

Cálculo de Programas

Resolução - Ficha 03

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Exercício 1

$\equiv \{\text{def. assocl, Fusão-}\times, \text{Reflexão-}\times, \text{Eq-}\times\}$
 $\equiv \{\text{Def-}\times, \text{Universal-}\times\}$
 $\equiv \{\text{associação à direita}\}$
 $\equiv \{\text{Universal-}\times\}$
 $\equiv \{\text{Natural-id, Universal-}\times\}$

Exercício 2

```
ghci> assocr = split (p1 . p1) (p2 >< id)
ghci> :t assocr
assocr :: ((b1, b2), d) -> (b1, (b2, d))
ghci> assocr ((3,4), 5)
(3,(4,5))
ghci> assocr ((x,y), z) = (x, (y,z))
ghci> :t assocr
assocr :: ((a1, a2), b) -> (a1, (a2, b))
ghci> assocr ((3,4), 5)
(3,(4,5))
```

Exercício 3

$$\frac{f : A \rightarrow B \quad g : C \rightarrow D}{f \times g : A \times C \rightarrow B \times D} \quad \frac{f : A \rightarrow B \quad g : A \rightarrow C}{\langle f, g \rangle : A \rightarrow B \times C} \quad \frac{f : A \rightarrow B \quad g : B \rightarrow C}{f \cdot g : A \rightarrow C}$$
$$\frac{\pi_2 : A \times B \rightarrow B \quad \pi_1 : A \times B \rightarrow A}{\langle \pi_2, \pi_1 \rangle : A \times B \rightarrow B \times A} \quad \frac{id : A \rightarrow A \quad swap : B \times C \rightarrow C \times B}{id \times swap : A \times (B \times C) \rightarrow A \times (C \times B)}$$
$$\frac{swap : D \times E \rightarrow E \times D \quad id \times swap : A \times (B \times C) \rightarrow A \times (C \times B)}{swap \cdot (id \times swap) : A \times (B \times C) \rightarrow (C \times B) \times A}$$
$$(F0) : \quad (g \times f) \cdot swap = swap \cdot (f \times g)$$

$$\begin{aligned}
& \beta \cdot (f \times (g \times h)) \\
= & \{\text{def. } \beta\} \\
& \text{swap} \cdot (\text{id} \times \text{swap}) \cdot (f \times (g \times h)) \\
= & \{(F0)\} \\
& (\text{id} \times \text{swap}) \cdot \text{swap} \cdot (f \times (g \times h)) \\
= & \{(F0)\} \\
& (\text{swap} \times \text{id}) \cdot ((g \times h) \times f) \cdot \text{swap} \\
= & \{\text{Functor-}\times\} \\
& ((\text{swap} \cdot (g \times h)) \times (\text{id} \cdot f)) \cdot \text{swap} \\
= & \{(F0)\} \\
& (((h \times g) \cdot \text{swap}) \times (f \cdot \text{id})) \cdot \text{swap} \\
= & \{\text{Functor-}\times\} \\
& ((h \times g) \times f) \cdot (\text{swap} \times \text{id}) \cdot \text{swap} \\
= & \{(F0)\} \\
& ((h \times g) \times f) \cdot \text{swap} \cdot (\text{id} \times \text{swap}) \\
= & \{\text{def. } \beta\} \\
& ((h \times g) \times f) \cdot \beta
\end{aligned}$$

Exercício 4

$$\underline{k} \ x = \underline{k} \ (x) = \underline{k} \ (\text{id} \ x) = \underline{k} \cdot \text{id} = \underline{k} = k$$

Exercício 5

```

ghci> data X = B Bool | P (Bool,Int)
ghci> :t B
B :: Bool -> X
ghci> :t P
P :: (Bool, Int) -> X
ghci> f = either B P
ghci> :t f
f :: Either Bool (Bool, Int) -> X

```

Exercício 6

$$\begin{array}{c}
\frac{\underline{False} : A \rightarrow \text{Bool} \quad \text{id} : A \rightarrow A}{\langle \underline{False}, \text{id} \rangle : A \rightarrow \text{Bool} \times A}
\end{array}
\qquad
\begin{array}{c}
\frac{f : A \rightarrow C \quad g : B \rightarrow C}{[f, g] : A + B \rightarrow C}
\end{array}$$

$$\frac{\langle \underline{False}, \text{id} \rangle : A \rightarrow \text{Bool} \times A \quad \langle \underline{True}, \text{id} \rangle : A \rightarrow \text{Bool} \times A}{[\langle \underline{False}, \text{id} \rangle, \langle \underline{True}, \text{id} \rangle] : A + A \rightarrow \text{Bool} \times A}$$

