1.1 vimec 1.2 Fast Integer Input 1.3 Increase stack size 1.4 Pragma optimization Plow, Matching 2.1 Dinic's Algorithm 2.2 Minimum-cost flow 2.3 Gomory-flu Tree 2.4 Stoer-Wagner Minimum Cut 2.5 Kuhn-Munkres Algorithm 2.6 Flow Model Data Structure 3.1 <ext phds=""> 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Mean Cycle 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.12 Virtual Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Dispurae Root 6.5 Multipoint Evaluation 6.6 Past Walsh-Hadmand Transform 6.7.1 XOR Convolution 6.7 Fast Walsh-Hadmand Transform 6.7.2 OR Convolution 6.7 Fast Walsh-Hadmand Transform 6.8 Simplex Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Missel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.23 Theorem 6.23 Theorem 6.23 Theorem 6.23 Little's Martix 6.23 Condition 6.23 Condition 6.25 Condition 6.25 Condition 6.25 Condition 6.25 Condition 6.25 Condition 6.27 Doptonial Concewe Hull 7.2 Doptonic Concewe Hull 7.2 Doptonic Concewe Convex  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 Doptonic Concewe Convex  Dynamic Programming 7.1 Totally Monotone (Concave/Convex)</ext>	Co	ntents
1.1 vimec   1.2 Fast Integer Input   1.3 Increases stack size   1.4 Pragma optimization   1.4 Pragma optimization   1.5 Plow, Matching   2.1 Dinic's Algorithm   2.2 Minimum-cost flow   2.3 Gomory-Hu Tree   2.4 Stoer-Wagner Minimum Cut   2.5 Kuhn-Munkres Algorithm   2.6 Flow Model   2.6 Flow M	Ва	sic
1.3   Increase stack size   1.4   Pragma optimization   Processing   2.1   Dinic's Algorithm   2.2   Minimum-cost flow   2.3   Gomory-Hu Tree   2.4   Stoer-Wagner Minimum Cut   2.5   Kuhn-Munkres Algorithm   2.6   Flow Model   Data Structure   3.1   <ext pubs =""> 3.2   Li Chao Tree   Craph   4.1   Link-Cut Tree   4.2   Heavy-Light Decomposition   4.3   Centroid Decomposition   4.4   Minimum Mean Cycle   4.5   Minimum Steiner Tree   4.6   Directed Minimum Spanning Tree   4.6   Directed Minimum Spanning Tree   4.7   Maximum Matching on General Graph   4.8   Maximum Weighted Matching on General Graph   4.9   Maximum Clique   4.10   Tarjan's Algorithm   4.11   Dominator Tree   4.12   Virtual Tree   4.13   System of Difference Constraints   String   5.1   Knuth-Morris-Pratt Algorithm   5.2   Algorithm   5.2   Algorithm   5.3   Manacher's Algorithm   5.4   Aho-Crasick Automaton   5.5   Suffix Automaton   5.5   Suffix Automaton   5.6   Suffix Automaton   6.7   Lexicographically Smallest Rotation   Math   6.1   Fast Fourier Transform   6.2   Number Theoretic Transform   6.2   Number Theoretic Transform   6.3   Polynomial Division   6.4   Polynomial Division   6.5   Multipoint Evaluation   6.7   Fast Walsh-Hadmand Transform   6.7   Fast Walsh-Hadmand Transform   6.7   And Convolution   6.7   Octoor   Convolution   6.7   Suffix Algorithm   6.8   Construction   6.9   Schreier-Sims Algorithm   6.10   Berlekamp-Massey Algorithm   6.10   Berlekamp-Massey Algorithm   6.12   Dulard's Rho   6.13   Millier Rabin   6.12   Dulard's Rho   6.13   Millier Rabin   6.12   Dulard's Rho   6.13   Millier Rabin   6.23   Construction   6.24   Construction   6.25   Construct</ext>		
1.4 Pragma optimization   Flow, Matching   2.1 Dinic's Algorithm   2.2 Minimum-cost flow   2.3 Gomory-Hu Tree   2.4 Stoer-Wagner Minimum Cut   2.5 Kuhn-Munkres Algorithm   2.6 Flow Model   2.6 Flow Model   2.6 Flow Model   2.7 Eval   2.6 Flow Model   2.8 Store-West   2.8 Eval   2.8 E	1.2	• •
Plow, Matching		
2.1 Dinic's Algorithm	1.4	Pragma optimization
2.2 Minimum-cost flow 2.3 Gomory-Hu Tree 2.4 Stoer-Wagner Minimum Cut. 2.5 Kuhn-Munkres Algorithm 2.6 Flow Model  Data Structure 3.1 ⟨ext/pbds⟩ 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Glique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XCR Convolution 6.8 Simplex Algorithm 6.8 Lonstruction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Genesamp-Massey Algorithm 6.10 Genesamp-Massey Algorithm 6.10 Genesamp-Massey Algorithm 6.10 Immediate Rabin 6.11 Pollard's Rho 6.12 Construction 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 µ function 6.18   The Immeration 6.19 Leuclidean Algorithms 6.23 Theorem 6.23 Theorem 6.23 Theorem 6.23 Theorem 6.24 Primes  Dynamic Convex Hull 7.1 Dynamic Convex Hull 7.1 Dynamic Convex Hull 7.2 ID/ID Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	Flo	
2.3 Gomory-Hu Tree 2.4 Stoer-Wagner Minimum Cut 2.5 Kuhn-Munkres Algorithm 2.6 Flow Model  Data Structure 3.1 <ext pbds=""> 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Weighted Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2 Number Theoretic Transform 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Missel-Lehmer Algorithm 6.14 Ulicer-Simple Algorithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 purple Elimination 6.19 Berlughese-Lehmer Algorithm 6.11 Ulicer Simple Algorithm 6.12 Dollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 purple Residue 6.20 Extended GCD 6.21 Euclidean Algorithm 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Trutte's Matrix 6.23.3 Cayley's Formula 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Totally Monotone (Concave/Convex)</ext>		9
2.4 Stoer-Wagner Minimum Cut 2.5 Kuhn-Munkres Algorithm 2.6 Flow Model  Data Structure 3.1 ⟨ext/pbds⟩ 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Metching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Aray 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.8 Simplex Algorithm 6.8 Simplex Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Ill Piller Rabin 6.11 Pollard's Rho 6.12 Euclidean Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 µ function 6.18 [‡] Enumeration 6.19 Light Enumeration 6.20 Extender GCD 6.21 Euclidean Algorithm 6.23 Theorem 6.23 Theorem 6.24 Primes  Dynamic Convex Hull 7.1 Dynamic Convex Hull 7.2 1D/LD Convex Optlmization 7.3.1 Totally Monotone (Concave/Convex)		
2.6 Flow Model  Data Structure 3.1 < ext/pbds> 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Hiller Rabin 6.11 µ function 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Ellmination 6.17 µ function 6.18 [a] Enumeration 6.19 De Bruign Sequence 6.20 Extended GCD 6.21 Euclidean Algorithm 6.3 Theorem 6.23.3 Theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		· ·
Data Structure   3.1		
3.1 ⟨ext/pbds⟩ 3.2 Li Chao Tree  Graph 4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manache's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 NCR Convolution 6.7.2 OR Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Gaussian Elimination 6.11 µ tunction 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 µ tunction 6.18 [½] Enumeration 6.29 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	2.6	Flow Model
3.2 Li Chao Tree	Da	ta Structure
Graph		
1.1 Link-Cut Tree	3.2	Li Chao Tree
4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Automaton 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \( \mu \) function 6.19 Let Leidlean Algorithm 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.2 Condition 7.3.1 Totally Monotone (Concave/Convex)		
4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Atnay 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8 I Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 µ function 6.19 Leptical Elemmeration 6.19 Leptical Elemmeration 6.19 Leptical Elemmeration 6.10 Polynomial Sequence 6.20 Extended GCD 6.21 Euclidean Algorithm 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.2 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2 Number Theoretic Transform 6.2 Number Theoretic Transform 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 µ function 6.19 Pinumeration 6.19 Brigin Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Clinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.2 Clyley's Formula 6.23.1 Totally Monotone (Concave/Convex)		* 9 *
4.6 Directed Minimum Spanning Tree 4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Automaton 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.19 Eruintion 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		*
4.7 Maximum Matching on General Graph 4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.3 AND Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Berlekamp-Massey Algorithm 6.10 Gludratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [♣] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Chiese Remainder Theorem 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
4.8 Maximum Weighted Matching on General Graph 4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Division 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \( \mu \) function 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.23 Tutte's Matrix 6.23 Cayley's Formula 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		1 9
4.9 Maximum Clique 4.10 Tarjan's Algorithm 4.11 Dominator Tree 4.12 Virtual Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 E Fujin Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.23.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		8
4.11 Dominator Tree 4.12 Virtual Tree 4.13 System of Difference Constraints  String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8.1 Construction 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \(\begin{array}{c} \mathbf{I} \text{purple interpolation} \mathbf{B} \text{cis} \text{due} \mathbf{C} \mathbf{I} \mat	4.9	
4.12 Virtual Tree   4.13 System of Difference Constraints		
### String  5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Array 5.7 Lexicographically Smallest Rotation    Math		
String		
5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.16 Gaussian Elimination 6.17 \(\mu\) function 6.18 \(\frac{1}{2}\) Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	G.	
5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Division 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton 5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7 Fast Walsh-Hadamard Transform 6.7.3 AND Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 μ function 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
5.5 Suffix Automaton 5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \( \mu\) function 6.18 \( \mu^2\) El numeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	5.3	=
5.6 Suffix Array 5.7 Lexicographically Smallest Rotation  Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.8.3 Simplex Algorithm 6.8.1 Construction 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \( \mu \) function 6.18 \( \frac{\pi}{\pi} \) Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoft's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 \( 1D/1D \) Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
Math         1           6.1 Fast Fourier Transform         1           6.2 Number Theoretic Transform         6.2.1 NTT Prime List           6.3 Polynomial Division         6.4 Polynomial Square Root           6.5 Multipoint Evaluation         6.6 Polynomial Interpolation           6.7 Fast Walsh-Hadamard Transform         6.7.1 XOR Convolution           6.7.2 OR Convolution         6.7.3 AND Convolution           6.8 Simplex Algorithm         6.8.1 Construction           6.9 Schreier-Sims Algorithm         6.10 Berlekamp-Massey Algorithm           6.11 Miller Rabin         6.12 Pollard's Rho           6.13 Meissel-Lehmer Algorithm         6.14 Discrete Logarithm           6.14 Discrete Logarithm         6.15 Quadratic Residue           6.16 Gaussian Elimination         6.17 \(\psi\$ function           6.18 \(\begin{array}{l} \frac{\pi}{l} \) Enumeration         6.19 De Bruijn Sequence           6.20 Extended GCD         6.21 Euclidean Algorithms           6.22 Chinese Remainder Theorem         6.23.2 Tutte's Matrix           6.23.2 Tutte's Matrix         6.23.2 Tutte's Matrix           6.23.2 Tutte's Matrix         6.23.4 Erdős-Gallai theorem           6.24 Primes         Dynamic Programming           7.1 Dynamic Convex Hull         7.2 1D/1D Convex Optimization           7.3 Conditon <t< td=""><td></td><td></td></t<>		
6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π/2] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π/2] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	ъл.	AL.
6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [ π / 1 Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [ ½ ] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)	6.2	
6.4 Polynomial Square Root 6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [ ½ ] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
6.5 Multipoint Evaluation 6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier–Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel–Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [ π		·
6.6 Polynomial Interpolation 6.7 Fast Walsh-Hadamard Transform 6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier–Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel–Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [ ½ ] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős–Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		ž ž
6.7.1 XOR Convolution 6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		÷
6.7.2 OR Convolution 6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)	6.7	
6.7.3 AND Convolution 6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.8 Simplex Algorithm 6.8.1 Construction 6.9 Schreier-Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 [π] Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3.1 Totally Monotone (Concave/Convex)		
6.9 Schreier–Sims Algorithm 6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel–Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 μ function 6.18 ⌊ ½ ∐ Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős–Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)	6.8	
6.10 Berlekamp-Massey Algorithm 6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel-Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 \(\mu\) function 6.18 \(\left \frac{n}{i}\right \text{ Enumeration}\) 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.11 Miller Rabin 6.12 Pollard's Rho 6.13 Meissel–Lehmer Algorithm 6.14 Discrete Logarithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 $\mu$ function 6.18 $\lfloor \frac{n}{4} \rfloor$ Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős–Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 $1D/1D$ Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		9
6.12 Pollard's Rho 6.13 Meissel–Lehmer Algorithm 6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 $\mu$ function 6.18 $\lfloor \frac{n}{i} \rfloor$ Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős–Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 $1D/1D$ Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.14 Discrete Logarithm 6.15 Quadratic Residue 6.16 Gaussian Elimination 6.17 $\mu$ function 6.18 $\lfloor \frac{n}{i} \rfloor$ Enumeration 6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 $1D/1D$ Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
$\begin{array}{c} 6.15 \text{ Quadratic Residue} \\ 6.16 \text{ Gaussian Elimination} \\ 6.17 \hspace{0.2cm} \mu \text{ function} \\ 6.18 \left \lfloor \frac{n}{i} \right \rfloor \text{ Enumeration} \\ 6.19 \text{ De Bruijn Sequence} \\ 6.20 \hspace{0.2cm} \text{ Extended GCD} \\ 6.21 \hspace{0.2cm} \text{ Euclidean Algorithms} \\ 6.22 \hspace{0.2cm} \text{ Chinese Remainder Theorem} \\ 6.23 \hspace{0.2cm} \text{ Theorem} \\ 6.23.1 \hspace{0.2cm} \text{ Kirchhoff's Theorem} \\ 6.23.2 \hspace{0.2cm} \text{ Tutte's Matrix} \\ 6.23.3 \hspace{0.2cm} \text{ Cayley's Formula} \\ 6.23.4 \hspace{0.2cm} \text{ Erdős-Gallai theorem} \\ 6.24 \hspace{0.2cm} \text{ Primes} \\ \\ \hline{\textbf{Dynamic Programming}} \\ 7.1 \hspace{0.2cm} \text{ Dynamic Convex Hull} \\ 7.2 \hspace{0.2cm} 1D/1D \hspace{0.2cm} \text{ Convex Optimization} \\ 7.3 \hspace{0.2cm} \text{ Conditon} \\ 7.3.1 \hspace{0.2cm} \text{ Totally Monotone (Concave/Convex)} \\ \end{array}$		
$\begin{array}{c} 6.16 \ {\rm Gaussian} \ {\rm Elimination} \\ 6.17 \ \mu \ {\rm function} \\ 6.18 \left \lfloor \frac{n}{i} \right \rfloor \ {\rm Enumeration} \\ 6.19 \ {\rm De} \ {\rm Bruijn} \ {\rm Sequence} \\ 6.20 \ {\rm Extended} \ {\rm GCD} \\ 6.21 \ {\rm Euclidean} \ {\rm Algorithms} \\ 6.22 \ {\rm Chinese} \ {\rm Remainder} \ {\rm Theorem} \\ 6.23 \ {\rm Theorem} \\ 6.23.1 \ {\rm Kirchhoff's} \ {\rm Theorem} \\ 6.23.2 \ {\rm Tutte's} \ {\rm Matrix} \\ 6.23.2 \ {\rm Tutte's} \ {\rm Matrix} \\ 6.23.3 \ {\rm Cayley's} \ {\rm Formula} \\ 6.23.4 \ {\rm Erdős-Gallai} \ {\rm theorem} \\ 6.24 \ {\rm Primes} \\ \\ \hline \begin{array}{c} {\rm Dynamic} \ {\rm Programming} \\ 7.1 \ {\rm Dynamic} \ {\rm Convex} \ {\rm Hull} \\ 7.2 \ 1D/1D \ {\rm Convex} \ {\rm Optimization} \\ 7.3 \ {\rm Conditon} \\ 7.3.1 \ {\rm Totally} \ {\rm Monotone} \ ({\rm Concave/Convex}) \\ \end{array}$		
$\begin{array}{c} 6.17~\mu~{\rm function} \\ 6.18~\lfloor\frac{n}{i}\rfloor~{\rm Enumeration} \\ 6.19~{\rm De~Bruijn~Sequence} \\ 6.20~{\rm Extended~GCD} \\ 6.21~{\rm Euclidean~Algorithms} \\ 6.22~{\rm Chinese~Remainder~Theorem} \\ 6.23~{\rm Theorem} \\ 6.23.1~{\rm Kirchhoff's~Theorem} \\ 6.23.2~{\rm Tutte's~Matrix} \\ 6.23.2~{\rm Tutte's~Matrix} \\ 6.23.3~{\rm Cayley's~Formula} \\ 6.23.4~{\rm Erdős-Gallai~theorem} \\ 6.24~{\rm Primes} \\ \\ \hline \begin{array}{c} {\bf Dynamic~Programming} \\ 7.1~{\rm Dynamic~Convex~Hull} \\ 7.2~{\rm 1}D/1D~{\rm Convex~Optimization} \\ 7.3~{\rm Conditon} \\ 7.3.1~{\rm Totally~Monotone~(Concave/Convex)} \\ \end{array}$		
6.19 De Bruijn Sequence 6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.20 Extended GCD 6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)	6.1	
6.21 Euclidean Algorithms 6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes   Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.22 Chinese Remainder Theorem 6.23 Theorem 6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.23.1 Kirchhoff's Theorem 6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes   Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.23.2 Tutte's Matrix 6.23.3 Cayley's Formula 6.23.4 Erdős-Gallai theorem 6.24 Primes   Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		3 Theorem
6.23.3 Cayley's Formula 6.23.4 Erdős–Gallai theorem 6.24 Primes   Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.23.4 Erdős–Gallai theorem 6.24 Primes  Dynamic Programming 7.1 Dynamic Convex Hull 7.2 1D/1D Convex Optimization 7.3 Conditon 7.3.1 Totally Monotone (Concave/Convex)		
6.24 Primes       1         Dynamic Programming       1         7.1 Dynamic Convex Hull       1         7.2 1D/1D Convex Optimization       1         7.3 Conditon       1         7.3.1 Totally Monotone (Concave/Convex)       1		
7.1 Dynamic Convex Hull  7.2 1D/1D Convex Optimization  7.3 Conditon  7.3.1 Totally Monotone (Concave/Convex)	6.2	
7.1 Dynamic Convex Hull  7.2 1D/1D Convex Optimization  7.3 Conditon  7.3.1 Totally Monotone (Concave/Convex)	D	namic Programming
7.2 1 <i>D</i> /1 <i>D</i> Convex Optimization	-	
7.3.1 Totally Monotone (Concave/Convex)		1D/1D Convex Optimization
	7.3	
		7.3.1 Totally Monotone (Concave/Convex)

```
8 Geometry
                          17
                          17
 8.1 Basic
 8.2 KD Tree .
 19
                          19
 8.11 Circle . . .
 8.12 Tangent of Circles and Points to Circle . . . . . . . . . .
                          20
 20
 21
                          21
 8.17 Minimum Enclosing Circle . . . . . . . . . . . . . . . .
                          22
 22
9 Miscellaneous
                          22
 9.1 Bitwise Hack . . . . . .
                          22
 9.2 Hilbert's Curve (faster Mo's algorithm) . . . . . . . . . . . . .
                          22
 9.3 Java . .
                          23
 9.5 Offline Dynamic MST
 9.7 IOI 2016 Alien trick . . . . . . . . . .
```

