19

24 24

Contents	8 Geometry 1 8.1 Basic 1 8.2 KD Tree 1
1 Basic 1.1 vimrc 1.2 Fast Integer Input 1.3 Increase stack size 1.4 Pragma optimization 2 Flow 2.1 Dinic's Algorithm 2.2 Minimum-cost flow 2.3 Gomory-Hu Tree 2.4 Stoer-Wagner Minimum Cut 2.5 Kuhn-Munkres Algorithm	8.3 Delaunay Triangulation 1 1 8.4 Sector Area 2 1 8.5 Half Plane Intersection 2 1 8.6 Rotating Sweep Line 2 1 8.7 Triangle Center 2 2 8.9 Maximum Triangle 2 8.10 Point in Polygon 2 2 8.11 Circle-Line Intersection 2 2 8.12 Circle-Triangle Intersection 2 2 8.13 Tangent from Point to Circle 2 2 8.14 Minimun Distance of 2 Polygons 2 2 8.15 2D Convex Hull 2 2 8.16 3D Convex Hull 2 3 8.17 Rotating Caliper 2
2.6 Flow Model	3 8.18 Minimum Enclosing Circle
3.1 Disjoint Set	3 9.1 Bitwise Hack 2 4 9.2 Hilbert's Curve (faster Mo's algorithm) 2 4 9.3 Java 2 9.4 Dancing Links 2 4 9.5 Offline Dynamic MST 2 9.6 Manhattan Distance MST 2
4.1 Link-Cut Tree 4.2 Heavy-Light Decomposition 4.3 Centroid Decomposition 4.4 Minimum Mean Cycle 4.5 Minimum Steiner Tree 4.6 Directed Minimum Spanning Tree	9.7 "Dynamic" Kth Element (parallel binary search)
4.7 Maximum Matching on General Graph	6 1 Basic 7 8 1.1 vimrc
4.11 Tarjan's Bridge 4.12 Dominator Tree 4.13 System of Difference Constraints	colo desert
5 String 5.1 Knuth-Morris-Pratt Algorithm 5.2 Z Algorithm 5.3 Manacher's Algorithm 5.4 Aho-Corasick Automaton	10 10
5.5 Suffix Automaton 5.6 Suffix Array	1.2 Past Integer Input
6 Math 6.1 Fast Fourier Transform 6.2 Number Theoretic Transform 6.2.1 NTT Prime List 6.3 Polynomial Division 6.4 Polynomial Square Root 6.5 Fast Walsh-Hadamard Transform	<pre>const int N = 4096; static char buffer[N]; static char *p = buffer, *end = buffer; if (p == end) { if ((end = buffer + fread(buffer, 1, N, stdin)) == buffer) return EOF;</pre>
6.5.1 XOR Convolution	14 return *p++; 14 }
6.6 Simplex Algorithm 6.6.1 Construction 6.7 Schreier-Sims Algorithm 6.8 Miller Rabin 6.9 Pollard's Rho 6.10 Meissel-Lehmer Algorithm 6.11 Discrete Logarithm 6.12 Gaussian Elimination	<pre>inline bool rit(T& x) { char c = 0; bool flag = false; while (c = getchar(), (c < '0' && c != '-') c > '9 ') if (c == -1) return false; c == '-' ? (flag = true, x = 0) : (x = c - '0'); while (c = getchar(), c >= '0' && c <= '9') x = x * 10 + c - '0';</pre>
$6.13 \text{ Linear Equations (full pivoting)} \\ 6.14 \mu \text{ function} \\ 6.15 \left\lfloor \frac{n}{i} \right\rfloor \text{ Enumeration} \\ 6.16 \text{ De Bruijn Sequence} \\ 6.17 \text{ Extended GCD} \\ 6.18 \text{ Euclidean Algorithms} \\ 6.19 \text{ Chinese Remainder Theorem}$	<pre>return true; return true; template <typename t,="" typenameargs=""> inline bool rit(T& x, Args&args) { return rit(x) && rit(args); }</typename></pre>
6.20 Theorem 6.20.1 Kirchhoff's Theorem 6.20.2 Tutte's Matrix 6.20.3 Cayley's Formula	$^{17}_{17}$ 1.3 Increase stack size
6.21 Primes	register long rsp asm("rsp"); char *p = (char*)malloc(size) + size, *bak = (char*)rsp
7.2 1D/1D Convex Optimization	18

1.4 Pragma optimization

2 Flow

2.1 Dinic's Algorithm

```
struct dinic {
  static const int inf = 1e9;
  struct edge {
     int dest, cap, rev;
     edge(int d, int c, int r): dest(d), cap(c), rev(r)
     {}
  vector<edge> g[maxn];
  int qu[maxn], ql, qr;
  int lev[maxn];
  void init() {
     for (int i = 0; i < maxn; ++i)
       g[i].clear();
  void add_edge(int a, int b, int c) {
  g[a].emplace_back(b, c, g[b].size() - 0);
  g[b].emplace_back(a, 0, g[a].size() - 1);
  bool bfs(int s, int t) {
  memset(lev, -1, sizeof(lev));
     lev[s] = 0;
    ql = qr = 0;

qu[qr++] = s;
     while (ql < qr) {</pre>
       int x = qu[ql++];
       for (edge &e : g[x]) if (lev[e.dest] == -1 && e.
     cap > 0) {
         lev[e.dest] = lev[x] + 1;
         qu[qr++] = e.dest;
     return lev[t] != -1;
  int dfs(int x, int t, int flow) {
     if (x == t) return flow;
     int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.dest]
     == lev[x] + 1) {
       int f = dfs(e.dest, t, min(e.cap, flow - res));
       res += f;
       e.cap -= f;
       g[e.dest][e.rev].cap += f;
     if (res == 0) lev[x] = -1;
    return res;
  int operator()(int s, int t) {
     int flow = 0;
     for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow;
};
```

2.2 Minimum-cost flow

```
struct mincost {
  struct edge {
    int dest, cap, w, rev;
    edge(int a, int b, int c, int d): dest(a), cap(b),
    w(c), rev(d) {}
  };
  vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
```

```
bool inq[maxn];
  void init() {
     for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
  void add_edge(int a, int b, int c, int d) {
    g[a].emplace\_back(b, c, +d, g[b].size() - 0);
     g[b].emplace_back(a, 0, -d, g[a].size() - 1);
  bool spfa(int s, int t, int &f, int &c) {
     for (int i = 0; i < maxn; ++i) {
       d[i] = inf;
p[i] = ed[i] = -1;
       inq[i] = false;
    d[s] = 0;
     queue<int> q;
     q.push(s);
     while (q.size()) {
       int x = q.front(); q.pop();
       inq[x] = false;
for (int i = 0; i < g[x].size(); ++i) {</pre>
         edge &e = g[x][i];
         if (e.cap > 0 && d[e.dest] > d[x] + e.w) {
  d[e.dest] = d[x] + e.w;
           p[e.dest] = x;
           ed[e.dest] = i;
           if (!inq[e.dest]) q.push(e.dest), inq[e.dest]
      = true;
         }
      }
     if (d[t] == inf) return false;
     int dlt = inf;
     for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[
    p[x]][ed[x]].cap);
for (int x = t; x != s; x = p[x]) {
       edge &e = g[p[x]][ed[x]];
       e.cap -= dlt;
       g[e.dest][e.rev].cap += dlt;
     f += dlt; c += d[t] * dlt;
    return true;
  pair<int, int> operator()(int s, int t) {
     int f = 0, c = 0;
    while (spfa(s, t, f, c));
return make_pair(f, c);
};
```

2.3 Gomory-Hu Tree

```
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;
  for(int i=2;i<=n;++i){
    int t=g[i];
    flow.reset(); // clear flows on all edge
    rt.push_back({i,t,flow(i,t)});
    flow.walk(i); // bfs points that connected to i (
    use edges not fully flow)
    for(int j=i+1;j<=n;++j){
        if(g[j]==t && flow.connect(j))g[j]=i; // check if
        i can reach j
    }
}
return rt;
}</pre>
```

2.4 Stoer-Wagner Minimum Cut

```
const int maxn = 500 + 5;
int w[maxn][maxn], g[maxn];
bool v[maxn], del[maxn];

void add_edge(int x, int y, int c) {
  w[x][y] += c;
```

```
w[y][x] += c;
pair<int, int> phase(int n) {
  memset(v, false, sizeof(v));
  memset(g, 0, sizeof(g));
  int s = -1, t = -1;
  while (true) {
     int c = -1;
     for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
       if (c = -1 | g[i] > g[c]) c = i;
     if (c == -1) break;
    v[c] = true;
     s = t, t = c;
    for (int i = 0; i < n; ++i) {
  if (del[i] || v[i]) continue;</pre>
       g[i] += w[c][i];
  return make_pair(s, t);
int mincut(int n) {
  int cut = 1e9;
  memset(del, false, sizeof(del));
  for (int i = 0; i < n - 1; ++i) {
    int s, t; tie(s, t) = phase(n);
del[t] = true;
     cut = min(cut, g[t]);
    for (int j = 0; j < n; ++j) {
    w[s][j] += w[t][j];
    w[j][s] += w[j][t];
  return cut;
```

2.5 Kuhn-Munkres Algorithm

```
int w[maxn][maxn], lx[maxn], ly[maxn];
int match[maxn], slack[maxn];
bool vx[maxn], vy[maxn];
bool dfs(int x) {
 vx[x] = true;
for (int i = 0; i < n; ++i) {</pre>
   if (vy[i]) continue;
   if (lx[x] + ly[i] > w[x][i]) {
     slack[i] = min(slack[i], lx[x] + ly[i] - w[x][i])
     continue:
   vy[i] = true;
   if (match[i] == -1 || dfs(match[i])) {
     match[i] = x;
     return true;
 return false;
int solve() {
 fill_n(match, n, -1);
 ][j]);
  for (int i = 0; i < n; ++i) {
   fill_n(slack, n, inf);
   while (true) {
      fill_n(vx, n, false);
     fill_n(vy, n, false);
if (dfs(i)) break;
     int dlt = inf;
      for (int j = 0; j < n; ++j) if (!vy[j]) dlt = min
    (dlt, slack[j]);
```

```
for (int j = 0; j < n; ++j) {
    if (vx[j]) lx[j] -= dlt;
    if (vy[j]) ly[j] += dlt;
    else slack[j] -= dlt;
    }
}
int res = 0;
for (int i = 0; i < n; ++i) res += w[match[i]][i];
return res;
}</pre>
```

2.6 Flow Model

- Maximum/Minimum flow with lower/upper bound from s to t
 - 1. Construct super source S and sink T
 - 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u l
 - 3. For each vertex v, denote in(v) as the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds
 - 4. If in(v)>0, connect $S\to v$ with capacity in(v), otherwise, connect $v\to T$ with capacity -in(v)
 - To maximize, connect $t \to s$ with capacity ∞ , and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, the maximum flow from s to t is the answer.
 - To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise, f' is the answer.
 - 5. The solution of each edge e is l_e+f_e , where f_e corresponds to the flow on the graph
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
 - 1. Redirect every edge $(y \to x \text{ if } (x,y) \in M, x \to y \text{ otherwise})$
 - 2. DFS from unmatched vertices in X
 - 3. $x \in X$ is chosen iff x is unvisited
 - 4. $y \in Y$ is chosen iff y is visited
- Minimum cost cyclic flow
 - 1. Consruct super source S and sink T
 - 2. For each edge (x,y,c), connect $x\to y$ with (cost,cap)=(c,1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1)
 - 3. For each edge with c<0, sum these cost as K, then increase d(y) by 1, decrease d(x) by 1
 - 4. For each vertex v with d(v)>0, connect $S\to v$ with (cost,cap)=(0,d(v))
 - 5. For each vertex v with d(v) < 0, connect $v \rightarrow T$ with (cost, cap) = (0, -d(v))
 - 6. Flow from S to T, the answer is the cost of the flow C + K
- Maximum density induced subgraph
 - 1. Binary search on answer, suppose we're checking answer T
 - 2. Construct a max flow model, let K be the sum of all weights
 - 3. Connect source $s \to v, v \in G$ with capacity K
 - 4. For each edge (u,v,w) in G, connect $u \to v$ and $v \to u$ with capacity w
 - 5. For $v \in G,$ connect it with sink $v \to t$ with capacity $K+2T-(\sum_{e \in E(v)} w(e)) 2w(v)$
 - 6. T is a valid answer if the maximum flow f < K|V|

3 Data Structure

3.1 Disjoint Set

```
struct DisjointSet {
  int p[maxn], sz[maxn], n, cc;
  vector<pair<int*, int>> his;
  vector<int> sh;
  void init(int _n) {
    n = _n; cc = n;
    for (int i = 0; i < n; ++i) sz[i] = 1, p[i] = i;
    sh.clear(); his.clear();
}</pre>
```

```
void assign(int *k, int v) {
    his.emplace_back(k, *k);
  void save() {
    sh.push_back((int)his.size());
  void undo() {
    int last = sh.back(); sh.pop_back();
    while (his.size() != last) {
  int *k, v;
       tie(k, v) = his.back(); his.pop_back();
    }
  int find(int x) {
    if (x == p[x]) return x;
    return find(p[x]);
  void merge(int x, int y) {
    x = find(x); y = find(y);
    if (x == y) return;
    if (sz[x] > sz[y]) swap(x, y);
assign(&sz[y], sz[x] + sz[y]);
    assign(&p[x], y);
    assign(\&cc, cc - 1);
} dsu;
```

