

# MARINA: Faster Non-Convex Distributed Learning with Compression

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KAUST

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KAUST



Federated  
Learning  
One  
World  
Seminar



IITP RAS



March 10, 2021





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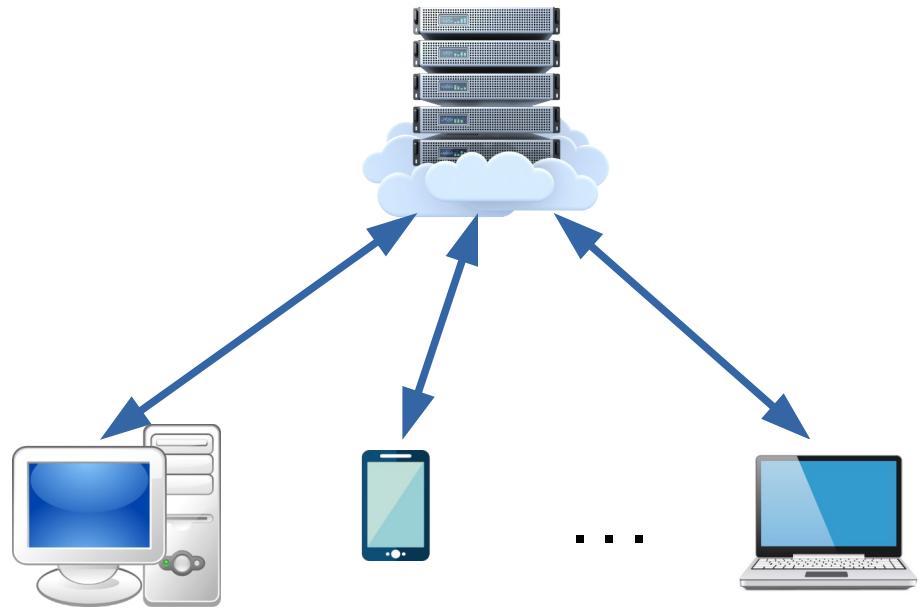


Peter Richtárik  
Professor of Computer Science  
KAUST

# Outline

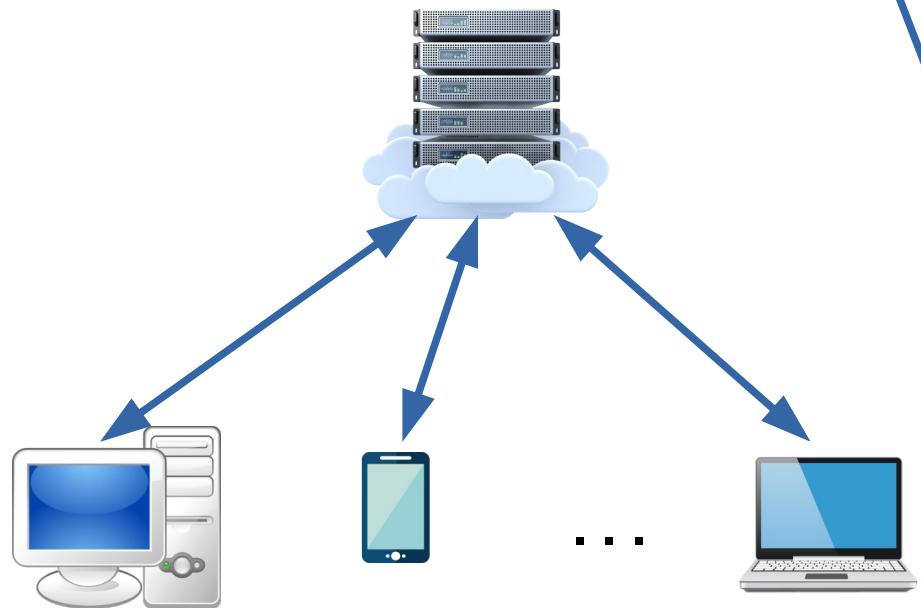
- 1 The problem
  - 2 Compressed communications
  - 3 Quantized Gradient Descent and DIANA
  - 4 MARINA
  - 5 MARINA and variance reduction
  - 6 MARINA and partial participation
  - 7 Experiments
- 
- The diagram features three vertical blue curly braces on the right side of the list. The first brace covers sections 1 through 3, labeled '10 minutes' in blue text. The second brace covers sections 4 through 6, labeled '25 minutes' in blue text. The third brace covers section 7, labeled '5 minutes' in blue text.

# 1. The Problem



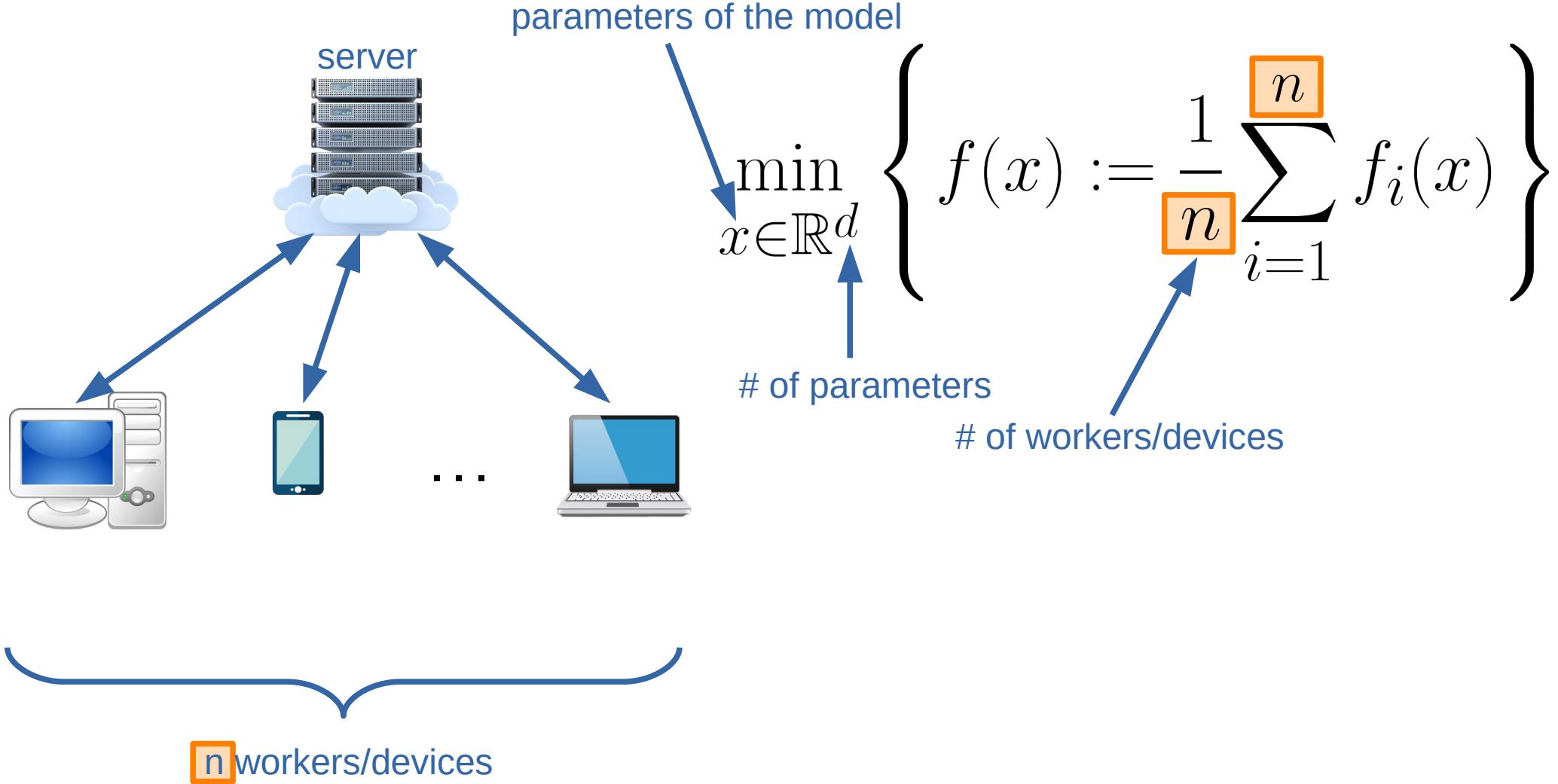
$$\min_{x \in \mathbb{R}^d} \left\{ f(x) := \frac{1}{n} \sum_{i=1}^n f_i(x) \right\}$$

