### 1 Basic

#### 1.1 vimrc

```
| se nu rnu bs=2 ru mouse=a cin et ts=4 sw=4 sts=4 syn on | colo desert | filetype indent on | inoremap {<CR> {<CR>}<Esc>0
```

### 1.2 Fast Integer Input

```
inline int gtx() {
  const int N = 4096;
  static char buffer[N];
  static char *p = buffer, *end = buffer;
  if (p == end) {
    if ((end = buffer + fread(buffer, 1, N, stdin)) == buffer)
     return EOF;
    p = buffer;
  return *p++;
template <typename T>
inline bool rit(T& x) {
  char c = 0; bool flag = false;
while (c = getchar(), (c < '0' && c != '-') || c > '9') if (c
  == -1) return false;
c == '-' ? (flag = true, x = 0) : (x = c - '0');
  while (c = getchar(), c >= '0' && c <= '9') x = x * 10 + c -
  if (flag) x = -x;
  return true;
```

#### 1.3 Increase stack size

```
const int size = 256 << 20;
register long rsp asm("rsp");
char *p = (char*)malloc(size) + size, *bak = (char*)rsp;
__asm__("movq %0, %%rsp\n"::"r"(p));
// main
__asm__("movq %0, %%rsp\n"::"r"(bak));</pre>
```

#### 1.4 Pragma optimization

# 2 Flows, Matching

#### 2.1 Dinic's Algorithm

```
struct dinic {
 static const int inf = 1e9;
  struct edge {
    int to, cap, rev;
    edge(int d, int c, int r): to(d), cap(c), rev(r) {}
 vector<edge> g[maxn];
 int qu[maxn], ql, qr;
  int lev[maxn];
 void init() {
   for (int i = 0; i < maxn; ++i)
      g[i].clear();
  void add_edge(int a, int b, int c) {
   g[a].emplace_back(b, c, g[b].size() - 0);
   g[b].emplace_back(a, 0, g[a].size() - 1);
 bool bfs(int s, int t) {
   memset(lev, -1, sizeof(lev));
   lev[s] = 0;
   ql = qr = 0;
   qu[qr++] = s;
   while (ql < qr) {</pre>
      int x = qu[ql++]
      for (edge &e : g[x]) if (lev[e.to] == -1 && e.cap > 0) {
       lev[e.to] = lev[x] + 1;
        qu[qr++] = e.to;
   }
```

```
return lev[t] != -1;
  int dfs(int x, int t, int flow) {
     if (x == t) return flow;
     int res = 0;
     for (edge &e : g[x]) if (e.cap > 0 && lev[e.to] == lev[x] +
      1) {
       int f = dfs(e.to, t, min(e.cap, flow - res));
       e.cap -= f;
       g[e.to][e.rev].cap += f;
     if (res == 0) lev[x] = -1;
     return res;
  int operator()(int s, int t) {
     int flow = 0;
     for (; bfs(s, t); flow += dfs(s, t, inf));
     return flow;
  }
};
      Minimum-cost flow
struct mincost {
  struct edge {
     int dest, cap, w, rev;
     edge(int a, int b, int c, int d): dest(a), cap(b), w(c),
     rev(d) {}
  vector<edge> g[maxn];
  int d[maxn], p[maxn], ed[maxn];
  bool inq[maxn];
  void init() {
    for (int i = 0; i < maxn; ++i) g[i].clear();</pre>
  void add_edge(int a, int b, int c, int d) {
    g[a].emplace_back(b, c, +d, g[b].size() - 0);
     g[b].emplace_back(a, 0, -d, g[a].size() - 1);
  bool spfa(int s, int t, int &f, int &c) {
     for (int i = 0; i < maxn; ++i) {
       d[i] = inf;
       p[i] = ed[i] = -1;
       inq[i] = false;
     d[s] = 0;
     queue<int> q;
     q.push(s);
     while (q.size()) {
       int x = q.front(); q.pop();
       inq[x] = false;
       for (int i = 0; i < g[x].size(); ++i) {</pre>
         edge &e = g[x][i];
         if (e.cap > 0 \& d[e.dest] > d[x] + e.w) {
           d[e.dest] = d[x] + e.w;
           p[e.dest] = x;
           ed[e.dest] = i;
           if (!inq[e.dest]) q.push(e.dest), inq[e.dest] = true;
        }
      }
     if (d[t] == inf) return false;
     int dlt = inf;
     for (int x = t; x != s; x = p[x]) dlt = min(dlt, g[p[x]][ed
      [x]].cap);
     for (int x = t; x != s; x = p[x]) {
       edge &e = g[p[x]][ed[x]];
e.cap -= dlt;
       g[e.dest][e.rev].cap += dlt;
     f += dlt; c += d[t] * dlt;
     return true;
  pair<int, int> operator()(int s, int t) {
    int f = 0, c = 0;
     while (spfa(s, t, f, c));
```

#### 2.3 Gomory-Hu Tree

return make\_pair(f, c);

```
int g[maxn];
vector<edge> GomoryHu(int n){
  vector<edge> rt;
  for(int i=1;i<=n;++i)g[i]=1;</pre>
```

```
for(int i=2;i<=n;++i){
   int t=g[i];
   flow.reset(); // clear flows on all edge
   rt.push_back({i,t,flow(i,t)});
   flow.walk(i); // bfs points that connected to i (use edges
   not fully flow)
   for(int j=i+1;j<=n;++j){
      if(g[j]==t && flow.connect(j))g[j]=i; // check if i can
      reach j
   }
}
return rt;
}</pre>
```

#### 2.4 Stoer-Wagner Minimum Cut

```
int w[kN][kN], g[kN], del[kN], v[kN];
void AddEdge(int x, int y, int c) {
  w[x][y] += c;
  w[y][x] += c;
pair<int, int> Phase(int n) {
  fill(v, v + n, 0), fill(g, g + n, 0);
  int s = -1, t = -1;
  while (true) {
    int c = -1;
    for (int i = 0; i < n; ++i) {
       if (del[i] || v[i]) continue;
       if (c == -1 || g[i] > g[c]) c = i;
    if (c == -1) break;
    v[c] = 1, s = t, t = c;
for (int i = 0; i < n; ++i) {
   if (del[i] || v[i]) continue;</pre>
      g[i] += w[c][i];
  }
  return make_pair(s, t);
int GlobalMinCut(int n) {
  int cut = kInf;
  fill(del, 0, sizeof(del));
  for (int i = 0; i < n - 1; ++i) {
    int s, t; tie(s, t) = Phase(n);
    del[t] = 1, cut = min(cut, g[t]);
    for (int j = 0; j < n; ++j) {
      w[s][j] += w[t][j];
w[j][s] += w[j][t];
    }
  return cut;
```

#### 2.5 Kuhn-Munkres Algorithm

```
int64_t KuhnMunkres(vector<vector<int>>> W) {
  int N = W.size();
  vector<int> fl(N, -1), fr(N, -1), hr(N), hl(N);
  for (int i = 0; i < N; ++i) {
   hl[i] = *max_element(W[i].begin(), W[i].end());
  auto Bfs = [&](int s) {
    vector<int> slk(N, kInf), pre(N);
    vector<bool> vl(N, false), vr(N, false);
    queue<int> que;
    que.push(s);
    vr[s] = true;
    auto Check = [\&](int x) \rightarrow bool {
      if (vl[x] = true, fl[x] != -1) {
        que.push(fl[x]);
        return vr[fl[x]] = true;
      while (x != -1) swap(x, fr[fl[x] = pre[x]]);
      return false;
    while (true) {
      while (!que.empty()) {
        int y = que.front(); que.pop();
        for (int x = 0, d = 0; x < N; ++x) {
          if (!vl[x] \&\& slk[x] >= (d = hl[x] + hr[y] - W[x][y])
            if (pre[x] = y, d) slk[x] = d;
            else if (!Check(x)) return;
        }
```

```
int d = kInf;
       for (int x = 0; x < N; ++x) {
         if (!vl[x] \&\& d > slk[x]) d = slk[x];
       for (int x = 0; x < N; ++x) {
         if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
       for (int x = 0; x < N; ++x) {
         if (!vl[x] && !slk[x] && !Check(x)) return;
    }
   };
   for (int i = 0; i < N; ++i) Bfs(i);
   int64_t res = 0;
   for (int i = 0; i < N; ++i) res += W[i][fl[i]];</pre>
   return res:
1 }
```

### 2.6 Maximum Matching on General Graph

```
namespace matching {
int fa[kN], pre[kN], match[kN], s[kN], v[kN];
vector<int> g[kN];
queue<int> q;
void Init(int n) {
  for (int i = 0; i <= n; ++i) match[i] = pre[i] = n;
for (int i = 0; i < n; ++i) g[i].clear();</pre>
void AddEdge(int u, int v) {
  g[u].push_back(v);
   g[v].push_back(u);
int Find(int u) {
  return u == fa[u] ? u : fa[u] = Find(fa[u]);
int LCA(int x, int y, int n) {
   static int tk = 0;
  tk++;
   x = Find(x), y = Find(y);
   for (; ; swap(x, y)) {
  if (x != n) {
       if (v[x] == tk) return x;
       v[x] = tk;
       x = Find(pre[match[x]]);
     }
  }
}
void Blossom(int x, int y, int l) {
  while (Find(x) != 1) {
    pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
     if (fa[x] == x) fa[x] = 1;
     if (fa[y] == y) fa[y] = 1;
     x = pre[y];
  }
bool Bfs(int r, int n) {
   for (int i = 0; i \le n; ++i) fa[i] = i, s[i] = -1;
   while (!q.empty()) q.pop();
   q.push(r);
   s[r] = 0;
   while (!q.empty()) {
     int x = q.front(); q.pop();
     for (int u : g[x]) {
       if (s[u] == -1) {
         pre[u] = x, s[u] = 1;
         if (match[u] == n) {
            for (int a = u, b = x, last; b != n; a = last, b =
      pre[a])
              last = match[b], match[b] = a, match[a] = b;
           return true;
         q.push(match[u]);
         s[match[u]] = 0;
       } else if (!s[u] && Find(u) != Find(x)) {
         int l = LCA(u, x, n);
Blossom(x, u, l);
         Blossom(u, x, 1);
    }
  }
  return false;
int Solve(int n) {
```

```
int res = 0;
for (int x = 0; x < n; ++x) {
   if (match[x] == n) res += Bfs(x, n);
}
return res;
}}</pre>
```

# 2.7 Maximum Weighted Matching on General Graph

```
struct WeightGraph {
 static const int inf = INT_MAX;
 static const int maxn = 514;
  struct edge {
   int u, v, w;
   edge(){}
   edge(int u, int v, int w): u(u), v(v), w(w) {}
 int n, n_x;
 edge g[maxn * 2][maxn * 2];
  int lab[maxn * 2];
 int match[maxn * 2], slack[maxn * 2], st[maxn * 2], pa[maxn *
 int flo_from[maxn * 2][maxn + 1], S[maxn * 2], vis[maxn * 2];
  vector<int> flo[maxn * 2];
 queue<int> q;
 int e_delta(const edge &e) { return lab[e.u] + lab[e.v] - g[e
     .u][e.v].w * 2; }
  void update_slack(int u, int x) { if (!slack[x] || e_delta(g[
    u][x]) < e_delta(g[slack[x]][x])) slack[x] = u; }
  void set_slack(int x) {
   slack[x] = 0;
    for (int u = 1; u \le n; ++u)
      if (g[u][x].w > 0 \&\& st[u] != x \&\& S[st[u]] == 0)
        update_slack(u, x);
  void q_push(int x) {
    if (x \le n) q.push(x);
   else for (size_t i = 0; i < flo[x].size(); i++) q_push(flo[</pre>
    x][i]);
 void set_st(int x, int b) {
   st[x] = b;
    if (x > n) for (size_t i = 0; i < flo[x].size(); ++i)
    set_st(flo[x][i], b);
  int get_pr(int b, int xr) {
    int pr = find(flo[b].begin(), flo[b].end(), xr) - flo[b].
    begin();
    if (pr % 2 == 1) {
      reverse(flo[b].begin() + 1, flo[b].end());
      return (int)flo[b].size() - pr;
   return pr;
 }
  void set_match(int u, int v) {
   match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    edge e = g[u][v];
    int xr = flo_from[u][e.u], pr = get_pr(u, xr)
    for (int i = 0; i < pr; ++i) set_match(flo[u][i], flo[u][i</pre>
    ^ 1]);
   set_match(xr, v);
   rotate(flo[u].begin(), flo[u].begin() + pr, flo[u].end());
  void augment(int u, int v) {
   for (; ; ) {
      int xnv = st[match[u]];
      set_match(u, v);
      if (!xnv) return;
      set_match(xnv, st[pa[xnv]]);
      u = st[pa[xnv]], v = xnv;
   }
 }
 int get_lca(int u, int v) {
   static int t = 0;
    for (++t; u || v; swap(u, v)) {
      if (u == 0) continue;
      if (vis[u] == t) return u;
     vis[u] = t;
      u = st[match[u]];
      if (u) u = st[pa[u]];
   return 0;
 void add_blossom(int u, int lca, int v) {
```

```
int b = n + 1;
  while (b <= n_x && st[b]) ++b;
  if (b > n_x) ++n_x;
  lab[b] = 0, S[b] = 0;
  match[b] = match[lca];
  flo[b].clear();
  flo[b].push_back(lca);
  for (int x = u, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y);
  reverse(flo[b].begin() + 1, flo[b].end())
  for (int x = v, y; x != lca; x = st[pa[y]])
    flo[b].push\_back(x), flo[b].push\_back(y = st[match[x]]),
   q_push(y)
  set_st(b, b);
  for (int x = 1; x <= n_x; ++x) g[b][x].w = g[x][b].w = 0; for (int x = 1; x <= n; ++x) flo_from[b][x] = 0;
  for (size_t i = 0; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    for (int x = 1; x <= n_x; ++x)
      if (g[b][x].w == 0 \mid \mid e_{delta}(g[xs][x]) < e_{delta}(g[b][x])
   (([x
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    for (int x = 1; x <= n; ++x)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slack(b);
}
void expand_blossom(int b) {
  for (size_t i = 0; i < flo[b].size(); ++i)</pre>
    set_st(flo[b][i], flo[b][i]);
  int xr = flo_from[b][g[b][pa[b]].u], pr = get_pr(b, xr);
  for (int i = 0; i < pr; i += 2) {
    int xs = flo[b][i], xns = flo[b][i + 1];
    pa[xs] = g[xns][xs].u;
    S[xs] = 1, S[xns] = 0;
slack[xs] = 0, set_slack(xns);
    q_push(xns);
  S[xr] = 1, pa[xr] = pa[b];
  for (size_t i = pr + 1; i < flo[b].size(); ++i) {</pre>
    int xs = flo[b][i];
    S[xs] = -1, set_slack(xs);
  st[b] = 0;
bool on_found_edge(const edge &e) {
  int u = st[e.u], v = st[e.v];
  if (S[v] == -1) {
    pa[v] = e.u, S[v] = 1;
    int nu = st[match[v]];
    slack[v] = slack[nu] = 0;
    S[nu] = 0, q_push(nu);
  } else if (S[v] == 0) {
    int lca = get_lca(u, v);
    if (!lca) return augment(u,v), augment(v,u), true;
    else add_blossom(u, lca, v);
  return false:
bool matching() {
  memset(S + 1, -1, sizeof(int) * n_x);
  memset(slack + 1, 0, sizeof(int) * n_x);
  q = queue<int>();
  for (int x = 1; x <= n_x; ++x)
    if (st[x] == x \&\& !match[x]) pa[x] = 0, S[x] = 0, q_push(
  if (q.empty()) return false;
  for (; ; ) {
  while (q.size()) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      for (int v = 1; v \le n; ++v)
        if (g[u][v].w > 0 \& st[u] != st[v]) {
           if (e_delta(g[u][v]) == 0) {
             if (on_found_edge(g[u][v])) return true;
          } else update_slack(u, st[v]);
    int d = inf;
    for (int b = n + 1; b \le n_x; ++b)
      if (st[b] == b \&\& S[b] == 1) d = min(d, lab[b] / 2);
    for (int x = 1; x <= n_x; ++x)
      if (st[x] == x \&\& slack[x]) {
        if (S[x] == -1) d = min(d, e_delta(g[slack[x]][x]));
```

```
else if (S[x] == 0) d = min(d, e_delta(g[slack[x]][x
     ]) / 2);
         }
       for (int u = 1; u \le n; ++u) {
         if (S[st[u]] == 0) {
           if (lab[u] <= d) return 0;</pre>
           lab[u] -= d;
         } else if (S[st[u]] == 1) lab[u] += d;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b) {
  if (S[st[b]] == 0) lab[b] += d * 2;
           else if (S[st[b]] == 1) lab[b] -= d * 2;
       q = queue<int>();
       for (int x = 1; x <= n_x; ++x)
if (st[x] == x && slack[x] && st[slack[x]] != x &&
     e_delta(g[slack[x]][x]) == 0)
           if (on_found_edge(g[slack[x]][x])) return true;
       for (int b = n + 1; b \le n_x; ++b)
         if (st[b] == b &\& S[b] == 1 &\& lab[b] == 0)
     expand_blossom(b);
    }
    return false;
  pair<long long, int> solve() {
  memset(match + 1, 0, sizeof(int) * n);
    n x = n:
    int n_matches = 0;
    long long tot_weight = 0;
     for (int u = 0; u <= n; ++u) st[u] = u, flo[u].clear();</pre>
     int w_max = 0;
    for (int u = 1; u \le n; ++u)
       for (int v = 1; v <= n; ++v) {
  flo_from[u][v] = (u == v ? u : 0);</pre>
         w_max = max(w_max, g[u][v].w);
    for (int u = 1; u <= n; ++u) lab[u] = w_max;</pre>
    while (matching()) ++n_matches;
     for (int u = 1; u \le n; ++u)
       if (match[u] && match[u] < u)</pre>
         tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
  void add_edge(int ui, int vi, int wi) { g[ui][vi].w = g[vi][
     ui].w = wi; }
  void init(int _n) {
    n = _n;
    for (int u = 1; u \le n; ++u)
       for (int v = 1; v \le n; ++v)
         g[u][v] = edge(u, v, 0);
  }
};
```

#### 2.8Minimum Cost Circulation

```
struct Edge { int to, cap, rev, cost; };
vector<Edge> g[kN];
int dist[kN], pv[kN], ed[kN];
bool mark[kN];
int NegativeCycle(int n) {
 memset(mark, false, sizeof(mark));
memset(dist, 0, sizeof(dist));
  int upd = -1;
  for (int i = 0; i \le n; ++i) {
    for (int j = 0; j < n; ++j) {
      int idx = 0:
      for (auto &e : g[j]) {
        if (e.cap > 0 && dist[e.to] > dist[j] + e.cost) {
          dist[e.to] = dist[j] + e.cost;
          pv[e.to] = j, ed[e.to] = idx;
          if (i == n) {
            upd = j;
            while (!mark[upd]) mark[upd] = true, upd = pv[upd];
            return upd;
          }
        idx++;
     }
   }
  return -1;
int Solve(int n) {
  int rt = -1, ans = 0;
  while ((rt = NegativeCycle(n)) >= 0) {
    memset(mark, false, sizeof(mark));
```

```
vector<pair<int, int>> cyc;
     while (!mark[rt]) {
       cyc.emplace_back(pv[rt], ed[rt]);
       mark[rt] = true;
       rt = pv[rt];
     }
     reverse(cyc.begin(), cyc.end());
     int cap = kInf;
for (auto &i : cyc) {
       auto &e = g[i.first][i.second];
       cap = min(cap, e.cap);
     for (auto &i : cyc) {
       auto &e = g[i.first][i.second];
       e.cap -= cap;
       g[e.to][e.rev].cap += cap;
ans += e.cost * cap;
   return ans;
}
```

#### 2.9 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x, y, l, u), connect  $x \to y$  with capacity u l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v) > 0, connect  $S \to v$  with capacity in(v), otherwise, connect  $v \to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution.
    - Otherwise, the maximum flow from s to t is the answer.

