== Abstract. ==

One of the difficulties faced when using a general purpose graphics processing on memory intensive tasks, is the considerable amount of time taken to transfer data from a CPU. Such is the case when one tries to upload a projection index from a CPU onto a GPU. One way to minimize the amount of data that needs to be transferred is through the use of compression. In this paper a Run Length Encoding (RLE) compression scheme is used to minimize the size of the data needed to be transferred. The idea is to compress a projection index using the RLE scheme and then decompress it within the GPU using a decompression algorithm. Two algorithms were proposed: one is load unbalanced and the other is load balanced. Both algorithms used the parallel prefix sum as a building block. The parallel prefix sum helps the algorithms determine how to allocate and copy the decompressed projection index within the GPU. To conclude, a benchmark test was performed comparing the two different algorithms suggested. It was determined from the results that the performance of the algorithms depends mostly on the GPU used and its computability, the percentage of compression of the projection index, and the nature of the data being processed. For a GPU with computability of 1.1, the algorithm showed no signs of performance improvement, mainly because of the number of cores and their clock rates to be too small to compete with the speed of the GPU’s bus. Moreover, under computability 1.1 accesses to global memory that are not sequential are not coalesced making the algorithms even slower. For a GPU with computability 1.3, an improvement in performance was observed, and the improvement was consistent for the load balanced algorithm, but it was not for the load unbalanced algorithm.

== Introduction ==

The projection index supports thread-level parallelism and therefore could potentially make good use of a GPU. However, most of the time spent when doing a query evaluation with a projection index, is spent in transferring data from the CPU to the GPU. Gosink et al [#x], improve on this bottleneck by reducing the size of the data that needs to be transferred; they do so by changing the encoding of the data that needs to be transferred. In contrast, in this paper, to reduce the size of the data that will be transferred from the CPU to the GPU, compression is used. Then after the index is transferred compressed, it is decompressed using a decompression algorithm within the GPU itself.

The compressed projection index is sent to the GPU as two separate arrays. One array is the frequencies, which represent the number of times an attribute value repeats, and the other array consists of the attribute values themselves. For the RLE compression scheme to be useful in reducing the size of the data, the projection index must be created on columns that are not unique and that allow themselves to be compressed somewhat. After the algorithm is sent to the GPU, it is decompressed using a parallel decompression algorithm. [You ought to give a figure here, showing a concrete example of the compressed input and the corresponding decompressed output]

Two algorithms were designed to perform this job in parallel, and both of them use Prefix sum.

== Parallel Prefix sum algorithm ==

The prefix sum algorithm is an essential building block for uncompressing a projection index that has been previously compressed in the RLE encoding format. #x et al [#x], present a method to calculate the prefix sum of an array in parallel. The article classifies two types of prefix sums, or scans as they are also called, inclusive and exclusive. The inclusive scan generates a new array in which every element j is the sum of all elements up to and including j. The exclusive scan, on the other hand, is an operation that contains the sum of all previous elements, but not j itself. Both types of scans are illustrated in figure 1.

Figure 1. Illustrates both the exclusive scan and inclusive scan of an arbitrary array of frequencies.

A scan can be performed sequentially to run on a single thread. Two arrays are kept for such scan, one is the input array, and the other is the output array. The input array contains the original elements before the scan and the output array is the array generated after the scan. To perform the scan, a loop is executed over the elements in the input array where the sum of the previous element of the input array and the output array is assigned to the current element in the output array. The algorithm is illustrated in the following listing (listing 1).

Listing 1. Sequential scan algorithm (Taken from Harris [#x])

void **scan**( float\* output, float\* input, int length)

{

output[0] = 0; *// since this is an exclusive scan*

for(int j = 1; j < length; ++j)

{

output[j] = input[j-1] + output[j-1];

}

}

To perform the algorithm in parallel, Harris started with a simple naïve algorithm and moved to one more complex but with better performance. The first algorithm presented in the paper is the naïve parallel scan. This algorithm assumes that there is one processor for each data element. For a GPU running CUDA this cannot be accomplished as the number of elements will often surpass the number of processors available. To work around this problem a double-buffer array is used, such that warps may work on arrays of 512 elements at a time. 512 elements are processed at a time, because this is the largest block size and data can only be synchronized within the block. The algorithm is illustrated in listing 2, and an illustration of how it performs the additions is illustrated in Figure 1. It is important to note that these operations are performed within the same array; the elements are added such that the distance between the elements increases each time by a power of 2.

Listing 2. Naïve parallel scan (Taken from Harris [#x])





Figure 1: Naïve parallel scan performed on 8 elements. (Taken from Harris [#x])

The naïve parallel scan has a work complexity equal to sum from d = 1 to log base 2 n n - 2^(d-1) = O(n log base 2 n ) addition operations. This scan's work complexity is even greater than the sequential scan which is of O(n) and therefore it is not work-efficient. The factor of Log base 2 n can significantly worsen the performance for the algorithm as n increases.

