# Part 2: JSON-Driven Domain Models in Practice

# Triangleness, Quantum Fields, and Beyond

Part 1: From Bits to Qubits with CMCC:Demonstrating Computational Universality through Triangles, Quantum Walks, the Ruliad and Multiway Systems

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### **Abstract**

Building upon the theoretical foundations established in Part 1, this second paper demonstrates the practical utility of CMCC through a straightforward, JSON-based rulebook approach. By encoding all conceptual knowledge in the same five declarative primitives and employing JSON as a universal, machine-readable mirror, developers can automatically generate consistent, cross-language implementations without resorting to specialized domain-specific languages or duplicated logic. We illustrate this practical approach through two distinct yet related examples: a geometric scenario defining and verifying properties of triangles, and a quantum simulation representing quantum walks and double-slit experiments.

Using template-based transformations (via Handlebars) from the JSON rulebook, we generate statically typed helper code across multiple programming environments, including Python, Golang, and Java. We demonstrate that the same declarative JSON file can robustly and simultaneously capture simple polygons as well as complex quantum phenomena, ensuring domain fidelity across diverse languages with minimal effort. This results in a robust, auditable, and easily maintainable system for modeling and evolving complex systems across domains. The paper closes by addressing performance considerations, scalability challenges, and schema evolution, reinforcing CMCC's potential as a powerful, practical, and universal declarative framework for computational modeling.

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### 1. Introduction

### 1.1 Context and Relationship to Part 1

In Part 1—From Bits to Qubits with CMCC—we established the Conceptual Model Completeness Conjecture (CMCC), showing how any finite computable concept can be represented purely via Schema (S), Data (D), Lookups (L), Aggregations (A), and Lambda Calculated Fields (F) in a snapshot-consistent (or ACID-compliant) environment. We explored emergent complexity in examples spanning geometry (triangleness) and quantum mechanics (quantum walks), along with theoretical ties to Wolfram's Multiway Systems and Wheeler's "It from Bit."

This second paper (**Part 2**) focuses on **practical implementation**: how you can store your entire CMCC rulebook in JSON, automatically generate cross-language code (e.g., Python, Golang, Java), and manage domain changes over time without re-coding every rule. Readers seeking the deeper theoretical underpinnings—Turing-completeness proofs, philosophical arguments, or Wolfram's multiway parallels—should refer back to **Part 1**.

### 1.2 Objectives: Hands-On CMCC Implementation

Here, we aim to demonstrate how:

- 1. A **single JSON rulebook** captures all relevant domain definitions—fields, relationships, aggregator formulas—using minimal, standardized syntax.
- 2. You can **auto-generate typed stubs** in Python, Golang, or Java (or your language of choice) via a simple template engine.
- 3. A small, **minimal imperative script** in each language suffices to run scenarios (e.g., verifying right triangles or simulating quantum interference).
- 4. **Updating** the JSON definitions seamlessly re-syncs logic across all language environments, eliminating DSL duplication or specialized parser overhead.

By the end, you'll see that the same approach supports both a straightforward geometry use case (triangleness) and a more advanced quantum scenario. This further underscores CMCC's promise: once domain logic is declared structurally, it can be shared, re-targeted, or extended with minimal friction.

### 1.3 JSON as a Universal "Rulebook" Format

A key part of CMCC is avoiding domain-specific languages (DSLs). Instead, we treat **JSON** as our universal "mirror" for:

- Machine-Readability: Any modern environment parses JSON out-of-the-box.
- Extensibility: Add aggregator fields or references on the fly.
- No Custom Syntax: You need not define or maintain new grammars—just structural, name-value pairs.

Throughout this paper, we'll detail how you can define everything from polygon edges to quantum amplitudes in JSON, ensuring that all emergent logic remains purely declarative.

#### 1.4 Roadmap of Practical Examples

After a brief recap of CMCC (Section 2) and the code-generation approach (Section 3), we dive into:

- 1. **Triangleness** (Section 4): JSON-based polygons and points, verifying right triangles with auto-generated code.
- 2. **Quantum Field / Double-Slit** (Section 5): Wavefunctions, coin operators, measurement, and interference patterns—again, auto-generated code with minimal hand-written logic.

We then explore advanced formula evaluation, best practices, scaling considerations, and schema evolution (Sections 6 and 7). Finally, we wrap up with next steps and an open invitation to experimentation or falsification.

### 1.5 HARDCORE FALSIFICATION CHECKLIST

As in Part 1, we emphasize that CMCC remains open to rigorous challenge. If you suspect a particular rule cannot be captured in (S, D, L, A, F), we welcome concrete counterexamples:

#### 1. DO NOT LIST 5 POSSIBLE PROBLEMS

(Hypotheticals like "could" or "might" are not real tests.)

#### 2. PICK A SPECIFIC RULE YOU THINK BREAKS CMCC

If it can be written in code or English, specify it clearly.

#### 3. **DECOMPOSE IT INTO CMCC PRIMITIVES**

Show how you tried using schema fields, data rows, lookups, aggregations, or formula fields.

### 4. CHECK IF THE RUNTIME ENGINE CAN HANDLE IT

CMCC only defines "what"—the "how" (e.g., code performance or concurrency) is a separate question.

### 5. IF YOU STILL THINK IT FAILS, TRY AGAIN

Many perceived "show-stoppers" yield to a purely declarative rethink.

**Still convinced you have a real counterexample?** Please reach out. We want to test any claim that a finite computable rule can't be expressed in CMCC.

### 1.6 Source Code and GitHub Repo

All examples in this paper—including fully working JSON files for both triangleness and the quantum walk—are available at:

https://github.com/eejai42/conceptual-model-completeness-conjecture-toe-meta-model

We encourage you to clone the repo, try out the code, tweak the JSON definitions, and see for yourself how easily the system adapts to domain changes.

# 2. Background: Recap of CMCC in Brief

### 2.1 Five Primitives (S, D, L, A, F) in Practice

The **Conceptual Model Completeness Conjecture (CMCC)** asserts that any finite computable domain can be fully encoded with only five declarative primitives—no extra syntax or imperative sidecar code required. These primitives are:

### Schema (S)

Describes what kinds of objects (or entities) exist in your domain and what attributes they have. Think of it as the meta-definition or blueprint.

Example: A "Polygon" schema with fields for "name," "edges," etc. A "Wavefunction" schema with fields for "TimeStep," "psi," etc.

### • Data (D)

Instantiates those schema definitions with actual records or objects.

Example: Multiple "Polygon" records for specific triangles, squares, or pentagons. Multiple
 "Wavefunction" entries for each discrete time step.

### Lookups (L)

Establish relationships between records. Conceptually, this is akin to a database foreign key or a pointer that says, "This object references that object."

Example: An "Edge" record references its two "Point" records as endpoints. A "Wavefunction" record might reference a "Grid" record to locate the amplitude array.

### Aggregations (A)

Summarize or roll up sets of data—like sums, counts, averages, or more sophisticated computations across multiple records.

 Example: Summing the squared magnitudes of amplitudes in a wavefunction to compute total probability. Counting the edges that belong to a polygon.