      To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to Tbe f'. If  $f + f' \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e + f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
  - 1. Redirect every edge:  $y \to x$  if  $(x, y) \in M$ ,  $x \to y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- $\bullet \quad \text{Maximum density induced subgraph} \\$ 
  - 1. Binary search on answer, suppose we're checking answer T
  - 2. Construct a max flow model, let K be the sum of all weights
  - 3. Connect source  $s \to v, v \in G$  with capacity K
  - 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
  - 5. For  $v \in G$ , connect it with sink  $v \to t$  with capacity K + 2T - $(\sum_{e \in E(v)} w(e)) - 2w(v)$
  - 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- · Project selection problem
  - 1. If  $p_v > 0$ , create edge (s, v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$
  - Create edge (u, v) with capacity w with w being the cost of choosing u without choosing v.
  - 3. The mincut is equivalent to the maximum profit of a subset of projects.
- $\bullet$  0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity  $c_x$  and create edge (s,y) with capacity  $c_y$ .
- 2. Create edge (x, y) with capacity  $c_{xy}$ .
- 3. Create edge (x, y) and edge (x', y') with capacity  $c_{xyx'y'}$ .

#### 3 Data Structure

#### <ext/pbds> 3.1

```
#include <bits/stdc++.h>
#include <bits/extc++.h>
#include <ext/rope>
using namespace __gnu_pbds;
using namespace __gnu_cxx;
#include <ext/pb_ds/assoc_container.hpp>
typedef tree<int, null_type, std::less<int>, rb_tree_tag,
     tree_order_statistics_node_update> tree_set;
typedef cc_hash_table<int, int> umap;
typedef priority_queue<int> heap;
int main() {
 tree set s:
  s.insert(71); s.insert(22);
  assert(*s.find_by\_order(0) == 22); assert(*s.find_by\_order(1))
      == 71):
  assert(s.order\_of\_key(22) == 0); assert(s.order\_of\_key(71) == 0);
      1):
  s.erase(22);
  assert(*s.find_by\_order(0) == 71); assert(s.order\_of\_key(71))
     == 0):
  // mergable heap
  heap a, b; a.join(b);
  // persistant
  rope<char> r[2];
  r[1] = r[0];
  std::string st = "abc";
  r[1].insert(0, st.c_str());
  r[1].erase(1, 1);
  std::cout << r[1].substr(0, 2) << std::endl;</pre>
```

#### Li Chao Tree

```
namespace lichao {
struct line {
   long long a, b;
  line(): a(0), b(0) {}
  line(long long a, long long b): a(a), b(b) {}
   long long operator()(int x) const { return a * x + b; }
line st[maxc * 4];
int sz, lc[maxc * 4], rc[maxc * 4];
int gnode() {
  st[sz] = line(1e9, 1e9);
   lc[sz] = -1, rc[sz] = -1;
  return sz++;
void init() {
  sz = 0;
void add(int l, int r, line tl, int o) {
  bool lcp = st[o](l) > tl(l);
   bool mcp = st[o]((l + r) / 2) > tl((l + r) / 2);
  if (mcp) swap(st[o], tl);
   if (r - l == 1) return;
  if (lcp != mcp) {
     if (lc[o] == -1) lc[o] = gnode();
     add(1, (1 + r) / 2, tl, lc[o]);
  } else {
     if (rc[o] == -1) rc[o] = gnode();
add((l + r) / 2, r, tl, rc[o]);
long long query(int l, int r, int x, int o) {
   if (r - l == 1) return st[o](x);
  if (x < (l + r) / 2) {
     if (lc[o] == -1) return st[o](x);
     return min(st[o](x), query(l, (l + r) / 2, x, lc[o]));
     if (rc[o] == -1) return st[o](x);
     return min(st[o](x), query((l + r) / 2, r, x, rc[o]));
| }}
```

#### 3.3Link-Cut Tree

```
struct node {
  node *ch[2], *fa, *pfa;
  int sum, v, rev, id;
```

```
node(int s, int id): id(id), v(s), sum(s), rev(0), fa(nullptr
     ), pfa(nullptr) {
    ch[0] = nullptr;
    ch[1] = nullptr;
  int relation() {
    return this == fa->ch[0] ? 0 : 1;
  void push() {
    if (!rev) return;
    swap(ch[0], ch[1]);
    if (ch[0]) ch[0]->rev ^= 1;
    if (ch[1]) ch[1]->rev ^= 1;
    rev = 0;
  void pull() {
    sum = v
    if (ch[0]) sum += ch[0]->sum;
    if (ch[1]) sum += ch[1]->sum;
  void rotate() -
    if (fa->fa) fa->fa->push();
    fa->push(), push(), swap(pfa, fa->pfa);
    int d = relation();
    node *t = fa;
    if (t->fa) t->fa->ch[t->relation()] = this;
    fa = t->fa, t->ch[d] = ch[d \land 1];
    if (ch[d ^ 1]) ch[d ^ 1]->fa = t;
ch[d ^ 1] = t, t->fa = this;
    t->pull(), pull();
  void splay() -
    while (fa) {
      if (!fa->fa) {
        rotate();
        continue:
      fa->fa->push(), fa->push();
      if (relation() == fa->relation()) fa->rotate();
      else rotate(), rotate();
    }
  }
  void evert() { access(), splay(), rev ^= 1; }
  void expose() {
    splay(), push();
    if (ch[1]) {
      ch[1]->fa = nullptr, ch[1]->pfa = this;
      ch[1] = nullptr, pull();
  }
  bool splice() {
    splay();
    if (!pfa) return false;
    pfa->expose(), pfa->ch[1] = this, fa = pfa;
    pfa = nullptr, fa->pull();
    return true:
  void access() {
    expose();
    while (splice());
  int query() { return sum; }
};
namespace lct {
node *sp[maxn];
void make(int u, int v) {
 // create node with id u and value v
  sp[u] = new node(v, u);
}
void link(int u, int v) {
  // u become v's parent
  sp[v]->evert();
  sp[v]->pfa = sp[u];
}
void cut(int u, int v) {
  // u was v's parent
  sp[u]->evert();
  sp[v]->access(), sp[v]->splay(), sp[v]->push();
  sp[v]->ch[0]->fa = nullptr;
  sp[v] -> ch[0] = nullptr;
  sp[v]->pull();
void modify(int u, int v) {
  sp[u]->splay();
  sp[u]->v = v
  sp[u]->pull();
```

```
}
int query(int u, int v) {
    sp[u]->evert(), sp[v]->access(), sp[v]->splay();
    return sp[v]->query();
}
int find(int u) {
    sp[u]->access();
    sp[u]->splay();
    node *p = sp[u];
    while (true) {
        p->push();
        if (p->ch[0]) p = p->ch[0];
        else break;
    }
    return p->id;
}
```

# 4 Graph

### 4.1 Heavy-Light Decomposition

```
void dfs(int x, int p) {
  dep[x] = \sim p ? dep[p] + 1 : dep[x];
  sz[x] = 1;
  to[x] = -1;
  fa[x] = p;
  for (const int &u : g[x]) {
    if (u == p) continue;
    dfs(u, x);
    sz[x] += sz[u];
    if (to[x] == -1 \mid | sz[to[x]] < sz[u]) to[x] = u;
void hld(int x, int t) {
 static int tk = 0;
  fr[x] = t;
  dfn[x] = tk++;
  if (!~to[x]) return;
 hld(to[x], t);
for (const int &u : g[x]) {
    if (u == fa[x] \mid \mid u == to[x]) continue;
    hld(u, u);
}
vector<pair<int, int>> get(int x, int y) {
 int fx = fr[x], fy = fr[y];
  vector<pair<int, int>> res;
  while (fx != fy) {
    if (dep[fx] < dep[fy]) {</pre>
      swap(fx, fy);
      swap(x, y);
    res.emplace_back(dfn[fx], dfn[x] + 1);
    x = fa[fx];
    fx = fr[x];
 }
  res.emplace_back(min(dfn[x], dfn[y]), max(dfn[x], dfn[y]) +
  int lca = (dep[x] < dep[y] ? x : y);
  return res;
```

#### 4.2 Centroid Decomposition

```
void get_center(int now) {
   v[now] = true; vtx.push_back(now);
   sz[now] = 1; mx[now] = 0;
   for (int u : G[now]) if (!v[u]) {
      get_center(u);
      mx[now] = max(mx[now], sz[u]);
      sz[now] += sz[u];
   }
}
void get_dis(int now, int d, int len) {
   dis[d][now] = cnt;
   v[now] = true;
   for (auto u : G[now]) if (!v[u.first]) {
      get_dis(u, d, len + u.second);
   }
}
void dfs(int now, int fa, int d) {
   get_center(now);
   int c = -1;
   for (int i : vtx) {
```

```
if (max(mx[i], (int)vtx.size() - sz[i]) <= (int)vtx.size()
    / 2) c = i;
    v[i] = false;
}
get_dis(c, d, 0);
for (int i : vtx) v[i] = false;
    v[c] = true; vtx.clear();
dep[c] = d; p[c] = fa;
for (auto u : G[c]) if (u.first != fa && !v[u.first]) {
    dfs(u.first, c, d + 1);
}
</pre>
```

### 4.3 Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
long long d[1003][1003],dp[1003][1003];
pair<long long,long long> MMWC(){
  memset(dp,0x3f,sizeof(dp))
  for(int i=1;i<=n;++i)dp[0][i]=0;</pre>
  for(int i=1;i<=n;++i){</pre>
    for(int j=1;j<=n;++j){</pre>
      for(int k=1;k<=n;++k){</pre>
        dp[i][k]=min(dp[i-1][j]+d[j][k],dp[i][k]);
    }
  long long au=1ll<<31,ad=1;</pre>
  for(int i=1;i<=n;++i){</pre>
    long long u=0,d=1;
    for(int j=n-1;j>=0;--j){
      if((dp[n][i]-dp[j][i])*d>u*(n-j)){
        u=dp[n][i]-dp[j][i];
        d=n-j;
      }
    if(u*ad<au*d)au=u,ad=d;
  long long g=__gcd(au,ad);
  return make_pair(au/g,ad/g);
```

#### 4.4 Minimum Steiner Tree

```
namespace steiner {
// Minimum Steiner Tree - O(N * 3^K + N^2 * 2^K)
// z[i] = the weight of the i-th vertex
const int maxn = 64, maxk = 10;
const int inf = 1e9;
int w[maxn][maxn], z[maxn], dp[1 << maxk][maxn], off[maxn];</pre>
void init(int n) {
  for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) w[i][j] = inf;
     z[i] = 0;
     w[i][i] = 0;
void add_edge(int x, int y, int d) {
  w[x][y] = min(w[x][y], d);
   w[y][x] = min(w[y][x], d);
void build(int n) {
  for (int i = 0; i < n; ++i) {
  for (int j = 0; j < n; ++j) {</pre>
       w[i][j] += z[i];
       if (i != j) w[i][j] += z[j];
   for (int k = 0; k < n; ++k) {
     for (int i = 0; i < n; ++i) {
       for (int j = 0; j < n; ++j) w[i][j] = min(w[i][j], w[i][k]
      ] + w[k][j] - z[k];
     }
  }
int solve(int n, vector<int> mark) {
  build(n);
  int k = (int)mark.size();
   assert(k < maxk);</pre>
   for (int s = 0; s < (1 << k); ++s) {
     for (int i = 0; i < n; ++i) dp[s][i] = inf;
   for (int i = 0; i < n; ++i) dp[0][i] = 0;
   for (int s = 1; s < (1 << k); ++s) {
```

```
if (__builtin_popcount(s) == 1) {
       int x = __builtin_ctz(s);
for (int i = 0; i < n; ++i) dp[s][i] = w[mark[x]][i];</pre>
       continue;
    for (int i = 0; i < n; ++i) {
       for (int sub = s & (s - 1); sub; sub = s & (sub - 1)) {
         dp[s][i] = min(dp[s][i], dp[sub][i] + dp[s \land sub][i] -
     z[i]);
      }
    }
    for (int i = 0; i < n; ++i) {
      off[i] = inf;
       for (int j = 0; j < n; ++j) off[i] = min(off[i], dp[s][j]
      + w[j][i] - z[j]);
    for (int i = 0; i < n; ++i) dp[s][i] = min(dp[s][i], off[i
     ]);
  int res = inf;
  for (int i = 0; i < n; ++i) res = min(res, dp[(1 << k) - 1][i
     ]);
  return res:
}}
```

#### Directed Minimum Spanning Tree

```
template <typename T> struct DMST {
 T g[maxn][maxn], fw[maxn];
 int n, fr[maxn];
 bool vis[maxn], inc[maxn];
  void clear() {
    for(int i = 0; i < maxn; ++i) {</pre>
      for(int j = 0; j < maxn; ++j) g[i][j] = inf;
      vis[i] = inc[i] = false;
   }
 void addedge(int u, int v, T w) {
    g[u][v] = min(g[u][v], w);
 T operator()(int root, int _n) {
    if (dfs(root) != n) return -1;
    T ans = 0;
    while (true) {
      for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;</pre>
      for (int i = 1; i <= n; ++i) if (!inc[i]) {
        for (int j = 1; j \le n; ++j) {
          if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
            fw[i] = g[j][i];
            fr[i] = j;
          }
        }
      for (int i = 1; i <= n; ++i) if (i != root && !inc[i]) {</pre>
        int j = i, c = 0;
        while (j != root \&\& fr[j] != i \&\& c <= n) ++c, j = fr[j]
        if (j == root || c > n) continue;
        else { x = i; break; }
      if (!~x) {
        for (int i = 1; i <= n; ++i) if (i != root && !inc[i])</pre>
    ans += fw[i];
        return ans;
      for (int i = 1; i <= n; ++i) vis[i] = false;
      do { ans += fw[y]; y = fr[y]; vis[y] = inc[y] = true; }
    while (y != x);
      inc[x] = false;
      for (int k = 1; k \le n; ++k) if (vis[k]) {
        for (int j = 1; j <= n; ++j) if (!vis[j]) {</pre>
          if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
          if (g[j][k] < \inf \&\& g[j][k] - fw[k] < g[j][x]) g[j][
    x] = g[j][k] - fw[k];
        }
     }
    return ans;
 int dfs(int now) {
    int r = 1;
    vis[now] = true;
    for (int i = 1; i \leftarrow n; ++i) if (g[now][i] \leftarrow inf \&\& !vis[i]
    ]) r += dfs(i);
```

```
};
4.6 Maximum Clique
```

return r;

}

```
| struct MaxClique {
   // change to bitset for n > 64.
   int n, deg[maxn];
   uint64_t adj[maxn], ans;
   vector<pair<int, int>> edge;
   void init(int n_) {
     n = n_{-}
     fill(adj, adj + n, 0ull);
     fill(deg, deg + n, 0);
     edge.clear();
   void add_edge(int u, int v) {
     edge.emplace_back(u, v);
     ++deg[u], ++deg[v];
   vector<int> operator()() {
     vector<int> ord(n);
     iota(ord.begin(), ord.end(), 0);
     sort(ord.begin(), ord.end(), [&](int u, int v) { return deg
      [u] < deg[v]; });
     vector<int> id(n);
     for (int i = 0; i < n; ++i) id[ord[i]] = i;</pre>
     for (auto e : edge) {
       int u = id[e.first], v = id[e.second];
       adj[u] = (1ull \ll v);
       adj[v] = (1ull \ll u);
    uint64_t r = 0, p = (1ull << n) - 1;
     ans = 0;
    dfs(r, p);
     vector<int> res;
     for (int i = 0; i < n; ++i) {
       if (ans >> i & 1) res.push_back(ord[i]);
     return res;
  }
#define pcount __builtin_popcountll
   void dfs(uint64_t r, uint64_t p) {
     if (p == 0) {
       if (pcount(r) > pcount(ans)) ans = r;
       return;
     if (pcount(r | p) <= pcount(ans)) return;</pre>
     int x = __builtin_ctzll(p & -p);
     uint64_t c = p & \simadj[x];
     while (c > 0) {
       // bitset._Find_first(); bitset._Find_next();
       x = __builtin_ctzll(c & -c);
r |= (1ull << x);</pre>
       dfs(r, p & adj[x]);
       r &= ~(1ull << x);
       p &= ~(1ull << x);
       c ^= (1ull << x);
  }
};
```

#### Tarjan's Algorithm 4.7

