TECNOLÓGICO DE MONTERREY

FUNDAMENTOS DE COMPUTACIÓN

Homework 6

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1 Problems

Solve the following problems:

- 1. For the selection algorithm, analyze and discuss the resulting complexity when the initial list is divided into groups of 19 elements (instead of 15). Derive the proper conclusions.
- 2. Given a set of n numbers, we want to find the i largest in sorted order using a comparison-based algorithm. Analyze and compare the following methods in terms of n and i:
 - (a) Sort the numbers, and list the *i* largest.
 - (b) Build a max-priority queue (like a heap) with the numbers and extract the minimum i items.
 - (c) Use the k-max (session 06) to find the i-th largest, partition around that number, and sort the i largest.
- 3. For n distinct elements $x_1, x_2, ..., x_n$ with positive weights $w_1, w_2, ..., w_n$ such that $\sum_{i=1}^n w_i = 1$, the weighted (lower) median is the element x_k satisfying $\sum_{x_i < x_k} w_i < \frac{1}{2}$ and $\sum_{x_i > x_k} w_i \leq \frac{1}{2}$

For example, if the elements are 0.1, 0.35, 0.05, 0.1, 0.15, 0.05, 0.2 and each element equals its weight then the median is 0.1, but the weighted median is 0.2.

- (a) Argue that the median of $x_1, x_2, ..., x_n$ is the weighted median of the x_i with weights $w_i = 1/n$ for i = 1, 2, ..., n
- (b) Show how to compute the weighted median of nelements in O(nlogn) worst-case using sorting.
- (c) Show how to compute the weighted median of nelements in O(n) worst-case.
- 4. Investigate on how the adversary argument concept can be used to determine the lower bound of merging two ordered lists.