LANGUAGE ENGINEERING WITH THE GEMOC STUDIO

Tutorial @ MODELS 2017, September 18th, 2017

Benoit Combemale (Univ. Toulouse)

http://www.combemale.fr benoit.combemale@irit.fr @bcombemale

Julien Deantoni (Univ. Nice-Sophia-Antipolis)

http://www.i3s.unice.fr/~deantoni/julien.deantoni@polytech.unice.fr



Add concurrency in your semantics!

https://github.com/gemoc/MODELS2017Tutorial



Reifying Concurrency in xDSML: Limitations

- Concurrency remains implicit and ad-hoc in language design and implementation:
 - Design: implicitly inherited from the meta-language used
 - Implementation: mostly embedded in the underlying execution environment
- The lack of an explicit concurrency specification in language design prevents:
 - leveraging the concurrency concern of a particular domain or platform
 - a complete understanding of the behavioral semantics
 - effective concurrency-aware analysis techniques
 - effective techniques for producing semantic variants
 - analysis of the deployment on parallel architectures



Reifying Concurrency in xDSML: Limitations

```
@Main
def public void main() {
  for(FSM sm : _self.ownedFsms){
    if (! sm.inputBuffer.isEmpty){
        sm.run()
        anFSMRan = true
    }
}
```

```
    System
    FSM fsm1
    FSM fsm2
    FSM fsm3
    FSM fsm0
```

Iterate over the collection

Scheduling order depends on the "ownedFsms" collection

Copy and paste changes the order and eventually the execution



Reifying Concurrency in xDSML: Proposition

```
→ System
→ FSM fsm1
→ FSM fsm2
→ FSM fsm3
```

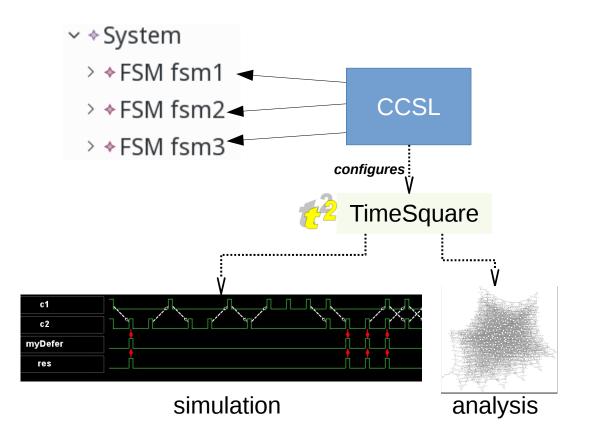
```
@Main
def public void main() {
  for(FSM sm : _self.ownedFsms){
    if (! sm.inputBuffer.isEmpty){
        sm.rup()
        anFsMRan = true
    }
}
```



We can use *Logical Time* to specify the partial ordering that are correct according to our semantics, in a platform independent point of view



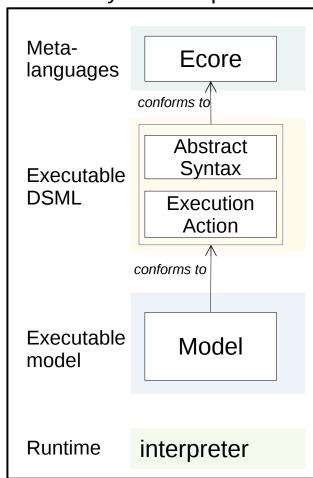
Reifying Concurrency in xDSML: Proposition



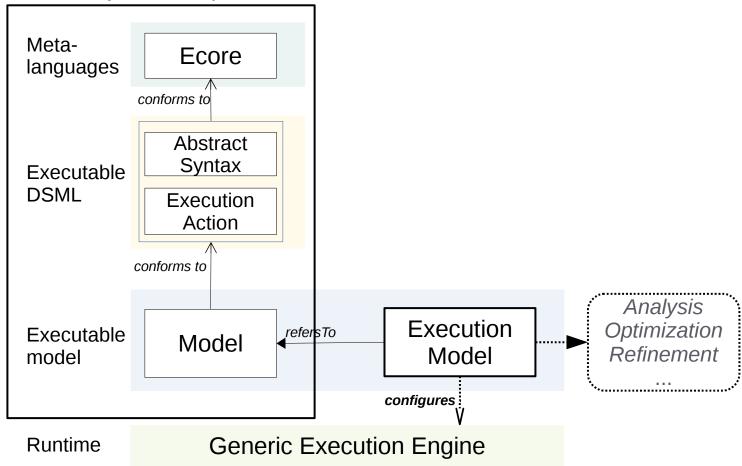


We can use *Logical Time* to specify the partial ordering that are correct according to our semantics, in a platform independent point of view