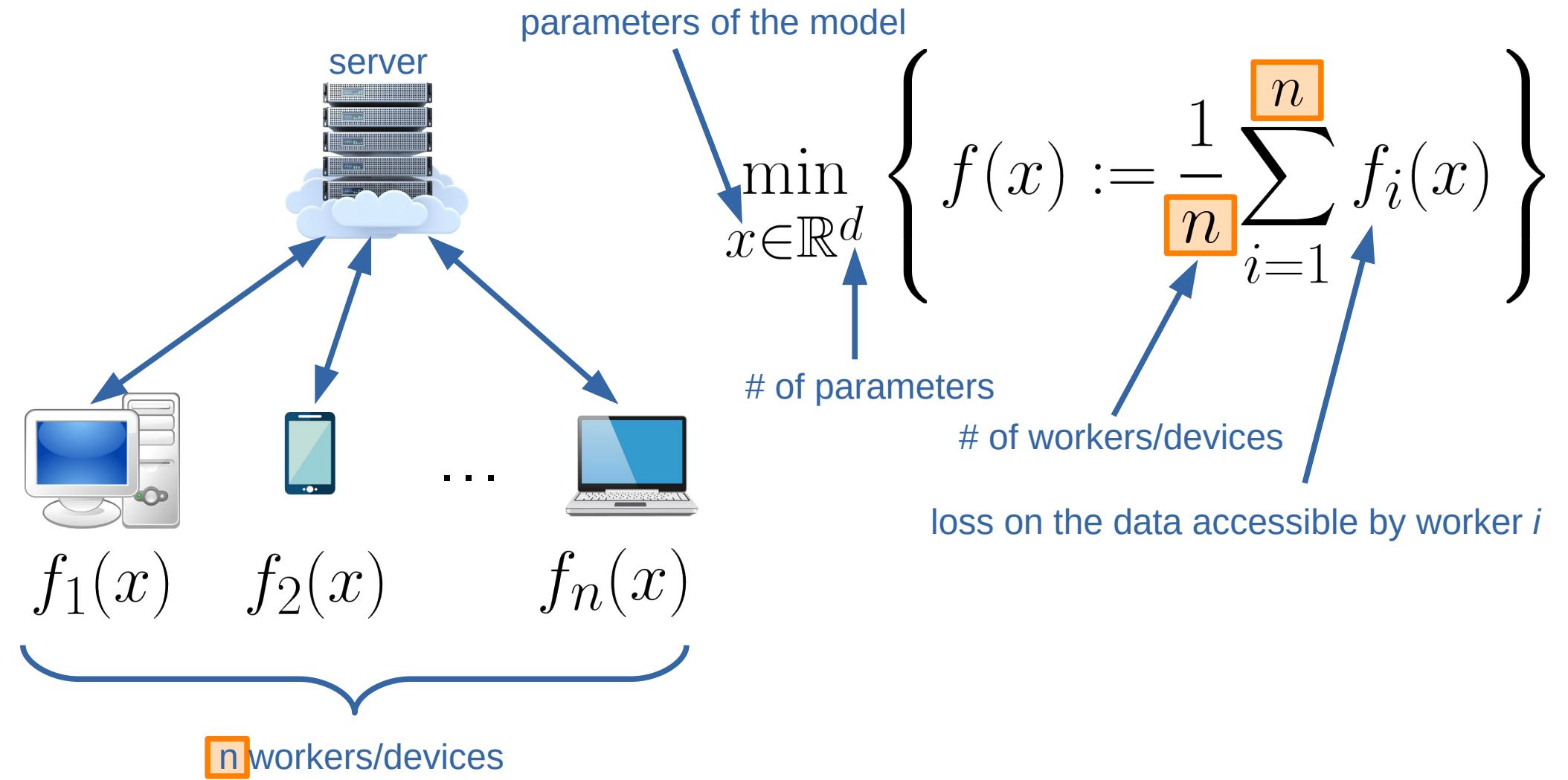
parameters of the model

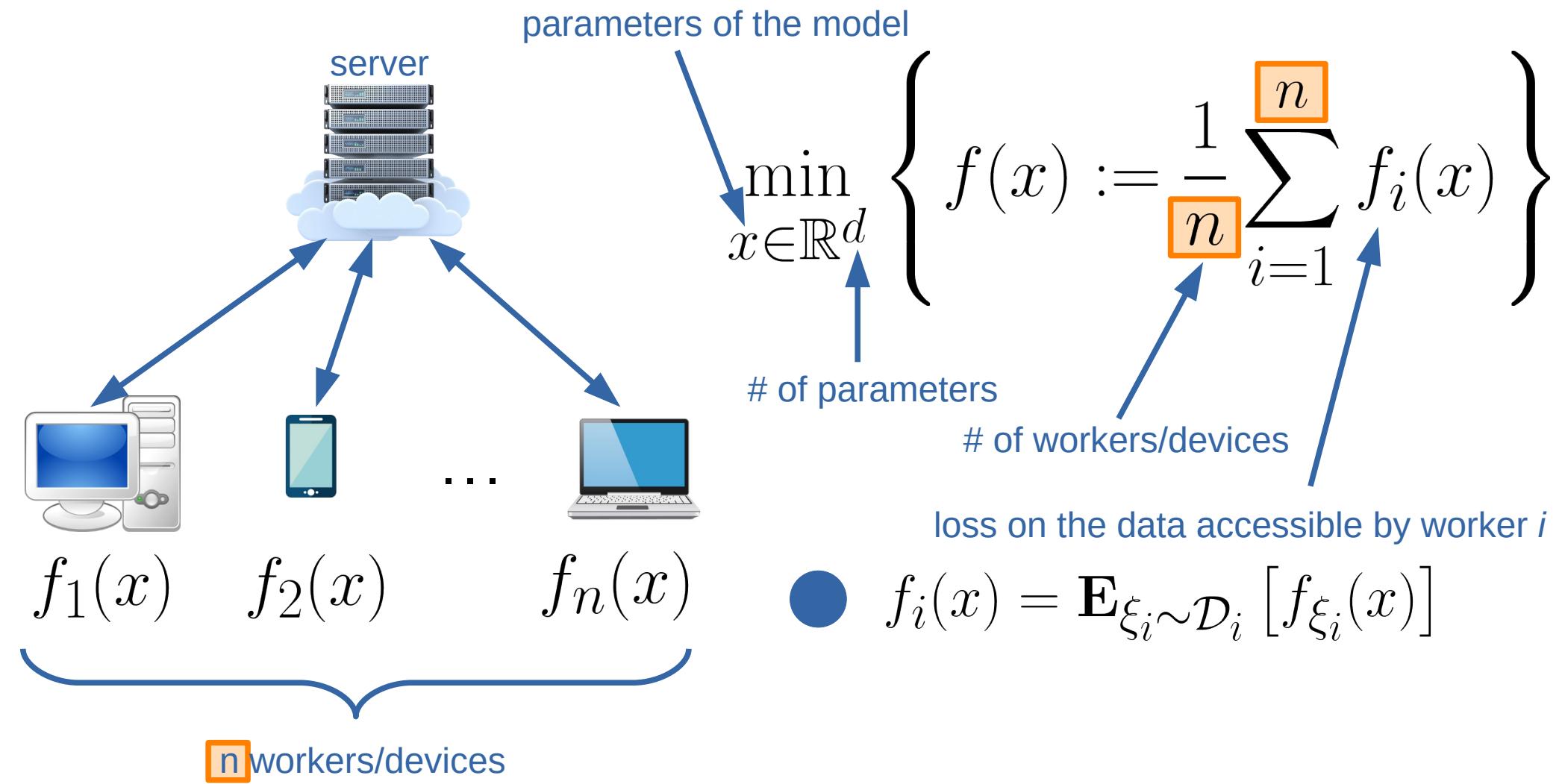


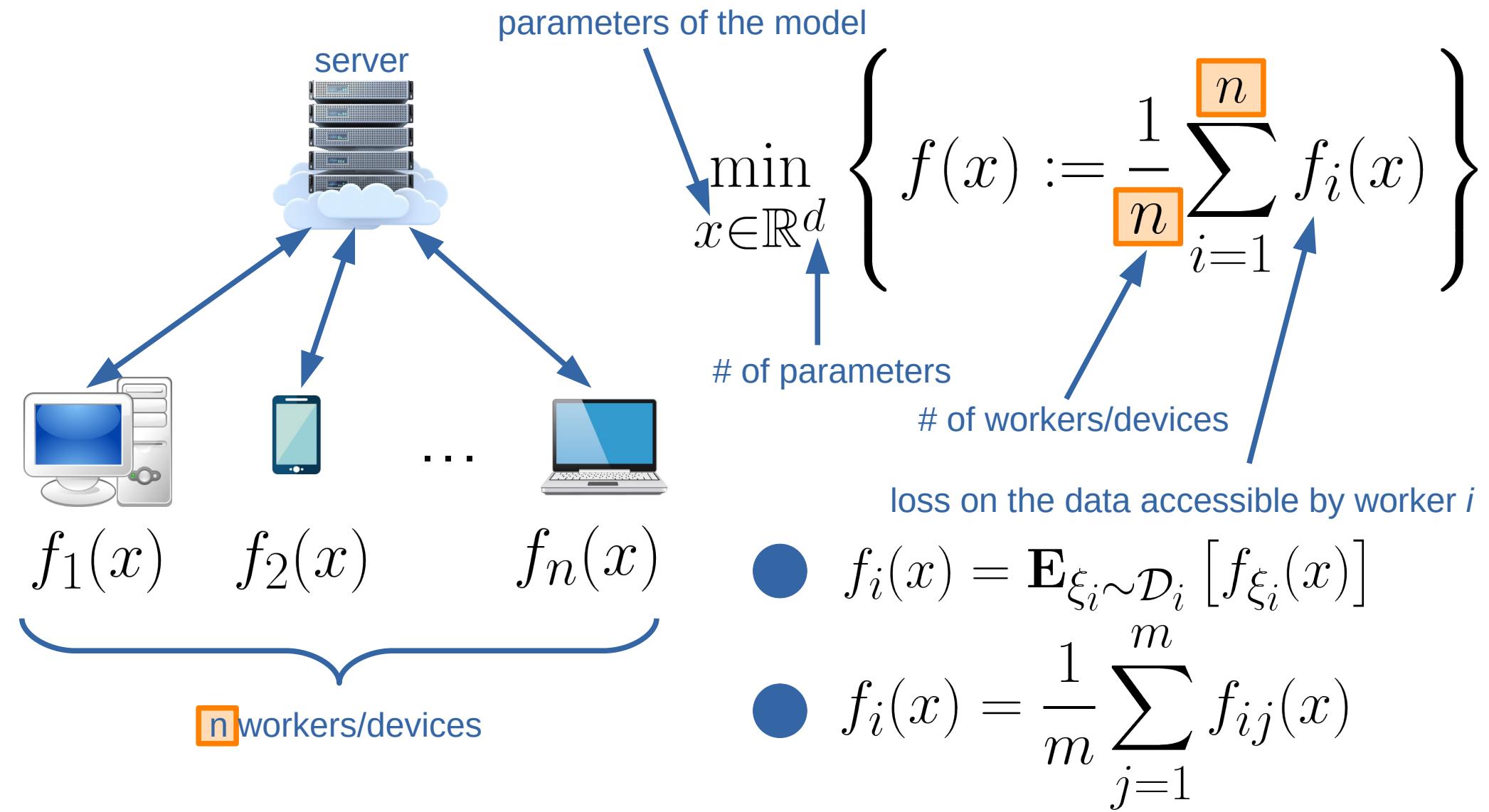
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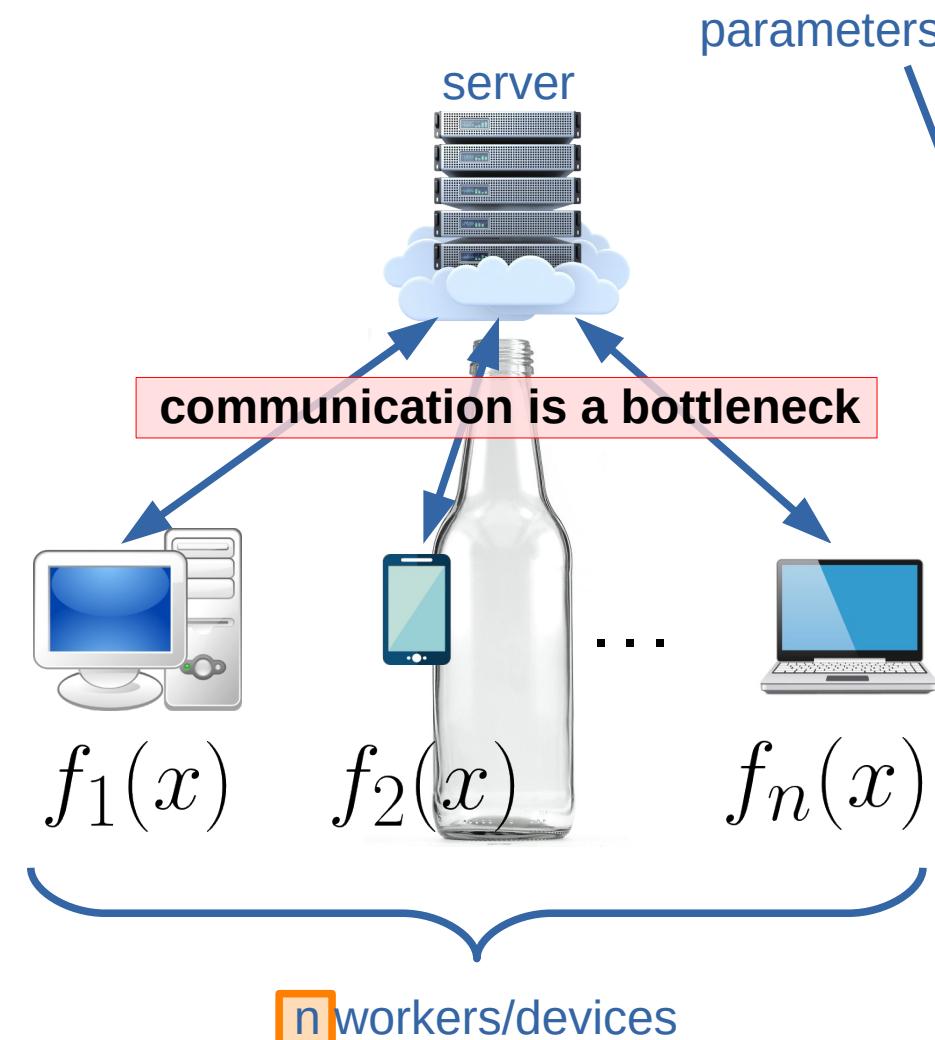
# of parameters