3.2 < ext/pbds >

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>,
    rb_tree_tag, tree_order_statistics_node_update>
    tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
 // rb tree
  tree_set s
  s.insert(71); s.insert(22);
  assert(*s.find_by_order(0) == 22); assert(*s.
    find_by_order(1) == 71);
  assert(s.order_of_key(22) == 0); assert(s.
order_of_key(71) == 1);
  s.erase(22);
  assert(*s.find_by\_order(0) == 71); assert(s.
     order_of_key(71) == 0);
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;</pre>
  return 0;
```

3.3 Li Chao Tree

```
namespace lichao {
struct line {
  long long a, b;
  line(): a(0), b(0) {}
  line(long long a, long long b): a(a), b(b) {}
  long long operator()(int x) const { return a * x + b;
    }
};
line st[maxc * 4];
int sz, lc[maxc * 4], rc[maxc * 4];
```

```
int gnode() {
  st[sz] = line(1e9, 1e9);
  lc[sz] = -1, rc[sz] = -1;
  return sz++:
void init() {
  sz = 0;
}
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
  bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
  if (mcp) swap(st[o], tl);
  if (r - l == 1) return;
  if (lcp != mcp) {
     if (lc[o] == -1) lc[o] = gnode();
     add(l, (l + r) / 2, tl, lc[o]);
  } else {
     if (rc[o] == -1) rc[o] = gnode();
     add((1 + r) / 2, r, tl, rc[o]);
long long query(int l, int r, int x, int o) {
  if (r - l == 1) return st[o](x);
  if (x < (l + r) / 2) {</pre>
     if (lc[o] == -1) return st[o](x);
     return min(st[o](x), query(l, (l + r) / 2, x, lc[o
     ]));
     if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((1 + r) / 2, r, x, rc[o
}}
```

4 Graph

4.1 Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
  int sum, v, rev, id;
node(int s, int id): id(id), v(s), sum(s), rev(0), fa
  (nullptr), pfa(nullptr) {
     ch[0] = nullptr;
     ch[1] = nullptr;
  int relation() {
     return this == fa \rightarrow ch[0] ? 0 : 1;
  void push() {
     if (!rev) return;
     swap(ch[0], ch[1]);
if (ch[0]) ch[0]->rev ^= 1;
     if (ch[1]) ch[1]->rev ^= 1;
     rev = 0;
  void pull() {
     if (ch[0]) sum += ch[0]->sum;
     if (ch[1]) sum += ch[1]->sum;
  void rotate() {
     if (fa->fa) fa->fa->push();
     fa->push(), push();
swap(pfa, fa->pfa);
     int d = relation();
     node *t = fa;
     if (t->fa) t->fa->ch[t->relation()] = this;
     fa = t \rightarrow fa;
     t->ch[d] = ch[d \land 1];
     if (ch[d \land 1]) ch[d \land 1] -> fa = t;
     ch[d \land 1] = t;
     t->fa = this;
     t->pull(), pull();
  void splay() {
     while (fa) {
       if (!fa->fa) {
```

```
rotate():
        continue;
      fa->fa->push(), fa->push();
      if (relation() == fa->relation()) fa->rotate(),
    rotate();
      else rotate(), rotate();
  void evert() {
    access();
    splay();
    rev ^= 1;
  void expose() {
    splay(), push();
    if (ch[1]) {
      ch[1]->fa = nullptr;
      ch[1]-pfa = this;
      ch[1] = nullptr;
      pull();
 bool splice() {
    splay();
    if (!pfa) return false;
    pfa->expose();
    pfa->ch[1] = this;
    fa = pfa;
    pfa = nullptr;
    fa->pull();
    return true;
  void access() {
    expose();
    while (splice());
  int query() {
    return sum;
namespace lct {
node *sp[maxn];
void make(int u, int v) {
  // create node with id u and value v
  sp[u] = new node(v, u);
void link(int u, int v) {
  // u become v's parent
  sp[v]->evert();
  sp[v]->pfa = sp[u];
void cut(int u, int v) {
 // u was v's parent
sp[u]->evert();
  sp[v]->access(), sp[v]->splay(), sp[v]->push();
  sp[v]->ch[0]->fa = nullptr;
  sp[v]->ch[0] = nullptr;
  sp[v]->pull();
void modify(int u, int v) {
  sp[u]->splay();
  sp[u]->v = v
  sp[u]->pull();
int query(int u, int v) {
  sp[u]->evert(), sp[v]->access(), sp[v]->splay();
  return sp[v]->query();
int find(int u) {
 sp[u]->access();
  sp[u]->splay();
 node *p = sp[u];
  while (true) {
    p->push()
    if (p->ch[0]) p = p->ch[0];
    else break;
 }
  return p->id;
```

4.2 Heavy-Light Decomposition

```
void dfs(int x, int p) {
  dep[x] = \sim p ? dep[p] + 1 : dep[x];
  sz[x] = 1;
  to[x] = -1;
  fa[x] = p;
  for (const int &u : g[x]) {
    if (u == p) continue;
    dfs(u, x);
    sz[x] += sz[u];
    if (to[x] == -1 \mid | sz[to[x]] < sz[u]) to[x] = u;
}
void hld(int x, int t) {
  static int tk = 0;
  fr[x] = t;
  dfn[x] = tk++;
  if (!~to[x]) return;
  hld(to[x], t);
for (const int &u : g[x]) {
    if (u == fa[x] || u == to[x]) continue;
    hld(u, u);
  }
}
vector<pair<int, int>> get(int x, int y) {
  int fx = fr[x], fy = fr[y];
vector<pair<int, int>> res;
while (fx != fy) {
    if (dep[fx] < dep[fy]) {</pre>
       swap(fx, fy);
       swap(x, y);
    res.emplace_back(dfn[fx], dfn[x] + 1);
    x = fa[fx];
    fx = fr[x];
  res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn
     [y]) + 1);
  int lca = (dep[x] < dep[y] ? x : y);
  return res;
```

4.3 Centroid Decomposition

```
void get_center(int now) {
  v[now] = true; vtx.push_back(now);
  sz[now] = 1; mx[now] = 0;
for (int u : G[now]) if (!v[u]) {
    get_center(u)
    mx[now] = max(mx[now], sz[u]);
    sz[now] += sz[u];
}
void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
  v[now] = true;
  for (auto u : G[now]) if (!v[u.first]) {
    get_dis(u, d, len + u.second);
}
void dfs(int now, int fa, int d) {
  get_center(now);
  int c = -1;
  for (int i : vtx) {
    if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx</pre>
     .size() / 2) c = i;
    v[i] = false;
  get_dis(c, d, 0);
for (int i : vtx) v[i] = false;
  v[c] = true; vtx.clear();
  dep[c] = d; p[c] = fa;
  for (auto u : G[c]) if (u.first != fa && !v[u.first])
```

```
dfs(u.first, c, d + 1);
}
```

4.4 Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
long long d[1003][1003],dp[1003][1003];
pair<long long,long long> MMWC(){
  memset(dp,0x3f,sizeof(dp))
  for(int i=1;i<=n;++i)dp[0][i]=0;</pre>
  for(int i=1;i<=n;++i){</pre>
    for(int j=1;j<=n;++j){</pre>
      for(int k=1; k<=n; ++k){</pre>
        dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
    }
  long long au=1ll<<31,ad=1;</pre>
  for(int i=1;i<=n;++i){</pre>
    if(dp[n][i]==0x3f3f3f3f3f3f3f3f)continue;
    long long \bar{u}=0, d=1;
    for(int j=n-1; j>=0; -- j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
  u=dp[n][i]-dp[j][i];
         d=n-j;
    if(u*ad<au*d)au=u,ad=d;
  long long g=__gcd(au,ad);
  return make_pair(au/g,ad/g);
```

4.5 Minimum Steiner Tree

```
namespace steiner {
const int maxn = 64, maxk = 10;
const int inf = 1e9;
int w[maxn][maxn], dp[1 << maxk][maxn], off[maxn];</pre>
void init(int n) {
  for (int i = 0; i < n; ++i) {</pre>
     for (int j = 0; j < n; ++j) w[i][j] = inf;
    w[i][i] = 0;
  }
void add_edge(int x, int y, int d) {
  w[x][y] = min(w[x][y], d);
  w[y][x] = min(w[y][x], d);
int solve(int n, vector<int> mark) {
  for (int k = 0; k < n; ++k) {
  for (int i = 0; i < n; ++i) {
       for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j]
     ], w[i][k] + w[k][j]);
  int k = (int)mark.size();
  assert(k < maxk);</pre>
  for (int s = 0; s < (1 << k); ++s) {
  for (int i = 0; i < n; ++i) dp[s][i] = inf;</pre>
  for (int i = 0; i < n; ++i) dp[0][i] = 0;
for (int s = 1; s < (1 << k); ++s) {
     if (__builtin_popcount(s) == 1) {
       int x = _{-}
                   _builtin_ctz(s);
       for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]]
     ]][i];
       continue;
     for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1);
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s ^
     sub][i]);
       }
```

```
for (int i = 0; i < n; ++i) {
    off[i] = inf;
    for (int j = 0; j < n; ++j) off[i] = min(off[i],
    dp[s][j] + w[j][i]);
    }
    for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i
    ], off[i]);
}
int res = inf;
for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k
    ) - 1][i]);
return res;
}}</pre>
```

4.6 Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
  T g[maxn] [maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {
  for(int j = 0; j < maxn; ++j) g[i][j] = inf;
  vis[i] = inc[i] = false;</pre>
    }
  }
  void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
  T operator()(int root, int _n) {
    n = n
     if (dfs(root) != n) return -1;
    T ans = 0;
     while (true) {
       for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] =</pre>
       for (int i = 1; i <= n; ++i) if (!inc[i]) {</pre>
         for (int j = 1; j <= n; ++j) {
  if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
              fw[i] = g[j][i];
              fr[i] = j;
           }
         }
       int x = -1;
       for (int i = 1; i <= n; ++i) if (i != root &&!
    while (j != root && fr[j] != i && c <= n) ++c,
     j = fr[j];
         if (j == root || c > n) continue;
         else { x = i; break; }
       if (!~x) {
         for (int i = 1; i <= n; ++i) if (i != root &&!
     inc[i]) ans += fw[i];
         return ans;
       for (int i = 1; i <= n; ++i) vis[i] = false;
do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] =</pre>
     true; } while (y != x);
       inc[x] = false;
for (int k = 1; k <= n; ++k) if (vis[k])</pre>
          for (int j = 1; j <= n; ++j) if (!vis[j]) {
            if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
if (g[j][k] < inf && g[j][k] - fw[k] < g[j][x]</pre>
     ]) g[j][x] = g[j][k] - fw[k];
       }
    }
    return ans;
  int dfs(int now) {
     int r = 1;
     vis[now] = true;
     for (int i = 1; i \le n; ++i) if (g[now][i] < inf &&
      !vis[i]) r += dfs(i);
    return r;
```

```
|};
```

4.7 Maximum Matching on General Graph

```
namespace matching {
int fa[maxn], pre[maxn], match[maxn], s[maxn], v[maxn];
vector<int> g[maxn];
queue<int> q;
void init(int n) {
  for (int i = 0; i \le n; ++i) match[i] = pre[i] = n;
  for (int i = 0; i < n; ++i) g[i].clear();</pre>
void add_edge(int u, int v) {
  g[u].push_back(v);
  g[v].push_back(u);
int find(int u) {
  if (u == fa[u]) return u;
  return fa[u] = find(fa[u]);
int lca(int x, int y, int n) {
  static int tk = 0;
  x = find(x), y = find(y);
  for (; ; swap(x, y)) {
  if (x != n) {
      if (v[x] == tk) return x;
      v[x] = tk;
      x = find(pre[match[x]]);
  }
void blossom(int x, int y, int l) {
  while (find(x) != 1) {
    pre[x] = y
    y = match[x];
    if (s[y] == 1) {
      q.push(y);
      s[y] = 0;
    if (fa[x] == x) fa[x] = 1;
    if (fa[y] == y) fa[y] = 1;
    x = pre[y];
bool bfs(int r, int n) {
  for (int i = 0; i <= n; ++i) {
    fa[i] = i;
    s[i] = -1;
  while (!q.empty()) q.pop();
  q.push(r);
  s[r] = 0;
  while (!q.empty()) {
    int x = q.front(); q.pop();
    for (int u : g[x]) {
      if (s[u] = -1) {
        pre[u] = x;
        s[u] = 1;
        if (match[u] == n) {
          for (int a = u, b = x, last; b != n; a = last
    , b = pre[a])
            last = match[b], match[b] = a, match[a] = b
          return true;
        q.push(match[u]);
        s[match[u]] = 0;
      lse if (!s[u] && find(u) != find(x)) {
int l = lca(u, x, n);
blossom(x, u, l);
        blossom(u, x, l);
    }
  return false;
int solve(int n) {
  int res = 0;
  for (int x = 0; x < n; ++x) {
```

```
if (match[x] == n) res += bfs(x, n);
}
return res;
}}
```