# 1 Basic

#### 1.1 vimrc

```
se nu rnu bs=2 ru mouse=a cin et ts=4 sw=4 sts=4 syn on colo desert filetype indent on inoremap {<CR> {<CR>}<Esc>0
```

# 1.2 Fast Integer Input

```
| inline int gtx() {
   const int N = 4096;
   static char buffer[N];
   static char *p = buffer, *end = buffer;
   if (p == end) {
     if ((end = buffer + fread(buffer, 1, N, stdin)) == buffer)
     return EOF;
    p = buffer;
  }
   return *p++;
 template <typename T>
inline bool rit(T& x) {
  char c = 0; bool flag = false;
while (c = getchar(), (c < '0' && c != '-') || c > '9') if (c
  if (flag) x = -x;
   return true;
13
```

### 1.3 Increase stack size

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size) + size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

#### 1.4 Pragma optimization

# 2 Flow, Matching

# 2.1 Dinic's Algorithm

```
struct dinic {
  static const int inf = 1e9;
  struct edge {
    int to, cap, rev;
    edge(int d, int c, int r): to(d), cap(c), rev(r) {}
  vector<edge> g[maxn];
  int qu[maxn], ql, qr;
  int lev[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i)
       g[i].clear();
  void add_edge(int a, int b, int c) {
    g[a].emplace_back(b, c, g[b].size() - 0);
    g[b].emplace_back(a, 0, g[a].size() - 1);
  bool bfs(int s, int t) {
    memset(lev, -1, sizeof(lev));
     lev[s] = 0;
    ql = qr = 0;
    qu[qr++] = s;
    while (ql < qr) {
       int x = qu[ql++];
       for (edge &e : g[x]) if (lev[e.to] == -1 && e.cap > 0) {
        lev[e.to] = lev[x] + 1;
         qu[qr++] = e.to;
    }
    return lev[t] != -1;
  int dfs(int x, int t, int flow) {
    if (x == t) return flow;
     int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.to] == lev[x] +
       int f = dfs(e.to, t, min(e.cap, flow - res));
      res += f;
e.cap -= f;
      g[e.to][e.rev].cap += f;
     if (res == 0) lev[x] = -1;
    return res;
  int operator()(int s, int t) {
    int flow = 0;
     for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow;
|};
```

#### 2.2 Minimum-cost flow

```
struct mincost {
  struct edge {
    int dest, cap, w, rev;
    edge(int a, int b, int c, int d): dest(a), cap(b), w(c),
     rev(d) {}
  vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
  bool inq[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
  void add_edge(int a, int b, int c, int d) {
    g[a].emplace_back(b, c, +d, g[b].size() - 0);
g[b].emplace_back(a, 0, -d, g[a].size() - 1);
  bool spfa(int s, int t, int &f, int &c) {
    for (int i = 0; i < maxn; ++i) {
      d[i] = inf;
      p[i] = ed[i] = -1;
       inq[i] = false;
    d\lceil s \rceil = 0;
    queue<int> q;
    q.push(s);
```

```
while (q.size()) {
      int_x_= q.front(); q.pop();
      inq[x] = false;
      for (int i = 0; i < g[x].size(); ++i) {
        edge &e = g[x][i];
        if (e.cap > 0 \& d[e.dest] > d[x] + e.w) {
          d[e.dest] = d[x] + e.w;
          p[e.dest] = x;
           ed[e.dest] = i;
          if (!inq[e.dest]) q.push(e.dest), inq[e.dest] = true;
        }
      }
    if (d[t] == inf) return false;
    int dlt = inf;
    for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[p[x]][ed
     [x]].cap);
    for (int x = t; x != s; x = p[x]) {
      edge &e = g[p[x]][ed[x]];
e.cap -= dlt;
      g[e.dest][e.rev].cap += dlt;
    f += dlt; c += d[t] * dlt;
return true;
  pair<int, int> operator()(int s, int t) {
    int f = 0, c = 0;
    while (spfa(s, t, f, c));
    return make_pair(f, c);
2.3 Gomory-Hu Tree
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;</pre>
  for(int i=2;i<=n;++i){</pre>
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
flow.walk(i); // bfs points that connected to i (use edges
     not fully flow)
    for(int j=i+1;j<=n;++j){</pre>
      if(g[j]==t && flow.connect(j))g[j]=i; // check if i can
     reach j
    }
```

# 2.4 Stoer-Wagner Minimum Cut

return rt;

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
 bool v[maxn], del[maxn];
void add_edge(int x, int y, int c) {
   w[x][y] += c;
   w[y][x] += c;
pair<int, int> phase(int n) {
  memset(v, false, sizeof(v));
memset(g, 0, sizeof(g));
   int s = -1, t = -1;
   while (true) {
      int c = -1;
     for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
        if (c == -1 || g[i] > g[c]) c = i;
     if (c == -1) break;
     v[c] = true;
s = t, t = c;
     for (int i = 0; i < n; ++i) {
        if (del[i] || v[i]) continue;
        g[i] += w[c][i];
   return make_pair(s, t);
int mincut(int n) {
```

```
int cut = 1e9;
memset(del, false, sizeof(del));
for (int i = 0; i < n - 1; ++i) {
   int s, t; tie(s, t) = phase(n);
   del[t] = true;
   cut = min(cut, g[t]);
   for (int j = 0; j < n; ++j) {
      w[s][j] += w[t][j];
      w[j][s] += w[j][t];
   }
}
return cut;
}</pre>
```

# 2.5 Kuhn–Munkres Algorithm

```
int w[maxn][maxn], lx[maxn], ly[maxn];
int match[maxn], slack[maxn];
bool vx[maxn], vy[maxn];
bool dfs(int x) {
 vx[x] = true;
  for (int i = 0; i < n; ++i) {
    if (vy[i]) continue;
    if (lx[x] + ly[i] > w[x][i]) {
      slack[i] = min(slack[i], lx[x] + ly[i] - w[x][i]);
      continue;
    vy[i] = true;
    if (match[i] == -1 || dfs(match[i])) {
      match[i] = x;
      return true;
 }
  return false;
int solve() {
 fill_n(match, n, -1);
 fill_n(lx, n, -inf);
  fill_n(ly, n, 0);
  for (int i = 0; i < n; ++i) {
    for (int j = 0; j < n; ++j) lx[i] = max(lx[i], w[i][j]);
 for (int i = 0; i < n; ++i) {
    fill_n(slack, n, inf);
    while (true) {
      fill_n(vx, n, false);
      fill_n(vy, n, false);
      if (dfs(i)) break;
      int dlt = inf;
      for (int j = 0; j < n; ++j) if (!vy[j]) dlt = min(dlt,
    slack[j]);
      for (int j = 0; j < n; ++j) {
        if (vx[j]) lx[j] -= dlt;
if (vy[j]) ly[j] += dlt;
        else slack[j] -= dlt;
     }
   }
 }
  int res = 0;
  for (int i = 0; i < n; ++i) res += w[match[i]][i];</pre>
  return res;
```

# 2.6 Flow Model

- Maximum/Minimum flow with lower/upper bound from s to t
  - 1. Construct super source S and sink T
  - 2. For each edge (x, y, l, u), connect  $x \to y$  with capacity u l
  - 3. For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v)
    - To maximize, connect  $t \to s$  with capacity  $\infty$ , and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, f' is the answer.

- 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$  corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge  $(y \to x \text{ if } (x,y) \in M, x \to y \text{ otherwise})$
  - 2. DFS from unmatched vertices in X
  - 3.  $x \in X$  is chosen iff x is unvisited
  - 4.  $y \in Y$  is chosen iff y is visited
- Minimum cost cyclic flow
  - 1. Construct super source S and sink T
  - 2. For each edge (x,y,c), connect  $x\to y$  with (cost,cap)=(c,1) if c>0, otherwise connect  $y\to x$  with (cost,cap)=(-c,1)
  - 3. For each edge with c < 0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost, cap)=(0,d(v))
  - 5. For each vertex v with d(v) < 0, connect  $v \rightarrow T$  with (cost, cap) = (0, -d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- · Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v, v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G,$  connect it with sink  $v \to t$  with capacity  $K+2T-(\sum_{e \in E(v)} w(e)) 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|

# 3 Data Structure

# 3.1 < ext/pbds >

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>, rb_tree_tag,
     tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
  // rb tree
  tree set s
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22); assert(*s.find_by_order(1)
     == 71);
  assert(s.order_of_key(22) == 0); assert(s.order_of_key(71) ==
      1);
  s.erase(22);
  assert(*s.find_by_order(0) == 71); assert(s.order_of_key(71)
     == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;
  return 0;
```

#### 3.2 Li Chao Tree

```
namespace lichao {
 struct line {
   long long a, b;
   line(): a(0), b(0) {}
   line(long long a, long long b): a(a), b(b) {}
  long long operator()(int x) const { return a * x + b; }
 line st[maxc * 4];
 int sz, lc[maxc * 4], rc[maxc * 4];
int gnode() {
  st[sz] = line(1e9, 1e9);
lc[sz] = -1, rc[sz] = -1;
   return sz++;
 void init() {
  sz = 0;
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
   bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
   if (mcp) swap(st[o], tl);
   if (r - l == 1) return;
  if (lcp != mcp) {
   if (lc[o] == -1) lc[o] = gnode();
}
     add(1, (1 + r) / 2, t1, lc[o]);
     if (rc[o] == -1) rc[o] = gnode();
     add((l + r) / 2, r, tl, rc[o]);
 long long query(int l, int r, int x, int o) {
   if (r - l == 1) return st[o](x);
  if (x < (l + r) / 2) {
  if (lc[o] == -1) return st[o](x);</pre>
     return min(st[o](x), query(l, (l + r) / 2, x, lc[o]));
     if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((l + r) / 2, r, x, rc[o]));
| }}
```

