Contents			8 Geometry       18         8.1 Basic       18
			8.2 KD Tree
1	Basic	1	1 8.4 Voronoi Diagram 20
	1.1 vimrc	1 1	20
	1.3 Increase stack size	1	0.C II-1f Dl I. d
	1.4 Pragma optimization	2	
_			8.8 Triangle Center
2	Flow, Matching 2.1 Dinic's Algorithm	<b>2</b> 2	o.b Tolygon Center
	2.2 Minimum-cost flow	2	0.10 Maximum friangic
	2.3 Gomory-Hu Tree	2	0.11 1 01110 111 1 017g011
	2.4 Stoer-Wagner Minimum Cut	2	2 8 13 Tangent of Circles and Points to Circle 22
	2.5 Kuhn–Munkres Algorithm	3	3 8.14 Area of Union of Circles 29
	2.6 Maximum Matching on General Graph	3 3	9 15 Minimum Distance of 9 Delimens
	2.8 Flow Model	5	0.10 OD Cl
			8.17 3D Convex Hull
3	Data Structure           3.1 <ext pbds=""></ext>	<b>5</b> 5	9
	3.2 Li Chao Tree	5	
	3.3 Link-Cut Tree	6	
		_	9.1 Bitwise Hack
4	Graph	6	9.2 Hilbert's Curve (laster Mo's algorithm)
	4.1 Heavy-Light Decomposition	6 7	9.3 Java
	4.3 Minimum Mean Cycle	7	7 9.4 Dancing Links
	4.4 Minimum Steiner Tree	7	0.6 Manhattan Distance MCT
	4.5 Directed Minimum Spanning Tree	7	0.7 101 0010 A1: 4:1
	4.6 Maximum Clique	8	
	4.8 Dominator Tree	8	
	4.9 Virtual Tree	9	9
	4.10 Vizing's Theorem	9	
	4.11 System of Difference Constraints	9	9 1 Basic
5	String	9	9
	5.1 Knuth-Morris-Pratt Algorithm	9	1 1 Trimano
	5.2 Z Algorithm	9	
		10 10	
	5.5 Suffix Automaton		
	5.6 Suffix Array		syn on
	5.7 Palindromic Tree		i Cil aluma dindant an
	5.8 Circular LCS	11 12	1. (CD (CD ) E O
	5.9 Lexicographically Smallest Rotation	12	Z Thoremap (tells) teles
6		<b>12</b>	
	6.1 Fast Fourier Transform		
	6.2 Number Theoretic Transform	13	
	6.3 Polynomial Division		· · · · · · · · · · · · · · · · · · ·
	6.4 Polynomial Square Root		
	6.5 Multipoint Evaluation		
	6.6 Polynomial Interpolation		
	6.7.1 XOR Convolution		
	6.7.2 OR Convolution	14	4   if (p == end) {
	6.7.3 AND Convolution		21 ((c 5a 1 5aa(5a ; 1,, 5ca))
	6.8 Simplex Algorithm		
	6.9 Schreier–Sims Algorithm		p 50,
	6.10 Berlekamp-Massey Algorithm		
		15	5  }
	6.12 Pollard's Rho		_   a surp and a supplemental
	6.14 Discrete Logarithm		The source of th
	6.15 Quadratic Residue		
	6.16 Gaussian Elimination		6 == -1) return false:
	6.17 $\mu$ function		c = '-'? (flag = true, x = 0) : (x = c - '0');
	6.18 Partition Function		$\frac{1}{2}$ write (c = getchar(), c >= 0 && c <= 9 ) x = x · 10 + c -
	6.20 De Bruijn Sequence		_   • • • • • • • • • • • • • • • • • •
	6.21 Extended GCD	17	return true:
	6.22 Euclidean Algorithms		17   3
	6.23 Chinese Remainder Theorem		
	6.24.1 Kirchhoff's Theorem		
	6.24.2 Tutte's Matrix	17	17
	6.24.3 Cayley's Formula		1.0 Increase stack size
	6.24.4 Erdős–Gallai Theorem		.8
	6.25 Primes	10	
7	v e e	18	register long rsn gsm("rsn").
	7.1 Dynamic Convex Hull		$-1$ Chan $\pi$ n $-1$ (Chan $\pi$ )malloc(Clash)   Clash $\pi$ hab $-1$ (Chan $\pi$ )mch
	7.2 1 <i>D</i> /1 <i>D</i> Convex Optimization		asm ("mova %0 %%nsn\n"··"n"(n)).
	7.3.1 Totally Monotone (Concave/Convex)		8 // main
		18	acm ("may a 0/0 0/0/ncm\n"··"n"(hak)).
	7.3.3 Optimal Split Point	18	.8

### 1.4 Pragma optimization

```
#pragma GCC optimize("Ofast", "no-stack-protector", "no-math-
             "unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4,sse4.2,popcnt,abm,
    mmx, avx, tune=native, arch=core-avx2, tune=core-avx2")
#pragma GCC ivdep
```

#### 2 Flow, Matching

### Dinic's Algorithm

```
struct dinic {
  static const int inf = 1e9;
   struct edge {
     int to, cap, rev;
     edge(int d, int c, int r): to(d), cap(c), rev(r) {}
  vector<edge> g[maxn];
  int qu[maxn], ql, qr;
  int lev[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i)
       g[i].clear();
   void add_edge(int a, int b, int c) {
     g[a].emplace_back(b, c, g[b].size() - 0);
     g[b].emplace_back(a, 0, g[a].size() - 1);
  bool bfs(int s, int t) {
     memset(lev, -1, sizeof(lev));
     lev[s] = 0;
     ql = qr = 0;
     qu[qr++] = s;
     while (ql < qr) {
       int x = qu[ql++];
       for (edge &e : g[x]) if (lev[e.to] == -1 && e.cap > 0) {
         lev[e.to] = lev[x] + 1;
         qu[qr++] = e.to;
     return lev[t] != -1;
  int dfs(int x, int t, int flow) {
     if (x == t) return flow;
     int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.to] == lev[x] +
      1) {
       int f = dfs(e.to, t, min(e.cap, flow - res));
      res += f;
e.cap -= f;
       g[e.to][e.rev].cap += f;
     if (res == 0) lev[x] = -1;
     return res;
  int operator()(int s, int t) {
     int flow = 0;
     for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow:
|};
```

### Minimum-cost flow

```
struct mincost {
 struct edge {
    int dest, cap, w, rev;
    edge(int a, int b, int c, int d): dest(a), cap(b), w(c),
    rev(d) {}
  vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
  bool inq[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
  void add_edge(int a, int b, int c, int d) {
    g[a].emplace\_back(b, c, +d, g[b].size() - 0);
    g[b].emplace_back(a, 0, -d, g[a].size() - 1);
```

```
bool spfa(int s, int t, int &f, int &c) {
    for (int i = 0; i < maxn; ++i) {
      d[i] = inf;
      p[i] = ed[i] = -1;
      inq[i] = false;
    d[s] = 0;
    queue<int> q;
    q.push(s);
    while (q.size()) {
      int x = q.front(); q.pop();
      inq[x] = false;
      for (int i = 0; i < g[x].size(); ++i) {
        edge &e = g[x][i];
        if (e.cap > 0 \&\& d[e.dest] > d[x] + e.w) {
          d[e.dest] = d[x] + e.w;
          p[e.dest] = x;
          ed[e.dest] = i;
          if (!inq[e.dest]) q.push(e.dest), inq[e.dest] = true;
        }
      }
    if (d[t] == inf) return false;
    int dlt = inf;
    for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[p[x]][ed
     [x]].cap);
    for (int x = t; x != s; x = p[x]) {
      edge &e = g[p[x]][ed[x]];
e.cap -= dlt;
      g[e.dest][e.rev].cap += dlt;
    f += dlt; c += d[t] * dlt;
return true;
  pair<int, int> operator()(int s, int t) {
    int f = 0, c = 0;
    while (spfa(s, t, f, c));
    return make_pair(f, c);
2.3
       Gomory-Hu Tree
int a[maxn];
```

```
vector<edge> GomoryHu(int n){
   vector<edge> rt;
   for(int i=1;i<=n;++i)g[i]=1;</pre>
   for(int i=2;i<=n;++i){</pre>
      int t=g[i];
     flow.reset(); // clear flows on all edge
     rt.push_back({i,t,flow(i,t)});
flow.walk(i); // bfs points that connected to i (use edges
      not fully flow)
     for(int j=i+1;j<=n;++j){</pre>
        if(g[j]==t && flow.connect(j))g[j]=i; // check if i can
      reach j
     }
   return rt;
}
```

### 2.4 Stoer-Wagner Minimum Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
  w[x][y] += c;
  w[y][x] += c;
pair<int, int> phase(int n) {
  memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
  int s = -1, t = -1;
  while (true) {
    int c = -1;
    for (int i = 0; i < n; ++i) {
          (del[i] || v[i]) continue;
       if (c == -1 || g[i] > g[c]) c = i;
    if (c == -1) break;
```

```
v[c] = true
     s = t, t = c;
    for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;</pre>
       g[i] += w[c][i];
  }
  return make_pair(s, t);
int mincut(int n) {
  int cut = 1e9;
  memset(del, false, sizeof(del));
  for (int i = 0; i < n - 1; ++i) {
    int s, t; tie(s, t) = phase(n);
    del[t] = true;
     cut = min(cut, g[t]);
    for (int j = 0; j < n; ++j) {
    w[s][j] += w[t][j];
    verilled
       w[j][s] += w[j][t];
    }
  return cut;
```

### 2.5 Kuhn-Munkres Algorithm

```
namespace km {
int w[kN][kN], h1[kN], hr[kN], s1k[kN];
int fl[kN], fr[kN], pre[kN], qu[kN], ql, qr;
bool vl[kN], vr[kN];
bool Check(int x) {
 if (vl[x] = true, fl[x] != -1) return vr[qu[qr++] = fl[x]] =
  while (x != -1) swap(x, fr[fl[x] = pre[x]]);
  return false;
}
void Bfs(int s, int n) {
  fill(slk, slk + n, kInf);
  fill(vl, vl + n, false);
  fill(vr, vr + n, false);
ql = qr = 0;
  qu[qr++] = s;
  vr[s] = true;
  while (true) {
    int d;
    while (ql < qr) {</pre>
      for (int x = 0, y = qu[ql++]; x < n; ++x) {
         if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] - w[x][y]))
          if (pre[x] = y, d) slk[x] = d;
          else if (!Check(x)) return;
      }
    }
    d = kInf;
    for (int x = 0; x < n; ++x) {
      if (!vl[x] \&\& d > slk[x]) d = slk[x];
    for (int x = 0; x < n; ++x) {
      if (vl[x]) hl[x] += d;
      else slk[x] -= d;
      if (vr[x]) hr[x] -= d;
    for (int x = 0; x < n; ++x) {
      if (!vl[x] && !slk[x] && !Check(x)) return;
 }
long long Solve(int n) {
 fill(fl, fl + n, -1);
fill(fr, fr + n, -1);
  fill(hr, hr + n, 0);
  for (int i = 0; i < n; ++i) hl[i] = *max_element(w[i], w[i] +
  for (int i = 0; i < n; ++i) Bfs(i, n);
long long res = 0;</pre>
  for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
  return res;
```

### 2.6 Maximum Matching on General Graph

```
namespace matching {
int fa[maxn], pre[maxn], match[maxn], s[maxn], v[maxn];
vector<int> g[maxn];
queue<int> q;
void init(int n) {
  for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
 void add_edge(int u, int v) {
  g[u].push_back(v);
   g[v].push_back(u);
int find(int u) {
   if (u == fa[u]) return u;
   return fa[u] = find(fa[u]);
int lca(int x, int y, int n) {
   static int tk = 0;
   tk++:
   x = find(x), y = find(y);
   for (; ; swap(x, y)) {
     if (x != n) {
       if (v[x] == tk) return x;
       v[x] = tk;
       x = find(pre[match[x]]);
  }
void blossom(int x, int y, int l) {
   while (find(x) != 1) {
     pre[x] = y;
     y = match[x];
     if (s[y] == 1) {
       q.push(y);
       s[y] = 0;
     if (fa[x] == x) fa[x] = 1;
     if (fa[y] == y) fa[y] = 1;
     x = pre[y];
  }
}
bool bfs(int r, int n) {
   for (int i = 0; i \le n; ++i) {
     fa[i] = i;
     s[i] = -1;
   while (!q.empty()) q.pop();
   q.push(r);
   s[r] = 0;
   while (!q.empty()) {
     int x = q.front(); q.pop();
     for (int u : g[x]) {
       if (s[u] == -1) {
         pre[u] = x;
         s[u] = 1;
         if (match[u] == n) {
           for (int a = u, b = x, last; b != n; a = last, b =
      pre[a])
             last = match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]);
         s[match[u]] = 0;
       } else if (!s[u] && find(u) != find(x)) {
         int l = lca(u, x, n);
         blossom(x, u, 1);
         blossom(u, x, 1);
       }
    }
  }
   return false;
}
 int solve(int n) {
   int res = 0;
   for (int x = 0; x < n; ++x) {
     if (match[x] == n) res += bfs(x, n);
   return res;
| }}
```

# 2.7 Maximum Weighted Matching on General Graph

