



Multimedia Networking

Transport Protocols Real-time Support RTP/RTCP

Main Sources:

- *Computer Networking: A Top Down Approach* 7th edition. Jim Kurose, Keith Ross, Addison-Wesley, 2017.
- RFC3550, RFC5506



Internet transport protocols

RTP motivation

- TCP not suited to real-time and interactive apps
 - point-to-point so not suitable for multicast
 - no way to associate **timing** with segments
 - causes arbitrary delays
- SCTP is multistream but has some of TCP drawbacks
- UDP does not include **timing** information nor any support for real-time apps. DCCP provides some support for buffer control but no media timing.
- Real-Time Protocol (RTP) aims to support the transport of *real-time* data, audio and video

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RTP



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- Goal:** provides mechanisms for end-to-end delivery of real-time data (audio, video)
- Typical apps: videoconferencing, IP telephony, etc.
 - Supports unicast or multicast sessions
 - Uses a data channel (RTP) and a control channel (RTCP)
 - RTCP (Real Time Control Protocol) reports the status of the data channel
 - RTP (RFC 3350)

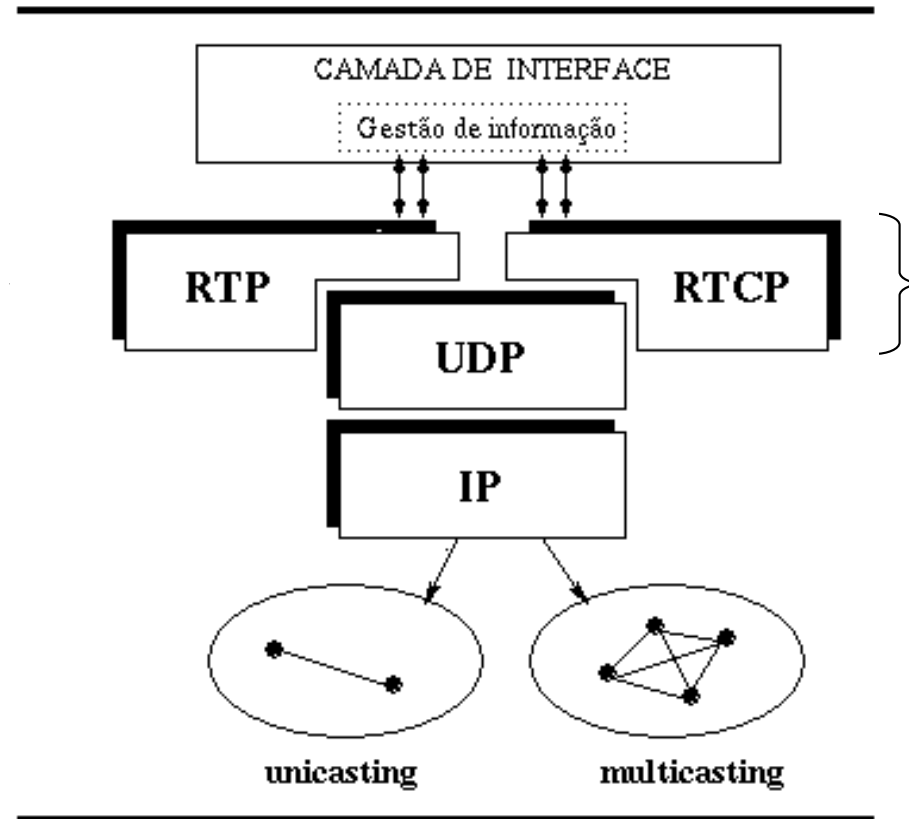
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RTP



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- Support for application real-time data delivery