# 4 Graph

#### 4.1 Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
  int sum, v, rev, id;
  node(int s, int id): id(id), v(s), sum(s), rev(0), fa(nullptr
     ), pfa(nullptr) {
    ch[0] = nullptr;
    ch[1] = nullptr;
  int relation() {
    return this == fa->ch[0] ? 0 : 1;
  void push() {
    if (!rev) return
    swap(ch[0], ch[1]);
if (ch[0]) ch[0]->rev ^= 1;
    if (ch[1]) ch[1]->rev ^= 1;
    rev = 0;
 }
  void pull() {
    sum = v
    if (ch[0]) sum += ch[0]->sum;
    if (ch[1]) sum += ch[1]->sum;
  void rotate() {
    if (fa->fa) fa->fa->push();
    fa->push(), push(), swap(pfa, fa->pfa);
    int d = relation();
    node *t = fa;
    if (t->fa) t->fa->ch[t->relation()] = this;
    fa = t->fa, t->ch[d] = ch[d ^ 1];
    if (ch[d \land 1]) ch[d \land 1] -> fa = t;
    ch[d \land 1] = t, t->fa = this;
    t->pull(), pull();
  void splay() {
```

```
while (fa) {
       if (!fa->fa) {
         rotate();
         continue;
       fa->fa->push(), fa->push();
       if (relation() == fa->relation()) fa->rotate();
       else rotate(), rotate();
   void evert() { access(), splay(), rev ^= 1; }
   void expose() {
     splay(), push();
     if (ch[1]) {
       ch[1]->fa = nullptr, ch[1]->pfa = this;
       ch[1] = nullptr, pull();
   bool splice() {
     splay();
     if (!pfa) return false;
     pfa->expose(), pfa->ch[1] = this, fa = pfa;
     pfa = nullptr, fa->pull();
     return true;
   void access() {
     expose();
     while (splice());
   int query() { return sum; }
};
namespace lct {
node *sp[maxn];
void make(int u, int v) {
  // create node with id u and value \boldsymbol{v}
   sp[u] = new node(v, u);
void link(int u, int v) {
  // u become v's parent
   sp[v]->evert();
   sp[v]->pfa = sp[u];
void cut(int u, int v) {
  // u was v's parent
   sp[u]->evert();
   sp[v]->access(), sp[v]->splay(), sp[v]->push();
   sp[v]->ch[0]->fa = nullptr;
   sp[v] -> ch[0] = nullptr;
   sp[v]->pull();
void modify(int u, int v) {
   sp[u]->splay();
   sp[u] -> v = v;
   sp[u]->pull();
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
   return sp[v]->query();
int find(int u) {
   sp[u]->access();
   sp[u]->splay();
   node *p = sp[u];
   while (true) {
     p->push();
     if (p->ch[0]) p = p->ch[0];
     else break;
  }
   return p->id;
| }}
```

#### 4.2 Heavy-Light Decomposition

```
void dfs(int x, int p) {
  dep[x] = ~p ? dep[p] + 1 : dep[x];
  sz[x] = 1;
  to[x] = -1;
  fa[x] = p;
  for (const int &u : g[x]) {
    if (u == p) continue;
    dfs(u, x);
    sz[x] += sz[u];
    if (to[x] == -1 || sz[to[x]] < sz[u]) to[x] = u;</pre>
```

```
}
}
void hld(int x, int t) {
  static int tk = 0;
  fr[x] = t;
  dfn[x] = tk++;
  if (!~to[x]) return;
  hld(to[x], t);
  for (const int &u : g[x]) {
     if (u == fa[x] || u == to[x]) continue;
    hld(u, u);
  }
}
vector<pair<int, int>> get(int x, int y) {
  int fx = fr[x], fy = fr[y];
  vector<pair<int, int>> res;
  while (fx != fy) {
    if (dep[fx] < dep[fy]) {</pre>
      swap(fx, fy);
       swap(x, y);
    res.emplace_back(dfn[fx], dfn[x] + 1);
    x = fa[fx];
    fx = fr[x];
  res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn[y]) +
     1);
  int lca = (dep[x] < dep[y] ? x : y);
  return res;
| }
```

# 4.3 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
   sz[now] = 1; mx[now] = 0;
   for (int u : G[now]) if (!v[u]) {
     get_center(u);
     mx[now] = max(mx[now], sz[u]);
     sz[now] += sz[u];
  }
 void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
   v[now] = true;
   for (auto u : G[now]) if (!v[u.first]) {
     get_dis(u, d, len + u.second);
 void dfs(int now, int fa, int d) {
   get_center(now);
   int c = -1;
   for (int i : vtx) {
     if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx.size()</pre>
      / 2) c = i;
     v[i] = false;
  get_dis(c, d, 0);
for (int i : vtx) v[i] = false;
  v[c] = true; vtx.clear();
  dep[c] = d; p[c] = fa;
for (auto u : G[c]) if (u.first != fa && !v[u.first]) {
     dfs(u.first, c, d + 1);
| }
```

#### 4.4 Minimum Mean Cycle

```
|// d[i][j] == 0 if {i,j} !in E
|long long d[1003][1003],dp[1003][1003];
|pair<long long,long long> MMWC(){
| memset(dp,0x3f,sizeof(dp));
| for(int i=1;i<=n;++i)dp[0][i]=0;
| for(int i=1;i<=n;++i){
| for(int j=1;j<=n;++j){
| for(int k=1;k<=n;++k){
| dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
| }
| }
| long long au=1ll<<31,ad=1;</pre>
```

```
for(int i=1;i<=n;++i){
   if(dp[n][i]==0x3f3f3f3f3f3f3f3f3f3f)continue;
   long long u=0,d=1;
   for(int j=n-1;j>=0;--j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
       u=dp[n][i]-dp[j][i];
      d=n-j;
    }
   }
   if(u*ad<au*d)au=u,ad=d;
}
long long g=__gcd(au,ad);
   return make_pair(au/g,ad/g);
}</pre>
```

# 4.5 Minimum Steiner Tree

```
namespace steiner {
// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
 // z[i] = the weight of the i-th vertex
const int maxn = 64, maxk = 10;
 const int inf = 1e9;
 int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[maxn];</pre>
 void init(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = inf;
     z[i] = 0;
     w[i][i] = 0;
  }
void add_edge(int x, int y, int d) {
   w[x][y] = min(w[x][y], d);
   w[y][x] = min(w[y][x], d);
}
void build(int n) {
   for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) {
       w[i][j] += z[i];
       if (i != j) w[i][j] += z[j];
   for (int k = 0; k < n; ++k) {
     for (int i = 0; i < n; ++i) {
       for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j], w[i][k]
      ] + w[k][j] - z[k]);
     }
  }
}
int solve(int n, vector<int> mark) {
   build(n);
   int k = (int)mark.size();
   assert(k < maxk);</pre>
   for (int s = 0; s < (1 << k); ++s) {
  for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
   for (int i = 0; i < n; ++i) dp[0][i] = 0;
   for (int s = 1; s < (1 << k); ++s) {
     if (__builtin_popcount(s) == 1) {
       int x = __builtin_ctz(s);
       for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]][i];
       continue;
     for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1)) {
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s \land sub][i] -
      z[i]);
       }
     for (int i = 0; i < n; ++i) {
       off[i] = inf;
       for (int j = 0; j < n; ++j) off[i] = min(off[i], dp[s][j]
       + w[j][i] - z[j]);
     for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i], off[i
     ]);
   for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k) - 1][i
     1);
   return res;
| }}
```

# 4.6 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
   T g[maxn][maxn], fw[maxn];
   int n, fr[maxn];
   bool vis[maxn], inc[maxn];
   void clear() {
     for(int i = 0; i < maxn; ++i) {</pre>
       for(int j = 0; j < maxn; ++j) g[i][j] = inf;
       vis[i] = inc[i] = false;
  }
   void addedge(int u, int v, T w) {
     g[u][v] = min(g[u][v], w);
   T operator()(int root, int _n) {
     if (dfs(root) != n) return -1;
     T ans = 0;
     while (true) {
       for (int i = 1; i \le n; ++i) fw[i] = inf, fr[i] = i;
       for (int i = 1; i <= n; ++i) if (!inc[i]) {
          for (int j = 1; j \ll n; ++j) {
           if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
              fw[i] = g[j][i];
fr[i] = j;
         }
       }
       int x = -1;
for (int i = 1; i <= n; ++i) if (i != root && !inc[i]) {</pre>
         int j = i, c = 0;
          while (j != root \&\& fr[j] != i \&\& c <= n) ++c, j = fr[j]
          if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
         for (int i = 1; i <= n; ++i) if (i != root && !inc[i])</pre>
      ans += fw[i];
         return ans;
       int y = x;
       for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
       do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true; }
      while (y != x);
       inc[x] = false;
       for (int k = 1; k \le n; ++k) if (vis[k]) {
          for (int j = 1; j <= n; ++j) if (!vis[j]) {
           if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x]) g[j][</pre>
      x] = g[j][k] - fw[k];
       }
     return ans;
   int dfs(int now) {
     int r = 1;
     vis[now] = true;
     for (int i = 1; i \le n; ++i) if (g[now][i] < inf && !vis[i]
      ]) r += dfs(i);
     return r;
| };
```

#### 4.7 Maximum Matching on General Graph

```
| namespace matching {
  int fa[maxn], pre[maxn], match[maxn], s[maxn], v[maxn];
  vector<int> g[maxn];
  queue<int> q;
  void init(int n) {
    for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
    for (int i = 0; i < n; ++i) g[i].clear();
  }
  void add_edge(int u, int v) {
    g[u].push_back(v);
    g[v].push_back(u);
  }
  int find(int u) {
    if (u == fa[u]) return u;
    return fa[u] = find(fa[u]);
  }
}</pre>
```

```
|int lca(int x, int y, int n) {
   static int tk = 0;
   tk++;
   x = find(x), y = find(y);
   for (; ; swap(x, y)) {
     if (x != n) {
       if (v[x] == tk) return x;
       v[x] = tk;
       x = find(pre[match[x]]);
  }
}
void blossom(int x, int y, int l) {
  while (find(x) != l) {
    pre[x] = y;
     y = match[x];
     if (s[y] == 1) {
       q.push(y);
       s[y] = 0;
     if (fa[x] == x) fa[x] = 1;
     if (fa[y] == y) fa[y] = 1;
     x = pre[y];
  }
bool bfs(int r, int n) {
   for (int i = 0; i \le n; ++i) {
     fa[i] = i;
     s[i] = -1;
   while (!q.empty()) q.pop();
  q.push(r);
   s[r] = 0;
   while (!q.empty()) {
     int x = q.front(); q.pop();
     for (int u : g[x]) {
       if (s[u] == -1) {
         pre[u] = x;
         s[u] = 1;
         if (match[u] == n) {
           for (int a = u, b = x, last; b != n; a = last, b =
     pre[a])
             last = match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]);
         s[match[u]] = 0;
       } else if (!s[u] && find(u) != find(x)) {
         int l = lca(u, x, n);
         blossom(x, u, l);
         blossom(u, x, 1);
    }
  }
   return false;
}
int solve(int n) {
   int res = 0;
   for (int x = 0; x < n; ++x) {
    if (match[x] == n) res += bfs(x, n);
   return res;
| }}
```

# 4.8 Maximum Weighted Matching on General Graph

```
struct WeightGraph {
    static const int inf = INT_MAX;
    static const int maxn = 514;
    struct edge {
        int u, v, w;
        edge(){}
        edge(int u, int v, int w): u(u), v(v), w(w) {}
    };
    int n, n_x;
    edge g[maxn * 2][maxn * 2];
    int lab[maxn * 2];
    int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa[maxn * 2];
    int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[maxn * 2];
    vector<int> flo[maxn * 2];
    queue<int> q;
    int e_delta(const edge &e) {
```

```
return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
void update_slack(int u, int x) {
  if (!slack[x] | l e_delta(g[u][x]) < e_delta(g[slack[x]][x])
  ) slack[x] = u;
void set_slack(int x) {
  slack[x] = 0;
  for (int u = 1; u <= n; ++u)
    if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
      update_slack(u, x);
void q_push(int x) {
  if (x \le n) q.push(x);
  else for (size_t i = 0; i < flo[x].size(); i++) q_push(flo[</pre>
  x][i]);
void set_st(int x, int b) {
  st[x] = b;
  if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
  set_st(flo[x][i], b);
int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) - flo[b].
  begin();
  if (pr % 2 == 1) {
    reverse(flo[b].begin() + 1, flo[b].end());
    return (int)flo[b].size() - pr;
  return pr;
void set_match(int u, int v) {
  match[u] = g[u][v].v;
  if (u <= n) return;</pre>
  edge e = g[u][v];
  int xr = flo_from[u][e.u], pr = get_pr(u, xr)
  for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo[u][i</pre>
  ^ 1]);
  set_match(xr, v);
  rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
void augment(int u, int v) {
 for (; ; ) {
  int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    u = st[pa[xnv]], v = xnv;
 }
int get_lca(int u, int v) {
  static int t = 0;
  for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
}
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[match[x]]),
   q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end());
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push_back(x), flo[b].push_back(y = st[match[x]]),
   q_push(y)
  set_st(b, b);
  for (int x = 1; x \le n_x; ++x) g[b][x].w = g[x][b].w = 0;
  for (int x = 1; x \le n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 \mid | e_delta(g[xs][x]) < e_delta(g[b][
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
```

