LB(\mathbf{q}, e.R), \widehat{ub} \leftarrow \widehat{ub} - UB(\mathbf{q}, e.R)$
13:         **if** e is leaf **then**
14:            $temp \leftarrow \sum_{\mathbf{P_i} \in e.R} w_i \, \mathcal{K}(\mathbf{q}, \mathbf{p_i})$
15:            $\widehat{lb} \leftarrow \widehat{lb} + temp, \widehat{ub} \leftarrow \widehat{ub} + temp$
16:         **else**
17:            **for** each child e_c in e **do**
18:                $\widehat{lb} \leftarrow \widehat{lb} + LB(\mathbf{q}, e_c.R)$
19:                $\widehat{ub} \leftarrow \widehat{ub} + UB(\mathbf{q}, e_c.R)$
20:                enheap e_c to H

Algorithm 1 is only used for the classification-based queries (types I-τ, II-τ and III-τ). For query type I-ϵ, the input threshold τ should be replaced by relative error ϵ. We also need to replace lines 7 to 10 by the following termination condition.

Algorithm 2 Terminiation condition for query type I-ϵ

1: **if** $\widehat{ub} \leq (1+\epsilon)\widehat{lb}$ **then**
2:     return $\frac{\widehat{lb}+\widehat{ub}}{2}$

There is another advance implementation for the termination condition for Algorithm 3, which is based on the unpublished paper of our work [3] in Earth Mover's Distance. This paper is now under submission to TKDE.

Algorithm 3 Advanced terminiation condition for query type I-ϵ [3]

1: **if** $\frac{\widehat{ub}-\widehat{lb}}{\widehat{ub}+\widehat{lb}} \leq \epsilon$ **then**
2:     $R = \frac{2 \times \widehat{lb} \times \widehat{ub}}{\widehat{lb}+\widehat{ub}}$
3:     return R

II. RELATIONSHIP BETWEEN SOTA AND KARL

Both SOTA and KARL utilize the same algorithm MSKP. However, compared with SOTA, KARL utilizes tighter bound functions to boost up the efficiency performance. During the implementation of KARL, we only replace LB and UB by our bounding functions [4].

III. AUTO-TUNING (OFFLLINE)

In Section III-C of our paper [4], we develop the auto-tuning method for obtaining the best index construction in the offline stage. First, we sample 1000 queries from the query dataset as the workload $\mathcal{WL}$. To ensure the fairness, these queries will not be used in the online phase. Then, our algorithm Auto chooses the index from either kd-tree [6] or ball-tree [8] and the most suitable leaf node capacity from the capacity list $CL = \{10, 20, 40, 80, 160, 320, 640\}$ in which the best setting bs provides the fastest running time f_{time} in the selected workload $\mathcal{WL}$.

Algorithm 4 shows the pseudocode of our implementation.

Algorithm 4 Auto-tuning (Offline)

1: **procedure** AUTO(Workload $\mathcal{WL}$, weights $\{w_1,...,w_n\}$, tree T, threshold τ, capacity list CL, tree list $TL =\{$kd,ball$\}$)
2:     $bs \leftarrow$ null
3:     $f_{time} \leftarrow \infty$
4:     **for** $t \in TL$ **do**
5:         **for** $c \in CL$ **do**
6:            $T \leftarrow$ Build tree t with capacity c
7:            $s \leftarrow timer()$
8:            **for** $q \in \mathcal{WL}$ **do**
9:                $MSKP(\mathbf{q}, weights, T, \tau)$
10:            $e \leftarrow timer()$
11:            $temp \leftarrow e - s$
12:            **if** $temp \leq f_{time}$ **then**
13:                $bs \leftarrow \{t, c\}$
14:                $f_{time} \leftarrow temp$
15:            Remove tree t
16:     return bs

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