- Support to control application real-time data

- RTP in the protocol stack

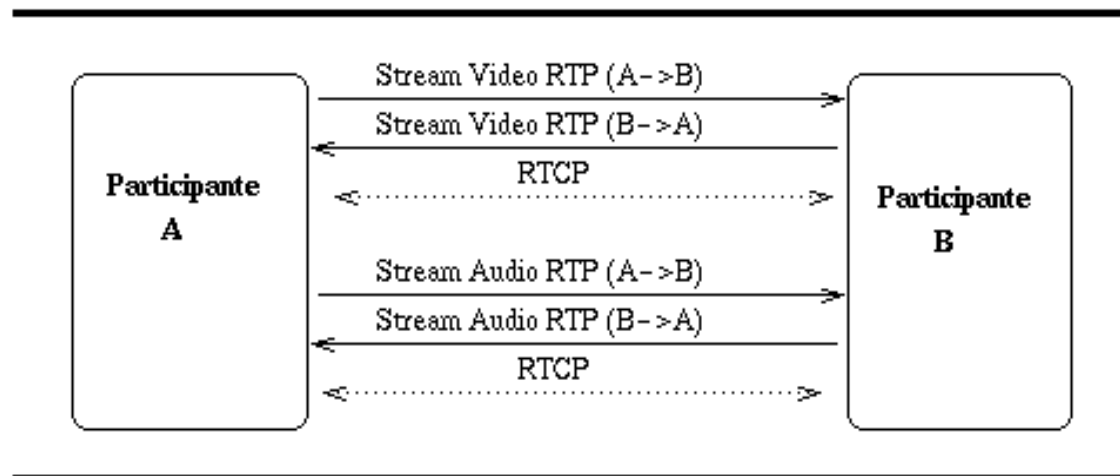
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- **RTP Session** = data channel + control channel
- Session address = network address (unicast, multicast), pair of ports
- RTCP channel = RTP port + 1
- In presence of multiple media streams, each one is transmitted in a separate session or mixed in a single session in each direction, depending on the codecs (e.g. MPEG-x mix audio and vídeo / session).



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- RTP issues
 - * **No** reservation, **No** Quality of Service (QoS) guarantees
 - * **No** guarantee of packet delivery
 - * Note that RTP **does not** provide any mechanism to ensure timely data delivery or other QoS guarantees.
 - * routers may provide differentiated service delivery, but there is no guarantees that RTP packets arrive at destination in a timely manner.
 - * RTP encapsulation is only seen at end-systems, not by intermediate routers.

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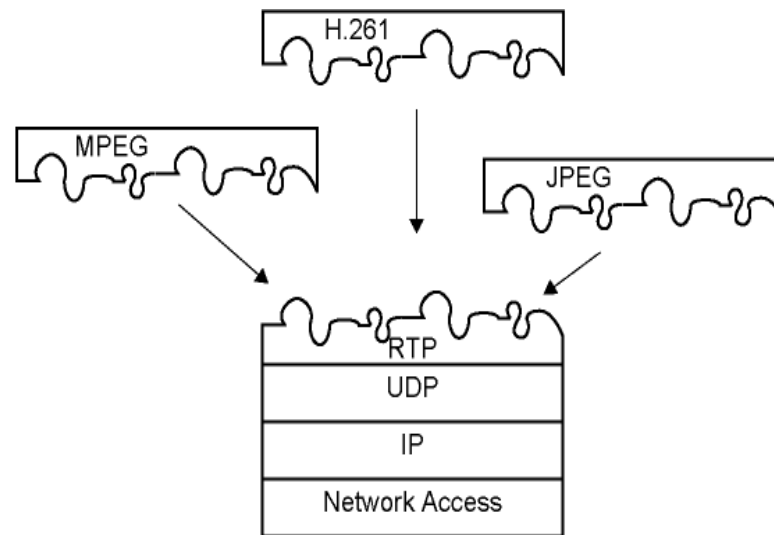
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- RTP issues

- * **Allows** timing reconstruction, loss detection and content identification



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- ❑ RTP specifies a packet structure for packets carrying audio, video data
- ❑ RTP packet provides
 - time stamping
 - packet sequence numbering
 - payload type identification
- ❑ RTP packets usually encapsulated in UDP segments (also in DCCP)
- ❑ Interoperability: if two Internet phone applications run RTP, then they may be able to work together (session control at RTP level)



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RTP example

- ❑ Consider sending 64 kbps PCM-encoded (8 bits) voice over RTP.
- ❑ application collects encoded data in chunks, e.g., every 20 msec = 160 samples in a chunk (160 bytes).
- ❑ an audio chunk + RTP header **form an RTP packet**, which is encapsulated in e.g. an UDP segment
- ❑ RTP header indicates type of audio encoding in each packet
 - sender can change encoding during conference.
- ❑ As mentioned, RTP header also contains sequence number and timestamp.