```
if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
for (size_t i = pr + 1; i < flo[b].size(); ++i) {
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  else\ if\ (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true;
    else add_blossom(u, lca, v);
  return false;
}
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  a = queue<int>();
  for (int x = 1; x <= n_x; ++x)
if (st[x] == x && !match[x]) pa[x] = 0, S[x] = 0, q_push(
  if (q.empty()) return false;
  for (; ; ) {
    while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v <= n; ++v)
         if (g[u][v].w > 0 && st[u] != st[v]) {
           if (e_delta(g[u][v]) == 0) {
             if (on_found_edge(g[u][v])) return true;
           } else update_slack(u, st[v]);
        }
    int d = inf;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x]));
         else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x
   ]) / 2);
    for (int u = 1; u <= n; ++u) {
      if (S[st[u]] == 0) {
        if (lab[u] <= d) return 0;</pre>
        lab[u] -= d;
      } else if (S[st[u]] == 1) lab[u] += d;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b) {
         if (S[st[b]] == 0) lab[b] += d * 2;
         else if (S[st[b]] == 1) lab[b] -= d * 2;
      }
    q = queue<int>();
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x] \&\& st[slack[x]] != x \&\&
   e_delta(g[slack[x]][x]) == 0)
        if (on_found_edge(g[slack[x]][x])) return true;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b && S[b] == 1 && lab[b] == 0)
   expand_blossom(b);
  return false;
}
```

```
pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
     n_x = n;
     int n_matches = 0;
     long long tot_weight = 0;
      for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear();</pre>
      int w_max = 0;
     for (int u = 1; u \le n; ++u)
        for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);
          w_max = max(w_max, g[u][v].w);
     for (int u = 1; u \leftarrow n; ++u) lab[u] = w_max;
     while (matching()) ++n_matches;
      for (int u = 1; u <= n; ++u)
        if (match[u] && match[u] < u)</pre>
          tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int ui, int vi, int wi) {
     g[ui][vi].w = g[vi][ui].w = wi;
   void init(int _n) {
     for (int u = 1; u <= n; ++u)
for (int v=1; v <= n; ++v)
          g[u][v] = edge(u, v, 0);
|};
```

# 4.9 Maximum Clique

struct MaxClique {

```
// change to bitset for n > 64.
  int n, deg[maxn];
 uint64_t adj[maxn], ans;
 vector<pair<int, int>> edge;
 void init(int n_) {
   fill(adj, adj + n, Oull);
    fill(deg, deg + n, 0);
   edge.clear();
 void add_edge(int u, int v) {
   edge.emplace_back(u, v);
    ++deg[u], ++deg[v];
 }
 vector<int> operator()() {
   vector<int> ord(n);
    iota(ord.begin(), ord.end(), 0);
   sort(ord.begin(), ord.end(), [&](int u, int v) { return deg
    [u] < deg[v]; });
    vector<int> id(n);
    for (int i = 0; i < n; ++i) id[ord[i]] = i;</pre>
    for (auto e : edge) {
      int u = id[e.first], v = id[e.second];
      adj[u] |= (1ull << v);
      adj[v] = (1ull \ll u);
   uint64_t r = 0, p = (1ull << n) - 1;
   dfs(r, p);
   vector<int> res;
for (int i = 0; i < n; ++i) {</pre>
      if (ans >> i & 1) res.push_back(ord[i]);
    return res;
#define pcount __builtin_popcountll
 void dfs(uint64_t r, uint64_t p) {
   if (p == 0) {
      if (pcount(r) > pcount(ans)) ans = r;
      return;
    if (pcount(r | p) <= pcount(ans)) return;</pre>
    int x = __builtin_ctzll(p & -p);
   uint64_t c = p & \sim adj[x];
   while (c > 0) {
      // bitset._Find_first(); bitset._Find_next();
      x = __builtin_ctzll(c & -c);
      r = (1ull \ll x);
      dfs(r, p & adj[x]);
      r &= ~(1ull << x);
      p &= ~(1ull << x);
```

```
c ^= (1ull << x);
}
| }
|};
```

# 4.10 Tarjan's Algorithm

```
void dfs(int x, int p) {
  dfn[x] = low[x] = tk++;
   int ch = 0;
   st.push(x); // bridge
   for (auto e : g[x]) if (e.first != p) {
     if (!ins[e.second]) { // articulation point
       st.push(e.second);
       ins[e.second] = true;
     if (~dfn[e.first]) {
       low[x] = min(low[x], dfn[e.first]);
       continue:
     dfs(u.first, x);
    if (low[u.first] >= low[x]) { // articulation point
       cut[x] = true;
       while (true) {
         int z = st.top(); st.pop();
         bcc[z] = sz;
         if (z == e.second) break;
       SZ++:
    }
  }
  if (ch == 1 \&\& p == -1) cut[x] = false;
  if (dfn[x] == low[x]) { // bridge
    while (true) {
       int z = st.top(); st.pop();
       bcc[z] = sz;
       if (z == x) break;
  }
}
```

# 4.11 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[maxn], val[
     maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1);
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
  fill(rp, rp + n, -1);
  fill(dom, dom + n, -1);
  for (int i = 0; i < n; ++i)
    g[i].clear();
}
void add_edge(int x, int y) {
  g[x].push_back(y);
}
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  for (int u : g[x]) {
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
 }
}
void merge(int x, int y) {
 fa[x] = y;
int find(int x, int c = 0) {
  if (fa[x] == x) return c? -1 : x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
```

```
vector<int> build(int s, int n) {
 // return the father of each node in the dominator tree
  // p[i] = -2 if i is unreachable from s
 for (int i = tk - 1; i >= 0; --i) {
    for (int u : r[i]) sdom[i] = min(sdom[i], sdom[find(u)]);
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
else dom[u] = p;
    if (i) merge(i, rp[i]);
 }
 vector<int> p(n, -2);
 p[s] = -1;
  for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i]) dom[i] =</pre>
    dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
```

#### 4.12 Virtual Tree

```
void VirtualTree(vector<int> v) {
  static int stk[kN];
  int sz = 0:
  sort(v.begin(), v.end(), [&](int i, int j) { return dfn[i] <</pre>
  dfn[j]; });
stk[sz++] = 0;
  for (int i = 0; i < v.size(); ++i) {</pre>
    int u = v[i];
    ck[u] = true;
    if (u == 0) continue;
    int p = LCA(u, stk[sz - 1]);
    if (p != stk[sz - 1]) {
      while (sz \ge 2 \& dfn[p] < dfn[stk[sz - 2]]) {
        AddEdge(stk[sz - 2], stk[sz - 1]);
      if (sz >= 2 && dfn[p] > dfn[stk[sz - 2]]) {
        AddEdge(p, stk[sz - 1]);
        stk[sz - 1] = p;
        AddEdge(p, stk[--sz]);
    stk[sz++] = u;
  for (int i = 0; i < sz - 1; ++i) AddEdge(stk[i], stk[i + 1]);</pre>
```

#### 4.13 System of Difference Constraints

Given m constrains on n variables  $x_1, x_2, \ldots, x_n$  of form  $x_i - x_j \leq w$  (resp,  $x_i - x_j \geq w$ ), connect  $i \to j$  with weight w. Then connect  $0 \to i$  for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to  $x_i$ .

# 5 String

# 5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
  vector<int> f(s.size(), 0);
  // f[i] = length of the longest prefix (excluding s[0:i])
  such that it coincides with the suffix of s[0:i] of the
  same length
  // i + 1 - f[i] is the length of the smallest recurring
  period of s[0:i]
  int k = 0;
  for (int i = 1; i < (int)s.size(); ++i) {
    while (k > 0 && s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
    f[i] = k;
  }
  return f;
}
vector<int> search(const string &s, const string &t) {
```

```
// return 0-indexed occurrence of t in s
vector<int> f = kmp(t), res;
int k = 0;
int k = 0;
for (int i = 0; i < (int)s.size(); ++i) {
   while (k > 0 && (k == (int)t.size() || s[i] != t[k])) k = f
   [k - 1];
   if (s[i] == t[k]) ++k;
   if (k == (int)t.size()) res.push_back(i - t.size() + 1);
}
return res;
```

# 5.2 Z Algorithm

```
int z[maxn];
// z[i] = LCP of suffix i and suffix 0
void z_function(const string& s) {
    memset(z, 0, sizeof(z));
    z[0] = (int)s.length();
    int l = 0, r = 0;
    for (int i = 1; i < s.length(); ++i) {
        z[i] = max(0, min(z[i - l], r - i + 1));
        while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
            l = i; r = i + z[i];
            ++z[i];
        }
    }
}</pre>
```

# 5.3 Manacher's Algorithm

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for (int i = 0; i < s.length(); ++i) t += s[i], t += '.';
   int l = 0, r = 0;
   for (int i = 1; i < t.length(); ++i) {
        z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
        while (i - z[i] >= 0 && i + z[i] < t.length() && t[i - z[i] ]] == t[i + z[i]]) ++z[i];
        if (i + z[i] > r) r = i + z[i], l = i;
    }
   int ans = 0;
   for (int i = 1; i < t.length(); ++i) ans = max(ans, z[i] - 1)
    ;
   return ans;
}</pre>
```

#### 5.4 Aho-Corasick Automaton

```
struct AC {
  static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn][26], f[
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
    cnt[sz] = 0;
    return sz++;
  void init() {
    root = gnode();
  int add(const string &s) {
    int now = root;
    for (int i = 0; i < s.length(); ++i) {</pre>
      if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a'] =
     gnode();
      now = ch[now][s[i] - 'a'];
    ed[now] = 1;
  void build_fail() {
    ql = qr = 0; q[qr++] = root;
    while (ql < qr) {</pre>
```

```
int now = q[ql++];
for (int i = 0; i < 26; ++i)_if (ch[now][i] != -1) {</pre>
         int p = ch[now][i], fp = f[now];
         while (fp != -1 && ch[fp][i] == -1) fp = f[fp];
         int pd = fp != -1 ? ch[fp][i] : root;
         f[p] = pd;
         el[p] = ed[pd] ? pd : el[pd];
         q[qr++] = p;
    }
  }
   void build(const string &s) {
     build_fail();
     int now = root;
     for (int i = 0; i < s.length(); ++i) {</pre>
      while (now != -1 \& ch[now][s[i] - 'a'] == -1) now = f[
       now = now != -1 ? ch[now][s[i] - 'a'] : root;
       ++cnt[now];
     for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] += cnt[q[i]]
     ]];
   long long solve(int n) {
     build_fail();
     vector<vector<long long>> dp(sz, vector<long long>(n + 1,
     0));
     for (int i = 0; i < sz; ++i) dp[i][0] = 1;
     for (int i = 1; i <= n; ++i) {
       for (int j = 0; j < sz; ++j) {
         for (int k = 0; k < 2; ++k) {
  if (ch[j][k] != -1) {
             if (!ed[ch[j][k]])
                dp[j][i] += dp[ch[j][k]][i - 1];
           } else {
             int z = f[j];
             while (z != root \&\& ch[z][k] == -1) z = f[z];
             int p = ch[z][k] == -1 ? root : ch[z][k];
             if (ch[z][k] == -1 || !ed[ch[z][k]]) dp[j][i] += dp
      [p][i - 1];
           }
         }
      }
     return dp[0][n];
  }
|};
```

#### 5.5 Suffix Automaton

struct SAM {

```
static const int maxn = 5e5 + 5;
int nxt[maxn][26], to[maxn], len[maxn];
int root, last, sz;
int gnode(int x) {
  for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
  to[sz] = -1;
  len[sz] = x;
  return sz++:
3
void init() {
  sz = 0;
  root = gnode(0);
last = root;
void push(int c) {
  int cur = last;
  last = gnode(len[last] + 1);
  for (; ~cur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
   ] = last;
  if (cur == -1) return to[last] = root, void();
  int link = nxt[cur][c];
  if (len[link] == len[cur] + 1) return to[last] = link, void
   ();
  int tlink = gnode(len[cur] + 1);
  for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
   ][c] = tlink;
  for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];</pre>
  to[tlink] = to[link];
  to[link] = tlink;
  to[last] = tlink;
void add(const string &s) {
  for (int i = 0; i < s.size(); ++i) push(s[i] - 'a');</pre>
```

```
bool find(const string &s) {
     int cur = root;
     for (int i = 0; i < s.size(); ++i) {</pre>
       cur = nxt[cur][s[i] - 'a'];
       if (cur == -1) return false;
     return true:
   int solve(const string &t) {
     int res = 0, cnt = 0;
     int cur = root;
     for (int i = 0; i < t.size(); ++i) {</pre>
       if (~nxt[cur][t[i] - 'a']) {
         ++cnt;
         cur = nxt[cur][t[i] - 'a'];
       } else {
         for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
         if (\simcur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a'
         else cnt = 0, cur = root;
       }
       res = max(res, cnt);
     return res;
};
```

# 5.6 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2], x[maxn], p[maxn],
     q[maxn * 2];
  sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
     suffix.
// hi[i]: longest common prefix of suffix sa[i] and suffix sa[i
      - 1].
void pre(int *sa, int *c, int n, int z) {
  memset(sa, 0, sizeof(int) * n);
  memcpy(x, c, sizeof(int) * z);
void induce(int *sa, int *c, int *s, bool *t, int n, int z) {
 memcpy(x + 1, c, sizeof(int) * (z - 1));

for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[x[
     s[sa[i] - 1]]++] = sa[i] - 1;
  memcpy(x, c, sizeof(int) * z);
  for (int i = n - 1; i \ge 0; --i) if (sa[i] && t[sa[i] - 1])
     sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int *c, int
      n, int z) {
  bool uniq = t[n - 1] = true;
  int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last = -1;
  memset(c, 0, sizeof(int) * z);
  for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
  for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
  if (uniq) {
    for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
  for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i + 1] ? t
     [i + 1] : s[i] < s[i + 1]);
  pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i - 1]) sa[--
     x[s[i]]] = p[q[i] = nn++] = i;
  induce(sa, c, s, t, n, z);
  for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[i]
      - 1]) {
    bool neq = last < 0 || memcmp(s + sa[i], s + last, (p[q[sa[
     i]] + 1] - sa[i]) * sizeof(int));
    ns[q[last = sa[i]]] = nmxz += neq;
  sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + 1);
  pre(sa, c, n, z);
  for (int i = nn - 1; i \ge 0; --i) sa[--x[s[p[nsa[i]]]] = p[
     nsa[i]];
  induce(sa, c, s, t, n, z);
void build(const string &s) {
  for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];</pre>
  _s[(int)s.size()] = 0; // s shouldn't contain 0
```