### Lambda Calculated Fields (F)

Encode computed logic or constraint checks in a purely declarative manner.

 $\circ$  Example: A formula for the length of an edge (x2-x1)2+(y2-y1)2\sqrt{(x\_2 - x\_1)^2 + (y\_2 - y\_1)^2}(x2-x1)2+(y2-y1)2. A unitarity check CoinOperator\*CoinOperator\*=I\text{CoinOperator} \times \text{CoinOperator}^\dagger = ICoinOperator\*CoinOperator\*=I.

By chaining these five primitives, you can represent any typical "imperative step" as a series of structural transformations: each new snapshot of data reflects the emergent outcome of aggregator fields and formulas. This eliminates the need to embed domain rules in a custom scripting or programming language.

#### 2.1.1 Additional Clarifications

- **Rulebook vs. Runtime**: As noted in Section 1.6, the five primitives form the *rulebook* of *what* must be true. *How* these rules get executed (e.g., in code or physical processes) is a separate concern.
- Formula Syntax and Adaptability: The potentially verbose function names (like SQRT and SUBTRACT) can be automatically mapped to any language or library call. See also Section 1.7 on how the entire model remains machine-readable, making it simple to adapt formulas across different toolchains.

### 2.2 Why We Don't Need a DSL

It's common for organizations or researchers to create specialized **Domain-Specific Languages (DSLs)** to describe complex domains—be it geometry, quantum simulations, or business logic. However, DSLs introduce extra overhead:

- Syntax Definition: You first have to invent the grammar or specialized keywords for your DSL.
- **Implementation**: You then parse or compile the DSL into something machine-executable.
- **Maintenance**: Over time, the DSL must evolve, introducing "DSL drift" if it's not perfectly synced with the domain.

Under CMCC, your data and rules live in standard, widely supported structures (e.g., JSON or relational tables) with no separate textual syntax. This "no-DSL" approach means:

- You immediately have a machine-readable format (JSON).
- Any code generation or runtime usage can parse JSON with off-the-shelf libraries.
- Updating a rule means editing the JSON—no re-tooling of a language specification.

### 2.2.1 Cross-Referencing the "Mirror" Concept

As described in Section 1.3, JSON acts as the universal mirror—storing the entire conceptual model. Because JSON is ubiquitous, we avoid writing or maintaining specialized DSLs that do nothing more than re-describe the same concepts.

### 2.2.2 Practical Tools and Extensibility

Later, in Section 3.5, we highlight the variety of simple transformations (e.g., json-to-xml, json-hbars-transform) that further reduce the friction of adopting CMCC for real projects. This set of plug-and-play tools accomplishes most tasks that DSL-based workflows might otherwise require.

### 2.3 Related Work & Prior Foundations

The present paper continues a broader research program on the **Conceptual Model Completeness Conjecture (CMCC)** and its applied variants. Over the past few years, we have published several papers that established:

### 1. Turing-Completeness & Theoretical Basis

- BRCC-Proof [1] introduces the Business Rule Completeness Conjecture (BRCC) and provides a proof sketch inspired by Turing-completeness arguments.
- CMCC: A Universal Declarative Computational Framework [2] formalizes CMCC's five primitives (Schema, Data, Lookups, Aggregations, and Calculated Fields) and outlines its alignment with well-known universal models (lambda calculus, cellular automata).

#### 2. Logical & Philosophical Extensions

 Formalizing Gödel's Incompleteness... [3] and Computational Paradoxes... [4] demonstrate how Gödelian self-reference and classical paradoxes fit into a fully declarative environment—showing that even purely declarative systems inherit certain logical constraints.

### 3. Domain-Specific Explorations

- Quantum CMCC [5] extends these ideas into quantum mechanics at design-time, emphasizing that the domain "rules" can be captured declaratively while actual quantum evolution remains an external runtime.
- CMCC-Driven Graph Isomorphism [6] reframes graph matching through CMCC's aggregator/calc logic, tackling structural and semantic alignment simultaneously.
- Applying CMCC to Model Theory... [7] illustrates how advanced mathematics (e.g., Zilber's pseudo-exponential fields) can be partially encoded via the same five declarative primitives.

Collectively, these works demonstrate that CMCC and its business-focused variant (BRCC) offer a unifying theory for capturing *any* computable rule in an Snapshot-Consistent schema—without requiring domain-specific languages or extraneous imperative code. For an in-depth treatment of theoretical foundations, readers can consult the above references.

The original CMCC and BRCC conjectures require strict ACID compliance to ensure foundational rigor, atomicity, and isolation for all rules and domain semantics [CMCC foundational paper ref.].

However, in practical modeling scenarios, particularly within the physics domain presented in this paper, snapshot-consistency is both sufficient and beneficial. This choice simplifies implementations, improves scalability, and reduces overhead, while still maintaining logically coherent domain states. Readers are referred to [CMCC foundational paper ref.] for an in-depth exploration of why ACID compliance is strictly required in foundational modeling scenarios

In contrast, this paper addresses a much more hands-on question: How can the same JSON-based "rulebook" approach be operationalized in day-to-day software practice across Python, Golang, Java, or other environments, with minimal overhead? Rather than revisiting the full theoretical apparatus from prior publications, we will illustrate practical code-generation workflows, highlight cross-language synchronization, and evaluate real-life examples—from simple polygons to quantum-like phenomena.

# 3. Practical Implementation Framework

### 3.1 JSON Domain Definitions

At the heart of our approach lies a single **JSON file** that "knows everything" about your domain. For instance, we might have:

```
"fields": [
        { "name": "edge_id", "type": "string" },
        { "name": "start_point", "type": "lookup", "references": "Point" },
        { "name": "end_point", "type": "lookup", "references": "Point" },
          "name": "length",
          "type": "calculated",
          "formula": "SQRT( POW(SUBTRACT(end_point.x, start_point.x),2) + POW(SUBTRACT(end_point.y,
start_point.y),2) )"
        }
      1
    },
      "name": "Point",
      "fields": [
        { "name": "x", "type": "number" },
        { "name": "y", "type": "number" }
      1
    }
  ]
}
```

- Schema (S): The entities array describes which tables or entity types exist.
- Data (D): In some scenarios, you might embed example or default data in the same JSON, or store it in a second file.
- Lookups (L): Fields like "type": "lookup" link an Edge to a Point.
- Aggregations (A): If you want, for example, a polygon to sum the lengths of its edges, you could define a field in Polygon with "type": "aggregation", "formula": "SUM(edges.length)".
- Calculated Fields (F): The length field in Edge is a perfect example.

You can obviously refine or restructure the JSON to your liking. The key is that it always references the same five primitives behind the scenes. We're *never* writing a DSL or embedding equations in code—these are just structured fields and formulas.

### 3.1.1 Referencing the Rulebook vs. Runtime Distinction

While the JSON stores **what** "an Edge" or "a Polygon" must do, it *does not* specify how to load or manipulate them at runtime. See Section 1.6 for how that interplay works in practice.