Harris also developed a work-efficient scan algorithm; to do this he employed an algorithmic pattern that is based on an algorithm used to build balanced binary trees in parallel. The algorithm consists of two phases: the reduce phase and the down-sweep phase. In the reduce phase, also called the up-sweep phase, the tree is traversed from the leaves to the root. As it traverses, it computes partial sums of neighboring nodes each time increasing the distance between them by a power of 2, until reaching the root of the tree. The root of the tree would hold the sum of all the nodes in the array. Pseudocode for this phase is listed in listing 3, and an illustration of the process is given in figure 2.

Listing 3.





Figure 2. Up-sweep or reduce phase on 8 elements. (Taken from Harris [#x])

Following the first phase of the algorithm, the second phase completes computing the scan by performing a down-sweep phase. The down-sweep phase starts from the root of the tree and uses the partial sums computed in the first phase. It discards the last sum of all elements, and replaces it with an element of value 0. A series of swap adds follows in which the sum of neighboring elements is assigned to the rightmost element. In this phase the distance between the neighboring elements decreases by powers of 2, starting from the last distance in the up-sweep phase. The pseudocode for the algorithm is listed in Listing 4, and an illustration of its process is given in Figure 3.

Listing 4.





Listing 4: The down-sweep phase of the work efficient parallel sum scan algorithm. It can be noted that the first step discards the last element of the array replacing it with a 0 (Taken from Harris [#x]).

== Design of Algorithms for Decompression ==

Two algorithm design approaches are taken for decompressing a projection index in the GPU (The projection index in these terms may also be referred to be a string). Both approaches use the prefix scan differently. The first algorithm which is called the unbalanced approach uses the prefix sum algorithm as an indicator for each thread to know from where to where to write the elements to allocate the uncompressed index. The algorithm is performed in two phases, and it is unbalanced because the workload of each thread is different, with threads handling elements that are heavily repeated doing most of the work.

The second algorithm, called the Load Balanced approach has five phases. In two of these phases parallel prefix sums are performed and the ending result of the last prefix sum is an array representing the uncompressed index. The algorithm is nicely load balanced because the amount of work done by each thread is the same. However, this algorithm uses almost twice as much memory as the unbalanced approach, and performs more than twice the number of kernel calls.

The input to both algorithms is a pair of arrays that represent the RLE compressed index. One array contains the symbols S, which are the different attribute values found on the projection index. The other array F contains the frequencies or repetitions of each index in the encoded algorithm. The length of both arrays is the same, as they correspond with each other; this length is denoted as C as it is the number of elements in compressed form. The length of the uncompressed index is denoted as U. The following figure (figure #x) illustrates the process where X3Y1Z7 is uncompressed sent as two arrays with length C = 3, and once uncompressed having a length of U = 11.

== The Load Unbalanced approach ==

The algorithm starts doing the decompression by appending an element of value 0 at the end of the array of frequencies, F. After doing that it obtains array X of length C+1 by performing an exclusive scan on the array of frequencies F, which also has been modified to have C+1. The last element of the exclusive scan X is the sum of all the frequencies, which is also the length of the decompressed array, U. It is used to allocate the amount of memory necessary in the GPU to hold the uncompressed array. The exclusive scan, array X, is then used to have each thread uncompress each element by writing it from one initial offset to where the element that is being repeated is changed. These two offsets are given by the exclusive scan and thus the decompression is performed. The workload of this algorithm is not balanced, as the amount of work a thread does depends directly on the frequency of the element it is decompressing. Threads decompressing elements with few repetitions will take considerably less time, than threads decompressing elements with various repetitions. The pseudocode for this algorithm is given in Listing 5.

Listing 5.

Phase 1: add element of value 0 at the end of array F.

get X as exclusive-scan of C+1 elements of F

Phase 2: for i:= 0 to C

forall k in parallel do

for j := X[i] to X[i+1]

result [j] = S[j]

== The Load Balanced approach ==

Taking in consideration that the first algorithm was not load balanced and that it depends greatly on the on the nature of the data it handles, a second approach was taken. This algorithm starts similar to the first one by obtaining an exclusive scan of the array of frequencies. However, in this algorithm the frequencies array is not modified, and thus it has a length of C, and so the exclusive scan, X will also have length of C. Adding the last element of the exclusive scan to the last frequency; one obtains the size of the uncompressed array, U. Having this size, memory for the decompressed array is allocated in the GPU, and phase two is on track. In the second phase, the uncompressed array A is initialized by having each thread assign a value 0 to all its positions. After the second phase is completed, Phase three has each thread write a 1 to positions given by the elements of the exclusive scan X. In other words, X represents the positions of the uncompressed array where there will be a change in symbol. Phase four follows, and an inclusive scan is performed on the now modified array A. The output of this scan is the positions of the corresponding symbols in the uncompressed array. In the final phase, phase five, these positions are used to write the actual symbols onto the final uncompressed index denoted B. The pseudocode for this algorithm is given in Listing 6, and an example illustration of the process is given in Figure #x.