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RTP



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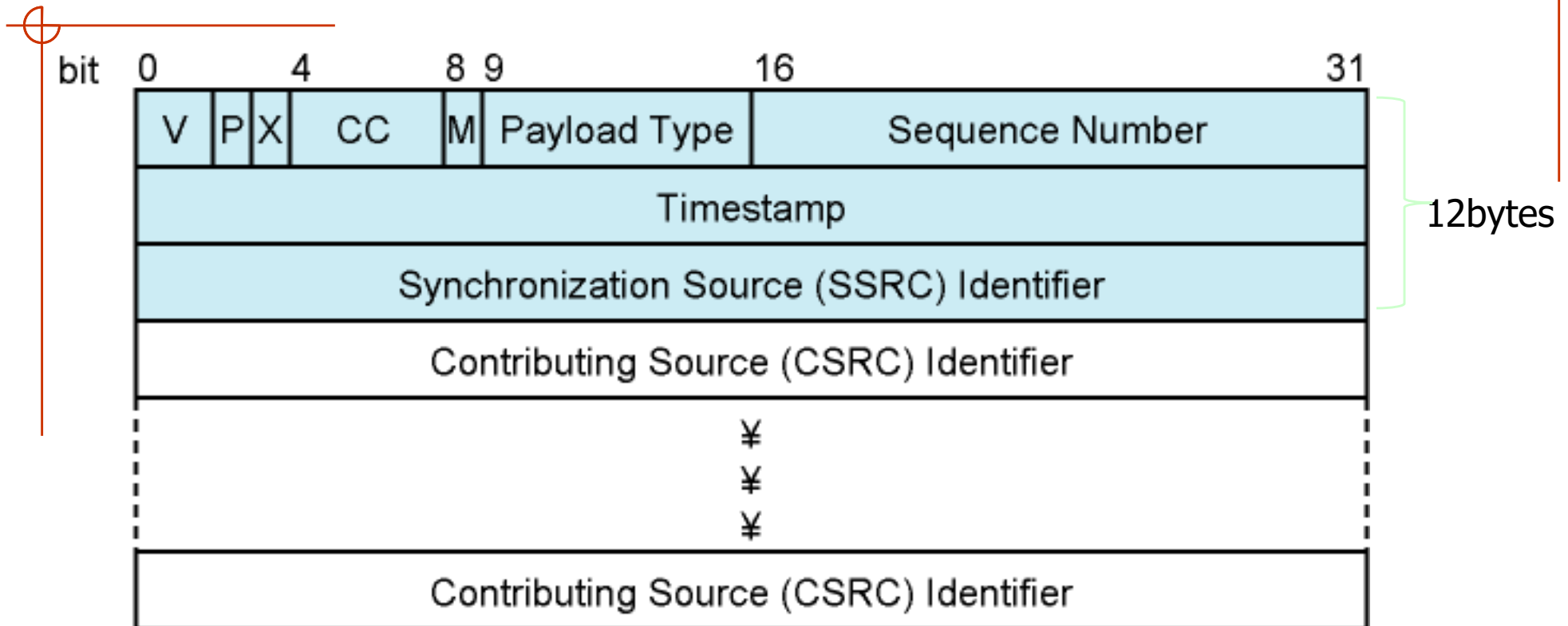
- RTP Timestamp
 - * **primary goal**: represent the inherent notion of **real-time** associated with media (sampling units)
 - * used to place packets in correct timing order
 - * may be used to evaluate jitter
- Sequence Number
 - * used to detect RTP packet loss
- If a video frame spans multiple RTP packets, then timestamp is unchanged and sequence number is incremented

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RTP header



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V = Version
P = Padding
X = Extension
CC = CSRC count
M = Marker

[W.Stallings, Computer Networks w/
Internet Prot. & Tech., 2004]

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RTP header



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Payload Type (7 bits): Indicates the type of encoding currently being used. If sender changes encoding in middle of a conference, sender informs receiver via payload type field.

- Payload type 0: PCM mu-law, 64 kbps
- Payload type 3, GSM, 13 kbps
- Payload type 7, LPC, 2.4 kbps
- Payload type 26, Motion JPEG
- Payload type 31. H.261
- Payload type 33, MPEG2 video
- see RFC 3551...

Seq. Number (16 bits): Increments by one for each RTP packet sent; may be used to detect packet loss and to restore packet sequence at receiver.

SSRC field (32 bits): Identifies the source of the RTP stream (allows independence of the network address), provides an ID for playout. Each stream in RTP session should have distinct SSRC (random init).