# 3.2 Generating Static Helpers with Handlebars

To make this JSON "come alive" in Python, Golang, and Java, we use a tool or script that reads the JSON and applies a Handlebars template to generate code. For example:

### Handlebars snippet (Python) might look like:

```
class {{name}}:
    def __init__(self, {{#each fields}}{{this.name}}{{#unless @last}},{{/unless}}{{/each}}):
        {{#each fields}}
```

```
self.{{this.name}} = {{this.name}}
   {{/each}}

{{#each fields}}
{{#if this.formula}}
@property
def {{this.name}}(self):
    # TODO: convert formula: "{{this.formula}}"
    return evaluate_formula("{{this.formula}}", self)
{{/if}}
{{/each}}
```

This template is then expanded for each "entity" in the JSON, producing Python classes with constructor fields and placeholders for formula evaluation.

• For Golang or Java, you'd have a different template—still referencing the same JSON definitions but generating structs or classes in the respective languages.

When you run:

#### bash

handlebars-codegen --input=domain.json --template=python\_class.hbs --output=domain\_generated.py

• you get a new domain\_generated.py that includes all the entity definitions, references, and so on.

### 3.2.1 Example: Adapting Function Names

If your JSON uses "SQRT(...)" but your target language expects "math.Sqrt(...)," a single pass in the Handlebars template can inject the correct local call. This is precisely why the verbose "SQRT" style is *not* a problem (see also Section 1.7).

# 3.3 Hand-Written Imperative Scripts: Minimal but Sufficient

We **do** need a small amount of code that orchestrates or executes domain logic at runtime—e.g., reading data from CSV or from the JSON, instantiating classes, and iterating a process. But this code is tiny and does **not** contain domain rules. Instead, it:

- 1. Parses the domain JSON (or references the newly generated classes).
- 2. Initializes the domain objects from external data or user input.
- 3. Invokes aggregator calculations or formula methods.
- 4. (Optionally) loops through time steps if we're simulating a dynamic phenomenon.

Critically, **any domain change** (like new fields, or a new formula) occurs in the JSON, not in this script. The script is stable—it just sets up a run. This aligns with CMCC's goal of letting you alter the "rulebook" or "blueprint" without touching code in every environment.

### 3.3.1 Tying Into the Larger System

In practice, these orchestrator scripts are your "runtime engine" if you're doing a purely software-based simulation. Alternatively, the same JSON definitions might inform a partially physical scenario (e.g., 3D printing

instructions or robotic step sequences), where the aggregator logic remains consistent but the actual runtime is hardware. Either way, the domain logic is *never* duplicated.

### 3.4 Keeping Everything in Sync

Because we have a single JSON file, all downstream artifacts (Python classes, Golang structs, Java code) come from the same source. If you tweak a formula in the JSON, you simply re-run:

#### bash

```
handlebars-codegen -i domain.json -t python.hbs -o domain_generated.py
handlebars-codegen -i domain.json -t golang.hbs -o domain_generated.go
handlebars-codegen -i domain.json -t java.hbs -o DomainGenerated.java
```

Then you use your existing orchestrator code in each language—and everything lines up automatically. No extra DSL, no rewriting the logic in multiple places.

#### 3.4.1 Links to Practical Tools

For real-world adoption, you might use one of the existing open-source pipelines that convert Airtable or Baserow configurations directly into JSON, or do advanced transformations with json-hbars-transform or xml-xlst-transform. We expand on these possibilities in Section 3.5.

### 3.5 Practical Tools and Export Pipelines

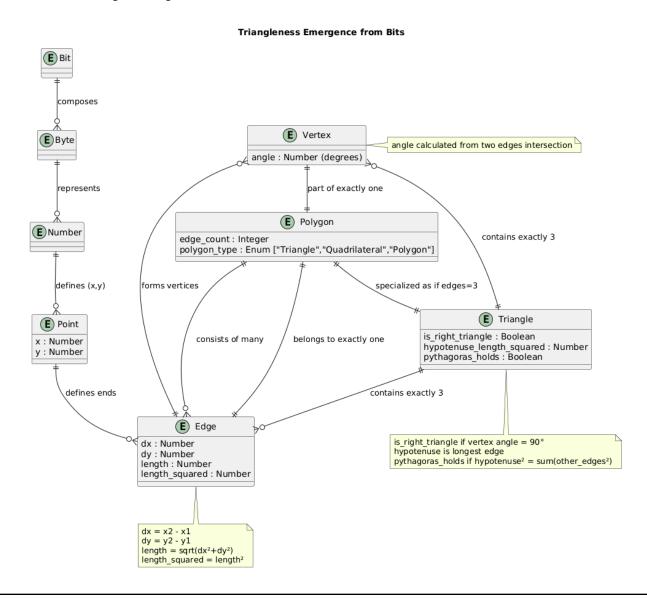
While the examples above show a *handwritten* JSON file, **CMCC** is equally comfortable with user-friendly modeling tools like Airtable or Baserow:

- **Airtable or Baserow**: Business analysts can define tables, formula fields, and relationships in a point-and-click interface. These platforms support snapshot-consistency (or isolation) and can export the entire structure as JSON.
- **NPM CLI Tools**: Simple commands like json-to-xml or json-to-jsd help convert your rulebook into different schemas or cross-checkers.
- **Downstream Code Generation**: Tools like json-hbars-transform apply your chosen Handlebars template to produce typed stubs in any language.
- **Version Control**: Because the JSON is plain text, you can commit your domain model to Git or other VCS, enabling collaborative, auditable evolution of the rulebook.

In other words, a robust open-source ecosystem already exists for **both** capturing the entire domain in a no-code environment **and** exporting that domain as machine-readable JSON for further transformations. Hence, no custom DSL or complex parser is required—just a consistent arrangement of the five CMCC primitives (S, D, L, A, F).

# 4. Case Study 1: Triangleness

This section demonstrates the approach with a *simple geometric domain*. We'll define minimal JSON for polygons, edges, and points. Then we'll generate Python, Golang, and Java helpers, culminating in a short script that checks for right triangles.



# 4.1 JSON Fields for Polygon Edges and Angles

Below is an illustrative snippet of the JSON:

```
{ "name": "max_edge_length", "type": "aggregation",
          "formula": "MAX(edges.length)" },
        { "name": "sum_of_squares", "type": "aggregation",
          "formula": "SUM( POW(edges.length,2) )"
        { "name": "is_triangle", "type": "calculated",
          "formula": "EQUAL(edge_count,3)"
        },
          "name": "is_right_triangle",
          "type": "calculated",
          "formula": "AND(is_triangle, EQUAL(sum_of_squares, POW(max_edge_length,2) * 2less))"
          // This might require referencing the "other two edges" in more detail,
          // but conceptually it's the same idea.
        }
      ]
   },
      "name": "Edge",
      "fields": [
        { "name": "edge_id", "type": "string" },
        { "name": "start_point", "type": "lookup", "references": "Point" },
        { "name": "end_point", "type": "lookup", "references": "Point" },
          "name": "length",
          "type": "calculated",
          "formula": "SQRT( POW(SUBTRACT(end_point.x, start_point.x),2) + POW(SUBTRACT(end_point.y,
start_point.y),2) )"
        }
      ]
   },
      "name": "Point",
     "fields": [
        { "name": "point_id", "type": "string" },
        { "name": "x", "type": "number" },
        { "name": "y", "type": "number" }
      1
   }
}
```

- The Polygon has an edge\_count aggregator, a max\_edge\_length aggregator, etc.
- A calculated field (is\_triangle) checks if edge\_count == 3.
- Another (is\_right\_triangle) tries to implement the Pythagorean check. (You might refine the
  actual formula to handle "largest edge vs. sum of squares of the other two.")