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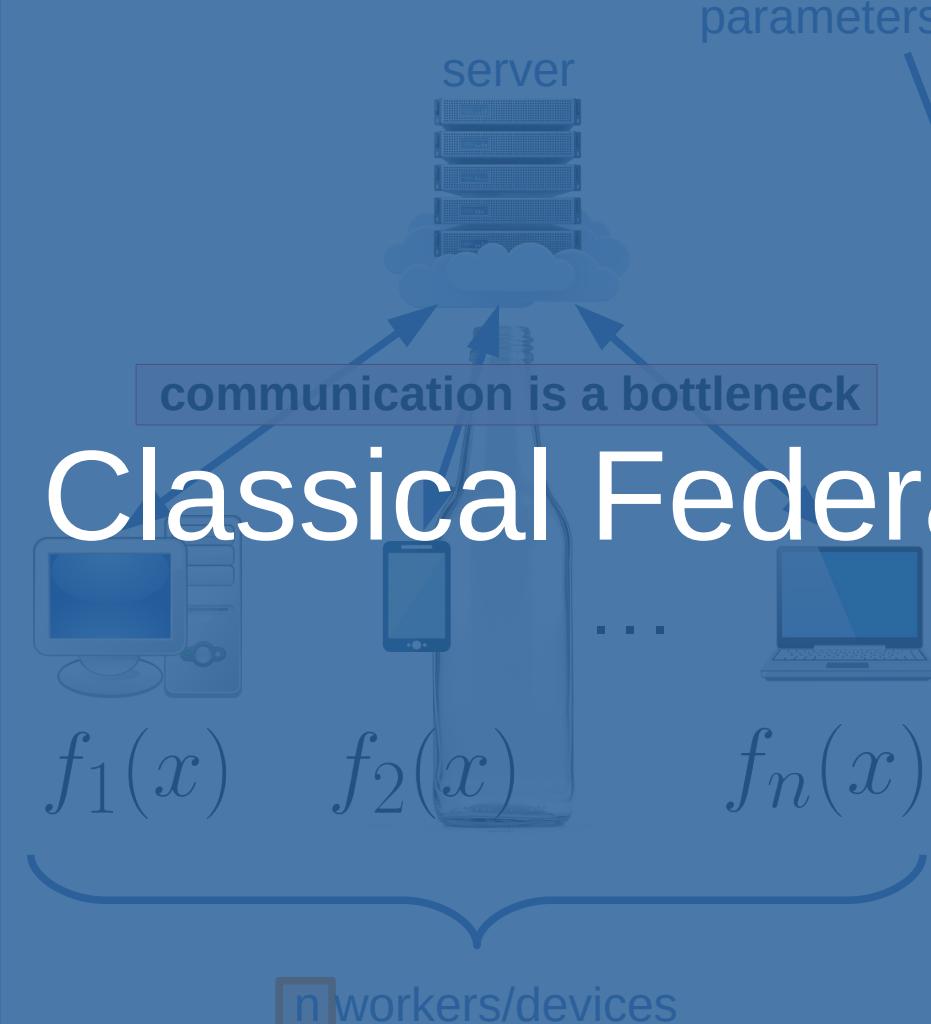
# of parameters      # of workers/devices

loss on the data accessible by worker  $i$

●  $f_i(x) = \mathbf{E}_{\xi_i \sim \mathcal{D}_i} [f_{\xi_i}(x)]$

●  $f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x)$

# Classical Federated Learning Setup



parameters of the model

$$\min_{x \in \mathbb{R}^d}$$

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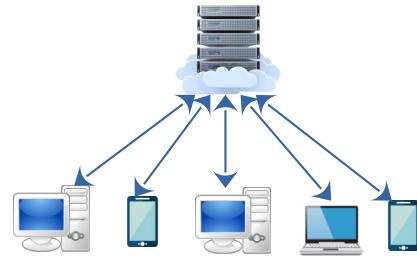
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- $f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x)$

# How to Handle Communication Bottleneck?

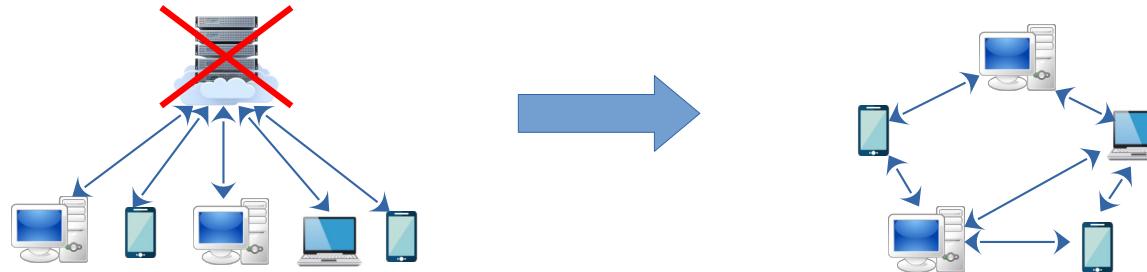
# How to Handle Communication Bottleneck?

- Change the topology of the network



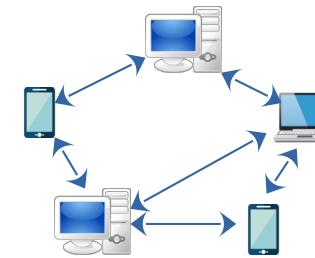
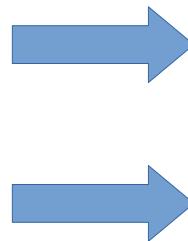
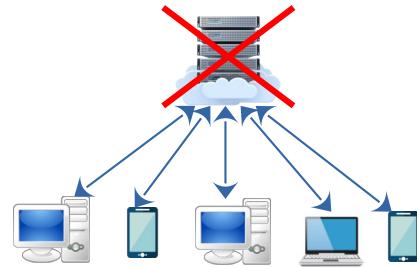
# How to Handle Communication Bottleneck?

- Change the topology of the network → Decentralized optimization



# How to Handle Communication Bottleneck?

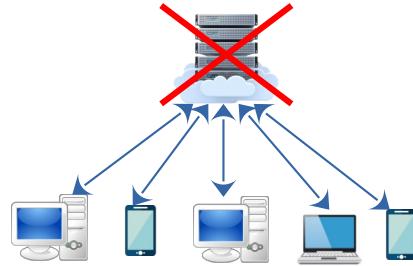
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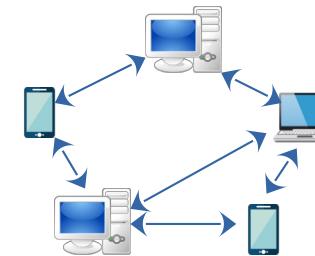
- Do more work on each worker in the hope of communicating less

# How to Handle Communication Bottleneck?

- Change the topology of the network



Decentralized optimization

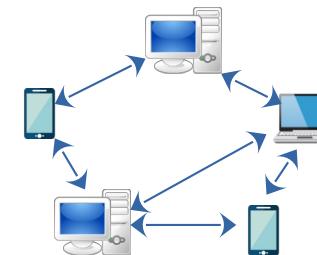
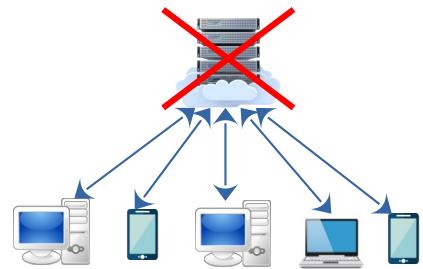


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Local-SGD/Federated Averaging

# How to Handle Communication Bottleneck?

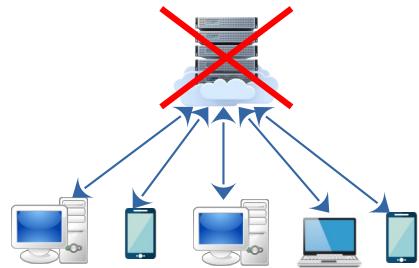
- Change the topology of the network → Decentralized optimization

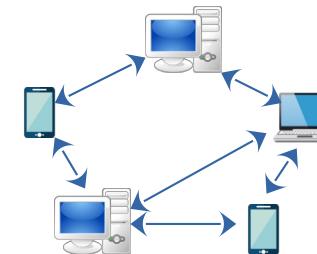


- Do more work on each worker in the hope of communicating less → Local-SGD/Federated Averaging
- Send less information to reduce the communication cost

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- Change the topology of the network → Decentralized optimization



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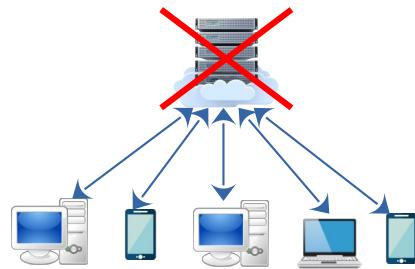
- Send less information to reduce the communication cost

Workers send dense vectors

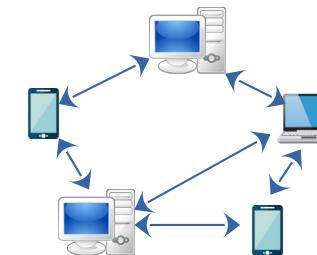
$$g = \begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix}$$

# How to Handle Communication Bottleneck?

- Change the topology of the network



Decentralized optimization



- Do more work on each worker in the hope of communicating less

Local-SGD/Federated Averaging

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$$g = \begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix}$$



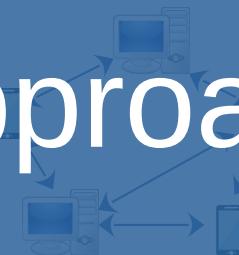
Workers send compressed/sparse vectors

$$\mathcal{Q}(g) = \frac{5}{2} \cdot \begin{pmatrix} 1 \\ 0 \\ 0 \\ -7 \\ 0 \end{pmatrix}$$

# How to Handle Communication Bottleneck?

- Change the topology of the network → Decentralized optimization

We study this approach



- Do more work on each worker in the hope of communicating less → Local-SGD/Federated Averaging

- Send less information to reduce the communication cost

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$$g = \begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix}$$

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Example: RandK (for K = 2)

$$\begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix} \longrightarrow \text{blue arrow}$$

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Pick K = 2 components uniformly at random

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Example: RandK (for K = 2)

$$\begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix} \xrightarrow{\text{for unbiasedness}} \begin{matrix} 5 \\ 2 \end{matrix} \cdot \begin{pmatrix} 1 \\ 0 \\ 0 \\ -7 \\ 0 \end{pmatrix}$$

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Example: RandK (for K = 2)

$$\begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix} \xrightarrow{\text{for unbiasedness}} 5 \cdot \begin{pmatrix} 1 \\ 0 \\ 0 \\ -7 \\ 0 \end{pmatrix}$$

$\omega = \frac{d}{K} - 1$

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Example: RandK (for K = 2)

$$d = 5 \left\{ \begin{pmatrix} 1 \\ -15 \\ 0.2 \\ -7 \\ 10 \end{pmatrix} \right. \xrightarrow{\text{for unbiasedness}} \left. \begin{matrix} 5 \\ 2 \end{matrix} \right\} \cdot \begin{pmatrix} 1 \\ 0 \\ 0 \\ -7 \\ 0 \end{pmatrix}$$

$$\omega = \frac{d}{K} - 1$$

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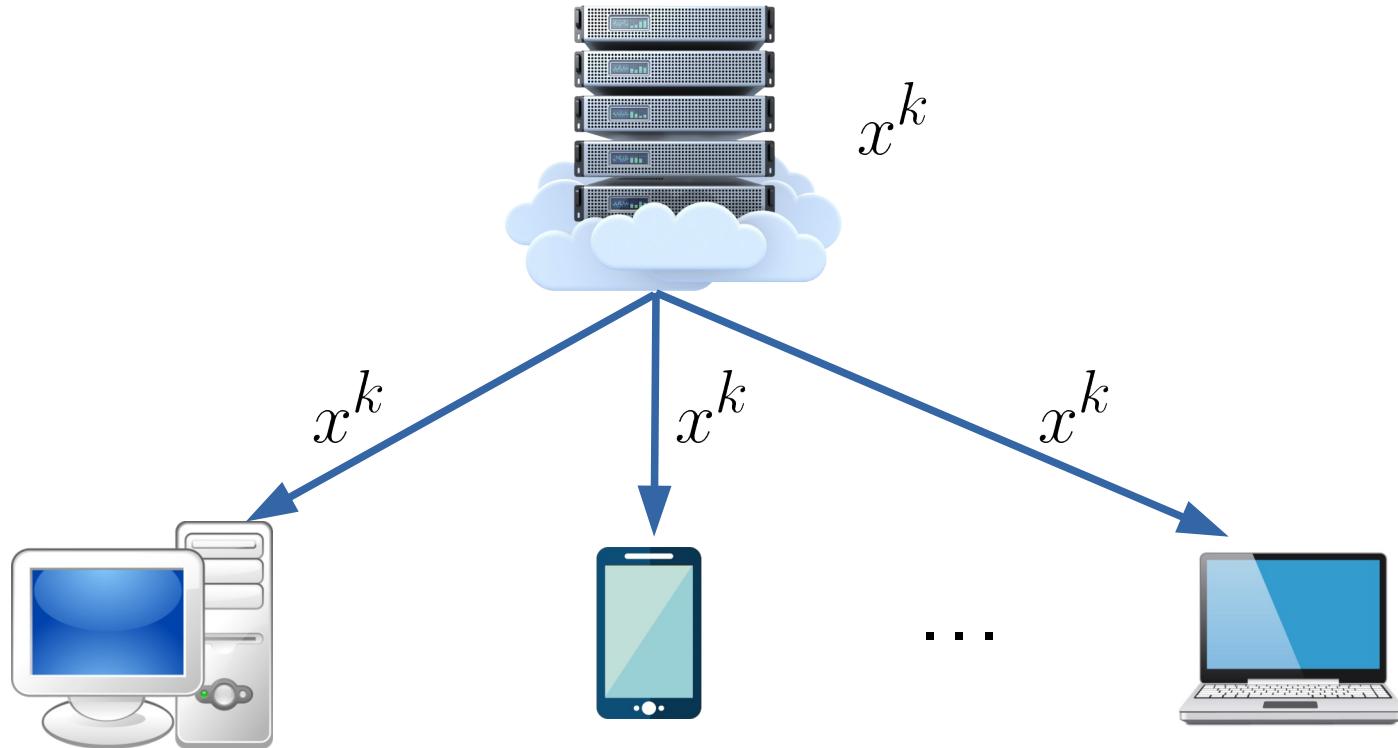
# 2. Quantized Gradient Descent (QGD)



Alistarh, Dan, Demjan Grubic, Jerry Li, Ryota Tomioka, and Milan Vojnovic. "**QSGD: Communication-efficient SGD via gradient quantization and encoding.**" *In Advances in Neural Information Processing Systems*, pp. 1709-1720. 2017.