```
void dfs(int x, int p) {
  dfn[x] = low[x] = tk++;
  int ch = 0;
  st.push(x); // bridge
  for (auto e : g[x]) if (e.first != p) {
    if (!ins[e.second]) { // articulation point
      st.push(e.second);
      ins[e.second] = true;
    if (~dfn[e.first]) {
      low[x] = min(low[x], dfn[e.first]);
      continue:
    dfs(u.first, x);
    if (low[u.first] >= low[x]) { // articulation point
      cut[x] = true;
      while (true) {
        int z = st.top(); st.pop();
        bcc[z] = sz;
        if (z == e.second) break;
      SZ++;
```

#### 4.8 Dominator Tree

```
namespace dominator {
vector<int> g[maxn], r[maxn], rdom[maxn];
int dfn[maxn], rev[maxn], fa[maxn], sdom[maxn], dom[maxn], val[
     maxn], rp[maxn], tk;
void init(int n) {
  // vertices are numbered from 0 to n - 1
  fill(dfn, dfn + n, -1);
  fill(rev, rev + n, -1);
fill(fa, fa + n, -1);
  fill(val, val + n, -1);
  fill(sdom, sdom + n, -1);
  fill(rp, rp + n, -1);
  fill(dom, dom + n, -1);
  tk = 0;
  for (int i = 0; i < n; ++i) {
    g[i].clear();
    r[i].clear();
    rdom[i].clear();
  }
void add_edge(int x, int y) { g[x].push_back(y); }
void dfs(int x) {
  rev[dfn[x] = tk] = x;
  fa[tk] = sdom[tk] = val[tk] = tk;
  tk++
  for (int u : g[x]) {
  if (dfn[u] == -1) dfs(u), rp[dfn[u]] = dfn[x];
    r[dfn[u]].push_back(dfn[x]);
void merge(int x, int y) { fa[x] = y; }
int find(int x, int c = 0) {
  if (fa[x] == x) return c ? -1 : x;
  int p = find(fa[x], 1);
if (p == -1) return c ? fa[x] : val[x];
  if (sdom[val[x]] > sdom[val[fa[x]]]) val[x] = val[fa[x]];
  fa[x] = p;
  return c ? p : val[x];
vector<int> build(int s, int n) {
  // return the father of each node in the dominator tree
  // p[i] = -2 if i is unreachable from s
  dfs(s);
  for (int i = tk - 1; i >= 0; --i) {
    for (int u : r[i]) sdom[i] = min(sdom[i], sdom[find(u)]);
    if (i) rdom[sdom[i]].push_back(i);
    for (int &u : rdom[i]) {
       int p = find(u);
      if (sdom[p] == i) dom[u] = i;
else dom[u] = p;
    if (i) merge(i, rp[i]);
  }
  vector<int> p(n, -2); p[s] = -1;
  for (int i = 1; i < tk; ++i) if (sdom[i] != dom[i]) dom[i] =</pre>
     dom[dom[i]];
  for (int i = 1; i < tk; ++i) p[rev[i]] = rev[dom[i]];</pre>
  return p;
```

#### 4.9 Virtual Tree

```
void VirtualTree(vector<int> v) {
   v.push_back(0);
   sort(v.begin(), v.end(), [&](int i, int j) { return dfn[i] <
      dfn[j]; });
   v.resize(unique(v.begin(), v.end()) - v.begin());
   vector<int> stk;
   for (int u : v) {
      if (stk.empty()) {
        stk.push_back(u);
      continue;
   }
}
```

```
int p = GetLCA(u, stk.back());
if (p != stk.back()) {
   while (stk.size() >= 2 && dep[p] <= dep[stk[stk.size() -
   2]]) {
    int x = stk.back();
    stk.pop_back();
   AddEdge(x, stk.back());
   }
   if (stk.back() != p) {
    AddEdge(stk.back(), p);
    stk.pop_back();
    stk.push_back(p);
   }
   stk.push_back(u);
}

for (int i = 0; i + 1 < stk.size(); ++i) AddEdge(stk[i], stk[i + 1]);
}</pre>
```

#### 4.10 Vizing's Theorem

```
\mbox{ namespace vizing } \{ \mbox{ // returns edge coloring in adjacent matrix } \mbox{ G. 1 - based }
int C[kN][kN], G[kN][kN];
void clear(int N) {
  for (int i = 0; i <= N; i++) {</pre>
                     for (int j = 0; j \le N; j++) C[i][j] = G[i][j] = 0;
         }
}
void solve(vector<pair<int, int>> &E, int N, int M) {
         int X[kN] = {}, a;
auto update = [&](int u) {
                   for (X[u] = 1; C[u][X[u]]; X[u]++);
          auto color = [&](int u, int v, int c) {
                     int p = G[u][v];
                     G[u][v] = G[v][u] = c;
                    C[u][c] = v, C[v][c] = u;
                    C[u][p] = C[v][p] = 0;
                     if (p) X[u] = X[v] = p
                     else update(u), update(v);
                     return p;
          auto flip = [&](int u, int c1, int c2) {
                    int p = C[u][c1];
                     swap(C[u][c1], C[u][c2]);
                     if (p) G[u][p] = G[p][u] = c2;
                     if (!C[u][c1]) X[u] = c1;
                     if (!C[u][c2]) X[u] = c2;
         };
          for (int i = 1; i <= N; i++) X[i] = 1;
for (int t = 0; t < E.size(); t++) {</pre>
                     int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c
                         = c0. d:
                     vector<pair<int, int>> L;
                     int vst[kN] = {};
                     while (!G[u][v0]) {
                              L.emplace_back(v, d = X[v]);
if (!C[v][c]) for (a = (int)L.size() - 1; a >= 0; a--) c
                         = color(u, L[a].first, c);
                              else if (!C[u][d]) for (a = (int)L.size() - 1; a >= 0; a
                           --) color(u, L[a].first, L[a].second);
                              else if (vst[d]) break
                               else vst[d] = 1, v = C[u][d];
                     if (!G[u][v0]) {
                              for (; v; v = flip(v, c, d), swap(c, d));
if (C[u][c0]) {
                                        for (a = (int)L.size() - 2; a >= 0 && L[a].second != c;
                                         for (; a \ge 0; a \longrightarrow 0; a \longrightarrow
                              } else t--;
         }
```

#### 4.11 System of Difference Constraints

Given m constrains on n variables  $x_1, x_2, \ldots, x_n$  of form  $x_i - x_j \leq w$  (resp,  $x_i - x_j \geq w$ ), connect  $i \to j$  with weight w. Then connect  $0 \to i$  for all i with weight 0 and find the shortest path (resp, longest path) on the graph. dis(i) will be the maximum (resp, minimum) solution to  $x_i$ .

# 5 String

### 5.1 Knuth-Morris-Pratt Algorithm

```
vector<int> kmp(const string &s) {
   vector<int> f(s.size(), 0);
   // f[i] = length of the longest prefix (excluding s[0:i])
      such that it coincides with the suffix of s[0:i] of the
      same length
   // i + 1 - f[i] is the length of the smallest recurring
     period of s[0:i]
   int k = 0:
   for (int i = 1; i < (int)s.size(); ++i) {</pre>
     while (k > 0 \& s[i] != s[k]) k = f[k - 1];
     if (s[i] == s[k]) ++k;
     f[i] = k;
   return f;
 vector<int> search(const string &s, const string &t) {
  // return 0-indexed occurrence of t in s
  vector<int> f = kmp(t), res;
for (int i = 0, k = 0; i < (int)s.size(); ++i) {</pre>
     while (k > 0 \&\& (k == (int)t.size() || s[i] != t[k])) k = f
     [k - 1];
     if (s[i] == t[k]) ++k;
     if (k == (int)t.size()) res.push_back(i - t.size() + 1);
   return res:
i}
```

### 5.2 Z Algorithm

```
| int z[maxn];
|// z[i] = LCP of suffix i and suffix 0
| void z_function(const string& s) {
| memset(z, 0, sizeof(z));
| z[0] = (int)s.length();
| int l = 0, r = 0;
| for (int i = 1; i < s.length(); ++i) {
| z[i] = max(0, min(z[i - l], r - i + 1));
| while (i + z[i] < s.length() && s[z[i]] == s[i + z[i]]) {
| l = i; r = i + z[i];
| ++z[i];
| }
| }
| }</pre>
```

#### 5.3 Manacher's Algorithm

```
int z[maxn];
int manacher(const string& s) {
   string t = ".";
   for (int i = 0; i < s.length(); ++i) t += s[i], t += '.';
   int l = 0, r = 0, ans = 0;
   for (int i = 1; i < t.length(); ++i) {
      z[i] = (r > i ? min(z[2 * l - i], r - i) : 1);
      while (i - z[i] >= 0 && i + z[i] < t.length() && t[i - z[i] = t[i + z[i]]) ++z[i];
      if (i + z[i] > r) r = i + z[i], l = i;
    }
   for (int i = 1; i < t.length(); ++i) ans = max(ans, z[i] - 1);
    return ans;
}</pre>
```

#### 5.4 Aho-Corasick Automaton

```
struct AC {
 static const int maxn = 1e5 + 5;
  int sz, ql, qr, root;
  int cnt[maxn], q[maxn], ed[maxn], el[maxn], ch[maxn][26], f[
    maxn];
  int gnode() {
    for (int i = 0; i < 26; ++i) ch[sz][i] = -1;
    f[sz] = -1;
    ed[sz] = 0;
    cnt[sz] = 0;
    return sz++;
  void init() {
    sz = 0;
    root = gnode();
  int add(const string &s) {
    int now = root;
```

```
for (int i = 0; i < s.length(); ++i) {
  if (ch[now][s[i] - 'a'] == -1) ch[now][s[i] - 'a'] =</pre>
      gnode();
       now = ch[now][s[i] - 'a'];
     ed[now] = 1;
     return now;
   void build_fail() {
     ql = qr = 0; q[qr++] = root;
     while (ql < qr) {</pre>
       int now = q[ql++];
       for (int i = 0; i < 26; ++i) if (ch[now][i] != -1) {
         int p = ch[now][i], fp = f[now];
         while (fp != -1 && ch[fp][i] == -1) fp = f[fp];
          int pd = fp != -1 ? ch[fp][i] : root;
         f[p] = pd;
         el[p] = ed[pd] ? pd : el[pd];
         q[qr++] = p;
       }
     }
   }
   void build(const string &s) {
     build fail():
     int now = root;
     for (int i = 0; i < s.length(); ++i) {</pre>
       while (now != -1 \&\& ch[now][s[i] - 'a'] == -1) now = f[
       now = now != -1 ? ch[now][s[i] - 'a'] : root;
       ++cnt[now];
     for (int i = qr - 1; i >= 0; --i) cnt[f[q[i]]] += cnt[q[i]]
      ]];
   long long solve(int n) {
     build_fail();
     vector<vector<long long>> dp(sz, vector<long long>(n + 1,
     for (int i = 0; i < sz; ++i) dp[i][0] = 1;
     for (int i = 1; i <= n; ++i) {
       for (int j = 0; j < sz; ++j)
         for (int k = 0; k < 2; ++k) {
           if (ch[j][k] != -1) {
              if (!ed[ch[j][k]])
                dp[j][i] += dp[ch[j][k]][i - 1];
           } else {
             int z = f[j];
              while (z != root \&\& ch[z][k] == -1) z = f[z];
              int p = ch[z][k] == -1 ? root : ch[z][k];
              if (ch[z][k] == -1 \mid \mid !ed[ch[z][k]]) dp[j][i] += dp
      [p][i - 1];
           }
         }
       }
     return dp[0][n];
|};
```

### 5.5 Suffix Automaton

```
struct SAM {
   static const int maxn = 5e5 + 5;
   int nxt[maxn][26], to[maxn], len[maxn];
   int root, last, sz;
   int gnode(int x) {
     for (int i = 0; i < 26; ++i) nxt[sz][i] = -1;
     to[sz] = -1;
     len[sz] = x:
     return sz++;
   void init() {
     sz = 0;
     root = gnode(0);
     last = root;
   void push(int c) {
     int cur = last:
     last = gnode(len[last] + 1);
     for (; ~cur && nxt[cur][c] == -1; cur = to[cur]) nxt[cur][c
      ] = last;
     if (cur == -1) return to[last] = root, void();
     int link = nxt[cur][c];
     if (len[link] == len[cur] + 1) return to[last] = link, void
     ();
     int tlink = gnode(len[cur] + 1);
```

```
for (; ~cur && nxt[cur][c] == link; cur = to[cur]) nxt[cur
     ][c] = tlink;
     for (int i = 0; i < 26; ++i) nxt[tlink][i] = nxt[link][i];</pre>
     to[tlink] = to[link];
    to[link] = tlink;
    to[last] = tlink;
  void add(const string &s) {
    for (int i = 0; i < s.size(); ++i) push(s[i] - 'a');</pre>
  bool find(const string &s) {
    int cur = root;
     for (int i = 0; i < s.size(); ++i) {</pre>
       cur = nxt[cur][s[i] - 'a'];
       if (cur == -1) return false;
    return true:
  int solve(const string &t) {
    int res = 0, cnt = 0;
     int cur = root;
     for (int i = 0; i < t.size(); ++i) {</pre>
       if (~nxt[cur][t[i] - 'a']) {
         ++cnt:
         cur = nxt[cur][t[i] - 'a'];
       } else {
         for (; ~cur && nxt[cur][t[i] - 'a'] == -1; cur = to[cur
         if (\simcur) cnt = len[cur] + 1, cur = nxt[cur][t[i] - 'a'
         else cnt = 0, cur = root;
       }
       res = max(res, cnt);
     return res:
∫};
```

### 5.6 Suffix Array

```
|// sa[i]: sa[i]-th suffix is the i-th lexigraphically smallest
     suffix.
   lcp[i]: longest common prefix of suffix sa[i] and suffix sa[
     i - 17.
namespace sfx {
vector<int> Build(const string &s) {
  int n = s.size();
  vector<int> str(n * 2), sa(n * 2), c(max(n, 256) * 2), x(max(n, 256)), p(n), q(n * 2), t(n * 2);
  for (int i = 0; i < n; ++i) str[i] = s[i];</pre>
  auto Pre = [&](int *sa, int *c, int n, int z) {
    memset(sa, 0, sizeof(int) * n);
    memcpy(x.data(), c, sizeof(int) * z);
  auto Induce = [&](int *sa, int *c, int *s, int *t, int n, int
      z) {
    memcpy(x.data() + 1, c, sizeof(int) * (z - 1));
    for (int i = 0; i < n; ++i) if (sa[i] && !t[sa[i] - 1]) sa[
     x[s[sa[i] - 1]]++] = sa[i] - 1;
    memcpy(x.data(), c, sizeof(int) * z);
for (int i = n - 1; i >= 0; --i) if (sa[i] && t[sa[i] - 1])
      sa[--x[s[sa[i] - 1]]] = sa[i] - 1;
  auto SAIS = [&](auto self, int *s, int *sa, int *p, int *q,
     int *t, int *c, int n, int z) -> void {
     bool uniq = t[n - 1] = true;
     int nn = 0, nmxz = -1, *nsa = sa + n, *ns = s + n, last =
    memset(c, 0, sizeof(int) * z);
     for (int i = 0; i < n; ++i) uniq &= ++c[s[i]] < 2;
     for (int i = 0; i < z - 1; ++i) c[i + 1] += c[i];
     if (uniq) {
       for (int i = 0; i < n; ++i) sa[--c[s[i]]] = i;
    for (int i = n - 2; i >= 0; --i) t[i] = (s[i] == s[i + 1] ?
t[i + 1] : s[i] < s[i + 1]);
    Pre(sa, c, n, z);
     for (int i = 1; i <= n - 1; ++i) if (t[i] && !t[i - 1]) sa
     [--x[s[i]]] = p[q[i] = nn++] = i;
    Induce(sa, c, s, t, n, z);
for (int i = 0; i < n; ++i) if (sa[i] && t[sa[i]] && !t[sa[
     i] - 1]) {
       bool neq = last < 0 \parallel memcmp(s + sa[i], s + last, (p[q[
     sa[i]] + 1] - sa[i]) * sizeof(int));
      ns[q[last = sa[i]]] = nmxz += neq;
```

```
self(self, ns, nsa, p + nn, q + n, t + n, c + z, nn, nmxz + n
       1):
     Pre(sa, c, n, z);
     for (int i = nn - 1; i >= 0; --i) sa[--x[s[p[nsa[i]]]]] = p
      [nsa[i]];
     Induce(sa, c, s, t, n, z);
   SAIS(SAIS, str.data(), sa.data(), p.data(), q.data(), t.data
      (), c.data(), n + 1, 256);
   return vector<int>(sa.begin() + 1, sa.begin() + n + 1);
vector<int> BuildLCP(const vector<int> &sa, const string &s) {
   int n = s.size();
   vector<int> lcp(n), rev(n);
   for (int i = 0; i < n; ++i) rev[sa[i]] = i;
for (int i = 0, ptr = 0; i < n; ++i) {
     if (!rev[i]) {
       ptr = 0;
       continue:
     while (i + ptr < n \&\& s[i + ptr] == s[sa[rev[i] - 1] + ptr]
      ]) ptr++
     lcp[rev[i]] = ptr ? ptr-- : 0;
   return lcp:
|}} // namespace sfx
```

#### 5.7 Palindromic Tree

