

# Cálculo de Programas

## Resolução - Ficha 11

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### Exercício 1

$$\begin{aligned} & \text{mirror} = \langle \text{in}_2 \cdot \alpha \rangle \\ \equiv & \quad \{ \text{def. in}_2, \text{def. } \alpha \} \\ & \text{mirror} = \langle [\text{Leaf}, \text{Fork}] \cdot (\text{id} + \text{swap}) \rangle \\ \equiv & \quad \{ \text{absorção-+} \} \\ & \text{mirror} = \langle [\text{Leaf}, \text{Fork} \cdot \text{swap}] \rangle \\ \equiv & \quad \{ \text{universal-cata} \} \\ & \text{mirror} \cdot \text{in} = [\text{Leaf}, \text{Fork} \cdot \text{swap}] \cdot (\text{id} + (\text{mirror} \times \text{mirror})) \\ \equiv & \quad \{ \text{def. in, fusão-+, absorção-+, eq-+} \} \\ & \begin{cases} \text{mirror} \cdot \text{Leaf} = \text{Leaf} \\ \text{mirror} \cdot \text{Fork} = \text{Fork} \cdot \text{swap} \cdot (\text{mirror} \times \text{mirror}) \end{cases} \\ \equiv & \quad \{ \text{pointwise, Def. comp} \} \\ & \begin{cases} \text{mirror} (\text{Leaf } x) = \text{Leaf } x \\ \text{mirror} (\text{Fork } (l, r)) = \text{Fork } (\text{swap } (\text{mirror } l, \text{mirror } r)) \end{cases} \\ \equiv & \quad \{ \text{def. swap} \} \\ & \begin{cases} \text{mirror} (\text{Leaf } x) = \text{Leaf } x \\ \text{mirror} (\text{Fork } (l, r)) = \text{Fork } (\text{mirror } r, \text{mirror } l) \end{cases} \end{aligned}$$

### Exercício 2

$$\begin{aligned} & \langle g \rangle \cdot \langle \text{in}_2 \cdot \alpha \rangle = \langle g \cdot \alpha \rangle \\ \Leftarrow & \quad \{ \text{fusão-cata} \} \\ & \langle g \rangle \cdot \text{in}_2 \cdot \alpha = g \cdot \alpha \cdot F \langle g \rangle \\ \equiv & \quad \{ (F1) \} \\ & g \cdot G \langle g \rangle \cdot \alpha = g \cdot \alpha \cdot F \langle g \rangle \\ \Leftarrow & \quad \{ \text{leibniz} \} \\ & G \langle g \rangle \cdot \alpha = \alpha \cdot F \langle g \rangle \\ \Leftarrow & \quad \{ \text{generalização de } \langle g \rangle \text{ em } f \} \\ & G f \cdot \alpha = \alpha \cdot F f \end{aligned}$$

### Exercício 3

$$\begin{aligned}
 & \text{mirror} = \langle g \rangle \\
 & \text{mirror} = \langle \text{in}_2 \cdot \alpha \rangle \\
 \\ 
 & \langle g \rangle = \langle \text{in}_2 \cdot \alpha \rangle \\
 \equiv & \quad \{ \text{def. in}_2, \text{def. } \alpha \} \\
 & \langle g \rangle = \langle [\text{Leaf}, \text{Fork}] \cdot (\text{id} + \text{swap}) \rangle \\
 \equiv & \quad \{ \text{absorção-+} \} \\
 & \langle g \rangle = \langle [\text{Leaf}, \text{Fork} \cdot \text{swap}] \rangle
 \end{aligned}$$

Podemos então dizer que  $g = [\text{Leaf}, \text{Fork} \cdot \text{swap}]$ . Precisamos também de provar que  $\text{id} = \langle g \cdot \alpha \rangle$ .

$$\begin{aligned}
 & \text{id} = \langle g \cdot \alpha \rangle \\
 \equiv & \quad \{ \text{def. } g, \text{def. } \alpha \} \\
 & \text{id} = \langle [\text{Leaf}, \text{Fork} \cdot \text{swap}] \cdot (\text{id} + \text{swap}) \rangle \\
 \equiv & \quad \{ \text{absorção-+} \} \\
 & \text{id} = \langle [\text{Leaf}, \text{Fork} \cdot \text{swap} \cdot \text{swap}] \rangle \\
 \equiv & \quad \{ \text{swap} \cdot \text{swap} = \text{id} \} \\
 & \text{id} = \langle [\text{Leaf}, \text{Fork}] \rangle \\
 \equiv & \quad \{ [\text{Leaf}, \text{Fork}] = \text{in}_{\text{LTree}} \} \\
 & \text{id} = \langle \text{in}_{\text{LTree}} \rangle \\
 \equiv & \quad \{ \text{reflexão-cata} \} \\
 & \text{id} = \text{id} \\
 & \square
 \end{aligned}$$