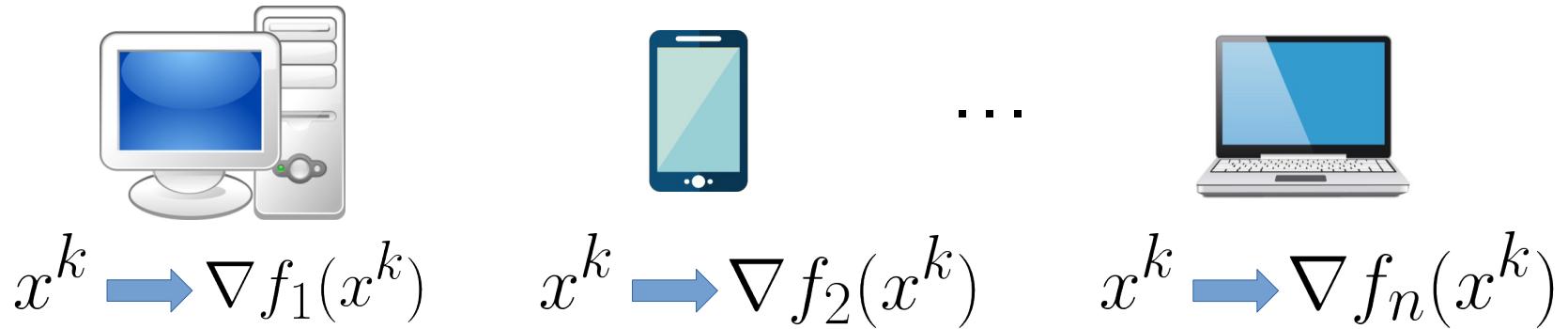
1

Server broadcasts the parameters



1 Server broadcasts the parameters

2 Devices compute the gradients



- 1 Server broadcasts the parameters
- 2 Devices compute the gradients
- 3 Devices quantize the gradients

 $x^k$ 

$$x^k \rightarrow \nabla f_1(x^k)$$



...



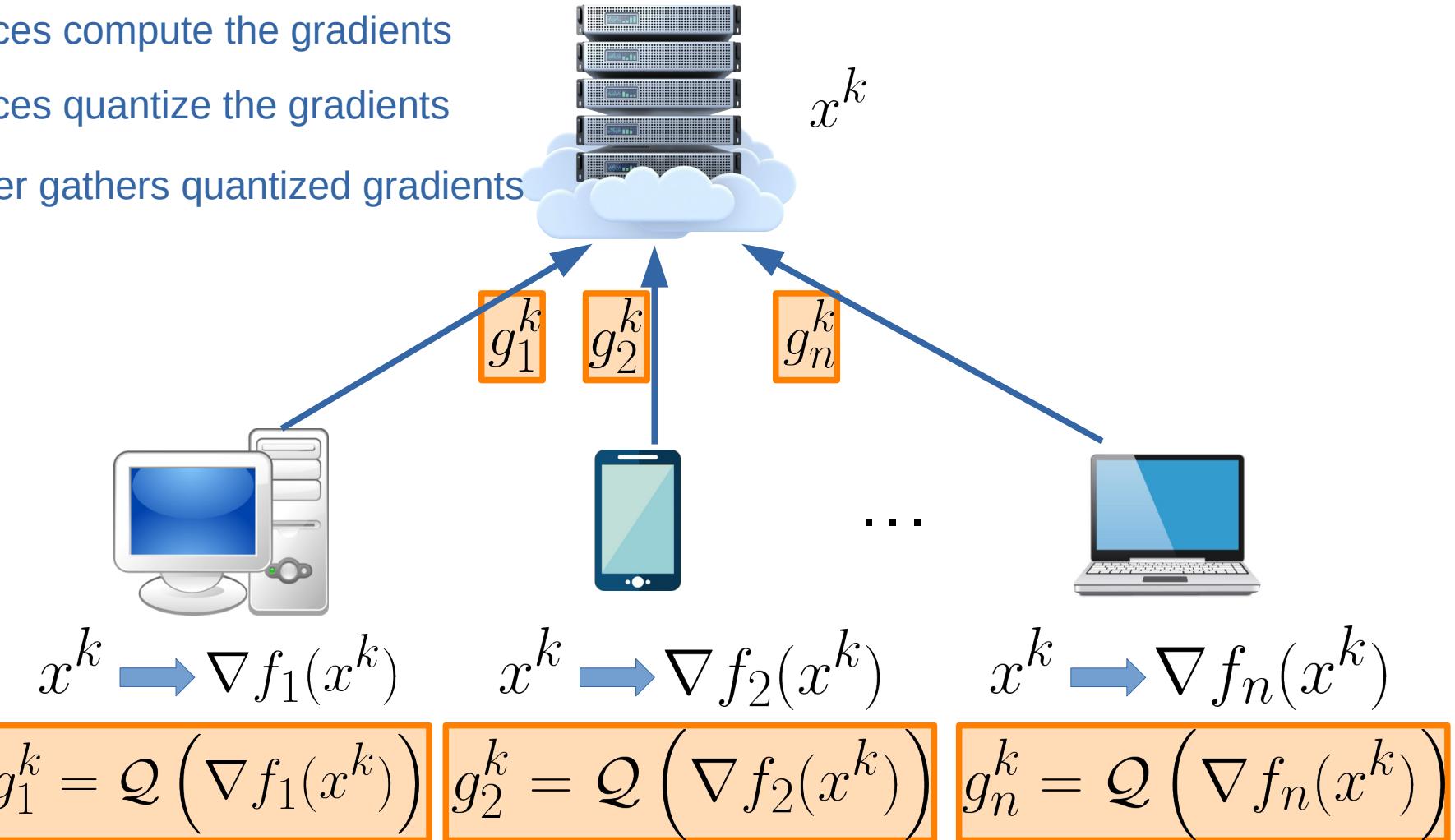
$$x^k \rightarrow \nabla f_n(x^k)$$

$$g_1^k = \mathcal{Q}(\nabla f_1(x^k))$$

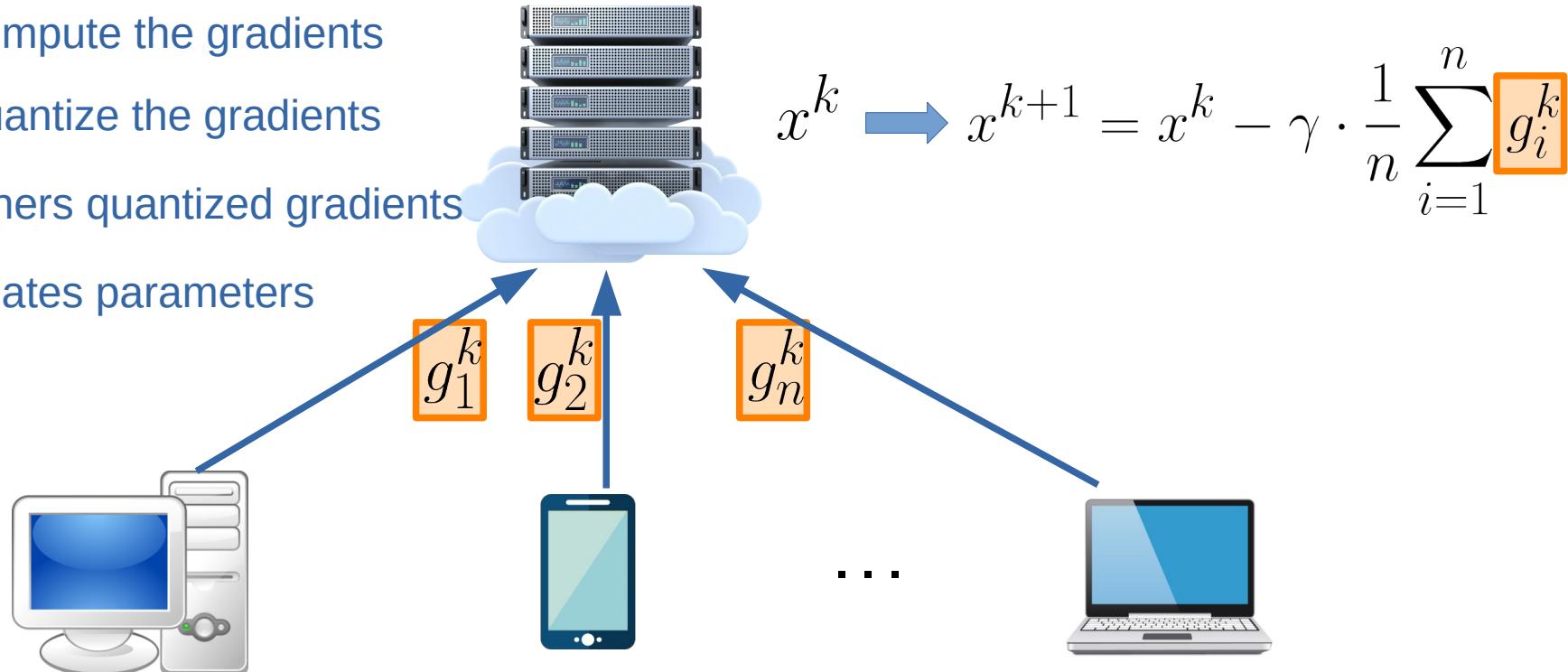
$$g_2^k = \mathcal{Q}(\nabla f_2(x^k))$$

$$g_n^k = \mathcal{Q}(\nabla f_n(x^k))$$

- 1 Server broadcasts the parameters
- 2 Devices compute the gradients
- 3 Devices quantize the gradients
- 4 Server gathers quantized gradients