```
static const int inf = INT_MAX;
static const int maxn = 514;
struct edge {
  int u, v, w;
  edge(){}
  edge(int u, int v, int w): u(u), v(v), w(w) {}
int n, n_x;
edge g[maxn * 2][maxn * 2];
int lab[maxn * 2];
int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa[maxn *
int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[maxn * 2];
vector<int> flo[maxn * 2];
queue<int> q;
int e_delta(const edge &e) { return lab[e.u] + lab[e.v] - g[e
   .u][e.v].w * 2; }
void update_slack(int u, int x) { if (!slack[x] || e_delta(g[
  u][x]) < e_delta(g[slack[x]][x])) slack[x] = u; }
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)</pre>
    if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
      update_slack(u, x);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++) q_push(flo[</pre>
  x][i]);
void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)</pre>
   set_st(flo[x][i], b);
int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) - flo[b].
  begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr);
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo[u][i</pre>
   ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
void augment(int u, int v) {
 for (; ; ) {
  int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
   if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[match[x]]),
   q_push(y):
  reverse(flo[b].begin() + 1, flo[b].end());
```

```
for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y)
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w = 0;
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 \mid \mid e_{delta}(g[xs][x]) < e_{delta}(g[b][x])
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
}
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  } else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true;
    else add_blossom(u, lca, v);
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \& !match[x]) pa[x] = 0, S[x] = 0, q_push(
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v \le n; ++v)
         if (g[u][v].w > 0 && st[u] != st[v]) {
           if (e_delta(g[u][v]) == 0) {
             if (on_found_edge(g[u][v])) return true;
           } else update_slack(u, st[v]);
        }
    int d = inf;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x]) {
         if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x]));
         else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x])
   1) / 2);
    for (int u = 1; u <= n; ++u) {</pre>
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
         if (S[st[b]] == 0) lab[b] += d * 2;
```

```
else if (S[st[b]] == 1) lab[b] -= d * 2;
          }
       q = queue<int>();
       for (int x = 1; x <= n_x; ++x)
  if (st[x] == x && slack[x] && st[slack[x]] != x &&</pre>
      e_delta(g[slack[x]][x]) == 0)
            if (on_found_edge(g[slack[x]][x])) return true;
        for (int b = n + 1; b \le n_x; ++b)
          if (st[b] == b \&\& S[b] == 1 \&\& lab[b] == 0)
      expand_blossom(b);
     return false;
   pair<long long, int> solve() {
     memset(match + 1, 0, sizeof(int) * n);
     n_x = n;
     int n_matches = 0;
     long long tot_weight = 0;
     for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear();</pre>
     int w_max = 0;
     for (int u = 1; u <= n; ++u)
       for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
          w_max = max(w_max, g[u][v].w);
     for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
     while (matching()) ++n_matches;
     for (int u = 1; u <= n; ++u)
  if (match[u] && match[u] < u)</pre>
          tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
  }
   void add_edge(int ui, int vi, int wi) { g[ui][vi].w = g[vi][
      ui].w = wi; }
   void init(int _n) {
     for (int u = 1; u <= n; ++u)</pre>
        for (int v = 1; v <= n; ++v)
          g[u][v] = edge(u, v, 0);
1 };
```

### 2.8 Flow Model

- Maximum/Minimum flow with lower/upper bound from s to t
  - 1. Construct super source S and sink T
  - 2. For each edge (x,y,l,u), connect  $x \to y$  with capacity u-l
  - For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
     If in(v) > 0, connect S > v with consists in(v) otherwise.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v)
    - To maximize, connect  $t \to s$  with capacity  $\infty$ , and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge  $(y \to x \text{ if } (x, y) \in M, x \to y \text{ otherwise})$
  - 2. DFS from unmatched vertices in X
  - 3.  $x \in X$  is chosen iff x is unvisited
  - 4.  $y \in Y$  is chosen iff y is visited
- Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x,y,c), connect  $x\to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v)<0, connect  $v\to T$  with (cost,cap)=(0,-d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K

- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v, v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G,$  connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', connected  $u' \to v'$  with weight w((u,v)).
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. The minimum weight perfect matching on the resulting graph is equivalent to the minimum weight edge cover of G by changing each (v,v') to the cheapest edge of v.

### 3 Data Structure

### 3.1 < ext/pbds >

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>, rb_tree_tag,
    tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree_set s
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22); assert(*s.find_by_order(1)
     == 71);
  assert(s.order_of_key(22) == 0); assert(s.order_of_key(71) ==
     1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71); assert(s.order_of_key(71)
     == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;</pre>
  return 0;
```

### 3.2 Li Chao Tree

```
namespace lichao {
struct line {
  long long a, b;
  line(): a(0), b(0) {}
  line(long long a, long long b): a(a), b(b) {}
  long long operator()(int x) const { return a * x + b; }
line st[maxc * 4];
int sz, lc[maxc * 4], rc[maxc * 4];
int gnode() {
  st[sz] = line(1e9, 1e9);
  lc[sz] = -1, rc[sz] = -1;
  return sz++;
void init() {
 sz = 0;
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
  bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
```

```
if (mcp) swap(st[o], tl);
  if (r - l == 1) return;
  if (lcp != mcp) {
    if (lc[o] == -1) lc[o] = gnode();
    add(1, (1 + r) / 2, tl, lc[o]);
  } else {
    if (rc[o] == -1) rc[o] = gnode();
    add((l + r) / 2, r, tl, rc[o]);
long long query(int l, int r, int x, int o) {
  if (r - l == 1) return st[o](x);
  if (x < (l + r) / 2) {
    if (lc[o] == -1) return st[o](x);
    return min(st[o](x), query(l, (l + r) / 2, x, lc[o]));
  } else {
    if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((l + r) / 2, r, x, rc[o]));
j }}
```

### 3.3 Link-Cut Tree

node \*ch[2], \*fa, \*pfa;

struct node {

```
int sum, v, rev, id;
node(int s, int id): id(id), v(s), sum(s), rev(0), fa(nullptr
   ), pfa(nullptr) {
  ch[0] = nullptr;
  ch[1] = nullptr;
int relation() {
 return this == fa->ch[0] ? 0 : 1;
}
void push() {
  if (!rev) return;
  swap(ch[0], ch[1]);
if (ch[0]) ch[0]->rev ^= 1;
  if (ch[1]) ch[1]->rev ^= 1;
  rev = 0:
void pull() {
  sum = v;
  if (ch[0]) sum += ch[0]->sum;
  if (ch[1]) sum += ch[1]->sum;
void rotate()
  if (fa->fa) fa->fa->push();
  fa->push(), push(), swap(pfa, fa->pfa);
  int d = relation();
  node *t = fa;
  if (t->fa) t->fa->ch[t->relation()] = this;
  fa = t->fa, t->ch[d] = ch[d ^ 1];
  if (ch[d \land 1]) ch[d \land 1] -> fa = t;
  ch[d \land 1] = t, t \rightarrow fa = this;
  t->pull(), pull();
}
void splay() {
  while (fa) -
    if (!fa->fa) {
      rotate();
      continue
    fa->fa->push(), fa->push();
    if (relation() == fa->relation()) fa->rotate();
    else rotate(), rotate();
void evert() { access(), splay(), rev ^= 1; }
void expose() {
  splay(), push();
  if (ch[1]) {
    ch[1]->fa = nullptr, ch[1]->pfa = this;
    ch[1] = nullptr, pull();
bool splice() {
  splay();
  if (!pfa) return false;
  pfa->expose(), pfa->ch[1] = this, fa = pfa;
  pfa = nullptr, fa->pull();
  return true;
void access() {
```

```
expose():
    while (splice());
  int query() { return sum; }
};
namespace lct {
node *sp[maxn];
void make(int u, int v) {
  // create node with id u and value v
  sp[u] = new node(v, u);
}
void link(int u, int v) {
  // u become v's parent
  sp[v]->evert();
  sp[v]->pfa = sp[u];
void cut(int u, int v) {
  // u was v's parent
  sp[u]->evert();
  sp[v]->access(), sp[v]->splay(), sp[v]->push();
  sp[v] \rightarrow ch[0] \rightarrow fa = nullptr;
  sp[v]->ch[0] = nullptr;
  sp[v]->pull();
}
void modify(int u, int v) {
  sp[u]->splay();
  sp[u] -> v = v
  sp[u]->pull();
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
  return sp[v]->query();
int find(int u) {
  sp[u]->access();
  sp[u]->splay();
  node *p = sp[u];
  while (true) {
    p->push();
    if (p->ch[0]) p = p->ch[0];
    else break;
  return p->id;
```

## 4 Graph

### 4.1 Heavy-Light Decomposition

```
| void dfs(int x, int p) {
| dep[x] = ~p ? dep[p] + 1 : dep[x];
   sz[x] = 1;
   to[x] = -1;
   fa[x] = p;
   for (const int &u : g[x]) {
     if (u == p) continue;
     dfs(u, x);
     sz[x] += sz[u];
     if (to[x] == -1 \mid \mid sz[to[x]] < sz[u]) to[x] = u;
void hld(int x, int t) {
   static int tk = 0;
   fr[x] = t;
   dfn[x] = tk++;
   if (!~to[x]) return;
   hld(to[x], t);
   for (const int &u : g[x]) {
     if (u == fa[x] || u == to[x]) continue;
     hld(u, u);
  }
vector<pair<int, int>> get(int x, int y) {
   int fx = fr[x], fy = fr[y];
   vector<pair<int, int>> res;
   while (fx != fy) {
     if (dep[fx] < dep[fy]) {</pre>
       swap(fx, fy);
       swap(x, y);
     res.emplace_back(dfn[fx], dfn[x] + 1);
     x = fa[fx];
```

```
fx = fr[x];
}
| res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn[y]) +
    1);
| int lca = (dep[x] < dep[y] ? x : y);
| return res;
|}</pre>
```

### 4.2 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
sz[now] = 1; mx[now] = 0;
   for (int u : G[now]) if (!v[u]) {
    get_center(u);
    mx[now] = max(mx[now], sz[u]);
    sz[now] += sz[u];
void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
  v[now] = true;
   for (auto u : G[now]) if (!v[u.first]) {
    get_dis(u, d, len + u.second);
}
void dfs(int now, int fa, int d) {
  get_center(now);
   int c = -1;
   for (int i : vtx) {
    if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx.size()</pre>
     / 2) c = i;
    v[i] = false;
  get_dis(c, d, 0);
  for (int i : vtx) v[i] = false;
  v[c] = true; vtx.clear();
  dep[c] = d; p[c] = fa;
  for (auto u : G[c]) if (u.first != fa && !v[u.first]) {
    dfs(u.first, c, d + 1);
į }
```

### 4.3 Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
long long d[1003][1003],dp[1003][1003];
pair<long long,long long> MMWC(){
 memset(dp,0x3f,sizeof(dp));
  for(int i=1;i<=n;++i)dp[0][i]=0;</pre>
  for(int i=1;i<=n;++i){</pre>
    for(int j=1;j<=n;++j){</pre>
      for(int k=1;k<=n;++k){</pre>
        dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
   }
  long long au=1ll<<31,ad=1;
  for(int i=1;i<=n;++i){</pre>
    if(dp[n][i]==0x3f3f3f3f3f3f3f3f)continue;
    long long u=0, d=1;
    for(int j=n-1; j>=0; -- j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
        u=dp[n][i]-dp[j][i];
        d=n-j;
      }
    if(u*ad<au*d)au=u,ad=d;</pre>
 long long g=__gcd(au,ad);
  return make_pair(au/g,ad/g);
```

### 4.4 Minimum Steiner Tree

```
| namespace steiner {
|// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
|// z[i] = the weight of the i-th vertex
| const int maxn = 64, maxk = 10;
```

```
const int inf = 1e9:
 int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[maxn];</pre>
void init(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = inf;
     z[i] = 0;
     w[i][i] = 0;
  }
void add_edge(int x, int y, int d) {
  w[x][y] = min(w[x][y], d);
  w[y][x] = min(w[y][x], d);
void build(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) {
       w[i][j] += z[i];
       if (i != j) w[i][j] += z[j];
   for (int k = 0; k < n; ++k) {
     for (int i = 0; i < n; ++i) {
       for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j], w[i][k]
      ] + w[k][j] - z[k]);
  }
}
int solve(int n, vector<int> mark) {
   build(n);
   int k = (int)mark.size();
   assert(k < maxk);</pre>
   for (int s = 0; s < (1 << k); ++s) {
     for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
   for (int i = 0; i < n; ++i) dp[0][i] = 0;
   for (int s = 1; s < (1 << k); ++s) {
     if (__builtin_popcount(s) == 1) {
       int x = __builtin_ctz(s);
       for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]][i];
       continue;
     for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1)) {
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s \land sub][i] -
     z[i]);
       }
     for (int i = 0; i < n; ++i) {
       off[i] = inf;
       for (int j = 0; j < n; ++j) off[i] = min(off[i], dp[s][j]
      + w[j][i] - z[j]);
     for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i], off[i
     1);
   int res = inf:
   for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k) - 1][i
     ]);
   return res;
| }}
```