```
struct PalindromicTree {
   int link[kN], len[kN], dp[kN], nxt[kN][26], sz, sf;
   int gnode(int 1, int fl = -1) {
     len[sz] = 1;
link[sz] = fl;
     fill(nxt[sz], nxt[sz] + 26, -1);
     return sz++;
   void Init() {
     sz = 0;
     sf = 1:
     gnode(-1, 0);
     gnode(0, 0);
   void Push(const string &s, int pos) {
     int cur = sf, z = s[pos] - 'a'
     while (pos - 1 - len[cur] < 0 || s[pos - 1 - len[cur]] != s</pre>
      [pos]) cur = link[cur];
     if (nxt[cur][z] != -1) {
       sf = nxt[cur][z];
     } else {
       int ch = gnode(len[cur] + 2);
       nxt[cur][z] = sf = ch;
if (len[ch] == 1) {
         link[ch] = 1;
       } else {
         cur = link[cur];
         while (pos - 1 - len[cur] < 0 || s[pos - 1 - len[cur]]
      != s[pos]) cur = link[cur];
         link[ch] = nxt[cur][z];
      }
     dp[sf] += 1;
   long long Build(const string &s) {
     for (int i = 0; i < s.size(); ++i) Push(s, i);</pre>
     for (int i = sz - 1; i >= 0; --i) dp[link[i]] += dp[i];
     long long res = 0;
     for (int i = 0; i < sz; ++i) res = max(res, 1LL * dp[i] *
      lenΓi]);
     return res;
  }
|} plt;
```

# 5.8 Circular LCS

```
| string s1, s2;
| int n, m;
| int dp[kN * 2][kN];
| int nxt[kN * 2][kN];
| void reroot(int px) {
| int py = 1;
| while (py <= m && nxt[px][py] != 2) py++;
| nxt[px][py] = 1;
| while (px < 2 * n && py < m) {
| if (nxt[px + 1][py] == 3) px++, nxt[px][py] = 1;
| else if (nxt[px + 1][py + 1] == 2) px++, py++, nxt[px][py] = 1;
```

```
else py++;
   while (px < 2 * n && nxt[px + 1][py] == 3) px++, nxt[px][py]
      = 1;
 int track(int x, int y, int e) { // use this routine to find
      LCS as string
   int ret = 0;
while (y != 0 && x != e) {
     if (nxt[x][y] == 1) y
     else if (nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
else if (nxt[x][y] == 3) x--;
   return ret;
 }
 int solve(string a, string b) {
   n = a.size(), m = b.size();
s1 = "#" + a + a, s1 = '#' + b;
   for (int i = 0; i \le 2 * n; i++) {
     for (int j = 0; j <= m; j++) {
  if (j == 0) { nxt[i][j] = 3; continue; }</pre>
        if (i == 0) { nxt[i][j] = 1; continue; }
        dp[i][j] = -1;
        if (dp[i][j] < dp[i][j - 1]) dp[i][j] = dp[i][j - 1], nxt</pre>
      [i][j] = 1;
        if (dp[i][j] < dp[i - 1][j - 1] + (s1[i] == s2[j])) dp[i]
      [j] = dp[i - 1][j - 1] + (s1[i] == s2[j]), nxt[i][j] = 2;
        if (dp[i][j] < dp[i - 1][j]) dp[i][j] = dp[i - 1][j], nxt</pre>
      [i][j] = 3;
   }
   int ret = dp[n][m];
   for (int i = 1; i < n; i++) reroot(i), ret = max(ret, track(n
        + i, m, i));
   return ret;
| }
```

#### 5.9 Lexicographically Smallest Rotation

```
| string rotate(const string &s) {
| int n = s.length();
| string t = s + s;
| int i = 0, j = 1;
| while (i < n && j < n) {
| int k = 0;
| while (k < n && t[i + k] == t[j + k]) ++k;
| if (t[i + k] <= t[j + k]) j += k + 1;
| else i += k + 1;
| if (i == j) ++j;
| }
| int pos = (i < n ? i : j);
| return t.substr(pos, n);
| }</pre>
```

#### 6 Math

#### 6.1 Fast Fourier Transform

```
namespace fft {
struct cplx {
  double re, im;
  cplx(): re(0), im(0) {}
  cplx(double r, double i): re(r), im(i) {}
  cplx operator+(const cplx &rhs) const { return cplx(re + rhs.
    re, im + rhs.im); }
  cplx operator-(const cplx &rhs) const { return cplx(re - rhs.
    re, im - rhs.im); }
  cplx operator*(const cplx &rhs) const { return cplx(re * rhs.
    re - im * rhs.im, re * rhs.im + im * rhs.re); }
  cplx conj() const { return cplx(re, -im); }
const double pi = acos(-1);
cplx omega[maxn + 1];
void prefft() {
  for (int i = 0; i \leftarrow maxn; ++i)
    omega[i] = cplx(cos(2 * pi * i / maxn), sin(2 * pi * i / maxn))
void bitrev(vector<cplx> &v, int n) {
  int z = __builtin_ctz(n) - 1;
  for (int i = 0; i < n; ++i) {
    for (int j = 0; (1 << j) < n; ++j) x ^= (i >> j & 1) << (z
     - j);
    if (x > i) swap(v[x], v[i]);
```

```
}
void fft(vector<cplx> &v, int n) {
  bitrev(v, n);
   for (int s = 2; s <= n; s <<= 1) {
     int z = s \gg 1;
     for (int i = 0; i < n; i += s) {
       for (int k = 0; k < z; ++k) {
         cplx x = v[i + z + k] * omega[maxn / s * k];
         v[i + z + k] = v[i + k] - x;
         v[i + k] = v[i + k] + x;
    }
}
void ifft(vector<cplx> &v, int n) {
  fft(v, n);
   reverse(v.begin() + 1, v.end());
   for (int i = 0; i < n; ++i) v[i] = v[i] * cplx(1. / n, 0);
vector<long long> convolution(const vector<int> &a, const
      vector<int> &b) {
   // Should be able to handle N <= 10^5, C <= 10^4
   int sz = 1;
   while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
   vector<cplx> v(sz);
   for (int i = 0; i < sz; ++i) {
     double re = i < a.size() ? a[i] : 0;</pre>
     double im = i < b.size() ? b[i] : 0;</pre>
     v[i] = cplx(re, im);
   fft(v, sz);
  for (int i = 0; i <= sz / 2; ++i) {
  int j = (sz - i) & (sz - 1);</pre>
     cplx x = (v[i] + v[j].conj()) * (v[i] - v[j].conj()) * cplx
      (0, -0.25);
     if (j != i) v[j] = (v[j] + v[i].conj()) * (v[j] - v[i].conj
()) * cplx(0, -0.25);
     v[i] = x;
   ifft(v, sz);
  vector<long long> c(sz);
   for (int i = 0; i < sz; ++i) c[i] = round(v[i].re);
   return c;
vector<int> convolution_mod(const vector<int> &a, const vector<</pre>
     int> &b, int p) {
   int sz = 1;
   while (sz < (int)a.size() + (int)b.size() - 1) sz <<= 1;</pre>
   vector<cplx> fa(sz), fb(sz);
   for (int i = 0; i < (int)a.size(); ++i)</pre>
     fa[i] = cplx(a[i] & ((1 << 15) - 1), a[i] >> 15);
   for (int i = 0; i < (int)b.size(); ++i)</pre>
     fb[i] = cplx(b[i] & ((1 << 15) - 1), b[i] >> 15);
  fft(fa, sz), fft(fb, sz);
double r = 0.25 / sz;
   cplx r2(0, -1), r3(r, 0), r4(0, -r), r5(0, 1);
   for (int i = 0; i \leftarrow (sz >> 1); ++i) {
     int j = (sz - i) & (sz - 1);
     cplx a1 = (fa[i] + fa[j].conj());
     cplx a2 = (fa[i] - fa[j].conj()) * r2;
     cplx \ b1 = (fb[i] + fb[j].conj()) * r3;
     cplx b2 = (fb[i] - fb[j].conj()) * r4;
     if (i != j) {
       cplx c1 = (fa[j] + fa[i].conj());
       cplx c2 = (fa[j] - fa[i].conj()) * r2;
       cplx d1 = (fb[j] + fb[i].conj()) * r3;
       cplx d2 = (fb[j] - fb[i].conj()) * r4;
       fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[j] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz), fft(fb, sz);
   vector<int> res(sz);
   for (int i = 0; i < sz; ++i) {
     long long a = round(fa[i].re);
     long long b = round(fb[i].re);
     long long c = round(fa[i].im);
     res[i] = (a + ((b \% p) << 15) + ((c \% p) << 30)) \% p;
   return res;
| }}
```

#### 6.2 Number Theoretic Transform

```
vector<int> omega;
void Init() {
  omega.resize(kN + 1);
  long long x = fpow(kRoot, (Mod - 1) / kN);
  omega[0] = 1;
  for (int i = 1; i <= kN; ++i) {
  omega[i] = 1LL * omega[i - 1] * x % kMod;</pre>
void Transform(vector<int> &v, int n) {
  BitReverse(v, n);
  for (int s = 2; s <= n; s <<= 1) {
     int z = s \gg 1;
     for (int i = 0; i < n; i += s) {
       for (int k = 0; k < z; ++k) {
         int x = 1LL * v[i + k + z] * omega[kN / s * k] % kMod;
         v[i + k + z] = (v[i + k] + kMod - x) % kMod;
         (v[i + k] += x) \% = kMod;
    }
  }
}
void InverseTransform(vector<int> &v, int n) {
  Transform(v, n);
  for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);
  const int kInv = fpow(n, kMod - 2);
for (int i = 0; i < n; ++i) v[i] = 1LL * v[i] * inv % kMod;</pre>
```

#### 6.3 NTT Prime List

```
Prime
             Root
                    Prime
                                   Root
                     167772161
7681
             17
12289
                     104857601
             11
                     985661441
65537
             3
                     998244353
786433
            10
                     1107296257
5767169
            3
                    2013265921
                                   31
7340033
                     2810183681
            3
                                   11
23068673
                     2885681153
469762049
                     605028353
```

#### 6.4 Formal Power Series

```
Poly Inverse(Poly f) {
  int n = f.size();
  Poly q(1, fpow(f[0], kMod - 2));
  for (int s = 2;; s <<= 1) {</pre>
     if (f.size() < s) f.resize(s);</pre>
    Poly fv(f.begin(), f.begin() + s);
    Poly fq(q.begin(), q.end());
     fv.resize(s + s);
    fq.resize(s + s);
    ntt::Transform(fv, s + s);
    ntt::Transform(fq, s + s);
    for (int i = 0; i < s + s; ++i) {
  fv[i] = 1LL * fv[i] * fq[i] % kMod * fq[i] % kMod;</pre>
    ntt::InverseTransform(fv, s + s);
    Poly res(s);
    for (int i = 0; i < s; ++i) {
       res[i] = kMod - fv[i];
       if (i < (s >> 1)) {
  int v = 2 * q[i] % kMod;
         (res[i] += v) >= kMod ? res[i] -= kMod : 0;
      }
    a = res;
    if (s >= n) break;
  q.resize(n);
  return q;
Poly Divide(const Poly &a, const Poly &b) {
  int n = a.size(), m = b.size(), k = 2;
  while (k < n - m + 1) k <<= 1;
  Poly ra(k), rb(k);
  for (int i = 0; i < min(n, k); ++i) ra[i] = a[n - 1 - i];
for (int i = 0; i < min(m, k); ++i) rb[i] = b[m - 1 - i];
  auto rbi = Inverse(rb);
auto res = Multiply(rbi, ra);
  res.resize(n - m + 1);
  reverse(res.begin(), res.end());
  return res;
Poly Modulo(const Poly &a, const Poly &b) {
  if (a.size() < b.size()) return a;</pre>
  auto dv = Multiply(Divide(a, b), b);
```

```
assert(dv.size() == a.size());
  for (int i = 0; i < dv.size(); ++i) {
  dv[i] = (a[i] + kMod - dv[i]) % kMod;</pre>
  while (!dv.empty() && dv.back() == 0) dv.pop_back();
Poly Derivative(const Poly &f) {
  int n = f.size();
  vector<int> res(n - 1);
  for (int i = 0; i < n - 1; ++i) {
  res[i] = 1LL * f[i + 1] * (i + 1) % kMod;</pre>
Poly Integral(const Poly &f) {
  int n = f.size();
  vector<int> res(n + 1);
  for (int i = 0; i < n; ++i) {
    res[i + 1] = 1LL * f[i] * fpow(i + 1, kMod - 2) % kMod;
  return res:
Poly Evaluate(const Poly &f, const vector<int> &x) {
  if (x.empty()) return Poly();
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {kMod - x[i], 1};
  for (int i = n - 1; i > 0; --i) up[i] = Multiply(up[i * 2],
     up[i * 2 + 1]);
  vector<Poly> down(n * 2);
  down[1] = Modulo(f, up[1]);
for (int i = 2; i < n * 2; ++i) down[i] = Modulo(down[i >>
     1], up[i]);
  vector<int> y(n);
  for (int i = 0; i < n; ++i) y[i] = down[i + n][0];
Poly Interpolate(const vector<int> &x, const vector<int> &y) {
  int n = x.size();
  vector<Poly> up(n * 2);
  for (int i = 0; i < n; ++i) up[i + n] = {kMod - x[i], 1}; for (int i = n - 1; i > 0; --i) up[i] = Multiply(up[i * 2],
     up[i * 2 + 1]);
  vector<int> a = Evaluate(Derivative(up[1]), x);
  for (int i = 0; i < n; ++i) {
    a[i] = 1LL * y[i] * fpow(a[i], kMod - 2) % kMod;
  vector<Poly> down(n * 2);
  for (int i = 0; i < n; ++i) down[i + n] = {a[i]};
  for (int i = n - 1; i > 0; --i) {
    auto lhs = Multiply(down[i * 2], up[i * 2 + 1]);
    auto rhs = Multiply(down[i * 2 + 1], up[i * 2]);
    assert(lhs.size() == rhs.size());
    down[i].resize(lhs.size());
    for (int j = 0; j < lhs.size(); ++j) {</pre>
       down[i][j] = (lhs[j] + rhs[j]) % kMod;
  }
  return down[1];
Poly Log(Poly f) {
  int n = f.size();
  if (n == 1) return {0};
auto d = Derivative(f);
  f.resize(n - 1);
  d = Multiply(d, Inverse(f));
  d.resize(n - 1)
  return Integral(d);
Poly Exp(Poly f) {
  int n = f.size();
  Poly q(1, 1);
  f[0] += 1;
  for (int s = 1; s < n; s <<= 1) {
     if (f.size() < s + s) f.resize(s + s);</pre>
    Poly g(f.begin(), f.begin() + s + s);
Poly h(q.begin(), q.end());
    h.resize(s + s);
    h = Log(h);
    for (int i = 0; i < s + s; ++i) {
       g[i] = (g[i] + kMod - h[i]) \% kMod;
    g = Multiply(g, q);
    g.resize(s + s);
```

```
assert(q.size() >= n);
  q.resize(n);
  return q;
Poly SquareRootImpl(Poly f) {
  if (f.empty()) return {0};
  int z = QuadraticResidue(f[0], kMod), n = f.size();
  constexpr int kInv2 = (kMod + 1) >> 1;
  if (z == -1) return {-1};
  vector<int> q(1, z);
  for (int s = 1; s < n; s <<= 1) {
    if (f.size() < s + s) f.resize(s + s);</pre>
    vector<int> fq(q.begin(), q.end());
    fq.resize(s + s);
    vector<int> f2 = Multiply(fq, fq);
    f2.resize(s + s);

for (int i = 0; i < s + s; ++i) {

f2[i] = (f2[i] + kMod - f[i]) % kMod;
    f2 = Multiply(f2, Inverse(fq));
    f2.resize(s + s);
for (int i = 0; i < s + s; ++i) {
   fq[i] = (fq[i] + kMod - 1LL * f2[i] * kInv2 % kMod) %</pre>
     kMod;
    q = fq;
  }
  q.resize(n);
  return q;
Poly SquareRoot(Poly f) {
  int n = f.size(), m = 0;
while (m < n && f[m] == 0) m++;</pre>
  if (m == n) return vector<int>(n);
  if (m & 1) return {-1};
  auto s = SquareRootImpl(vector<int>(f.begin() + m, f.end()));
  if (s[0] == -1) return {-1};
  vector<int> res(n);
  for (int i = 0; i < s.size(); ++i) res[i + m / 2] = s[i];</pre>
```

#### 6.5 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0\pmod{x^{2^k}}$ , then

$$Q_{k+1} = Q_k - \frac{F(Q_k)}{F'(Q_k)} \pmod{x^{2^{k+1}}}$$

### 6.6 General Purpose Numbers

#### 6.6.1 Bernoulli numbers

EGF of Bernoulli numbers is  $B(x) = \frac{x}{e^x - 1}$ .

#### 6.6.2 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$
 
$$S(n,1) = S(n,n) = 1$$
 
$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k} (-1)^{k-i} {k \choose i} i^{n}$$

# 6.7 Fast Walsh-Hadamard Transform

1. XOR Convolution

 $\begin{array}{l} \bullet \quad f(A) = (f(A_0) + f(A_1), f(A_0) - f(A_1)) \\ \bullet \quad f^{-1}(A) = (f^{-1}(\frac{A_0 + A_1}{2}), f^{-1}(\frac{A_0 - A_1}{2})) \end{array}$ 

2. OR Convolution

•  $f(A) = (f(A_0), f(A_0) + f(A_1))$ •  $f^{-1}(A) = (f^{-1}(A_0), f^{-1}(A_1) - f^{-1}(A_0))$ 

3. AND Convolution

•  $f(A) = (f(A_0) + f(A_1), f(A_1))$ •  $f^{-1}(A) = (f^{-1}(A_0) - f^{-1}(A_1), f^{-1}(A_1))$ 

### 6.8 Simplex Algorithm