All geometry logic is in the JSON. There is no geometry code in the imperative scripts (aside from reading or writing data).

### 4.2 Showing How "Triangleness" Is an Emergent, Declarative Truth

From a CMCC perspective, "Triangleness" is *not* an imperative routine. Instead, it **emerges** from relationships:

- Schema (S) + Data (D) define polygons that have edges, which have endpoints.
- Aggregations (A) compute how many edges belong to a polygon, or sum up edge lengths squared.
- Lookups (L) link edges to points.
- Calculated Fields (F) compare those aggregator outputs against the 3-edge requirement or the Pythagorean sum.

Hence, the property "triangle" (or "right triangle") *automatically* holds whenever the aggregator logic aligns with the geometry. This stands in contrast to a typical procedural approach, where you'd write code that says "if edge\_count == 3, then do X." Under CMCC, that rule is purely declarative in JSON—the orchestrator or physical system that uses it (see Section 1.6) just interprets these constraints as part of the domain's truth.

### 4.3 Auto-Generated Helpers in Python, Golang, and Java

Using Handlebars (or your preferred engine), we produce classes or structs like:

### Python (domain\_generated.py excerpt):

```
class Polygon:
    def __init__(self, polygon_id, edges):
       self.polygon_id = polygon_id
        self.edges = edges # This might be a list of Edge objects
    @property
    def edge_count(self):
       # aggregator: COUNT(edges)
        return len(self.edges)
    @property
    def max_edge_length(self):
        # aggregator: MAX(edges.length)
        return max(edge.length for edge in self.edges)
    @property
    def sum_of_squares(self):
       # aggregator: SUM( POW(edges.length, 2) )
        return sum(edge.length**2 for edge in self.edges)
    @property
    def is_triangle(self):
       # is_triangle = EQUAL(edge_count, 3)
       return (self.edge_count == 3)
    @property
    def is_right_triangle(self):
       # placeholder for formula
        # let's do a naive check for brevity:
       if not self.is_triangle:
            return False
        sorted_edges = sorted((e.length for e in self.edges))
        return abs(sorted_edges[0]**2 + sorted_edges[1]**2 - sorted_edges[2]**2) < 1e-6
# ... similarly for Edge, Point classes ...
```

(Exact code structure depends on your template. The key is that each property's logic is "pulled in" from the JSON definitions—so if the formula changes in JSON, these stubs adapt automatically at generation time.)

### Golang (domain\_generated.go excerpt):

```
type Polygon struct {
   PolygonID string
   Edges
                 []*Edge
}
func (p *Polygon) EdgeCount() int {
   return len(p.Edges)
}
func (p *Polygon) MaxEdgeLength() float64 {
   maxVal := 0.0
   for _, e := range p.Edges {
       if e.Length() > maxVal {
           maxVal = e.Length()
    }
   return maxVal
}
func (p *Polygon) SumOfSquares() float64 {
    sum := 0.0
   for _, e := range p.Edges {
       length := e.Length()
        sum += (length * length)
    }
    return sum
}
func (p *Polygon) IsTriangle() bool {
   return p.EdgeCount() == 3
}
func (p *Polygon) IsRightTriangle() bool {
    if !p.IsTriangle() {
        return false
   // naive approach
   // ...
   return true
}
Java (DomainGenerated. java excerpt):
public class Polygon {
   private String polygonId;
   private List<Edge> edges;
```

```
public Polygon(String polygonId, List<Edge> edges) {
    this.polygonId = polygonId;
    this.edges = edges;
}
public int getEdgeCount() {
    return edges.size();
}
public double getMaxEdgeLength() {
    return edges.stream().mapToDouble(Edge::getLength).max().orElse(0.0);
}
public double getSumOfSquares() {
    return edges.stream().mapToDouble(e -> Math.pow(e.getLength(),2)).sum();
public boolean isTriangle() {
    return getEdgeCount() == 3;
}
public boolean isRightTriangle() {
    if (!isTriangle()) return false;
    // Implementation details
    // ...
    return true;
}
```

For more on how function names like SQRT or SUBTRACT become Math.sqrt or - in the generated code, see Section 1.7.

# 4.4 Running the Demo: Checking Right Triangles

We then write a **tiny** imperative script in Python, Golang, or Java to demonstrate usage:

#### python

}

```
import domain_generated as dg

# Suppose we manually create Points and Edges:
p1 = dg.Point(point_id="P1", x=0, y=0)
p2 = dg.Point(point_id="P2", x=3, y=0)
p3 = dg.Point(point_id="P3", x=3, y=4)

e1 = dg.Edge(edge_id="E1", start_point=p1, end_point=p2)
e2 = dg.Edge(edge_id="E2", start_point=p2, end_point=p3)
e3 = dg.Edge(edge_id="E3", start_point=p3, end_point=p1)

triangle = dg.Polygon(polygon_id="Triangle1", edges=[e1,e2,e3])

print(f"Is it a triangle? {triangle.is_triangle}")
print(f"Is it a right triangle? {triangle.is_right_triangle}")
```

That's it—no geometry logic is in this script. The logic all came from the auto-generated classes, which in turn came from the JSON definitions.

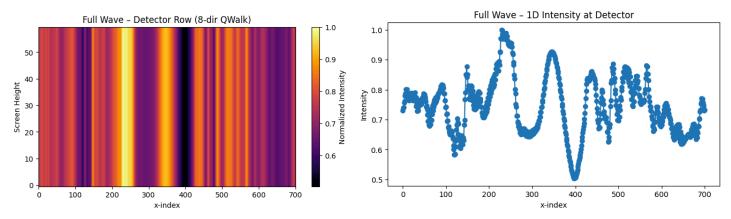
- If we decide to add a new field, say "area" or "circumcircle," we edit the JSON, re-run our Handlebars generator, and the updated classes automatically appear in domain\_generated.py.
- The user script remains the same, or might call the new property if desired.

(For large-scale or concurrent geometry tasks, see Section 7.2 on how the blueprint remains the same even if advanced HPC strategies are used.)

