

Distributed Storage Tuesday afternoon



Outline

- Common data patterns in HTC Applications
- Storage architecture for HTC sites
- Strategies for Managing Data
- Reliability and the Cost of Complexity
- Prestaging data on the OSG
- Advanced Data Management Architectures



Common Storage Patterns

 There are many variations in how storage is handled.

 Not only variations from different workflows, but even within a workflow.

 Important: Pick a proper storage strategy for your application!



Common Storage Patterns

 I'm going to highlight a few of the storage patterns seen at HTC sites.

 It's important to note that moving files around is relatively easy. Error cases are the hard part of storage management.



Classifying Usage

- Ask yourself, per job:
 - Events in a job's life:
 - INPUT: What should be available when the job is starting?
 - RUNTIME: What is needed while the job runs?
 - OUTPUT: What output is produced?
 - Important quantities:
 - FILES: How many files?
 - DATA: How big is the "working set" of data? How big is the sum of all the files?
 - RATES: How much data is consumed on average? At peak?



Why we care about quantities?

- FILES: How many files?
 - Many file systems cannot handle lots of small files

DATA:

- Working set is how much you need on a worker node for 1 'unit' of work
- How much space do you need on the file server to store the input? Output?



Why we care about quantities?

- RATES: How much data is consumed on average? At peak?
 - Even professionals have a hard time figuring this out
 - Rates determine how you should stage your data (talk about staging shortly)



Examples of usage

- Simulation
 - Small input, big out
- Searches
 - Big input, small output
- Data processing
 - ~same input and output
- Analysis
 - Big input, small output



Simulation

- Based on different input configurations, generate the physical response of a system.
 - Input: The HTC application must manage many input files; one per job.
 - Runtime: An executable reads the input and later produces output. Sometimes, temporary scratch space is necessary for intermediate results.
 - Output: These outputs can be small [KB] (total energy and configuration of a single molecule) or large [GB] (electronic readout of a detector for hundreds of simulated particle collisions).
 - "Huge" outputs [TB] are currently not common in HTC.



Searches

- Given a database, calculate a solution to a given problem.
 - Input: A database (possibly several GB)
 shared for all jobs, and an input file per job.
 - Runtime: Job reads the configuration file at startup, and accesses the database throughout the job's runtime.
 - Output: Typically small (order of a few MB); the results of a query/search.



Data Processing

- Transform dataset(s) into new dataset(s).
 The input dataset might be re-partitioned into new logical groupings, or changed into a different data format.
 - Input: Configuration file. Input dataset
 - Runtime: Configuration is read at startup;
 input dataset is read through, one file at a time. Scratch space used for staging output.
 - Output: Output dataset; similar in size to the input dataset.



Analysis

- Given some dataset, analyze and summarize its contents.
 - Input: Configuration file and data set.
 - Runtime: Configuration file is read, then the process reads through the files in the dataset, one at a time (approximately constant rate). Intermediate output written to scratch area.
 - Output: Summary of dataset; smaller than input dataset.



OSG Storage Patterns

- Just as we want to identify common successful patterns, we want to identify common patterns that are unsuccessful on OSG.
 - Files larger than >5GB.
 - Requiring a (POSIX) shared file system.
 - Lots of small files (more than 1 file per minute of computation time).
 - Jobs consuming more than 10GB per hour, or needing scratch space more than 5GB.
 - Locking, appending, or updating files.
- When using OSG resources opportunistically, the effective limitations may be more restrictive than above!



Exercise to help identify problems

- Remember, we want to identify:
 - Input data How much data is needed for the entire workflow?
 - Working set How much data is needed to run 1 unit of work? Including possible temporary files.
 - Output data How much data is output to the workflow.

Open Science Grid

Questions?

- Questions? Comments?
 - Feel free to ask me questions: Horst Severini <severini@ou.edu>
- Upcoming sessions
 - 14:30-15:30
 - Hands-on exercises
 - -15:30 16:00
 - Break
 - -16:00 18:00
 - More!



