

CSC196: Great Ideas in Computing

Tutorial 3

5th October 2022

Birthday paradox

- Write your DOBs on a piece of paper

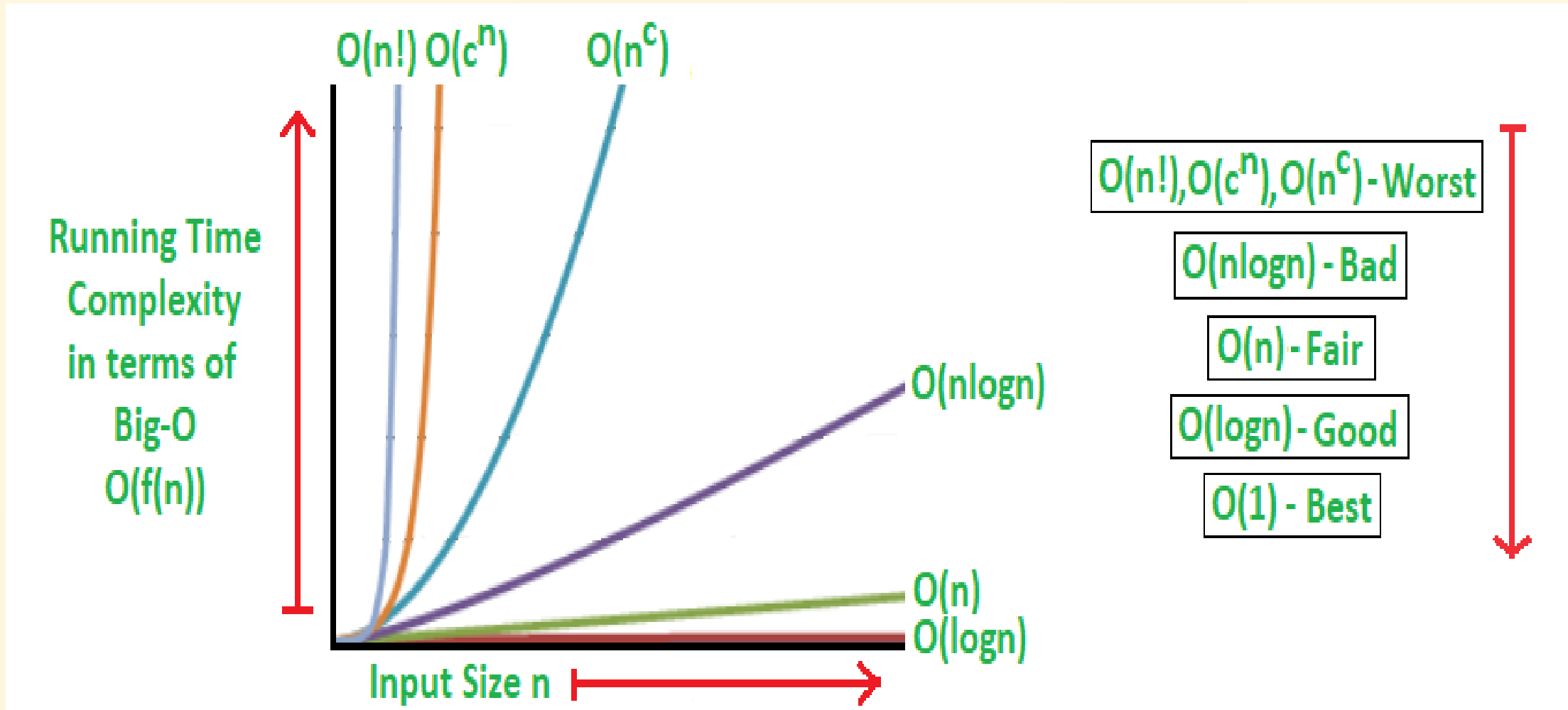
Q2. Asymptotic analysis of algorithms

- Big O (Upper bound)
 - $f(n)$ is $O(g(n))$
 - a.k.a $f(n)$ is Big oh of $g(n)$
 - a.k.a $f(n)$ is order of $g(n)$
 - $f(n) \leq c * g(n)$ for all $n > n_0$ for an arbitrary n_0 and c

- Examples

- $O(1)$: Constant time
- $O(\log N)$: Logarithmic time
- $O(N)$: Linear time
- $O(N^2)$: Quadratic time
- $O(N^2 + \log N) = O(N^2) + O(\log N)$

- Other bounds
 - Big Omega (lower bound)
 - Big Theta (tight bound)



- Reference: [https://www.geeksforgeeks.org/analysis-algorithms-](https://www.geeksforgeeks.org/analysis-algorithms-big-o-analysis/)