```
sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for (int i = 0; i < (int)s.size(); ++i) sa[i] = sa[i + 1];
for (int i = 0; i < (int)s.size(); ++i) rev[sa[i]] = i;
int ind = 0; hi[0] = 0;
for (int i = 0; i < (int)s.size(); ++i) {
   if (!rev[i]) {
      ind = 0;
      continue;
   }
   while (i + ind < (int)s.size() && s[i + ind] == s[sa[rev[i] - 1] + ind]) ++ind;
   hi[rev[i]] = ind ? ind-- : 0;
}
</pre>
```

# 5.7 Lexicographically Smallest Rotation

```
| string rotate(const string &s) {
    int n = s.length();
    string t = s + s;
    int i = 0, j = 1;
    while (i < n && j < n) {
        int k = 0;
        while (k < n && t[i + k] == t[j + k]) ++k;
        if (t[i + k] <= t[j + k]) j += k + 1;
        else i += k + 1;
        if (i == j) ++j;
    }
    int pos = (i < n ? i : j);
    return t.substr(pos, n);
}</pre>
```

# 6 Math

#### 6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(re + rhs.
    re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(re - rhs.
    re, im - rhs.im); }
  cplx operator*(const cplx &rhs) const { return cplx(re * rhs.
    re - im * rhs.im, re * rhs.im + im * rhs.re); }
  cplx conj() const { return cplx(re, -im); }
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
  for (int i = 0; i \leftarrow maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi * i /
    maxn));
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    int x = 0:
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j & 1) << (z
      j);
    if (x > i) swap(v[x], v[i]);
  }
void fft(vector<cplx> &v, int n) {
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
      for (int k = 0; k < z; ++k) {
        cplx x = v[i + z + k] * omega[maxn / s * k];
        v[i + z + k] = v[i + k] - x;
        v[i + k] = v[i + k] + x;
    }
 }
void ifft(vector<cplx> &v, int n) {
 fft(v, n);
```

```
for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n, 0);
}
vector<long long> convolution(const vector<int> &a, const
      vector<int> &b) {
   // Should be able to handle N <= 10^5, C <= 10^4
   int sz = 1;
   while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
   vector<cplx> v(sz);
   for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;</pre>
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
   fft(v, sz);
   for (int i = 0; i \le sz / 2; ++i) {
     int j = (sz - i) & (sz - 1);
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()) * cplx
      (0, -0.25);
     if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i].conj
      ()) * cplx(0, -0.25);
     v[i] = x;
   ifft(v, sz);
   vector<long long> c(sz);
   for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);</pre>
vector<int> convolution mod(const vector<int> &a. const vector<
      int> &b, int p) {
   int sz = 1;
   while (sz < (int)a.size() + (int)b.size() - 1) sz <<= 1;</pre>
   vector<cplx> fa(sz), fb(sz);
for (int i = 0; i < (int)a.size(); ++i)</pre>
     fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
   for (int i = 0; i < (int)b.size(); ++i)</pre>
     fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
   fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1); for (int i = 0; i <= (sz >> 1); ++i) {
     int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx b1 = (fb[i] + fb[j].conj()) * r3;
     cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
       cplx c1 = (fa[j] + fa[i].conj());
       cplx c2 = (fa[j] - fa[i].conj()) * r2;
       cplx d1 = (fb[j] + fb[i].conj()) * r3;
       cplx d2 = (fb[j] - fb[i].conj()) * r4;
fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \% p;
   }
   return res;
1 } }
```

reverse(v.begin() + 1, v.end());

#### 6.2 Number Theoretic Transform

```
template <long long mod, long long root>
struct NTT {
  vector<long long> omega;
  NTT() {
    omega.resize(maxn + 1);
    long long x = fpow(root, (mod - 1) / maxn);
    omega[0] = 111;
    for (int i = 1; i <= maxn; ++i)
        omega[i] = omega[i - 1] * x % mod;
  }
  void bitrev(vector<long long> &v, int n) {
    int z = __builtin_ctz(n) - 1;
    for (int i = 0; i < n; ++i) {
        int x = 0;
    }
}</pre>
```

```
for (int j = 0; j \le z; ++j) x = (i >> j & 1) << (z - j)
      if (x > i) swap(v[x], v[i]);
   }
 }
  void ntt(vector<long long> &v, int n) {
    bitrev(v, n);
    for (int s = 2; s <= n; s <<= 1) {
      int z = s \gg 1;
      for (int i = 0; i < n; i += s) {
        for (int k = 0; k < z; ++k) {
          long long x = v[i + k + z] * omega[maxn / s * k] %
          v[i + k + z] = (v[i + k] + mod - x) \% mod;
          (v[i + k] += x) \%= mod;
     }
   }
 }
 void intt(vector<long long> &v, int n) {
    ntt(v, n);
    for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
    long long inv = fpow(n, -1);
    for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;
 vector<long long> operator()(vector<long long> a, vector<long</pre>
    int sz = 1;
    while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
    while (a.size() < sz) a.push_back(0);</pre>
    while (b.size() < sz) b.push_back(0);</pre>
    ntt(a, sz), ntt(b, sz);
    vector<long long> c(sz);
    for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] % mod;
    intt(c, sz);
    return c;
 }
vector<long long> convolution(vector<long long> a, vector<long
    long> b) {
 NTT<mod1, root1> conv1;
NTT<mod2, root2> conv2;
  vector<long long> pa(a.size()), pb(b.size());
 for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] % mod1</pre>
    + mod1) % mod1;
 for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] \% mod1
    + mod1) % mod1;
  vector<long long> c1 = conv1(pa, pb);
 for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i] \% mod2)
    + mod2) % mod2;
 for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i] % mod2</pre>
    + mod2) % mod2;
  vector<long long> c2 = conv2(pa, pb);
 long long x = conv2.fpow(mod1, -1);
 long long y = conv1.fpow(mod2, -1);
long long prod = mod1 * mod2;
 vector<long long> res(c1.size());
 for (int i = 0; i < c1.size(); ++i) {</pre>
    long long z = ((ull)fmul(c1[i] * mod2 % prod, y, prod) + (
    ull)fmul(c2[i] * mod1 % prod, x, prod)) % prod;
    if (z >= prod / 2) z -= prod;
   res[i] = z;
  return res;
```

### 6.2.1 NTT Prime List

Prime	Root	Prime	Root
7681	17	167772161	3
12289	11	104857601	3
40961	3	985661441	3
65537	3	998244353	3
786433	10	1107296257	10
5767169	3	2013265921	31
7340033	3	2810183681	11
23068673	3	2885681153	3
469762049	3	605028353	3

#### 6.3 Polynomial Division

```
vector<int> inverse(const vector<int> &v, int n) {
   vector<int> q(1, fpow(v[0], mod - 2));
   for (int i = 2; i <= n; i <<= 1) {
      vector<int> fv(v.begin(), v.begin() + i);
   }
}
```

```
vector<int> fq(q.begin(), q.end());
     fv.resize(2 * i), fq.resize(2 * i);
ntt(fq, 2 * i), ntt(fv, 2 * i);
for (int j = 0; j < 2 * i; ++j) {
    fv[j] = fv[j] * 1ll * fq[j] % mod * 1ll * fq[j] % mod;</pre>
      intt(fv, 2 * i);
      vector<int> res(i);
      for (int j = 0; j < i; ++j) {
  res[j] = mod - fv[j];</pre>
        if (j < (i >> 1)) (res[j] += 2 * q[j] % mod) %= mod;
     q = res;
   return q;
}
vector<int> divide(const vector<int> &a, const vector<int> &b)
   // leading zero should be trimmed
   int n = (int)a.size(), m = (int)b.size();
   int k = 2;
   while (k < n - m + 1) k <<= 1;
   vector<int> ra(k), rb(k);
   for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - i - 1];
   for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - i - 1];
   vector<int> rbi = inverse(rb, k);
   vector<int> res = convolution(rbi, ra);
   res.resize(n - m + 1);
   reverse(res.begin(), res.end());
   return res;
1 }
```

# 6.4 Polynomial Square Root

```
// Find G(x) such that G^2(x) = F(x) (mod x^{N+1})
vector<int> solve(vector<int> b, int n) {
   if (n == 1) return {sqr[b[0]]};
   vector<int> h = solve(b, n >> 1); h.resize(n);
   vector<int> c = inverse(h, n);
   h.resize(n << 1); c.resize(n << 1);
   vector<int> res(n << 1);
   conv.ntt(h, n << 1);
   for (int i = n; i < (n << 1); ++i) b[i] = 0;
   conv.ntt(b, n << 1);
   conv.ntt(c, n << 1);
   for (int i = 0; i < (n << 1); ++i) res[i] = 111 * (h[i] + 111
        * c[i] * b[i] % mod) % mod * inv2 % mod;
   conv.intt(res, n << 1);
   for (int i = n; i < (n << 1); ++i) res[i] = 0;
   return res;
}</pre>
```

# 6.5 Multipoint Evaluation

```
struct MultiEval {
  MultiEval *lc, *rc;
  vector<int> p, ml;
// v is the points to be queried
  MultiEval(const vector<int> &v, int l, int r) : lc(nullptr),
     rc(nullptr) {
    if (r - 1 \le 64) {
      p = vector<int>(v.begin() + l, v.begin() + r);
      ml.resize(1, 1);
      for (int x : p) ml = Multiply(ml, {kMod - x, 1});
      return;
    int m = (l + r) >> 1;
    lc = new MultiEval(v, l, m), rc = new MultiEval(v, m, r);
    ml = Multiply(lc->ml, rc->ml);
  // poly is the polynomial to be evaluated
  void Query(const vector<int> &poly, vector<int> &res, int l,
     int r) const {
    if (r - 1 <= 64) {
      for (int x : p) {
        int s = 0, bs = 1;
         for (int i = 0; i < poly.size(); ++i) {</pre>
          (s += 1LL * bs * poly[i] % kMod) %= kMod;
bs = 1LL * bs * x % kMod;
        res.push_back(s);
```

```
return;
}
auto pol = Modulo(poly, ml);
int m = (l + r) >> 1;
lc->Query(pol, res, l, m), rc->Query(pol, res, m, r);
}
};
```

# 6.6 Polynomial Interpolation

```
vector<int> Interp(const vector<int> &x, const vector<int> &y)
  vector<vector<int>>> v;
  int n = x.size();
  v.emplace_back(n);
  for (int i = 0; i < n; ++i) v[0][i] = \{\{kMod - x[i], 1\}\};
  while (v.back().size() > 1) {
    int n2 = v.back().size();
    vector<vector<int>>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) f[i] = Multiply(v.back ()[2 * i], v.back()[2 * i + 1]);
    if (n2 & 1) f.back() = v.back().back();
    v.push_back(f);
  }
  vector<int> df(v.back()[0].size() - 1);
  for (int i = 0; i < df.size(); ++i) df[i] = 1LL * v.back()
  [0][i + 1] * (i + 1) % kMod;</pre>
  vector<int> s;
  MultiEval(x, 0, n).Query(df, s, 0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) res[i] = {1LL * y[i] * fpow(s[i],
      kMod - 2) % kMod};
  for (int p = 1; p < v.size(); ++p) {
  int n2 = v[p - 1].size();</pre>
    vector<vector<int>> f((n2 + 1) >> 1);
    for (int i = 0; i < (n2 >> 1); ++i) {
      auto a = Multiply(res[i * 2], v[p - 1][2 * i + 1]);
      auto b = Multiply(res[i * 2 + 1], v[p - 1][2 * i]);
      assert(a.size() == b.size());
      f[i].resize(a.size());
      for (int j = 0; j < a.size(); ++j) f[i][j] = (a[j] + b[j]
     ]) % kMod;
    if (n2 & 1) f.back() = res.back();
    res = f;
  return res[0];
```

#### 6.7 Fast Walsh-Hadamard Transform

#### 6.7.1 XOR Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_0) tf(A_1))$
- $utf(A) = (utf(\frac{A_0 + A_1}{2}), utf(\frac{A_0 A_1}{2}))$

#### 6.7.2 OR Convolution

- $tf(A) = (tf(A_0), tf(A_0) + tf(A_1))$
- $utf(A) = (utf(A_0), utf(A_1) utf(A_0))$

#### 6.7.3 AND Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_1))$
- $utf(A) = (utf(A_0) utf(A_1), utf(A_1))$

# 6.8 Simplex Algorithm

```
| namespace simplex {
|// maximize c^Tx under Ax <= B
|// return vector<double>(n, -inf) if the solution doesn't exist
|// return vector<double>(n, +inf) if the solution is unbounded
const double eps = 1e-9;
| const double inf = 1e+9;
| int n, m;
| vector<vector<double>>> d;
| vector<int> p, q;
```

```
void pivot(int r, int s) {
   double inv = 1.0 / d[r][s];
   for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {</pre>
       if (i != r && j != s) d[i][j] -= d[r][j] * d[i][s] * inv;
     }
   for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s] *= -inv;
   for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j] *= +inv;
   d[r][s] = inv;
   swap(p[r], q[s]);
bool phase(int z) {
   int x = m + z;
   while (true) {
     int s = -1;
     for (int i = 0; i \le n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1 \mid | d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;</pre>
       if (r == -1 \mid | d[i][n + 1] / d[i][s] < d[r][n + 1] / d[r]
      ][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
}
vector<double> solve(const vector<vector<double>> &a, const
      vector<double> &b, const vector<double> &c) {
   m = b.size(), n = c.size();
   d = vector<vector<double>>(m + 2, vector<double>(n + 2));
   for (int i = 0; i < m; ++i) {
  for (int j = 0; j < n; ++j) d[i][j] = a[i][j];</pre>
   p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][
      n + 1] = b[i];
   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   int r = 0;
   for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r
      = i;
   if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) \mid \mid d[m + 1][n + 1] < -eps) return vector<
     double>(n, -inf);
for (int i = 0; i < m; ++i) if (p[i] == -1) {</pre>
       int s = min_element(d[i].begin(), d[i].end() - 1) - d[i].
      begin();
       pivot(i, s);
   if (!phase(0)) return vector<double>(n, inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d[i][n +
      17;
   return x;
}}
```

#### 6.8.1 Construction

Standard form: maximize  $\sum_{1 \leq i \leq n} c_i x_i$  such that for all  $1 \leq j \leq m$ ,  $\sum_{1 < i < n} A_{ji} x_i \leq b_j$  and  $x_i \geq 0$  for all  $1 \leq i \leq n$ .