Exercise

 https://twiki.grid.iu.edu/bin/view/ Education/AfricaGridSchoolStorageEx1



Lecture 2: Using Remote Storage Systems

Horst Severini < severini@ou.edu > Slides courtesy of Derek Weitzel < dweitzel@cse.unl.edu >



Storage Architectures on the OSG

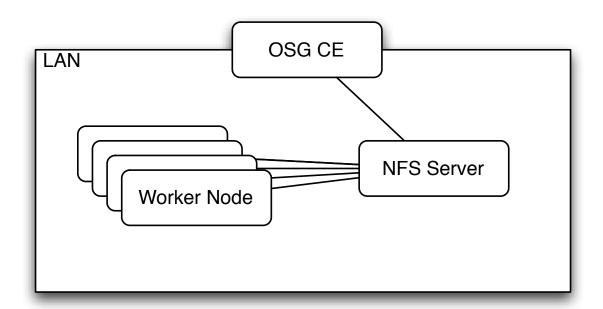
 Lets look at how storage is structured on the OSG.

 This will give you an idea of where the bottlenecks can be.

 Also can help when forming data solutions on your own.



Storage at OSG CEs



- All OSG sites have some kind of shared, POSIX-mounted storage (typically NFS).* This is almost never a distributed or high-performance file system
- This is mounted and writable on the CE.*
- This is readable (sometimes read-only) from the OSG worker nodes.

*Exceptions apply; Sites ultimately decide



Why Not?

- This setup is called the "classic SE" setup, because this is how the grid worked circa 2003.
 - Why didn't this work?
- High-performance filesystems not reliable or cheap enough.
- Scalability issues.
- Difficult to manage space.

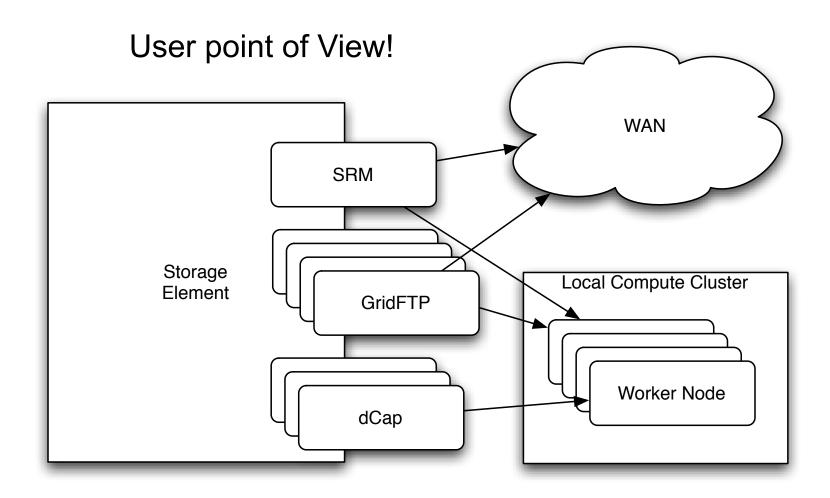


Storage Elements

- In order to make storage and transfers scalable, sites set up a separate system for storage (the Storage Element).
- Most sites have an attached SE, but there's a wide range of scalability.
- These are separated from the compute cluster; normally, you interact it via a get or put of the file.
 - Not necessarily POSIX.