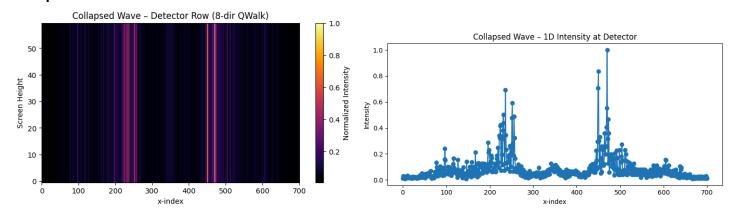
# 5. Case Study 2: Quantum Field Experiment (QFT)

While Triangleness showcases CMCC's declarative modeling for geometry, we now turn to a more **physics-intensive** example: a simplified quantum walk or double-slit experiment. This scenario involves wavefunctions, coin (or spin) operators, and measurements—an ideal test of how far "pure data + aggregator logic" can go.

#### Full Wave Visualization at the Screen & 1d Profile



### Collapsed Wave Visualization at the Screen & 1d Profile



This interference pattern, and collapsed state are purely emergent from the json code below, and don't "solve" the Schrodinger equation at any point in the process. Instead, the quantum interference and collapse outcomes emerge purely from structural transformations defined by CMCC primitives, without explicitly solving Schrödinger's equation. Both interference patterns and measurement-induced collapse are thus natural, predicted consequences within the CMCC framework.

You can find a short primer for non-physicists in Section 5.1, explaining the conceptual background behind wave interference and measurement. Essentially, we illustrate how no specialized quantum DSL is needed; the entire domain can be declaratively expressed in JSON.

### **5.1 QFT Primer for Non-Physicists**

A **quantum walk** can be seen as a discrete analog to how a quantum particle's wavefunction spreads and interferes. At each time step:

- A "coin operator" rotates or mixes spin/directional states.
- Amplitudes propagate (shift) to neighboring grid cells.
- Interference emerges when multiple paths converge with differing phases.

Our aim is *not* to teach quantum mechanics in depth but to show how these rules (wavefunction shape, spin flips, interference) can be *fully captured* in a JSON-based declarative model. Domain experts can verify the coin operator's unitarity or the wavefunction's norm, while non-specialists appreciate that *no special quantum DSL or imperative PDE solver is required*.

### 5.2 JSON Layout for Fields, Wavefunction, Coin Operator

Below is a truncated JSON snippet (inspired by the *From Bits to Qubits* paper) showing three entities: **Grid**, **CoinOperator**, and **Wavefunction**. Each has fields referencing the CMCC primitives (S, D, L, A, F).

```
json
{
  "entities": [
      "name": "Grid",
      "fields": [
        { "name": "grid_id", "type": "string" },
        { "name": "nx", "type": "number" },
        { "name": "ny", "type": "number" },
        { "name": "barrier_y_phys", "type": "number" },
        { "name": "detector_y_phys", "type": "number" },
          "name": "barrier_row",
          "type": "calculated",
          "formula": "FLOOR(DIVIDE(ADD(barrier_y_phys, DIVIDE(Ly,2)), dy))"
        // etc. (slit spacing, dx, dy, etc.)
      1
    },
      "name": "CoinOperator",
      "fields": [
        { "name": "matrix", "type": "tensor", "tensor_shape": "(8,8)" },
          "name": "unitarity_check",
          "type": "calculated",
          "formula": "EQUAL(MULTIPLY(matrix, CONJUGATE_TRANSPOSE(matrix)), IDENTITY(8))"
        },
```

```
{ "name": "seed", "type": "number" }
      1
    },
      "name": "Wavefunction",
      "fields": [
        { "name": "timestep", "type": "number" },
          "name": "psi",
          "type": "tensor",
          "tensor_shape": "(ny,nx,8)"
        },
          "name": "total_norm",
          "type": "aggregation",
          "formula": "SUM( POW( ABS(psi), 2 ) )"
      1
    }
}
```

- **Grid** might store domain geometry: how large the 2D lattice is, where the barrier row is, etc.
- CoinOperator keeps a unitary matrix that flips spin states (like the "coin toss" in a quantum walk). The
  unitarity\_check field is a calculated property that verifies if matrix\*matrix†=I\text{matrix} \times
  \text{matrix}^\dagger = Imatrix\*matrix†=I.
- Wavefunction references a 3D tensor psi with shape (ny,nx,8)(ny, nx, 8)(ny,nx,8) to accommodate multiple spin or direction states. Its total\_norm aggregator sums | ψx,y,spin|2|\psi {x,y,spin}|^2|ψx,y,spin|2 across the entire grid.

Though this is just a high-level snippet, it demonstrates how **every piece** of quantum logic can reside declaratively in JSON: no special quantum DSL or Schrödinger-equation text. You simply define the domain's structure, relationships, and basic aggregator/calculated constraints.

# 5.4 Example Imperative Script for a Quantum Walk/Double-Slit

Imagine we have a minimal script in Python (or another language). For clarity:

### python

```
import domain_generated as dg
import numpy as np

def main():
    # Step 1: Initialize the domain from some config/data
    coin = dg.CoinOperator(matrix=np.eye(8), seed=42)
    wave = dg.Wavefunction(timestep=0, psi=np.zeros((200,200,8), dtype=np.complex64))
    wave.psi[100,100,0] = 1.0 # set initial amplitude at center

# Step 2: Check unitarity
    print("Coin operator is unitary?", coin.unitarity_check)
```

```
# Step 3: Evolve wavefunction in discrete steps
for t in range(50):
    # - apply coin operator spin transform
    # - shift psi to neighbors (like a quantum walk)
    wave.timestep = t
    # pseudo 'update_psi' call
    wave.psi = do_quantum_walk_step(wave.psi, coin.matrix)
    # barrier or slit logic to zero out amplitude in blocked cells
    # ...
    print(f"Timestep {t}, total_norm = {wave.total_norm}")

if __name__ == "__main__":
    main()
```

#### Here:

- The script references the generated CoinOperator and Wavefunction classes (plus aggregator logic).
- The actual "update" method (do\_quantum\_walk\_step) might be partially hand-coded, but it **does not** redefine domain rules—it's simply orchestrating how to apply the matrix and shift amplitudes.
- If you expand your JSON to add more fields (e.g., a "barrier\_mask" or "slit\_mask"), you can regenerate domain\_generated.py or .go or .java—and your script can incorporate those fields with minimal change.

(See Section 1.6 on how the JSON "rulebook" dictates the truth, while the script is just one possible "runtime engine" that obeys it.)

# 5.5 Observing Interference Patterns

In a typical quantum walk or double-slit scenario, you'd see an evolving amplitude distribution that forms interference fringes at the detector row. The key takeaway is that **no quantum DSL or sidecar language** is necessary: everything from grid geometry to unitarity checks is **encoded** as JSON-based aggregator/calc fields. By referencing these fields in your minimal script, you watch interference "emerge" from the data-driven transformations—fully consistent with the CMCC approach.

As the simulation proceeds, you might notice constructive or destructive interference building at certain grid coordinates, eventually forming wave-like patterns. If you implement a measurement step, "collapse" is likewise just a matter of zeroing amplitude outside the measured region—again, no special imperative logic. All these transformations naturally follow from your aggregator fields and formulas (see Section 2.1.1 for how you chain them).

In this manner, quantum behaviors—just like "triangleness"—become a matter of *declarative structure*, not coded instructions. For performance considerations when scaling to large grids or many time steps, refer to Section 7.2. And if you're concerned about dynamic expression security (e.g., passing user-supplied formulas), see Section 7.5.