- 1. In case of minimization, let  $c'_i = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
  - $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

# 6.9 Schreier-Sims Algorithm

```
namespace schreier {
int n:
vector<vector<int>>> bkts, binv;
vector<vector<int>> lk;
vector<int> operator*(const vector<int> &a, const vector<int> &
    b) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[i] = b[a[i]];
  return res;
vector<int> inv(const vector<int> &a) {
 vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[a[i]] = i;</pre>
  return res;
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
  vector<int> p = g;
  for (int i = 0; i < n; ++i) {
    assert(p[i] >= 0 && p[i] < (int)lk[i].size());
    int res = lk[i][p[i]];
    if (res == -1) {
      if (add) {
        bkts[i].push_back(p);
        binv[i].push_back(inv(p));
        lk[i][p[i]] = (int)bkts[i].size() - 1;
      return i;
    p = p * binv[i][res];
  return -1;
bool inside(const vector<int> &g) { return filter(g, false) ==
     -1; }
void solve(const vector<vector<int>> &gen, int _n) {
  bkts.clear(), bkts.resize(n);
  binv.clear(), binv.resize(n);
  lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
for (int i = 0; i < n; ++i) {</pre>
    lk[i].resize(n, -1);
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i) filter(gen[i]);</pre>
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
      for (int k = 0; k < (int)bkts[i].size(); ++k) {</pre>
        for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
          upd.emplace(make_pair(i, k), make_pair(j, l));
   }
 }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.first][b.
    second]);
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].size() -
    1);
    for (int i = 0; i < n; ++i) {
      for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
        if (i <= res) upd.emplace(make_pair(i, j), pr);</pre>
        if (res <= i) upd.emplace(pr, make_pair(i, j));</pre>
   }
 }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i) res = res * bkts[i].size();</pre>
  return res;
```

```
Berlekamp-Massey Algorithm
```

```
vector<int> BerlekampMassey(vector<int> x) {
   vector<int> cur, ls;
    int lf = 0, ld = 0;
    for (int i = 0; i < (int)x.size(); ++i) {</pre>
      for (int j = 0; j < (int)cur.size(); ++j)
  (t += 1LL * cur[j] * x[i - j - 1] % P) %= P;</pre>
      if (t == x[i]) continue;
      if (cur.empty()) {
        cur.resize(i + 1);
        lf = i, ld = (t + P - x[i]) % P;
        continue:
      int k = 1LL * fpow(ld, P - 2, P) * (t + P - x[i]) % P;
      vector<int> c(i - lf - 1);
      c.push_back(k);
      for (int j = 0; j < (int)ls.size(); ++j)
  c.push_back(1LL * k * (P - ls[j]) % P);</pre>
      if (c.size() < cur.size()) c.resize(cur.size());</pre>
      for (int j = 0; j < (int)cur.size(); ++j)
      c[j] = (c[j] + cur[j]) % P;
if (i - lf + (int)ls.size() >= (int)cur.size()) {
        ls = cur, lf = i;
        ld = (t + P - x[i]) \% P;
      cur = c;
   }
   return cur;
1 }
```

#### Miller Rabin 6.11

```
|// n < 4759123141
// n < 2^64
                 chk = [2, 325, 9375, 28178, 450775, 9780504,
     17952650227
//
vector<long long> chk = { 2, 325, 9375, 28178, 450775, 9780504,
      1795265022 };
bool check(long long a, long long u, long long n, int t) {
  a = fpow(a, u, n);
  if (a == 0) return true;
  if (a == 1 \mid \mid a == n - 1) return true;
  for (int i = 0; i < t; ++i) {
    a = fmul(a, a, n);
    if (a == 1) return false;
    if (a == n - 1) return true;
  return false:
}
bool is_prime(long long n) {
  if (n < 2) return false;
  if (n % 2 == 0) return n == 2;
  long long u = n - 1; int t = 0;
  for (; !(u & 1); u >>= 1, ++t);
for (long long i : chk) {
    if (!check(i, u, n, t)) return false;
  return true;
```

#### 6.12Pollard's Rho

```
map<long long, int> cnt;
long long f(long long x, long long n, int p) { return (fmul(x, y))
     x, n) + p) % n; }
void pollard_rho(long long n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2 == 0) return pollard_rho(n / 2), ++cnt[2], void();
long long x = 2, y = 2, d = 1, p = 1;
  while (true) {
    if (d != n && d != 1) {
       pollard_rho(n / d);
       pollard_rho(d);
       return;
    if (d == n) ++p;
    x = f(x, n, p); y = f(f(y, n, p), n, p);
    d = \_gcd(abs(x - y), n);
```

```
| }
|}
```

# 6.13 Meissel-Lehmer Algorithm

```
int prc[maxn];
long long phic[msz][nsz];
void sieve() {
  bitset<maxn> v
  pr.push_back(0);
   for (int i = 2; i < maxn; ++i) {
     if (!v[i]) pr.push_back(i);
     for (int j = 1; i * pr[j] < maxn; ++j) {
       v[i * pr[j]] = true;
       if (i % pr[j] == 0) break;
    }
  for (int i = 1; i < pr.size(); ++i) prc[pr[i]] = 1;</pre>
   for (int i = 1; i < maxn; ++i) prc[i] += prc[i - 1];</pre>
long long p2(long long, long long);
long long phi(long long m, long long n) {
  if (m < msz && n < nsz && phic[m][n] != -1) return phic[m][n</pre>
      ];
  if (n == 0) return m;
   if (pr[n] >= m) return 1;
  long long ret = phi(m, n - 1) - phi(m / pr[n], n - 1);
   if (m < msz && n < nsz) phic[m][n] = ret;</pre>
  return ret;
long long pi(long long m) {
   if (m < maxn) return prc[m];</pre>
  long long n = pi(cbrt(m));
return phi(m, n) + n - 1 - p2(m, n);
long long p2(long long m, long long n) {
   long long ret = 0;
   long long lim = sqrt(m);
   for (int i = n + 1; pr[i] <= lim; ++i) ret += pi(m / pr[i]) -</pre>
       pi(pr[i]) + 1;
   return ret;
| }
```

# 6.14 Discrete Logarithm

```
// to solve discrete x for x^a = b \pmod{p} with p is prime
// let c = primitive root of p
// find k such that c^k = b \pmod{p} by bsgs
// solve fa = k \pmod{p-1} by euclidean algorithm
// x = c^f
int bsgs(int a, int b, int p) {
  // return L such that a^L = b \pmod{p}
  if (p == 1) {
    if (!b) return a != 1;
return -1;
  if (b == 1) {
    if (a) return 0;
    return -1;
  if (a % p == 0) {
    if (!b) return 1;
return -1;
  int num = 0, d = 1;
  while (true) {
    int r = __gcd(a, p);
if (r == 1) break;
    if (b % r) return -1;
    ++num;
b /= r, p /= r;
d = (1ll * d * a / r) % p;
  for (int i = 0, now = 1; i < num; ++i, now = 1ll * now * a %
    p) {
    if (now == b) return i;
  int m = ceil(sqrt(p)), base = 1;
  map<int, int> mp;
  for (int i = 0; i < m; ++i) {
    if (mp.find(base) == mp.end()) mp[base] = i;
    else mp[base] = min(mp[base], i);
```

```
base = 111 * base * a % p;
}
for (int i = 0; i < m; ++i) {
    // can be modified to fpow if p is prime
    int r, x, y; tie(r, x, y) = extgcd(d, p);
    x = (111 * x * b % p + p) % p;
    if (mp.find(x) != mp.end()) return i * m + mp[x] + num;
    d = 111 * d * base % p;
}
return -1;
}</pre>
```

# 6.15 Quadratic Residue

```
| int Jacobi(int a, int m) {
   int s = 1;
   for (; m > 1; ) {
   a %= m;
    if (a == 0) return 0;
     const int r = __builtin_ctz(a);
     if ((r \& 1) \& ((m + 2) \& 4)) s = -s;
     if (a \& m \& 2) s = -s;
    swap(a, m);
  }
   return s;
}
int QuadraticResidue(int a, int p) {
   if (p == 2) return a & 1;
   const int jc = Jacobi(a, p);
   if (jc == 0) return 0;
   if (jc == -1) return -1;
   int b, d;
   for (; ; ) {
    b = rand() % p;
d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
   for (int e = (p + 1) >> 1; e; e >>= 1) {
     if (e & 1) {
       tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p
       q1 = (1LL * q0 * f1 + 1LL * q1 * f0) % p;
       g0 = tmp;
     tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
     f1 = (2LL * f0 * f1) % p;
     f0 = tmp;
   return g0;
}
```

# 6.16 Gaussian Elimination

```
double gauss(vector<vector<double>> &d) {
   int n = d.size(), m = d[0].size();
   double det = 1;
   for (int i = 0; i < m; ++i) {
     int p = -1;
     for (int j = i; j < n; ++j) {
       if (fabs(d[j][i]) < eps) continue;</pre>
       if (p == -1 \mid | fabs(d[j][i]) > fabs(d[p][i])) p = j;
     if (p == -1) continue;
     if (p != i) det *= -1;
     for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
     for (int j = 0; j < n; ++j) {
       if (i == j) continue;
       double z = d[j][i] / d[i][i];
       for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
   for (int i = 0; i < n; ++i) det *= d[i][i];</pre>
   return det;
}
```

### 6.17 $\mu$ function

```
int mu[maxn], pi[maxn];
 vector<int> prime;
 void sieve() {
   mu[1] = pi[1] = 1;
   for (int i = 2; i < maxn; ++i) {
      if (!pi[i]) {
        pi[i] = i;
        prime.push_back(i);
        mu[i] = -1;
     for (int j = 0; i * prime[j] < maxn; ++j) {
  pi[i * prime[j]] = prime[j];
  mu[i * prime[j]] = -mu[i];</pre>
         if (i % prime[j] == 0) {
           mu[i * prime[j]] = 0;
           break;
     }
   }
|}
```

# 6.18 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

# 6.19 De Bruijn Sequence

```
int res[maxn], aux[maxn], a[maxn], sz;
void db(int t, int p, int n, int k) {
  if (t > n) {
     if (n \% p == 0) {
       for (int i = 1; i \le p; ++i) res[sz++] = aux[i];
  } else {
     aux[t] = aux[t - p];
     db(t + 1, p, n, k);
     for (int i = aux[t - p] + 1; i < k; ++i) {
  aux[t] = i;</pre>
       db(t + 1, t, n, k);
  }
}
int de_bruijn(int k, int n) {
  // return cyclic string of length k^n such that every string
     of length n using k character appears as a substring.
   if (k == 1) {
     res[0] = 0;
  for (int i = 0; i < k * n; i++) aux[i] = 0;
  sz = 0;
  db(1, 1, n, k);
return sz;
į }
```

#### 6.20 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
  if (!b) return make_tuple(a, 1, 0);
  T d, x, y;
  tie(d, x, y) = extgcd(b, a % b);
  return make_tuple(d, y, x - (a / b) * y);
}
```

# 6.21 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^{n} i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

# 6.22 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
  long long mult = mod[0];
  int n = (int)mod.size();
  long long res = a[0];
  for (int i = 1; i < n; ++i) {
    long long d, x, y;
    tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
    if ((a[i] - res) % d) return -1;
  long long new_mult = mult / __gcd(mult, 1ll * mod[i]) * mod
    [i];
  res += x * ((a[i] - res) / d) % new_mult * mult % new_mult;
  mult = new_mult;
  ((res %= mult) += mult) %= mult;
}
return res;
}</pre>
```

#### 6.23 Theorem

# 6.23.1 Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i), \, L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 6.23.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniform randomly) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 6.23.3 Cayley's Formula

- Given a degree sequence  $d_1,d_2,\ldots,d_n$  for each labeled vertices, there're  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1, 2, \ldots, k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .

#### 6.23.4 Erdős-Gallai theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1 + d_2 + \ldots + d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \le k \le n$ .

#### 6.24 Primes

# 7 Dynamic Programming

# 7.1 Dynamic Convex Hull

```
struct Line {
  mutable int64_t a, b, p;
bool operator<(const Line &rhs) const { return a < rhs.a; }</pre>
   bool operator<(int64_t x) const { return p < x; }</pre>
 struct DynamicHull : multiset<Line, less<>>> {
   static const int64_t kInf = 1e18;
   int64_t Div(int64_t a, int64_t b) { return a / b - ((a \land b) <
       0 && a % b); }
   bool Isect(iterator x, iterator y) {
     if (y == end()) { x->p = kInf; return false; }
     if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
     else x->p = Div(y->b - x->b, x->a - y->a);
     return x->p >= y->p;
  }
   void Insert(int64_t a, int64_t b) {
     auto z = insert(\{a, b, 0\}), y = z++, x = y;
     while (Isect(y, z)) z = erase(z);
     if (x != begin() && Isect(--x, y)) Isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p) Isect(x,
      erase(y));
   int64_t Query(int64_t x) {
     auto l = *lower_bound(x);
     return l.a * x + l.b;
|};
```

# 7.2 1D/1D Convex Optimization

```
struct segment {
   int i, l, r;
   segment() {}
   segment(int a, int b, int c): i(a), l(b), r(c) {}
 inline long long f(int l, int r) {
  return dp[l] + w(l + 1, r);
 void solve() {
   dp[0] = 011;
   deque<segment> deq; deq.push_back(segment(0, 1, n));
   for (int i = 1; i \le n; ++i) {
     dp[i] = f(deq.front().i, i);
     while (deq.size() && deq.front().r < i + 1) deq.pop_front()</pre>
     deq.front().l = i + 1;
     segment seg = segment(i, i + 1, n);
     while (deq.size() && f(i, deq.back().1) < f(deq.back().i,</pre>
     deq.back().1)) deq.pop_back();
     if (dea.size()) {
       int d = 1048576, c = deq.back().1;
       while (d >>= 1) if (c + d <= deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c += d;
       deq.back().r = c; seg.l = c + 1;
     if (seg.l <= n) deq.push_back(seg);</pre>
1}
```

#### 7.3 Condition

#### 7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

#### 7.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \ B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

# 8 Geometry

#### 8.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps; }</pre>
struct P {
  double x, y;
  P() : x(0), y(0) {}
  P(double x, double y): x(x), y(y) {}