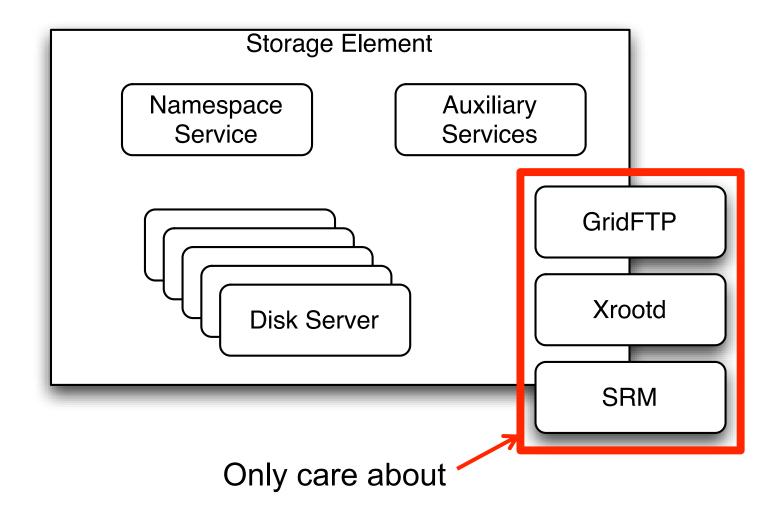


Storage Elements on the OSG





SE Internals





Complicated?



GridFTP in One Slide

- An set of extensions to the classic FTP protocol.
- Two most important extensions:
 - Security: Authentication, encryption, and integrity provided by GSI/X509. Use proxies instead of username/password.
 - Improved Scalability: Instead of transferring a file over one TCP connection, multiple TCP connections may be used.



SRM

- SRM = Storage Resource Management
- SRM is a web-services-based protocol for doing:
 - Metadata Operations.
 - Load-balancing.
 - Space management.
- This allows us to access storage remotely, and to treat multiple storage implementations in a homogeneous manner.



Data Access Methods

- To go with our archetypical problems, we'll have a few common approaches to implementing solutions:
 - Job Sandbox
 - Prestaging
 - Caching
 - Remote I/O
- A realistic solution will combine multiple methods.
- Descriptions to follow discuss the common case; there are always exceptions or alternates.



Job Sandbox

- Data sent with jobs.
- The user generates the files on the submit machines.
- These files are copied from submit node to worker node.
 - With Condor, a job won't even start if the sandbox can't be copied first. You can always assume it is found locally.



Prestaging

- Data placed into some place "near" to where the jobs are run.
- By increasing the number of locations of the data and doing intelligent placement, we increase scalability.
 - In some cases, prestaging = making a copy at each site where the job is run.
 - Not always true; a single storage element may be able to support the load of many sites.



Caching

- A cache transparently stores data close to the user process. Future requests can be served from the cache.
 - User must place file/object at a source site responsible for keeping the file permanently. Cache must be able to access the file at its source.
 - Requesting a file will bring it into the cache.
 - Once the cache becomes full, an eviction policy is invoked to decide what files to remove.
 - The cache, not the user, decides what files will be kept in the cache.



Remote I/O

- Application data is streamed from a remote site upon demand. I/O calls are transformed into network transfers.
 - Typically, only the requested bytes are moved.
 - Often, data is streamed from a common source site.
 - Can be done transparently so the user application doesn't know the I/O isn't local.
- Can be effectively combined with caching to scale better.



Scalability

- For each of the four previous methods (sandboxes, prestaging, caching, remote I/O), how well do they scale?
- Depends on:
 - What are the limitations on data size?
 - What resources are consumed per userprocess?
 - What are the scalability bottlenecks?



The Cost of Reliability

- For each of the four previous methods, what must be done to create a reliable system? Depends on:
 - What's the cost for tracking the location of files?
 - What recovery must be done on failures?
 - What must be done manually, or require extra infrastructure?
 - What is the most critical bottlenecks or points of failure?
 - How do these tie into the job submission system?



Comparing Compute and Storage

- The "shared resource" for computing is the gatekeeper.
 - One badly behaving user can overload the gatekeeper, preventing new jobs from starting.
 - However, once jobs are started, they are mostly independent.



Comparing Compute and Storage

- Storage is different:
 - It is often used throughout the job's lifetime, especially for remote I/O.
 - One badly behaved user can crash the storage resource – or at least severely degrade.
 - Opportunistic usage of storage via prestaging does not automatically have a limited lifetime.
 - Most users assume that data, once written to a SE, will be retrievable.