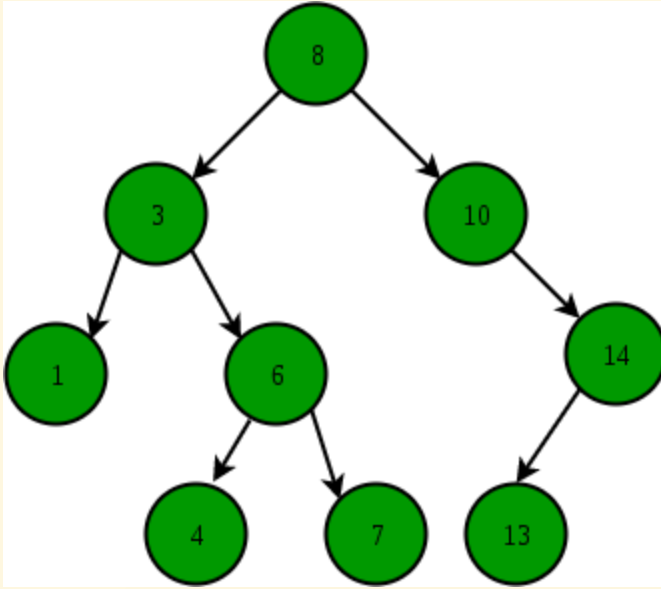
n	Constant $O(1)$	Logarithmic $O(\log n)$	Linear $O(n)$	Linear Logarithmic $O(n \log n)$	Quadratic $O(n^2)$	Cubic $O(n^3)$
1	1	1	1	1	1	1
2	1	1	2	2	4	8
4	1	2	4	8	16	64
8	1	3	8	24	64	512
16	1	4	16	64	256	4,096
1,024	1	10	1,024	10,240	1,048,576	1,073,741,824

- Reference: <https://dzone.com/articles/learning-big-o-notation-with-on-complexity>

Q2. Binary Search Tree

- Binary Tree
 - There is a root node which has in degree = 0 (no incoming edges)
 - There are leaf nodes which has out degree = 0 (no outgoing edges)
 - Each node has at most two outgoing pointers and at most one incoming pointer
 - Number of levels = $O(\log(\text{number_of_nodes}))$

- Binary Search Tree (BST)
 - All nodes in the left subtree have lesser value than root node
 - All nodes in the right subtree have greater value than root node
 - Left and right subtrees are BST as well
- Used to implement ordered dictionary (keys are sorted)



- Reference: <https://www.geeksforgeeks.org/binary-search-tree-data-structure/>

Q2. Sorted Array

- Array can be implemented using linked list
 - Doubly linked list for traversal from either end
- Searching (reading) is faster in sorted array, however writing is slower
 - Reading: $O(\log(n))$
 - Writing: $O(\log(n))$
- Unsorted array has faster writes but slower reads
 - Reading: $O(n)$
 - Writing: $O(1)$

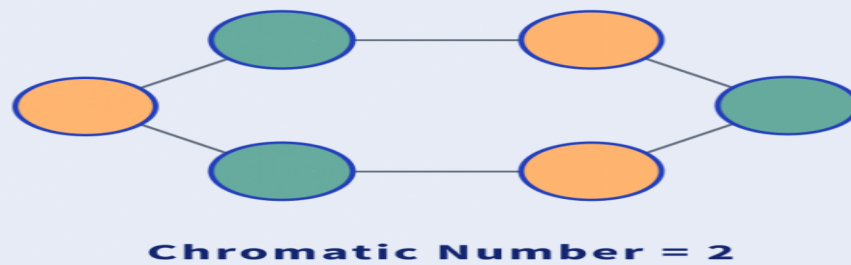
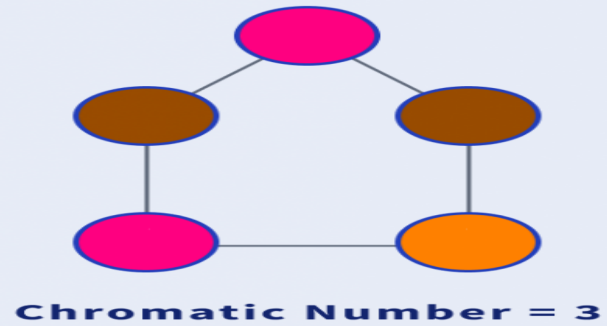
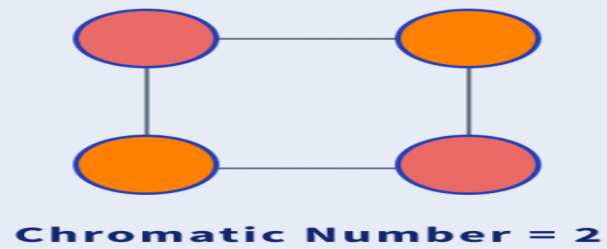
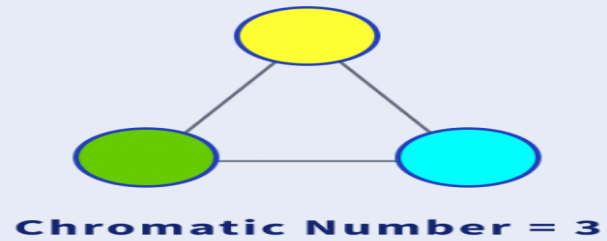
Memory Location									
200	201	202	203	204	205	206	▪	▪	▪
U	B	F	D	A	E	C	▪	▪	▪
0	1	2	3	4	5	6	▪	▪	▪
Index									

