



POP Performance analysis methodology

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HORIZON-EUROHPC-JU-2023-COE



EuroHPC
Joint Undertaking

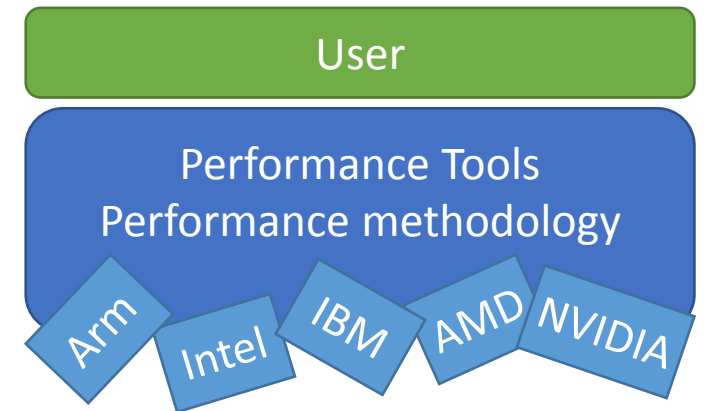
1 January 2024– 31 December 2026

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Context and history



- **POP-CoE** starts in 2015
- Objectives:
 - To offer performance analysis services to European HPC users.
 - To develop a common methodology for performance analysis
- **POP1** runs from October 2015 to March 2018
- **POP2** runs from December 2018 to June 2022
- **POP3** will run between January 2023 and December 2026



• Who?

- BSC (coordinator), ES
- HLRS, DE
- JSC, DE
- INES-ID, PT
- RWTH Aachen, IT Center, DE
- TERATEC, FR
- IT4I, CZ
- UVSQ, FR



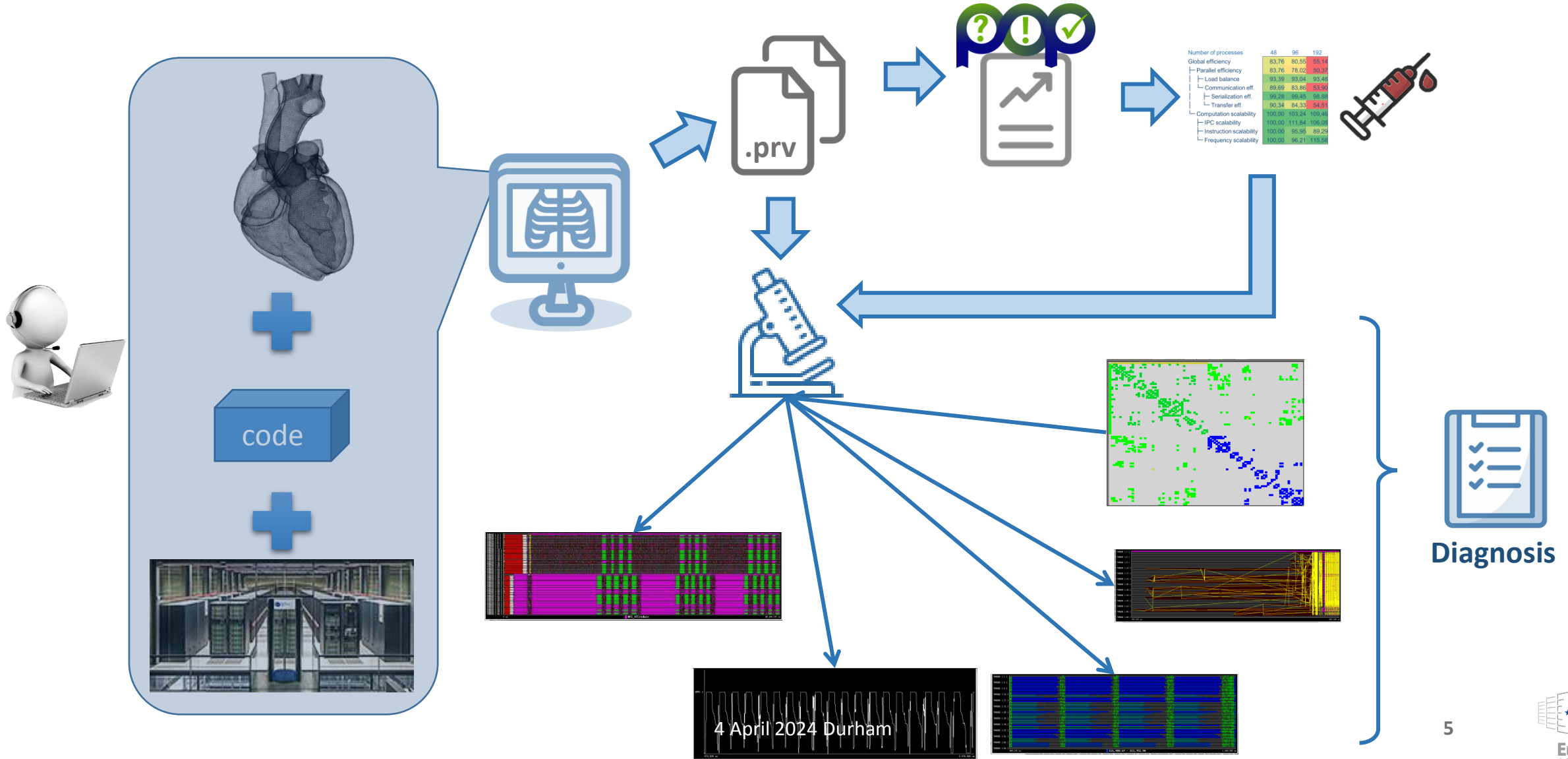
A team with

- Excellence in performance tools and tuning
- Excellence in programming models and practices
- Research and development background AND proven commitment in application to real academic and industrial use cases

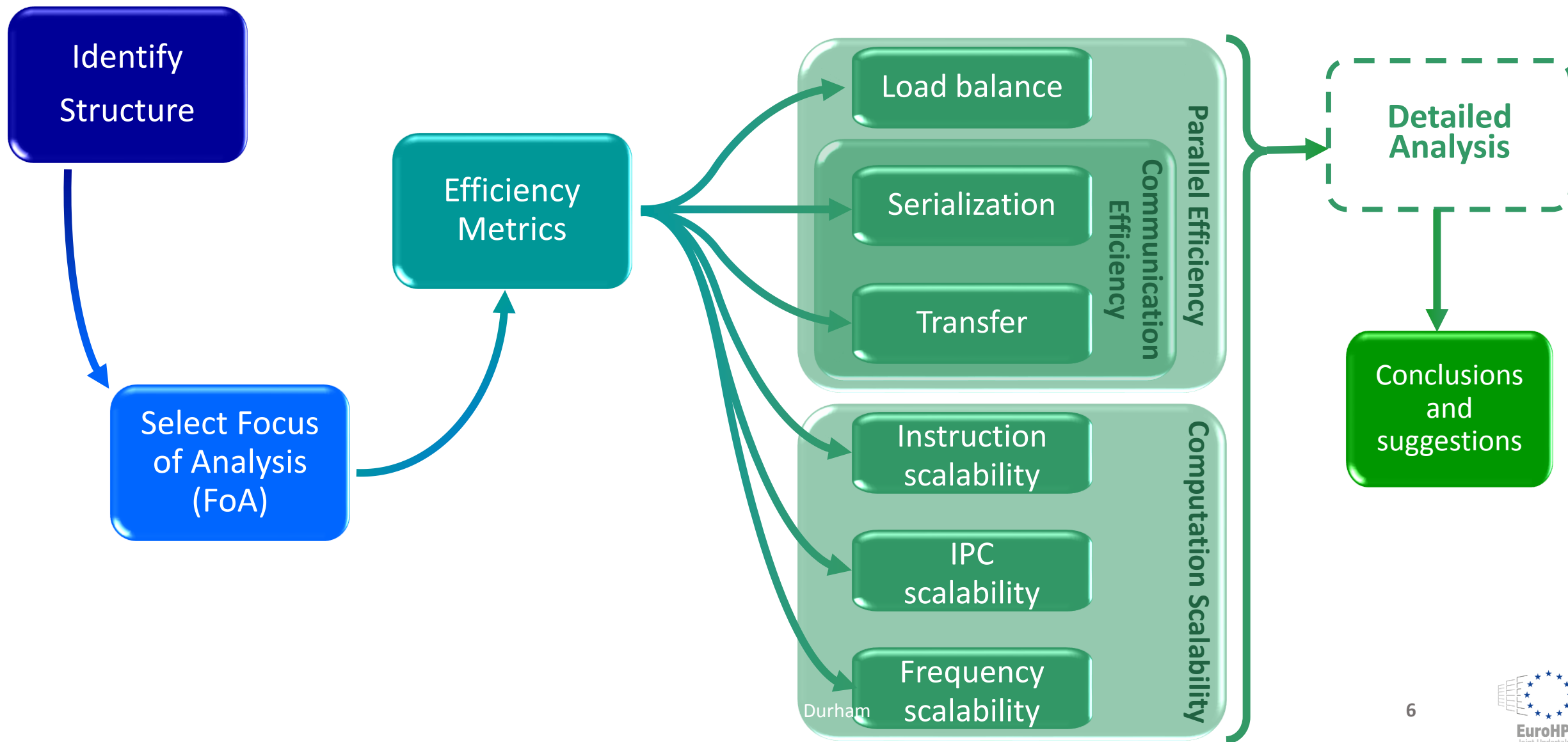


- POP methodology
 - A methodology for performance analysis
 - Agnostic of the tools
 - But this presentation is based on BSC tools (Extrae, Paraver, Dimemas and ModelFactors)
 - You will see examples of the methodology applied using JSC tools on 25th April 2024
- MPI only applications
 - But POP methodology is being extended to:
 - Hybrid models: OpenMP, GPUs...
 - I/O
 - Vectorization

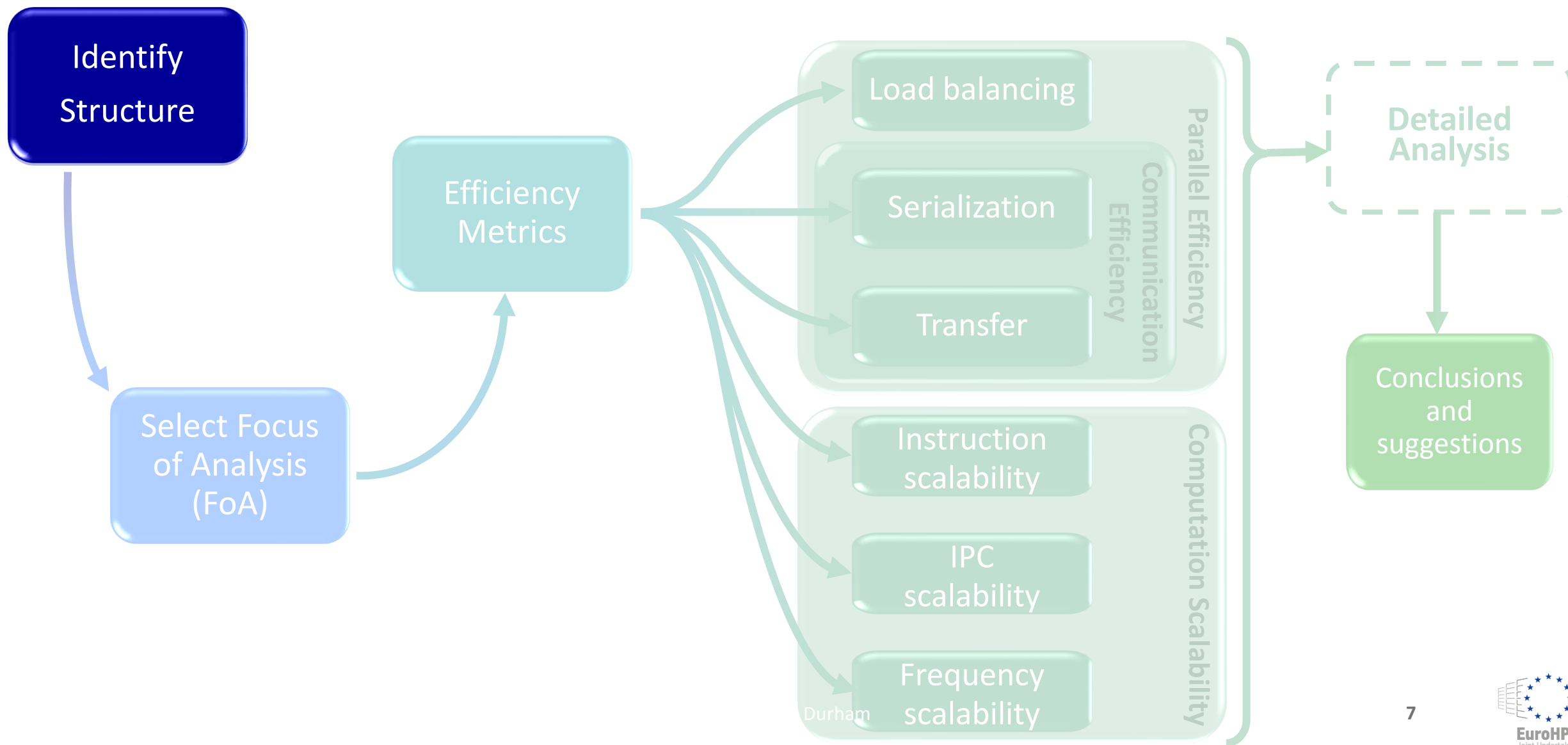
Performance analysis



The journey



The journey



Identify structure

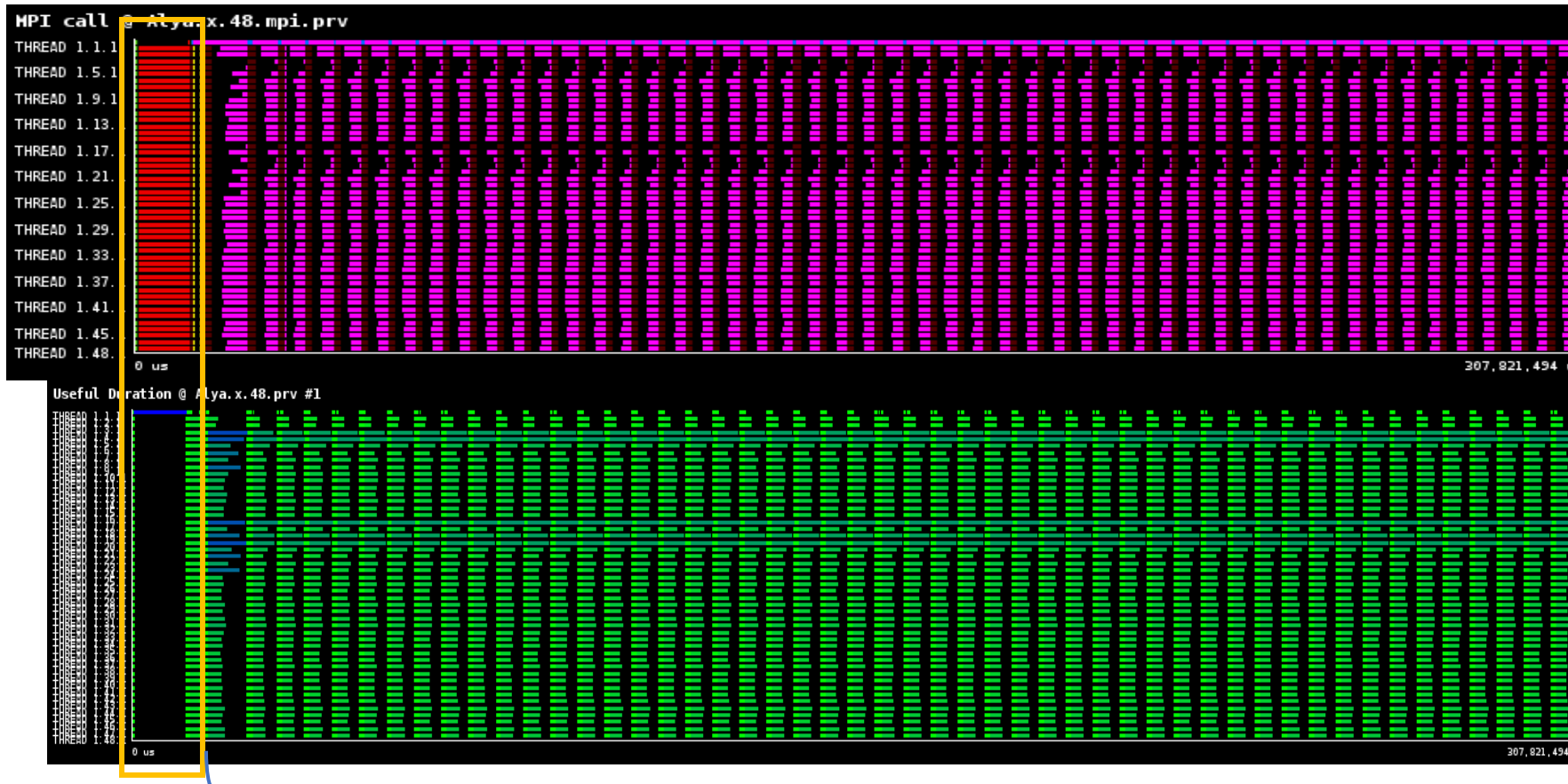


- Objectives:
 - Understand general structure
 - Identify initialization/finalization phases
 - Detect iterative pattern
 - Granularity
- Different levels of “difficulty” → Different levels of knowledge