P operator + (P b) { return P(x + b.x, y + b.y); }

P operator - (P b) { return P(x - b.x, y - b.y); }

P operator * (double b) { return P(x * b, y * b); }
  P operator / (double b) { return P(x / b, y / b); } double operator * (P b) { return x * b.x + y * b.y; } double operator ^ (P b) { return x * b.y - y * b.x; }
  double abs() { return hypot(x, y); }
  P unit() { return *this / abs(); }
  P spin(double o) {
    double c = cos(o), s = sin(o);
     return P(c * x - s * y, s * x + c * y);
  double angle() { return atan2(y, x); }
};
struct L {
  // ax + by + c = 0
  double a, b, c, o;
  P pa, pb;
  L(): a(0), b(0), c(0), o(0), pa(), pb() {}
  L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa \land pb), o
      (atan2(-a, b)), pa(pa), pb(pb) {}
  P project(P p) { return pa + (pb - pa).unit() * ((pb - pa) *
      (p - pa) / (pb - pa).abs()); }
  P reflect(P p) { return p + (project(p) - p) * 2; }
  double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
     pa).abs() * (pb - pa).abs()); }
}:
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
  if (max(p1.x, p2.x) < min(p3.x, p4.x) | | max(p3.x, p4.x) < |
     min(p1.x, p2.x)) return false
  if (max(p1.y, p2.y) < min(p3.y, p4.y) || max(p3.y, p4.y) <
     min(p1.y, p2.y)) return false
  return sign((p3 - p1) ^ (p4 - p1)) * sign((p3 - p2) ^ (p4 -
     p2)) <= 0 &&
       sign((p1 - p3) \land (p2 - p3)) * sign((p1 - p4) \land (p2 - p4))
}
bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b, x.a *
     y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
```

# 8.2 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[
    maxn];
point p[maxn];
```

```
int build(int 1, int r, int dep = 0) {
  if (l == r) return -1;
   function<br/><br/>bool(const point &, const point &)> f = [dep](const
      point &a, const point &b) {
     if (dep & 1) return a.x < b.x;</pre>
     else return a.y < b.y;</pre>
  int m = (l + r) >> 1;
nth_element(p + l, p + m, p + r, f);
   xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
   lc[m] = build(l, m, dep + 1);
  if (\simlc[m]) {
     xl[m] = min(xl[m], xl[lc[m]]);
     xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
     xl[m] = min(xl[m], xl[rc[m]]);
     xr[m] = max(xr[m], xr[rc[m]]);
     yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  }
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
   if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
  q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
       (a.y - b.y) * 111 * (a.y - b.y);
void dfs(const point &q, long long &d, int o, int dep = 0) {
   if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
   if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x || !(dep & 1) && q.y < p[o].y)</pre>
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
     if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
}
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
return res;
j }}
```

# 8.3 Delaunay Triangulation

```
namespace triangulation {
static const int maxn = 1e5 + 5;
vector<point> p;
set<int> g[maxn];
int o[maxn];
set<int> s;
void add_edge(int x, int y) {
  s.insert(x), s.insert(y);
  g[x].insert(y);
  g[y].insert(x);
bool inside(point a, point b, point c, point p) {
  if (((b - a) ^ (c - a)) < 0) swap(b, c);</pre>
  function<long(int)> sqr = [](int x) { return x * 1ll * x}
     ; };
  long long k11 = a.x - p.x, k12 = a.y - p.y, k13 = sqr(a.x) -
     sqr(p.x) + sqr(a.y) - sqr(p.y);
  long long k21 = b.x - p.x, k22 = b.y - p.y, k23 = sqr(b.x) - b.y
     sqr(p.x) + sqr(b.y) - sqr(p.y);
  long long k31 = c.x - p.x, k32 = c.y - p.y, k33 = sqr(c.x) - c.y
  sqr(p.x) + sqr(c.y) - sqr(p.y);
long long det = k11 * (k22 * k33 - k23 * k32) - k12 * (k21 *
     k33 - k23 * k31) + k13 * (k21 * k32 - k22 * k31);
```

```
return det > 0:
bool intersect(const point &a, const point &b, const point &c,
     const point &d) {
  return ((b - a) \land (c - a)) * ((b - a) \land (d - a)) < 0 &&
      ((d - c) \wedge (a - c)) * ((d - c) \wedge (b - c)) < 0;
void dfs(int 1, int r) {
  if (r - 1 \le 3) {
    for (int i = 1; i < r; ++i) {
      for (int j = i + 1; j < r; ++j) add_edge(i, j);
    return:
  }
  int m = (l + r) >> 1;
  dfs(l, m), dfs(m, r);
  int pl = l, pr = r - 1;
  while (true) {
    int z = -1;
    for (int u : g[pl]) {
      long long c = ((p[pl] - p[pr]) \wedge (p[u] - p[pr]));
      if (c > 0 \mid | c == 0 \& abs(p[u] - p[pr]) < abs(p[pl] - p[
    pr])) {
    z = u;
        break;
      }
    if (z != -1) {
      pl = z;
      continue;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) \wedge (p[u] - p[pl]));
      if (c < 0 \mid l \mid c == 0 \& abs(p[u] - p[pl]) < abs(p[pr] - p[
     pl])) {
        z = u:
        break;
      }
    if (z != -1) {
      pr = z;
      continue;
    break;
  }
  add_edge(pl, pr);
  while (true) {
    int z = -1;
    bool b = false;
    for (int u : g[pl]) {
      long long c = ((p[pl] - p[pr]) \wedge (p[u] - p[pr]));
      if (c < 0 \& (z == -1 || inside(p[pl], p[pr], p[z], p[u])
     )) z = u;
    for (int u : g[pr]) {
      long long c = ((p[pr] - p[pl]) \land (p[u] - p[pl]));
      if (c > 0 && (z == -1 \mid \mid inside(p[pl], p[pr], p[z], p[u])
     )) z = u, b = true;
    if (z == -1) break;
    int x = pl, y = pr;
    if (b) swap(x, y);
    for (auto it = g[x].begin(); it != g[x].end(); ) {
      int u = *it;
      if (intersect(p[x], p[u], p[y], p[z])) {
        it = g[x].erase(it);
        g[u].erase(x);
      } else {
        ++it;
    if (b) add_edge(pl, z), pr = z;
    else add_edge(pr, z), pl = z;
vector<vector<int>> solve(vector<point> v) {
  int n = v.size();
  for (int i = 0; i < n; ++i) g[i].clear();</pre>
  for (int i = 0; i < n; ++i) o[i] = i;
  sort(o, o + n, [&](int i, int j) { return v[i] < v[j]; });</pre>
  p.resize(n);
  for (int i = 0; i < n; ++i) p[i] = v[o[i]];
  dfs(0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i) {
    for (int j : g[i]) res[o[i]].push_back(o[j]);
```

#### 8.4 Sector Area

```
// calc area of sector which include a, b
| double SectorArea(P a, P b, double r) {
| double o = atan2(a.y, a.x) - atan2(b.y, b.x);
| while (o <= 0) o += 2 * pi;
| while (o >= 2 * pi) o -= 2 * pi;
| o = min(o, 2 * pi - o);
| return r * r * o / 2;
| }
```

# 8.5 Half Plane Intersection

```
bool jizz(L l1,L l2,L l3){
  P p=intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
bool cmp(const L &a,const L &b){
  return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;</pre>
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
   vector<L> pls(1,ls[0]);
  for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls.back().</pre>
      o))pls.push_back(ls[i]);
   deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],pls[b],
     pls[c]))
  for(int i=2;i<(int)pls.size();++i){</pre>
    meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
    meow(i,dq[0],dq[1])dq.pop_front();
    dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back();
  meow(dq.back(),dq[0],dq[1])dq.pop_front();
  if(dq.size()<3u)return vector<P>(); // no solution or
     solution is not a convex
   vector<P> rt;
  for(int i=0;i<(int)dq.size();++i)rt.push_back(intersect(pls[</pre>
     dq[i]],pls[dq[(i+1)%dq.size()]]));
   return rt;
| }
```

#### 8.6 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
 int n=int(ps.size());
 vector<int> id(n),pos(n);
 vector<pair<int,int>> line(n*(n-1)/2);
  int m=-1;
  for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=make_pair</pre>
    (i,j); ++m;
  sort(line.begin(),line.end(),[&](const pair<int,int> &a,const
     pair<int, int> &b)->bool{
    if(ps[a.first].first==ps[a.second].first)return 0;
    if(ps[b.first].first==ps[b.second].first)return 1;
   return (double)(ps[a.first].second-ps[a.second].second)/(ps
    [a.first].first-ps[a.second].first) < (double)(ps[b.first</pre>
    ].second-ps[b.second].second)/(ps[b.first].first-ps[b.
    second].first);
 });
 for(int i=0;i<n;++i)id[i]=i;</pre>
  sort(id.begin(),id.end(),[&](const int &a,const int &b){
    return ps[a]<ps[b]; });</pre>
  for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
  for(int i=0;i<m;++i){</pre>
   auto l=line[i];
   tie(pos[l.first],pos[l.second],id[pos[l.first]],id[pos[l.
    second]])=make_tuple(pos[l.second],pos[l.first],l.second,l
     .first);
```

# 8.7 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay)) / (
     sin(a1) * cos(a2) - sin(a2) * cos(a1));
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
     TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);

res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

# 8.8 Polygon Center

```
| Point BaryCenter(vector<Point> &p, int n) {
    Point res(0, 0);
    double s = 0.0, t;
    for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
    }
    res.x /= (3 * s);
    res.y /= (3 * s);
    return res;
}
```

# 8.9 Maximum Triangle

# 8.10 Point in Polygon

```
int pip(vector<P> ps, P p) {
  int c = 0;
  for (int i = 0; i < ps.size(); ++i) {
    int a = i, b = (i + 1) % ps.size();
    L l(ps[a], ps[b]);
    P q = l.project(p);
    if ((p - q).abs() < eps && l.inside(q)) return 1;
    if (same(ps[a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
    if (ps[a].y > ps[b].y) swap(a, b);
    if (ps[a].y <= p.y && p.y < ps[b].y && p.x <= ps[a].x + (ps
        [b].x - ps[a].x) / (ps[b].y - ps[a].y) * (p.y - ps[a].y))
        ++c;
    }
    return (c & 1) * 2;
}</pre>
```

#### 8.11 Circle

```
struct C {
  P c;
  double r;
  C(P \ c = P(0, 0), double \ r = 0) : c(c), r(r) \{\}
vector<P> Intersect(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  vector<P> p;
  if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c).unit()
     * a.r);
  else if (a.r + b.r > d \&\& d + a.r >= b.r) {
    double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
    P i = (b.c - a.c).unit();
    p.push_back(a.c + i.spin(o) * a.r);
    p.push_back(a.c + i.spin(-o) * a.r);
  return p;
double IntersectArea(C a, C b) {
  if (a.r > b.r) swap(a, b);
  double d = (a.c - b.c).abs();
  if (d \ge a.r + b.r - eps) return 0;
  if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
double q = acos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
  return p * sq(a.r) + q * sq(b.r) - a.r * d * <math>sin(p);
// remove second level if to get points for line (defalut:
     segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
  double C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
  vector<P> t;
  if (d >= -eps) {
    d = \max(0., d);
    double i = (-B - sqrt(d)) / (2 * A);
double j = (-B + sqrt(d)) / (2 * A);
    if (i - 1.0 \le eps \&\& i \ge -eps) t.emplace_back(a.x + i * x)
      a.y + i * y);
    if (j - 1.0 \le eps \&\& j \ge -eps) t.emplace_back(a.x + j * x)
     , a.y + j * y);
  }
  return t:
// calc area intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
  auto p = CircleCrossLine(a, b, P(0, 0), r);
  if (ina) {
    if (inb) return abs(a ^ b) / 2;
    return SectorArea(b, p[0], r) + abs(a \land p[0]) / 2;
  if (inb) return SectorArea(p[0], a, r) + abs(p[0] ^b / 2;
  if (p.size() == 2u) return SectorArea(a, p[0], r) +
     SectorArea(p[1], b, r) + abs(p[0] \land p[1]) / 2;
  else return SectorArea(a, b, r);
// for any trianale
double AreaOfCircleTriangle(vector<P> ps, double r) {
```

```
double ans = 0;
for (int i = 0; i < 3; ++i) {
   int j = (i + 1) % 3;
   double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y, ps[j].x
   );
   if (o >= pi) o = o - 2 * pi;
   if (o <= -pi) o = o + 2 * pi;
   ans += AreaOfCircleTriangle(ps[i], ps[j], r) * (o >= 0 ? 1
        : -1);
}
return abs(ans);
}
```

# 8.12 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
 #define Pij \
   P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
   z.emplace_back(a.c + i, a.c + i + j);
 #define deo(I,J) \
   double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos(e / d)
  P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-o);\
z.emplace_back(a.c + j * a.r, b.c J j * b.r);\
z.emplace_back(a.c + k * a.r, b.c J k * b.r);
   if (a.r < b.r) swap(a, b);
   vector<L> z;
   if ((a.c - b.c).abs() + b.r < a.r) return z;</pre>
   else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
   else {
     deo(-,+);
     if (same(d, a.r + b.r)) { Pij; }
     else if (d > a.r + b.r) \{ deo(+,-); \}
   return z;
}
vector<L> tangent(C c, P p) {
   vector<L> z;
double d = (p - c.c).abs();
   if (same(d, c.r)) {
     P i = (p - c.c).spin(pi / 2);
     z.emplace_back(p, p + i);
   } else if (d > c.r) {
     double o = acos(c.r / d);
     P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.spin(-o) * c.r;
     z.emplace_back(c.c + j, p);
     z.emplace_back(c.c + k, p);
   return z;
}
```

#### 8.13 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
  vector<pair<double, double>> res;
if (same(a.r + b.r, d));
  else if (d \le abs(a.r - b.r) + eps) {
     if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
  } else if (d < abs(a.r + b.r) - eps) {
     double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
     ), z = (b.c - a.c).angle();
     if (z < 0) z += 2 * pi;
double l = z - o, r = z + o;
     if (l < 0) l += 2 * pi;
     if (r > 2 * pi) r -= 2 * pi;
     if (l > r) res.emplace_back(l, 2 * pi), res.emplace_back(0,
      r);
    else res.emplace_back(l, r);
  return res;
double CircleUnionArea(vector<C> c) { // circle should be
     identical
   int n = c.size();
  double a = 0, w;
  for (int i = 0; w = 0, i < n; ++i) {
     vector<pair<double, double>> s = {{2 * pi, 9}}, z;
     for (int j = 0; j < n; ++j) if (i != j) {
       z = CoverSegment(c[i], c[j]);
       for (auto &e : z) s.push_back(e);
```

