

Cálculo de Programas

Resolução - Ficha 06

Eduardo Freitas Fernandes

2026

Exercício 1

$$\begin{aligned} & ap \cdot (\bar{f} \times id) = f \\ \equiv & \quad \{ \text{ pointwise } \} \\ & ap \cdot (\bar{f} \times id) (a, b) = f (a, b) \\ \equiv & \quad \{ \text{ def. comp, def-}\times \} \\ & \bar{f} a b = f (a, b) \\ \equiv & \quad \{ \text{ def. curry } \} \\ & f (a, b) = f (a, b) \\ \square & \end{aligned}$$

Exercício 2

$$\begin{aligned} & ap \cdot (\bar{f} \times id) = f \\ \equiv & \quad \{ \text{ pointwise } \} \\ & ap \cdot (\bar{f} \times id) (a, b) = f (a, b) \\ \equiv & \quad \{ \text{ def. comp, def-}\times \} \\ & \bar{f} a b = f (a, b) \\ \equiv & \quad \{ f = \hat{g} \} \\ & \bar{\hat{g}} a b = \hat{g} (a, b) \\ \equiv & \quad \{ \text{ def. curry } \} \\ & \hat{g} (a, b) = \hat{g} (a, b) \\ \square & \end{aligned}$$

Exercício 3

$$\begin{aligned} & \bar{f} \cdot (g \times h) = \overline{ap \cdot (id \times h)} \cdot \bar{f} \cdot g \\ \equiv & \quad \{ \text{ universal-exp } \} \\ & f \cdot (g \times h) = ap \cdot ((\overline{ap \cdot (id \times h)} \cdot \bar{f} \cdot g) \times id) \\ \equiv & \quad \{ \text{ natural-id, functor-}\times \} \\ & f \cdot (g \times h) = ap \cdot (\overline{ap \cdot (id \times h)} \times id) \cdot (\bar{f} \cdot g \times id) \end{aligned}$$

$$\begin{aligned}
&\equiv \{ \text{cancelamento-exp} \} \\
&f \cdot (g \times h) = ap \cdot (id \times h) \cdot ((\bar{f} \cdot g) \times id) \\
&\equiv \{ \text{functor-}\times \} \\
&f \cdot (g \times h) = ap \cdot ((\bar{f} \cdot g) \times h) \\
&\equiv \{ \text{natural-id, functor-}\times \} \\
&f \cdot (g \times h) = ap \cdot (\bar{f} \times id) \cdot (g \times h) \\
&\equiv \{ \text{cancelamento-exp} \} \\
&f \cdot (g \times h) = f \cdot (g \times h) \\
&\square
\end{aligned}$$

Exercício 4

$$\begin{aligned}
&flip(flip f) = f \\
&\equiv \{ \text{def. } flip \} \\
&\overline{(flip f) \cdot swap} = f \\
&\equiv \{ \text{pointwise} \} \\
&\overline{(flip f) \cdot swap} a b = f a b \\
&\equiv \{ \text{def. curry} \} \\
&((\widehat{flip f}) \cdot swap)(a, b) = f a b \\
&\equiv \{ \text{def. comp, def. swap} \} \\
&(\widehat{flip f})(b, a) = f a b \\
&\equiv \{ \text{def. uncurry} \} \\
&flip f b a = f a b \\
&\equiv \{ \dots \} \\
&\dots
\end{aligned}
\qquad
\begin{aligned}
&flip f x y = f y x \\
&\equiv \{ \text{def. } flip \} \\
&\widehat{f \cdot swap} x y = f y x \\
&\equiv \{ \text{def. curry} \} \\
&(\widehat{f \cdot swap})(x, y) = f y x \\
&\equiv \{ \text{def. comp, def. swap} \} \\
&\widehat{f}(y, x) = f y x \\
&\equiv \{ \text{def. uncurry} \} \\
&f y x = f y x \\
&\square
\end{aligned}$$

