

Ascending order of General Time Complexities

In []: Ascending order of General Time Complexities --> $O(1)$, $O(\log n)$, $O(n)$, $O(n\log n)$, $O(n^2)$, $O(n^3)$, $O(2^n)$

Python Program to find the sum of N natural Number

```
In [1]: n=int(input())
sum=0
for i in range(0,n+1):
    sum=sum+i
print(sum)

#Time Complexity for the above code is O(n) which is more we can also optimize it

10
55
```

Optimized Code

```
In [ ]: n=int(input())
sum=(n*(n+1))//2    #O(1)
print(sum)

#Time Complexity of this code is O(1) which is the least one and furthur we cannot optimized it.
```

Searching

In []: --> The process of finding the desired information **from** the set of items stored **in** the form of elements **in** the computer memory **is** referred to **as** '**searching in data structure**'.

Types of Searching Algorithms

In []: Two types of searching algorithms:
1.Linear Search --> $O(N)$ --> Sorted **as** well **as** Unsorted
2.Binary Search --> $O(\log N)$ --> Sorted Array

Linear Search Algorithm

In []: --> Linear search **is** the simplest method **for** searching.
--> In Linear search technique of searching; the element to be found **in** searching the elements to be found **is** searched sequentially **in** the list.
--> This method can be performed on a sorted **or** an unsorted list (usually arrays).

Implementation of Linear Search Algorithm

```
In [2]: def linear_search(array,key):    #5
        for i in range(len(array)):    #O(n)
            if array[i]==key:
                print("Element found at index :",i)
                break
            else:
                return "Element not found"

array=[10,20,30,40,50]
key=20
linear_search(array,key)

Element found at index : 1
```

Discussion on Time Complexity of Linear Search

Best Case(Big Omega) of Linear Search

In []: --> The element being searched could be found **in** the first position.
--> In this case, the search ends **with** a single successful comparison.
--> Thus, **in** the best-case scenario, the linear search algorithm performs $O(1)$ operations.

Average Case(Big Theta) of Linear Search

In []: --> When the element to be searched **is in** the middle of the array, the average case of the Linear Search Algorithm **is** $O(n)$.

Worst Case(Big O) of Linear Search

In []: --> The element being searched may be at the last position **in** the array **or not** at all.
--> In the first case, the search succeeds **in** 'n' comparisons.
--> In the next case, the search fails after 'n' comparisons.
--> Thus, **in** the worst-case scenario, the linear search algorithm performs $O(n)$ operations.

Binary Search Algorithm

In []: --> Binary searches are efficient algorithms based on the concept of "**divide and conquer**" that improves the search by recursively dividing the array **in** half until you either find the element **or** the list gets narrowed down to one piece that doesn't match the needed element.

--> Binary searches work under the principle of using the sorted information **in** the array to reduce the time complexity to zero ($\log n$).

The binary search approach's basic steps:

In []: --> Sort the array **in** ascending order.
--> Set the low index to the first element of the array **and** the high index to the last element.
--> Set the middle index to the average of the low **and** high indices.
--> If the element at the middle index **is** the target element, **return** the middle index.
--> If the target element **is** less than the element at the middle index, set the high index to the middle index - 1.
--> If the target element **is** greater than the element at the middle index, set the low index to the middle index + 1.
--> Repeat steps 3-6 until the element **is** found **or** it **is** clear that the element **is not** present **in** the array.

Implementation of Binary Search

```
In [4]: def binary_search(array,key):
        low=0
        high = len(array)-1
        mid=0
        while low<=high:
            mid=(high+low)//2
            if array[mid]<key:
                low=mid+1
            elif array[mid]>key:
                high=mid-1
            else:
                return mid
        return -1

array = [1,2,3,4,5,6]
key=6
x= binary_search(array,key)
if x!= -1:
    print("Element present at index :",x)
else:
    print("Element not present")

Element present at index : 5
```

Analysis of input size at each iteration of Binary Search:

In []: At Iteration 1:

Length of array = n

At Iteration 2:

Length of array = $n/2$

At Iteration 3:

Length of array = $(n/2)/2 = n/2^2$

Therefore, after Iteration k:

Length of array = $n/2^k$

Also, we know that after k iterations, the length of the array becomes 1 Therefore, the Length of the array

 $n/2^k = 1$
 $\Rightarrow n = 2^k$

Applying log function on both sides:

 $\Rightarrow \log 2n = \log 2^k$

 $\Rightarrow \log 2n = k * \log 2$

As $(\log_2 (a) = 1)$ Therefore, $k = \log_2(n)$

Discussion of time Complexity of Binary Search

Best Case Time Complexity(Big Omega)

In []: --> The best time complexity of binary search occurs when the required element **is** found **in** the first comparison itself, **and** only one iteration occurs. Therefore we use $O(1)$.
--> Essentially **for** this case, the element needs to be **in** the exact middle of the list because, **in** binary search, the first comparison occurs **with** the middle element. Once the middle element does **not return** the correct answer, the iteration begins **for** the lesser half of the greater half.

