

LiveArchive 3506

4 values whose sum is 0

Prof. Edson Alves – UnB/FGA

The SUM problem can be formulated as follows: given four lists A, B, C, D of integer values, compute how many quadruplet (a, b, c, d) belongs to $A \times B \times C \times D$ are such that $a + b + c + d = 0$. In the following, we assume that all lists have the same size n .

Input

The input begins with a single positive integer on a line by itself indicating the number of the cases following, each of them as described below. This line is followed by a blank line, and there is also a blank line between two consecutive inputs.

The first line of the input file contains the size of the lists n (this value can be as large as 4000). We then have n lines containing four integer values (with absolute value as large as 2^{28}) that belong respectively to A, B, C and D .

Output

For each test case, your program has to write the number quadruplets whose sum is zero.

The outputs of two consecutive cases will be separated by a blank line.

Exemplo de entradas e saídas

Sample Input

1

6

-45 22 42 -16

-41 -27 56 30

-36 53 -37 77

-36 30 -75 -46

26 -38 -10 62

-32 -54 -6 45

Sample Output

5

Solução com complexidade $O(N^2 \log N)$

- O produto cartesiano $A \times B \times C \times D$ tem $4000^4 = 256 \times 10^{12}$, o que inviabiliza uma solução *naive* $O(N^4)$
- É possível utilizar a técnica *meet in the middle*, observando que $a + b = -(c + d)$
- Assim, é preciso computar as somas xs e ys dos pares $A \times B$ e $C \times D$, respectivamente
- As somas registradas em ys devem ser ordenadas, de modo que seja possível utilizar a busca binária
- Para cada valor $x \in xs$, a resposta será incrementada em $L - R$, onde $[L, R)$ é o intervalo de índices de elementos y_i em ys tais que $y_i = -x$
- Este intervalo pode ser computado através da função `equal_range()` da STL da linguagem C++

Solução com complexidade $O(N^2 \log N)$

```
1 #include <bits/stdc++.h>
2
3 using namespace std;
4 using ll = long long;
5
6 const int MAX { 4010 };
7
8 ll as[MAX], bs[MAX], cs[MAX], ds[MAX], xs[MAX*MAX], ys[MAX*MAX];
9
10 ll solve(int N)
11 {
12     ll ans = 0;
13
14     for (int i = 0; i < N; ++i)
15         for (int j = 0; j < N; ++j)
16             xs[j + i*N] = as[i] + bs[j];
17
18     for (int i = 0; i < N; ++i)
19         for (int j = 0; j < N; ++j)
20             ys[j + i*N] = cs[i] + ds[j];
```

Solução com complexidade $O(N^2 \log N)$

```
22  sort(ys, ys + N*N);
23
24  for (int i = 0; i < N*N; i++)
25  {
26      auto p = equal_range(ys, ys + N*N, -xs[i]);
27      ans += (p.second - p.first);
28  }
29
30  return ans;
31 }
32
33 int main()
34 {
35     ios::sync_with_stdio(false);
36
37     int T;
38     cin >> T;
```

Solução com complexidade $O(N^2 \log N)$

```
40  for (int test = 0; test < T; ++test)
41  {
42      int N;
43      cin >> N;
44
45      for (int i = 0; i < N; ++i)
46          cin >> as[i] >> bs[i] >> cs[i] >> ds[i];
47
48      if (test)
49          cout << endl;
50
51      cout << solve(N) << endl;
52  }
53
54  return 0;
55 }
```