# 6. "Lambda-to-X" Runtime Libraries

The examples so far assume our JSON formulas are either hard-coded into the generated classes (e.g., def total\_norm()) or replaced with a short snippet of Python or Java. But what if we want to evaluate **any arbitrary expression** from the JSON at runtime—without generating new code for every formula change?

This is where **in-language expression evaluation libraries** come in.

### 6.1 Python Expression Evaluation

Python offers multiple ways to dynamically evaluate strings as expressions, such as eval() or the safer ast.literal\_eval() (though it's limited). There are also more advanced libraries like <a href="mailto:numexpr">numexpr</a> or asteval.

- In principle, your generated code could pass the JSON formula string to numexpr.evaluate(...) at runtime.
- This means you don't have to recompile or regenerate classes if you alter a formula, because the runtime engine will parse and evaluate whatever the JSON says.

However, see Section 7.5 for **important security** considerations. If untrusted users can modify formulas, "eval" can introduce serious risks. Often, generating a statically typed library from the rulebook is more robust.

### 6.2 Golang Expression Libraries

A popular library is <u>govaluate</u>, which parses a string expression (like "SQRT(x\*x + y\*y)") and evaluates it dynamically over variables in a map[string]interface{}.

Another approach is <u>expr</u>, providing a small JIT-like engine for expressions.

• If you embed govaluate or expr calls in your generated code, you can change the JSON formula at runtime, and the library will parse & run it **without** a separate compilation step.

Again, weigh performance and security carefully. For large or mission-critical computations, a design-time approach may offer safer, faster code.

### **6.3 Java Expression Interpreters**

Java has libraries like MVEL or Janino, which can compile or interpret expression strings on the fly.

- With Janino, you can compile a snippet of Java code at runtime to a class, then invoke it.
- MVEL is simpler for quick expression evaluation.

These allow mid-run formula changes in a Java application, but at the cost of dynamic compilation overhead and the need for a safe sandbox (see Section 7.5).

### 6.4 Toward a Truly Dynamic Architecture

By pairing your JSON formulas with **in-language interpreters**, you allow changes even while the system is running—no re-gen or recompile. The trade-off is **potentially less performance** and the **need for sandboxing** if you accept untrusted expressions. Nonetheless, this synergy with dynamic interpreters cements the idea that **the domain logic is not embedded in code**—the code simply executes whatever logic the JSON prescribes.

- 1. CMCC Model: "What is the rule?"
- 2. Runtime Engine: "Let me parse/evaluate that rule string right now."

Hence, you remain syntax-free at the "rule" level, while preserving a pragmatic way to run those rules in your environment of choice. For many production workflows, Section 7.5 argues that generating derived SDKs or libraries at design time is safer, *especially* when formulas originate from external or unverified sources.

# 7. Discussion

### 7.1 Advantages and Possible Caveats

### Advantages:

### 1. No Drift Across Languages

Because every domain rule (e.g., geometry constraints, wavefunction definitions) lives in the JSON, changes propagate automatically to Python, Golang, Java, or any other environment with the same code-generation pipeline.

## 2. Declarative Clarity

The logic (e.g., "A polygon is a triangle if it has exactly three edges," or "unitarity\_check = CoinMatrix × CoinMatrixt - I must be zero") is easy to inspect in the JSON definitions—no hidden assumptions in custom DSL syntax.

### 3. Scalable to Larger Domains

You can add more fields, aggregator logic, or references without needing to create or modify a specialized parser or DSL. Standard JSON-based solutions also plug into countless existing tools for indexing, versioning, or distributed storage.

### 4. Bridging Theory and Practice

The same minimal set of CMCC primitives (S, D, L, A, F) that previously handled pure "paper-level" examples—triangles, quantum phenomena—now appear in real code. This helps unify a conceptual foundation with an actual development workflow.

### **Possible Caveats:**

#### 1. Performance Overheads

- Generating code from large JSON schemas can be slow if you have extremely large or deeply nested models.
- Dynamic expression evaluation libraries (e.g., MVEL, govaluate, or Python's eval) may introduce runtime overhead or security concerns if expressions come from untrusted sources.

### 2. Complexity in Very Large Models

- For massive multi-entity domains (e.g., thousands of tables and relationships), raw JSON can become unwieldy to maintain manually.
- Solutions include splitting the model into multiple JSON files, employing a database-based schema editor, or layering a simple GUI on top.

### 3. Runtime vs. Design-Time

- The difference between "blueprint" (the JSON definitions) and "runtime engine" can create confusion if developers expect the JSON itself to "run."
- Clarify that the JSON is the *rulebook*, and the actual stepping or evolution requires a consistent update cycle (transaction or snapshot logic) in an application or database.

### 7.2 Performance and Scalability Notes

Several database and big-data approaches can optimize the aggregator logic or handle concurrency seamlessly—e.g., Fully scalable relational systems (PostgreSQL, MySQL, Oracle) or distributed DBs with snapshot isolation. If you store these JSON definitions in a table and rely on standard SQL queries for aggregator expressions, you leverage decades of database optimizations for scaling.

Additionally, for HPC (High-Performance Computing) contexts (like large grid-based simulations of wavefunctions), you might store the large data arrays separately (in HDF5, for example) while still referencing them in your JSON-based "schema" definitions. That way, your aggregator fields or formula logic can call into optimized native libraries (NumPy, BLAS, etc.) for the heavy lifting.

(Recall, Section 5.5 illustrates how quantum-level models can be quite large but remain valid under a purely declarative definition.)

# 7.3 Future Directions for Real-Time or Large-Scale Systems

- Real-Time Updates: With incremental or streaming data, you can re-apply aggregator calculations at
  fixed intervals. Using a dynamic expression-evaluation library, you can tweak formulas while the system
  runs.
- Distributed Systems: If multiple nodes share a "master" JSON schema, code generation or formula
  evaluation can be done per node. Ensuring consistent snapshots across distributed nodes calls for
  well-known concurrency protocols (e.g., <u>Raft</u> or <u>Paxos</u>).

(See also Section 7.4 on how schemas can evolve in large or long-lived deployments, and Section 7.5 on securing dynamic evaluations.)

# 7.4 Schema Evolution and Long-Lived Systems

Real-world systems often require updating domain definitions over time—adding new fields, changing formulas, or splitting out large aggregates. Because CMCC treats the **model** as self-describing data (the JSON), these changes can be versioned in Git or a comparable VCS:

 Add or Remove Fields: You can insert or remove schema fields, letting aggregator or calculated properties adjust accordingly.

- 2. **Generate Code**: Re-run the generation step for all target languages, ensuring each environment "follows" the new rules.
- 3. **Agent or Human Check**: Any hand-written references to old fields can be updated by developers or Al-based updaters, ensuring no orphaned references linger.

In database-backed environments (Airtable, Baserow, etc.), these schema changes are typically atomic or transactional, making it trivial to expand your model without breaking old data. The synergy between a no-code platform's consistent snapshots and CMCC's purely declarative definitions simplifies long-lived model evolution.

