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1 Basic

1.1 vimrc

```
se nu rnu bs=2 ru mouse=a cin et ts=4 sw=4 sts=4 syn on colo desert filetype indent on inoremap {<CR> {<CR>}<Esc>0
```

1.2 Fast Integer Input

```
| inline int gtx() {
   const int N = 4096;
   static char buffer[N];
   static char *p = buffer, *end = buffer;
   if (p == end) {
     if ((end = buffer + fread(buffer, 1, N, stdin)) == buffer)
     return EOF;
    p = buffer;
  }
   return *p++;
 template <typename T>
inline bool rit(T& x) {
  char c = 0; bool flag = false;
while (c = getchar(), (c < '0' && c != '-') || c > '9') if (c
  if (flag) x = -x;
   return true;
13
```

1.3 Increase stack size

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size) + size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

1.4 Pragma optimization

2 Flow, Matching

2.1 Dinic's Algorithm

```
struct dinic {
  static const int inf = 1e9;
  struct edge {
    int to, cap, rev;
    edge(int d, int c, int r): to(d), cap(c), rev(r) {}
  vector<edge> g[maxn];
  int qu[maxn], ql, qr;
  int lev[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i)
       g[i].clear();
  void add_edge(int a, int b, int c) {
    g[a].emplace_back(b, c, g[b].size() - 0);
    g[b].emplace_back(a, 0, g[a].size() - 1);
  bool bfs(int s, int t) {
    memset(lev, -1, sizeof(lev));
     lev[s] = 0;
    ql = qr = 0;
    qu[qr++] = s;
    while (ql < qr) {
       int x = qu[ql++];
       for (edge &e : g[x]) if (lev[e.to] == -1 && e.cap > 0) {
        lev[e.to] = lev[x] + 1;
         qu[qr++] = e.to;
    }
    return lev[t] != -1;
  int dfs(int x, int t, int flow) {
    if (x == t) return flow;
     int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.to] == lev[x] +
       int f = dfs(e.to, t, min(e.cap, flow - res));
      res += f;
e.cap -= f;
      g[e.to][e.rev].cap += f;
     if (res == 0) lev[x] = -1;
    return res;
  int operator()(int s, int t) {
    int flow = 0;
     for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow;
|};
```

2.2 Minimum-cost flow

```
struct mincost {
  struct edge {
    int dest, cap, w, rev;
    edge(int a, int b, int c, int d): dest(a), cap(b), w(c),
     rev(d) {}
  vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
  bool inq[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
  void add_edge(int a, int b, int c, int d) {
    g[a].emplace_back(b, c, +d, g[b].size() - 0);
g[b].emplace_back(a, 0, -d, g[a].size() - 1);
  bool spfa(int s, int t, int &f, int &c) {
    for (int i = 0; i < maxn; ++i) {
      d[i] = inf;
      p[i] = ed[i] = -1;
       inq[i] = false;
    d\lceil s \rceil = 0;
    queue<int> q;
    q.push(s);
```

```
while (q.size()) {
      int_x_= q.front(); q.pop();
      inq[x] = false;
      for (int i = 0; i < g[x].size(); ++i) {
        edge &e = g[x][i];
        if (e.cap > 0 \& d[e.dest] > d[x] + e.w) {
          d[e.dest] = d[x] + e.w;
          p[e.dest] = x;
           ed[e.dest] = i;
          if (!inq[e.dest]) q.push(e.dest), inq[e.dest] = true;
        }
      }
    if (d[t] == inf) return false;
    int dlt = inf;
    for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[p[x]][ed
     [x]].cap);
    for (int x = t; x != s; x = p[x]) {
      edge &e = g[p[x]][ed[x]];
e.cap -= dlt;
      g[e.dest][e.rev].cap += dlt;
    f += dlt; c += d[t] * dlt;
return true;
  pair<int, int> operator()(int s, int t) {
    int f = 0, c = 0;
    while (spfa(s, t, f, c));
    return make_pair(f, c);
2.3 Gomory-Hu Tree
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;</pre>
  for(int i=2;i<=n;++i){</pre>
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
flow.walk(i); // bfs points that connected to i (use edges
     not fully flow)
    for(int j=i+1;j<=n;++j){</pre>
      if(g[j]==t && flow.connect(j))g[j]=i; // check if i can
     reach j
    }
```

2.4 Stoer-Wagner Minimum Cut

return rt;

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
 bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
   w[x][y] += c;
   w[y][x] += c;
pair<int, int> phase(int n) {
  memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
   int s = -1, t = -1;
   while (true) {
      int c = -1;
     for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
        if (c == -1 || g[i] > g[c]) c = i;
     if (c == -1) break;
     v[c] = true;
s = t, t = c;
     for (int i = 0; i < n; ++i) {
        if (del[i] || v[i]) continue;
        g[i] += w[c][i];
   return make_pair(s, t);
int mincut(int n) {
```

```
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {
   int s, t; tie(s, t) = phase(n);
   del[t] = true;
   cut = min(cut, g[t]);
   for (int j = 0; j < n; ++j) {
      w[s][j] += w[t][j];
      w[j][s] += w[j][t];
   }
}
return cut;
}</pre>
```

2.5 Kuhn–Munkres Algorithm

```
int w[maxn][maxn], lx[maxn], ly[maxn];
int match[maxn], slack[maxn];
bool vx[maxn], vy[maxn];
bool dfs(int x) {
 vx[x] = true;
  for (int i = 0; i < n; ++i) {
    if (vy[i]) continue;
    if (lx[x] + ly[i] > w[x][i]) {
      slack[i] = min(slack[i], lx[x] + ly[i] - w[x][i]);
      continue;
    vy[i] = true;
    if (match[i] == -1 || dfs(match[i])) {
      match[i] = x;
      return true;
 }
  return false;
int solve() {
 fill_n(match, n, -1);
 fill_n(lx, n, -inf);
  fill_n(ly, n, 0);
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) lx[i] = max(lx[i], w[i][j]);
 for (int i = 0; i < n; ++i) {
    fill_n(slack, n, inf);
    while (true) {
      fill_n(vx, n, false);
      fill_n(vy, n, false);
      if (dfs(i)) break;
      int dlt = inf;
      for (int j = 0; j < n; ++j) if (!vy[j]) dlt = min(dlt,
    slack[j]);
      for (int j = 0; j < n; ++j) {
        if (vx[j]) lx[j] -= dlt;
if (vy[j]) ly[j] += dlt;
        else slack[j] -= dlt;
     }
   }
 }
  int res = 0;
  for (int i = 0; i < n; ++i) res += w[match[i]][i];</pre>
  return res;
```

2.6 Flow Model

- Maximum/Minimum flow with lower/upper bound from s to t
 - 1. Construct super source S and sink T
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l
 - 3. For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v)
 - To maximize, connect $t \to s$ with capacity ∞ , and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.

- 5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge $(y \to x \text{ if } (x,y) \in M, x \to y \text{ otherwise})$
 - 2. DFS from unmatched vertices in X
 - 3. $x \in X$ is chosen iff x is unvisited
 - 4. $y \in Y$ is chosen iff y is visited
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost, cap)=(0,d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer T
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G,$ connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e)) 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|

3 Data Structure

3.1 < ext/pbds >

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>, rb_tree_tag,
     tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree set s
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22); assert(*s.find_by_order(1)
     == 71);
  assert(s.order_of_key(22) == 0); assert(s.order_of_key(71) ==
      1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71); assert(s.order_of_key(71)
     == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;
  return 0;
```

3.2 Li Chao Tree

```
namespace lichao {
 struct line {
   long long a, b;
   line(): a(0), b(0) {}
   line(long long a, long long b): a(a), b(b) {}
  long long operator()(int x) const { return a * x + b; }
 line st[maxc * 4];
 int sz, lc[maxc * 4], rc[maxc * 4];
int gnode() {
  st[sz] = line(1e9, 1e9);
lc[sz] = -1, rc[sz] = -1;
   return sz++;
 void init() {
  sz = 0;
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
   bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
   if (mcp) swap(st[o], tl);
   if (r - l == 1) return;
  if (lcp != mcp) {
   if (lc[o] == -1) lc[o] = gnode();
}
     add(1, (1 + r) / 2, t1, lc[o]);
     if (rc[o] == -1) rc[o] = gnode();
     add((l + r) / 2, r, tl, rc[o]);
 long long query(int l, int r, int x, int o) {
   if (r - l == 1) return st[o](x);
  if (x < (l + r) / 2) {
  if (lc[o] == -1) return st[o](x);</pre>
     return min(st[o](x), query(l, (l + r) / 2, x, lc[o]));
     if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((l + r) / 2, r, x, rc[o]));
| }}
```

4 Graph

4.1 Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
  int sum, v, rev, id;
  node(int s, int id): id(id), v(s), sum(s), rev(0), fa(nullptr
     ), pfa(nullptr) {
    ch[0] = nullptr;
    ch[1] = nullptr;
  int relation() {
    return this == fa->ch[0] ? 0 : 1;
  void push() {
    if (!rev) return
    swap(ch[0], ch[1]);
if (ch[0]) ch[0]->rev ^= 1;
    if (ch[1]) ch[1]->rev ^= 1;
    rev = 0;
 }
  void pull() {
    sum = v
    if (ch[0]) sum += ch[0]->sum;
    if (ch[1]) sum += ch[1]->sum;
  void rotate() {
    if (fa->fa) fa->fa->push();
    fa->push(), push(), swap(pfa, fa->pfa);
    int d = relation();
    node *t = fa;
    if (t->fa) t->fa->ch[t->relation()] = this;
    fa = t->fa, t->ch[d] = ch[d ^ 1];
    if (ch[d \land 1]) ch[d \land 1] -> fa = t;
    ch[d \land 1] = t, t->fa = this;
    t->pull(), pull();
  void splay() {
```

```
while (fa) {
       if (!fa->fa) {
         rotate();
         continue;
       fa->fa->push(), fa->push();
       if (relation() == fa->relation()) fa->rotate();
       else rotate(), rotate();
   void evert() { access(), splay(), rev ^= 1; }
   void expose() {
     splay(), push();
     if (ch[1]) {
       ch[1]->fa = nullptr, ch[1]->pfa = this;
       ch[1] = nullptr, pull();
   bool splice() {
     splay();
     if (!pfa) return false;
     pfa->expose(), pfa->ch[1] = this, fa = pfa;
     pfa = nullptr, fa->pull();
     return true;
   void access() {
     expose();
     while (splice());
   int query() { return sum; }
};
namespace lct {
node *sp[maxn];
void make(int u, int v) {
  // create node with id u and value \boldsymbol{v}
   sp[u] = new node(v, u);
void link(int u, int v) {
  // u become v's parent
   sp[v]->evert();
   sp[v]->pfa = sp[u];
void cut(int u, int v) {
  // u was v's parent
   sp[u]->evert();
   sp[v]->access(), sp[v]->splay(), sp[v]->push();
   sp[v]->ch[0]->fa = nullptr;
   sp[v] -> ch[0] = nullptr;
   sp[v]->pull();
void modify(int u, int v) {
   sp[u]->splay();
   sp[u] -> v = v;
   sp[u]->pull();
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
   return sp[v]->query();
int find(int u) {
   sp[u]->access();
   sp[u]->splay();
   node *p = sp[u];
   while (true) {
     p->push();
     if (p->ch[0]) p = p->ch[0];
     else break;
  }
   return p->id;
| }}
```

4.2 Heavy-Light Decomposition