Prestaging Data on the OSG

- Prestaging is currently the most popular data management method on the OSG, but requires the most work.
 - Requires a transfer system to move a list of files between two sites. Example systems: Globus Online, Stork, FTS.
 - Requires a bookkeeping system to record the location of datasets. Examples: LFC, DBS.
 - Requires a mechanism to verify the current validity of file locations in the bookkeeping systems. Most systems are ad-hoc or manual.



Previous example

 In the first exercise, you formulated a data plan for BLAST.

 Now, you are going to implement the data movement of blast using prestaging.



Exercise Time!

https://twiki.grid.iu.edu/bin/view/
 Education/AfricaGridSchoolStorageEx2



Lecture 3: Caching and Remote I/O

Horst Severini < severini@ou.edu > Slides courtesy of Derek Weitzel < dweitzel@cse.unl.edu >



HTTP Caching is easy on the OSG

- Widely deployed caching servers (required by other experiments)
- Tools are easy to use
 - curl, wget
- Servers are easy to setup
 - Just need http server, somewhere



HTTP Caching

 In practice, you can cache almost any URL

 Sites are typically configured to cache any URL, but only accept requests from WN's

 Sites control the size and policy for caches, keep this in mind



HTTP Caching - Pitfalls

- curl Have to add special argument to use caching:
 - -curl -H "Pragma:" <url>
- Services like Dropbox and Google Docs explicitly disable caching
 - Dropbox: cache-control: max-age=0
 - Google Docs: Cache-Control: private, max-age=0



HTTP Caching

- The HTTP protocol is perhaps the most popular application-layer protocol on the planet.
- Built-in to HTTP 1.1 are elaborate mechanisms to use HTTP through a proxy and for caching.
- There are mature, widely-used commandline clients for the HTTP protocol on Linux.
 - Curl and wget are the most popular; we will use wget for the exercises.
 - Oddly enough, curl disables caching by default.



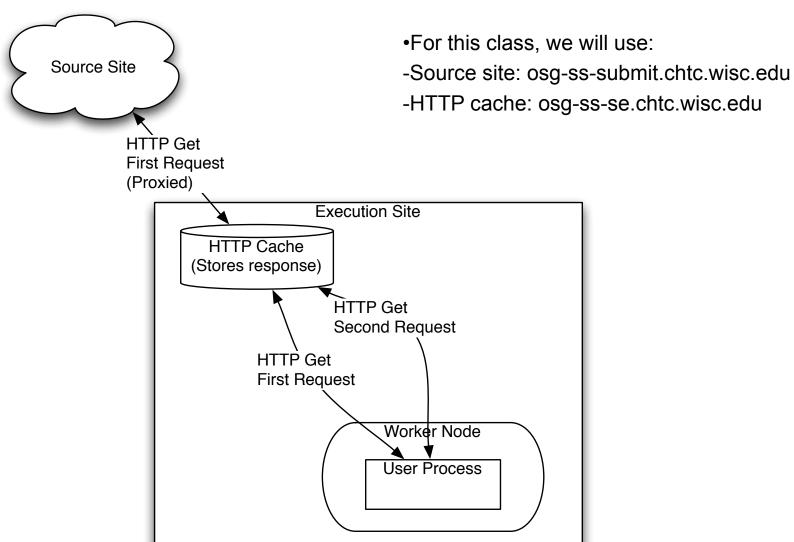
HTTP Caching: What do you need?

All you need is a public webserver...
 somewhere.

 Will pull down files to the worker node through squid cache.