- 1 Server broadcasts the parameters
- 2 Devices compute the gradients
- 3 Devices quantize the gradients
- 4 Server gathers quantized gradients
- 5 Server updates parameters



$$x^k \rightarrow \nabla f_1(x^k)$$

$$x^k \rightarrow \nabla f_2(x^k)$$

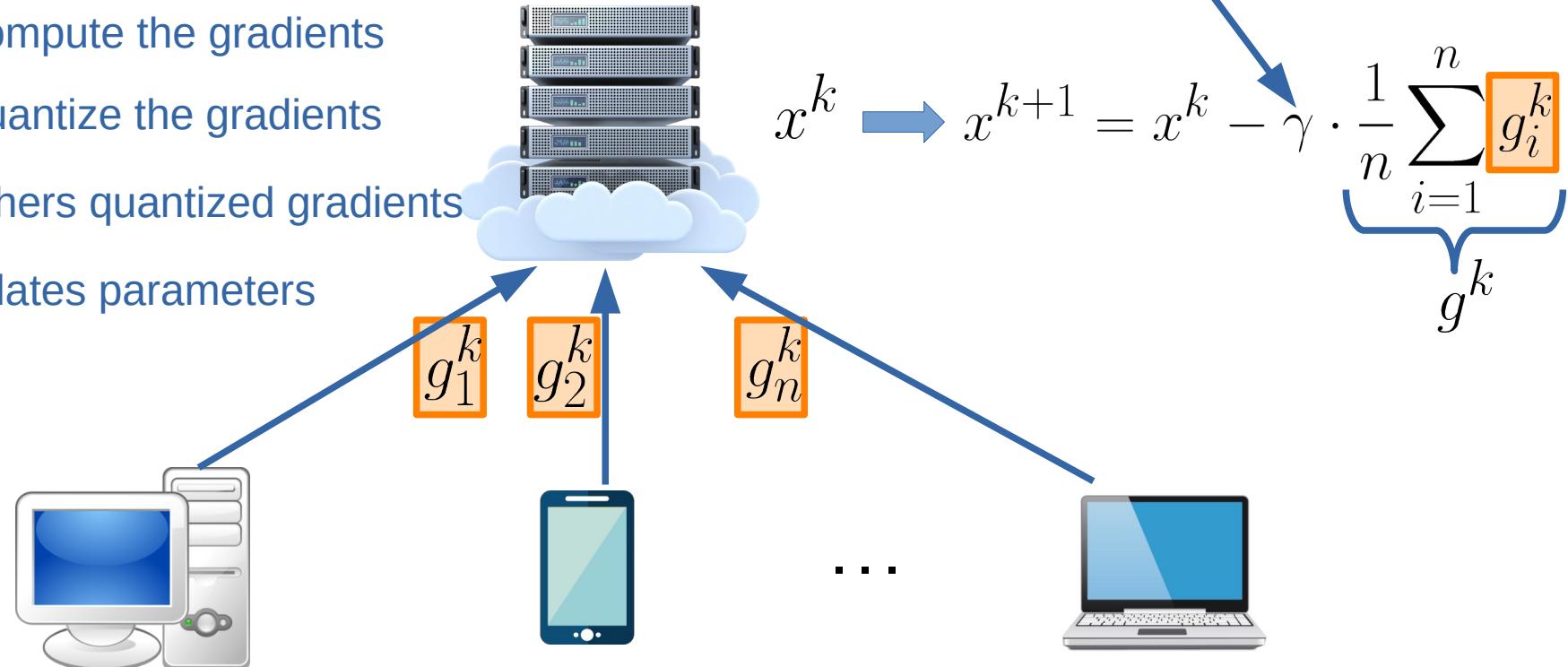
$$x^k \rightarrow \nabla f_n(x^k)$$

$$g_1^k = \mathcal{Q}(\nabla f_1(x^k))$$

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$$x^k \rightarrow \nabla f_1(x^k)$$

$$x^k \rightarrow \nabla f_2(x^k)$$

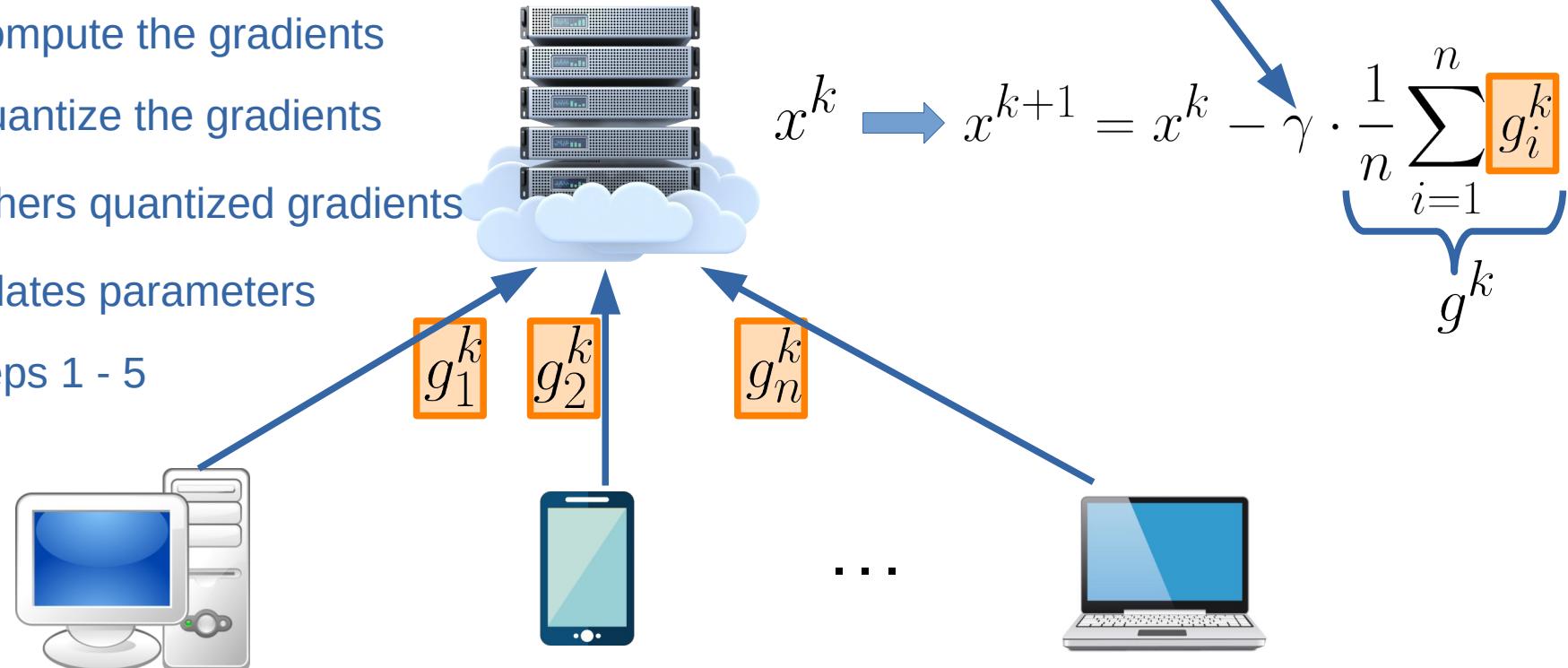
$$x^k \rightarrow \nabla f_n(x^k)$$

$$g_1^k = Q(\nabla f_1(x^k))$$

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- 1 Server broadcasts the parameters
- 2 Devices compute the gradients
- 3 Devices quantize the gradients
- 4 Server gathers quantized gradients
- 5 Server updates parameters
- 6 Repeat steps 1 - 5



$$x^k \rightarrow \nabla f_1(x^k)$$

$$x^k \rightarrow \nabla f_2(x^k)$$

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$$g_1^k = Q(\nabla f_1(x^k))$$

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# Assumptions

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1 Uniform lower bound:

$$\exists f_* \in \mathbb{R} : \forall x \in \mathbb{R}^d \quad f(x) \geq f_*$$

2 Smoothness:

$$\|\nabla f_i(x) - \nabla f_i(y)\| \leq L_i \|x - y\|$$

# Complexity Bound for QGD



Khaled, Ahmed, and Peter Richtárik. "Better theory for SGD in the nonconvex world." arXiv preprint arXiv:2002.03329 (2020).

QGD finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \varepsilon^2$  after

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$$\mathcal{O} \left( \frac{\Delta_0}{\varepsilon^2} + \frac{(1+\omega)\Delta_0^2}{\varepsilon^4 n} + \frac{(1+\omega)\Delta_0\Delta_f^*}{\varepsilon^4 n} \right)$$

communication  
rounds

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communication rounds

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$$\mathcal{O} \left( \frac{\Delta_0}{\varepsilon^2} + \frac{(1 + \omega) \Delta_0^2}{\varepsilon^4 n} + \frac{(1 + \omega) \Delta_0 \Delta_f^*}{\varepsilon^4 n} \right)$$

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$$\Delta_0 = f(x^0) - f_*$$

# Complexity Bound for QGD



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communication rounds

$$\mathbb{E} \| \mathcal{Q}(x) - x \|^2 \leq \omega \| x \|^2$$

$$\Delta_0 = f(x^0) - f_*$$

$$\Delta_f^* = f_* - \frac{1}{n} \sum_{i=1}^n f_{i,*}$$

# 3. DIANA



Mishchenko, Konstantin, Eduard Gorbunov, Martin Takáč, and Peter Richtárik.  
**"Distributed learning with compressed gradient differences."** arXiv preprint  
arXiv:1901.09269 (2019).



Horváth, Samuel, Dmitry Kovalev, Konstantin Mishchenko, Sebastian Stich, and Peter  
Richtárik. **"Stochastic distributed learning with gradient quantization and variance  
reduction."** arXiv preprint arXiv:1904.05115 (2019).

$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

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QGD:  $g_i^k = \mathcal{Q}\left(\nabla f_i(x^k)\right)$

$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

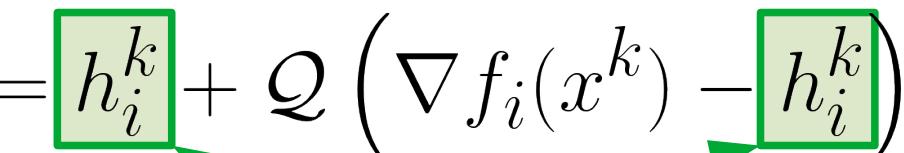
**QGD:**  $g_i^k = \mathcal{Q}\left(\nabla f_i(x^k)\right)$

**DIANA:**  $g_i^k = h_i^k + \mathcal{Q}\left(\nabla f_i(x^k) - h_i^k\right)$

$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

**QGD:**  $g_i^k = \mathcal{Q}\left(\nabla f_i(x^k)\right)$

**DIANA:**  $g_i^k = h_i^k + \mathcal{Q}\left(\nabla f_i(x^k) - h_i^k\right)$



learnable local shifts

$$h_i^{k+1} = h_i^k + \alpha \mathcal{Q}\left(\nabla f_i(x^k) - h_i^k\right)$$

$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

**QGD:**  $g_i^k = \mathcal{Q}(\nabla f_i(x^k))$  vectors that devices have to send

**DIANA:**  $g_i^k = h_i^k + \mathcal{Q}(\nabla f_i(x^k) - h_i^k)$  learnable local shifts

$$h_i^{k+1} = h_i^k + \alpha \mathcal{Q}(\nabla f_i(x^k) - h_i^k)$$

# Complexity Bound for DIANA

DIANA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \varepsilon^2$  after

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DIANA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \varepsilon^2$  after

$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + (1 + \omega) \sqrt{\omega/n} \right)}{\varepsilon^2} \right)$$

communication  
rounds

# Complexity Bound for DIANA

DIANA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \varepsilon^2$  after

Hides  
numerical  
factors and  
smoothness  
constants

$$\xrightarrow{\text{Hides numerical factors and smoothness constants}} O \left( \frac{\Delta_0 \left( 1 + (1 + \omega) \sqrt{\omega/n} \right)}{\varepsilon^2} \right)$$

communication  
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$$\Delta_0 = f(x^0) - f_*$$

# Complexity Bounds for DIANA and QGD

QGD:  $\mathcal{O} \left( \frac{\Delta_0}{\varepsilon^2} + \frac{(1+\omega)\Delta_0^2}{\varepsilon^4 n} + \frac{(1+\omega)\Delta_0\Delta_f^*}{\varepsilon^4 n} \right)$

DIANA:  $\mathcal{O} \left( \frac{\Delta_0 \left( 1 + (1+\omega)\sqrt{\omega/n} \right)}{\varepsilon^2} \right)$

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QGD:  $\mathcal{O} \left( \frac{\Delta_0}{\varepsilon^2} + \frac{(1 + \omega)\Delta_0^2}{\varepsilon^4 n} + \frac{(1 + \omega)\Delta_0\Delta_f^*}{\varepsilon^4 n} \right)$

Is it possible to get better rates?

