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6.7 Schreier-Sims Algorithm 14 template <typename t=""> 6.8 Miller Rabin 15 inline bool rit(T& x) { 6.9 Pollard's Rho 15 char c = 0; bool flag = false; 6.10 Meissel-Lehmer Algorithm 15 while (c = getchar(), (c < '0' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' ' '</typename>				17	return 'p++,
6.8 Miller Rabin 15 inline bool rit(T& x) { char c = 0; bool flag = false; while (c = getchar(), (c < '0' or '0' or '0' or '1' or '1				1 t	template <typename t=""></typename>
6.9 Pollard's Rho 15 6.10 Meissel-Lehmer Algorithm 15 6.11 Discrete Logarithm 15 6.12 Gaussian Elimination 16 6.13 μ function 16 6.14 $\lfloor \frac{n}{i} \rfloor$ Enumeration 16 6.15 De Bruijn Sequence 16 6.17 Euclidean Algorithms 16 6.18 Chinese Remainder Theorem 17 6.19 Theorem 17 6.19.1 Kirchhoff's Theorem 17	9				
6.11 Discrete Logarithm 15 6.12 Gaussian Elimination 16 6.13 μ function 16 6.14 $\lfloor \frac{n}{i} \rfloor$ Enumeration 16 6.15 De Bruijn Sequence 16 6.16 Extended GCD 16 6.17 Euclidean Algorithms 16 6.18 Chinese Remainder Theorem 17 6.19 Theorem 17 6.19.1 Kirchhoff's Theorem 17				i	
6.12 Gaussian Elimination 16 6.13 μ function 16 6.14 $\lfloor \frac{n}{i} \rfloor$ Enumeration 16 6.15 De Bruijn Sequence 16 6.16 Extended GCD 16 6.17 Euclidean Algorithms 16 6.18 Chinese Remainder Theorem 17 6.19 Theorem 17 6.19.1 Kirchhoff's Theorem 17	6.10 Meissel-Lehmer Alge	orithm	. 15	5	while (c = getchar(), (c < '0' && c != '-') c > '9
6.13 μ function	6.11 Discrete Logarithm		. 15	5	
$ \begin{array}{cccccccccccccccccccccccccccccccccccc$				i	c == '-' ? (flag = true, x = 0) : (x = c - '0');
6.15 De Bruijn Sequence 16 6.16 Extended GCD 16 6.17 Euclidean Algorithms 16 6.18 Chinese Remainder Theorem 17 6.19 Theorem 17 6.19.1 Kirchhoff's Theorem 17	•			i i	
6.16 Extended GCD	- : -				
6.17 Euclidean Algorithms					
6.18 Chinese Remainder Theorem 17 6.19 Theorem 17 6.19.1 Kirchhoff's Theorem 17				17	
6.19.1 Kirchhoff's Theorem	_				'
6.19.1 Kirchhoff's Theorem					
10 T 1 •					1.9 Impropaga ata al- al-a
6.19.2 Tutte's Matrix					1.5 Increase stack size
6.19.3 Cayley's Formula					
6.19.4 Erdős-Gallai theorem					const int size = 256 << 20;
<pre>register long rsp asm("rsp");</pre>	0.20 I IIIIO5		11	ľ	<pre>register long rsp asm("rsp");</pre>
	7 Dynamic Programmin	ng	17		<pre>char *p = (char*)malloc(size) + size, *bak = (char*)rsp</pre>
7.1 Convex Hull Optimization				7	; (Human 0/0 0/0/) - Hard Hall (-)
	, -				asm("movq %0, %%rsp\n"::"r"(p));
7.3 Conditon				1	asm("movq %0, %%rsp\n"::"r"(bak));
7.3.1 totally monotone (concave/convex)	· ·	` '			

1.4 Pragma optimization

```
#pragma GCC optimize("Ofast", "no-stack-protector", "no
    -math-errno", "unroll-loops")
#pragma GCC target("sse,sse2,sse3,sse3,sse4,sse4.2,
   popcnt,abm,mmx,avx,tune=native,arch=core-avx2,tune=
    core-avx2"
#pragma GCC ivdep
```

2 Flow

2.1Dinic's Algorithm

```
struct dinic {
  static const int inf = 1e9;
  struct edge {
    int to, cap, rev;
    edge(int d, int c, int r): to(d), cap(c), rev(r) {}
  vector<edge> g[maxn];
  int qu[maxn], ql, qr;
  int lev[maxn];
  void init() {
  for (int i = 0; i < maxn; ++i)</pre>
       g[i].clear();
  void add_edge(int a, int b, int c) {
  g[a].emplace_back(b, c, g[b].size() - 0);
  g[b].emplace_back(a, 0, g[a].size() - 1);
  bool bfs(int s, int t) {
    memset(lev, -1, sizeof(lev));
    lev[s] = 0;
    ql = qr = 0;
    qu[qr++] = s;
    while (ql < qr) {
       int x = qu[ql++];
       for (edge &e : g[x]) if (lev[e.to] == -1 && e.cap
      > 0) {
         lev[e.to] = lev[x] + 1;
         qu[qr++] = e.to;
       }
    return lev[t] != -1;
  int dfs(int x, int t, int flow) {
    if (x == t) return flow;
    int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.to] ==
     lev[x] + 1) {
       int f = dfs(e.to, t, min(e.cap, flow - res));
      res += f;
       e.cap -= f;
       g[e.to][e.rev].cap += f;
    if (res == 0) lev[x] = -1;
    return res;
  int operator()(int s, int t) {
    int flow = 0;
    for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow;
};
```

Minimum-cost flow

```
struct mincost {
  struct edge {
    int dest, cap, w, rev;
    edge(int a, int b, int c, int d): dest(a), cap(b),
    w(c), rev(d) {}
 vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
 bool inq[maxn];
```

```
void init() {
     for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
   void add_edge(int a, int b, int c, int d) {
    g[a].emplace_back(b, c, +d, g[b].size() - 0);
    g[b].emplace_back(a, 0, -d, g[a].size() - 1);
   bool spfa(int s, int t, int &f, int &c) {
     for (int i = 0; i < maxn; ++i) {
       d[i] = inf;
p[i] = ed[i] = -1;
        inq[i] = false;
     d[s] = 0;
     queue<int> q;
     q.push(s);
     while (q.size()) {
  int x = q.front(); q.pop();
       inq[x] = false;
for (int i = 0; i < g[x].size(); ++i) {</pre>
          edge &e = g[x][i];
          if (e.cap > 0 \&\& d[e.dest] > d[x] + e.w) {
            d[e.dest] = d[x] + e.w;
            p[e.dest] = x;
            ed[e.dest] = i
            if (!inq[e.dest]) q.push(e.dest), inq[e.dest]
      = true;
          }
       }
     if (d[t] == inf) return false;
     int dlt = inf;
     for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[
     p[x]][ed[x]].cap);
     for (int x = t; x != s; x = p[x]) {
       edge &e = g[p[x]][ed[x]];
       e.cap -= dlt;
       g[e.dest][e.rev].cap += dlt;
     f += dlt; c += d[t] * dlt;
     return true;
   pair<int, int> operator()(int s, int t) {
     int f = 0, c = 0;
     while (spfa(s, t, f, c));
     return make_pair(f, c);
};
2.3 Gomory-Hu Tree
```

```
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;</pre>
  for(int i=2;i<=n;++i){</pre>
     int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
flow.walk(i); // bfs points that connected to i (
    use edges not fully flow)
     for(int j=i+1; j<=n;++j){</pre>
       if(g[j]==t && flow.connect(j))g[j]=i; // check if
      i can reach j
    }
  return rt;
```

2.4 Stoer-Wagner Minimum Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
 w[x][y] += c;
 w[y][x] += c;
```

```
pair<int, int> phase(int n) {
  memset(v, false, sizeof(v));
  memset(g, 0, sizeof(g));
  int s = -1, t = -1;
  while (true) {
     int c = -1;
     for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
       if (c == -1 || g[i] > g[c]) c = i;
     if (c == -1) break;
     v[c] = true;
     s = t, t = c;
     for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;</pre>
       g[i] += w[c][i];
  return make_pair(s, t);
int mincut(int n) {
  int cut = 1e9;
  memset(del, false, sizeof(del));
  for (int i = 0; i < n - 1; ++i)
     int s, t; tie(s, t) = phase(n);
     del[t] = true;
     cut = min(cut, g[t]);
for (int j = 0; j < n; ++j) {</pre>
       w[s][j] += w[t][j];
       w[j][s] += w[j][t];
    }
  return cut;
```

2.5 Kuhn-Munkres Algorithm

```
int w[maxn][maxn], lx[maxn], ly[maxn];
int match[maxn], slack[maxn];
bool vx[maxn], vy[maxn];
bool dfs(int x) {
  vx[x] = true;
  for (int i = 0; i < n; ++i) {
    if (vy[i]) continue;
if (lx[x] + ly[i] > w[x][i]) {
       slack[i] = min(slack[i], lx[x] + ly[i] - w[x][i])
       continue;
    vy[i] = true;
    if (match[i] == -1 || dfs(match[i])) {
       match[i] = x;
       return true;
  }
  return false;
int solve() {
  fill_n(match, n, -1);
  fill_n(lx, n, -inf);
fill_n(ly, n, 0);
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) lx[i] = max(lx[i], w[i])
     ][j]);
  for (int i = 0; i < n; ++i) {
    fill_n(slack, n, inf);
    while (true) {
       fill_n(vx, n, false);
       fill_n(vy, n, false);
       if (dfs(i)) break;
       int dlt = inf;
for (int j = 0; j < n; ++j) if (!vy[j]) dlt = min</pre>
     (dlt, slack[j]);
       for (int j = 0; j < n; ++j) {
  if (vx[j]) lx[j] -= dlt;
  if (vy[j]) ly[j] += dlt;
         else slack[j] -= dlt;
```

```
fint res = 0;
for (int i = 0; i < n; ++i) res += w[match[i]][i];
return res;
}</pre>
```

2.6 Flow Model

- Maximum/Minimum flow with lower/upper bound from s to t
 - 1. Construct super source S and sink T
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l
 - 3. For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v)
 - To maximize, connect $t \to s$ with capacity ∞ , and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge $(y \to x \text{ if } (x,y) \in M, x \to y \text{ otherwise})$
 - 2. DFS from unmatched vertices in X
 - 3. $x \in X$ is chosen iff x is unvisited
 - 4. $y \in Y$ is chosen iff y is visited
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x\to y$ with $(\cos t, cap)=(c,1)$ if c>0, otherwise connect $y\to x$ with $(\cos t, cap)=(-c,1)$
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=(0,d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C + K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer ${\cal T}$
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u\to v$ and $v\to u$ with capacity w
 - 5. For $v \in G,$ connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e)) 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|

3 Data Structure

3.1 < ext/pbds >

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>,
    rb_tree_tag, tree_order_statistics_node_update>
    tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
    // rb tree
    tree_set s;
```

```
s.insert(71); s.insert(22);
assert(*s.find_by\_order(0) == 22); assert(*s.
   find_by_order(1) == 71);
assert(s.order_of_key(22) == 0); assert(s.
   order_of_key(71) == 1);
s.erase(22);
assert(*s.find_by_order(0) == 71); assert(s.
   order_of_key(71) == 0);
// mergable heap
heap a, b; a.join(b);
// persistant
rope<char> r[2];
r[1] = r[0];
std::string st = "abc";
r[1].insert(0, st.c_str());
r[1].erase(1, 1);
std::cout << r[1].substr(0, 2) << std::endl;
return 0;
```

