

Cálculo de Programas

Resolução - Ficha 06

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Exercício 1

$$\begin{aligned} & ap \cdot (\bar{f} \times id) = f \\ \equiv & \quad \{ \text{pointwise} \} \\ & ap \cdot (\bar{f} \times id) (a, b) = f (a, b) \\ \equiv & \quad \{ \text{def. comp, def-}\times \} \\ & \bar{f} a b = f (a, b) \\ \equiv & \quad \{ \text{def. curry} \} \\ & f (a, b) = f (a, b) \\ & \square \end{aligned}$$

Exercício 2

$$\begin{aligned} & ap \cdot (\bar{f} \times id) = f \\ \equiv & \quad \{ \text{pointwise} \} \\ & ap \cdot (\bar{f} \times id) (a, b) = f (a, b) \\ \equiv & \quad \{ \text{def. comp, def-}\times \} \\ & \bar{f} a b = f (a, b) \\ \equiv & \quad \{ f = \widehat{g} \} \\ & \widehat{\bar{g}} a b = \widehat{g} (a, b) \\ \equiv & \quad \{ \text{def. curry} \} \\ & \widehat{g} (a, b) = \widehat{g} (a, b) \\ & \square \end{aligned}$$

Exercício 3

$$\begin{aligned} & \overline{f \cdot (g \times h)} = \overline{ap \cdot (id \times h) \cdot \bar{f} \cdot g} \\ \equiv & \quad \{ \text{universal-exp} \} \\ & f \cdot (g \times h) = ap \cdot (\overline{ap \cdot (id \times h) \cdot \bar{f} \cdot g} \times id) \\ \equiv & \quad \{ \text{natural-id, functor-}\times \} \\ & f \cdot (g \times h) = ap \cdot (\overline{ap \cdot (id \times h)} \times id) \cdot (\bar{f} \cdot g \times id) \end{aligned}$$

$$\begin{aligned}
&\equiv \{ \text{cancelamento-exp} \} \\
&\quad f \cdot (g \times h) = ap \cdot (id \times h) \cdot ((\bar{f} \cdot g) \times id) \\
&\equiv \{ \text{functor-}\times \} \\
&\quad f \cdot (g \times h) = ap \cdot ((\bar{f} \cdot g) \times h) \\
&\equiv \{ \text{natural-id, functor-}\times \} \\
&\quad f \cdot (g \times h) = ap \cdot (\bar{f} \times id) \cdot (g \times h) \\
&\equiv \{ \text{cancelamento-exp} \} \\
&\quad f \cdot (g \times h) = f \cdot (g \times h) \\
&\square
\end{aligned}$$

Exercício 4

$$\begin{aligned}
&\quad flip (flip f) = f \\
&\equiv \{ \text{def. flip} \} \\
&\quad \overline{(flip f) \cdot swap} = f \\
&\equiv \{ \text{pointwise} \} \\
&\quad \overline{(flip f) \cdot swap a b} = f a b \\
&\equiv \{ \text{def. curry} \} \\
&\quad ((\widehat{flip f}) \cdot swap) (a, b) = f a b \\
&\equiv \{ \text{def. comp, def. swap} \} \\
&\quad (\widehat{flip f}) (b, a) = f a b \\
&\equiv \{ \text{def. uncurry} \} \\
&\quad flip f b a = f a b \\
&\equiv \{ \dots \} \\
&\dots
\end{aligned}$$

$$\begin{aligned}
&\quad flip f x y = f y x \\
&\equiv \{ \text{def. flip} \} \\
&\quad \overline{\widehat{f} \cdot swap x y} = f y x \\
&\equiv \{ \text{def. curry} \} \\
&\quad (\widehat{f} \cdot swap) (x, y) = f y x \\
&\equiv \{ \text{def. comp, def. swap} \} \\
&\quad \widehat{f} (y, x) = f y x \\
&\equiv \{ \text{def. uncurry} \} \\
&\quad f y x = f y x \\
&\square
\end{aligned}$$

