ALPHABET I STRINGS on I LANGUAGE: set of strings on I

$$L_{\varepsilon} = \{\varepsilon\}$$

$$L_{\infty} = \{\xi\}$$

$$L_{\alpha} = \{\alpha\} \text{ for any } \alpha \in \Sigma$$

$$L_{\Sigma} = \{\alpha \mid \alpha \in \Sigma\}$$

{"a" | 'a' \ \]

L, ELz are languages, L, ULz bethe normal set-theoretic union. {abc, ab} v {ca, cd} = {abc, ab, ca, cd}

L, L2 is the concatenation of L, and L2 {aB|aEL, ^BEL2}

{ab, ac}{cd, be} = {abcd, abbe, accd, acbe}

EXPONENTI ATION k times {aaa, aab, aba, abb, baa, bab, bba, bbb} {a,b} {aaa} 9033 10 = {E} 1' = L KLEENE CLOSURE

 $L^* = L^0 U L^1 U L^2 U L^3 U \dots$ $\{a\}^* = \{\epsilon, a, aa, aaa, aaaa, \dots\}$

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REGULAR LANGUAGE:
Start with La aEZi
apply union (alternation), concatenation, Kleene closure
         -> Regular Language
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Regular Expressions
R.E. Language
a Lauls «B
     Lalp
2B
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-use parentheses for grouping

$$a(b+c) = \{ab,ac\}$$

$$a(b+c)^2 = \{abb,abc,acb,acc\}$$

$$a(bc+c) = \{abc,ac\}$$

$$\{a\}\{\{b\}\{c\}\}\}\{c\}\}$$

$$\{a\}\{\{bc\}\}\{c\}\}$$

$$\{a\}\{\{bc,c\}\}$$

 $ab* = \{a, ab, abb, abbb, \dots\}$

A Regular Expression & matches a string s if s is in the language denoted by &.

can also say s matches ox. ab* matches abb

Use regular expressions in many places.
-editors

- -shells -lexical analyzers, grammars.

Notational Extensions

"optional 2" d? (8+2) "one or more 2" とく* \propto^{+} (a|b|-...12)

every character except a, b, c. [abdf] [ab...2] "abc] (E, x, x Ba, a BaBa, ...) "B-seperated list of Dor more [^abc] & MB

floating literal:

[0.9] [0.9] ((e|E)(+|-)[0.9])?

Oncormore digits. one or more digits.

179

3.9 109.725

2.3e+21 34.22E-7

optional part!

an'e', sign, at least one digit.

E' tor