```
void dfs(int x, int p) {
  dep[x] = ~p ? dep[p] + 1 : dep[x];
  sz[x] = 1;
  to[x] = -1;
  fa[x] = p;
  for (const int &u : g[x]) {
    if (u == p) continue;
    dfs(u, x);
    sz[x] += sz[u];
    if (to[x] == -1 || sz[to[x]] < sz[u]) to[x] = u;</pre>
```

```
}
}
void hld(int x, int t) {
  static int tk = 0;
  fr[x] = t;
  dfn[x] = tk++;
  if (!~to[x]) return;
  hld(to[x], t);
  for (const int &u : g[x]) {
     if (u == fa[x] || u == to[x]) continue;
    hld(u, u);
  }
}
vector<pair<int, int>> get(int x, int y) {
  int fx = fr[x], fy = fr[y];
  vector<pair<int, int>> res;
  while (fx != fy) {
    if (dep[fx] < dep[fy]) {</pre>
      swap(fx, fy);
       swap(x, y);
    res.emplace_back(dfn[fx], dfn[x] + 1);
    x = fa[fx];
    fx = fr[x];
  res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn[y]) +
     1);
  int lca = (dep[x] < dep[y] ? x : y);
  return res;
| }
```

4.3 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
   sz[now] = 1; mx[now] = 0;
   for (int u : G[now]) if (!v[u]) {
     get_center(u);
     mx[now] = max(mx[now], sz[u]);
     sz[now] += sz[u];
  }
 void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
   v[now] = true;
   for (auto u : G[now]) if (!v[u.first]) {
     get_dis(u, d, len + u.second);
 void dfs(int now, int fa, int d) {
   get_center(now);
   int c = -1;
   for (int i : vtx) {
     if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx.size()</pre>
      / 2) c = i;
     v[i] = false;
  get_dis(c, d, 0);
for (int i : vtx) v[i] = false;
  v[c] = true; vtx.clear();
  dep[c] = d; p[c] = fa;
for (auto u : G[c]) if (u.first != fa && !v[u.first]) {
     dfs(u.first, c, d + 1);
| }
```

4.4 Minimum Mean Cycle

```
|// d[i][j] == 0 if {i,j} !in E
|long long d[1003][1003],dp[1003][1003];
|pair<long long,long long> MMWC(){
| memset(dp,0x3f,sizeof(dp));
| for(int i=1;i<=n;++i)dp[0][i]=0;
| for(int i=1;i<=n;++i){
| for(int j=1;j<=n;++j){
| for(int k=1;k<=n;++k){
| dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
| }
| }
| long long au=1ll<<31,ad=1;</pre>
```

```
for(int i=1;i<=n;++i){
   if(dp[n][i]==0x3f3f3f3f3f3f3f3f3f3f)continue;
   long long u=0,d=1;
   for(int j=n-1;j>=0;--j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
       u=dp[n][i]-dp[j][i];
      d=n-j;
    }
   }
   if(u*ad<au*d)au=u,ad=d;
}
long long g=__gcd(au,ad);
   return make_pair(au/g,ad/g);
}</pre>
```

4.5 Minimum Steiner Tree

```
namespace steiner {
// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
 // z[i] = the weight of the i-th vertex
const int maxn = 64, maxk = 10;
 const int inf = 1e9;
 int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[maxn];</pre>
 void init(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = inf;
     z[i] = 0;
     w[i][i] = 0;
  }
void add_edge(int x, int y, int d) {
   w[x][y] = min(w[x][y], d);
   w[y][x] = min(w[y][x], d);
}
void build(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) {
       w[i][j] += z[i];
       if (i != j) w[i][j] += z[j];
   for (int k = 0; k < n; ++k) {
     for (int i = 0; i < n; ++i) {
       for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j], w[i][k]
      ] + w[k][j] - z[k]);
     }
  }
}
int solve(int n, vector<int> mark) {
   build(n);
   int k = (int)mark.size();
   assert(k < maxk);</pre>
   for (int s = 0; s < (1 << k); ++s) {
  for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
   for (int i = 0; i < n; ++i) dp[0][i] = 0;
   for (int s = 1; s < (1 << k); ++s) {
     if (__builtin_popcount(s) == 1) {
       int x = __builtin_ctz(s);
       for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]][i];
       continue;
     for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1)) {
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s \land sub][i] -
      z[i]);
       }
     for (int i = 0; i < n; ++i) {
       off[i] = inf;
       for (int j = 0; j < n; ++j) off[i] = min(off[i], dp[s][j]
       + w[j][i] - z[j]);
     for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i], off[i
     ]);
   for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k) - 1][i
     1);
   return res;
| }}
```

4.6 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
   T g[maxn][maxn], fw[maxn];
   int n, fr[maxn];
   bool vis[maxn], inc[maxn];
   void clear() {
     for(int i = 0; i < maxn; ++i) {</pre>
       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
       vis[i] = inc[i] = false;
  }
   void addedge(int u, int v, T w) {
     g[u][v] = min(g[u][v], w);
   T operator()(int root, int _n) {
     if (dfs(root) != n) return -1;
     T ans = 0;
     while (true) {
       for (int i = 1; i \le n; ++i) fw[i] = inf, fr[i] = i;
       for (int i = 1; i <= n; ++i) if (!inc[i]) {
          for (int j = 1; j \ll n; ++j) {
           if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
              fw[i] = g[j][i];
fr[i] = j;
         }
       }
       int x = -1;
for (int i = 1; i <= n; ++i) if (i != root && !inc[i]) {</pre>
         int j = i, c = 0;
          while (j != root && fr[j] != i && c <= n) ++c, j = fr[j]
          if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
         for (int i = 1; i <= n; ++i) if (i != root && !inc[i])</pre>
      ans += fw[i];
         return ans;
       int y = x;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
       do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true; }
      while (y != x);
       inc[x] = false;
       for (int k = 1; k \le n; ++k) if (vis[k]) {
          for (int j = 1; j <= n; ++j) if (!vis[j]) {
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x]) g[j][</pre>
      x] = g[j][k] - fw[k];
       }
     return ans;
   int dfs(int now) {
     int r = 1;
     vis[now] = true;
     for (int i = 1; i \le n; ++i) if (g[now][i] < inf && !vis[i]
      ]) r += dfs(i);
     return r;
| };
```

4.7 Maximum Matching on General Graph

```
| namespace matching {
  int fa[maxn], pre[maxn], match[maxn], s[maxn], v[maxn];
  vector<int> g[maxn];
  queue<int> q;
  void init(int n) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
    for (int i = 0; i < n; ++i) g[i].clear();
  }
  void add_edge(int u, int v) {
    g[u].push_back(v);
    g[v].push_back(u);
  }
  int find(int u) {
    if (u == fa[u]) return u;
    return fa[u] = find(fa[u]);
  }
}</pre>
```

```
|int lca(int x, int y, int n) {
   static int tk = 0;
   tk++;
   x = find(x), y = find(y);
   for (; ; swap(x, y)) {
     if (x != n) {
       if (v[x] == tk) return x;
       v[x] = tk;
       x = find(pre[match[x]]);
  }
}
void blossom(int x, int y, int l) {
  while (find(x) != l) {
    pre[x] = y;
     y = match[x];
     if (s[y] == 1) {
       q.push(y);
       s[y] = 0;
     if (fa[x] == x) fa[x] = 1;
     if (fa[y] == y) fa[y] = 1;
     x = pre[y];
  }
bool bfs(int r, int n) {
   for (int i = 0; i \le n; ++i) {
     fa[i] = i;
     s[i] = -1;
   while (!q.empty()) q.pop();
  q.push(r);
   s[r] = 0;
   while (!q.empty()) {
     int x = q.front(); q.pop();
     for (int u : g[x]) {
       if (s[u] == -1) {
         pre[u] = x;
         s[u] = 1;
         if (match[u] == n) {
           for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
             last = match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]);
         s[match[u]] = 0;
       } else if (!s[u] && find(u) != find(x)) {
         int l = lca(u, x, n);
         blossom(x, u, l);
         blossom(u, x, 1);
    }
  }
   return false;
}
int solve(int n) {
   int res = 0;
   for (int x = 0; x < n; ++x) {
    if (match[x] == n) res += bfs(x, n);
   return res;
| }}
```

4.8 Maximum Weighted Matching on General Graph

```
struct WeightGraph {
    static const int inf = INT_MAX;
    static const int maxn = 514;
    struct edge {
        int u, v, w;
        edge(){}
        edge(int u, int v, int w): u(u), v(v), w(w) {}
    };
    int n, n_x;
    edge g[maxn * 2][maxn * 2];
    int lab[maxn * 2];
    int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa[maxn * 2];
    int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[maxn * 2];
    vector<int> flo[maxn * 2];
    queue<int> q;
    int e_delta(const edge &e) {
```

```
return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
void update_slack(int u, int x) {
  if (!slack[x] | l e_delta(g[u][x]) < e_delta(g[slack[x]][x])
  ) slack[x] = u;
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)
    if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
      update_slack(u, x);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++) q_push(flo[</pre>
  x][i]);
void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
  set_st(flo[x][i], b);
int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) - flo[b].
  begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo[u][i</pre>
  ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
void augment(int u, int v) {
 for (; ; ) {
  int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
}
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y)
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w = 0;
  for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 \mid | e_delta(g[xs][x]) < e_delta(g[b][
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
```

```
if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
for (size_t i = pr + 1; i < flo[b].size(); ++i) {
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  else\ if\ (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true;
    else add_blossom(u, lca, v);
  return false;
}
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  a = queue<int>();
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0, q_push(
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
         if (g[u][v].w > 0 && st[u] != st[v]) {
           if (e_delta(g[u][v]) == 0) {
             if (on_found_edge(g[u][v])) return true;
           } else update_slack(u, st[v]);
        }
    int d = inf;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x]));
         else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x
   ]) / 2);
    for (int u = 1; u <= n; ++u) {
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
         if (S[st[b]] == 0) lab[b] += d * 2;
         else if (S[st[b]] == 1) lab[b] -= d * 2;
      }
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x] \&\& st[slack[x]] != x \&\&
   e_delta(g[slack[x]][x]) == 0)
        if (on_found_edge(g[slack[x]][x])) return true;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
  return false;
}
```

