

### 1.4.1 Proof by simple induction

- a) Use simple induction to show that for every positive integer  $n$ ,  $5^n - 1$  is divisible by 4  
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 Base case:

$$\begin{aligned} 5^1 - 1 &= 4j \\ &= 4 = 4j \\ 5^k - 1 &= 4j \\ 5^k &= 4j + 1 \end{aligned}$$

Induction step:

$$\begin{aligned} &= 5^{k+1} - 1 \\ &= 5^k \cdot 5 - 1 \\ &= (4j + 1) \cdot 5 - 1 \\ &= (20j + 5) - 1 \\ &= 20j + 4 \\ &= 4 \cdot (5j + 1) \\ &\implies \text{true for } n = k + 1 \end{aligned}$$

- b) Use simple induction to show that for every positive integer  $n$ ,  $n^3 - n$  is divisible by 3. (Hint: In the induction step, you will need to make use of the arithmetic fact that  $(k+1)^3 = k^3 + 3k^2 + 3k + 1$ )  
 $n^3 - n$  is divisible by 3  
 Base case:

$$\begin{aligned} 1^3 - 1 &= 3j \\ &= 0 = 3j \\ k^3 - k &= 3j \\ k^3 &= 3j + k \end{aligned}$$

Induction step:

$$\begin{aligned} &= (k+1)^3 - (k+1) \\ &= (k^3 + 3k^2 + 3k + 1) - (k+1) \\ &= k^3 + 3k^2 + 2k \\ &= (3j + k) + 3k^2 + 2k \\ &= 3j + 3k^2 + 3k \\ &= 3 \cdot (j + k^2 + k) \\ &\implies \text{true for } n = k + 1 \end{aligned}$$

- c) Show by simple induction that for every natural number  $n$ ,  $\sum_{i=0}^n 2^i = 2^{n+1} - 1$   
 $f(n) = 2^{n+1} - 1$   
 Base case:

$$\begin{aligned} f(1) &= 2^{1+1} - 1 \\ f(1) &= 4 - 1 \\ f(1) &= 3 \\ f(k) &= 2^{k+1} - 1 \end{aligned}$$

Induction step:

$$\begin{aligned}
 f(k+1) &= 2^{k+2} - 1 \\
 f(k+1) &= 2^1 + 2^2 \dots + 2^k + 2^{k+1} \\
 &= f(k) + 2^{k+1} \\
 &= 2^{k+1} - 1 + 2^{k+1} \\
 &= 2(2^{k+1}) - 1 \\
 &= 2^1 \cdot 2^{k+1} - 1 \\
 &= 2^{k+1+1} - 1 \\
 &= 2^{k+2} - 1 \\
 &\implies \text{true for } n = k + 1
 \end{aligned}$$

### 1.4.2: Definition by simple recursion

a) Let  $f : N \rightarrow N$  be the function defined by putting  $f(0) = 0$  and  $f(n+1) = n$  for all  $n \in N$ .

i) Evaluate this function bottom-up for all arguments 0-5.

$$\begin{aligned}
 f(0) &= 0 \\
 f(1) &= 0 \\
 f(2) &= 1 \\
 f(3) &= 2 \\
 f(4) &= 3 \\
 f(5) &= 4
 \end{aligned}$$

ii) Explain what  $f$  does by expressing it in explicit terms (i.e. without a recursion).

This function is not recursive, it only returns the antecesor of the current value

b) Let  $f : N^+ \rightarrow N$  be the function that takes each positive integer  $n$  to the greatest natural number  $p$  with  $2^p \leq n$ . Define this function by a simple recursion. (Hint: You will need to divide the recursion step into two cases.)

$$\begin{aligned}
 f(n) &= 0 && , \text{when } n = 1 \\
 f(n) &= f(n-1) && , \text{when } n > 1, \log_2(n) \notin N \\
 f(n) &= f(n-1) + 1 && , \text{when } n > 1, \log_2(n) \in N
 \end{aligned}$$

c) Let  $g : NXN \rightarrow N$  be defined by putting  $g(m, 0) = m$  for all  $m \in N$  and  $g(m, n+1) = f(g(m, n))$  where  $f$  is the function defined in part (a) of this exercise.

i) Evaluate  $g(3, 4)$  top-down.

$$\begin{aligned}
 g(3, 4) &= f(g(3, 3)) \\
 &= f(f(g(3, 2))) \\
 &= f(f(f(g(3, 1)))) \\
 &= f(f(f(f(g(3, 0))))) \\
 &= f(f(f(f(f(3))))) \\
 &= f(f(f(f(2)))) \\
 &= f(f(f(1))) \\
 &= f(f(0)) \\
 &= f(0) \\
 &= 0
 \end{aligned}$$

- ii) Explain what  $g$  does by expressing it in explicit terms (i.e. without a recursion).  
It subtracts the right item from the left, but if the remain is negative, it returns 0

### 1.4.3: Proof by cumulative induction

- a) Use cumulative induction to show that any postage cost of four or more pence can be covered by two-pence and five-pence stamps.

