

# Bargaining Games

An Application of Sequential Move Games

## The Bargaining Problem

- The “Bargaining Problem” arises in economic situations where there are **gains from trade**, for example, when a buyer values an item more than a seller.
- The problem is *how to divide the gains, for example, what price should be charged?*
- Bargaining problems arise when the size of the market is *small and there are no obvious price standards* because the good is *unique, e.g., a house at a particular location. A custom contract to develop a web page, etc.*
- We can describe **bargaining games (in extensive form)** that allow us to better understand the bargaining problem in various economic settings.

## Bargaining Games



- A bargaining game is one in which two (or more) players bargain over how to divide the gains from trade.
- The gains from trade are represented by a sum of money,  $M$ , that is “on the table.”
- Players move sequentially, making alternating offers.
- Examples:
  - A Seller and a Buyer bargain over the price of a house.
  - A Labor Union and Firm bargain over wages & benefits.
  - Two countries, e.g., the U.S. and Canada bargain over the terms of a trade agreement.

## The *Disagreement Value*

- If both players in a 2-player bargaining game disagree as to how to divide the sum of money  $M$ , (and walk away from the game) then each receives their **disagreement value**.
- Let  $a$ =the disagreement value to the first player and let  $b$ =the disagreement value to the second player.
- In many cases,  $a=b=0$ , e.g., if a movie star and film company cannot come to terms, the movie star doesn't get the work and the film company doesn't get the movie star.
- The disagreement value is known by some other terms, e.g., the best alternative to negotiated agreement “BATNA.”
- By **gains from trade we mean that  $M > a+b$** .

## Take it or Leave it Bargaining Games

- “Take-it-or-leave-it” is the simplest sequential move bargaining game between two players; each player makes one move.
- Player 1 moves first and proposes a division of  $M$ .
  - For example,  $x$  for player 1 and  $M-x$  for player 2.
- Player 2 moves second and must decide whether to accept or reject Player 1’s proposal.
- If Player 2 accepts, the proposal is implemented: Player 1 gets  $x$  and Player 2 gets  $M-x$ .
- If Player 2 rejects, then both players receive their disagreement values,  $a$  for Player 1 and  $b$  for Player 2.
- This game has a simple “rollback” equilibrium:
  - Player 2 accepts if  $M-x \geq b$ , her disagreement value.
- Often we can give an even more precise solution.

## Used Car Example



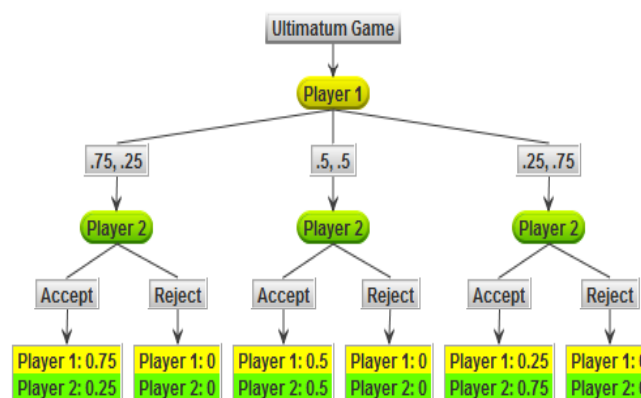
- Buyer is willing to pay a maximum price of \$8,500.
- Seller will not sell for a price less than \$8,000.
- $M = \$8,500 - \$8,000 = \$500$ ,  $a = b = 0$ .
- Suppose the seller moves first and there is *perfect information*: the seller knows the maximum value the buyer attaches to the car. Then the seller knows the buyer will reject any price  $p > \$8,500$ , and will accept any price  $p \leq \$8,500$ .
- The seller maximizes his profits by proposing  $p = \$8,500$ , or  $x = \$500$ . The buyer accepts, since  $M - x \geq b$ .
- The seller gets the entire amount,  $M = \$500$ .
- What happens if the buyer moves first?