4.8 Maximum Weighted Matching on General Graph

```
struct WeightGraph {
  static const int inf = INT_MAX;
  static const int maxn = 514;
  struct edge {
    int u, v, w;
    edge(){}
    edge(int u, int v, int w): u(u), v(v), w(w) {}
  int n, n_x;
  edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
  int match[maxn * 2], slack[maxn * 2], st[maxn * 2],
    pa[maxn * 2];
  int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[
    maxn * 2];
  vector<int> flo[maxn * 2];
  queue<int> q;
  int e_delta(const edge &e) {
    return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2;
  void update_slack(int u, int x) {
    if (!slack[x] | | e_delta(g[u][x]) < e_delta(g[slack])
    [x]][x])) slack[x] = u;
  void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; ++u)
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x \le n) q.push(x);
    else for (size_t i = 0; i < flo[x].size(); i++)</pre>
    q_push(flo[x][i]);
  void set_st(int x, int b) {
    st[x] = b;
    if(x > n) for (size_t i = 0; i < flo[x].size(); ++
    i) set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
  int pr = find(flo[b].begin(), flo[b].end(), xr) -
    flo[b].begin();
if (pr % 2 == 1) {
      reverse(flo[b].begin() + 1, flo[b].end());
      return (int)flo[b].size() - pr;
    return pr;
  }
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;
    edge e = g[u][v];
int xr = flo_from[u][e.u], pr = get_pr(u, xr)
    for (int i = 0; i < pr; ++i) set_match(flo[u][i],</pre>
    flo[u][i ^ 1]);
    set_match(xr, v);
rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].
    end());
  void augment(int u, int v) {
    for (; ; ) {
  int xnv = st[match[u]];
      set_match(u, v);
      if (!xnv) return
      set_match(xnv, st[pa[xnv]]);
      u = st[pa[xnv]], v = xnv;
  int get_lca(int u, int v) {
    static int t = 0;
```

```
for (++t; u || v; swap(u, v)) {
    if (u == 0) continue;
    if (vis[u] == t) return u;
    vis[u] = t;
    u = st[match[u]];
    if (u) u = st[pa[u]];
  return 0;
void add_blossom(int u, int lca, int v) {
  int b = n + 1;
  while (b <= n_x && st[b]) ++b;</pre>
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
for (int x = u, y; x != lca; x = st[pa[y]])
     flo[b].push_back(x), flo[b].push_back(y = st[
  match[x]]), q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end())
  for (int x = v, y; x != lca; x = st[pa[y]])
     flo[b].push_back(x), flo[b].push_back(y = st[
  match[x]]), q_push(y);
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].
  for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
     int xs = flo[b][i];
     for (int x = 1; x <= n_x; ++x)
       if (g[b][x].w == 0 \mid | e_delta(g[xs][x]) <
  e_{delta(g[b][x])}
         g[b][x] = g[xs][x], g[x][b] = g[x][xs];
     for (int x = 1; x \le n; ++x)
       if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)
  set_st(flo[b][i], flo[b][i]);</pre>
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b)
  for (int i = 0; i < pr; i += 2) {
  int xs = flo[b][i], xns = flo[b][i + 1];</pre>
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
    slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
pa[v] = e.u, S[v] = 1;
     int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  else\ if\ (S[v] == 0) {
     int lca = get_lca(u, v);
     if (!lca) return augment(u,v), augment(v,u), true
    else add_blossom(u, lca, v);
  }
  return false;
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0,
   q_push(x);
  if (q.empty()) return false;
  for (; ; ) {
```

```
while (q.size()) {
          int u = q.front(); q.pop();
          if (S[st[u]] == 1) continue;
          for (int v = 1; v <= n; ++v)
            if (g[u][v].w > 0 && st[u] != st[v]) {
              if (e_delta(g[u][v]) == 0) {
                if (on_found_edge(g[u][v])) return true;
              } else update_slack(u, st[v]);
       int d = inf;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b])
     / 2);
       for (int x = 1; x <= n_x; ++x)
          if (st[x] == x \&\& slack[x]) {
            if (S[x] == -1) d = min(d, e_delta(g[slack[x]]))
            else if (S[x] == 0) d = min(d, e_delta(g[
     slack[x]][x]) / \overline{2};
       for (int u = 1; u \ll n; ++u) {
         if (S[st[u]] == 0) {
   if (lab[u] <= d) return 0;</pre>
            lab[u] -= d;
         } else if (S[st[u]] == 1) lab[u] += d;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b) {
            if (S[st[b]] == 0) lab[b] += d * 2;
            else if (S[st[b]] == 1) lab[b] -= d * 2;
       q = queue<int>();
       for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x
      && e_delta(g[slack[x]][x]) == 0)
            if (on_found_edge(g[slack[x]][x])) return
       for (int b = n + 1; b <= n_x; ++b)
if (st[b] == b && S[b] == 1 && lab[b] == 0)
     expand_blossom(b);
     return false;
   pair<long long, int> solve() {
     memset(match + 1, 0, sizeof(int) * n);
     n x = n:
     int n_matches = 0;
     long long tot_weight = 0;
     for (int u = 0; u \le n; ++u) st[u] = u, flo[u].
     clear();
     int w_max = 0;
     for (int u = 1; u \le n; ++u)
       for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
         w_max = max(w_max, g[u][v].w);
     for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
     while (matching()) ++n_matches;
     for (int u = 1; u \le n; ++u)
       if (match[u] && match[u] < u)</pre>
          tot_weight += g[u][match[u]].w;
     return make_pair(tot_weight, n_matches);
   void add_edge(int ui, int vi, int wi) {
     g[ui][vi].w = g[vi][ui].w = wi;
   void init(int _n) {
     n = _n;
     for (int u = 1; u <= n; ++u)</pre>
       for (int_v=1; v <= n; ++v)</pre>
         g[u][v] = edge(u, v, 0);
  }
};
4.9 Maximum Clique
```

```
struct MaxClique {
  // change to bitset for n > 64.
  int n, deg[maxn];
```

```
uint64_t adj[maxn], ans;
  vector<pair<int, int>> edge;
  void init(int n_) {
    n = n_{-}
    fill(adj, adj + n, 0ull);
    fill(deg, deg + n, 0);
    edge.clear();
  void add_edge(int u, int v) {
    edge.emplace_back(u, v);
    ++deg[u], ++deg[v];
  vector<int> operator()() {
    vector<int> ord(n);
    iota(ord.begin(), ord.end(), 0)
    sort(ord.begin(), ord.end(), [&](int u, int v) {
    return deg[u] < deg[v]; });</pre>
    vector<int> id(n);
     for (int i = 0; i < n; ++i) id[ord[i]] = i;</pre>
    for (auto e : edge) {
       int u = id[e.first], v = id[e.second];
       adj[u] |= (1ull << v);
       adj[v] = (1ull \ll u);
    uint64_t r = 0, p = (1ull << n) - 1;
    ans = 0;
    dfs(r, p);
    vector<int> res;
    for (int i = 0; i < n; ++i) {
       if (ans >> i & 1) res.push_back(ord[i]);
    return res;
#define pcount __builtin_popcountll
  void dfs(uint64_t r, uint64_t p) {
    if (p == 0) {
       if (pcount(r) > pcount(ans)) ans = r;
       return;
    if (pcount(r | p) <= pcount(ans)) return;</pre>
    int x = __builtin_ctzll(p & -p);
    uint64_t c = p_{\alpha} \sim adj[x];
    while (c > 0) {
   // bitset._Find_first(); bitset._Fint
      x = __builtin_ctzll(c & -c);
r |= (1ull << x);</pre>
      dfs(r, p & adj[x]);
r &= ~(1ull << x);</pre>
       p &= ~(1ull << x);
       c ^= (1ull << x);
  }
};
```

4.10 Tarjan's Articulation Point

```
vector<pair<int, int>> g[maxn];
int low[maxn], tin[maxn], t;
int bcc[maxn], sz;
int a[maxn], b[maxn], deg[maxn];
bool cut[maxn], ins[maxn];
vector<int> ed[maxn];
stack<int> st;
void dfs(int x, int p) {
 tin[x] = low[x] = ++t;
  int ch = 0;
  for (auto u : g[x]) if (u.first != p) {
    if (!ins[u.second]) st.push(u.second), ins[u.second
    if (tin[u.first]) {
      low[x] = min(low[x], tin[u.first]);
      continue;
    }
    ++ch;
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
```

```
if (low[u.first] >= tin[x]) {
    cut[x] = true;
    ++sz;
    while (true) {
        int e = st.top(); st.pop();
        bcc[e] = sz;
        if (e == u.second) break;
    }
}
if (ch == 1 && p == -1) cut[x] = false;
}
```

4.11 Tarjan's Bridge

```
vector<pair<int, int>> g[maxn];
int tin[maxn], low[maxn], t;
int a[maxn], b[maxn];
int bcc[maxn], sz;
bool br[maxn];
stack<int> st;
void dfs(int x, int p) {
  tin[x] = low[x] = ++t;
  st.push(x);
  for (auto u : g[x]) if (u.first != p) {
    if (tin[u.first]) {
      low[x] = min(low[x], tin[u.first]);
      continue;
    dfs(u.first, x);
    low[x] = min(low[x], low[u.first]);
    if (low[u.first] == tin[u.first]) br[u.second] =
  if (tin[x] == low[x]) {
    ++SZ;
    while (st.size()) {
      int u = st.top(); st.pop();
      bcc[u] = sz;
      if (u == x) break;
  }
}
```

4.12 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[
     maxn], val[maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1);
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
  fill(rp, rp + n, -1)
  fill(dom, dom + n, -1);
  tk = 0;
  for (int i = 0; i < n; ++i)
    g[i].clear();
void add_edge(int x, int y) {
  g[x].push_back(y);
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  for (int &u : g[x])
    if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
     r[dfn[u]].push_back(dfn[x]);
}
```

```
void merge(int x, int y) {
 fa[x] = y;
int find(int x, int c = 0) {
  if (fa[x] == x) return x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[
    x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
 // return the father of each node in the dominator
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int &u : r[i]) sdom[i] = min(sdom[i], sdom[
    find(u)]);
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
      int p = find(u);
      if (sdom[p] == i) dom[u] = i;
      else dom[u] = p;
    if (i) merge(i, rp[i]);
 vector<int> p(n, -1);
  for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i])</pre>
    dom[i] = dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
```

4.13 System of Difference Constraints

Given m constrains on n variables x_1, x_2, \ldots, x_n of form $x_i - x_j \leq w$ (resp, $x_i - x_j \geq w$), connect $i \to j$ with weight w. Then connect $0 \to i$ for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to x_i .

5 String

5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
 vector<int> f(s.size(), 0);
 // f[i] = length of the longest prefix (excluding s)
    [0:i]) such that it coincides with the suffix of s
    [0:i] of the same length
 // i + 1 - f[i] is the length of the smallest
    recurring period of s[0:i]
  int k = 0:
 for (int i = 1; i < (int)s.size(); ++i) {
    while (k > 0 \& s[i] != s[k]) k = f[k - 1];
    if (s[i] == s[k]) ++k;
    f[i] = k;
 return f;
vector<int> search(const string &s, const string &t) {
 // return 0-indexed occurrence of t in s
 vector<int> f = kmp(t), res;
  int k = 0;
 for (int i = 0; i < (int)s.size(); ++i) {
   while (k > 0 \&\& (k == (int)t.size() || s[i] != t[k]
    ])) k = f[k - 1];
    if (s[i] == t[k]) ++k;
   if (k == (int)t.size()) res.push_back(i - t.size()
    + 1);
 }
  return res;
```

5.2 Z Algorithm

```
int z[maxn];
// z[i] = longest common prefix of suffix i and suffix

void z_function(const string& s) {
    memset(z, 0, sizeof(z));
    z[0] = (int)s.length();
    int l = 0, r = 0;
    for (int i = 1; i < s.length(); ++i) {
        z[i] = max(0, min(z[i - l], r - i + 1));
        while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
            l = i; r = i + z[i];
            ++z[i];
        }
    }
}</pre>
```

5.3 Manacher's Algorithm

5.4 Aho-Corasick Automaton

```
struct AC {
  static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn
    ][26], f[maxn];
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
    cnt[sz] = 0;
    return sz++;
  void init() {
    sz = 0;
    root = gnode();
  int add(const string &s) {
    int now = root;
    for (int i = 0; i < s.length(); ++i) {</pre>
      if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a']
    ] = gnode();
      now = ch[now][s[i] - 'a'];
    ed[now] = 1;
    return now;
  void build_fail() {
    ql = qr = 0; q[qr++] = root;
    while (ql < qr) {</pre>
      int now = q[ql++];
      for (int i = 0; i < 26; ++i) if (ch[now][i] !=
        int p = ch[now][i], fp = f[now];
while (fp != -1 && ch[fp][i] == -1) fp = f[fp];
         int pd = fp != -1 ? ch[fp][i] : root;
         f[p] = pd;
         el[p] = ed[pd] ? pd : el[pd];
```

```
q[qr++] = p;
    }
  }
  void build(const string &s) {
    build_fail();
    int now = root;
    for (int i = 0; i < s.length(); ++i) {</pre>
      while (now != -1 && ch[now][s[i] - 'a'] == -1)
    now = f[now];
      now = now != -1 ? ch[now][s[i] - 'a'] : root;
      ++cnt[now];
    for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] +=
    cnt[q[i]];
};
```

Suffix Automaton

```
struct SAM {
  static const int maxn = 5e5 + 5;
  int nxt[maxn][26], to[maxn], len[maxn];
  int root, last, sz;
  int gnode(int x) {
    for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
    to[sz] = -1;
    len[sz] = x;
    return sz++;
  void init() {
    sz = 0;
    root = gnode(0);
    last = root;
  void push(int c) {
    int cur = last;
    last = gnode(len[last] + 1);
for (; ~cur && nxt[cur][c] == -1; cur = to[cur])
    nxt[cur][c] = last;
    if (cur == -1) return to[last] = root, void();
    int link = nxt[cur][c];
    if (len[link] == len[cur] + 1) return to[last] =
    link, void();
    int tlink = gnode(len[cur] + 1);
    for (; ~cur && nxt[cur][c] == link; cur = to[cur])
    nxt[cur][c] = tlink;
    for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[</pre>
    link][i];
    to[tlink] = to[link];
    to[link] = tlink;
    to[last] = tlink;
  void add(const string &s) {
    for (int i = 0; i < s.size(); ++i) push(s[i] - 'a')
  bool find(const string &s) {
    int cur = root;
for (int i = 0; i < s.size(); ++i) {
   cur = nxt[cur][s[i] - 'a'];
</pre>
      if (cur == -1) return false;
    }
    return true;
  int solve(const string &t) {
    int res = 0, cnt = 0;
    int cur = root;
for (int i = 0; i < t.size(); ++i) {</pre>
      if (~nxt[cur][t[i] - 'a']) {
        ++cnt;
        cur = nxt[cur][t[i] - 'a'];
      = to[cur]);
        if (~cur) cnt = len[cur] + 1, cur = nxt[cur][t[
       - 'a<sup>†</sup>7;
        else cnt = 0, cur = root;
      res = max(res, cnt);
```

```
return res;
};
```