3.2 Li Chao Tree

```
namespace lichao {
struct line {
  long long a, b:
  line(): a(0), b(0) {}
  line(long long a, long long b): a(a), b(b) {}
  long long operator()(int x) const { return a * x + b;
line st[maxc * 4];
int sz, lc[maxc * 4], rc[maxc * 4];
int gnode() {
  st[sz] = line(1e9, 1e9);
  lc[sz] = -1, rc[sz] = -1;
  return sz++;
void init() {
  sz = 0;
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
  bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
  if (mcp) swap(st[o], tl);
  if (r - l == 1) return;
if (lcp != mcp) {
   if (lc[o] == -1) lc[o] = gnode();
}
    add(l, (l + r) / 2, tl, lc[o]);
  } else {
    if (rc[o] == -1) rc[o] = gnode();
    add((l + r) / 2, r, tl, rc[o]);
  }
long long query(int 1, int r, int x, int o) {
  if (r - l == 1) return st[o](x);
if (x < (l + r) / 2) {</pre>
    if (lc[o] == -1) return st[o](x);
    return min(st[o](x), query(l, (l + r) / 2, x, lc[o
    ]));
  } else {
    if (rc[o] == -1) return st[o](x);
    return min(st[o](x), query((l + r) / 2, r, x, rc[o
    J));
}}
```

4 Graph

4.1 Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
  int sum, v, rev, id;
  node(int s, int id): id(id), v(s), sum(s), rev(0), fa
    (nullptr), pfa(nullptr) {
    ch[0] = nullptr;
```

```
ch[1] = nullptr;
  int relation() {
    return this == fa->ch[0] ? 0 : 1;
  void push() {
    if (!rev) return
    swap(ch[0], ch[1]);
if (ch[0]) ch[0]->rev ^= 1;
    if (ch[1]) ch[1]->rev ^= 1;
    rev = 0:
  }
  void pull() {
    sum = v;
    if (ch[0]) sum += ch[0]->sum;
    if (ch[1]) sum += ch[1]->sum;
  void rotate()
    if (fa->fa) fa->fa->push();
    fa->push(), push();
swap(pfa, fa->pfa);
    int d = relation();
    node *t = fa;
    if (t->fa) t->fa->ch[t->relation()] = this;
    fa = t -> fa;
    t->ch[d] = ch[d \land 1];
    if (ch[d \land 1]) ch[d \land 1] -> fa = t;
    ch[d \land 1] = t;
    t->fa = this;
    t->pull(), pull();
  void splay() {
    while (fa) {
      if (!fa->fa) {
        rotate();
        continue:
      fa->fa->push();
      if (relation() == fa->relation()) fa->rotate(),
    rotate();
      else rotate(), rotate();
  void evert() {
    access();
    splay();
    rev ^= 1;
  void expose() {
    splay(), push();
    if (ch[1]) {
      ch[1]->fa = nullptr;
      ch[1]->pfa = this;
      ch[1] = nullptr;
      pull();
  }
  bool splice() {
    splay();
    if (!pfa) return false;
    pfa->expose();
    pfa->ch[1] = this;
    fa = pfa;
    pfa = nullptr;
    fa->pull();
    return true;
  void access() {
    expose():
    while (splice());
  int query() {
    return sum;
};
namespace lct {
node *sp[maxn];
void make(int u, int v) {
  // create node with id u and value v
  sp[u] = new node(v, u);
```

```
void link(int u, int v) {
  // u become v's parent
  sp[v]->evert();
  sp[v]->pfa = sp[u];
void cut(int u, int v) {
  // u was v's parent
  sp[u]->evert();
  sp[v] \rightarrow access(), sp[v] \rightarrow splay(), sp[v] \rightarrow push();
  sp[v]->ch[0]->fa = nullptr;
sp[v]->ch[0] = nullptr;
  sp[v]->pull();
void modify(int u, int v) {
  sp[u]->splay();
  sp[u] -> v = v
  sp[u]->pull();
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
  return sp[v]->query();
int find(int u) {
  sp[u]->access();
  sp[u]->splay();
  node *p = sp[u];
  while (true) {
    p->push():
    if (p->ch[0]) p = p->ch[0];
    else break;
  return p->id;
```

4.2 Heavy-Light Decomposition

```
void dfs(int x, int p) {
  dep[x] = ~p ? dep[p] + 1 : dep[x];
  sz[x] = 1;
  to[x] = -1;
  fa[x] = p;
  for (const int &u : g[x]) {
    if (u == p) continue;
    dfs(u, x);
    sz[x] += sz[u];
if (to[x] == -1 || sz[to[x]] < sz[u]) to[x] = u;
void hld(int x, int t) {
  static int tk = 0;
  fr[x] = t;
  dfn[x] = tk++;
  if (!~to[x]) return;
  hld(to[x], t);
  for (const int &u : g[x]) {
    if (u == fa[x] || u == to[x]) continue;
    hld(u, u);
vector<pair<int, int>> get(int x, int y) {
  int fx = fr[x], fy = fr[y];

  vector<pair<int, int>> res;
  while (fx != fy) {
    if (dep[fx] < dep[fy]) {
  swap(fx, fy);</pre>
       swap(x, y);
    }
    res.emplace_back(dfn[fx], dfn[x] + 1);
    x = fa[fx];
    fx = fr[x];
  res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn
    [y]) + 1)
  int lca = (dep[x] < dep[y] ? x : y);
  return res;
```

4.3 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
  sz[now] = 1; mx[now] = 0;
for (int u : G[now]) if (!v[u]) {
    get_center(u)
    mx[now] = max(mx[now], sz[u]);
    sz[now] += sz[u];
}
void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
  v[now] = true;
  for (auto u : G[now]) if (!v[u.first]) {
    get_dis(u, d, len + u.second);
}
void dfs(int now, int fa, int d) {
  get_center(now);
  int c = -1;
  for (int i : vtx) {
    if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx</pre>
     .size() / 2) c = i;
    v[i] = false;
  get_dis(c, d, 0);
  for (int i : vtx) v[i] = false;
  v[c]
       = true; vtx.clear();
  dep[c] = d; p[c] = fa;
  for (auto u : G[c]) if (u.first != fa && !v[u.first])
    dfs(u.first, c, d + 1);
}
```

4.4 Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
long long d[1003][1003],dp[1003][1003];
pair<long long,long long> MMWC(){
  memset(dp,0x3f,sizeof(dp))
  for(int i=1;i<=n;++i)dp[0][i]=0;</pre>
  for(int i=1;i<=n;++i){</pre>
    for(int j=1; j<=n; ++j){
  for(int k=1; k<=n; ++k){</pre>
         dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
    }
  long long au=1ll<<31,ad=1;</pre>
  for(int i=1;i<=n;++i)</pre>
     if(dp[n][i]==0x3f3f3f3f3f3f3f3f3f)continue;
     long long u=0,d=1;
     for(int j=n-1;j>=0;--j){
  if((dp[n][i]-dp[j][i])*d>u*(n-j)){
         u=dp[n][i]-dp[j][i];
         d=n-j;
       }
     if(u*ad<au*d)au=u,ad=d;
  long long g=\_gcd(au,ad)
  return make_pair(au/g,ad/g);
}
```

4.5 Minimum Steiner Tree

```
namespace steiner {
// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
// z[i] = the weight of the i-th vertex
const int maxn = 64, maxk = 10;
const int inf = 1e9;
int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[
    maxn];
void init(int n) {
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) w[i][j] = inf;
    z[i] = 0;</pre>
```

```
w[i][i] = 0;
void add_edge(int x, int y, int d) {
  w[x][y] = min(w[x][y], d);
  w[y][x] = min(w[y][x], d);
void build(int n) {
  for (int i = 0; i < n; ++i) {
  for (int j = 0; j < n; ++j) {
    w[i][j] += z[i];
       if (i != j) w[i][j] += z[j];
    }
  for (int k = 0; k < n; ++k) {
     for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j
], w[i][k] + w[k][j] - z[k]);</pre>
  }
int solve(int n, vector<int> mark) {
  build(n);
  int k = (int)mark.size();
  assert(k < maxk);</pre>
  for (int s = 0; s < (1 << k); ++s) {
  for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
  for (int i = 0; i < n; ++i) dp[0][i] = 0;
  for (int s = 1; s < (1 << k); ++s) {
  if (__builtin_popcount(s) == 1) {</pre>
       int x = __builtin_ctz(s);
        for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x
     ]][i];
       continue;
     for (int i = 0; i < n; ++i) {
       for (int sub = s \& (s - 1); sub; sub = s \& (sub - 1);
      1)) {
          dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s ^
     sub][i] - z[i]);
       }
     for (int i = 0; i < n; ++i) {
       off[i] = inf;
     for (int j = 0; j < n; ++j) off[i] = min(off[i],
dp[s][j] + w[j][i] - z[j]);</pre>
     for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i]
     ], off[i]);
  int res = inf;
  for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k
     ) - 1][i]);
  return res;
}}
```

4.6 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
   T g[maxn][maxn], fw[maxn];
   int n, fr[maxn];
   bool vis[maxn], inc[maxn];
   void clear() {
      for(int i = 0; i < maxn; ++i) {
            for(int j = 0; j < maxn; ++j) g[i][j] = inf;
            vis[i] = inc[i] = false;
      }
   }
   void addedge(int u, int v, T w) {
      g[u][v] = min(g[u][v], w);
   }
   T operator()(int root, int _n) {
      n = _n;
      if (dfs(root) != n) return -1;
      T ans = 0;
   while (true) {
      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =
      i;
      for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
```

```
for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
               fw[i] = g[j][i];
               fr[i] = j;
        int x = -1;
        for (int i = 1; i <= n; ++i) if (i != root &&!
      inc[i]) {
          int j = i, c = 0;
          while (j != root && fr[j] != i && c <= n) ++c,
      j = fr[j];
          if (j == root || c > n) continue;
else { x = i; break; }
        if (!~x) {
          for (int i = 1; i <= n; ++i) if (i != root &&!
      inc[i]) ans += fw[i];
          return ans;
        int y = x;
        for (int i = 1; i <= n; ++i) vis[i] = false;
do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =</pre>
      true; } while (y != x);
        inc[x] = false;
        for (int k = 1; k \le n; ++k) if (vis[k])
          for (int j = 1; j <= n; ++j) if (!vis[j]) {
             if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x]</pre>
     ]) g[j][x] = g[j][k] - fw[k];
       }
     }
     return ans;
   int dfs(int now) {
     int r = 1;
     vis[now] = true;
      for (int i = 1; i <= n; ++i) if (g[now][i] < inf &&
       !vis[i]) r += dfs(i);
     return r;
  }
};
```