DIANA:  $\mathcal{O} \left( \frac{\Delta_0 \left( 1 + (1 + \omega) \sqrt{\omega/n} \right)}{\varepsilon^2} \right)$

# 4. MARINA

$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

**DIANA:**  $g_i^k = h_i^k + \mathcal{Q}(\nabla f_i(x^k) - h_i^k)$

$$h_i^{k+1} = h_i^k + \alpha \mathcal{Q}(\nabla f_i(x^k) - h_i^k)$$

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typically small

$p$

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w.p.  $p$

w.p.  $1-p$

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$$x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k = x^k - \gamma g^k$$

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$$\mathbb{E}[g^k | x^k] \neq \nabla f(x^k)$$

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$$p = \frac{1}{\omega + 1} = \Theta \left( \frac{\zeta_{\mathcal{Q}}}{d} \right)$$

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$$\Delta_0 = f(x^0) - f_*$$

$$p = \frac{1}{\omega + 1} = \Theta \left( \frac{\zeta_Q}{d} \right)$$

assumption (holds for RandK, l2-quantization)

$$\zeta_Q = \sup_{x \in \mathbb{R}^d} \mathbb{E} [\|Q(x)\|_0]$$

expected density

# Complexity Bounds for MARINA and DIANA

DIANA:

$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + (1 + \omega) \sqrt{\omega/n} \right)}{\varepsilon^2} \right)$$

MARINA:

$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + \omega / \sqrt{n} \right)}{\varepsilon^2} \right)$$

# Complexity Bounds for MARINA and DIANA

DIANA:

$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + (1 + \boxed{\omega}) \sqrt{\boxed{\omega}/n} \right)}{\varepsilon^2} \right)$$

MARINA:

$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + \boxed{\omega}/\sqrt{n} \right)}{\varepsilon^2} \right)$$

# 5. MARINA and Variance Reduction

$$f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x)$$

$$f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x) \quad x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

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**VR-MARINA:**  $g_i^k = \begin{cases} \nabla f_i(x^k) & \text{w.p. } p \\ g^{k-1} + \mathcal{Q}(\nabla f_{ij}(x^k) - \nabla f_{ij}(x^{k-1})) & \text{w.p. } 1-p \end{cases}$

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$\boxed{j} \sim \{1, \dots, m\}$  uniformly at random

$$f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x) \quad x^{k+1} = x^k - \gamma \cdot \frac{1}{n} \sum_{i=1}^n g_i^k$$

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$$\mathcal{O} \left( \frac{\Delta_0 \left( 1 + \max\{\omega, \sqrt{(1+\omega)m}\} / \sqrt{n} \right)}{\varepsilon^2} \right)$$

communication  
rounds/oracle  
calls per node

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VR-MARINA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \varepsilon^2$  after

Hides numerical factors and smoothness constants

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communication rounds/oracle calls per node

$$\mathbb{E} \|Q(x) - x\|^2 \leq \omega \|x\|^2$$

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$$\mathbb{E} \|Q(x) - x\|^2 \leq \omega \|x\|^2$$

$$\Delta_0 = f(x^0) - f_*$$

$$p = \min \left\{ \frac{1}{\omega + 1}, \frac{1}{m + 1} \right\}$$

# Complexity Bounds for MARINA and VR-MARINA

$$f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x)$$

communication rounds

oracle calls per node

MARINA:

$$\mathcal{O}\left(\frac{\Delta_0 (1 + \boxed{\omega}/\sqrt{n})}{\varepsilon^2}\right) \quad \mathcal{O}\left(\frac{\boxed{m}\Delta_0(1 + \boxed{\omega}/\sqrt{n})}{\varepsilon^2}\right)$$

VR-MARINA:

$$\mathcal{O}\left(\frac{\Delta_0 \left(1 + \max\{\boxed{\omega}, \sqrt{(1 + \boxed{\omega})\boxed{m}}\} / \sqrt{n}\right)}{\varepsilon^2}\right)$$

# Complexity Bounds for VR-DIANA and VR-MARINA

$$f_i(x) = \frac{1}{m} \sum_{j=1}^m f_{ij}(x)$$

communication rounds

oracle calls per node

VR-DIANA:

$$\mathcal{O}\left(\frac{\Delta_0 \left(m^{2/3} + \omega\right) \sqrt{1 + \omega/n}}{\varepsilon^2}\right)$$

VR-MARINA:

$$\mathcal{O}\left(\frac{\Delta_0 \left(1 + \max\{\omega, \sqrt{(1 + \omega)m}/\sqrt{n}\}\right)}{\varepsilon^2}\right)$$

# 6. MARINA and Partial Participation

# PP-MARINA

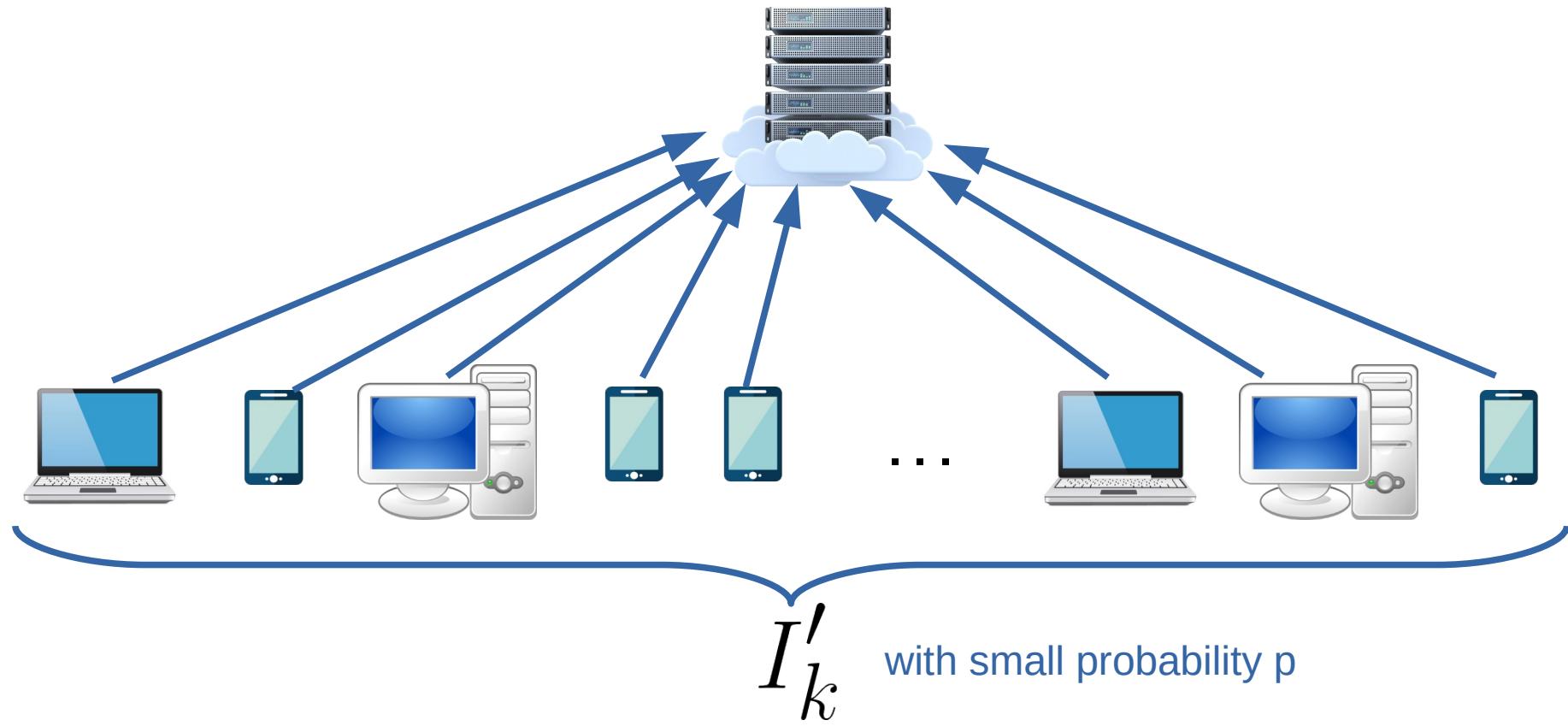


# PP-MARINA

 $I'_k$ 

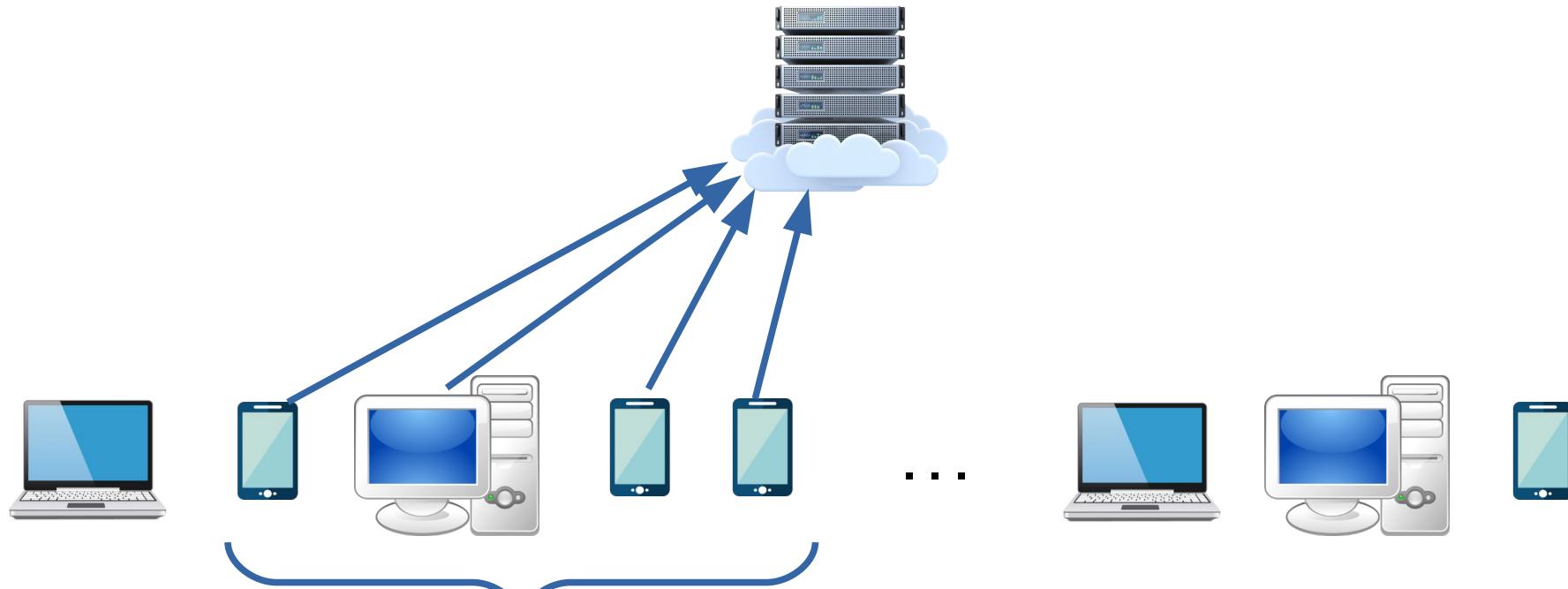
- the set of clients participating in communication on step k

# PP-MARINA



$I'_k$  - the set of clients participating in communication on step k

# PP-MARINA



$I'_k'$  - the set of clients participating in communication on step k

# PP-MARINA

$$x^{k+1} = x^k - \gamma g^k$$

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# PP-MARINA

$$x^{k+1} = x^k - \gamma g^k$$

$$c_k \sim \text{Be}(p)$$

$$g^k = \begin{cases} \nabla f\left(x^k\right) & \text { if } c_k=1 \\ g^{k-1}+\frac{1}{r} \sum_{i_k \in I'_k} \mathcal{Q}\left(\nabla f_{i_k}\left(x^k\right)-\nabla f_{i_k}\left(x^{k-1}\right)\right) & \text { if } c_k=0 \end{cases}$$

# PP-MARINA

$$x^{k+1} = x^k - \gamma g^k$$

full participation & uncompressed communication

$$c_k \sim \text{Be}(p)$$

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$$x^{k+1} = x^k - \gamma g^k$$

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if  $c_k = 1$

if  $c_k = 0$

partial participation & compressed communication

$$|I'_k| = r$$

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communication  
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PP-MARINA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \boxed{\varepsilon^2}$  after