```
}
sort(s.begin(), s.end());
auto F = [&] (double t) { return c[i].r * (c[i].r * t + c[i].c.x * sin(t) - c[i].c.y * cos(t)); };
for (auto &e : s) {
    if (e.first > w) a += F(e.first) - F(w);
    w = max(w, e.second);
}
return a * 0.5;
}
```

# 8.14 Minimun Distance of 2 Polygons

# 8.15 2D Convex Hull

```
bool operator < (const P &a, const P &b) { return same(a.x, b.x
     ) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(a.x, b.x
     ) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) \wedge (c - a))
vector<P> convex(vector<P> ps) {
  vector<P> p;
  sort(ps.begin(), ps.end(), [&] (P a, P b) { return same(a.x,
     b.x) ? a.y < b.y : a.x < b.x; });
  for (int i = 0; i < ps.size(); ++i)</pre>
    while (p.size() \ge 2 \& crx(p[p.size() - 2], ps[i], p[p.
     size() - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  int t = p.size();
  for (int i = (int)ps.size() - 2; i >= 0; --i) {
    while (p.size() > t && crx(p[p.size() - 2], ps[i], p[p.size
     () - 1]) >= 0) p.pop_back();
    p.push_back(ps[i]);
  p.pop_back();
  return p;
int sgn(double x) { return same(x, \emptyset) ? \emptyset : x > \emptyset ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
  double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
  int n;
  vector<P> p, u, d;
  CH() {}
  CH(vector<P> ps) : p(ps) {
    n = ps.size();
    rotate(p.begin(), min_element(p.begin(), p.end()), p.end())
    auto t = max_element(p.begin(), p.end());
    d = vector<P>(p.begin(), next(t));
    u = \text{vector} < P > (t, p.end()); u.push_back(p[0]);
  int find(vector<P> &v, P d) {
```

```
int l = 0, r = v.size();
     while (l + 5 < r) {
  int L = (l * 2 + r) / 3, R = (l + r * 2) / 3;
       if (v[L] * d > v[R] * d) r = R;
       else l = L;
     int x = 1;
     for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x] * d) x
      = i;
     return x:
   int findFarest(P v) {
     if (v.y > 0 \mid | v.y == 0 \& v.x > 0) return ((int)d.size() -
      1 + find(u, v)) % p.size();
     return find(d, v);
   P get(int 1, int r, P a, P b) {
     int s = sgn(crx(a, b, p[l % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(a, b, p[m % n])) == s) l = m;
       else r = m;
     return isLL(a, b, p[l % n], p[(l + 1) % n]);
   vector<P> getIS(P a, P b) {
     int X = findFarest((b - a).spin(pi / 2));
int Y = findFarest((a - b).spin(pi / 2));
     if (X > Y) swap(X, Y);
     if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0) return
{get(X, Y, a, b), get(Y, X + n, a, b)};</pre>
     return {};
   void update_tangent(P q, int i, int &a, int &b) {
     if (sgn(crx(q, p[a], p[i])) > 0) a = i;
     if (sgn(crx(q, p[b], p[i])) < 0) b = i;
   void bs(int l, int r, P q, int &a, int &b) {
     if (l == r) return;
     update_tangent(q, 1 % n, a, b);
     int s = sgn(crx(q, p[l % n], p[(l + 1) % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l = m;
     update_tangent(q, r % n, a, b);
   bool contain(P p) {
     if (p.x < d[0].x \mid | p.x > d.back().x) return 0;
     auto it = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
     if (it->x == p.x) {
       if (it->y > p.y) return 0;
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     it = lower_bound(u.begin(), u.end(), P(p.x, 1e12), greater<
      P>());
     if (it->x == p.x) {
       if (it->y < p.y) return 0;</pre>
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
   bool get_tangent(P p, int &a, int &b) { // b -> a
     if (contain(p)) return 0;
     int i = lower_bound(d.begin(), d.end(), p) - d.begin();
     bs(0, i, p, a, b);
bs(i, d.size(), p, a, b);
     i = lower_bound(u.begin(), u.end(), p, greater<P>()) - u.
      begin();
     bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a, b);
     bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.size(), p,
     a, b);
return 1;
  }
};
```

#### 8.16 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d){
  return abs(((b-a)^(c-a))*(d-a))/6;
}
struct convex3D{
  static const int maxn=1010;
```

```
struct T{
      int a,b,c;
bool res;
       T(){}
      T(int a,int b,int c,bool res=1):a(a),b(b),c(c),res(res){}
  int n,m;
  P p[maxn];
  T f[maxn*8];
  int id[maxn][maxn];
  bool on(T &t,P &q){
      return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
  void meow(int q,int a,int b){
       int f2=id[a][b];
       if(f[f2].res){
            if(on(f[f2],p[q]))dfs(q,f2);
           else{
                id[q][b]=id[a][q]=id[b][a]=m;
                 f[m++]=T(b,a,q,1);
      }
  }
  void dfs(int p,int i){
      f[i].res=0;
       meow(p,f[i].b,f[now].a);
      meow(p,f[i].c,f[now].b);
       meow(p,f[i].a,f[now].c);
  void operator()(){
       if(n<4)return
       if([&]()->int{
            for(int i=1; i< n; ++i)if(abs(p[0]-p[i])>eps)return swap(p[1],
             p[i]),0;
            return 1;
       }())return;
       if([&]()->int{
            for(int i=2; i< n; ++i)if(abs((p[0]-p[i])^(p[1]-p[i]))>eps)
             return swap(p[2],p[i]),0;
            return 1;
       }())return;
       if([&]()->int{
            for(int i=3;i<n;++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))*(p[i]-</pre>
             p[0]))>eps)return swap(p[3],p[i]),0;
               eturn 1;
       }())return;
       for(int i=0;i<4;++i){</pre>
           T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
            if(on(t,p[i]))swap(t.b,t.c)
            id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
            f[m++]=t;
       for(int i=4;i< n;++i) for(int j=0;j< m;++j) if(f[j].res \&\& on(f[j]) for(int j=0;j< m;++j) for(int j=0;j< m;+
             ],p[i])){
            dfs(i,j);
           break:
       int mm=m; m=0;
       for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
  bool same(int i,int j){
       return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].a])>eps
              || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].b])>eps ||
              absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
  }
  int faces(){
       for(int i=0;i<m;++i){</pre>
            int iden=1;
            for(int j=0;j<i;++j)if(same(i,j))iden=0;</pre>
           r+=iden;
       return r;
|} tb;
```

# 8.17 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 ^ p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
```

```
circle min_enclosing(vector<pt> &p) {
   random_shuffle(p.begin(), p.end());
   double r = 0.0:
   pt cent;
   for (int i = 0; i < p.size(); ++i) {
  if (norm2(cent - p[i]) <= r) continue;</pre>
     r = 0.0:
     for (int j = 0; j < i; ++j) {
        if (norm2(cent - p[j]) <= r) continue;</pre>
       cent = (p[i] + p[j]) / 2;
        r = norm2(p[j] - cent);
        for (int k = 0; k < j; ++k) {
          if (norm2(cent - p[k]) <= r) continue;</pre>
          cent = center(p[i], p[j], p[k]);
          r = norm2(p[k] - cent);
       }
     }
   }
   return circle(cent, sqrt(r));
1 }
```

#### 8.18 Closest Pair

```
double closest_pair(int l, int r) {
   // p should be sorted increasingly according to the x-
      coordinates.
   if (l == r) return 1e9;
   if (r - l == 1) return dist(p[l], p[r]);
   int m = (l + r) >> 1;
   double d = min(closest_pair(l, m), closest_pair(m + 1, r));
   vector<int> vec;
for (int i = m; i >= l && fabs(p[m].x - p[i].x) < d; --i) vec</pre>
      .push_back(i);
   for (int i = m + 1; i \le r \& fabs(p[m].x - p[i].x) < d; ++i)
      vec.push_back(i);
   sort(vec.begin(), vec.end(), [&](int a, int b) { return p[a].
     y < p[b].y; \});
   for (int i = 0; i < vec.size(); ++i) {</pre>
     for (int j = i + 1; j < vec.size() && fabs(p[vec[j]].y - p[
      vec[i]].y) < d; ++j) {
       d = min(d, dist(p[vec[i]], p[vec[j]]));
    }
   return d;
1}
```

# 9 Miscellaneous

# 9.1 Bitwise Hack

```
| long long next_perm(long long v) {
| long long t = v | (v - 1);
| return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1))
| ;
| }
| void subset(long long s) {
| long long sub = s;
| while (sub) sub = (sub - 1) & s;
| }
```

# 9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 111 * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) {
            x = s - 1 - x;
            y = s - 1 - y;
        }
        swap(x, y);
    }
}
```

return res:

```
9.3
        Java
import java.io.*;
import java.util.*;
import java.lang.*;
import java.math.*;
public class filename{
  static Scanner in = new Scanner(System.in);
   public static void main(String[] args) throws Exception {
     Scanner fin = new Scanner(new File("infile"));
PrintWriter fout = new PrintWriter("outfile", "UTF-8");
     fout.println(fin.nextLine());
     fout.close():
     while (in.hasNext()) {
       String str = in.nextLine(); // getline
       String stu = in.next(); // string
     System.out.println("Case #" + t);
     System.out.printf("%d\n", 7122);
     int[][] d = \{\{7,1,2,2\},\{8,7\}\};
     int g = Integer.parseInt("-123");
     long f = (long)d[0][2];
     List<Integer> l = new ArrayList<>();
     Random rg = new Random();
     for (int i = 9; i >= 0; --i) {
       l.add(Integer.valueOf(rg.nextInt(100) + 1));
       l.add(Integer.valueOf((int)(Math.random() * 100) + 1));
     Collections.sort(l, new Comparator<Integer>() {
       public int compare(Integer a, Integer b) { return a - b;
     });
     for (int i = 0; i < l.size(); ++i)</pre>
       System.out.print(l.get(i));
     Set<String> s = new HashSet<String>(); // TreeSet
s.add("jizz");
     System.out.println(s);
     System.out.println(s.contains("jizz"));
    Map<String, Integer> m = new HashMap<String, Integer>();
m.put("lol", 7122);
     System.out.println(m);
     for(String key: m.keySet())
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
     System.out.println(Math.acos(-1));
     BigInteger bi = in.nextBigInteger(), bj = new BigInteger("
      -7122"), bk = BigInteger.value0f(17171);
     int sgn = bi.signum(); // sign(bi)
     bi = bi.subtract(BigInteger.ONE).multiply(bj).divide(bj).
     and(bj).gcd(bj).max(bj).pow(87);
     int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
|}
```

# 9.4 Dancing Links

```
| namespace dlx {
| int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn], rw[maxn],
| bt[maxn], s[maxn], head, sz, ans;
| void init(int c) {
| for (int i = 0; i < c; ++i) {
| up[i] = dn[i] = bt[i] = i;
| lt[i] = i == 0 ? c : i - 1;
| rg[i] = i == c - 1 ? c : i + 1;
| s[i] = 0;
| }
| rg[c] = 0, lt[c] = c - 1;</pre>
```

```
up[c] = dn[c] = -1;
  head = c, sz = c + 1;
}
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
   int f = sz;
  for (int i = 0; i < (int)col.size(); ++i) {</pre>
     int c = col[i], v = sz++;
     dn[bt[c]] = v;
     up[v] = bt[c], bt[c] = v;
     rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r, cl[v] = c;
     ++s[c];
    if (i > 0) lt[v] = v - 1;
  lt[f] = sz - 1;
void remove(int c) {
  lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
   for (int i = dn[c]; i != c; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j]) {
  up[dn[j]] = up[j], dn[up[j]] = dn[j];
       --s[cl[j]];
  }
fvoid restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
    for (int j = lt[i]; j != i; j = lt[j]) {
}
       ++s[cl[j]];
       up[dn[j]] = j, dn[up[j]] = j;
  lt[rg[c]] = c, rg[lt[c]] = c;
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i)
     dn[bt[i]] = i, up[i] = bt[i];
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
   int w = c;
  for (int x = c; x != head; x = rg[x]) {
     if (s[x] < s[w]) w = x;
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
     for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j]);
     dfs(dep + 1);
     for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j]);
  }
  restore(w);
}
int solve() {
 ans = 1e9, dfs(0);
 return ans;
```

#### 9.5 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second =
     weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i such
     that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int> &x,
     vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
    if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  }):
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(st[qr[i].first], ed[qr</pre>
     [i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
```

```
djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[x[i]],</pre>
    ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
    }
  djs.undo();
}
void solve(int l, int r, vector<int> v, long long c) {
 if (l == r) {
  cost[qr[l].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first]) {
      printf("%lld\n", c);
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(minv,
    cost[v[i]]);
    printf("%lld\n", c + minv);
    return;
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first);
 contract(l, m, lv, x, y);
long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
 solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

#### 9.6 Manhattan Distance MST

```
void solve(int n) {
       init();
         vector<int> v(n), ds;
        for (int i = 0; i < n; ++i) {
               v[i] = i;
                ds.push_back(x[i] - y[i]);
        sort(ds.begin(), ds.end());
        ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
         sort(v.begin(), v.end(), [\&](int i, int j) { return x[i] == x}
                   [j] ? y[i] > y[j] : x[i] > x[j]; y[j] : x[i] > x[j]; y[j] : y[j] : y[i] > y[j] : y[i] > y[i]
         int j = 0;
        for (int i = 0; i < n; ++i) {
                int p = lower_bound(ds.begin(), ds.end(), x[v[i]] - y[v[i]]) - ds.begin() + 1;
                pair<int, int> q = query(p);
                 // query return prefix minimum
                 if (~q.second) add_edge(v[i], q.second);
                add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
       }
```

```
| }
| void make_graph() {
| solve(n);
| for (int i = 0; i < n; ++i) swap(x[i], y[i]);
| solve(n);
| for (int i = 0; i < n; ++i) x[i] = -x[i];
| solve(n);
| for (int i = 0; i < n; ++i) swap(x[i], y[i]);
| solve(n);
| }</pre>
```

# 9.7 IOI 2016 Alien trick

```
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative as well, depending on the problem.
    if (c - (1LL << d) < 0) continue;
    long long ck = c - (1LL << d);
    pair<long long, int> r = check(ck);
    if (r.second == k) return r.first - ck * k;
    if (r.second < k) c = ck;
  }
  pair<long long, int> r = check(c);
  return r.first - c * k;
}
```