# Even Simpler Streaming and Deterministic Matrix Sketching

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## **Abstract**

This paper provides a one-line proof of Frequent Directions (FD) for sketching streams of matrices. It simplifies the main results in [1] and [2]. The simpler proof arises from sketching the *covariance* of the stream of matrices rather than the stream itself.

### Introduction

Let  $X_t \in \mathbb{R}^{d \times n_t}$  be a stream of matrices. Let  $C = \sum_{t=1}^T X_t X_t^T \in \mathbb{R}^{d \times d}$  be their covariance matrix. Frequent Directions [1] maintains a rank deficient approximate covariance matrix  $\tilde{C}_t \in \mathbb{R}^{d \times d}$  using Algorithm 1. Set  $\tilde{C}_0 \in \mathbb{R}^{d \times d}$  to be the all zeros matrix. Then, at time  $t = 1, \ldots, T$  compute  $\tilde{C}_t = \text{Update}(\tilde{C}_{t-1}, X, \ell)$ .

#### Algorithm 1 Frequent Directions (FD) Update

- 1: **function** UPDATE $(\tilde{C}, X, \ell)$
- 2:  $U\Lambda U^T = \tilde{C} + XX^T$
- 3: **return**  $U \cdot \max(\Lambda I \cdot \lambda_{\ell}, 0) \cdot U^T$
- 4: end function

Above,  $U\Lambda U^T$  is the eigen-decomposition of  $\tilde{C}+XX^T$  and  $\lambda_\ell$  is the its  $\ell$ 'th largest eigenvalue. Note that the rank of  $\tilde{C}_t$  is at most  $\ell-1$  for all t by construction. It can therefore be stored in  $O(d\ell)$  space. Assuming  $n_t < \ell$ , the update operation itself also consumes at most  $O(d\ell)$  space.

**Lemma 1** (simplified from [2] and [1]). Let  $\tilde{C}$  denote the approximated covariance produced by FD and  $\lambda_i$  be the eigenvalues of the exact covariance C in descending order. For any  $\ell$  and simultaneously for all  $k < \ell$  we have

$$||C - \tilde{C}|| \le \frac{1}{\ell - k} \sum_{i=k+1}^{d} \lambda_i$$

## Short proof Lemma 1

Define  $\Delta_t = X_t X_t^T - \tilde{C}_t + \tilde{C}_{t-1}$ . Then  $\sum_{t=1}^T \Delta_t = \sum_{t=1}^T X_t X_t^T - \sum_{t=1}^T (\tilde{C}_t - \tilde{C}_{t-1}) = C - \tilde{C}$  where  $\tilde{C}$  stands for  $\tilde{C}_T$ , the final sketch.

Moreover, note that the top  $\ell$  eigenvalues of  $\Delta_t$  are all equal to one another because  $\Delta_t = U_t \cdot \min(\Lambda_t, I \cdot \lambda_\ell^t) \cdot U_t^T$ . As a result  $\|\Delta_t\| < \frac{1}{\ell-k} \operatorname{tr}(\bar{P}_k \Delta_t \bar{P}_k)$  for any projection  $\bar{P}_k$  having a null space of dimension at most k. Specifically, this holds for  $\bar{P}_k$  whose null space contains the eigenvectors of C corresponding to its largest eigenvalues.

$$||C - \tilde{C}|| = ||\sum_{t=1}^{T} \Delta_t|| \le \sum_{t=1}^{T} ||\Delta_t||$$

$$\le \frac{1}{\ell - k} \operatorname{tr} \left( \bar{P}_k \left( \sum_{t=1}^{T} \Delta_t \right) \bar{P}_k \right)$$

$$\le \frac{1}{\ell - k} \operatorname{tr} \left( \bar{P}_k C \bar{P}_k \right) = \frac{1}{\ell - k} \sum_{i=k+1}^{d} \lambda_i$$

Here we used that  $\operatorname{tr}(\bar{P}_k\tilde{C}\bar{P}_k) \geq 0$  because  $\tilde{C}$  (and therefore  $\bar{P}_k\tilde{C}\bar{P}_k$ ) is positive semidefinite. This completes the proof.

#### References

- [1] Edo Liberty. Simple and deterministic matrix sketching. In *The 19th ACM SIGKDD International Conference on Knowledge Discovery and Data Mining, KDD 2013, Chicago, IL, USA, August 11-14, 2013*, pages 581–588. ACM, 2013.
- [2] Mina Ghashami and Jeff M. Phillips. Relative errors for deterministic low-rank matrix approximations. In Chandra Chekuri, editor, Proceedings of the Twenty-Fifth Annual ACM-SIAM Symposium on Discrete Algorithms, SODA 2014, Portland, Oregon, USA, January 5-7, 2014, pages 707-717. SIAM, 2014.

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