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$$\mathbb{E} \|Q(x) - x\|^2 \leq \omega \|x\|^2$$

$$\Delta_0 = f(x^0) - f_*$$

$$p = \frac{\boxed{r}}{n(\omega + 1)} = \Theta \left( \frac{\boxed{r} \zeta_Q}{nd} \right)$$

# Complexity Bound for PP-MARINA

PP-MARINA finds such  $\hat{x}$  that  $\mathbb{E} \left[ \|\nabla f(\hat{x})\|^2 \right] \leq \boxed{\varepsilon^2}$  after

Hides numerical factors and smoothness constants  $\rightarrow \mathcal{O} \left( \frac{\Delta_0 (1 + (1 + \boxed{\omega}) \boxed{\sqrt{n/r}})}{\varepsilon^2} \right)$  communication rounds

$$\mathcal{O} \left( \frac{\Delta_0 (1 + (1 + \boxed{\omega}) \boxed{\sqrt{n/r}})}{\varepsilon^2} \right)$$

$$\mathbb{E} \|Q(x) - x\|^2 \leq \omega \|x\|^2$$

$$\Delta_0 = f(x^0) - f_*$$

$$p = \frac{r}{n(\omega + 1)} = \Theta \left( \frac{r \zeta_Q}{nd} \right)$$

$\uparrow$  assumption (holds for RandK, l2-quantization)

$\uparrow$  expected density

$$\zeta_Q = \sup_{x \in \mathbb{R}^d} \mathbb{E} [\|Q(x)\|_0]$$

# Complexity Bounds for PP-MARINA and MARINA

MARINA:  $\mathcal{O} \left( \frac{\Delta_0 (1 + \omega / \sqrt{n})}{\varepsilon^2} \right)$

PP-MARINA:  $\mathcal{O} \left( \frac{\Delta_0 (1 + (1 + \omega) \sqrt{n/r})}{\varepsilon^2} \right)$

# 7. Experiments

# The Problem

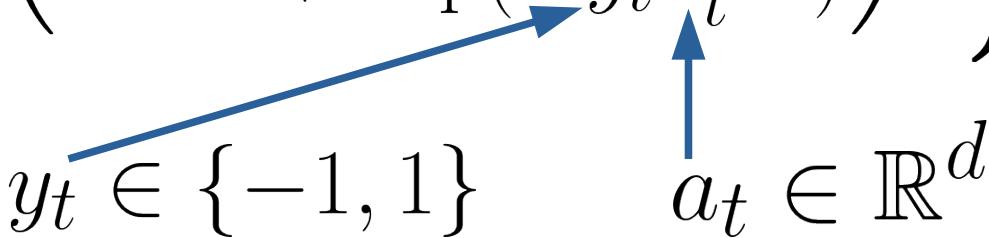
$$\min_{x \in \mathbb{R}^d} \left\{ f(x) = \frac{1}{N} \sum_{t=1}^N \left( 1 - \frac{1}{1 + \exp(-y_t a_t^\top x)} \right)^2 \right\}$$

# The Problem

$$\min_{x \in \mathbb{R}^d} \left\{ f(x) = \frac{1}{N} \sum_{t=1}^N \left( 1 - \frac{1}{1 + \exp(-y_t a_t^\top x)} \right)^2 \right\}$$

A diagram illustrating the components of the loss function. Two blue arrows point upwards from the labels  $y_t \in \{-1, 1\}$  and  $a_t \in \mathbb{R}^d$  to the terms  $y_t$  and  $a_t^\top x$  in the equation respectively.

# The Problem

$$\min_{x \in \mathbb{R}^d} \left\{ f(x) = \frac{1}{N} \sum_{t=1}^N \left( 1 - \frac{1}{1 + \exp(-y_t a_t^\top x)} \right)^2 \right\}$$


- The dataset was split into 5 equal parts among 5 clients

# The Problem

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$y_t \in \{-1, 1\}$        $a_t \in \mathbb{R}^d$

- The dataset was split into 5 equal parts among 5 clients
- Theoretical stepsizes

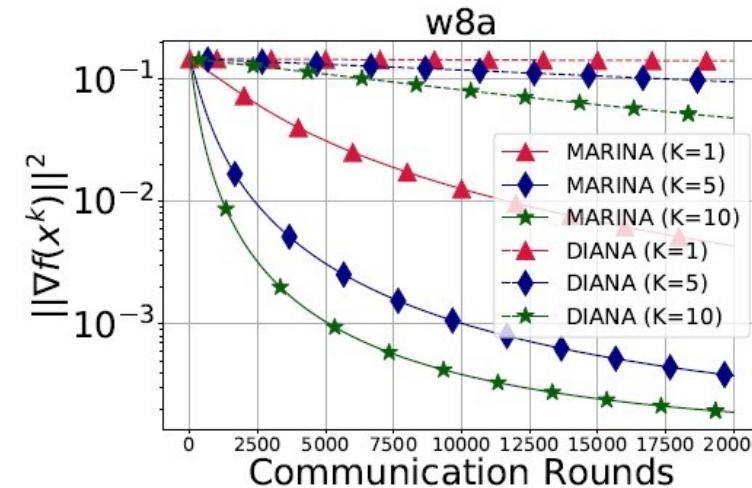
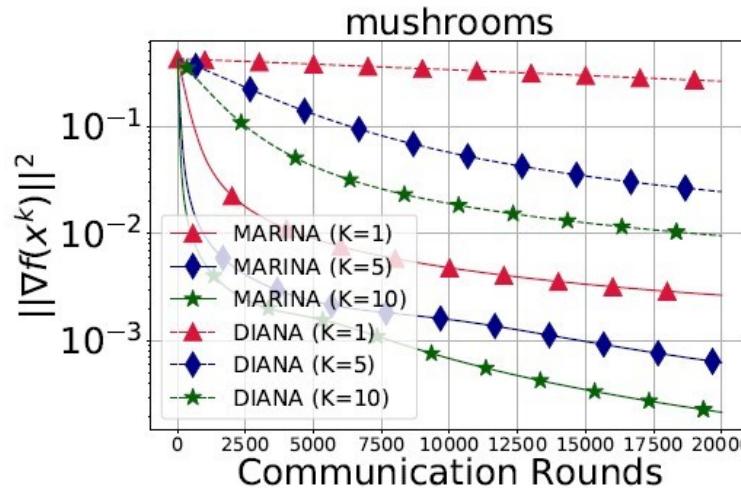
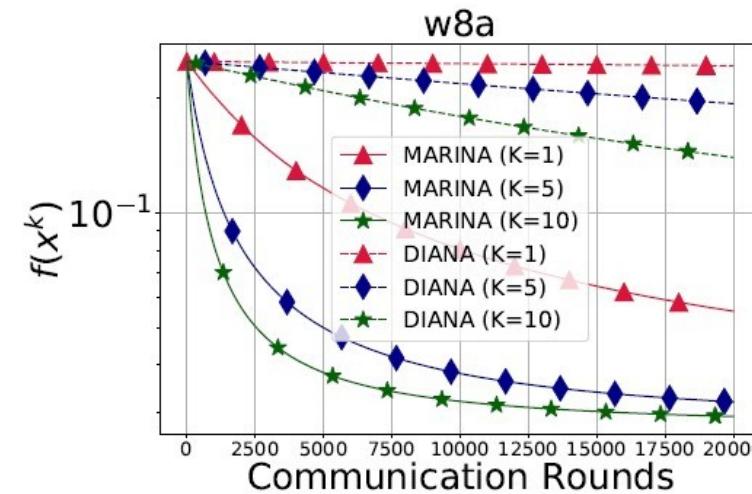
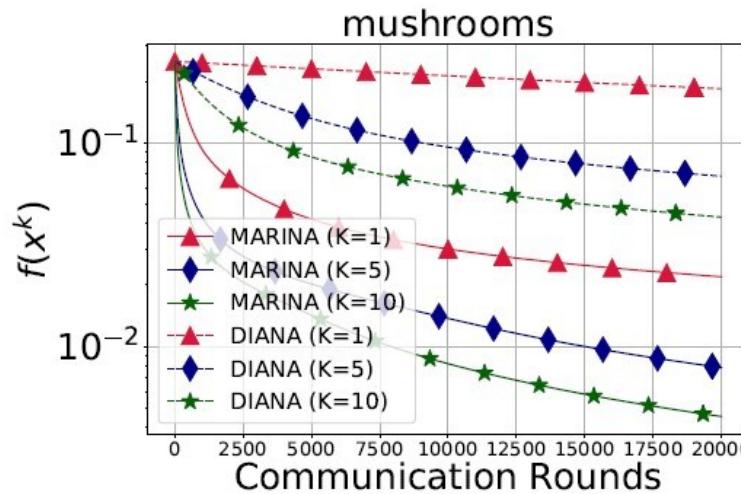
# The Problem

$$\min_{x \in \mathbb{R}^d} \left\{ f(x) = \frac{1}{N} \sum_{t=1}^N \left( 1 - \frac{1}{1 + \exp(-y_t a_t^\top x)} \right)^2 \right\}$$

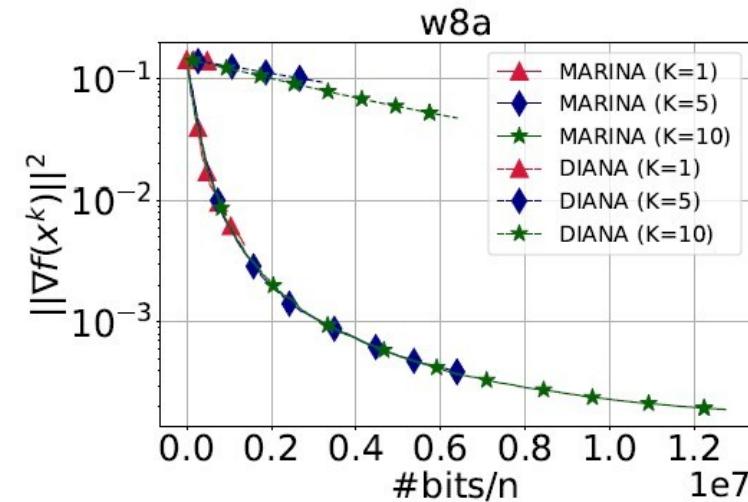
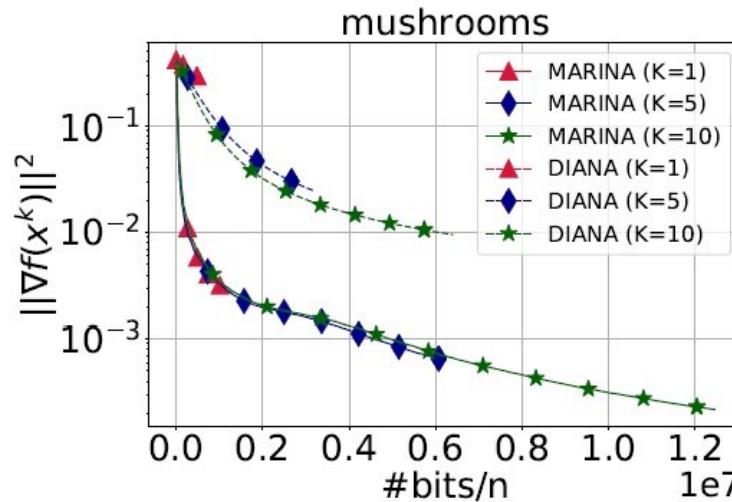
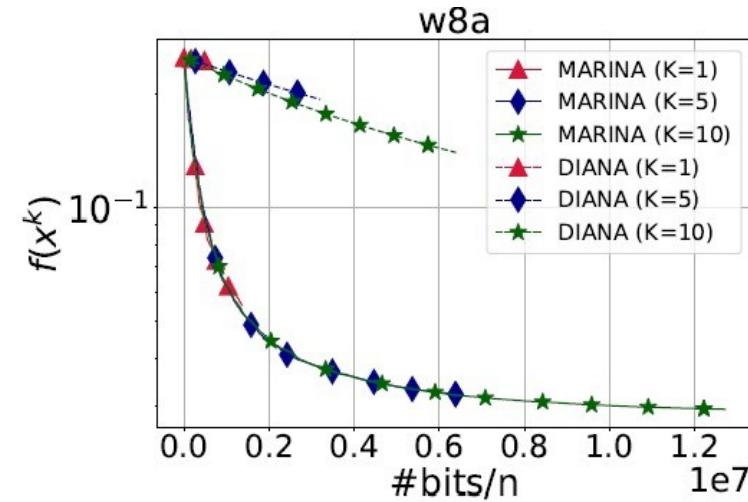
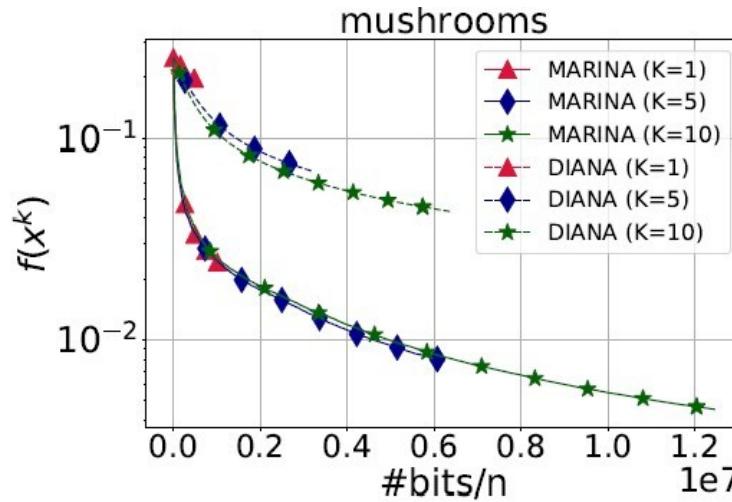
$y_t \in \{-1, 1\}$        $a_t \in \mathbb{R}^d$