5.6 Suffix Array

```
namespace sfxarray {
bool t[maxn * 2];
int hi[maxn], rev[maxn];
int _s[maxn * 2], sa[maxn * 2], c[maxn * 2], x[maxn], p
      [maxn], q[maxn * 2];
    sa[i]: sa[i]-th suffix is the i-th lexigraphically
      smallest suffix.
 // hi[i]: longest common prefix of suffix sa[i] and
      suffix sa[i - 1].
void pre(int *sa, int *c, int n, int z) {
   memset(sa, 0, sizeof(int) * n);
   memcpy(x, c, sizeof(int) * z);
void induce(int *sa, int *c, int *s, bool *t, int n,
     int z) {
   memcpy(x + 1, c, sizeof(int) * (z - 1));
for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] -
     1]) sa[x[s[sa[i] - 1]]++] = sa[i] - 1;
   memcpy(x, c, sizeof(int) * z);
   for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i]
       - 1]) sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
void sais(int *s, int *sa, int *p, int *q, bool *t, int
       *c, int n, int z) {
   bool uniq = t[n - 1] = true;
   int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n,
     last = -1;
   memset(c, 0, sizeof(int) * z);
for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];</pre>
   if (uniq) {
     for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
   for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i]
      + 1] ? t[i + 1] : s[i] < s[i + 1]);
   pre(sa, c, n, z);
for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i -</pre>
      1]) sa[--x[s[i]]] = p[q[i] = nn++] = i;
   induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] &&
      !t[sa[i] - 1]) {
     bool neq = last < 0 \mid \mid memcmp(s + sa[i], s + last,
      (p[q[sa[i]] + 1] - sa[i]) * sizeof(int));
     ns[q[last = sa[i]]] = nmxz += neq;
   sais(ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz +
      1);
   pre(sa, c, n, z);
   for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i
     ]]]]] = p[nsa[i]];
   induce(sa, c, s, t, n, z);
void build(const string &s) {
   for (int i = 0; i < (int)s.size(); ++i) _s[i] = s[i];
   _s[(int)s.size()] = 0; // s shouldn't contain 0
sais(_s, sa, p, q, t, c, (int)s.size() + 1, 256);
for (int i = 0; i < (int)s.size(); ++i) sa[i] = sa[i
     + 1];
   for (int i = 0; i < (int)s.size(); ++i) rev[sa[i]] =</pre>
   int ind = 0; hi[0] = 0;
   for (int i = 0; i < (int)s.size(); ++i) {</pre>
     if (!rev[i]) {
        ind = 0;
        continue;
     while (i + ind < (int)s.size() \&\& s[i + ind] == s[
      sa[rev[i] - 1] + ind]) ++ind;
     hi[rev[i]] = ind ? ind-- : 0;
}}
```

5.7 Lexicographically Smallest Rotation

```
string rotate(const string &s) {
  int n = s.length();
  string t = s + s;
  int i = 0, j = 1;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && t[i + k] == t[j + k]) ++k;
    if (t[i + k] <= t[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
  }
  int pos = (i < n ? i : j);
  return t.substr(pos, n);
}</pre>
```

6 Math

6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
cplx operator+(const cplx &rhs) const { return cplx(
  re + rhs.re, im + rhs.im); }
cplx operator-(const cplx &rhs) const { return cplx(
  re - rhs.re, im - rhs.im); }
  cplx operator*(const cplx &rhs) const { return cplx(
    re * rhs.re - im * rhs.im, re * rhs.im + im * rhs.
    re); }
  cplx conj() const { return cplx(re, -im); }
const int maxn = 262144;
const double pi = acos(-1);
cplx omega[maxn + 1];
bool init;
void prefft() {
  for (int i = 0; i <= maxn; ++i)</pre>
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi
      i / maxn));
void bitrev(vector<cplx> &v, int n) {
  int z = builtin_ctz(n) - 1;
  int z = __builtin_ctz(n) -
  for (int i = 0; i < n; ++i) {
    int x = 0;
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j \& 
    1) << (z - j);
    if (x > i) swap(v[x], v[i]);
void fft(vector<cplx> &v, int n) {
  if (!init) {
    init = true;
    prefft();
  bitrev(v, n);
  for (int s = 2; s <= n; s <<= 1) {
    int z = s \gg 1;
    for (int i = 0; i < n; i += s) {
       for (int k = 0; k < z; ++k) {
         cplx x = v[i + z + k] * omega[maxn / s * k];

v[i + z + k] = v[i + k] - x;
         v[i + k] = v[i + k] + x;
 }
void ifft(vector<cplx> &v, int n) {
  fft(v, n)
  reverse(v.begin() + 1, v.end())
  for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n
vector<long long> convolution(const vector<int> &a,
    const vector<int> &b) {
```

```
// Should be able to handle N \leq 10^5, C \leq 10^4
  int sz = 1;
  while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
  vector<cplx> v(sz);
  for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
  fft(v, sz);
  for (int i = 0; i <= sz / 2; ++i) {
     int j = (sz - i) & (sz - 1);
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()
     ) * cplx(0, -0.25);
if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v
[i].conj()) * cplx(0, -0.25);
     v[i] = x;
  ifft(v, sz);
  vector<long long> c(sz);
  for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
  return c;
vector<int> convolution_mod(const vector<int> &a, const
      vector<int> &b, int p) {
  int sz = 1;
  while (sz < (int)a.size() + (int)b.size() - 1) sz <<=</pre>
  vector<cplx> fa(sz), fb(sz);
  for (int i = 0; i < (int)a.size(); ++i) {
  int x = (a[i] % p + p) % p;
  fa[i] = cplx(x & ((1 << 15) - 1), x >> 15);
  for (int i = 0; i < (int)b.size(); ++i) {</pre>
     int x = (b[i] \% p + p) \% p;
     fb[i] = cp[x(x \& ((1 << 15) - 1), x >> 15);
  fft(fa, sz), fft(fb, sz);
  double r = 0.25 / sz; cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
  for (int i = 0; i <= (sz >> 1); ++i) {
  int j = (sz - i) & (sz - 1);

     cplx a1 = (fa[i] + fa[j].conj());

cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx b1 = (fb[i] + fb[j].conj()) * r3;
cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
  cplx c1 = (fa[j] + fa[i].conj());
       cplx c2 = (fa[j] - fa[i].conj()) * r2;
cplx d1 = (fb[j] + fb[i].conj()) * r3;
        cplx_d2 = (fb[j] - fb[i].conj()) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
fb[j] = a1 * b2 + a2 * b1;
  fft(fa, sz), fft(fb, sz);
  vector<int> res(sz);
  for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \%
     р;
  return res;
}}
6.2 Number Theoretic Transform
```

```
template <long long mod, long long root>
struct NTT {
  vector<long long> omega;
  NTT() {
    omega.resize(maxn + 1);
    long long x = fpow(root, (mod - 1) / maxn);
    omega[0] = 1ll;
    for (int i = 1; i <= maxn; ++i)
        omega[i] = omega[i - 1] * x % mod;</pre>
```

```
long long fpow(long long a, long long n) {
    (n += mod - 1) \%= mod - 1;
    long long r = 1;
    for (; n; n >>= 1) {
      if (n & 1) (r *= a) %= mod;
      (a *= a) \%= mod;
    return r:
  void bitrev(vector<long long> &v, int n) {
    int z = __builtin_ctz(n) - 1;
    for (int i = 0; i < n; ++i) {
      int x = 0;
      for (int j = 0; j \le z; ++j) x \stackrel{\wedge}{=} (i >> j \& 1) <<
     (z - j);
      if (x > i) swap(v[x], v[i]);
  void ntt(vector<long long> &v, int n) {
    bitrev(v, n);
    for (int s = 2; s <= n; s <<= 1) {
      int z = s >> 1;
for (int i = 0; i < n; i += s) {</pre>
        for (int k = 0; k < z; ++k) {
  long long x = v[i + k + z] * omega[maxn / s *</pre>
     k] % mod;
           v[i + k + z] = (v[i + k] + mod - x) \% mod;
           (v[i + k] += x) \% = mod;
      }
    }
  void intt(vector<long long> &v, int n) {
    ntt(v, n);
    for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i
    long long inv = fpow(n, -1);
    for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;
  vector<long long> operator()(vector<long long> a,
    vector<long long> b) {
    int sz = 1;
    while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
    while (a.size() < sz) a.push_back(0);</pre>
    while (b.size() < sz) b.push_back(0);</pre>
    ntt(a, sz), ntt(b, sz);
    vector<long long> c(sz);
    for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] %
    mod:
    intt(c, sz);
    return c;
vector<long long> convolution(vector<long long> a,
    vector<long long> b) {
  NTT<mod1, root1> conv1;
 NTT<mod2, root2> conv2;
  vector<long long> pa(a.size()), pb(b.size());
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i]
     % mod1 + mod1) % mod1;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i])
     % mod1 + mod1) % mod1;
  vector<long long> c1 = conv1(pa, pb);
  for (int i = 0; i < (int)a.size(); ++i) pa[i] = (a[i]
     % mod2 + mod2) % mod2;
  for (int i = 0; i < (int)b.size(); ++i) pb[i] = (b[i])
     % mod2 + mod2) % mod2;
  vector<long long> c2 = conv2(pa, pb);
  long long x = conv2.fpow(mod1, -1);
long long y = conv1.fpow(mod2, -1);
  long long prod = mod1 * mod2;
  vector<long long> res(c1.size());
  for (int i = 0; i < c1.size(); ++i) {
  long long z = ((ull)fmul(c1[i] * mod2 % prod, y</pre>
    prod) + (ull)fmul(c2[i] * mod1 % prod, x, prod)) %
    if (z \ge prod / 2) z = prod;
    res[i] = z;
  return res;
```

6.2.1 NTT Prime List

}

```
Prime
             Root
                     Prime
                                   Root
                     167772161
7681
             17
12289
             11
                     104857601
40961
                     985661441
65537
                     998244353
786433
             10
                     1107296257
                                   10
5767169
             3
                     2013265921
                                   31
7340033
                     2810183681
                                   11
             3
23068673
                     2885681153
469762049
```

6.3 Polynomial Division

```
vector<int> inverse(const vector<int> &v, int n) {
  vector<int> q(1, fpow(v[0], mod - 2));
for (int i = 2; i <= n; i <<= 1) {
     vector<int> fv(v.begin(), v.begin() + i);
     vector<int> fq(q.begin(), q.end());
fv.resize(2 * i), fq.resize(2 * i);
     ntt(fq, 2 * i), ntt(fv, 2 * i);

for (int j = 0; j < 2 * i; ++j) {

    fv[j] = fv[j] * 1ll * fq[j] % mod * 1ll * fq[j] %
      mod:
     intt(fv, 2 * i);
     vector<int> res(i);
     for (int j = 0; j < i; ++j) {
  res[j] = mod - fv[j];</pre>
       if (j < (i >> 1)) (res[j] += 2 * q[j] % mod) %=
     mod:
     q = res;
  }
  return q;
vector<int> divide(const vector<int> &a, const vector<</pre>
     int> &b) {
   // leading zero should be trimmed
  int n = (int)a.size(), m = (int)b.size();
  int k = 2;
  while (k < n - m + 1) k <<= 1;
  vector<int> ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - i -
  for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - i -
      1];
  vector<int> rbi = inverse(rb, k);
  vector<int> res = convolution(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
```

6.4 Polynomial Square Root

```
// Find G(x) such that G^2(x) = F(x) (mod x^{N+1})
vector<int> solve(vector<int> b, int n) {
   if (n == 1) return {sqr[b[0]]};
   vector<int> h = solve(b, n >> 1); h.resize(n);
   vector<int> c = inverse(h, n);
   h.resize(n << 1); c.resize(n << 1);
   vector<int> res(n << 1);
   conv.ntt(h, n << 1);
   for (int i = n; i < (n << 1); ++i) b[i] = 0;
   conv.ntt(b, n << 1);
   conv.ntt(c, n << 1);
   for (int i = 0; i < (n << 1); ++i) res[i] = 111 * (h[
        i] + 111 * c[i] * b[i] % mod) % mod * inv2 % mod;
   conv.intt(res, n << 1);
   for (int i = n; i < (n << 1); ++i) res[i] = 0;
   return res;
}</pre>
```

6.5 Fast Walsh-Hadamard Transform

6.5.1 XOR Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_0) tf(A_1))$
- $utf(A) = (utf(\frac{A_0 + A_1}{2}), utf(\frac{A_0 A_1}{2}))$

6.5.2 OR Convolution

- $tf(A) = (tf(A_0), tf(A_0) + tf(A_1))$
- $utf(A) = (utf(A_0), utf(A_1) utf(A_0))$

6.5.3 AND Convolution

- $tf(A) = (tf(A_0) + tf(A_1), tf(A_1))$
- $utf(A) = (utf(A_0) utf(A_1), utf(A_1))$

6.6 Simplex Algorithm