### 7.5 Security and Deployment (eval vs. Derived SDK)

**Dynamic evaluation** of JSON formulas (as per Section 6) can be powerful but poses potential risks:

- Malicious Input: An attacker could inject OS commands or memory exploits if "eval" is not properly sandboxed.
- Performance Surprises: If a user-provided formula tries to do nested loops or large allocations, you
  may see unexpected CPU or memory use.

### **Recommended Approach:**

- Generate a Domain-Specific Library from the JSON at design time, rather than evaluating expressions on the fly.
- Only allow live formula editing when you trust the users or have a robust sandbox (and still be mindful of performance overhead).
- For extremely large or specialized domains, you might compile the rulebook into a custom HPC or GPU pipeline—still referencing the same aggregator logic from the JSON, but ensuring safe, optimized code.

In short, "eval" is a **last-resort** technique. CMCC's machine-readability and structural approach make it easy to produce statically typed "SDKs" that handle all the aggregator or calculated fields at high speed, with minimal risk.

# 8. Conclusions and Future Work

# 8.1 Key Takeaways for Cross-Language Code Generation

### 1. JSON as Single Source of Truth

By modeling your entire domain (S, D, L, A, F) in a single JSON, you eliminate the need to re-implement domain logic in each target language.

### 2. Minimal "Glue Code"

The handful of imperative lines—loading data, instantiating objects, stepping through timesteps—remains nearly identical across Python, Golang, or Java, with **no embedded domain logic**.

### 3. CMCC in Action

Triangleness (geometry) and quantum walks (wavefunction evolution) demonstrate that even *diverse* domains can be handled under the same method. This approach systematically shows the *practical* dimension of earlier theoretical claims from *From Bits to Qubits with CMCC*.

### 8.2 Next Steps: Expanding to More Domains

- Biology and Systems Modeling: Use aggregator fields to represent gene regulatory networks, with references to enzymes, promoters, and dynamic concentration fields.
- **Real-Time Finance**: Capture trading rules or compliance constraints in JSON aggregator logic, automatically generate code for multiple high-level languages.
- **Al/ML Model Introspection**: Store model hyperparameters and transformations in JSON as aggregator or lambda fields. Generate code for data preprocessing in Python, Java, and Golang.

(And for advanced concurrency or schema evolution in these domains, see Sections 7.3 and 7.4.)

### 8.3 Final Remarks

This paper illustrates one possible route for turning CMCC's structural declarations into actual multi-language runtime code using widely available tools (JSON, Handlebars, expression-evaluation libraries). The approach is far from the only option, but it clearly highlights the power of storing the "what" (domain logic) as data, removing the friction of DSL design or ad hoc duplication across languages.

CMCC remains open for extension, testing, and real-world adoption. We look forward to collaborations that push these ideas into even broader, more complex arenas—and, of course, any sincere falsification attempts that might sharpen or challenge the conjecture's boundaries (see the "Hardcore Falsification Checklist" in Section 2.3).

# 9. References

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## Additional Background and Research on the CMCC Framework