Worst Case Time Complexity(Big O)

In []: --> The worst time complexity of binary search occurs when the element **is** found **in** the very first index **or** the very last index of the array. For this scenario, the number of comparisons **and** iterations required **is** $\log n$, where n **is** the number of elements **in** the array. It **is** called the worst time complexity because it consumes a lot of time **for** large arrays containing hundreds **and** thousands of values. Accordingly, hundreds **and** thousands of iterations must occur.

Average Case Time Complexity(Big Theta)

In []: The average case Time Complexity **in** Binary Search **is** $n * \log n / (n+1)$, which **is** approximately $\log n$.

Summary of Data Structures and Algorithms

In []: Data Structures: **is** a way of storing the data so that we can use that data **in** an efficient manner **for** our further use.

Types of Data Structures:

1. Primitive Data Structure(Hold a single value) --> int,float,long,bool complex.....
2. Non Primitive Data Structure(Hold multiple values)

--> Based on Implementation

- > Physical --> Array linkedlist
- > logical(ADT) --> Stack **and** Queue, Graph **and** Trees

--> Based on Storage

- > Linear Data Structures --> Array, linkedlist, Stack **and** Queue
- > Non Linear Data Structures --> Trees **and** Graphs

Linkedlist: Linkedlist **is** a linear data structure. It **is** collection of nodes **and** each node **is** having data **and** next (Memory location of next element)

Time Complexity of Linkedlist:

- accessing any element of linkedlist --> $O(n)$
- Traversing over linkedlist --> $O(N)$
- Insertion at first position --> $O(1)$
- Insertion at given position --> $O(N)$
- Insertion at last --> $O(N)$
- Deletion at first position --> $O(1)$
- Deletion at given position --> $O(N)$
- Deletion at last --> $O(N)$

Operations on linkedlist:

- 1.Insertion --> Insertion at begin , Insertion at end , Insertion at given pos
- 2.Deletion --> Deletion at begin , Deletion at end , Deletion at given pos
- 3.Traversal

Array: Array **is** a linear Data Structure **in** which elements are stored at contiguous memory location.

- Accessing any element of array --> $O(1)$
- Traversing over array --> $O(N)$
- Insertion at first position --> $O(1)$
- Insertion at given position --> $O(N)$
- Insertion at last --> $O(N)$
- Deletion at first pos --> $O(1)$
- Deletion at given pos --> $O(N)$
- deletion at last --> $O(n)$

Operations on Array:

- 1.Insertion --> Insertion at begin , Insertion at end , Insertion at given pos
- 2.Deletion --> Deletion at begin , Deletion at end , Deletion at given pos
- 3.Traversal

Stack : Stack **is** a linear data structure which **is** following the principle of LIFO(Last **in** first out)
: Insertion **and** Deletion **in** stack will be done **from** one end(Top of stack)

Terminologies **and** Operations of stack:

- 1.Push --> insertion at top
- 2.Pop --> Deletion **from** top
- 3.Peek --> Getting Top Value

Time Complexities of stack:

- push --> $O(1)$ --> $O(N)$
- pop --> $O(1)$
- peek --> $O(1)$

Queue : Queue **is** a linear data structure which is following the principle of FIFO(First **in** first out).
: Insertion will be done **from** Rear end **and** deletion will be done **from** front end

Terminologies **and** Operations of queue:

- 1.Enqueue --> insertion at Rear --> $O(1)$
- 2.Dequeue --> Deletion **from** Front --> $O(1)$
- 3.PeekinQueue --> Getting front Value

Time Complexity of Queue:

- 1.Enqueue --> $O(1)$
- 2.Dequeue --> $O(1)$
- 3.PeekinQueue --> $O(1)$

Never ending Queue(Circular Queue) -->By linkedlist only