- The dataset was split into 5 equal parts among 5 clients
- Theoretical stepsizes
- For stochastic methods batchsize was 1% of local data

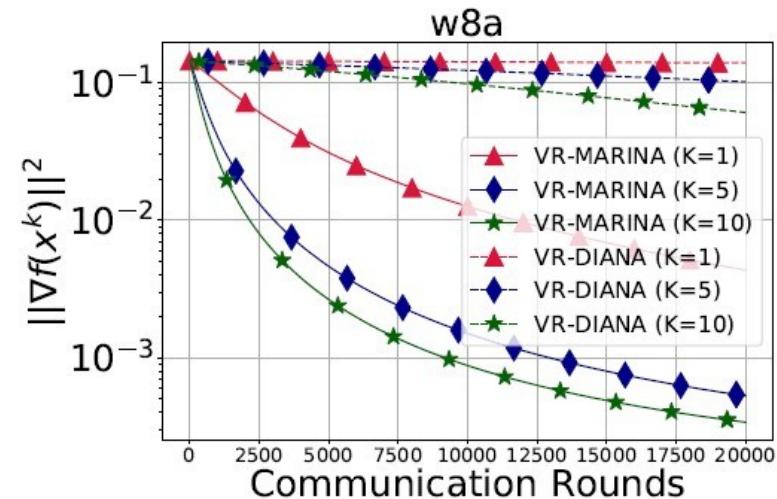
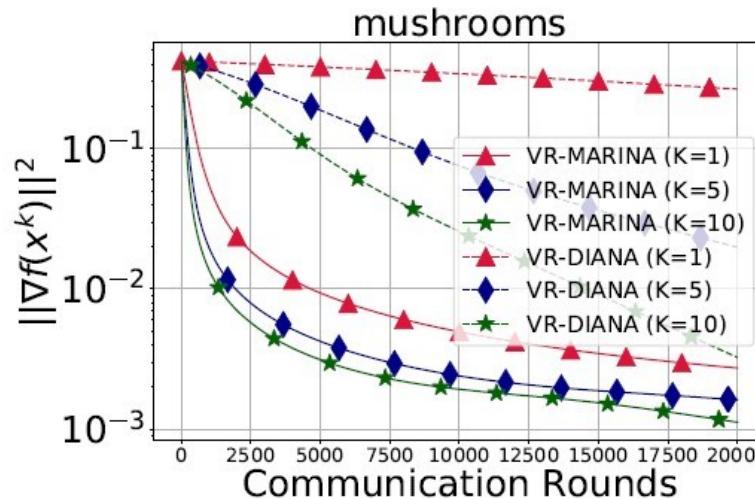
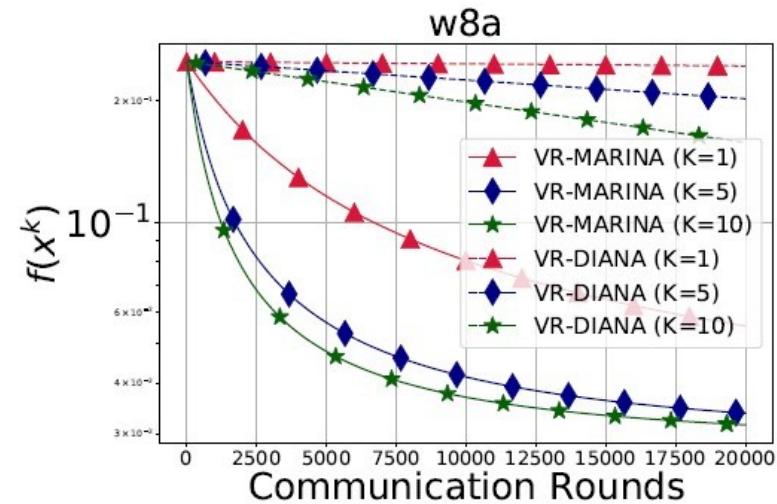
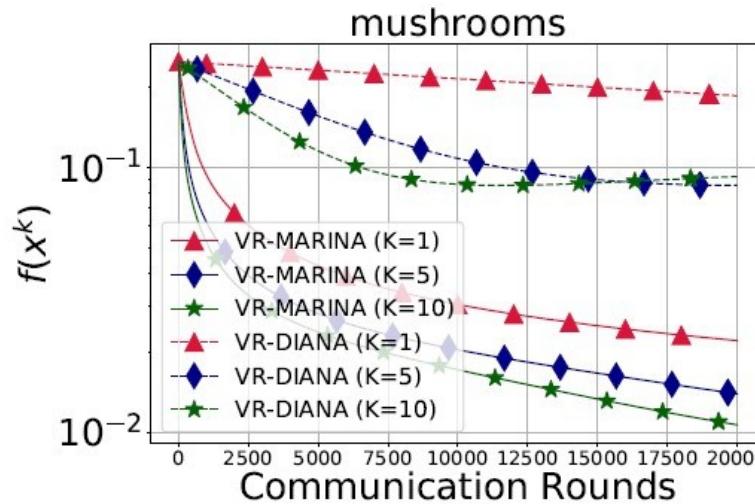
# MARINA vs DIANA (RandK)



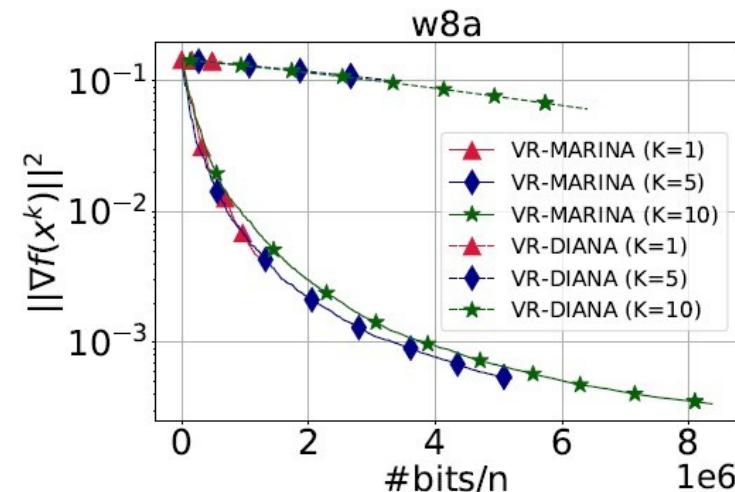
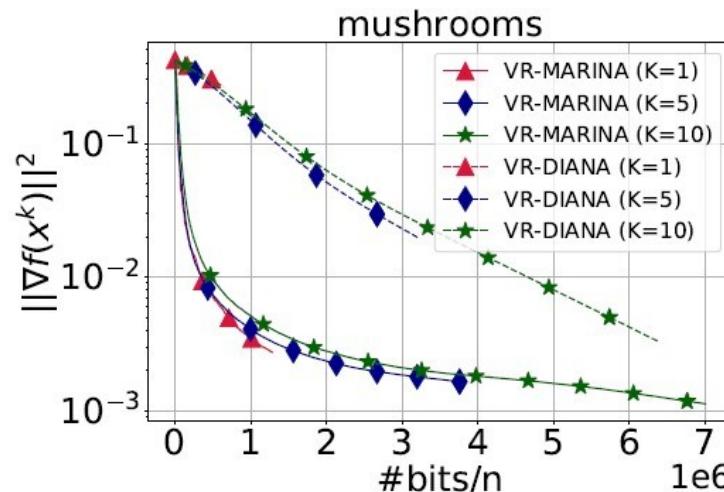
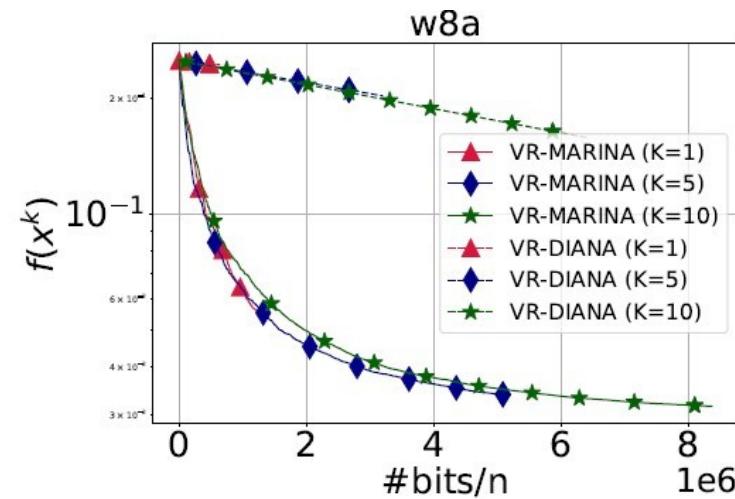
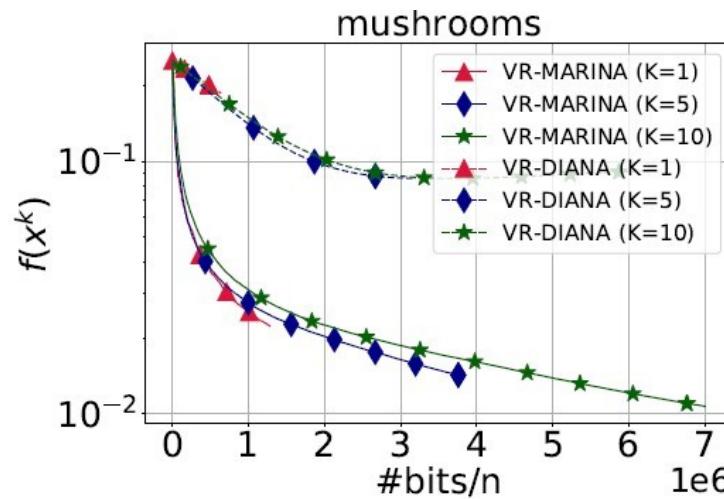
# MARINA vs DIANA (RandK)



# VR-MARINA vs VR-DIANA (RandK)



# VR-MARINA vs VR-DIANA (RandK)



# 8. Extra Results

In the paper, we also have:

- Mini-batched version of Variance Reduced MARINA

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- Variance Reduced MARINA for the problems with  $f_i(x) = \mathbf{E}_{\xi_i \sim \mathcal{D}_i} [f_{\xi_i}(x)]$

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- Mini-batched version of Variance Reduced MARINA
- Variance Reduced MARINA for the problems with  $f_i(x) = \mathbf{E}_{\xi_i \sim \mathcal{D}_i} [f_{\xi_i}(x)]$
- Rates under Polyak- Lojasiewicz Condition

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- Rates under Polyak- Lojasiewicz Condition
- Explicit dependencies on smoothness constants, non-uniform sampling
- Simple proofs