```
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
     n_x = n;
     int n_matches = 0;
     long long tot_weight = 0;
      for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear();</pre>
      int w_max = 0;
     for (int u = 1; u \le n; ++u)
        for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);
          w_max = max(w_max, g[u][v].w);
     for (int u = 1; u \leftarrow n; ++u) lab[u] = w_max;
     while (matching()) ++n_matches;
      for (int u = 1; u <= n; ++u)
        if (match[u] && match[u] < u)</pre>
          tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int ui, int vi, int wi) {
     g[ui][vi].w = g[vi][ui].w = wi;
   void init(int _n) {
     for (int u = 1; u <= n; ++u)
for (int v=1; v <= n; ++v)
          g[u][v] = edge(u, v, 0);
|};
```

4.9 Maximum Clique

struct MaxClique {

```
// change to bitset for n > 64.
  int n, deg[maxn];
 uint64_t adj[maxn], ans;
 vector<pair<int, int>> edge;
 void init(int n_) {
   fill(adj, adj + n, Oull);
    fill(deg, deg + n, 0);
   edge.clear();
 void add_edge(int u, int v) {
   edge.emplace_back(u, v);
    ++deg[u], ++deg[v];
 }
 vector<int> operator()() {
   vector<int> ord(n);
    iota(ord.begin(), ord.end(), 0);
   sort(ord.begin(), ord.end(), [&](int u, int v) { return deg
    [u] < deg[v]; });
    vector<int> id(n);
    for (int i = 0; i < n; ++i) id[ord[i]] = i;</pre>
    for (auto e : edge) {
      int u = id[e.first], v = id[e.second];
      adj[u] |= (1ull << v);
      adj[v] = (1ull \ll u);
   uint64_t r = 0, p = (1ull << n) - 1;
   dfs(r, p);
   vector<int> res;
for (int i = 0; i < n; ++i) {</pre>
      if (ans >> i & 1) res.push_back(ord[i]);
    return res;
#define pcount __builtin_popcountll
 void dfs(uint64_t r, uint64_t p) {
   if (p == 0) {
      if (pcount(r) > pcount(ans)) ans = r;
      return;
    if (pcount(r | p) <= pcount(ans)) return;</pre>
    int x = __builtin_ctzll(p & -p);
   uint64_t c = p & \sim adj[x];
   while (c > 0) {
      // bitset._Find_first(); bitset._Find_next();
      x = __builtin_ctzll(c & -c);
      r = (1ull \ll x);
      dfs(r, p & adj[x]);
      r &= ~(1ull << x);
      p &= ~(1ull << x);
```

```
c ^= (1ull << x);
}
| }
|};
```

4.10 Tarjan's Algorithm

```
void dfs(int x, int p) {
  dfn[x] = low[x] = tk++;
   int ch = 0;
   st.push(x); // bridge
   for (auto e : g[x]) if (e.first != p) {
     if (!ins[e.second]) { // articulation point
       st.push(e.second);
       ins[e.second] = true;
     if (~dfn[e.first]) {
       low[x] = min(low[x], dfn[e.first]);
       continue:
     dfs(u.first, x);
    if (low[u.first] >= low[x]) { // articulation point
       cut[x] = true;
       while (true) {
         int z = st.top(); st.pop();
         bcc[z] = sz;
         if (z == e.second) break;
       SZ++:
    }
  }
  if (ch == 1 \&\& p == -1) cut[x] = false;
  if (dfn[x] == low[x]) { // bridge
    while (true) {
       int z = st.top(); st.pop();
       bcc[z] = sz;
       if (z == x) break;
  }
}
```

4.11 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[maxn], val[
     maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1);
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
  fill(rp, rp + n, -1);
  fill(dom, dom + n, -1);
  for (int i = 0; i < n; ++i)
    g[i].clear();
}
void add_edge(int x, int y) {
  g[x].push_back(y);
}
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
 }
}
void merge(int x, int y) {
 fa[x] = y;
int find(int x, int c = 0) {
  if (fa[x] == x) return c? -1 : x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
```

```
|}
|vector<int> build(int s, int n) {
    // return the father of each node in the dominator tree
    // p[i] = -2 if i is unreachable from s
    dfs(s);
    for (int i = tk - 1; i >= 0; --i) {
        for (int u : r[i]) sdom[i] = min(sdom[i], sdom[find(u)]);
        if (i) rdom[sdom[i]].push_back(i);
        for (int &u : rdom[i]) {
            int p = find(u);
            if (sdom[p] == i) dom[u] = i;
            else dom[u] = p;
        }
        if (i) merge(i, rp[i]);
    }
    vector<int> p(n, -2);
    p[s] = -1;
    for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i]) dom[i] =
        dom[dom[i]];
    for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];
    return p;
}</pre>
```

4.12 Virtual Tree

```
void VirtualTree(vector<int> v) {
  static int stk[kN];
  int sz = 0:
  sort(v.begin(), v.end(), [&](int i, int j) { return dfn[i] <</pre>
  dfn[j]; });
stk[sz++] = 0;
 for (int u : v) {
  if (u == 0) continue;
    int p = LCA(u, stk[sz - 1]);
    if (p != stk[sz - 1]) {
      while (sz \ge 2 \&\& dfn[p] < dfn[stk[sz - 2]]) {
        AddEdge(stk[sz - 2], stk[sz - 1]);
         --sz:
      if (sz >= 2 && dfn[p] > dfn[stk[sz - 2]]) {
        AddEdge(p, stk[sz - 1]);
        stk[sz - 1] = p;
      } else {
        AddEdge(p, stk[--sz]);
      }
    stk[sz++] = u;
  for (int i = 0; i < sz - 1; ++i) AddEdge(stk[i], stk[i + 1]);
```

4.13 System of Difference Constraints

Given m constrains on n variables x_1, x_2, \ldots, x_n of form $x_i - x_j \leq w$ (resp, $x_i - x_j \geq w$), connect $i \to j$ with weight w. Then connect $0 \to i$ for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to x_i .

5 String

5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  // f[i] = length of the longest prefix (excluding s[0:i])
    such that it coincides with the suffix of s[0:i] of the
     same length
  // i + 1 - f[i] is the length of the smallest recurring
    period of s[0:i]
  for (int i = 1; i < (int)s.size(); ++i) {</pre>
    while (k > 0 \& s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
    f[i] = k;
 }
  return f;
vector<int> search(const string &s, const string &t) {
  // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), res;
```

```
int k = 0;
for (int i = 0; i < (int)s.size(); ++i) {
  while (k > 0 && (k == (int)t.size() || s[i] != t[k])) k = f
    [k - 1];
  if (s[i] == t[k]) ++k;
  if (k == (int)t.size()) res.push_back(i - t.size() + 1);
}
return res;
```

5.2 Z Algorithm

```
int z[maxn];
// z[i] = LCP of suffix i and suffix 0
void z_function(const string& s) {
    memset(z, 0, sizeof(z));
    z[0] = (int)s.length();
    int l = 0, r = 0;
    for (int i = 1; i < s.length(); ++i) {
        z[i] = max(0, min(z[i - l], r - i + 1));
        while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
        l = i; r = i + z[i];
        ++z[i];
    }
}
</pre>
```

5.3 Manacher's Algorithm

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for (int i = 0; i < s.length(); ++i) t += s[i], t += '.';
   int l = 0, r = 0;
   for (int i = 1; i < t.length(); ++i) {
      z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
      while (i - z[i] >= 0 && i + z[i] < t.length() && t[i - z[i] ]
      if (i + z[i] > r) r = i + z[i], l = i;
    }
   int ans = 0;
   for (int i = 1; i < t.length(); ++i) ans = max(ans, z[i] - 1)
    ;
   return ans;
}</pre>
```

5.4 Aho-Corasick Automaton

```
struct AC {
  static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn][26], f[
     maxn1:
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
cnt[sz] = 0;
    return sz++:
  void init() {
    sz = 0;
    root = gnode();
  int add(const string &s) {
    int now = root;
    for (int i = 0; i < s.length(); ++i) {
      if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a'] =
     gnode();
      now = ch[now][s[i] - 'a'];
    ed[now] = 1;
    return now;
  void build_fail() {
    ql = qr = 0; q[qr++] = root;
    while (ql < qr) {</pre>
      int now = q[ql++];
      for (int i = 0; i < 26; ++i) if (ch[now][i] != -1) {
```

```
int p = ch[now][i], fp = f[now];
while (fp != -1 && ch[fp][i] == -1) fp = f[fp];
           int pd = fp != -1 ? ch[fp][i] : root;
          f[p] = pd;
          el[p] = ed[pd] ? pd : el[pd];
          q[qr++] = p;
     }
   void build(const string &s) {
     build_fail();
      int now = root;
      for (int i = 0; i < s.length(); ++i) {</pre>
        while (now != -1 \& ch[now][s[i] - 'a'] == -1) now = f[
       now = now != -1 ? ch[now][s[i] - 'a'] : root;
        ++cnt[now];
      for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] += cnt[q[i]]
      ]];
   long long solve(int n) {
     build_fail();
      vector<vector<long long>> dp(sz, vector<long long>(n + 1,
      0));
      for (int i = 0; i < sz; ++i) dp[i][0] = 1; for (int i = 1; i <= n; ++i) {
        for (int j = 0; j < sz; ++j) {
          for (int k = 0; k < 2; ++k) {
  if (ch[j][k] != -1) {</pre>
               if (!ed[ch[j][k]])
                  dp[j][i] += dp[ch[j][k]][i - 1];
             } else {
               int z = f[j];
               while (z != root \&\& ch[z][k] == -1) z = f[z];
               int p = ch[z][k] == -1 ? root : ch[z][k];
if (ch[z][k] == -1 || !ed[ch[z][k]]) dp[j][i] += dp
      [p][i - 1];
          }
       }
     }
      return dp[0][n];
|};
```

5.5 Suffix Automaton

struct SAM {

```
static const int maxn = 5e5 + 5;
int nxt[maxn][26], to[maxn], len[maxn];
int root, last, sz;
int gnode(int x) {
  for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
  to[sz] = -1;
  len[sz] = x;
  return sz++;
void init() {
  sz = 0;
  root = gnode(0);
last = root;
void push(int c) {
  int cur = last;
  last = gnode(len[last] + 1);
  for (; ~cur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
  ] = last;
  if (cur == -1) return to[last] = root, void();
  int link = nxt[cur][c];
  if (len[link] == len[cur] + 1) return to[last] = link, void
  ();
  int tlink = gnode(len[cur] + 1);
  for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
  ][c] = tlink;
  for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];</pre>
  to[tlink] = to[link];
  to[link] = tlink;
  to[last] = tlink;
void add(const string &s) {
  for (int i = 0; i < s.size(); ++i) push(s[i] - 'a');
bool find(const string &s) {
```

```
int cur = root:
     for (int i = 0; i < s.size(); ++i) {</pre>
       cur = nxt[cur][s[i] - 'a'];
       if (cur == -1) return false;
     return true;
  }
   int solve(const string &t) {
     int res = 0, cnt = 0;
     int cur = root;
     for (int i = 0; i < t.size(); ++i) {</pre>
       if (~nxt[cur][t[i] - 'a']) {
         cur = nxt[cur][t[i] - 'a'];
       } else {
         for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
         if (\simcur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a'
         else cnt = 0, cur = root;
       res = max(res, cnt);
     return res:
1};
```