HTTP Cache Dataflow





Proxy Request

 Here's the HTTP headers for the request that goes to the proxy osg-ssse.cs.wisc.edu:

```
GET http://osg-ss-submit.chtc.wisc.edu/~dweitzel/stuff.html HTTP/
1.0
User-Agent: Wget/1.11.4 Red Hat modified
Accept: */*
Host: osg-ss-submit.chtc.wisc.edu
```



Proxy Request

Response:

```
HTTP/1.0 200 OK

Date: Tue, 19 Jun 2012 22:53:26 GMT

Server: Apache/2.2.3 (Scientific Linux)

Last-Modified: Tue, 19 Jun 2012 21:26:56 GMT

ETag: "136d0f-59-4c2d9f120e800"

Accept-Ranges: bytes

Content-Length: 89

Content-Type: text/html; charset=UTF-8

X-Cache: HIT from osg-ss-se.chtc.wisc.edu

X-Cache-Lookup: HIT from osg-ss-se.chtc.wisc.edu:3128

Via: 1.0 osg-ss-se.chtc.wisc.edu:3128 (squid/2.6.STABLE21)

Proxy-Connection: close
```



HTTP Proxy

HTTP Proxy is very popular for small VO's

Easy to manage, easy to implement

 Highly recommend HTTP Proxy for files less than 100MB



Now on to Remote I/O

 Do you only read parts of the data (skip around in the file?)

 Do you have data that is only available on the submit host?

 Do you have an application that requires a specific directory structure?

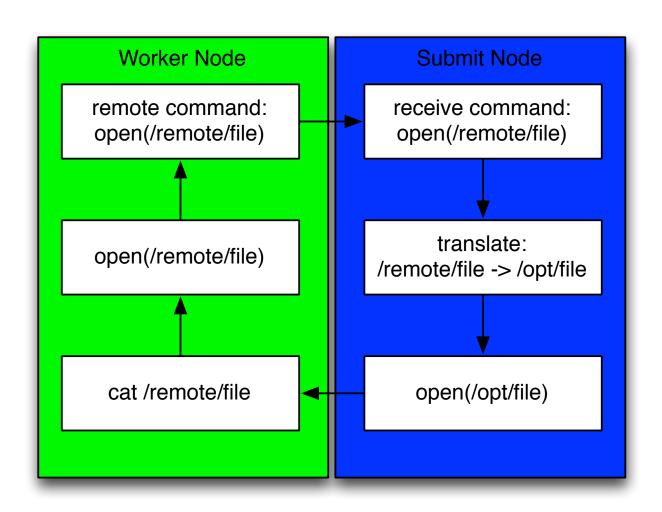




Remote I/O is for you



Remote I/O





Remote I/O



The Cooperative Computing Lab

- We will perform Remote I/O with the help of Parrot from ND.
- It redirects I/O requests back to the submit host.
- Can 'pretend' that a directory is locally accessible.



Examples of Remote I/O

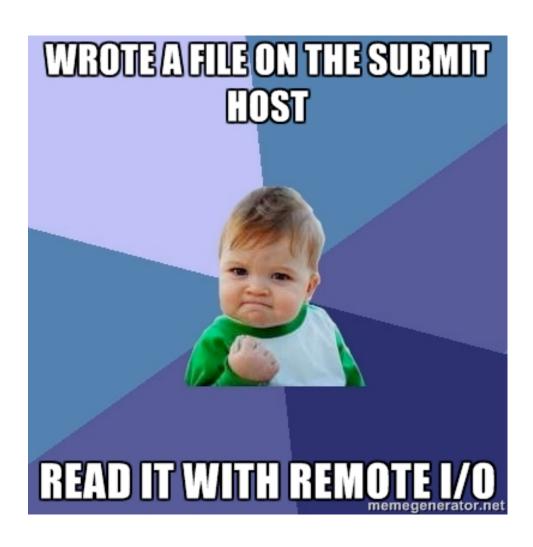
What does Remote IO look like in practice?

```
submit host:
$ echo "hello from submit host" > newfile

script on the worker node (in Nebraska?)
$ cat /chirp/CONDOR/home/dweitzel/newfile
hello from submit host
$
```



Celebrate!





Downsides of Remote I/O

 What are some downsides to Remote I/O?

All I/O goes through the submit host.

 Distributed remote I/O is hard (though can be done, see XrootD)



Now lets do it!

 https://opensciencegrid.org/bin/view/ Education/AfricaGridSchoolStorageEx3