```
Description: maximize \mathbf{c}^T \mathbf{x} subject to A\mathbf{x} \leq \mathbf{b} and \mathbf{x} \geq 0. Returns -\infty if infeasible and \infty if unbounded.
```

```
const double eps = 1e-9;
const double inf = 1e+9;
int n, m;
vector<vector<double>> d;
vector<int> p, q;
void pivot(int r, int s) {
   double inv = 1.0 / d[r][s];
for (int i = 0; i < m + 2; ++i) {
  for (int j = 0; j < n + 2; ++j) {</pre>
       if (i != r && j != s) d[i][j] -= d[r][j] * d[i][s] * inv;
   for (int i = 0; i < m + 2; ++i) if (i != r) d[i][s] *= -inv;
   for (int j = 0; j < n + 2; ++j) if (j != s) d[r][j] *= +inv;
   d[r][s] = inv;
   swap(p[r], q[s]);
bool phase(int z) {
   int x = m + z;
   while (true) {
     int s = -1;
     for (int i = 0; i \le n; ++i) {
       if (!z && q[i] == -1) continue;
       if (s == -1 || d[x][i] < d[x][s]) s = i;
     if (d[x][s] > -eps) return true;
     int r = -1;
     for (int i = 0; i < m; ++i) {
       if (d[i][s] < eps) continue;
if (r == -1 || d[i][n + 1] / d[i][s] < d[r][n + 1] / d[r</pre>
      ][s]) r = i;
     if (r == -1) return false;
     pivot(r, s);
  }
}
vector<double> solve(const vector<vector<double>> &a, const
      vector<double> &b, const vector<double> &c) {
   m = b.size(), n = c.size();
   d = vector<vector<double>>(m + 2, vector<double>(n + 2));
   for (int i = 0; i < m; ++i) {
     for (int j = 0; j < n; ++j) d[i][j] = a[i][j];
  p.resize(m), q.resize(n + 1);
   for (int i = 0; i < m; ++i) p[i] = n + i, d[i][n] = -1, d[i][
      n + 1] = b[i];
   for (int i = 0; i < n; ++i) q[i] = i, d[m][i] = -c[i];
   q[n] = -1, d[m + 1][n] = 1;
   for (int i = 1; i < m; ++i) if (d[i][n + 1] < d[r][n + 1]) r
   if (d[r][n + 1] < -eps) {
     pivot(r, n);
     if (!phase(1) || d[m + 1][n + 1] < -eps) return vector<</pre>
      double>(n, -inf);
     for (int i = 0; i < m; ++i) if (p[i] == -1) {
       int s = min_element(d[i].begin(), d[i].end() - 1) - d[i].
      begin();
       pivot(i, s);
     }
   if (!phase(0)) return vector<double>(n, inf);
   vector<double> x(n);
   for (int i = 0; i < m; ++i) if (p[i] < n) \times [p[i]] = d[i][n +
  1];
return x;
}
6.9 Subset Convolution
```

```
Description: h(s) = \sum_{s' \subset s} f(s')g(s \setminus s')
```

```
if (s >> j & 1) {
    a[i][s] += a[i][s ^ (1 << j)];
    b[i][s] += b[i][s ^ (1 << j)];</pre>
  }
}
vector<vector<int>>> c(n + 1, vector<int>(m));
for (int s = 0; s < m; ++s) {
   for (int i = 0; i <= n; ++i) {
     for (int j = 0; j \le i; ++j) c[i][s] += a[j][s] * b[i - j]
   ][s];
  }
for (int i = 0; i <= n; ++i) {
  for (int j = 0; j < n; ++j) {
  for (int s = 0; s < m; ++s) {</pre>
       if (s >> j & 1) c[i][s] -= c[i][s ^ (1 << j)];</pre>
  }
}
vector<int> res(m);
for (int i = 0; i < m; ++i) res[i] = c[__builtin_popcount(i)</pre>
   ][i];
return res:
```

#### 6.9.1 Construction

Standard form: maximize  $\mathbf{c}^T\mathbf{x}$  subject to  $A\mathbf{x} \leq \mathbf{b}$  and  $\mathbf{x} \geq 0$ . Dual LP: minimize  $\mathbf{b}^T\mathbf{y}$  subject to  $A^T\mathbf{y} \geq \mathbf{c}$  and  $\mathbf{y} \geq 0$ .  $\bar{\mathbf{x}}$  and  $\bar{\mathbf{y}}$  are optimal if and only if for all  $i \in [1,n]$ , either  $\bar{x}_i = 0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j = c_i$  holds and for all  $i \in [1,m]$  either  $\bar{y}_i = 0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j = b_j$  holds.

- 1. In case of minimization, let  $c'_i = -c_i$
- 2.  $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j \to \sum_{1 \le i \le n} -A_{ji} x_i \le -b_j$
- 3.  $\sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$   $\sum_{1 \le i \le n} A_{ji} x_i \ge b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.10 Schreier-Sims Algorithm

```
namespace schreier {
vector<vector<int>>> bkts, binv;
vector<vector<int>> lk:
vector<int> operator*(const vector<int> &a, const vector<int> &
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[i] = b[a[i]];</pre>
  return res;
vector<int> inv(const vector<int> &a) {
  vector<int> res(a.size());
  for (int i = 0; i < (int)a.size(); ++i) res[a[i]] = i;</pre>
  return res;
int filter(const vector<int> &g, bool add = true) {
 n = (int)bkts.size();
  vector < int > p = g;
  for (int i = 0; i < n; ++i) {
    assert(p[i] >= 0 && p[i] < (int)lk[i].size());</pre>
    int res = lk[i][p[i]];
if (res == -1) {
      if (add) {
        bkts[i].push_back(p);
        binv[i].push_back(inv(p));
        lk[i][p[i]] = (int)bkts[i].size() - 1;
      return i;
    }
    p = p * binv[i][res];
bool inside(const vector<int> &g) { return filter(g, false) ==
     -1; }
void solve(const vector<vector<int>> &gen, int _n) {
  bkts.clear(), bkts.resize(n);
  binv.clear(), binv.resize(n);
 lk.clear(), lk.resize(n);
vector<int> iden(n);
  iota(iden.begin(), iden.end(), 0);
```

```
for (int i = 0; i < n; ++i) {
  lk[i].resize(n, -1);</pre>
    bkts[i].push_back(iden);
    binv[i].push_back(iden);
    lk[i][i] = 0;
  for (int i = 0; i < (int)gen.size(); ++i) filter(gen[i]);</pre>
  queue<pair<pair<int, int>, pair<int, int>>> upd;
  for (int i = 0; i < n; ++i) {
    for (int j = i; j < n; ++j) {
      for (int k = 0; k < (int)bkts[i].size(); ++k) {</pre>
        for (int l = 0; l < (int)bkts[j].size(); ++l)</pre>
           upd.emplace(make_pair(i, k), make_pair(j, l));
    }
  }
  while (!upd.empty()) {
    auto a = upd.front().first;
    auto b = upd.front().second;
    upd.pop();
    int res = filter(bkts[a.first][a.second] * bkts[b.first][b.
     second]);
    if (res == -1) continue;
    pair<int, int> pr = make_pair(res, (int)bkts[res].size() -
     1);
    for (int i = 0; i < n; ++i) {
  for (int j = 0; j < (int)bkts[i].size(); ++j) {</pre>
         if (i <= res) upd.emplace(make_pair(i, j), pr);</pre>
         if (res <= i) upd.emplace(pr, make_pair(i, j));</pre>
      }
    }
 }
long long size() {
  long long res = 1;
  for (int i = 0; i < n; ++i) res = res * bkts[i].size();</pre>
  return res;
```

#### 6.11 Berlekamp-Massey Algorithm

```
template <int P>
vector<int> BerlekampMassey(vector<int> x) {
  vector<int> cur, ls;
  int lf = 0, ld = 0;
for (int i = 0; i < (int)x.size(); ++i) {
    int t = 0;
    for (int j = 0; j < (int)cur.size(); ++j)
      (t += 1LL * cur[j] * x[i - j - 1] % P) %= P;
    if (t == x[i]) continue;
    if (cur.empty()) {
       cur.resize(i + 1);
       lf = i, ld = (t + P - x[i]) % P;
      continue:
    int k = 1LL * fpow(ld, P - 2, P) * (t + P - x[i]) % P;
    vector<int> c(i - lf - 1);
    c.push_back(k);
    for (int j = 0; j < (int)ls.size(); ++j)
  c.push_back(1LL * k * (P - ls[j]) % P);</pre>
    if (c.size() < cur.size()) c.resize(cur.size());</pre>
    for (int j = 0; j < (int)cur.size(); ++j)
      c[j] = (c[j] + cur[j]) % P;
    if (i - lf + (int)ls.size() >= (int)cur.size()) {
      ls = cur, lf = i;
      ld = (t + P - x[i]) \% P;
    cur = c:
  return cur;
```

#### 6.12 Fast Linear Recurrence

#### 6.13 Miller Rabin

```
l// n < 4759123141
                    chk = [2, 7, 61]
 // n < 1122004669633 chk = [2, 13, 23, 1662803]
 // n < 2^64
                 chk = [2, 325, 9375, 28178, 450775, 9780504,
     17952650227
 vector<long long> chk =
     {2,325,9375,28178,450775,9780504,1795265022};
 bool Check(long long a, long long u, long long n, int t) {
  a = fpow(a, u, n);
   if (a == 0 \mid \mid a == 1 \mid \mid a == n - 1) return true;
  for (int i = 0; i < t; ++i) {
     a = fmul(a, a, n);
     if (a == 1) return false;
     if (a == n - 1) return true;
  return false;
bool IsPrime(long long n) {
  if (n < 2) return false;
  if (n % 2 == 0) return n == 2;
   long long u = n - 1; int t = 0;
   for (; !(u & 1); u >>= 1, ++t);
  for (long long i : chk) {
    if (!Check(i, u, n, t)) return false;
   return true;
| }
```

#### 6.14 Pollard's Rho

```
map<long long, int> cnt;
 void PollardRho(long long n) {
   if (n == 1) return;
   if (prime(n)) return ++cnt[n], void();
   if (n % 2 == 0) return PollardRho(n / 2), ++cnt[2], void(); long long x = 2, y = 2, d = 1, p = 1;
   auto f = [\&](auto x, auto n, int p) { return <math>(fmul(x, x, n) +
       p) % n; }
   while (true) {
     if (d != n && d != 1) {
       PollardRho(n / d);
       PollardRho(d);
       return;
     if (d == n) ++p;
     x = f(x, n, p); y = f(f(y, n, p), n, p);
     d = \_gcd(abs(x - y), n);
| }
```

#### 6.15 Meissel-Lehmer Algorithm

```
int64_t PrimeCount(int64_t n) {
   if (n <= 1) return 0;
   const int v = sqrt(n);
   vector<int> smalls(v + 1);
   for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;
   int s = (v + 1) / 2;
   vector<int> roughs(s);
   for (int i = 0; i < s; ++i) roughs[i] = 2 * i + 1;
   vector<int64_t> larges(s);
   for (int i = 0; i < s; ++i) larges[i] = (n / (2 * i + 1) + 1) / 2;
   vector<br/>
   vector<br/>
   int pc = 0;
   for (int p = 3; p <= v; ++p) {
      if (smalls[p] > smalls[p - 1]) {
        int q = p * p;
        pc++;
      if (1LL * q * q > n) break;
      }
}
```

```
skip[p] = true;
                  for (int i = q; i <= v; i += 2 * p) skip[i] = true;</pre>
                  int ns = 0;
                  for (int k = 0; k < s; ++k) {
                          int i = roughs[k];
                          if (skip[i]) continue;
                          int64_t d = 1LL * i * p;
                          larges[ns] = larges[k] - (d \ll v ? larges[smalls[d] - v ] larges[sm
             pc] : smalls[n / d]) + pc;
                          roughs[ns++] = i;
                 }
                 s = ns;
                 for (int j = v / p; j >= p; --j) {
                          int c = smalls[j] - pc;
                          for (int i = j * p, e = min(i + p, v + 1); i < e; ++i)
             smalls[i] -= c;
        }
for (int k = 1; k < s; ++k) {
  const int64_t m = n / roughs[k];</pre>
         int64_t = larges[k] - (pc + k - 1);
         for (int l = 1; l < k; ++l) {
                 int p = roughs[l];
                 if (1LL * p * p > m) break;
                 s = smalls[m / p] - (pc + l - 1);
         larges[0] -= s;
}
return larges[0];
```

#### 6.16 Discrete Logarithm

Description: to find x such that  $x^a \equiv b \pmod{p}$ , let g be the primitive root of p, find k such that  $g^k \equiv b \pmod{p}$  and x can be found by  $g^d$  where  $ad \equiv k \pmod{p-1}$ .

```
int DiscreteLog(int s, int x, int y, int m) {
   constexpr int kStep = 32000;
   unordered_map<int, int> p;
   int b = 1;
   for (int i = 0; i < kStep; ++i) {</pre>
     p[y] = i;
y = 1LL * y * x % m;
b = 1LL * b * x % m;
   for (int i = 0; i < m + 10; i += kStep) {
   s = 1LL * s * b % m;</pre>
      if (p.find(s) != p.end()) return i + kStep - p[s];
   return -1:
int DiscreteLog(int x, int y, int m) {
   if (m == 1) return 0;
   int s = 1;
   for (int i = 0; i < 100; ++i) {
     if (s == y) return i;
s = 1LL * s * x % m;
   if (s == y) return 100;
   int p = 100 + DiscreteLog(s, x, y, m);
   if (fpow(x, p, m) != y) return -1;
```

#### 6.17 Quadratic Residue

```
int Jacobi(int a, int m) {
  int s = 1;
  for (; m > 1; ) {
    a %= m;
     if (a == 0) return 0;
     const int r = __builtin_ctz(a);
     if ((r \& 1) \&\& ((m + 2) \& 4)) s = -s;
    a \gg r
     if (a \& m \& 2) s = -s;
    swap(a, m);
  return s;
int QuadraticResidue(int a, int p) {
  if (p == 2) return a & 1;
   const int jc = Jacobi(a, p);
   if (jc == 0 || jc == -1) return jc;
  int b, d;
  for (; ; ) {
    b = rand() \% p;
```

```
d = (1LL * b * b + p - a) % p;
if (Jacobi(d, p) == -1) break;
}
int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
for (int e = (p + 1) >> 1; e; e >>= 1) {
   if (e & 1) {
     tmp = (1LL * g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
     g0 = tmp;
   }
   tmp = (1LL * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
   f1 = (2LL * f0 * f1) % p;
   f0 = tmp;
}
return g0;
}
```

#### 6.18 Gaussian Elimination

```
| double Gauss(vector<vector<double>>> &d) {
    int n = d.size(), m = d[0].size();
    double det = 1;
    for (int i = 0; i < m; ++i) {
        int p = -1;
        for (int j = i; j < n; ++j) {
            if (fabs(d[j][i]) < kEps) continue;
            if (p == -1 || fabs(d[j][i]) > fabs(d[p][i])) p = j;
        }
        if (p == -1) continue;
        if (p != i) det *= -1;
        for (int j = 0; j < m; ++j) swap(d[p][j], d[i][j]);
        for (int j = 0; j < n; ++j) {
            if (i == j) continue;
            double z = d[j][i] / d[i][i];
            for (int k = 0; k < m; ++k) d[j][k] -= z * d[i][k];
        }
    }
    for (int i = 0; i < n; ++i) det *= d[i][i];
    return det;
}</pre>
```

#### 6.19 $\mu$ function

```
int mu[kC], dv[kC];
vector<int> prime;
void Sieve() {
    mu[1] = dv[1] = 1;
    for (int i = 2; i < kC; ++i) {
        if (!dv[i]) {
            dv[i] = i, mu[i] = -1;
            prime.push_back(i);
    }
    for (int j = 0; i * prime[j] < kC; ++j) {
            dv[i * prime[j]] = prime[j];
            mu[i * prime[j]] = -mu[i];
            if (i % prime[j] == 0) {
                mu[i * prime[j]] = 0;
                 break;
          }
     }
}</pre>
```

#### 6.20 Partition Function

# 6.21 $\lfloor \frac{n}{i} \rfloor$ Enumeration

```
T_0 = 1, T_i = \lfloor \frac{n}{\lfloor \frac{n}{T_{i-1}+1} \rfloor} \rfloor
```

# 6.22 De Bruijn Sequence

```
int res[kN], aux[kN], a[kN], sz;
void Rec(int t, int p, int n, int k) {
   if (t > n) {
     if (n \% p == 0)
       for (int i = 1; i <= p; ++i) res[sz++] = aux[i];</pre>
  } else {
     aux[t] = aux[t - p];
     Rec(t + 1, p, n, k);
     for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t]) Rec(t +
       1, t, n, k);
  }
int DeBruijn(int k, int n) {
   // return cyclic string of length k^n such that every string
     of length n using k character appears as a substring.
   if (k == 1) return res[0] = 0, 1;
fill(aux, aux + k * n, 0);
   return sz = 0, Rec(1, 1, n, k), sz;
```

#### 6.23 Extended GCD

```
template <typename T> tuple<T, T, T> extgcd(T a, T b) {
   if (!b) return make_tuple(a, 1, 0);
   T d, x, y;
   tie(d, x, y) = extgcd(b, a % b);
   return make_tuple(d, y, x - (a / b) * y);
}
```

# 6.24 Euclidean Algorithms

```
• m = \lfloor \frac{an+b}{c} \rfloor
```

• Time complexity:  $O(\log n)$ 

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ + f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm - f(c, c - b - 1, a, m - 1), & \text{otherwise} \end{cases} \end{split}$$

$$\begin{split} g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c, b \bmod c, c, n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ \frac{1}{2} \cdot (n(n+1)m - f(c, c-b-1, a, m-1) \\ -h(c, c-b-1, a, m-1)), & \text{otherwise} \end{cases} \end{split}$$