### 4.5 Directed Minimum Spanning Tree

```
template <tvpename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {
      for(int j = 0; j < maxn; ++j) g[i][j] = inf;
      vis[i] = inc[i] = false;
    }
  }
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
  T operator()(int root, int _n) {
    if (dfs(root) != n) return -1;
    T ans = 0;
    while (true) {
      for (int i = 1; i \le n; ++i) fw[i] = inf, fr[i] = i;
      for (int i = 1; i <= n; ++i) if (!inc[i]) {
        for (int j = 1; j \le n; ++j) {
```

```
if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
             fw[i] = g[j][i];
             fr[i] = j;
       }
       int x = -1;
       for (int i = 1; i <= n; ++i) if (i != root && !inc[i]) {
         int j = i, c = 0;
         while (j != root && fr[j] != i && c <= n) ++c, j = fr[j
         if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
         for (int i = 1; i <= n; ++i) if (i != root && !inc[i])</pre>
      ans += fw[i];
         return ans;
       int y = x;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
       do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true; }
      while (y != x);
       inc[x] = false;
       for (int k = 1; k \le n; ++k) if (vis[k]) {
         for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
           if (g[j][k] < \inf \&\& g[j][k] - fw[k] < g[j][x]) g[j][
      x] = g[j][k] - fw[k];
         }
       }
     return ans;
   int dfs(int now) {
     int r = 1;
     vis[now] = true;
     for (int i = 1; i <= n; ++i) if (g[now][i] < inf && !vis[i
     ]) r += dfs(i);
     return r;
|};
```

### 4.6 Maximum Clique

```
struct MaxClique {
 // change to bitset for n > 64.
 int n, deg[maxn];
 uint64_t adj[maxn], ans;
  vector<pair<int, int>> edge;
  void init(int n_) {
    fill(adj, adj + n, 0ull);
   fill(deg, deg + n, 0);
   edge.clear();
  void add_edge(int u, int v) {
   edge.emplace_back(u, v);
   ++deg[u], ++deg[v];
 vector<int> operator()() {
   vector<int> ord(n);
   iota(ord.begin(), ord.end(), 0);
   sort(ord.begin(), ord.end(), [&](int u, int v) { return deg
    [u] < deg[v]; });
   vector<int> id(n);
    for (int i = 0; i < n; ++i) id[ord[i]] = i;
    for (auto e : edge) {
      int u = id[e.first], v = id[e.second];
      adj[u] |= (1ull << v);
      adj[v] = (1ull \ll u);
   uint64_t r = 0, p = (1ull << n) - 1;
   ans = 0;
   dfs(r, p);
   vector<int> res;
    for (int i = 0; i < n; ++i) {
      if (ans >> i & 1) res.push_back(ord[i]);
    return res;
#define pcount __builtin_popcountll
 void dfs(uint64_t r, uint64_t p) {
   if (p == 0) {
```

```
if (pcount(r) > pcount(ans)) ans = r;
       return;
     }
     if (pcount(r | p) <= pcount(ans)) return;</pre>
     int x = __builtin_ctzll(p & -p);
     uint64_t c = p & \simadj[x];
     while (c > 0) {
       // bitset._Find_first(); bitset._Find_next();
       x = \__builtin_ctzll(c \& -c);
       r = (1ull \ll x);
       dfs(r, p & adj[x]);
       r \&= \sim (1ull << x);
       p \&= \sim (1ull << x);
       c \wedge = (1ull \ll x);
  }
};
```

### 4.7 Tarjan's Algorithm

```
void dfs(int x, int p) {
  dfn[x] = low[x] = tk++;
   int ch = 0;
   st.push(x); // bridge
   for (auto e : g[x]) if (e.first != p) {
     if (!ins[e.second]) { // articulation point
       st.push(e.second);
       ins[e.second] = true;
     if (~dfn[e.first]) {
       low[x] = min(low[x], dfn[e.first]);
     dfs(u.first, x);
     if (low[u.first] >= low[x]) { // articulation point
       cut[x] = true;
       while (true) {
         int z = st.top(); st.pop();
         bcc[z] = sz;
         if (z == e.second) break;
       SZ++;
     }
  if (ch == 1 \&\& p == -1) cut[x] = false;
  if (dfn[x] == low[x]) { // bridge
     while (true) {
       int z = st.top(); st.pop();
       bcc[z] = sz;
       if (z == x) break;
  }
1}
```

### 4.8 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[maxn], val[
     maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1);
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
  fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
fill(rp, rp + n, -1);
  fill(dom, dom + n, -1);
  tk = 0;
  for (int i = 0; i < n; ++i) {
    g[i].clear();
    r[i].clear();
    rdom[i].clear();
 }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  tk++:
 for (int u : g[x]) {
```

```
if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
 }
}
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
 int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
 fa[x] = p;
return c ? p : val[x];
vector<int> build(int s, int n) {
  // return the father of each node in the dominator tree
  // p[i] = -2 if i is unreachable from s
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int u : r[i]) sdom[i] = min(sdom[i], sdom[find(u)]);
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
 }
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i]) dom[i] =
    dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
```

### 4.9 Virtual Tree

```
void VirtualTree(vector<int> v) {
  static int stk[kN];
  int sz = 0;
 sort(v.begin(), v.end(), [&](int i, int j) { return dfn[i] <</pre>
    dfn[j]; });
 stk[sz++] = 0;
  for (int u : v) {
    if (u == 0) continue;
    int p = LCA(u, stk[sz - 1]);
    if (p != stk[sz - 1]) {
      while (sz \ge 2 \&\& dfn[p] < dfn[stk[sz - 2]]) {
        AddEdge(stk[sz - 2], stk[sz - 1]);
      if (sz \ge 2 \& dfn[p] > dfn[stk[sz - 2]]) {
        AddEdge(p, stk[sz - 1]);
        stk[sz - 1] = p;
      } else {
        AddEdge(p, stk[--sz]);
   stk[sz++] = u;
  for (int i = 0; i < sz - 1; ++i) AddEdge(stk[i], stk[i + 1]);</pre>
```

#### 4.10 Vizing's Theorem

```
namespace vizing { // returns edge coloring in adjacent matrix
    G. 1 - based
int C[kN][kN], G[kN][kN];
void clear(int N) {
    for (int i = 0; i <= N; i++) {
        for (int j = 0; j <= N; j++) C[i][j] = G[i][j] = 0;
    }
}
void solve(vector<pair<int, int>> &E, int N, int M) {
    int X[kN] = {}, a;
    auto update = [&](int u) {
        for (X[u] = 1; C[u][X[u]]; X[u]++);
    };
    auto color = [&](int u, int v, int c) {
        int p = G[u][v];
        G[u][v] = G[v][u] = c;
        C[u][c] = v, C[v][c] = u;
        C[u][p] = C[v][p] = 0;
```

```
if (p) X[u] = X[v] = p;
      else update(u), update(v);
      return p;
   };
   auto flip = [&](int u, int c1, int c2) {
      int p = C[u][c1];
     swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
if (!C[u][c1]) X[u] = c1;
      if (!C[u][c2]) X[u] = c2;
      return p;
   for (int i = 1; i <= N; i++) X[i] = 1;
   for (int t = 0; t < E.size(); t++) {</pre>
      int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c
      = c0, d;
      vector<pair<int, int>> L;
      int vst[kN] = {}
      while (!G[u][v0]) {
        L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a = (int)L.size() - 1; a >= 0; a--) c
      = color(u, L[a].first, c);
else if (!C[u][d]) for (a = (int)L.size() - 1; a >= 0; a
      --) color(u, L[a].first, L[a].second);
        else if (vst[d]) break;
        else vst[d] = 1, v = C[u][d];
      if (!G[u][v0]) {
        for (; v; v = flip(v, c, d), swap(c, d));
        if (C[u][c0]) {
          for (a = (int)L.size() - 2; a >= 0 && L[a].second != c;
       a--);
          for (; a \ge 0; a - -) color(u, L[a].first, L[a].second);
        } else t--:
   }
}}
```

### 4.11 System of Difference Constraints

Given m constrains on n variables  $x_1, x_2, \ldots, x_n$  of form  $x_i - x_j \leq w$  (resp,  $x_i - x_j \geq w$ ), connect  $i \to j$  with weight w. Then connect  $0 \to i$  for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to  $x_i$ .

### 5 String

#### 5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  // f[i] = length of the longest prefix (excluding s[0:i])
     such that it coincides with the suffix of s[0:i] of the
     same length
   // i + 1 - f[i] is the length of the smallest recurring
     period of s[0:i]
   int k = 0;
   for (int i = 1; i < (int)s.size(); ++i) {</pre>
    while (k > 0 \& s[i] != s[k]) k = f[k - 1];
     if (s[i] == s[k]) ++k;
    f[i] = k;
  }
  return f;
}
vector<int> search(const string &s, const string &t) {
  // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), res;
   for (int i = 0, k = 0; i < (int)s.size(); ++i) {
    while (k > 0 \& (k == (int)t.size() || s[i] != t[k])) k = f
     [k - 1];
     if (s[i] == t[k]) ++k;
     if (k == (int)t.size()) res.push_back(i - t.size() + 1);
   return res;
}
```

### 5.2 Z Algorithm

```
|int z[maxn];
|// z[i] = LCP of suffix i and suffix 0
```

```
void z_function(const string& s) {
    memset(z, 0, sizeof(z));
    z[0] = (int)s.length();
    int l = 0, r = 0;
    for (int i = 1; i < s.length(); ++i) {
        z[i] = max(0, min(z[i - l], r - i + 1));
        while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
            l = i; r = i + z[i];
            ++z[i];
        }
    }
}</pre>
```

### 5.3 Manacher's Algorithm

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for (int i = 0; i < s.length(); ++i) t += s[i], t += '.';
   int l = 0, r = 0, ans = 0;
   for (int i = 1; i < t.length(); ++i) {
        z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
        while (i - z[i] >= 0 && i + z[i] < t.length() && t[i - z[i] ] = t[i + z[i]]) ++z[i];
        if (i + z[i] > r) r = i + z[i], l = i;
        }
        for (int i = 1; i < t.length(); ++i) ans = max(ans, z[i] - 1)
        ;
        return ans;
}</pre>
```

#### 5.4 Aho-Corasick Automaton

```
struct AC {
  static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn][26], f[
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
    cnt[sz] = 0;
    return sz++;
  void init() {
    root = gnode();
  int add(const string &s) {
    int now = root;
for (int i = 0; i < s.length(); ++i) {</pre>
      if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a'] =
     gnode();
      now = ch[now][s[i] - 'a'];
    ed[now] = 1;
    return now;
  void build_fail() {
    ql = qr = 0; q[qr++] = root;
    while (ql < qr) {</pre>
      int now = q[ql++];
      for (int i = 0; i < 26; ++i) if (ch[now][i] != -1) {
        int p = ch[now][i], fp = f[now];
        while (fp != -1 \& ch[fp][i] == -1) fp = f[fp];
        int pd = fp != -1 ? ch[fp][i] : root;
        f[p] = pd;
        el[p] = ed[pd] ? pd : el[pd];
        q[qr++] = p;
      }
   }
  void build(const string &s) {
    build_fail();
    int now = root:
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (now != -1 && ch[now][s[i] - 'a'] == -1) now = f[
     nowl:
      now = now != -1 ? ch[now][s[i] - 'a'] : root;
      ++cnt[now];
```

```
for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] += cnt[q[i]]
      ]];
   long long solve(int n) {
     build_fail();
     vector<vector<long long>> dp(sz, vector<long long>(n + 1,
      0));
     for (int i = 0; i < sz; ++i) dp[i][0] = 1;
     for (int i = 1; i <= n; ++i) {
       for (int j = 0; j < sz; ++j) {
         for (int k = 0; k < 2; ++k) {
  if (ch[j][k] != -1) {</pre>
              if (!ed[ch[j][k]])
                dp[j][i] += dp[ch[j][k]][i - 1];
           } else {
              int z = f[j];
              while (z != root \&\& ch[z][k] == -1) z = f[z];
              int p = ch[z][k] == -1 ? root : ch[z][k];
              if (ch[z][k] == -1 \mid \mid !ed[ch[z][k]]) dp[j][i] += dp
      [p][i - 1];
           }
         }
       }
     return dp[0][n];
  }
};
```

### 5.5 Suffix Automaton

```
struct SAM {
  static const int maxn = 5e5 + 5;
  int nxt[maxn][26], to[maxn], len[maxn];
  int root, last, sz;
  int gnode(int x) {
    for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
    to[sz] = -1;
    len[sz] = x;
    return sz++;
  void init() {
    root = gnode(0);
    last = root;
  void push(int c) {
    int cur = last;
    last = gnode(len[last] + 1);
    for (; ~cur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
     1 = last;
    if (cur == -1) return to[last] = root, void();
    int link = nxt[cur][c];
    if (len[link] == len[cur] + 1) return to[last] = link, void
     ();
    int tlink = gnode(len[cur] + 1);
    for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
     ][c] = tlink;
    for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];</pre>
    to[tlink] = to[link];
    to[link] = tlink;
    to[last] = tlink;
  void add(const string &s) {
    for (int i = 0; i < s.size(); ++i) push(s[i] - 'a');</pre>
  bool find(const string &s) {
    int cur = root;
    for (int i = 0; i < s.size(); ++i) {</pre>
      cur = nxt[cur][s[i] - 'a'];
      if (cur == -1) return false;
    return true:
  }
  int solve(const string &t) {
    int res = 0, cnt = 0;
    int cur = root;
    for (int i = 0; i < t.size(); ++i) {</pre>
      if (~nxt[cur][t[i] - 'a']) {
        ++cnt;
        cur = nxt[cur][t[i] - 'a'];
      } else {
        for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
        if (\simcur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a'
     ];
```

