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#!/usr/bin/env/stack
-- stack --install-ghc runghc

-- CS 381 HW 3
-- Prof Erwig

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main :: IO()
main = putStrLn("Program Executed.")

--exercise 1

--- part a ---

data Cmd = LD Int | ADD | MULT
          | DUP | INC | SWAP
          | POP Int
          deriving Show

type Stack = [Int]

type Prog = [Cmd]

type D = Stack -> Stack

sem :: Prog -> D
sem [] a = a
sem (x:xs) a = sem xs (semCmd x a)

semCmd :: Cmd -> D
semCmd (LD a) xs = ([a] ++ xs)

type Rank = Int
type CmdRank = (Int, Int)

rankC :: Cmd -> CmdRank
rankC (LD a) = (0, 1)
rankC ADD = (2, 1)
rankC MULT = (2, 1)
rankC DUP = (1, 2)
rankC INC = (1, 1)
rankC SWAP = (2, 1)
rankC (POP a) = (a, 0)

rankP :: Prog -> Maybe Rank
rankP xs = rank xs 0

rank :: Prog -> Rank -> Maybe Rank
rank [] s | s >= 0 = Just s

```

```
rank (x:xs) s | under >= 0 = rank xs (under+push)
              where (pop, push) = rankC x
                  under         = s - pop
rank _ _ = Nothing
```

---- Part b ----

```
data Type = S Stack | TypeError deriving Show
```

```
typeSafe :: Prog -> Bool
typeSafe p = (rankP p) /= Nothing
```

```
semStatTC :: Prog -> Type
semStatTC p | typeSafe p = S(sem p [])
            | otherwise   = TypeError
```

-- Part 2 --

```
data Shape = X
            | TD Shape Shape
            | LR Shape Shape
            deriving Show
```

```
type BBox = (Int, Int)
```

--- Define a type checker for the shape language ---

```
bbox :: Shape -> BBox
bbox (TD i j) -- width is that of the wider one; height is sum of heights
  | ix >= jx = (ix, iy + jy)
  | ix < jx = (jx, iy + jy)
  where (ix, iy) = bbox i
        (jx, jy) = bbox j
bbox (LR i j) -- width is sum of widths; height is that of the taller one
  | iy >= jy = (ix + jx, iy)
  | iy < jy = (ix + jx, jy)
  where (ix, iy) = bbox i
        (jx, jy) = bbox j
bbox X = (1, 1)
```

```
rect :: Shape -> Maybe BBox
rect X = Just(1, 1)
rect (TD i j) =
  case rect i of
    Nothing -> Nothing          --- If i is not a rectangle then return nothing ---
    Just (ix, iy) -> case rect j of --- If i is a rectangle then check j ---
      Nothing -> Nothing
      Just (jx, jy) -> case (ix == jx) of --- If j is a rectangle check if
        i is the same width ---
          True -> Just (ix, iy + jy) --- If it is same width,
            place j on top of i. ---
```

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False -> Nothing

rect (LR i j) =
  case rect i of
    -- with width the sum of widths. Else Nothing.
    Nothing -> Nothing
    Just (ix, iy) -> case rect j of
      Nothing -> Nothing
      Just (jx, jy) -> case (iy == jy) of
        True -> Just (ix + jx, iy)
        False -> Nothing

-- Part 3 --

-- 1 --
-- a. What types are f and g?
-- Function f is of type [a] -> [a] -> [a]
-- Function g is of type [a] -> [b] -> [b]

-- b. Why do the functions have these types?
-- f has this type because it takes two lists and outputs one, and both lists must have the same
-- type because either x or y can be returned and a Haskell function can have at most 1 return
type.
-- g has this type because it can return either an untyped list (empty list) or a list of type y.
-- the type of x does not matter because it will not be returned no matter what.

-- c. Which type is more general?
-- g has a more general type, because there are less restrictions. in f, all three lists must
be of the same type,
-- but in g, the type of x doesnt matter, allowing for more flexibility

-- d. Why do f and g have different types?
-- They need to. Although any valid input to f can also be a valid input to g, we still must
specify the
-- restrictions of f. Mostly this means denoting that all list types in f must be the same, but
the only
-- two that need to be the same in g are parameters 2 and 3.

-- 2 --

h :: [b] -> [(a,b)] -> [b]
h b (x:xs) = [snd(x)] ++ b;
h b _ = b

-- 3 --
--k :: (a -> b) -> ((a -> b) -> a) -> b    --Order of parentheses should not matter, I think?
--k f x = func(f x)

-- 4 --

-- No, you cannot make this definition for a function.
-- To convert one type to another, you must know all possible values of the type you are
converting to, in this case b.
-- Unless you know all those values and develop case statements that only allow valid output,

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you will end

-- up restricting type b with some operation related to a. For example, if the function added 5 to everything, then a would

-- have to be an numerical type, but then so would b because by adding to a, you restrict the output, b, to numerical types.

-- In any event, b would be the same type as a and the definition would fail, as it stipulates a and b are different types.