4.7 Maximum Matching on General Graph

```
namespace matching {
int fa[maxn], pre[maxn], match[maxn], s[maxn], v[maxn];
vector<int> g[maxn];
queue<int> q;
void init(int n) {
  for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
  for (int i = 0; i < n; ++i) g[i].clear();</pre>
void add_edge(int u, int v) {
  g[u].push_back(v);
  g[v].push_back(u);
int find(int u) {
  if (u == fa[u]) return u;
  return fa[u] = find(fa[u]);
int lca(int x, int y, int n) {
  static int tk = 0;
  tk++
  x = find(x), y = find(y);
  for (; ; swap(x, y)) {
  if (x != n) {
       if (v[x] == tk) return x;
       v[x] = tk;
       x = find(pre[match[x]]);
  }
void blossom(int x, int y, int l) {
  while (find(x) != l) {
     pre[x] = y
     y = match[x]
     if (s[y] == 1) {
```

```
q.push(y);
       s[y] = 0;
    if (fa[x] == x) fa[x] = l;
if (fa[y] == y) fa[y] = l;
    x = pre[y];
bool bfs(int r, int n) {
  for (int i = 0; i <= n; ++i) {</pre>
    fa[i] = i;
    s[i] = -1;
  while (!q.empty()) q.pop();
  q.push(r);
  s[r] = 0;
  while (!q.empty()) {
  int x = q.front(); q.pop();
     for (int u : g[x]) {
       if (s[u] = -1) {
         pre[u] = x;
         s[u] = 1;
         if (match[u] == n) {
            for (int a = u, b = x, last; b != n; a = last
      b = pre[a]
              last = match[b], match[b] = a, match[a] = b
            return true;
         }
         q.push(match[u]);
         s[match[u]] = 0
       else\ if\ (!s[u]\ \&\&\ find(u)\ !=\ find(x))\ \{
         int l = lca(u, x, n);
blossom(x, u, l);
         blossom(u, x, 1);
    }
  return false;
int solve(int n) {
  int res = 0;
for (int x = 0; x < n; ++x) {
    if (match[x] == n) res += bfs(x, n);
  return res;
```

4.8 Maximum Weighted Matching on General Graph

```
struct WeightGraph {
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
  int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) {
   return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
  void update_slack(int u, int x) {
    if (!slack[x] | | e_delta(g[u][x]) < e_delta(g[slack])
    [x]][x])) slack[x] = u;
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
      if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
```

```
update_slack(u, x);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++)
  q_push(flo[x][i]);
void set_st(int x, int b) {
  st[x] = b;
  if(x > n) for (size_t i = 0; i < flo[x].size(); ++
  i) set_st(flo[x][i], b);
int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) -
  flo[b].begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
}
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
  flo[u][i ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
  end());
void augment(int u, int v) {
  for (; ; ) {
  int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
}
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b \leftarrow n_x && st[b]) ++b;
  if (b > n_x) ++n_x;
lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[
  match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end())
  for (int x = v, y; x != lca; x = st[pa[y]])
  flo[b].push_back(x), flo[b].push_back(y = st[
  match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
       if (g[b][x].w == 0 \mid \mid e_{delta}(g[xs][x]) <
  e_delta(g[b][x])
    g[b][x] = g[xs][x], g[x][b] = g[x][xs];
for (int x = 1; x <= n; ++x)
       if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
```

```
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, y)
  for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];</pre>
     pa[xs] = g[xns][xs].u;
     S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
     q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
     int xs = flo[b][i];
     S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
   int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
     pa[v] = e.u, S[v] = 1;
     int nu = st[match[v]];
     slack[v] = slack[nu] = 0;
  S[nu] = 0, q_push(nu);
} else if (S[v] == 0) {
     int lca = get_lca(u, v);
     if (!lca) return augment(u,v), augment(v,u), true
     else add_blossom(u, lca, v);
  }
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0,
    q_push(x);
   if (q.empty()) return false;
  for (; ; ) {
  while (q.size()) {
       int u = q.front(); q.pop();
if (S[st[u]] == 1) continue;
for (int v = 1; v <= n; ++v)</pre>
          if (g[u][v].w > 0 && st[u] != st[v]) {
             if (e_delta(g[u][v]) == 0) {
  if (on_found_edge(g[u][v])) return true;
            } else update_slack(u, st[v]);
          }
     int d = inf;
     for (int b = n + 1; b <= n_x; ++b)
if (st[b] == b && S[b] == 1) d = min(d, lab[b]
   / 2);
     for (int x = 1; x <= n_x; ++x)
  if (st[x] == x && slack[x]) {</pre>
          if (S[\bar{x}] == -1) d = min(d, e_delta(g[slack[x]])
   ]][x]));
  else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x]) / 2);
     for (int u = 1; u <= n; ++u) {
       if (S[st[u]] == 0) {
          if (lab[u] <= d) return 0;</pre>
       lab[u] -= d;
} else if (S[st[u]] == 1) lab[u] += d;
     for (int b = n + 1; b \ll n_x; ++b)
       if (st[b] == b) {
          if (S[st[b]] == 0) lab[b] += d * 2;
          else if (S[st[b]] == 1) lab[b] -= d * 2;
     q = queue<int>();
     for (int x = 1; x <= n_x; ++x)
        if (st[x] == x \&\& slack[x] \&\& st[slack[x]] != x
    && e_delta(g[slack[x]][x]) == 0)
          if (on_found_edge(g[slack[x]][x])) return
     for (int b = n + 1; b \le n_x; ++b)
```

```
if (st[b] == b \&\& S[b] == 1 \&\& lab[b] == 0)
     expand_blossom(b);
     return false;
   pair<long long, int> solve() {
     memset(match + 1, 0, sizeof(int) * n);
     n_x = n;
     int n_matches = 0;
     long long tot_weight = 0;
     for (int u = 0; u \le n; ++u) st[u] = u, flo[u].
     clear();
     int w_max = 0;
     for (int u = 1; u <= n; ++u)
for (int v = 1; v <= n; ++v) {
          flo_from[u][v] = (u == v ? u : 0);
          w_max = max(w_max, g[u][v].w);
     for (int u = 1; u <= n; ++u) lab[u] = w_max;
     while (matching()) ++n_matches;
     for (int u = 1; u <= n; _++u)
       if (match[u] && match[u] < u)</pre>
          tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int ui, int vi, int wi) {
     g[ui][vi].w = g[vi][ui].w = wi;
   void init(int _n) {
     n = _n;
for (int u = 1; u <= n; ++u)
       for (int v=1; v <= n; ++v)
         g[u][v] = edge(u, v, 0);
};
4.9 Maximum Clique
struct MaxClique {
   // change to bitset for n > 64.
   int n, deg[maxn];
   uint64_t adj[maxn], ans;
   vector<pair<int, int>> edge;
   void init(int n_) {
     fill(adj, adj + n, 0ull);
     fill(deg, deg + n, 0);
     edge.clear();
   void add_edge(int u, int v) {
  edge.emplace_back(u, v);
     ++deg[u], ++deg[v];
   vector<int> operator()() {
     vector<int> ord(n);
     iota(ord.begin(), ord.end(), 0);
sort(ord.begin(), ord.end(), [&](int u, int v) {
  return deg[u] < deg[v]; });</pre>
     vector<int> id(n);
for (int i = 0; i < n; ++i) id[ord[i]] = i;
for (auto e : edge) {</pre>
       int u = id[e.first], v = id[e.second];
       adj[u] |= (1ull << v);
adj[v] |= (1ull << u);
     uint64_t r = 0, p = (1ull << n) - 1;
     ans = 0;
     dfs(r, p);
     vector<int> res;
     for (int i = 0; i < n; ++i) {
       if (ans >> i & 1) res.push_back(ord[i]);
     return res;
#define pcount __builtin_popcountll
   void dfs(uint64_t r, uint64_t p) {
```

 $if (p == 0) {$

return;

if (pcount(r) > pcount(ans)) ans = r;

```
if (pcount(r | p) <= pcount(ans)) return;
int x = __builtin_ctzll(p & -p);
uint64_t c = p & ~adj[x];
while (c > 0) {
    // bitset._Find_first(); bitset._Find_next();
    x = __builtin_ctzll(c & -c);
    r != (1ull << x);
    dfs(r, p & adj[x]);
    r &= ~(1ull << x);
    p &= ~(1ull << x);
    c ^= (1ull << x);
    c ^= (1ull << x);
}</pre>
```

4.10 Tarjan's Articulation Point

```
void dfs(int x, int p) {
  tin[x] = low[x] = ++t;
  int ch = 0;
  for (auto u : g[x]) if (u.first != p) {
    if (!ins[u.second]) st.push(u.second), ins[u.second
    ] = true:
    if (tin[u.first])
      low[x] = min(low[x], tin[u.first]);
      continue;
    }
    ++ch;
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
    if (low[u.first] >= tin[x]) {
      cut[x] = true;
      ++SZ;
      while (true) {
        int e = st.top(); st.pop();
        bcc[e] = sz;
        if (e == u.second) break;
   }
  if (ch == 1 && p == -1) cut[x] = false;
```

4.11 Tarjan's Bridge

```
void dfs(int x, int p) {
  tin[x] = low[x] = ++t;
  st.push(x);
  for (auto u : g[x]) if (u.first != p) {
    if (tin[u.first]) {
       low[x] = min(low[x], tin[u.first]);
       continue;
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
if (low[u.first] == tin[u.first]) br[u.second] =
  if (tin[x] == low[x]) {
    while (st.size()) {
       int u = st.top(); st.pop();
       bcc[u] = sz;
       if (u == x) break;
  }
}
```

4.12 Dominator Tree

```
namespace dominator {
vectorsint> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[
    maxn], val[maxn], rp[maxn], tk;
void init(int n) {
```

```
^{\prime} vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1)
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
  fill(rp, rp + n, -1)
  fill(dom, dom + n, -1);
  tk = 0;
  for (int i = 0; i < n; ++i)
     g[i].clear();
void add_edge(int x, int y) {
  g[x].push_back(y);
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  for (int &u : g[x]) {
     if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
     r[dfn[u]].push_back(dfn[x]);
}
void merge(int x, int y) {
  fa[x] = y;
int find(int x, int c = 0) {
  if (fa[x] == x) return x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[
     x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
  // return the father of each node in the dominator
     tree
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
     for (int &u : r[i]) sdom[i] = min(sdom[i], sdom[
     find(u)]);
     if (i) rdom[sdom[i]].push_back(i);
     for (int &u : rdom[i]) {
       int p = find(u);
if (sdom[p] == i) dom[u] = i;
       else dom[u] = p;
     if (i) merge(i, rp[i]);
  vector<int> p(n, -1);
for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i])</pre>
     dom[i] = dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
}}
```

4.13 System of Difference Constraints

Given m constrains on n variables x_1, x_2, \ldots, x_n of form $x_i - x_j \leq w$ (resp, $x_i - x_j \geq w$), connect $i \to j$ with weight w. Then connect $0 \to i$ for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to x_i .