```
else cnt = 0, cur = root;

res = max(res, cnt);

return res;
}

};
```

### 5.6 Suffix Array

```
namespace sfxarray {
 bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2], x[maxn], p[maxn],
     q[maxn * 2];
   sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
     suffix.
 // hi[i]: longest common prefix of suffix sa[i] and suffix sa[i
       - 1].
 void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
 void induce(int *sa, int *c, int *s, bool *t, int n, int z) {
  memcpy(x + 1, c, sizeof(int) * (z - 1));
   for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[x[
     s[sa[i] - 1]]++] = sa[i] - 1;
   memcpy(x, c, sizeof(int) * z);
   for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
}
void sais(int *s, int *sa, int *p, int *q, bool *t, int *c, int
      n, int z) {
   bool uniq = t[n - 1] = true;
   int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
   memset(c, 0, sizeof(int) * z);
   for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
   for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
   for (int i = n - 2; i \ge 0; --i) t[i] = (s[i] == s[i + 1] ? t
     [i + 1] : s[i] < s[i + 1]);
   pre(sa, c, n, z);
   for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i - 1]) sa[--
     x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
   for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[i]
     bool neq = last < 0 | | memcmp(s + sa[i], s + last, (p[q[sa[
     i]] + 1] - sa[i]) * sizeof(int));
     ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p[
     nsa[i]];
   induce(sa, c, s, t, n, z);
 void build(const string &s) {
   for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
   _s[(int)s.size()] = 0; // s shouldn't contain 0
   sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
   for (int i = 0; i < (int)s.size(); ++i) sa[i] = sa[i + 1];</pre>
   for (int i = 0; i < (int)s.size(); ++i) rev[sa[i]] = i;</pre>
   int ind = 0; hi[0] = 0;
   for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
       ind = 0;
       continue;
     while (i + ind < (int)s.size() && s[i + ind] == s[sa[rev[i]]]
       - 1] + ind]) ++ind;
     hi[rev[i]] = ind ? ind-- : 0;
| }}
```

### 5.7 Palindromic Tree

```
|struct PalindromicTree {
| int link[kN], len[kN], dp[kN], nxt[kN][26], sz, sf;
```

```
int gnode(int l, int fl = -1) {
     len[sz] = l;
link[sz] = fl;
     fill(nxt[sz], nxt[sz] + 26, -1);
     return sz++;
   void Init() {
     sz = 0:
     sf = 1;
     gnode(-1, 0);
     gnode(0, 0);
   void Push(const string &s, int pos) {
     int cur = sf, z = s[pos] - 'a'
     while (pos - 1 - len[cur] < 0 | | s[pos - 1 - len[cur]] != s
      [pos]) cur = link[cur];
     if (nxt[cur][z] != -1) {
       sf = nxt[cur][z];
     } else {
       int ch = gnode(len[cur] + 2);
       nxt[cur][z] = sf = ch;
       if (len[ch] == 1) {
          link[ch] = 1;
       } else {
          cur = link[cur];
while (pos - 1 - len[cur] < 0 | | s[pos - 1 - len[cur]]</pre>
      != s[pos]) cur = link[cur];
          link[ch] = nxt[cur][z];
     dp[sf] += 1;
   long long Build(const string &s) {
     for (int i = 0; i < s.size(); ++i) Push(s, i);
for (int i = sz - 1; i >= 0; --i) dp[link[i]] += dp[i];
     long long res = 0;
     for (int i = 0; i < sz; ++i) res = max(res, 1LL * dp[i] *
      len[i]);
     return res;
} plt;
```

#### 5.8 Circular LCS

```
string s1, s2;
int n, m;
int dp[kN * 2][kN];
int nxt[kN * 2][kN];
void reroot(int px) {
  int py = 1;
  while (py <= m && nxt[px][py] != 2) py++;</pre>
  nxt[px][py] = 1;
  while (px < 2 * n \&\& py < m) {
    if (nxt[px + 1][py] == 3) px++, nxt[px][py] = 1;
    else if (nxt[px + 1][py + 1] == 2) px++, py++, nxt[px][py]
     = 1;
    else py++;
  while (px < 2 * n \& nxt[px + 1][py] == 3) px++, nxt[px][py]
int track(int x, int y, int e) { // use this routine to find
     LCS as string
  int ret = 0;
  while (y != 0 && x != e) {
    if (nxt[x][y] == 1) y--
    else if (nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
else if (nxt[x][y] == 3) x--;
  return ret;
int solve(string a, string b) {
  n = a.size(), m = b.size();
s1 = "#" + a + a, s1 = '#' + b;
  for (int i = 0; i \le 2 * n; i++) {
    for (int j = 0; j <= m; j++) {
       if (j == 0) { nxt[i][j] = 3; continue;
       if (i == 0) { nxt[i][j] = 1; continue; }
       dp[i][j] = -1;
       if (dp[i][j] < dp[i][j - 1]) dp[i][j] = dp[i][j - 1], nxt
     [i][j] = 1;
       if (dp[i][j] < dp[i - 1][j - 1] + (s1[i] == s2[j])) dp[i
     [j] = dp[i - 1][j - 1] + (s1[i] == s2[j]), nxt[i][j] = 2;
       if (dp[i][j] < dp[i - 1][j]) dp[i][j] = dp[i - 1][j], nxt</pre>
     [i][j] = 3;
```

```
}
}
int ret = dp[n][m];
for (int i = 1; i < n; i++) reroot(i), ret = max(ret, track(n + i, m, i));
return ret;
}</pre>
```

### 5.9 Lexicographically Smallest Rotation

```
| string rotate(const string &s) {
| int n = s.length();
| string t = s + s;
| int i = 0, j = 1;
| while (i < n && j < n) {
| int k = 0;
| while (k < n && t[i + k] == t[j + k]) ++k;
| if (t[i + k] <= t[j + k]) j += k + 1;
| else i += k + 1;
| if (i == j) ++j;
| }
| int pos = (i < n ? i : j);
| return t.substr(pos, n);
|}</pre>
```

### 6 Math

### 6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(re + rhs.
    re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(re - rhs.
    re, im - rhs.im); }
  cplx operator*(const cplx &rhs) const { return cplx(re * rhs.
    re - im * rhs.im, re * rhs.im + im * rhs.re); }
  cplx conj() const { return cplx(re, -im); }
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
  for (int i = 0; i \leftarrow maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi * i / maxn))
    maxn));
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    int x = 0:
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j \& 1) << (z
     - j);
    if (x > i) swap(v[x], v[i]);
 }
}
void fft(vector<cplx> &v, int n) {
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
        cplx x = v[i + z + k] * omega[maxn / s * k];
        v[i + z + k] = v[i + k] - x;
        v[i + k] = v[i + k] + x;
   }
 }
}
void ifft(vector<cplx> &v, int n) {
  fft(v, n);
  reverse(v.begin() + 1, v.end());
  for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n, 0);
vector<long long> convolution(const vector<int> &a, const
    vector<int> &b) {
  // Should be able to handle N <= 10^5, C <= 10^4
  int sz = 1;
```

```
while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
   vector<cplx> v(sz);
   for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;</pre>
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
   fft(v, sz);
   for (int i = 0; i \le sz / 2; ++i) {
     int j = (sz - i) & (sz - 1);
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()) * cplx
      (0, -0.25);
     if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i].conj
()) * cplx(0, -0.25);
     v[i] = x;
  }
  ifft(v, sz);
  vector<long long> c(sz);
   for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
vector<int> convolution_mod(const vector<int> &a, const vector<</pre>
     int> &b, int p) {
   int sz = 1;
   while (sz < (int)a.size() + (int)b.size() - 1) sz <<= 1;</pre>
   vector<cplx> fa(sz), fb(sz);
   for (int i = 0; i < (int)a.size(); ++i)</pre>
     fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
   for (int i = 0; i < (int)b.size(); ++i)</pre>
     fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
   fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
   for (int i = 0; i \leftarrow (sz >> 1); ++i) {
     int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx b1 = (fb[i] + fb[j].conj()) * r3;
     cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
       cplx c1 = (fa[j] + fa[i].conj());
       cplx c2 = (fa[j] - fa[i].conj()) * r2;
       cplx d1 = (fb[j] + fb[i].conj()) * r3;
       cplx d2 = (fb[j] - fb[i].conj()) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \% p;
   return res;
}}
```

#### 6.2 Number Theoretic Transform

```
template <long long mod, long long root>
struct NTT {
  vector<long long> omega;
  NTT() {
    omega.resize(maxn + 1);
    long long x = fpow(root, (mod - 1) / maxn);
    omega[0] = 111;
    for (int i = 1; i <= maxn; ++i)
  omega[i] = omega[i - 1] * x % mod;</pre>
  void bitrev(vector<long long> &v, int n) {
    int z = __builtin_ctz(n) - 1;
    for (int i = 0; i < n; ++i) {
      int x = 0;
       for (int j = 0; j \ll z; ++j) x \sim (i \gg j \& 1) \ll (z - j)
       if (x > i) swap(v[x], v[i]);
    }
  }
  void ntt(vector<long long> &v, int n) {
    bitrev(v, n);
```

```
for (int s = 2; s <= n; s <<= 1) {
       int z = s \gg 1;
       for (int i = 0; i < n; i += s) {
         for (int k = 0; k < z; ++k) {
           long long x = v[i + k + z] * omega[maxn / s * k] %
           v[i + k + z] = (v[i + k] + mod - x) \% mod;
           (v[i + k] += x) \% = mod;
         }
       }
    }
  }
  void intt(vector<long long> &v, int n) {
     ntt(v, n);
     for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
     long long inv = fpow(n, -1);
     for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;</pre>
  vector<long long> operator()(vector<long long> a, vector<long</pre>
       long> b) {
     int sz = 1;
     while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
     while (a.size() < sz) a.push_back(0);
while (b.size() < sz) b.push_back(0);</pre>
     ntt(a, sz), ntt(b, sz);
     vector<long long> c(sz);
     for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] % mod;
     intt(c, sz);
     return c:
  }
};
vector<long long> convolution(vector<long long> a, vector<long</pre>
      long> b) {
  NTT<mod1, root1> conv1;
NTT<mod2, root2> conv2;
   vector<long long> pa(a.size()), pb(b.size());
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] \% mod1)
     + mod1) % mod1;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] \% mod1)
     + mod1) % mod1;
   vector<long long> c1 = conv1(pa, pb);
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] \% mod2)
     + mod2) % mod2;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] \% mod2)
     + mod2) % mod2;
   vector<long long> c2 = conv2(pa, pb);
  long long x = conv2.fpow(mod1, -1);
  long long y = conv1.fpow(mod2, -1);
  long long prod = mod1 * mod2;
   vector<long long> res(c1.size());
  for (int i = 0; i < c1.size(); ++i) {</pre>
     long long z = ((ull)fmul(c1[i] * mod2 % prod, y, prod) + (
     ull)fmul(c2[i] * mod1 % prod, x, prod)) % prod;
     if (z >= prod / 2) z -= prod;
     res[i] = z;
   return res;
| }
```

### 6.2.1 NTT Prime List

Prime	Root	Prime	Root
7681	17	167772161	3
12289	11	104857601	3
40961	3	985661441	3
65537	3	998244353	3
786433	10	1107296257	10
5767169	3	2013265921	31
7340033	3	2810183681	11
23068673	3	2885681153	3
469762049	3	605028353	3

#### 6.3 Polynomial Division

```
vector<int> inverse(const vector<int> &v, int n) {
  vector<int> q(1, fpow(v[0], mod - 2));
  for (int i = 2; i <= n; i <<= 1) {
    vector<int> fv(v.begin(), v.begin() + i);
    vector<int> fq(q.begin(), q.end());
    fv.resize(2 * i), fq.resize(2 * i);
    ntt(fq, 2 * i), ntt(fv, 2 * i);
    for (int j = 0; j < 2 * i; ++j) {
        fv[j] = fv[j] * 1ll * fq[j] % mod * 1ll * fq[j] % mod;
    }
    intt(fv, 2 * i);</pre>
```

```
vector<int> res(i);
      for (int j = 0; j < i; ++j) {
  res[j] = mod - fv[j];</pre>
        if (j < (i >> 1)) (res[j] += 2 * q[j] % mod) %= mod;
     q = res;
   }
   return q;
}
vector<int> divide(const vector<int> &a. const vector<int> &b)
   // leading zero should be trimmed
   int n = (int)a.size(), m = (int)b.size();
   int k = 2;
   while (k < n - m + 1) k <<= 1;
   vector<int> ra(k), rb(k);
   for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - i - 1];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - i - 1];
vector<int> rbi = inverse(rb, k);
   vector<int> res = convolution(rbi, ra);
   res.resize(n - m + 1);
   reverse(res.begin(), res.end());
   return res;
13
```

### 6.4 Polynomial Square Root

```
// Find G(x) such that G^2(x) = F(x) \pmod{x^{N+1}} vector<int> solve(vector<int> b, int n) {
   if (n == 1) return {sqr[b[0]]};
   vector<int> h = solve(b, n >> 1); h.resize(n);
   vector<int> c = inverse(h, n);
   h.resize(n \ll 1); c.resize(n \ll 1);
   vector<int> res(n << 1);</pre>
   conv.ntt(h, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) b[i] = 0;
   conv.ntt(b, n << 1);
   conv.ntt(c, n << 1);</pre>
   for (int i = 0; i < (n << 1); ++i) res[i] = 1ll * (h[i] + 1ll
       * c[i] * b[i] % mod) % mod * inv2 % mod;
   conv.intt(res, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) res[i] = 0;
   return res;
i }
```

