

```
fun append (xs,ys) =  
  if xs=[]  
  then ys  
  else (hd xs)::append(tl xs,ys)  
  
fun map (f,xs) =  
  case xs of  
    [] => []  
  | x::xs' => (f x)::(map(f,xs'))  
  
val a = map (increment, [4,8,12,16])  
val b = map (hd, [[8,6],[7,5],[3,0,9]])
```

Programming Languages

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2013

The Subtype Relation

Keeping subtyping separate

A programming language already has a lot of typing rules and we do not want to change them

- Example: The type of an actual function argument must ***equal*** the type of the function parameter

We can do this by adding “just two things to our language”

- *Subtyping*: Write $t1 <: t2$ for $t1$ is a subtype of $t2$
- One new typing rule that uses subtyping:
If e has type $t1$ and $t1 <: t2$,
then e (also) has type $t2$

Now all we need to do is define $t1 <: t2$

Subtyping is not a matter of opinion

- Misconception: If we are making a new language, we can have whatever typing and subtyping rules we want
- Not if you want to prevent what you claim to prevent [soundness]
 - Here: No accessing record fields that do not exist
- Our typing rules were *sound* before we added subtyping
 - We should keep it that way
- Principle of *substitutability*: If $\tau_1 <: \tau_2$, then any value of type τ_1 must be usable in every way a τ_2 is
 - Here: Any value of subtype needs all fields any value of supertype has

Four good rules

For our record types, these rules all meet the substitutability test:

1. “Width” subtyping: A supertype can have a subset of fields with the same types
2. “Permutation” subtyping: A supertype can have the same set of fields with the same types in a different order
3. Transitivity: If $t1 <: t2$ and $t2 <: t3$, then $t1 <: t3$
4. Reflexivity: Every type is a subtype of itself

(4) may seem unnecessary, but it composes well with other rules in a full language and “does no harm”