5 String

5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  // f[i] = length of the longest prefix (excluding s
      [0:i]) such that it coincides with the suffix of s
      [0:i] of the same length
  // i + 1 - f[i] is the length of the smallest
      recurring period of s[0:i]
  int k = 0;
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
```

```
while (k > 0 && s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
    f[i] = k;
}
return f;
}
vector<int> search(const string &s, const string &t) {
    // return 0-indexed occurrence of t in s
    vector<int> f = kmp(t), res;
    int k = 0;
    for (int i = 0; i < (int)s.size(); ++i) {
        while (k > 0 && (k == (int)t.size() || s[i] != t[k]
        ])) k = f[k - 1];
        if (s[i] == t[k]) ++k;
        if (k == (int)t.size()) res.push_back(i - t.size() + 1);
    }
    return res;
}
```

5.2 Z Algorithm

5.3 Manacher's Algorithm

5.4 Aho-Corasick Automaton

```
struct AC {
    static const int maxn = 1e5 + 5;
    int sz, ql, qr, root;
    int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn
        ][26], f[maxn];
    int gnode() {
        for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
        f[sz] = -1;
        ed[sz] = 0;
        cnt[sz] = 0;
        return sz++;
    }
    void init() {
        sz = 0;
        root = gnode();
    }
}</pre>
```

```
int add(const string &s) {
     int now = root;
for (int i = 0; i < s.length(); ++i) {</pre>
       if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a']
     ] = gnode();
       now = ch[now][s[i] - 'a'];
     ed[now] = 1;
     return now;
   void build_fail() {
     ql = qr = 0; q[qr++] = root;
     while (ql < qr) {
       int now = q[ql++];
       for (int i = 0; i < 26; ++i) if (ch[now][i] !=</pre>
     -1) {
          int p = ch[now][i], fp = f[now];
          while (fp != -1 && ch[fp][i] == -1) fp = f[fp];
          int pd = fp != -1 ? ch[fp][i] : root;
          f[p] = pd;
          el[p] = ed[pd] ? pd : el[pd];
          q[qr++] = p;
       }
     }
   void build(const string &s) {
     build_fail();
     int now = root;
     for (int i = 0; i < s.length(); ++i) {
  while (now != -1 && ch[now][s[i] - 'a'] == -1)</pre>
     now = f[now];
       now = now != -1 ? ch[now][s[i] - 'a'] : root;
       ++cnt[now];
     for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] +=
     cnt[q[i]];
   long long solve(int n) {
     build_fail();
     vector<vector<long long>> dp(sz, vector<long long>(
     n + 1, 0));
     for (int i = 0; i < sz; ++i) dp[i][0] = 1;
     for (int i = 1; i <= n; ++i) {
       for (int j = 0; j < sz; ++j) {
  for (int k = 0; k < 2; ++k) {
    if (ch[j][k]]!= -1) {
              if (!ed[ch[j][k]])
                 dp[j][i] += dp[ch[j][k]][i - 1];
            } else {
              int z = f[j];
              while (z != root \&\& ch[z][k] == -1) z = f[z]
     ];
              int p = ch[z][k] == -1 ? root : ch[z][k];
              if (ch[z][k] == -1 \mid \mid \cdot \mid ed[ch[z][k]]) dp[j][
     i] += dp[p][i - 1];
            }
         }
       }
     return dp[0][n];
};
```

5.5 Suffix Automaton

```
struct SAM {
    static const int maxn = 5e5 + 5;
    int nxt[maxn][26], to[maxn], len[maxn];
    int root, last, sz;
    int gnode(int x) {
        for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
        to[sz] = -1;
        len[sz] = x;
        return sz++;
    }
    void init() {
        sz = 0;
        root = gnode(0);
        last = root;</pre>
```

```
void push(int c) {
     int cur = last;
     last = gnode(len[last] + 1);
     for (; ~cur && nxt[cur][c] == -1; cur = to[cur])
     nxt[cur][c] = last;
     if (cur == -1) return to[last] = root, void();
     int link = nxt[cur][c];
     if (len[link] == len[cur] + 1) return to[last] =
     link, void();
     int tlink = gnode(len[cur] + 1);
for (; ~cur && nxt[cur][c] == link; cur = to[cur])
     nxt[cur][c] = tlink;
     for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[</pre>
     link][i];
     to[tlink] = to[link];
     to[link] = tlink;
to[last] = tlink;
   void add(const string &s) {
     for (int i = 0; i < s.size(); ++i) push(s[i] - 'a')</pre>
  bool find(const string &s) {
     int cur = root;
     for (int i = 0; i < s.size(); ++i) {
  cur = nxt[cur][s[i] - 'a'];</pre>
       if (cur == -1) return false;
     return true;
   int solve(const string &t) {
     int res = 0, cnt = 0;
     int cur = root;
     for (int i = 0; i < t.size(); ++i) {</pre>
       if (~nxt[cur][t[i] - 'a']) {
         ++cnt;
         cur = nxt[cur][t[i] - 'a'];
       } else {
  for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur
     = to[cur]);
     if (~cur) cnt = len[cur] + 1, cur = nxt[cur][t[
i] - 'a'];
         else cnt = 0, cur = root;
       res = max(res, cnt);
     return res;
|};
```

5.6 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2], x[maxn], p
     [maxn], q[maxn * 2];
// sa[i]: sa[i]-th suffix is the i-th lexigraphically
     smallest suffix.
// hi[i]: longest common prefix of suffix sa[i] and
     suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa, int *c, int *s, bool *t, int n,
     int z) {
  memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -
     1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
       - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
     *c, int n, int z) {
  bool uniq = t[n - 1] = true;
   int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
     last = -1;
```

```
memset(c, 0, sizeof(int) * z);
   for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
     return:
   for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i]
     + 1] ? t[i + 1] : s[i] < s[i + 1]);
   pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i -</pre>
      1]) sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] &&
      !t[sa[i] - 1]) {
     bool neq = last < 0 \mid \mid memcmp(s + sa[i], s + last,
      (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
     ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
   pre(sa, c, n, z);
   for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i
      ]]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
void build(const string &s) {
   for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
   _s[(int)s.size()] = 0; // s shouldn't contain 0
   sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for (int i = 0; i < (int)s.size(); ++i) sa[i] = sa[i</pre>
     + 1];
   for (int i = 0; i < (int)s.size(); ++i) rev[sa[i]] =</pre>
     i;
   int ind = 0; hi[0] = 0;
   for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
        ind = 0;
        continue;
     while (i + ind < (int)s.size() && s[i + ind] == s[
sa[rev[i] - 1] + ind]) ++ind;</pre>
     hi[rev[i]] = ind ? ind-- : 0;
}}
```

5.7 Lexicographically Smallest Rotation

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
  int i = 0, j = 1;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && t[i + k] == t[j + k]) ++k;
    if (t[i + k] <= t[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

6 Math

6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(
    re + rhs.re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(
    re - rhs.re, im - rhs.im); }
```

```
cplx operator*(const cplx &rhs) const { return cplx(
  re * rhs.re - im * rhs.im, re * rhs.im + im * rhs.
     re); }
  cplx conj() const { return cplx(re, -im); }
const int maxn = 262144;
const double pi = acos(-1);
cplx omega[maxn + 1];
bool init:
void prefft() {
  for (int i = 0; i \le maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi
     * i / maxn));
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    int x = 0:
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j &
    1) << (z - j);
    if (x > i) swap(v[x], v[i]);
  }
void fft(vector<cplx> &v, int n) {
  if (!init) {
    init = true;
    prefft();
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
  int z = s >> 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
         cplx x = v[i + z + k] * omega[maxn / s * k];
         v[i + z + k] = v[i + k] - x;
         v[i + k] = v[i + k] + x;
    }
 }
void ifft(vector<cplx> &v, int n) {
  fft(v, n);
  reverse(v.begin() + 1, v.end());
  for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n
    , 0);
vector<long long> convolution(const vector<int> &a,
    const vector<int> &b) {
  // Should be able to handle N <= 10^5, C <= 10^4
  int sz = 1;
  while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
  vector<cplx> v(sz);
  for (int i = 0; i < sz; ++i) {
    double re = i < a.size() ? a[i] : 0;
    double im = i < b.size() ? b[i] : 0;
    v[i] = cplx(re, im);
  fft(v, sz);
  for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);
  rely ((sz - 1);
    cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()
    ) * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v
[i].conj()) * cplx(0, -0.25);
    v[i] = x;
  ifft(v, sz);
  vector<long long> c(sz);
  for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
vector<int> convolution_mod(const vector<int> &a, const
     vector<int> &b, int p) {
  int sz = 1;
  while (sz < (int)a.size() + (int)b.size() - 1) sz <<=</pre>
  vector<cplx> fa(sz), fb(sz);
  for (int i = 0; i < (int)a.size(); ++i) {
  int x = (a[i] % p + p) % p;</pre>
    fa[i] = cp[x(x \& ((1 << 15) - 1), x >> 15);
  for (int i = 0; i < (int)b.size(); ++i) {
```

```
int x = (b[i] \% p + p) \% p
     fb[i] = cplx(x & ((1 << 15) - 1), x >> 15);
  fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
   for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
cplx b1 = (fb[i] + fb[j].conj()) * r3;
     cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
        cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
        cplx d1 = (fb[j] + fb[i].conj()) * r3;
        cplx d2 = (fb[j] - fb[i].conj()) * r4;
fa[i] = c1 * d1 + c2 * d2 * r5;
        fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \%
  }
   return res;
}}
```

6.2 Number Theoretic Transform

```
template <long long mod, long long root>
struct NTT {
  vector<long long> omega;
  NTT() {
    omega.resize(maxn + 1);
    long long x = fpow(root, (mod - 1) / maxn);
    omega[0] = 111;
for (int i = 1; i <= maxn; ++i)
      omega[i] = omega[i - 1] * x % mod;
  long long fpow(long long a, long long n) {
    (n += mod - 1) \% = mod - 1;
    long long r = 1;
    for (; n; n >>= 1) {
  if (n & 1) (r *= a) %= mod;
      (a *= a) \% = mod;
    return r;
  void bitrev(vector<long long> &v, int n) {
    int z = __builtin_ctz(n) - 1;
    for (int i = 0; i < n; ++i) {
       int x = 0;
      for (int j = 0; j \ll z; ++j) x ^= (i >> j & 1) <<
     (z - j);
       if (x > i) swap(v[x], v[i]);
  void ntt(vector<long long> &v, int n) {
    bitrev(v, n);
    for (int s = 2; s <= n; s <<= 1) {
      int z = s \gg 1;
       for (int i = 0; i < n; i += s) {
         for (int k = 0; k < z; ++k) {
  long long x = v[i + k + z] * omega[maxn / s *</pre>
           v[i + k + z] = (v[i + k] + mod - x) \% mod;
           (v[i + k] += x) \% = mod;
      }
    }
  void intt(vector<long long> &v, int n) {
```

```
ntt(v, n);
    for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i
    ]);
    long long inv = fpow(n, -1);
    for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;
  vector<long long> operator()(vector<long long> a,
    vector<long long> b) {
    int sz = 1;
    while (sz < a.size() + b.size() - 1) sz <<= 1;
    while (a.size() < sz) a.push_back(0);
while (b.size() < sz) b.push_back(0);</pre>
    ntt(a, sz), ntt(b, sz);
    vector<long long> c(sz);
    for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] %
    mod:
    intt(c, sz);
    return c;
};
vector<long long> convolution(vector<long long> a,
    vector<long long> b) {
  NTT<mod1, root1> conv1;
  NTT<mod2, root2> conv2;
  vector<long long> pa(a.size()), pb(b.size())
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i])
     % mod1 + mod1) % mod1;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i]
     % mod1 + mod1) % mod1;
  vector<long long> c1 = conv1(pa, pb);
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i]
     % mod2 + mod2) % mod2;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i]</pre>
     % \mod 2 + \mod 2) \% \mod 2;
  vector<long long> c2 = conv2(pa, pb);
  long long x = conv2.fpow(mod1, -1);
  long long y = conv1.fpow(mod2, -1);
  long long prod = mod1 * mod2;
  vector<long long> res(c1.size());
  for (int i = 0; i < c1.size(); ++i) {</pre>
    long long z = ((ull)fmul(c1[i] * mod2 % prod, y,
prod) + (ull)fmul(c2[i] * mod1 % prod, x, prod)) %
    prod;
    if (z \ge prod / 2) z = prod;
    res[i] = z;
  }
  return res;
```

6.2.1 NTT Prime List

Prime	Root	Prime	Root
7681	17	167772161	3
12289	11	104857601	3
40961	3	985661441	3
65537	3	998244353	3
786433	10	1107296257	10
5767169	3	2013265921	31
7340033	3	2810183681	11
23068673	3	2885681153	3
469762049	3	605028353	3

6.3 Polynomial Division