5.6 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2], x[maxn], p[maxn],
     q[maxn * 2];
  sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
     suffix.
// hi[i]: longest common prefix of suffix sa[i] and suffix sa[i
      - 17
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa, int *c, int *s, bool *t, int n, int z) {
  memcpy(x + 1, c, sizeof(int) * (z - 1));
  for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[x[
     s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i] - 1])
     sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int *c, int
      n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    return;
  for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i + 1] ? t
     [i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i - 1]) sa[--</pre>
     x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[i]
      - 1]) {
    bool neq = last < 0 || memcmp(s + sa[i], s + last, (p[q[sa[
     i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]] = p[
     nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
  _s[(int)s.size()] = 0; // s shouldn't contain 0
  sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for (int i = 0; i < (int)s.size(); ++i) sa[i] = sa[i + 1];</pre>
```

```
for (int i = 0; i < (int)s.size(); ++i) rev[sa[i]] = i;
int ind = 0; hi[0] = 0;
for (int i = 0; i < (int)s.size(); ++i) {
   if (!rev[i]) {
      ind = 0;
      continue;
   }
   while (i + ind < (int)s.size() && s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
   hi[rev[i]] = ind ? ind-- : 0;
}
</pre>
```

5.7 Lexicographically Smallest Rotation

```
| string rotate(const string &s) {
| int n = s.length();
| string t = s + s;
| int i = 0, j = 1;
| while (i < n && j < n) {
| int k = 0;
| while (k < n && t[i + k] == t[j + k]) ++k;
| if (t[i + k] <= t[j + k]) j += k + 1;
| else i += k + 1;
| if (i == j) ++j;
| }
| int pos = (i < n ? i : j);
| return t.substr(pos, n);
| }</pre>
```

6 Math

6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(re + rhs.
     re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(re - rhs.
    re, im - rhs.im); }
 cplx operator*(const cplx &rhs) const { return cplx(re * rhs.
    re - im * rhs.im, re * rhs.im + im * rhs.re); }
  cplx conj() const { return cplx(re, -im); }
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
  for (int i = 0; i \leftarrow maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi * i / maxn))
     maxn));
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    int x = 0;
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j & 1) << (z
     - j);
    if (x > i) swap(v[x], v[i]);
 }
void fft(vector<cplx> &v, int n) {
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
    cplx x = v[i + z + k] * omega[maxn / s * k];
        v[i + z + k] = v[i + k] - x;
        v[i + k] = v[i + k] + x;
   }
 }
void ifft(vector<cplx> &v, int n) {
 fft(v, n);
  reverse(v.begin() + 1, v.end());
  for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n, 0);
```

```
vector<long long> convolution(const vector<int> &a, const
      vector<int> &b) {
   // Should be able to handle N <= 10^5, C <= 10^4
   int sz = 1;
   while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
   vector<cplx> v(sz);
   for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;</pre>
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
   fft(v, sz);
   for (int i = 0; i \le sz / 2; ++i) {
     int j = (sz - i) & (sz - 1);
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()) * cplx
      (0, -0.25);
     if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i].conj
()) * cplx(0, -0.25);
     v[i] = x;
   ifft(v, sz);
   vector<long long> c(sz);
   for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
   return c:
}
vector<int> convolution_mod(const vector<int> &a, const vector<</pre>
      int> &b, int p) {
   int sz = 1:
  while (sz < (int)a.size() + (int)b.size() - 1) sz <<= 1;
   vector<cplx> fa(sz), fb(sz);
   for (int i = 0; i < (int)a.size(); ++i)</pre>
     fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
   for (int i = 0; i < (int)b.size(); ++i)</pre>
     fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
   fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
   for (int i = 0; i \leftarrow (sz >> 1); ++i) {
     int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx b1 = (fb[i] + fb[j].conj()) * r3;
     cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
       cplx c1 = (fa[j] + fa[i].conj());
cplx c2 = (fa[j] - fa[i].conj()) * r2;
       cplx d1 = (fb[j] + fb[i].conj()) * r3;
       cplx d2 = (fb[j] - fb[i].conj()) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \% p;
   return res;
}}
```

6.2 Number Theoretic Transform

```
if (x > i) swap(v[x], v[i]);
    }
  }
   void ntt(vector<long long> &v, int n) {
     bitrev(v, n);
     for (int s = 2; s <= n; s <<= 1) {
       int z = s \gg 1;
       for (int i = 0; i < n; i += s) {
         for (int k = 0; k < z; ++k) {
           long long x = v[i + k + z] * omega[maxn / s * k] %
           v[i + k + z] = (v[i + k] + mod - x) \% mod;
           (v[i + k] += x) \% = mod;
         }
      }
    }
  }
  void intt(vector<long long> &v, int n) {
     ntt(v, n);
     for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
     long long inv = fpow(n, -1);
     for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;
  vector<long long> operator()(vector<long long> a, vector<long</pre>
      long> b) {
     int sz = 1;
     while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
     while (a.size() < sz) a.push_back(0);
while (b.size() < sz) b.push_back(0);</pre>
     ntt(a, sz), ntt(b, sz);
     vector<long long> c(sz);
     for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] % mod;
     intt(c, sz);
     return c;
  }
};
vector<long long> convolution(vector<long long> a, vector<long</pre>
     long> b) {
  NTT<mod1, root1> conv1;
NTT<mod2, root2> conv2;
   vector<long long> pa(a.size()), pb(b.size());
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] \% mod1)
     + mod1) % mod1;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] \% mod1)
     + mod1) % mod1;
   vector<long long> c1 = conv1(pa, pb);
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] % mod2</pre>
     + mod2) % mod2;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] % mod2</pre>
     + mod2) % mod2;
   vector<long long> c2 = conv2(pa, pb);
  long long x = conv2.fpow(mod1, -1);
  long long y = conv1.fpow(mod2, -1);
  long long prod = mod1 * mod2;
   vector<long long> res(c1.size());
   for (int i = 0; i < c1.size(); ++i) {
     long long z = ((ull)fmul(c1[i] * mod2 % prod, y, prod) + (
     ull)fmul(c2[i] * mod1 % prod, x, prod)) % prod;
     if (z >= prod / 2) z -= prod;
     res[i] = z;
   return res:
| }
6.2.1 NTT Prime List
```

| Prime | Root | Prime | Root |
|-----------|------|------------|------|
| 7681 | 17 | 167772161 | 3 |
| 12289 | 11 | 104857601 | 3 |
| 40961 | 3 | 985661441 | 3 |
| 65537 | 3 | 998244353 | 3 |
| 786433 | 10 | 1107296257 | 10 |
| 5767169 | 3 | 2013265921 | 31 |
| 7340033 | 3 | 2810183681 | 11 |
| 23068673 | 3 | 2885681153 | 3 |
| 469762049 | 3 | 605028353 | 3 |

6.3 Polynomial Division

```
| vector<int> inverse(const vector<int> &v, int n) {
| vector<int> q(1, fpow(v[0], mod - 2));
| for (int i = 2; i <= n; i <<= 1) {
| vector<int> fv(v.begin(), v.begin() + i);
| vector<int> fq(q.begin(), q.end());
| fv.resize(2 * i), fq.resize(2 * i);
```

```
ntt(fq, 2 * i), ntt(fv, 2 * i);
for (int j = 0; j < 2 * i; ++j) {
    fv[j] = fv[j] * 1ll * fq[j] % mod * 1ll * fq[j] % mod;</pre>
      intt(fv, 2 * i);
      vector<int> res(i);
      for (int j = 0; j < i; ++j) {
  res[j] = mod - fv[j];</pre>
         if (j < (i >> 1)) (res[j] += 2 * q[j] % mod) %= mod;
      q = res;
   return q;
}
vector<int> divide(const vector<int> &a. const vector<int> &b)
    // leading zero should be trimmed
    int n = (int)a.size(), m = (int)b.size();
   int k = 2;
   while (k < n - m + 1) k <<= 1;
   vector<int> ra(k), rb(k);
   for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - i - 1];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - i - 1];
vector<int> rbi = inverse(rb, k);
   vector<int> res = convolution(rbi. ra):
   res.resize(n - m + 1);
   reverse(res.begin(), res.end());
    return res;
13
```

6.4 Polynomial Square Root

```
// Find G(x) such that G^2(x) = F(x) \pmod{x^{N+1}}
vector<int> solve(vector<int> b, int n) {
   if (n == 1) return {sqr[b[0]]};
   vector<int> h = solve(b, n >> 1); h.resize(n);
   vector<int> c = inverse(h, n);
   h.resize(n << 1); c.resize(n << 1);
   vector<int> res(n << 1);</pre>
   conv.ntt(h, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) b[i] = 0;
   conv.ntt(b, n << 1);</pre>
   conv.ntt(c, n << 1);</pre>
   for (int i = 0; i < (n << 1); ++i) res[i] = 1ll * (h[i] + 1ll
       * c[i] * b[i] % mod) % mod * inv2 % mod;
   conv.intt(res, n << 1);</pre>
   for (int i = n; i < (n << 1); ++i) res[i] = 0;
   return res;
}
```

6.5 Multipoint Evaluation

```
struct MultiEval {
  MultiEval *lc, *rc;
  vector<int> p, ml;
  // v is the points to be queried
  MultiEval(const vector<int> &v, int l, int r) : lc(nullptr),
     rc(nullptr) {
    if (r - 1 <= 64) {
      p = vector<int>(v.begin() + 1, v.begin() + r);
      ml.resize(1, 1);
      for (int x : p) ml = Multiply(ml, {kMod - x, 1});
      return;
    int m = (l + r) >> 1;
    lc = new MultiEval(v, l, m), rc = new MultiEval(v, m, r);
    ml = Multiply(lc->ml, rc->ml);
  // poly is the polynomial to be evaluated
  void Query(const vector<int> &poly, vector<int> &res, int l,
    int r) const {
if (r - l <= 64) {
      for (int x : p) {
        int s = 0, bs = 1;
        for (int i = 0; i < poly.size(); ++i) {</pre>
          (s += 1LL * bs * poly[i] % kMod) %= kMod;
          bs = 1LL * bs * x % kMod;
        res.push_back(s);
      return:
```

```
| auto pol = Modulo(poly, ml);
| int m = (l + r) >> 1;
| lc->Query(pol, res, l, m), rc->Query(pol, res, m, r);
| }
| };
```