Listing 6.

Phase 1: get X as Exclusive-scan of F

Phase 2: for i = 0 to U

forall k in parallel do

write 0 to item i in array A

Phase 3: for i = 0 to C

forall k in parallel do

write a 1 to item X[i] in array A

Phase 4: overwrite A as an Inclusive-scan of array A

Phase 5: for i = 0 to U

forall k in parallel do

write item S[A[i]] to Uncompressed Index B

To improve further on the performance of this algorithm, one could possibly do away with the 5th phase by sending the query one wants to perform in the projection index in terms of the positions of the elements in S. Not only would this save time, but also memory, as array B of size U would no longer need to be created.

== Performance Analysis==

== Data Analyzed ==

To test the algorithms, synthetic data was generated in the CPU that would represent a projection index after being sorted and loaded from disk. This is to simplify the interpretation on performance tests. Two different arrays were created, one with characters or integers representing the attribute values, and one with the number of times each of those attribute values repeat themselves. The distribution of this data was simulated in three different distributions: In the first distribution, the data was generated such that the next element repeats itself once more than its previous element. This is done first for 500 elements, and it is increased by 500 until reaching 5000 elements. This type of distribution did not favor the load unbalanced algorithm. In the second case, the attribute values are repeated a 1024 times initially, and in the next iteration that amount is doubled, so essentially the same number of elements would repeat themselves more times in each execution. Finally, the last distribution of data consisted of an uncompressed index of fixed size, 16777216 elements in total. The number of different repeated elements was then varied from 1024 to 8’388,608 by doubling on each iteration. An example of the data distributions is shown in Figure #x.

Both algorithms suffer mostly from the amount of time it takes to actually move a poorly compressed index from the CPU to the GPU, such as when an index has each element repeat only two times. This problem is inherent of the data, as it is not always compressible under the RLE scheme and thus moving to the GPU and decompressing would be a waste of time. However there were good cases where compressing and decompressing the projection index proved to save time.

== Performance on a GPU with Computability 1.1 ==

== Problems ==

Initially the tests were performed on an NVIDIA 9400m GPU, however the results obtained from this machine were …

Initially, tests were performed in a 9400m NVidia GPU with 16 cores and 256 Mb of video main memory. Both algorithms were tested against copying the uncompressed index directly to the GPU. The data distribution of this first test was the sequentially incremented one (Owen.. How would you call that type of distribution??). The outcome of the test in this GPU was that neither the Load Unbalanced (LU) nor Load Balanced (LB) algorithm was a good approach to improve the time it takes to transfer a projection index. Transferring the uncompressed index (UC) proved to be a better option (See Figure #x). Presumably, the GPU did not have sufficient cores to make the computations necessary quickly enough to uncompress faster than transferring the index (Rate of transfer from CPU to GPU). Additionally this GPU has a computability of 1.1, and so writes are not coalesced if the positions were writes on an array are not organized sequentially. There are such writes in the Load Balanced Algorithm.

It was also noticed that there was not much difference between the Load Balanced algorithm and the unbalanced algorithm. Each Phase of the load balanced algorithm was analyzed in a pie chart (see Figure #x) to determine which phases took the majority of time to do the decompression. From the pie chart, it can be determined that Phase 4 takes the most time of the algorithm. Most of this is attributed (blamed?) to the fact that the graphics processor only had 16 cores, and the sheer size of the array. Phases 2 and Phase 5 follow Phase 4 in amount of time taken. Phase two takes a long time due to the sheer size of the array which could only improve with a greater number of processing elements. The writes in this phase should be coalesced because they are sequential. The long time taken in Phase 5 is attributed to the computability issue. This is issues is shared by Phase 3, but it is only significant in Phase 5 as the amount of computation in Phase 3 is much less.

Figure #x. Time taken to transfer a sequence of characters and decompressing using the Load Unbalanced Algorithm, and the Load Balanced Algorithm (LB) versus transferring the uncompressed index.

Figure #x. Performance analysis of the different phases of the algorithm. Notice that only Phases 2, 4, and 5 were the most significant.

== A small improvement to Phase 5 ==

The fifth phase of the Load balanced algorithm involves a lot of read/write operations. The read operations could be accelerated by bringing the array of symbols in texture memory. This would have the effect of caching this constant array and thus it would improve performance.