#### 6.5 Multipoint Evaluation

```
struct MultiEval {
   MultiEval *lc, *rc;
   vector<int> p, ml;
   // v is the points to be queried
   MultiEval(const vector<int> &v, int l, int r) : lc(nullptr),
      rc(nullptr) {
     if (r - 1 <= 64) {
       p = vector<int>(v.begin() + l, v.begin() + r);
       ml.resize(1, 1);
       for (int x : p) ml = Multiply(ml, {kMod - x, 1});
       return;
     int m = (l + r) >> 1;
     lc = new MultiEval(v, l, m), rc = new MultiEval(v, m, r);
     ml = Multiply(lc->ml, rc->ml);
   // poly is the polynomial to be evaluated
   void Query(const vector<int> &poly, vector<int> &res, int l,
     int r) const {
if (r - l <= 64) {
       for (int x : p) {
         int s = 0, bs = 1;
         for (int i = 0; i < poly.size(); ++i) {
   (s += 1LL * bs * poly[i] % kMod) %= kMod;</pre>
            bs = 1LL * bs * x % kMod;
         }
         res.push_back(s);
       return;
     auto pol = Modulo(poly, ml);
     int m = (l + r) >> 1;
     lc->Query(pol, res, l, m), rc->Query(pol, res, m, r);
|};
```

### 6.6 Polynomial Interpolation

```
vector<int> Interp(const vector<int> &x, const vector<int> &y)
  vector<vector<int>>> v;
  int n = x.size();
  v.emplace_back(n);
  for (int i = 0; i < n; ++i) v[0][i] = \{\{kMod - x[i], 1\}\};
  while (v.back().size() > 1) {
    int n2 = v.back().size();
    vector<vector<int>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) f[i] = Multiply(v.back ()[2 * i], v.back()[2 * i + 1]); if (n2 & 1) f.back() = v.back().back();
    v.push_back(f);
  }
  vector<int> df(v.back()[0].size() - 1);
  for (int i = 0; i < df.size(); ++i) df[i] = 1LL * v.back()
  [0][i + 1] * (i + 1) % kMod;</pre>
  vector<int> s;
  MultiEval(x, 0, n).Query(df, s, 0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) res[i] = {1LL * y[i] * fpow(s[i],
       kMod - 2) % kMod};
  for (int p = 1; p < v.size(); ++p) {</pre>
    int n2 = v[p - 1].size();
    vector<vector<int>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) {
  auto a = Multiply(res[i * 2], v[p - 1][2 * i + 1]);
  auto b = Multiply(res[i * 2 + 1], v[p - 1][2 * i]);
       assert(a.size() == b.size());
       f[i].resize(a.size());
       for (int j = 0; j < a.size(); ++j) f[i][j] = (a[j] + b[j]
     ]) % kMod;
     if (n2 & 1) f.back() = res.back();
    res = f;
  return res[0];
```

### 6.7 Fast Walsh-Hadamard Transform

#### 6.7.1 XOR Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_0) tf(A_1))$
- $utf(A) = (utf(\frac{A_0 + A_1}{2}), utf(\frac{A_0 A_1}{2}))$

#### 6.7.2 OR Convolution

- $tf(A) = (tf(A_0), tf(A_0) + tf(A_1))$
- $utf(A) = (utf(A_0), utf(A_1) utf(A_0))$

### 6.7.3 AND Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_1))$
- $utf(A) = (utf(A_0) utf(A_1), utf(A_1))$

### 6.8 Simplex Algorithm

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn't exist
// return vector<double>(n, +inf) if the solution is unbounded
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
vector<double>> d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {
      if (i != r && j != s) d[i][j] -= d[r][j] * d[i][s] * inv;
    }
  }
  for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s] *= -inv;
  for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j] *= +inv;
```

```
d[r][s] = inv;
   swap(p[r], q[s]);
bool phase(int z) {
   int x = m + z;
   while (true) {
     int s = -1;
     for (int i = 0; i <= n; ++i) {</pre>
       if (!z && q[i] == -1) continue;
       if (s == -1 || d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;</pre>
       if (r == -1 \mid \mid d[i][n + 1] / d[i][s] < d[r][n + 1] / d[r]
     ][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
}
 vector<double> solve(const vector<vector<double>> &a, const
     vector<double> &b, const vector<double> &c) {
   m = b.size(), n = c.size();
   d = vector<vector<double>>(m + 2, vector<double>(n + 2));
   for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][
     n + 1] = b[i];
   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
   for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r
   if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) \mid | d[m + 1][n + 1] < -eps) return vector<
      double>(n, -inf);
     for (int i = 0; i < m; ++i) if (p[i] == -1) {
       int s = min_element(d[i].begin(), d[i].end() - 1) - d[i].
     begin();
       pivot(i, s);
    }
   if (!phase(0)) return vector<double>(n, inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] < n) x[p[i]] = d[i][n +
  1];
return x;
| }}
```

#### 6.8.1 Construction

Standard form: maximize  $\sum_{1 \leq i \leq n} c_i x_i$  such that for all  $1 \leq j \leq m$ ,  $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$  and  $x_i \geq 0$  for all  $1 \leq i \leq n$ .

- 1. In case of minimization, let  $c'_i = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \leq i \leq n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
  - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

### 6.9 Schreier-Sims Algorithm

```
| namespace schreier {
| int n;
| vector<vector<vector<int>>> bkts, binv;
| vector<vector<int>> lk;
| vector<int> operator*(const vector<int> &a, const vector<int> &
| b) {
| vector<int> res(a.size());
| for (int i = 0; i < (int)a.size(); ++i) res[i] = b[a[i]];
| return res;
| }
| vector<int> inv(const vector<int> &a) {
```

```
vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[a[i]] = i;
  return res;
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
  vector<int> p = g;
for (int i = 0; i < n; ++i) {
   assert(p[i] >= 0 && p[i] < (int)lk[i].size());</pre>
     int res = lk[i][p[i]];
     if (res == -1) {
       if (add) {
         bkts[i].push_back(p);
         binv[i].push_back(inv(p));
         lk[i][p[i]] = (int)bkts[i].size() - 1;
       return i:
      = p * binv[i][res];
    р
  return -1;
bool inside(const vector<int> &g) { return filter(g, false) ==
      -1; }
void solve(const vector<vector<int>> &gen, int _n) {
  n = _n;
  bkts.clear(), bkts.resize(n);
  binv.clear(), binv.resize(n);
  lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
  for (int i = 0; i < n; ++i) {
    lk[i].resize(n, -1)
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i) filter(gen[i]);</pre>
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
  for (int j = i; j < n; ++j) {</pre>
       for (int k = 0; k < (int)bkts[i].size(); ++k) {</pre>
         for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l));
       }
    }
  }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.first][b.
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].size() -
     1);
     for (int i = 0; i < n; ++i) {
       for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
         if (i <= res) upd.emplace(make_pair(i, j), pr);</pre>
         if (res <= i) upd.emplace(pr, make_pair(i, j));</pre>
    }
  }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i) res = res * bkts[i].size();</pre>
  return res;
į }}
```

### 6.10 Berlekamp-Massey Algorithm

```
| template <int P>
| vector<int> BerlekampMassey(vector<int> x) {
| vector<int> cur, ls;
| int lf = 0, ld = 0;
| for (int i = 0; i < (int)x.size(); ++i) {
| int t = 0;
| for (int j = 0; j < (int)cur.size(); ++j)
| (t += 1LL * cur[j] * x[i - j - 1] % P) %= P;
| if (t == x[i]) continue;
| if (cur.empty()) {
| cur.resize(i + 1);
| lf = i, ld = (t + P - x[i]) % P;</pre>
```

```
continue;
}
int k = 1LL * fpow(ld, P - 2, P) * (t + P - x[i]) % P;
vector<int> c(i - lf - 1);
c.push_back(k);
for (int j = 0; j < (int)ls.size(); ++j)
    c.push_back(1LL * k * (P - ls[j]) % P);
if (c.size() < cur.size()) c.resize(cur.size());
for (int j = 0; j < (int)cur.size(); ++j)
    c[j] = (c[j] + cur[j]) % P;
if (i - lf + (int)ls.size() >= (int)cur.size()) {
    ls = cur, lf = i;
    ld = (t + P - x[i]) % P;
}
cur = c;
}
return cur;
}
```

#### 6.11 Miller Rabin

```
// n < 4759123141
                      chk = [2, 7, 61]
// n < 1122004669633 chk = [2, 13, 23, 1662803]
// n < 2^64 chk = [2, 325, 9375, 28178, 450775, 9780504,
      17952650227
vector<long long> chk = { 2, 325, 9375, 28178, 450775, 9780504,
       1795265022 }:
bool check(long long a, long long u, long long n, int t) {
  a = fpow(a, u, n);
   if (a == 0) return true;
   if (a == 1 \mid \mid a == n - 1) return true;
   for (int i = 0; i < t; ++i) {
     a = fmul(a, a, n);
     if (a == 1) return false;
     if (a == n - 1) return true;
  }
  return false:
}
bool is_prime(long long n) {
   if (n < 2) return false;
   if (n % 2 == 0) return n == 2;
   long long u = n - 1; int t = 0;
   for (; !(u & 1); u >>= 1, ++t);
   for (long long i : chk) {
     if (!check(i, u, n, t)) return false;
   return true;
```

#### 6.12 Pollard's Rho

```
map<long long, int> cnt;
 long long f(long long x, long long n, int p) { return (fmul(x,
      x, n) + p) % n; }
 void pollard_rho(long long n) {
   if (n == 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2 == 0) return pollard_rho(n / 2), ++cnt[2], void(); long long x = 2, y = 2, d = 1, p = 1;
   while (true) {
     if (d != n && d != 1) {
       pollard_rho(n / d);
       pollard_rho(d);
       return;
     if (d == n) ++p;
     x = f(x, n, p); y = f(f(y, n, p), n, p);
     d = \_gcd(abs(x - y), n);
  }
}
```

### 6.13 Meissel-Lehmer Algorithm

```
|int prc[maxn];
|long long phic[msz][nsz];
|void sieve() {
```

```
bitset<maxn> v
   pr.push_back(0);
   for (int i = 2; i < maxn; ++i) {</pre>
     if (!v[i]) pr.push_back(i);
     for (int j = 1; i * pr[j] < maxn; ++j) {
       v[i * pr[j]] = true;
       if (i % pr[j] == 0) break;
     }
   for (int i = 1; i < pr.size(); ++i) prc[pr[i]] = 1;</pre>
   for (int i = 1; i < maxn; ++i) prc[i] += prc[i - 1];</pre>
 long long p2(long long, long long);
 long long phi(long long m, long long n) {
  if (m < msz && n < nsz && phic[m][n] != -1) return phic[m][n</pre>
   if (n == 0) return m;
  if (pr[n] >= m) return 1;
   long long ret = phi(m, n - 1) - phi(m / pr[n], n - 1);
   if (m < msz && n < nsz) phic[m][n] = ret;</pre>
   return ret;
long long pi(long long m) {
   if (m < maxn) return prc[m];</pre>
   long long n = pi(cbrt(m));
   return phi(m, n) + n - 1 - p2(m, n);
long long p2(long long m, long long n) {
  long long ret = 0;
   long long lim = sqrt(m);
   for (int i = n + 1; pr[i] <= lim; ++i) ret += pi(m / pr[i]) -</pre>
       pi(pr[i]) + 1;
   return ret;
1}
```

### 6.14 Discrete Logarithm

```
// to solve discrete x for x^a = b \pmod{p} with p is prime
// let c = primitive root of p
// find k such that c^k = b \pmod{p} by bsgs
// solve fa = k \pmod{p-1} by euclidean algorithm
// x = c^f
int bsgs(int a, int b, int p) {
  // return L such that a^L = b \pmod{p}
  if (p == 1) {
    if (!b) return a != 1;
return -1;
  if (b == 1) {
    if (a) return 0;
    return -1;
  if (a \% p == 0) {
    if (!b) return 1;
    return -1;
  }
  int num = 0, d = 1;
  while (true) {
    int r = \__gcd(a, p);
    if (r == 1) break;
    if (b % r) return -1;
    ++num;
b /= r, p /= r;
    d = (111 * d * a / r) \% p;
  for (int i = 0, now = 1; i < num; ++i, now = 111 * now * a %
     if (now == b) return i;
  int m = ceil(sqrt(p)), base = 1;
  map<int, int> mp;
  for (int i = 0; i < m; ++i) {
     if (mp.find(base) == mp.end()) mp[base] = i;
    else mp[base] = min(mp[base], i);
base = 111 * base * a % p;
  for (int i = 0; i < m; ++i) {
    \ensuremath{/\!/} can be modified to fpow if p is prime
    int r, x, y; tie(r, x, y) = extgcd(d, p);
x = (111 * x * b % p + p) % p;
    if (mp.find(x) != mp.end()) return i * m + mp[x] + num;
d = 111 * d * base % p;
  return -1;
```

### 6.15 Quadratic Residue

```
int Jacobi(int a, int m) {
   int s = 1;
   for (; m > 1; ) {
   a %= m;
     if (a == 0) return 0;
     const int r = __builtin_ctz(a);
     if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
     if (a \& m \& 2) s = -s;
     swap(a, m);
   return s;
int QuadraticResidue(int a, int p) {
   if (p == 2) return a & 1;
   const int jc = Jacobi(a, p);
if (jc == 0) return 0;
   if (jc == -1) return -1;
   int b, d;
   for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
   for (int e = (p + 1) >> 1; e; e >>= 1) {
     if (e & 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p
       g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
       g0 = tmp;
     tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
     f1 = (2LL * f0 * f1) % p;
     f0 = tmp;
   return q0;
```

#### 6.16 Gaussian Elimination