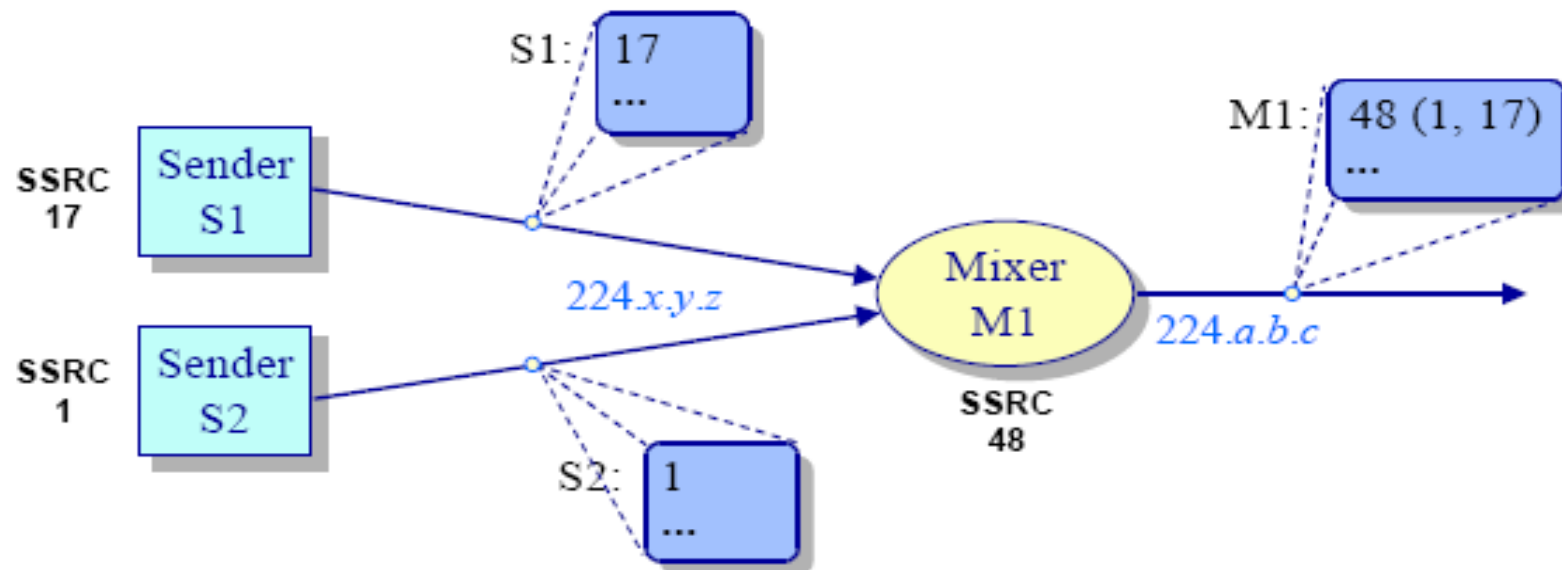
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RTP header – SSRCs vs. CSRCs



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RTP mixer - an intermediate system that receives and combines RTP PDUs of one or more RTP sessions into a new RTP PDU



Synchronization Sources(SSRC) become Contributing Sources (CSRC)

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RTP header - timestamp



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- ❑ Timestamp field (32 bits): reflects the **sampling instant** of first octet in the RTP data packet
 - initial value random for each RTP stream
 - for audio, timestamp clock typically increments by one for each sampling period (e.g., each 125 usecs for 8 kHz sampling clock)
 - if application generates chunks of 160 encoded samples each 20ms, then timestamp increases by 160 units for each RTP packet when source is active.
 - timestamp clock continues to increase at constant rate when source is inactive; it increments monotonically and linearly
 - if the sampling rate is variable (e.g. software codecs), the system clock may be used (e.g. gettimeofday())



Internet transport protocols

RTCP – Real-Time Control Protocol

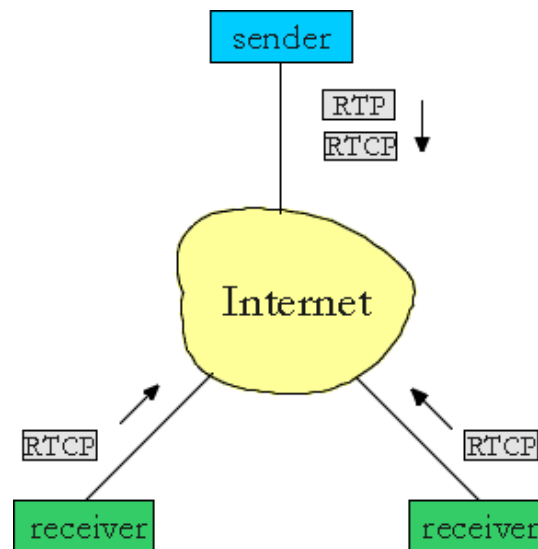
- ❑ RTCP works in conjunction with RTP.
- ❑ each participant in RTP session periodically transmits RTCP control packets to all other participants.
- ❑ each RTCP packet contains **sender and/or receiver reports**
- ❑ minimal overhead (e.g. 1 RTCP pkt every 5 sec)
- ❑ feedback can be used to control application performance
 - sender may modify its transmissions based on feedback
 - report statistics useful to application: #packets sent, #packets lost, interarrival jitter, etc.
- ❑ allows to identify and keep track of participants