6.6 Polynomial Interpolation

```
vector<int> Interp(const vector<int> &x, const vector<int> &y)
  vector<vector<int>>> v;
  int n = x.size();
  v.emplace_back(n);
  for (int i = 0; i < n; ++i) v[0][i] = \{\{kMod - x[i], 1\}\};
  while (v.back().size() > 1) {
    int n2 = v.back().size();
    vector<vector<int>>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) f[i] = Multiply(v.back ()[2 * i], v.back()[2 * i + 1]);
    if (n2 & 1) f.back() = v.back().back();
    v.push_back(f);
  vector<int> df(v.back()[0].size() - 1);
  for (int i = 0; i < df.size(); ++i) df[i] = 1LL * v.back() 
 [0][i + 1] * (i + 1) % kMod;
  vector<int> s;
  MultiEval(x, 0, n).Query(df, s, 0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) res[i] = {1LL * y[i] * fpow(s[i],
      kMod - 2) % kMod};
  for (int p = 1; p < v.size(); ++p) {</pre>
    int n2 = v[p - 1].size();
    vector<vector<int>>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) {
    auto a = Multiply(res[i * 2], v[p - 1][2 * i + 1]);
    auto b = Multiply(res[i * 2 + 1], v[p - 1][2 * i]);
       assert(a.size() == b.size());
       f[i].resize(a.size());
       for (int j = 0; j < a.size(); ++j) f[i][j] = (a[j] + b[j]
     ]) % kMod;
    if (n2 & 1) f.back() = res.back();
    res = f;
  }
  return res[0];
```

6.7 Fast Walsh-Hadamard Transform

6.7.1 XOR Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_0) tf(A_1))$
- $utf(A) = (utf(\frac{A_0 + A_1}{2}), utf(\frac{A_0 A_1}{2}))$

6.7.2 OR Convolution

- $tf(A) = (tf(A_0), tf(A_0) + tf(A_1))$
- $utf(A) = (utf(A_0), utf(A_1) utf(A_0))$

6.7.3 AND Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_1))$
- $utf(A) = (utf(A_0) utf(A_1), utf(A_1))$

6.8 Simplex Algorithm

```
| namespace simplex {
    // maximize c^Tx under Ax <= B
    // return vector<double>(n, -inf) if the solution doesn't exist
    // return vector<double>(n, +inf) if the solution is unbounded
    const double eps = 1e-9;
    const double inf = 1e+9;
    int n, m;
    vector<vector<double>> d;
    vector<int> p, q;
    void pivot(int r, int s) {
        double inv = 1.0 / d[r][s];
    }
}
```

```
for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {</pre>
       if (i != r \&\& j != s) d[i][j] -= d[r][j] * d[i][s] * inv;
   for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s] *= -inv;
for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j] *= +inv;
   d[r][s] = inv;
   swap(p[r], q[s]);
bool phase(int z) {
   int x = m + z;
   while (true) {
     for (int i = 0; i <= n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;</pre>
       if (r == -1 \mid | d[i][n + 1] / d[i][s] < d[r][n + 1] / d[r]
      ][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
 vector<double> solve(const vector<vector<double>> &a, const
      vector<double> &b, const vector<double> &c) {
   m = b.size(), n = c.size();
   d = vector<vector<double>>(m + 2, vector<double>(n + 2));
   for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
   p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][
      n + 1] = b[i];
   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   int r = 0:
   for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r
   if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps) return vector<</pre>
      double>(n, -inf);
     for (int i = 0; i < m; ++i) if (p[i] == -1) {
       int s = min_element(d[i].begin(), d[i].end() - 1) - d[i].
      begin();
       pivot(i, s);
   if (!phase(0)) return vector<double>(n, inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d[i][n +
      1];
   return x;
| }}
```

6.8.1 Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$ and $x_i \geq 0$ for all $1 \leq i \leq n$.

- 1. In case of minimization, let $c'_i = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3. $\sum_{1 < i < n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.9 Schreier-Sims Algorithm

```
namespace schreier {
int n;
vector<vector<int>>>> bkts, binv;
vector<vector<int>>> lk;
```

```
vector<int> operator*(const vector<int> &a, const vector<int> &
     b) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[i] = b[a[i]];</pre>
  return res;
vector<int> inv(const vector<int> &a) {
  vector<int> res(a.size());
for (int i = 0; i < (int)a.size(); ++i) res[a[i]] = i;</pre>
  return res:
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
 vector<int> p = g;
for (int i = 0; i < n; ++i) {
   assert(p[i] >= 0 && p[i] < (int)lk[i].size());</pre>
    int res = lk[i][p[i]];
    if (res == -1) {
      if (add) {
        bkts[i].push_back(p);
         binv[i].push_back(inv(p));
         lk[i][p[i]] = (int)bkts[i].size() - 1;
      return i:
    }
    p = p * binv[i][res];
  return -1;
bool inside(const vector<int> &g) { return filter(g, false) ==
     -1; }
void solve(const vector<vector<int>> &gen, int _n) {
  bkts.clear(), bkts.resize(n);
binv.clear(), binv.resize(n);
  lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
  for (int i = 0; i < n; ++i) {
    lk[i].resize(n, -1);
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i) filter(gen[i]);</pre>
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
      for (int k = 0; k < (int)bkts[i].size(); ++k) {</pre>
         for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l));
    }
  }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.first][b.
     second]);
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].size() -
     1);
    for (int i = 0; i < n; ++i) {
      for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
         if (i <= res) upd.emplace(make_pair(i, j), pr);</pre>
         if (res <= i) upd.emplace(pr, make_pair(i, j));</pre>
    }
  }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i) res = res * bkts[i].size();</pre>
  return res;
```

6.10 Berlekamp-Massey Algorithm

```
| template <int P>
| vector<int> BerlekampMassey(vector<int> x) {
| vector<int> cur, ls;
| int lf = 0, ld = 0;
| for (int i = 0; i < (int)x.size(); ++i) {</pre>
```

```
int t = 0;
for (int j = 0; j < (int)cur.size(); ++j)
  (t += 1LL * cur[j] * x[i - j - 1] % P) %= P;
      if (t == x[i]) continue;
      if (cur.empty()) {
         cur.resize(i + 1);
         lf = i, ld = (t + P - x[i]) % P;
         continue:
      int k = 1LL * fpow(ld, P - 2, P) * (t + P - x[i]) % P;
      vector<int> c(i - lf - 1);
      c.push_back(k);
      for (int j = 0; j < (int)ls.size(); ++j)
  c.push_back(1LL * k * (P - ls[j]) % P);</pre>
      if (c.size() < cur.size()) c.resize(cur.size());</pre>
      for (int j = 0; j < (int)cur.size(); ++j)</pre>
        c[j] = (c[j] + cur[j]) % P;
      if (i - lf + (int)ls.size() >= (int)cur.size()) {
  ls = cur, lf = i;
        ld = (t + P - x[i]) \% P;
      cur = c;
   }
   return cur:
}
```

6.11 Miller Rabin

```
// n < 4759123141
                     chk = [2, 7, 61]
// n < 1122004669633 chk = [2, 13, 23, 1662803]
                   chk = [2, 325, 9375, 28178, 450775, 9780504,
// n < 2^64
     17952650227
//
vector<long long> chk = { 2, 325, 9375, 28178, 450775, 9780504,
       1795265022 };
bool check(long long a, long long u, long long n, int t) {
  a = fpow(a, u, n);
   if (a == 0) return true;
   if (a == 1 \mid \mid a == n - 1) return true;
   for (int i = 0; i < t; ++i) {
     a = fmul(a, a, n);
     if (a == 1) return false;
     if (a == n - 1) return true;
  }
   return false;
bool is_prime(long long n) {
   if (n < 2) return false;</pre>
   if (n % 2 == 0) return n == 2;
   long long u = n - 1; int t = 0;
   for (; !(u & 1); u >>= 1, ++t);
for (long long i : chk) {
     if (!check(i, u, n, t)) return false;
   return true;
```

6.12 Pollard's Rho

```
map<long long, int> cnt;
 long long f(long long x, long long n, int p) { return (fmul(x, y))
      x, n) + p) % n; }
 void pollard_rho(long long n) {
   if (n == 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2 == 0) return pollard_rho(n / 2), ++cnt[2], void(); long long x = 2, y = 2, d = 1, p = 1;
   while (true) {
     if (d != n && d != 1) {
        pollard_rho(n / d);
        pollard_rho(d);
        return:
     if (d == n) ++p;
     x = f(x, n, p); y = f(f(y, n, p), n, p);
     d = \_gcd(abs(x - y), n);
}
```

6.13 Meissel-Lehmer Algorithm

```
int prc[maxn];
long long phic[msz][nsz];
void sieve() {
 bitset<maxn> v:
  pr.push_back(0);
  for (int i = 2; i < maxn; ++i) {
    if (!v[i]) pr.push_back(i);
    for (int j = 1; i * pr[j] < maxn; ++j) {
  v[i * pr[j]] = true;</pre>
      if (i % pr[j] == 0) break;
    }
  for (int i = 1; i < pr.size(); ++i) prc[pr[i]] = 1;</pre>
  for (int i = 1; i < maxn; ++i) prc[i] += prc[i - 1];
long long p2(long long, long long);
long long phi(long long m, long long n) {
 if (m < msz && n < nsz && phic[m][n] != -1) return phic[m][n</pre>
     1;
  if (n == 0) return m;
  if (pr[n] >= m) return 1;
  long long ret = phi(m, n - 1) - phi(m / pr[n], n - 1);
  if (m < msz && n < nsz) phic[m][n] = ret;</pre>
  return ret;
long long pi(long long m) {
  if (m < maxn) return prc[m];</pre>
  long long n = pi(cbrt(m));
  return phi(m, n) + n - 1 - p2(m, n);
long long p2(long long m, long long n) {
 long long ret = 0;
long long lim = sqrt(m);
  for (int i = n + 1; pr[i] <= lim; ++i) ret += pi(m / pr[i]) -
      pi(pr[i]) + 1;
  return ret;
```

6.14 Discrete Logarithm

```
// to solve discrete x for x^a = b \pmod{p} with p is prime
// let c = primitive root of p
// find k such that c^k = b \pmod{p} by bsgs
// solve fa = k \pmod{p-1} by euclidean algorithm
// x = c^f
int bsgs(int a, int b, int p) {
  // return L such that a^L = b \pmod{p}
  if (p == 1) {
    if (!b) return a != 1;
    return -1;
  if (b == 1) {
    if (a) return 0;
    return -1;
  if (a % p == 0) {
    if (!b) return 1;
    return -1;
  }
  int num = 0, d = 1;
  while (true) {
    int r = \__gcd(a, p);
    if (r == 1) break;
    if (b % r) return -1;
    ++num;
    b /= r, p /= r;
    d = (111 * d * a / r) \% p;
  for (int i = 0, now = 1; i < num; ++i, now = 1ll * now * a %
    if (now == b) return i;
  int m = ceil(sqrt(p)), base = 1;
  map<int, int> mp;
  for (int i = 0; i < m; ++i) {
    if (mp.find(base) == mp.end()) mp[base] = i;
    else mp[base] = min(mp[base], i);
base = 111 * base * a % p;
  for (int i = 0; i < m; ++i) {
    \ensuremath{\text{//}} can be modified to fpow if p is prime
```

```
int r, x, y; tie(r, x, y) = extgcd(d, p);
x = (1ll * x * b % p + p) % p;
if (mp.find(x) != mp.end()) return i * m + mp[x] + num;
d = 1ll * d * base % p;
}
return -1;
}
```