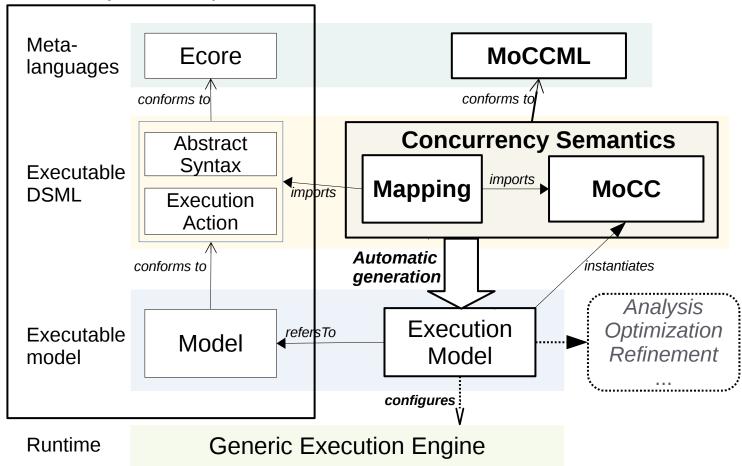




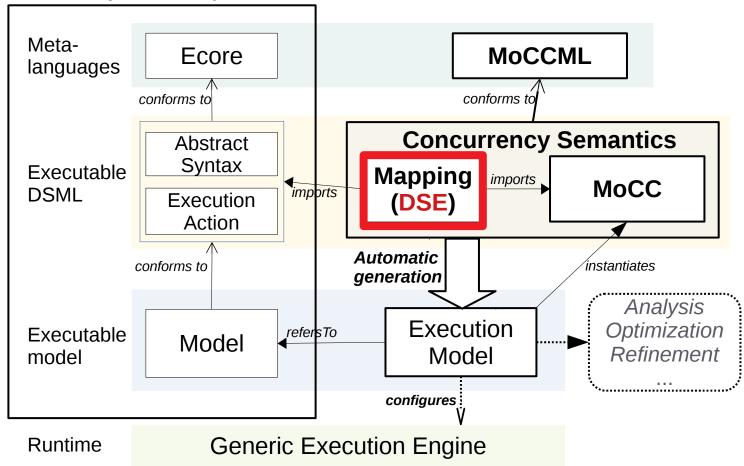




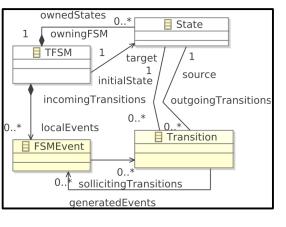


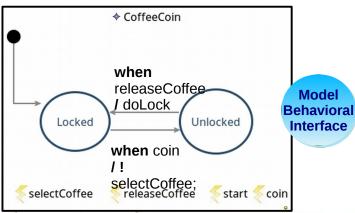




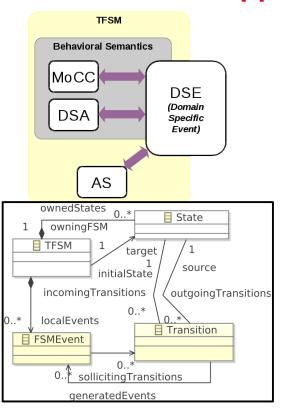


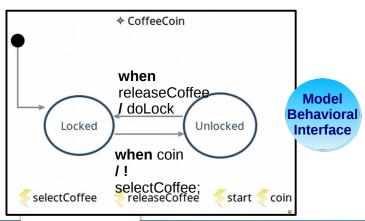




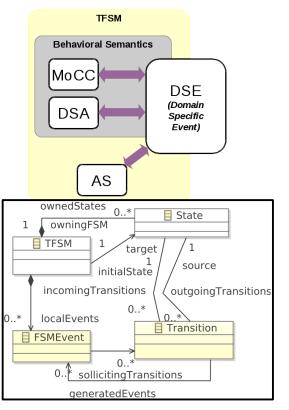


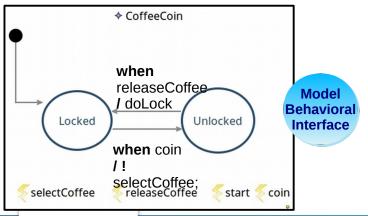








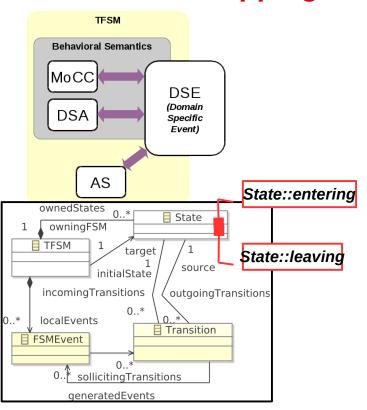


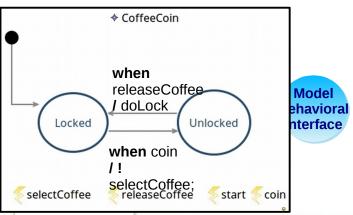




This is mainly the OCL syntax which is used...

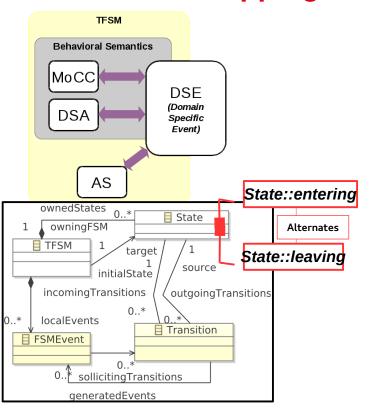


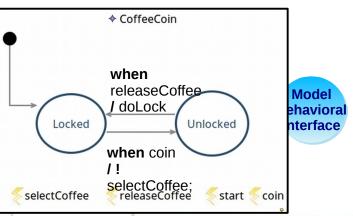




```
import 'http://org.gemoc.models17.fsm.xfsm/model/'