```
vector<int> inverse(const vector<int> &v, int n) {
  vector<int> q(1, fpow(v[0], mod - 2));
  for (int i = 2; i <= n; i <<= 1) {
     vector<int> fv(v.begin(), v.begin() + i);
     vector<int> fq(q.begin(), q.end());
     fv.resize(2 * i), fq.resize(2 * i);
     ntt(fq, 2 * i), ntt(fv, 2 * i);
     for (int j = 0; j < 2 * i; ++j) {
        fv[j] = fv[j] * 1ll * fq[j] % mod * 1ll * fq[j] %
        mod;
     }
     intt(fv, 2 * i);
     vector<int> res(i);
     for (int j = 0; j < i; ++j) {
        res[j] = mod - fv[j];
     }
}</pre>
```

```
if (j < (i >> 1)) (res[j] += 2 * q[j] % mod) %=
    mod;
    q = res;
  return q;
vector<int> divide(const vector<int> &a, const vector<</pre>
    int> &b) {
  // leading zero should be trimmed
  int n = (int)a.size(), m = (int)b.size();
  int k = 2;
  while (k < n - m + 1) k <<= 1;
  vector<int> ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - i -
     1];
  for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - i -
     1];
  vector<int> rbi = inverse(rb, k);
  vector<int> res = convolution(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
```

6.4 Polynomial Square Root

```
// Find G(x) such that G^2(x) = F(x) \pmod{x^{N+1}}
vector<int> solve(vector<int> b, int n) {
   if (n == 1) return \{sqr[b[0]]\}
   vector<int> h = solve(b, n >> 1); h.resize(n);
   vector<int> c = inverse(h, n);
   h.resize(n << 1); c.resize(n << 1);
   vector<int> res(n << 1);</pre>
   conv.ntt(h, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) b[i] = 0;
   conv.ntt(b, n << 1);
   conv.ntt(c, n << 1);</pre>
   for (int i = 0; i < (n << 1); ++i) res[i] = 1ll * (h[
   i] + 1ll * c[i] * b[i] % mod) % mod * inv2 % mod;</pre>
   conv.intt(res, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) res[i] = 0;
   return res;
}
```

6.5 Fast Walsh-Hadamard Transform

6.5.1 XOR Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_0) tf(A_1))$
- $utf(A) = (utf(\frac{A_0 + A_1}{2}), utf(\frac{A_0 A_1}{2}))$

6.5.2 OR Convolution

- $tf(A) = (tf(A_0), tf(A_0) + tf(A_1))$
- $utf(A) = (utf(A_0), utf(A_1) utf(A_0))$

6.5.3 AND Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_1))$
- $utf(A) = (utf(A_0) utf(A_1), utf(A_1))$

6.6 Simplex Algorithm

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
    't exist
// return vector<double>(n, +inf) if the solution is
    unbounded
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
vector<vector<double>>> d;
```

```
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {
    if (i != r && j != s) d[i][j] -= d[r][j] * d[i][s
       * inv;
  for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s]
  for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j]
     *= +inv;
  d[r][s] = inv;
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
while (true) {
    int s = -1;
for (int i = 0; i <= n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1 || d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;</pre>
       if (r == -1 \mid | d[i][n + 1] / d[i][s] < d[r][n +
     1] / d[r][s]) r = i;
     if (r == -1) return false;
    pivot(r, s);
vector<double> solve(const vector<vector<double>> &a,
     const vector<double> &b, const vector<double> &c) {
  m = b.size(), n = c.size();
  d = vector<vector<double>>(m + 2, vector<double>(n +
  for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
  for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] =
   -1, d[i][n + 1] = b[i];
for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i]</pre>
  q[n] = -1, d[m + 1][n] = 1;
  int r = 0;
  for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n]
     + 1]) r = i;
  if (d[r][n + 1] < -eps) {</pre>
    pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps) return
     vector<double>(n, -inf);
     for (int i = 0; i < m; ++i) if (p[i] == -1) {
      int s = min_element(d[i].begin(), d[i].end() - 1)
      - d[i].begin();
       pivot(i, s);
  if (!phase(0)) return vector<double>(n, inf);
  vector<double> x(n);
  for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d
     [i][n + 1];
  return x;
}}
```

6.6.1 Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$ and $x_i \geq 0$ for all $1 \leq i \leq n$.

- 1. In case of minimization, let $c'_i = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$
 - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$

```
• \sum_{1 \le i \le n} A_{ji} x_i \ge b_j
```

4. If x_i has no lower bound, replace x_i with $x_i - x_i'$

6.7 Schreier–Sims Algorithm

```
namespace schreier {
int n;
vector<vector<int>>> bkts, binv;
vector<vector<int>> lk;
vector<int> operator*(const vector<int> &a, const
    vector<int> &b) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i)
    res[i] = b[a[i]];
  return res;
vector<int> inv(const vector<int> &a) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i)
    res[a[i]] = i;
  return res;
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
  vector<int> p = g;
for (int i = 0; i < n; ++i) {</pre>
    assert(p[i] >= 0 && p[i] < (int)lk[i].size());
    int res = lk[i][p[i]];
     if (res == -1) {
      if (add) {
  bkts[i].push_back(p);
         binv[i].push_back(inv(p));
         lk[i][p[i]] = (int)bkts[i].size() - 1;
       return i;
    p = p * binv[i][res];
  }
  return -1;
bool inside(const vector<int> &g) {
  return filter(g, false) == -1;
void solve(const vector<vector<int>> &gen, int _n) {
  bkts.clear(), bkts.resize(n);
binv.clear(), binv.resize(n);
  lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
  for (int i = 0; i < n; ++i) {
    lk[i].resize(n, -1);
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i)</pre>
    filter(gen[i]);
  queue<pair<pair<int, int>, pair<int, int>>> upd;
for (int i = 0; i < n; ++i) {</pre>
    for (int j = i; j < n; ++j) {
  for (int k = 0; k < (int)bkts[i].size(); ++k) {
    for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l))
     ;
      }
    }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
int res = filter(bkts[a.first][a.second] * bkts[b.
     first][b.second]);
     if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].
     size() - 1);
    for (int i = 0; i < n; ++i) {
       for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
```

6.8 Miller Rabin

```
9780504, 1795265022
vector < long long > chk = \{ 2, 325, 9375, 28178, 450775, 
    9780504, 1795265022 };
bool check(long long a, long long u, long long n, int t
  a = fpow(a, u, n);
  if (a == 0) return true;
  if (a == 1 \mid \mid a == n - 1) return true;
  for (int i = 0; i < t; ++i) {
    a = fmul(a, a, n);
    if (a == 1) return false;
    if (a == n - 1) return true;
  return false;
bool is_prime(long long n) {
  if (n < 2) return false;
  if (n % 2 == 0) return n == 2;
  long long u = n - 1; int t = 0;
 for (; !(u & 1); u >>= 1, ++t);
for (long long i : chk) {
   if (!check(i, u, n, t)) return false;
  return true;
}
```

6.9 Pollard's Rho

```
map<long long, int> cnt;
long long f(long long x, long long n, int p) { return (
fmul(x, x, n) + p) % n; }
void pollard_rho(long long n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
if (n % 2 == 0) return pollard_rho(n / 2), ++cnt[2],
     void();
  long long x = 2, y = 2, d = 1, p = 1;
  while (true) {
     if (d != n && d != 1) {
       pollard_rho(n / d);
       pollard_rho(d);
       return;
    if (d == n) ++p;
    x = f(x, n, p); y = f(f(y, n, p), n, p);
     d = \underline{gcd(abs(x - y), n)};
  }
}
```

6.10 Meissel-Lehmer Algorithm

```
int prc[maxn];
long long phic[msz][nsz];
void sieve() {
  bitset<maxn> v
   pr.push_back(0);
   for (int i = 2; i < maxn; ++i) {
     if (!v[i]) pr.push_back(i);
     for (int j = 1; i * pr[j] < maxn; ++j) {
       v[i * pr[j]] = true;
       if (i % pr[j] == 0) break;
  for (int i = 1; i < pr.size(); ++i) prc[pr[i]] = 1;
for (int i = 1; i < maxn; ++i) prc[i] += prc[i - 1];</pre>
long long p2(long long, long long);
long long phi(long long m, long long n) {
   if (m < msz && n < nsz && phic[m][n] != -1) return</pre>
     phic[m][n];
   if (n == 0) return m;
   if (pr[n] >= m) return 1;
   long long ret = phi(m, n-1) - phi(m / pr[n], n-1)
   if (m < msz && n < nsz) phic[m][n] = ret;</pre>
   return ret;
long long pi(long long m) {
   if (m < maxn) return prc[m];</pre>
   long long n = pi(cbrt(m));
   return phi(m, n) + n - 1 - p2(m, n);
long long p2(long long m, long long n) {
   long long ret = 0;
   long long lim = sqrt(m);
   for (int i = n + 1; pr[i] <= lim; ++i) ret += pi(m /</pre>
     pr[i]) - pi(pr[i]) + 1;
   return ret;
}
```

6.11 Discrete Logarithm

```
// to solve discrete x for x^a = b \pmod{p} with p is
    prime
// let c = primitive root of p
// find k such that c^k = b \pmod{p} by bsgs
// solve fa = k \pmod{p-1} by euclidean algorithm
// x = c^f
int bsgs(int a, int b, int p) {
  // return L such that a^L = b \pmod{p}
  if (p == 1) {
    if (!b) return a != 1;
    return -1;
  if (b == 1) {
    if (a) return 0;
    return -1;
  if (a \% p == 0) {
    if (!b) return 1;
    return -1;
  int num = 0, d = 1;
  while (true) {
    int r = __gcd(a, p);
if (r == 1) break;
if (b % r) return -1;
    ++num;
    b /= r, p /= r;
d = (111 * d * a / r) % p;
  for (int i = 0, now = 1; i < num; ++i, now = 1ll *
    now * a % p) {
    if (now == b) return i;
  int m = ceil(sqrt(p)), base = 1;
  map<int, int> mp;
  for (int i = 0; i < m; ++i) {
    if (mp.find(base) == mp.end()) mp[base] = i;
    else mp[base] = min(mp[base], i);
```

```
base = 111 * base * a % p;
}
for (int i = 0; i < m; ++i) {
    // can be modified to fpow if p is prime
    int r, x, y; tie(r, x, y) = extgcd(d, p);
    x = (111 * x * b % p + p) % p;
    if (mp.find(x) != mp.end()) return i * m + mp[x] +
    num;
    d = 111 * d * base % p;
}
return -1;
}</pre>
```

6.12 Gaussian Elimination

```
void gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   for (int i = 0; i < m; ++i) {
      int p = -1;
      for (int j = i; j < n; ++j) {
        if (fabs(d[j][i]) < eps) continue;
        if (p == -1 || fabs(d[j][i]) > fabs(d[p][i])) p =
        j;
      }
      if (p == -1) continue;
      for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
      for (int j = 0; j < n; ++j) {
        if (i == j) continue;
        double z = d[j][i] / d[i][i];
        for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
      }
    }
}</pre>
```

6.13 μ function

6.14 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

6.15 De Bruijn Sequence

```
int res[maxn], aux[maxn], a[maxn], sz;

void db(int t, int p, int n, int k) {
   if (t > n) {
      if (n % p == 0) {
       for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
      }
   } else {
      aux[t] = aux[t - p];
}</pre>
```

```
db(t + 1, p, n, k);
  for (int i = aux[t - p] + 1; i < k; ++i) {
    aux[t] = i;
    db(t + 1, t, n, k);
  }
}
int de_bruijn(int k, int n) {
  // return cyclic string of length k^n such that every
    string of length n using k character appears as a
    substring.
  if (k == 1) {
    res[0] = 0;
    return 1;
}
for (int i = 0; i < k * n; i++) aux[i] = 0;
    sz = 0;
    db(1, 1, n, k);
    return sz;
}</pre>
```

6.16 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
  if (!b) return make_tuple(a, 1, 0);
  T d, x, y;
  tie(d, x, y) = extgcd(b, a % b);
  return make_tuple(d, y, x - (a / b) * y);
}
```

6.17 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ + g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.18 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
  long long mult = mod[0];
  int n = (int)mod.size();
  long long res = a[0];
  for (int i = 1; i < n; ++i) {
    long long d, x, y;
    tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
    if ((a[i] - res) % d) return -1;
    long long new_mult = mult / __gcd(mult, 1ll * mod[i]) * mod[i];</pre>
```

```
res += x * ((a[i] - res) / d) % new_mult * mult %
    new_mult;
    mult = new_mult;
    ((res %= mult) += mult) %= mult;
}
return res;
}
```

6.19 Theorem

6.19.1 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

6.19.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

6.19.3 Cayley's Formula

- Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

6.19.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \leq k \leq n$.