## “Ultimatum Game”

### Discrete Version of Take it or Leave it Bargaining

- Player 1 moves first and proposes a division of \$1.00. Suppose there are just 3 possible discrete divisions, limited to \$0.25 increments:
  - Player 1 can propose  $x = \$0.25$ ,  $x = \$0.50$ , or  $x = \$0.75$  for himself, with the remainder,  $1-x$  going to Player 2.
- Player 2 moves second and can either *accept* or *reject* Player 1’s proposal.
- If Player 2 *accepts*, the proposal is implemented.
- If Player 2 *rejects*, both players get \$0 each. The \$1.00 gains from trade *vanish*.

## Computer Screen View



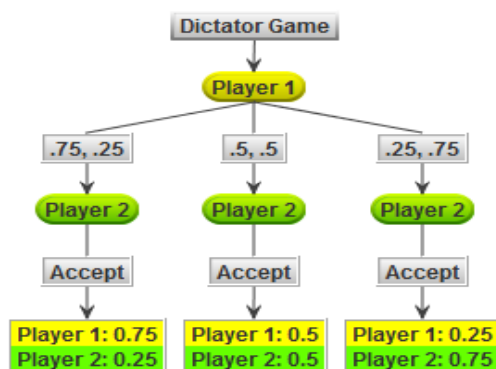
## Problems with Take-it-or-Leave-it

- Take-it-or-leave-it games are too trivial; there is no back-and-forth bargaining.
- Another problem is the credibility of take-it-or-leave-it proposals.
  - If player 2 rejects player 1's offer, *is it really believable* that both players walk away even though there are potential gains from trade?
  - Or do they continue bargaining? Recall that  $M > a+b$ .
- What about fairness? Is it really likely that Player 1 will keep as much of  $M$  as possible for himself?

## The Dictator Game



- Are Player 1's concerned about *fairness*, or are they concerned that Player 2's will reject their proposals? The Dictator Game gets at this issue.



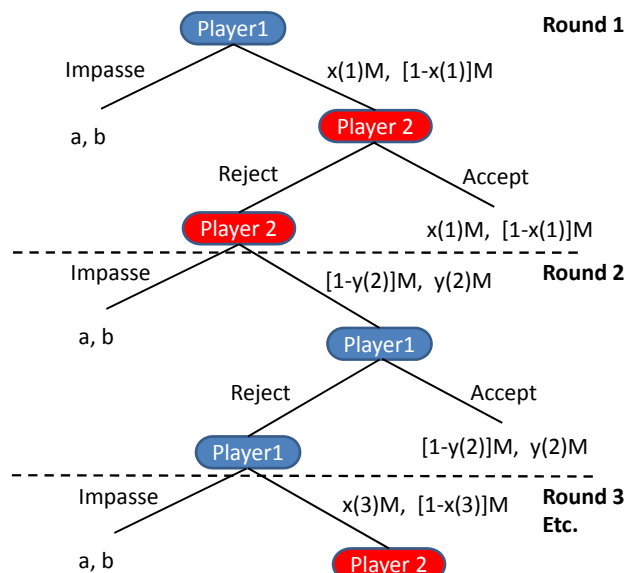
## The Alternating Offers Model of Bargaining

- A sequential move game where players have perfect information at each move.
- Players take turns making alternating offers, with one offer per round, i.e., this is real back-and-forth bargaining.
- Round numbers  $t = 1, 2, 3, \dots$
- Let  $0 \leq x(t) \leq 1$  be the fraction of  $M$  that player 1 asks for in bargaining round  $t$ , and let  $0 \leq y(t) \leq 1$  be the fraction of  $M$  that player 2 asks for in bargaining round  $t$ .

## Alternating Offer Rules

- Player 1 begins in the first round by proposing to keep  $x(1)M$  for himself and giving Player 2  $[1-x(1)]M$ .
- If Player 2 accepts, the deal is struck. If Player 2 rejects, another bargaining round may be played. In round 2, player 2 proposes to keep  $y(2)M$  for herself, and giving  $[1-y(2)]M$  to player 1.
- If Player 1 accepts, the deal is struck, otherwise, it is round 3 and Player 1 gets to make another proposal.
- Bargaining continues in this manner until a deal is struck or no agreement is reached (an *impasse* is declared by one player – a “holdout”).
- If no agreement is reached, Player 1 earns  $a$ , and Player 2 earns  $b$  (the disagreement values).