```
double gauss(vector<vector<double>> &d) {
  int n = d.size(), m = d[0].size();
double det = 1;
  for (int i = 0; i < m; ++i) {
    int p = -1;
    for (int j = i; j < n; ++j) {
      if (fabs(d[j][i]) < eps) continue;</pre>
      if (p == -1 \mid | fabs(d[j][i]) > fabs(d[p][i])) p = j;
    if (p == -1) continue;
    if (p != i) det *= -1;
    for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
    for (int j = 0; j < n; ++j) {
      if (i == j) continue;
      double z = d[j][i] / d[i][i];
      for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
    }
  }
  for (int i = 0; i < n; ++i) det *= d[i][i];</pre>
  return det;
```

### 6.17 $\mu$ function

```
int mu[maxn], pi[maxn];
vector<int> prime;

void sieve() {
    mu[1] = pi[1] = 1;
    for (int i = 2; i < maxn; ++i) {
        if (!pi[i]) {
            pi[i] = i;
            prime.push_back(i);
            mu[i] = -1;
        }
        for (int j = 0; i * prime[j] < maxn; ++j) {
            pi[i * prime[j]] = prime[j];
            mu[i * prime[j]] = -mu[i];
        }
</pre>
```

```
if (i % prime[j] == 0) {
    mu[i * prime[j]] = 0;
    break;
}
}
}
```

#### 6.18 Partition Function

```
void Build(int n) {
    // P[i] = the number of ways of representing i as the sum of
    a non-decreasing sequence.
    vector<pair<int, int>> g = {{0, 0}};
    for (int i = 1; g.back().second <= n; ++i) {
        g.emplace_back(i % 2 ? 1 : kMod - 1, i * (3 * i - 1) / 2);
        g.emplace_back(i % 2 ? 1 : kMod - 1, i * (3 * i + 1) / 2);
    }
    P[0] = P[1] = 1;
    for (int i = 2; i <= n; ++i) {
        for (auto it : g) {
            if (i < it.second) continue;
            P[i] += 1LL * P[i - it.second] * it.first % kMod;
            P[i] %= kMod;
        }
    }
}</pre>
```

### 6.19 $\left| \frac{n}{i} \right|$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

### 6.20 De Bruijn Sequence

```
int res[maxn], aux[maxn], a[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
    if (n \% p == 0) {
      for (int i = 1; i <= p; ++i) res[sz++] = aux[i];</pre>
 } else {
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
    for (int i = aux[t - p] + 1; i < k; ++i) {
      aux[t] = i;
      db(t + 1, t, n, k);
    }
 }
}
int de_bruijn(int k, int n) {
  // return cyclic string of length k^n such that every string
    of length n using k character appears as a substring.
  if (k == 1) {
    res[0] = 0;
    return 1:
  for (int i = 0; i < k * n; i++) aux[i] = 0;
  sz = 0;
  db(1, 1, n, k);
  return sz;
```

#### 6.21 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
  if (!b) return make_tuple(a, 1, 0);
  T d, x, y;
  tie(d, x, y) = extgcd(b, a % b);
  return make_tuple(d, y, x - (a / b) * y);
}
```

### 6.22 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

#### 6.23 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
  long long mult = mod[0];
  int n = (int)mod.size();
  long long res = a[0];
  for (int i = 1; i < n; ++i) {
    long long d, x, y;
    tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
    if ((a[i] - res) % d) return -1;
    long long new_mult = mult / __gcd(mult, 1ll * mod[i]) * mod
    [i];
    res += x * ((a[i] - res) / d) % new_mult * mult % new_mult;
    mult = new_mult;
    ((res %= mult) += mult) %= mult;
  }
  return res;
}</pre>
```

### 6.24 Theorem

#### 6.24.1 Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i), \ L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 6.24.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

### 6.24.3 Cayley's Formula

- Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \ldots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

#### 6.24.4 Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

#### 6.25 Primes

```
\begin{array}{l} 97,101,131,487,593,877,1087,1187,1487,1787,3187,12721,\\ 13331,14341,75577,123457,222557,556679,999983,\\ 1097774749,1076767633,100102021,999997771,\\ 1001010013,1000512343,987654361,999991231,\\ 999888733,98789101,987777733,999991921,1000000007,\\ 1000000087,1000000123,1010101333,1010102101,\\ 100000000039,10000000000037,2305843009213693951,\\ 4611686018427387847,9223372036854775783,18446744073709551557\end{array}
```

### 7 Dynamic Programming

### 7.1 Dynamic Convex Hull

```
struct Line {
   mutable int64_t a, b, p;
   bool operator<(const Line &rhs) const { return a < rhs.a; }</pre>
   bool operator<(int64_t x) const { return p < x; }</pre>
 struct DynamicHull : multiset<Line, less<>> {
    static const int64_t kInf = 1e18;
   int64_t Div(int64_t a, int64_t b) { return a / b - ((a \land b) <
       0 && a % b); }
   bool Isect(iterator x, iterator y) {
     if (y == end()) { x->p = kInf; return false; }
     if (x->a == y->a) x->p = x->b > y->b? kInf : -kInf;
      else x->p = Div(y->b - x->b, x->a - y->a);
      return x->p >= y->p;
   void Insert(int64_t a, int64_t b) {
     auto z = insert({a, b, 0}), y = z++, x = y;
      while (Isect(y, z)) z = erase(z);
     if (x != begin() && Isect(--x, y)) Isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p) Isect(x,
      erase(y));
   int64_t Query(int64_t x) {
     auto l = *lower_bound(x);
return l.a * x + l.b;
|};
```

### 7.2 1D/1D Convex Optimization

```
struct segment {
  int i, l, r;
  segment(int a, int b, int c): i(a), l(b), r(c) {}
inline long long f(int l, int r) { return dp[l] + w(l + 1, r);
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
  dp[i] = f(deq.front().i, i);</pre>
    while (deq.size() && deq.front().r < i + 1) deq.pop_front()</pre>
    deq.front().l = i + 1;
    segment seg = segment(i, i + 1, n);
    while (deq.size() && f(i, deq.back().l) < f(deq.back().i,</pre>
     deq.back().1)) deq.pop_back();
    if (deq.size()) {
      int d = 1048576, c = deq.back().1;
while (d >>= 1) if (c + d <= deq.back().r) {</pre>
         if (f(i, c + d) > f(deq.back().i, c + d)) c += d;
      deq.back().r = c; seq.l = c + 1;
    if (seg.l <= n) deq.push_back(seg);</pre>
```

### 7.3 Condition

}

#### 7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{ll} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

#### 7.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \, B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \, B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

| bool same(double a, double b) { return abs(a - b) < eps; }

### 8 Geometry

#### 8.1 Basic

```
struct P {
  double x
  P() : x(0), y(0) {}
  P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
  P operator * (double b) { return P(x / b, y / b); } double operator * (P b) { return x * b.x + y * b.y; } double operator * (P b) { return x * b.y - y * b.x; }
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P spin(double o) {
     double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
  double angle() { return atan2(y, x); }
};
struct L {
   // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
  L(): a(0), b(0), c(0), o(0), pa(), pb() {}
  L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa ^ pb), o
     (atan2(-a, b)), pa(pa), pb(pb) {}
  P reflect(P p) { return p + (project(p) - p) * 2; }
  double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
     pa).abs() * (pb - pa).abs()); }
};
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
  if (max(p1.x, p2.x) < min(p3.x, p4.x) | | max(p3.x, p4.x) < |
     min(p1.x, p2.x)) return false;
  if (max(p1.y, p2.y) < min(p3.y, p4.y) | | max(p3.y, p4.y) <
     min(p1.y, p2.y)) return false:
  return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^ (p4 -
     p2)) <= 0 &&
       sign((p1 - p3) \land (p2 - p3)) * sign((p1 - p4) \land (p2 - p4))
}
bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b, x.a *
     y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

#### 8.2 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[
point p[maxn];
int build(int 1, int r, int dep = 0) {
  if (l == r) return -1;
  function<bool(const point &, const point &)> f = [dep](const
    point &a, const point &b) {
    if (dep \& 1) return a.x < b.x;
    else return a.y < b.y;
  };
  int m = (l + r) >> 1;
  nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
   xr[m] = max(xr[m], xr[rc[m]]);
yl[m] = min(yl[m], yl[rc[m]]);
   yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
      (a.y - b.y) * 111 * (a.y - b.y);
void dfs(const point &q, long long &d, int o, int dep = 0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x || !(dep & 1) && q.y < p[o].y)
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
 } else {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void init(const vector<point> &v) {
 for (int i = 0; i < v.size(); ++i) p[i] = v[i];
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
 dfs(q, res, root);
return res;
```

#### 8.3 Delaunay Triangulation

```
|/* Delaunay Triangulation:
| Given a sets of points on 2D plane, find a
| triangulation such that no points will strictly
| inside circumcircle of any triangle.
| find : return a triangle contain given point
| add_point : add a point into triangulation
| A Triangle is in triangulation iff. its has_chd is 0.
| Region of triangle u: iterate each u.edge[i].tri,
| each points are u.p[(i+1)%3], u.p[(i+2)%3]
| calculation involves O(|V|^6) */
| const double inf = 1e9;
| double eps = 1e-6; // 0 when integer
| // return p4 is in circumcircle of tri(p1,p2,p3)
| bool in_cc(P &p1, P &p2, P &p3, P &p4) {
| int o1 = (abs(p1.x) >= inf * 0.99 || abs(p1.y) >= inf * 0.99);
| int o2 = (abs(p2.x) >= inf * 0.99 || abs(p2.y) >= inf * 0.99);
```

```
int o3 = (abs(p3.x) >= inf * 0.99 || abs(p3.y) >= inf * 0.99);
 int rtrue = o1 + o2 + o3;
 int rfalse = abs(p4.x) \Rightarrow inf * 0.99 | | abs(p4.y) \Rightarrow inf *
     0.99;
 if (rtrue == 3) return true;
 if (rtrue) {
  P in(0, 0), out(0, 0);
  if (o1) out = out + p1; else in = in + p1;
  if (o2) out = out + p2; else in = in + p2;
  if (o3) out = out + p3; else in = in + p3;
  return (p4 - in) * (out - in) > 0;
 if (rfalse) return false;
// ^ ?
 double u11 = p1.x - p4.x, u12 = p1.y - p4.y;
 double u21 = p2.x - p4.x, u22 = p2.y - p4.y;
double u31 = p3.x - p4.x, u32 = p3.y - p4.y;
 double u13 = sq(p1.x) - sq(p4.x) + sq(p1.y) - sq(p4.y);
 double u23 = sq(p2.x) - sq(p4.x) + sq(p2.y) - sq(p4.y);
 double u33 = sq(p3.x) - sq(p4.x) + sq(p3.y) - sq(p4.y);
double det = -u13 * u22 * u31 + u12 * u23 * u31 + u13 * u21 *
u32 - u11 * u23 * u32 - u12 * u21 * u33 + u11 * u22 * u33;
 return det > eps;
double side(P &a, P &b, P &p) { return (b - a) ^ (p - a); }
struct Tri:
struct Edge {
 Tri *tri;
 int side;
 Edge() : tri(0), side(0) {}
Edge(Tri *_tri, int _side) : tri(_tri), side(_side) {}
struct Tri {
 P p[3];
 Edge edge[3];
Tri *ch[3];
 Tri() {}
 Tri(P p0, P p1, P p2) {
  p[0] = p0; p[1] = p1; p[2] = p2;
  ch[0] = ch[1] = ch[2] = 0;
 bool has_ch() { return ch[0] != 0; }
 int num_ch() {
  return ch[0] == 0 ? 0 : ch[1] == 0 ? 1 : ch[2] == 0 ? 2 : 3;
 bool contains(P &q) {
  for (int i = 0; i < 3; ++i)
   if (side(p[i], p[(i + 1) % 3], q) < -eps) return false;
} pool[maxn * 10], *tris;
void edge(Edge a, Edge b) {
  if (a.tri) a.tri->edge[a.side] = b;
if (b.tri) b.tri->edge[b.side] = a;
struct Trig {
 Trig() {
  the_root = new (tris++) Tri(P(-inf, -inf), P(inf * 2, -inf),
   P(-inf, inf * 2));
   // all p should in
 Tri *find(P p) { return find(the_root, p); }
 void add_point(P &p) { add_point(find(the_root, p), p); }
 Tri *the_root;
 static Tri *find(Tri *root, P &p) {
  while (true) {
   if (!root->has_ch()) return root;
   for (int i = 0; i < 3 \&\& root->ch[i]; ++i)
    if (root->ch[i]->contains(p)) {
      root = root->ch[i];
     break;
  assert(false); // "point not found"
 void add_point(Tri *root, P &p) {
  Tri *tab, *tbc, *tca;
  tab = new (tris++) Tri(root->p[0], root->p[1], p);
  tbc = new (tris++) Tri(root->p[1], root->p[2], p);
  tca = new (tris++) Tri(root->p[2], root->p[0], p);
  edge(Edge(tab, 0), Edge(tbc, 1));
  edge(Edge(tbc, 0), Edge(tca, 1));
  edge(Edge(tca, 0), Edge(tab, 1));
  edge(Edge(tab, 2), root->edge[2]);
  edge(Edge(tbc, 2), root->edge[0]);
  edge(Edge(tca, 2), root->edge[1]);
  root->ch[0] = tab; root->ch[1] = tbc; root->ch[2] = tca;
  flip(tab, 2); flip(tbc, 2); flip(tca, 2);
```