```
namespace simplex {
// maximize c^Tx under Ax <= B
// return vector<double>(n, -inf) if the solution doesn
     't exist
// return vector<double>(n, +inf) if the solution is
    unbounded
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
vector<vector<double>> d;
vector<int> p, q;
void pivot(int r, int s) {
  double inv = 1.0 / d[r][s];
  for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {
    if (i != r && j != s) d[i][j] -= d[r][j] * d[i][s]
      * inv;
  for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s]
    *= -inv;
  for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j]
    *= +inv;
  d[r][s] = inv;
  swap(p[r], q[s]);
bool phase(int z) {
  int x = m + z;
while (true) {
    int s = -1;
    for (int i = 0; i \le n; ++i) {
      if (!z && q[i] == -1) continue;
      if (s == -1) | d[x][i] < d[x][s]) s = i;
    if (d[x][s] > -eps) return true;
    int r = -1;
    for (int i = 0; i < m; ++i) {
      if (d[i][s] < eps) continue;
if (r == -1 || d[i][n + 1] / d[i][s] < d[r][n +</pre>
    1] / d[r][s]) r = i;
    if (r == -1) return false;
    pivot(r, s);
  }
vector<double> solve(const vector<vector<double>> &a,
    const vector<double> &b, const vector<double> &c) {
  m = b.size(), n = c.size();
  d = vector<vector<double>>(m + 2, vector<double>(n +
    2));
  for (int i = 0; i < m; ++i) {
    for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] =</pre>
    -1, d[i][n + 1] = b[i];
  for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i]
    ];
```

```
q[n] = -1, d[m + 1][n] = 1;
  int r = 0;
  for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n
    + 1]) r = i;
  if (d[r][n + 1] < -eps) {
    pivot(r, n);
    if (!phase(1) \mid | d[m + 1][n + 1] < -eps) return
    vector<double>(n, -inf);
    for (int i = 0; i < m; ++i) if (p[i] == -1) {
      int s = min_element(d[i].begin(), d[i].end() - 1)
      - d[i].begin();
      pivot(i, s);
    }
  if (!phase(0)) return vector<double>(n, inf);
  vector<double> x(n);
  for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d
    [i][n + 1];
  return x:
}}
```

6.6.1 Construction

Standard form: maximize $\sum_{1 \leq i \leq n} c_i x_i$ such that for all $1 \leq j \leq m$, $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$ and $x_i \geq 0$ for all $1 \leq i \leq n$.

- 1. In case of minimization, let $c'_i = -c_i$
- 2. $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3. $\sum_{1 \le i \le n} A_{ji} x_i = b_j$
 - $\sum_{1 \leq i \leq n} A_{ji} x_i \leq b_j$
 - $\sum_{1 < i < n} A_{ji} x_i \ge b_j$
- 4. If x_i has no lower bound, replace x_i with $x_i x_i'$

6.7 Schreier-Sims Algorithm

```
namespace schreier {
int n;
vector<vector<int>>> bkts, binv;
vector<vector<int>> lk;
vector<int> operator*(const vector<int> &a, const
     vector<int> &b) {
  vector<int> res(a.size());
for (int i = 0; i < (int)a.size(); ++i)</pre>
    res[i] = b[a[i]];
  return res;
vector<int> inv(const vector<int> &a) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i)
    res[a[i]] = i;
  return res;
int filter(const vector<int> &g, bool add = true) {
  n = (int)bkts.size();
  vector<int> p = g;
for (int i = 0; i < n; ++i) {</pre>
    assert(p[i] >= 0 \&\& p[i] < (int)lk[i].size());
    int res = lk[i][p[i]];
if (res == -1) {
       if (add) {
         bkts[i].push_back(p);
         binv[i].push_back(inv(p))
         lk[i][p[i]] = (int)bkts[i].size() - 1;
       return i;
    p = p * binv[i][res];
  return -1;
bool inside(const vector<int> &g) {
  return filter(g, false) == -1;
void solve(const vector<vector<int>> &gen, int _n) {
 n = _n;
```

```
bkts.clear(), bkts.resize(n);
binv.clear(), binv.resize(n);
lk.clear(), lk.resize(n);
  vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
  for (int i = 0; i < n; ++i) {
    lk[i].resize(n, -1)
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i)
    filter(gen[i]);
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
      for (int k = 0; k < (int)bkts[i].size(); ++k) {
  for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l))
      }
    }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.
    first][b.second]);
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].
    size() - 1);
    for (int i = 0; i < n; ++i) {
  for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
         if (i <= res)
           upd.emplace(make_pair(i, j), pr);
         if (res <= i)
           upd.emplace(pr, make_pair(i, j));
      }
    }
 }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i)
    res = res * bkts[i].size();
  return res;
```

6.8 Miller Rabin

```
// n < 2^64
                chk = [2, 325, 9375, 28178, 450775,
    9780504, 17952650227
vector<long long> chk = { 2, 325, 9375, 28178, 450775,
    9780504, 1795265022 };
bool check(long long a, long long u, long long n, int t
  a = fpow(a, u, n);
  if (a == 0) return true;
  if (a == 1 \mid \mid a == n - 1) return true;
  for (int i = 0; i < t; ++i) {
   a = fmul(a, a, n);
if (a == 1) return false;
    if (a == n - 1) return true;
  return false;
bool is_prime(long long n) {
  if (n < 2) return false;</pre>
  if (n % 2 == 0) return n == 2;
  long long u = n - 1; int t = 0;
  for (; !(u & 1); u >>= 1, ++t);
  for (long long i : chk) {
   if (!check(i, u, n, t)) return false;
```

Pollard's Rho

return true;

6.9

```
long long f(long long x, long long n, int p) { return (
    fmul(x, x, n) + p) % n; }
map<long long, int> cnt;
void pollard_rho(long long n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n \% 2 == 0) return pollard_rho(n / 2), ++cnt[2],
  long long x = 2, y = 2, d = 1, p = 1;
  while (true) {
    if (d != n && d != 1) {
      pollard_rho(n / d);
      pollard_rho(d);
      return;
    if (d == n) ++p;
    x = f(x, n, p); y = f(f(y, n, p), n, p);
    d = \_gcd(abs(x - y), n);
}
```

6.10 Meissel-Lehmer Algorithm

```
int prc[maxn];
long long phic[msz][nsz];
void sieve() {
  bitset<maxn> v:
  pr.push_back(0);
for (int i = 2; i < maxn; ++i) {</pre>
     if (!v[i]) pr.push_back(i);
     for (int j = 1; i * pr[j] < maxn; ++j) {
       v[i * pr[j]] = true;
       if (i % pr[j] == 0) break;
  for (int i = 1; i < pr.size(); ++i) prc[pr[i]] = 1;
for (int i = 1; i < maxn; ++i) prc[i] += prc[i - 1];</pre>
long long p2(long long, long long);
long long phi(long long m, long long n) {
  if (m < msz && n < nsz && phic[m][n] != -1) return</pre>
     phic[m][n];
  if (n == 0) return m;
if (pr[n] >= m) return 1;
  long long ret = phi(m, n - 1) - phi(m / pr[n], n - 1)
  if (m < msz && n < nsz) phic[m][n] = ret;</pre>
  return ret;
long long pi(long long m) {
  if (m < maxn) return prc[m];</pre>
  long long n = pi(cbrt(m));
  return phi(m, n) + n - 1 - p2(m, n);
long long p2(long long m, long long n) {
  long long ret = 0;
  long long lim = sqrt(m);
  for (int i = n + 1; pr[i] <= lim; ++i) ret += pi(m /</pre>
     pr[i]) - pi(pr[i]) + 1;
  return ret;
}
```

6.11 Discrete Logarithm

```
// to solve discrete x for x^a = b \pmod{p} with p is
// let c = primitive root of p
// find k such that c^k = b \pmod{p} by bsgs
// solve fa = k (mod p - 1) by euclidean algorithm
// x = c^f
int bsgs(int a, int b, int p) {
   // return L such that a^L = b \pmod{p}
  if (p == 1) {
     if (!b) return a != 1;
     return -1:
  if (b == 1) {
     if (a) return 0;
     return -1;
  if (a \% p == 0) {
     if (!b) return 1;
     return -1;
  int num = 0, d = 1;
  while (true) {
    int r = __gcd(a, p);
if (r == 1) break;
    if (b % r) return -1;
     ++num;
    b /= r, p /= r;
    d = (111 * d * a / r) % p;
  for (int i = 0, now = 1; i < num; ++i, now = 1ll *
    now * a % p) {</pre>
     if (now == b) return i;
  int m = ceil(sqrt(p)), base = 1;
  map<int, int> mp;
for (int i = 0; i < m; ++i) {</pre>
     if (mp.find(base) == mp.end()) mp[base] = i;
     else mp[base] = min(mp[base], i);
    base = \overline{111} * base * \overline{a} % p;
  for (int i = 0; i < m; ++i) {
    \ensuremath{\text{//}} can be modified to fpow if p is prime
    int r, x, y; tie(r, x, y) = extgcd(d, p);
x = (111 * x * b % p + p) % p;
    if (mp.find(x) != mp.end()) return i * m + mp[x] +
    d = 111 * d * base % p;
  }
  return -1;
}
```

6.12 Gaussian Elimination

6.13 Linear Equations (full pivoting)

```
void linear_equation(vector<vector<double>> &d, vector<
   double> &aug, vector<double> &sol) {
  int n = d.size(), m = d[0].size();
```

```
vector<int> r(n), c(m);
iota(r.begin(), r.end(), 0);
iota(c.begin(), c.end(), 0);
for (int i = 0; i < m; ++i) {</pre>
  int p = -1, z = -1;
  for (int j = i; j < n; ++j) {
  for (int k = i; k < m; ++k) {
    if (fabs(d[r[j]][c[k]]) < eps) continue;</pre>
       if (p == -1 \mid | fabs(d[r[j]][c[k]]) > fabs(d[r[p]
  ]][c[z]])) p = j, z = k;
  if (p == -1) continue;
  swap(r[p], r[i]), swap(c[z], c[i]);
for (int j = 0; j < n; ++j) {</pre>
     if (i == j) continue
     double z = d[r[j]][c[i]] / d[r[i]][c[i]];
     for (int k = 0; k < m; ++k) d[r[j]][c[k]] -= z *
  d[r[i]][c[k]];
     aug[r[j]] -= z * aug[r[i]];
}
vector<vector<double>> fd(n, vector<double>(m));
vector<double> faug(n), x(n);
for (int i = 0; i < n; ++i) {
  for (int j = 0; j < m; ++j) fd[i][j] = d[r[i]][c[j
  11;
  faug[i] = aug[r[i]];
d = fd, aug = faug;
for (int i = n - 1; i >= 0; --i) {
  double p = 0.0;
  for (int j = i + 1; j < n; ++j) p += d[i][j] * x[j]
  x[i] = (aug[i] - p) / d[i][i];
for (int i = 0; i < n; ++i) sol[c[i]] = x[i];
```

6.14 μ function

6.15 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1} + 1} \rfloor} \rfloor
```

6.16 De Bruijn Sequence

```
int res[maxn], aux[maxn], a[maxn], sz;

void db(int t, int p, int n, int k) {
   if (t > n) {
      if (n % p == 0) {
       for (int i = 1; i <= p; ++i) res[sz++] = aux[i];
      }
}</pre>
```

```
} else {
    aux[t] = aux[t - p];
    db(t + 1, p, n, k);
    for (int i = aux[t - p] + 1; i < k; ++i) {
        aux[t] = i;
        db(t + 1, t, n, k);
    }
}

int de_bruijn(int k, int n) {
    // return cyclic string of length k^n such that every
        string of length n using k character appears as a
        substring.
    if (k == 1) {
        res[0] = 0;
        return 1;
    }
    for (int i = 0; i < k * n; i++) aux[i] = 0;
    sz = 0;
    db(1, 1, n, k);
    return sz;
}</pre>
```

6.17 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
  if (!b) return make_tuple(a, 1, 0);
  T d, x, y;
  tie(d, x, y) = extgcd(b, a % b);
  return make_tuple(d, y, x - (a / b) * y);
}
```

6.18 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity: $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \end{cases} \\ = \frac{\frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1)}{-h(c, c-b-1, a, m-1)),} & \text{otherwise} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

6.19 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
  long long mult = mod[0];
  int n = (int)mod.size();
  long long res = a[0];
  for (int i = 1; i < n; ++i) {
    long long d, x, y;
    tie(d, x, y) = extgcd(mult, mod[i] * 1ll);
    if ((a[i] - res) % d) return -1;</pre>
```

```
long long new_mult = mult / __gcd(mult, 1ll * mod[i
]) * mod[i];
  res += x * ((a[i] - res) / d) % new_mult * mult %
    new_mult;
  mult = new_mult;
  ((res %= mult) += mult) %= mult;
}
return res;
}
```

6.20 Theorem

6.20.1 Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

6.20.2 Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniform randomly) if i < j and $(i,j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

6.20.3 Cayley's Formula

- Given a degree sequence d_1,d_2,\ldots,d_n for each labeled vertices, there're $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees.
- Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1, 2, \ldots, k$ belong to different components. Then $T_{n,k} = kn^{n-k-1}$.