```
\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ + \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ + h(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ + 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c-b-1, a, m-1) \\ - 2f(c, c-b-1, a, m-1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}
```

### 6.25 Chinese Remainder Theorem

```
long long crt(vector<int> mod, vector<int> a) {
   long long mult = mod[0];
   int n = (int)mod.size();
   long long res = a[0];
   for (int i = 1; i < n; ++i) {
      long long d, x, y;
      tie(d, x, y) = extgcd(mult, mod[i] * 111);
      if ((a[i] - res) % d) return -1;
      long long new_mult = mult / __gcd(mult, 111 * mod[i]) * mod
      [i];
      res += x * ((a[i] - res) / d) % new_mult * mult % new_mult;
      mult = new_mult;
      ((res %= mult) += mult) %= mult;
   }
   return res;
}</pre>
```

#### 6.26 Theorem

#### 6.26.1 Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i, j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

#### 6.26.2 Tutte's Matrix

Let D be a  $n \times n$  matrix, where  $d_{ij} = x_{ij}$   $(x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}$ .  $\frac{rank(D)}{2}$  is the maximum matching on G.

#### 6.26.3 Cayley's Formula

- Given a degree sequence  $d_1, d_2, \ldots, d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
- Let  $T_{n,k}$  be the number of labeled forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k}=kn^{n-k-1}$ .

#### 6.26.4 Erdős-Gallai Theorem

A sequence of non-negative integers  $d_1 \geq d_2 \geq \ldots \geq d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if  $d_1+d_2+\ldots+d_n$  is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all  $1 \leq k \leq n$ .

#### 6.27 Primes

 $\begin{array}{l} 97,101,131,487,593,877,1087,1187,1487,1787,3187,12721,\\ 13331,14341,75577,123457,222557,556679,999983,\\ 1097774749,1076767633,100102021,999997771,\\ 1001010013,1000512343,987654361,999991231,\\ 999888733,98789101,987777733,999991921,1000000007,\\ 10000000087,1000000123,1010101333,1010102101,\\ 100000000039,10000000000037,2305843009213693951,\\ 4611686018427387847,9223372036854775783,\\ 18446744073709551557\end{array}$ 

# 7 Dynamic Programming

#### 7.1 Dynamic Convex Hull

```
struct Line {
          mutable int64_t a, b, p;
          bool operator<(const Line &rhs) const { return a < rhs.a; }</pre>
          bool operator<(int64_t x) const { return p < x; }</pre>
struct DynamicHull : multiset<Line, less<>>> {
          static const int64_t kInf = 1e18;
           int64_t Div(int64_t a, int64_t b) { return a / b - ((a \land b) < (a \land b) < (a
                               0 && a % b); }
          bool Isect(iterator x, iterator y) {
                     if (y == end()) { x->p = kInf; return false; }
                     if (x->a == y->a) x->p = x->b > y->b ? kInf : -kInf;
                      else x->p = Div(y->b - x->b, x->a - y->a);
                     return x->p >= y->p;
           void Insert(int64_t a, int64_t b) {
                   auto z = insert({a, b, 0}), y = z++, x = y;
while (Isect(y, z)) z = erase(z);
if (x != begin() && Isect(--x, y)) Isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p) Isect(x,
                          erase(y));
         int64_t Query(int64_t x) {
  auto l = *lower_bound(x);
  return l.a * x + l.b;
```

### 7.2 1D/1D Convex Optimization

```
struct segment {
  int i, 1, r;
  segment(int a, int b, int c): i(a), l(b), r(c) {}
inline long long f(int l, int r) { return dp[l] + w(l + 1, r);
void solve() {
  dp[0] = 011;
  deque<segment> deq; deq.push_back(segment(0, 1, n));
  for (int i = 1; i <= n; ++i) {
    dp[i] = f(deq.front().i, i);
    while (deq.size() && deq.front().r < i + 1) deq.pop_front()</pre>
    deq.front().l = i + 1;
    segment seg = segment(i, i + 1, n);
while (deq.size() && f(i, deq.back().l) < f(deq.back().i,</pre>
     deq.back().1)) deq.pop_back();
    if (deq.size()) {
      int d = 1048576, c = deq.back().1;
      while (d \gg 1) if (c + d \ll deq.back().r) {
         if (f(i, c + d) > f(deq.back().i, c + d)) c += d;
      deq.back().r = c; seq.l = c + 1;
    if (seg.l <= n) deq.push_back(seg);</pre>
```

#### 7.3 Condition

#### 7.3.1 Totally Monotone (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \ B[i][j] \leq B[i'][j] \implies B[i][j'] \leq B[i'][j'] \\ \forall i < i', j < j', \ B[i][j] \geq B[i'][j] \implies B[i][j'] \geq B[i'][j'] \end{array}
```

#### 7.3.2 Monge Condition (Concave/Convex)

```
\begin{array}{l} \forall i < i', j < j', \, B[i][j] + B[i'][j'] \geq B[i][j'] + B[i'][j] \\ \forall i < i', j < j', \, B[i][j] + B[i'][j'] \leq B[i][j'] + B[i'][j] \end{array}
```

#### 7.3.3 Optimal Split Point

```
If B[i][j] + B[i+1][j+1] \geq B[i][j+1] + B[i+1][j] then H_{i,j-1} \leq H_{i,j} \leq H_{i+1,j}
```

# 8 Geometry

#### 8.1 Basic

```
bool same(double a, double b) { return abs(a - b) < eps; }</pre>
struct P {
   double x, y;
   P() : x(0), y(0) {}
  P(double x, double y): x(x), y(y) {}
P operator + (P b) { return P(x + b.x, y + b.y); }
P operator - (P b) { return P(x - b.x, y - b.y); }
P operator * (double b) { return P(x * b, y * b); }
P operator / (double b) { return P(x / b, y / b); }
   double operator * (P b) { return x * b.x + y * b.y; }
double operator ^ (P b) { return x * b.y - y * b.x; }
   double abs() { return hypot(x, y); }
   P unit() { return *this / abs(); }
   P spin(double o) {
      double c = cos(o), s = sin(o);
      return P(c * x - s * y, s * x + c * y);
   double angle() { return atan2(y, x); }
struct L {
   // ax + by + c = 0
  double a, b, c, o; P pa, pb;
   L(): a(0), b(0), c(0), o(0), pa(), pb() {}
   L(P pa, P pb) : a(pa.y - pb.y), b(pb.x - pa.x), c(pa ^ pb), o
       (atan2(-a, b)), pa(pa), pb(pb) {}
     project(P p) { return pa + (pb - pa).unit() * ((pb - pa) * (p - pa) / (pb - pa).abs()); }
  P reflect(P p) { return p + (project(p) - p) * 2; }
```

```
double get_ratio(P p) { return (p - pa) * (pb - pa) / ((pb -
                                                                                                   /* Delaunay Triangulation:
       pa).abs() * (pb - pa).abs()); }
                                                                                                        Given a sets of points on 2D plane, find a
};
                                                                                                        triangulation such that no points will strictly
                                                                                                        inside circumcircle of any triangle.
bool SegmentIntersect(P p1, P p2, P p3, P p4) {
                                                                                                     find : return a triangle contain given point
  if (\max(p1.x, p2.x) < \min(p3.x, p4.x) \mid | \max(p3.x, p4.x) <
                                                                                                     add_point : add a point into triangulation
       min(p1.x, p2.x)) return false;
                                                                                                     A Triangle is in triangulation iff. its has_chd is 0.
   if (max(p1.y, p2.y) < min(p3.y, p4.y) | | max(p3.y, p4.y) < 
                                                                                                     Region of triangle u: iterate each u.edge[i].tri,
       min(p1.y, p2.y)) return false
                                                                                                     each points are u.p[(i+1)%3], u.p[(i+2)%3]
   return sign((p3 - p1) \land (p4 - p1)) * sign((p3 - p2) \land (p4 - p4)) * sign((p3 - p2) \land (p4 - p4)) * sign((p3 - p4)) * sign((p4 - p4)) * sig
                                                                                                     calculation involves O(|V|^6) */
       p2)) <= 0 &&
                                                                                                     const double inf = 1e9;
        sign((p1 - p3) \land (p2 - p3)) * sign((p1 - p4) \land (p2 - p4))
                                                                                                     double eps = 1e-6; // 0 when integer
                                                                                                     // return p4 is in circumcircle of tri(p1,p2,p3)
}
                                                                                                     bool in_cc(P &p1, P &p2, P &p3, P &p4) {
                                                                                                      int o1 = (abs(p1.x) >= inf * 0.99 | | abs(p1.y) >= inf * 0.99);
bool parallel(L x, L y) { return same(x.a * y.b, x.b * y.a); }
                                                                                                       int o2 = (abs(p2.x) >= inf * 0.99 | | abs(p2.y) >= inf * 0.99);
                                                                                                       int o3 = (abs(p3.x) >= inf * 0.99 | | abs(p3.y) >= inf * 0.99);
P Intersect(L x, L y) { return P(-x.b * y.c + x.c * y.b, x.a *
                                                                                                       int rtrue = 01 + 02 + 03;
       y.c - x.c * y.a) / (-x.a * y.b + x.b * y.a); }
                                                                                                       int rfalse = abs(p4.x) >= inf * 0.99 || abs(p4.y) >= inf *
                                                                                                            0.99;
8.2 KD Tree
                                                                                                       if (rtrue == 3) return true;
                                                                                                       if (rtrue) {
namespace kdt {
                                                                                                        P in(0, 0), out(0, 0);
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn], yl[maxn], yr[
                                                                                                        if (o1) out = out + p1; else in = in + p1;
       maxn];
                                                                                                        if (o2) out = out + p2; else in = in + p2;
point p[maxn];
                                                                                                        if (o3) out = out + p3; else in = in + p3;
int build(int 1, int r, int dep = 0) {
                                                                                                        return (p4 - in) * (out - in) > 0;
   if (l == r) return -1;
   function<bool(const point &, const point &)> f = [dep](const
                                                                                                       if (rfalse) return false;
      point &a, const point &b) {
      if (dep & 1) return a.x < b.x;</pre>
                                                                                                      double u11 = p1.x - p4.x, u12 = p1.y - p4.y;
double u21 = p2.x - p4.x, u22 = p2.y - p4.y;
double u31 = p3.x - p4.x, u32 = p3.y - p4.y;
      else return a.y < b.y;
   int m = (l + r) >> 1;
                                                                                                       double u13 = sq(p1.x) - sq(p4.x) + sq(p1.y) -
                                                                                                                                                                             sq(p4.v);
   nth_element(p + l, p + m, p + r, f);
                                                                                                      double u23 = sq(p2.x) - sq(p4.x) + sq(p2.y) - sq(p4.y);
                                                                                                      double u33 = sq(p3.x) - sq(p4.x) + sq(p3.y) - sq(p4.y);
double det = -u13 * u22 * u31 + u12 * u23 * u31 + u13 * u21 *
u32 - u11 * u23 * u32 - u12 * u21 * u33 + u11 * u22 * u33;
   xl[m] = xr[m] = p[m].x;
   yl[m] = yr[m] = p[m].y;
   [c[m]] = build(l, m, dep + 1);
                                                                                                       return det > eps;
   if (~lc[m]) {
      xl[m] = min(xl[m], xl[lc[m]]);
                                                                                                     double side(P &a, P &b, P &p) { return (b - a) ^ (p - a); }
      xr[m] = max(xr[m], xr[lc[m]]);
                                                                                                     struct Tri;
      yl[m] = min(yl[m], yl[lc[m]]);
                                                                                                     struct Edge {
      yr[m] = max(yr[m], yr[lc[m]]);
                                                                                                      Tri *tri;
                                                                                                       int side;
   rc[m] = build(m + 1, r, dep + 1);
                                                                                                      Edge() : tri(0), side(0) {}
   if (~rc[m]) {
      xl[m] = min(xl[m], xl[rc[m]]);
                                                                                                      Edge(Tri *_tri, int _side) : tri(_tri), side(_side) {}
                                                                                                     };
      xr[m] = max(xr[m], xr[rc[m]]);
                                                                                                     struct Tri {
      yl[m] = min(yl[m], yl[rc[m]]);
                                                                                                      P p[3];
      yr[m] = max(yr[m], yr[rc[m]]);
                                                                                                      Edge edge[3];
                                                                                                      Tri *ch[3];
   return m;
                                                                                                       Tri() {}
                                                                                                      Tri(P p0, P p1, P p2) {
bool bound(const point &q, int o, long long d) {
                                                                                                        p[0] = p0; p[1] = p1; p[2] = p2;
   double ds = sqrt(d + 1.0);
   if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
                                                                                                        ch[0] = ch[1] = ch[2] = 0;
      q.y < yl[o] - ds || q.y > yr[o] + ds) return false;
   return true;
                                                                                                       bool has_ch() { return ch[0] != 0; }
                                                                                                       int num_ch() {
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
                                                                                                        return ch[0] == 0 ? 0 : ch[1] == 0 ? 1 : ch[2] == 0 ? 2 : 3;
         (a.y - b.y) * 111 * (a.y - b.y);
                                                                                                       bool contains(P &q) {
                                                                                                        for (int i = 0; i < 3; ++i)
void dfs(const point &q, long long &d, int o, int dep = 0) {
                                                                                                          if (side(p[i], p[(i + 1) % 3], q) < -eps) return false;
  if (!bound(q, o, d)) return;
long long cd = dist(p[o], q);
                                                                                                        return true
                                                                                                     } pool[maxn * 10], *tris;
   if (cd != 0) d = min(d, cd);
   if ((dep & 1) && q.x < p[o].x || !(dep & 1) && <math>q.y < p[o].y)
                                                                                                     void edge(Edge a, Edge b) {
                                                                                                      if (a.tri) a.tri->edge[a.side] = b;
                                                                                                       if (b.tri) b.tri->edge[b.side] = a;
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                                                     }
  } else {
                                                                                                     struct Trig {
      if (~rc[o]) dfs(q, d, rc[o], dep + 1);
                                                                                                      Trig() {
                                                                                                        the_root = new (tris++) Tri(P(-inf, -inf), P(inf * 2, -inf),
   P(-inf, inf * 2));
      if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
                                                                                                         // all p should in
                                                                                                       Tri *find(P p) { return find(the_root, p); }
void init(const vector<point> &v) {
                                                                                                       void add_point(P &p) { add_point(find(the_root, p), p); }
   for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
   root = build(0, v.size());
                                                                                                      Tri *the_root;
                                                                                                      static Tri *find(Tri *root, P &p) {
                                                                                                        while (true) {
long long nearest(const point &q) {
  long long res = 1e18;
                                                                                                          if (!root->has_ch()) return root;
   dfs(q, res, root);
                                                                                                          for (int i = 0; i < 3 \&\& root->ch[i]; ++i)
   return res;
                                                                                                           if (root->ch[i]->contains(p)) {
```

root = root->ch[i];

break;

#### 8.3 Delaunay Triangulation

```
assert(false); // "point not found"
 }
  void add_point(Tri *root, P &p) {
  Tri *tab, *tbc, *tca;
  tab = new (tris++) Tri(root->p[0], root->p[1], p);
   tbc = new (tris++) Tri(root->p[1], root->p[2], p);
  tca = new (tris++) Tri(root->p[2], root->p[0], p);
   edge(Edge(tab, 0), Edge(tbc, 1));
   edge(Edge(tbc, 0), Edge(tca, 1));
  edge(Edge(tca, 0), Edge(tab, 1));
  edge(Edge(tab, 2), root->edge[2]);
   edge(Edge(tbc, 2), root->edge[0]);
   edge(Edge(tca, 2), root->edge[1]);
  root->ch[0] = tab; root->ch[1] = tbc; root->ch[2] = tca;
  flip(tab, 2); flip(tbc, 2); flip(tca, 2);
 void flip(Tri *tri, int pi) {
   Tri *trj = tri->edge[pi].tri;
  int pj = tri->edge[pi].side;
   if (!trj) return;
  if (!in_cc(tri->p[0], tri->p[1], tri->p[2], trj->p[pj]))
     return:
   /* flip edge between tri,trj */
  Tri *trk = new (tris++) Tri(tri->p[(pi + 1) % 3], trj->p[pj],
      tri->p[pi]);
   Tri *trl = new (tris++) Tri(trj->p[(pj + 1) % 3], tri->p[pi],
      trj->p[pj]);
   edge(Edge(trk, 0), Edge(trl, 0));
   edge(Edge(trk, 1), tri->edge[(pi + 2) % 3]);
   edge(Edge(trk, 2), trj->edge[(pj + 1) % 3]);
  edge(Edge(trl, 1), trj->edge[(pj + 2) % 3]);
  edge(Edge(trl, 2), tri->edge[(pi + 1) % 3]);
tri->ch[0] = trk; tri->ch[1] = trl; tri->ch[2] = 0;
  trj->ch[0] = trk; trj->ch[1] = trl; trj->ch[2] = 0;
   flip(trk, 1); flip(trk, 2);
   flip(trl, 1); flip(trl, 2);
vector<Tri *> triang;
set<Tri *> vst;
void go(Tri *now) {
 if (vst.find(now) != vst.end()) return;
 vst.insert(now);
 if (!now->has_ch()) {
  triang.push_back(now);
 for (int i = 0; i < now->num_ch(); ++i) go(now->ch[i]);
void build(int n, P *ps) {
 tris = pool
 random\_shuffle(ps, ps + n);
 Trig tri;
 for (int i = 0; i < n; ++i) tri.add_point(ps[i]);</pre>
 go(tri.the_root);
8.4 Voronoi Diagram
int gid(P &p) {
  auto it = ptoid.find(p);
   if (it == ptoid.end()) return -1;
  return it->second;
}
```

```
L make_line(P p, L l) {
  P d = 1.pb - 1.pa; d = d.spin(pi / 2);
  P m = (1.pa + 1.pb) / 2;
  l = L(m, m + d);
  if (((l.pb - l.pa) \land (p - l.pa)) < 0) l = L(m + d, m);
  return 1;
double calc_ans(int i) {
  vector<P> ps = HPI(ls[i]);
  double rt = 0;
  for (int i = 0; i < (int)ps.size(); ++i) {</pre>
    rt += (ps[i] ^ ps[(i + 1) % ps.size()]);
  return abs(rt) / 2;
}
void solve() {
  for (int i = 0; i < n; ++i) ops[i] = ps[i], ptoid[ops[i]] = i
  random\_shuffle(ps, ps + n);
  build(n, ps);
for (auto *t : triang) {
    int z[3] = \{gid(t-p[0]), gid(t-p[1]), gid(t-p[2])\};
```

```
for (int i = 0; i < 3; ++i) for (int j = 0; j < 3; ++j) if
    (i != j && z[i] != -1 && z[j] != -1) {
        L l(t->p[i], t->p[j]);
        ls[z[i]].push_back(make_line(t->p[i], l));
    }
}
vector<P> tb = convex(vector<P>(ps, ps + n));
for (auto &p : tb) isinf[gid(p)] = true;
for (int i = 0; i < n; ++i) {
    if (isinf[i]) cout << -1 << '\n';
    else cout << fixed << setprecision(12) << calc_ans(i) << '\n';
}
}</pre>
```