Podemos então provar que *mirror* é o seu próprio isomorfismo:

$$\begin{aligned}
 & \langle g \rangle \cdot \langle \text{in}_2 \cdot \alpha \rangle = \langle g \cdot \alpha \rangle \\
 \Longleftarrow & \quad \{ \text{(F3)} \} \\
 & G f \cdot \alpha = \alpha \cdot F f \\
 \equiv & \quad \{ G f = F f = \text{id} + f \times f, \text{def. } \alpha \text{ (twice)} \} \\
 & (\text{id} + f \times f) \cdot (\text{id} + \text{swap}) = (\text{id} + \text{swap}) \cdot (\text{id} + f \times f) \\
 \equiv & \quad \{ \text{functor-+} \} \\
 & (\text{id} + ((f \times f) \cdot \text{swap})) = (\text{id} + (\text{swap} \cdot (f \times f))) \\
 & \square
 \end{aligned}$$

Através da propriedade grátis da função *swap* (i.e:  $\text{swap} \cdot (f \times g) = (g \times f) \cdot \text{swap}$ ), podemos garantir a veracidade desta propriedade.

### Exercício 4

$$\begin{aligned}
 & \langle g \rangle \cdot T f = \langle g \cdot B(f, \text{id}) \rangle \\
 \equiv & \quad \{ \text{def-map-cata} \}
 \end{aligned}$$

$$\begin{aligned}
& \llbracket g \rrbracket \cdot \llbracket \text{id} \cdot B(f, \text{id}) \rrbracket = \llbracket g \cdot B(f, \text{id}) \rrbracket \\
\Leftarrow & \quad \{ \text{(F3)} \} \\
& Gf \cdot B(f, \text{id}) = B(f, \text{id}) \cdot Ff \\
\equiv & \quad \{ Gf = B(\text{id}, f) \} \\
& B(\text{id}, f) \cdot B(f, \text{id}) = B(f, \text{id}) \cdot B(\text{id}, f) \\
\equiv & \quad \{ \text{functor-id-F (para bi-funtores)} \} \\
& B(f, f) = B(f, f) \\
& \square
\end{aligned}$$

### Exercício 5

$$\begin{aligned}
& \text{while } p \text{ f } g = \text{tailr } ((g + f) \cdot (\neg \cdot p)?) \\
\equiv & \quad \{ \text{def. tailr} \} \\
& \text{while } p \text{ f } g = \llbracket \text{join}, (g + f) \cdot (\neg \cdot p)? \rrbracket \\
\equiv & \quad \{ \llbracket f, g \rrbracket = f \cdot F \llbracket f, g \rrbracket \cdot g \} \\
& \text{while } p \text{ f } g = \text{join} \cdot F(\text{while } p \text{ f } g) \cdot ((g + f) \cdot (\neg \cdot p)?) \\
\equiv & \quad \{ \text{def. join, def. functor F} \} \\
& \text{while } p \text{ f } g = [\text{id}, \text{id}] \cdot (\text{id} + (\text{while } p \text{ f } g)) \cdot ((g + f) \cdot (\neg \cdot p)?) \\
\equiv & \quad \{ \text{absorção-+ (twice)} \} \\
& \text{while } p \text{ f } g = [g, (\text{while } p \text{ f } g) \cdot f] \cdot (\neg \cdot p)? \\
\equiv & \quad \{ \text{def. condicional de McCarthy} \} \\
& \text{while } p \text{ f } g = (\neg \cdot p) \rightarrow g, (\text{while } p \text{ f } g) \cdot f \\
\equiv & \quad \{ \text{pointwise} \} \\
& \text{while } p \text{ f } g \ x = \text{if } \neg (p \ x) \text{ then } g \ x \text{ else while } p \text{ f } g (f \ x)
\end{aligned}$$

### Exercício 6

$$\begin{aligned}
& (\text{tailr } g) \cdot f = \text{tailr } h \\
\equiv & \quad \{ \text{def. tailr, def. hilomorfismo} \} \\
& \llbracket \Delta \rrbracket \cdot \llbracket [g] \rrbracket = \llbracket \Delta \rrbracket \cdot \llbracket [h] \rrbracket \\
\Leftarrow & \quad \{ \text{leibniz} \} \\
& \llbracket [g] \rrbracket \cdot f = \llbracket [h] \rrbracket \\
\Leftarrow & \quad \{ \text{fusão-ana, def. functor} \} \\
& g \cdot f = (\text{id} + f) \cdot h \\
& \square
\end{aligned}$$

### Exercício 7

$$\begin{aligned}
& f \bullet [g, h] = [f \bullet g, f \bullet h] \\
\equiv & \quad \{ \text{(F9) (3*)} \} \\
& \mu \cdot T f \cdot [g, h] = [\mu \cdot T f \cdot g, \mu \cdot T f \cdot h]
\end{aligned}$$

$$\equiv \{ \text{ fus\~{a}o} + \}$$

$$[\mu \cdot T f \cdot g, \mu \cdot T f \cdot h] = [\mu \cdot T f \cdot g, \mu \cdot T f \cdot h]$$

$$\square$$