## Alternating Offers in Extensive Form



## When Does it End??



- Alternating offer bargaining games could continue indefinitely. In reality they do not.
- Why not?
  - Both sides have agreed to a deadline in advance (or  $M=0$  at a certain date).
  - The gains from trade,  $M$ , *diminish* in value over time, and may fall below  $a+b$ .
  - The players are *impatient* (time is money!).
- Take-it-or-leave-it has 1 round deadline. Let's focus on the last two possibilities.

## Decreasing Gains from Trade

- Suppose there are 2 rounds, Player 1 proposes a division of  $M$  first, player 2 accepts/rejects.
- If accepted the proposal is implemented.
- If rejected, player 2 gets to make a counter-proposal for how to split a reduced amount of money  $\alpha M$ , where  $0 < \alpha < 1$  is the rate of decay.
- Player 1 can then accept or reject this final proposal. The bargaining game is then over with certainty.
- Example: Suppose that  $M = \$12$   $\alpha = 1/3$ .

## Shrinking Gains Equilibrium

- What is the equilibrium?
- Note first that  $M = \$12$  in round 1, but only  $\$4$  in round 2.
- Working backward, player 2 knows s/he can get  $\$3.99$  in round 2 as the first mover since player 1 will strictly prefer  $\$0.01$  to nothing in round 2, and will accept that division (this follows from our assumption that players are forward looking, rational, payoff maximizers).
- Knowing this, player 1 must offer player 2  $\$4.00 > \$3.99$  in round 1 keeping  $\$12 - 4 = 8$  for himself. As player 2 recognizes this as a better payoff than can be had by waiting, she accept player 1's first round proposal.
- Player 1's first stage offer equals the amount at stake at the start of the final round,  $\alpha M$ . This generalizes to any finite  $n$ -round bargaining game, where  $2 < n < \infty$ .



### *Impatience as a Reason for Ending Bargaining:* The Period Discount Factor, $\delta$

- The period discount factor,  $0 < \delta < 1$ , provides a means of evaluating future money amounts in terms of current equivalent money amounts.
- Suppose a player values a \$1 offer now as equivalent to  $\$1(1+r)$  one period later. The discount factor in this case is:  $\delta = 1/(1+r)$ , since  $\delta \times \$1 = \$1/(1+r)$  now = \$1 later.
- If  $r$  is high,  $\delta$  is low: players discount future money amounts heavily, and are therefore very impatient.
- If  $r$  is low,  $\delta$  is high; players regard future money almost the same as current amounts of money and are more patient (less impatient).

### Example: Bargaining over a House



- Suppose the minimum price a seller will sell her house for is \$150,000, and the maximum price the buyer will pay for the house is \$160,000. Therefore,  $M = \$10,000$ .
- Suppose both players have the exact same discount factor,  $\delta = .80$ . (This implies that  $r = .25$ ).
- Suppose that there are just two rounds of bargaining. Why? The Seller has to sell by a certain date (buying another house or the Buyer has to start a new job and needs a house).
- Suppose the buyer makes a proposal in the first round, and the seller makes a proposal in the second round.
- Work backward starting in the second (last) round of bargaining and apply backward induction.

## Two period bargaining over a house, continued

- Work backward. From the perspective of today, the value of the gain from trade in the second and final round is  $gM$ . In that round, the seller has to make a counterproposal.
- In that second and final round, the seller's proposal will be to keep  $gM$  for herself and the buyer must accept or reject. Since he is indifferent between accepting and rejecting (he gets 0 in either case), let's suppose he will accept the proposal.
- Knowing this, the buyer must offer the seller  $gM$  in the first period and, since in this case the seller is made indifferent between waiting and accepting, the seller accepts immediately.
- In our example, where  $g=.8$  and  $M=\$10,000$ , the buyer offers  $.8M$  to the seller, or  $\$8,000$ , keeping  $\$2,000$  for himself. The sale price of the home is thus  $\$150,000+\$8,000=\$158,000$ .
- So, the same logic as in the decreasing gain case.

