

Question 1

//bubble sort

```

Algorithm beautiful(A, n)
    for int i <- 0 to n do
        for int j <- 0 to n-i do
            temp <- arr[j]
            arr[j] <- arr[j+1]
            arr[j+1] <- temp

```

In every case, in bubble sort, we still have to traverse through all the elements twice for a loop. So that is n^2 complexity for best and worst-case same.

Question 2

ascending: 2^n , $2^{(n+1)}$, $2^{(2n)}$, $2^{(2^n)}$

Question 3

$O(1)$ - add element to array
 $O(\log n)$ - binary search
 $O(n)$ - find min element
 $O(n \log n)$ - merge sort
 $O(n^2)$ - bubble sort
 $O(n^3)$ - 3 variables equation solver
 $O(2^n)$ - Find all subsets

Question 4

We can't use the master theorem for Fibonacci n . So in this case we can use Counting self-calls:

$S(0) = 0$, $S(1) = 0$

So the self calls of Fibonacci n will be:

$S(n) = 2 + S(n-1) + S(n-2)$

Question 5

$T(1) = 1$

$T(n) = 2T(n/2) + c$

$N = 2^k$

Counting self-calls:

$S(n^k) = 2 + S(n^{k-1}) + S(n^{k-2})$

Question 5-2

Master theorem examples:

1) $T(n) = T + (2/2) + cn^2$

$a = 1$ $b = 2$ $c = 2$

$1 < 2^{1/2}$: That means the time complexity will be: $\Theta(n^k)$

2) $T(n) = 2T + (n/2) + n$

$a = 2$ $b = 2$ $c = 2$

$2 < 2^{1/2}$: That means the time complexity will be: $\Theta(n^k)$

3) $T(n) = 2T + (2/2) + n$
 $a = 2$ $b = 4$, $d = 1/2$

$2 = 4^{(1/2)}$: That means the time complexity will be: $\Theta(n^k \log n)$