Exercício 7

$$\begin{aligned}
\alpha &= [\langle \underline{False}, id \rangle, \langle \underline{True}, id \rangle] \\
&\equiv \{\text{Universal-+}\} \\
&\quad \begin{cases} \alpha \cdot i_1 = \langle \underline{False}, id \rangle \\ \alpha \cdot i_2 = \langle \underline{True}, id \rangle \end{cases} \\
&\equiv \{\text{point wise}\} \\
&\quad \begin{cases} (\alpha \cdot i_1) a = \langle \underline{False}, id \rangle a \\ (\alpha \cdot i_2) a = \langle \underline{True}, id \rangle a \end{cases} \\
&\equiv \{\text{def. composição, def. split}\} \\
&\quad \begin{cases} \alpha(i_1 a) = (\underline{False} a, id a) \\ \alpha(i_2 a) = (\underline{True} a, id a) \end{cases} \\
&\equiv \{\text{def. const, def. id}\} \\
&\quad \begin{cases} \alpha(i_1 a) = (False, a) \\ \alpha(i_2 a) = (True, a) \end{cases}
\end{aligned}$$

```
ghci> alpha (Left a) = (False, a)
ghci> alpha (Right a) = (True, a)
```

Exercício 8

$$\begin{array}{c}
\begin{array}{c} \pi_1 : A \times B \rightarrow A \\ id : C \rightarrow C \end{array} \\
\hline
\pi_1 \times id : (A \times B) \times C \rightarrow A \times C
\end{array}
\qquad
\begin{array}{c}
\begin{array}{c} \pi_2 : A \times B \rightarrow B \\ \pi_1 : (A \times B) \times C \rightarrow A \times B \end{array} \\
\hline
\pi_2 \cdot \pi_1 : (A \times B) \times C \rightarrow B
\end{array}$$

$$\begin{aligned}
&\frac{\pi_1 \times id : (A \times B) \times C \rightarrow A \times C \quad \pi_2 \cdot \pi_1 : (A \times B) \times C \rightarrow B}{\langle \pi_1 \times id, \pi_2 \cdot \pi_1 \rangle : (A \times B) \times C \rightarrow (A \times C) \times B} \\
&xr \cdot \langle \langle f, g \rangle, h \rangle = \langle \langle f, h \rangle, g \rangle \\
&\equiv \{\text{Universal-}\times\} \\
&\quad \begin{cases} \pi_1 \cdot xr \cdot \langle \langle f, g \rangle, h \rangle = \langle f, h \rangle \\ \pi_2 \cdot xr \cdot \langle \langle f, g \rangle, h \rangle = g \end{cases} \\
&\equiv \{\text{def. xr, Cancelamento}\} \\
&\quad \begin{cases} (\pi_1 \times id) \cdot \langle \langle f, g \rangle, h \rangle = \langle f, h \rangle \\ \pi_2 \cdot \pi_1 \cdot \langle \langle f, g \rangle, h \rangle = g \end{cases} \\
&\equiv \{\text{Absorção-}\times, \text{Cancelamento-}\times\} \\
&\quad \begin{cases} \langle \pi_1 \cdot \langle f, g \rangle, id \cdot h \rangle = \langle f, h \rangle \\ \pi_2 \cdot \langle f, g \rangle = g \end{cases} \\
&\equiv \{\text{Cancelamento-}\times\} \\
&\quad \begin{cases} \langle f, h \rangle = \langle f, h \rangle \\ g = g \end{cases}
\end{aligned}$$

Exercício 9

```

mkInd :: (Bib, Aux) -> Ind
mkInd = map (id >< sort)
      . uncurry applyMapping
      . swap
      . (invertRelation >< invertRelation)

invertRelation :: (Ord a, Ord b) => [(a, [b])] -> [(b, [a])]
invertRelation = groupPairs . sortOn p1 . map swap . expandPairs

expandPairs :: [(a, [b])] -> [(a, b)]
expandPairs = concat . map (\(x, l) -> map (x,) l)

groupPairs :: (Eq a) => [(a, b)] -> [(a, [b])]
groupPairs = uncurry zip . split
            (nub . map p1)
            (map (map p2)
              . groupBy (curry (uncurry (==)
                                   . (p1 >< p1))))))

applyMapping :: (Eq b) => [(b, [c])] -> [(a, [b])] -> [(a, [c])]
applyMapping a = map (applyMappingOne a)

applyMappingOne :: (Eq b) => [(b, [c])] -> (a, [b]) -> (a, [c])
applyMappingOne d = id >< (concat . map ( \x -> case lookup x d of
                                           Just xs -> xs
                                           _ -> []))

```

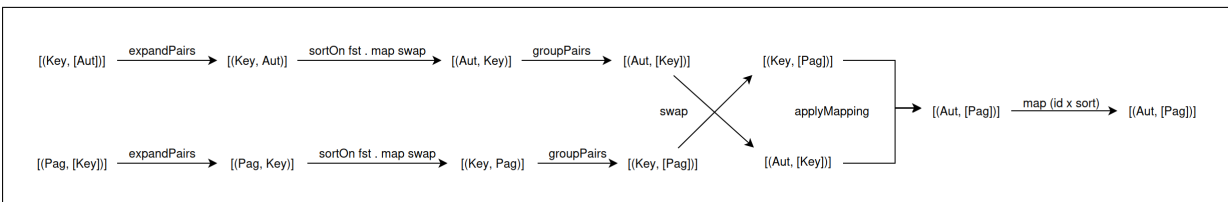


Figura 1: mkInd