## Infinitely Repeated Analysis

- Now suppose there is no end to the number of bargaining rounds; bargaining can go on forever (an infinitely repeated game)
- If the Buyer's moves first, the amount he proposes to keep for himself,  $x(1)M$ , must leave the Seller an amount that is equivalent to that which the Seller can get in the next round, 2, by rejecting and proposing  $y(2)M$  for herself next round. The equivalent amount now, in period 1, has value to the seller of  $gy(2)M$ , where  $g$  is the period discount factor.
- Dropping the time indexes, the Buyer offers  $(1-x)M = gyM$  to the Seller, »  $x=1-gy$
- By a similar argument, the Seller must offer  $(1-y)M = gxM$  to the buyer, »  $y=1-gx$ .
- $x=1-g(1-gx)$ ,  $x=(1-g)/(1-g^2)=1/(1+\delta)$ .
- $y=1-g(1-gy)$ ,  $y=(1-g)/(1-g^2)=1/(1+\delta)$ .

### Infinitely Repeated Analysis (Continued)

- $x=y=1/(1+\delta)$ . (Note that  $x+y > 1$ !) What is  $x$  and  $y$ ?
- $x$  is the amount the Buyer gets if he makes the first proposal in the very first round; the Seller would get  $(1-x)=\delta/(1+\delta)$ , because the shares must add to 100%
- $y=1/(1+\delta)$  is the amount the Seller gets if she makes the first proposal in the very first round; Buyer would get  $(1-y)=\delta/(1+\delta)$  in this case.
- If the Buyer is the first proposer, he gets  $xM$ , and the Seller gets  $(1-x)M$ . Price is  $\$150,000+(1-x)M$ .
- If the Seller is the first proposer, she gets  $yM$  and the Buyer gets  $(1-y)M$ . Price is  $\$150,000+yM$ .
- In our example, the Buyer was the first proposer:  
 $x=1/(1+.8)=1/1.8=.556M$ . The Seller gets  $(1-x)M=(1-.556)M=.444M$ . Since  $M=10,000$ , the price of the house is  $\$154,440$ . ( $\$150,000+.444*10,000$ ).

### First Mover Advantage and Limiting Outcomes

- Summarizing, the first mover gets  $1/(1+\delta)$  and the second mover gets  $\delta/(1+\delta)$ .
- Note that so long as  $0<\delta<1$  there is a first mover advantage, as  $1/(1+\delta)>\delta/(1+\delta)$ .
- What happens to the first mover advantage as  $\delta\rightarrow 1$ ? We have  $x=1/2$ ,  $1-x=1/2$ ; the first mover advantage disappears!
- As  $\delta\rightarrow 0$ ? We get  $x=1$ ,  $1-x=0$ ; strong first mover advantage.
- The extent of the first mover advantage depends on the discount factor.

## Differing Discount Factors

- Suppose the two players have different discount factors, for example the buyer's discount factor  $\alpha_b$  is less than the seller's discount factor  $\alpha_s$ .
  - Buyer is less patient than the seller. Who gets more in this case?
  - When Buyer is the first mover, he now offers  $(1-x)M = \alpha_s yM$  to the Seller, and when Seller is the first mover she offers  $(1-y)M = \alpha_b xM$  to the Buyer.
  - $x = 1 - \alpha_s y$  and  $y = 1 - \alpha_b x$ .  $x^{new} = (1 - \alpha_s) / (1 - \alpha_s \alpha_b)$ .
  - It is easy to show that  $x^{new} < (1 - g) / (1 - g^2) = x$  when both buyer and seller had the same discount factors.
  - Example:  $\alpha_b = .5$ ,  $\alpha_s = .8$ ;  $x = (1 - \alpha_s) / (1 - \alpha_s \alpha_b) = .2 / .6 = .333$ . Recall that when  $\alpha_b = \alpha_s = g = .8$  that  $x = (1 - g) / (1 - g^2) = .2 / .36 = .556$ .

## Practical Lessons



- In reality, bargainers do not know one another's discount factors,  $g$  (or their relative levels of patience), but may try to guess these values.
- Signal that you are patient, even if you are not. For example, do not respond with counteroffers right away. Act unconcerned that time is passing- have a "poker face."
- Remember that our bargaining model indicates that the more patient player gets the higher fraction of the amount  $M$  that is on the table.