#### 8.5 Sector Area

```
// calc area of sector which include a, b
double SectorArea(P a, P b, double r) {
  double o = atan2(a.y, a.x) - atan2(b.y, b.x);
  while (o <= 0) o += 2 * pi;
  while (o >= 2 * pi) o -= 2 * pi;
  o = min(o, 2 * pi - o);
  return r * r * o / 2;
}
```

#### 8.6 Half Plane Intersection

```
bool jizz(L l1,L l2,L l3){
  P p=Intersect(12,13);
  return ((l1.pb-l1.pa)^(p-l1.pa))<-eps;</pre>
}
bool cmp(const L &a,const L &b){
  return same(a.o,b.o)?(((b.pb-b.pa)^(a.pb-b.pa))>eps):a.o<b.o;</pre>
// availble area for L l is (l.pb-l.pa)^(p-l.pa)>0
vector<P> HPI(vector<L> &ls){
  sort(ls.begin(),ls.end(),cmp);
   vector<L> pls(1,ls[0]);
   for(int i=0;i<(int)ls.size();++i)if(!same(ls[i].o,pls.back().</pre>
      o))pls.push_back(ls[i]);
   deque<int> dq; dq.push_back(0); dq.push_back(1);
#define meow(a,b,c) while(dq.size()>1u && jizz(pls[a],pls[b],
     pls[c]))
   for(int i=2;i<(int)pls.size();++i){</pre>
     meow(i,dq.back(),dq[dq.size()-2])dq.pop_back();
     meow(i,dq[0],dq[1])dq.pop_front();
     dq.push_back(i);
  meow(dq.front(),dq.back(),dq[dq.size()-2])dq.pop_back();
   meow(dq.back(),dq[0],dq[1])dq.pop_front();
   if(dq.size()<3u)return vector<P>(); // no solution or
      solution is not a convex
   vector<P> rt;
   for(int i=0;i<(int)dq.size();++i)rt.push_back(Intersect(pls[</pre>
      dq[i]],pls[dq[(i+1)%dq.size()]]));
   return rt:
}
```

#### 8.7 Rotating Sweep Line

```
void rotatingSweepLine(vector<pair<int,int>> &ps){
  int n=int(ps.size());
  vector<int> id(n),pos(n);
  vector<pair<int,int>> line(n*(n-1)/2);
  int m=-1;
  for(int i=0;i<n;++i)for(int j=i+1;j<n;++j)line[++m]=make_pair</pre>
     (i,j); ++m;
  sort(line.begin(),line.end(),[&](const pair<int,int> &a,const
    pair<int, int> &b)->bool{
if(ps[a.first].first==ps[a.second].first)return 0;
    if(ps[b.first].first==ps[b.second].first)return 1;
    return (double)(ps[a.first].second-ps[a.second].second)/(ps
     [a.first].first-ps[a.second].first) < (double)(ps[b.first
     ].second-ps[b.second].second)/(ps[b.first].first-ps[b.
     second].first);
  for(int i=0;i<n;++i)id[i]=i;</pre>
  sort(id.begin(),id.end(),[&](const int &a,const int &b){
     return ps[a]<ps[b]; });</pre>
  for(int i=0;i<n;++i)pos[id[i]]=i;</pre>
  for(int i=0;i<m;++i){</pre>
```

```
auto l=line[i];
     // meow
    tie(pos[l.first],pos[l.second],id[pos[l.first]],id[pos[l.
     second]])=make_tuple(pos[l.second],pos[l.first],l.second,l
i}
```

#### Triangle Center 8.8

```
Point TriangleCircumCenter(Point a, Point b, Point c) {
  Point res;
  double a1 = atan2(b.y - a.y, b.x - a.x) + pi / 2;
  double a2 = atan2(c.y - b.y, c.x - b.x) + pi / 2;
  double ax = (a.x + b.x) / 2;
  double ay = (a.y + b.y) / 2;
  double bx = (c.x + b.x) / 2;
  double by = (c.y + b.y) / 2;
double r1 = (sin(a2) * (ax - bx) + cos(a2) * (by - ay)) / (
    sin(a1) * cos(a2) - sin(a2) * cos(a1));
  return Point(ax + r1 * cos(a1), ay + r1 * sin(a1));
Point TriangleMassCenter(Point a, Point b, Point c) {
  return (a + b + c) / 3.0;
Point TriangleOrthoCenter(Point a, Point b, Point c) {
  return TriangleMassCenter(a, b, c) * 3.0 -
     TriangleCircumCenter(a, b, c) * 2.0;
Point TriangleInnerCenter(Point a, Point b, Point c) {
  Point res;
  double la = len(b - c);
  double lb = len(a - c);
  double lc = len(a - b);

res.x = (la * a.x + lb * b.x + lc * c.x) / (la + lb + lc);
  res.y = (la * a.y + lb * b.y + lc * c.y) / (la + lb + lc);
  return res;
```

### Polygon Center

```
Point BaryCenter(vector<Point> &p, int n) {
 Point res(0, 0);
 double s = 0.0, t;
for (int i = 1; i < p.size() - 1; i++) {
    t = Cross(p[i] - p[0], p[i + 1] - p[0]) / 2;
    res.x += (p[0].x + p[i].x + p[i + 1].x) * t;
    res.y += (p[0].y + p[i].y + p[i + 1].y) * t;
 res.x /= (3 * s);
 res.y = (3 * s);
 return res;
```

#### 8.10 Maximum Triangle

```
double ConvexHullMaxTriangleArea(Point p[], int res[], int
    chnum) {
  double area = 0, tmp;
  res[chnum] = res[0];
  for (int i = 0, j = 1, k = 2; i < chnum; i++) {
   while (fabs(Cross(p[res[j]] - p[res[i]], p[res[(k + 1) %
    chnum]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
    p[res[k]] - p[res[i]])) k = (k + 1) % chnum;
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
    11));
   if (tmp > area) area = tmp;
while (fabs(Cross(p[res[(j + 1) % chnum]] - p[res[i]], p[
    res[k]] - p[res[i]])) > fabs(Cross(p[res[j]] - p[res[i]],
    p[res[k]] - p[res[i]]))) j = (j + 1) % chnum;
    tmp = fabs(Cross(p[res[j]] - p[res[i]], p[res[k]] - p[res[i
    11));
   if (tmp > area) area = tmp;
  return area / 2;
```

#### 8.11 Point in Polygon

```
int pip(vector<P> ps, P p) {
 int c = 0;
  for (int i = 0; i < ps.size(); ++i) {</pre>
    int a = i, b = (i + 1) \% ps.size();
```

```
20
     L l(ps[a], ps[b]);
     P q = l.project(p);
     if ((p - q).abs() < eps && l.inside(q)) return 1;</pre>
     if (same(ps[a].y, ps[b].y) && same(ps[a].y, p.y)) continue;
     if (ps[a].y > ps[b].y) swap(a, b);
     if (ps[a].y \le p.y \&\& p.y < ps[b].y \&\& p.x \le ps[a].x + (ps
      [b].x - ps[a].x) / (ps[b].y - ps[a].y) * (p.y - ps[a].y))
      ++C:
  }
   return (c & 1) * 2;
}
        Circle
8.12
struct C {
  Рс;
   double r;
   C(P \ c = P(0, 0), double \ r = 0) : c(c), r(r) \{\}
vector<P> Intersect(C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
   vector<P> p;
   if (same(a.r + b.r, d)) p.push_back(a.c + (b.c - a.c).unit()
      * a.r);
   else if (a.r + b.r > d \&\& d + a.r >= b.r) {
     double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
     P i = (b.c - a.c).unit();
     p.push_back(a.c + i.spin(o) * a.r);
     p.push_back(a.c + i.spin(-o) * a.r);
  }
   return p:
double IntersectArea(C a, C b) {
   if (a.r > b.r) swap(a, b);
   double d = (a.c - b.c).abs();
   if (d >= a.r + b.r - eps) return 0;
   if (d + a.r \le b.r + eps) return sq(a.r) * acos(-1);
  double p = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d));
double q = acos((sq(b.r) + sq(d) - sq(a.r)) / (2 * b.r * d));
   return p * sq(a.r) + q * sq(b.r) - a.r * d * sin(p);
}
// remove second level if to get points for line (defalut:
      segment)
vector<P> CircleCrossLine(P a, P b, P o, double r) {
  double x = b.x - a.x, y = b.y - a.y, A = sq(x) + sq(y), B = 2 * x * (a.x - o.x) + 2 * y * (a.y - o.y);
   double C = sq(a.x - o.x) + sq(a.y - o.y) - sq(r), d = B * B - 4 * A * C;
   vector<P> t;
   if (d >= -eps) {
     d = max(0., d);
     double i = (-B - sqrt(d)) / (2 * A);
     double j = (-B + sqrt(d)) / (2 * A);
     if (i - 1.0 \le eps \&\& i \ge -eps) t.emplace_back(a.x + i * x)
       a.y + i * y);
     if (j - 1.0 \le eps \& j \ge -eps) t.emplace_back(a.x + j * x)
     , a.y + j * y);
   return t;
}
// calc area intersect by circle with radius r and triangle OAB
double AreaOfCircleTriangle(P a, P b, double r) {
  bool ina = a.abs() < r, inb = b.abs() < r;
   auto p = CircleCrossLine(a, b, P(0, 0), r);
   if (ina) {
     if (inb) return abs(a ^ b) / 2;
     return SectorArea(b, p[0], r) + abs(a p[0]) / 2;
   if (inb) return SectorArea(p[0], a, r) + abs(p[0] \land b) / 2;
   if (p.size() == 2u) return SectorArea(a, p[0], r) +
    SectorArea(p[1], b, r) + abs(p[0] ^ p[1]) / 2;
   else return SectorArea(a, b, r);
// for any triangle
double AreaOfCircleTriangle(vector<P> ps, double r) {
   double ans = 0;
   for (int i = 0; i < 3; ++i) {
     int j = (i + 1) \% 3;
     double o = atan2(ps[i].y, ps[i].x) - atan2(ps[j].y, ps[j].x
     if (o >= pi) o = o - 2 * pi;
```

if  $(o \le -pi) o = o + 2 * pi;$ 

: -1);

ans += AreaOfCircleTriangle(ps[i], ps[j], r) \* (o >= 0 ? 1

```
| }
| return abs(ans);
|}
```

#### 8.13 Tangent of Circles and Points to Circle

```
vector<L> tangent(C a, C b) {
 #define Pij \
   P i = (b.c - a.c).unit() * a.r, j = P(i.y, -i.x); 
   z.emplace_back(a.c + i, a.c + i + j);
 #define deo(I,J) \
   double d = (a.c - b.c).abs(), e = a.r I b.r, o = acos(e / d)
   P i = (b.c - a.c).unit(), j = i.spin(o), k = i.spin(-o);\
z.emplace_back(a.c + j * a.r, b.c J j * b.r);\
   z.emplace_back(a.c + k * a.r, b.c J k * b.r);
   if (a.r < b.r) swap(a, b);
   vector<L> z;
   if ((a.c - b.c).abs() + b.r < a.r) return z;</pre>
   else if (same((a.c - b.c).abs() + b.r, a.r)) { Pij; }
   else {
     deo(-,+);
     if (same(d, a.r + b.r)) { Pij; }
     else if (d > a.r + b.r) \{ deo(+,-); \}
 }
 vector<L> tangent(C c, P p) {
   vector<L> z;
double d = (p - c.c).abs();
   if (same(d, c.r)) {
  P i = (p - c.c).spin(pi / 2);
     z.emplace_back(p, p + i);
   } else if (d > c.r) {
     double o = acos(c.r / d);
     P i = (p - c.c).unit(), j = i.spin(o) * c.r, k = i.spin(-o)
       * c.r;
     z.emplace_back(c.c + j, p);
     z.emplace_back(c.c + k, p);
   return z;
1}
```

#### 8.14 Area of Union of Circles

```
vector<pair<double, double>> CoverSegment(C &a, C &b) {
  double d = (a.c - b.c).abs();
       vector<pair<double, double>> res;
       if (same(a.r + b.r, d));
      else if (d \le abs(a.r - b.r) + eps) {
             if (a.r < b.r) res.emplace_back(0, 2 * pi);</pre>
      } else if (d < abs(a.r + b.r) - eps) {
             double o = acos((sq(a.r) + sq(d) - sq(b.r)) / (2 * a.r * d)
               ), z = (b.c - a.c).angle();
            if (z < 0) z += 2 * pi;
double l = z - o, r = z + o;
if (l < 0) l += 2 * pi;
if (r > 2 * pi) r -= 2 * pi;
             if (l > r) res.emplace_back(l, 2 * pi), res.emplace_back(0,
             else res.emplace_back(l, r);
      return res;
}
double CircleUnionArea(vector<C> c) { // circle should be
               identical
       int n = c.size();
       double a = 0, w;
       for (int i = 0; w = 0, i < n; ++i) {
             vector<pair<double, double>> s = \{\{2 * pi, 9\}\}, z; for (int j = 0; j < n; ++j) if (i != j) {
                   z = CoverSegment(c[i], c[j]);
                    for (auto &e : z) s.push_back(e);
             sort(s.begin(), s.end());
auto F = [&] (double t) { return c[i].r * (c[i].r * t + c[i].r * 
               ].c.x * sin(t) - c[i].c.y * cos(t)); };
              for (auto &e : s) {
                    if (e.first > w) a += F(e.first) - F(w);
                    w = max(w, e.second);
      return a * 0.5;
```

### 8.15 Minimun Distance of 2 Polygons

#### 8.16 2D Convex Hull

```
| bool operator < (const P &a, const P &b) { return same(a.x, b.x
      ) ? a.y < b.y : a.x < b.x; }
bool operator > (const P &a, const P &b) { return same(a.x, b.x
      ) ? a.y > b.y : a.x > b.x; }
#define crx(a, b, c) ((b - a) \land (c - a))
vector<P> convex(vector<P> ps) {
   vector<P> p;
   sort(ps.begin(), ps.end(), [&] (P a, P b) { return same(a.x,
      b.x) ? a.y < b.y : a.x < b.x; });
   for (int i = 0; i < ps.size(); ++i) {
     while (p.size() \ge 2 \& crx(p[p.size() - 2], ps[i], p[p.
      size() - 1]) >= 0) p.pop_back();
     p.push_back(ps[i]);
   int t = p.size();
   for (int i = (int)ps.size() - 2; i >= 0; --i) {
  while (p.size() > t && crx(p[p.size() - 2], ps[i], p[p.size
      () - 1]) >= 0) p.pop_back();
     p.push_back(ps[i]);
   p.pop_back();
   return p;
}
int sgn(double x) { return same(x, 0) ? 0 : x > 0 ? 1 : -1; }
P isLL(P p1, P p2, P q1, P q2) {
   double a = crx(q1, q2, p1), b = -crx(q1, q2, p2);
return (p1 * b + p2 * a) / (a + b);
struct CH {
   int n;
   vector<P> p, u, d;
   CH() {}
   CH(vector<P> ps) : p(ps) {
     n = ps.size();
     rotate(p.begin(), min_element(p.begin(), p.end()), p.end())
     auto t = max_element(p.begin(), p.end());
     d = vector<P>(p.begin(), next(t));
     u = vector < P > (t, p.end()); u.push_back(p[0]);
   int find(vector<P> &v, P d) {
     int l = 0, r = v.size();
     int L = (1 + 5 < r) {
  int L = (1 * 2 + r) / 3, R = (1 + r * 2) / 3;
  if (v[L] * d > v[R] * d) r = R;
  else l = L;
     int x = 1;
     for (int i = l + 1; i < r; ++i) if (v[i] * d > v[x] * d) x
      = i:
     return x;
   int findFarest(P v) {
     if (v.y > 0 \mid \mid v.y == 0 \& v.x > 0) return ((int)d.size() - v.y == 0 & v.x > 0)
       1 + find(u, v)) % p.size();
     return find(d, v);
   P get(int 1, int r, P a, P b) {
```

```
int s = sgn(crx(a, b, p[l % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(a, b, p[m % n])) == s) l = m;
     return isLL(a, b, p[l % n], p[(l + 1) % n]);
   }
   vector<P> getIS(P a, P b) {
     int X = findFarest((b - a).spin(pi / 2));
     int Y = findFarest((a - b).spin(pi / 2));
     if (X > Y) swap(X, Y);
     if (sgn(crx(a, b, p[X])) * sgn(crx(a, b, p[Y])) < 0) return
       \{get(X, Y, a, b), get(Y, X + n, a, b)\};
     return {};
   void update_tangent(P q, int i, int &a, int &b) {
     if (sgn(crx(q, p[a], p[i])) > 0) a = i;
if (sgn(crx(q, p[b], p[i])) < 0) b = i;
   void bs(int l, int r, P q, int &a, int &b) {
     if (l == r) return;
     update_tangent(q, 1 % n, a, b);
     int s = sgn(crx(q, p[l % n], p[(l + 1) % n]));
     while (l + 1 < r) {
       int m = (l + r) >> 1;
       if (sgn(crx(q, p[m % n], p[(m + 1) % n])) == s) l = m;
     update_tangent(q, r % n, a, b);
   bool contain(P p) {
     if (p.x < d[0].x \mid | p.x > d.back().x) return 0;
     auto it = lower_bound(d.begin(), d.end(), P(p.x, -1e12));
     if (it->x == p.x) {
       if (it->y > p.y) return 0;
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     it = lower_bound(u.begin(), u.end(), P(p.x, 1e12), greater<</pre>
     P>());
     if (it->x == p.x) {
       if (it->y < p.y) return 0;</pre>
     } else if (crx(*prev(it), *it, p) < -eps) return 0;</pre>
     return 1;
   bool get_tangent(P p, int &a, int &b) { // b -> a
     if (contain(p)) return 0;
     a = b = 0;
     int i = lower_bound(d.begin(), d.end(), p) - d.begin();
     bs(0, i, p, a, b);
     bs(i, d.size(), p, a, b);
     i = lower_bound(u.begin(), u.end(), p, greater < P > ()) - u.
     beain():
     bs((int)d.size() - 1, (int)d.size() - 1 + i, p, a, b);
     bs((int)d.size() - 1 + i, (int)d.size() - 1 + u.size(), p,
     a, b);
     return 1;
|};
```