6.20 Primes

 $\begin{array}{l} 97,101,131,487,593,877,1087,1187,1487,1787,3187,12721,\\ 13331,14341,75577,123457,222557,556679,999983,\\ 1097774749,1076767633,100102021,999997771,\\ 1001010013,1000512343,987654361,999991231,\\ 999988733,98789101,987777733,999991921,1000000007,\\ 1000000087,1000000123,1010101333,1010102101,\\ 100000000039,100000000000037,2305843009213693951,\\ 4611686018427387847,9223372036854775783,18446744073709551557,\\ \end{array}$

7 Dynamic Programming

7.1 Convex Hull Optimization

```
struct line {
  int m, y;
  int l, r;
  line(int m = 0,int y = 0, int l = -5, int r =
     1000000009): m(m), y(y), l(l), r(r) {}
  int get(int x) const { return m * x + y; }
  int useful(line le) const {
    return (int)(get(l) >= le.get(l)) + (int)(get(r) >=
     le.get(r));
  }
};
int magic;
bool operator < (const line &a, const line &b) {
  if (magic) return a.m < b.m;
  return a.l < b.l;
}</pre>
```

```
set<line> st:
void addline(line 1) {
  magic = 1;
  auto it = st.lower_bound(1);
  if (it != st.end() && it->useful(l) == 2) return;
  while (it != st.end() && it->useful(l) == 0) it = st.
    erase(it);
  if (it != st.end() && it->useful(l) == 1) {
    int L = it->l, R = it->r, M;
while (R > L) {
      M = (L + R + 1) >> 1;
       if (it->get(M) >= l.get(M)) R = M - 1;
      else L = M;
    line cp = *it;
    st.erase(it);
    cp.l = L + 1;
    if (cp.l <= cp.r) st.insert(cp);</pre>
    l.r = L;
  else if (it != st.end()) l.r = it->l - 1;
  it = st.lower_bound(1);
while (it != st.begin() && prev(it)->useful(1) == 0)
    it = st.erase(prev(it));
  if (it != st.begin() && prev(it)->useful(l) == 1) {
    int \hat{L} = it \rightarrow l, R = it \rightarrow r, M;
    while (R > L) {
      M = (L + R) >> 1;
       if (it->get(M) >= l.get(M)) L = M + 1;
       else R = M;
    line cp = *it;
    st.erase(it);
    cp.r = L - 1;
    if (cp.l <= cp.r) st.insert(cp);</pre>
  else if (it != st.begin()) l.l = prev(it)->r + 1;
  if (l.l <= l.r) st.insert(l);</pre>
int getval(int d) {
  magic = 0;
  return (--st.upper_bound(line(0, 0, d, 0)))->get(d);
```

7.2 1D/1D Convex Optimization

```
struct segment {
  int i, l, r;
  segment() {}
  segment(int a, int b, int c): i(a), l(b), r(c) {}
};
inline long long f(int l, int r) {
  return dp[l] + w(l + 1, r);
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i \le n; ++i) {
    dp[i] = f(deq.front().i, i);
    while (deq.size() \&\& deq.front().r < i + 1) deq.
    pop_front();
    deq.front().l = i + 1;
    segment seg = segment(i, i + 1, n);
    while (deq.size() && f(i, deq.back().l) < f(deq.back().i, deq.back().l)) deq.pop_back();</pre>
    if (deq.size()) {
  int d = 1048576, c = deq.back().l;
      while (d \gg 1) if (c + d \ll deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c +=
    d;
      deq.back().r = c; seg.l = c + 1;
    if (seg.l <= n) deq.push_back(seg);</pre>
```

Condition

}

7.3

7.3.1 totally monotone (concave/convex)

```
\begin{array}{ll} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

7.3.2 monge condition (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', \, B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \, B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8 Geometry

8.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps</pre>
struct P {
  double x, y;
 double x, y;
P() : x(0), y(0) {}
P(double x, double y) : x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }

double connector * (P b) { return P(x / b, y / b); }
  double operator * (P b) { return x * b.x + y * b.y; double operator ^ (P b) { return x * b.y - y * b.x;
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P spin(double o) {
     double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
   double angle() { return atan2(y, x); }
struct L {
  // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
L(): a(0), b(0), c(0), o(0), pa(), pb() {}
  L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa
  ^ pb), o(atan2(-a, b)), pa(pa), pb(pb) {}
P project(P p) { return pa + (pb - pa).unit() * ((pb
      - pa) * (p - pa) / (pb - pa).abs()); }
  };
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
  if (\max(p1.x, p2.x) < \min(p3.x, p4.x) \mid | \max(p3.x, p4.x) |
       (x) < min(p1.x, p2.x)) return false;
   if (max(p1.y, p2.y) < min(p3.y, p4.y) | | max(p3.y, p4
  .y) < min(p1.y, p2.y)) return false;
return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^
       (p4 - p2)) \le 0 \&\&
        sign((p1 - p3) ^ (p2 - p3)) * sign((p1 - p4) ^ (
     p2 - p4)) <= 0;
bool parallel(L x, L y) { return same(x.a * y.b, x.b *
     y.a); }
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b , x.a * y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a)
```

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[
     maxn], yr[maxn];
point p[maxn];
int build(int 1, int r, int dep = 0) {
   if (l == r) return -1;
   function<bool(const point &, const point &)> f = [dep
     ](const point &a, const point &b) {
     if (dep \& 1) return a.x < b.x;
     else return a.y < b.y;</pre>
   int m = (l + r) >> 1;
   nth_element(p + l, p + m, p + r, f);
   xl[m] = xr[m] = p[m].x;
   yl[m] = yr[m] = p[m].y;
   [c[m] = build(l, m, dep + 1);
   if (~lc[m]) {
     xl[m] = min(xl[m], xl[lc[m]]);
     xr[m] = max(xr[m], xr[lc[m]]);
     yl[m] = min(yl[m], yl[lc[m]]);
    yr[m] = max(yr[m], yr[lc[m]]);
   rc[m] = build(m + 1, r, dep + 1);
   if (~rc[m]) {
     xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
yl[m] = min(yl[m], yl[rc[m]));
yr[m] = max(yr[m], yr[rc[m]));
   return m;
bool bound(const point &q, int o, long long d) {
   double ds = sqrt(d + 1.0);
   if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
     q.y < yl[o] - ds | | q.y > yr[o] + ds) return false;
   return true;
void dfs(const point &q, long long &d, int o, int dep =
      0) {
   if (!bound(q, o, d)) return;
   long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
   if ((dep & 1) && q.x < p[o].x || !(dep & 1) && q.y <
     p[o].y) {
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
   } else {
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
}
void init(const vector<point> &v) {
   for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
   root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
   dfs(q, res, root);
   return res:
| }}
8.3
      Delaunay Triangulation
namespace triangulation {
```

```
namespace triangulation {
static const int maxn = 1e5 + 5;
vector<point> p;
set<int> g[maxn];
int o[maxn];
set<int> s;
void add_edge(int x, int y) {
   s.insert(x), s.insert(y);
   g[x].insert(y);
   g[y].insert(x);
}
bool inside(point a, point b, point c, point p) {
   if (((b - a) ^ (c - a)) < 0) swap(b, c);</pre>
```

```
function<long long(int)> sqr = [](int x) { return x *
     111 * x;
  long long k11 = a.x - p.x, k12 = a.y - p.y, k13 = sqr
    (a.x) - sqr(p.x) + sqr(a.y) - sqr(p.y);
  long long k21 = b.x - p.x, k22 = b.y - p.y, k23 = sqr
    (b.x) - sqr(p.x) + sqr(b.y) - sqr(p.y);
  long long k31 = c.x - p.x, k32 = c.y - p.y, k33 = sqr
  (c.x) - sqr(p.x) + sqr(c.y) - sqr(p.y);
long long det = k11 * (k22 * k33 - k23 * k32) - k12 *
     (k21 * k33 - k23 * k31) + k13 * (k21 * k32 - k22 *
     k31);
  return det > 0;
bool intersect(const point &a, const point &b, const
  point &c, const point &d) {
return ((b - a) ^ (c - a)) * ((b - a) ^ (d - a)) < 0
      ((d - c) \wedge (a - c)) * ((d - c) \wedge (b - c)) < 0;
void dfs(int 1, int r) {
  if (r - l <= 3) {</pre>
    for (int i = 1; i < r; ++i) {
      for (int j = i + 1; j < r; ++j) add_edge(i, j);
    return:
  int m = (l + r) >> 1;
 dfs(l, m), dfs(m, r);
  int pl = l, pr = r - 1;
 while (true) {
    int z = -1;
    for (int u : g[pl]) {
      long long c = ((p[pl] - p[pr]) ^ (p[u] - p[pr]));
if (c > 0 || c == 0 && abs(p[u] - p[pr]) < abs(p[</pre>
    pl] - p[pr])) {
         z = u;
        break;
      }
    if (z != -1) {
      pl = z;
      continue;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) \wedge (p[u] - p[pl]));
       if (c < 0 \mid | c == 0 \& abs(p[u] - p[pl]) < abs(p[
    pr] - p[pl])) {
         z = u;
        break;
      }
    if (z != -1) {
      pr = z;
      continue;
    break;
  add_edge(pl, pr);
  while (true) {
    int z = -1;
    bool b = false;
    for (int u : g[pl]) {
      long long c = (p[pl] - p[pr]) \wedge (p[u] - p[pr]);
      if (c < 0 \&\& (z == -1 | l inside(p[pl], p[pr], p[z])
    ], p[u]))) z = u;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) ^ (p[u] - p[pl]));
if (c > 0 && (z == -1 || inside(p[pl], p[pr], p[z
    ], p[u]))) z = u, b = true;
    if (z == -1) break;
    int x = pl, y = pr;
    if (b) swap(x, y);
    for (auto it = g[x].begin(); it != g[x].end(); ) {
      int u = *it;
      if (intersect(p[x], p[u], p[y], p[z])) {
         it = g[x].erase(it);
         g[u].erase(x);
      } else {
         ++it;
```

```
if (b) add_edge(pl, z), pr = z;
    else add_edge(pr, z), pl = z;
vector<vector<int>> solve(vector<point> v) {
  int n = v.size();
  for (int i = 0; i < n; ++i) g[i].clear();
  for (int i = 0; i < n; ++i) o[i] = i;
  sort(o, o + n, [\&](int i, int j) \{ return v[i] < v[j] \}
    ]; });
  p.resize(n);
  for (int i = 0; i < n; ++i) p[i] = v[o[i]];</pre>
  dfs(0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) {
    for (int j : g[i]) res[o[i]].push_back(o[j]);
  return res;
}}
```

8.4 Sector Area