6.21 Primes

```
97, 101, 131, 487, 593, 877, 1087, 1187, 1487, 1787, 3187, 12721, \\ 13331, 14341, 75577, 123457, 222557, 556679, 999983, \\ 1097774749, 1076767633, 100102021, 999997771, \\ 1001010013, 1000512343, 987654361, 999991231, \\ 999888733, 98789101, 987777733, 999991921, 1000000007, \\ 1000000087, 1000000123, 101010133, 1010102101, \\ 1000000000039, 10000000000037, 2305843009213693951, \\ 4611686018427387847, 9223372036854775783, 18446744073709551557
```

7 Dynamic Programming

7.1 Convex Hull Optimization

```
struct line {
  int m, y;
  int l, r;
  line(int m = 0, int y = 0, int l = -5, int r = 0
  1000000009): m(m), y(y), l(l), r(r) {} int get(int x) const { return m * x + y; }
  int useful(line le) const {
    return (int)(get(l) >= le.get(l)) + (int)(get(r) >=
      le.get(r));
};
int magic;
bool operator < (const line &a, const line &b) {</pre>
  if (magic) return a.m < b.m;</pre>
  return a.l < b.l;</pre>
set<line> st;
void addline(line l) {
  magic = 1;
  auto it = st.lower_bound(1);
  if (it != st.end() && it->useful(l) == 2) return;
  while (it != st.end() && it->useful(1) == 0) it = st.
    erase(it);
  if (it != st.end() && it->useful(l) == 1) {
```

```
int L = it -> l, R = it -> r, M;
    while (R > L) {
      M = (L + R + 1) >> 1;
      if (it->get(M) >= l.get(M)) R = M - 1;
      else L = M;
    line cp = *it;
    st.erase(it);
    cp.l = L + 1;
    if (cp.l <= cp.r) st.insert(cp);</pre>
 }
  else if (it != st.end()) l.r = it->l - 1;
  it = st.lower_bound(l);
 while (it != st.begin() && prev(it)->useful(l) == 0)
    it = st.erase(prev(it));
  if (it != st.begin() && prev(it)->useful(l) == 1) {
    --it;
    int L = it \rightarrow l, R = it \rightarrow r, M;
    while (R > L) {
      M = (L + R) >> 1;
      if (it->get(M) >= l.get(M)) L = M + 1;
      else R = M;
    line cp = *it;
    st.erase(it);
    cp.r = L - 1;
    if (cp.l <= cp.r) st.insert(cp);</pre>
   1.1 = L;
  else if (it != st.begin()) l.l = prev(it)->r + 1;
  if (l.l <= l.r) st.insert(l);
int getval(int d) {
 magic = 0;
  return (--st.upper_bound(line(0, 0, d, 0)))->get(d);
```

7.2 1D/1D Convex Optimization

```
struct segment {
   int i, l, r;
  segment() {}
   segment(int a, int b, int c): i(a), l(b), r(c) {}
inline long long f(int l, int r) {
  return dp[l] + w(l + 1, r);
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
  dp[i] = f(deq.front().i, i);</pre>
     while (deq.size() && deq.front().r < i + 1) deq.</pre>
     pop_front();
     deq.front().l = i + 1;
     segment seg = segment(i, i + 1, n);
     while (deq.size() && f(i, deq.back().1) < f(deq.</pre>
     back().i, deq.back().l)) deq.pop_back();
     if (deq.size()) {
       int d = 1048576, c = deq.back().1;
       while (d \gg 1) if (c + d \ll deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c +=
     d;
       deq.back().r = c; seg.l = c + 1;
     if (seg.l <= n) deq.push_back(seg);</pre>
  }
}
```

7.3 Condition

7.3.1 totally monotone (concave/convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

7.3.2 monge condition (concave/convex)

bool same(const double a, const double b){ return abs(a-

```
\begin{array}{l} \forall i < i', j < j', \, B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \, B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

8 Geometry

8.1 Basic

```
b)<1e-9; }
struct Point{
  double x,y
  Point():\hat{x}(0),y(0){}
  Point(double x, double y):x(x),y(y){}
Point operator+(const Point a,const Point b){ return
    Point(a.x+b.x,a.y+b.y); }
Point operator-(const Point a, const Point b){ return
    Point(a.x-b.x,a.y-b.y); }
Point operator*(const Point a,const double b){ return
    Point(a.x*b,a.y*b); }
Point operator/(const Point a,const double b){ return
    Point(a.x/b,a.y/b); }
double operator^(const Point a,const Point b){ return a
     .x*b.y-a.y*b.x;
double abs(const Point a){ return sqrt(a.x*a.x+a.y*a.y)
    ; }
struct Line{
  // ax+by+c=0
  double a,b,c;
  double angle;
  Point pa,pb;
  Line():a(0),b(0),c(0),angle(0),pa(),pb(){}
  Line(Point pa,Point pb):a(pa.y-pb.y),b(pb.x-pa.x),c(
    pa^pb), angle(atan2(-a,b)), pa(pa), pb(pb){}
};
Point intersect(Line la,Line lb){
  if(same(la.a*lb.b,la.b*lb.a))return Point(7122,7122);
  double bot=-la.a*ib.b+la.b*ib.a;
  return Point(-la.b*lb.c+la.c*lb.b,la.a*lb.c-la.c*lb.a
    )/bot;
}
bool intersect(Point p1, Point p2, Point p3, Point p4)
  if (\max(p1.x, p2.x) < \min(p3.x, p4.x) \mid \max(p3.x, p4.x) \mid
  .x) < min(p1.x, p2.x)) return false;
if (max(p1.y, p2.y) < min(p3.y, p4.y) || max(p3.y, p4
  .y) < min(p1.y, p2.y)) return false;
return sign((p3 - p1) % (p4 - p1)) * sign((p3 - p2) %
     (p4 - p2)) \le 0 \&\&
      sign((p1 - p3) \% (p2 - p3)) * sign((p1 - p4) % (
    p2 - p4)) <= 0;
int contain(const vector<Point> &ps, Point p) {
  // ps is not necessarily convex.
  int n = (int)ps.size();
  for (int i = 0; i < n; ++i) {
    Point a = ps[i], b = ps[(i + 1) % n];
    // on segment
    if ((p - a) * (b - a) >= 0 && (p - b) * (a - b) >=
    0 \& (p - a) \% (b - a) == 0) return 1;
  // infinity
  Point q = Point(1000000000, p.y);
  int res = 0;
```

```
for (int i = 0; i < n; ++i) {
   Point a = ps[i], b = ps[(i + 1) % n];
   if (intersect(a, b, p, q) && p.y >= min(a.y, b.y)
   && p.y < max(a.y, b.y)) res ^= 1;
}
// ps contains p.
if (res == 1) return 2;
return 0;
}</pre>
```

8.2 KD Tree

```
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[
     maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
  if (l == r) return -1;
  function<bool(const point &, const point &)> f = [dep
     ](const point &a, const point &b) {
     if (dep & 1) return a.x < b.x;
     else return a.y < b.y;</pre>
  };
  int m = (l + r) >> 1;
  nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
yl[m] = min(yl[m], yl[rc[m]]);
yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
     q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 111 * (a.x - b.x) +
       (a.y - b.y) * 111 * (a.y - b.y);
void dfs(const point &q, long long &d, int o, int dep =
      0) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x || !(dep & 1) && q.y <
    p[o].y) {
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
     if (~rc[o]) dfs(q, d, rc[o], dep + 1);
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res;
```

8.3 Delaunay Triangulation

```
namespace triangulation {
static const int maxn = 1e5 + 5;
vector<point> p;
set<int> g[maxn];
int o[maxn];
set<int> s;
void add_edge(int x, int y) {
  s.insert(x), s.insert(y);
  g[x].insert(y);
  g[y].insert(x);
bool inside(point a, point b, point c, point p) {
  if (((b - a) \land (c - a)) < 0) swap(b, c);
  function<long long(int)> sqr = [](int x) { return x *
      111 * x; };
  long long k11 = a.x - p.x, k12 = a.y - p.y, k13 = sqr
  (a.x) - sqr(p.x) + sqr(a.y) - sqr(p.y);
long long k21 = b.x - p.x, k22 = b.y - p.y, k23 = sqr
  (b.x) - sqr(p.x) + sqr(b.y) - sqr(p.y);
long long k31 = c.x - p.x, k32 = c.y - p.y, k33 = sqr
  (c.x) - sqr(p.x) + sqr(c.y) - sqr(p.y);
long long det = k11 * (k22 * k33 - k23 * k32) - k12 *
      (k21 * k33 - k23 * k31) + k13 * (k21 * k32 - k22 *
      k31);
  return det > 0;
bool intersect(const point &a, const point &b, const
     point &c, const point &d) {
  return ((b - a) \wedge (c - a)) * ((b - a) \wedge (d - a)) < 0
       ((d - c) \wedge (a - c)) * ((d - c) \wedge (b - c)) < 0;
void dfs(int l, int r) {
  if (r - l <= 3) {</pre>
     for (int i = 1; i < r; ++i) {</pre>
       for (int j = i + 1; j < r; ++j) add_edge(i, j);
     return;
  int m = (l + r) >> 1;
  dfs(l, m), dfs(m, r);
int pl = l, pr = r - 1;
  while (true) {
     int z = -1;
     for (int u : g[pl]) {
       long long c = ((p[p1] - p[pr]) \wedge (p[u] - p[pr]));
if (c > 0 \mid | c == 0 \&\& abs(p[u] - p[pr]) < abs(p[u] - p[pr])
     pl] - p[pr])) {
         z = u;
         break;
       }
     if (z != -1) {
       pl = z;
       continue;
     for (int u : g[pr]) {
       long long c = ((p[pr] - p[pl]) ^ (p[u] - p[pl]));
if (c < 0 || c == 0 && abs(p[u] - p[pl]) < abs(p[
     pr] - p[pl])) {
         z = u;
         break;
       }
     if (z != -1) {
       pr = z;
       continue;
     break;
  add_edge(pl, pr);
  while (true) {
     int z = -1;
     bool b = false;
     for (int u : g[pl]) {
       long long c = ((p[pl] - p[pr]) \land (p[u] - p[pr]));
       if (c < 0 \&\& (z == -1 | l inside(p[pl], p[pr], p[z])
     ], p[u]))) z = u;
     for (int u : g[pr]) {
       long long c = ((p[pr] - p[pl]) \wedge (p[u] - p[pl]));
```

```
if (c > 0 \& (z == -1 || inside(p[pl], p[pr], p[z
    if (z == -1) break;
    int x = pl, y = pr;
    if (b) swap(x, y);
    for (auto it = g[x].begin(); it != g[x].end(); ) {
  int u = *it;
      if (intersect(p[x], p[u], p[y], p[z])) {
        it = g[x].erase(it);
        g[u].erase(x);
      } else {
        ++it;
      }
    if (b) add_edge(pl, z), pr = z;
    else add_edge(pr, z), pl = z;
vector<vector<int>> solve(vector<point> v) {
  int n = v.size();
  for (int i = 0; i < n; ++i) g[i].clear();
for (int i = 0; i < n; ++i) o[i] = i;</pre>
  sort(o, o + n, [&](int i, int j) { return v[i] < v[j</pre>
    ]; });
  p.resize(n);
  for (int i = 0; i < n; ++i) p[i] = v[o[i]];
  dfs(0, n);
  vector<vector<int>> res(n);
  for (int i = 0; i < n; ++i)
    for (int j : g[i]) res[o[i]].push_back(o[j]);
  return res;
}}
```

8.4 Sector Area

```
// calc area of sector which include a, b
double SectorArea(Point a, Point b, double r) {
  double theta = atan2(a.y, a.x) - atan2(b.y, b.x);
  while (theta <= 0) theta += 2 * pi;
  while (theta >= 2 * pi) theta -= 2 * pi;
  theta = min(theta, 2 * pi - theta);
  return r * r * theta / 2;
}
```

8.5 Half Plane Intersection

```
bool jizz(Line l1,Line l2,Line l3){
  Point p=intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
bool cmp(const Line &a,const Line &b){
  return same(a.angle,b.angle)?(((b.pb-b.pa)^(a.pb-b.pa
    ))>eps):a.angle<b.angle;</pre>
// availble area for Line l is (l.pb-l.pa)^(p-l.pa)>0
vector<Point> HPI(vector<Line> &ls){
  sort(ls.begin(),ls.end(),cmp);
  vector<Line> pls(1,ls[0]);
for(unsigned int i=0;i<ls.size();++i)if(!same(ls[i].
    angle,pls.back().angle))pls.push_back(ls[i]);</pre>
  deque<int> dq; dq.push_back(0); dq.push_back(1);
  for(unsigned int i=2u;i<pls.size();++i){</pre>
    while(dq.size()>1u && jizz(pls[i],pls[dq.back()],
    pls[dq[dq.size()-2]]))dq.pop_back()
    while(dq.size()>1u && jizz(pls[i],pls[dq[0]],pls[dq
    [1]]))dq.pop_front();
    dq.push_back(i);
  while(dq.size()>1u && jizz(pls[dq.front()],pls[dq.
    back()],pls[dq[dq.size()-2]]))dq.pop_back();
  while(dq.size()>1u && jizz(pls[dq.back()],pls[dq[0]],
    pls[dq[1]]))dq.pop_front();
```

```
if(dq.size()<3u)return vector<Point>(); // no
    solution or solution is not a convex
vector<Point> rt;
for(unsigned int i=0u;i<dq.size();++i)rt.push_back(
    intersect(pls[dq[i]],pls[dq[(i+1)%dq.size()]]));
return rt;</pre>
```

8.6 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
  int n=int(ps.size());
  vector<int> id(n),pos(n);
  vector<pair<int,int>> line(n*(n-1)/2);
  int m=-1:
  for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=</pre>
  make_pair(i,j); ++m;
sort(line.begin(),line.end(),[&](const pair<int,int>
     &a,const pair<int,int> &b)->bool{
     if(ps[a.first].first==ps[a.second].first)return 0;
    if(ps[b.first].first==ps[b.second].first)return 1;
    return (double)(ps[a.first].second-ps[a.second].
     second)/(ps[a.first].first-ps[a.second].first) <</pre>
     double)(ps[b.first].second-ps[b.second].second)/(ps
     [b.first].first-ps[b.second].first);
  });
  for(int i=0;i<n;++i)id[i]=i</pre>
  sort(id.begin(),id.end(),[&](const int &a,const int &
    b){ return ps[a]<ps[b]; });
  for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
  for(int i=0;i<m;++i){</pre>
    auto l=line[i];
    // meow
    tie(pos[l.first],pos[l.second],id[pos[l.first]],id[
    pos[l.second]])=make_tuple(pos[l.second],pos[l.
     first],l.second,l.first);
}
```