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
package model
 * DSE - MoCCML mapping
context State
    def: entering : Event = self
    def: leaving : Event = self
endpackage
```

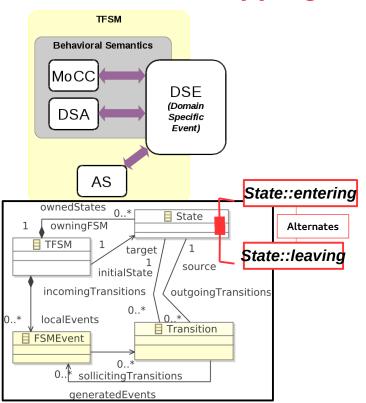


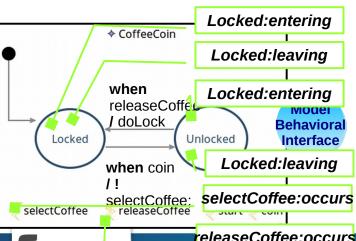




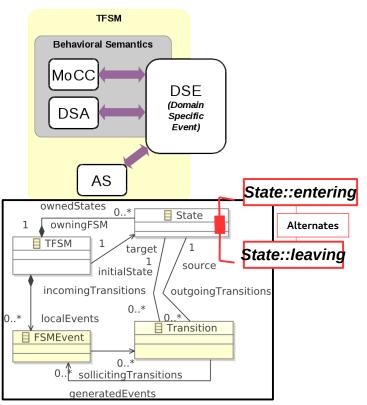
```
import 'http://org.gemoc.models17.fsm.xfsm/model/'
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
package model
 * DSE - MoCCML mapping
context State
    def: entering : Event = self
    def: leaving : Event = self
/**
 * Constraints
context State
    inv enterThenLeaveNoReEntrance:
        Relation Alternates(self.entering, leaving)
endpackage
```