```
// calc area of sector which include a, b
double SectorArea(P a, P b, double r) {
  double theta = atan2(a.y, a.x) - atan2(b.y, b.x);
  while (theta <= 0) theta += 2 * pi;
  while (theta >= 2 * pi) theta -= 2 * pi;
  theta = min(theta, 2 * pi - theta);
  return r * r * theta / 2;
}
```

8.5 Half Plane Intersection

```
bool jizz(L l1,L l2,L l3){
  P p=intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
bool cmp(const L &a,const L &b){
  return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):
    a.o<b.o;
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
vector<L> pls(1,ls[0]);
  for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls</pre>
     .back().o))pls.push_back(ls[i]);
  deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],
    pls[b],pls[c]))
  for(int i=2;i<(int)pls.size();++i){</pre>
    meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
    meow(i,dq[0],dq[1])dq.pop_front();
    dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back
    ();
  meow(dq.back(),dq[0],dq[1])dq.pop_front();
  if(dq.size()<3u)return vector<P>(); // no solution or
     solution is not a convex
  vector<P> rt
  for(int i=0;i<(int)dq.size();++i)rt.push_back(</pre>
    intersect(pls[dq[i]],pls[dq[(i+1)%dq.size()]]));
  return rt;
```

8.6 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
  int n=int(ps.size());
  vector<int> id(n),pos(n);
```

```
vector<pair<int,int>> line(n*(n-1)/2);
  int m=-1:
  for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=</pre>
  make_pair(i,j); ++m;
sort(line.begin(),line.end(),[&](const pair<int,int>
    &a,const pair<int,int> &b)->bool{
    if(ps[a.first].first==ps[a.second].first)return 0;
    if(ps[b.first].first==ps[b.second].first)return 1;
    return (double)(ps[a.first].second-ps[a.second].
    second)/(ps[a.first].first-ps[a.second].first) <</pre>
     double)(ps[b.first].second-ps[b.second].second)/(ps
    [b.first].first-ps[b.second].first);
  for(int i=0;i<n;++i)id[i]=i;</pre>
  sort(id.begin(),id.end(),[&](const int &a,const int &
    b){ return ps[a]<ps[b]; })</pre>
  for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
  for(int i=0;i<m;++i){</pre>
    auto l=line[i];
    // meow
    tie(pos[l.first],pos[l.second],id[pos[l.first]],id[
    pos[l.second]])=make_tuple(pos[l.second],pos[l.
     first], l. second, l. first);
}
```

8.7 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2
double bx = (c.x + b.x) / 2
  double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)
  )) / (sin(a1) * cos(a2) - sin(a2) * cos(a1));
return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
     TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
   double lb = len(a - c);
  double lc = len(a - b);
  res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb +
      lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb +
      lc);
   return res;
}
```

8.8 Polygon Center

```
Point BaryCenter(vector<Point> &p, int n) {
   Point res(0, 0);
   double s = 0.0, t;
   for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
   }
   res.x /= (3 * s);
   res.y /= (3 * s);
   return res;
}</pre>
```

8.9 Maximum Triangle

```
double ConvexHullMaxTriangleArea(Point p[], int res[],
    int chnum) {
  double area = 0,
  res[chnum] = res[0];
  1) % chnum]] - p[res[i]])) > fabs(Cross(p[res[j]]
    - p[res[i]], p[res[k]] - p[res[i]]))) k = (k + 1) %
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] -
     p[res[i]]));
    if (tmp > area) area = tmp;
    while (fabs(Cross(p[res[(j + 1) % chnum]] - p[res[i
]], p[res[k]] - p[res[i]])) > fabs(Cross(p[res[j]])
    - p[res[i]], p[res[k]] - p[res[i]]))) j = (j + 1) %
     chnum;
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] -
     p[res[i]]));
    if (tmp > area) area = tmp;
  return area / 2;
}
```

8.10 Point in Polygon

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) % ps.size();
    L l(ps[a], ps[b]);
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
    if (same(ps[a].y, ps[b].y) && same(ps[a].y, p.y))
    continue;
    if (ps[a].y > ps[b].y) swap(a, b);
    if (ps[a].y <= p.y && p.y < ps[b].y && p.x <= ps[a].x + (ps[b].x - ps[a].x) / (ps[b].y - ps[a].y) * (
        p.y - ps[a].y)) ++c;
    }
    return (c & 1) * 2;
}</pre>
```

8.11 Circle

```
struct C {
  Pc;
  double r:
  C(P \ c = P(0, 0), double \ r = 0) : c(c), r(r) \{\}
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c))
    .unit() * a.r);
  else if (a.r + b.r > d & d + a.r >= b.r) {
    double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 *
    a.r * d));
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.spin(o) * a.r);
    p.push_back(a.c + i.spin(-o) * a.r);
  return p;
double IntersectArea(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d >= a.r + b.r - eps) return 0;
  if (d + a.r \leftarrow b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.
    r * d));
  double q = acos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.
  r * d));
return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
```

```
// remove second level if to get points for line (
     defalut: segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r)
  vector<P> t;
  if (d >= -eps) {
     double i = (-B - sqrt(abs(d))) / (2 * A);
double j = (-B + sqrt(abs(d))) / (2 * A);
     if (i - 1.0 \le eps \& i \ge -eps) t.emplace_back(a.x
     + i * x, a.y + i * y);
if (j - 1.0 <= eps && j >= -eps) t.emplace_back(a.x
       + j * x, a.y + j * y);
  return t;
// calc area intersect by circle with radius r and
     triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() - r < 0, inb = b.abs() - r < 0;
  if (ina) {
     if (inb) return abs(a ^ b) / 2;
     else {
        P p = CircleCrossLine(a, b, P(0, 0), r)[0];
        return SectorArea(b, p, r) + abs(a \land p) / 2;
  } else {
     auto p = CircleCrossLine(a, b, P(0, 0), r);
     if (inb) return SectorArea(p[0], a, r) + abs(p[0] \land
      b) / 2;
     else {
        if (p.size() == 2u) return SectorArea(a, p[0], r)
       + SectorArea(p[1], b, r) + abs(p[0] \land p[1]) / 2;
        else return SectorArea(a, b, r);
     }
  }
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
  double ans = 0;
  for (int i = 0; i < 3; ++i) {
     int j = (i + 1) \% 3;
     double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y,
      ps[j].x);
     if (o >= pi) o = o - 2 * pi;
     if (o \le -pi) o = o + 2 * pi;
     ans += AreaOfCircleTriangle(ps[i], ps[j], r) * (o
     >= 0 ? 1 : -1);
  return abs(ans);
```

8.12 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
#define Pij \
  P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x);
  z.emplace_back(a.c + i, a.c + i + j);
#define deo(I,J) \
  double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos
    (e / d);\
  P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-a.c)
    0);\
  z.emplace\_back(a.c + j * a.r, b.c J j * b.r);
  z.emplace_back(a.c + k * a.r, b.c J k * b.r);
  if (a.r < b.r) swap(a, b);
  vector<L> z;
  if ((a.c - b.c).abs() + b.r < a.r) return z;
  else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
    deo(-,+);
    if (same(d, a.r + b.r)) { Pij; }
else if (d > a.r + b.r) { deo(+,-); }
 }
  return z;
```

```
vector<L> tangent(C c, P p) {
  vector<L> z;
  double d = (p - c.c).abs();
  if (same(d, c.r)) {
    P i = (p - c.c).spin(pi / 2);
    z.emplace_back(p, p + i);
} else if (d > c.r) {
    double o = acos(c.r / d);
    P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.
    spin(-o) * c.r;
    z.emplace_back(c.c + j, p);
    z.emplace_back(c.c + k, p);
}
return z;
}
```

8.13 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
   double d = (a.c - b.c).abs();
   vector<pair<double, double>> res;
   if (same(a.r + b.r, d));
  else if (d <= abs(a.r - b.r) + eps) {
  if (a.r < b.r) res.emplace_back(0, 2 * pi);
} else if (d < abs(a.r + b.r) - eps) {</pre>
      double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 *
     a.r * d), z = (b.c - a.c).angle();
if (z < 0) z += 2 * pi;
      double l = z - o, r = z + o;
     if (l < 0) l += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
      if (l > r) res.emplace_back(l, 2 * pi), res.
      emplace_back(0, r);
     else res.emplace_back(l, r);
   return res;
}
double CircleUnionArea(vector<C> c) { // circle should
      be identical
   int n = c.size();
   double a = 0, w;
   for (int i = 0; w = 0, i < n; ++i) {
     vectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvectorvector
        z = CoverSegment(c[i], c[j]);
        for (auto &e : z) s.push_back(e);
     sort(s.begin(), s.end());
      auto F = [\&] (double t) { return c[i].r * (c[i].r *
       t + c[i].c.x * sin(t) - c[i].c.y * cos(t)); };
      for (auto &e : s) {
        if (e.first > w) a += F(e.first) - F(w);
        w = max(w, e.second);
   return a * 0.5;
}
```

8.14 Minimun Distance of 2 Polygons

bool operator < (const P &a, const P &b) { return same(</pre>

8.15 2D Convex Hull

```
a.x, b.x) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(
    a.x, b.x) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) \wedge (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return
    same(a.x, b.x) ? a.y < b.y : a.x < b.x; });
  for (int i = 0; i < ps.size(); ++i) {</pre>
    while (p.size() >= 2 && crx(p[p.size() - 2], ps[i],
p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() - 2], ps[i],
p[p.size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double x) { return same(x, 0) ? 0 : x > 0 ? 1 :
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector<P> p, u, d;
  CH() \{\}
  CH(vector<P> ps) : p(ps) {
    n = ps.size();
    rotate(p.begin(), min_element(p.begin(), p.end()),
    p.end());
    auto t = max_element(p.begin(), p.end());
    d = vector<P>(p.begin(), next(t));
    u = vector<P>(t, p.end()); u.push_back(p[0]);
  int find(vector<P> &v, P d) {
    int l = 0, r = v.size();
    while (l + 5 < r) {
  int L = (l * 2 + r) / 3, R = (l + r * 2) / 3;
  if (v[L] * d > v[R] * d) r = R;
      else l = L:
    int x = 1;
    for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x]
     * d) x = i;
    return x;
  int findFarest(P v) {
    if (v.y > 0 | | v.y == 0 && v.x > 0) return ((int)d.
    size() - 1 + find(u, v)) % p.size();
    return find(d, v);
  P get(int 1, int r, P a, P b) {
    int s = sgn(crx(a, b, p[l % n]));
    while (l + 1 < r) {
       int m = (l + r) >> 1;
      if (sgn(crx(a, b, p[m % n])) == s) l = m;
    return isLL(a, b, p[l % n], p[(l + 1) % n]);
```

```
vector<P> getIS(P a, P b) {
     int X = findFarest((b - a).spin(pi / 2));
     int Y = findFarest((a - b).spin(pi / 2));
     if (X > Y) swap(X, Y)
     if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) <
     0) return {get(X, Y, a, b), get(Y, X + n, a, b)};
    return {};
  void update_tangent(P q, int i, int &a, int &b) {
    if (sgn(crx(q, p[a], p[i])) > 0) a = i;
if (sgn(crx(q, p[b], p[i])) < 0) b = i;</pre>
  void bs(int l, int r, P q, int &a, int &b) {
    if (l == r) return
     update_tangent(q, 1 % n, a, b);
    int s = sgn(crx(q, p[l \% n], p[(l + 1) \% n]));
while (l + 1 < r) {
       int m = (l + r) >> 1;
      if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l
      = m;
      else r = m;
     update_tangent(q, r % n, a, b);
  bool contain(P p) {
     if (p.x < d[0].x | | p.x > d.back().x) return 0;
     auto it = lower_bound(d.begin(), d.end(), P(p.x, -1
     e12));
     if (it->x == p.x) {
       if (it->y > p.y) return 0;
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     it = lower_bound(u.begin(), u.end(), P(p.x, 1e12),
     greater<P>());
     if (it->x == p.x) {
       if (it->y < p.y) return 0;</pre>
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     return 1:
  bool get_tangent(P p, int &a, int &b) { // b -> a
    if (contain(p)) return 0;
     a = b = 0;
     int i = lower_bound(d.begin(), d.end(), p) - d.
     begin();
    bs(0, i, p, a, b);
bs(i, d.size(), p, a, b);
     i = lower_bound(u.begin(), u.end(), p, greater<P>()
     ) - u.begin();
     bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a,
     bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.
     size(), p, a, b);
     return 1;
};
```