6.15 Quadratic Residue

```
int Jacobi(int a, int m) {
   int s = 1;
   for (; m > 1; ) {
    a %= m;
     if (a == 0) return 0;
     const int r = __builtin_ctz(a);
     if ((r \& 1) \& \& ((m + 2) \& 4)) s = -s;
     a >>= r;
     if (a \& m \& 2) s = -s;
     swap(a, m);
   return s;
}
int QuadraticResidue(int a, int p) {
   if (p == 2) return a & 1;
   const int jc = Jacobi(a, p);
   if (jc == 0) return 0;
   if (jc == -1) return -1;
   int b, d;
for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
   for (int e = (p + 1) >> 1; e; e >>= 1) {
     if (e & 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p
       g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
     tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
     f1 = (2LL * f0 * f1) % p;
     f0 = tmp;
   return g0;
1}
```

6.16 Gaussian Elimination

```
double gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   double det = 1;
   for (int i = 0; i < m; ++i) {
     int p = -1;
     for (int j = i; j < n; ++j) {
       if (fabs(d[j][i]) < eps) continue;</pre>
       if (p == -1 \mid | fabs(d[j][i]) > fabs(d[p][i])) p = j;
     if (p == -1) continue;
     if (p != i) det *= -1;
     for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
     for (int j = 0; j < n; ++j) {
       if (i == j) continue;
       double z = d[j][i] / d[i][i];
       for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
   for (int i = 0; i < n; ++i) det *= d[i][i];</pre>
   return det;
}
```

6.17 μ function

```
int mu[maxn], pi[maxn];
vector<int> prime;

void sieve() {
  mu[1] = pi[1] = 1;
```

```
for (int i = 2; i < maxn; ++i) {
    if (!pi[i]) {
        pi[i] = i;
        prime.push_back(i);
        mu[i] = -1;
    }
    for (int j = 0; i * prime[j] < maxn; ++j) {
        pi[i * prime[j]] = prime[j];
        mu[i * prime[j]] = -mu[i];
        if (i % prime[j] == 0) {
            mu[i * prime[j]] = 0;
            break;
        }
    }
}</pre>
```

6.18 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

6.19 De Bruijn Sequence

```
int res[maxn], aux[maxn], a[maxn], sz;
void db(int t, int p, int n, int k) {
 if (t > n) {
    if (n \% p == 0) {
      for (int i = 1; i <= p; ++i) res[sz++] = aux[i];</pre>
 } else {
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
    for (int i = aux[t - p] + 1; i < k; ++i) {
      aux[t] = i;
      db(t + 1, t, n, k);
 }
}
int de_bruijn(int k, int n) {
  // return cyclic string of length k^n such that every string
    of length n using k character appears as a substring.
  if (k == 1) {
    res[0] = 0;
    return 1;
  for (int i = 0; i < k * n; i++) aux[i] = 0;
  sz = 0;
  db(1, 1, n, k);
  return sz;
```

6.20 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
  if (!b) return make_tuple(a, 1, 0);
  T d, x, y;
  tie(d, x, y) = extgcd(b, a % b);
  return make_tuple(d, y, x - (a / b) * y);
}
```

6.21 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{a} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ -2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.22 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
   long long mult = mod[0];
   int n = (int)mod.size();
   long long res = a[0];
   for (int i = 1; i < n; ++i) {
      long long d, x, y;
      tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
      if ((a[i] - res) % d) return -1;
      long long new_mult = mult / __gcd(mult, 1ll * mod[i]) * mod
      [i];
      res += x * ((a[i] - res) / d) % new_mult * mult % new_mult;
      mult = new_mult;
      ((res %= mult) += mult) %= mult;
   }
   return res;
}</pre>
```

6.23 Theorem

6.23.1 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i), \, L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

6.23.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

6.23.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

6.23.4 Erdős-Gallai theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1+d_2+\ldots+d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \le k \le n$.

6.24 Primes

7 Dynamic Programming

7.1 Dynamic Convex Hull

```
struct Line {
  mutable int64_t a, b, p;
bool operator<(const Line &rhs) const { return a < rhs.a; }</pre>
   bool operator<(int64_t x) const { return p < x; }</pre>
 struct DynamicHull : multiset<Line, less<>>> {
   static const int64_t kInf = 1e18;
   int64_t Div(int64_t a, int64_t b) { return a / b - ((a \land b) <
       0 && a % b); }
   bool Isect(iterator x, iterator y) {
     if (y == end()) { x->p = kInf; return false; }
     if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
     else x->p = Div(y->b - x->b, x->a - y->a);
     return x->p >= y->p;
  }
   void Insert(int64_t a, int64_t b) {
     auto z = insert(\{a, b, 0\}), y = z++, x = y;
     while (Isect(y, z)) z = erase(z);
     if (x != begin() && Isect(--x, y)) Isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p) Isect(x,
      erase(y));
   int64_t Query(int64_t x) {
     auto l = *lower_bound(x);
     return l.a * x + l.b;
|};
```

7.2 1D/1D Convex Optimization

```
struct segment {
  int i, l, r;
  segment() {}
   segment(int a, int b, int c): i(a), l(b), r(c) {}
 inline long long f(int l, int r) {
  return dp[l] + w(l + 1, r);
 void solve() {
  dp[0] = 011;
   deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
     dp[i] = f(deq.front().i, i);
     while (deq.size() && deq.front().r < i + 1) deq.pop_front()</pre>
     deq.front().l = i + 1;
     segment seg = segment(i, i + 1, n);
     while (deq.size() && f(i, deq.back().1) < f(deq.back().i,</pre>
     deq.back().1)) deq.pop_back();
     if (dea.size()) {
       int d = 1048576, c = deq.back().1;
       while (d >>= 1) if (c + d <= deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c += d;
       deq.back().r = c; seg.l = c + 1;
     if (seg.l <= n) deq.push_back(seg);</pre>
1}
```

7.3 Condition

7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

7.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \ B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

8 Geometry

8.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps; }</pre>
struct P {
  double x, y;
  P() : x(0), y(0) {}
  P(double x, double y): x(x), y(y) {}

P operator + (P b) { return P(x + b.x, y + b.y); }

P operator - (P b) { return P(x - b.x, y - b.y); }

P operator * (double b) { return P(x * b, y * b); }
  P operator / (double b) { return P(x / b, y / b); } double operator * (P b) { return x * b.x + y * b.y; } double operator ^ (P b) { return x * b.y - y * b.x; }
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P spin(double o) {
    double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
  double angle() { return atan2(y, x); }
};
struct L {
  // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
  L(): a(0), b(0), c(0), o(0), pa(), pb() {}
  L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa ^ pb), o
      (atan2(-a, b)), pa(pa), pb(pb) {}
  P project(P p) { return pa + (pb - pa).unit() * ((pb - pa) *
      (p - pa) / (pb - pa).abs()); }
  P reflect(P p) { return p + (project(p) - p) * 2; }
  double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
     pa).abs() * (pb - pa).abs()); }
}:
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
  if (max(p1.x, p2.x) < min(p3.x, p4.x) | | max(p3.x, p4.x) < |
     min(p1.x, p2.x)) return false
  if (max(p1.y, p2.y) < min(p3.y, p4.y) || max(p3.y, p4.y) <
     min(p1.y, p2.y)) return false
  return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^ (p4 -
     p2)) <= 0 &&
       sign((p1 - p3) \land (p2 - p3)) * sign((p1 - p4) \land (p2 - p4))
}
bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b, x.a *
     y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

8.2 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[
    maxn];
point p[maxn];
```

```
int build(int 1, int r, int dep = 0) {
  if (l == r) return -1;
   function<br/><br/>bool(const point &, const point &)> f = [dep](const
      point &a, const point &b) {
     if (dep & 1) return a.x < b.x;</pre>
     else return a.y < b.y;</pre>
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
   xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
   lc[m] = build(l, m, dep + 1);
  if (\simlc[m]) {
     xl[m] = min(xl[m], xl[lc[m]]);
     xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
     xl[m] = min(xl[m], xl[rc[m]]);
     xr[m] = max(xr[m], xr[rc[m]]);
     yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  }
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
   if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
  q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
       (a.y - b.y) * 111 * (a.y - b.y);
void dfs(const point &q, long long &d, int o, int dep = 0) {
   if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x || !(dep & 1) && q.y < p[o].y)</pre>
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
}
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
return res;
j }}
```

8.3 Delaunay Triangulation

```
namespace triangulation {
static const int maxn = 1e5 + 5;
vector<point> p;
set<int> g[maxn];
int o[maxn];
set<int> s;
void add_edge(int x, int y) {
  s.insert(x), s.insert(y);
  g[x].insert(y);
  g[y].insert(x);
bool inside(point a, point b, point c, point p) {
  if (((b - a) ^ (c - a)) < 0) swap(b, c);</pre>
  function<long(int)> sqr = [](int x) { return x * 1ll * x}
     ; };
  long long k11 = a.x - p.x, k12 = a.y - p.y, k13 = sqr(a.x) -
     sqr(p.x) + sqr(a.y) - sqr(p.y);
  long long k21 = b.x - p.x, k22 = b.y - p.y, k23 = sqr(b.x) - b.y
     sqr(p.x) + sqr(b.y) - sqr(p.y);
  long long k31 = c.x - p.x, k32 = c.y - p.y, k33 = sqr(c.x) - c.y
  sqr(p.x) + sqr(c.y) - sqr(p.y);
long long det = k11 * (k22 * k33 - k23 * k32) - k12 * (k21 *
     k33 - k23 * k31) + k13 * (k21 * k32 - k22 * k31);
```

```
return det > 0:
bool intersect(const point &a, const point &b, const point &c,
     const point &d) {
  return ((b - a) \land (c - a)) * ((b - a) \land (d - a)) < 0 &&
      ((d - c) \wedge (a - c)) * ((d - c) \wedge (b - c)) < 0;
void dfs(int 1, int r) {
  if (r - 1 \le 3) {
    for (int i = 1; i < r; ++i) {
      for (int j = i + 1; j < r; ++j) add_edge(i, j);
    return:
  }
  int m = (l + r) >> 1;
  dfs(l, m), dfs(m, r);
  int pl = l, pr = r - 1;
  while (true) {
    int z = -1;
    for (int u : g[pl]) {
      long long c = ((p[pl] - p[pr]) \wedge (p[u] - p[pr]));
      if (c > 0 | | c == 0 \& abs(p[u] - p[pr]) < abs(p[pl] - p[
    pr])) {
    z = u;
        break;
      }
    if (z != -1) {
      pl = z;
      continue;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) \wedge (p[u] - p[pl]));
      if (c < 0 \mid l \mid c == 0 \& abs(p[u] - p[pl]) < abs(p[pr] - p[
     pl])) {
        z = u:
        break;
      }
    if (z != -1) {
      pr = z;
      continue;
    break;
  }
  add_edge(pl, pr);
  while (true) {
    int z = -1;
    bool b = false;
    for (int u : g[pl]) {
      long long c = ((p[pl] - p[pr]) \wedge (p[u] - p[pr]));
      if (c < 0 \& (z == -1 || inside(p[pl], p[pr], p[z], p[u])
     )) z = u;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) \land (p[u] - p[pl]));
      if (c > 0 && (z == -1 \mid \mid inside(p[pl], p[pr], p[z], p[u])
     )) z = u, b = true;
    if (z == -1) break;
    int x = pl, y = pr;
    if (b) swap(x, y);
    for (auto it = g[x].begin(); it != g[x].end(); ) {
      int u = *it;
      if (intersect(p[x], p[u], p[y], p[z])) {
        it = g[x].erase(it);
        g[u].erase(x);
      } else {
        ++it;
    if (b) add_edge(pl, z), pr = z;
    else add_edge(pr, z), pl = z;
vector<vector<int>>> solve(vector<point> v) {
  int n = v.size();
  for (int i = 0; i < n; ++i) g[i].clear();</pre>
  for (int i = 0; i < n; ++i) o[i] = i;
  sort(o, o + n, [&](int i, int j) { return v[i] < v[j]; });</pre>
  p.resize(n);
  for (int i = 0; i < n; ++i) p[i] = v[o[i]];
  dfs(0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) {
    for (int j : g[i]) res[o[i]].push_back(o[j]);
```

8.4 Sector Area