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- ❑ each RTP session: typically uses a single multicast address; all RTP/RTCP packets belonging to that session **use multicast** address (IP class D).
- ❑ RTP, RTCP packets distinguished from each other via distinct port numbers.
- ❑ to limit traffic, each participant reduces RTCP traffic as the number of conference participants increases



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RTCP packets

Sender Report (SR) packets:

- ❑ SSRC of RTP stream, current time, number of packets sent, number of octets sent
- ❑ allows receiver to:
 - estimate mean data rate;
 - estimate mean packet size;
 - distinguish reception breaks (pauses in transmission vs. network problems)

Receiver Report (RR) packets:

- ❑ fraction of packets lost, last sequence number, average interarrival jitter

Source Description packets:

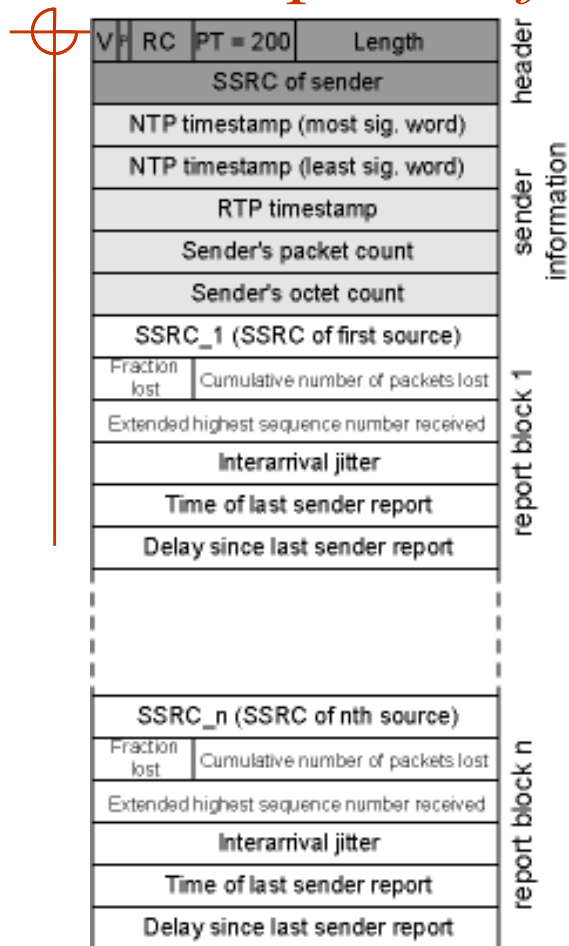
- ❑ e-mail address of sender, sender's name, SSRC of associated RTP stream
- ❑ provide mapping between the SSRC and the user/host name

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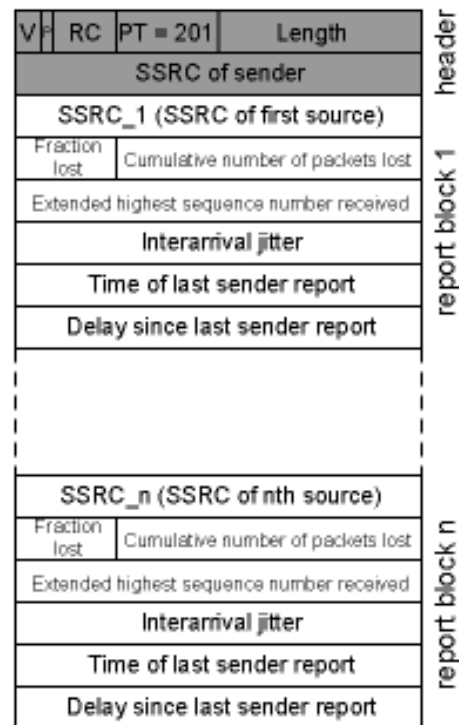
RTCP packet formats