8.7 Triangle Center

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2
  double by = (c.y + b.y) / 2;
double r1 = (\sin(a2) * (ax - bx) + \cos(a2) * (by - ay) / (\sin(a1) * \cos(a2) - \sin(a2) * \cos(a1));
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
}
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
}
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);
res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb +
      lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb +
      lc);
  return res;
```

8.8 Polygon Center

```
Point BaryCenter(vector<Point> &p, int n) {
   Point res(0, 0);
   double s = 0.0, t;
   for (int i = 1; i < p.size() - 1; i++) {
        t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
        s += t;
        res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
        res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
   }
   res.x /= (3 * s);
   res.y /= (3 * s);
   return res;
}</pre>
```

8.9 Maximum Triangle

```
double ConvexHullMaxTriangleArea(Point p[], int res[],
    int chnum) {
  double area = 0, tmp;
  res[chnum] = res[0];
  for (int i = 0, j = 1, k = 2; i < chnum; i++) {
   while (fabs(Cross(p[res[j]] - p[res[i]], p[res[(k +
     1) % chnum]] - p[res[i]])) > fabs(Cross(p[res[j]]
    - p[res[i]], p[res[k]] - p[res[i]]))) k = (k + 1) %
     chnum:
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] -
     p[res[i]]));
    if (tmp > area) area = tmp;
    while (fabs(Cross(p[res[(j + 1) % chnum]] - p[res[i
    ]], p[res[k]] - p[res[i]])) > fabs(Cross(p[res[j]]
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] -
     p[res[i]]));
    if (tmp > area) area = tmp;
  return area / 2;
}
```

8.10 Point in Polygon

```
bool on(point a, point b, point c) {
  if (a.x == b.x) {
    if (c.x != a.x) return false;
    if (c.y >= min(a.y, b.y) \&\& c.y <= max(a.y, b.y))
    return true;
    return false;
  if (((a - c) ^ (b - c)) != 0) return false;
  if (a.x > b.x) swap(a, b);
  if (c.x < min(a.x, b.x) | | c.x > max(a.x, b.x))
    return false;
  return ((a - b) \wedge (a - c)) == 0;
int sgn(long long x) {
  if (x > 0) return 1;
  if (x < 0) return -1;
  return 0;
bool in(const vector<point> &c, point p) {
  int last = -2;
  int n = c.size();
  for (int i = 0; i < c.size(); ++i) {
  if (on(c[i], c[(i + 1) % n], p)) return true;</pre>
    int g = sgn((c[i] - p) \wedge (c[(i + 1) % n] - p));
    if (last == -2) last = g;
    else if (last != g) return false;
  return true;
bool in(point a, point b, point c, point p) {
 return in({ a, b, c }, p);
```

```
}
bool inside(const vector<point> &ch, point t) {
  point p = ch[1] - ch[0];
  point q = t - ch[0];
  if ((p ^ q) < 0) return false;
  if ((p \land q) == 0) {
     if (p * q < 0) return false;</pre>
     if (q.len() > p.len()) return false;
     return true;
  p = ch[ch.size() - 1] - ch[0];
  if ((p \land q) > 0) return false;
  if ((p \land q) == 0) {
    if (p * q < 0) return false;</pre>
     if (q.len() > p.len()) return false;
    return true;
  p = ch[1] - ch[0];
  double ang = acos(1.0 * (p * q) / p.len() / q.len());
  int d = 20, z = ch.size() - 1;
  while (d--) {
    if (z - (1 << d) < 1) continue;
point p1 = ch[1] - ch[0];</pre>
    point p2 = ch[z - (1 << d)] - ch[0];
double tang = acos(1.0 * (p1 * p2) / p1.len() / p2.</pre>
     len());
     if (tang >= ang) z -= (1 << d);
  return in(ch[0], ch[z - 1], ch[z], t);
```

8.11 Circle-Line Intersection

```
// remove second level if to get points for line (
             defalut: segment)
void CircleCrossLine(Point a, Point b, Point o, double
             r, Point ret[], int &num) {
      double x0 = o.x, y0 = o.y
      double x1 = a.x, y1 = a.y;
     double x2 = b.x, y2 = b.y;

double dx = x2 - x1, dy = y2 - y1;

double A = dx * dx + dy * dy;

double B = 2 * dx * (x1 - x0) + 2 * dy * (y1 -
      double C = (x1 - x0) * (x1 - x0) + (y1 - y0) * (y1 - y0)
      double delta = B * B - 4 * A * C;
      num = 0;
      if (epssgn(delta) >= 0) {
             double t1 = (-B - sqrt(fabs(delta))) / (2 * A);
             double t2 = (-B + sqrt(fabs(delta))) / (2 * A);
             if (epssgn(t1 - 1.0) \le 0 \& epssgn(t1) \ge 0) ret[
             num++] = Point(x1 + t1 * dx, y1 + t1 * dy);
            if (epssgn(t2 - 1.0) <= 0 && epssgn(t2) >= 0) ret[num++] = Point(x1 + t2 * dx, y1 + t2 * dy);
}
vector<Point> CircleCrossLine(Point a, Point b, Point o
                  double r) {
      double x0 = o.x, y0 = o.y;
      double x1 = a.x, y1 = a.y;
      double x2 = b.x, y2 = b.y;
double dx = x2- x1, dy = y2 - y1;
      double A = dx * dx + dy * dy;
      double B = 2 * dx * (x1 - x0) + 2 * dy * (y1 - y0);
double C = (x1 - x0) * (x1 - x0) + (y1 - y0) * (y1 - y0) *
             y0) - r * r;
      double delta = B * B - 4 * A * C:
      vector<Point> ret;
      if (epssgn(delta) >= 0) {
            double t1 = (-B - sqrt(fabs(delta))) / (2 * A);
double t2 = (-B + sqrt(fabs(delta))) / (2 * A);
            if (epssgn(t1 - 1.0) \le 0 \& epssgn(t1) >= 0) ret.

emplace\_back(x1 + t1 * dx, y1 + t1 * dy);
            if (epssgn(t2 - 1.0) \le 0 \& epssgn(t2) >= 0) ret.
emplace_back(x1 + t2 * dx, y1 + t2 * dy);
      return ret;
```

8.12Circle-Triangle Intersection

```
// calc area intersect by circle with radius r and
    triangle OAB
double Calc(Point a, Point b, double r) {
  Point p[2];
  int num = 0;
  bool ina = epssqn(len(a) - r) < 0, inb = epssqn(len(b) - r)
    ) - r) < 0;
  if (ina) {
    if (inb) return fabs(Cross(a, b)) / 2.0; //
    triangle in circle
    else { // a point inside and another outside: calc
    sector and triangle area
CircleCrossLine(a, b, Point(0, 0), r, p, num);
       return SectorArea(b, p[0], r) + fabs(Cross(a, p
     [0])) / 2.0;
  } else {
    CircleCrossLine(a, b, Point(0, 0), r, p, num)
    if (inb) return SectorArea(p[0], a, r) + fabs(Cross
    (p[0], b)) / 2.0;
    else {
  if (num == 2) return SectorArea(a, p[0], r) +
    SectorArea(p[1], b, r) + fabs(Cross(p[\overline{0}], p[1])) /
    2.0; // segment ab has 2 point intersect with
    circle
      else return SectorArea(a, b, r); // segment has
    no intersect point with circle
  }
}
```

8.13Tangent from Point to Circle

```
array<Point, 2> tangent(Point o, double r, Point p) {
  double dist = sqrt((p - o) * (p - o));
  double len = sqrt(dist * dist - r * r);
  double ang = acos(len / dist);
  Point vec = (o - p) / dist * len;
  array<Point, 2> res;
  for (int i = 0; i < 2; ++i) {
    int z = i == 0 ? 1 : -1;
    Point v(vec.x * cos(z * ang) - vec.y * sin(z * ang)
      vec.x * sin(z * ang) + vec.y * cos(z * ang));
     res[i] = p + v;
  return res;
}
```

8.14 Minimum Distance of 2 Polygons

```
// p, q is convex
double TwoConvexHullMinDist(Point P[], Point Q[], int n
        int m) {
  int YMinP = 0, YMaxQ = 0;
  double tmp, ans = 999999999;
  for (i = 0; i < n; ++i) if (P[i].y < P[YMinP].y) YMinP
      = i;
  for (i = 0; i < m; ++i) if (Q[i].y > Q[YMaxQ].y) YMaxQ
       = i;
  P[n] = P[0], Q[m] = Q[0];
for (int i = 0; i < n; ++i) {
   while (tmp = Cross(Q[YMaxQ + 1] - P[YMinP + 1], P[</pre>
     YMinP] - P[YMinP + 1]) > Cross(Q[YMaxQ] - P[YMinP + 1], P[YMinP] - P[YMinP + 1]) YMaxQ = (YMaxQ + 1)
     if (tmp < 0) ans = min(ans, PointToSegDist(P[YMinP</pre>
     ], P[YMinP + 1], Q[YMaxQ]));
    else ans = min(ans, TwoSegMinDist(P[YMinP], P[YMinP
+ 1], Q[YMaxQ], Q[YMaxQ + 1]));
YMinP = (YMinP + 1) % n;
  }
  return ans;
```

8.15 2D Convex Hull

```
vector<point> convex(vector<point> p) {
   sort(p.begin(), p.end());
   vector<point> ch;
   for (int i = 0; i < n; ++i) {
   while (ch.size() >= 2 && ((p[i] - ch[ch.size() -
   2]) ^ (ch[ch.size() - 1] - ch[ch.size() - 2])) >=
       0) ch.pop_back();
      ch.push_back(p[i]);
   int t = ch.size();
   for (int i = n - 2; i >= 0; --i) {
  while (ch.size() > t && ((p[i] - ch[ch.size() - 2])
    ^ (ch[ch.size() - 1] - ch[ch.size() - 2])) >= 0)
       ch.pop_back();
      ch.push_back(p[i]);
   ch.pop_back();
   return ch;
```

8.16 3D Convex Hull

```
double absvol(const Point a,const Point b,const Point c
    ,const Point d){
  return abs(((b-a)^{(c-a)})^*(d-a))/6;
struct convex3D{
static const int maxn=1010;
struct Triangle{
  int a,b,c;
  bool res;
  Triangle(){}
  Triangle(int a,int b,int c,bool res=1):a(a),b(b),c(c)
    ,res(res){}
};
int n,m;
Point p[maxn];
Triangle f[maxn*8];
int id[maxn][maxn];
bool on(Triangle &t,Point &pt){
  return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(pt-p[t.a])>
void meow(int pi,int a,int b){
  int f2=id[a][b];
  if(f[f2].res){
    if(on(f[f2],p[pi]))dfs(pi,f2);
      id[pi][b]=id[a][pi]=id[b][a]=m;
      f[m++]=Triangle(b,a,pi,1);
    }
  }
void dfs(int pi,int now){
  f[now].res=0;
  meow(pi,f[now].b,f[now].a);
  meow(pi,f[now].c,f[now].b);
  meow(pi,f[now].a,f[now].c);
void operator()(){
  if(n<4)return;
  if([&]()->int{
    for(int i=1;i<n;++i){</pre>
      if(abs(p[0]-p[i])>eps){
        swap(p[1],p[i]);
        return 0;
      }
    }
    return 1;
  }())return;
  if([&]()->int{
    for(int i=2;i<n;++i){</pre>
      if(abs((p[0]-p[i])^{(p[1]-p[i]))>eps){
        swap(p[2],p[i]);
        return 0;
```

```
return 1;
  }())return;
  if([&]()->int{
    for(int i=3;i<n;++i){</pre>
      if(abs(((p[1]-p[0])^{p[2]-p[0]))*(p[i]-p[0]))>eps
        swap(p[3],p[i]);
        return 0;
      }
    return 1;
  }())return;
  for(int i=0;i<4;++i){</pre>
    Triangle tmp((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
    if(on(tmp,p[i]))swap(tmp.b,tmp.c);
    id[tmp.a][tmp.b]=id[tmp.b][tmp.c]=id[tmp.c][tmp.a]=
    f[m++]=tmp;
  for(int i=4;i<n;++i){</pre>
    for(int j=0;j<m;++j){</pre>
      if(f[j].res && on(f[j],p[i])){
        dfs(i,j);
        break;
      }
    }
  int mm=m; m=0;
  for(int i=0;i<mm;++i){</pre>
    if(f[i].res)f[m++]=f[i];
bool same(int i,int j){
  return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j]
    a])>eps || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f
    [j].b])>eps || absvol(p[f[i].a],p[f[i].b],p[f[i].c
    ],p[f[j].c])>eps);
int faces(){
  int rt=0;
  for(int i=0;i<m;++i){</pre>
    int iden=1;
    for(int j=0;j<i;++j){</pre>
      if(same(i,j))iden=0;
    rt+=iden;
  }
  return rt;
} tb;
```

8.17 Rotating Caliper

```
struct pnt {
  int x, y
  pnt(): x(0), y(0) {};
  pnt(int xx, int yy): x(xx), y(yy) {};
} p[maxn];
pnt operator-(const pnt &a, const pnt &b) { return pnt(
    b.x - a.x, b.y - a.y); }
int operator^(const pnt &a, const pnt &b) { return a.x
     b.y - a.y * b.x; } //cross
int operator*(const pnt &a, const pnt &b) { return (a -
     b).x * (a - b).x + (a - b).y * (a - b).y; } //
    distance
int tb[maxn], tbz, rsd;
int dist(int n1, int n2){
  return p[n1] * p[n2];
int cross(int t1, int t2, int n1){
 return (p[t2] - p[t1]) ^ (p[n1] - p[t1]);
bool cmpx(const pnt &a, const pnt &b) { return a.x == b
    .x ? a.y < b.y : a.x < b.x; }
void RotatingCaliper() {
 sort(p, p + n, cmpx);
```