```
// calc area of sector which include a, b
| double SectorArea(P a, P b, double r) {
| double o = atan2(a.y, a.x) - atan2(b.y, b.x);
| while (o <= 0) o += 2 * pi;
| while (o >= 2 * pi) o -= 2 * pi;
| o = min(o, 2 * pi - o);
| return r * r * o / 2;
| }
```

8.5 Half Plane Intersection

```
bool jizz(L l1,L l2,L l3){
  P p=intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
bool cmp(const L &a,const L &b){
  return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;</pre>
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
   vector<L> pls(1,ls[0]);
  for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls.back().</pre>
      o))pls.push_back(ls[i]);
   deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],pls[b],
     pls[c]))
  for(int i=2;i<(int)pls.size();++i){</pre>
    meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
    meow(i,dq[0],dq[1])dq.pop_front();
    dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop_front();
  if(dq.size()<3u)return vector<P>(); // no solution or
     solution is not a convex
   vector<P> rt;
  for(int i=0;i<(int)dq.size();++i)rt.push_back(intersect(pls[</pre>
     dq[i]],pls[dq[(i+1)%dq.size()]]));
   return rt;
| }
```

8.6 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
 int n=int(ps.size());
 vector<int> id(n),pos(n);
 vector<pair<int,int>> line(n*(n-1)/2);
  int m=-1;
  for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=make_pair</pre>
    (i,j); ++m;
  sort(line.begin(),line.end(),[&](const pair<int,int> &a,const
     pair<int,int> &b)->bool{
    if(ps[a.first].first==ps[a.second].first)return 0;
    if(ps[b.first].first==ps[b.second].first)return 1;
   return (double)(ps[a.first].second-ps[a.second].second)/(ps
    [a.first].first-ps[a.second].first) < (double)(ps[b.first</pre>
    ].second-ps[b.second].second)/(ps[b.first].first-ps[b.
    second].first);
 });
 for(int i=0;i<n;++i)id[i]=i;</pre>
  sort(id.begin(),id.end(),[&](const int &a,const int &b){
    return ps[a]<ps[b]; });</pre>
  for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
  for(int i=0;i<m;++i){</pre>
   auto l=line[i];
   tie(pos[l.first],pos[l.second],id[pos[l.first]],id[pos[l.
    second]])=make_tuple(pos[l.second],pos[l.first],l.second,l
     .first);
```

8.7 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)) / (
     sin(a1) * cos(a2) - sin(a2) * cos(a1));
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
     TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);

res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

8.8 Polygon Center

```
| Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

8.9 Maximum Triangle

8.10 Point in Polygon

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) % ps.size();
    L l(ps[a], ps[b]);
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
    if (same(ps[a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
    if (ps[a].y > ps[b].y) swap(a, b);
    if (ps[a].y <= p.y && p.y < ps[b].y && p.x <= ps[a].x + (ps
        [b].x - ps[a].x) / (ps[b].y - ps[a].y) * (p.y - ps[a].y))
        ++c;
    }
    return (c & 1) * 2;
}</pre>
```

8.11 Circle

```
struct C {
  P c;
  double r;
  C(P \ c = P(0, 0), double \ r = 0) : c(c), r(r) \{\}
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c).unit()
     * a.r);
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
    double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.spin(o) * a.r);
    p.push_back(a.c + i.spin(-o) * a.r);
  return p;
double IntersectArea(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d \ge a.r + b.r - eps) return 0;
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
double q = acos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
  return p * sq(a.r) + q * sq(b.r) - a.r * d * <math>sin(p);
// remove second level if to get points for line (defalut:
     segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2
* x * (a.x - o.x) + 2 * y * (a.y - o.y);
  double C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
  if (d >= -eps) {
    d = \max(0., d);
    double i = (-B - sqrt(d)) / (2 * A);
double j = (-B + sqrt(d)) / (2 * A);
    if (i - 1.0 \le eps \&\& i \ge -eps) t.emplace_back(a.x + i * x)
      a.y + i * y);
    if (j - 1.0 \le eps \&\& j \ge -eps) t.emplace_back(a.x + j * x)
     , a.y + j * y);
  }
  return t:
// calc area intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
    return SectorArea(b, p[0], r) + abs(a \land p[0]) / 2;
  if (inb) return SectorArea(p[0], a, r) + abs(p[0] ^b / 2;
  if (p.size() == 2u) return SectorArea(a, p[0], r) +
     SectorArea(p[1], b, r) + abs(p[0] \land p[1]) / 2;
  else return SectorArea(a, b, r);
// for any trianale
double AreaOfCircleTriangle(vector<P> ps, double r) {
```

```
double ans = 0;
for (int i = 0; i < 3; ++i) {
   int j = (i + 1) % 3;
   double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y, ps[j].x
   );
   if (o >= pi) o = o - 2 * pi;
   if (o <= -pi) o = o + 2 * pi;
   ans += AreaOfCircleTriangle(ps[i], ps[j], r) * (o >= 0 ? 1
        : -1);
}
return abs(ans);
}
```

8.12 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
 #define Pij \
   P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
   z.emplace_back(a.c + i, a.c + i + j);
 #define deo(I,J) \
   double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos(e / d)
  P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-o);\
z.emplace_back(a.c + j * a.r, b.c J j * b.r);\
z.emplace_back(a.c + k * a.r, b.c J k * b.r);
   if (a.r < b.r) swap(a, b);
   vector<L> z;
   if ((a.c - b.c).abs() + b.r < a.r) return z;</pre>
   else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
   else {
     deo(-,+);
     if (same(d, a.r + b.r)) { Pij; }
     else if (d > a.r + b.r) \{ deo(+,-); \}
   return z;
}
vector<L> tangent(C c, P p) {
   vector<L> z;
double d = (p - c.c).abs();
   if (same(d, c.r)) {
     P i = (p - c.c).spin(pi / 2);
     z.emplace_back(p, p + i);
   } else if (d > c.r) {
     double o = acos(c.r / d);
     P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.spin(-o) * c.r;
     z.emplace_back(c.c + j, p);
     z.emplace_back(c.c + k, p);
   return z;
}
```

8.13 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
  vector<pair<double, double>> res;
if (same(a.r + b.r, d));
  else if (d \le abs(a.r - b.r) + eps) {
     if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
  } else if (d < abs(a.r + b.r) - eps) {
     double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
     ), z = (b.c - a.c).angle();
     if (z < 0) z += 2 * pi;
double l = z - o, r = z + o;
     if (l < 0) l += 2 * pi;
     if (r > 2 * pi) r -= 2 * pi;
     if (l > r) res.emplace_back(l, 2 * pi), res.emplace_back(0,
      r);
    else res.emplace_back(l, r);
  return res;
double CircleUnionArea(vector<C> c) { // circle should be
     identical
   int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
     vector<pair<double, double>> s = {{2 * pi, 9}}, z;
     for (int j = 0; j < n; ++j) if (i != j) {
       z = CoverSegment(c[i], c[j]);
       for (auto &e : z) s.push_back(e);
```

```
}
sort(s.begin(), s.end());
auto F = [&] (double t) { return c[i].r * (c[i].r * t + c[i].c.x * sin(t) - c[i].c.y * cos(t)); };
for (auto &e : s) {
    if (e.first > w) a += F(e.first) - F(w);
    w = max(w, e.second);
}
return a * 0.5;
}
```

8.14 Minimun Distance of 2 Polygons

8.15 2D Convex Hull

```
bool operator < (const P &a, const P &b) { return same(a.x, b.x
     ) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(a.x, b.x
     ) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) \wedge (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return same(a.x,
     b.x) ? a.y < b.y : a.x < b.x; });
  for (int i = 0; i < ps.size(); ++i)</pre>
    while (p.size() \ge 2 \& crx(p[p.size() - 2], ps[i], p[p.
     size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() - 2], ps[i], p[p.size
     () - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double x) { return same(x, \emptyset) ? \emptyset : x > \emptyset ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector<P> p, u, d;
  CH() {}
  CH(vector<P> ps) : p(ps) {
    n = ps.size();
    rotate(p.begin(), min_element(p.begin(), p.end()), p.end())
    auto t = max_element(p.begin(), p.end());
    d = vector<P>(p.begin(), next(t));
    u = \text{vector} < P > (t, p.end()); u.push_back(p[0]);
  int find(vector<P> &v, P d) {
```

```
int l = 0, r = v.size();
     while (l + 5 < r) {
  int L = (l * 2 + r) / 3, R = (l + r * 2) / 3;
       if (v[L] * d > v[R] * d) r = R;
       else l = L;
     int x = 1;
     for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x] * d) x
      = i;
     return x:
   int findFarest(P v) {
     if (v.y > 0 \mid | v.y == 0 \& v.x > 0) return ((int)d.size() -
      1 + find(u, v)) % p.size();
     return find(d, v);
   P get(int 1, int r, P a, P b) {
     int s = sgn(crx(a, b, p[l % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(a, b, p[m % n])) == s) l = m;
       else r = m;
     return isLL(a, b, p[l % n], p[(l + 1) % n]);
   vector<P> getIS(P a, P b) {
     int X = findFarest((b - a).spin(pi / 2));
int Y = findFarest((a - b).spin(pi / 2));
     if (X > Y) swap(X, Y);
     if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0) return
{get(X, Y, a, b), get(Y, X + n, a, b)};</pre>
     return {};
   void update_tangent(P q, int i, int &a, int &b) {
     if (sgn(crx(q, p[a], p[i])) > 0) a = i;
     if (sgn(crx(q, p[b], p[i])) < 0) b = i;
   void bs(int l, int r, P q, int &a, int &b) {
     if (l == r) return;
     update_tangent(q, 1 % n, a, b);
     int s = sgn(crx(q, p[l % n], p[(l + 1) % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l = m;
     update_tangent(q, r % n, a, b);
   bool contain(P p) {
     if (p.x < d[0].x \mid | p.x > d.back().x) return 0;
     auto it = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
     if (it->x == p.x) {
       if (it->y > p.y) return 0;
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     it = lower_bound(u.begin(), u.end(), P(p.x, 1e12), greater<
      P>());
     if (it->x == p.x) {
       if (it->y < p.y) return 0;</pre>
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
   bool get_tangent(P p, int &a, int &b) { // b -> a
     if (contain(p)) return 0;
     int i = lower_bound(d.begin(), d.end(), p) - d.begin();
     bs(0, i, p, a, b);
bs(i, d.size(), p, a, b);
     i = lower_bound(u.begin(), u.end(), p, greater<P>()) - u.
      begin();
     bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a, b);
     bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.size(), p,
     a, b);
return 1;
  }
};
```