Identify structure



- Usually use “MPI calls” view or “Useful Duration”



- Clear iterative pattern
- With an initialization phase
- All iterations are similar
 - We can select a few to analyze

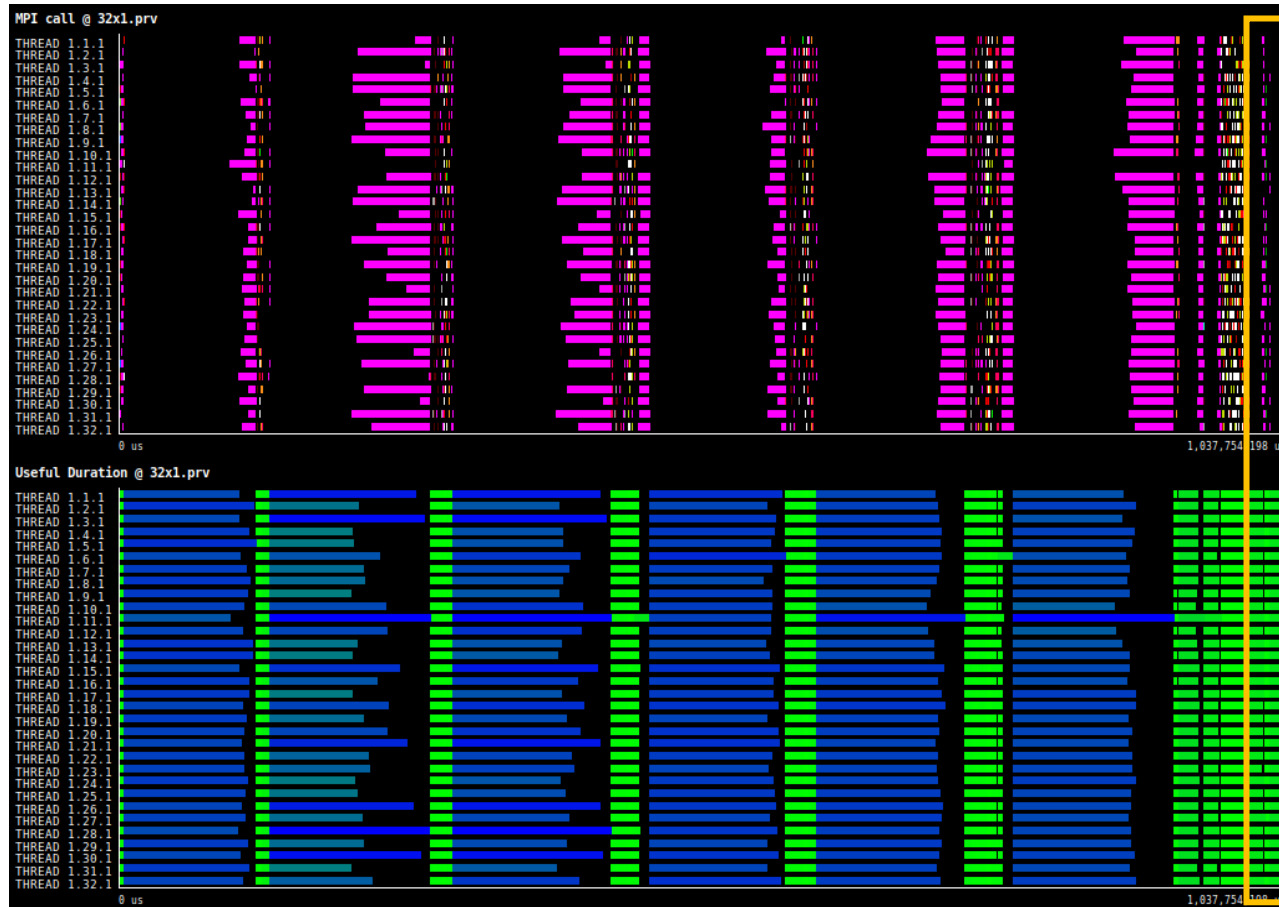
Initialization

Iterative phase

Identify structure



- Not always easy



- There are 6 iterations and one finalization phase
- Iterations are not regular along time
 - Different pattern of load balance
 - Different pattern of durations

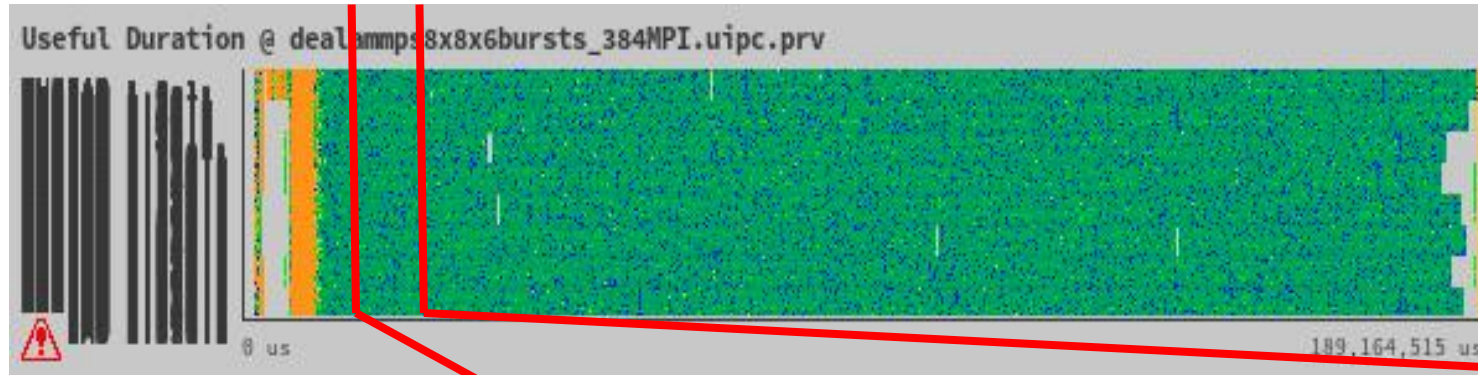
6 iterations

Finalization

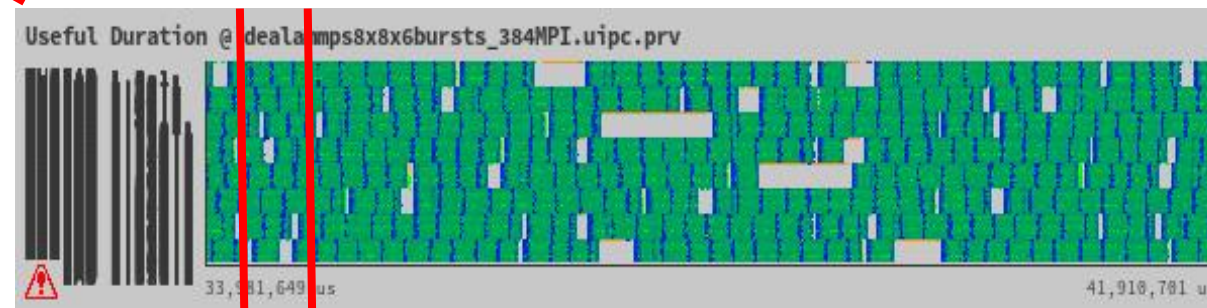
Identify structure



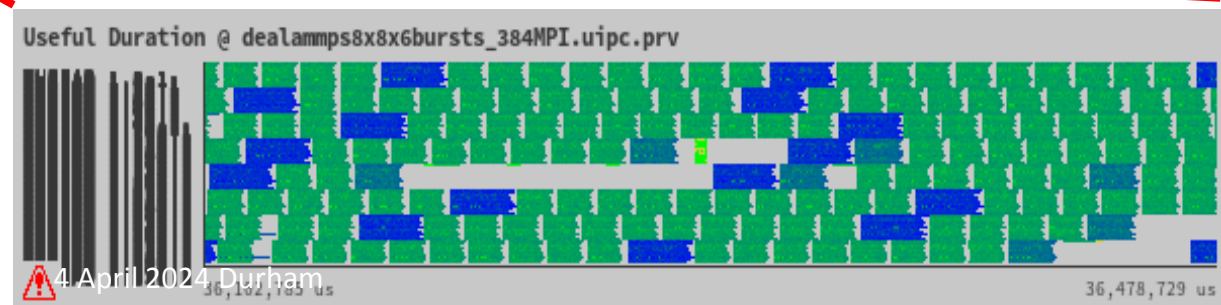
- Not always easy



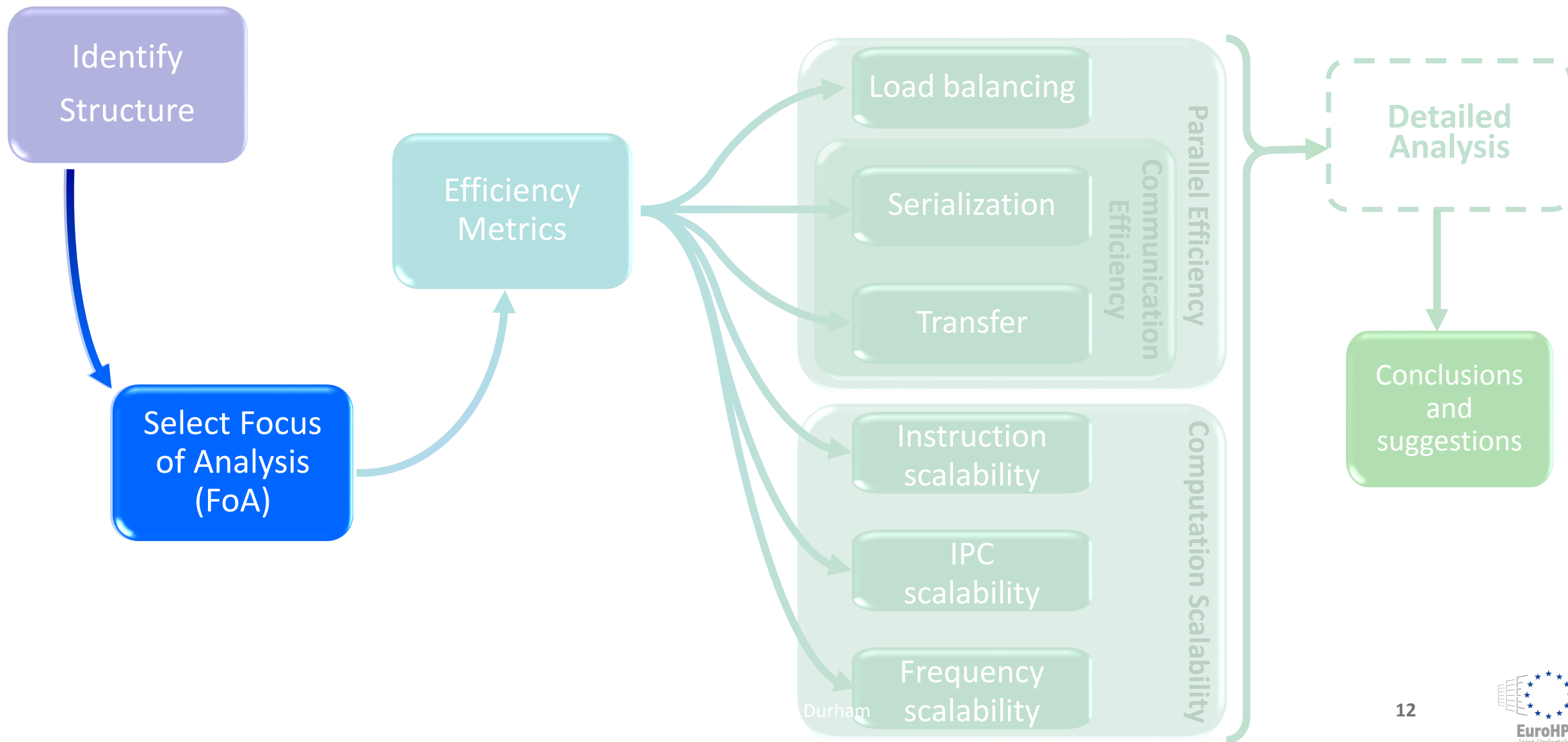
- It is not easy to detect an iterative structure
- No global synchronizations



- Zooming in we detect that synchronizations only happen at node level
- Not possible to determine iterations