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# **Appendix A: Concrete example in Python**

```
#!/usr/bin/env python3
import numpy as np
import matplotlib.pyplot as plt
# CLASSES
class Grid:
   11 11 11
   Holds geometry info: nx, ny, domain size, barrier row, slit geometry, etc.
   def __init__(self, nx, ny, Lx, Ly, barrier_y_phys, detector_y_phys,
              slit width, slit spacing):
       self.nx = nx
      self.ny = ny
      self.Lx = Lx
      self.Ly = Ly
      self.dx = Lx / nx
       self.dy = Ly / ny
       # Convert physical coords to grid indices
       self.barrier row = int((barrier y phys + Ly/2) / self.dy)
       self.detector row = int((detector y phys + Ly/2) / self.dy)
       # Slit geometry
       center x = nx // 2
       self.slit width = slit width
       self.slit spacing = slit spacing
       self.slit1 xstart = center x - slit spacing // 2
       self.slit1 xend = self.slit1 xstart + slit width
       self.slit2_xstart = center_x + slit_spacing // 2
       self.slit2 xend = self.slit2 xstart + slit width
       # For logging/demonstration
       print(f"Barrier row={self.barrier_row}, Detector row={self.detector_row}")
       print(f"Slit1=({self.slit1 xstart}:{self.slit1 xend}),
Slit2=({self.slit2_xstart}:{self.slit2_xend})")
class CoinOperator:
   Stores the NxN (in this case 8x8) matrix for the local coin operation.
   def init (self, seed=42):
       self.matrix = self._make_coin_8(seed)
   @staticmethod
   def _make_coin_8(seed):
       rng = np.random.default rng(seed=seed)
```

```
mat = np.ones((8,8), dtype=np.complex128)
       alpha = 2.0
        for i in range(8):
           mat[i,i] -= alpha
       rnd = 0.05*(rng.random((8,8)) + 1j*rng.random((8,8)))
       mat += rnd
        # Force unitarity via SVD
       U, s, Vh = np.linalg.svd(mat, full_matrices=True)
       return U @ Vh
   def apply(self, spin in):
       spin in shape=(8,), returns spin out shape=(8,)
       return self.matrix @ spin in
class Wavefunction:
   An immutable snapshot of the wavefunction at a given time:
     psi.shape = (ny, nx, 8)
   We'll have a method evolve one step(...) that returns a NEW Wavefunction.
   DIRECTION OFFSETS = [
        (-1, 0), # up
        (+1, 0), \# down
        (0, -1), # left
        (0, +1), # right
        (-1, -1), # up-left
        (-1, +1), # up-right
        (+1, -1), # down-left
        (+1, +1), # down-right
   ]
    def init (self, grid: Grid, array psi: np.ndarray):
       array psi is shape=(ny,nx,8), complex
        ....
       self.grid = grid
       self.psi = array_psi # store the array as immutable
        # no direct assignment to self.psi[...] from outside
    @classmethod
    def initial condition(cls, grid: Grid):
       Build the wavefunction at t=0, as a Gaussian in y near the bottom,
       wide in x, same amplitude in all directions.
       psi0 = np.zeros((grid.ny, grid.nx, 8), dtype=np.complex128)
       \# let's pick a src y \sim 15% from bottom:
       src y = int(grid.ny * 0.15)
```

```
sigma y = 5.0
    for y in range(grid.ny):
       dy = y - src y
       amp = np.exp(-0.5*(dy/sigma y)**2)
       for d in range(8):
           psi0[y,:,d] = amp
   return cls(grid, psi0)
def evolve one step(self, coin: CoinOperator, measure barrier=False):
   Return a NEW Wavefunction at time t+1, applying:
     1) coin step
     2) shift step
     3) barrier or measure (if measure barrier=True)
   ny, nx, ndir = self.psi.shape
    # 1) Coin step
   psi coin = np.zeros like(self.psi, dtype=np.complex128)
   for y in range(ny):
       for x in range(nx):
           spin in = self.psi[y,x,:] \# shape=(8,)
           spin out = coin.apply(spin in)
           psi_coin[y,x,:] = spin_out
    # 2) Shift step
   psi shift = np.zeros like(psi coin, dtype=np.complex128)
    for d, (ofy, ofx) in enumerate(self.DIRECTION OFFSETS):
        shifted dir = np.roll(psi coin[:,:,d], shift=ofy, axis=0)
       shifted dir = np.roll(shifted dir, shift=ofx, axis=1)
        psi shift[:,:,d] = shifted dir
    # 3) Barrier or measurement
    if measure_barrier:
        # measure collapse barrier logic
       psi_out = self._collapse_barrier(psi_shift)
   else:
       # normal barrier
       psi out = self. apply barrier(psi shift)
   return Wavefunction(self.grid, psi out)
def apply barrier(self, psi in):
   Normal barrier => zero out barrier row except slit columns.
   psi out = psi in.copy()
   br = self.grid.barrier row
   psi_out[br,:,:] = 0
   s1s, s1e = self.grid.slit1 xstart, self.grid.slit1 xend
    s2s, s2e = self.grid.slit2 xstart, self.grid.slit2 xend
```

```
psi out[br, s1s:s1e, :] = psi in[br, s1s:s1e, :]
        psi out[br, s2s:s2e, :] = psi in[br, s2s:s2e, :]
        return psi out
    def collapse_barrier(self, psi_in):
        Collapsing amplitude in barrier row => sum intensities across directions,
        keep only slit columns, sqrt(keep / max), put in direction=0
       psi_out = psi_in.copy()
       ny, nx, ndir = psi in.shape
       br = self.grid.barrier row
        # sum intensities across directions
        row intens = np.sum(np.abs(psi in[br,:,:])**2, axis=-1) # shape=(nx,)
        keep = np.zeros_like(row_intens)
        s1s, s1e = self.grid.slit1 xstart, self.grid.slit1 xend
        s2s, s2e = self.grid.slit2 xstart, self.grid.slit2 xend
        keep[s1s:s1e] = row intens[s1s:s1e]
        keep[s2s:s2e] = row_intens[s2s:s2e]
       m = np.max(keep)
        if m > 1e-30:
           keep /= m
       amps = np.sqrt(keep)
        psi out[br,:,:] = 0
       psi out[br,:,0] = amps # put amplitude in direction=0
        return psi out
    def total norm(self):
       .....
       Returns the sum of |psi|^2 over all y, x, d.
       return np.sum(np.abs(self.psi)**2)
   def detector_row_intensity(self):
        11 11 11
       Summation over directions at 'detector_row', returns shape=(nx,).
       dr = self.grid.detector row
       row amp = self.psi[dr, :, :] # shape=(nx,8)
        row_intens = np.sum(np.abs(row_amp)**2, axis=-1) # shape=(nx,)
       return row intens
class QWalkRunner:
   Orchestrates the layer-by-layer evolution in an immutable, functional style.
   - We keep a list of Wavefunction objects, wave[t].
   - wave[t+1] = wave[t].evolve one step(...)
```

```
We can optionally do a measurement collapse at t=steps to barrier.
   .....
   def init (self, grid: Grid, coin: CoinOperator, steps to barrier, steps after barrier):
       self.grid = grid
       self.coin = coin
       self.steps_to_barrier = steps_to_barrier
       self.steps after barrier = steps after barrier
   def run experiment(self, collapse=False):
       Return the final Wavefunction after steps to barrier + steps after barrier.
       t final = self.steps to barrier + self.steps after barrier
       # We'll keep each wavefunction in a list for demonstration
       wave = [None] * (t final+1)
       wave[0] = Wavefunction.initial condition(self.grid)
       # Evolve up to barrier
       for t in range(self.steps to barrier):
          wave[t+1] = wave[t].evolve one step(self.coin, measure barrier=False)
       # If collapse => measure at t=steps_to_barrier
       if collapse:
          wave[self.steps to barrier] = wave[self.steps to barrier-1].evolve one step(self.coin,
measure barrier=True)
      else:
          # else wave[self.steps to barrier] was already created with measure=False above
       # Evolve remainder
       for t in range(self.steps to barrier, t final):
          wave[t+1] = wave[t].evolve one step(self.coin, measure barrier=False)
       return wave[t final] # final wavefunction
# MAIN
def main():
   # 1) Build the Grid
   nx = 201
   ny = 201
   Lx, Ly = 16.0, 16.0
   steps to barrier = 80
   steps after barrier = 200
   barrier y phys = -2.0
   detector y phys = 5.0
   slit width = 3
   slit_spacing = 12
   grid = Grid(nx, ny, Lx, Ly, barrier y phys, detector y phys,
```

```
slit width, slit spacing)
# 2) Create the Coin
coin = CoinOperator(seed=42)
# 3) Create a QWalkRunner
runner = QWalkRunner(grid, coin, steps to barrier, steps after barrier)
# 4) Run the "FULL" wave (no measurement)
print("Running FULL wave (no barrier measurement)...")
wave full = runner.run experiment(collapse=False)
norm_full = wave_full.total_norm()
print(f"Final norm (full)={norm full:.3g}")
# 5) Run the "COLLAPSED" wave (with measurement)
print("Running COLLAPSED wave (with barrier measurement)...")
wave_coll = runner.run_experiment(collapse=True)
norm coll = wave coll.total norm()
print(f"Final norm (collapsed) = {norm coll:.3g}")
# 6) Measure intensity at the detector row => sum over directions => shape=(nx,)
int full = wave full.detector row intensity()
int_coll = wave_coll.detector_row_intensity()
# 7) Normalize each
mf = np.max(int full)
if mf > 1e-30:
   int full /= mf
mc = np.max(int coll)
if mc > 1e-30:
    int coll /= mc
# 8) Build 2D "screen" => tile 1D intensity
screen height = 60
screen_full = np.tile(int_full, (screen_height,1))
screen coll = np.tile(int coll, (screen height,1))
# 9) Plot
plt.figure(figsize=(8,4))
plt.imshow(screen full, origin="lower", aspect="auto", cmap="inferno")
plt.title("Full Wave - Detector Row (OO, Functional layering)")
plt.xlabel("x-index")
plt.ylabel("Screen Height")
plt.colorbar(label="Normalized Intensity")
plt.figure(figsize=(8,4))
plt.plot(int full, 'o-')
plt.title("Full Wave - 1D Intensity at Detector")
plt.xlabel("x-index")
plt.ylabel("Intensity")
plt.figure(figsize=(8,4))
```

```
plt.imshow(screen_coll, origin="lower", aspect="auto", cmap="inferno")
plt.title("Collapsed Wave - Detector Row (OO, Functional layering)")
plt.xlabel("x-index")
plt.ylabel("Screen Height")
plt.colorbar(label="Normalized Intensity")

plt.figure(figsize=(8,4))
plt.plot(int_coll, 'o-')
plt.title("Collapsed Wave - 1D Intensity at Detector")
plt.xlabel("x-index")
plt.ylabel("Intensity")

plt.tight_layout()
plt.show()
```