$$2x + 5y = n, \text{ when } n \geq 4$$

Base case:

$$n = 4$$

$$4 = 2(2) + 5(0)$$

Induction step:

$$\text{hypothesis} \rightarrow \forall j < k, j = 2x + 5y$$

$$\text{goal} \rightarrow k = 2x + 3y$$

$$\text{case(1)} \rightarrow k \text{ is multiple of two} \implies y = 0 \text{ and } x \in N$$

$$\text{case(2)} \rightarrow k \text{ is multiple of five} \implies x = 0 \text{ and } y \in N$$

$$\text{case(3)} \rightarrow k \text{ is not multiple of any} \implies y \in N \text{ and } x \in N$$

- b) Use cumulative induction to show that for every natural number  $n$ ,  $F(n) \leq 2^n - 1$ , where  $F$  is the Fibonacci function.

$$F(n) \leq 2^n - 1,$$

Base case:

$$n \leq 1$$

$$1 \leq 2^1 - 1$$

Induction step:

$$\text{hypothesis} \rightarrow \forall j < k, F(j) \leq 2^j - 1$$

$$\text{goal} \rightarrow F(k) \leq 2^k - 1$$

$$\text{case(1)} \rightarrow F(k) \text{ is greater than } 2^k - 1 \implies 2^{k+1} - 1 < F(j+1)$$

*Impossible given that  $F$  grows at  $F(k)$  and  $2^k - 1$  grows at double rate*

$$\text{case(2)} \rightarrow F(k) \text{ is less or equal than } 2^k - 1 \implies \text{This must be true then}$$

- c) Calculate  $F(5)$  top-down, and then again bottom-up, where again  $F$  is the Fibonacci function

Top-down:

$$F(5) = F(4) + F(3)$$

$$F(5) = (F(3) + F(2)) + (F(2) + F(1))$$

$$F(5) = ((F(2) + F(1)) + (F(1) + F(0))) + ((F(1) + F(0)) + 1)$$

$$F(5) = (((F(1) + F(0)) + 1) + (1 + 0)) + ((1 + 0) + 1)$$

$$F(5) = (((1 + 0) + 1) + (1 + 0)) + ((1 + 0) + 1)$$

$$F(5) = ((1 + 1) + (1)) + (1 + 1)$$

$$F(5) = (2 + 1) + (2)$$

$$F(5) = 3 + 2$$

$$F(5) = 5$$

Bottom-up

$$\begin{aligned}
 F(0) &= 0 \\
 F(1) &= 1 \\
 F(2) &= F(1) + F(0) = 1 + 0 = 1 \\
 F(3) &= F(2) + F(1) = 1 + 1 = 2 \\
 F(4) &= F(3) + F(2) = 2 + 1 = 3 \\
 F(5) &= F(4) + F(3) = 3 + 2 = 5
 \end{aligned}$$

- d) Express each of the numbers 14, 15 and 16 as a sum of  $3s$  and/or  $8s$ . Using this fact in your basis, show by cumulative induction that every positive integer  $n \leq 14$  may be expressed as a sum of  $3s$  and/or  $8s$ .

$$n = 3x + 8y, \text{ when } n > 14$$

Base case:

$$\begin{array}{ll}
 n = 14 & 14 = 3(2) + 8(1) \\
 n = 15 & 15 = 3(5) + 8(0) \\
 n = 16 & 16 = 3(0) + 8(2)
 \end{array}$$

Induction step:

$$\begin{aligned}
 \text{hypothesis} &\rightarrow \forall j < k, j = 3x + 8y \\
 \text{goal} &\rightarrow k = 3x + 8y \\
 \text{case(1)} &\rightarrow k \text{ is multiple of three} \implies y = 0 \text{ and } x \in N \\
 \text{case(2)} &\rightarrow k \text{ is multiple of eight} \implies x = 0 \text{ and } y \in N \\
 \text{case(3)} &\rightarrow k \text{ is not multiple of any} \implies y \in N \text{ and } x \in N
 \end{aligned}$$

- e) Show by induction that for every natural number  $n$ ,  $A(1, n) = n + 2$ , where  $A$  is the Ackermann function.

$$A(m, n) = \begin{cases} n + 1 & \text{if } m = 0 \\ A(m - 1, 1) & \text{if } m > 0 \text{ and } n = 0 \\ A(m - 1, A(m, n - 1)) & \text{if } m > 0 \text{ and } n > 0 \end{cases}$$

Base case:

$$\begin{array}{ll}
 n = 0 & A(1, 0) = A(0, 1) \\
 & = 2 \\
 & = (0) + 2 \\
 n = 1 & A(1, 1) = A(0, A(1, 0)) \\
 & = A(0, 2) \\
 & = 3 \\
 & = (1) + 2
 \end{array}$$

Induction step:

$$\begin{aligned}
 \text{hypothesis} &\rightarrow \forall j < k, A(1, j) = j + 2 \\
 \text{goal} &\rightarrow A(1, k) = k + 2 \\
 A(1, k + 1) &= (k + 1) + 2 \\
 &= A(0, A(1, k)) \\
 &= A(1, k) + 1 \\
 &= (k + 2) + 1 \\
 A(1, k + 1) &= k + 3 \\
 \therefore A(1, k) &= k + 2
 \end{aligned}$$