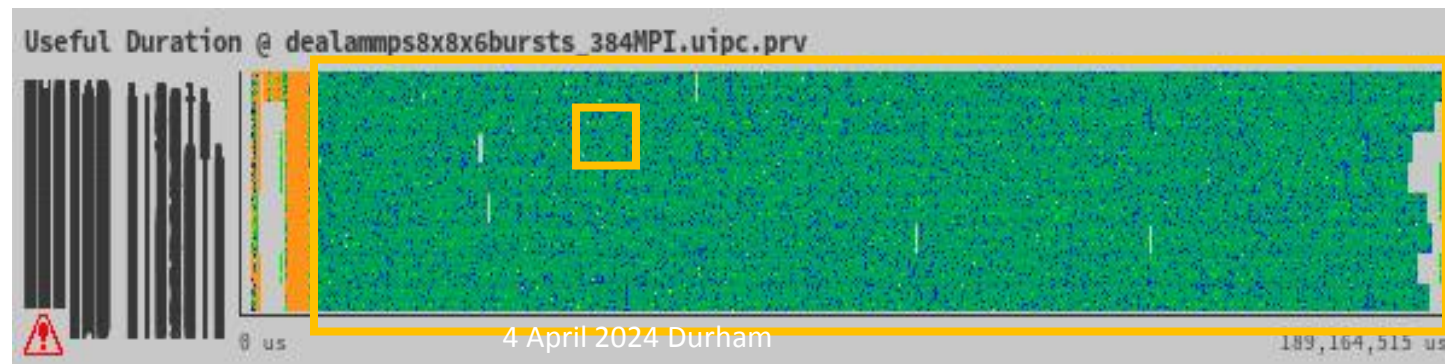
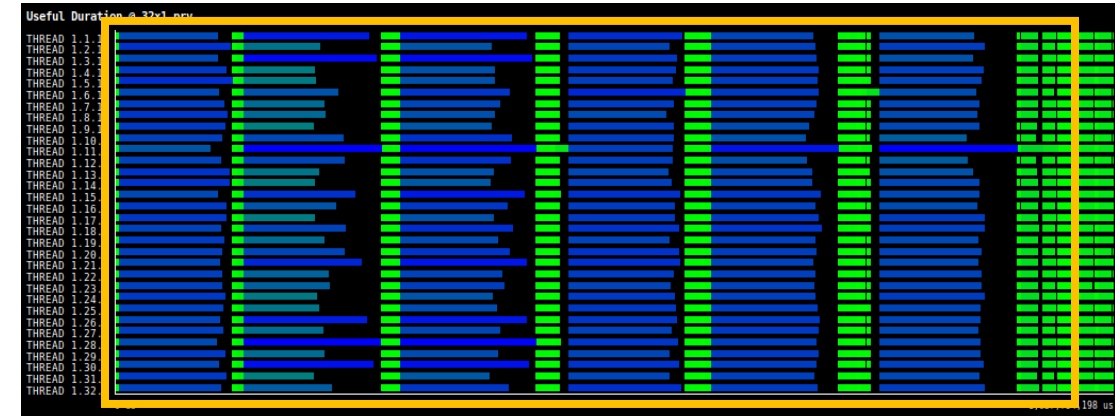
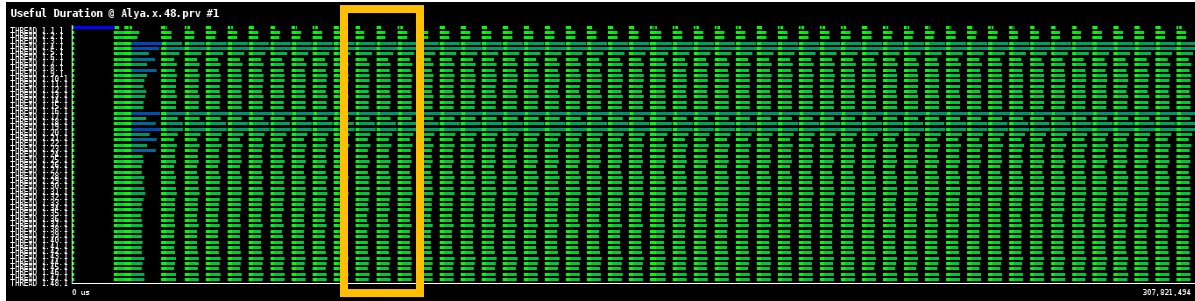
The journey



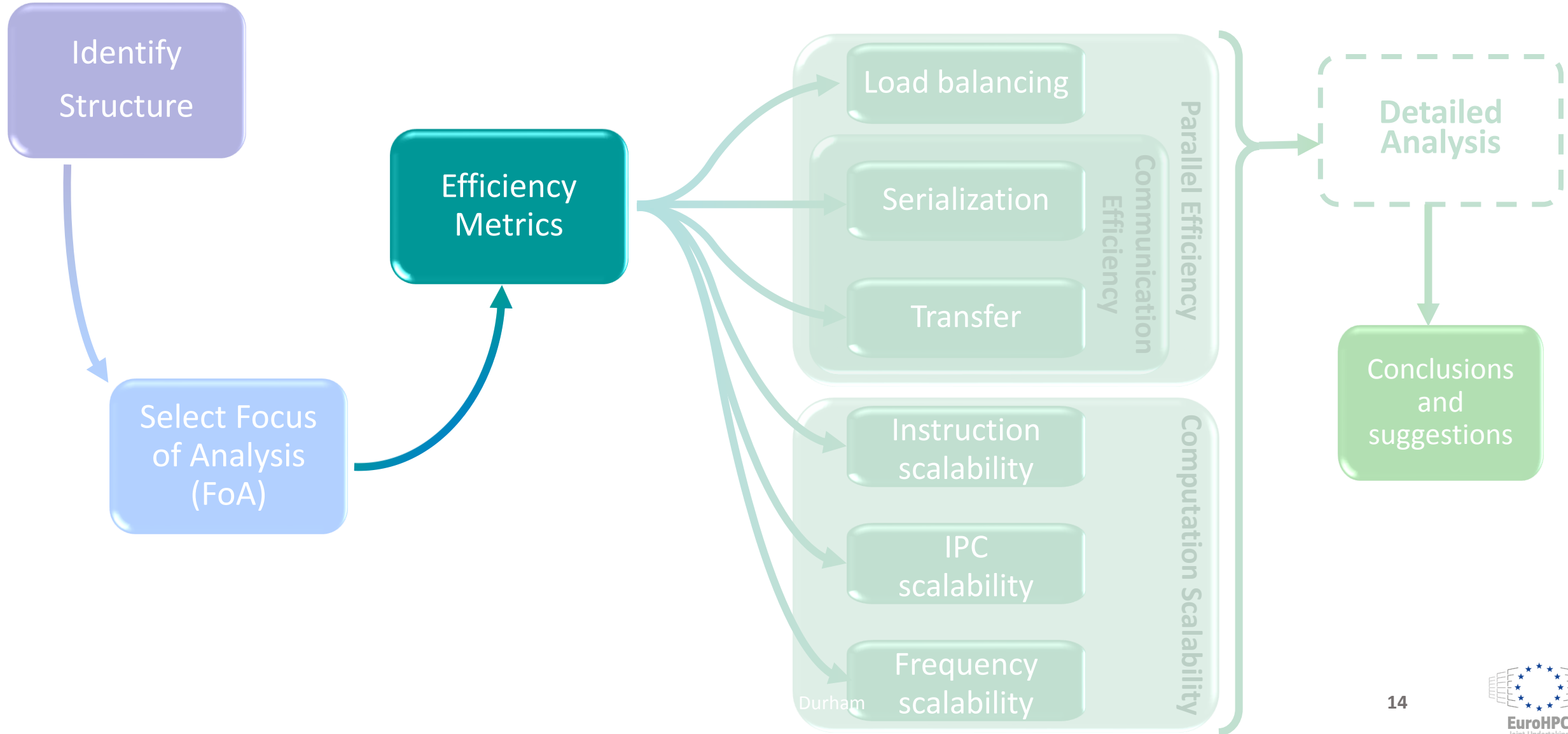
Select Focus of Analysis (FoA)



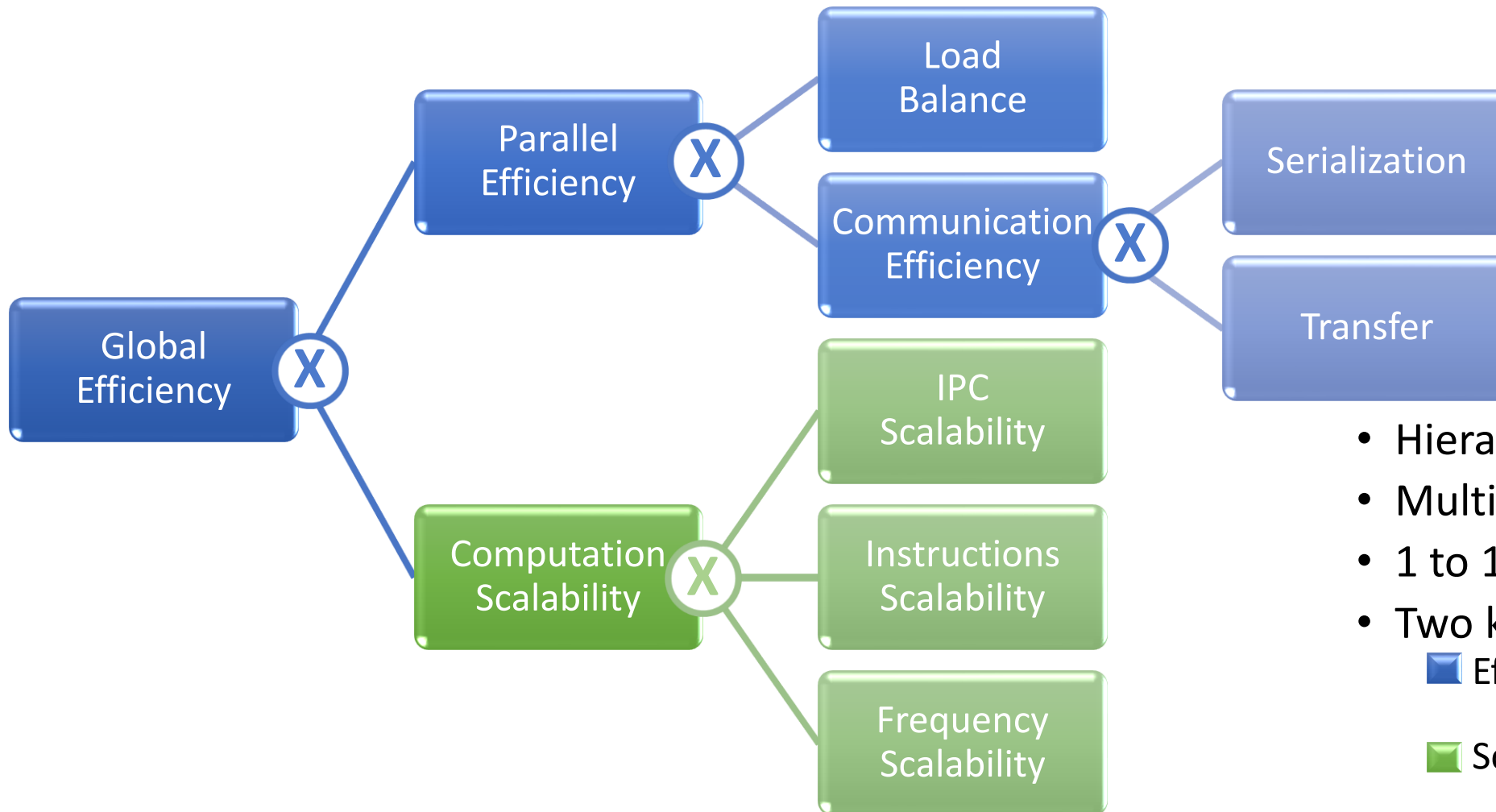
- **Objective:** Select the region we want to analyze
 - Not a correct answer, depends on the context of the analysis
 - For the same trace we may select two different FoA to perform two different studies with two different objectives.



The journey



The Efficiency Metrics



- Hierarchical model
- Multiplicative metrics
- 1 to 100% scale
- Two kind of metrics:
 - Efficiency metrics
 - Absolute metrics
 - Scalability metrics
 - Relative to a base case
 - 100% for the base case

The Efficiency Metrics

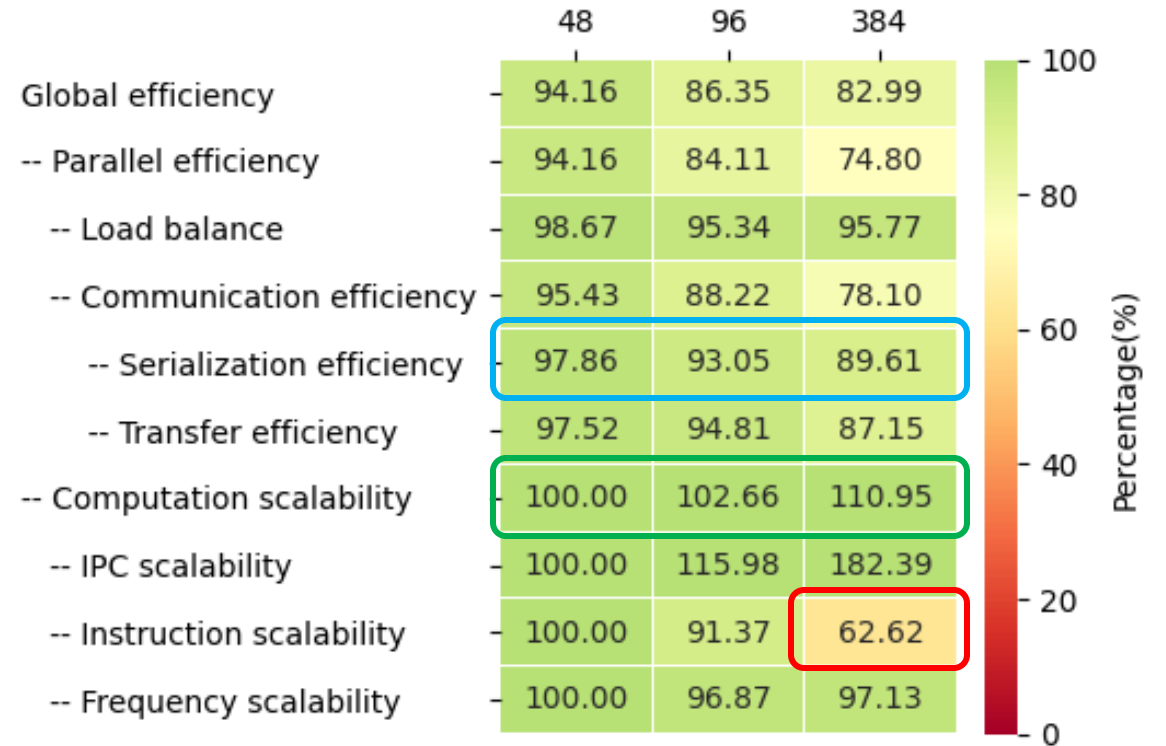


- What they are NOT
 - The end of the journey
- What they are...
 - ... a general mechanism to describe the fundamental concepts of parallelism
 - ... a hint that tells where to look
 - ... a way to quantify efficiency loss
 - ... a fair comparison between different...
 - ... code versions
 - ... architectures
 - ... core counts (scalability)
 - ... applications
- For each metric we are going to see:
 - What it quantifies
 - How it is computed
 - Formula
 - Graphical
 - Interpretation
 - Where to look next

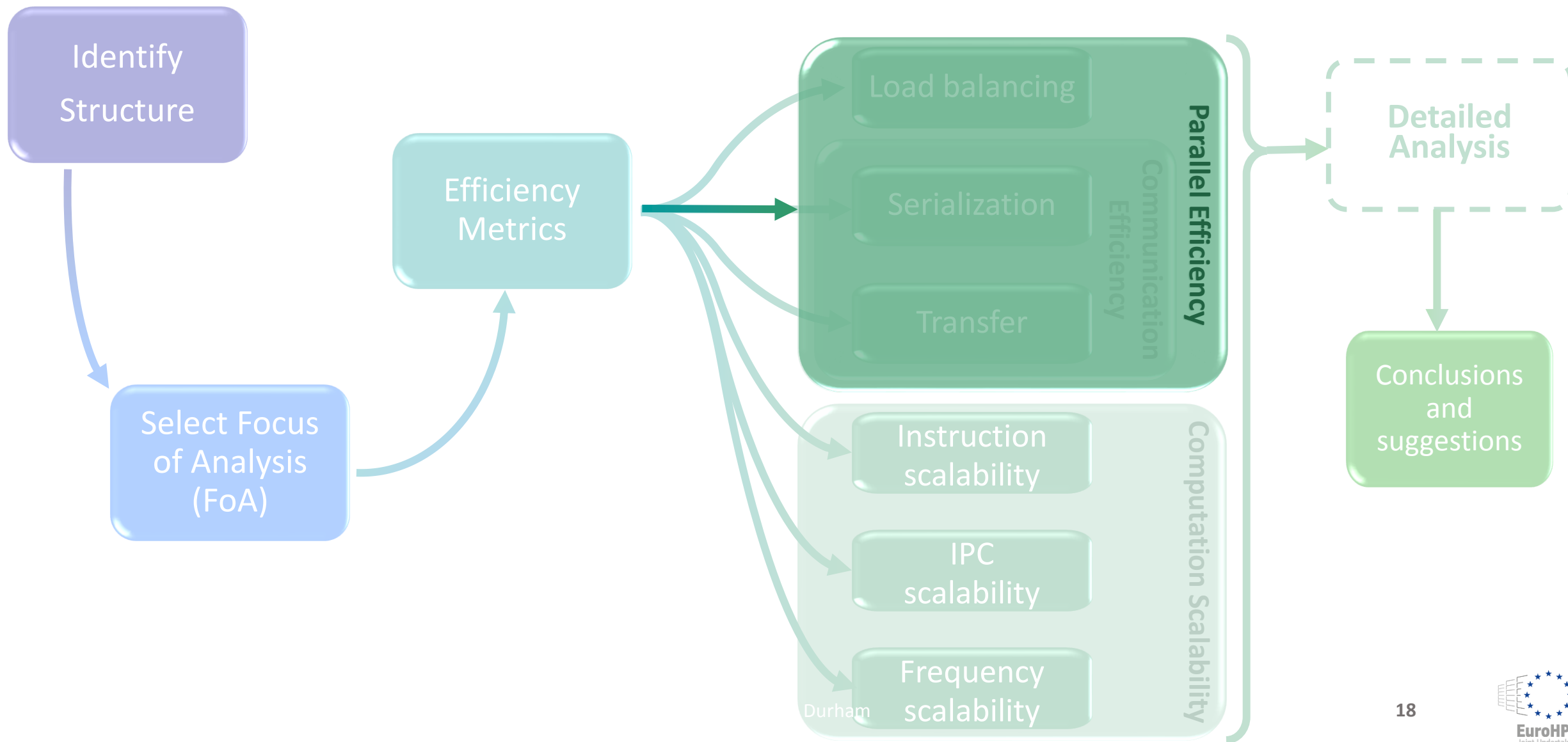
The Efficiency Metrics



- Usually shown in a table
 - Rows: Metrics
 - Columns: Different traces
- With colored cells as a heat map
- What to look for?
 - Low values
 - Trend
 - High values