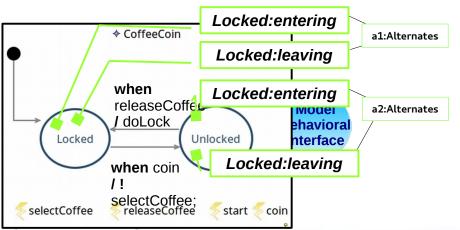




```
import 'http://org.gemoc.models17.fsm.xfsm/model/'
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
package model
 * DSE - MoCCML mapping
context State
    def: entering : Event = self
    def: leaving : Event = self
/**
 * Constraints
context State
    inv enterThenLeaveNoReEntrance:
        Relation Alternates(self.entering, leaving)
endpackage
```

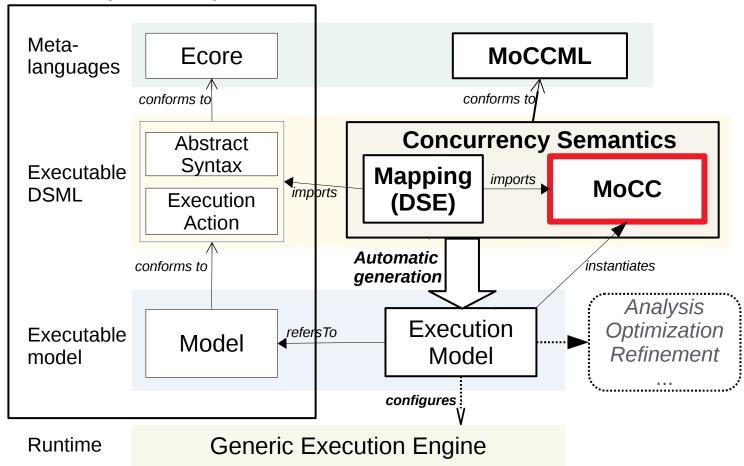


```
import 'http://org.gemoc.models17.fsm.xfsm/model/'
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
ECLimport "platform:/plugin/fr.inria.aoste.timesquare.ccslkernel
package model
 * DSE - MoCCML mapping
context State
    def: entering : Event = self
    def: leaving : Event = self
/**
 * Constraints
context State
    inv enterThenLeaveNoReEntrance:
        Relation Alternates(self.entering, leaving)
endpackage
```



In green, this is the generated execution model, which is a symbolic formal representation of the scheduling state space!







MoCCML Domain Specific Constraints

MoCCML allows for the definition of possibly complex constraints related to a specific domain :

- PeriodicTask, SporadicTask, Deadline in real time,
- Fork, Join, ActivityTriggering in fUML
- Specific Communication Protocol in Distributed Systems
 It can be defined either
- 1°) based on a conjunction of primitive Event Constraints as defined in the CCSL language (e.g., Precedes, Coincides, Excludes, DelayedFor, Alternates)
- 2°) based on a MoCCML automata, i.e., a constraint automata



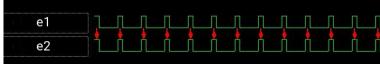
MoccML Domain Specific Constraints

1°) based on a conjunction of primitive Event Constraints as defined in the CCSL language (e.g., Precedes, Coincides, Excludes, DelayedFor, Alternates)

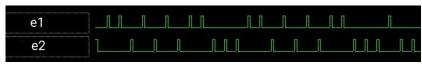
Precedes: e1 precedes e2 means that the ith occurrence of e1 arrives before the ith of e2



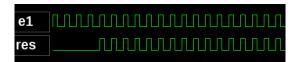
Coincides: **e1 coincides with e2** means that the ith occurrence of e1 arrives synchronously with the ith of e2



Excludes: **e1 excludes** with **e2** means that none of the occurrences of e1 arrives synchronously with one of e2



DelayedFor: res = e1 DelayedFor N on e2 means that for each occurrences of e1 between two occurrences of e2, there is an occurrence of res after N occurrences of e2. A special case is res = e1 delayedFor N on e1. In this case the N first occurrences of e1 are removed in res.

