

Cálculo de Programas

Resolução - Ficha 11

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Exercício 1

$$\begin{aligned}
 & \text{mirror} = \langle \langle in_2 \cdot \alpha \rangle \rangle \\
 \equiv & \{ \text{Def. } in_2, \text{Def. } \alpha \} \\
 & \text{mirror} = \langle \langle [Leaf, Fork] \cdot (id + swap) \rangle \rangle \\
 \equiv & \{ \text{Absorção-+} \} \\
 & \text{mirror} = \langle \langle [Leaf, Fork \cdot swap] \rangle \rangle \\
 \equiv & \{ \text{Universal-cata} \} \\
 & \text{mirror} \cdot in_{LTree} = [Leaf, Fork \cdot swap] \cdot (id + (\text{mirror} \times \text{mirror})) \\
 \equiv & \{ \text{Def. } in_{LTree}, \text{Fusão-+}, \text{Absorção-+}, \text{Eq-+} \} \\
 & \begin{cases} \text{mirror} \cdot Leaf = Leaf \\ \text{mirror} \cdot Fork = Fork \cdot swap \cdot (\text{mirror} \times \text{mirror}) \end{cases} \\
 \equiv & \{ \text{pointwise, Def. composição} \} \\
 & \begin{cases} \text{mirror} (Leaf \ x) = Leaf \ x \\ \text{mirror} (Fork \ (l, r)) = Fork \ (swap \ (\text{mirror} \ l, \text{mirror} \ r)) \end{cases} \\
 \equiv & \{ \text{Def. swap} \} \\
 & \begin{cases} \text{mirror} (Leaf \ x) = Leaf \ x \\ \text{mirror} (Fork \ (l, r)) = Fork \ (\text{mirror} \ r, \text{mirror} \ l) \end{cases}
 \end{aligned}$$

Exercício 2

$$\begin{aligned}
 & \dots\dots\dots \\
 \Leftarrow & \{ \text{Fusão-cata} \} \\
 & \dots\dots\dots \\
 \equiv & \{ \langle \langle g \rangle \rangle \cdot in_2 = g \cdot G \ \langle \langle g \rangle \rangle \text{ (F1)} \} \\
 & \dots\dots\dots \\
 \Leftarrow & \{ \text{Leibniz} \} \\
 & \dots\dots\dots \\
 \Leftarrow & \{ \text{Generalização de } \langle \langle g \rangle \rangle \text{ em f} \} \\
 & \dots\dots\dots
 \end{aligned}$$

Exercício 3

$$\begin{aligned} mirror &= \langle\!\langle g \rangle\!\rangle \\ mirror &= \langle\!\langle in_2 \cdot \alpha \rangle\!\rangle \end{aligned}$$

$$\begin{aligned} \langle\!\langle g \rangle\!\rangle &= \langle\!\langle in_2 \cdot \alpha \rangle\!\rangle \\ &\equiv \{\text{Def. } in_2, \text{Def. } \alpha\} \\ \langle\!\langle g \rangle\!\rangle &= \langle\!\langle [Leaf, Fork] \cdot (id + swap) \rangle\!\rangle \\ &\equiv \{\text{Absorção-+}\} \\ \langle\!\langle g \rangle\!\rangle &= \langle\!\langle [Leaf, Fork \cdot swap] \rangle\!\rangle \end{aligned}$$

Podemos então dizer que $g = [Leaf, Fork \cdot swap]$. Precisamos também de provar também que $id = \langle\!\langle g \cdot \alpha \rangle\!\rangle$.

$$\begin{aligned} id &= \langle\!\langle g \cdot \alpha \rangle\!\rangle \\ &\equiv \{\text{Def. } g, \text{Def. } \alpha\} \\ id &= \langle\!\langle [Leaf, Fork \cdot swap] \cdot (id + swap) \rangle\!\rangle \\ &\equiv \{\text{Absorção-+}\} \\ id &= \langle\!\langle [Leaf, Fork \cdot swap \cdot swap] \rangle\!\rangle \\ &\equiv \{swap \cdot swap = id\} \\ id &= \langle\!\langle [Leaf, Fork] \rangle\!\rangle \\ &\equiv \{[Leaf, Fork] = in_{LTree}\} \\ id &= \langle\!\langle in_{LTree} \rangle\!\rangle \\ &\equiv \{\text{Reflexão-cata}\} \\ &True \end{aligned}$$