The journey

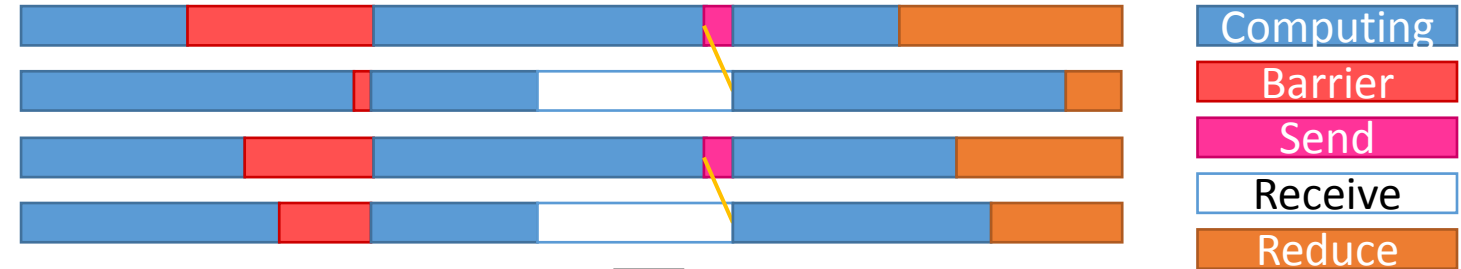


Some semantics first



- The state of a processes is simplified to two values:

- Useful == Computing
- Not useful == Otherwise

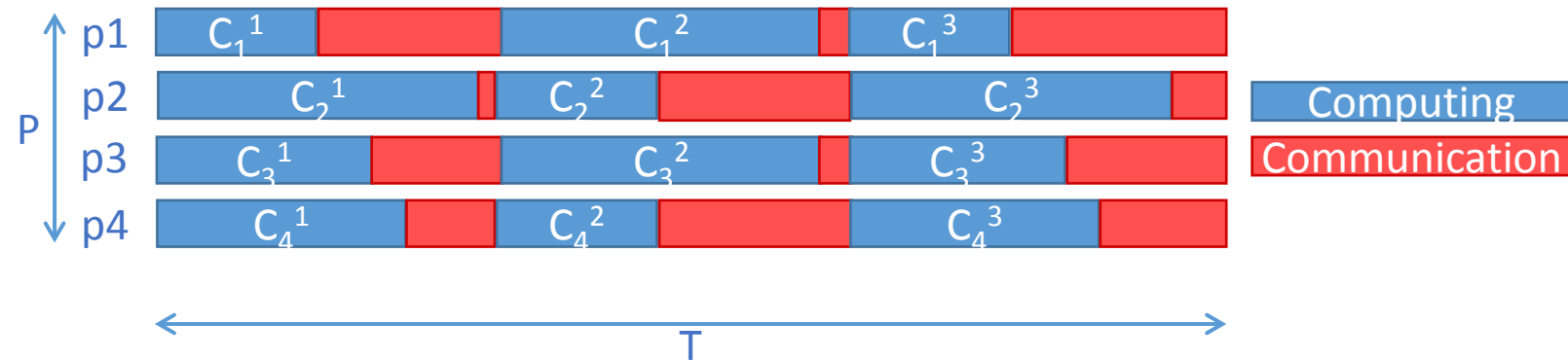


- T = Elapsed time
- P = Number of processes
- c_i = Compute time of process i

- $c_i = \sum_{j=1}^n c_i^j$

- C = Total compute time

- $C = \sum_{i=1}^P c_i$



Parallel Efficiency



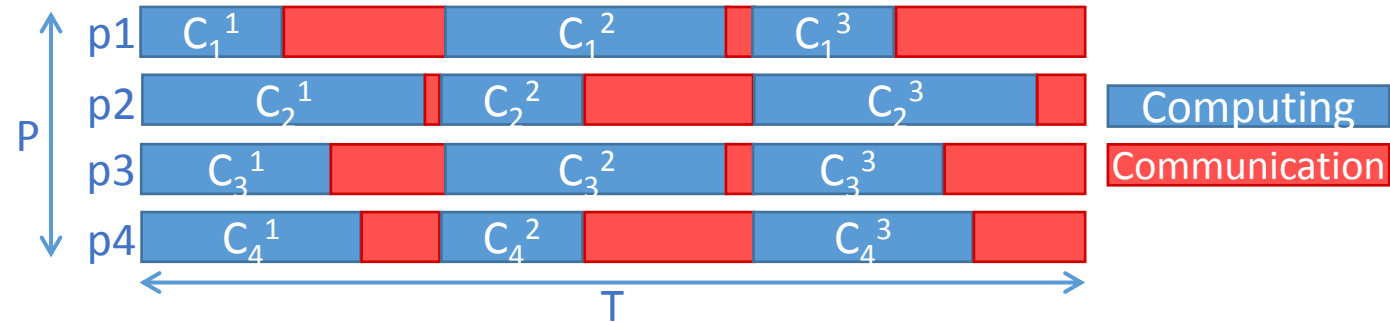
Quantifies: The extent to which all resources in the system are kept active doing useful work

How it is computed: Ratio between time used to do useful computation and consumed cpu time

$$\text{Parallel Efficiency} = \frac{\text{Computing}}{\text{Computing} + \text{Communication}}$$

How it is computed:

$$\text{Parallel Efficiency} = \frac{\sum_{i=1}^P c_i}{P * T}$$

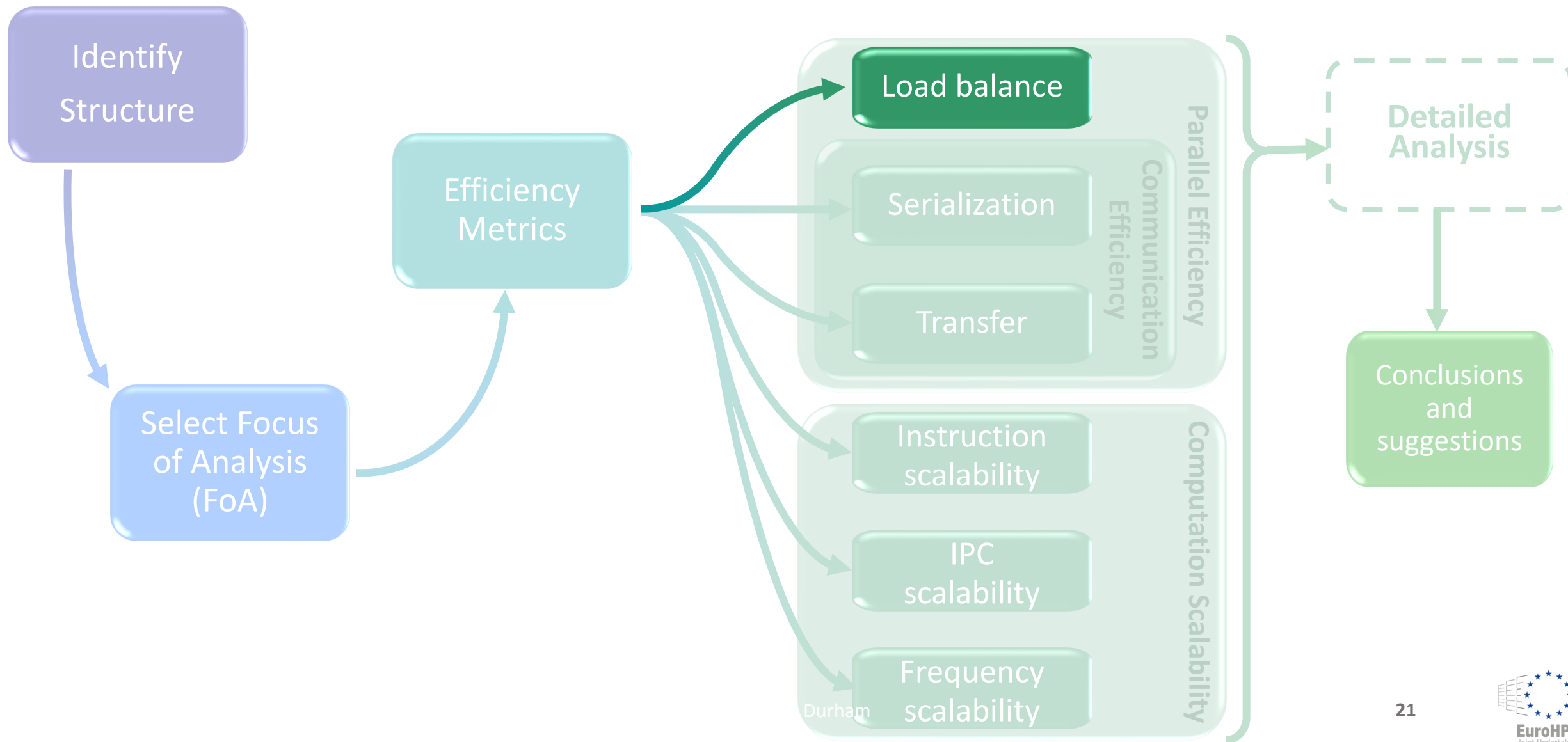


Interpretation: A low value indicates that a low fraction of the time consumed is used to do useful computation.

Where to look next?

- Look at its child metrics

The journey



Load Balance



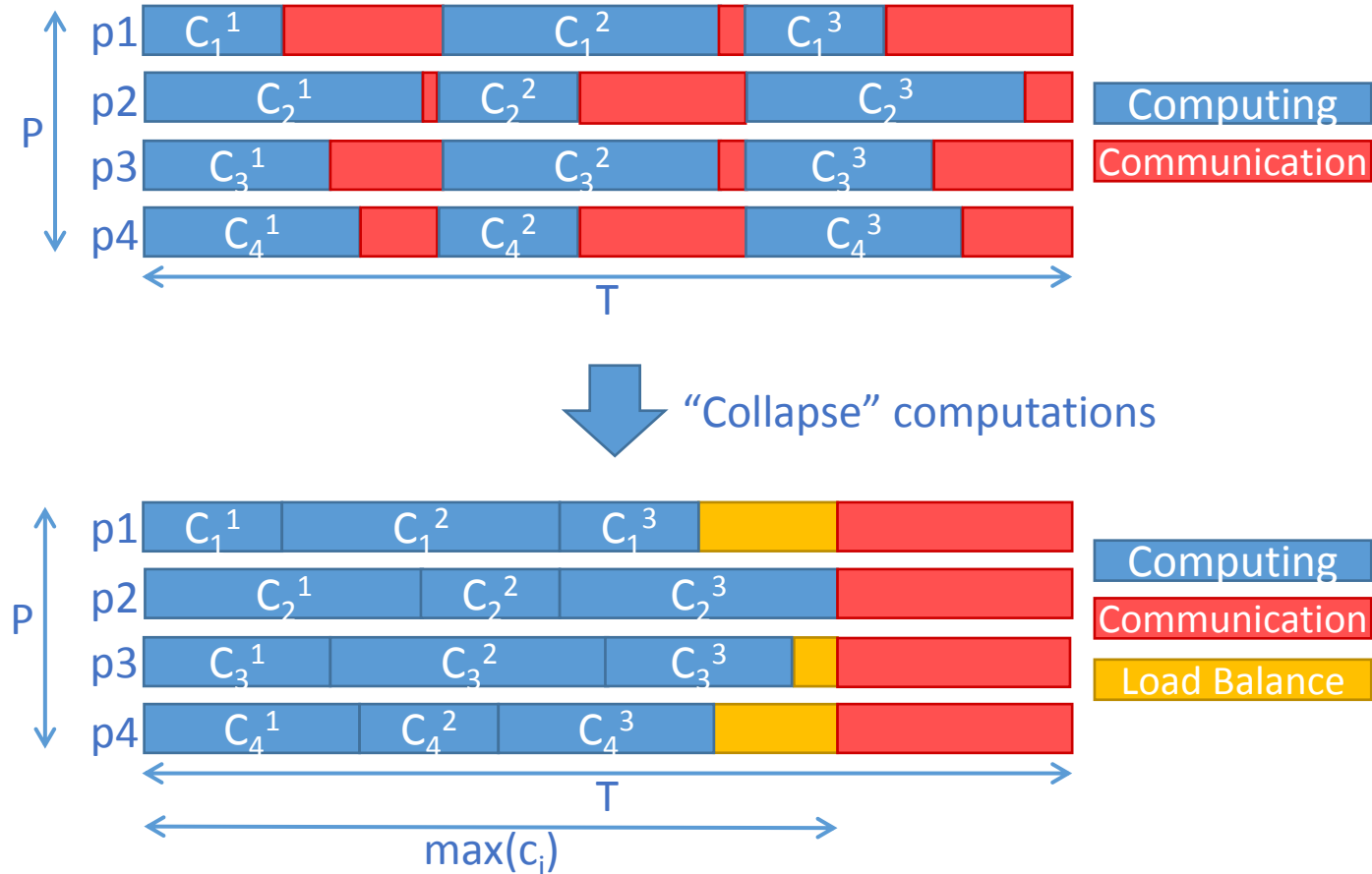
Quantifies: The efficiency loss due to the global distribution of work among processes.

How it is computed: Ratio between time used to do useful computation and the useful computation of the most loaded process multiplied by the number of processes

$$\text{Load Balance} = \frac{\text{blue}}{(\text{blue} + \text{yellow})}$$

How it is computed:

$$\text{Load Balance} = \frac{\sum_{i=1}^P C_i}{\max_{i=1}^P (c_i) * P}$$



Load Balance



Interpretation: A low value in this metric indicates that more highly loaded processes keep other processes idle for a significant amount of time.