8.16 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d){
  return abs(((b-a)^(c-a))*(d-a))/6;
}
struct convex3D{
  static const int maxn=1010;
```

```
struct T{
      int a,b,c;
bool res;
       T(){}
      T(int a,int b,int c,bool res=1):a(a),b(b),c(c),res(res){}
  int n,m;
  P p[maxn];
  T f[maxn*8];
  int id[maxn][maxn];
  bool on(T &t,P &q){
      return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
       int f2=id[a][b];
       if(f[f2].res){
            if(on(f[f2],p[q]))dfs(q,f2);
           else{
                id[q][b]=id[a][q]=id[b][a]=m;
                 f[m++]=T(b,a,q,1);
      }
  }
  void dfs(int p,int i){
      f[i].res=0;
       meow(p,f[i].b,f[now].a);
      meow(p,f[i].c,f[now].b);
       meow(p,f[i].a,f[now].c);
  void operator()(){
       if(n<4)return
       if([&]()->int{
            for(int i=1; i< n; ++i)if(abs(p[0]-p[i])>eps)return swap(p[1],
             p[i]),0;
            return 1;
       }())return;
       if([&]()->int{
            for(int i=2; i< n; ++i)if(abs((p[0]-p[i])^(p[1]-p[i]))>eps)
             return swap(p[2],p[i]),0;
            return 1;
       }())return;
       if([&]()->int{
            for(int i=3;i<n;++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))*(p[i]-</pre>
             p[0]))>eps)return swap(p[3],p[i]),0;
               eturn 1;
       }())return;
       for(int i=0;i<4;++i){</pre>
           T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
            if(on(t,p[i]))swap(t.b,t.c)
            id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
            f[m++]=t;
       for(int i=4;i< n;++i) for(int j=0;j< m;++j) if(f[j].res \&\& on(f[j]) for(int j=0;j< m;++j) for(int j=0;j< m;+
             ],p[i])){
            dfs(i,j);
           break:
       int mm=m; m=0;
       for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
  bool same(int i,int j){
       return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].a])>eps
              || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].b])>eps ||
              absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
  }
  int faces(){
       for(int i=0;i<m;++i){</pre>
            int iden=1;
            for(int j=0;j<i;++j)if(same(i,j))iden=0;</pre>
           r+=iden;
       return r;
|} tb;
```

8.17 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^ p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
```

```
circle min_enclosing(vector<pt> &p) {
   random_shuffle(p.begin(), p.end());
   double r = 0.0:
   pt cent;
   for (int i = 0; i < p.size(); ++i) {
  if (norm2(cent - p[i]) <= r) continue;</pre>
     r = 0.0:
     for (int j = 0; j < i; ++j) {
        if (norm2(cent - p[j]) <= r) continue;</pre>
       cent = (p[i] + p[j]) / 2;
        r = norm2(p[j] - cent);
        for (int k = 0; k < j; ++k) {
          if (norm2(cent - p[k]) <= r) continue;</pre>
          cent = center(p[i], p[j], p[k]);
          r = norm2(p[k] - cent);
       }
     }
   }
   return circle(cent, sqrt(r));
1 }
```

8.18 Closest Pair

```
double closest_pair(int l, int r) {
   // p should be sorted increasingly according to the x-
      coordinates.
   if (l == r) return 1e9;
   if (r - l == 1) return dist(p[l], p[r]);
   int m = (l + r) >> 1;
   double d = min(closest_pair(l, m), closest_pair(m + 1, r));
   vector<int> vec;
for (int i = m; i >= l && fabs(p[m].x - p[i].x) < d; --i) vec</pre>
      .push_back(i);
   for (int i = m + 1; i \le r \& fabs(p[m].x - p[i].x) < d; ++i)
      vec.push_back(i);
   sort(vec.begin(), vec.end(), [&](int a, int b) { return p[a].
     y < p[b].y; \});
   for (int i = 0; i < vec.size(); ++i) {</pre>
     for (int j = i + 1; j < vec.size() && fabs(p[vec[j]].y - p[
      vec[i]].y) < d; ++j) {
       d = min(d, dist(p[vec[i]], p[vec[j]]));
    }
   return d;
1}
```

9 Miscellaneous

9.1 Bitwise Hack

```
| long long next_perm(long long v) {
| long long t = v | (v - 1);
| return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1))
| ;
| }
| void subset(long long s) {
| long long sub = s;
| while (sub) sub = (sub - 1) & s;
| }
```

9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) {
            x = s - 1 - x;
            y = s - 1 - y;
        }
        swap(x, y);
    }
}
```

return res:

```
9.3
        Java
import java.io.*;
import java.util.*;
import java.lang.*;
import java.math.*;
public class filename{
  static Scanner in = new Scanner(System.in);
   public static void main(String[] args) throws Exception {
     Scanner fin = new Scanner(new File("infile"));
PrintWriter fout = new PrintWriter("outfile", "UTF-8");
     fout.println(fin.nextLine());
     fout.close():
     while (in.hasNext()) {
       String str = in.nextLine(); // getline
       String stu = in.next(); // string
     System.out.println("Case #" + t);
     System.out.printf("%d\n", 7122);
     int[][] d = \{\{7,1,2,2\},\{8,7\}\};
     int g = Integer.parseInt("-123");
     long f = (long)d[0][2];
     List<Integer> l = new ArrayList<>();
     Random rg = new Random();
     for (int i = 9; i >= 0; --i) {
       l.add(Integer.valueOf(rg.nextInt(100) + 1));
       l.add(Integer.valueOf((int)(Math.random() * 100) + 1));
     Collections.sort(l, new Comparator<Integer>() {
       public int compare(Integer a, Integer b) { return a - b;
     });
     for (int i = 0; i < l.size(); ++i)</pre>
       System.out.print(l.get(i));
     Set<String> s = new HashSet<String>(); // TreeSet
s.add("jizz");
     System.out.println(s);
     System.out.println(s.contains("jizz"));
    Map<String, Integer> m = new HashMap<String, Integer>();
m.put("lol", 7122);
     System.out.println(m);
     for(String key: m.keySet())
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
     System.out.println(Math.acos(-1));
     BigInteger bi = in.nextBigInteger(), bj = new BigInteger("
      -7122"), bk = BigInteger.value0f(17171);
     int sgn = bi.signum(); // sign(bi)
     bi = bi.subtract(BigInteger.ONE).multiply(bj).divide(bj).
     and(bj).gcd(bj).max(bj).pow(87);
     int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
|}
```

9.4 Dancing Links

```
| namespace dlx {
| int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn], rw[maxn],
| bt[maxn], s[maxn], head, sz, ans;
| void init(int c) {
| for (int i = 0; i < c; ++i) {
| up[i] = dn[i] = bt[i] = i;
| lt[i] = i == 0 ? c : i - 1;
| rg[i] = i == c - 1 ? c : i + 1;
| s[i] = 0;
| }
| rg[c] = 0, lt[c] = c - 1;</pre>
```

```
up[c] = dn[c] = -1;
  head = c, sz = c + 1;
}
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
   int f = sz;
  for (int i = 0; i < (int)col.size(); ++i) {</pre>
     int c = col[i], v = sz++;
     dn[bt[c]] = v;
     up[v] = bt[c], bt[c] = v;
     rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r, cl[v] = c;
     ++s[c];
    if (i > 0) lt[v] = v - 1;
  lt[f] = sz - 1;
void remove(int c) {
  lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
   for (int i = dn[c]; i != c; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j]) {
  up[dn[j]] = up[j], dn[up[j]] = dn[j];
       --s[cl[j]];
  }
fvoid restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
    for (int j = lt[i]; j != i; j = lt[j]) {
}
       ++s[cl[j]];
       up[dn[j]] = j, dn[up[j]] = j;
  lt[rg[c]] = c, rg[lt[c]] = c;
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i)
     dn[bt[i]] = i, up[i] = bt[i];
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
   int w = c;
  for (int x = c; x != head; x = rg[x]) {
     if (s[x] < s[w]) w = x;
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
     for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j]);
     dfs(dep + 1);
     for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j]);
  }
  restore(w);
}
int solve() {
 ans = 1e9, dfs(0);
 return ans;
```

9.5 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second =
     weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i such
     that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int> &x,
     vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
    if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  }):
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(st[qr[i].first], ed[qr</pre>
     [i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
```

```
djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[x[i]],</pre>
    ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
}
void solve(int l, int r, vector<int> v, long long c) {
 if (l == r) {
  cost[qr[l].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first]) {
      printf("%lld\n", c);
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(minv,
    cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first);
 contract(l, m, lv, x, y);
long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
 solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

9.6 Manhattan Distance MST

```
void solve(int n) {
       init();
         vector<int> v(n), ds;
        for (int i = 0; i < n; ++i) {
               v[i] = i;
                ds.push_back(x[i] - y[i]);
        sort(ds.begin(), ds.end());
        ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
         sort(v.begin(), v.end(), [\&](int i, int j) { return x[i] == x}
                   [j] ? y[i] > y[j] : x[i] > x[j]; y[j] : x[i] > x[j]; y[j] : y[j] : y[i] > y[j] : y[i] > y[i]
         int j = 0;
        for (int i = 0; i < n; ++i) {
                int p = lower_bound(ds.begin(), ds.end(), x[v[i]] - y[v[i]]) - ds.begin() + 1;
                pair<int, int> q = query(p);
                 // query return prefix minimum
                 if (~q.second) add_edge(v[i], q.second);
                add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
       }
```

```
| }
| void make_graph() {
| solve(n);
| for (int i = 0; i < n; ++i) swap(x[i], y[i]);
| solve(n);
| for (int i = 0; i < n; ++i) x[i] = -x[i];
| solve(n);
| for (int i = 0; i < n; ++i) swap(x[i], y[i]);
| solve(n);
| }</pre>
```

9.7 IOI 2016 Alien trick

```
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative as well, depending on the problem.
    if (c - (1LL << d) < 0) continue;
    long long ck = c - (1LL << d);
    pair<long long, int> r = check(ck);
    if (r.second == k) return r.first - ck * k;
    if (r.second < k) c = ck;
  }
  pair<long long, int> r = check(c);
  return r.first - c * k;
}
```