```
void flip(Tri *tri, int pi) {
  Tri *trj = tri->edge[pi].tri;
  int pj = tri->edge[pi].side;
  if (!trj) return;
  if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[pj]))
  /* flip edge between tri,trj */
  Tri *trk = new (tris++) Tri(tri->p[(pi + 1) % 3], trj->p[pj],
       tri->p[pi]);
  Tri *trl = new (tris++) Tri(trj->p[(pj + 1) % 3], tri->p[pi],
       trj->p[pj]);
  edge(Edge(trk, 0), Edge(trl, 0));
edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
  edge(Edge(trk, 1), trj->edge[(pj + 1) % 3]);
edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
tri->ch[0] = trk; tri->ch[1] = trl; tri->ch[2] = 0;
trj->ch[0] = trk; trj->ch[1] = trl; trj->ch[2] = 0;
  flip(trk, 1); flip(trk, 2);
  flip(trl, 1); flip(trl, 2);
vector<Tri *> triang;
set<Tri *> vst;
void go(Tri *now) {
if (vst.find(now) != vst.end()) return;
 vst.insert(now);
 if (!now->has_ch()) {
  triang.push_back(now);
  return;
 for (int i = 0; i < now->num\_ch(); ++i) go(now->ch[i]);
void build(int n, P *ps) {
tris = pool
 random\_shuffle(ps, ps + n);
for (int i = 0; i < n; ++i) tri.add_point(ps[i]);</pre>
 go(tri.the_root);
```

### 8.4 Voronoi Diagram

```
int gid(P &p) {
   auto it = ptoid.find(p);
   if (it == ptoid.end()) return -1;
   return it->second;
 L make_line(P p, L l) {
   P d = 1.pb - 1.pa; d = d.spin(pi / 2);
   P m = (1.pa + 1.pb) / 2;
   l = L(m, m + d);
   if (((1.pb - 1.pa) \land (p - 1.pa)) < 0) l = L(m + d, m);
   return 1;
 double calc_ans(int i) {
   vector<P> ps = HPI(ls[i]);
double rt = 0;
   for (int i = 0; i < (int)ps.size(); ++i) {
     rt += (ps[i] ^ ps[(i + 1) % ps.size()]);
   return abs(rt) / 2;
 }
 void solve() {
   for (int i = 0; i < n; ++i) ops[i] = ps[i], ptoid[ops[i]] = i
   random_shuffle(ps, ps + n);
   build(n, ps);
   for (auto *t : triang) {
     int z[3] = {gid(t->p[0]), gid(t->p[1]), gid(t->p[2])};
for (int i = 0; i < 3; ++i) for (int j = 0; j < 3; ++j) if
(i != j && z[i] != -1 && z[j] != -1) {
       L l(t->p[i], t->p[j]);
        ls[z[i]].push_back(make_line(t->p[i], l));
     }
   }
   vector<P> tb = convex(vector<P>(ps, ps + n));
   for (auto &p : tb) isinf[gid(p)] = true;
   for (int i = 0; i < n; ++i) {
     if (isinf[i]) cout << -1 << '\n';</pre>
     else cout << fixed << setprecision(12) << calc_ans(i) << '</pre>
   }
| }
```

### 8.5 Sector Area

```
// calc area of sector which include a, b
double SectorArea(P a, P b, double r) {
  double o = atan2(a.y, a.x) - atan2(b.y, b.x);
  while (o <= 0) o += 2 * pi;
  while (o >= 2 * pi) o -= 2 * pi;
  o = min(o, 2 * pi - o);
  return r * r * o / 2;
}
```

### 8.6 Half Plane Intersection

```
bool jizz(L l1,L l2,L l3){
  P p=Intersect(12,13);
   return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
bool cmp(const L &a,const L &b){
  return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;</pre>
}
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
   sort(ls.begin(),ls.end(),cmp);
   vector<L> pls(1,ls[0]);
   for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls.back().</pre>
      o))pls.push_back(ls[i]);
   deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],pls[b],
     pls[c]))
   for(int i=2;i<(int)pls.size();++i){</pre>
     meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
     meow(i,dq[0],dq[1])dq.pop_front();
     dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop_front();
   if(dq.size()<3u)return vector<P>(); // no solution or
      solution is not a convex
   vector<P> rt;
   for(int i=0;i<(int)dq.size();++i)rt.push_back(Intersect(pls[</pre>
      dq[i]],pls[dq[(i+1)%dq.size()]]));
13
```

### 8.7 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
   int n=int(ps.size());
   vector<int> id(n),pos(n);
   vector<pair<int,int>> line(n*(n-1)/2);
   for(int i=0;i< n;++i)for(int j=i+1;j< n;++j)line[++m]=make_pair
     (i,j); ++m;
   sort(line.begin(),line.end(),[&](const pair<int,int> &a,const
      pair<int, int> &b)->bool{
     if(ps[a.first].first==ps[a.second].first)return 0;
     if(ps[b.first].first==ps[b.second].first)return 1;
     return (double)(ps[a.first].second-ps[a.second].second)/(ps
      [a.first].first-ps[a.second].first) < (double)(ps[b.first</pre>
      ].second-ps[b.second].second)/(ps[b.first].first-ps[b.
      second].first);
  });
   for(int i=0;i<n;++i)id[i]=i;</pre>
   sort(id.begin(),id.end(),[&](const int &a,const int &b){
      return ps[a]<ps[b]; });</pre>
   for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
   for(int i=0;i<m;++i){</pre>
     auto l=line[i];
     // meow
     tie(pos[l.first],pos[l.second],id[pos[l.first]],id[pos[l.
      second]])=make_tuple(pos[l.second],pos[l.first],l.second,l
      .first);
  }
}
```

### 8.8 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by - ay)) / (
    sin(a1) * cos(a2) - sin(a2) * cos(a1));
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
}
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
     TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);

res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

### 8.9 Polygon Center

```
| Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
    res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
    res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

#### 8.10 Maximum Triangle

```
double ConvexHullMaxTriangleArea(Point p[], int res[], int
     chnum) {
  double area = 0, tmp;
  res[chnum] = res[0];
  for (int i = 0, j = 1, k = 2; i < chnum; i++) {
    while (fabs(Cross(p[res[j]] - p[res[i]], p[res[(k + 1) %
     chnum]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
     p[res[k]] - p[res[i]])) k = (k + 1) % chnum;
     tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
     ]]));
    if (tmp > area) area = tmp;
    while (fabs(Cross(p[res[(j + 1) % chnum]] - p[res[i]], p[
     res[k]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
     p[res[k]] - p[res[i]])) j = (j + 1) % chnum;
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
     ]]));
    if (tmp > area) area = tmp;
  return area / 2;
| }
```

### 8.11 Point in Polygon

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) % ps.size();
    L l(ps[a], ps[b]);</pre>
```

#### 8.12 Circle

```
struct C {
  P c;
double r;
  C(P c = P(0, 0), double r = 0) : c(c), r(r) {}
};
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c).unit()
      * a.r);
  else if (a.r + b.r > d & d + a.r >= b.r) {
     double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.spin(o) * a.r);
    p.push_back(a.c + i.spin(-o) * a.r);
  return p;
double IntersectArea(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d >= a.r + b.r - eps) return 0;
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));

double q = acos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));

return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
// remove second level if to get points for line (defalut:
     segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
  double C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
  if (d >= -eps) {
    d = max(0., d);
double i = (-B - sqrt(d)) / (2 * A);
    double j = (-B + sqrt(d)) / (2 * A);
     if (i - 1.0 \leftarrow eps \& i \rightarrow eps) t.emplace_back(a.x + i * x)
       a.y + i * y);
     if (j - 1.0 \le eps \&\& j \ge -eps) t.emplace_back(a.x + j * x)
     , a.y + j * y);
  return t:
}
\ensuremath{/\!/} calc area intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
     if (inb) return abs(a ^ b) / 2;
     return SectorArea(b, p[0], r) + abs(a \land p[0]) / 2;
  if (inb) return SectorArea(p[0], a, r) + abs(p[0] ^b b) / 2; if (p.size() == 2u) return SectorArea(a, p[0], r) +
     SectorArea(p[1], b, r) + abs(p[0] \land p[1]) / 2;
  else return SectorArea(a, b, r);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
     int j = (i + 1) \% 3;
     double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y, ps[j].x
     if (o >= pi) o = o - 2 * pi;
     if (o <= -pi) o = o + 2 * pi;
```

### 8.13 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
#define Pij ∖
  P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x);
  z.emplace_back(a.c + i, a.c + i + j);
#define deo(I,J) \
  double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos(e / d)
  P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-o);\
z.emplace_back(a.c + j * a.r, b.c J j * b.r);\
z.emplace_back(a.c + k * a.r, b.c J k * b.r);
  if (a.r < b.r) swap(a, b);
  vector<L> z;
  if ((a.c - b.c).abs() + b.r < a.r) return z;
   else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
  else {
     deo(-,+);
     if (same(d, a.r + b.r)) { Pij; }
     else if (d > a.r + b.r) \{ deo(+,-); \}
}
vector<L> tangent(C c, P p) {
   vector<L> z;
   double d = (p - c.c).abs();
   if (same(d, c.r)) {
     P i = (p - c.c).spin(pi / 2);
     z.emplace_back(p, p + i);
  } else if (d > c.r) {
     double o = acos(c.r / d);
     P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.spin(-o) * c.r;
     z.emplace_back(c.c + j, p);
     z.emplace_back(c.c + k, p);
   return z;
į }
```

#### 8.14 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
 double d = (a.c - b.c).abs();
 vector<pair<double, double>> res;
  if (same(a.r + b.r, d));
  else if (d \le abs(a.r - b.r) + eps) {
    if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
 } else if (d < abs(a.r + b.r) - eps) {
    double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
    ), z = (b.c - a.c).angle();
   if (z < 0) z += 2 * pi;
double l = z - o, r = z + o;
    if (l < 0) l += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
    if (l > r) res.emplace_back(l, 2 * pi), res.emplace_back(0,
     r);
    else res.emplace_back(l, r);
  return res;
double CircleUnionArea(vector<C> c) { // circle should be
    identical
  int n = c.size():
  double a = 0, w;
 for (int i = 0; w = 0, i < n; ++i) {
    vector<pair<double, double>> s = \{\{2 * pi, 9\}\}, z;
    for (int j = 0; j < n; ++j) if (i != j) {
      z = CoverSegment(c[i], c[j]);
      for (auto &e : z) s.push_back(e);
    sort(s.begin(), s.end());
    auto F = [\&] (double t) { return c[i].r * (c[i].r * t + c[i])
     ].c.x * sin(t) - c[i].c.y * cos(t)); };
    for (auto &e : s) {
      if (e.first > w) a += F(e.first) - F(w);
      w = max(w, e.second);
```