- A single highly loaded process will make this metric report a low value



$$LB = \frac{2 + 1 + 1 + 1 + 1}{2 * 5} = \frac{6}{10} = 0.6$$

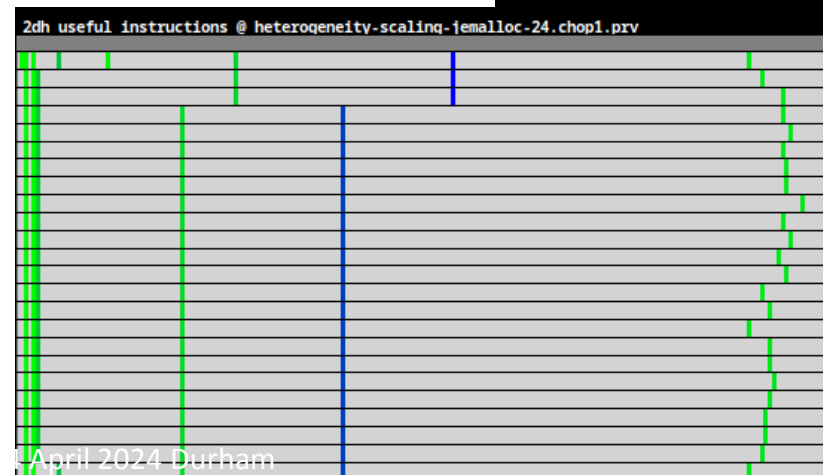
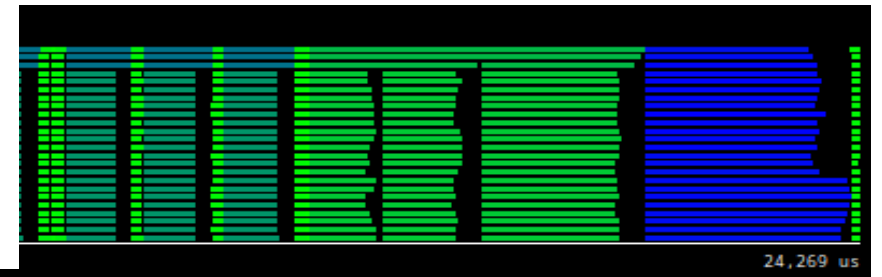
- A single low loaded process won't have an effect on this metric



$$LB = \frac{2 + 2 + 2 + 2 + 1}{2 * 5} = \frac{9}{10} = 0.9$$

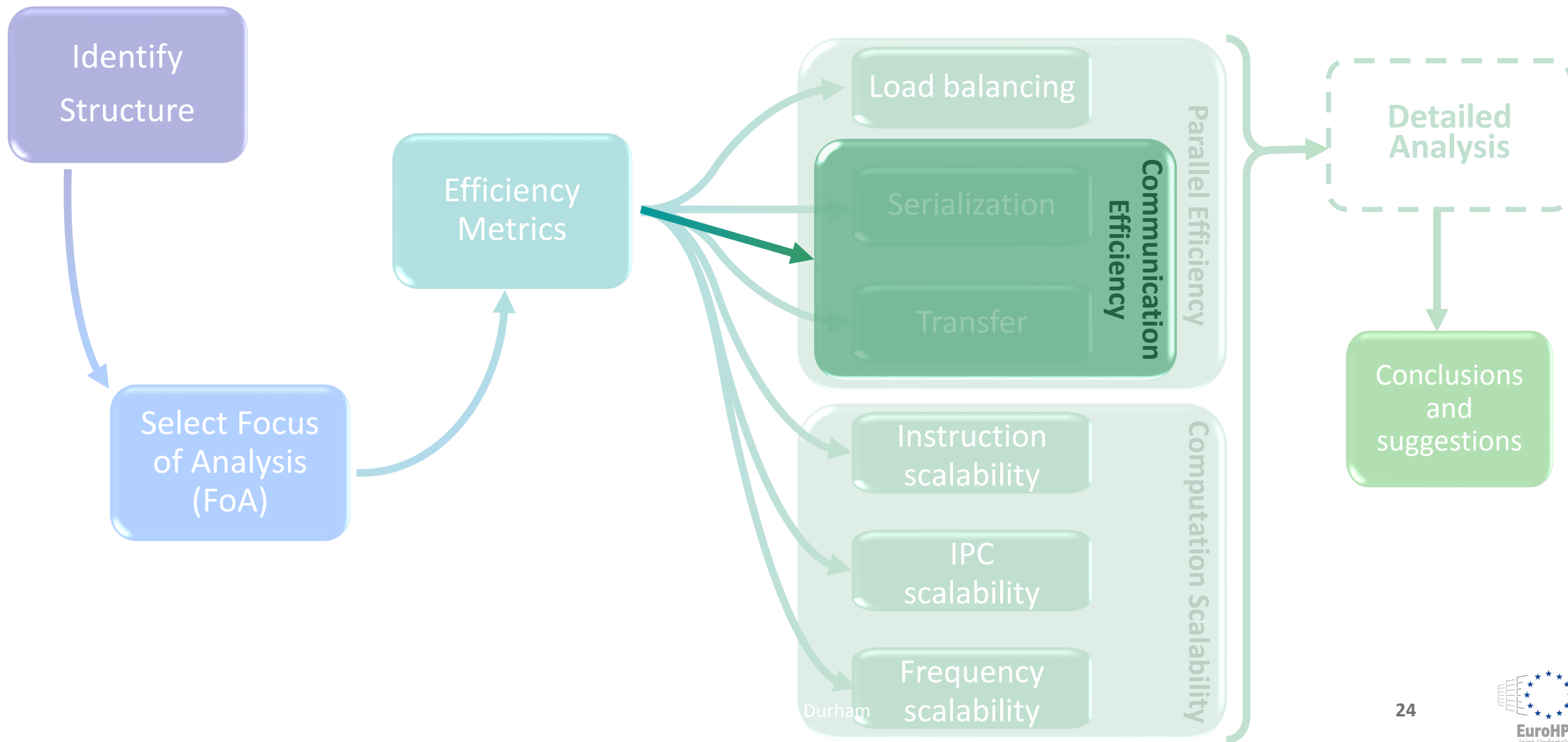
• What to look next?

- We try to understand the cause of the Load Imbalance
- In general can have 3 sources
 - Number of instructions → Histogram of useful instructions
 - IPC → Histogram of IPC
 - Frequency → Histogram of cycles per us



April 2024 Durham

The journey



Communication Efficiency



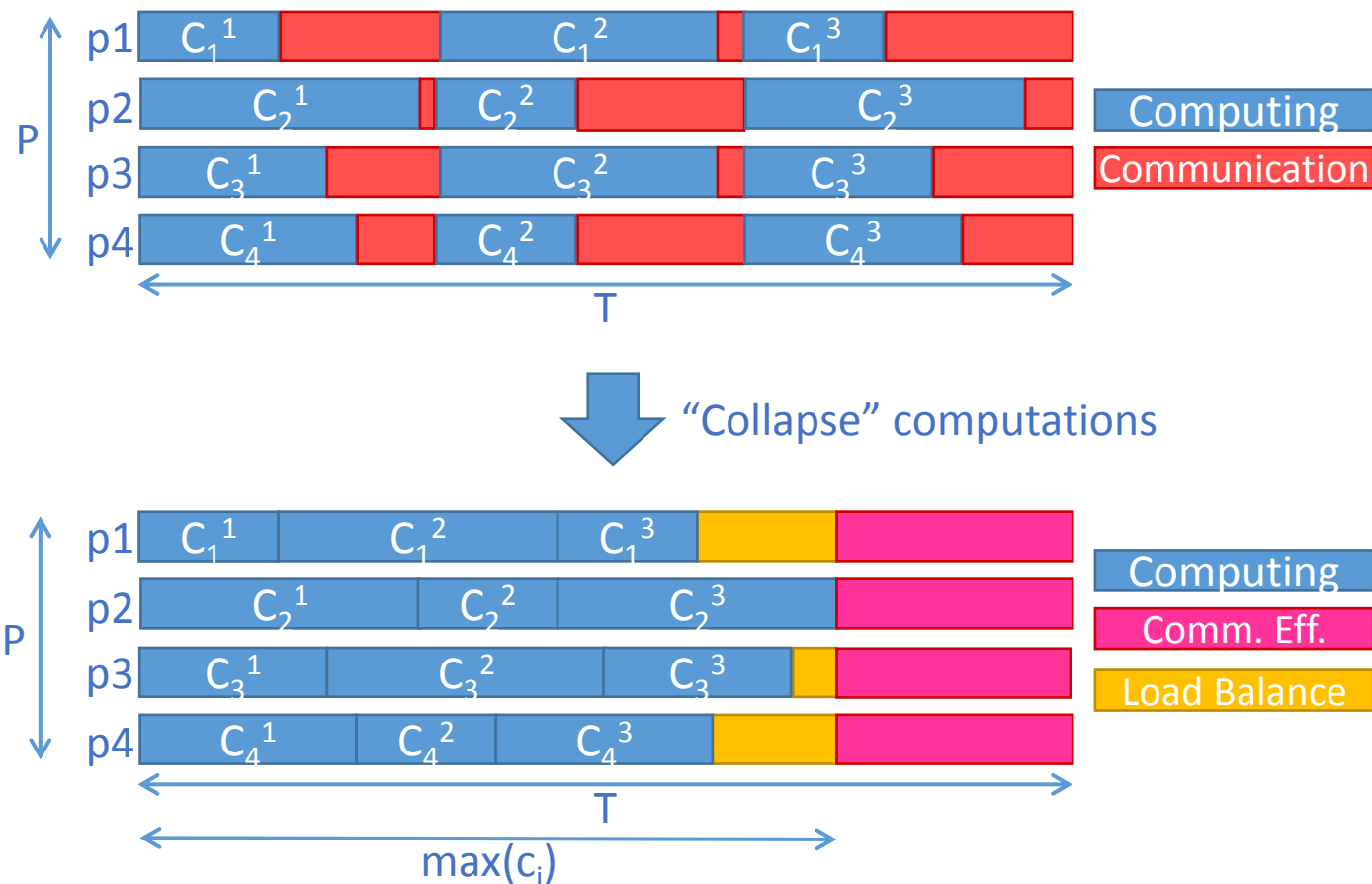
Quantifies: The efficiency loss due to the communication of data. Be it due to synchronizations between processes or to the overhead introduced by the communication itself. Excluding time loss due to global load imbalance

How it is computed: Ratio between the useful computation time of the most loaded processes and the total elapsed time

Communication Eff.= ( + ) / ( +  + )

How it is computed:

$$\text{Communication Eff.} = \frac{\max_{i=1}^P(c_i)}{T}$$

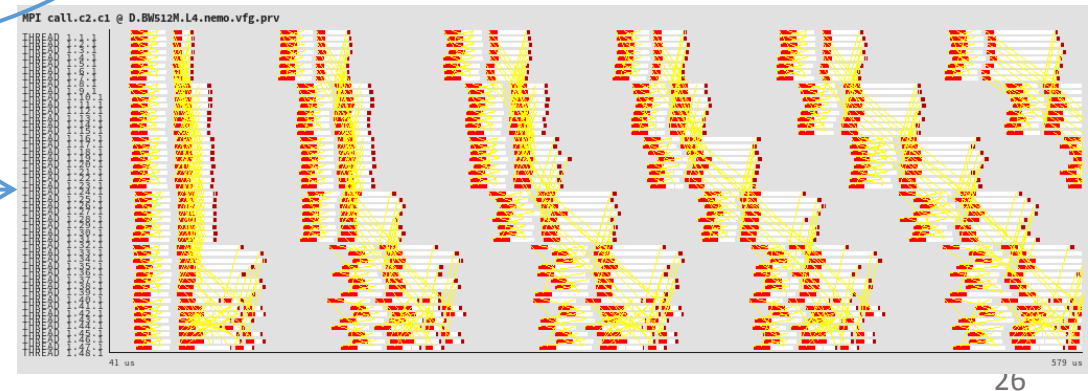
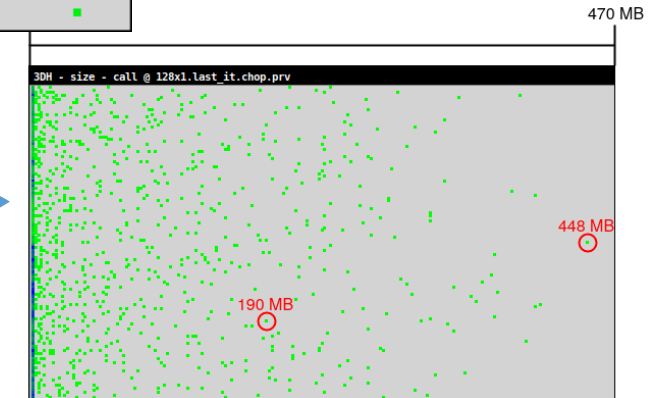
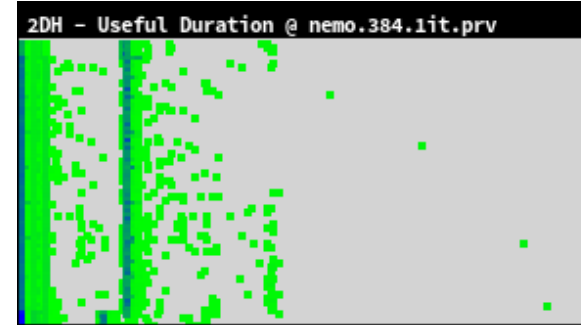


Communication Efficiency

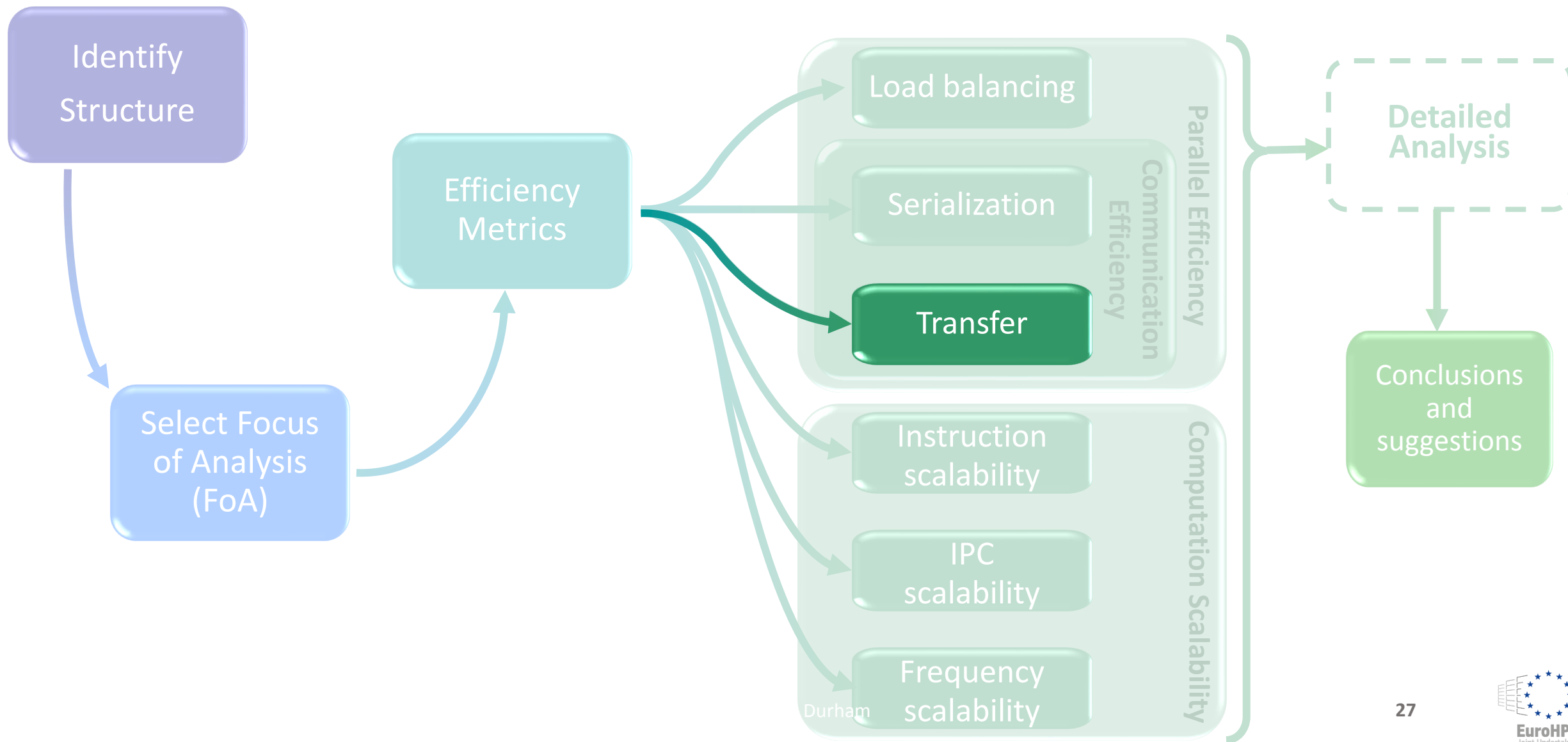


Interpretation: A low value in this metric indicates that the interaction between processes is impacting the performance.