Exercício 5

$$\begin{aligned}
& \text{junc} \cdot \text{unjunc} = \text{id} \\
\equiv & \quad \{ \text{pointwise} \} \\
& (\text{junc} \cdot \text{unjunc}) \, l = \text{id} \, k \\
\equiv & \quad \{ \text{def. id, def. comp, def. unjunc} \} \\
& \text{junc} (k \cdot i_1, k \cdot i_2) = k \\
\equiv & \quad \{ \text{def. junc} \} \\
& [k \cdot i_1, k \cdot i_2] = k \\
\equiv & \quad \{ \text{fusão-+} \} \\
& k \cdot [i_1, i_2] = k \\
\equiv & \quad \{ \text{reflexão-+, natural-id} \} \\
& k = k
\end{aligned}$$

□

$$\begin{aligned}
& \text{unjunc} \cdot \text{junc} = \text{id} \\
\equiv & \quad \{ \text{pointwise} \} \\
& (\text{unjunc} \cdot \text{junc}) (f, g) = \text{id} (f, g) \\
\equiv & \quad \{ \text{def. id, def. comp, def. junc} \} \\
& \text{unjunc} [\cdot, \cdot] f g = (f, g) \\
\equiv & \quad \{ \text{def. unjunc} \} \\
& ([f, g] \cdot i_1, [f, g] \cdot i_2) = (f, g) \\
\equiv & \quad \{ \text{cancelamento-+} \} \\
& (f, g) = (f, g)
\end{aligned}$$

□

Exercício 6

$$\begin{aligned}
& \text{for } b \, i \cdot \text{in} = [g_1, g_2] \cdot (\text{id} + \text{for } b \, i) \\
\equiv & \quad \{ \text{def. in, fusão-+, absorção-+} \} \\
& [\text{for } b \, i \cdot \underline{0}, \text{for } b \, i \cdot \text{succ}] = [g_1 \cdot \text{id}, g_2 \cdot \text{for } b \, i] \\
\equiv & \quad \{ \text{natural-id, eq-+} \} \\
& \left\{ \begin{array}{l} \text{for } b \, i \cdot \underline{0} = g_1 \\ \text{for } b \, i \cdot \text{succ} = g_2 \cdot \text{for } b \, i \end{array} \right. \\
\equiv & \quad \{ \text{pointwise, def. comp, def. } \underline{0}, \text{def. succ} \} \\
& \left\{ \begin{array}{l} \text{for } b \, i \, 0 = g_1 \, () \\ \text{for } b \, i \, (n+1) = g_2 \, \text{for } b \, i \, n \end{array} \right. \\
\equiv & \quad \{ (\text{F8}) \} \\
& \left\{ \begin{array}{l} g_1 = i \\ g_2 \, \text{for } b \, i \, n = b \, \text{for } b \, i \, n \end{array} \right. \\
\equiv & \quad \{ \text{igualdade extensional (73)} \} \\
& \left\{ \begin{array}{l} g_1 = i \\ g_2 = b \end{array} \right.
\end{aligned}$$

Exercício 8

$$\begin{aligned}
& \left\{ \begin{array}{l} \text{for } b \, i \, 0 = i \\ \text{for } b \, i \, (n+1) = b \, \text{for } b \, i \, n \end{array} \right. \\
\equiv & \quad \{ \text{pointfree} \} \\
& \left\{ \begin{array}{l} \text{for } b \, i \cdot \underline{0} = \underline{i} \\ \text{for } b \, i \cdot \text{succ} = b \cdot \text{for } b \, i \end{array} \right.
\end{aligned}$$

$$\begin{aligned}
& \left\{ \begin{array}{l} a + 0 = a \\ a + (n+1) = 1 + (a + n) \end{array} \right. \\
\equiv & \quad \{ \text{pointfree} \} \\
& \left\{ \begin{array}{l} (a+) \cdot \underline{0} = \underline{a} \\ (a+) \cdot \text{succ} = \text{succ} \cdot (a+) \end{array} \right.
\end{aligned}$$

Deduzimos então que $(a+) = \text{for succ } a$.

Exercício 9

```
int k(int n, int a) {  
    int r = 0;  
    int j;  
    for (j = 1; j < n + 1; j++) {  
        r = a + r;  
    }  
    return r;  
};
```

Exercício 10

func b = (maybe b id.) · flip lookup

a = [(140999000, "Manuel"), (200100300, "Mary"), (000111222, "Teresa")]

b = [(140999000, "PT"), (200100300, "UK")]

c = [(140999000, "Braga"), (200100300, "Porto"), (151999000, "Lisbon")]