### 8.15 Minimun Distance of 2 Polygons

#### 8.16 2D Convex Hull

```
bool operator < (const P &a, const P &b) { return same(a.x, b.x</pre>
     ) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(a.x, b.x
     ) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) ^ (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return same(a.x,
     b.x) ? a.y < b.y : a.x < b.x; });
  for (int i = 0; i < ps.size(); ++i) {
    while (p.size() >= 2 && crx(p[p.size() - 2], ps[i], p[p.
     size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() - 2], ps[i], p[p.size
     () - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double x) { return same(x, 0) ? 0 : x > 0 ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector<P> p, u, d;
  CH() {}
  CH(vector<P> ps) : p(ps) {
    n = ps.size();
    rotate(p.begin(), min_element(p.begin(), p.end()), p.end())
    auto t = max_element(p.begin(), p.end());
    d = vector<P>(p.begin(), next(t));
    u = vector<P>(t, p.end()); u.push_back(p[0]);
  int find(vector<P> &v, P d) {
    int l = 0, r = v.size();
    while (l + 5 < r) {
      int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;
if (v[L] * d > v[R] * d) r = R;
else l = L;
    }
    int x = 1;
```

```
for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x] * d) x
                                                                       P p[maxn]:
                                                                        T f[maxn*8];
      = i:
     return x;
                                                                        int id[maxn][maxn]
                                                                        bool on(T &t,P &q){
   int findFarest(P v) {
                                                                          return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
     if (v.y > 0 \mid | v.y == 0 \& v.x > 0) return ((int)d.size() - (int)d.size() = 0
                                                                        }
       1 + find(u, v)) % p.size();
                                                                        void meow(int q,int a,int b){
     return find(d, v);
                                                                          int f2=id[a][b];
                                                                          if(f[f2].res){
   P get(int l, int r, P a, P b) {
                                                                            if(on(f[f2],p[q]))dfs(q,f2);
     int s = sgn(crx(a, b, p[l % n]));
                                                                            else{
     while (l + 1 < r) {
                                                                              id[q][b]=id[a][q]=id[b][a]=m;
       int m = (l + r) >> 1;
                                                                              f[m++]=T(b,a,q,1);
       if (sgn(crx(a, b, p[m % n])) == s) l = m;
       else r = m:
                                                                         }
                                                                        }
     return isLL(a, b, p[l % n], p[(l + 1) % n]);
                                                                        void dfs(int p,int i){
                                                                          f[i].res=0;
   vector<P> getIS(P a, P b) {
                                                                          meow(p,f[i].b,f[now].a);
     int X = findFarest((b - a).spin(pi / 2));
                                                                          meow(p,f[i].c,f[now].b);
     int Y = findFarest((a - b).spin(pi / 2));
                                                                          meow(p,f[i].a,f[now].c);
     if (X > Y) swap(X, Y)
     if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0) return
                                                                        void operator()(){
       \{get(X, Y, a, b), get(Y, X + n, a, b)\};
                                                                          if(n<4)return;
     return {};
                                                                          if([&]()->int{
   }
                                                                            for(int i=1;i< n;++i)if(abs(p[0]-p[i])>eps)return swap(p[1],
   void update_tangent(P q, int i, int &a, int &b) {
                                                                             p[i]),0;
     if (sgn(crx(q, p[a], p[i])) > 0) a = i;
if (sgn(crx(q, p[b], p[i])) < 0) b = i;</pre>
                                                                            return 1:
                                                                          }())return;
                                                                          if([&]()->int{
   void bs(int l, int r, P q, int &a, int &b) {
                                                                            for(int i=2;i<n;++i)if(abs((p[0]-p[i])^(p[1]-p[i]))>eps)
     if (l == r) return;
                                                                             return swap(p[2],p[i]),0;
     update_tangent(q, 1 % n, a, b);
                                                                            return 1;
     int s = sgn(crx(q, p[l % n], p[(l + 1) % n]));
                                                                          }())return;
     while (l + 1 < r) {
                                                                          if([&]()->int{
       int m = (l + r) >> 1;
                                                                            for(int i=3; i< n; ++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))*(p[i]-
       if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l = m;
                                                                             p[0]))>eps)return swap(p[3],p[i]),0;
                                                                            return 1;
                                                                          }())return;
     update_tangent(q, r % n, a, b);
                                                                          for(int i=0;i<4;++i){</pre>
                                                                            T t((i+1)%4,(i+2)%4,(i+3)%4,1);
   bool contain(P p) {
                                                                            if(on(t,p[i]))swap(t.b,t.c);
     if (p.x < d[0].x \mid | p.x > d.back().x) return 0;
                                                                            id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
     auto it = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
                                                                            f[m++]=t;
     if (it->x == p.x) {
     if (it->y > p.y) return 0;
} else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
                                                                          for(int i=4;i< n;++i)for(int j=0;j< m;++j)if(f[j].res && on(f[j
                                                                             ],p[i])){
                                                                            dfs(i,j);
     it = lower_bound(u.begin(), u.end(), P(p.x, 1e12), greater<
     P>());
                                                                            break;
     if (it->x == p.x) {
       if (it->y < p.y) return 0;</pre>
                                                                          int mm=m: m=0:
                                                                          for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     return 1:
                                                                        bool same(int i,int j){
   bool get_tangent(P p, int &a, int &b) { // b -> a
                                                                          return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].a])>eps
                                                                             || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].b])>eps ||
     if (contain(p)) return 0;
                                                                             absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
     int i = lower_bound(d.begin(), d.end(), p) - d.begin();
     bs(0, i, p, a, b);
                                                                        int faces(){
     bs(i, d.size(), p, a, b);
                                                                          int r=0;
     i = lower_bound(u.begin(), u.end(), p, greater<P>()) - u.
                                                                          for(int i=0;i<m;++i){</pre>
                                                                            int iden=1;
     bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a, b);
                                                                            for(int j=0;j<i;++j)if(same(i,j))iden=0;</pre>
     bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.size(), p,
                                                                            r+=iden;
     a, b);
return 1;
                                                                          return r;
  }
|};
                                                                       |} tb;
```

#### 8.17 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d){
   return abs(((b-a)^(c-a))*(d-a))/6;
}

struct convex3D{
   static const int maxn=1010;
   struct T{
      int a,b,c;
      bool res;
      T(){}
      T(int a,int b,int c,bool res=1):a(a),b(b),c(c),res(res){}
};
int n,m;
```

### 8.18 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^ p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
}

circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {</pre>
```

```
if (norm2(cent - p[i]) <= r) continue;
cent = p[i];
r = 0.0;
for (int j = 0; j < i; ++j) {
    if (norm2(cent - p[j]) <= r) continue;
    cent = (p[i] + p[j]) / 2;
    r = norm2(p[j] - cent);
    for (int k = 0; k < j; ++k) {
        if (norm2(cent - p[k]) <= r) continue;
        cent = center(p[i], p[j], p[k]);
        r = norm2(p[k] - cent);
    }
}
return circle(cent, sqrt(r));
}</pre>
```

#### 8.19 Closest Pair

```
double closest_pair(int 1, int r) {
 // p should be sorted increasingly according to the x-
    coordinates.
  if (l == r) return 1e9;
 if (r - l == 1) return dist(p[l], p[r]);
 int m = (l + r) >> 1;
 double d = min(closest_pair(l, m), closest_pair(m + 1, r));
  vector<int> vec;
 for (int i = m; i >= l \&\& fabs(p[m].x - p[i].x) < d; --i) vec
    .push_back(i);
 for (int i = m + 1; i \le r \& fabs(p[m].x - p[i].x) < d; ++i)
     vec.push_back(i);
  sort(vec.begin(), vec.end(), [&](int a, int b) { return p[a].
    y < p[b].y; \});
  for (int i = 0; i < vec.size(); ++i) {</pre>
   for (int j = i + 1; j < vec.size() && fabs(p[vec[j]].y - p[
    vec[i]].y) < d; ++j) {
      d = min(d, dist(p[vec[i]], p[vec[j]]));
   }
  return d;
```

### 9 Miscellaneous

### 9.1 Bitwise Hack

```
long long next_perm(long long v) {
  long long t = v | (v - 1);
  return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1))
   ;
  }
  void subset(long long s) {
  long long sub = s;
  while (sub) sub = (sub - 1) & s;
}
```

### 9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) x = s - 1 - x, y = s - 1 - y;
        swap(x, y);
    }
  }
  return res;
}
```

### 9.3 Java

```
import java.io.*
import java.util.*;
import java.lang.*
import java.math.*;
public class filename{
  static Scanner in = new Scanner(System.in);
  public static void main(String[] args) throws Exception {
   Scanner fin = new Scanner(new File("infile"));
     PrintWriter fout = new PrintWriter("outfile", "UTF-8");
     fout.println(fin.nextLine());
     fout.close();
    while (in.hasNext()) {
       String str = in.nextLine(); // getline
       String stu = in.next(); // string
     System.out.println("Case #" + t);
    System.out.printf("%d\n", 7122);
int[][] d = {{7,1,2,2},{8,7}};
     int g = Integer.parseInt("-123");
     long f = (long)d[0][2];
     List<Integer> l = new ArrayList<>();
     Random rg = new Random();
     for (int i = 9; i >= 0; --i) {
       l.add(Integer.valueOf(rg.nextInt(100) + 1));
       l.add(Integer.valueOf((int)(Math.random() * 100) + 1));
     Collections.sort(l, new Comparator<Integer>() {
       public int compare(Integer a, Integer b) { return a - b;
     });
     for (int i = 0; i < l.size(); ++i)</pre>
       System.out.print(l.get(i));
     Set<String> s = new HashSet<String>(); // TreeSet
     s.add("jizz");
     System.out.println(s);
     System.out.println(s.contains("jizz"));
    Map<String, Integer> m = new HashMap<String, Integer>();
m.put("lol", 7122);
     System.out.println(m);
     for(String key: m.keySet())
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol"));
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
    System.out.println(Math.acos(-1));
    BigInteger bi = in.nextBigInteger(), bj = new BigInteger("
-7122"), bk = BigInteger.valueOf(17171);
int sgn = bi.signum(); // sign(bi)
     bi = bi.subtract(BigInteger.ONE).multiply(bj).divide(bj).
     and(bj).gcd(bj).max(bj).pow(87);
    int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
```

### 9.4 Dancing Links

}

```
namespace dlx {
int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn], rw[maxn],
      bt[maxn], s[maxn], head, sz, ans;
void init(int c) {
  for (int i = 0; i < c; ++i) {
    up[i] = dn[i] = bt[i] = i;
lt[i] = i == 0 ? c : i - 1;
    rg[i] = i == c - 1 ? c : i + 1;
    s[i] = 0;
  rg[c] = 0, lt[c] = c - 1;
  up[c] = dn[c] = -1;
  head = c, sz = c + 1;
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
  int f = sz;
  for (int i = 0; i < (int)col.size(); ++i) {</pre>
    int c = col[i], v = sz++;
```

```
dn[bt[c]] = v;
    up[v] = bt[c], bt[c] = v;
    rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
                                                                           void solve(int 1, int r, vector<int> v, long long c) {
    rw[v] = r, cl[v] = c;
                                                                             if (l == r) {
                                                                                cost[qr[l].first] = qr[l].second;
    ++s[c];
    if (i > 0) lt[v] = v - 1;
                                                                                if (st[qr[l].first] == ed[qr[l].first]) {
                                                                                  printf("%lld\n", c);
                                                                                  return:
  lt[f] = sz - 1;
                                                                                int minv = qr[l].second;
void remove(int c) {
                                                                                for (int i = 0; i < (int)v.size(); ++i) minv = min(minv,
 lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
for (int i = dn[c]; i != c; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j])
                                                                                cost[v[i]]);
                                                                               printf("%lld\n", c + minv);
      up[dn[j]] = up[j], dn[up[j]] = dn[j], --s[cl[j]];
  }
                                                                             int m = (l + r) >> 1;
}
                                                                             vector < int > lv = v, rv = v;
void restore(int c) {
                                                                             vector<int> x, y;
  for (int i = up[c]; i != c; i = up[i]) {
  for (int j = lt[i]; j != i; j = lt[j])
                                                                             for (int i = m + 1; i \ll r; ++i) {
                                                                                cnt[qr[i].first]--
      ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
                                                                                if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first);
  lt[rg[c]] = c, rg[lt[c]] = c;
                                                                             contract(l, m, lv, x, y);
long long lc = c, rc = c;
// Call dlx::make after inserting all rows.
                                                                             djs.save();
void make(int c) {
                                                                             for (int i = 0; i < (int)x.size(); ++i) {</pre>
  for (int i = 0; i < c; ++i)
                                                                                lc += cost[x[i]];
    dn[bt[i]] = i, up[i] = bt[i];
                                                                               djs.merge(st[x[i]], ed[x[i]]);
void dfs(int dep) {
                                                                             solve(l, m, y, lc);
  if (dep >= ans) return;
                                                                             djs.undo();
  if (rg[head] == head) return ans = dep, void();
                                                                             x.clear(), y.clear();
for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;
for (int i = l; i <= m; ++i) {</pre>
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
  int w = c;
                                                                                cnt[qr[i].first]--;
  for (int x = c; x != head; x = rg[x]) if (s[x] < s[w]) w = x;
                                                                                if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first);
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
                                                                             contract(m + 1, r, rv, x, y);
    for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j]);
                                                                             djs.save();
    dfs(dep + 1);
                                                                             for (int i = 0; i < (int)x.size(); ++i) {</pre>
    for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j]);
                                                                                rc += cost[x[i]];
  }
                                                                               djs.merge(st[x[i]], ed[x[i]]);
  restore(w);
                                                                             solve(m + 1, r, y, rc);
int solve() {
                                                                             djs.undo();
 ans = 1e9, dfs(0);
                                                                             for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
 return ans;
```

### 9.5 Offline Dynamic MST

djs.undo();

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second =
    weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i such
    that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int> &x,
    vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
   if (cost[i] == cost[j]) return i < j;</pre>
   return cost[i] < cost[j];</pre>
 djs.save();
 for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first], ed[qr</pre>
    [i].first]);
 for (int i = 0; i < (int)v.size(); ++i) {</pre>
   if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
 djs.undo();
  djs.save();
 for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[x[i]],
    ed[x[i]]);
 for (int i = 0; i < (int)v.size(); ++i) {</pre>
   if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
```

#### 9.6 Manhattan Distance MST

```
void solve(int n) {
   init();
   vector<int> v(n), ds;
   for (int i = 0; i < n; ++i) {
     v\Gamma i = i:
     ds.push_back(x[i] - y[i]);
   sort(ds.begin(), ds.end());
   ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
   sort(v.begin(), v.end(), [&](int i, int j) { return x[i] == x
   [j] ? y[i] > y[j] : x[i] > x[j]; });
   int j = 0;
   for (int i = 0; i < n; ++i) {
     int p = lower_bound(ds.begin(), ds.end(), x[v[i]] - y[v[i
      ]]) - ds.begin() + 1;
     pair<int, int> q = query(p);
     // query return prefix minimum
     if (~q.second) add_edge(v[i], q.second)
     add(p, \ make\_pair(x[v[i]] + y[v[i]], \ v[i])); \\
}
 void make_graph() {
   solve(n);
   for (int i = 0; i < n; ++i) swap(x[i], y[i]);
   solve(n);
   for (int i = 0; i < n; ++i) x[i] = -x[i];
   solve(n);
   for (int i = 0; i < n; ++i) swap(x[i], y[i]);
   solve(n);
1}
```

### 9.7 IOI 2016 Alien trick

```
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative as well, depending on the problem.
    if (c - (1LL << d) < 0) continue;
    long long ck = c - (1LL << d);
    pair<long long, int> r = check(ck);
    if (r.second == k) return r.first - ck * k;
    if (r.second < k) c = ck;
  }
  pair<long long, int> r = check(c);
  return r.first - c * k;
}
```

### 9.8 Matroid Intersection

Start from  $S = \emptyset$ . In each iteration, let

- $Y_1 = \{x \not\in S \mid S \cup \{x\} \in I_1\}$
- $Y_2 = \{x \not\in S \mid S \cup \{x\} \in I_2\}$

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \not\in S$ , connect

- $x \to y$  if  $S \{x\} \cup \{y\} \in I_1$ .
- $y \to x \text{ if } S \{x\} \cup \{y\} \in I_2.$

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration.