

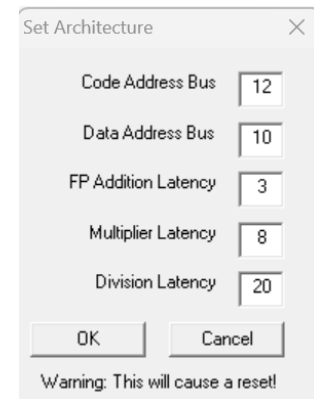
**Laboratory**  
**2**

Expected delivery of lab\_02.zip must include:

- **program\_1.s** and **program\_2.s**
- This file, filled with information and possibly compiled in a pdf format.

Please, configure the winMIPS64 simulator with the *Base Configuration* provided in the following (*in italics not user controllable configuration*):

- Code address bus: 12
- Data address bus: 10
- Pipelined FP arithmetic unit (latency): 3 stages
- Pipelined multiplier unit (latency): 8 stages
- divider unit (latency): not pipelined unit, 20 clock cycles
- Forwarding is enabled
- Branch prediction is disabled
- Branch delay slot is disabled
- *Integer ALU: 1 clock cycle*
- *Data memory: 1 clock cycle*
- *Branch delay slot: 1 clock cycle.*



- 1) Write an assembly program (**program\_1.s**) for the *winMIPS64* architecture described before able to implement the following piece of code described at high-level:

```
for (i = 0; i < 64; i++) {
    v5[i] = ((v1[i]* v2[i]) + v3[i]) + v4[i];
    v6[i] = v5[i]/(v4[i]+v1[i]);
    v7[i] = v6[i]*(v2[i]+v3[i]);
}
```

Assume that the vectors **v1[]**, **v2[]**, **v3[]**, and **v4[]** are allocated previously in memory and contain 64 double precision **floating point** values; assume also that **v1[]** and **v4[]** do not contain 0 values. Additionally, the vectors **v5[]**, **v6[]**, **v7[]** are empty vectors also allocated in memory.

**Calculate** the data memory footprint of your program:

Data	Number of Bytes
V1	512
V2	512
V3	512
V4	512
V5	512
V6	512
V7	512
Total	3584

Are there any issues? Yes, where and why? No ? Do you need to change something?

Your answer:

We need to modify the data bus size, when set to 10 it can only store  $2^{10}$  bytes which is 1024, store the 3584 bytes we need at least a bus size of 12 ( $2^{12} == 4096$  bytes)

**ATTENTION:** winMIPS64 has a limitation due to the underlying software. There is a limitation in the string length when declaring a vector. Split the vectors elements in multiple lines (it also increases the readability) .

Example: my\_fancy\_vector: .byte 4, 5 ,7, 8  
.byte 5,77, 8  
.byte .....

a. Compute the CPU performance equation (CPU time) of the previous program following the next directions, assume a clock frequency of 1MHz:

CPU time =  $\frac{IC_i}{CPI_i}$

- Count manually, the number of the different instructions ( $IC_i$ ) executed in the program
- Assume that the  $CPI_i$  for every type of instructions equals the number of clock cycles in the instruction EXE stage, for example:
  - integer instructions  $CPI = 1$
  - LD/SD instructions  $CPI = 1$
  - FP MUL instructions  $CPI = 8$
  - FP DIV instructions  $CPI = 20$
  - ...

b. Compute by hand again the CPU performance equation assuming that you can improve the FP Multiplier or the FP Divider by speeding up by 2 only one of the units at a time:

- Pipelined FP multiplier unit (latency):  $8 \rightarrow 4$  stages  
Or
- FP Divider unit (latency): not pipelined unit,  $20 \rightarrow 10$  clock cycles

Table 1: CPU time by hand

	CPU Time initial (a)	CPU Time (b – MUL speed up)	CPU Time (b – DIV speed up)
program_1.S	3.777 ms	3.265 ms	3.137

c. Using the simulator calculate again the CPU time and complete the following table:

Table 2: CPU time using the simulator

	CPU Time initial (a)	CPU Time (b – MUL speed up)	CPU Time (b – DIV speed up)
program_1.S	3.845 ms	3.333 ms	3.205 ms

Are there any difference? If yes, where and why? If Not, provide some comments in the following:

Your answer:

It is slightly slower.

Maybe some load instructions are slower in practice because of stalls.

- d. Using the simulator and the *Base Configuration*, disable the Forwarding option and compute how many clock cycles the program takes to execute.