Figure #x. Performance of Phase 5 with texture memory (P5T), and without it (P5). Notice that there is a small improvement to Phase 5 when using Texture memory.

== Performance in GPU with Computability 1.3 ==

Tests were also performed in a GeForceGTX285 NVidia GPU with 240 cores and 1 GB of RAM memory. This GPU has a higher computability, greater number of cores, faster cores, and more memory. The original test using the data distribution of sequentially incremented attribute values was also used, but this time the attribute values were integers. The outcome of this test was different, both Load Unbalanced (LU) and Load Balanced (LB) algorithms were faster approaches to make the index available rather than transferring the uncompressed projection index (see figure #x). Furthermore, it was noticed that the Load Balanced algorithm was more efficient than the load unbalanced approach for this set of data. The data is not friendly to the Load Unbalanced algorithm because the last element will repeat itself more than all the other, thus the work is not well distributed among all threads (see figure #x).



Figure #x. #TODO

Figure #x. Comparison between the Load Balanced and Load Unbalanced Algorithm, as the number of elements in the uncompressed index is increased.

The test was also performed by maintaining the size of the compressed index constant while doubling the frequency of each element on every iteration. This way, the threads in the load unbalanced algorithm would have the same amount of work. At first the Load Unbalanced algorithm was faster than the Load Balanced algorithm, but as the amount of repetitions increased the Load Balanced algorithm took less time to decompress. The Load Balanced algorithm (LB) was only influenced by the increasing size of the uncompressed array within the GPU. On the other hand, the Load Unbalanced algorithm (LU) not only is slowed down by the fact that it has to write on a bigger array, but it also does not distribute the work among all the possible threads. Since there are only 1024 different elements, only 1024 threads do work at a time. The LU algorithm also cannot make as good use of the GPU as the LB algorithm when the elements have too many repetitions. Sending the uncompressed index (UC) was only the better option for the first cases up to where the frequency was eight. At this frequency the sequence was still fairly small with 8192 elements in its uncompressed form. Notice also that there was no compression for a frequency of only two (See Table #x).



Finally, a test using the third data distribution was performed on the two algorithms. For this test the size of the index was fixed to 16777216 elements in uncompressed form, and the number of different elements in the index was varied, achieving different levels of compression. In other words, the fewer different elements the more compression achieved and the more different elements the lesser the compression. Using the previous experiment as reference, the amount of time taken to transfer the uncompressed index without an algorithm was taken to be 335.80 milliseconds. The outcome of this experiment shows that obviously when no compression is achieved, it is best to send the uncompressed index without decompressing. The compression percentage measure in our results is taken as the percentage in size of the compressed index as a fraction of the original uncompressed index. In the test results, notice that if the compression is of at least 50% on a large index, such as this one, both Load Balanced and Load Unbalanced algorithms performed better than sending the uncompressed index (See Table #x).



On a deeper look, one can also notice that the speedup for the load balanced algorithm is steady as the index becomes more and more compressed. The speed up for the load unbalanced algorithm, however, is very inconsistent. This is due to two competing factors: on is the decreasing time to copy a more compressed index on each iteration, and the other is the increasing time threads take by having less different elements that are repeated more frequently.

/\*On a deeper look, the amount of compression achieved with the first and second data distributions is very high extremely high; the data is reduced to in the lowest iteration 0.79% to 0.079% of the original data. However a lot of the real world data will not allow itself to be compressed that much. The second distribution was also \*/

== Conclusions and Future work ==

Many things were concluded for this work:

From the results, it was determined that the performance of the algorithms depends mostly on the GPU used and its computability, the percentage of compression of the index, and the nature of the data being processed. Both algorithms performed badly on the GPU with computability of 1.1, and performed well on the GPU with computability 1.2.

The load balanced algorithm was mostly limited by its fourth phase in which an inclusive-scan is performed to obtain the position of corresponding elements in the decompressed index

The load balanced and unbalanced algorithms both are dependent on how well compressed is a projection index. This problem is inherent from the RLE.

One possible avenue for future work is to compare the GPU's decompression against the CPU, as it may not be a good algorithm for transferring an index quickly in the GPU, but it may be a good way to perform decompression when using a GPU.

It may be noticed that one of the problems that was not addressed was the fact that it is possible to have a compressed index that will not fit in the GPUs memory once it is uncompressed within the GPU.

/\* Much analysis may still be done on the data as the percentage of \*/

/\* There are other possible performance enhancements related to use of different memories \*/

/\*Compressibility of the projection index matters very much\*/

To conclude the project a benchmark test will compare and find the cases where a compressed index can be more readily available to the GPU by uncompressing as opposed to loading it as an uncompressed index.

Projection index with 1024 different elements and then doubles the amount of elements.

Projection index with fixed size of elements and then increasing the number of different elements from 2 different to having all elements with a frequency of 3.