Exercício 5

$$\begin{aligned}
 & junc \cdot unjunc = id \\
 \equiv & \{ \text{ pointwise } \} & unjunc \cdot junc = id \\
 & (junc \cdot unjunc) l = id k \\
 \equiv & \{ \text{ def. } id, \text{ def. comp, def. } unjunc \} & (unjunc \cdot junc) (f, g) = id (f, g) \\
 & junc (k \cdot i_1, k \cdot i_2) = k \\
 \equiv & \{ \text{ def. } junc \} & \{ \text{ def. } id, \text{ def. comp, def. } junc \} \\
 & [k \cdot i_1, k \cdot i_2] = k & unjunc [\cdot, \cdot] f g = (f, g) \\
 \equiv & \{ \text{ fusão-+ } \} & \{ \text{ def. } unjunc \} \\
 & k \cdot [i_1, i_2] = k & ([f, g] \cdot i_1, [f, g] \cdot i_2) = (f, g) \\
 \equiv & \{ \text{ reflexão-+, natural-}id \} & \{ \text{ cancelamento-+ } \} \\
 & k = k & (f, g) = (f, g) \\
 \square & & \square
 \end{aligned}$$

Exercício 6

$$\begin{aligned}
 & \text{for } b i \cdot \text{in} = [g_1, g_2] \cdot (id + \text{for } b i) \\
 \equiv & \{ \text{ def. in, fusão-+, absorção-+ } \} \\
 & [\text{for } b i \cdot \underline{0}, \text{for } b i \cdot \text{succ}] = [g_1 \cdot id, g_2 \cdot \text{for } b i] \\
 \equiv & \{ \text{ natural-}id, \text{eq-+ } \} \\
 & \left\{ \begin{array}{l} \text{for } b i \cdot \underline{0} = g_1 \\ \text{for } b i \cdot \text{succ} = g_2 \cdot \text{for } b i \end{array} \right. \\
 \equiv & \{ \text{ pointwise, def. comp, def. } \underline{0}, \text{def. succ } \} \\
 & \left\{ \begin{array}{l} \text{for } b i \ 0 = g_1 () \\ \text{for } b i \ (n+1) = g_2 \text{ for } b i \ n \end{array} \right. \\
 \equiv & \{ (\text{F8}) \} \\
 & \left\{ \begin{array}{l} g_1 = i \\ g_2 \text{ for } b i \ n = b \text{ for } b i \ n \end{array} \right. \\
 \equiv & \{ \text{ igualdade extensional (73) } \} \\
 & \left\{ \begin{array}{l} g_1 = i \\ g_2 = b \end{array} \right.
 \end{aligned}$$

Exercício 8

$$\begin{aligned}
 & \left\{ \begin{array}{l} \text{for } b i \ 0 = i \\ \text{for } b i \ (n+1) = b \text{ for } b i \ n \end{array} \right. & \left\{ \begin{array}{l} a + 0 = a \\ a + (n+1) = 1 + (a+n) \end{array} \right. \\
 \equiv & \{ \text{ pointfree } \} & \equiv \{ \text{ pointfree } \} \\
 & \left\{ \begin{array}{l} \text{for } b i \cdot \underline{0} = i \\ \text{for } b i \cdot \text{succ} = b \cdot \text{for } b i \end{array} \right.
 \end{aligned}$$

Deduzimos então que $(a+) = \text{for succ } a$.

Exercício 9

```
int k(int n, int a) {
    int r = 0;
    int j;
    for (j = 1; j < n + 1; j++) {
        r = a + r;
    }
    return r;
};
```

Exercício 10

func b = (maybe b id) · flip lookup
 $a = [(140999000, "Manuel"), (200100300, "Mary"), (000111222, "Teresa")]$
 $b = [(140999000, "PT"), (200100300, "UK")]$
 $c = [(140999000, "Braga"), (200100300, "Porto"), (151999000, "Lisbon")]$