- This metric can report good values in codes where a profile reports significant time in MPI and that would be reported by a bad load balance efficiency
- What to look next?
- If possible child metrics
- If not...
 - How many MPI calls are made?
 - Histogram of MPI calls
 - How often? Granularity of computations
 - Useful duration
 - How much data is sent?
 - Bytes sent per MPI call
 - Which are the semantics of the communication?
 - Chains of dependences?
 - MPI calls



The journey



Transfer Efficiency



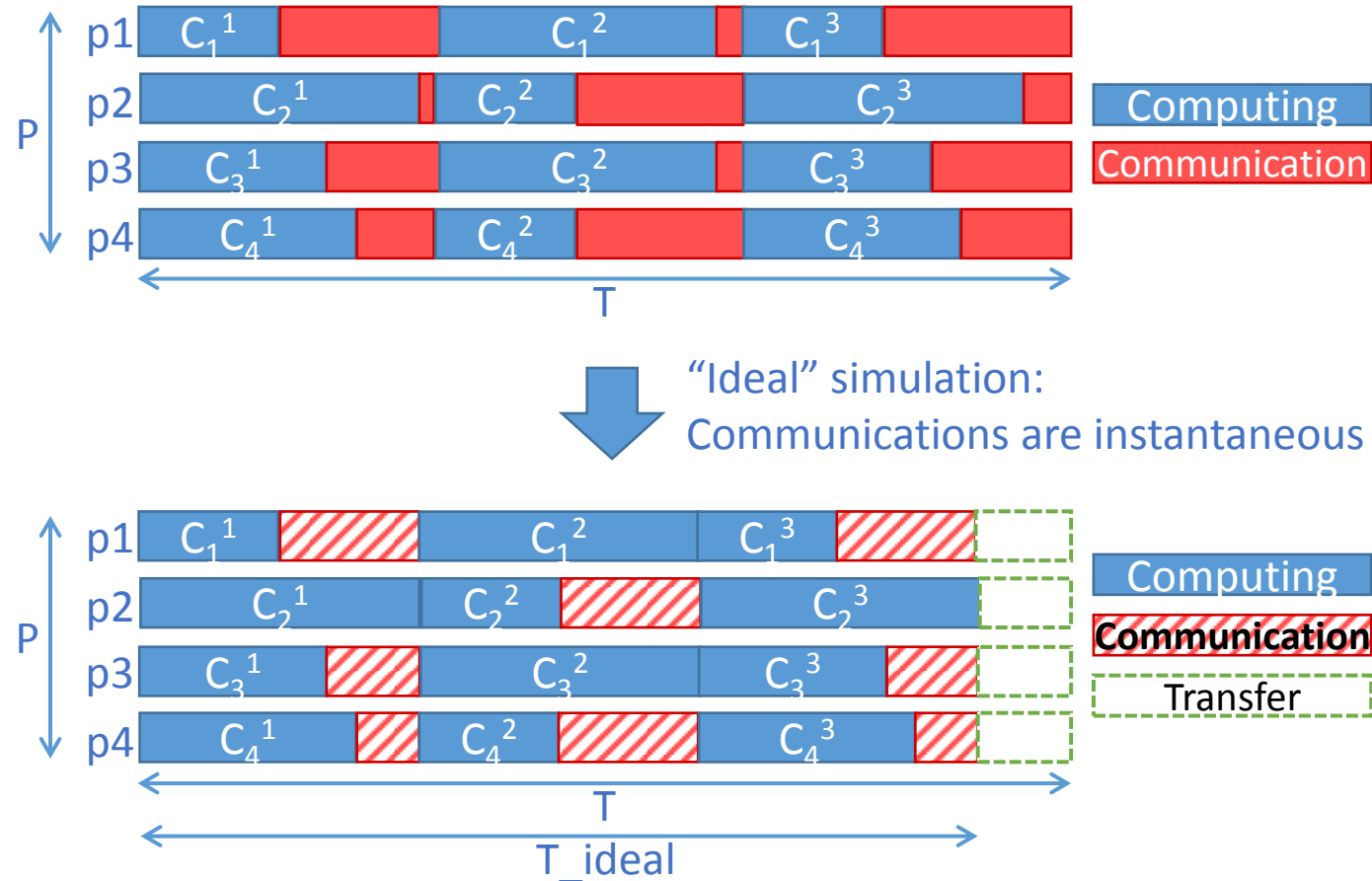
Quantifies: Efficiency loss related to the non instantaneous nature of communication mechanisms. Includes time to transmit the data over the physical channel and the overhead in the libraries.

How it is computed: Ratio between elapsed time in the ideal simulation and the elapsed time in the real execution

$$\text{Transfer Eff.} = (\text{blue} + \text{red}) / (\text{blue} + \text{red} + \text{green})$$

How it is computed:

$$\text{Transfer Eff.} = \frac{T_{\text{ideal}}}{T}$$

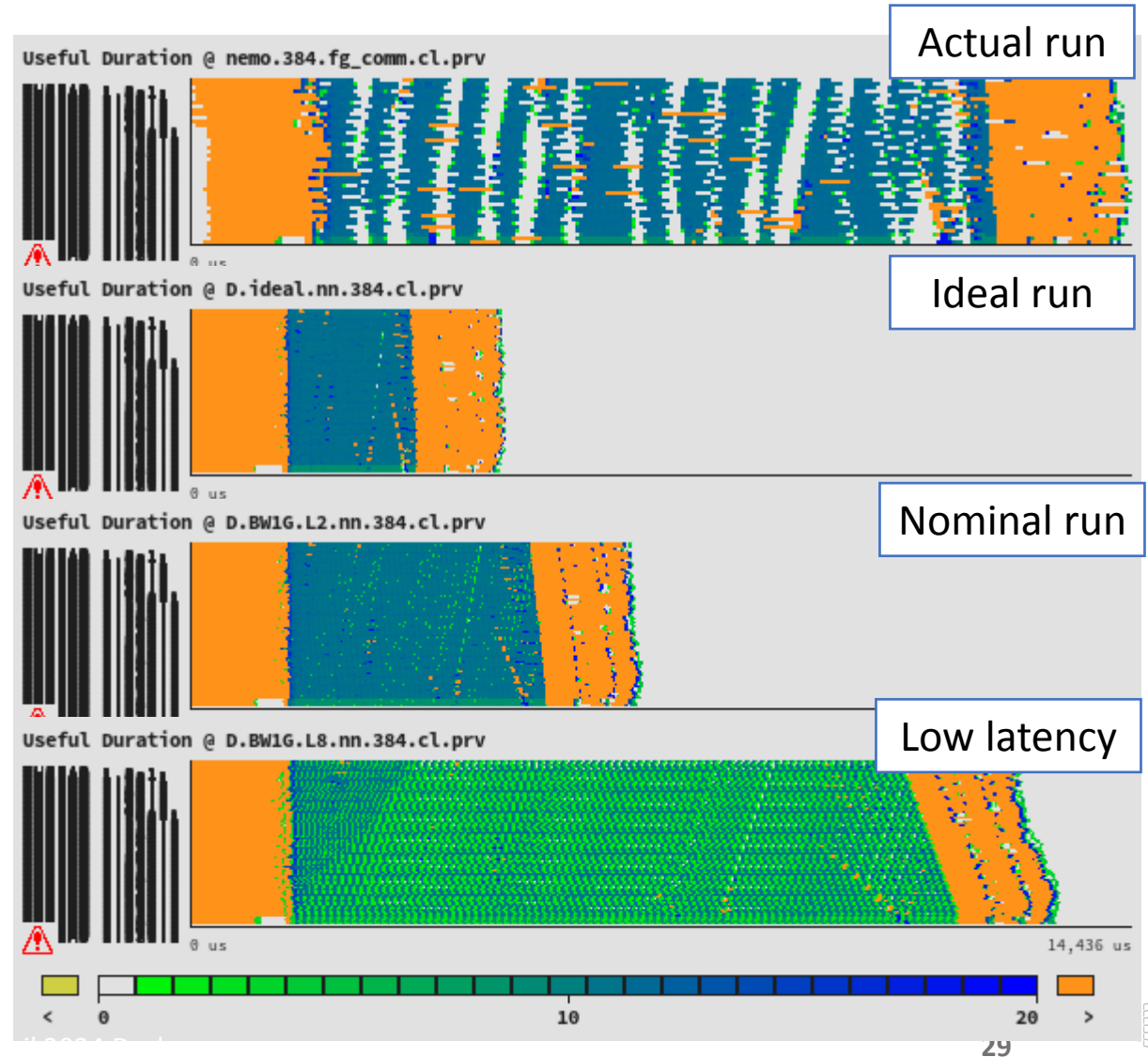


Transfer Efficiency

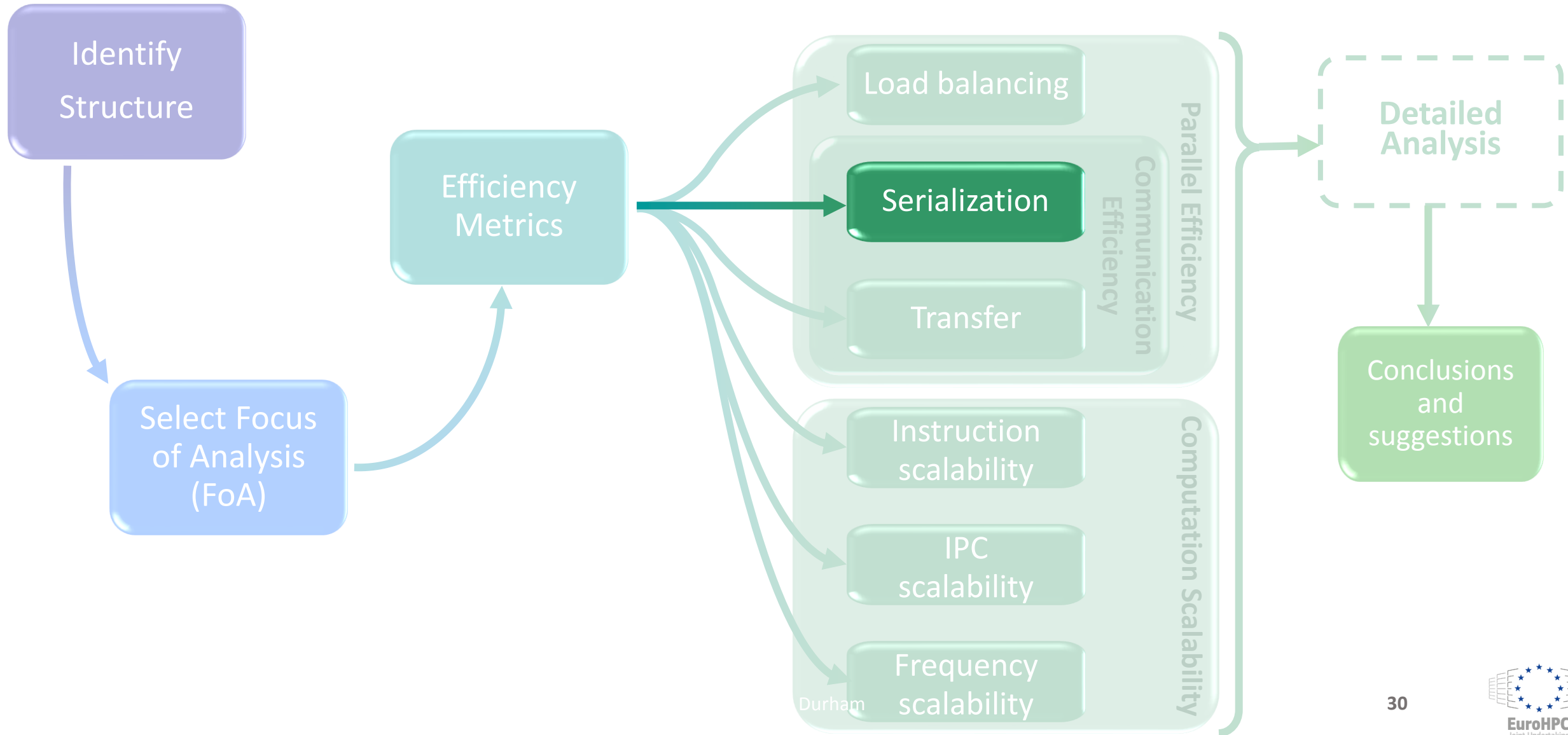


Interpretation: Low values indicate that the execution is suffering from a high overhead of the runtime or a poor latency or bandwidth of the network.

- What to look next?
 - Determine if the transfer problem is Bandwidth or Latency
 - Note that in Latency we include the overhead of the communication library
 - Use Simulations of Dimemas
 - Different BW and latency



The journey



Serialization Efficiency



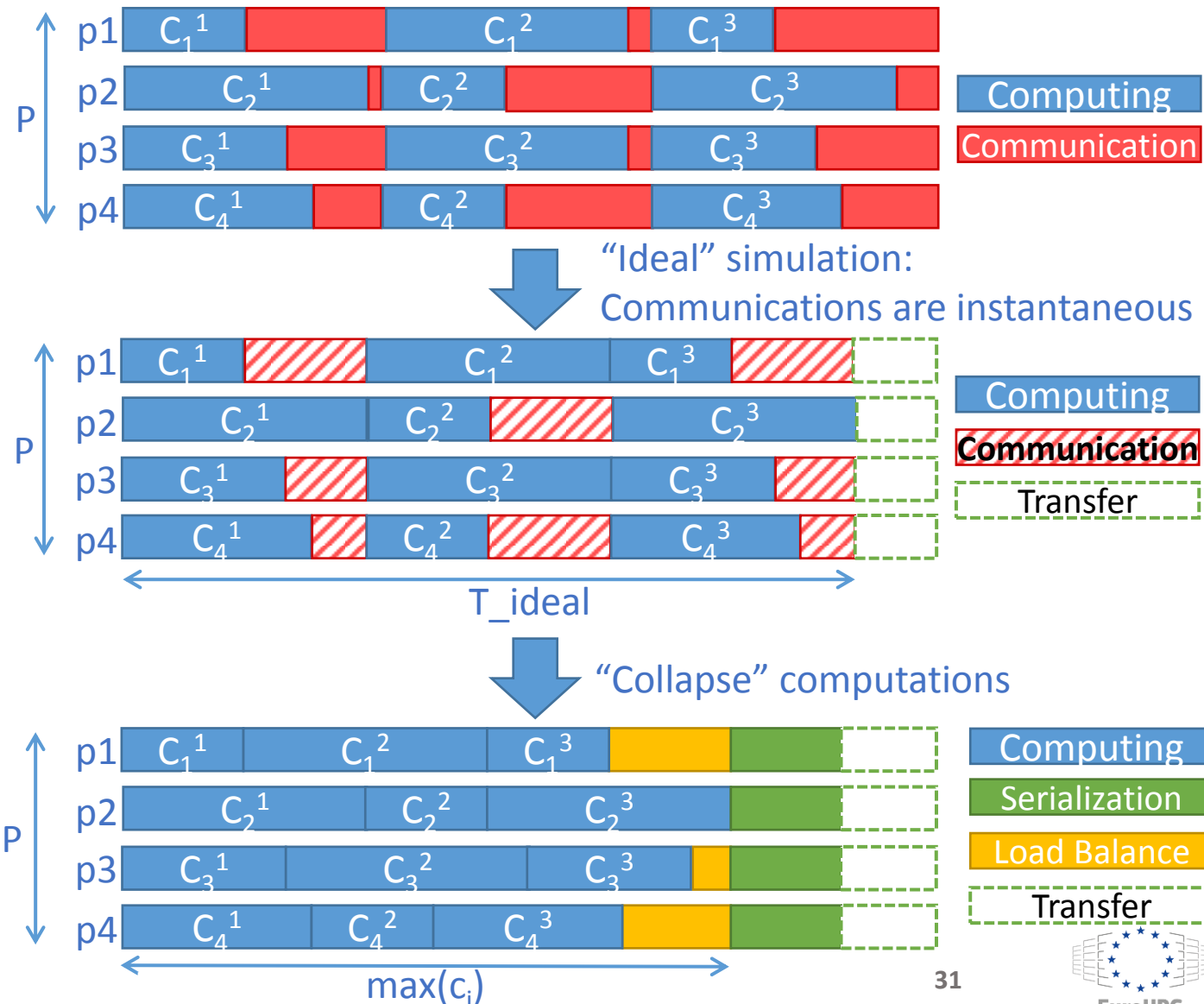
Quantifies: Inefficiencies caused by circular dependences or non-uniform imbalances