# 8.17 3D Convex Hull

```
double absvol(const P a,const P b,const P c,const P d) {
 return abs(((b-a)^{(c-a)})*(d-a))/6;
struct convex3D {
static const int maxn=1010;
struct T{
  int a,b,c;
  bool res;
 T(){}
 T(int a, int b, int c, bool res=1):a(a), b(b), c(c), res(res){}
int n,m;
P p[maxn];
T f[maxn*8];
int id[maxn][maxn];
bool on(T &t,P &q){
 return ((p[t.c]-p[t.b])^(p[t.a]-p[t.b]))*(q-p[t.a])>eps;
void meow(int q,int a,int b){
  int g=id[a][b];
  if(f[g].res){
    if(on(f[g],p[q]))dfs(q,g);
      id[q][b]=id[a][q]=id[b][a]=m;
      f[m++]=T(b,a,q,1);
```

```
}
}
void dfs(int p,int i){
   f[i].res=0;
   meow(p,f[i].b,f[i].a);
   meow(p,f[i].c,f[i].b);
   meow(p,f[i].a,f[i].c);
void operator()(){
   if(n<4)return;</pre>
   if([&](){
     for(int i=1;i<n;++i)if(abs(p[0]-p[i])>eps)return swap(p[1],
      p[i]),0;
     return 1
   }() || [&](){
     for(int i=2;i<n;++i)if(abs((p[0]-p[i])^(p[1]-p[i]))>eps)
      return swap(p[2],p[i]),0;
     return 1:
   }() || [&](){
     for(int i=3; i< n; ++i)if(abs(((p[1]-p[0])^(p[2]-p[0]))*(p[i]-
      p[0]))>eps)return swap(p[3],p[i]),0;
      eturn 1;
   }())return;
   for(int i=0;i<4;++i){</pre>
     T t((i+1)\%4,(i+2)\%4,(i+3)\%4,1);
     if(on(t,p[i]))swap(t.b,t.c);
     id[t.a][t.b]=id[t.b][t.c]=id[t.c][t.a]=m;
     f\lceil m++\rceil=t;
   for(int i=4; i< n; ++i)for(int j=0; j< m; ++j)if(f[j].res && on(f[j])
      ],p[i])){
     dfs(i,j);
     break:
   int mm=m; m=0;
   for(int i=0;i<mm;++i)if(f[i].res)f[m++]=f[i];</pre>
bool same(int i,int j){
   return !(absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].a])>eps
      || absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].b])>eps ||
      absvol(p[f[i].a],p[f[i].b],p[f[i].c],p[f[j].c])>eps);
int faces(){
   int r=0;
   for(int i=0;i<m;++i){</pre>
     int iden=1;
     for(int j=0;j<i;++j)if(same(i,j))iden=0;</pre>
     r+=iden;
   return r;
} tb;
```

#### 8.18 Minimum Enclosing Circle

```
pt center(const pt &a, const pt &b, const pt &c) {
  pt p0 = b - a, p1 = c - a;
double c1 = norm2(p0) * 0.5, c2 = norm2(p1) * 0.5;
  double d = p0 \land p1;
  double x = a.x + (c1 * p1.y - c2 * p0.y) / d;
  double y = a.y + (c2 * p0.x - c1 * p1.x) / d;
  return pt(x, y);
circle min_enclosing(vector<pt> &p) {
  random_shuffle(p.begin(), p.end());
  double r = 0.0;
  pt cent;
  for (int i = 0; i < p.size(); ++i) {</pre>
    if (norm2(cent - p[i]) <= r) continue;</pre>
    cent = p[i];
     r = 0.0;
    for (int j = 0; j < i; ++j) {
       if (norm2(cent - p[j]) <= r) continue;</pre>
      cent = (p[i] + p[j]) / 2;
r = norm2(p[j] - cent);
       for (int k = 0; k < j; ++k) {
         if (norm2(cent - p[k]) <= r) continue;</pre>
         cent = center(p[i], p[j], p[k]);
         r = norm2(p[k] - cent);
    }
  }
  return circle(cent, sqrt(r));
```

#### 8.19 Closest Pair

```
double closest_pair(int l, int r) {
  // p should be sorted increasingly according to the x-
      coordinates.
   if (l == r) return 1e9;
  if (r - l == 1) return dist(p[l], p[r]);
  int m = (l + r) >> 1;
   double d = min(closest_pair(l, m), closest_pair(m + 1, r));
   vector<int> vec;
  for (int i = m; i >= 1 && fabs(p[m].x - p[i].x) < d; --i) vec
      .push_back(i);
   for (int i = m + 1; i \le r \&\& fabs(p[m].x - p[i].x) < d; ++i)
       vec.push_back(i);
   sort(vec.begin(), vec.end(), [&](int a, int b) { return p[a].
     y < p[b].y; \});
  for (int i = 0; i < vec.size(); ++i) {
    for (int j = i + 1; j < vec.size() && fabs(p[vec[j]].y - p[
     vec[i]].y) < d; ++j) {
       d = min(d, dist(p[vec[i]], p[vec[j]]));
   return d:
į }
```

### 9 Miscellaneous

#### 9.1 Bitwise Hack

```
long long next_perm(long long v) {
  long long t = v | (v - 1);
  return (t + 1) | (((~t & -~t) - 1) >> (__builtin_ctz(v) + 1))
    ;
  }
  void subset(long long s) {
    long long sub = s;
    while (sub) sub = (sub - 1) & s;
}
```

#### 9.2 Hilbert's Curve (faster Mo's algorithm)

```
long long hilbert(int n, int x, int y) {
  long long res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x & s) > 0;
    int ry = (y & s) > 0;
    res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) x = s - 1 - x, y = s - 1 - y;
        swap(x, y);
    }
  }
  return res;
}
```

#### 9.3 Mo's Algorithm on Tree

```
void MoAlgoOnTree() {
   Dfs(0, -1);
   vector<int> euler(tk);
   for (int i = 0; i < n; ++i) {
  euler[tin[i]] = i;</pre>
      euler[tout[i]] = i;
   vector<int> l(q), r(q), qr(q), sp(q, -1);
   for (int i = 0; i < q; ++i) {
  if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
      int z = GetLCA(u[i], v[i]);
      sp[i] = z[i];
      if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
      else l[i] = tout[u[i]], r[i] = tin[v[i]];
      qr[i] = i;
   sort(qr.begin(), qr.end(), [&](int i, int j) {
   if (l[i] / kB == l[j] / kB) return r[i] < r[j];</pre>
      return l[i] / kB < l[j] / kB;</pre>
   vector<bool> used(n);
   // Add(v): add/remove v to/from the path based on used[v]
   for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
      while (tl < l[qr[i]]) Add(euler[tl++]);
while (tl > l[qr[i]]) Add(euler[--tl]);
      while (tr > r[qr[i]]) Add(euler[tr--]);
      while (tr < r[qr[i]]) Add(euler[++tr]);</pre>
      // add/remove LCA(u, v) if necessary
   }
| }
```

#### 9.4 Java

```
import java.io.*;
 import java.util.*;
 import java.lang.*;
 import java.math.*;
 public class filename{
   static Scanner in = new Scanner(System.in);
   public static void main(String[] args) throws Exception {
   Scanner fin = new Scanner(new File("infile"));
     PrintWriter fout = new PrintWriter("outfile",
     fout.println(fin.nextLine());
     fout.close();
     while (in.hasNext()) {
        String str = in.nextLine(); // getline
        String stu = in.next(); // string
     System.out.println("Case #" + t);
     System.out.printf("%d\n", 7122);
     int[] d = \{\{7,1,2,2\},\{8,7\}\};
     int g = Integer.parseInt("-123");
     long f = (long)d[0][2];
     List<Integer> l = new ArrayList<>();
     Random rg = new Random();
for (int i = 9; i >= 0; --i) {
        l.add(Integer.valueOf(rg.nextInt(100) + 1));
        l.add(Integer.valueOf((int)(Math.random() * 100) + 1));
     Collections.sort(l, new Comparator<Integer>() {
        public int compare(Integer a, Integer b) { return a - b;
     for (int i = 0; i < l.size(); ++i)</pre>
        System.out.print(l.get(i));
     Set<String> s = new HashSet<String>(); // TreeSet
     s.add("jizz");
     System.out.println(s);
     System.out.println(s.contains("jizz"));
     Map<String, Integer> m = new HashMap<String, Integer>();
m.put("lol", 7122);
     System.out.println(m);
     for(String key: m.keySet())
  System.out.println(key + " : " + m.get(key));
     System.out.println(m.containsKey("lol"));
     System.out.println(m.containsValue(7122));
     System.out.println(Math.PI);
     System.out.println(Math.acos(-1));
     BigInteger bi = in.nextBigInteger(), bj = new BigInteger("
    -7122"), bk = BigInteger.valueOf(17171);
     int sgn = bi.signum(); // sign(bi)
     bi = bi.subtract(BigInteger.ONE).multiply(bj).divide(bj).
      and(bj).gcd(bj).max(bj).pow(87);
     int meow = bi.compareTo(bj); // -1 0 1
String stz = "f5abd69150";
     BigInteger b16 = new BigInteger(stz, 16);
     System.out.println(b16.toString(2));
}
```

#### 9.5 Dancing Links

```
namespace dlx {
int lt[maxn], rg[maxn], up[maxn], dn[maxn], cl[maxn], rw[maxn],
    bt[maxn], s[maxn], head, sz, ans;

void init(int c) {
    for (int i = 0; i < c; ++i) {
        up[i] = dn[i] = bt[i] = i;
        lt[i] = i == 0 ? c : i - 1;
        rg[i] = i == c - 1 ? c : i + 1;
        s[i] = 0;
    }
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
}

void insert(int r, const vector<int> &col) {
    if (col.empty()) return;
    int f = sz;
    for (int i = 0; i < (int)col.size(); ++i) {
        int c = col[i], v = sz++;
    }
}</pre>
```

```
dn[bt[c]] = v;
    up[v] = bt[c], bt[c] = v;
    rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
    rw[v] = r, cl[v] = c;
    ++s[c];
    if (i > 0) lt[v] = v - 1;
  lt[f] = sz - 1;
void remove(int c) {
 lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
for (int i = dn[c]; i != c; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j])
       up[dn[j]] = up[j], dn[up[j]] = dn[j], --s[cl[j]];
  }
}
void restore(int c) {
  for (int i = up[c]; i != c; i = up[i]) {
  for (int j = lt[i]; j != i; j = lt[j])
       ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
  lt[rg[c]] = c, rg[lt[c]] = c;
// Call dlx::make after inserting all rows.
void make(int c) {
  for (int i = 0; i < c; ++i)
  dn[bt[i]] = i, up[i] = bt[i];</pre>
void dfs(int dep) {
  if (dep >= ans) return;
  if (rg[head] == head) return ans = dep, void();
  if (dn[rg[head]] == rg[head]) return;
  int c = rg[head];
  int w = c;
  for (int x = c; x != head; x = rg[x]) if (s[x] < s[w]) w = x;
  remove(w);
  for (int i = dn[w]; i != w; i = dn[i]) {
    for (int j = rg[i]; j != i; j = rg[j]) remove(cl[j]);
    dfs(dep + 1);
    for (int j = lt[i]; j != i; j = lt[j]) restore(cl[j]);
  }
  restore(w);
int solve() {
 ans = 1e9, dfs(0);
 return ans;
```

### 9.6 Offline Dynamic MST

```
int cnt[maxn], cost[maxn], st[maxn], ed[maxn];
pair<int, int> qr[maxn];
// qr[i].first = id of edge to be changed, qr[i].second =
     weight after operation
// cnt[i] = number of operation on edge i
// call solve(0, q - 1, v, 0), where v contains edges i such
     that cnt[i] == 0
void contract(int 1, int r, vector<int> v, vector<int> &x,
    vector<int> &y) {
  sort(v.begin(), v.end(), [&](int i, int j) {
    if (cost[i] == cost[j]) return i < j;</pre>
    return cost[i] < cost[j];</pre>
  });
  djs.save();
  for (int i = 1; i <= r; ++i) djs.merge(st[qr[i].first], ed[qr</pre>
    [i].first]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      x.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
 }
  djs.undo();
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) djs.merge(st[x[i]],
    ed[x[i]]);
  for (int i = 0; i < (int)v.size(); ++i) {</pre>
    if (djs.find(st[v[i]]) != djs.find(ed[v[i]])) {
      y.push_back(v[i]);
      djs.merge(st[v[i]], ed[v[i]]);
   }
  dis.undo();
```

```
void solve(int 1, int r, vector<int> v, long long c) {
  if (l == r) {
    cost[qr[i].first] = qr[l].second;
    if (st[qr[l].first] == ed[qr[l].first]) {
      printf("%lld\n", c);
      return:
    int minv = qr[l].second;
    for (int i = 0; i < (int)v.size(); ++i) minv = min(minv,
     cost[v[i]]);
    printf("%lld\n", c + minv);
  int m = (l + r) >> 1;
  vector<int> lv = v, rv = v;
  vector<int> x, y;
  for (int i = m + 1; i \le r; ++i) {
    cnt[qr[i].first]--
    if (cnt[qr[i].first] == 0) lv.push_back(qr[i].first);
  contract(l, m, lv, x, y);
long long lc = c, rc = c;
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    lc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(l, m, y, lc);
  djs.undo();
  x.clear(), y.clear();
  for (int i = m + 1; i <= r; ++i) cnt[qr[i].first]++;</pre>
  for (int i = l; i <= m; ++i) {</pre>
    cnt[qr[i].first]--;
    if (cnt[qr[i].first] == 0) rv.push_back(qr[i].first);
  contract(m + 1, r, rv, x, y);
  djs.save();
  for (int i = 0; i < (int)x.size(); ++i) {</pre>
    rc += cost[x[i]];
    djs.merge(st[x[i]], ed[x[i]]);
  solve(m + 1, r, y, rc);
  djs.undo();
  for (int i = l; i <= m; ++i) cnt[qr[i].first]++;</pre>
```

#### 9.7 Manhattan Distance MST

```
void solve(int n) {
  init();
  vector<int> v(n), ds;
  for (int i = 0; i < n; ++i) {
    v[i] = i;
    ds.push_back(x[i] - y[i]);
  sort(ds.begin(), ds.end());
  ds.resize(unique(ds.begin(), ds.end()) - ds.begin());
sort(v.begin(), v.end(), [&](int i, int j) { return x[i] == x
     [j] ? y[i] > y[j] : x[i] > x[j]; y[j] : x[i] > x[j]; y[i]
  int j = 0;
  for (int i = 0; i < n; ++i) {
    int p = lower_bound(ds.begin(), ds.end(), x[v[i]] - y[v[i]]
     ]]) - ds.begin() + 1;
    pair<int, int> q = query(p);
     // query return prefix minimum
    if (~q.second) add_edge(v[i], q.second);
    add(p, \ make\_pair(x[v[i]] + y[v[i]], \ v[i])); \\
  }
}
void make_graph() {
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);
  solve(n);
  for (int i = 0; i < n; ++i) x[i] = -x[i];
  solve(n);
  for (int i = 0; i < n; ++i) swap(x[i], y[i]);
  solve(n);
```

#### 9.8 IOI 2016 Aliens Trick

```
long long Alien() {
  long long c = kInf;
  for (int d = 60; d >= 0; --d) {
    // cost can be negative as well, depending on the problem.
    if (c - (1LL << d) < 0) continue;</pre>
```

```
long long ck = c - (1LL << d);
pair<long long, int> r = check(ck);
if (r.second == k) return r.first - ck * k;
if (r.second < k) c = ck;
}
pair<long long, int> r = check(c);
return r.first - c * k;
}
```

#### 9.9 Matroid Intersection

Start from  $S = \emptyset$ . In each iteration, let

```
 \begin{array}{ll} \bullet & Y_1 = \{x \not \in S \mid S \cup \{x\} \in I_1\} \\ \bullet & Y_2 = \{x \not \in S \mid S \cup \{x\} \in I_2\} \end{array}
```

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \not\in S$ , create edges

```
• x \to y \text{ if } S - \{x\} \cup \{y\} \in I_1.
• y \to x \text{ if } S - \{x\} \cup \{y\} \in I_2.
```

Find a shortest path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of S will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \notin S$ . Find the path with the minimum number of edges among all minimum length paths and alternate it.