Table 3: forwarding disabled

	Number of clock cycles	IPC (Instructions Per Clock)
program_1.S	4871	0.237

Enable one at a time the **optimization features** that were initially disabled and collect statistics to fill the following table (fill all required data in the table before exporting this file to pdf format to be delivered).

Table 4: **Program performance for different processor configurations**

Program	Forwarding		Branch Target Buffer		Delay Slot		Forwarding + Branch Target Buffer	
	IPC	CC	IPC	CC	IPC	CC	IPC	CC
Program_1.S	0.3	3845	0.240	4812	N/A		0.305	3786

- 2) Using the WinMIPS64 simulator, validate experimentally the Amdahl's law, defined as follows:

$$\text{speedup}_{\text{overall}} = \frac{\text{execution time}_{\text{old}}}{\text{execution time}_{\text{new}}} = \frac{1}{(1 - \text{fraction}_{\text{enhanced}}) + \frac{\text{fraction}_{\text{enhanced}}}{\text{speedup}_{\text{enhanced}}}}$$

- a. Using the program developed before: **program\_1.s**
- b. Modify the processor architectural parameters related with multicycle instructions (Menu→Configure→Architecture) in the following way:
  - 1) Configuration 1
    - Starting from the *Base Configuration*, change only the FP addition latency to 6
  - 2) Configuration 2
    - Starting from the *Base Configuration*, change only the Multiplier latency to 4
  - 3) Configuration 3
    - Starting from the *Base Configuration*, change only the division latency to 10

Compute by hand (using the Amdahl's Law) and using the simulator the speed-up for any one of the previous processor configurations. Compare the obtained results and complete the following table.

Table 5: **program\_1.s** speed-up computed by hand and by simulation

Proc. Config.	Base config. [c.c.]	Config. 1	Config. 2	Config. 3
Speed-up comp.				
By hand	1	$\frac{1/((1-0.203) + 0.203/0.5)}{0.83} =$	$\frac{1/((1-0.271) + 0.271/2)}{1.16} =$	$\frac{1/((1-0.339) + 0.339/2)}{1.20} =$
By simulation	1	0.83	1.15	1.20

- 3) Write an assembly program (**program\_2.s**) for the winMIPS64 architecture able to compute the output (y) of a **neural computation** (see the Fig. below):

$$x = \sum_{j=0}^{K-1} i_j * w_j + b$$

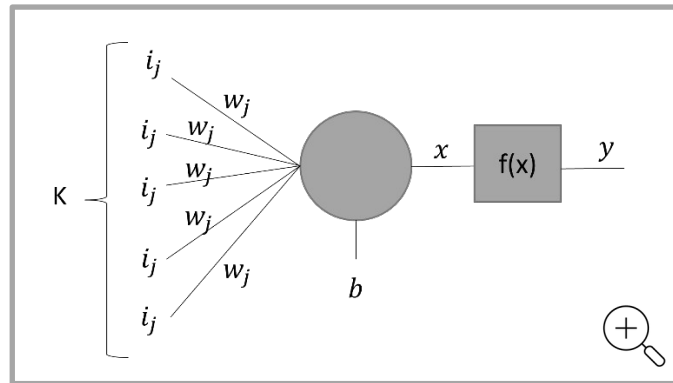
$$y = f(x)$$

where, to prevent the propagation of NaN (Not a Number), the activation function  $f$  is defined as:

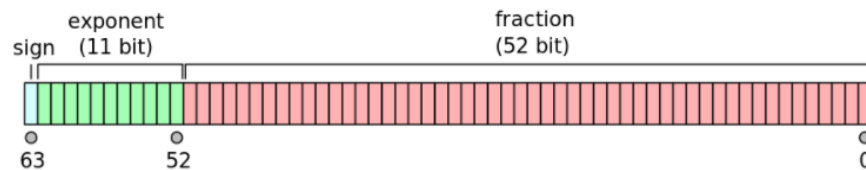
$$f(x) = \begin{cases} 0, & \text{if the exponent part of } x \text{ is equal } 0x7f \\ x, & \text{otherwise} \end{cases}$$

Assume the vectors  $i$  and  $w$  respectively store the inputs entering the neuron and the weights of the connections. They contain  $K=30$  double precision **floating point**

elements. Assume that  $b$  is a double precision **floating point** constant and is equal to  $0xab$ , and  $y$  is a double precision **floating point** value stored in memory. Compute  $y$ .



Below is reported the encoding of IEEE 754 double-precision binary floating-point format:



Given the *Base Configuration*, run your program and extract the following information.

	Number of clock cycles	IPC (Instructions Per Clock)	CPI (Clock per Instructions)
program_2.S	792	0.651	1.535