How it is computed: Ratio between useful computation of most loaded process and elapsed time in the ideal simulation

$$\text{Serialization Eff.} = (\text{blue} + \text{yellow}) / (\text{blue} + \text{red})$$

How it is computed:

$$\text{Serialization Eff.} = \frac{\max_{i=1}^P(c_i)}{T_{ideal}}$$

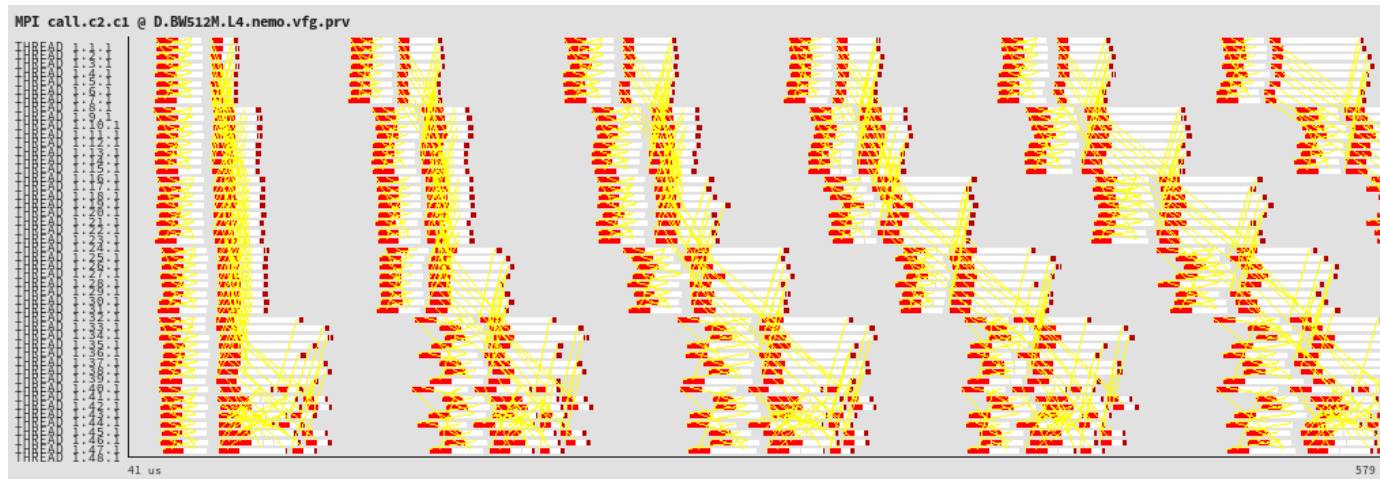


Serialization Efficiency

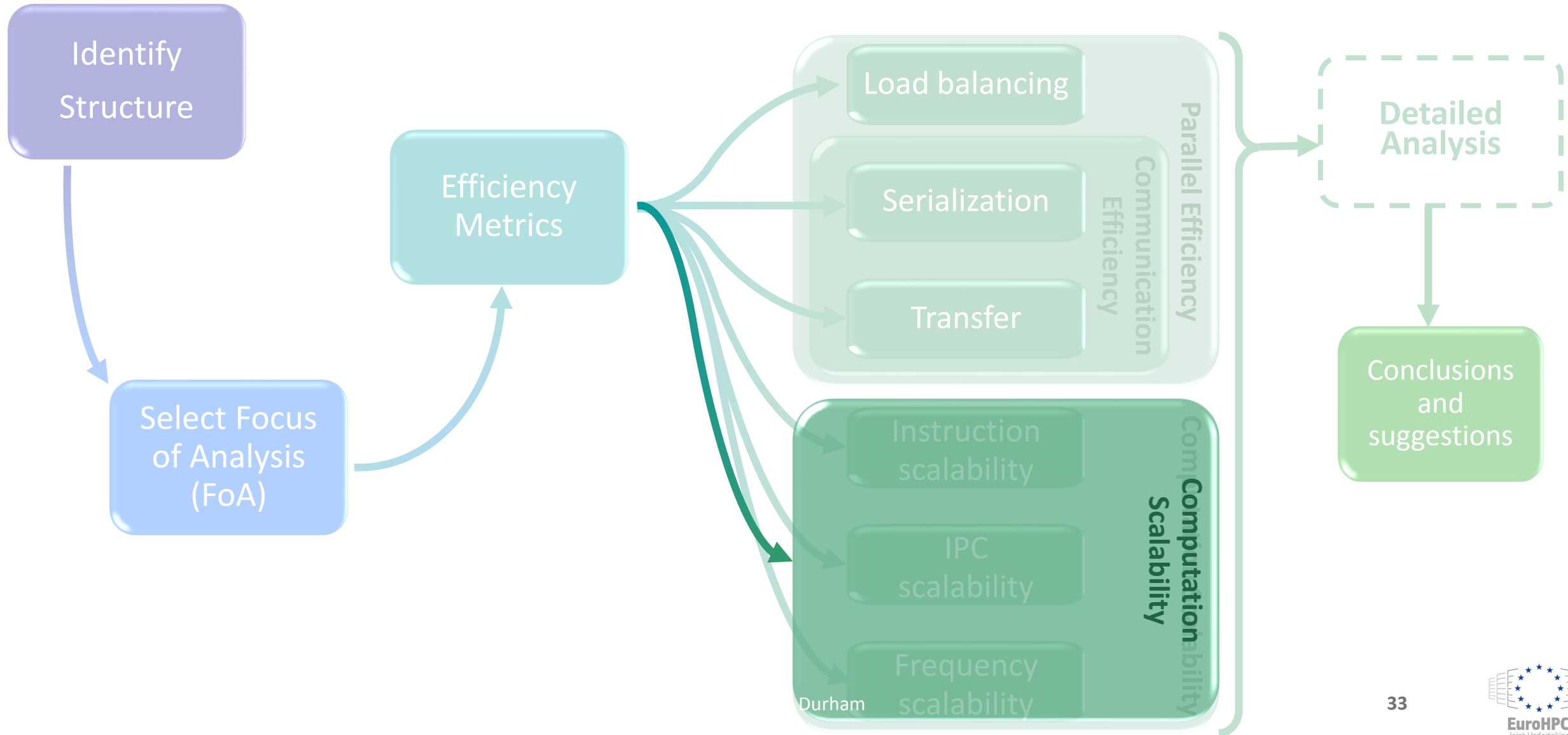


Interpretation: A low value indicates the existence of circular dependences. They can be caused by actual algorithmic serialization, irregularities in the load of processes during the execution or noise.

- What to look next?
 - Discard noise
 - Cycles per us
 - Try to identify causes for circular waits.
 - Understand semantics of the communication
 - MPI Calls



The journey



Computation Scalability



Quantifies: How the time spent computing scales with respect to the reference case.

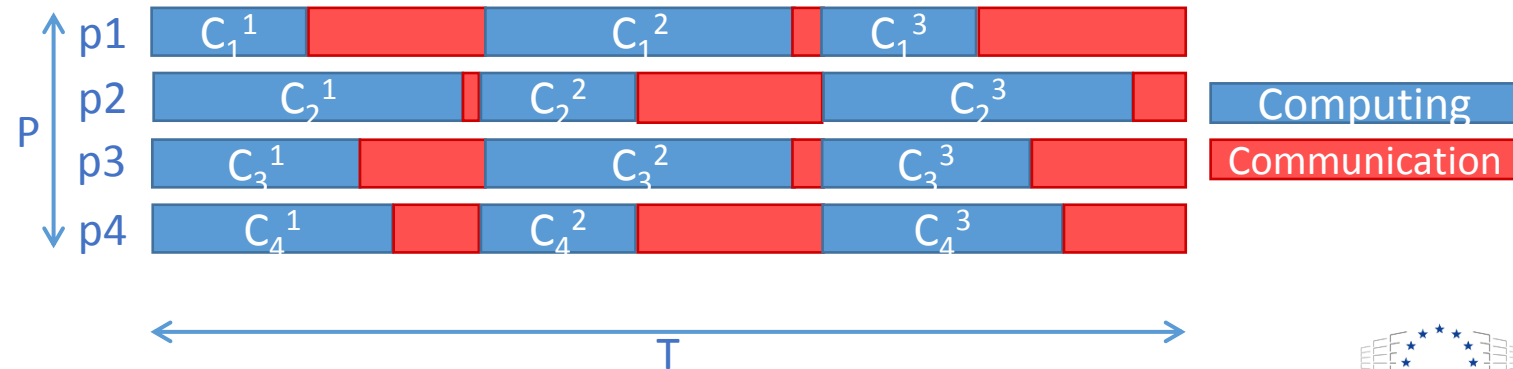
How it is computed: Ratio between time spent on useful computation in the reference case and the time spent on useful computation on current run.

How it is computed:

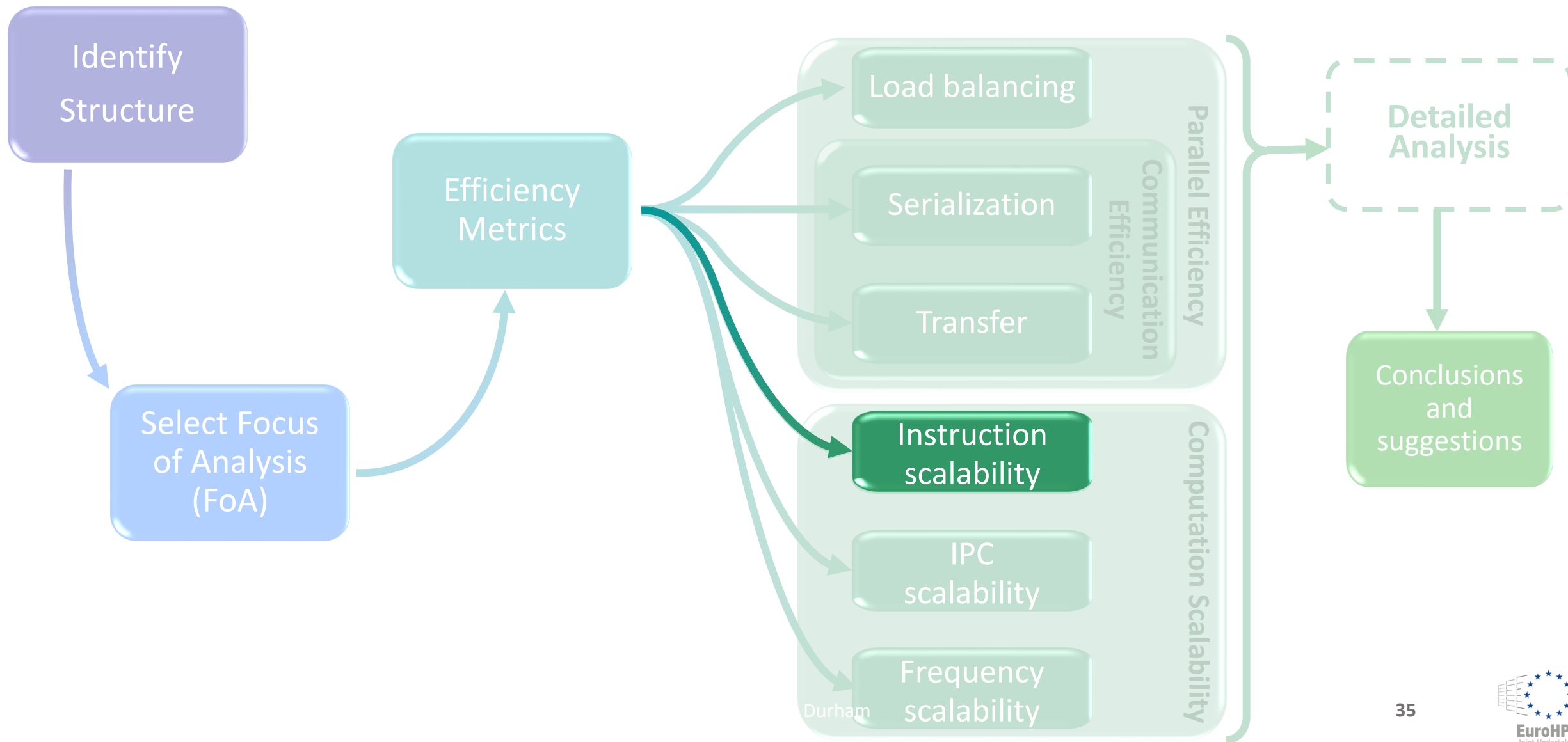
$$\text{Computation Sc.} = \frac{C_{\text{ref}}}{C_{\text{current}}}$$

Interpretation: A low value indicates time grows per core count. Ideally the total compute time to solve a problem should be constant, independent of core count (in strong scaling).

- Relative metric, based on a reference case.
 - C_{ref} = Useful Computation of the reference run
 - C_{current} = Useful computation current run
- 100% for the reference case
- Can be weak or strong scaling
 - Strong scaling for all the formulas presented
 - In strong scaling we assume total compute time and total number of instructions should remain constant as we increase the number of processes
 - In an analogous way weak scaling can be computed
- High level metric composed by three child metrics based on:
 - $T = \# \text{ instr.} / (\text{IPC} * \text{freq})$



The journey



Instructions Scalability



Quantifies: How the number of instructions scales with respect to the reference case.

How it is computed: Ratio between number of instructions executed on the reference case and the number of instructions executed on current run.