Podemos então provar que *mirror* é o seu próprio isomorfismo:

$$\begin{aligned} \langle\!\langle g \rangle\!\rangle \cdot \langle\!\langle in_2 \cdot \alpha \rangle\!\rangle &= \langle\!\langle g \cdot \alpha \rangle\!\rangle \\ \iff \{(F3)\} \\ G \ f \cdot \alpha &= \alpha \cdot F \ f \\ &\equiv \{G \ f = F \ f = id + f \times f\} \\ (id + f \times f) \cdot (id + swap) &= (id + swap) \cdot (id + f \times f) \\ &\equiv \{\text{Functor-+}\} \\ (id + (f \times f) \cdot swap) &= (id + swap \cdot (f \times f)) \end{aligned}$$

Através da propriedade grátis da função *swap* (i.e: $swap \cdot (f \times g) = (g \times f) \cdot swap$), podemos garantir a veracidade desta propriedade.

Exercício 4

$$\begin{aligned} & \llbracket g \rrbracket \cdot T f = \llbracket g \cdot B(f, id) \rrbracket \\ & \equiv \{\text{Def-map-cata}\} \\ & \llbracket g \rrbracket \cdot \llbracket in \cdot B(f, id) \rrbracket = \llbracket g \cdot B(f, id) \rrbracket \\ & \Leftarrow \{(F3)\} \\ & G f \cdot B(f, id) = B(f, id) \cdot F f \\ & \equiv \{\text{????}\} \\ & B(id, f) \cdot B(f, id) = B(f, id) \cdot B(id, f) \\ & \equiv \{\text{Functor-id-F (para Bi Functores)}\} \\ & B(f, f) = B(f, f) \end{aligned}$$

Exercício 5

$$\begin{aligned} & while\ p\ f\ g = tailr\ ((g + f) \cdot (\neg \cdot p)?) \\ & \equiv \{\text{Def. tailr}\} \\ & while\ p\ f\ g = \llbracket join, ((g + f) \cdot (\neg \cdot p)?) \rrbracket \\ & \equiv \{\llbracket f, g \rrbracket = f \cdot F \llbracket f, g \rrbracket \cdot g\} \\ & while\ p\ f\ g = join \cdot F (while\ p\ f\ g) \cdot ((g + f) \cdot (\neg \cdot p)?) \\ & \equiv \{\text{Def. join, Def. Functor}\} \\ & while\ p\ f\ g = [id, id] \cdot (id + (while\ p\ f\ g)) \cdot ((g + f) \cdot (\neg \cdot p)?) \\ & \equiv \{\text{Absorção-+ (2×)}\} \\ & while\ p\ f\ g = [g, (while\ p\ f\ g) \cdot f] \cdot (\neg \cdot p)? \\ & \equiv \{\text{Def. Condicional de McCarthy}\} \\ & while\ p\ f\ g = (\neg \cdot p) \rightarrow g, (while\ p\ f\ g) \cdot f \\ & \equiv \{\text{pointwise}\} \\ & while\ p\ f\ g\ x = \text{if not } (p\ x) \text{ then } g\ x \text{ else } while\ p\ f\ g\ (f\ x) \end{aligned}$$

Exercício 6

$$\begin{aligned} & \dots\dots\dots \\ & \equiv \{\text{Def. tailr, Def. hylomorfismo}\} \\ & \dots\dots\dots \\ & \Leftarrow \{\text{Leibniz}\} \\ & \dots\dots\dots \\ & \Leftarrow \{\text{Fusão-ana, Def. Functor}\} \\ & \dots\dots\dots \end{aligned}$$

Exercício 7

$$\begin{aligned} & f \bullet [g, h] = [f \bullet g, f \bullet h] \\ & \equiv \{(F9) (3\times)\} \\ & \mu \cdot T f \cdot [g, h] = [\mu \cdot T f \cdot g, \mu \cdot T f \cdot h] \\ & \equiv \{\text{Fusão-+}\} \\ & [\mu \cdot T f \cdot g, \mu \cdot T f \cdot h] = [\mu \cdot T f \cdot g, \mu \cdot T f \cdot h] \end{aligned}$$