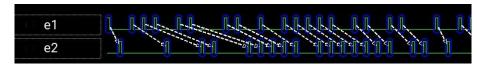




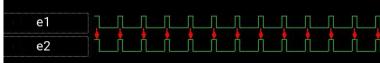
MoccML Domain Specific Constraints

1°) based on a conjunction of primitive Event Constraints as defined in the CCSL language (e.g., Precedes, Coincides, Excludes, DelayedFor, Alternates)

Precedes: e1 precedes e2 means that the ith occurrence of e1 arrives before the ith of e2



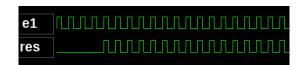
Coincides: **e1 coincides with e2** means that the ith occurrence of e1 arrives synchronously with the ith of e2



Excludes: **e1 excludes** with **e2** means that none of the occurrences of e1 arrives synchronously with one of e2



DelayedFor: res = e1 DelayedFor N on e2 means that for each occurrences of e1 between two occurrences of e2, there is an occurrence of res after N occurrences of e2. A special case is res = e1 delayedFor N on e1. In this case the N first occurrences of e1 are removed in res.







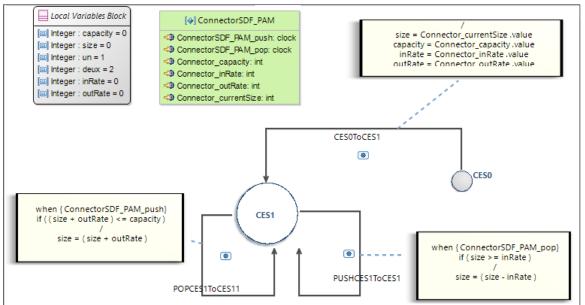
MoCCML language

MoCCML automata intuitive semantics

- Clocks as transition triggers
- Condition expressions on integer variables
- Assignment and arithmetic operators on variables

Operational Semantics of the Model of Concurrency and Communication Language :

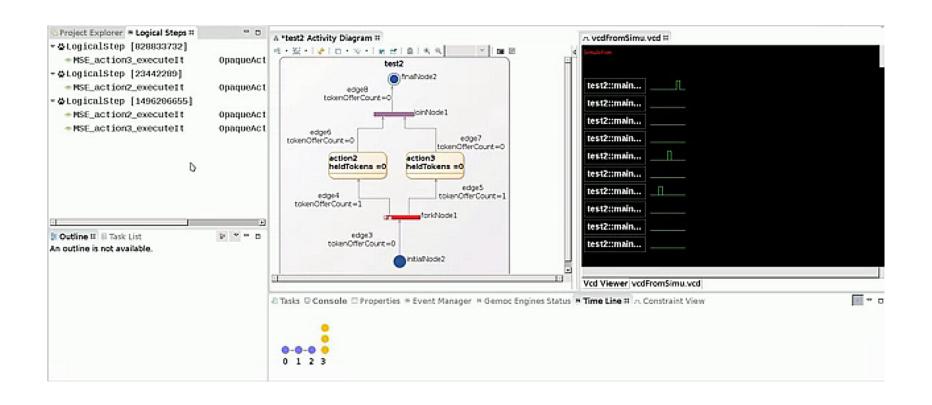
https://hal.inria.fr/hal-01060601v2





Activity Diagram Debugger





Benoit Combemale, Julien Deantoni, Matias Vara Larsen, Frédéric Mallet, Olivier Barais, Benoit Baudry, Robert France, "Reifying Concurrency for Executable Metamodeling" In Software Language Engineering (SLE), 2013







Program outline (see README)

Language workbench: create your xFSM DSL

- Changing the language specification from sequential to concurrent
- Creation of the DSE and MoCCML mapping project
- Modification of the ECL file
- Create Domain Specific MoCC constraints



Wrap-up and discussion

- Today you have
 - executed models of the communicating FSM language
 - Completed the operational semantics by using Kermeta
 - extended its concrete graphical syntax for animation using Sirius
 - Added a explicit and formal concurrency model to your semantics by using MoCCML
 - Saw hox to coordinate heterogeneous executable modeling languages to support concurrent model execution by using BCOoL
 - Had fun ?

