



Data Structures and Algorithms



Week-2

Greedy Algorithms

Greedy is an algorithmic paradigm that builds up a solution piece by piece, always choosing the next piece that offers the most obvious and immediate benefit.

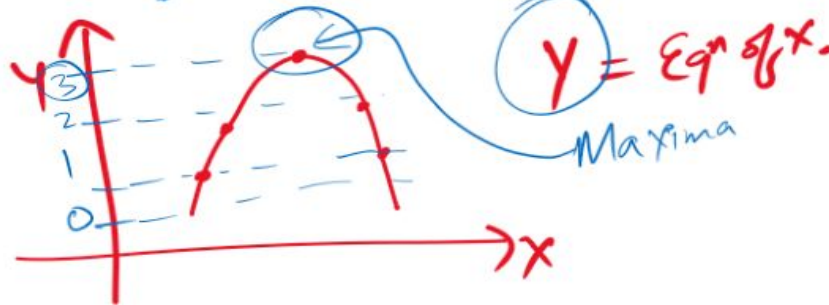
Used to solve Optimisation Problems

Examples: Fractional Knapsack, Maximum Number Problem, Car Fueling , etc

Optimisation Problems and the ways to solve them:

Optimisation Problems

→ Finding best solution



Ways of Solving

- ① Greedy Algo (Method)
- ② Branch & Bound
- ③ Dynamic Programming

Greedy Strategy:

- Make a Greedy choice
- Reduce to a subproblem
- Iterate

Links to develop Intuition: <http://dm.compsicclub.ru/app/quiz-largest-number>

What is the Greedy Choice here? Choosing the largest number OR choosing the number such that leftmost position should be the greatest one possible?

Largest Concatenate Problem

Compile the largest number by concatenating the given numbers.

Level 1

Level 2

Level 3

Level 4

2

3

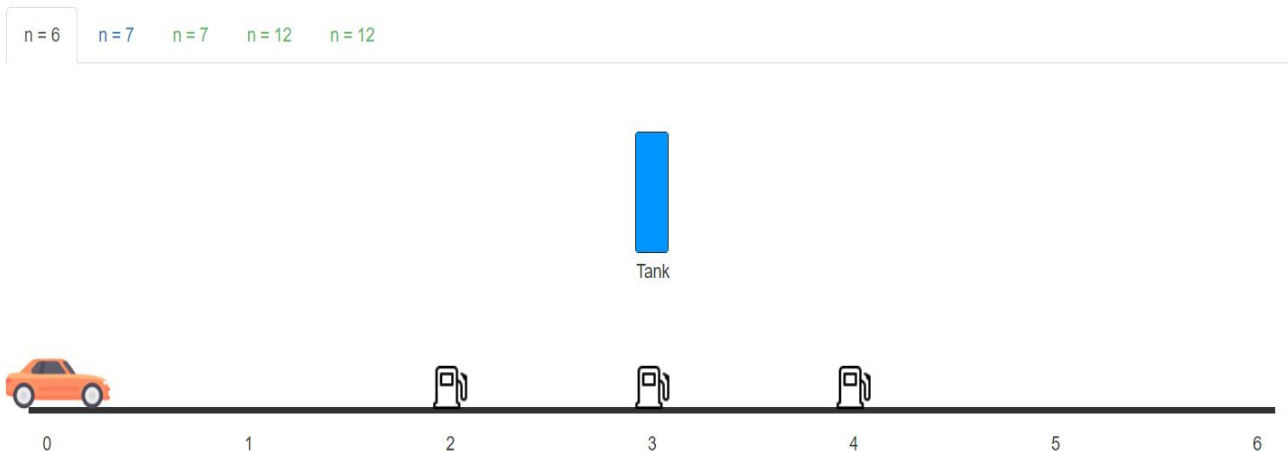
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Links to develop Intuition:

<http://dm.compsclub.ru/app/quiz-car-fueling>

Car Fueling

A car can travel at most 3 miles on a full tank. You want to make as few refills as possible while getting from A to B. Select gas stations where you would like to refuel and press "Start journey" button.



Greedy Choice here?

1. Refill at the closest gas station.
2. Refill at the farthest reachable gas station
3. Go until there is no fuel.

So option 2 is correct.

Problem in the code format

Car Fueling

Input: A car which can travel at most L kilometers with full tank, a source point A , a destination point B and n gas stations at distances $x_1 \leq x_2 \leq x_3 \leq \dots \leq x_n$ in kilometers from A along the path from A to B .

Output: The minimum number of refills to get from A to B , besides refill at A .

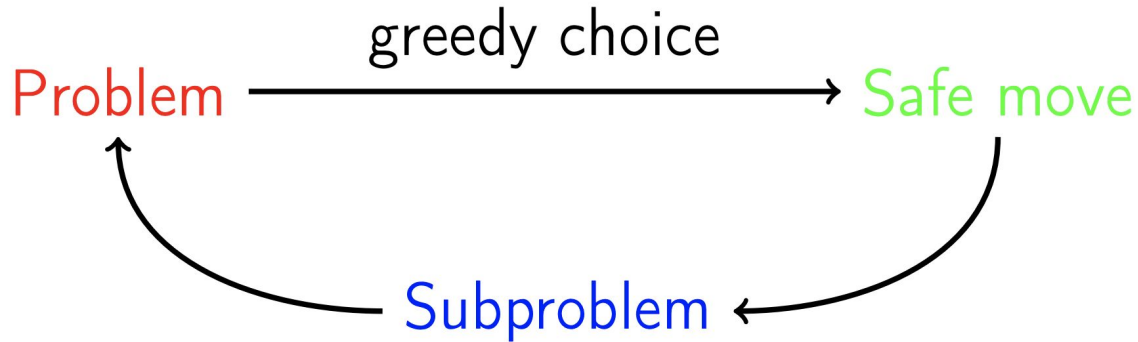
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$$A = x_0 \leq x_1 \leq x_2 \leq \dots \leq x_n \leq x_{n+1} = B$$

MinRefills(*x*, *n*, *L*)

```
numRefills  $\leftarrow$  0, currentRefill  $\leftarrow$  0
while currentRefill  $\leq$  n:
    lastRefill  $\leftarrow$  currentRefill
    while (currentRefill  $\leq$  n and
            $x[\textit{currentRefill} + 1] - x[\textit{lastRefill}] \leq L$ ):
        currentRefill  $\leftarrow$  currentRefill + 1
    if currentRefill == lastRefill:
        return IMPOSSIBLE
    if currentRefill  $\leq$  n:
        numRefills  $\leftarrow$  numRefills + 1
return numRefills
```


General Strategy



- Make a greedy choice
- **Prove** that it is a **safe move**
- Reduce to a **subproblem**
- Solve the **subproblem**

Fractional Knapsack

