

```
fun append (xs,ys) =  
  if xs=[]  
  then ys  
  else (hd xs)::append(tl xs,ys)  
  
fun map (f,xs) =  
  case xs of  
    [] => []  
  | x::xs' => (f x)::(map(f,xs'))  
  
val a = map (increment, [4,8,12,16])  
val b = map (hd, [[8,6],[7,5],[3,0,9]])
```

# Programming Languages

Dan Grossman  
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*Pattern-Matching So Far: Precisely*

# Careful definitions

When a language construct is “new and strange,” there is *more* reason to define the evaluation rules precisely...

... so let's review datatype bindings and case expressions “so far”

- *Extensions* to come but won't invalidate the “so far”

# *Datatype bindings*

```
datatype t = C1 of t1 | C2 of t2 | ... | Cn of tn
```

Adds type  $t$  and constructors  $C_i$  of type  $t_i \rightarrow t$

- $C_i \ v$  is a value, i.e., the result “includes the tag”

Omit “of  $t$ ” for constructors that are just tags, no underlying data

- Such a  $C_i$  is a value of type  $t$

Given an expression of type  $t$ , use *case expressions* to:

- See which variant (tag) it has
- Extract underlying data once you know which variant

# *Datatype bindings*

```
case e of p1 => e1 | p2 => e2 | ... | pn => en
```

- As usual, can use a case expressions anywhere an expression goes
  - Does not need to be whole function body, but often is
- Evaluate **e** to a value, call it **v**
- If **p<sub>i</sub>** is the first *pattern* to *match* **v**, then result is evaluation of **e<sub>i</sub>** in environment “extended by the match”
- Pattern **C<sub>i</sub> (x<sub>1</sub>, ..., x<sub>n</sub>)** matches value **C<sub>i</sub> (v<sub>1</sub>, ..., v<sub>n</sub>)** and extends the environment with **x<sub>1</sub>** to **v<sub>1</sub>** ... **x<sub>n</sub>** to **v<sub>n</sub>**
  - For “no data” constructors, pattern **C<sub>i</sub>** matches value **C<sub>i</sub>**