8.16 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d){
  return abs(((b-a)^{(c-a)})^*(d-a))/6;
}
struct convex3D{
static const int maxn=1010;
struct T{
  int a,b,c;
  bool res;
  T(){}
  T(int a, int b, int c, bool res=1):a(a),b(b),c(c),res(
    res){}
int n,m;
P p[maxn]
T f[maxn*8];
int id[maxn][maxn];
bool on(T \&t, P \&q){
  return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>
}
```

```
void meow(int q,int a,int b){
  int f2=id[a][b];
  if(f[f2].res){
    if(on(f[f2],p[q]))dfs(q,f2);
      id[q][b]=id[a][q]=id[b][a]=m;
      f[m++]=T(b,a,q,1);
  }
void dfs(int p,int i){
  f[i].res=0;
  meow(p,f[i].b,f[now].a);
  meow(p,f[i].c,f[now].b);
  meow(p,f[i].a,f[now].c);
void operator()(){
  if(n<4)return
  if([&]()->int{
    for(int i=1;i<n;++i)if(abs(p[0]-p[i])>eps)return
    swap(p[1],p[i]),0;
    return 1;
  }())return;
  if([&]()->int{
    for(int i=2;i<n;++i)if(abs((p[0]-p[i])^(p[1]-p[i]))</pre>
    >eps)return swap(p[2],p[i]),0;
    return 1;
  }())return;
  if([&]()->int{
    for(int i=3; i< n; ++i) if (abs(((p[1]-p[0])^{p[2]-p[0]})
    )*(p[i]-p[0]))>eps)return swap(p[3],p[i]),0;
  }())return;
  for(int i=0; i<4; ++i){
    T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
    if(on(t,p[i]))swap(t.b,t.c);
    id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
    f[m++]=t;
  for(int i=4;i<n;++i)for(int j=0;j<m;++j)if(f[j].res</pre>
    && on(f[j],p[i])){
    dfs(i,j);
    break;
  int mm=m; m=0;
  for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
bool same(int i,int j){
  return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].
    a])>eps || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f
     [j].b])>eps || absvol(p[f[i].a],p[f[i].b],p[f[i].c
    ],p[f[j].c])>eps);
int faces(){
  int r=0;
for(int i=0;i<m;++i){</pre>
    int iden=1;
    for(int j=0;j<i;++j)if(same(i,j))iden=0;</pre>
    r+=iden;
  return r;
} tb;
```

8.17 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^ p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
}

circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {</pre>
```

```
if (norm2(cent - p[i]) <= r) continue;
cent = p[i];
r = 0.0;
for (int j = 0; j < i; ++j) {
    if (norm2(cent - p[j]) <= r) continue;
    cent = (p[i] + p[j]) / 2;
    r = norm2(p[j] - cent);
    for (int k = 0; k < j; ++k) {
        if (norm2(cent - p[k]) <= r) continue;
        cent = center(p[i], p[j], p[k]);
        r = norm2(p[k] - cent);
    }
}
return circle(cent, sqrt(r));
</pre>
```

8.18 Closest Pair

```
double closest_pair(int 1, int r) {
  // p should be sorted increasingly according to the x
     -coordinates.
  if (l == r) return 1e9;
  if (r - l == 1) return dist(p[l], p[r]);
  int m = (l + r) >> 1;
  double d = min(closest_pair(l, m), closest_pair(m +
    1, r));
  vector<int> vec;
  for (int i = m; i >= 1 && fabs(p[m].x - p[i].x) < d;
     --i) vec.push_back(i);
  for (int i = m + 1; i \le r \& fabs(p[m].x - p[i].x) <
     d; ++i) vec.push_back(i);
  sort(vec.begin(), vec.end(), [&](int a, int b) {
    return p[a].y < p[b].y; });</pre>
  for (int i = 0; i < vec.size(); ++i) {
  for (int j = i + 1; j < vec.size() && fabs(p[vec[j ]].y - p[vec[i]].y) < d; ++j) {
      d = min(d, dist(p[vec[i]], p[vec[j]]));
  }
  return d;
```

9 Miscellaneous / Problems

9.1 Bitwise Hack

9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) {
            x = s - 1 - x;
            y = s - 1 - y;
        }
    }
```

```
swap(x, y);
}

BigInteger b16 = new BigInteger(stz, 16);
System.out.println(b16.toString(2));
}
return res;
}
```

9.3 Java

```
import java.io.*;
import java.util.*;
import java.lang.*
import java.math.*;
public class filename{
  static Scanner in = new Scanner(System.in);
  public static void main(String[] args) {
    int t = 7122;
    while(in.hasNext()) {
      t = in.nextInt();
      float b = in.nextFloat();
      String str = in.nextLine(); // getline
      String stu = in.next(); // string
    System.out.println("Case #" + t);
System.out.printf("%d\n", 7122);
    int[] c = new int[5];
    int[][] d = {{7,1,2,2},{8,7}};
int g = Integer.parseInt("-123");
    long f = (long)d[0][2];
    List<Integer> l = new ArrayList<>();
    Random rg = new Random();
    for (int i = 9; i >= 0; --i) {
      1.add(Integer.valueOf(rg.nextInt(100) + 1));
      1.add(Integer.valueOf((int)(Math.random() * 100)
    + 1));
    Collections.sort(l, new Comparator<Integer>() {
      public int compare(Integer a, Integer b) {
        return a - b;
    });
    for (int i = 0; i < l.size(); ++i) {</pre>
      System.out.print(l.get(i));
    Set<String> s = new HashSet<String>(); // TreeSet
s.add("jizz");
    System.out.println(s);
    System.out.println(s.contains("jizz"));
    Map<String, Integer> m = new HashMap<String,</pre>
    Integer>();
m.put("lol", 7122);
    System.out.println(m);
    for(String key: m.keySet()) {
   System.out.println(key + "
                                    : " + m.get(key));
    System.out.println(m.containsKey("lol"));
    System.out.println(m.containsValue(7122));
    System.out.println(Math.PI);
    System.out.println(Math.acos(-1));
    BigInteger bi = in.nextBigInteger(), bj = new
    BigInteger("-7122"), bk = BigInteger.value0f(17171)
    int sgn = bi.signum(); // -1 if bi is negative, 0
    if bi is zero and 1 if bi is positive.
    bi = bi.add(bj);
    bi = bi.subtract(BigInteger.ONE);
    bi = bi.multiply(bj);
    bi = bi.divide(bj);
    bi = bi.and(bj);
    bi = bi.gcd(bj);
    bi = bi.max(bj);
    bi = bi.pow(10);
    int meow = bi.compareTo(bj); // -1 0 1
    String stz = "f5abd69150";
```

9.4 Dancing Links

```
namespace dlx {
int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn],
  rw[maxn], bt[maxn], s[maxn], head, sz, ans;
void init(int c) {
  for (int i = 0; i < c; ++i) {
  up[i] = dn[i] = bt[i] = i;</pre>
    lt[i] = i == 0 ? c : i - 1;
    rg[i] = i == c - 1 ? c : i' + 1;
    s[i] = 0;
  rg[c] = 0, lt[c] = c - 1;
  up[c] = dn[c] = -1;
  head = c, sz = c + 1;
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
  int f = sz;
  for (int i = 0; i < (int)col.size(); ++i) {</pre>
    int c = col[i], v = sz++;
    dn[bt[c]] = v;
    up[v] = bt[c], bt[c] = v;
rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r, cl[v] = c;
    ++s[c];
if (i > 0) lt[v] = v - 1;
  lt[f] = sz - 1;
void remove(int c) {
  lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
  for (int i = dn[c]; i != c; i = dn[i]) {
  for (int j = rg[i]; j != i; j = rg[j]) {
       up[dn[j]] = up[j], dn[up[j]] = dn[j];
       --s[cl[j]];
  }
}
void restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
    for (int j = lt[i]; j != i; j = lt[j]) {
    ++s[cl[j]];
       up[dn[j]] = j, dn[up[j]] = j;
  lt[rg[c]] = c, rg[lt[c]] = c;
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i)
    dn[bt[i]] = i, up[i] = bt[i];
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
  int w = c;
  for (int x = c; x != head; x = rg[x]) {
    if (s[x] < s[w]) w = x;
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j
     1);
    dfs(dep + 1);
    for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j
    ]);
  restore(w);
}
int solve() {
 ans = 1e9, dfs(0);
 return ans;
```

9.5 Offline Dynamic MST

| } }

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
       weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
      such that cnt[i] == 0
void contract(int l, int r, vector<int> v, vector<int>
    &x, vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
  if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  djs.save();
  for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first</pre>
    ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {
  if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {</pre>
      x.push_back(v[i])
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
  x[i]], ed[x[i]]);
for (int i = 0; i < (int)v.size(); ++i) {
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i])
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
}
void solve(int l, int r, vector<int> v, long long c) {
  if (l == r) {
    cost[qr[l].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first]) {
  printf("%lld\n", c);
      return:
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(</pre>
    minv, cost[v[i]]);
printf("%lld\n", c + minv);
    return;
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
  contract(l, m, lv, x, y);
  long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
    );
  }
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
```

```
rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
}
solve(m + 1, r, y, rc);
djs.undo();
for (int i = l; i <= m; ++i) cnt[qr[i].first]++;
}</pre>
```

9.6 Manhattan Distance MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {
     v[i] = i;
     ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
sort(v.begin(), v.end(), [&](int i, int j) { return x
  [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
for (int i = 0; i < n; ++i) {
     int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
     y[v[i]]) - ds.begin() + 1;
pair<int, int> q = query(p);
     // query return prefix minimum
     if (~q.second) add_edge(v[i], q.second);
     add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
}
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);
  solve(n):
  for (int i = 0; i < n; ++i) x[i] = -x[i];
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);
  solve(n);
```

9.7 IOI 2016 Alien trick

```
struct result {
  long long m; int v
  result(): m(0), v(0) {}
  result(long long a, int b): m(a), v(b) {}
  result operator+(const result &r) const { return
    result(m + r.m, v + r.v);
  bool operator<(const result &r) const { return m == r</pre>
    .m ? v < r.v : m < r.m; }
  bool operator>(const result &r) const { return m == r
    .m ? v > r.v : m > r.m; }
} dp[maxn];
result check(int p);
long long alien() {
  long long c = inf;
  for (int d = 60; d >= 0; --d) {
    if (c - (111 \ll d) < 0) continue;
    result r = check(c - (1ll << d));
    if (r.v == k) return r.m - (c - (111 << d)) * k;
    if (r.v < k) c -= (1ll << d);
  result r = check(c);
  return r.m - c * k;
```