

CSC196: Great Ideas in Computing

Tutorial 5

16th November 2022

Assignment 3

- Deadline: Wednesday, November 18, 8 AM EST
- Please try to export your solutions to .pdf before submitting, it makes grading easier
- Assignment created on MarkUs with 48 hours penalty decay of 0.5% per hour
- You have one week from the deadline to ask for re-grading
 - You will get the assignment grades within 4 days from the deadline

Q1

- Write search engine query to get some "complex" information
 - "Complex" information is a mixture of multiple "simple" information
 - "Simple" information: List of countries in the world
 - "Complex" information: List of executive leaders in the European Union
 - You will probably not get the desired result in the first time
 - Refine query **if necessary** and try again
 - Use incognito mode so that search information is not cached

Q1 (cont.)

- Record the following for each query:
 - Search Engine of choice (Google, DuckDuckGo etc...)
 - Query string ("Who is the president of South Africa")
 - Is this a refined query? If so, how?
 - Time at which search was executed
 - $\max(k, 10)$ titles, where k is the number of results, ignore sponsored results and subtitles
- Repeat the same set of queries after a day
 - Use incognito mode so that search information is not cached
 - Search engines might learn to produce different (e.g. more relevant) results

Q1 (cont.)

- Are the titles recorded the same for each query? If it is different, then why is it the case?
 - Is it more relevant or less?
 - Are the sources more authentic than before?
 - These are some of the questions you can ask yourselves
 - Again, there are no right or wrong answers here
- **To the point answers are appreciated, don't write a long essay to get more marks.**

Q2

- Find allocations for each agent based on their value density function
 - Allocation A_1 for agent 1 is a range of rational numbers
 - $A_1 = [a, b] \subset [0, 1]$ such that $a, b \in Q$
 - Value function v_1 for agent 1 maps allocation A_1 to value y_1
 - $v_1(A_1) = y_1$
- Show that allocations are **equitable** and **envy-free**
 - Equitable: $y_1 = y_2 = y_3$
 - Envy-Free: $v_i(A_i) > v_i(A_j) \forall i, j \in [1, 2, 3]$
- Hint: **Use symmetry of the given value density functions**

Q3

- Assume $P \neq NP$
- It is a proven fact that 3-coloring problem is NP-complete
 - NP-complete: cannot be decided in polynomial time
- For graph $G = (V, E)$, whether the following languages are decidable in polynomial time? Justify.
 - \mathcal{L}_1 : Is graph G , 100 colorable?
 - \mathcal{L}_2 : Is graph G , $|V| - 1$ colorable?
 - \mathcal{L}_3 : Is graph G , $|V|/2$ colorable such that $|V|$ is even?

Q3 (cont.)

- If a NP-complete problem A can be reduced in polynomial time to problem B , then B is NP-complete
- Come up with a transformation to reduce 3-coloring problem to 100-coloring problem?
 - Use a graph G which is 3 colorable and has $|V'|$ vertices, then transform it
 - Hint: Clique K_n has n vertices and it is n -colorable
- What should be the size of the clique to reduce 3-coloring problem to $|V|/2$ colorable problem?
 - Here $|V|$ is the number of vertices in the transformed graph
 - Use a graph G which is 3 colorable and has $|V'|$ vertices, then transform it

Q3 (cont.)

- What about $|V| - 1$ colorable problem?
 - Every graph is $|V| - 1$ colorable if it is not a clique, **why?**
 - You can't use above transformations for this problem, **why?**

Best of luck for the quiz on 25th November!

- If you want to prove a problem A is NP, either show
 - Some NP problem B is reducible to A , note the direction of reducibility
 - Show exponential time solution and show that there is no better solution
- Polynomial vs Exponential
 - n^k for constant k is polynomial
 - n^k where $k = f(n)$ is exponential
- Revise previous assignments