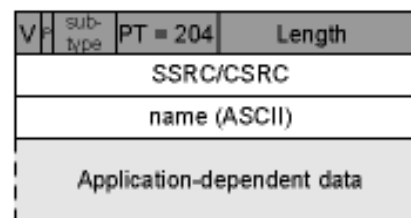
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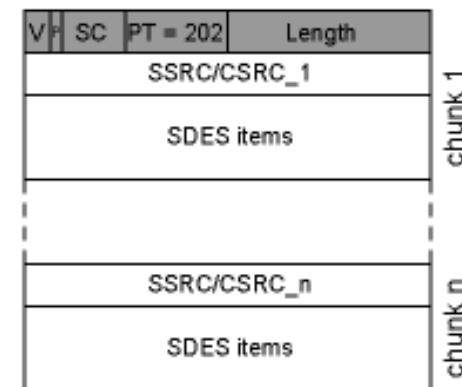
(a) RTCP Sender Report



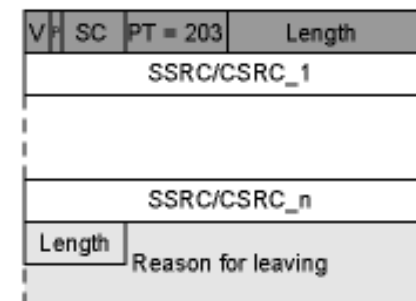
(b) RTCP Receiver Report



(c) RTCP Application-defined packet



(d) RTCP Source Description



(e) RTCP BYE

[W.Stallings, Computer Networks w/
Internet Prot. & Tech., 2004]

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RTCP Reports



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- ❑ RRs and SRs indicate reception statistics from each receiver and for every sender. These statistics include:
 - ❑ **packet loss ratio** since the last SR or RR was sent;
 - ❑ **total # packets lost since the beginning of the session**
 - ❑ **highest sequence number received** so far, which allows to estimate how much data is in flight (when used together with the SR and RR timestamps);
 - ❑ **moving average of the interarrival jitter** of media packets; which gives the sender an indirect view of the size of any adaptive playout buffer used at the receiver ([RFC3611]), e.g. useful for VoIP sessions.

[RFC 5968., 2010]



Internet transport protocols

RTCP synchronization of streams

- ❑ RTCP can **synchronize different media** streams within a RTP session
- ❑ consider videoconferencing app for which each sender generates one RTP stream for video, one for audio.
- ❑ timestamps in RTP packets are **tied to** the video, audio **sampling clocks**
 - **not** tied to wall-clock time (NTP gives an absolute time)
- ❑ each RTCP sender-report packet contains (for most recently generated packet in associated RTP stream):
 - timestamp of RTP packet
 - wall-clock time for when packet was created.
- ❑ receivers uses association to synchronize playout of audio, video

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RTCP bandwidth scaling



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- ❑ RTCP has a potential bandwidth scaling problem as the number of receivers R increases
- ❑ RTCP aggregated bandwidth consumption increases linearly with R
- ❑ To minimize that, RTCP receivers adjusts the sending rate to the multicast tree according to R



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RTCP bandwidth scaling

- ❑ RTCP attempts to limit its traffic to 5% of session bandwidth.

Example

- ❑ Suppose one sender, sending video at 2 Mbps. Then RTCP attempts to limit its traffic to 100 kbps.
- ❑ RTCP gives 75% of rate to receivers; remaining 25% to sender
- ❑ 75 kbps is equally shared among receivers:
 - with R receivers, each receiver gets to send RTCP traffic at $75/R$ kbps.
- ❑ sender gets to send RTCP traffic at 25 kbps.
- ❑ participant determines RTCP packet transmission period by calculating avg RTCP packet size (across entire session) and dividing by its allocated rate

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RTP/RTCP references



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- Lots of good material at Henning Schulzrinne's RTP page
<http://www.cs.columbia.edu/~hgs/rtp/>
- RFCs on RTP/RTCP
 - G.Ott, C.Perkins, *Guidelines for Extending the RTP Control Protocol (RTCP)*, RFC 5968, Sep. 2010.
 - C.Perkins, T. Schierl, *Rapid Synchronisation of RTP Flows*, RFC 6051, Nov. 2010
 - A.Begen, *RTP Control Protocol (RTCP) Port for Source-Specific Multicast (SSM) Sessions*, RFC 6128, Feb. 2011.
 - A. Began, C. Perkins, D.Wing, *Guidelines for Choosing RTP Control Protocol (RTCP) Canonical Names (CNAMEs)*, RFC 6222, Apr. 2011.
 - ...