```
for (int i = 0; i < n; ++i) {
  while (tbz > 1 && cross(tb[tbz - 2], tb[tbz - 1], i
      ) <= 0) --tbz;
     tb[tbz++] = i;
   rsd = tbz - 1;
  for (int i = n - 2; i >= 0; --i) {
  while (tbz > rsd + 1 && cross(tb[tbz - 2], tb[tbz -
       1], i) <= 0) --tbz;
     tb[tbz++] = i;
   --tbz;
  int lpr = 0, rpr = rsd;
// tb[lpr], tb[rpr]
   while (lpr < rsd || rpr < tbz - 1) {</pre>
     if (lpr < rsd && rpr < tbz - 1) {</pre>
        pnt rvt = p[tb[rpr + 1]] - p[tb[rpr]];
pnt lvt = p[tb[lpr + 1]] - p[tb[lpr]];
        if ((lvt ^{\prime} rvt) < 0) ++\overline{l}pr;
        else ++rpr;
     else if (lpr == rsd) ++rpr;
     else ++lpr;
      }
```

8.18 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c -
                               a:
  double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 \land p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {</pre>
     if (norm2(cent - p[i]) <= r) continue;</pre>
     cent = p[i];
     r = 0.0;
     for (int j = 0; j < i; ++j) {
  if (norm2(cent - p[j]) <= r) continue;</pre>
       cent = (p[i] + p[j]) / 2;
       r = norm2(p[j] - cent);
for (int k = 0; k < j;
                                   ++k) {
          if (norm2(cent - p[k]) <= r) continue;</pre>
          cent = center(p[i], p[j], p[k]);
          r = norm2(p[k] - cent);
       }
    }
  return circle(cent, sqrt(r));
```

8.19 Closest Pair

```
double closest_pair(int l, int r) {
    // p should be sorted increasingly according to the x
        -coordinates.
    if (l == r) return 1e9;
    if (r - l == 1) return dist(p[l], p[r]);
    int m = (l + r) >> 1;
    double d = min(closest_pair(l, m), closest_pair(m +
        1, r));
    vector<int> vec;
    for (int i = m; i >= l && fabs(p[m].x - p[i].x) < d;
        --i) vec.push_back(i);
    for (int i = m + 1; i <= r && fabs(p[m].x - p[i].x) <
        d; ++i) vec.push_back(i);
    sort(vec.begin(), vec.end(), [&](int a, int b) {
        return p[a].y < p[b].y; });
    for (int i = 0; i < vec.size(); ++i) {</pre>
```

```
for (int j = i + 1; j < vec.size() && fabs(p[vec[j
]].y - p[vec[i]].y) < d; ++j) {
    d = min(d, dist(p[vec[i]], p[vec[j]]));
    }
}
return d;
}</pre>
```

9 Miscellaneous / Problems

9.1 Bitwise Hack

9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) {
            x = s - 1 - x;
            y = s - 1 - y;
        }
        swap(x, y);
    }
   return res;
}
```

9.3 Java

```
import java.io.*;
import java.util.*;
import java.lang.*
import java.math.*;
public class filename{
  static Scanner in = new Scanner(System.in);
  public static void main(String[] args) {
    int t = 7122;
    while(in.hasNext()) {
      t = in.nextInt();
       float b = in.nextFloat();
      String str = in.nextLine(); // getline
      String stu = in.next(); // string
    System.out.println("Case #" + t);
System.out.printf("%d\n", 7122);
    int[] c = new int[5];
    int[][] d = {{7,1,2,2},{8,7}};
int g = Integer.parseInt("-123");
    long f = (long)d[0][2];
    List<Integer> l = new ArrayList<>();
    Random rg = new Random();
for (int i = 9; i >= 0; --i) {
      l.add(Integer.valueOf(rg.nextInt(100) + 1));
```

```
1.add(Integer.valueOf((int)(Math.random() * 100)
     + 1));
     Collections.sort(l, new Comparator<Integer>() {
       public int compare(Integer a, Integer b) {
         return a - b;
     });
     for (int i = 0; i < l.size(); ++i) {</pre>
       System.out.print(l.get(i));
     Set<String> s = new HashSet<String>(); // TreeSet
     s.add("jizz");
     System.out.println(s);
     System.out.println(s.contains("jizz"));
     Map<String, Integer> m = new HashMap<String,
     Integer>();
m.put("lol"
                 ', 7122);
     System.out.println(m);
     for(String key: m.keySet()) {
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol")):
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
     System.out.println(Math.acos(-1));
     BigInteger bi = in.nextBigInteger(), bj = new
     BigInteger("-7122"), bk = BigInteger.valueOf(17171)
     bi = bi.add(bj);
     bi = bi.subtract(BigInteger.ONE);
     bi = bi.multiply(bj);
     bi = bi.divide(bj);
     bi = bi.and(bj);
     bi = bi.gcd(bj);
     bi = bi.max(bj);
     bi = bi.pow(10);
     int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
}
```

9.4 Dancing Links

```
namespace dlx {
const int maxn = 1000 * 1000 + 5;
int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn],
   rw[maxn], bt[maxn], s[maxn], head, sz, ans;
lt[i] = i == 0 ? c : i - 1;
    rg[i] = i == c - 1 ? c : i + 1;
    s[i] = 0;
  rg[c] = 0;
  lt[c] = c' - 1;
  up[c] = dn[c] = -1;
  head = c;
  sz = c + 1;
void insert(int r, const vector<int> &col) {
  if (col.empty()) return;
  int f = sz;
for (int i = 0; i < (int)col.size(); ++i) {</pre>
    int c = col[i], v = sz++;
    dn[bt[c]] = v;
    up[v] = bt[c];
    bt[c] = v;
    rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r;
    cl[v] = c;
    ++s[c];
    if (i > 0) lt[v] = v - 1;
```

```
lt[f] = sz - 1;
void remove(int c) {
  lt[rg[c]] = lt[c];
  rg[lt[c]] = rg[c];
  for (int i = dn[c]; i != c; i = dn[i]) {
  for (int j = rg[i]; j != i; j = rg[j]) {
    up[dn[j]] = up[j];
    dn[up[j]] = dn[j];
    --s[c][i]]
       --s[cl[j]];
    }
  }
void restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
    for (int j = lt[i]; j != i; j = lt[j]) {
    ++s[cl[j]];
       up[dn[j]] =
       dn[up[j]] = j;
    }
  lt[rg[c]] = c;
rg[lt[c]] = c;
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i) {
    dn[bt[i]] = i;
    up[i] = bt[i];
  }
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
  int w = c;
  for (int x = c; x != head; x = rg[x]) {
  if (s[x] < s[w]) w = x;</pre>
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
     for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j
     dfs(dep + 1);
     for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j
  restore(w);
int solve() {
 ans = 1e9;
 dfs(0);
 return ans;
```

9.5 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second
     = weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i
     such that cnt[i] == 0
sort(v.begin(), v.end(), [&](int i, int j) {
  if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  djs.save();
  for (int i = l; i <= r; ++i) djs.merge(st[qr[i].first</pre>
    ], ed[qr[i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i])
      djs.merge(st[v[i]], ed[v[i]]);
```

```
}
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[</pre>
    x[i]], ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {
  if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {</pre>
       y.push_back(v[i])
       djs.merge(st[v[i]], ed[v[i]]);
  djs.undo();
}
void solve(int 1, int r, vector<int> v, long long c) {
  if (l == r) {
  cost[qr[l].first] = qr[l].second;
     if (st[qr[l].first] == ed[qr[l].first]) {
      printf("%lld\n", c);
       return;
    int minv = qr[l].second;
for (int i = 0; i < (int)v.size(); ++i) minv = min(
minv, cost[v[i]]);</pre>
    printf("%lld\n", c + minv);
    return;
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \ll r; ++i) {
    cnt[qr[i].first]-
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first
  contract(l, m, lv, x, y);
  long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = 1; i <= m; ++i) {</pre>
    cnt[qr[i].first]--:
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first
  }
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merqe(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = 1; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

9.6 Manhattan Distance MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  }
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
  sort(v.begin(), v.end(), [&](int i, int j) { return x
    [i] == x[j] ? y[i] > y[j] : x[i] > x[j]; });
  int j = 0;
  for (int i = 0; i < n; ++i) {
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] -
    y[v[i]]) - ds.begin() + 1;</pre>
```

```
pair<int, int> q = query(p);
    // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, make_pair(x[v[i]] + y[v[i]], v[i]));
}

void make_graph() {
    solve(n);
    for (int i = 0; i < n; ++i) swap(x[i], y[i]);
    solve(n);
    for (int i = 0; i < n; ++i) x[i] = -x[i];
    solve(n);
    for (int i = 0; i < n; ++i) swap(x[i], y[i]);
    solve(n);
}</pre>
```

9.7 "Dynamic" Kth Element (parallel binary search)

```
struct query { int op, l, r, k, qid; };
// op = 1: insertion (l = pos, r = val)
// op = 2: deletion (l = pos, r = val)
// op = 3: query
void bs(vector<query> &qry, int 1, int r) {
  // answer to queries in ary are from l to r
  if (l == r) {
    for (int i = 0; i < qry.size(); ++i) {</pre>
      if (qry[i].op == 3) ans[qry[i].qid] = 1;
    return:
  if (qry.size() == 0) return;
  int m = 1 + r >> 1;
  for (int i = 0; i < qry.size(); ++i) {</pre>
    if (qry[i].op == 1 \& qry[i].r \Leftarrow m) bit.add(qry[i])
    ].l, 1);
else if (qry[i].op == 2 && qry[i].r <= m) bit.add(
    qry[i].l, -1)
    else if (qry[i].op == 3) tmp[qry[i].qid] += bit.qry
    (qry[i].r) - bit.qry(qry[i].l - 1);
 vector<query> ql, qr;
  for (int i = 0; i < qry.size(); ++i) {</pre>
    if (qry[i].op == 3) {
      if (qry[i].k - tmp[qry[i].qid] > 0) qry[i].k -=
    tmp[qry[i].qid], qr.push_back(qry[i]);
      else ql.push_back(qry[i]);
      tmp[qry[i].qid] = 0;
      continue;
    if (qry[i].r <= m) ql.push_back(qry[i]);</pre>
    else qr.push_back(qry[i]);
  for (int i = 0; i < qry.size(); ++i) {</pre>
    if (qry[i].op == 1 && qry[i].r <= m) bit.add(qry[i</pre>
    [].l, -1);
else if (qry[i].op == 2 && qry[i].r <= m) bit.add(</pre>
 bs(ql, l, m), bs(qr, m + 1, r);
```

9.8 Dynamic Kth Element (persistent segment tree)

```
// segtree: persistant segment tree which supports
    range sum query

void init(int n) {
    segtree::sz = 0;
    bit[0] = segtree::build(0, ds.size());
    for (int i = 1; i <= n; ++i) bit[i] = bit[0];
}

void add(int p, int n, int x, int v) {</pre>
```

```
for (; p \le n; p += p \& -p)
    bit[p] = segtree::modify(0, ds.size(), x, v, bit[p
     ]);
}
vector<int> query(int p) {
  vector<int> z;
  for (; p; p -= p & -p)
    z.push_back(bit[p]);
  return z;
int dfs(int 1, int r, vector<int> lz, vector<int> rz,
  int k) {
if (r - l == 1) return l;
  int ls = 0, rs = 0;
  for (int i = 0; i < lz.size(); ++i) ls += segtree::st
  [segtree::lc[lz[i]]];</pre>
  for (int i = 0; i < rz.size(); ++i) rs += segtree::st
     [segtree::lc[rz[i]]];
  if (rs - ls >= k)
     for (int i = 0; i < lz.size(); ++i) lz[i] = segtree</pre>
     ::lc[lz[i]];
     for (int i = 0; i < rz.size(); ++i) rz[i] = segtree</pre>
     ::lc[rz[i]];
     return dfs(l, (l + r) / 2, lz, rz, k);
  } else {
   for (int i = 0; i < lz.size(); ++i) lz[i] = segtree</pre>
     ::rc[lz[i]];
     for (int i = 0; i < rz.size(); ++i) rz[i] = segtree
     ::rc[rz[i]]
     return dfs((l + r) / 2, r, lz, rz, k - (rs - ls));
  }
}
void solve() {
  init(n);
  for (int i = 1; i <= n; ++i) add(i, n, a[i], 1);
for (int i = 0; i < q; ++i) {
  if (qr[i][0] == 1) {</pre>
       vector<int> lz = query(qr[i][1] - 1);
       vector<int> rz = query(qr[i][2]);
int ans = dfs(0, ds.size(), lz, rz, qr[i][3]);
       printf("%d\n", ds[ans]);
    } else {
       add(qr[i][1], n, a[qr[i][1]], -1);
add(qr[i][1], n, qr[i][2], 1);
       a[qr[i][1]] = qr[i][2];
}
```

9.9 IOI 2016 Alien trick

```
struct result {
  long long m; int v;
  result(): m(0), v(0) {}
  result(long long a, int b): m(a), v(b) {}
  result operator+(const result &r) const { return
    result(m + r.m, v + r.v); }
  bool operator<(const result &r) const { return m == r</pre>
     .m ? v < r.v : m < r.m; }
  bool operator>(const result &r) const { return m == r
     .m ? v > r.v : m > r.m; }
} dp[maxn];
result check(int p);
long long alien() {
  long long c = inf;
  for (int d = 60; d >= 0; --d) {
    if (c - (111 << d) < 0) continue;
    result r = \text{check}(c - (111 \ll d));
if (r.v == k) return r.m - (c - (111 \ll d)) * k;
    if (r.v < k) c -= (111 << d);
  result r = check(c);
  return r.m - c * k;
```