# Companion slides for The Art of Multiprocessor Programming by Maurice Herlihy & Nir Shavit

- Provides First-Come-First-Served
- · How?
  - Take a "number"
  - Wait until lower numbers have been served
- · Lexicographic order
  - -(a,i) > (b,j)
    - If a > b, or a = b and i > j

```
class Bakery implements Lock {
   boolean[] flag;
   Label[] label;
  public Bakery (int n) {
    flag = new boolean[n];
    label = new Label[n];
    for (int i = 0; i < n; i++) {
       flag[i] = false; label[i] = 0;
```

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class Bakery implements Lock {
   boolean[] flag;
   Label[] label;
  public Bakery (int n) {
    flag = new boolean[n];
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    for (int i = 0; i < n; i++) {
   flag[i] = false; label[i] = 0;</pre>
```

```
class Bakery implements Lock {
    in some arbitrary

public void lock() {
        flag[i] = true;
    label[i] = max(label[0], ..., label[n-1])+1;

while (∃k flag[k]
        && (label[i],i) > (label[k],k));
}
```

Take increasing

```
class Bakery implements Lock {
  boolean flag[n];
                       Someone is
  int label[n];
                       interested
 public void lock() {
  flag[i] = true;
             max(label[0], ..., label[n-1])+1;
 while (3k flag[k]
          && (label[i],i) > (label[k],k));
```

## With lower (label,i) in lexicographic order

```
class Bakery implements Lock {
    No longer interested

    public void an lock {
        flag[i] = false;
    }
}

labels are always increasing
```

#### No Deadlock

- There is always one thread with earliest label
- Ties are impossible (why?)

- Suppose A and B in CS together
- Suppose A has earlier label
- When B entered, it must have seen
  - flag[A] is false, or
  - label[A] > label[B]

- · Labels are strictly increasing so
- B must have seen flag[A] == false

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- · Labels are strictly increasing so
- B must have seen flag[A] == false
- · Labeling<sub>B</sub>  $\rightarrow$  read<sub>B</sub>(flag[A])  $\rightarrow$  write<sub>A</sub>(flag[A])  $\rightarrow$  Labeling<sub>A</sub>
- Which contradicts the assumption that A has an earlier label



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