How it is computed:

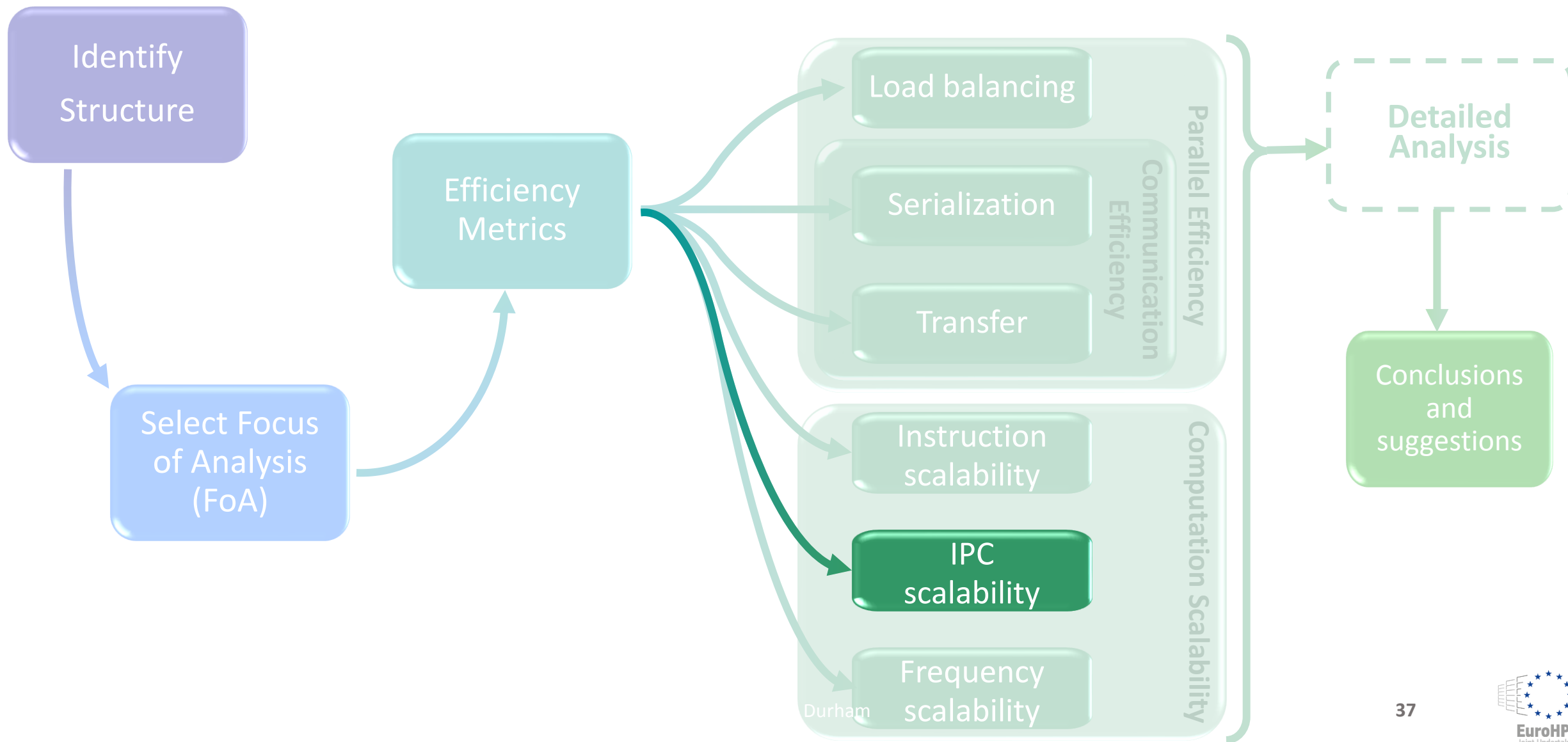
$$\text{Instructions Sc.} = \frac{I_{\text{ref}}}{I_{\text{current}}}$$

- Relative metric, based on a reference case.
 - I_{ref} = Total number of instructions executed on the reference run
 - I_{current} = Total number of instructions executed on the current run
- 100% for reference case

Interpretation: A value less than 100 indicate that the total number of instructions to solve the problem grows with core count, which ideally should not be the case.

May be caused by code replication (computed by all processes, by an increase in the surface to volume ratio when computations on the surface/boundary are "replicated", ...

The journey



Quantifies: How the IPC scales with respect to the reference case.

How it is computed: Ratio between average IPC on the current run and the average IPC on the reference run.

How it is computed:

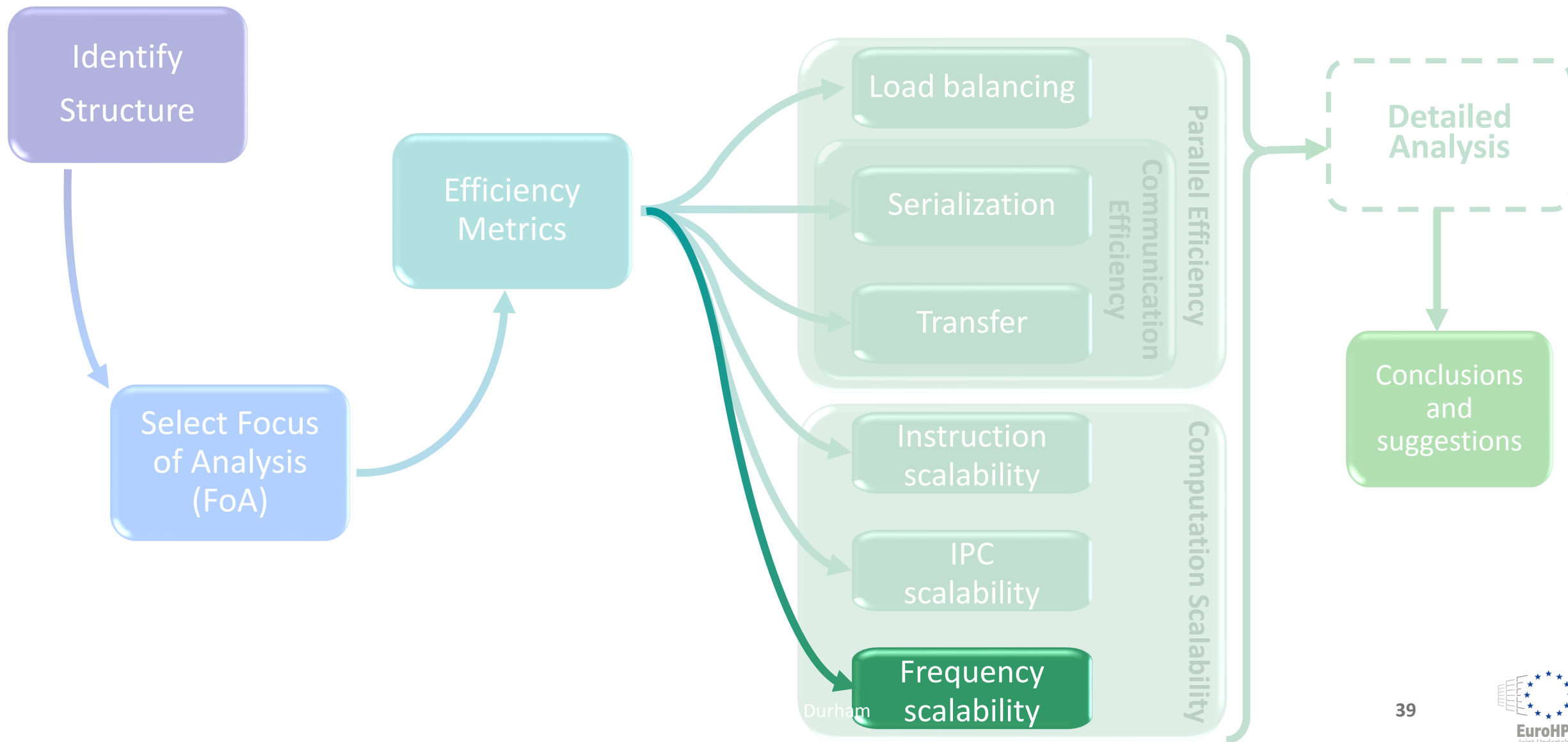
$$IPC_{Sc.} = \frac{IPC_{current}}{IPC_{ref}}$$

- Relative metric, based on a reference case.
 - IPC_{ref} = Average IPC of the reference run
 - $IPC_{current}$ = Average IPC current run
- 100% otherwise

Interpretation: A value less than 100 indicates that the IPC for the specific core count is worse than that of the reference case. May be caused by different locality behavior, contention on resources such as memory bandwidth, ...

A value above 100 indicates a higher IPC than the reference case this can be produced by cache effects for example.

The journey



Frequency Scalability



Quantifies: How the frequency scales with respect to the reference case.

How it is computed: Ratio between average frequency on the current run and the average frequency on the reference run.

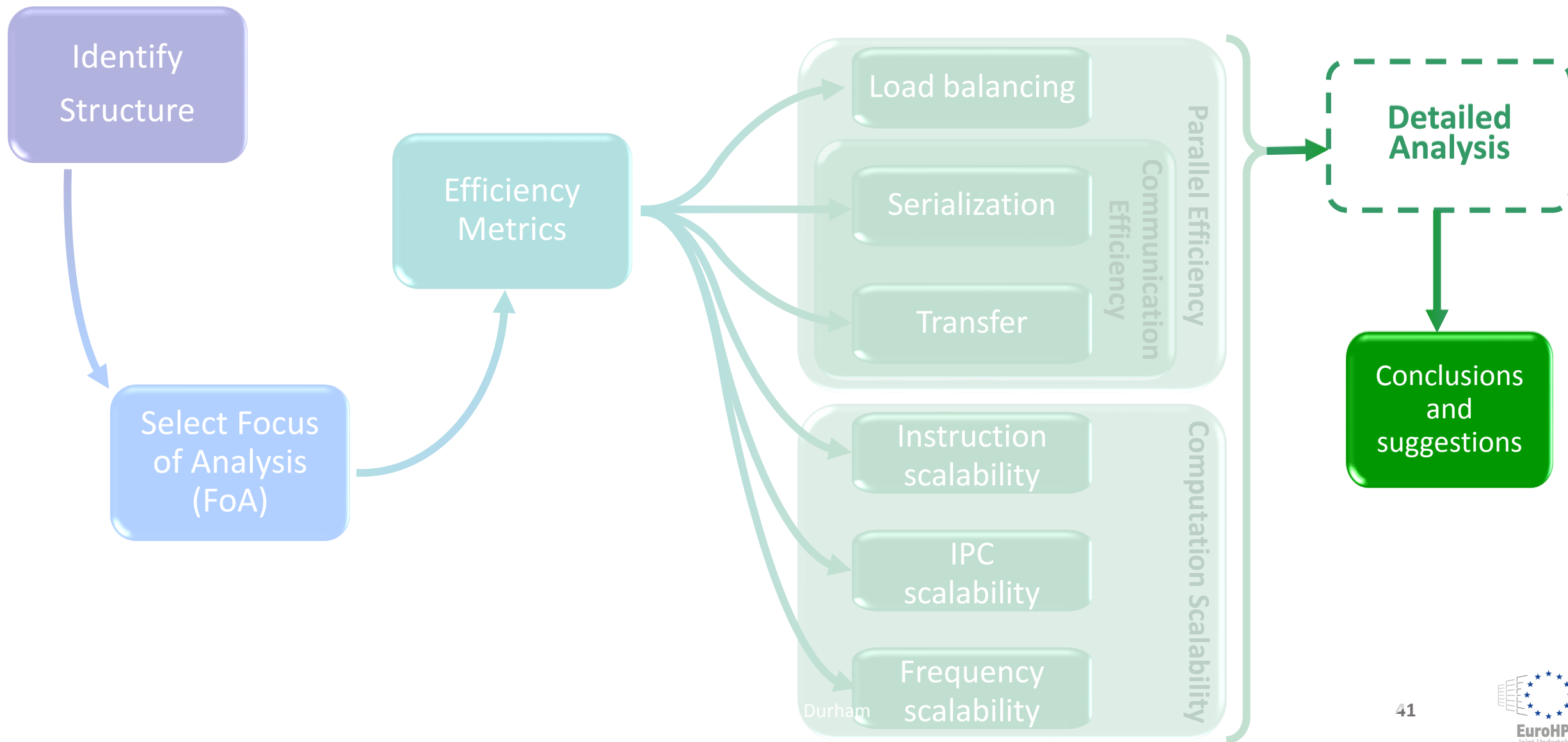
How it is computed:

$$Frequency\ Sc. = \frac{Freq_{current}}{Freq_{ref}}$$

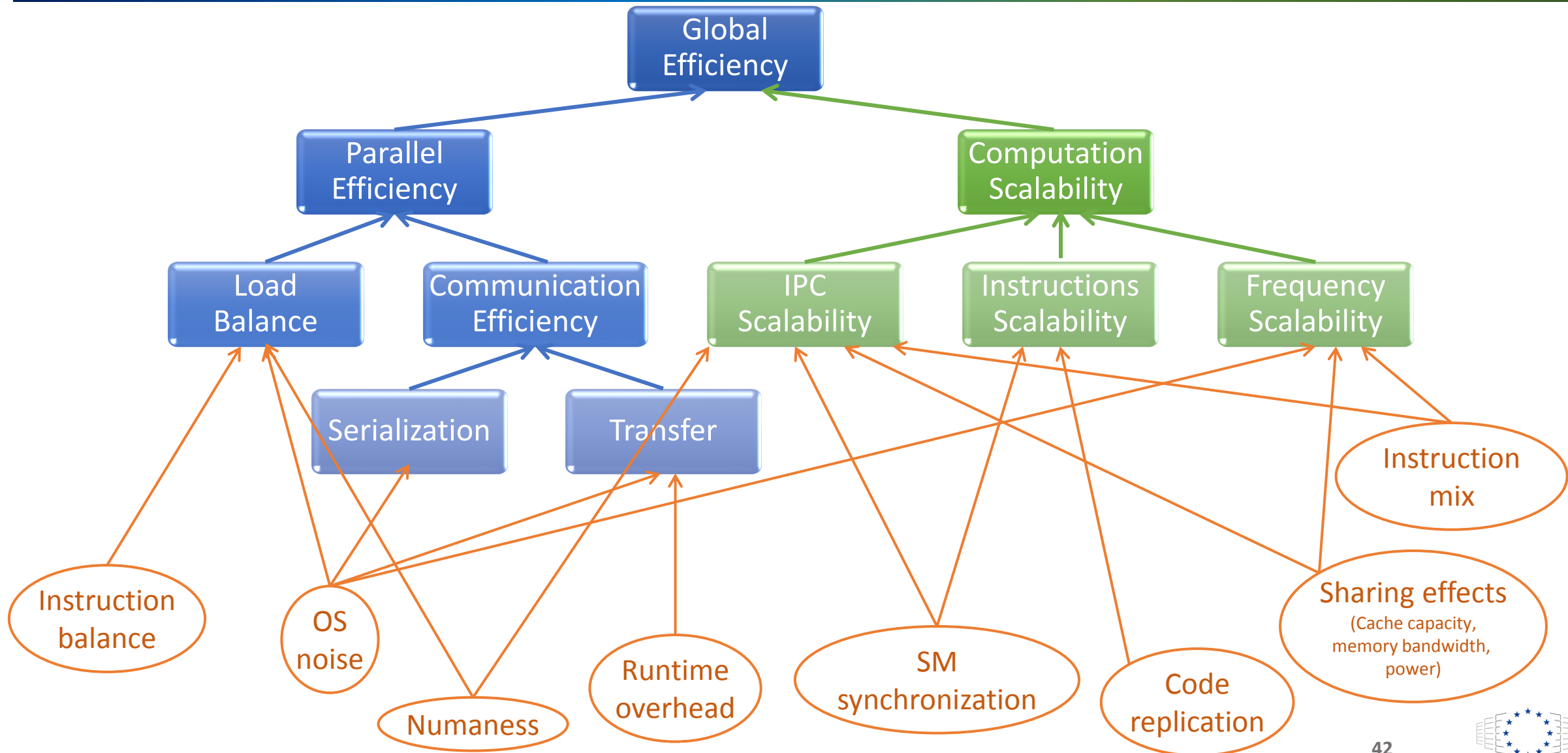
- Relative metric, based on a reference case.
 - $Freq_{ref}$ = Average frequency on the reference run
 - $Freq_{current}$ = Average frequency on the current run
- 100% for the reference case

Interpretation: A value less than 100 indicates that the frequency is lower than the reference frequency. This may be caused by "preemptions", power management measures, ...

The journey



Factors and causes





Performance Optimisation and Productivity 3

A Centre of Excellence in HPC

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 [@POP_HPC](#)

 [youtube.com/POPHPC](https://www.youtube.com/POPHPC)



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