Operating Systems II

Process Management

Process Management

- We already talked about the task_struct data structure that holds all of the information about the task.
- Also talked about
 - the states a task can exist in.
 - the queues the task can be held in.
- What about the other aspects of process management.

Task Creation

- □ Every task is created from another task.
 - Maybe from the system idle task?
- □ A new task is created when a parent process invokes either a fork() or clone() system call.
 - clone() inherits more from the parent task.
- The system calls eventually call the do_fork() function that creates a new task.

do fork()

- Allocate a new task_struct data structure.
- Link the task_struct into the process table.
- Create a new kernel space stack for execution when inside the kernel.
- Copy the fields from the parent's task_struct into the child's.

do fork() (cont.)

- Modify the child's task_struct fields specific to itself.
 - New process identifier
 - Links to the tasks parent and siblings
 - initialize the process specific timers. (creation, time quantum, and so on)

do fork() (cont.)

- Copy other data structures that are in parent and should be replicated for the child task.
 - File table and new file descriptor for each open file.
 - Create a new user data segment and copy the data.
 - Copy signal and signal handling information.
 - Copy virtual memory tables.
- Change the child's state to RUNNING.

Task Termination

- □ Several ways to terminate a task
 - task making the exit() system call.
 - delivered a signal with the signal handler disposition to die.
 - Being forced to die under certain exceptions.
- Termination work is done by do exit().
- ☐ What does do exit() do?

do exit()

- Sets a global kernel lock.
- Sets the task state to ZOMBIE
- Notifies any child of process termination.
- Notifies the parent with an exit signal (usually SIGCHILD).
- Releases any resources allocated by do fork() (such as open files, ...)
- Calls schedule() (never returns).

Task Scheduling

- ☐ Tasks are scheduled via the schedule() function.
- What does schedule() do?
 - Exactly what happens is dependent on the scheduling algorithm.

Schedule() - Overview

- Performs various periodic work (such as running the bottom halves of interrupt handlers)
- Inspect the set of tasks in the running state.
- Chose one task to execute based on the scheduling policy.
- Dispatch the task to run on the CPU until an interrupt occurs.

Schedule()

- □ Release the global kernel lock.
- □ Do any software interrupts (2nd stage).
- ☐ Grab current process and current CPU.
- Set state appropriately.
- Find next process to schedule.
- Use switch_to() macro to perform the transfer (saves state of old task, loads state of new task).
 - Architecture Specific!!!

Process Table Support

- Some common functions found in a kernel to assist process management
 - init_task The first task in the process table.
 - find_task_by_pid(x) given a pid will return the task_struct pointer.
 - for_each_task(p) For loop construct that will allow you to iterate over each task in the process table.