- Reference: <https://www.geeksforgeeks.org/array-data-structure/>

Q3. Graph coloring problem

- Vertex coloring is the most common graph coloring problem
- Color the vertices such that no two adjacent vertices are colored the same
 - Graph is k colorable if you can do this with k colors
 - Chromatic number: Least possible value of k
 - Use Breadth first search (BFS) for two coloring
- Reference: <https://www.geeksforgeeks.org/graph-coloring-applications/>

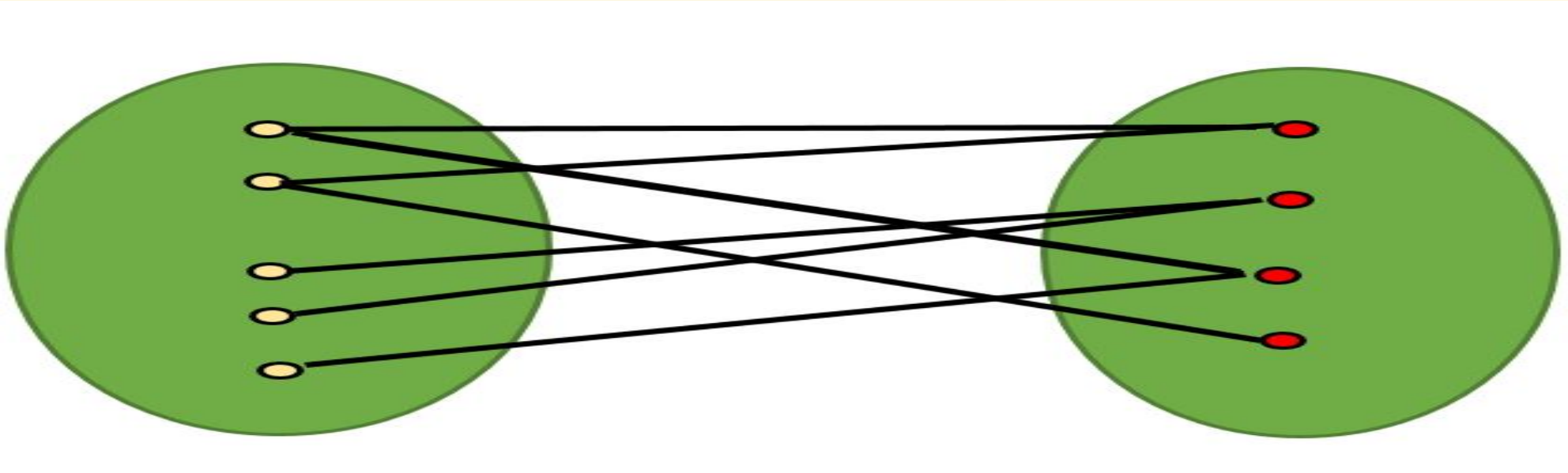
**Chromatic
Number of
Cycle Graph
Examp^ls**



- Finding chromatic number ($N > 3$) is NP complete
- Reference: <https://www.interviewbit.com/blog/graph-coloring-problem/>

Q3. Bipartite graph

- Vertices can be divided into two disjoint sets
 - All edges are in between the sets



Assignment 1

- Deadline 7th October (Friday 8AM)

Birthday Paradox

- Pigeon hole principle: Same birthday guaranteed for $n > 365$ people
- Probability of same birthday is not $n/365$ but
 - $P = 1 - \frac{365*364*...*(365-(n-1))}{365^n} = 1 - \frac{P_n^{365}}{365^n}$
 - For $n = 20$, $P = 41.1\%$
 - For $n = 30$, $P = 70.6\%$
- Reference:
https://en.wikipedia.org/wiki/Birthday_problem#Calculating_the_probability

Example of computable transformation

- Equivalence of Vertex cover problem and Independent set problem
 - The solution to one problem can be computed using solution to the other problem
- Vertex Cover problem
 - Identify a set of vertices such that all edges are connected to these vertices
- Independent Set problem
 - Identify a set of vertices such that there are no edges between them
- Removing the vertices of a given vertex cover from the graph leaves us with independent set
 - And vice versa